(OS) Day7(3)

[Ch2] Process Homework(Simulation)

Exercise 1

1. Runprocess-run.py with the following flags: -1 5:100, 5:100. What should the CPU utilization be (e.g., the percent of time the CPU is in use?) Why do you know this? Use the -c and -p flags to see if you were right.

The CPU executes process0 first. When it finishes executing process0, it will start to execute process1. The total time cost is 10 time units.

```
PS D:\ostep-homework\cpu-intro> python .\process-run.py -1 5:100,5:100 -c -p
              PID: 0
                              PID: 1
                                                   CPU
                               READY
 1
2
3
4
5
6
7
8
9
10
             RUN:cpu
             RUN:cpu
                                READY
             RUN:cpu
                                READY
             RUN:cpu
                                READY
                                READY
             RUN:cpu
                DONE
                             RUN:cpu
                DONE
                             RUN:cpu
                DONE
                             RUN:cpu
                DONE
                             RUN:cpu
                DONE
                             RUN:cpu
Stats: Total Time 10
Stats: CPU Busy 10 (100.00%)
Stats: IO Busy 0 (0.00%)
```

Exercise 2

2. Now run with these flags: ./process-run.py -1 4:100,1:0. These flags specify one process with 4 instructions (all to use the CPU), and one that simply issues an I/O and waits for it to be done. How long does it take to complete both processes? Use -c and -p to find out if you were right.

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CPU will execute process0 first, which has 4 CPU instructions and cost 4 time units.

After finishing process0, it will execute process1, which have one I/O instruction(cost 2 time units) and keep the I/O busy for 5 time units.

Thus, the execution of the two processes will cost 11 time units.

```
PS D:\ostep-homework\cpu-intro> {	t python} .\process-run.py -1 4:100,1:0 -c -p
Time
              PID: 0
                              PID: 1
                                                                    I0s
                                READY
             RUN:cpu
  123456789
             RUN:cpu
                                READY
             RUN:cpu
                                READY
             RUN:cpu
                                READY
                DONE
                              RUN:io
                DONE
                             WAITING
                DONE
                             WAITING
                             WAITING
                DONE
                             WAITING
                DONE
 10
                DONE
                             WAITING
 11*
                DONE
                         RUN:io_done
Stats: Total Time 11
Stats: CPU Busy 6 (54.55%)
Stats: IO Busy 5 (45.45%)
```

Exercise 3

3. Switch the order of the processes: -1 1:0, 4:100. What happens now? Does switching the order matter? Why? (As always, use -c and -p to see if you were right)

Switch order matters.

In this case, the I/O instruction is executed first. CPU can run process1 when process0 is waiting for I/O to be done.

After process1 is finished, there is still 1 time unit left for the I/O to do its job and 1 time unit cost for the CPU to end the I/O.

Thus, the total time cost is 7 time units.

Exercise 4

4. We'll now explore some of the other flags. One important flag is -S, which determines how the system reacts when a process issues an I/O. With the flag set to SWITCH_ON_END, the system will NOT switch to another process while one is doing I/O, instead waiting until the process is completely finished. What happens when you run the following two processes (-1 1:0, 4:100 -c -S SWITCH_ON_END), one doing I/O and the other doing CPU work?

With the -s flag set to SWITCH_ON_END, the CPU is not doing work when the I/O instruction is being executed.

PS C+1	osten-homework\	rnu-intro> nvtho	n \nrocess-rur	n.py -1 1:0,4:100 -c -5 SWITCH_ON	I END
Time	PID: 0	PID: 1	CPU	10s	
1	RUN:io	READY	1		
2	WAITING	READY		1	
3	WAITING	READY		1	
4	WAITING	READY		1	
5	WAITING	READY		1	
6	WAITING	READY		1	
7*	RUN:io_done	READY	1		
8	DONE	RUN:cpu	1		
9	DONE	RUN:cpu	1		
10	DONE	RUN:cpu	1		
11	DONE	RUN:cpu	1		

Exercise 5

5. Now, run the same processes, but with the switching behavior set to switch to another process whenever one is WAITING for I/O (-1 1:0, 4:100 -c -S SWITCH_ON_IO). What happens now? Use -c and -p to confirm that you are right.

Now, process1(the 4 CPU instructions) will be executed while waiting for I/O. T he I/O instruction executes for 1 time unit, then passes the control to process1 and have it execute for 4 time units.

Then, another time unit is spent for I/O. Finally, the control was passed to CPU when the I/O is done.

Exercise 6

6. One other important behavior is what to do when an I/O completes. With -I IO_RUN_LATER, when an I/O completes, the process that issued it is not necessarily run right away; rather, whatever was running at the time keeps running. What happens when you run this combination of processes? (Run ./process-run.py -1 3:0,5:100,5:100,5:100 -S SWITCH_ON_IO -I IO_RUN_LATER -c -p) Are system resources being effectively utilized?

The system resources are not being effectively utilized. After the first I/O finishes, instead of executing the second I/O instruction and then the second CPU instruction, the second CPU instruction was executed.

This caused a waste of time later. During the second I/O instruction is being executed, the CPU just waits without doing any work.

