

# A comparison between functional and object oriented programming approaches in JavaScript

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# 1 Abbreviations

ECMAScript 2015 = ES6

Functional Programming = FP

Object-Oriented Programming = OOP

# 2 Background and Related work

Object oriented or imperative languages are dominating with languages such as C, C++, Java and C#, however functional languages seem to be of interest in the industry for handling big data and concurrency with languages such as Scala, Erlang or Haskell [1, 4, 9, 11]. Some are also arguing that a more functional approach will give you less code, more readable code, code that is easier to maintain and easier to test, and that learning it will give you better experience as a programmer [1, 2, 3, 9].

Here we will describe the functional programming paradigm, FP, and also the object oriented programming paradigm, OOP.

## 2.1 Functional programming

Functional programming, FP, is a programming paradigm that at its core is based on lambda-calculus [35]. Programs are constructed using functions and by avoiding changing the state. By not modifying the state side effects are avoided. Computation is done by changing the environment, rewriting the functions, rather than changing variables. Multiple functions can be composed into larger and more complex functions, and should be reduced to its simplest state following mathematical rules. The following are concepts used in FP.

**Lazy evaluation** - Lazy evaluation is when a value is not calculated until it is needed [5].

**Static type checking** - Pure functional languages usually have static type checking [35]. This means that variables are of a certain type, for example int or char. In such languages it is not allowed to use functions with the wrong types. So if there is for example a function taking an int as parameter it is illegal to call that function with a String as its parameter. In dynamic type checked-languages this would be legal, since variables are not specified as types, and might cause an unexpected error. For example in JavaScript:

```

1 //JavaScript has dynamic type checking
2 function double(nr) {
3   return nr * 2;
4 }
5
6 var word = "string";
7 double(word); //Will return NaN but still continue
               running

```

A similar program in go lang, that is a statically type checked language:

```

1 func main() {
2   var word string = "string"
3   fmt.Printf("%d", double(word)) //Error: cannot use
   word (type string) as type int in argument to double
4 }
5
6 func double(nr int) int {
7   return nr * 2
8 }

```

This will give an error when compiling.

There are however pure functional languages with dynamic type checking, such as Clean. There are also languages that are mostly functional with dynamic type checking, such as Lisp or Scheme [52].

**Side effects** - Side effects are for example changing a variable or any interaction outside the function [1]. This is avoided in functional programming since it may result in incorrect and unexpected behaviour.

**Pure functions** - A pure functions always returns the same result, given the same input, and does not have side effects [1]. See the following example:

```

1 var sum = 0;
2
3 //Impure
4 function add(a, b) {
5   sum = a + b;
6 }
7
8 //Pure
9 function add(a, b) {
10  var tmp = a + b;
11  return tmp;
12 }

```

**Higher order functions** - Higher order functions are an important concept in functional programming [5]. First class functions mean that functions are treated as values, which means that they can be stored in variables,

stored in arrays or created if needed. A higher order function is a first class function that either takes a function as a parameter, returns a function as a result or both. This makes it possible to compose larger and more complex functions.

```
1 function add(a, b) {
2   return a + b;
3 }
4
5 //Functions can be placed in variables.
6 var thisFunc = add;
7
8 //And also put in other variables.
9 var sameFunc = thisFunc;
10 sameFunc(1, 2); //Will give the output 3.
11
12 //Functions can also be used as parameters or return
   values.
13 function applyFunc(f, a, b) {
14   return f(a, b);
15 }
16
17 applyFunc(function(a, b) {
18   return a * b;
19 }, 3, 2); //Will give output of 6
20
21 //Returns a function that returns a string
22 function getStringCreator(category, unit) {
23   return function(value) {
24     return category + ': ' + value + unit;
25   }
26 }
27
28 var weightStringCreator = getStringCreator('Weight', 'kg
   ');
29 weightStringCreator(5); //Outputs   Weight: 5 kg
```

**Recursion** - Recursive functions are functions that call themselves and are used as loops [5]. It is important in functional programming since it can hide mutable state and also implement laziness. In functional programming recursion is used rather than loops.

```
1 //Adds all numbers from start to end with a loop
2 function iterativeAdd(start, end) {
3   var sum = 0;
4   while(start <= end) {
5     sum += start;
6     start++;
7   }
8 }
```

```

9     return sum;
10 }
11
12 //Adds all numbers from start to end recursively
13 function recursiveAdd(start, end) {
14     if (start == end) {
15         return end;
16     }
17     else {
18         return start + recursiveAdd(start + 1, end);
19     }
20 }
21 iterativeAdd(1, 6); //Outputs 21
22 recursiveAdd(1, 6); //Also outputs 21

```

**Currying** - Currying means a function can be called with fewer arguments than it expects and it will return a function that takes the remaining arguments [1].

```

1 //Returns a new function that takes the remaining
   argument or the sum of a and b if both are provided.
2 function addWithCurrying(a, b) {
3     if(b) {
4         return a + b;
5     }
6     else {
7         return function(b) {
8             return a + b;
9         }
10    }
11 }
12
13 var curryAdd = addWithCurrying(4);
14 curryAdd(6); //Outputs 10
15 addWithCurrying(4, 6); //Also outputs 10

```

**Immutable data structures** - In functional programming mutations, that are side effects, are avoided [5]. Hidden side effects can result in a chains of unpredicted behaviour in large systems. Instead of mutating the data itself a local copy is mutated and returned as a result, as seen in the pure functions example.

