

1) Grammar:

$$\begin{aligned} \langle S \rangle &\rightarrow \langle A \rangle a \langle B \rangle b \\ \langle A \rangle &\rightarrow \langle A \rangle b \mid b \\ \langle B \rangle &\rightarrow a \langle B \rangle \mid a \end{aligned}$$
a) baab

$$\begin{aligned} \langle S \rangle &\rightarrow \langle A \rangle a \langle B \rangle b \\ &\rightarrow ba \langle B \rangle b \\ &\rightarrow baab \end{aligned}$$

$$\begin{aligned} \langle A \rangle &\rightarrow b \\ \langle B \rangle &\rightarrow a \end{aligned}$$
 $\therefore$  baab is valid!b) bbbab

$$\begin{aligned} \langle S \rangle &\rightarrow \langle A \rangle a \langle B \rangle b \\ &\rightarrow \langle A \rangle ba \langle B \rangle b \\ &\rightarrow \langle A \rangle bba \langle B \rangle b \\ &\rightarrow bbaa \langle B \rangle b \end{aligned}$$

$$\begin{aligned} \langle A \rangle &\rightarrow \langle A \rangle b \\ \langle A \rangle &\rightarrow \langle A \rangle b \\ \langle A \rangle &\rightarrow b \end{aligned}$$

$$\therefore \text{bbbab is not generated by the grammar, so is } \underline{\text{invalid!}}$$
c) bbbaaaaa

$$\begin{aligned} \langle S \rangle &\rightarrow \langle A \rangle a \langle B \rangle b \\ &\rightarrow \langle A \rangle ba \langle B \rangle b \\ &\rightarrow bba \langle B \rangle b \\ &\rightarrow bbaa \langle B \rangle b \\ &\rightarrow bbaaa \langle B \rangle b \\ &\rightarrow bbaaaa \langle B \rangle b \\ &\rightarrow bbaaaaa \langle B \rangle b \\ &\rightarrow bbaaaaaa \langle B \rangle b \end{aligned}$$

$$\begin{aligned} \langle A \rangle &\rightarrow \langle A \rangle b \\ \langle A \rangle &\rightarrow b \\ \langle B \rangle &\rightarrow a \langle B \rangle \\ \langle B \rangle &\rightarrow a \langle B \rangle \\ \langle B \rangle &\rightarrow a \langle B \rangle \\ \langle B \rangle &\rightarrow a \langle B \rangle \\ \langle B \rangle &\rightarrow a \end{aligned}$$

$$\therefore \text{sentence is } \underline{\text{invalid}}$$
d) bbaab

$$\begin{aligned} \langle S \rangle &\rightarrow \langle A \rangle a \langle B \rangle b \\ &\rightarrow \langle A \rangle ba \langle B \rangle b \\ &\rightarrow bba \langle B \rangle b \\ &\rightarrow bbaab \end{aligned}$$

$$\begin{aligned} \langle A \rangle &\rightarrow \langle A \rangle b \\ \langle A \rangle &\rightarrow b \\ \langle B \rangle &\rightarrow b \end{aligned}$$

$$\therefore \text{bbaab is derived from the grammar and is therefore } \underline{\text{valid!}}$$

2.) Grammar

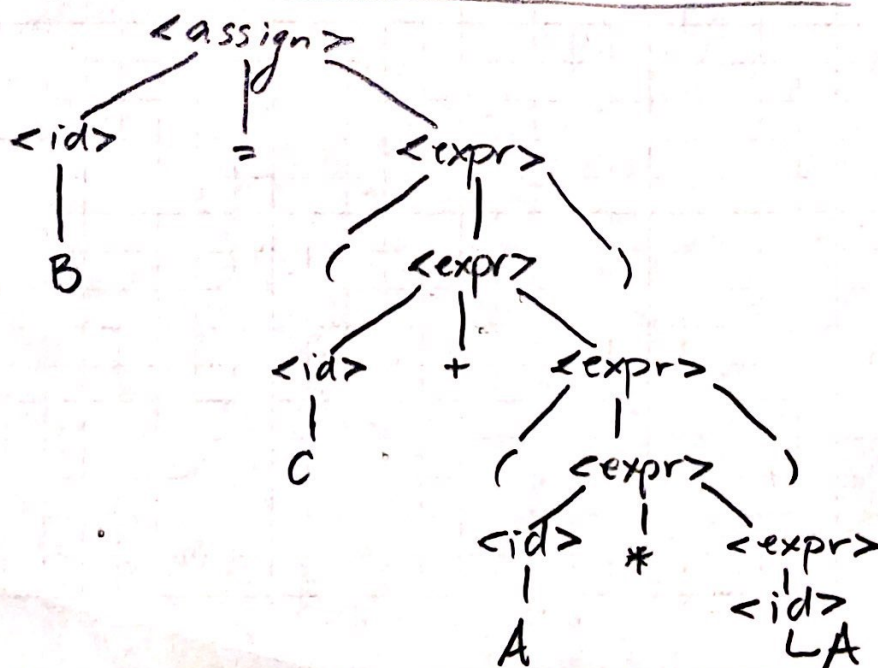
$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$   
 $\langle \text{id} \rangle \rightarrow A | B | C$   
 $\langle \text{expr} \rangle \rightarrow \langle \text{id} \rangle \pm \langle \text{expr} \rangle$   
 $| \langle \text{id} \rangle * \langle \text{expr} \rangle$   
 $| ( \langle \text{expr} \rangle )$   
 $| \langle \text{id} \rangle$

