SOLUTIONS AND MARKING SCHEME



ECS629U-759P: ARTIFICIAL INTELLIGENCE

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Coursework 1 (12%): Search Problems
Winter 2020

DUE DATE: PLEASE REFER TO THE EXACT DATE FROM QMPLUS NO LATE submission is allowed: submission link will disappear upon deadline!

Instructions: The coursework is composed of three questions. The first two involve coding exercises. The last one is a writing exercise about adversarial search. You are expected to submit two files ONLY through the QM+ submission link:

- 1. A single **PDF** file that must be **at most 5 pages long**. Any file submitted under the wrong format (e.g. doc, docx, odt, etc) or exceeding the length limit may be penalised. The PDF can include scans of handwritten material as long as it is legible but a typewritten version is encouraged (even better if you use LATEX). **ATTENTION: Do not include the text of the question in your answers. Failure to do so may disqualify your submission!**
- 2. A *single compressed folder* containing all your (well documented and sufficiently commented) codes. The codes can be in any language, although Python/java/C/julia/JS/Lisp are preferred in that order. Your codes should run (no bugs!). There must be an easy to follow readme.txt file inside the folder.

It is allowed to use codes from online resources. However, this has to be clearly cited with reference in your report. Collaboration is <u>NOT</u> permitted when attempting the answers. <u>High level</u> discussion when not attempting the questions may be fine, but discouraged. A better means is using our discussion forum on QM+. There is zero tolerance policy for cases of plagiarism. If in doubt, ask! Please be aware that systems can be busy and slow to respond shortly before deadlines. So you should aim to submit at least one hour before the announced deadline.

Question 1: Let's Travel! (Uniform Cost Search).....

A friend of yours gave you a map of UK, but encoded as a json file: UK_cities.json. The file encodes all the roads between major cities of the UK with their length of the road (in km). After numerous arguments, you both decided to go from London to Aberdeen. Although you would like to go there while visiting as many cities as possible, neither you nor your friend can afford the time! You are then trying to reach Aberdeen from London taking the shortest path (the total driving distance).

(a) (15 points) Following the UCS algorithm, manually execute it for three iterations for this problem. As a reminder, the pseudo-code of the UCS is provided in Figure 1, which matches Figure 3.13 in the textbook (3rd edition).

```
function UNIFORM-COST-SEARCH(problem) returns a solution, or failure
  node \leftarrow a node with State = problem.Initial-State, Path-Cost = 0
  frontier \leftarrow a priority queue ordered by PATH-COST, with node as the only element
  explored \leftarrow an empty set
  loop do
      if EMPTY?(frontier) then return failure
      node \leftarrow Pop(frontier) /* chooses the lowest-cost node in frontier */
      if problem.GOAL-TEST(node.STATE) then return SOLUTION(node)
      add node.STATE to explored
      for each action in problem.ACTIONS(node.STATE) do
          child \leftarrow \text{CHILD-NODE}(problem, node, action)
         if child.State is not in explored or frontier then
             frontier \leftarrow Insert(child, frontier)
         else if child.State is in frontier with higher Path-Cost then
             replace that frontier node with child
```

Figure 1: Pseudo-code for the USC (Uniform-Cost Search) algorithm on a graph. This is borrowed from [1, Fig. 3.13].

- By iteration, we are referring to the outer "loop". So, in your trace, at least three nodes should be selected for expansion.
- In particular, you should provide (only) the content of the following variables, with the following specific data structure:
 - (i) the "frontier": a queue of "nodes" ordered by their path-cost
 - (ii) the "explored": a set of "states"
- You should represent each "node" in the frontier as a tuple of the following information: ("state", "path-cost", "best-path-to-reach-that-state"). Note: the "state" in our problem is just the name of the city you are at! The "path-cost" is just a positive real number; and the "best-path-to-reach-that-state" is a list of states starting from the initial state to that node's state.
- For the trace, you should provide only one value per each line. Write down only the new value each time any of the above two variables change. The order definitely matters.
- If you encounter any tie-breaking situation in the executing, e.g. in adding the nodes to the frontier, go with the alphabetical (lexicographical) order based on the name of the city.
- No explanation or comments are necessary. For instance, you don't even need to tell which line/part of the pseudo-code was responsible for an update. We only check whether the trace is correct (including the order of updates).

