

COMENIUS UNIVERSITY IN BRATISLAVA
FACULTY OF MATHEMATICS, PHYSICS AND INFORMATICS

TRIANGULATION OF IMPLICIT SURFACE WITH
SINGULARITIES
MASTER'S THESIS

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BC. KRISTÍNA KORECOVÁ

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FACULTY OF MATHEMATICS, PHYSICS AND INFORMATICS

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Department: Department of Algebra and Geometry
Supervisor: doc. RNDr. Pavel Chalmovianský, PhD.

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Bc. Kristína Korecová



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Introduction

In mathematics, implicit surface is a set of points where the implicit function $F : \mathbb{R}^3 \rightarrow \mathbb{R}$ is zero. Thanks to unification of constructive solid geometry modelling and implicit surfaces, we can now model and represent very complex surfaces using a single function.

In computer graphics, triangular mesh is the most common way of surface representation and visualization. Calculating the intersection of a ray and a triangular mesh is far easier task as calculating the intersection of a ray and an implicit surface.

In this thesis we follow up on our effort in bachelor's thesis [10] to present an algorithm for creating the triangular mesh of regular implicit surfaces.

First, we reimplement the algorithm to be more effective, by using advanced data structures for mesh representation and 3D point search.

Next, we extend this algorithm to include the triangulation of certain types of isolated singularities - ADE singularities. We analyze the geometry of these singularities and propose an approach to triangulate these singularities.

Lastly, we improve the adaptive technique, which changes the size of the triangles based on the local curvature of the surface.

Chapter 1

Theoretical background

1.1 Implicit surfaces

Implicit functions are a tool for surface representation and manipulation. In computer graphics, they can be used for modelling of complex surfaces using boolean operations, realistic animations, rendering and other.

Implicit functions do not define the surface explicitly, instead the surface is defined as a zero set of a function.

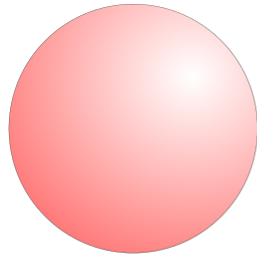
Definition 1 *Given a function $F : \mathbb{R}^3 \rightarrow \mathbb{R}$, one can define an implicit surface as a set of points $(x, y, z) \in \mathbb{R}^3$ that satisfy $F(x, y, z) = 0$.*

Some examples of implicit surfaces and their equations can be seen on the Figure 1.1.

Unit normal vector of the implicit surface in point (x_0, y_0, z_0) is normalized gradient of the implicit function in that point, provided it is non-zero.

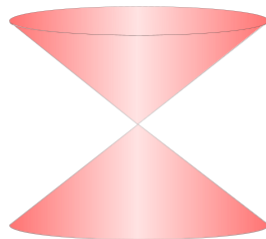
Definition 2 *Gradient vector of a function $F : \mathbb{R}^3 \rightarrow \mathbb{R}$ is defined as*

$$\nabla F(x, y, z) = \left(\frac{\partial F(x, y, z)}{\partial x}, \frac{\partial F(x, y, z)}{\partial y}, \frac{\partial F(x, y, z)}{\partial z} \right).$$



$$x^2 + y^2 + z^2 - 1 = 0$$

a) sphere



$$x^2 + y^2 - z^2 = 0$$

b) cone



$$z = 0$$

c) plane

Figure 1.1: Implicit surfaces with corresponding equations.

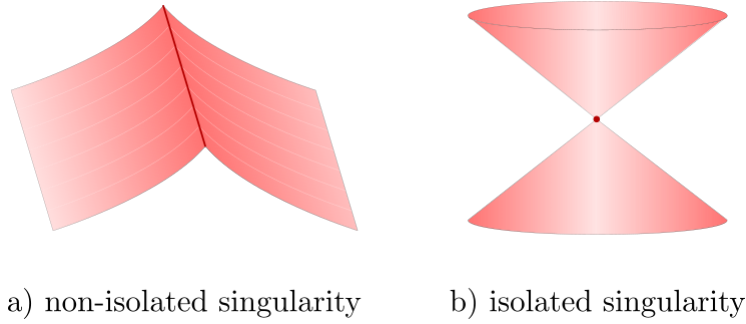


Figure 1.2: Isolated and non-isolated singularity.

If $\nabla F(x, y, z) \neq 0$, we can define the unit normal vector of F as a normalized gradient vector

$$N(F(x, y, z)) = \frac{\nabla F(x, y, z)}{\|\nabla F(x, y, z)\|}.$$

Points lying on the implicit surface can be classified as regular or singular based on the value of the gradient vector in that point.

Definition 3 Point $P = (x, y, z)$ lying on the implicit surface is said to be regular, if $\nabla F(x, y, z) \neq 0$. On the contrary, point P is said to be singular, if $\nabla F(x, y, z) = 0$.

Singular points can be further classified as isolated or non-isolated.

Definition 4 Singular point P of a surface S is said to be isolated, if there exists an open ball $B_\varepsilon(P)$, which does not contain any other singular point of S . Singular point P is said to be non-isolated if it is not isolated.

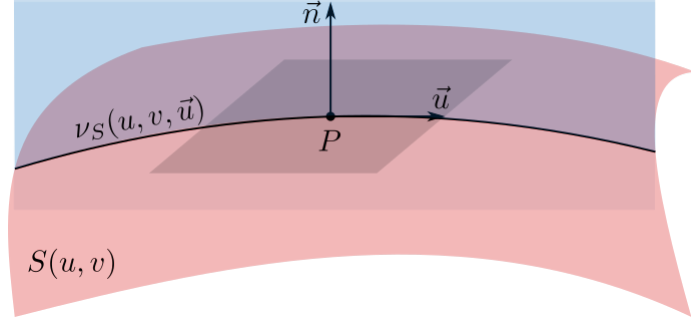
On the Figure 1.2, one can see an example of isolated and non-isolated singularities.

Curvature of a surface

Curvature is a fundamental concept in differential geometry of curves and surfaces. In case of curves, curvature is a measure of how much does the curve differ from a straight line. It is defined as the inverse of the radius of the osculating circle, which is the second order approximation of the curve.

For surfaces, curvature is a measure of how much does the surface differ from a plane. The definition of the curvature of a surface is not as straightforward as in the case of curves. The curvature of a curve on a surface depends on the choice of the tangent direction of a measured point.

The idea of measuring the curvature of a surface has a long history in mathematics. One of the first contributors was a mathematician Carl Friedrich Gauss, who developed

Figure 1.3: Normal cut of the parametric surface $S(u, v)$.

the idea of the Gaussian curvature of surfaces. In this subsection, we are drawing from the summary presented by Tiago Novello et al.[13].

Normal curvature of the surface

Let S be a parametric surface

$$S : D \subseteq \mathbb{R}^2 \rightarrow \mathbb{R}^3$$

$$S(u, v) = (S_x(u, v), S_y(u, v), S_z(u, v)), \quad (u, v) \in \mathbb{R}^2$$

Let the unit normal vector of the surface S at the point $S(u, v)$ be $\vec{n}(u, v)$.

We define the normal curvature of the surface as a function of the location of the point on the surface given by parameters u and v and the unit tangent vector in that point \vec{u} .

Definition 5 *Normal cut of a surface S in the regular point P in the direction of the unit tangent vector \vec{u} is defined as an intersection of the surface S and a plane given by the vectors \vec{u} and $\vec{n}(u, v)$ shifted to the point $S(u, v)$.*

The visualisation of the normal cut is shown on the Figure 1.3. It is clear, that the normal cut is a plane curve lying on the surface, we denote this normal cut as $\nu_S(u, v, \vec{u})$.

Definition 6 *Oriented normal curvature of the surface at the regular point P in the direction of the unit tangent vector \vec{u} is defined as the oriented curvature of the normal cut $\nu_S(u, v, \vec{u})$. Non-oriented normal curvature is defined as an absolute value of the oriented normal curvature.*

Definition 7 *Minimal and maximal curvature at the point $P = S(u, v)$ are defined as*

$$\kappa_{min}(u, v) = \min_{\vec{u} \in T_P S} \nu_S(u, v, \vec{u}),$$

$$\kappa_{max}(u, v) = \max_{\vec{u} \in T_P S} \nu_S(u, v, \vec{u}),$$

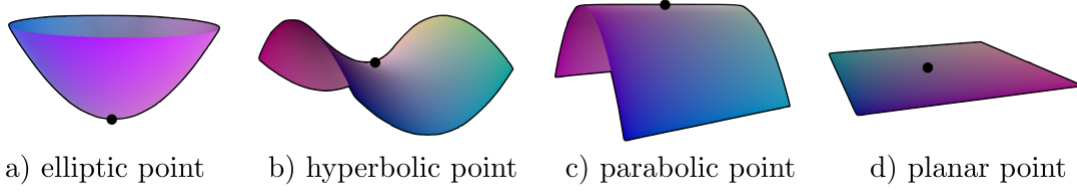


Figure 1.4: Categorization of the surface points based on Gaussian curvature.

where $T_P S$ is a tangent plane of the surface S in the point P .

Minimal and maximal curvatures are called principal.

Definition 8 Gaussian curvature at a point $S(u, v)$ is defined as a product of principal curvatures at that point, i.e.

$$\kappa_G(u, v) = \kappa_{min}(u, v)\kappa_{max}(u, v).$$

Gaussian curvature describes the shape of the surface in the local neighborhood of the point. Points of the surface S , where the Gaussian curvature is positive, are called elliptic. Points of the surface S , where the Gaussian curvature is negative are called hyperbolic. Points of the surface S , where only one of $\kappa_{min}, \kappa_{max}$ is zero, are called parabolic and points, where both κ_{min} and κ_{max} are zero, are called planar. The shape of the surface in the local neighborhoods of the points, visualized on the Figure 1.4, is as follows:

- elliptic points \longrightarrow surface is curved like an ellipsoid,
- hyperbolic points \longrightarrow surface is curved like a saddle,
- parabolic points \longrightarrow surface is curved like a parabolic cylinder,
- planar points \longrightarrow surface is flat.

Gaussian curvature is an intrinsic property, which means that it is also independent of the isometric image of the surface.

Definition 9 Mean curvature of a surface S at a point $P = S(u, v)$ is defined as an arithmetic mean of corresponding principal curvatures:

$$\kappa_M(u, v) = \frac{\kappa_{min}(u, v) + \kappa_{max}(u, v)}{2}.$$

Minimal, maximal, Gaussian and mean curvature are visualized in the Figure 1.5, where blue color indicates positive value of the corresponding visualized curvature, red color indicates negative value of the corresponding visualized curvature and white color indicates zero value of the corresponding visualized curvature.

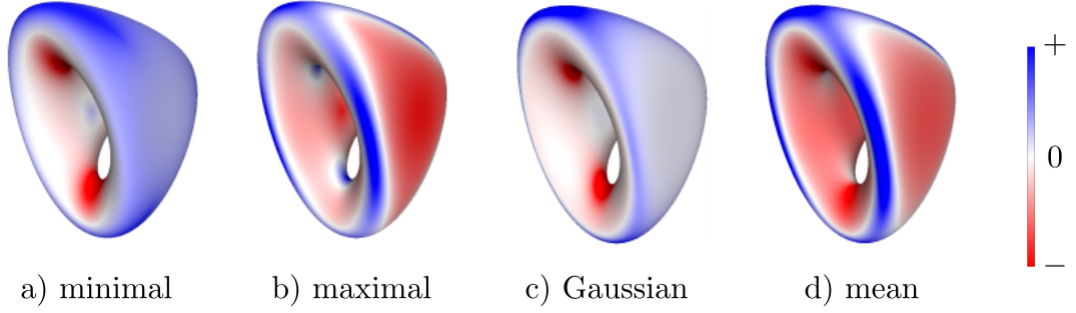


Figure 1.5: Visualisation of minimal, maximal, Gaussian and mean curvatures of the double-torus [13]. Blue/white/red color correspond to positive/zero/negative value of the corresponding curvature.

