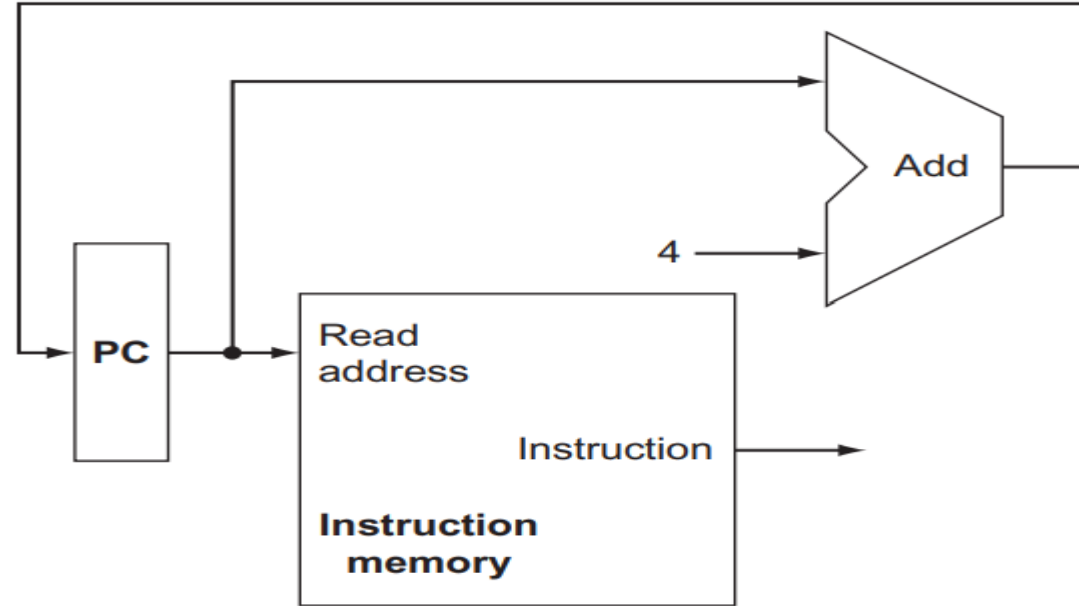


Computer Architecture (ENE1004)

