

Motivation

- Most of existing file systems only address cloud systems
 - NFS, IPFS, Ceph
- Edge scenario
 - IoT cameras and sensors continuously producing data
 - Operators want to store and observe the stream of data with low latency

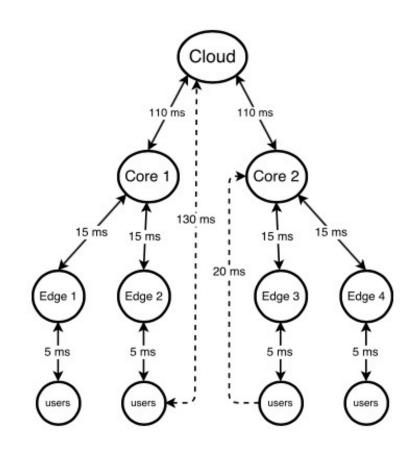
Motivation

Solution:

- A file system that addresses the heterogeneity of resources
- Provides fast data propagation to support streaming
- Building on top of existing systems
 - CloudPath paper: an edge computing framework with routing, execution, and storage facilities
 - PathStore: CloudPath's storage system
 - Based on Cassandra
 - Data represented by CQL queries

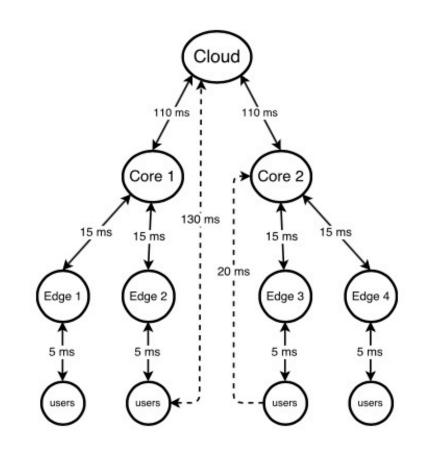
PathStore Architecture

- Hierarchy of nodes
- Each node store a subset of parent's data
- Moving down the hierarchy
 - Resources are more limited
 - Delay and bandwidth usage decrease
- Data replication policy
 - Lazy propagation of updates
 - Push and Pull servers
 - Interest subscription



Event-driven System

- Updates will be propagated eagerly
 - Upward propagation
 - Downward propagation
- Downward propagation based on interest subscription in PathStore
 - Use a private table to check query matching



File System Interface

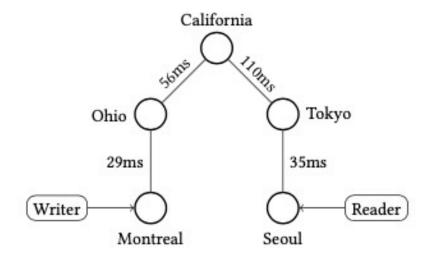
- Entry vs Filedata tables
- Implemented FUSE interface
 - File System for User Space
 - FUSE Syscalls
 - Read
 - Write
 - Readdir
 - Open
 - Create
 - Symlink

```
Entry (parent_id, name):
       parent id: uuid
       name: varchar
       id: uuid
       mode: int (type & permissions)
       userld: int
       groupld: int
       size: bigint
       created: timestamp
       modified: timestamp
       changed: timestamp
       blocks: int
       follow: varchar
Filedata (file_id, insert_time)
       file id: uuid
       insert_time: timestamp
       raw data: blob
```