Curvature formulas for implicit surface

A version of curvature formulas for implicit surfaces appeared in [16] and were reformulated, summarized and proved by Ron Goldman [6]. In this subsection, we point out these formulas.

Let $F : \mathbb{R}^3 \rightarrow R$ be an implicit function which defines surface by the equation $F(x, y, z) = 0$. Let us denote $F_t = \frac{\partial F}{\partial t}$ and $F_{ts} = \frac{\partial^2 F}{\partial t \partial s}$. Hessian matrix – the matrix of the second order derivatives is defined as

$$H(F) = \begin{pmatrix} F_{xx} & F_{xy} & F_{xz} \\ F_{yx} & F_{yy} & F_{yz} \\ F_{zx} & F_{zy} & F_{zz} \end{pmatrix},$$

and the adjoint of the Hessian is defined as a transpose of its cofactor matrix.

$$H^*(F) = \begin{pmatrix} \begin{vmatrix} F_{yy} & F_{yz} \\ F_{zy} & F_{zz} \end{vmatrix} & -\begin{vmatrix} F_{xy} & F_{xz} \\ F_{zy} & F_{zz} \end{vmatrix} & \begin{vmatrix} F_{xy} & F_{xz} \\ F_{yy} & F_{yz} \end{vmatrix} \\ -\begin{vmatrix} F_{yx} & F_{yz} \\ F_{zx} & F_{zz} \end{vmatrix} & \begin{vmatrix} F_{xx} & F_{xz} \\ F_{zx} & F_{zz} \end{vmatrix} & -\begin{vmatrix} F_{xx} & F_{xz} \\ F_{yx} & F_{yz} \end{vmatrix} \\ \begin{vmatrix} F_{yx} & F_{yy} \\ F_{zx} & F_{zy} \end{vmatrix} & -\begin{vmatrix} F_{xx} & F_{xy} \\ F_{zx} & F_{zy} \end{vmatrix} & \begin{vmatrix} F_{xx} & F_{xy} \\ F_{yx} & F_{yy} \end{vmatrix} \end{pmatrix}.$$

Now, one can find the formulas of Gaussian, mean, minimal and maximal curvature for an implicit surface $F = 0$ at a regular point.

Gaussian curvature is given by

$$\kappa_G = \frac{\nabla F \cdot H^*(F) \cdot \nabla F^T}{|\nabla F|^4}.$$

Mean curvature of the implicit surface defined by function F is given by

$$\kappa_M = \frac{\nabla F \cdot H^*(F) \cdot \nabla F^T - |\nabla F|^2 \text{Tr}(H)}{2|\nabla F|^3}.$$

The principal curvatures κ_{min} and κ_{max} can be calculated from Gaussian curvature and mean curvature as

$$\kappa_{min}, \kappa_{max} = \kappa_M \pm \sqrt{\kappa_M^2 - \kappa_G}.$$

1.2 ADE singularities

ADE singularities, also referred to as du Val singularities are a specific class of simple, isolated surface singularities. They were classified by Arnold's [1] according to ADE classification [7] based on correspondence of these singularities to simply laced Dynkin diagrams [4]. We know infinitely many A singularities – A_1, A_2, \dots , infinitely many D singularities – D_4, D_5, \dots and three E singularities – E_6, E_7 and E_8 . ADE singularities are specified by their normal forms. When working in complex space, each singularity has a single normal form:

- A_n $F(x, y, z) = x^{n+1} + y^2 + z^2,$
- D_n $F(x, y, z) = yx^2 + y^{n-1} + z^2,$
- E_6 $F(x, y, z) = x^3 + y^4 + z^2,$
- E_7 $F(x, y, z) = x^3 + xy^3 + z^2,$
- E_8 $F(x, y, z) = x^3 + y^5 + z^2.$

Each ADE singularity on a surface can be locally expressed by their normal form.

In the real case, changing the signs in these equations produces different surfaces and therefore, ADE singularities can be further classified by their signature.

Definition 10 *Let real surface singularities based on their signature be denoted as follows:*

- $A_{n\pm\pm}$ $F(x, y, z) = x^{n+1} \pm y^2 \pm z^2,$
- $D_{n\pm\pm}$ $F(x, y, z) = yx^2 \pm y^{n-1} \pm z^2,$
- $E_{6\pm\pm}$ $F(x, y, z) = x^3 \pm y^4 \pm z^2,$
- $E_{7\pm\pm}$ $F(x, y, z) = x^3 \pm xy^3 \pm z^2,$
- $E_{8\pm\pm}$ $F(x, y, z) = x^3 \pm y^5 \pm z^2.$

The most common example of a surface with ADE singularity is an ordinary cone. Given as the zero set of the function $F(x, y, z) = x^2 - y^2 - z^2$, cone has a singular point $P = (0, 0, 0)$. This singular point is an example of A_{1--} singularity.

Correspondence between $SO(3, \mathbb{R})$ group and ADE singularities

$SO(3, \mathbb{R})$ is special orthogonal group over the field of real numbers in three dimensions. It is also called 3D rotation group, as it is a group of all rotations around axes passing through the origin in \mathbb{R}^3 .

Definition 11 $SO(3, \mathbb{R})$ is a group of 3×3 orthogonal matrices of real numbers with determinant 1.

$$SO(3, \mathbb{R}) = \left\{ A \in \mathbb{R}^{3 \times 3} \mid AA^T = I, \det(A) = 1 \right\}.$$

Simply laced Dynkin diagrams correspond to all finite subgroups of $SO(3, \mathbb{R})$. Finite subgroups of $SO(3, \mathbb{R})$ are the rotational symmetry groups of

- pyramid with n vertices (cyclic subgroup \overline{C}_n) – Figure 1.6 a),
- double pyramid with n vertices (dihedral subgroup \overline{D}_n) – Figure 1.6 b),
- platonic solids:
 - tetrahedron (tetrahedral subgroup \overline{T}) – Figure 1.6 c),
 - octahedron (octahedral subgroup \overline{O}) – Figure 1.6 d),
 - icosahedron (icosahedral subgroup \overline{I}) – Figure 1.6 e).

The correspondence is as follows:

- $A_n \iff \overline{C}_{n+1}$,
- $D_n \iff \overline{D}_{n+2}$,
- $E_6 \iff \overline{T}$,
- $E_7 \iff \overline{O}$,
- $E_8 \iff \overline{I}$.

The conclusion is that ADE singularities correspond to finite subgroups of $SO(3, \mathbb{R})$, which represent certain types of symmetries in \mathbb{R}^3 .

1.3 CSG modelling for implicit surfaces

In this section we will discuss how constructive solid geometry can be used for modelling complex implicit surfaces. First major publication regarding constructive solid geometry was published in 1977 by Requicha and Voelcker [15]. More detailed mathematical foundations were published a year later, in 1978 by Requicha and Tilove [14].

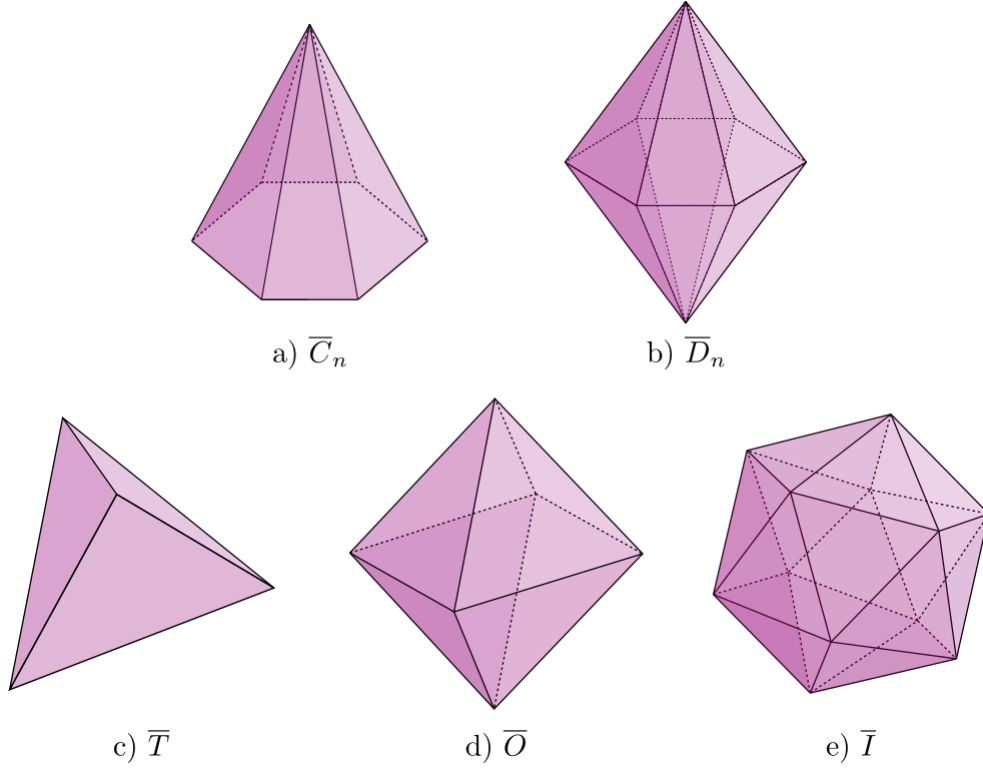


Figure 1.6: Shapes representing the $SO(3, \mathbb{R})$ group.

Constructive solid geometry (CSG)

Constructive solid geometry is a technique used for modelling complex geometric objects using boolean operations on sets of points – union, intersection and difference. It gained popularity in 1980s as a powerful tool to create complex shapes from sets of geometrical primitives, such as cylinders, spheres and cones. The resulting model can be represented as a CSG tree, which contains geometrical primitives in its leaves and boolean operations in its internal nodes [5]. An example of such CSG tree can be seen on the Figure 1.7.

CSG for implicit surfaces

As the implicit surface divides space into inside ($F(x, y, z) < 0$) and outside ($F(x, y, z) > 0$), one can easily combine the idea of implicit surfaces and CSG modelling. Given two implicit surfaces represented by implicit functions F and G , the surface of the intersection of the interiors can be defined by implicit function $H = \min(F, G)$. Similarly, the surface of the union of the interiors can be defined by implicit function $H = \max(F, G)$ and lastly, the surface of the difference of the interiors can be defined by implicit function $H = \max(F, -G)$. Two dimensional example of this idea can be seen on the Figure 1.8.