• An example few first steps of the trace with the above description would be as follows:

```
1 frontier = [('london', 0, ['london'])]
2 \text{ explored} = \{\}
3 \text{ frontier} = []
4 explored = { 'london '}
5 frontier = [('birmingham', 110,['london','birmingham'])]
6
7
  . . .
8 ...
```

Solution: Here is a trace at the outer loop:

```
Frontier is: [[0, 'london', None]]
Frontier len is: 1
Explored nodes are: {'london'}
Explored nodes len is: 1
Current node is: london
Frontier is: [[52, 'brighton', ['london', None]],
\hbox{\tt [54, 'cambridge', ['london', None]],}\\
[57, 'oxford', ['london', None]],
[71, 'dover', ['london', None]],
[110, 'birmingham', ['london', None]],
      'bristol', ['london', None]],
[161, 'cardiff', ['london', None]],
[172, 'exeter', ['london', None]],
[172, \ 'hull', \ ['london', \ None]],
[198, 'leeds', ['london', None]],
      'liverpool', ['london', None]],
[302, 'carlisle', ['london', None]], [396, 'glasgow', ['london', None]]]
Frontier len is: 13
Explored nodes are: {'brighton', 'london'}
Explored nodes len is: 2
Current node is: brighton
Frontier is: [[54, 'cambridge', ['london', None]],
[57, 'oxford', ['london', None]],
[71, 'dover', ['london', None]],
[101, 'portsmouth', ['brighton', ['london', None]]],
[110, 'birmingham', ['london', None]],
      'bristol', ['london', None]],
[116,
[161, 'cardiff', ['london', None]],
[172, 'exeter', ['london', None]],
       'hull', ['london', None]],
       'leeds', ['london', None]],
       'liverpool', ['london', None]],
       'nottingham', ['brighton', ['london', None]]],
'sheffield', ['brighton', ['london', None]]],
[230,
[268,
       'swansea', ['brighton', ['london', None]]],
       'aberystwyth', ['brighton', ['london', None]]],
[302, 'carlisle', ['london', None]],
[302, 'york', ['brighton', ['london', None]]],
```

```
[329, 'penzance', ['brighton', ['london', None]]],
[396, 'glasgow', ['london', None]]]
Frontier len is: 19
Explored nodes are: {'cambridge', 'brighton', 'london'}
Explored nodes len is: 3
Current node is: cambridge
```

Marking scheme: 15 points. You could instead provide the trace with the inner loop. Both were accepted.

- (b) (10 points) Implement the UCS for this problem. Run it with the provided data. You don't need to include the code in your report (you submit it separately). Instead:
 - (b-1) Provide the optimal path found by the algorithm (both the path and its length);
 - (b-2) Make the code print the trace (in the same format as requested in part (a) but only provide the *last two* iterations (of the outer loop) in your report.

Note: you don't have to follow the provided pseudo-code in your implementation necessarily. However, if the logic of your code deviates significantly from this version, then you need to provide your pseudo-code as well.

Solution: An example implementation of the UCS is provided in ucs.py (find it in the same place you found this solution sheet on QM+).

It turns out that you have multiple optimal paths: [london, cambridge, york, aberdeen], [london, cambridge, york, edinburgh, aberdeen] and [london, cambridge, york, newcastle, edinburgh, aberdeen], all of which share the same total cost of 502.

The trace of the last two epochs would look something like the following:

```
Frontier is: [[387, 'edinburgh',
['york', ['cambridge', ['london', None]]]],
[396, 'glasgow', ['london', None]]]
Frontier len is: 2
Explored nodes are: {'exeter', 'manchester', 'newcastle', 'penzance',
'cambridge', 'bristol', 'swansea', 'aberystwyth', 'liverpool', 'leeds', 'brighton', 'birmingham', 'carlisle', 'cardiff', 'hull', 'sheffield',
'london',
             'dover', 'york', 'edinburgh', 'portsmouth', 'nottingham',
'oxford'}
Explored nodes len is: 23
Current node is: edinburgh
Found solution for [502, 'aberdeen',
['edinburgh', \ ['york', \ ['cambridge', \ ['london', \ None]]]]]\\
Frontier is: [[396, 'glasgow', ['london', None]]]
Frontier len is: 1
Explored nodes are: {'exeter', 'manchester', 'glasgow', 'newcastle', 'penzance', 'cambridge', 'bristol', 'swansea', 'aberystwyth', 'liverpool', 'leeds', 'brighton', 'birmingham', 'carlisle', 'cardiff', 'hull', 'sheffield', 'london', 'dover', 'york',
'edinburgh', 'portsmouth', 'nottingham', 'oxford'}
Explored nodes len is: 24
Current node is: glasgow
```

Marking scheme: 10 points. Providing one of the optimal path was enough to get full mark for the first part of the question.