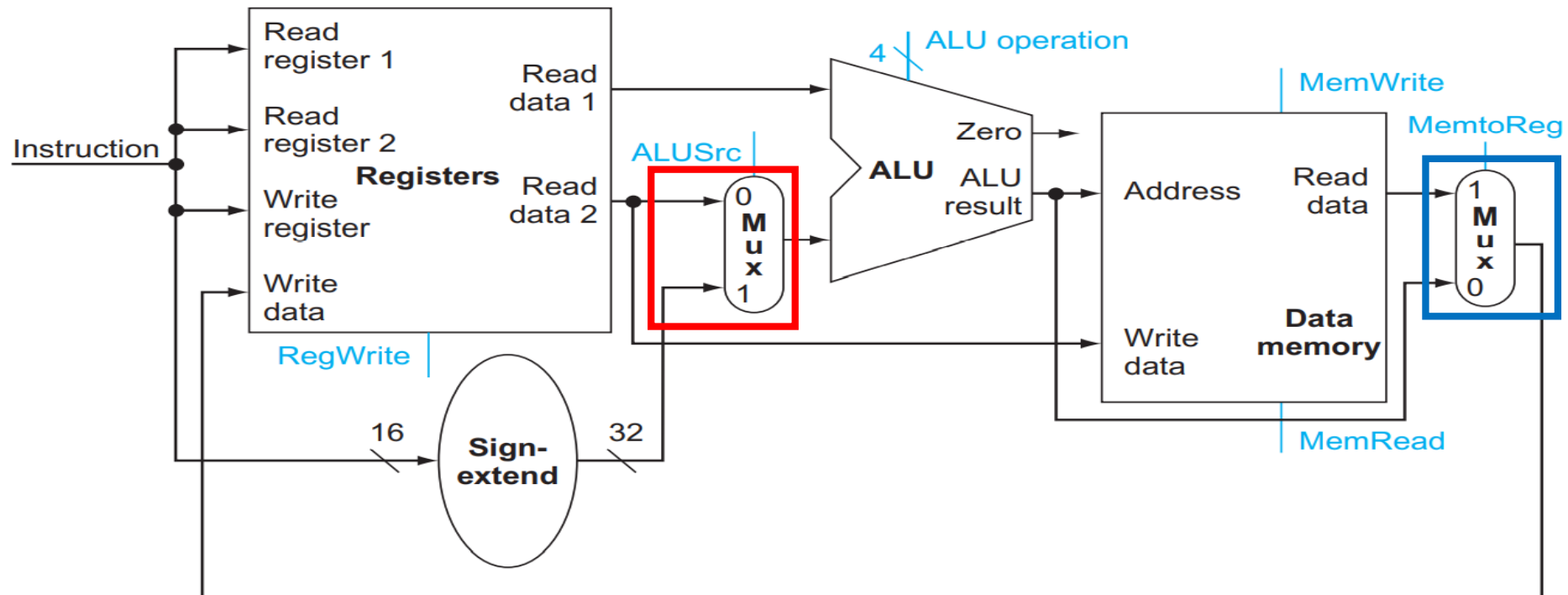
Lec - 11: The Processor (Chapter 4) - 2

A Datapath for Fetching Instructions



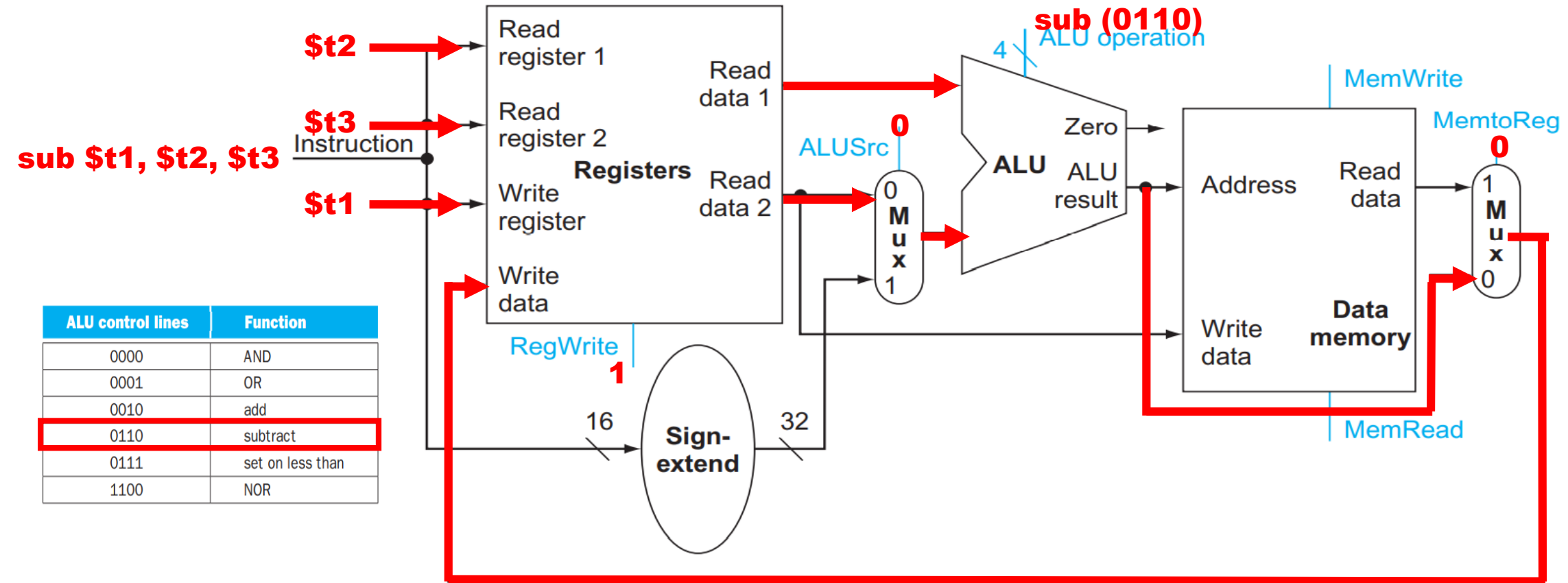
- (1) To execute an instruction, we must “fetch” the instruction from memory
 - The processor obtains an instruction when the address of PC is given to the instruction memory
- (2) To prepare for the execution of the next instruction, we must also “update (increment) the PC” so that it points to the next (sequential) instruction
 - The size of an instruction is 4 bytes
 - For jump and branch instructions, PC will point at the target address (instruction)

