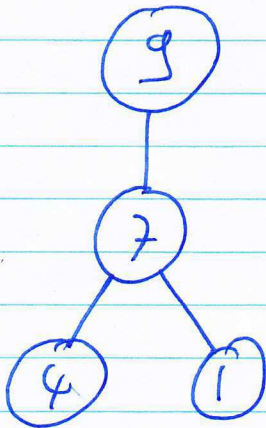
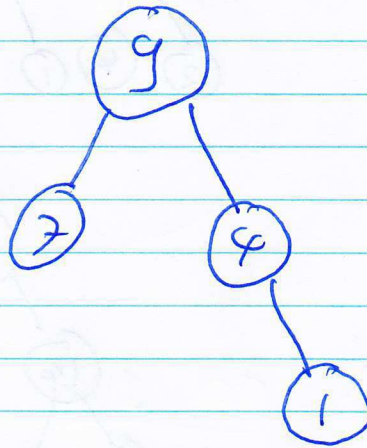
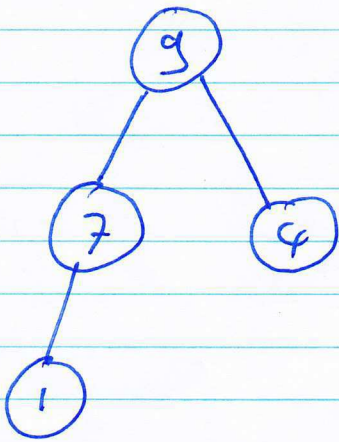
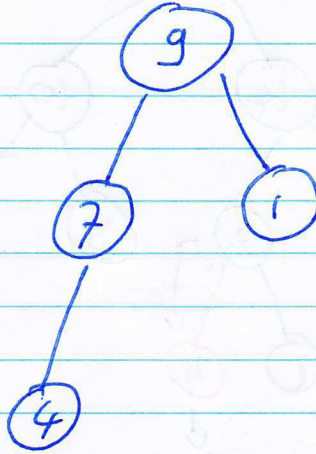
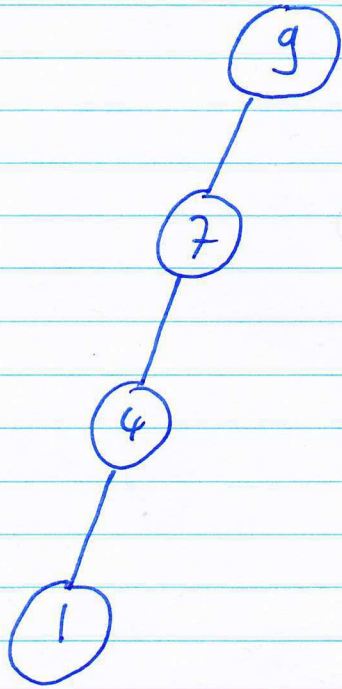


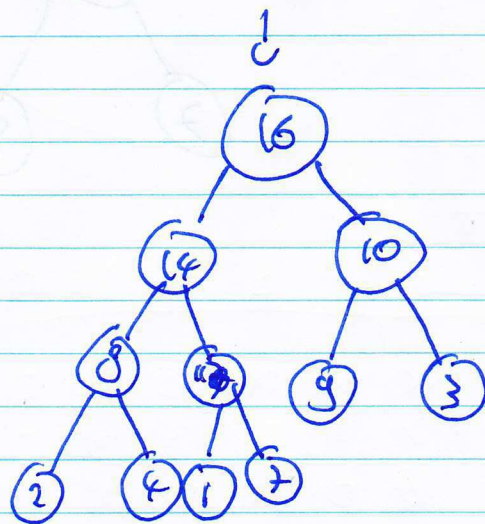
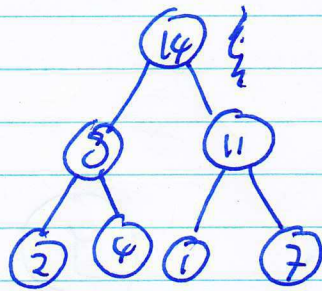
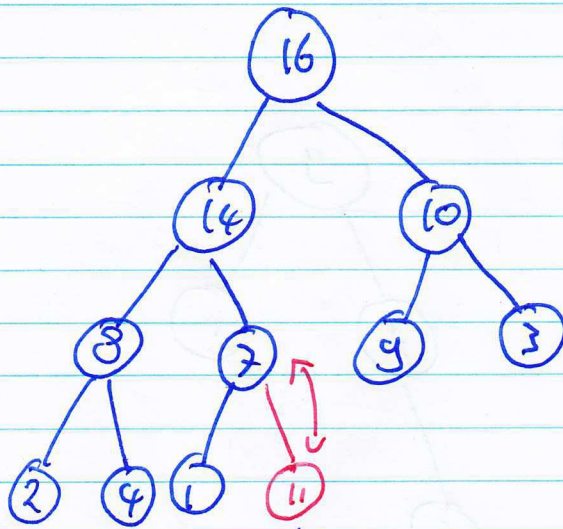
IV

1.

