

Discussion 10: November 11th

Operational Semantics

Formal semantics of a PL: Mathematical description of meaning of programs in a PL.

Terminologies in Operational Semantics

$$A; e1 \Rightarrow v1 \quad A; e2 \Rightarrow v2 \quad v3 \text{ is } v1 + v2$$
$$-----$$
$$A; e1 + e2 \Rightarrow v3$$

- **Expression:** A program that evaluates to a value
- **Value:** A result of an expression
- **Environment:** A mapping from variables to values
- **Hypothesis:** A set of rules that describe the meaning of expressions
- **Judgement:** A statement with expressions and values ($e \Rightarrow v$). The expression e evaluates to the value v .

Operational Semantics Exercise

Solve the following problems using operational semantics:

Evaluate $1 + (2 + 3) = 6$ using the following hypotheses:

$$-----$$
$$n \Rightarrow n$$
$$e1 \Rightarrow v1 \quad e2 \Rightarrow v2 \quad v3 \text{ is } v1 + v2$$
$$-----$$
$$e1 + e2 \Rightarrow v3$$

Evaluate the expression using the given hypotheses: $A; \text{let } y = 1 \text{ in let } x = 2 \text{ in } x \Rightarrow 2$

$\frac{}{n \Rightarrow n}$	$\frac{A(x) = v}{A; x \Rightarrow v}$
$\frac{A; e1 \Rightarrow v1 \quad A, x: v1; e2 \Rightarrow v2}{A; \text{let } x = e1 \text{ in } e2 \Rightarrow v2}$	$\frac{A; e1 \Rightarrow v1 \quad A; e2 \Rightarrow v2 \quad v3 \text{ is } v1 + v2}{A; e1 + e2 \Rightarrow v3}$

Lambda Calculus

Lambda Calculus Terminologies

$\lambda x. \lambda y. x + y + a$

- **Lambda Expression:** A lambda expression is a function that takes an argument and returns a value.
- **Free Variable:** A variable that is not bound by a lambda expression.
- **Lambda Abstraction:** A lambda abstraction is a function that takes an argument and returns a function.
- **Alpha Conversion:** A process of renaming a variable in a lambda expression to avoid name conflict. Does not change the meaning of the expression. **Do not rename free variables**
- **Beta Reduction:** A process of substituting a lambda expression for a variable in a lambda abstraction.

Lambda Calculus Exercise

Things to keep in mind: 1. Alpha conversion: Do not rename free variables 2. Explicit Parentheses: Scope of a variable extends to far right or the first `)` seen. 3. Lambda Calculus is left-associative. 4. Beta reduction: Keep applying the functions until you can't anymore.

Solve the following problems using lambda calculus:

1. Make parentheses explicit in the following lambda expressions:

1. $a\ b\ c$

2. $\lambda a. \lambda b. a\ b$

3. $\lambda a. a\ b\ \lambda a. a\ b$

2. Identify the free variables in the following lambda expressions:

1. $\lambda a. a\ b\ a$

2. $\lambda a. \lambda b. a\ b$

3. $\lambda a. (\lambda b. a\ b)\ a\ b$

3. Perform alpha conversion on the following lambda expressions:

1. $\lambda a. \lambda a. a$

2. $(\lambda a. a)\ a\ b$

3. $(\lambda a. (\lambda a. (\lambda a. a)\ a)\ a)\ a$

4. Perform beta reduction on the following lambda expressions:

1. $(\backslash a. a b) \times b$

2. $(\backslash a. b) (\backslash a. \backslash b. \backslash c. a b c)$

3. $(\backslash a. a a) (\backslash a. a a)$