Theorem 1 *The minimum function $\min(F, G)$ of two functions $F : \mathbb{R}^3 \rightarrow \mathbb{R}$ and*

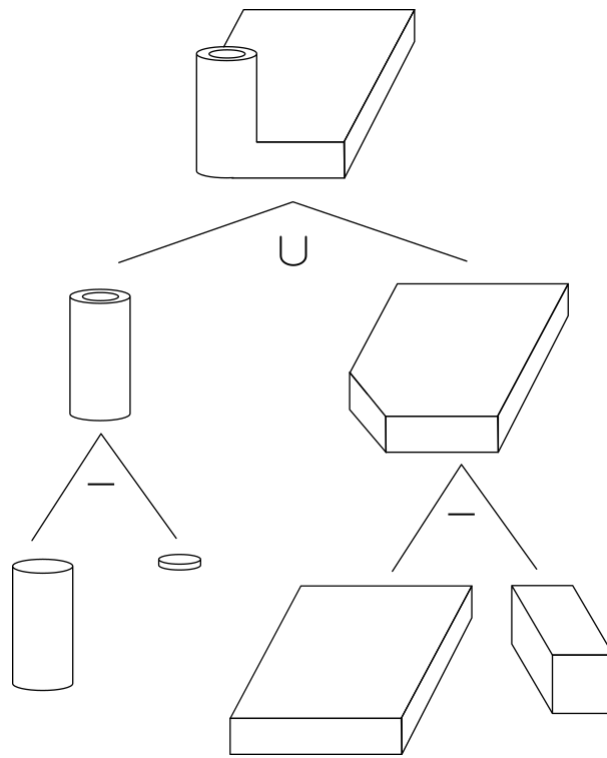


Figure 1.7: Example of CSG tree [5].

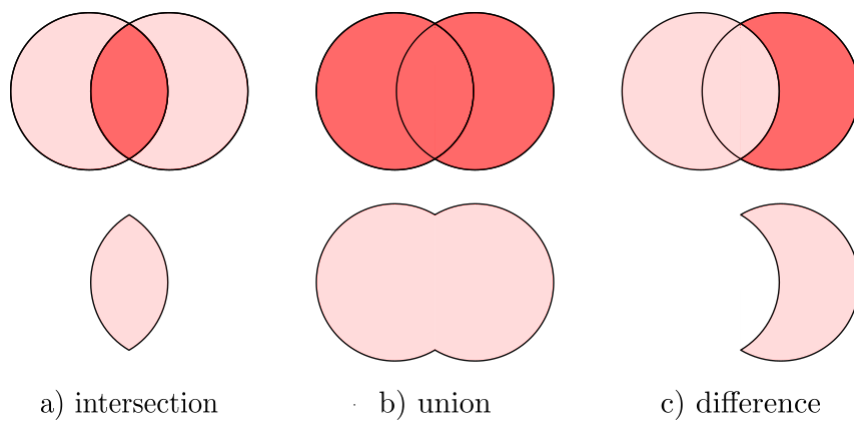


Figure 1.8: Boolean operations on implicit curves.

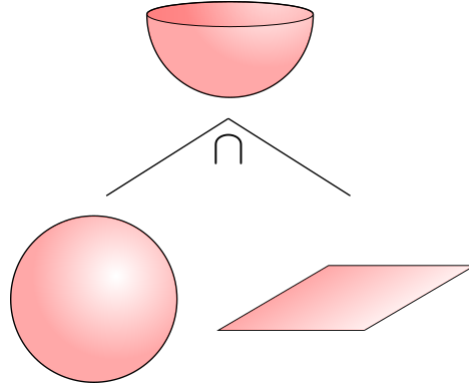


Figure 1.9: Intersection of a sphere and a plane.

$G : \mathbb{R}^3 \rightarrow \mathbb{R}$ is defined as

$$\min(F, G) = F + G - \sqrt{F^2 + G^2}$$

. The maximum function $\max(F, G)$ of two functions $F : \mathbb{R}^3 \rightarrow \mathbb{R}$ and $G : \mathbb{R}^3 \rightarrow \mathbb{R}$ is defined as

$$\max(F, G) = F + G + \sqrt{F^2 + G^2}$$

.

These formulas allow us to model implicit surfaces which are the result of performing finite number of operations union, intersection and difference on arbitrary implicit functions.

An easy example is creating an implicit equation representing a half of a sphere. Let us have an implicit function $F(x, y, z) = x^2 + y^2 + z^2 - 1$, which represents a unit sphere and another implicit function $G(x, y, z) = z$ which represents the xy plane. The minimum function of these two functions

$$\min(F, G) = x^2 + y^2 + z^2 - 1 + z - \sqrt{(x^2 + y^2 + z^2 - 1)^2 + z^2}$$

represents the surface of a half sphere. The visualisation of these surfaces can be seen on the Figure 1.9.

1.4 Non-isolated surface singularities

As we already mentioned, non-isolated singular point of the surface S is the singular point, for which every open ball $B_\epsilon(P)$ contains other singular points of the surface S . Non-isolated singular points are also called singular curves.

In the CSG modelling, singular curves most commonly arise as a result of intersection, union and difference operations on the interiors of these surfaces. When performing the intersection, union or difference operation on the interiors of two regular

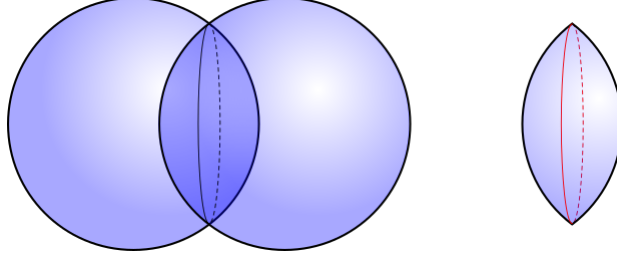


Figure 1.10: Intersection of a two spheres.

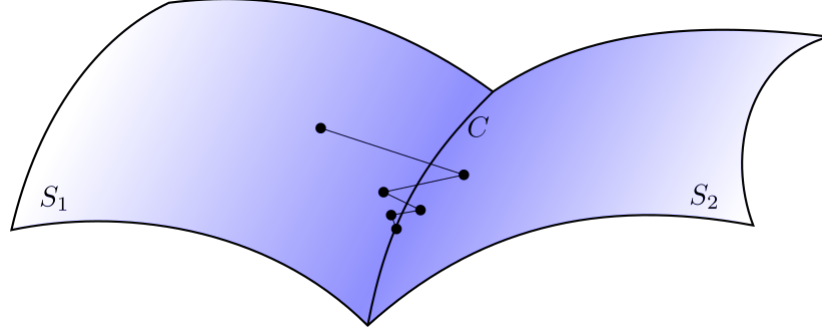


Figure 1.11: Projecting a point to the implicit curve.

surfaces, the singular curve is the curve given by the intersection of these surfaces. An example of an intersection of two spherical interiors is shown on the Figure 1.10, the resulting singular curve given by an intersection of the two spheres is displayed by the red color.

Projecting point to the implicit curve

Implicit curve C is given by two implicit equations - the equations of two surfaces which intersect in the curve C . To project a point close to the surface to the given surface, we use the method where we project the point in the direction of the normal vector using the iterative Newton-Raphson method. An extension of this method can be used to project a point to the implicit curve, by alternately projecting to the two surfaces until the point lies on both with required precision. A visualization of this approach can be seen on the image 1.11. We denote the projection function $proj_C$.

Differential geometry of implicit curves

An implicit curve C is given by two implicit equations $F_1, F_2 : \mathbb{R}^3 \rightarrow \mathbb{R}$ representing two implicit surfaces S_1 and S_2 , respectively which are intersecting in the implicit curve C . The unit tangent vector \vec{t}_C of the curve C in the point $P \in C$ is given by the cross product of unit normal vectors of the surfaces S_1 and S_2

$$\vec{t}_C(P) = \vec{n}_{S_1}(P) \times \vec{n}_{S_2}(P).$$

Chapter 2

Our contributions

2.1 Triangulation adaptive to the local curvature

As we explained in the beginning of chapter 1.1, curvature of the surface is a measure of how much the surface bends.

The triangulation of the surface should be accurate enough, but also memory efficient. This can be achieved by creating a triangulation which is locally adaptive to the curvature of the surface. Therefore having smaller triangles in the places where the surface is curved and having bigger triangles where surface is flatter.

In this section we present our implementation of the triangulation adaptive to the local curvature.

In the original algorithm, the height of the triangle which is projected to the surface is set to the constant value $\frac{\sqrt{3}}{2}e$, where e is the required length of the side of the triangle. To achieve the adaptivity of the triangles size, we set the height of the triangle to depend on the curvature in the given point, as shown in the Figure 2.1.

To identify the curved areas, we decided to use the maximal curvature. As the minimal and maximal curvature are both signed, we want to identify the areas depending on the absolute value of these curvatures.

We do not allow arbitrary height of the triangle to avoid edge-cases. We decided to restrict the allowed height to $\frac{1}{4}\frac{\sqrt{3}}{2}e > h > 4\frac{\sqrt{3}}{2}e$. We create a variable m -multiplier, depending on κ_T and set the height of the new triangle to $h = m\frac{\sqrt{3}}{2}e$.

As κ_T has values in range $\langle 0, \infty \rangle$.

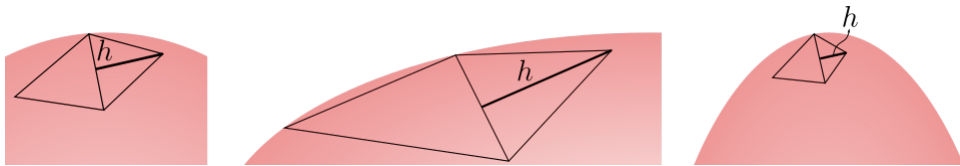


Figure 2.1: Adaptive height of the new triangle.

Definition 12 *Let us define triangulation curvature of the surface S in the point $P = S(u, v)$ as $\kappa_T(u, v) = \max(|\kappa_{\min}|, |\kappa_{\max}|)$.*

NOT FINISHED YET

2.2 Triangulation of ADE singularities

Analysis of the geometry of ADE singularities

ADE singularities are simple, isolated surface singularities, which can be expressed by corresponding implicit equations.

We already know, that A_{1--} singularity is locally represented as a cone. In this section we discuss geometric structure of other ADE surface singularities.

Definition 13 *(TODO rewrite) Let us define branch of ADE singularity as the part of the surface, which is connected to the rest only by the singular point.*

For our needs, we pick one triangulation vector for each branch of each ADE singularity. This triangulation vector is normalized vector either in the direction of rotation symmetry axis or an intersection of reflection symmetry planes of the corresponding branch. If the branch has only one reflection symmetry plane, the triangulation vector is picked to lie in the reflection symmetry plane.

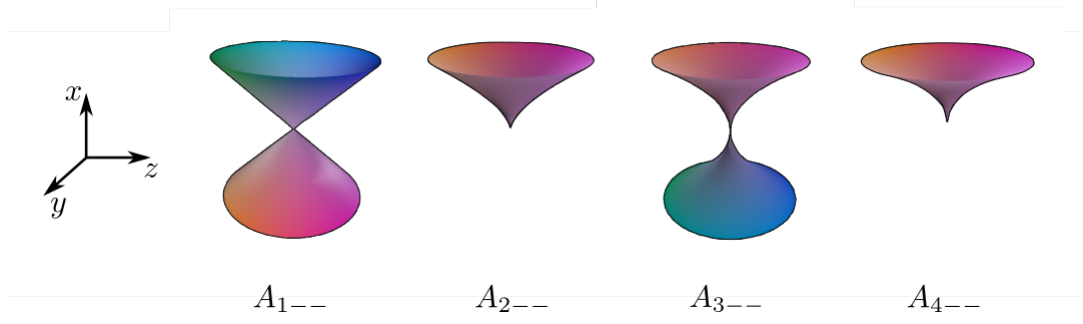
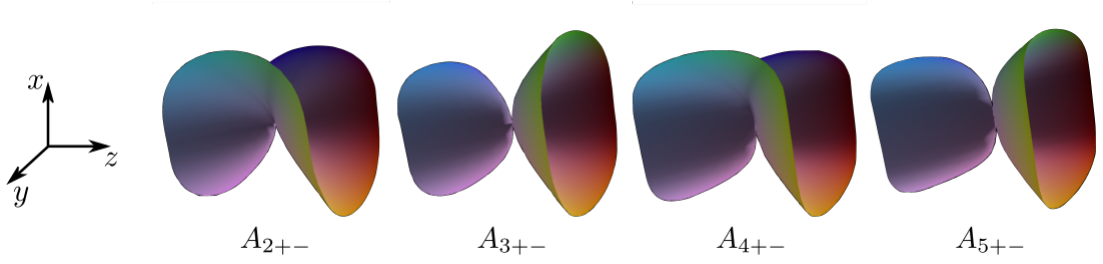
In the general case, triangulation vectors serve us as a partial information about the orientation of a singularity with respect to its normal form.

