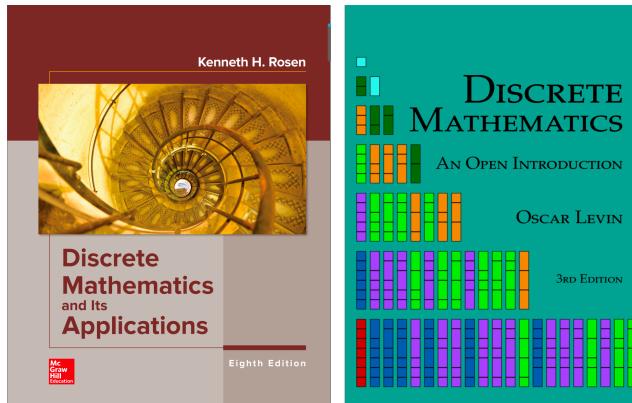




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**Session 9
Counting (part 3), October, 2022**

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Introduction

- Part 1: The basic rules of counting
 - ❖ Product rules
 - ❖ Sum rule
 - ❖ The Inclusion-Exclusion rule
 - ❖ The Pigeonhole Principle
- Part 2: Permutations and Combinations
 - ❖ Permutation
 - ❖ Combination
 - ❖ Formulas of combinations
 - ❖ The extended combination principles
- Part 3: Enumeration

Part 3: Enumeration

- A. Recursive method**
- B. Loop method**

Recursive method

- ❖ General rules
- ❖ Problem: eight queens
- ❖ Enumerate the simple combination

General rule

In order to enumerate all cases $S=s_1s_2..s_k$

The algorithm suppose having the sub-case $s_1s_2..s_{i-1}$. In the tried step of finding s_i , **Try(i)**, we loop for all values of j that are voted for s_i and perform the following steps:

General rule

For each j (voted for s_i), execute four steps:

- 1) $s_i = j;$
- 2) *<change the status>;*
- 3) If ($i == k$) Print(S); else Try($i+1$);
- 4) *<return the old state>;*

The main function call $Try(1)$

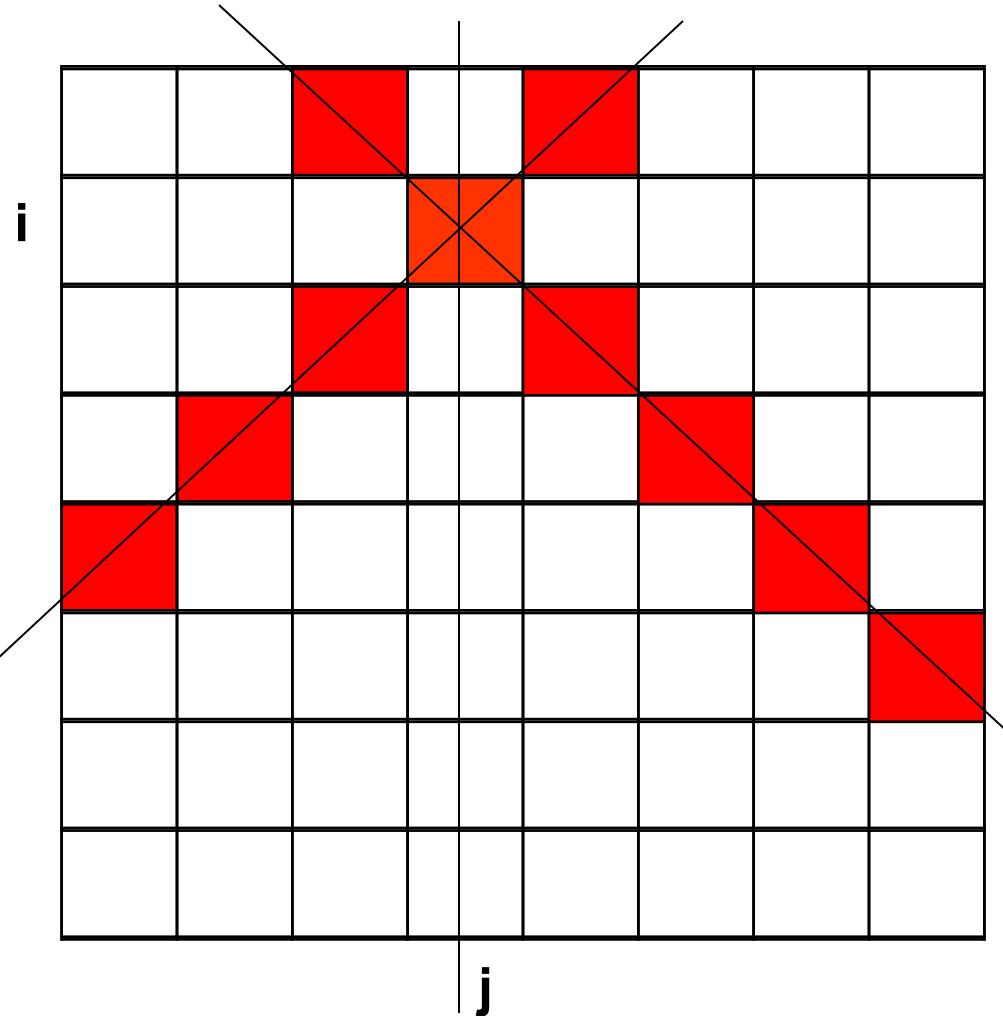
General rule

Note: in some cases, there is no step 2) and 4)

The above enumerated algorithm, the case of simple combination are always assumed that $X = \{1..n\}$.