## 2.2 Object oriented programming

Object oriented programming is a programming paradigm which is built around objects. These objects may contain data, in form of fields, often referred as attributes and code in form of procedures, often referred to as methods [26].

These objects are created by the programmer to represent something with the help of its attributes and methods. For example, a class can represent an employee. An employee has the attributes assignment and salary. The employee also has a method for doing work. Then the employee has to work somewhere, so we create another object for a company where our employee can work. The company has the attributes income and number of employees. It also has methods to hire employees and fire employees. Since there are more employees who work at this company, we can add more employees by creating new objects of the information in employee class.

By building objects together like this, you can build programs by the object oriented paradigm.

**Class** - A class is a model for a set of objects, which clearly the object oriented paradigm is built around[35]. The class establishes what the object will contain, for example variables and functions, and signatures and visibility of these. To create a object of any kind, a class must be present.

```
1 //Creates a class Animal with the properties name and
   age and a function for logging the properties to the
   screen
2 class Animal {
3     constructor (name, age) {
4         this.name = name;
5         this.age = age;
6     }
7
8     logAnimal() {
9         console.log('Name:␣' + this.name + '\nAge:␣' + this.
            age);
10    }
11 }
12
13 var animal = new Animal('Buster', '9');
14 animal.logAnimal();
15 //Outputs Name: Buster
16 //Age: 9
```

**Object** - An object is a capsule that contains the variables and functions established in the class. All data and functions can be accessible from outside the object.

**Inheritance** - When an object acquires all properties and functionality of another object [41], it is called inheritance. This provides code reusability.

```
1 //Based upon the Animal class with an added property
   race
2 //The original logAnimal() is overridden so the new
   property is also logged to the screen.
```



```

3  class Dog extends Animal {
4      constructor(name, age, race) {
5          super(name, age);
6          this.race = race;
7      }
8
9      logAnimal() {
10         super.logAnimal();
11         console.log('Race:␣' + this.race);
12     }
13 }
14
15 var dog = new Dog('Buster', '9', 'Shitzu');
16 dog.logAnimal();
17 //Outputs Name: Buster
18 //Age: 9
19 //Race: Shitzu

```

**Encapsulation** - Encapsulation can be used to refer to two different things [35, 41]. A mechanism to restrict direct access to some object components and the language construct that facilitates the bundling of data with methods. The access part is done by making the different parts of an object public or private and the bundling is made with objects.

## 2.3 JavaScript

The support for the different paradigms in JavaScript

### 2.3.1 Functional programming

JavaScript is not a functional language, but it is possible to write functional code with it [1]. There are also libraries that makes functional programming in JavaScript easier, such as Underscore.js [5], however we will use ECMA 2015, ES6, that already provide many of the functions provided by Underscore.js. This is also to use similar environment for our different implementations in this experiment.

In JavaScript it is possible to treat functions as any other variable, pass them as function parameters or store them in arrays, so called first class functions [1]. There are also higher order functions such as `map()`, `filter()` and `reduce()` that might replace loops [2]:

**map()** - Takes calls a provided function on every element in an array and returns an array with the outputs [27].

**filter()** - Returns a new array with all the elements that passes a test in a provided function.

**reduce()** - Reduces a array to a single value.

However there is no automatic currying or immutable data structures in JavaScript. JavaScript is also dynamically type checked. All of these can be added to JavaScript with libraries.

### 2.3.2 Object oriented programming

JavaScript has always had support for OOP. But with ES6, OOP starts to look more like classical OOP languages like for example Java [30].

JavaScript do not fully support encapsulation, since you cant make variables and functions private,

## 2.4 Related work

In Curse of the excluded middle [3] Erik Meijer arguments with multiple examples that the industry should not focus on a combination between functional and objected oriented methods to counter handling big data with concurrency and parallelism. He concludes that it is not good enough avoid side effects in imperative or object oriented languages. It is also not good enough to try to ignore side effects in pure functional languages. Instead he thinks that people should either accept side effects or think more seriously about using the functional paradigm.

In Functional at scale [4] by Marius Eriksen he is explaining why Twitter uses methods from functional programming to handle concurrent events that arises in large distributed systems in cloud environments. In functional programming it is possible build complex parts out of simple building blocks, thus making systems more modular. He concludes that the functional paradigm has multiple tools for handling the complexity present in modern software.

Eriksson and rleryd are looking at how to use functional practises when developing front end applications by taking inspiration from Elm in their master's thesis [11]. They have researched each practise in Elm to see if it is possible to use these practises in JavaScript together with tools and libraries. Their conclusion was that it is possible to replicate functional practises from Elm in JavaScript, but that they prefer working with Elm. In JavaScript multiple libraries would have to be used to use the same practises. They also concluded

that even though functional programming is not widely used within the industry, functional practises can still be used in all projects.

In Improving Testability and Reuse by Transitioning to Functional Programming [12], Benton and Radziwill state that functional programming is better suited for test driven development (TDD) and concludes that a shift toward the functional paradigm benefits reuse and testability of cloud-based applications.

Alic, Omanovic and Giedrimas has made a comparative analysis of functional and object-oriented programming languages[15]. They have compared four languages, C#, F#, Java and Haskell based on performance, runtime and memory usage. Their conclusion is that Java is the fastest while Haskell uses much less memory, and that programming paradigms should be combined to increase execution efficiency.

