# Beweren en Bewijzen Leertaak 9

#### 28 april 2017

### Opgave 1

a) Mijn studentnummer is dus ik maak de even afleidingsbomen.

b) (2) Stelling:  $P \to Q, Q \to R, P \vdash R$ 

Afleidingsboom:

$$\frac{P \rightarrow Q, Q \rightarrow R, P \vdash Q \rightarrow R}{P \rightarrow Q, Q \rightarrow R, P \vdash P \rightarrow Q} \frac{hyp}{P \rightarrow Q, Q \rightarrow R, P \vdash P} \frac{hyp}{P \rightarrow Q, Q \rightarrow R, P \vdash P} \rightarrow E}{P \rightarrow Q, Q \rightarrow R, P \vdash R} \rightarrow E$$

(4) Stelling:  $P \to (Q \to R) \vdash (P \land Q) \to R$ 

Afkorting(en):  $\Sigma = P \to (Q \to R)$ 

Afleidingsboom:

$$\frac{\sum, (P \land Q) \vdash P \to (Q \to R)}{\sum, (P \land Q) \vdash P \to Q} \overset{hyp}{\wedge E1} \xrightarrow{\sum, (P \land Q) \vdash P \to Q} \overset{hyp}{\wedge E1} \xrightarrow{\sum, (P \land Q) \vdash P \to Q} \overset{hyp}{\wedge E2} \xrightarrow{\sum, (P \land Q) \vdash Q \to R} \xrightarrow{\sum, (P \land Q) \vdash R} \xrightarrow{} I$$

(6) Stelling:  $P \wedge (Q \vee R) \vdash (P \wedge Q) \vee (P \wedge R)$ 

Afkorting(en):  $\Sigma = P \wedge (Q \vee R)$ ) Afleidingsboom:

$$\frac{\sum . Q \vdash P \land (Q \lor R)}{\sum . Q \vdash P \land (Q \lor R)} \stackrel{hyp}{\land E1} \frac{\sum . Q \vdash P \land (Q \lor R)}{\sum . Q \vdash Q} \stackrel{hyp}{\land I} \frac{\sum . Q \vdash P \land (Q \lor R)}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\land I} \frac{\sum . Q \vdash P \land (Q \lor R)}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\land I} \frac{\sum . Q \vdash P \land (Q \lor R)}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\land I} \frac{\sum . Q \vdash P \land (Q \lor R)}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\land I} \frac{\sum . Q \vdash P \land (Q \lor R)}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\land I} \frac{\sum . Q \vdash P \land (Q \lor R)}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\land I} \frac{\sum . Q \vdash P \land (Q \lor R)}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\land I} \frac{\sum . Q \vdash P \land (Q \lor R)}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\land I} \frac{\sum . Q \vdash P \land (Q \lor R)}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\land I} \frac{\sum . Q \vdash P \land (Q \lor R)}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\land I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\land I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash P \land Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash Q} \stackrel{hyp}{\lor I} \frac{\sum . Q \vdash P \land Q}{\sum . Q \vdash Q} \stackrel{hyp}{\lor I} \stackrel$$

(8) Stelling:  $P \to R, Q \to R \quad \vdash \quad (P \lor Q) \to R$ 

Afkorting(en):  $\Sigma = P \to R, Q \to R$ )

Afleidingsboom:

$$\frac{\sum_{...}P\vee Q\vdash P\vee Q}{\sum_{...}P\vee Q,P\vdash P} \stackrel{hyp}{\longrightarrow} \frac{\overline{\sum_{...}P\vee Q,P\vdash P\to R}}{\sum_{...}P\vee Q,P\vdash R} \stackrel{hyp}{\longrightarrow} \frac{\overline{\sum_{...}P\vee Q,Q\vdash Q}}{\sum_{...}P\vee Q,Q\vdash Q} \stackrel{hyp}{\longrightarrow} \frac{\overline{\sum_{...}P\vee Q,Q\vdash Q\to R}}{\sum_{...}P\vee Q,Q\vdash R} \stackrel{hyp}{\longrightarrow} E$$

$$\frac{\overline{\sum_{...}P\vee Q\vdash R}}{\sum_{...}P\vee Q\to R} \rightarrow I$$

c)