- (c) (10 points) Suppose both you and your friend are environmentally aware. So you would like to find the path that has the lowest overall cost, which is the sum of the overall driving time plus the overall cost of the air pollution due to your driving. Here are the details:
 - You are free to choose your driving speed on each road of the path (but assume that you don't change your speed on a given road). So e.g. if you choose to drive at 40 km/h on the road between London to Birmingham, you maintain that speed on this road. But you are free to choose a different speed on the next road, etc.).
 - If you drive at speed v km/h (on any road), you are responsible for an air pollution cost of $(0.00001v^2)$ units *per hour*.

Using your answer to the previous part, provide the best path in this scenario (both the optimal cost and the path). Again, you don't need to provide the code in your report, but you need to briefly explain your approach.

Solution: The overall cost of a path here is the total time plus the total air pollution cost. But this is equal to the driving times plus the air pollution cost summed over each edge of the path independently. Mathematically, if we consider a path from an initial node to a goal node to be (e_1, e_2, \ldots, e_n) , were e_i is the i'th node, then cost(path) can be broken down to $cost(e_1) + cost(e_2) + \ldots + cost(e_n)$. Therefore, instead of having to search in the domain of the entire paths, we can firs find the optimal edge weights for each edge in the network (more accurately, for each edge that the algorithm needs), and then find the optimal path using those optimal weights.

So now consider a general edge e, that is a road between two cities. Let the original weight of that edge be denoted by l, representing the length of that road. Suppose we choose to drive on the road with velocity v. The new cost of that edge would be:

$$cost(e) = \frac{l}{v} + \alpha v^2 \times \frac{l}{v} = (\frac{1}{v} + \alpha v)l \tag{1}$$

Some of you overlooked the fact that αv^2 (where $\alpha=0.00001=10^{-5}$) was given as the *per hour* cost of air pollution; so to get the air pollution cost of that edge, we need to multiply that by driving time l/v. Note that l here is not a variable, that is just the distance between those two cities at the both ends of that edge, what is a variable here is v, the driving speed, which the problems says we can choose per each edge.

There are no limits set for us (except that of course the speed cannot be negative, and here, it cannot be zero either!, so v>0). So in order to find the optimal speed, we can take the derivative of the function $\frac{1}{v}+\alpha v$ with respect to v and set it equal to zero:

$$\frac{d}{dv}\left(v^{-1} + \alpha v\right) = -v^{-2} + \alpha = 0$$

The valid solution to this is $v^* = \sqrt{1/\alpha}$, which if you replace for $\alpha = 10^{-5}$, gives you $\sqrt{10^5} \approx 316.2$ km/h. Technically, to be sure that what we have found is indeed a minimiser value (and not a maximiser), we should also check the second derivative to make sure it is positive (so the function around the zero-derivative point looks

like a cup, not an upside-down cup!). The second derivative is v^{-3} , which is always positive for any v > 0, including our v^* . So, what is the optimal weight of that edge? All we need to do is to plug in v^* in (1):

optimal
$$cost(e) = (\frac{1}{v^*} + \alpha v^*)l = (2\sqrt{\alpha})l \approx (0.006325)l$$

At this point, you can just re-run the UCS with the difference that instead of the previous weights of each edge, you scale it with 0.006325. However, if you are observant, you can notice that an interesting thing has just happened(!): the optimal weight of each edge has became just a constant scaling of l, which was the previous weight of that edge. So, we in fact don't even need to re-run the UCS as we already know what the optimal solution is! The optimal path will be just as before! In fact, we also know what the overall cost is going to be: it will be just a re-scaling of the previous overall cost, it will be $0.006325 \times 502 \approx 3.175$ units.

Fun-fact/disclaimer(!): you should of course just follow the assumptions of the question, which is no speed limit implied on the roads. That said, here the optimal speed of 316.2 km/h on each road is obviously impractical. This is because of a typo (mi scuzi!): the α was supposed to be 0.0001 as opposed to 0.00001, which gives an optimal speed of 100 km/h, a more realistic/practicable value!