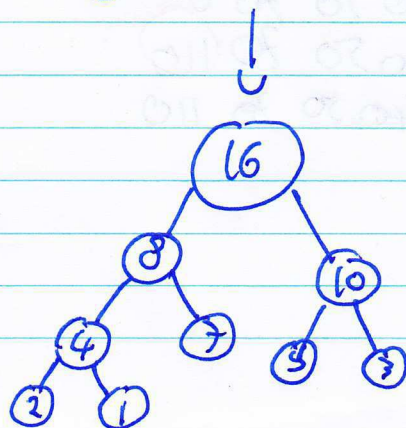
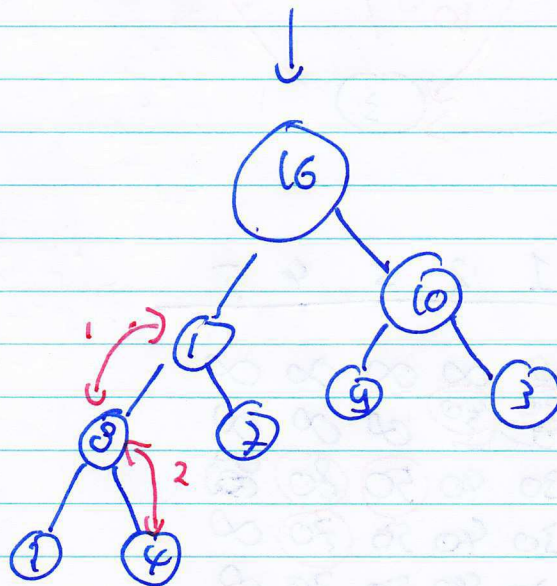
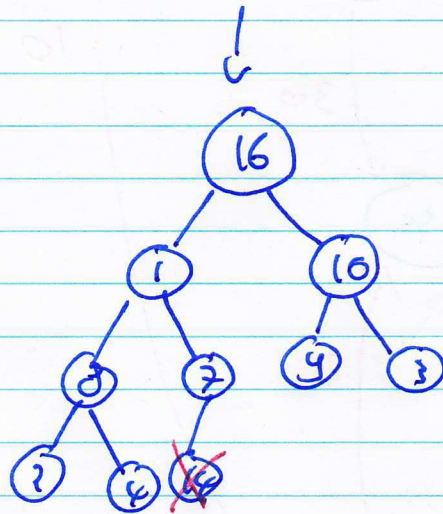
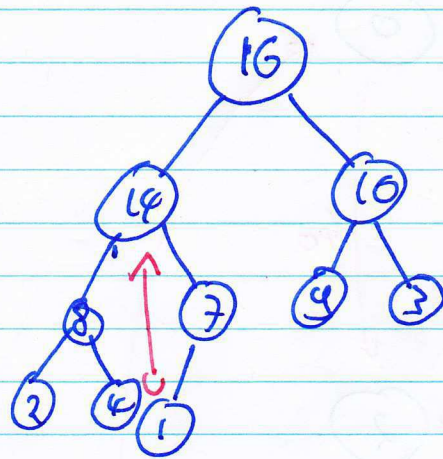


