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Answers for COMP2012 A1

Question 1 and 2

Assignment 1 Zhu JinShun 22101071D

1) Proof by contradiction

Assume ~~if~~ every number in a_1, a_2, \dots, a_n is greater than m .
The sum of all the numbers will be
 $S = a_1 + a_2 + a_3 + \dots + a_n$
Because every number of a_1, a_2, \dots is greater than m .
 $S > m + m + \dots + m = nxm$
Divide both sides by n , we get
 $S/n > m$
But in question $S/n = (a_1 + a_2 + \dots + a_n)/n$ which leads
to $m = S/n$, so a_1, a_2, \dots, a_n can't be all greater than m .
Which means there exist some number in a_1, a_2, \dots, a_n
is such ~~gt~~ smaller or equal to m to make $m = (a_1 + a_2 + \dots + a_n)/n$

2) a) Truth table b) Truth table

p	q	$p \wedge q$	$\neg p \wedge \neg q$	$(p \wedge q) \vee (\neg p \wedge \neg q)$	$p \leftrightarrow q$
T	T	T	F	T	T
T	F	F	F	F	F
F	T	F	F	F	F
F	F	F	T	T	T

By evaluating truth tables of both sides,
we could see all possible combination of q and p . And $(p \wedge q) \vee (\neg p \wedge \neg q)$
and $p \leftrightarrow q$ are the same. We could conclude $p \leftrightarrow q = (p \wedge q) \vee (\neg p \wedge \neg q)$
is true.

b) Truth table

p	q	r	$p \wedge q \wedge r$	$\neg(p \wedge q \wedge r)$	$\neg p$	$\neg q$	$\neg r$	$\neg p \vee \neg q \vee \neg r$
T	T	T	T	F	F	F	F	F
T	T	F	F	T	F	F	T	T
T	F	T	F	T	F	T	F	T
T	F	F	F	T	F	T	T	T
F	T	T	F	T	T	F	F	T
F	T	F	F	T	T	F	T	T
F	F	T	F	T	T	T	F	T
F	F	F	F	T	T	T	T	T

By evaluating truth tables of both sides, we could see all possible
combination of p, q, r . ~~$(p \wedge q) \wedge r$~~ $\neg(p \wedge q \wedge r)$ has the same cases with
 $\neg p \vee \neg q \vee \neg r$. We could conclude $\neg(p \wedge q \wedge r) = \neg p \vee \neg q \vee \neg r$.

Question 3

(Question 3)

a) We set two proposed function: $P(x)$ and $Q(y)$

For $P(x)$: " x is a prime number that is greater than or equal to 3 and less than or equal to 7"

For $Q(y)$: " y is a prime number greater than or equal to 10 and less than or equal to 25"

Proof: (1) Statement 1: $\exists x \in D(P(x))$

This is true because there are prime numbers between 3 and 7 in the domain, such as, 3, 5 and 7.

Proof: (2) Statement 2: $\exists y \in D(Q(y))$

This is true because there are prime numbers between 10 and 25 in the domain, such as 11, 13, 17, 19, 23

Proof: (3) Statement 3: $\neg \exists z \in D(P(z) \wedge Q(z))$

Because $P(x)$ is in range $3 \leq x \leq 7$, $Q(y)$ is in range $10 \leq y \leq 25$, So it is not true there exist a D so $P(x), Q(y)$ both true. But with the "not" sign, not false is true makes it satisfy the condition

With all three conditions true, the proposed $P(x), Q(y)$ satisfy the above conditions

b) For $L_1: x+y=0$, we can conclude $x=-y$ or $-x=y$

Then we have values from Domain, we get points $(-3,3) (-2,2) (-1,1) (0,0) (1,-1) (2,-2) (3,-3)$ in Set A that contains fulfill requirement

For $L_2: x-y=0$, we can conclude $x=y$

From values from Domain, we get points $(-3,-3) (-2,-2) (-1,-1) (0,0) (1,1) (2,2) (3,3)$ in Set B

Expressing the findings with set builder notation

$$A = \{(x,y) \mid x,y \in D, x+y=0\} \quad B = \{(x,y) \mid x,y \in D, x-y=0\}$$