Dobre and Xhafa writes in Parallel Programming Paradigms and Frameworks in Big Data Era that we now are in a big data era[47]. They also review different frameworks, programming paradigms and more in a big data perspective. Around paradigms they state, functional programming (FP) is actually considered today to be the most prominent Programming Paradigm, as it allows actually more flexibility in defining Big Data distributed processing workflows.

In Comparing programming paradigms: an evaluation of functional and object-oriented programs compares R. Harrison, L.G. Samaraweera, M.R. Dobie and P.H. Lewis the quality of code in functional and object oriented programming[51]. To compare these, they use C++ and SML. While their discussion states that they would probably use OOP, since C++ is a better language, the standard list functions was of great help, the debugging was better and reusability was much higher in SML.

Mention concurrency, scala osv? Argument why this subject might be relevant. Functional programming can give more experience, more popular in the industry blablabla. Find studies!

## 3 Method

What we are doin.

### 3.1 Algorithms

In this experiment we will implement four different algorithms. We have chosen algorithms that are well known, so that our focus will be on the different implementations, rather than on how the algorithms work. We will implement

a search tree, the shellsort algorithm, the tower of hanoi algorithm and Dijkstras algorithm. In the search tree we will implement tree traversal, which is a recursive algorithm, as is the tower of hanoi algorithm. These algorithms fit well for FP since it uses recursion. The implementation of the binary tree structure can however be implemented by using classes and objects, that is used in OOP. Shellsort uses state and iteration which is used in OOP and avoided in FP. Dijkstras algorithm also uses state and iteration, and the graph can also be implemented using classes and objects.

We chose algorithms that uses both recursion and iteration to not give advantage to any paradigm. When implementing these algorithms and data structures we can also make use of OOP methods, such as classes and objects. The algorithms purpose are also different to avoid for example implementing four different sorting algorithms.

### 3.1.1 Tree search algorithms

A binary tree is a tree consisting of nodes, where a node can have a maximum of two children [38]. These children can be described as the left and the right subtree and the parent is the root. The property that differs a binary search tree from a standard binary tree, is the order of the nodes. In a binary tree the order can be undecided, but in a search tree the nodes are stored in an order based on some property. For example in our binary search tree the left subtree of a root will contain smaller numbers than the root, and the right subtree will contain larger numbers, see example in figure 1. Using tree traversal it is possible to find certain nodes or get a list of sorted objects. Our binary tree will contain random numbers and the functions:

**findNode(comparable, rootNode)** - Finds a specific key in the tree.

**inOrderTraversal(root)** - Returns an array of all numbers in the tree, sorted smallest to biggest.

**insert(comparable, rootNode)** - Inserts number into the tree. Returns true or false depending on success. If comparable is already in the tree, false should be returned. Note that in the functional implementation this function will return the resulting tree instead of mutating the tree and returning true or false.

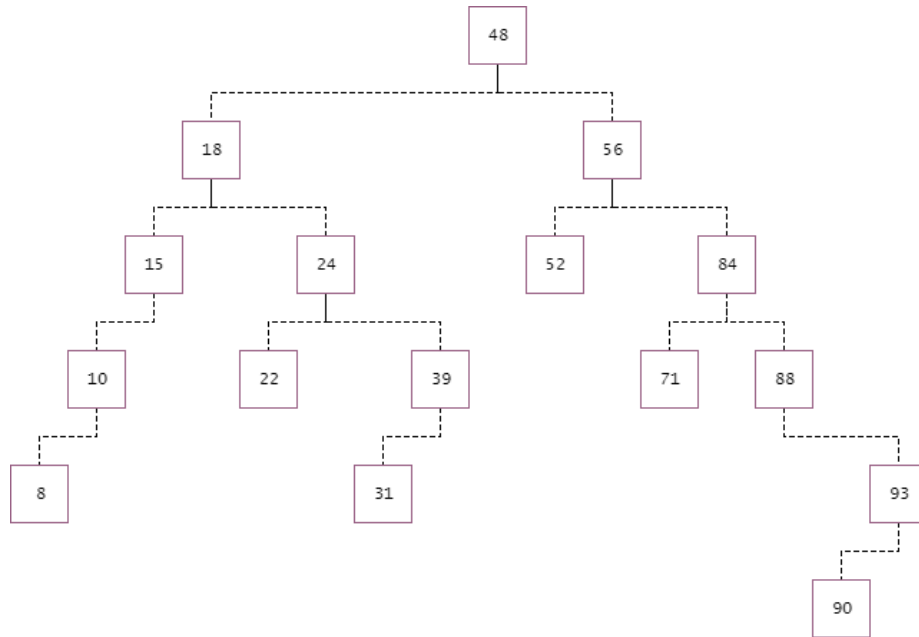


Figure 1: Binary tree example

#### Algorithm descriptions:

**algorithm** findNode(comparable, root)

```

1  if comparable is equal to comparable in root then
2    return root
3  else if comparable is larger than comparable in root
    then
4    return findNode(comparable, rightSubtree in root)
5  else if comparable is smaller than comparable in root
    then
6    return findNode(comparable, leftSubtree in root)
7  end if

```

**algorithm** inOrderTraversal(root)

```

1  set array to empty array
2  if comparable in root is undefined then
3    return array
4  else if comparable in root is not undefined then
5    add inOrderTraversal(leftSubtree in root) to end of
      array
6    add comparable in root at end of array
7    add inOrderTraversal(rightSubtree in root) at end of
      array

```

```

8   return array
9 end if

```

Running this in the tree in Figure 1 would return the array [8, 10, 15, 18, 24, 22, 31, 39, 48, 52, 56, 71, 84, 88, 90, 93].

