

Inferbo: Infer 기반 버퍼 오버런 오류 분석기

허기홍 조성근 이광근
서울대학교
2017. 4. 14 @ 고려대학교



Infer

- Facebook 의 오픈소스 정적 분석기
 - 개별 분석 (modular analysis) 기반
 - 산업 현장에서 쓸수 있을 정도로 빠르고 정확
 - 대상: memory/resource leak, null dereference, buffer-overflow 등

└ Infer

[Docs](#) [Support](#) [Blog](#) [Twitter](#) [Facebook](#) [GitHub](#)

A tool to detect bugs in Android and iOS apps before they ship

Facebook Infer is a static analysis tool - If you give Infer some Objective-C, Java, or C code, it produces a list of potential bugs. Anyone can use Infer to intercept critical bugs before they have shipped to people's phones, and help prevent crashes or poor performance.

GET STARTED

TRY INFER IN YOUR BROWSER



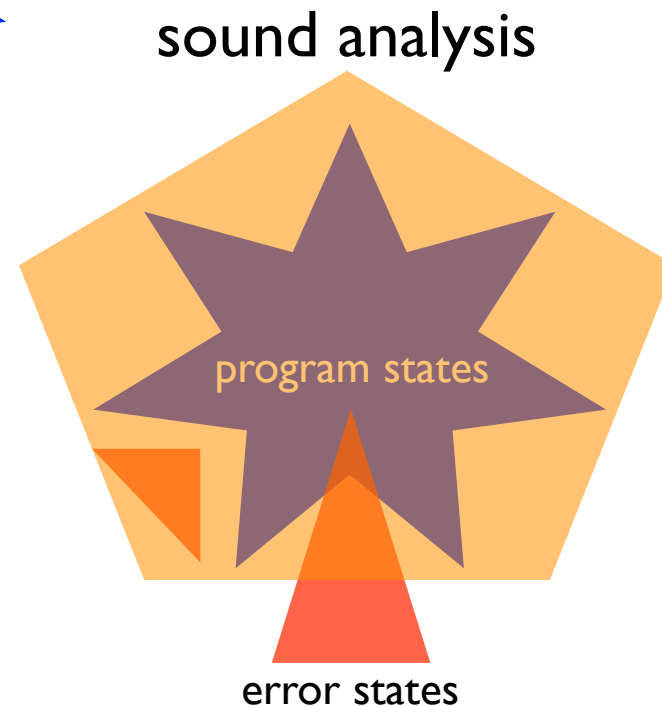
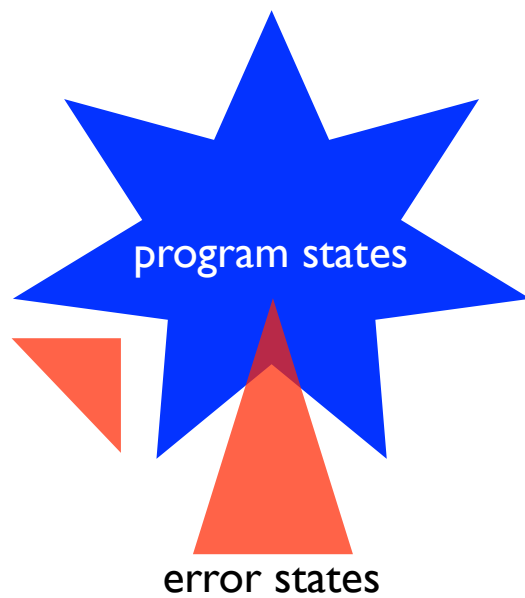
Star

