# Team DMS

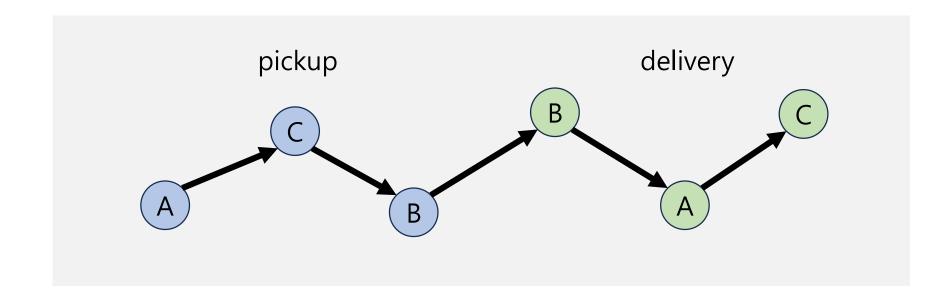
Wonjae Jang, Seoyong Lee

jangwj1026@snu.ac.kr, sylee2685@snu.ac.kr

**Optimization Grand Challenge 2024** 

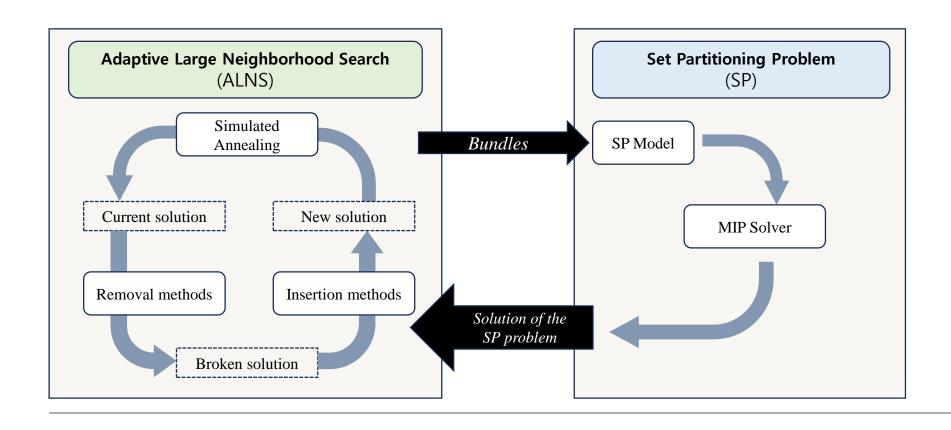
## Problem Summary

A Variant of the <u>Pickup and Delivery Problem</u> with <u>Time Windows</u> for <u>Capacitated</u> and <u>Heterogeneous</u> Vehicle Fleets



## Algorithm: High level View

**Our approach**: Iteratively apply the <u>Adaptive Large Neighborhood Search</u> (ALNS) and <u>Set Partitioning (SP) models</u> in succession.

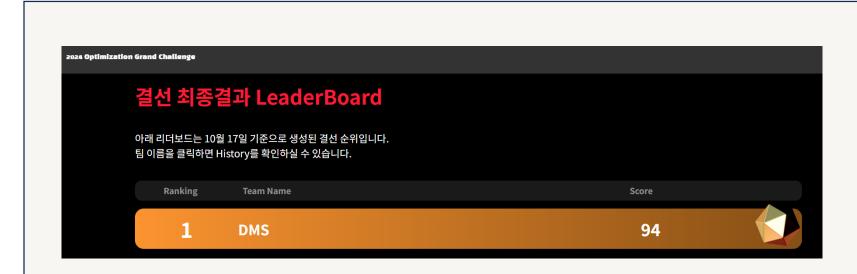


## Algorithm: High level View

#### Algorithm 1 Main Algorithm $\triangleright$ Time list L1: procedure SOLVE(L) $\mathbf{S} \leftarrow \phi$ ▶ Solution S $\mathbf{B} \leftarrow \phi$ ▶ Bundles **B** 3: for $(T_{ALNS}, T_{SP}) \in L$ do 4: $(\mathbf{S}, \mathbf{B}) \leftarrow \text{ALNS}(\mathbf{S}, \mathbf{B}, T_{ALNS})$ 5: $\mathbf{S} \leftarrow \mathrm{SP}(\mathbf{S}, \, \mathbf{B}, \, T_{SP})$ 6: end for 7: return S 9: end procedure

```
10: procedure ALNS(S, B, T)
                                                                                    \triangleright Time limit T
         if S = \phi then
11:
             S \leftarrow generateInitialSolution()
12:
         end if
13:
         \mathbf{p} \leftarrow \text{Uniform}
                                               > prob dist p for choosing removal method
14:
         S_{best} \leftarrow S
15:
         while spent time < T do
16:
             Randomly select op from p and n_{\text{remove}}
                                                                           ⊳ destroy method op
17:
             \hat{\mathbf{S}} \leftarrow \text{APPLYDELETION}(\text{op, } n_{\text{remove}}, \mathbf{S})
18:
             \tilde{\mathbf{S}} \leftarrow \text{APPLYINSERTION}(\tilde{\mathbf{S}})
19:
             Store the routes of \tilde{\mathbf{S}} in \mathbf{B}
20:
             if acceptanceCriteria(S, S) then
                                                                          ▶ Simulated Annealing
21:
                  \mathbf{S} \leftarrow \tilde{\mathbf{S}}
             end if
             if cost(\tilde{S}) < cost(S_{best}) then
24:
                  S_{best} \leftarrow S
             end if
26:
             Update p
                                                        ▶ based on performance : Adaptive!
27:
         end while
28:
         return (S<sub>best</sub>, B)
30: end procedure
31: procedure SP(S, B, T)
         model \leftarrow buildSetCoveringModel(S,B)
         sol_{SC} \leftarrow solveMIP(model, timelimit)
         sol_{SC} \mapsto sol_{SP}
34:
         return sol_{SP}
36: end procedure
```