A Datapath for R-Type and Load/Store



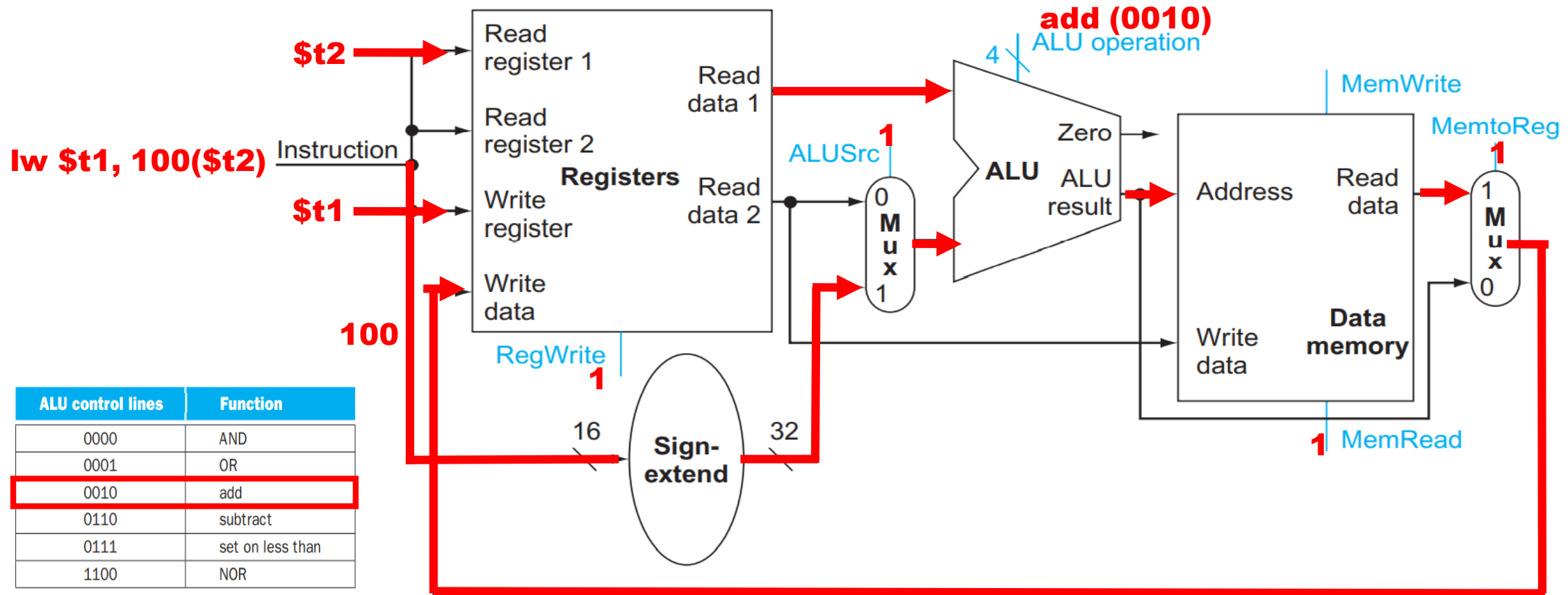
- We combine (II) + (III) into a single datapath to execute any of the two types
 - R-type (e.g., **add \$t1, \$t2, \$t3**): register file, ALU
 - Load/store (e.g., **lw \$t1, 100(\$t2)**): register file, ALU, sign extension unit, data memory
- We add multiplexer (**Mux**) to select the required path for a given instruction
 - **Mux** for an input of the ALU: data from a register? or immediate from the instruction?
 - **Mux** for the data input to the register file: data from the data memory? or result of the ALU?