6,371

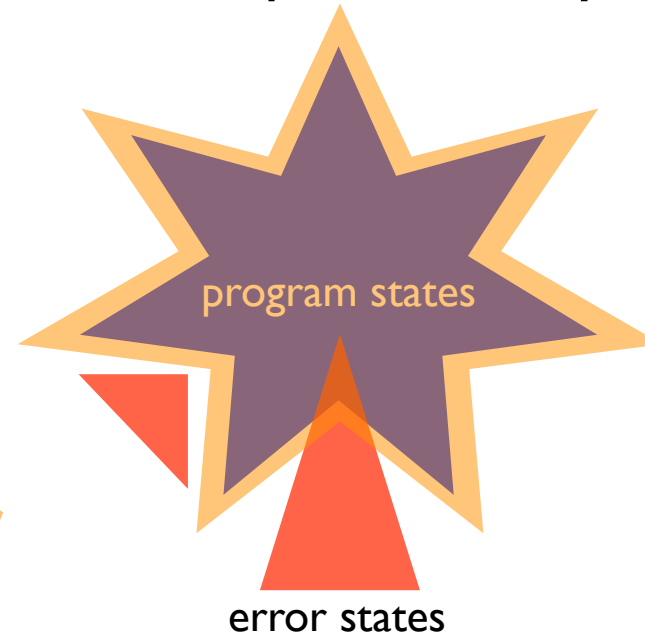


정적 분석

- 자동으로 SW 의 동작을 미리 어림잡는 일반적인 방법
- 목적에 따라 다양하게 요약



sound & precise analysis

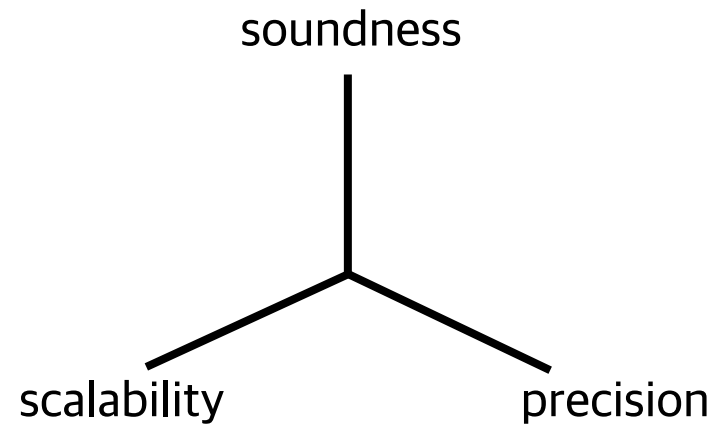


unsound analysis

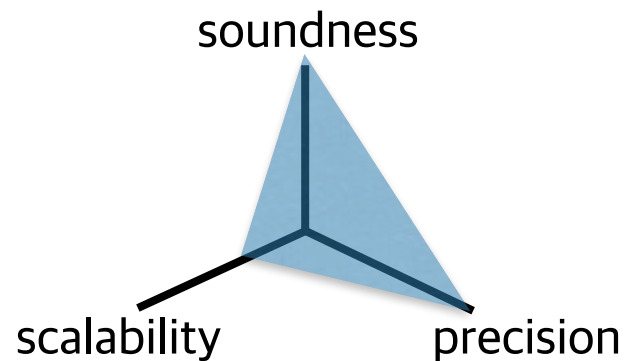


도전 과제

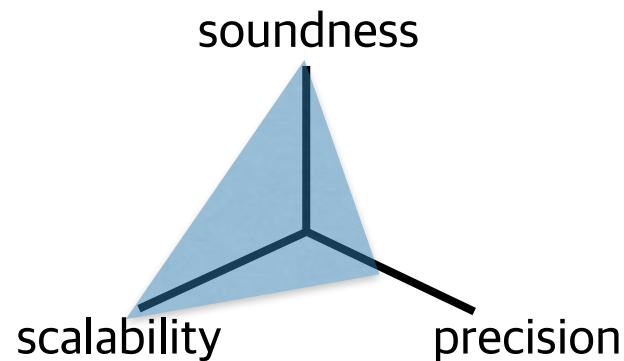
- 성능의 세가지 축: 모두 달성하는 것은 불가능



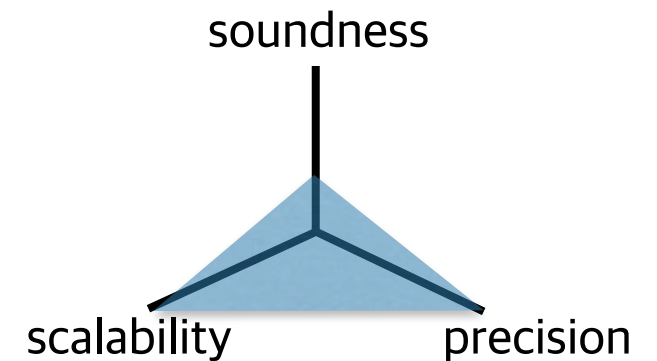
- 전통적인 분류,



무결성 검증용



코드 최적화용

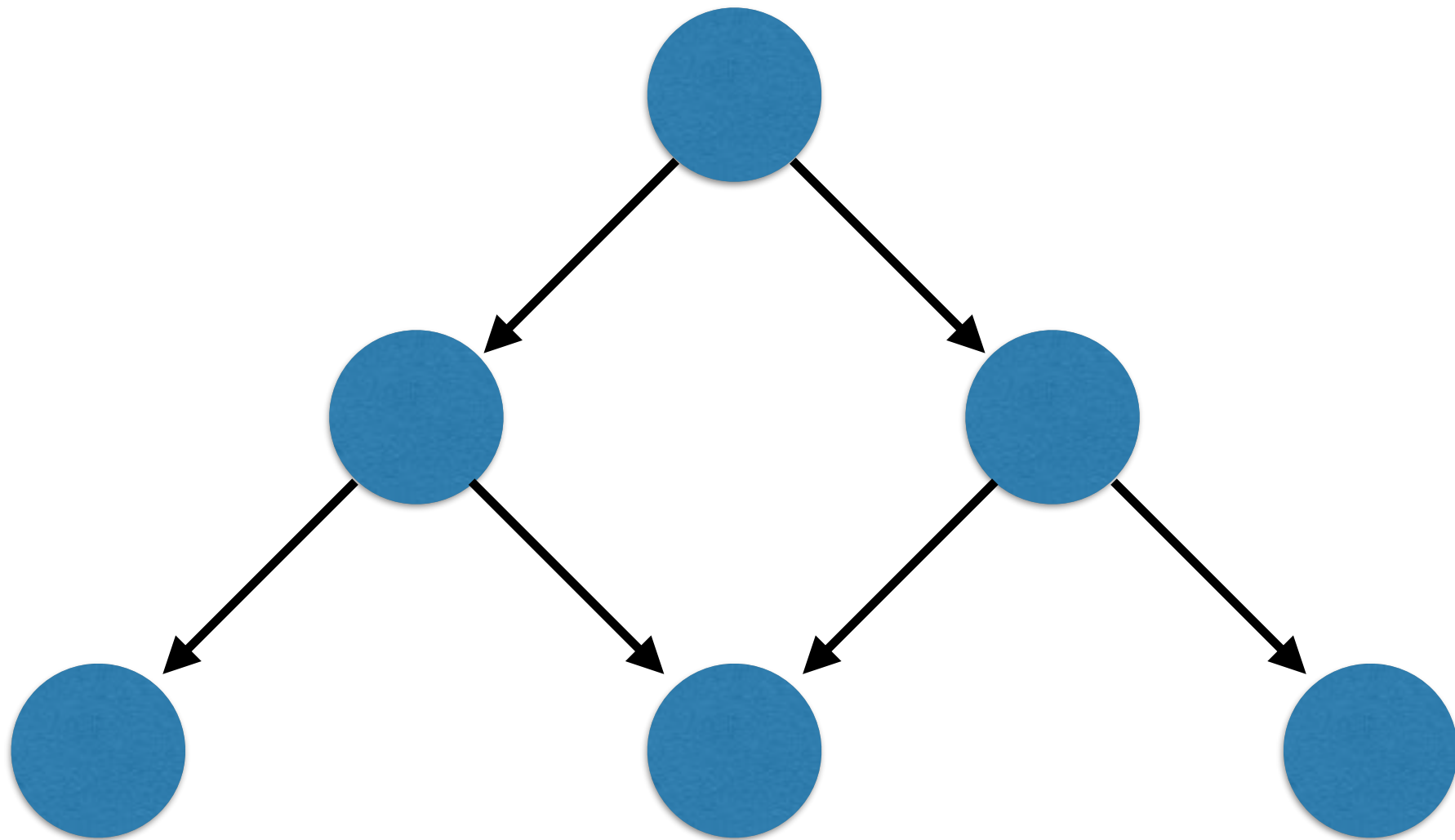


오류 검출용

개별 분석 vs 전체 분석?
(modular) (global)

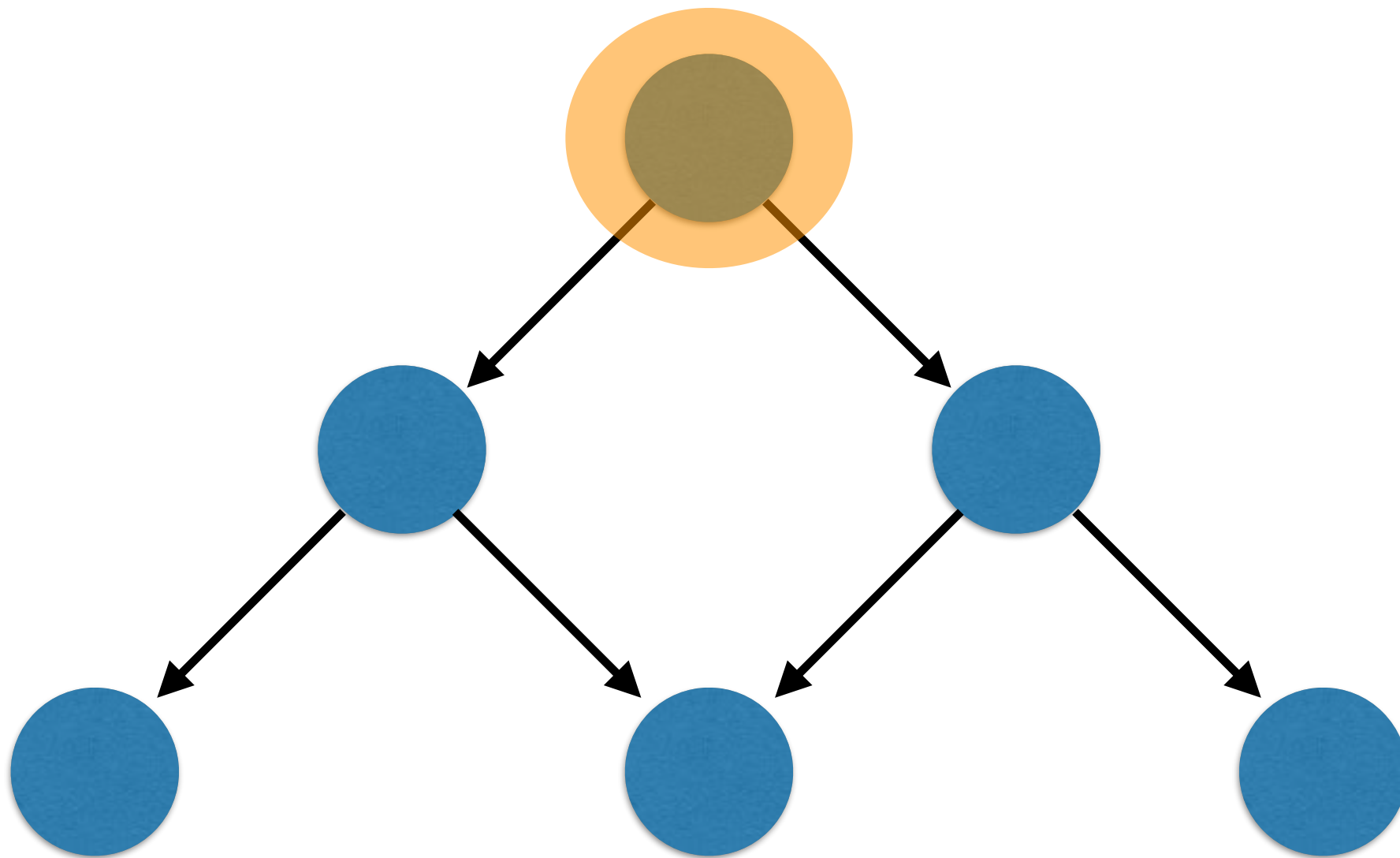
전체 분석 (global analysis)