A_n singularities

As we can see from the equations $F(x, y, z) = x^{n+1} \pm y^2 \pm z^2$, A_{n-+} singularities are just rotated A_{n+-} singularities and A_{n++} singularities are a single point if n is odd and reflected A_{n--} singularities if n is even. We therefore only discuss geometry of A_{n--} and A_{n+-} singularities.

A_{n--} singularities are topologically equivalent to a cone if n is odd, therefore they have two branches. If n is even, they are topologically equivalent to a half cone or a plane, therefore they have a single branch. As n gets bigger, the tip of the cone gets sharper. As A_{n--} singularities are rotationally symmetrical, we pick the direction of axis of symmetry as triangulation vector. For a normal form, the triangulation vectors are $(1, 0, 0)$ (and $(-1, 0, 0)$ if n is odd). First four A_{n--} singularities can be seen on the Figure 2.2.

A_{n+-} singularities are topologically equivalent to a cone if n is odd, therefore they have two branches. In the contrary with the previous singularities, as n gets bigger, the tip of the cone gets less sharp and flatter. Branches of these singularities have

Figure 2.2: A_{n--} singularities. [12]Figure 2.3: A_{n+-} singularities. [12]

reflection symmetry planes $x = 0$ and $y = 0$, therefore we pick the vectors $(0, 0, 1)$ and $(0, 0, -1)$ as the triangulation vectors.

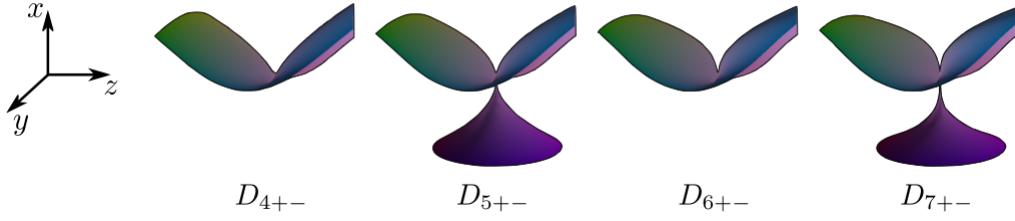
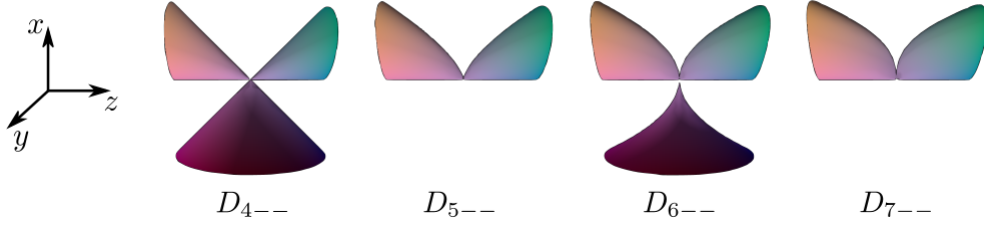
If n is even, A_{n+-} singularities are topologically equivalent to a plane with shape similar to hyperbolic paraboloid, therefore they have a single branch. First four A_{n+-} singularities can be seen on the Figure 2.3. For this case, we pick the vector $(1, 0, 0)$ as a triangulation vector as these singularities have reflection symmetry planes $y = 0$ and $z = 0$.

D_n singularities

Given by equations $F(x, y, z) = yx^2 \pm y^{n-1} \pm z^2$, we consider 8 categories. For given sign combination and parity of n , the singularities are topologically equivalent, with sharper(or flatter) features around the singularities for increasing value of n similar to A_n singularities.

We can therefore say that D_n singularities can be classified into 8 categories locally represented by the following equations:

- D_{4++} $yx^2 + y^3 + z^2$
- D_{5++} $yx^2 + y^4 + z^2$
- D_{4+-} $yx^2 + y^3 - z^2$
- D_{5+-} $yx^2 + y^4 - z^2$

Figure 2.4: D_{n+-} singularities. [12]Figure 2.5: D_{n--} singularities. [12]

- $D_{4-+} \quad yx^2 - y^3 + z^2$
- $D_{5-+} \quad yx^2 - y^4 + z^2$
- $D_{4--} \quad yx^2 - y^3 - z^2$
- $D_{5--} \quad yx^2 - y^4 - z^2.$

Now we look at some equivalences between these 8 categories. D_{4++} singularity is reflected D_{4+-} singularity. D_{5++} singularity is reflected D_{5--} singularity. D_{5-+} singularity is reflected D_{5+-} singularity. D_{4-+} singularity is reflected D_{4--} singularity.

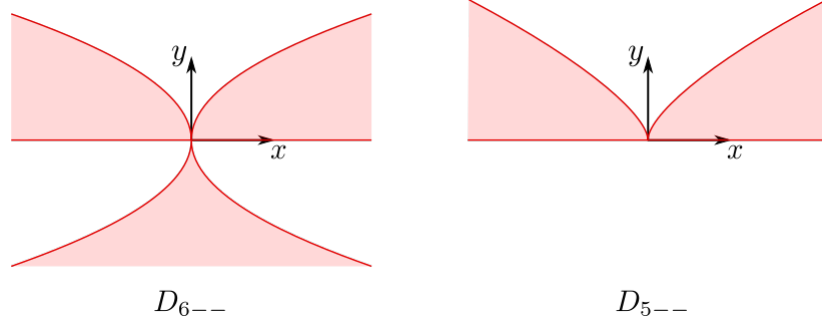
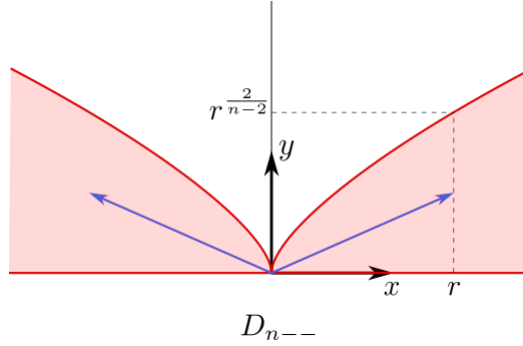
We therefore only analyze geometry of D_{n+-} singularities and D_{n--} singularities.

D_{n+-} singularities are topologically equivalent to a plane when n is even and to a cone when n is odd. Again, as n gets bigger, the features around singularities get sharper. Symmetry planes of these singularities are $x = 0$ and $z = 0$, therefore we pick $(0, 1, 0)$ (and $(0, -1, 0)$ when n is odd) as triangulation vectors. First four D_{n+-} singularities can be seen on the Figure 2.4.

D_{n--} singularities are topologically equivalent to a cone when n is odd and to a 3 halfcones connected in the singular point when n is even. First four D_{n--} singularities can be seen on the Figure 2.5.

Symmetry plane for all branches of these singularities is $z = 0$. the intersection of the surface and plane $z = 0$ is displayed on the Figure 2.6.

For D_{n--} singularity, the intersections of the two branches where $y \geq 0$ are bounded by curves $y = 0$ and $x^2 = y^{n-2}$. For given r , we pick the triangulation vectors as

Figure 2.6: Intersection of D_{n--} singularities with plane $z = 0$.Figure 2.7: Triangulation vectors for two branches of D_{n--} singularities.

$(r, \frac{1}{2}r^{\frac{2}{n-2}}, 0)$ and $(-r, \frac{1}{2}r^{\frac{2}{n-2}}, 0)$. The resulting vectors are displayed on the Figure 2.7 by blue arrow. Parameter r is changed based on the length of the edge of triangulation triangle. displayed on the Figure 2.7 by blue arrow. Parameter r is changed based on the length of the edge of triangulation triangle.

The third branch where $y \leq 0$ has another plane of symmetry $x = 0$, therefore triangulation vector for this branch is chosen as $(0, -1, 0)$.

E_6, E_7 and E_8 singularities

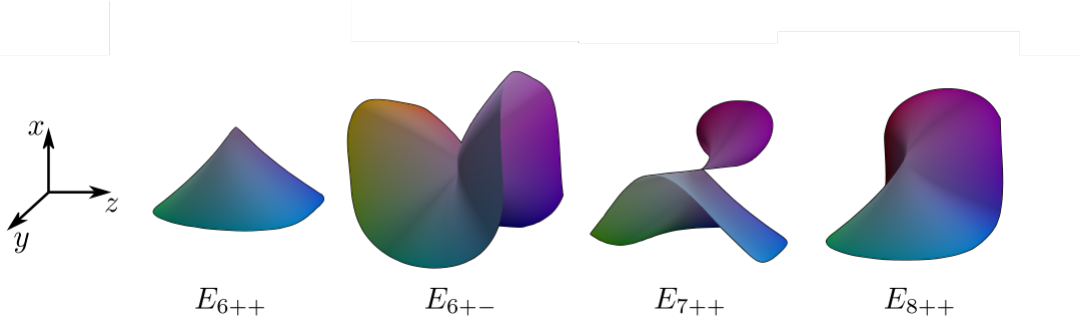
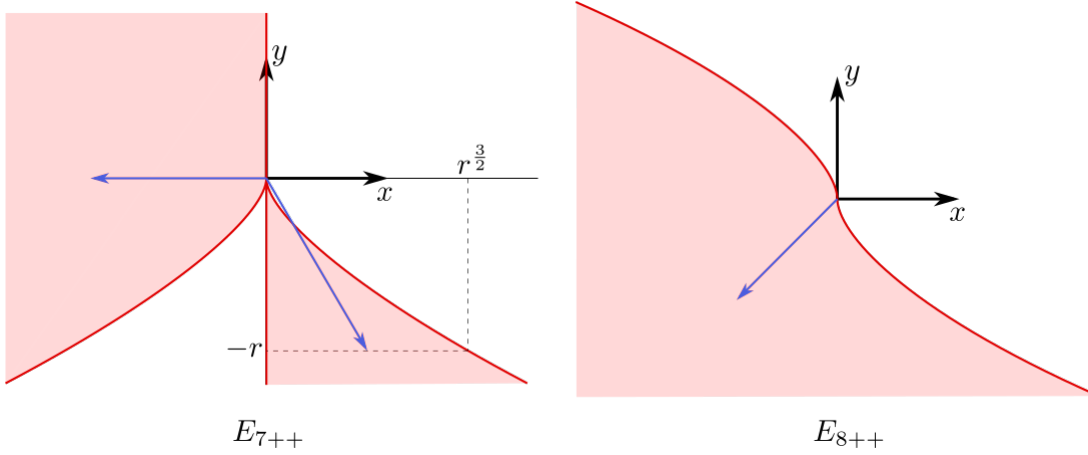
Given by equations $F(x, y, z) = x^3 \pm y^4 \pm z^2$, $F(x, y, z) = x^3 \pm xy^3 \pm z^2$ and $F(x, y, z) = x^3 \pm y^5 \pm z^2$, we can see the following equivalences: E_{6++} singularity is reflected E_{6--} singularity. E_{6+-} singularity is reflected E_{6-+} singularity. E_{7+-} , E_{7-+} and E_{7--} are all reflected E_{7++} singularity. E_{8+-} , E_{8-+} and E_{8--} are all reflected E_{8++} singularity.

