

# Hybrid Cost Model for Input Row Activation

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## 1 Introduction

Accurately estimating DRAM Row Activations for Input Feature Maps is critical for performance optimization. Simple analytical models (e.g., assuming contiguous access) fail to capture complex interactions between:

- **Tile Alignment:** The relative offset of a tile within a DRAM Row.
- **Loop Order:** The specific traversal order (e.g., Row-Major vs. Block-Major) within a tile.
- **Row Thrashing:** The penalty incurred when a single tile repeatedly crosses DRAM Row boundaries.

This document describes a **Hybrid Cost Model** that combines **Number Theory (GCD)** for efficient quantity estimation and **Micro-Trace Simulation** for accurate unit cost calculation.

## 2 The Hybrid Approach: Categorization & Sampling

The total Input Row Activation cost is derived by categorizing tiles into two states: **Safe** (aligned within a row) and **Crossing** (straddling a row boundary).

$$\text{TotalCost} = N_{\text{safe}} \cdot C_{\text{safe}} + N_{\text{crossing}} \cdot C_{\text{crossing}} \quad (1)$$

### 2.1 1. Cost Determination ( $C$ )

We use **Micro-Trace Simulation** to determine the unit costs.

- **Safe Cost ( $C_{\text{safe}}$ ):** The cost when the tile is perfectly aligned (Offset = 0).

$$C_{\text{safe}} = \text{MicroTrace}(0) \quad (2)$$

- **Crossing Cost ( $C_{\text{crossing}}$ ):** The expected cost when the tile is in the “Danger Zone”. We sample offsets specifically from the crossing region.

$$C_{\text{crossing}} = \frac{1}{|\mathcal{O}_{\text{cross}}|} \sum_{o \in \mathcal{O}_{\text{cross}}} \text{MicroTrace}(o) \quad (3)$$

### 2.2 2. Quantity Determination ( $N$ )

Instead of simulating every tile to check if it crosses, we use **Number Theory** to calculate the exact counts.

Let  $P$  be the Period (Row Buffer Size) and  $S$  be the Step (Stride in linear memory). The “Danger Zone”  $Z$  is the set of offsets that cause a tile of width  $W_{\text{tile}}$  to cross a boundary:

$$Z = [P - W_{\text{tile}} + 1, P - 1] \quad (4)$$

The sequence of offsets is  $o_i = (i \cdot S) \pmod{P}$ . This sequence visits the set of values  $\{0, g, 2g, \dots\}$  where  $g = \text{gcd}(S, P)$ . The probability of landing in the Danger Zone is:

$$\text{Prob}_{\text{crossing}} = \frac{\text{Count of multiples of } g \text{ in } Z}{P/g} \quad (5)$$

Thus, the exact quantities are:

$$N_{crossing} = N_{total} \times Prob_{crossing} \quad (6)$$

$$N_{safe} = N_{total} - N_{crossing} \quad (7)$$

### 3 Micro-Trace Simulation

To determine the cost of a single tile at a specific offset  $o$ , we use a lightweight simulation.

#### 3.1 Strict Ordering (Realistic Hardware)

We simulate the exact memory access sequence generated by the hardware loop nest (e.g., Row-Major  $H, W$ ).

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**Algorithm 1** Micro-Trace Simulation (Strict Order)

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ActiveRow  $\leftarrow -1$ 
Activations  $\leftarrow 0$ 
for  $h \in [0, TileH - 1]$  do
  for  $w \in [0, TileW - 1]$  do
    Addr  $\leftarrow BaseAddr + o + h \cdot W_{total} + w$ 
    RowID  $\leftarrow Addr // RowBufferSize$ 
    if  $RowID \neq ActiveRow$  then
      Activations  $\leftarrow Activations + 1$ 
      ActiveRow  $\leftarrow RowID$ 
    end if
  end for
end for
return Activations

```

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This captures **Row Thrashing** where a tile straddles a row boundary and the access pattern causes repeated switching (e.g.,  $A \rightarrow B \rightarrow A \rightarrow B$ ).

### 4 Validation Results

We compared the Hybrid Model against a full exhaustive simulation of all tiles.

| Scenario     | Params (H, W, Tile, Stride) | Ground Truth | Hybrid Model | Error |
|--------------|-----------------------------|--------------|--------------|-------|
| Aligned      | 100, 1024, 3x3, 1           | 3.0000       | 3.0000       | 0.00% |
| Misaligned   | 224, 224, 3x3, 1            | 1.4324       | 1.4375       | 0.35% |
| Large Stride | 224, 224, 3x3, 2            | 1.4324       | 1.4375       | 0.35% |
| Large Tile   | 224, 224, 7x7, 2            | 2.3119       | 2.3125       | 0.02% |

Table 1: Validation of Hybrid Cost Model vs. Exhaustive Simulation

### 5 Implementation in ILP

The ILP optimizer will use this model to precompute a lookup table or cost function:

1. For a candidate mapping  $(TileH, TileW)$ , calculate  $E[Cost]$ .
2. Use  $E[Cost]$  as the coefficient for the Input Row Activation term in the objective function.
3. This ensures the optimizer penalizes mappings that result in high thrashing probabilities.