2

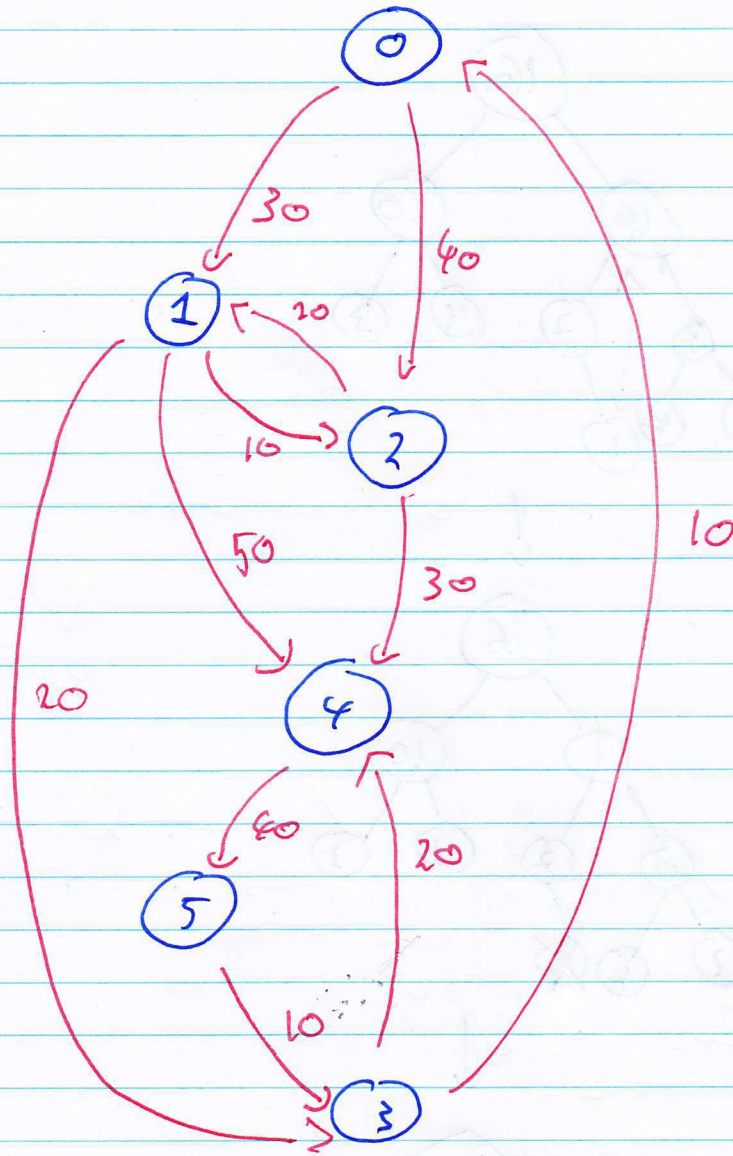
i



ii



3.



Q: 0 1 2 3 4 5

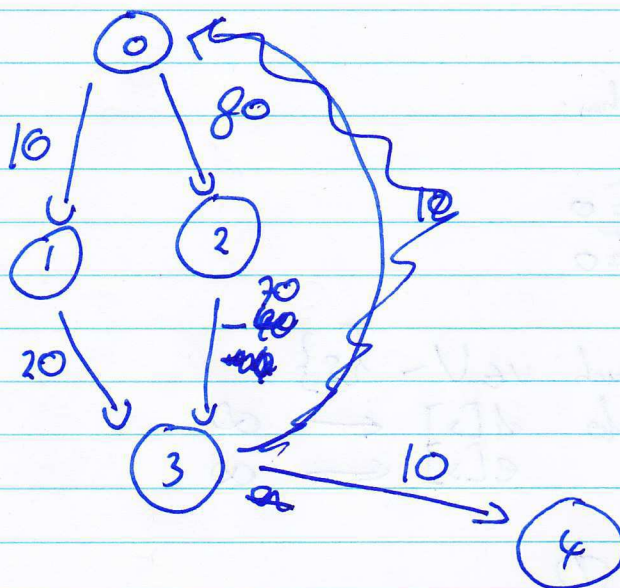
0	∞	∞	∞	∞	∞
0	30	40	∞	∞	∞
0	30	40	50	80	110
0	30	40	50	70	∞
0	30	40	50	70	∞
0	30	40	50	70	110
0	30	40	50	70	110

~~S: {0}~~
 S: {0}
 S: {0, 1}
 S: {0, 1, 2}
 S: {0, 1, 2, 3}
 S: {0, 1, 2, 3, 4}
 S: {0, 1, 2, 3, 4, 5}

(x) = update

4.

Graph



Q:

	0	1	2	3	4
0	0				
1	10	0	∞	∞	∞
2	80	80	0	∞	∞
3	10	80	30	0	∞
4	10	80	30	40	0

 $S: \{0\}$ $S: \{0, 1\}$ $S: \{0, 1, 2\}$ $S: \{0, 1, 3, 4\}$ $S: \{0, 1, 3, 4, 2\}$ $S: \{0, 1, 3, 4, 2\}$

We see that the update with the length of $0 \rightsquigarrow 3$ updates to a lower value but the path $0 \rightsquigarrow 4$ does not update because it is already in S

5. For this problem, we also Dijkstra's algorithm a little:

Algorithm:

$d[s] \leftarrow 0$
 $e[s] \leftarrow 0$

for each $v \in V - \{s\}$
do $d[v] \leftarrow \infty$
 $e[v] \leftarrow \infty$

$S \leftarrow \emptyset$
 $Q \leftarrow V$

while $Q \neq \emptyset$
do $u \leftarrow \text{Extract_Min}(Q)$
 $S \leftarrow S \cup \{u\}$
for each $v \in \text{Adj}[u]$
do if $d[v] > d[u] + w(u, v)$
then $d[v] \leftarrow d[u] + w(u, v)$
 $e[v] \leftarrow e[u] + 1$
else if $d[v] == d[u] + w(u, v)$
then if $e[u] + 1 < e[v]$
then $d[v] \leftarrow d[u] + w(u, v)$
 $e[v] \leftarrow e[u] + 1$