- 프로그램을 구조를 따라 전체를 한꺼번에 분석
(예, Sparrow)



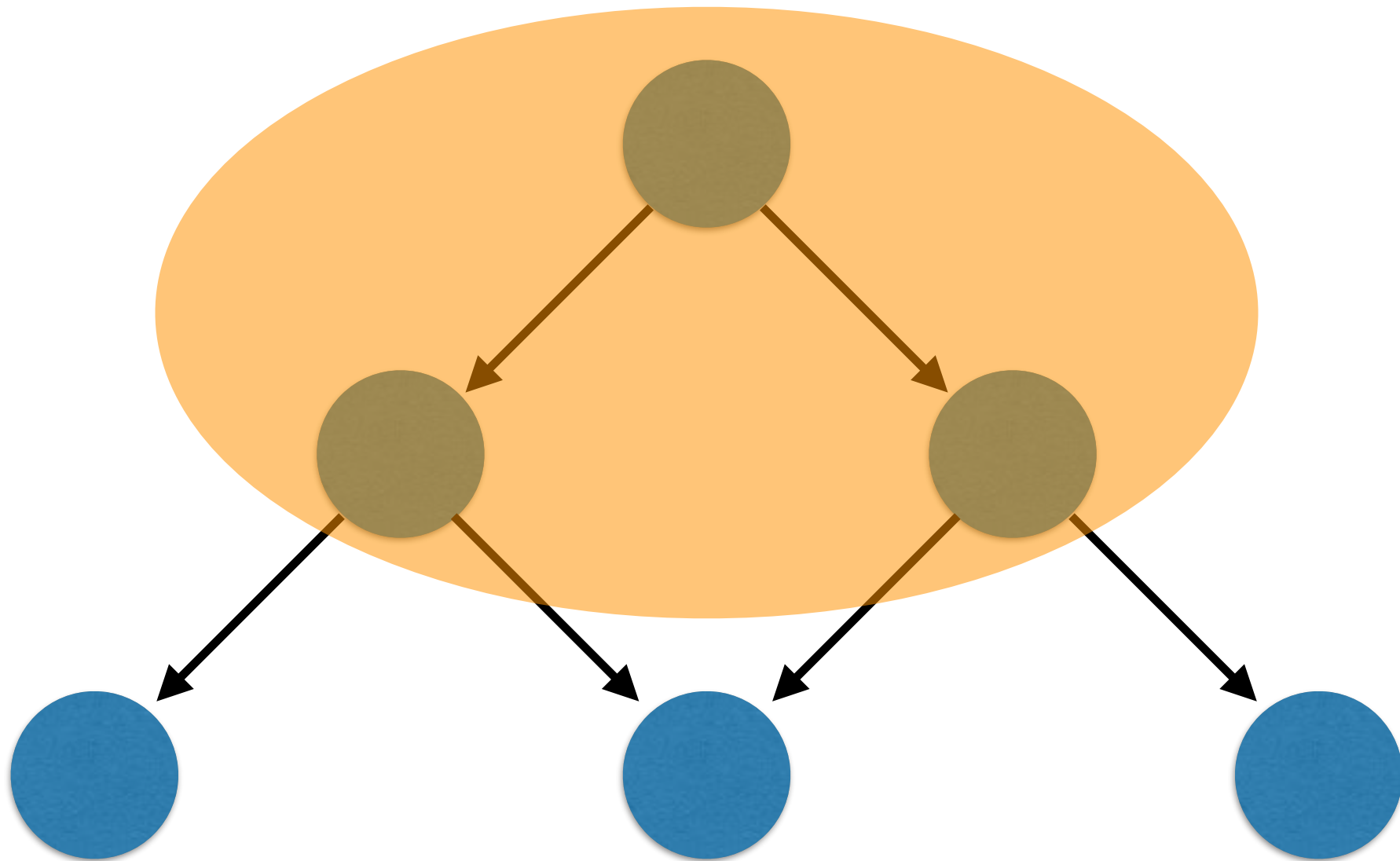
전체 분석 (global analysis)

- 프로그램을 구조를 따라 전체를 한꺼번에 분석
(예, Sparrow)



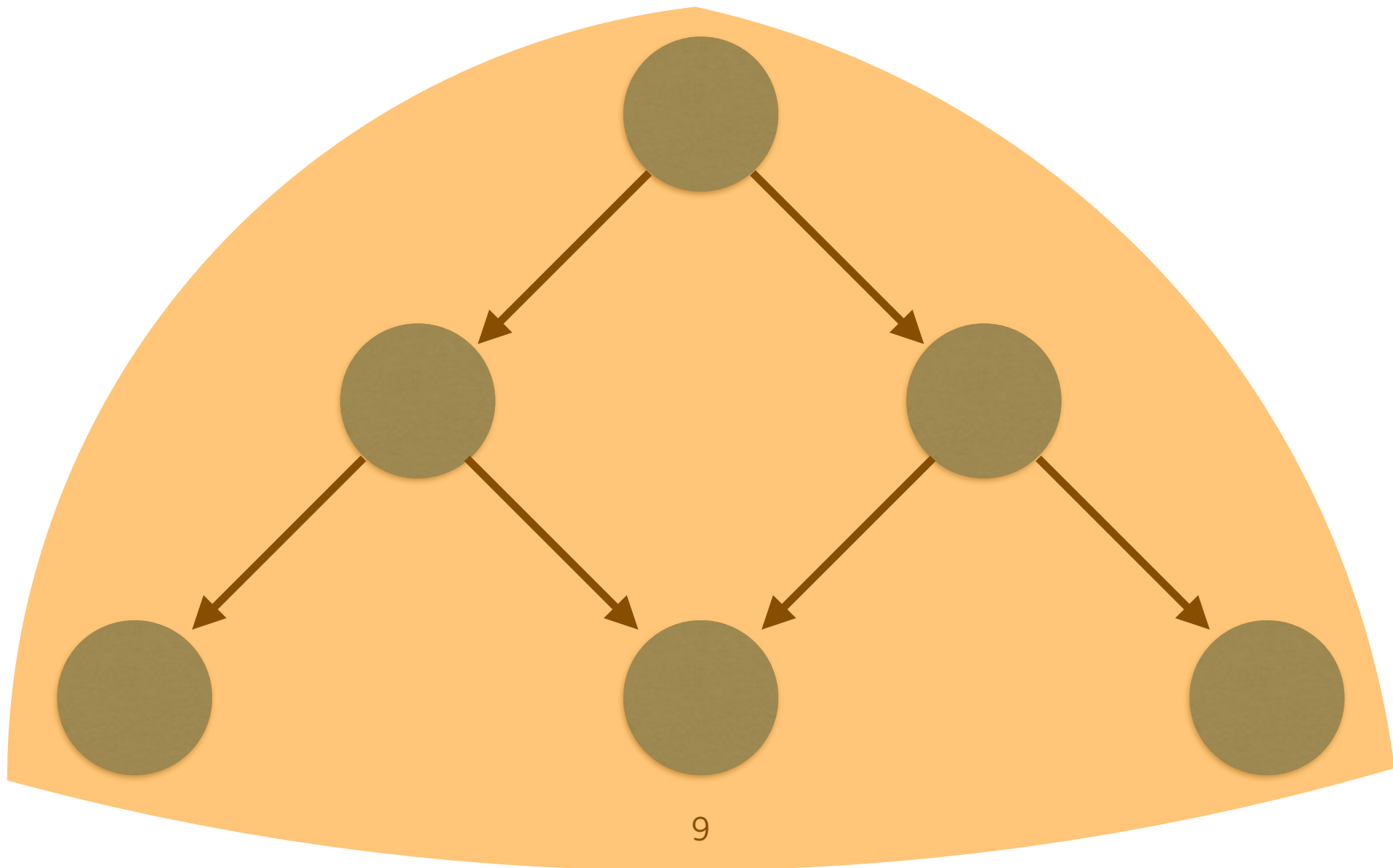
전체 분석 (global analysis)