Stelling	#hyp	$\# \to E$	#  o I	$\# \wedge E1$	$\# \wedge E2$	$\# \wedge I$	$\# \vee E$	$\# \vee I1$	$\# \vee I2$	#totaal
(2)	3	2	0	0	0	0	0	0	0	5

Stelling	#hyp	$\# \to E$	$\# \to I$	$\# \wedge E1$	$\# \wedge E2$	$\# \wedge I$	$\# \vee E$	$\# \vee I1$	$\# \lor I2$	#totaal
(4)	3	2	1	1	1	0	0	0	0	8

ſ	Stelling	#hyp	$\# \to E$	$\# \to I$	$\# \wedge E1$	$\# \wedge E2$	$\# \wedge I$	$\# \vee E$	$\# \vee I1$	$\# \vee I2$	#totaal
ĺ	(6)	5	0	0	2	1	2	1	1	1	13

Stelling	#hyp	$\# \to E$	$\# \to I$	$\# \wedge E1$	$\# \wedge E2$	$\# \wedge I$	$\# \vee E$	$\# \vee I1$	$\# \lor I2$	#totaal
(8)	5	2	0	0	0	0	1	0	0	8

#### Opgave 2

Dit is het alternatieve bewijs voor het feit dat uit de aannames volgt dat Kees gelukkig is.

Stelling:  $\Sigma \vdash KG$ 

Afkorting(en): 
$$\Sigma = JV \land PK, JV \rightarrow ML, PK \rightarrow WT, KT \lor \neg KT, WT \rightarrow \neg KT, KT \land ML \rightarrow KG, \neg KT \rightarrow KG$$

Afleidingsboom:

$$\underbrace{\frac{\sum_{,KT\vdash JV \to ML} hyp}{\sum_{,KT\vdash KT} hyp} \frac{\frac{\sum_{,KT\vdash JV \to ML} hyp}{\sum_{,KT\vdash JV} \to ML} \bigwedge^{hyp}}_{\sum_{,KT\vdash KT} hyp} \frac{\frac{\sum_{,KT\vdash JV \to ML} hyp}{\sum_{,KT\vdash KT} \land ML}}{\sum_{,KT\vdash KT} \land ML} \to E} \underbrace{\frac{\sum_{,KT\vdash KT} hyp}{\sum_{,KT\vdash KG} \to E}}_{\sum_{,KT\vdash KG}} \underbrace{\frac{\sum_{,KT\vdash JV \to ML} hyp}{\sum_{,KT\vdash KG} \land KG}} \bigwedge^{hyp}}_{\sum_{,KT\vdash KG}} \underbrace{\frac{\sum_{,KT\vdash KT \to ML} hyp}{\sum_{,KT\vdash KG} \land KG}} \bigwedge^{hyp}}_{\sum_{,KT\vdash KG}} \underbrace{\frac{\sum_{,KT\vdash KT \to ML} hyp}{\sum_{,KT\vdash KG} \land KG}} \bigwedge^{hyp}}_{\sum_{,KT\vdash KG}} \underbrace{\frac{\sum_{,KT\vdash KT \to ML} hyp}{\sum_{,KT\vdash KG} \land KG}} \bigwedge^{hyp}}_{\sum_{,KT\vdash KG}} \underbrace{\frac{\sum_{,KT\vdash KT \to ML} hyp}{\sum_{,KT\vdash KG} \land KG}} \bigwedge^{hyp}}_{\sum_{,KT\vdash KG}} \underbrace{\frac{\sum_{,KT\vdash KT \to ML} hyp}{\sum_{,KT\vdash KG} \land KG}} \bigwedge^{hyp}}_{\sum_{,KT\vdash KG}} \underbrace{\frac{\sum_{,KT\vdash KT \to ML} hyp}{\sum_{,KT\vdash KG} \land KG}} \bigwedge^{hyp}}_{\sum_{,KT\vdash KG}} \underbrace{\frac{\sum_{,KT\vdash KT \to ML} hyp}{\sum_{,KT\vdash KG}}}_{\sum_{,KT\vdash KG}} \underbrace{\frac{\sum_{,KT\vdash KT \to ML} hyp}{\sum_{,KT\vdash KG}}}_{\sum_{,KT\vdash KG}} \underbrace{\frac{\sum_{,KT\vdash KG} hyp}{\sum_{,KT\vdash KG}}}_{\sum_{,KT\vdash KG}}$$