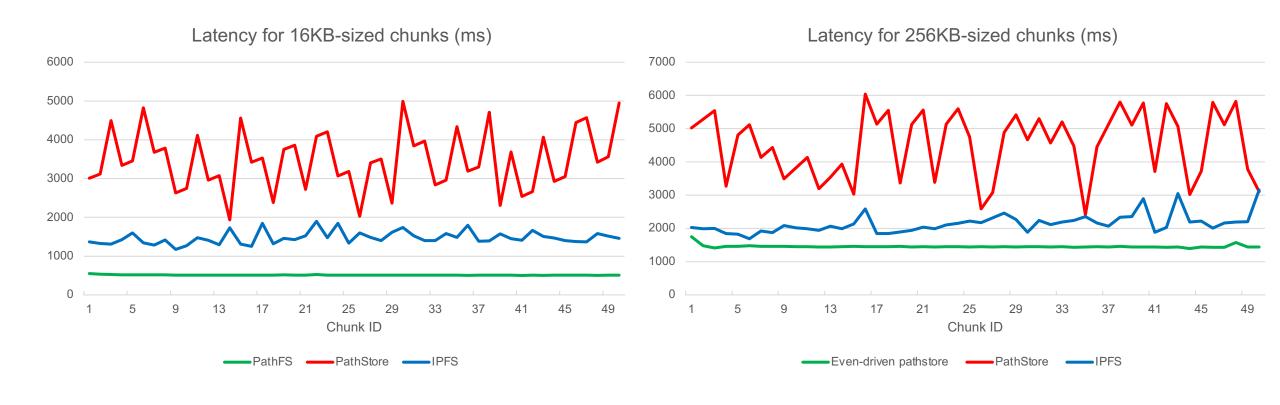


Evaluation

- Evaluating the event-driven system
- 5-node topology with simulated latency
- Writer repeatedly inserts 16KB and 256KBsized data chunks
- Measured how long it takes to reach reader
- Compared with PathStore and IPFS



Event-driven system results



FUSE Evaluation

- Meaursing single-node performance of the schema and FUSE implementation
- Used FIO benchmark
 - Repeatedly reading/writing data with 4K-sized chunks
- Compared with the native file system, nfs, and Cassandra

FUSE Performance

FIO Benchmark (random read/ writes)	PathFS	Cassandra pathfs	native	nfs (1 host)
read	629 KB/s	1.2 MB/s	1.8 MB/s	621 KB/s
write	230 KB/s	444 KB/s	666 KB/s	227 KB/s

Summary

- PathFS: File system for edge environments
 - Heterogeneity of resources
 - Support for streaming
- Built on top of PathStore: A key-value edge storage using Cassandra
- Contributions:
 - Event-driven update propagation
 - FUSE interface support
- 60% Improvement in propagation latency

Future Work

- More rigorous testing of FUSE
 - Finding out performance bottlenecks
- PathFS writes are append-only
 - Implement support for edits
- Conflict handling



System Overview



Experiments

- Measure each layer's overhead
- Experiments with:
 - Cassandra simple insert/update queries
 - Cassandra + our schema (fuseless)
 - Cassandra + our schema + fuse
 - Cassandra + our schema + fuse + FIO (the benchmark)

New Swift Nodes!

- Infrastructural changes:
 - Moved the swift nodes to SSD drives
 - Cool and fast nodes!
 - Unit tests for PathFS

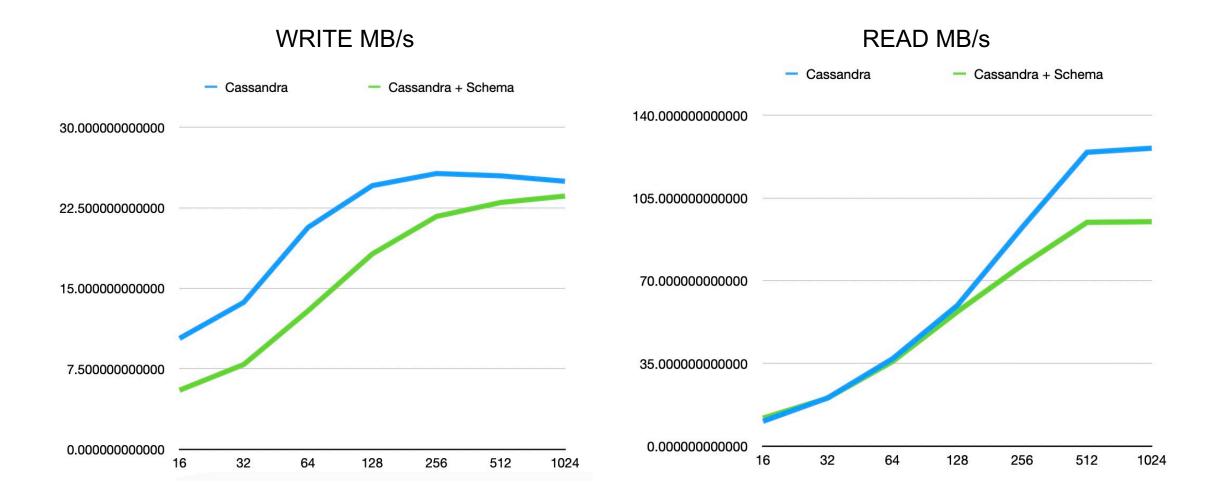
Experiments

- PathFS now assumes fixed chunk sizes
- In all experiments:
 - We create a 128MB-sized file
 - Add chunks one by one
 - Close and and re-open
 - Read chunks one by one
- Chunk sizes vary!