- 프로그램을 구조를 따라 전체를 한꺼번에 분석
(예, Sparrow)



전체 분석 (global analysis)

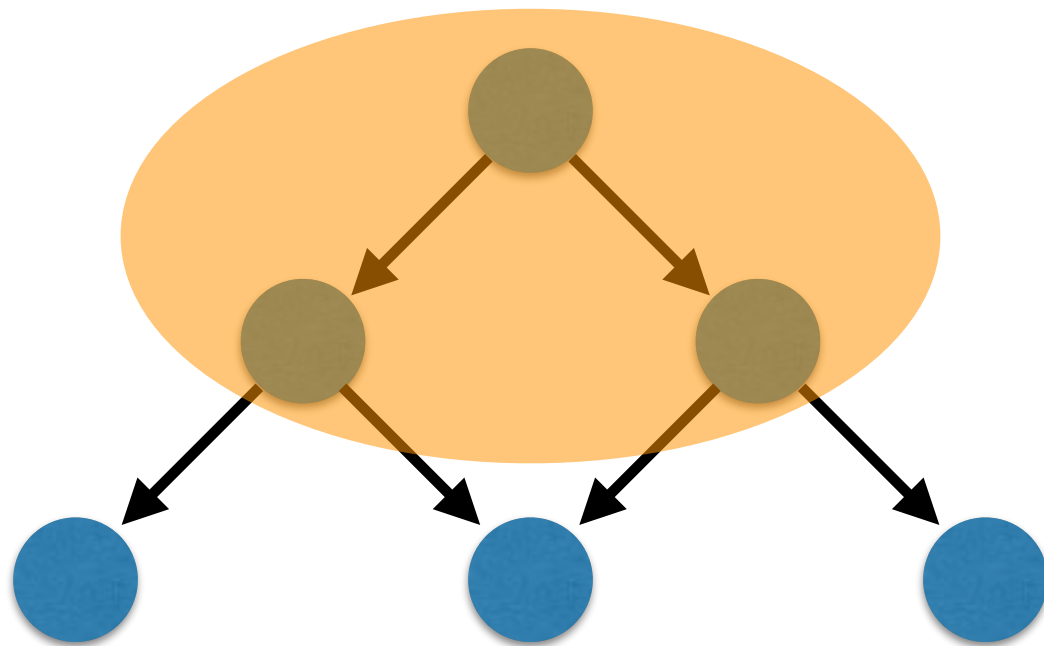
- 프로그램을 구조를 따라 전체를 한꺼번에 분석
(예, Sparrow)



전체 분석 (global analysis)

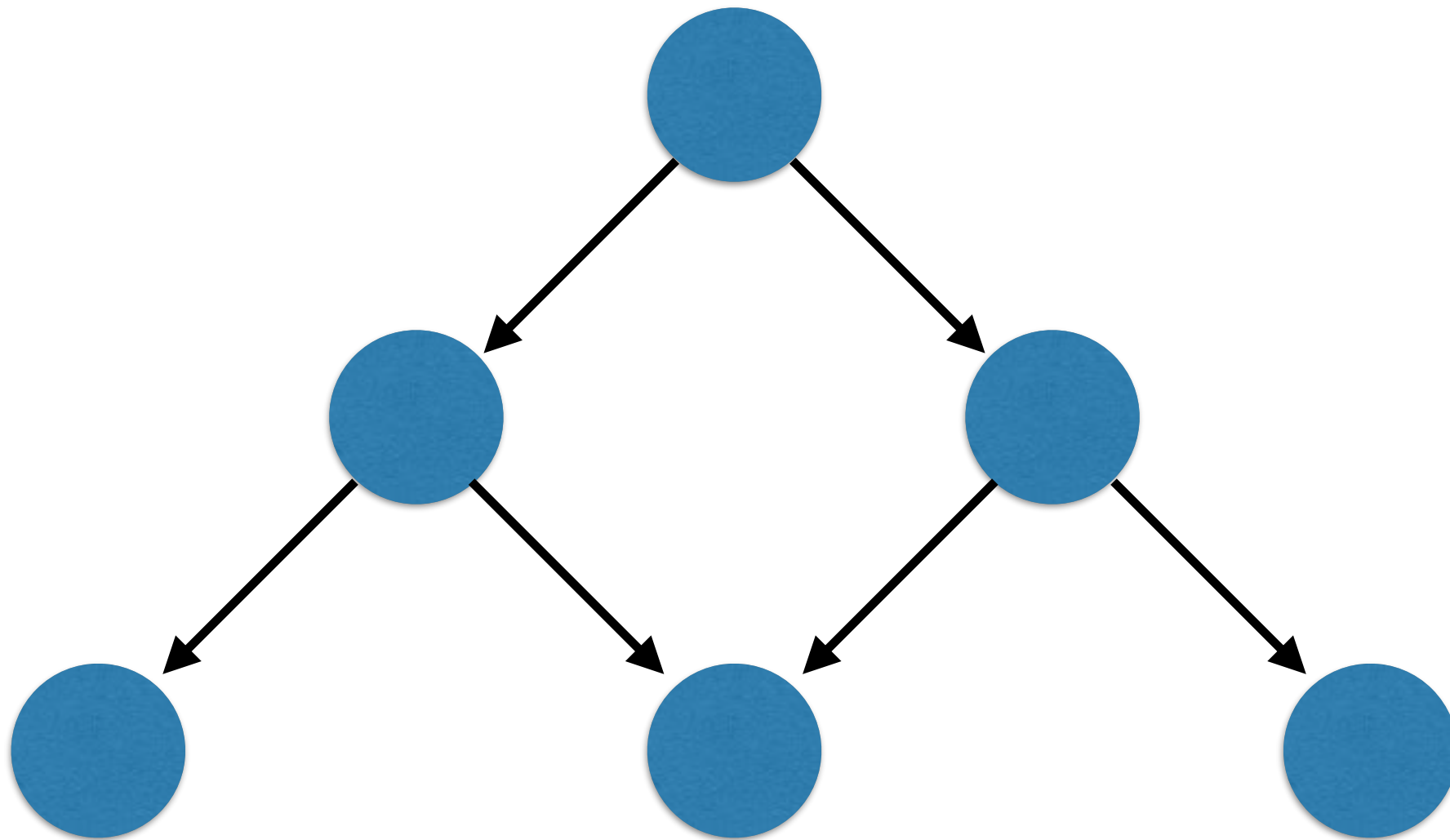
(일반적으로)