**algorithm** insert(comparable, root)

```

1  if comparable is equal to comparable in root then
2    return false
3  else if comparable is larger than comparable in root
    then
4    if rightSubtree in root is undefined
5      create newNode
6      set comparable in newNode to comparable
7      set rightSubtree of root to newNode
8      return true
9    else if rightSubtree in root is not undefined then
10     return insert(comparable, rightSubtree in root)
11   end if
12 else if comparable is smaller than comparable in root
    then
13   if leftSubtree in root is undefined then
14     create newNode
15     set comparable in newNode to comparable
16     set leftSubtree in root to newNode
17     return true
18   else if leftSubtree in root is not undefined then
19     return insert(comparable, leftSubtree in root)
20   end if
21 end if

```

### 3.1.2 Shellsort algorithm

The shellsort algorithm is named after its creator Donald Shell [13]. It was one of the first algorithms to break the quadratic time barrier. The algorithm sorts by sorting items using insertion sort with a gap. For each run the gap is decreased until the gap is 1 and the items are sorted. How the gap is decreased is decided with a gap sequence. Different gap sequences gives shellsort a different worst-case running time.

We will use one of Sedgewicks gap sequences that has one of the fastest worst-case running times  $O(n^3/4)$ .

The sequence is

{1, 5, 19, 41, 109}, where the terms are of the form

$4^k - 3 * 2^k + 1$ , or

$9 * 4^k - 9 * 2^k + 1$

In our implementation the sequence is put in an array and not calculated during execution, since we do not want to get different results between our implementation because of the calculation of the gap sequence.

This algorithm will sort an array filled with randomized numbers. Our implementation uses this algorithm in a function:

**shellsort(array)** - Takes an unsorted array as an input and returns a sorted copy of the array.

**Algorithm description:**

**algorithm algorithm shellsort(array of size n)**

```
1  set sortedArray to array
2  set gapSequence to Sedgewick s gap sequence.
3  set currentGapIndex to 0
4  set currentGap to the largest gap i gapSequence where
   gap is smaller than n divided by 2
5  set currentGapIndex to the index of currentGap in
   gapSequence
6  while currentGap is larger than 0 do
7    for i = currentGap to n
8      set currentValue to array[i]
9      set currentIndex to i
10     while currentIndex - currentGap is larger or equal
        to 0 and sortedArray[currentIndex - currentGap] is
        larger than currentValue do
11       set sortedArray[currentIndex] to sortedArray[
        currentIndex - currentGap]
12       set currentIndex to currentIndex - currentGap
13     end while
14     set sortedArray[currentIndex] to currentValue
15   end for
16   set currentGapIndex to currentGapIndex - 1
17   set currentGap to gapSequence[currentGapIndex]
18 end while
19 return sortedArray
```

### 3.1.3 The Tower of Hanoi algorithm

The Tower of Hanoi is a game invented by mathematician douard Lucas in 1883 [14]. The game consists of three pegs and a number of disks stacked in

decreasing order on one of the pegs, see figure 2. The goal is to move the tower from one peg to another by moving one disk at a time to one of the other pegs. A disk can not be placed on a peg on top a smaller disk.

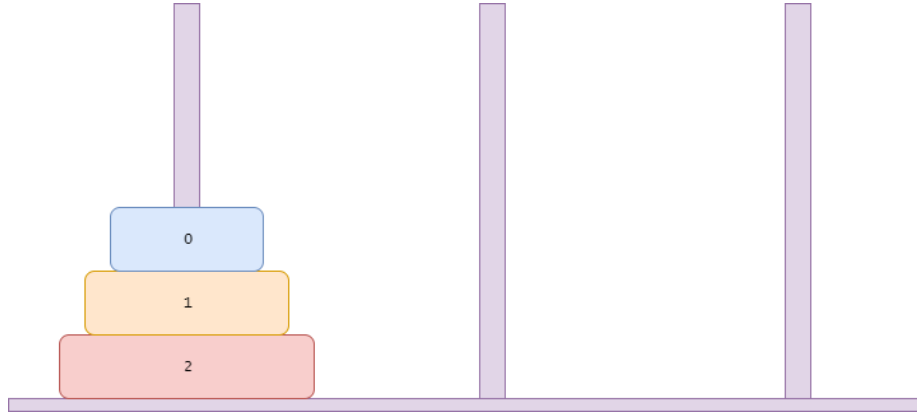


Figure 2: Tower of hanoi image

This algorithm will move a tower from one peg to another in the function:

**hanoi(tower, start, dest, aux)** - Moves tower from start to dest following the rules of the tower of hanoi problem. Returns the start, dest and aux pegs with the repositioned tower.