Execution of R-Type on the Datapath

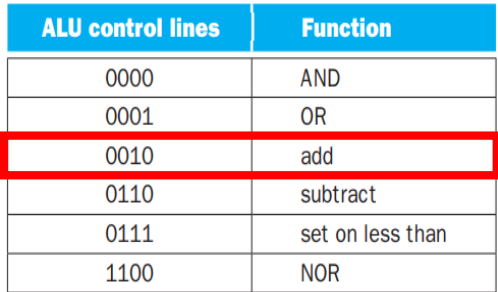


- **ALUSrc = 0** to select **Read data 2** of the register file to an input of ALU
- **ALU operation = 0110** to perform **sub** instruction on ALU
- **MemtoReg = 0** to select the data value of **ALU result** to supply **Write data** of the register file
- When writing into a register, **RegWrite = 1**

Execution of Load Instruction on the Datapath



- **ALUSrc = 1** to select a **sign-extended immediate value** to an input of ALU
- **ALU operation = 0010** to perform **add** instruction on ALU
- **MemRead = 1** to read data from the data memory
- **MemtoReg = 1** to select the data read from the data memory to supply **Write data** of the register file
- When writing a register, **RegWrite = 1**



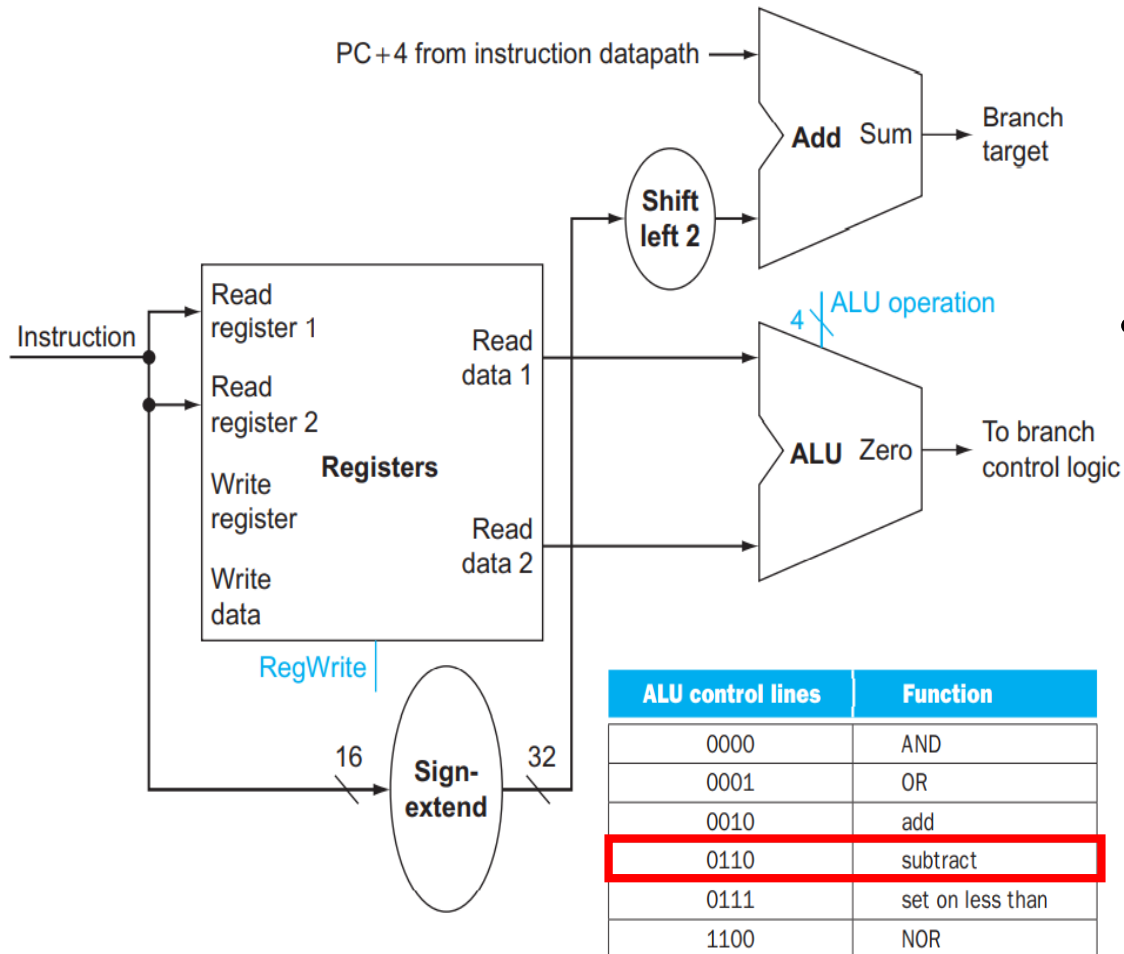
- **ALUSrc = 1** to select a **sign-extended immediate value** to an input of ALU
- **ALU operation = 0010** to perform **add** instruction on ALU
- **MemWrite = 1** to write data into the data memory