## "So, How Did We Clinch the Victory?"



	P1	P2	Р3	P4	P5	P6	P7	P8	Р9	P10
Ours	2574.61	2597.28	1826.98	1803.32	2761.76	2124.65	2977.45	2140.61	2176.57	2692.06
Others	2577.69	2608.17	1862.10	1908.47	2754.71	2118.11	2972.85	2136.69	2231.24	2692.34
Diff	-0.12%	-0.42%	-1.89%	-5.51%	0.26%	0.31%	0.15%	0.18%	-2.45%	-0.01%

## "We need to build the best system, not just the best algorithm!"

Python to C++ ~ 10x speedup

Multithreading (via OpenMP) ~ 2x speedup

Novel Insertion Method ~ 260x speedup¹

Performance Engineering ~3x speedup

The common insertion methods widely used in literature are

- k-regret insertion
- Greedy insertion

However, they fail to scale for large instances due to their  $O(K^3)$  complexity.

- Calculating regret/best value of a single order takes O(K).
- To insert a single order, need to calculate for all O(K) orders that is removed.
- We have O(K) orders to insert.

### Thus, their effectiveness is significantly limited

The next method we can think of is...

- Random Order best insertion
  - Randomly choose the order to be inserted
  - Insert the order in the best location

At first glance, the algorithm seems more foolish. However, due to the  $O(K^2)$  nature, it allows more iterations leading to better solutions.

An improvement over before, but still not great.

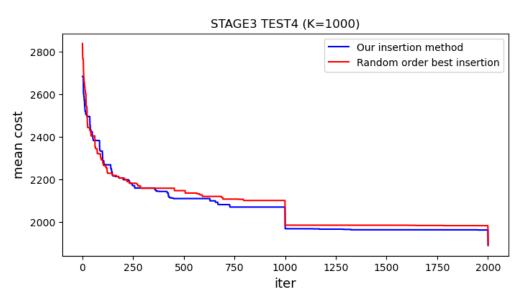
**Intuition:** "Evaluating for all bundles during insertion is unnecessary"

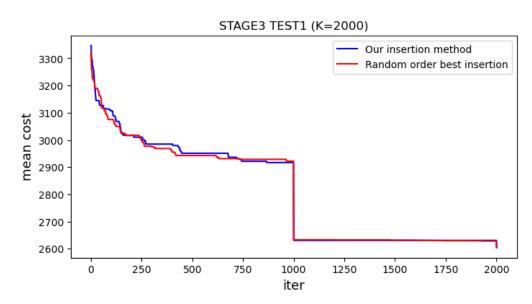
**Idea:** "Evaluate only for the *promising* bundles"

- 0. Preprocess score[i][j] array before ALNS
  - The value increases as the difference in time and distance becomes greater<sup>[1]</sup>

$$c_{ij} + \gamma^{WT} \max\{e_j - \tau_i - \delta_{ij} - l_i, 0\} + \gamma^{TW} \max\{e_i + \tau_i + \delta_{ij} - l_j, 0\}$$
 distance minimum waiting time minimum penalty

- 1. For a given order to insert, calculate how promising the bundle is
  - promising( $order_i$ ,  $bundle_B$ ) :=  $-\frac{\sum_{order_j \in B}(score[i][j])}{|bundle_B|}$
- 2. Select the 10-20 most promising bundles and evaluate only those.





#### No deduction in score for the same number of iterations!

	Regret	Random order best	Ours
Avg time per iteration (K=2000)	7.007s ( <b>260x</b> )	0.1229s ( <b>4.6x</b> )	0.0269s

Speedup gained → More iterations → Better ALNS result + More bundles → Better SP results

### **ALNS**: Removal Methods

```
# orders to remove = (0.05 \sim 0.15) \times K
```

- 1. Random removal
- 2. Distance oriented removal (pickup, delivery & pickup+delivery)
- 3. Route removal
- 4. Shaw removal
- 5. Historical action pair removal
- 6. Worst removal
- 7. <u>Semi-worst removal</u>

"Even a suboptimal removal method improves the solution, and using more diverse methods yields better results."

## ALNS-SP: Additional Insights

- •Time distribution between ALNS and SP plays a critical role.
- •Annealing factor (temperature decay) is a crucial component.
- •Gurobi NoRel heuristics significantly impact performance.
- Acceptance criteria need to be adjusted

	First ALNS & SP	Latter ALNS & SP	
Acceptance Criteria	Simulated Annealing	Hill Climbing Search	
Solver Search Method	Full Exact	Full NoRel Heuristic	

## Proposal for Next Year...

1. Unexplored problems in the literature

2. Commercial solver