**Algorithm description:**

**algorithm hanoi(tower, start, dest, aux)**

```

1  (*Pseudo code based on that tower is the number of the
   largest disk, where 0 is the smallest disk in the
   tower.*)
2  if tower is equal to 0
3    move tower from start to dest
4  else
5    hanoi(tower - 1, start, aux, dest)
6    move tower from start to dest
7    hanoi(tower - 1, aux, dest, start)
8  end if

```

### 3.1.4 Dijkstras algorithm

Dijkstras algorithm is an algorithm for finding the shortest path in a graph consisting of a number of nodes connected by edges [18], where the weight of



the edges is known. The algorithm will find the shortest path from the start node to the end node in a graph. It is initiated by:

1. setting the distance to the start node to 0.
2. setting the distance to all other nodes to infinity.
3. mark all nodes as unvisited.

See figure 3. The algorithm will then find the shortest path using the following steps:

1. Set current node to the node with the smallest distance that has not already been visited.
2. For all neighbors to current node that has not already been visited, check if their distance is smaller than the distance of current node + the distance to the neighbor. If so, update the distance of the neighbor.
3. Mark current node as visited.
4. Repeat until all nodes have been visited or the end node has been visited.

See figure 4.

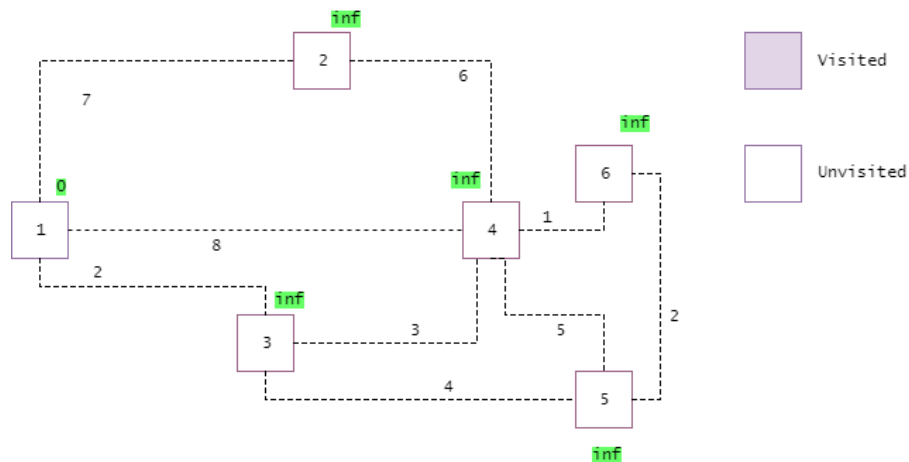


Figure 3: After initiation of Dijkstra's algorithm

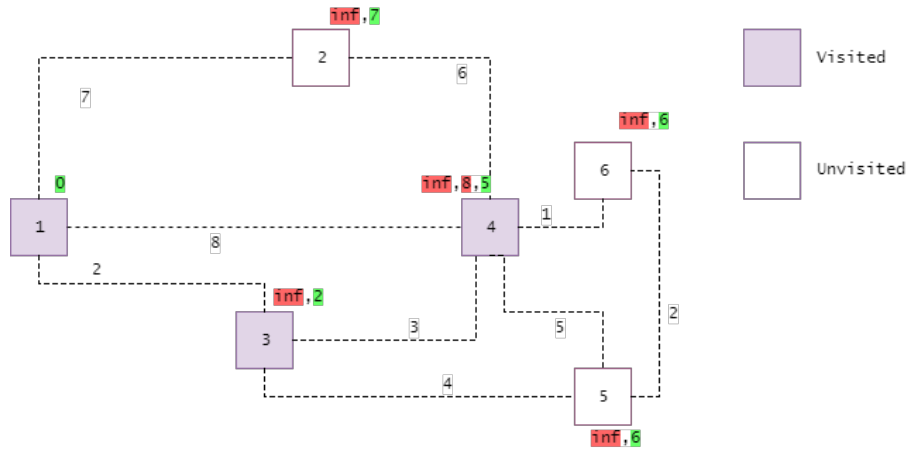


Figure 4: After three nodes have been visited with Dijkstras algorithm

This algorithm will be used in a function:

**dijkstras(graph, startNode, endNode)** - That takes graph, startNode and endNode and returns an array with the shortest path from startNode to destNode.

**Algorithm description:**

**algorithm dijkstras(graph, startNode, endNode)**

```

1  for each node in graph
2    set dist[node] to infinity
3    set path[node] to undefined
4  end for
5  set dist[startNode] to 0
6  set unvisitedNodes to nodes in graph
7  while unvisitedNodes is not empty or endNode is not in
    unvisitedNodes do
8    set current to node where dist[node] is smallest and
    node is in unvisitedNodes
9    remove current from unvisitedNodes
10   for each neighbor of current where neighbor is in
    unvisited do
11     set temp to dist[current] + weight of edge in
    graph, where edge is from current to neighbor
    distance between current and neighbor node
12   if temp is smaller than dist[neighbor] then
13     set dist[neighbor] to temp
14     set path[neighbor] to path[current] + current

```

```

15         end if
16     end for
17 end while
18 return path[endNode]

```

## 3.2 Environment

We will implement the algorithms using Atom, a text editor, and compile and run our tests and our code with Node.js. To manage our dependencies we are using npm. For automation we are using Grunt.

(Add image describing the build of our files)

Since node.js does not support ES6-modules we are using babel to transpile our ES6-modules to node-modules that are supported. We are doing this since we have prior experience with ES6 modules and not with node-modules. Since node.js support almost everything else in ES6 we will only be transpiling the modules, leaving our code in an ES6 standard.