Marking scheme: 10 points. If you don't argue why the optimal path does not change from before, if everything else is perfect, you will get 7/10. (Note: it is not trivial that this would be the case: since the weight of each edge is changing, then the optimal path could have been different! It just happened (fortuitously!) in this specific problem that the new weights were just a constant scaling of the previous weights, with the same scaling for all edges.) If you don't make the observation that you don't have to re-run the UCS but everything else is perfect, you get 9/10. If you omit the multiplication by time in computation of the air pollution, but everything else is perfect, you get 6/10. If you don't do any analysis and just use a brute-force to find the best speed per each link, but everything else is perfect, you get 7/10. If you don't find the best speed per each link, but rather choose an arbitrary or random value, then mark is capped at 4.

- (d) (10 points (bonus)) Now, for a bonus 10 points!, suppose you have a super-car and you are environmentally ignorant! Suppose:
 - As in the previous part, you can choose your driving speed per each road.
 - However: suppose each road now has a speed-limit. For simplicity, assume that the value of the speed limit is the same as the value provided for the length of that road in the json file.
 - On each road, you can even choose to go over the speed limit! However, there is a likelihood that you get caught on a speed-camera and get a fine of 1000 units. In particular, if you drive on a road with speed limit of $v_{\rm lim}$ at speed v, then the likelihood (i.e., probability) of getting the fine is as follows:

$$\begin{cases} 1 - e^{-(v - v_{\text{lim}})}, & \text{if } v > v_{\text{lim}} \\ 0, & \text{if } v \le v_{\text{lim}} \end{cases}$$

 You may only be fined at most once on a given road, but fines on different roads accumulate. Again, using your UCS code, find the best path in this problem, assuming the overall cost is the total travel time in hours plus your total fines (if any).

Solution: As in part c, the overall cost of a path can be broken down as the sum of costs per each edge that constitutes that path. So again, we can find the optimal speed per each edge, and then apply our UCS algorithm with the new weights to find the optimal path (or get lucky again, and somehow find the optimal path without having to run UCS again).

The cost of an edge with length l when we drive at speed v on that road (where $0 < v \le 300$ km/h), would be:

optimal
$$cost(e) = 100 \times \frac{1}{v} + 1000 \times \begin{cases} 1 - e^{-(v - v_{lim})}, & \text{if } v > v_{lim} \\ 0, & \text{if } v \le v_{lim} \end{cases}$$

This is because $100 \times (l/v)$ is the rental fee of the car pertaining to that road, and the second part is the *expected value* of the fine we would receive on that road. Note that whether or not we receive a fine on a road (if we drive above speed limit) is a random variable. In order to be able to meaningfully compare different options, we have to convert it to a non-random value. The most obvious candidate for this is to use the expected value (the average value if we would have to repeat this scenario many times). The expected value of the fine, is just 1000 (the ticket value) times the probability that we get a ticket.

This is also a function of just v, and we should just find a speed that minimises this per each road. First, we can make the observation that there is no point in driving below the speed limit on the road (unless of course the speed limit is above 300 km/h, which is the maximum speed of our super-car). This is because up to the speed limit, we will be reducing the first part (the rental fee of the car) without adding any expected fine (which is zero). So the only question is, will it be worth it to go above the speed limit, as it will reduce the first part, but starts to increase the second part.

Now either analytically, or through brute-force search, or by using any numerical optimisation library like <code>scipy.optimize</code>, you can run the USC with the weight of each edge being passed as the result of this single-variable optimisation (and that would have scored you a perfect mark if you indeed get the correct answer, which is 300 units).

However, with a little bit of analysis, you can also show that with the given parameters, it is never worth going over the speed limit (the reduction in the rental fee never outweighs the expected fine for any road). So for the given parameters of the problem, it is always optimal to just stick with the speed limit. Now, because we are told to assume the speed limit is the same as the distance, this means the cost per each edge is just 100×1 , i.e., 100. So the optimal solution becomes any path that can get us from London to Aberdeen with taking as few roads as possible. The lowest number of roads is three, which gets you the overall cost of 300.

Marking scheme: 10 points. The perfect score is reserved for a combination of the correct answer and correct rich explanation and/or code. Partial scores possible depending on the quality of the answer and/or code. If instead of using the expected value, you are generating the fines randomly (according to the probabilities), then your score is capped at 4 (because such analyses would not make much sense!).

Question 2: It's on the tip of my tongue! (Genetic Algorithm).....

Becoming old is certainly difficult. And one of the signs of growing older is some form of amnesia. As a matter of fact, I cannot exactly remember my two favourite passwords, but still, I have a hazy idea about what it looked like!