- 장점: 문맥을 아는 상태로 분석 (파라미터, 전역변수 등) => 정확
- 단점: 같은 부분을 여러번 분석 => 느림



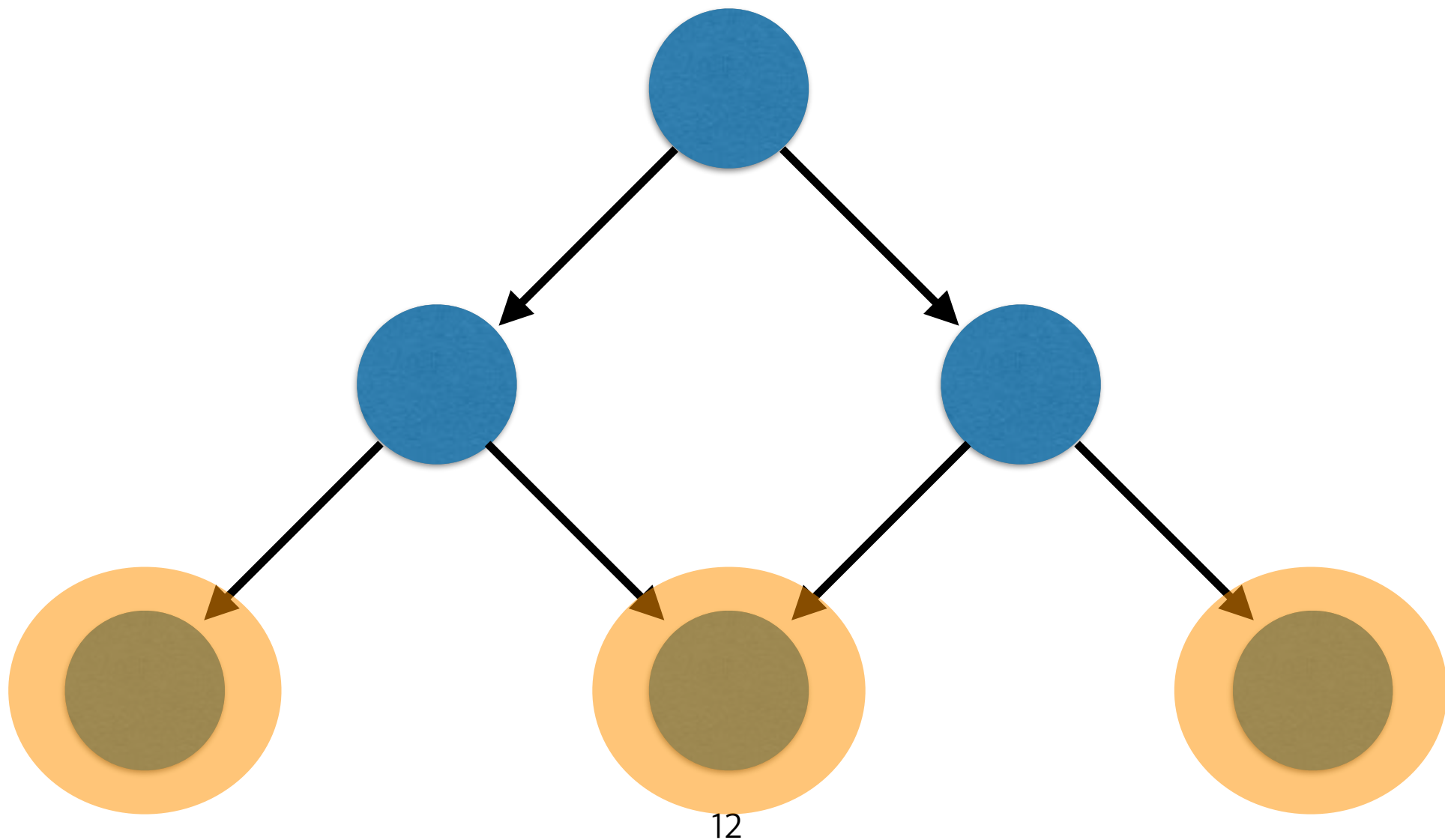
개별 분석 (modular analysis)

- 큰 프로그램의 각 부분(주로 함수)을 따로 분석 + 조합



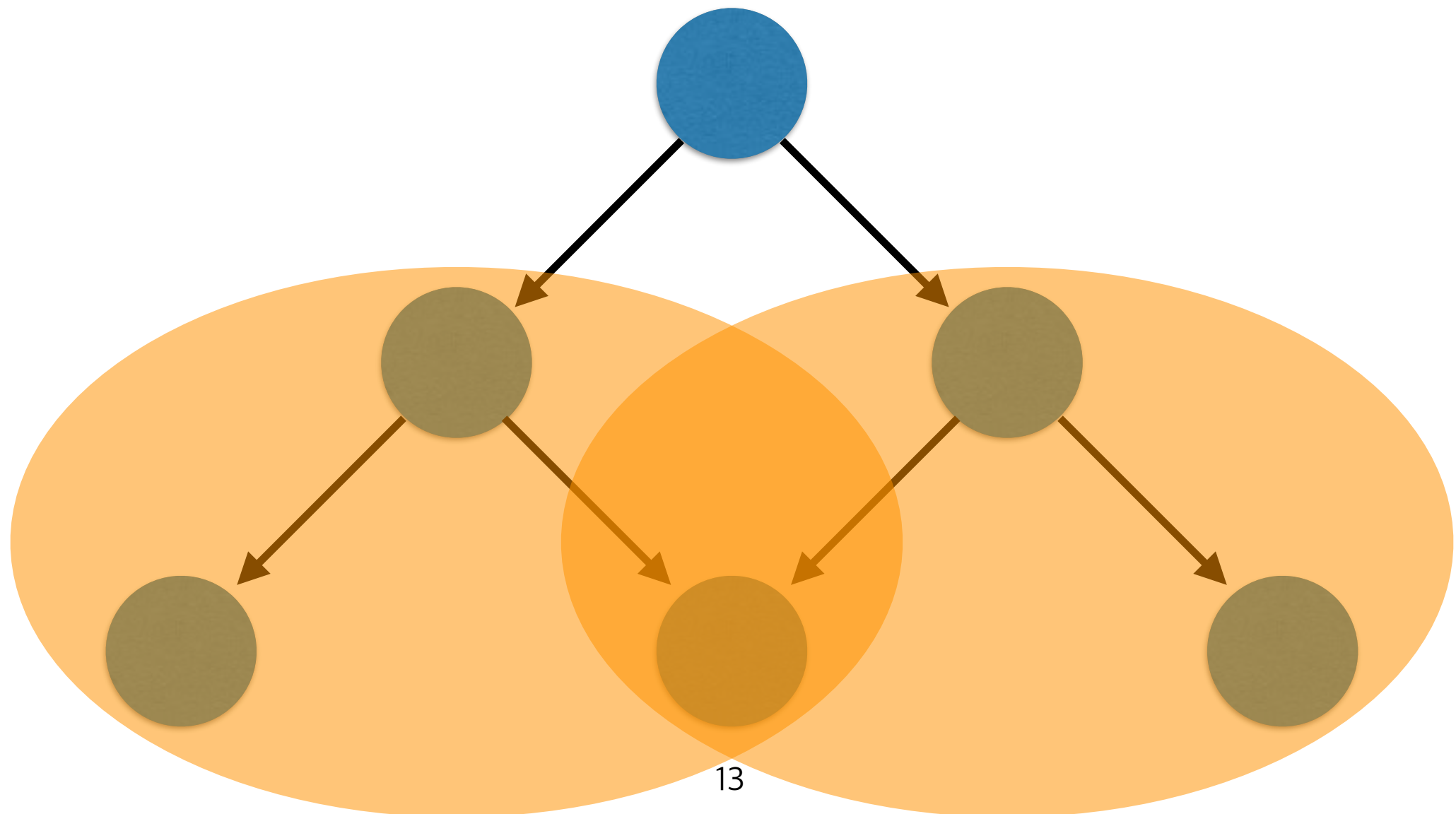
개별 분석 (modular analysis)