We only analyze geometry of E_{6++} , E_{6+-} , E_{7++} and E_{8++} singularities. These singularities are displayed on the the Figure 2.8.

Both E_{6++} and E_{6+-} are topologically equivalent to a plane, thus they each have only one branch. The planes of symmetry of both of these branches are $y = 0$ and $z = 0$, therefore we pick $(-1, 0, 0)$ as the triangulation vector.

E_{7++} singularity is topologically equivalent to a cone, therefore it has two branches. The plane of symmetry of this singularity is $z = 0$.

E_{8++} singularity is also topologically equivalent to a plane, therefore it has only

Figure 2.8: E_n singularities. [12]Figure 2.9: Intersection of E_{7++} and E_{8++} singularities with plane $z = 0$.

one branch. This branch has only one plane of symmetry $z = 0$.

We again look at the intersection of the surfaces with the plane of symmetry, this is displayed on the Figure 2.9.

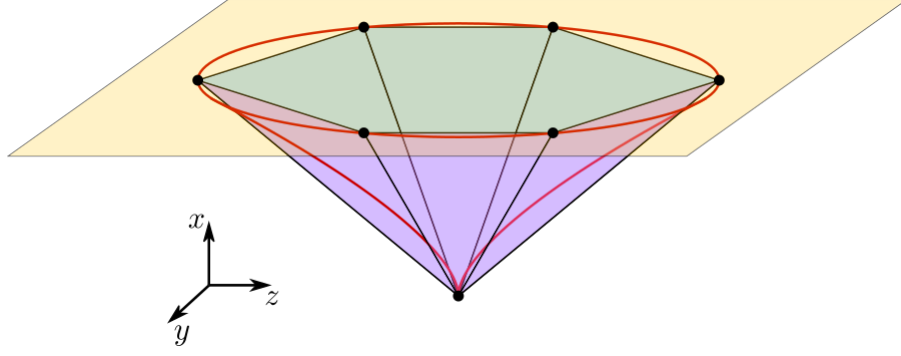
For E_{7++} singularity, we pick $(-1, 0, 0)$ and $(\frac{1}{2}r^{\frac{3}{2}}, -r, 0)$ as triangulation vectors. For E_{8++} singularity, we pick $(-1, -1, 0)$ as a triangulation vector. These vectors are displayed on the Figure 2.9 as blue arrows.

Analytical calculation of local triangulation of some ADE singularities

For given edge size e , we want to calculate the local triangulation of ADE singularities, such that edges on the border of the local triangulation have length e .

A_{n--} singularities

For A_{n--} singularities, we create a disc of six isosceles triangles with vertex in the singular point. The bases of these triangles create regular hexagon in the plane P parallel to the plane $x = 0$, as showed on the Figure 2.10. Given by equation $x^{n+1} -$

Figure 2.10: Triangulation of A_{n--} singularity.

$y^2 - z^2 = 0$, we find the distance of the plane P from the plane $x = 0$ for the given length e of the sides of the hexagon.

Let e be the length of the side of the hexagon, then the circumscribed circle has radius e . This circle is identical with the intersection of the surface and the plane $x = h$. The equation of the intersecting circle is $y^2 + z^2 = h^{n+1}$ therefore, the radius can be also expressed as $r = h^{\frac{n+1}{2}}$, which emerges $h = e^{\frac{2}{n+1}}$. Knowing the distance of the plane, one can easily calculate the length of the arms of the triangles using Pythagorean theorem:

$$a^2 = h^2 + e^2 \implies a = \sqrt{e^{\frac{4}{n+1}} + e^2}$$

D_n singularities

Some D_n singularities have branches with elliptical intersection with a plane parallel to the plane $y = 0$. As ellipses have two axes of symmetry, we create eight triangles with apex in the singular point for these branches. The other points of the triangles lie on the ellipse and they have the same length of the base.

Let us have an ellipse E with semi-major axis a , semi-minor axis b and the center in the point $(0, 0)$.

$$E : \frac{x^2}{a^2} + \frac{y^2}{b^2} = 1$$

As displayed on the Figure 2.11, we pick the leftmost, the rightmost, the top and the bottom points. As shown on the figure 2.11, the coordinates of these points are $P_1 = (a, 0)$, $P_2 = (0, b)$, $P_3 = (-a, 0)$, $P_4 = (0, -b)$. Then we can calculate the point P on ellipse equidistant from points P_1 and P_2 . We calculate this point by taking the point P_{12} in the middle of a line segment P_1P_2 .

$$P_{12} = \frac{1}{2}(a, b)$$

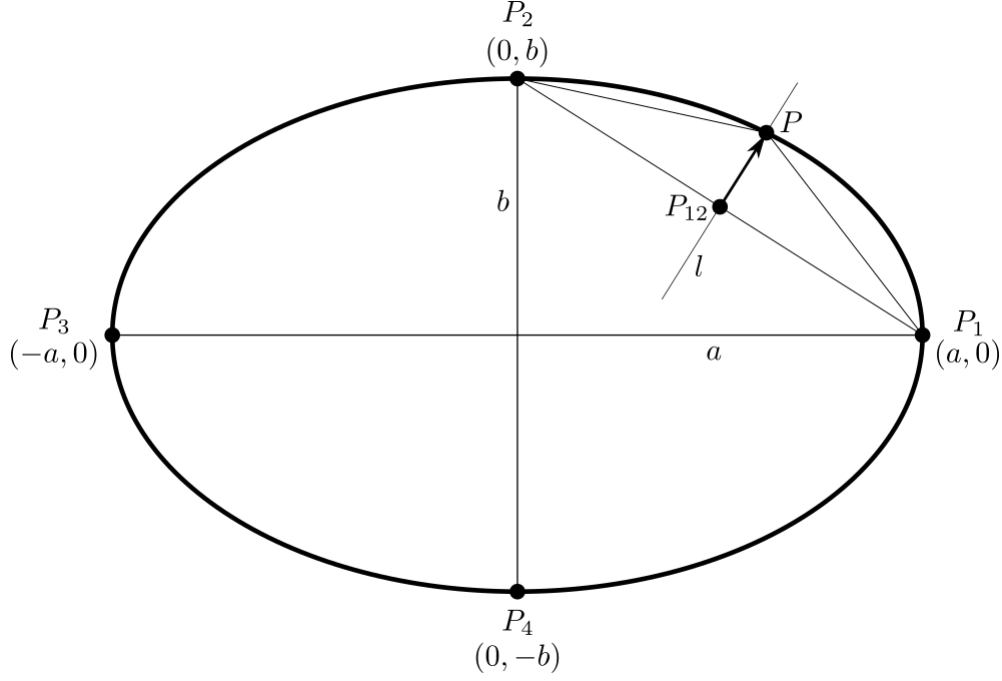


Figure 2.11: Equidistant points on ellipse.

Then, the point P is lying on the intersection of the ellipse and a line l passing through the point P_{12} , perpendicular to the line segment P_1P_2 .

$$l : \frac{1}{2}(a, b) + \frac{t}{2}(b, a), \quad t \in \mathbb{R}$$

Given the ellipse with semi-major axis a , semi-minor axis b and the center in the point $(0, 0)$, the point P can be calculated as follows:

$$P \in l \cap E \implies \frac{(a + tb)^2}{4a^2} + \frac{(b + ta)^2}{4b^2} = 1$$

$$t = \frac{ab(\sqrt{3a^4 + 2a^2b^2 + 3b^4} - a^2 - b^2)}{a^4 + b^4},$$

therefore

$$P = \frac{1}{2}(a, b) + \frac{ab(\sqrt{3a^4 + 2a^2b^2 + 3b^4} - a^2 - b^2)}{2(a^4 + b^4)}(b, a).$$

TODO binary search for height

Given edge length e , we are not able to calculate the height in which the distance between points P_1 and P is e . The visualization showing this is on the Figure 2.12.

We use binary search to find such height. Given the height and the singularity class, we can calculate the semi-major axis and semi-minor axis as

$$D_{n+-} : -hx^2 + h^{n-1} - z^2 = 0 \quad h > 0$$

$$x^2 + \frac{z^2}{h} = h^{n-2}$$

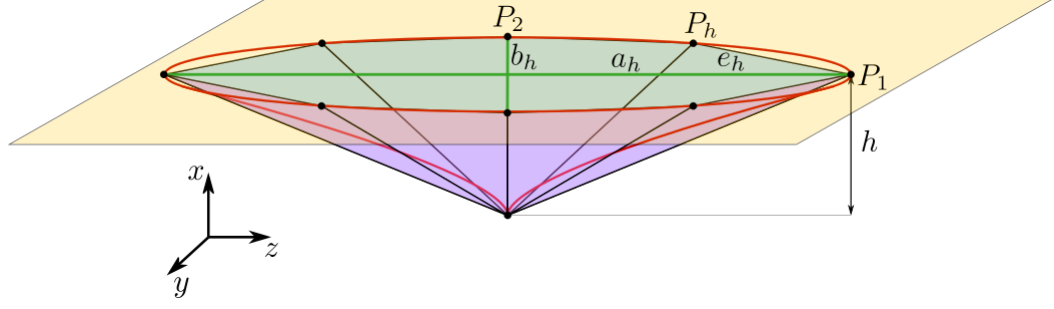


Figure 2.12: Calculating the point P_h – point on the ellipse equidistant from P_1 and P_2 .

$$\frac{x^2}{h^{n-2}} + \frac{z^2}{h^{n-1}} = 1 \implies a_h = \max(h^{\frac{n-2}{2}}, h^{\frac{n-1}{2}}) \wedge b_h = \min(h^{\frac{n-2}{2}}, h^{\frac{n-1}{2}}).$$

As we can see, we get the same ellipse for D_{n--} singularities:

$$\begin{aligned} D_{n--} &: -hx^2 - h^{n-1} - z^2 = 0 \quad h > 0 \\ 2|n \wedge x^2 + \frac{z^2}{h} &= -h^{n-2} \implies x^2 + \frac{z^2}{h} = h^{n-2}. \end{aligned}$$

Then

$$\begin{aligned} P_h &= \frac{1}{2}(h^{\frac{n-2}{2}}, h^{\frac{n-1}{2}}) + \frac{h^{\frac{2n-3}{2}}(\sqrt{3h^{2n-4} + 2h^{2n-3} + 3h^{2n-2}} - h^{n-2} - h^{n-1})}{2(h^{2n-4} + h^{2n-2})}(h^{\frac{n-1}{2}}, h^{\frac{n-2}{2}}) \\ P_h &= \frac{1}{2}(h^{\frac{n-2}{2}}, h^{\frac{n-1}{2}}) + \frac{h^{\frac{1}{2}}(\sqrt{3 + 2h + 3h^2} - 1 - h)}{2(1 + h^2)}(h^{\frac{n-1}{2}}, h^{\frac{n-2}{2}}) \end{aligned}$$

and we can calculate $e_h = \|P_h - P_1\|$.

As $e \leq a_h$, we can start the binary search on the interval $\langle 0, a^{\frac{2}{n-2}} \rangle$ or $\langle 0, a^{\frac{2}{n-1}} \rangle$ and finish, when required precision is reached.