Your goal is to help me remember these two passwords using a genetic algorithm: It turns our even though I am amnesiac, I can still tell how close to my password a candidate is. Let me give you the few details I remember with certainty: the password was exactly 10 characters long and could only contain upper-case English letters, numbers between 0 and 9, and the underscore symbol _. For instance, 1_PASSWORD and MYPA55WORD are legitimate candidates.

So here is the deal: For each password you show me, I am going to give a score between 0.0 and 1.0 indicating how close to my forgotten password I feel it is (where 1.0 means that the password is exactly correct) and despite my memory, I have a trustworthy intuition!

To be more precise, you have access to a function called blurry_memory as follows:

input: a list of candidate passwords, e.g. ["___TEST____", "IS_I7_L0V3"], along with your own student number (e.g 190123123), and the index of the password you are trying to find: 0 or 1 (recall: there are two passwords to recover).

output an ordered dictionary containing the passwords as keys and the associated scores as values (e.g "____TEST____": 0.045, "IS_I7_L0V3": 0.247)

To access this function, you can add the following line to the header of your python3 code:

```
from amnesiac import blurry_memory
```

Implement a Genetic Algorithm to recover the two forgotten passwords.

What should you provide in your report:

(a) (10 points) What did you get for the two passwords? (recall: you should use your own student ID number whenever invoking blurry_memory).

Solution: A vanilla version of what has been asked can be found in algo_gen.py. You can run it by changing the student number.