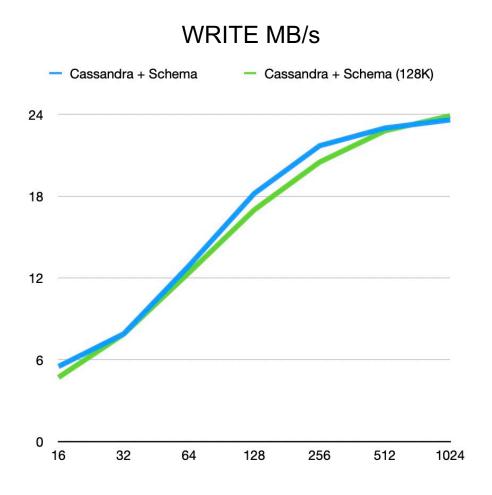
Wow!

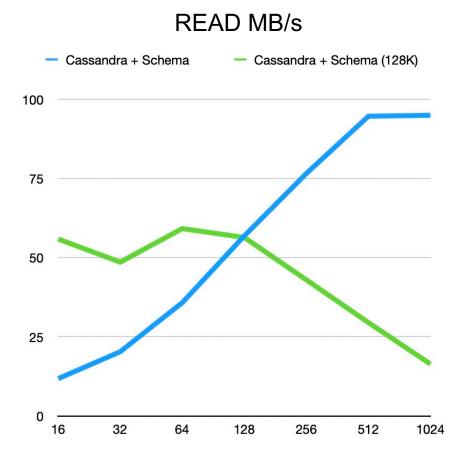
- Always seen performance degradation in fuse for large chunks
- The reason: Fuse modifies the requested chunk size!
 - 0 4K -> 4K
 - 4 16K -> 16K
 - 16 64K -> 64K
 - 64 infinity -> 128
- Also, Fuse is multi-threaded

Schema Overhead



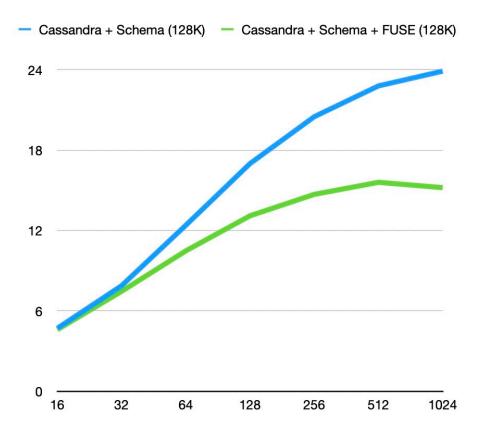
The Impact of Fixed Sizes



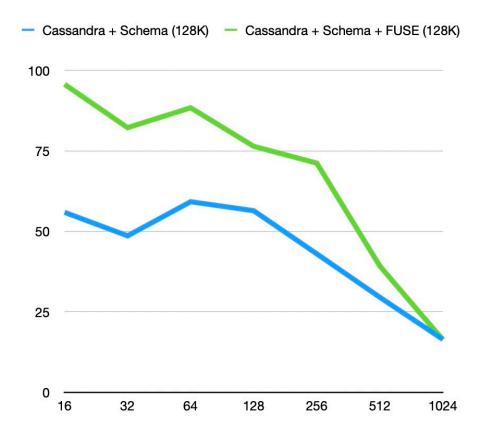


FUSE Overhead!

WRITE MB/s

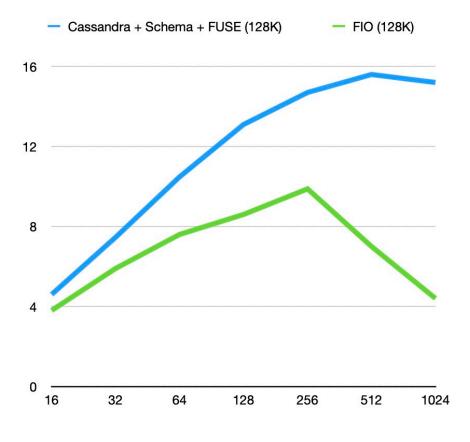


READ MB/s



FIO Overhead

WRITE MB/s



READ MB/s