Datapath Elements for Branch

- A branch instruction computes its target addresses using PC-relative addressing
 - **beq \$t1, \$t2, offset_value**
 - Target address = ① address of the subsequent instruction + ② branch offset in bytes
 - Target address = ① (PC + 4) + ② (offset_value * 4)
- Let us consider what datapath elements branch instructions need
 - ① (PC + 4) can be obtained from the datapath for fetching an instruction
 - ② (offset_value * 4) can be done by shifting left the offset_value by 2
- Branch instructions can have two different scenarios depending on the condition
- (1) If it is true, the next instruction is the instruction at the target address
 - We say that the branch is “taken”
 - $PC \leftarrow \text{target address}$
- (2) If it is not true, the next instruction is the instruction that follows sequentially
 - We say that the branch is “not taken”
 - $PC \leftarrow PC + 4$

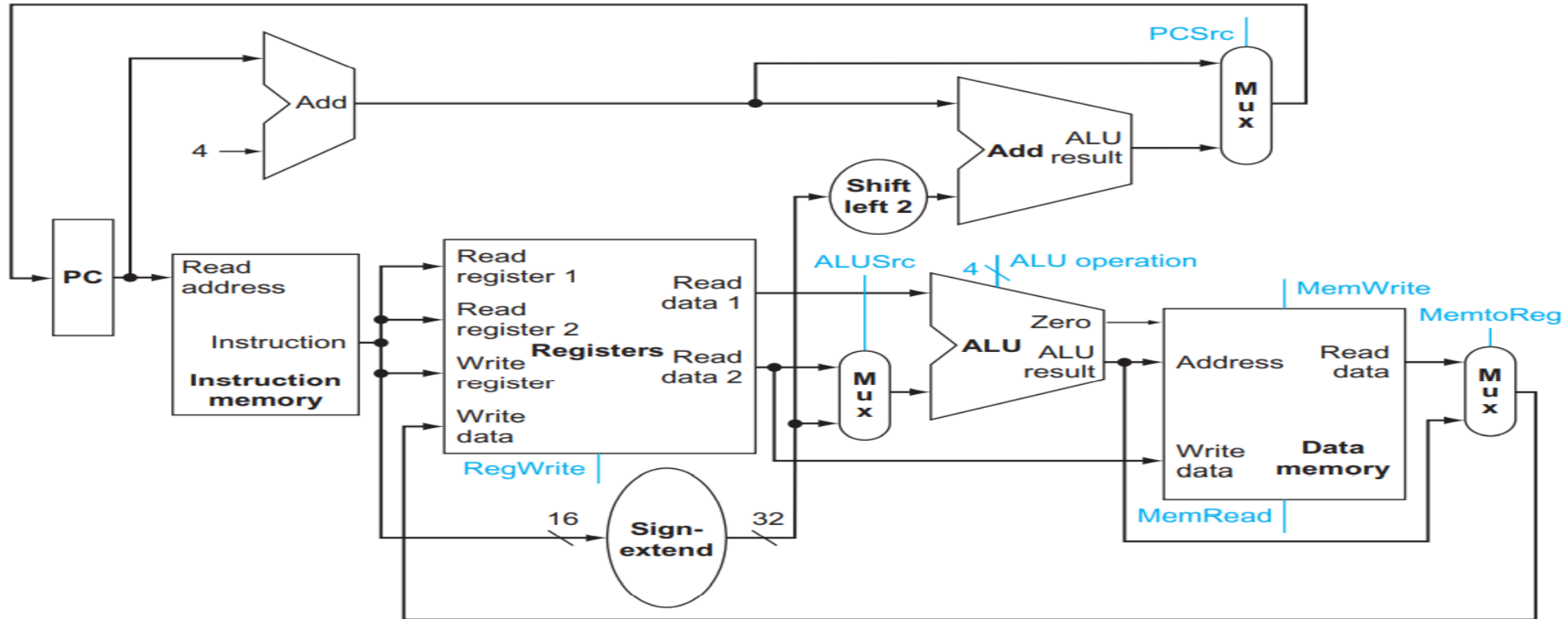
Datapath Elements for Branch

beq \$t1, \$t2, offset_value



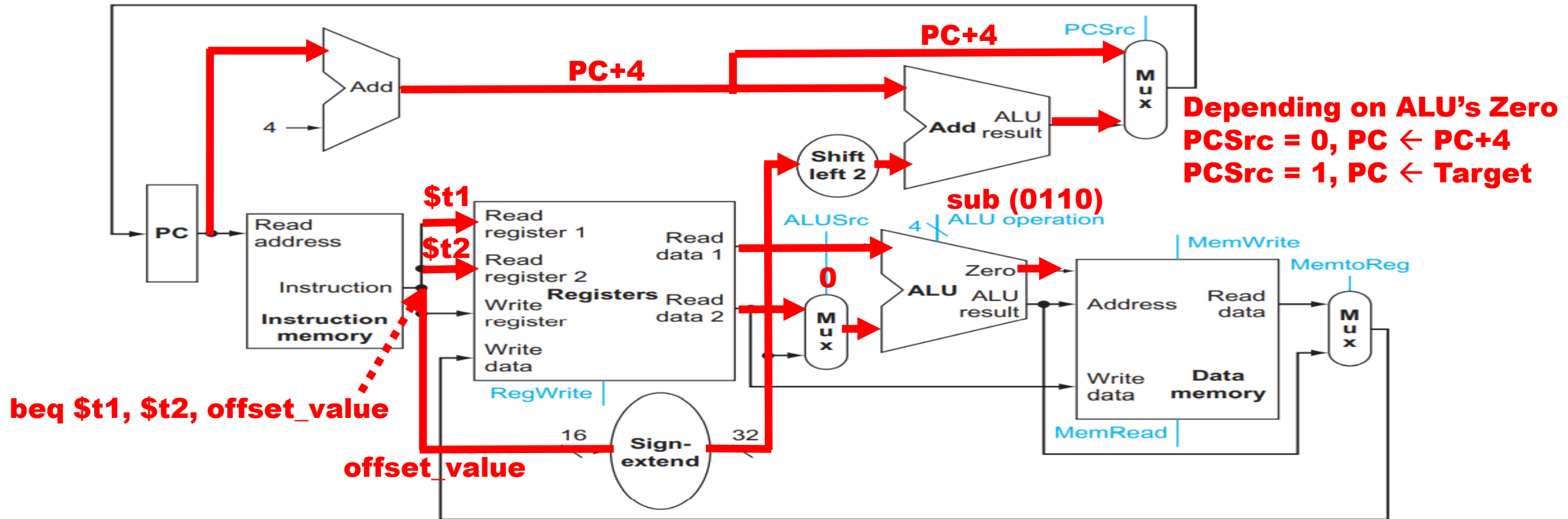
- (1) Computing the target address
 - Target address = $(PC + 4) + (\text{offset_value} * 4)$
 - **Sign extension unit** for (offset_value)
 - **Shift left 2 unit** for (offset_value * 4)
 - **Adder** for $(PC + 4) + (\text{offset_value} * 4)$
- (2) Comparing the register contents
 - **Register file** to supply two operands (**\$t1** & **\$t2**)
 - **ALU** for comparing the two operands (**\$t1** vs **\$t2**)
 - **ALU operation = 0110** for a subtract
 - If **Zero = 1** (two values are equal), the instruction in the target address should be executed next
 - If **Zero = 0** (two values are not equal), the instruction that follows sequentially should be executed next

A Single Datapath that Combines R, I, and Branch



- For fetching instructions, we need PC, Instruction memory, Adder
- For R-type instructions, we need Register file and ALU
- For load/store instructions, we need Register file, ALU, Sign extension unit, Data memory
- For branch instructions, we need Register file, ALU, Sign extension unit, Shift left 2 unit, Adder

Execution of Branch Instruction on the Datapath



- (1) Computing the target address: $(PC+4) + (\text{offset_value} * 4)$ using sign extension/shift left 2/adder
- (2) At the same time, comparing the two register operands: (**\$t1** vs **\$t2**) using register file/ALU
- The result of (2) is evaluated whether it is zero or not using **ALU's Zero**
- Based on the evaluation, either $(PC+4)$ or $\{(PC+4) + (\text{offset_value} * 4)\}$ is selected using **PCSrc = 0/1**