Here is the table containing the passwords of each student:

```
+----+
 Student | Password1 | Password2 |
| 170928388 | 6RUMPY_C47 | _C0NGR47S5 |
170368517 | Y0UDON754Y | 9_FR33D0M8 |
| 181122056 | 1_L1K3TH4T | J0HN__5NOW |
| 190084873 | _4_3VER4I_ | 9_FR33D0M8 |
| 150047755 | F1ND_NE55Y | 6RUMPY_C47 |
| 170343949 | M1N_I_M4X_ | F1ND_NE55Y |
| 190583055 | 1_H4T3_DF5 | F1ND_NE55Y |
| 190028305 | 1_L1K3TH4T | WE1L_D0N3_ |
| 190517522 | CHUKN088IS | 1_H4T3_DF5 |
| 160205075 | J0HN__5NOW | D0GW0OWFUN |
| 190734613 | _4_3VER4I_ | CHUKN088IS |
```

```
| 160486423 | 9_FR33D0M8 | CHUKN088IS |
| 170434328 | _CONGR47S5 | 6RUMPY_C47 |
| 170247197 | CHUKN088IS | JOHN__5NOW |
| 190262046 | 1_H4T3_DF5 | WE1L_D0N3_ |
| 190778145 | 1_L1K3TH4T | 9_FR33D0M8 |
| 170049926 | K4NG_OUF__ | WE1L_D0N3_ |
| 160580134 | M1N_I_M4X_ | B3L1K3B1LL |
| 170709031 | _Y0UM4DBRO | F1ND_NE55Y |
| 170980700 | M1N I M4X | Y0UDON754Y |
| 170172202 | _CONGR47S5 | 1_H4T3_DF5 |
| 190572587 | 1_H4T3_DF5 | 7ROLL_F4CE |
| 170074157 | 9_FR33D0M8 | B3L1K3B1LL |
| 161125684 | B3L1K3B1LL | 9_FR33D0M8 |
| 170639668 | _Y0UM4DBRO | 9_FR33D0M8 |
| 190536501 | 6RUMPY_C47 | M1N_I_M4X_ |
| 160365449 | 7ROLL_F4CE | Y0UDON754Y |
| 190347321 | Y0UDON754Y | J0HN__5NOW |
| 190203706 | 6RUMPY_C47 | F1ND_NE55Y |
| 160794683 | 9_FR33D0M8 | 1_H4T3_DF5 |
| 170407562 | K4NG_OUF__ | F1ND_NE55Y |
| 170219839 | _9CE_TW1CK | 1_L1K3TH4T |
| 190759775 | DOGWOOWFUN | M1N I M4X |
| 170488387 | F1ND_NE55Y | B3L1K3B1LL |
| 170954310 | 7ROLL_F4CE | _Y0UM4DBRO |
| 130785761 | _CONGR47S5 | 9_FR33D0M8 |
| 170180425 | _YOUM4DBRO | M1N_I_M4X_ |
| 170506315 | _4_3VER4I_ | B3L1K3B1LL |
| 130052685 | 9_FR33D0M8 | _4_3VER4I_ |
| 170732110 | K4NG_OUF__ | _9CE_TW1CK |
| 170984786 | _4_3VER4I_ | M1N_I_M4X_ |
| 190766164 | _4_3VER4I_ | WE1L_D0N3_ |
| 170496854 | R1DD7E4U___ | B3S7_M0DUL |
| 190766681 | 7ROLL_F4CE | B3L1K3B1LL |
| 170213468 | 1_H4T3_DF5 | _4_3VER4I_ |
| 170434362 | 6RUMPY_C47 | 9_FR33D0M8 |
| 190649694 | DOGWOOWFUN | B3S7_MODUL |
| 160554847 | CHUKN088IS | F1ND_NE55Y |
| 190789473 | DOGWOOWFUN | _9CE_TW1CK |
| 160557413 | R1DD7E4U__ | 1_H4T3_DF5 |
| 160392294 | _CONGR47S5 | _4_3VER4I_ |
| 161107303 | 7ROLL_F4CE | 1_L1K3TH4T |
| 190677099 | CHUKN088IS | Y0UDON754Y |
| 170946670 | 1_L1K3TH4T | D0GW0OWFUN |
| 190172273 | DOGWOOWFUN | K4NG_OUF__ |
```

```
| 161051763 | _4_3VER4I_ | _Y0UM4DBRO |
| 170367093 | 9_FR33D0M8 | J0HN__5NOW |
| 190838391 | WE1L_D0N3_ | 7ROLL_F4CE |
| 190188920 | 9_FR33D0M8 | B3L1K3B1LL |
| 140368985 | CHUKN088IS | M1N_I_M4X_ |
| 190635387 | M1N_I_M4X_ | _Y0UM4DBRO |
| 170346364 | K4NG_OUF__ | 1_H4T3_DF5 |
| 170465045 | CHUKN088IS | R1DD7E4U___ |
| 170788737 | B3S7 MODUL | WE1L D0N3 |
| 170459778 | JOHN__5NOW | YOUDON754Y |
| 190649476 | WE1L_DON3_ | F1ND_NE55Y |
| 180273942 | CHUKN088IS | _9CE_TW1CK |
| 190760070 | _YOUM4DBRO | JOHN__5NOW |
| 170216584 | DOGWOOWFUN | 9_FR33D0M8 |
| 190561673 | F1ND_NE55Y | 1_L1K3TH4T |
| 140869770 | JOHN__5NOW | _9CE_TW1CK |
| 170467212 | 7ROLL_F4CE | B3S7_M0DUL |
| 170276290 | _CONGR47S5 | M1N_I_M4X_ |
| 170111887 | JOHN__5NOW | CHUKN088IS |
| 150442640 | R1DD7E4U___ | F1ND_NE55Y |
| 160497430 | DOGWOOWFUN | 6RUMPY_C47 |
| 190735252 | K4NG_OUF__ | B3S7_MODUL |
| 170146926 | WE1L_D0N3_ | 6RUMPY_C47 |
| 170224279 | K4NG_OUF__ | 1_L1K3TH4T |
| 170404505 | B3S7_MODUL | 9_FR33D0M8 |
| 160446874 | 1_L1K3TH4T | 6RUMPY_C47 |
| 150411934 | 9_FR33D0M8 | B3L1K3B1LL |
| 190583712 | 7ROLL_F4CE | _C0NGR47S5 |
| 160951714 | 6RUMPY_C47 | CHUKN088IS |
| 190758309 | JOHN__5NOW | DOGWOOWFUN |
| 190573735 | CHUKN088IS | 1 L1K3TH4T |
| 160166824 | M1N_I_M4X_ | 7ROLL_F4CE |
| 190189237 | R1DD7E4U__ | 7ROLL_F4CE |
| 170434889 | _CONGR47S5 | _9CE_TW1CK |
| 170285753 | 7ROLL_F4CE | D0GW0OWFUN |
| 170992828 | CHUKN088IS | B3S7_M0DUL |
| 190620349 | 6RUMPY_C47 | 9_FR33D0M8 |
| 170366144 | M1N_I_M4X_ | _Y0UM4DBRO |
| 170182337 | 1_H4T3_DF5 | D0GW0OWFUN |
| 170226527 | 9_FR33D0M8 | J0HN__5NOW |
| 170747723 | JOHN__5NOW | _4_3VER4I_ |
| 170442529 | K4NG_OUF___ | CHUKN088IS |
| 170219976 | 9 FR33D0M8 | M1N I M4X |
| 170253006 | Y0UDON754Y | CHUKN088IS |
| 181004877 | 9_FR33D0M8 | WE1L_D0N3_ |
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190602194
             | R1DD7E4U___ | B3L1K3B1LL |
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            | 6RUMPY_C47 | K4NG_0UF___ |
 170364117
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            | K4NG_OUF__ | _YOUM4DBRO |
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            | _C0NGR47S5 | WE1L_D0N3_ |
 190385633
            | CHUKN088IS | _9CE_TW1CK |
 181044334
            | YOUDON754Y | WE1L DON3 |
 190620132
            | 9_FR33D0M8 | F1ND_NE55Y |
 170252238
            | 6RUMPY_C47
                          | JOHN___5NOW
 151006445
            | 7ROLL F4CE | 1 H4T3 DF5 |
 190875374
            | 1_H4T3_DF5 | 9_FR33D0M8 |
 170793199
            | R1DD7E4U___ | __4_3VER4I__ |
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            | J0HN__5NOW | F1ND_NE55Y |
 170469881
            | _4_3VER4I_ | _Y0UM4DBRO |
            | Y0UDON754Y | R1DD7E4U_
160426234
            | K4NG_OUF__ | F1ND_NE55Y |
 190005502
```