- 큰 프로그램의 각 부분(주로 함수)을 따로 분석 + 조합



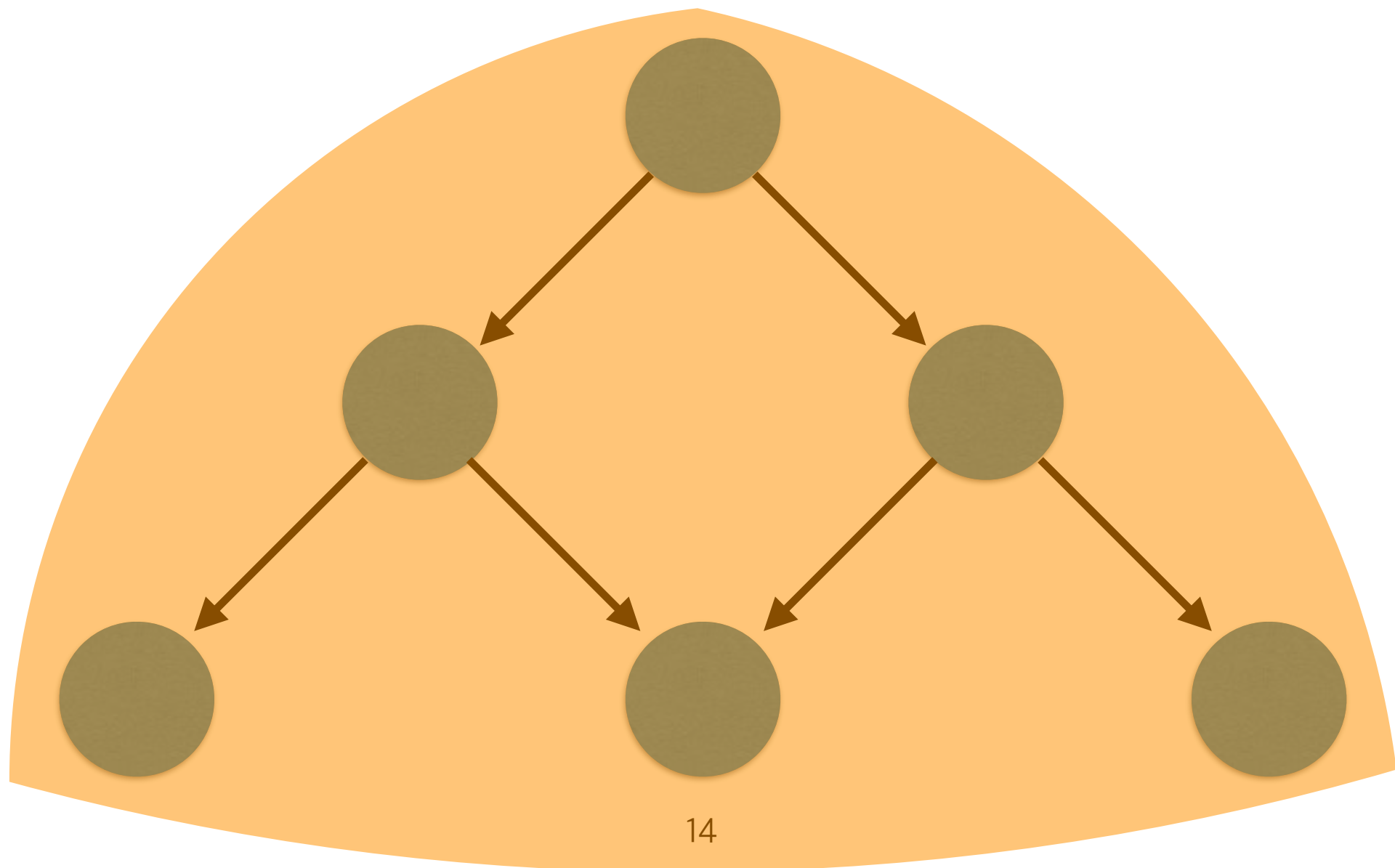
개별 분석 (modular analysis)

- 큰 프로그램의 각 부분(주로 함수)을 따로 분석 + 조합



개별 분석 (modular analysis)

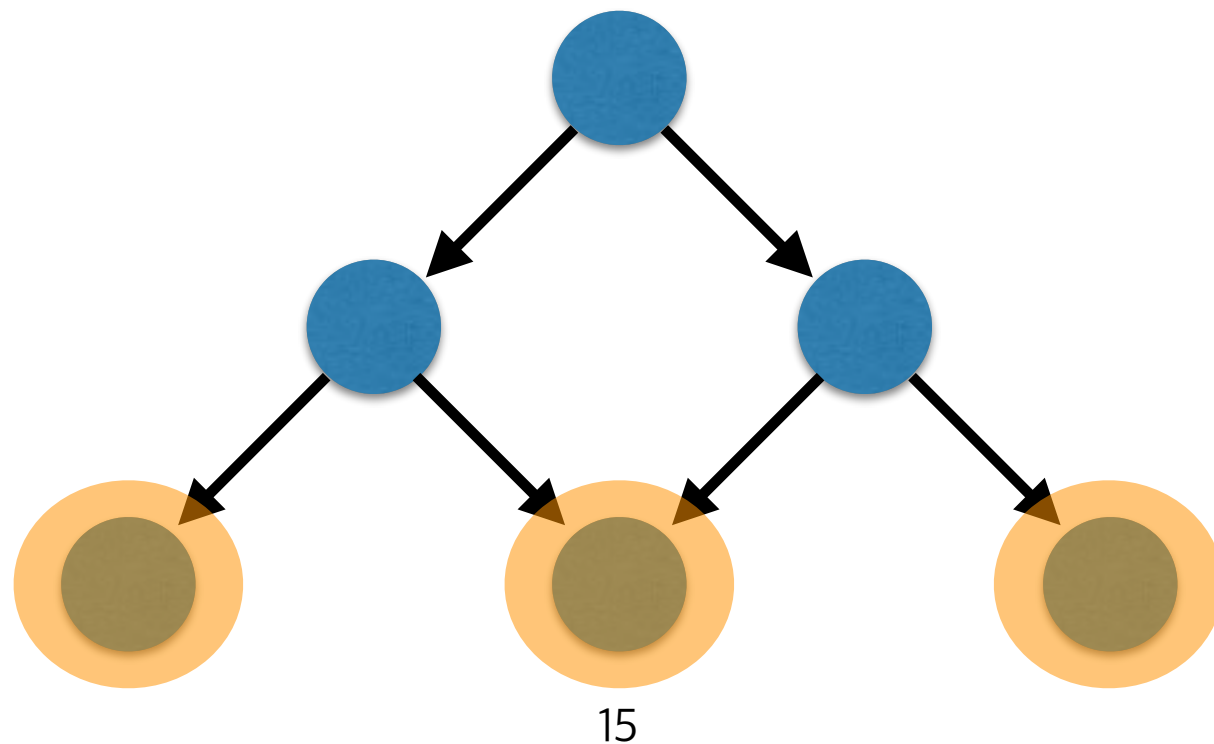
- 큰 프로그램의 각 부분(주로 함수)을 따로 분석 + 조합



개별 분석 (modular analysis)

(일반적으로)

- 장점: 각 함수를 한 번씩만 분석 => 빠름
- 단점: 문맥을 모름 => 부정확 혹은 불안전



개별 분석 (modular analysis)

- 장점 극대화, 단점 최소화하는 핵심 기술:
부분 분석 결과를 재사용 가능하도록 잘 가공
- 전통적으로, 간단하거나 귀납적으로 잘 정의되는 성질 분석 (tree, list, 등)

• 예)

```
void foo(void* p){  
    *p = 0;  
}
```

p 가 null 인경우 위험한 함수

```
void bar(list* p){  
    while(!p)  
        p = p->next;  
}
```

p 가 올바른 list 인경우 잘 순환하는 함수

Inferbo

- 개별분석 + 인터벌 도메인 기반 버퍼오버런 분석기
 - 인터벌: 복잡, 귀납적 x
- 도메인 디자인 원칙:
 - 재사용 가능한 인터벌 분석 결과 => Symbolic bound
 - Facebook 내부 코드에서 발견된 버그 패턴에 집중
- Infer 에 통합되어 Facebook 에서 사용 중

반응



오바다

2월 6일 · 🌐

"Infer is an amazing robot suit for static analysis designers....Infer's modular analysis engine is a rocket that shoots us to the mesmerizing modular analysis orbit"

From Yi Kwangkeun's report on his sabbatical visit with the Infer team@Facebook, and the amazing work he did with his PhD students Kihong Heo and Sungkeun Cho



Inferbo: Infer-based buffer overrun analyzer

This is a report of my experience during a 2-month sabbatical I recently completed with the Facebook Infer static analysis team in London. During my sabbatical with my PhD students, we developed Inferbo, a static analyzer to detect C-like...