To measure memory and run time we are using node.js process, that is a global and therefore always available within node.js applications [50]. We are using the following functions:

**process.hrtime(time)** - Returns the current high resolution time in a [seconds, nanoseconds] tuple Array. If the optional time-parameter, an earlier hrtime, is used, it will return the difference between that time and the current time.

**process.memoryUsage()** - Returns an object describing the memory usage of the Node.js process measured in bytes.

## 3.3 Testing

Our code will be tested through code reviews and unit testing.

### 3.3.1 Code reviews

We will review each others implementations and our implementations will also be reviewed by a third party. The reviews should help us to find bugs and also to confirm that we have used FP or OOP methods.

### 3.3.2 Unit testing

Unit testing will be done using the JavaScript libraries Mocha [45] for writing tests, Chai [46] for evaluating expressions, and Karma [44] for automated tests.

Our implementations will have to pass the tests in appendix A to be accepted as done.

## 3.4 Implementation

We will proceed from a template with the following structure:

```
|---dist
|   \---<Files generated from babel>
|---src
|   |---js
|       |---<function and class files>
|       \---index.js
|---es6-test
|   \---<Test files>
|---test
|   \---<Test files generated from babel>
| Gruntfile.js
| package.json
```

To run tests use the command:

**npm test** - Will transpile the test files in the es6-test-folder, place the generated files in the test-folder and run the tests.

To run the algorithm use the command:

**npm start** - Will transpile the JavaScript-files in src-folder, place the generated files in the dist-folder and run index.js in the dist-folder.

For an implementation of an algorithm to be considered done the following has to be implemented and provided:

- Tests for the algorithm that are defined in AppendixA placed in `jprojectName;/es6-test`
- A complete implementation of the algorithm that has passed the tests placed in `jprojectName;/src/js`

- A measurement implementation that must run the algorithm and measure time and memory usage using node.js process. For the object-oriented implementation a UML diagram shall also be provided.
- A timelog describing how long the implementation took.

### **3.4.1 Guidelines for writing the different paradigm**

Functional programming guidelines:

- Do not use classes
- Treat functions as variables
- Compose functions to build programs
- Use already provided pure functions when possible
- Write pure functions
- Do not change variables
- Simplify the code

Object-Oriented Programming guidelines:

- Use classes and objects to build programs
- Use inheritance when possible
- Use encapsulation
- Each class should have only one job

## A Appendix A

### A.1 Test cases

ID	T1
Algorithm	Binary search tree algorithms
Function	insert(comparable, rootNode)
Description	Inserted items should be added at the correct place in the tree. If number already exists in the binary search tree, it should not be inserted.
Preconditions	See figure 5
Input	9, node(13)
Expected values	FP: None, OOP: See figure 5
Expected output	FP: See figure 5, OOP: false
Input	8, node(13)
Expected values	FP: None, OOP: See figure 6
Expected output	FP: See figure 6, OOP: true
Input	1, node(13)
Expected values	FP: None, OOP: See figure 7
Expected output	FP: See figure 7, OOP: true
Input	33, node(13)
Expected values	FP: None, OOP: See figure 8
Expected output	FP: See figure 8, OOP: true

ID	T2
Algorithm	Binary search tree algorithms
Function	findNode(comparable, rootNode)
Description	findNode should return the node containing comparable. If comparable is not in the tree undefined should be returned.
Preconditions	See figure 5
Input	5, node(13)
Expected values	None
Expected output	undefined
Input	13, node(13)
Expected values	None
Expected output	node(13)
Input	2, node(13)
Expected values	None
Expected output	node(2)
Input	32, node(13)
Expected values	None
Expected output	node(32)
Input	20, node(13)
Expected values	None
Expected output	node(20)

ID	T3
Algorithm	Binary search tree algorithms
Function	inOrderTraversal()
Description	Should return a sorted array of the elements in the tree. If the tree is it should return an empty array.
Preconditions	See figure 5
Input	FP: node(13), OOP: none
Expected values	None
Expected output	[2, 3, 6, 7, 9, 13, 16, 20, 24, 32]
Preconditions	Empty tree
Input	FP: none, OOP: none
Expected values	None
Expected output	[]

ID	T4
Algorithm	Shellsort
Function	shellsort(array)
Description	If an array of numbers is input a sorted array should be returned. If an empty array is input an empty array should be returned.
Preconditions	None
Input	[]
Expected values	None
Expected output	[]
Input	[ 9, 8, 1, 15, 3, 4, 11, 2, 7, 6]
Expected values	None
Expected output	[ 9, 8, 1, 15, 3, 4, 11, 2, 7, 6]

ID	T5
Algorithm	The Tower of Hanoi
Function	hanoi(tower, start, dest, aux)
Description	Should return start, dest and aux pegs with moved tower. If tower has zero discs it should return empty start, dest and aux. The algorithm should take $2^n - 1$ moves.
Preconditions	None
Input	tower size 8, peg1 with tower, peg3 empty, peg2 empty
Expected values	FP: none, OOP: nrOfMoves=255
Expected output	peg1 empty, peg2 empty, peg3 with tower
Input	tower size 0, peg1 empty, peg3 empty, peg2 empty
Expected values	FP: none, OOP: nrOfMoves=0
Expected output	peg1 empty, peg2 empty, peg3 empty