In the real case, take elements of X to index for array, and not use $S[0]$

Problem: eight queens



There are no two queens that share the same row, column, or diagonal.

Problem: eight queens

Declare:

- The chess table: `int S[8];`
- The arrays that need to note:
 - column: `int a[8];`
 - parallel to the main diagonal: `int b[15];`
 - parallel to the sub diagonal: `int c[15];`

Before calling Try(0) need to full-fill a, b, c with value 1 (TRUE)

Problem: eight queens

```
void Try(int i)
{
    int j;
    for (j=0; j<8; j++)
        if (a[j] && b[i-j+7] && c[i+j])
    {
        S[i]=j;
        a[j]=0; b[i-j+7]=0; c[i+j]=0;
        if (i==7)
            print( );
        else
            Try(i+1);
        a[j]=1; b[i-j+7]=1; c[i+j]=1;
    }
}
```

See: *8queens.cpp*

Enumerate arrangement with repetition

```
void Try(int i)
{
    int j;
    for (j=1; j<=n; j++){
        S[i]=j;
        if (i==k) print();
        else Try(i+1);
    }
}
```

See: *EnumArrRepeat.c*

Enumerate arrangement without repetition

We need an array a to note values j that have been used

```
void Try(int i)
{
    int j;
    for (j=1; j<=n; j++)
        if (a[j])
    {
        S[i]=j;    a[j]=0;
        if (i==k) print( ); else Try(i+1);
        a[j]=1;
    }
}
```

See: *EnumArrNonRepeat.cpp*

Enumerate combination

*Enumerate the increasing orders. Value j (voted for s_i) is:
 $1+s_{i-1} \leq j \leq n - k + i$. Let Try(1): array S with $S[0]=0$.*

```
void Try(int i)
{
    int j;
    for (j=1+S[i-1]; j<=n-k+i; j++){
        S[i]=j;
        if (i==k) print( ); else Try(i+1);
    }
}
```

*See: *EnumCom.c**

Loop method

- ❖ General rule
- ❖ Enumeration of simple combinations

General rule

Enumerate all $S=s_1s_2..s_k$ following the lexical ordering, including the following steps:

- Find S smallest; $\text{Print}(S)$;
- Repeat until end of S and satisfying the condition:
Find the next S ; $\text{Print}(S)$;

or while S is not the largest:

Find the next S ; $\text{Print}(S)$;

General rule

Next step: find the next S depending on the condition that making S change as follows:

- *Find $S[i]$ on the right hand side that can be increased.*
- *Increase $S[i]$ with the smallest increasing value.*
- *Decrease subset $S[i+1..k]$ to smallest.*

Or you can also find $S[i]$ while as the same time of checking stop condition.

Enumerate the repetition arrangement

```
for (i=1; i<=k; i++)  
S[i]=1; print(); c=1;  
while (c<nk)  
{  
    i=k; while (S[i]==n) i--;  
    S[i]++;  
    for (j=i+1; j<=k; j++) S[j]=1;  
    print();  
}
```

See: *EnumErrNonRepeatER.cpp*

Enumerate the combination

Combination is a non-order set. Enumeration follows increasing order: $s_i < s_{i+1}$. The largest of s_i is $n-k+i$.

```
for (i=1; i<=k; i++) S[i]=i; print(); c=1;
while (c<C(n,k))
{
    i=k; while (S[i]==n-k+i) i--;
    S[i]++;
    for (j=i+1; j<=k; j++) S[j]=S[j-1]+1;
    print( );
}
```

Enumerate the permutation

The smallest follows increasing order and the largest follows decreasing order:

```
for (i=1; i<=n; i++) S[i]=i; print( ); c=1;
    while (c<n!)
    {
        i=n-1; while (S[i]>S[i+1]) i--;
        j=n; while (S[j]<S[i]) j--;
        temp=S[i]; S[i]=S[j]; S[j]= temp;
        j=i+1; k=n; while (j<k){
            temp =S[j]; S[j]=S[k]; S[k]= temp; j++; k--;}
        print( );
    }
```

Enumerate the permutation

Putting the algorithm of permutation-generation into each step of generation of combination in order to obtain the generation of non-repetition arrangement.

Homework (G3)

- ❖ Doing the following exercises:
6, 14, 16, 20, 24, 28, 32, 34 (pages: 435, 436)
- ❖ Implement the following exercises using C/C++:
8 Queens (the size is entered from user: slide 10)
- ❖ Enumeration (see algorithms in slides: 11-13)
 - Combination (repeat and not repeat)
 - Arrangement (repeat and not repeat)

Submit on the BB before 10:00 PM, 26/10/2022