RESEARCH.FB.COM

👍 좋아요 💬 댓글 달기 ➦ 공유하기

👍❤️ 김동선님, Woosuk Lee님 외 41명

공유 4회



Infer @fbinfer · 2월 6일

An analyzer as realistic and modular as Inferbo in about 5 weeks? I would not have believed it". Prof Kwangkeun Yi research.fb.com/inferbo-infer-...

🌐 영어 번역하기

FEBRUARY 6, 2017

Inferbo: Infer-based buffer overrun analyzer

by: Professor Kwangkeun Yi, Dept. of Computer Science and Engineering, Seoul National University



Sam Blackshear님이 링크를 공유했습니다.

2월 7일

FB ❤️ AbsInt: "[Infer's] extensible abstract interpretation framework is super-convenient to add a new kind of analysis like Inferbo, and its multi-lingual front end frees us from worrying about targeting varied source languages."



Facebook Academics

2월 7일 · 🌐

A great overview of Inferbo: Infer-based buffer overrun analyze developed by Facebook visiting professor Professor Kwangkeun Yi, from the Dept. of Computer Science and Engineering at Seoul National University.

예제

```
char * malloc_wrapper(int n) {  
    return malloc(n);  
}
```

Memory:

n |-> [s\$0, s\$1]

ret |-> (offset: [0, 0], size: [s\$0, s\$1])

예제

```
void set_i(int* arr, int index){  
    arr[index] = 0;  
}
```

Memory:

arr \mapsto (offset: [s\$4, s\$5], size: [s\$6, s\$7])

index \mapsto [s\$8, s\$9]

Safety Condition:

$[s\$4 + s\$8, s\$5 + s\$9] < [s\$6, s\$7]$

예제

```
char * maximum(int x) {  
    if (x < 9) return x;  
    else return 8  
}
```

Memory:

x \mapsto [s\$10, s\$11]

ret \mapsto [min(8, s\$10), 8]

예제

```
void interprocedural() {  
  char *arr = malloc_wrapper(9);  
  int i;  
  for (i = 0; i < 9; i+=1) {  
    set_i(arr, i); // safe  
    set_i(arr, i + 1); // alarm  
    set_i(arr, maximum(i+1)); //safe  
  }  
}
```

Summary of malloc_wrapper:

n |-> [s\$0, s\$1]

ret |-> (offset: [0, 0], size: [s\$0, s\$1])

Memory:

arr |-> (offset: [0, 0], size: [9, 9])

예제

```
void interprocedural() {  
  char *arr = malloc_wrapper(9);  
  int i;  
  for (i = 0; i < 9; i+=1) {  
    set_i(arr, i); // safe  
    set_i(arr, i + 1); // alarm  
    set_i(arr, maximum(i+1)); //safe  
  }  
}
```

Summary of set_i:

Memory:

arr \mapsto (offset: [s\$4, s\$5], size: [s\$6, s\$7])

index \mapsto [s\$8, s\$9]

Safety Condition:

$[s\$4 + s\$8, s\$5 + s\$9] < [s\$6, s\$7]$

Memory:

arr \mapsto (offset: [0, 0], size: [9, 9])

i \mapsto [0, 8]

Safety Condition

$[0 + 0, 0 + 8] < [9, 9]$

예제

```
void interprocedural() {  
  char *arr = malloc_wrapper(9);  
  int i;  
  for (i = 0; i < 9; i+=1) {  
    set_i(arr, i); // safe  
    set_i(arr, i + 1); // alarm  
    set_i(arr, maximum(i+1)); //safe  
  }  
}
```

Summary of set_i:

Memory:

arr \mapsto (offset: [s\$4, s\$5], size: [s\$6, s\$7])

index \mapsto [s\$8, s\$9]

Safety Condition:

$[s\$4 + s\$8, s\$5 + s\$9] < [s\$6, s\$7]$

Memory:

arr \mapsto (offset: [0, 0], size: [9, 9])

i \mapsto [1, 9]

Safety Condition

$[0 + 1, 0 + 9] < [9, 9]$

예제

```
void interprocedural() {  
    char *arr = malloc_wrapper(9);  
    int i;  
    for (i = 0; i < 9; i+=1) {  
        set_i(arr, i); // safe  
        set_i(arr, i + 1); // alarm  
        set_i(arr, maximum(i+1)); //safe  
    }  
}
```

Summary of maximum:

Memory:

x |-> [s\$10, s\$11]

ret |-> [min(8, s\$10), 8]

Memory:

arr |-> (offset: [0, 0], size: [9, 9])

i |-> [1, 9]

예제

```
void interprocedural() {  
    char *arr = malloc_wrapper(9);  
    int i;  
    for (i = 0; i < 9; i+=1) {  
        set_i(arr, i); // safe  
        set_i(arr, i + 1); // alarm  
        set_i(arr, maximum(i+1)); //safe  
    }  
}
```

Summary of maximum:

Memory:

x \mapsto [s\$10, s\$11]

ret \mapsto [min(8, s\$10), 8]

Memory:

arr \mapsto (offset: [0, 0], size: [9, 9])

i \mapsto [1, 9]

Safety Condition

$[0 + 1, 0 + 8] < [9, 9]$

실험 결과

- 오픈 소스 프로그램에 수동으로 버그 삽입후 테스트

Software	KLoC	#injected bugs	#true alarms	#false alarms	Notes
spell	2	7	7	0	caught all, no false alarms
unhtml-2.3.9	2	11	6	0	missed 5/11, no false alarms
spell++-2.0.2	7	7	6	0	missed 1/7, no false alarms
bc-1.06	16	25	24	0	missed 1/25, no false alarms
gzip-1.8	61	5	5	6	caught all, 6 false alarms

도메인

$$State = Node \rightarrow Mem$$

$$Mem = Loc \rightarrow Val$$

$$Val = \hat{\mathbb{Z}} \times 2^{AllocSite \times \hat{\mathbb{Z}} \times \hat{\mathbb{Z}}}$$

$$\hat{\mathbb{Z}} = \{ [l, u] \mid l, u \in Bound \} \cup \{ \perp \}$$

$$Bound = SymLinear \cup SymMinMax \cup \{ -\infty, +\infty \}$$

$$SymLinear = \{ c_0 + \sum c_i s_i \mid c_k \in \mathbb{Z}, s_i \in Symbols \}$$

$$SymMinMax = \{ \min(c, s) \mid c \in \mathbb{Z}, s \in Symbols \} \cup \{ \max(c, s) \mid c \in \mathbb{Z}, s \in Symbols \}$$

도메인

- Order

$$[l_1, u_1] \sqsubseteq [l_2, u_2] \quad \text{iff} \quad l_2 \hat{\leq} l_1 \wedge u_1 \hat{\leq} u_2$$

$$c_0 + \sum c_i s_i \hat{\leq} d_0 + \sum d_i s_i \quad \text{iff} \quad c_0 \leq d_0 \wedge (\forall i \geq 1. c_i = d_i)$$

$$\min(c, s) \hat{\leq} \min(d, s) \quad \text{iff} \quad c \leq d$$

$$\max(c, s) \hat{\leq} \max(c, s) \quad \text{iff} \quad c \leq d$$

$$\min(c, s) \hat{\leq} d_0 + d_1 s \quad \text{iff} \quad (c \leq d_0 \wedge d_1 = 0) \vee (d_0 = 0 \wedge d_1 = 1)$$

$$c_0 + c_1 s \hat{\leq} \max(d, s) \quad \text{iff} \quad (c_0 \leq d \wedge c_1 = 0) \vee (c_0 = 0 \wedge c_1 = 1)$$

$$\min(c, s_1) \hat{\leq} \max(d, s_2) \quad \text{iff} \quad c \leq d \vee s_1 = s_2$$