ID	T6
Algorithm	Dijkstra's algorithm
Function	dijkstras(graph, startNode, endNode)
Description	Should return the shortest path from startNode to endNode. If startNode is same as endNode it should return empty path. If graph is empty it should return empty path.
Preconditions	None
Input	See figure 9, node1, node6
Expected values	None
Expected output	See figure 10
Input	See figure 9, node1, node1
Expected values	None
Expected output	empty path
Input	nodes=[], edges[], noNode, noNode
Expected values	None
Expected output	empty path

## A.2 Test case images

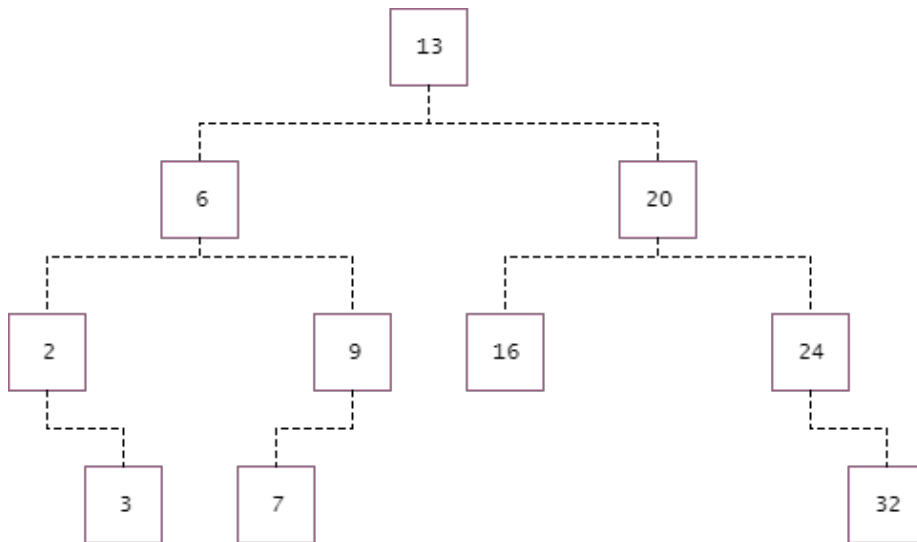


Figure 5: Input tree



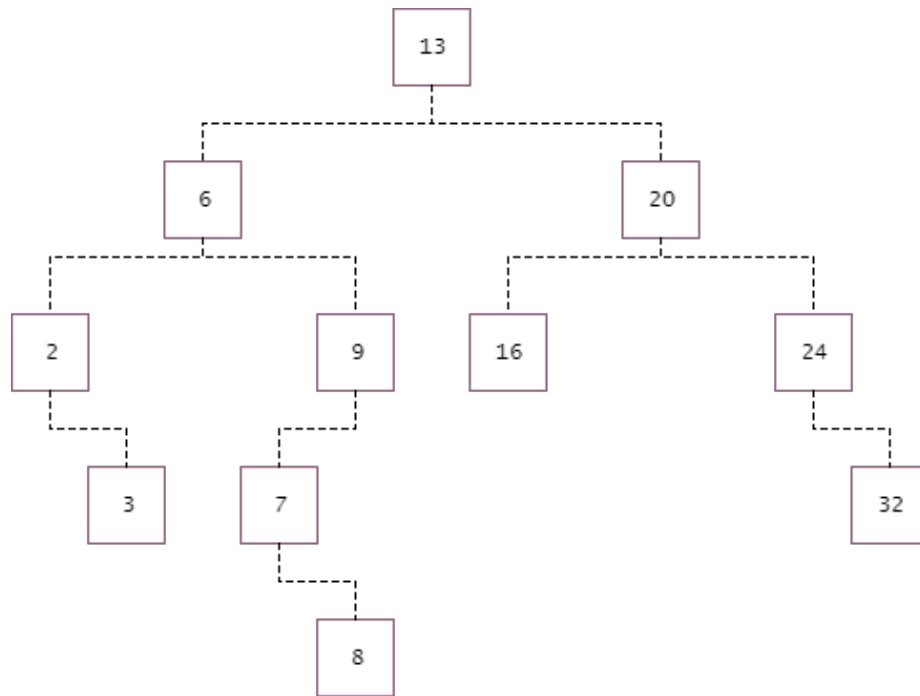


Figure 6: Output tree after inserting 8

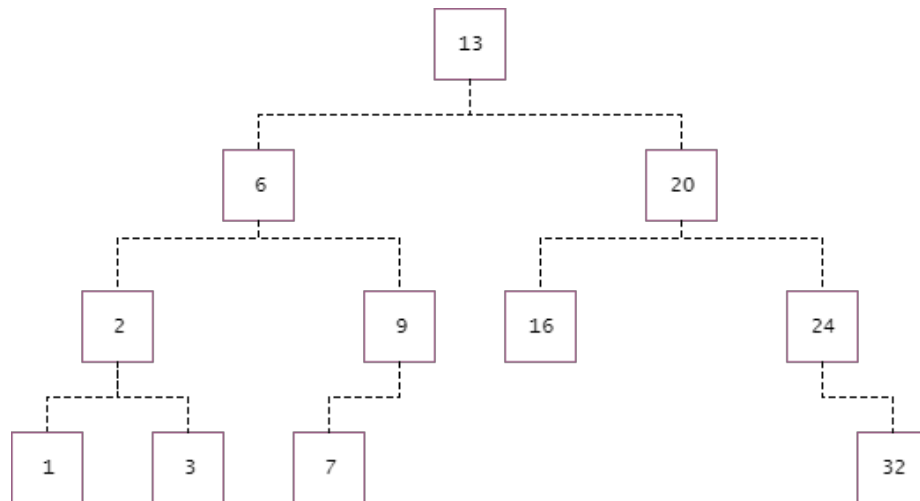


Figure 7: Output tree after inserting 1

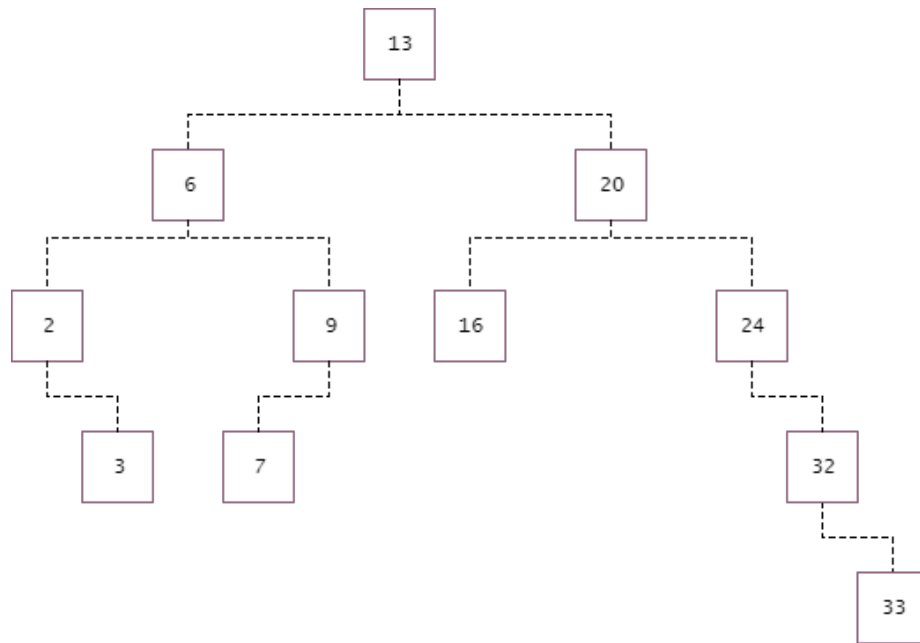


Figure 8: Output tree after inserting 33

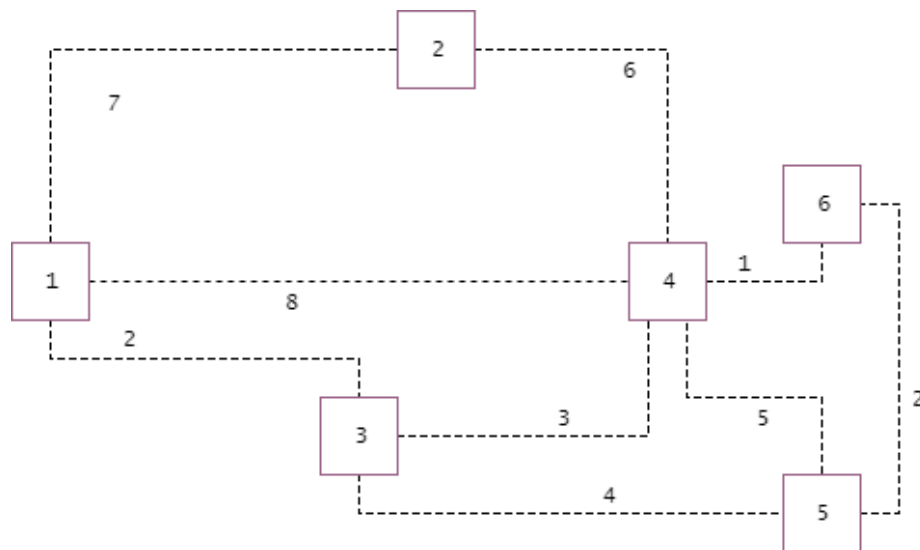


Figure 9: Input graph

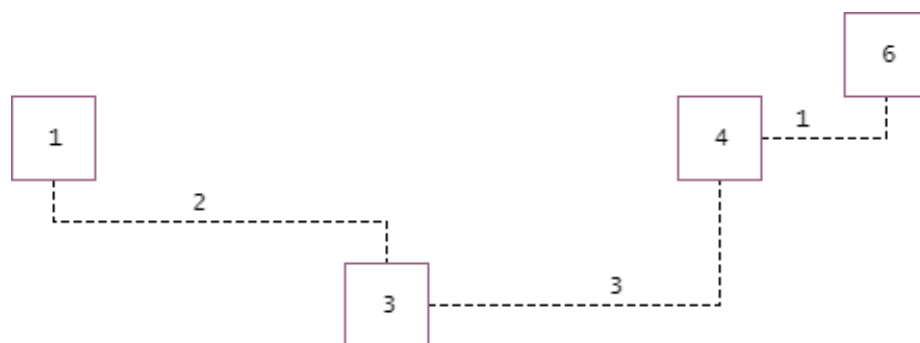


Figure 10: Output graph