TODO want to try to prove that $\|P_h - P_1\|$ is monotone in h .

Numerical calculation of local triangulation of ADE singularities

For other types of ADE singularities, the exact analytical calculations become more complicated. In this section we present an approach for triangulation of all types of ADE singularities using numerical algorithms such as binary search.

We begin by dividing ADE singularities into three categories.

1. Singularities topologically equivalent to a plane:

- A_{n--} , where $2|n$,
- A_{n+-} , where $2|n$,
- D_{n+-} , where $2|n$,

- E_{6++} ,
- E_{6+-} ,
- E_{8++} .

2. Singularities topologically equivalent to a cone except E_{7++} :

- A_{n--} , where $2 \nmid n$,
- A_{n+-} , where $2 \nmid n$,
- D_{n+-} , where $2 \nmid n$,
- D_{n--} , where $2 \nmid n$

3. Other singularities:

- D_{n--} , where $2|n$,
- E_{7++} .

We propose a general solution for first and second categories and we treat the singularities in the third category as a special cases.

Singularities topologically equivalent to a plane

Singularities topologically equivalent to a cone except E_{7++}

Other singularities - special cases

Triangulation of a plane with multiple A_{n--} singularities

In this section, we present an approach for creating an implicit equation of a surface which consists of a plane and arbitrary many A_{n--} singularities C^1 smoothly connected to this plane.

Input and output

In this section, the following data are provided on the input:

1. the number of singularities - m ,
2. m discrete points on a plane - $(x_1, y_1), \dots, (x_m, y_m)$,
3. m degrees of the singularities - n_1, \dots, n_m ,
4. m heights at which each singularity is connected - h_1, \dots, h_m .

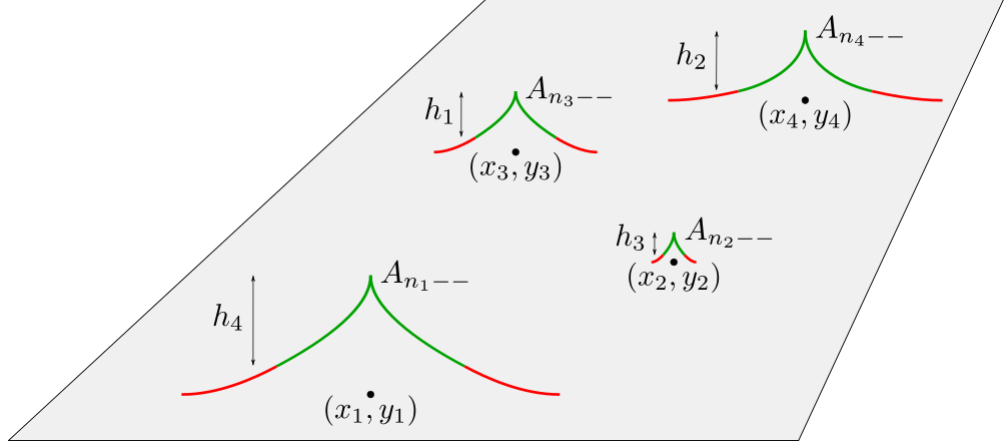


Figure 2.13: Plane with singularities.

The visualisation of desired output function can be seen on the Figure 2.13. On this figure, the singularity is displayed by green color, the red color is used to display the function which connects the singularity to a plane - the bump function. There are some limitations on the input data. As we do not want the singularities or the bump functions to intersect, we require that each pair of input points is distanced d_{ij} from each other. We specify the value of d_{ij} in the section TODO.

Bump function

Definition 14 *The support $\text{supp}(f)$ of a function $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is a set of points where f is not zero:*

$$\text{supp}(f) = \{x \in \mathbb{R}^n : f(x) \neq 0\}.$$

The closed support of the function f is defined as a closure of $\text{supp}(f)$.

Bump function is a function $f : \mathbb{R}^n \rightarrow \mathbb{R}$ which is smooth (C^∞) and compactly supported (the closed support of the function f is a compact subset of \mathbb{R}^n).

The most common example of such bump function is the function

$$f(x) = \begin{cases} e^{-\frac{1}{1-x^2}}, & x \in (-1, 1) \\ 0, & \text{otherwise,} \end{cases}$$

which is both C^∞ and compactly supported.

The bump function can be used to smoothly connect a curve to a line or a surface to a plane in higher dimension. If we only need to connect a curve and a line C^n smoothly, we only need a bump function which connects C^n smoothly to a line.

In our work, we connect two surfaces with C^1 continuity and for this purpose we use the function

$$f(x) = \begin{cases} -q \cdot \cos(k \cdot x) - q, & x \in (-\frac{\pi}{k}, \frac{\pi}{k}) \\ 0, & \text{otherwise,} \end{cases}$$

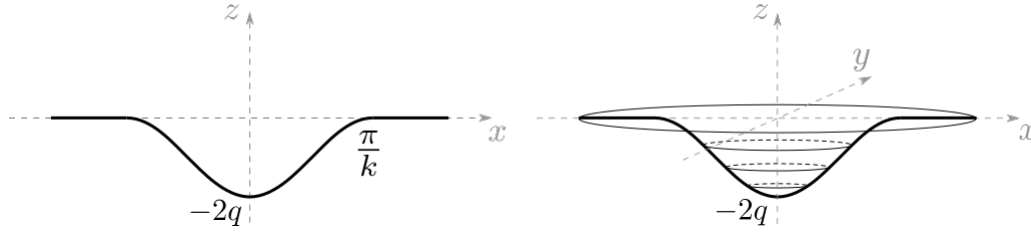
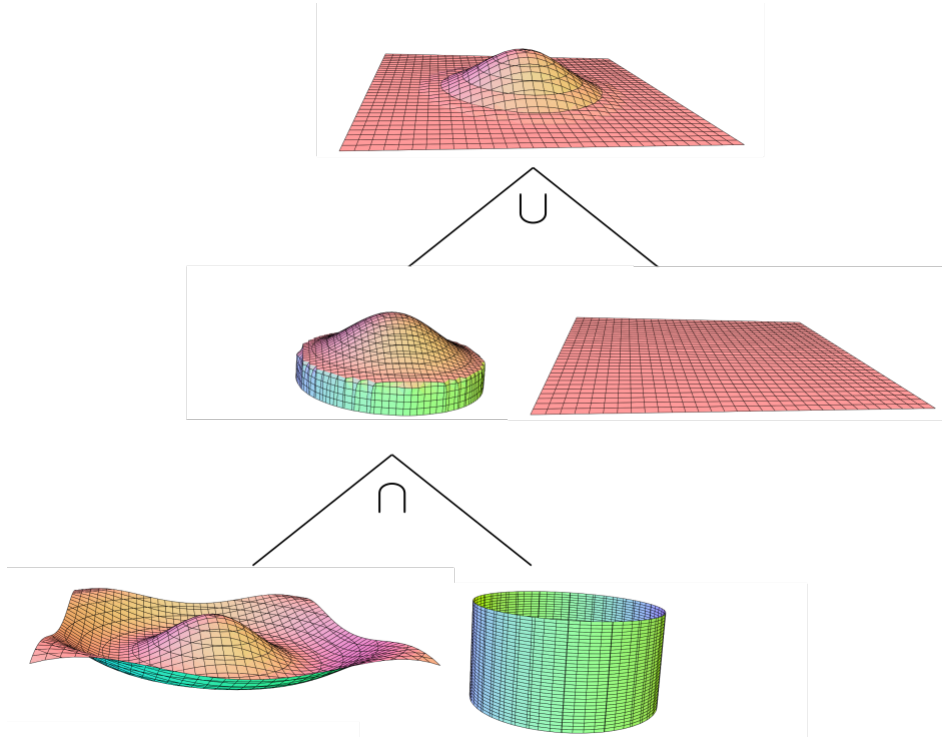
Figure 2.14: C^1 cosine bump function.

Figure 2.15: Construction of the cosine bump function using CSG.

rotated about z -axis. The result of the rotation is the following function:

$$f(x, y) = \begin{cases} -q \cdot \cos(k \cdot (x^2 + y^2)) - q, & x^2 + y^2 \leq (\frac{\pi}{k})^2 \\ 0, & \text{otherwise.} \end{cases}$$

Both of these functions can be seen on the Figure 2.14.

Implicit equation of the cosine bump function

To construct the implicit equation of the cosine bump function, we use CSG - constructive solid geometry, which is described in the section 1.3.

First, we cut out the part of the rotated cosine, where $x^2 + y^2 < (\frac{\pi}{k})^2$ using cylinder and intersection operation. Next, we use a plane and the union operation to *glue* the bump to the plane. The described process is displayed on the Figure 2.15 in the form of CSG tree. We use the following equations of the surfaces to model the cosine bump function:

Function name	Implicit equation
Rotated cosine function	$x + q \cdot \cos(k \cdot \sqrt{(y - p_y)^2 + (z - p_z)^2}) + q = 0$
Cylinder	$(y - p_y)^2 + (z - p_z)^2 - (\frac{\pi}{k})^2 = 0$
Plane	$x=0$

Table 2.1: Implicit equations for bump function modeling.

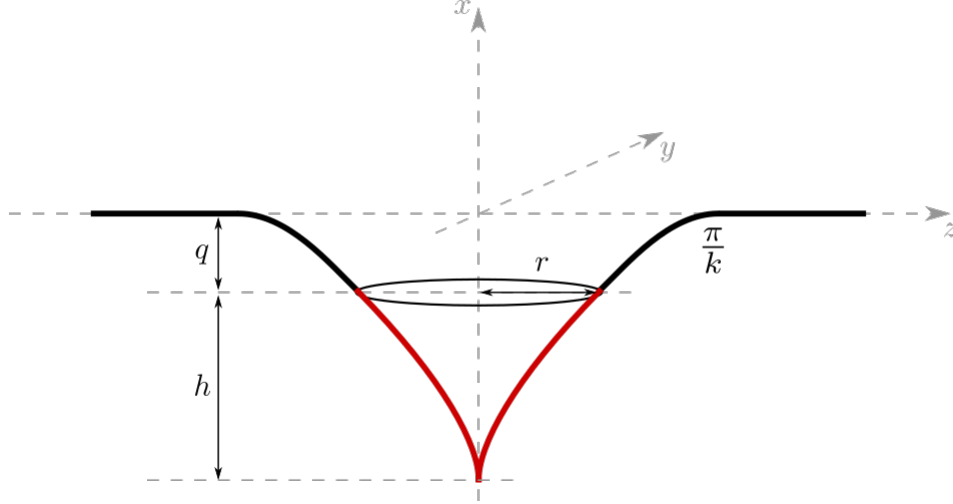


Figure 2.16: Attaching the singularity to a plane using the cosine bump function.

Parameters q and k allow us to change the amplitude and the frequency of the cosine function, parameters p_y and p_z are used to move the bump function to the given point (p_y, p_z) .

Attaching singularities to the plane using the cosine bump function

Given the type of the singularity - n and given height - h , we calculate the constants of the cosine bump function to connect C^1 smoothly to the given singularity.

The singularity given by the implicit equation $x^{n+1} - y^2 - z^2$ intersected with the plane $x = h$ produces a circle with the radius $r = \sqrt{h^{n+1}}$. The cosine bump function is scaled using q and k to smoothly connect the singularity in the middle of the cosine bump function. This approach is displayed on the Figure 2.16.