도메인

- Order

$$[l_1, u_1] \sqsubseteq [l_2, u_2] \quad \text{iff} \quad l_2 \hat{\leq} l_1 \wedge u_1 \hat{\leq} u_2$$

$$c_0 + \sum c_i s_i \hat{\leq} d_0 + \sum d_i s_i \quad \text{iff} \quad c_0 \leq d_0 \wedge (\forall i \geq 1. c_i = d_i)$$

$$\min(c, s) \leq \min(d, s) \quad \text{iff} \quad c \leq d$$

$$\max(c, s) \hat{\leq} \max(c, s) \quad \text{iff} \quad c \leq d$$

$$\min(c, s) \leq d_0 + d_1 s \quad \text{iff} \quad (c \leq d_0 \wedge d_1 = 0) \vee (d_0 = 0 \wedge d_1 = 1)$$

$$c_0 + c_1 s \hat{\leq} \max(d, s) \quad \text{iff} \quad (c_0 \leq d \wedge c_1 = 0) \vee (c_0 = 0 \wedge c_1 = 1)$$

$$\min(c, s_1) \hat{\leq} \max(d, s_2) \quad \text{iff} \quad c \leq d \vee s_1 = s_2$$

도메인

- Join

$$[l_1, u_1] \sqcup [l_2, u_2] = [\min(l_1, l_2), \max(u_1, u_2)]$$

$$\min(x, y) = \begin{cases} x & \text{if } x \hat{\leq} y \\ y & \text{if } y \hat{\leq} x \\ \min(c, s) & \text{if } x = c \wedge y = s \\ \min(c, s) & \text{if } y = c \wedge x = s \\ -\infty & \text{o.w.} \end{cases}$$

연산

$$[l_1, u_1] + [l_2, u_2] = [l_1 \hat{+}_l l_2, u_1 \hat{+}_u u_2]$$

$$x \hat{+}_l y = \begin{cases} (c_0 + d_0) + \Sigma(c_i + d_i)s_i & \text{if } x = c_0 + \Sigma c_i s_i \wedge y = d_0 + \Sigma d_i s_i \\ c + d & \text{if } x = c \wedge y = \max(d, s) \\ c + d & \text{if } x = \max(c, s) \wedge y = d \\ -\infty & \text{o.w.} \end{cases}$$

연산

$$[l_1, u_1] + [l_2, u_2] = [l_1 \hat{+}_l l_2, u_1 \hat{+}_u u_2]$$

$$x \hat{+}_l y = \begin{cases} (c_0 + d_0) + \Sigma(c_i + d_i)s_i & \text{if } x = c_0 + \Sigma c_i s_i \wedge y = d_0 + \Sigma d_i s_i \\ c + d & \text{if } x = c \wedge y = \max(d, s) \\ c + d & \text{if } x = \max(c, s) \wedge y = d \\ -\infty & \text{o.w.} \end{cases}$$

연산

$$prune : \hat{\mathbb{Z}} \times \hat{\mathbb{Z}} \rightarrow \hat{\mathbb{Z}}$$

$$prune_{\leq}([l_1, u_1], [l_2, u_2]) = \begin{cases} [l_1, u_2] & \text{if } u_1 = +\infty \\ [l_1, \min(c_0, d_0) + \sum c_i s_i] & \text{if } u_1 = c_0 + \sum c_i s_i \wedge u_2 = d_0 + \sum d_i s_i \\ & \wedge \forall i \geq 1. c_i = d_i \\ [l_1, \min(c, s)] & \text{if } (u_1 = c \wedge u_2 = s) \vee (u_1 = s \wedge u_2 = c) \\ [l_1, \min(\min(c, d), s)] & \text{if } (u_1 = c \wedge u_2 = \min(d, s)) \\ & \vee (u_1 = \min(c, s) \wedge u_2 = d) \\ & \vee (u_1 = \min(c, s) \wedge u_2 = \min(d, s)) \\ [l_1, u_1] & \text{o.w.} \end{cases}$$

연산

$$\text{prune} : \hat{\mathbb{Z}} \times \hat{\mathbb{Z}} \rightarrow \hat{\mathbb{Z}}$$

$$\text{prune}_{\leq}([l_1, u_1], [l_2, u_2]) = \begin{cases} [l_1, u_2] & \text{if } u_1 = +\infty \\ [l_1, \min(c_0, d_0) + \sum c_i s_i] & \text{if } u_1 = c_0 + \sum c_i s_i \wedge u_2 = d_0 + \sum d_i s_i \\ & \wedge \forall i > 1. c_i = d_i \\ [l_1, \min(c, s)] & \text{if } (u_1 = c \wedge u_2 = s) \vee (u_1 = s \wedge u_2 = c) \\ [l_1, \min(\min(c, d), s)] & \text{if } (u_1 = c \wedge u_2 = \min(d, s)) \\ & \vee (u_1 = \min(c, s) \wedge u_2 = d) \\ & \vee (u_1 = \min(c, s) \wedge u_2 = \min(d, s)) \\ [l_1, u_1] & \text{o.w.} \end{cases}$$

함수 서머리

- 정의

$$\begin{aligned}\sigma &\rightarrow \lambda[s_1, s_2]. \langle i, \kappa \rangle \\ \kappa &\rightarrow \text{True} \mid \kappa \wedge (i < i)\end{aligned}$$

- 예) 함수 f 의 서머리

$$\sigma_f = \lambda[s_1, s_2]. \langle i_f, \kappa_f \rangle$$

where

$$\begin{aligned}i_f &= \bigsqcup_{i \in \text{RetVal}_f} i \\ \kappa_f &= \bigwedge_{\langle i_{idx}, i_{size} \rangle \in \text{BufAccess}_f} (i_{idx} < i_{size}) \wedge \bigwedge_{g([l, u]) \in \text{CallSite}_f} (\sigma_g[l, u]).2\end{aligned}$$

함수 서머리

- 정의

$$\sigma \rightarrow \lambda[s_1, s_2]. \langle i, \kappa \rangle \quad \boxed{\text{<리턴값, 조건식>}}$$

$$\kappa \rightarrow \text{True} \mid \kappa \wedge (i < i)$$

- 예) 함수 f 의 서머리

$$\sigma_f = \lambda[s_1, s_2]. \langle i_f, \kappa_f \rangle$$

where

$$i_f = \bigsqcup_{i \in \text{RetVal}_f} i \quad \boxed{\text{모든 리턴값 포섭}}$$

$$\kappa_f = \bigwedge_{\langle i_{idx}, i_{size} \rangle \in \text{BufAccess}_f} (i_{idx} < i_{size}) \wedge \bigwedge_{g([l, u]) \in \text{CallSite}_f} (\sigma_g[l, u]).2$$

$\boxed{\text{f 의 조건식}}$

$\boxed{\text{모든 callee 의 조건식}}$

경험

- 잘 정리된 분석기 제작틀
 - multi-lingual & modular analysis framework
- OCaml 프로그래밍의 미학: 가독성 + 성능
- 전 세계와 함께 숨쉬는 즐거움 (공동 개발, 활발한 질답)

결론

- 개별분석 + 인터벌 도메인 기반 버퍼오버런 분석기
- 페이스북 내부 테스트를 통해 실용성 입증
- 계획: 여러 언어로 만든 프로그램에 테스트/튜닝 (C++, Java, Obj-C, etc)

THE FACEBOOK WALL

What's on your mind?



고맙습니다