Tak3:

$$\frac{\overline{\Sigma, \neg KT \vdash \neg KT \to KG} \ \underset{\Sigma, \neg KT \vdash \neg KT}{hyp} \ \overline{\Sigma, \neg KT \vdash \neg KT} \ \underset{\Sigma}{hyp}} \to E$$

## Opgave 3

- a) Dit is de definitie van het onderdeel Knop in propositielogica:
- b) En dit is het bewijs van de correctheidsstelling.

Stelling: 
$$\Sigma \vdash \text{KnopIn} \leftrightarrow \text{Rinkel}$$

Afkorting(en):  $\Sigma = \text{SpanningA}, \text{SpanningB} \leftrightarrow \text{SpanningC}, \text{Rinkel} \leftrightarrow \text{SpanningC}, \text{SpanningA} \rightarrow (\text{KnopIn} \leftrightarrow \text{SpanningB})$ 

Afleidingsboom:

$$\frac{\frac{\sum_{i} \text{Rinkel} \vdash \text{SpanningB}}{\sum_{i} \text{Rinkel} \vdash \text{SpanningB}} hyp}{\frac{\sum_{i} \text{Rinkel} \vdash \text{SpanningB}}{\sum_{i} \text{Rinkel} \vdash \text{SpanningB}}}{\frac{\sum_{i} \text{Rinkel} \vdash \text{SpanningB}}{\sum_{i} \text{Rinkel} \vdash \text{KnopIn}}} \leftrightarrow E2} \leftrightarrow E2$$

$$\frac{Tak1}{\sum_{i} \vdash \text{KnopIn} \leftrightarrow \text{Rinkel}} \leftrightarrow I$$

Tak1:

$$\frac{\sum_{\mathsf{K}} \mathsf{KnopIn} \vdash \mathsf{Rinkel} \leftrightarrow \mathsf{SpanningC}}{\mathsf{E}, \mathsf{KnopIn} \vdash \mathsf{Rinkel}} \overset{hyp}{\longleftarrow} \underbrace{\frac{\mathsf{T}ak3}{\mathsf{\Sigma}, \mathsf{KnopIn} \vdash \mathsf{KnopIn}}_{\mathsf{\Sigma}, \mathsf{KnopIn} \vdash \mathsf{SpanningB}}_{\mathsf{\Sigma}, \mathsf{KnopIn} \vdash \mathsf{SpanningC}}}_{\mathsf{\Sigma}, \mathsf{KnopIn} \vdash \mathsf{SpanningC}} \overset{hyp}{\longleftrightarrow} E1}$$

Tak2:

$$\frac{\overline{\Sigma, \operatorname{Rinkel} \vdash \operatorname{Rinkel}} \leftrightarrow \operatorname{SpanningC}}{\Sigma, \operatorname{Rinkel} \vdash \operatorname{SpanningC}} \xrightarrow{L} \frac{hyp}{\Sigma, \operatorname{Rinkel} \vdash \operatorname{Rinkel}} \leftrightarrow E1$$

Tak3:

$$\frac{\overline{\Sigma, \operatorname{Rinkel\vdash SpanningA}} \ hyp}{\Sigma, \operatorname{Rinkel\vdash SpanningA} \to (\operatorname{KnopIn} \leftrightarrow \operatorname{SpanningB})} \ \underset{\Sigma, \operatorname{Rinkel\vdash KnopIn} \leftrightarrow \operatorname{SpanningB}}{} \ hyp}$$