RUN:io MAITING MAITING MAITING MAITING MAITING MAITING READY	READY RUN:cpu RUN:cpu RUN:cpu RUN:cpu RUN:cpu DONE DONE DONE DONE DONE	READY READY READY READY READY READY READY RUN:cpu RUN:cpu RUN:cpu RUN:cpu	READY	1 1 1 1 1 1	1 1 1 1		
MAITING MAITING MAITING MAITING READY	RUN: cpu RUN: cpu RUN: cpu DONE DONE DONE DONE DONE DONE DONE	READY READY READY READY RUN:cpu RUN:cpu RUN:cpu	READY READY READY READY READY READY	1 1 1 1 1	1 1 1		
MAITING MAITING READY	RUN: cpu RUN: cpu RUN: cpu DONE DONE DONE DONE DONE DONE	READY READY READY RUN:cpu RUN:cpu RUN:cpu	READY READY READY READY READY	1 1 1 1 1	1 1		
MAITING MAITING READY	RUN: cpu RUN: cpu DONE DONE DONE DONE DONE DONE	READY READY RUN:cpu RUN:cpu RUN:cpu	READY READY READY READY	1 1 1 1	1		
NAITING READY	RUN: cpu DONE DONE DONE DONE DONE DONE	READY RUN:cpu RUN:cpu RUN:cpu	READY READY READY	1 1			
READY READY READY READY READY READY READY	DONE DONE DONE DONE DONE	RUN:cpu RUN:cpu RUN:cpu	READY READY	1 1	1		
READY READY READY READY READY READY	DONE DONE DONE DONE	RUN:cpu RUN:cpu	READY	1			
READY READY READY READY READY	DONE DONE DONE	RUN:cpu					
READY READY READY READY	DONE DONE		READY				
READY READY READY	DONE	RUN:cpu		1			
READY READY			READY	1			
READY	DOME	RUN:cpu	READY	1			
	DONE	DONE	RUN:cpu	1			
	DONE	DONE	RUN:cpu	1			
READY	DONE	DONE	RUN:cpu	1			
READY	DONE	DONE	RUN:cpu	1			
READY	DONE	DONE	RUN:cpu	1			
io_done	DONE	DONE	DONE	1			
RUN:io	DONE	DONE	DONE	1			
WAITING	DONE	DONE	DONE		1		
WAITING	DONE	DONE	DONE		1		
WAITING	DONE	DONE	DONE		1		
WAITING	DONE	DONE	DONE		1		
WAITING	DONE	DONE	DONE		1		
io done	DONE	DONE	DONE	1			
RUN:io	DONE	DONE	DONE	1			
WAITING	DONE	DONE	DONE		1		
WAITING	DONE	DONE	DONE		1		
WAITING	DONE	DONE	DONE		1		
WAITING	DONE	DONE	DONE		1		
WAITING	DONE	DONE	DONE		1		
				1			
	O_done RUN:io IAITING	AC_done DONE RUN:io DONE RUN:io DONE RAITING DONE ROUGHOUSE	ACO_DONE DONE DONE RUN:io DONE DONE DONE DONE DONE DONE DONE DONE	RUN:io DONE DONE DONE ALTING DONE DONE DONE BALTING DONE DONE DONE	DONE DONE	DONE DONE	DONE DONE

Exercise 7

7. Now run the same processes, but with -I IO_RUN_IMMEDIATE set, which immediately runs the process that issued the I/O. How does this behavior differ? Why might running a process that just completed an I/O again be a good idea?

Now, the CPU is fully utilized.

me	PID: 0	PID: 1	PID: 2	PID: 3	CPU	I0s	
1	RUN:io	READY	READY	READY	1		
2	WAITING	RUN:cpu	READY	READY	1	1	
3	WAITING	RUN:cpu	READY	READY	1	1	
4	WAITING	RUN:cpu	READY	READY	1	1	
5	WAITING	RUN:cpu	READY	READY	1	1	
6	WAITING	RUN:cpu	READY	READY	1	1	
7*	RUN:io_done	DONE	READY	READY	1		
8	RUN:io	DONE	READY	READY	1		
9	WAITING	DONE	RUN:cpu	READY	1	1	
0	WAITING	DONE	RUN:cpu	READY	1	1	
1	WAITING	DONE	RUN:cpu	READY	1	1	
2	WAITING	DONE	RUN:cpu	READY	1	1	
3	WAITING	DONE	RUN:cpu	READY	1	1	
4*	RUN:io_done	DONE	DONE	READY	1		
5	RUN:io	DONE	DONE	READY	1		
6	WAITING	DONE	DONE	RUN:cpu	1	1	
7	WAITING	DONE	DONE	RUN:cpu	1	1	
8.	WAITING	DONE	DONE	RUN:cpu	1	1	
9	WAITING	DONE	DONE	RUN:cpu	1	1	
20	WAITING	DONE	DONE	RUN:cpu	1	1	
1*	RUN:io done	DONE	DONE	DONE	1		

Exercise 8

8. Now run with some randomly generated processes: -s 1 -1 3:50, 3:50 or -s 2 -1 3:50, 3:50 or -s 3 -1 3:50, 3:50. See if you can predict how the trace will turn out. What happens when you use the flag -I IO_RUN_IMMEDIATE vs. -I IO_RUN_LATER? What happens when you use -S SWITCH_ON_IO vs. -S SWITCH_ON_END?

When IO_RUN_IMMEDIATE is set, the CPU will run another I/O instruction right after the current instruction finishes, and thus fully utilizes the time waiting for the data to be fetched.

When switch_on_io is set, the CPU transfers control to another process when an I/O instruction is invoked.