Marking scheme: 5pt if password is correct. 4pt if student only submitted 1 correct password, (5 for correct password 4 for the other one).âĂć2-3pt if password is incorrect (non converging algorithm) depending on closeness of password. âĂć1pt if password does not converge and looks random (less than 0.50 fitness). 0pt if None

(b) (5 points) Explain *briefly* how specifically you chose to do the state encoding, selection, crossover, and mutation.

Solution: You don't have a unique solution since it depends on how you have implemented it. Here is what was expected:

- Initialisation/encoding: Explain clearly how the population was created and which format was used to create population. Some possible answers for this stage are: Representing population as a list of strings, list of tuples (string + fitness), custom class, etc. Strings are created randomly or from a sample of pre-defined passwords.
- Selection: Given an existing population, explain how the blurry_memory function has been run for every single specimen to determine individual fitness. Explain the method to select a number of specimens that will create a new, superior population. Some possible answers for this stage are: Roulette wheel method (probabilistic chance of passing genes over, with more chance assigned to fitter specimens), elitism method (the n-best specimens pass their genes and the other specimens are forgotten), tournament method (create tournaments of n specimens where the fittest wins and passes their genes).
- Crossover: Explain how Crossover has been performed after selecting fittest individuals. Possible answers for crossover methods include: Single-point crossover, n-point crossover, uniform crossover

• Mutation: Explain how the generation have been mutated. Answer for mutation stage: Given a x% mutation chance, each specimen may be mutated. Mutation method may include changing a single character in the specimen, randomising whole specimen, change a few characters, etc.

Novel ideas for any of these stages are allowed and not penalised, as long as you have explained it in the report.

Marking scheme: 5 points: 1pt for correctly explaining representation/encoding, 1pt for correctly explaining selection stage, 1pt for correctly explaining crossover, 1pt for correctly explaining mutation, 1pt for correctly explaining when mutation is triggered (chance of mutation) and/or if you show a good understanding of election/mutation/crossover.

(c) (10 points) Report on the number of reproductions you had to do to converge to the password? Ideally, run your code multiple times and report on the average and variance instead of a single run (since Genetic Algorithm involves randomisation).

Solution: Of course this question can be answered in a lot of different ways. What was expected it for you to explain how many times you ran it, give the average number of iterations to converge to each password. A possible answer for this stage may be:After running the genetic algorithm code N times, I found that the correct password is found after an average of X iterations with a standard deviation of Y. The results can be visualised on graph or table Y

Marking scheme: Result (5pt) - Explanation/times (3pt) - avg/var given (1.5pt) -> avg (1pt) + var (0.5pt) - Code or/and table (0.5pt)

(d) (10 points) Explore the effect of at least one hyper-parameter on the performance of your algorithm. This could for instance be the choice of the mutation rate, or the selection/pairing scheme, the population number, etc. Ideally, provide a table or a figure with a very concise discussion.