As the singularity is attached in the middle of the bump function, we get the equality $r = \frac{\pi}{2k}$ and therefore $\sqrt{h^{n+1}} = \frac{\pi}{2k} \implies k = \pi/(2\sqrt{h^{n+1}})$. The parameter q is calculated from C^1 continuity requirement. We require the gradients to be linearly dependent in the points of connection. Due to the rotation symmetry, we check it only for the intersection with the plane $z = 0$ for the point $(-q, \frac{\pi}{2k})$.

$$F = (x + h + q)^{n+1} - y^2 \implies \nabla F = [(n+1)(x + h + q)^n, -2y]$$

$$\nabla F \left(-q, \frac{\pi}{2k} \right) = \left[(n+1)h^n, -\frac{\pi}{k} \right]$$

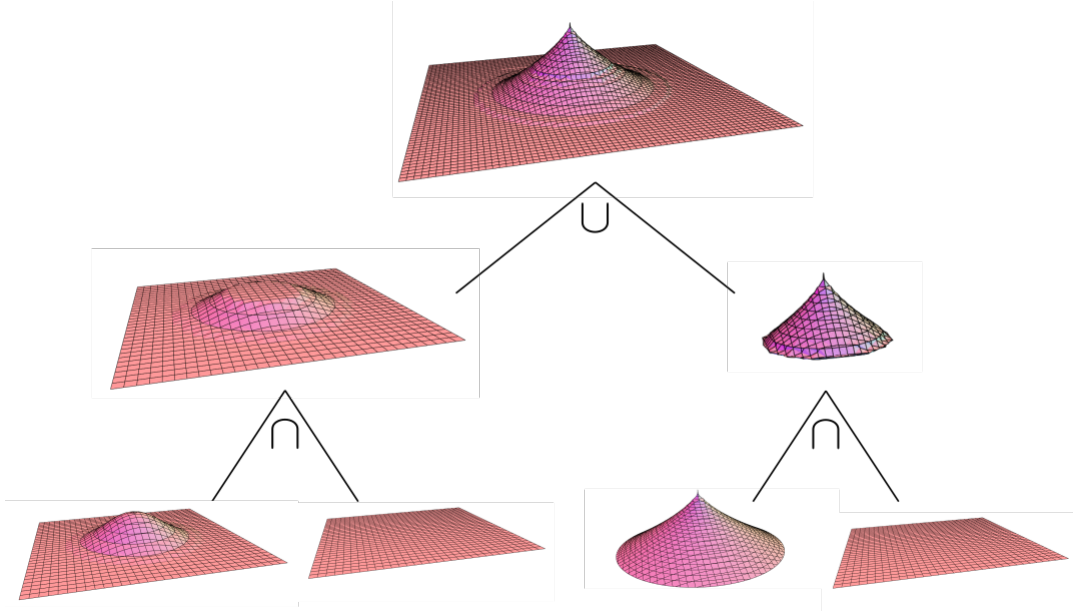


Figure 2.17: Attaching the singularity to a plane using CSG.

$$G = x + q \cdot \cos(ky) + q \implies \nabla G = [1, -qk \cdot \sin(ky)]$$

$$\nabla G \left(-q, \frac{\pi}{2k} \right) = [1, -qk]$$

Requiring $\nabla F(-q, \frac{\pi}{2k}) = s \cdot \nabla G(-q, \frac{\pi}{2k})$ and knowing $k = \pi/(2\sqrt{h^{n+1}})$, we get $s = (n+1)h^n$ and therefore $q = 4h/(\pi(n+1))$.

After calculating the parameters q and k of the cosine bump function, we proceed to connect the singularity to the bump function. We use intersection with the plane $x = -q$ to get the sections of the singularity and the section of the bump function and lastly, we use union of these two surfaces. The described process is displayed on the Figure 2.17. Curious reader can find the detailed calculation of the implicit equation in the appendix 4.

To connect multiple singularities to the same plane, we construct the implicit equation for each of these singularities and then use the union operation to create a surface with multiple singularities. This procedure is displayed on the Figure 2.18.

Limitations on the input data

As we already mentioned, we require that each pair of input points is distanced d_{ij} from each other. As the radius of the closed support of the cosine bump function is $r = \frac{\pi}{k} = 2\sqrt{h^{n+1}}$, the distance between two input points p_i, p_j must be at least $d_{ij} = 2\sqrt{h_i^{n_i+1}} + 2\sqrt{h_j^{n_j+1}}$. This way, both singularities and the corresponding bump functions do not intersect.

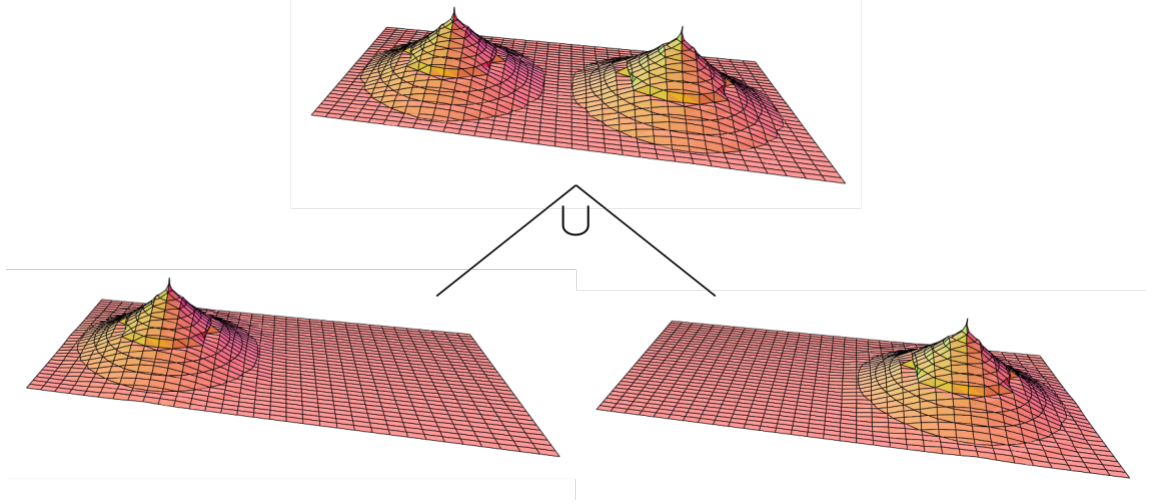


Figure 2.18: Plane with multiple attached singularities.

Visualization of results of the triangulated surfaces

2.3 Triangulation of non-isolated singularities

We present an approach for triangulation of implicit surfaces with singular curves which are a result of performing intersection operation on the interiors of two regular implicit surfaces.

We start by creating a local mesh for the surroundings of these singular curves and finish the mesh in the regular parts.

2.3.1 Creating the local mesh around the singular curves

We start by approximating the singular curve by a polyline. On the input, one point close to the singular curve is given. This point serves as a starting point P_0 . The tangent vector $\vec{t}_C(P_0)$ of the curve is computed as the cross product of the unit normal vectors of the two surfaces S_1 and S_2 . Given the required approximate edge length e , a point Q_1 is computed as $Q_1 = P_0 + e \cdot \vec{t}_C(P_0)$. We obtain the point P_1 by projecting the point Q_1 to the curve C using an approach described in the section 1.4. Other points are then created iteratively by the same approach:

$$Q_{n+1} = P_n + e \cdot \vec{t}_C(P_n),$$

$$P_{n+1} = proj_C(Q_{n+1}),$$

while checking if the new point is inside the axis-aligned bounding box. We stop once the new point is outside of the axis-aligned bounding box or if the new point is close to the starting point P_0 , which means the approximated curve is a closed curve.

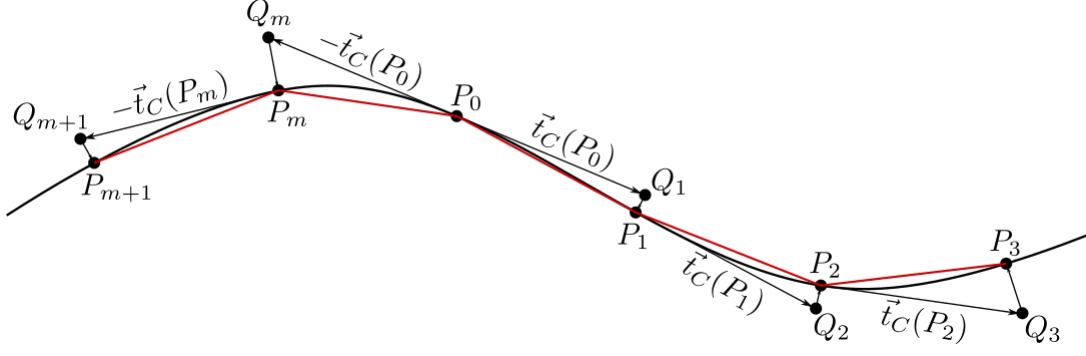


Figure 2.19: Approximation of the implicit curve by polyline.

In case of the open curve, to complete the whole polyline, we also approximate the second part of the singular curve by iteratively creating points in the opposite direction, starting from point P_0 . Let m be number of points in the polyline, the second part of the polyline consists of the points P_m, P_{m+1}, \dots , where

$$Q_m = P_0 - e \cdot \vec{T}_C(P_0),$$

$$P_m = \text{proj}_C(Q_m),$$

and again, iteratively

$$Q_{n+1} = P_n - e \cdot \vec{T}_C(P_n), \quad n = m, m+1, \dots$$

$$P_{n+1} = \text{proj}_C(Q_{n+1}), \quad n = m, m+1, \dots$$

The presented approach is visualized on the Figure 2.19.

TODO too big image

Chapter 3

Results

3.1 Quality criteria

3.2 Comparison with TODO

Chapter 4

Future work

Conclusion

TODO Conclusion

Bibliography

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Appendix A

$$f_1 = (x - p_x)^{n+1} - (y - p_y)^2 - (z - p_z)^2$$

$$f_2 = -x - q$$

$$f_3 = f_1 \cap f_2$$

$$f_3 = f_1 + f_2 - \sqrt{f_1^2 + f_2^2}$$

$$f_4 = x + q + q \cdot \cos \left(k \sqrt{(y - p_y)^2 + (z - p_z)^2} \right)$$

$$f_5 = 4h^{n+1} - (z - p_z)^2 - (y - p_y)^2$$

$$f_6 = f_4 \cap f_5$$

$$f_6 = f_4 + f_5 - \sqrt{f_4^2 + f_5^2}$$

$$f_7 = x$$

$$f_8 = f_6 \cup f_7$$

$$f_8 = f_6 + f_7 + \sqrt{f_6^2 + f_7^2}$$

$$f_9 = x + q$$

$$f_{10} = f_8 \cap f_9$$

$$f_{10} = f_8 + f_9 - \sqrt{f_8^2 + f_9^2}$$

$$f_{11} = f_{10} \cup f_3$$

$$f_{11} = f_{10} + f_3 + \sqrt{f_{10}^2 + f_3^2}$$

$$f_{11} =$$

The corresponding surfaces are displayed on the figure 4.1

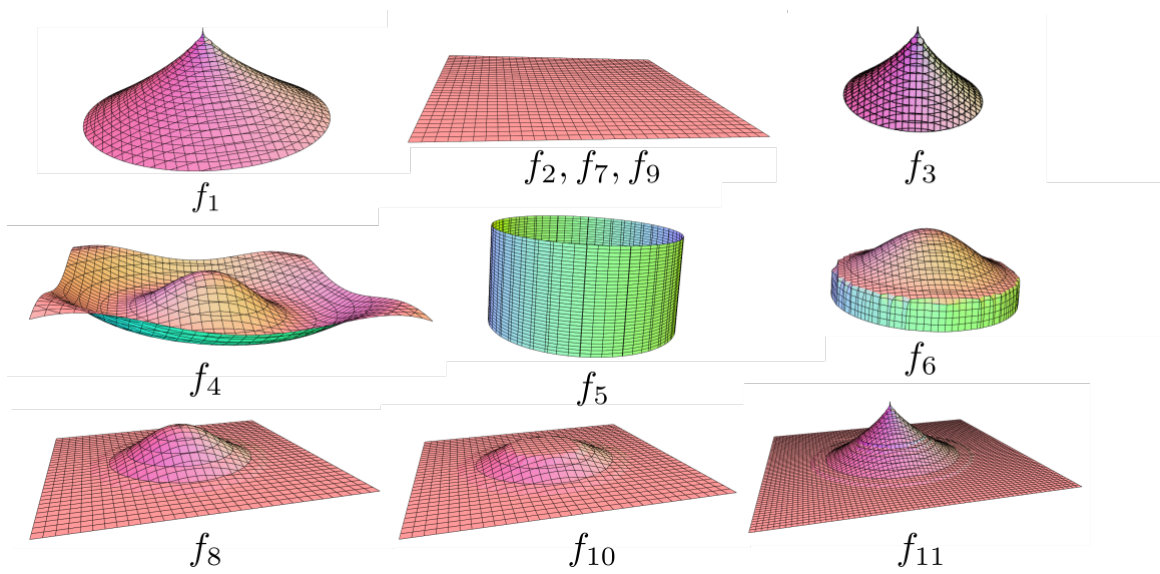


Figure 4.1: Attaching the singularity to a plane using CSG.