Tokens	Lexemes
variables	A, B, C
Arith. ops	+, *
Assign. ops	=
Special chars	(, )

3.) Left-most derivation for:  $B = B + (C + (A * A))$

$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$   
 $\rightarrow B = \langle \text{expr} \rangle$   
 $\rightarrow B = \langle \text{id} \rangle + \langle \text{expr} \rangle$   
 $\rightarrow B = B + \langle \text{expr} \rangle$   
 $\rightarrow B = B + ( \langle \text{expr} \rangle )$   
 $\rightarrow B = B + ( \langle \text{id} \rangle + \langle \text{expr} \rangle )$   
 $\rightarrow B = B + ( C + \langle \text{expr} \rangle )$   
 $\rightarrow B = B + ( C + ( \langle \text{expr} \rangle ) )$   
 $\rightarrow B = B + ( C + ( \langle \text{id} \rangle * \langle \text{expr} \rangle ) )$   
 $\rightarrow B = B + ( C + ( A * \langle \text{expr} \rangle ) )$   
 $\rightarrow B = B + ( C + ( A * \langle \text{id} \rangle ) )$   
 $\rightarrow B = B + ( C + ( A * A ) ) \checkmark$

Parse Tree





4.) Grammar

$S \rightarrow Aa | Bb$   
 $A \rightarrow Aa | AbC | C$   
 $B \rightarrow S | bb$   
 $C \rightarrow c$

By removing the left recursion!

$$A = A\alpha / B \Rightarrow A = BA'$$

$$A' = \alpha A' / \epsilon$$

Grammar after removing the recursion!

$S \rightarrow Aa | Bb$   
 $A \rightarrow CA'$   
 $A' \rightarrow aA' | bCA' | \epsilon$   
 $B \rightarrow S | bb$   
 $C \rightarrow c$

5.) Grammar

$A \rightarrow aBc | ac | a$   
 $B \rightarrow b | aB$

$\therefore 'a'$  is common thru out grammar...

$A \rightarrow aA'$   
 $A' \rightarrow Bc | C | \epsilon$

So... using left factoring for grammar results in:

$A \rightarrow aA'$   
 $A' \rightarrow Bc | C | \epsilon$   
 $B \rightarrow b | aB$