Solution:

Once again, it exists a lot of different answers. The point is for you to experiment by changing the value of some hyper parameter and showcase your findings. Findings don't need to be positive, just showcase that the workings/dynamics of the GA change with the hyper-parameters. Possible parameters to change include mutation chance, cross chance, amount of characters mutated, population number, number of iterations, crossover/mutation/selection method, etc. A possible answer for this stage may be: I decided to experiment with varying the mutation chance of my algorithm. I found out that if the mutation chance was too low, password would not always converge to a correct password, but if it was too high, the amount of iterations needed to create valid passwords is higher. My optimal mutation chance is X%. You must show understanding of the importance of hyper-parameter changed and correctly represent results using a graph or table.

Marking scheme: Explanation quality (8pt) - (4pt if password not found) - Table/- Fig (2pt)

Question 3: Let's play a game! (Adversarial search)

You are now engaged in a thrilling two-player adversarial game which is represented by the following tree. Suppose you are the MAX player (who goes first).

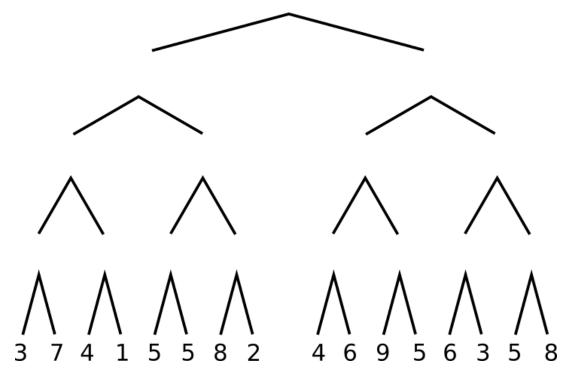


Figure 2: Tree of the game

(a) (10 points) Play optimally (MINIMAX algorithm)

Knowing that you start (top of the tree is you), complete the blank nodes in the provided tree using the MINIMAX approach.

Solution: Using the MINIMAX algorithm, the obtained tree is represented in Figure 3.

Marking scheme: The more correct nodes the more points

(b) (10 points) Play optimally, quicker thinking! (α - β pruning)

Suppose there is a mean arbiter that doesn't like too much hesitation at the beginning. Perform α - β pruning on the tree you filled, searching from left to right. Which branches are pruned and why?

Solution: The red branches of Figure 4 correspond to the pruned one. You also have the values of alpha and beta for each node.

Marking scheme: 7 points on pruning and 3 points on explanation. Each branch wrongly pruned results in 1 point subtracted.

- (c) (10 points) Suppose the entry fee to play the game is x units (for either players).
 - c-(1) What are the ranges of values for x that will be still worth playing the game? (to enter the game at all).
 - c-(2) For a fixed value of x, do you prefer to be the first player (MAXimiser) or the second player (MINimiser)? Very briefly explain your answer.

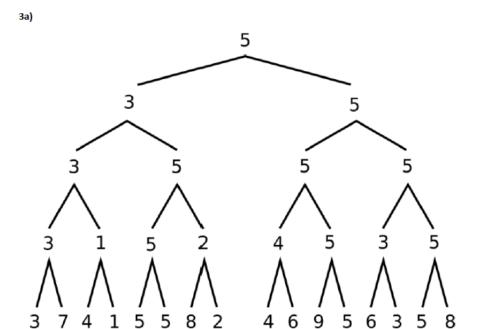


Figure 3: Filled tree of the game

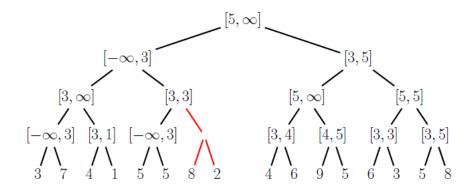


Figure 4: Result of α - β pruning

Solution: 1) The output of MINIMAX allows to know the sequence of actions that are going to be performed if the other player is playing optimally. We also know that this will end up by the MAX player winning 5 units. So we are looking for all values x such as 5 - x >= 0 which means the acceptable range of values is if the entry fee is between 0 and 5.

2) We just showed that if playing optimally it would be better to be MAX player for all entry fee between 0 and 5 since the gain is positive. In a similar way if you play as the MIN player your gain will be x - 5, which means that if the gain is bigger than 5, it's better to play as MIN.

References

[1] Stuart J Russell and Peter Norvig. Artificial Intelligence: A Modern Approach. Pearson, 2009.