Appendix B

4.1 Triangulation of regular implicit surfaces

In this section, we shortly describe the algorithm for triangulation of the regular parts of implicit surfaces introduced in bachelor's thesis [10].

The algorithm creates triangles iteratively, one at a time. The new triangle is created at the edge of the existing triangle. The new point is projected on the surface using the Newton-Raphson method. This method finds the root of the implicit function lying on the line, which is in the direction of the gradient of the implicit function. This approach is displayed on the Figure 4.2.

After the new point is projected on the surface, conditions are checked. Some of these conditions are based on the Delaunay triangulation introduced by Hilton [8]. The Delaunay condition checks, if some other points are in the proximity of circumcenter of the new triangle. The conditions minimize the chances of triangle intersection. If the new triangle intersects with some existing triangles, the algorithm tries to connect the new triangle to the existing triangles in its proximity.

The algorithm is enriched with the possibility to triangulate adaptively to the curvature of the surface using approach presented by Akkouche [2]. It can also triangulate surfaces in the bounded volume - axis aligned bounding box, given by six numbers - minimal and maximal value for each of the three axes.

As a part of our work, we reimplement the algorithm to be more effective by using advanced data structures, such as the half-edge data structure [9] and the range tree [11].

4.2 Data structures for triangulation algorithm

Half-edge data structure

Triangular mesh is given by a set of vertices, non-oriented edges and triangular faces. There are multiple methods for mesh representation. The straightforward one - list of vertices, edges and faces does not provide any information about the local surroundings of the vertices, edges and faces and therefore the searching for incident faces or

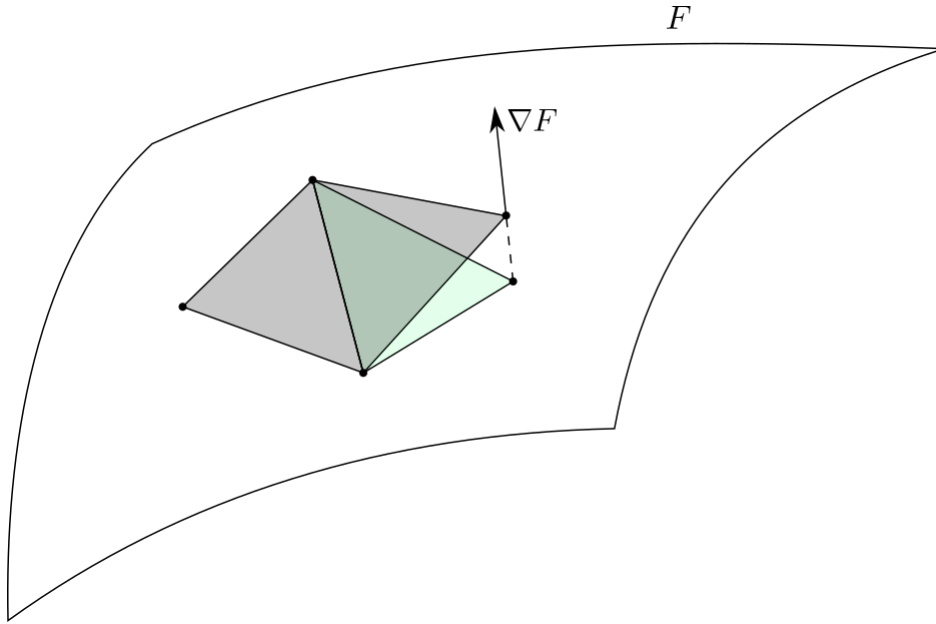


Figure 4.2: TODO

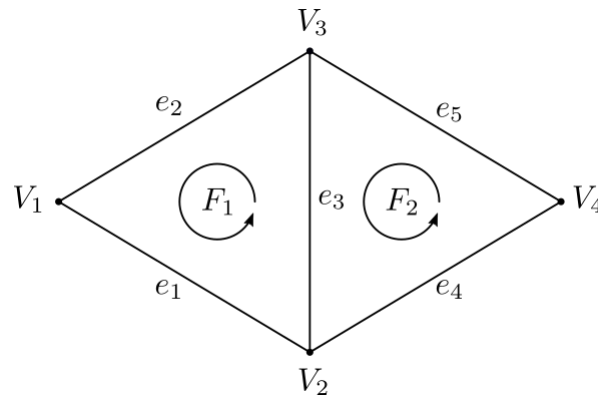


Figure 4.3: Example of winged edge representation.

incident edges is complicated and inefficient.

In 1975, Baumgart [3] presented a representation using winged edges, which was further improved in 1985 by Weiler [17] who presented the modification called half-edge data structure. Both of these representations are edge-based representations, each edge stores references (pointers) to the surrounding vertices, edges and faces. One can easily extract the information of the surrounding vertices, edges and faces.

In the winged edge representation, each edge is represented as oriented line segment and stores references to its vertices, both incident faces, left and right traverse. Each vertex stores a reference to one of the outgoing edges and each face stores a reference to one of its incident edge. An example of such representation is shown on the Figure 4.3 and tables 4.1, 4.2, 4.3.

In the half-edge representation, edges are split into two halves. Each half-edge stores reference to its initial vertex, left face, opposite half-edge and left traverse -

Edge	Vertex		Face		Left traverse		Right traverse	
	Initial	Terminal	Left	Right	Pred.	Succ.	Pred.	Succ.
e_1	V_1	V_2	F_1		e_2	e_3		
e_2	V_3	V_1	F_1		e_3	e_1		
e_3	V_2	V_3	F_1	F_2	e_1	e_2	e_5	e_4
e_4	V_2	V_4	F_2		e_3	e_5		
e_5	V_4	V_3	F_2		e_4	e_3		

Table 4.1: Edge table of a winged edge data structure for the Figure 4.3.

Vertex	Coordinates			Edge
	x	y	z	Outgoing edge
V_1	x_1	y_1	z_1	e_1
V_2	x_2	y_2	z_2	e_4
V_3	x_3	y_3	z_3	e_2
V_4	x_4	y_4	z_4	e_5

Table 4.2: Vertex table of a winged edge data structure for the Figure 4.3.

Face	Edge
F_1	e_1
F_2	e_4

Table 4.3: Face table of a winged edge data structure for the Figure 4.3.

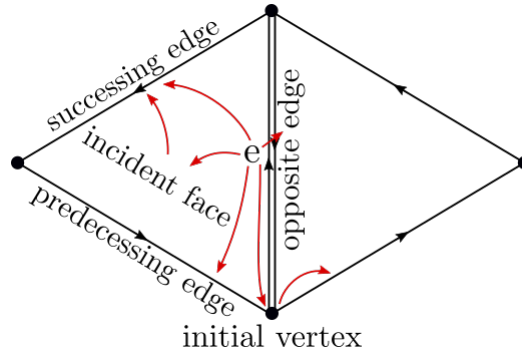


Figure 4.4: Visualisation of the half-edge data structure.

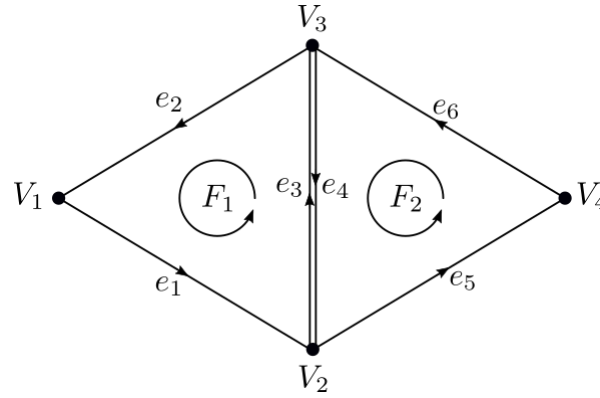


Figure 4.5: Example of half-edge representation.

predecessing and successing half-edge. A visualisation of the half-edge is displayed on the Figure 4.4. An example of half-edge representation is shown on the Figure 4.5 and tables 4.4, 4.5, 4.6.

Range tree

Mesh structure

Triangular mesh consists of vertices, edges and faces. We use the half-edge data structure for mesh representation and the range tree for time efficient search of vertices in three dimensional interval. The mesh structure used for maintaining and modifying the triangular mesh consists of:

- list of vertices,
- list of half-edges,
- list of faces,
- set of active edges,
- set of checked edges,
- set of bounding edges,
- range tree of all vertices.

Edge	Vertex	Face	Edges		
	Initial	Incident	Pred.	Succ.	Opposite
e_1	V_1	F_1	e_2	e_3	
e_2	V_3	F_1	e_3	e_1	
e_3	V_2	F_1	e_1	e_2	e_4
e_4	V_3	F_2	e_6	e_5	e_3
e_5	V_2	F_2	e_4	e_6	
e_6	V_4	F_2	e_5	e_4	

Table 4.4: Edge table of a half-edge data structure for the Figure 4.5.

Vertex	Coordinates			Edge
	x	y	z	Outgoing edge
V_1	x_1	y_1	z_1	e_1
V_2	x_2	y_2	z_2	e_3
V_3	x_3	y_3	z_3	e_2
V_4	x_4	y_4	z_4	e_6

Table 4.5: Vertex table of a half-edge data structure for the Figure 4.5.

Face	Edge
F_1	e_1
F_2	e_4

Table 4.6: Face table of a half-edge data structure for the Figure 4.5.

Vertex consists of

- three coordinates,
- index of itself in the list of vertices,
- list of indices to all outgoing edges.

Half-edge consists of

- six coordinates,
- index of itself in the list of half-edges,
- index of initial vertex in the list of vertices,
- index of terminal vertex in the list of vertices,
- index of opposite edge in the list of half-edges,
- index of predeccessing edge in the list of half-edges,
- index of successing edge in the list of half-edges,
- index of incident face in the list of faces.

Face consists of

- nine coordinates,
- index of incident half-edge in the list of half-edges.

Implementation

TODO