6.) Grammar

$E \rightarrow E + T | E * T | T$   
 $T \rightarrow (E) | id$

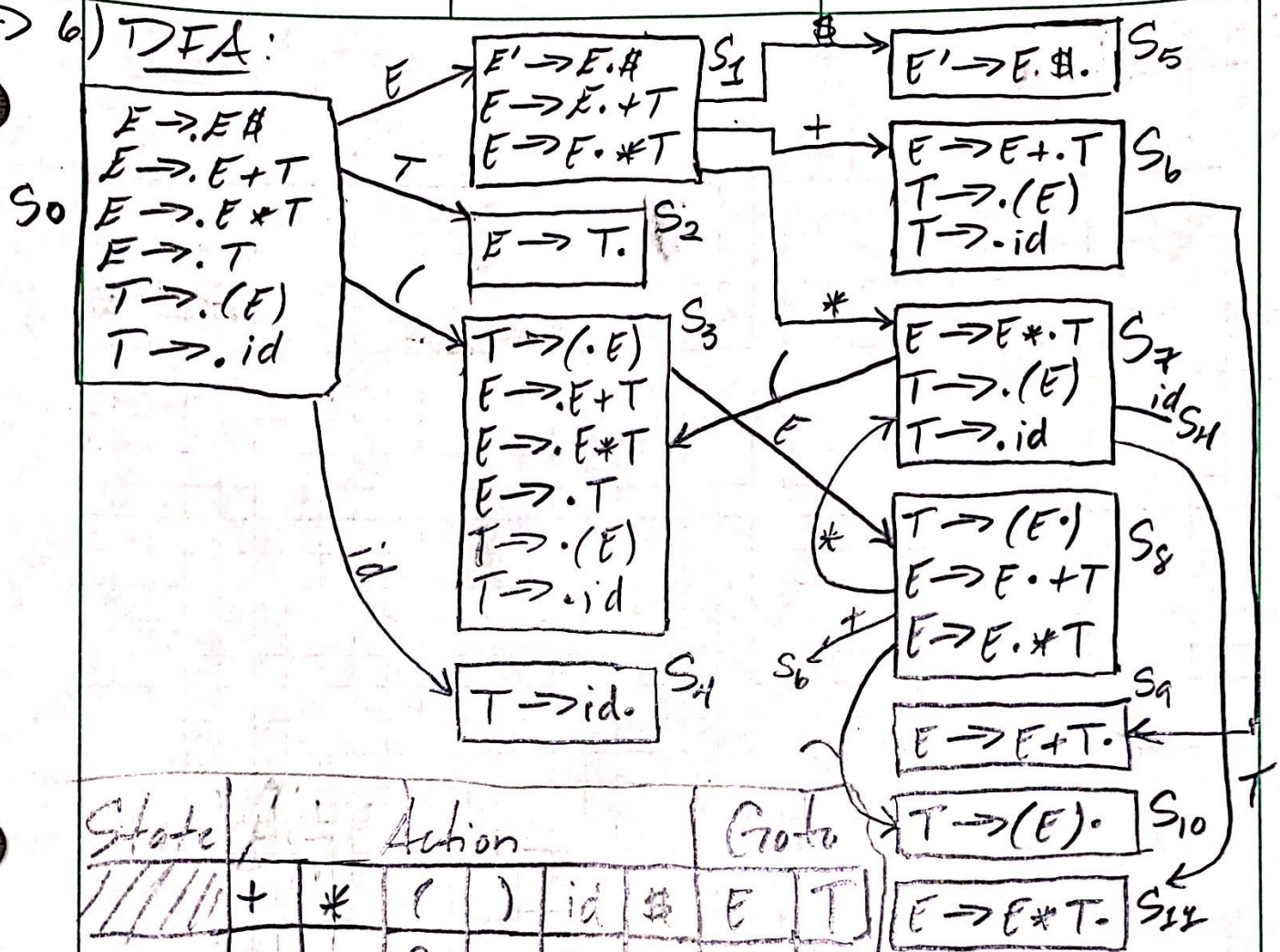
① Closure:  $E' \rightarrow E$

$E \rightarrow E + T$   
 $E \rightarrow E * T$   
 $E \rightarrow T$   
 $T \rightarrow (E)$   
 $T \rightarrow id$

(continued on next page)



6) DFA:



State	Action						Goto	
	+	*	(	)	id	\$	E	T
0			S <sub>3</sub>		S <sub>4</sub>		1	2
1	S <sub>6</sub>	S <sub>7</sub>				S <sub>5</sub>		
2	R <sub>3</sub>	R <sub>3</sub>	R <sub>3</sub>	R <sub>3</sub>	R <sub>3</sub>	R <sub>3</sub>		
3					S <sub>4</sub>			
4	R <sub>5</sub>	R <sub>5</sub>	R <sub>5</sub>	R <sub>5</sub>	R <sub>5</sub>	R <sub>5</sub>		
5						Acc		
6			S <sub>3</sub>		S <sub>4</sub>			
7			S <sub>3</sub>		S <sub>4</sub>			
8	S <sub>6</sub>	S <sub>7</sub>		S <sub>10</sub>				
9	R <sub>1</sub>	R <sub>1</sub>	R <sub>1</sub>	R <sub>1</sub>	R <sub>1</sub>	R <sub>1</sub>		
10	R <sub>4</sub>	R <sub>4</sub>	R <sub>4</sub>	R <sub>4</sub>	R <sub>4</sub>	R <sub>4</sub>		
11	R <sub>2</sub>	R <sub>2</sub>	R <sub>2</sub>	R <sub>2</sub>	R <sub>2</sub>	R <sub>2</sub>		