Final result in set form:

$$A = \{(-3,3), (-2,2), (-1,1), (0,0), (1,-1), (2,-2), (3,-3)\}$$

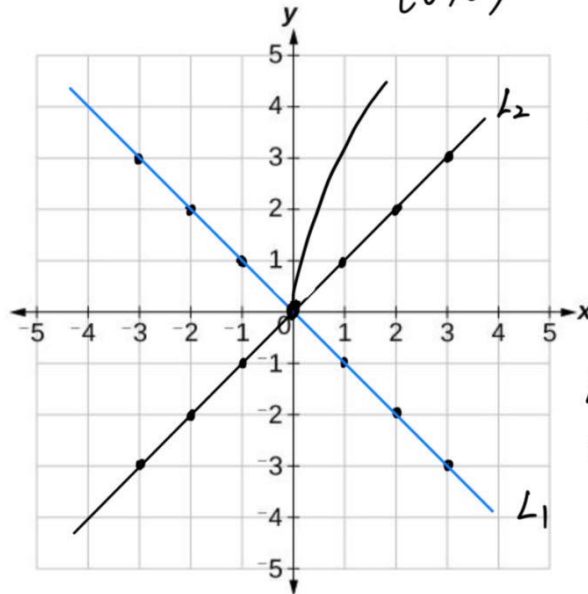
$$B = \{(-3,-3), (-2,-2), (-1,-1), (0,0), (1,1), (2,2), (3,3)\} \quad x=0, y=0$$

The Final answer for L_1, L_2 will be $(0,0)$ $x=y$

Solution guide:

- List out all the points (x, y) of each line in the given domain D ;
- Solve the equations diagrammatically with the aid of the *coordinate grid* given below. Explain with the concept of set.
- Express your findings with the aid of set builder representation, and write down the final answer in set form.

$(0,0)$ - solution $\begin{cases} x=0 \\ y=0 \end{cases}$



L_1 : Each point's
x-coordinate added
it's y-coordinate
equals 0, x y has
opposite value

L_2 : Each point's
x-coordinate minus
it's y-coordinate
equals 0
x and y has same
value

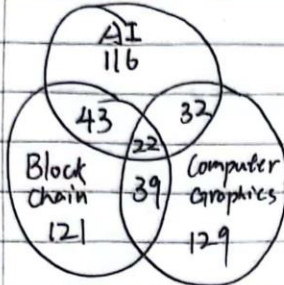
Set $A(L_1)$: $\{(-3,3), (-2,2), (-1,1), (0,0), (1,-1), (2,-2), (3,-3)\}$

Set $B(L_2)$: $\{(-3,-3), (-2,-2), (-1,-1), (0,0), (1,1), (2,2), (3,3)\}$

Question 4

Question 4)

1) Venn Diagram



In the Venn Diagram, We set students taking AI as set A, students taking blockchain as set B, students taking Computer Graphics as set C

Then we get

$$|A| = 116 \quad |B| = 121 \quad |C| = 129$$

$$|A \cap B| = 43 \quad |B \cap C| = 39 \quad |A \cap C| = 32$$

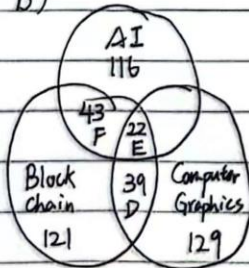
$$|A \cap B \cap C| = 22$$

To calculate the number of students of Comp. according to the Principle of Inclusion-Exclusion

$$\begin{aligned} \text{We have } |A \cup B \cup C| &= |A| + |B| + |C| - |A \cap B| - |B \cap C| - |A \cap C| + |A \cap B \cap C| \\ &= 116 + 121 + 129 - 43 - 39 - 32 + 22 \\ &= 274 \end{aligned}$$

Number of COMP Students is 274

b)



From above Venn Diagram, and question, we know a student has to complete at least 2 courses and one of the course has to be Blockchain to receive FinTech major achievement

The students would be the result of D, E, F

$$D = 39 - 22 = 17$$

$$E = 22$$

$$F = 43 - 22 = 21$$

$$\text{Students that receive FinTech major achievement} = D + E + F = 17 + 22 + 21 = 60$$

60 students will be eligible to receive a degree certificate with a FinTech major

Question 5

(Question 5

a) Proof: Truth Table

P	q	$P \rightarrow q$	$P \wedge (P \rightarrow q)$	$P \wedge (P \rightarrow q) \rightarrow q$
T	T	T	T	T
T	F	F	F	F
F	T	T	F	T
F	F	T	F	T

Therefore, Modus ponens is proved

b) Proof:

Simplification: $P \wedge q = (P \wedge q) \rightarrow P$

Then, addition: $P = P \rightarrow (P \vee q)$ pvs)

Since $(P \vee q) \rightarrow \neg r$ and $\neg r$ can be implied

Also, $r \vee t$

then Disjunctive syllogism $(C(r \vee t) \wedge \neg r) \rightarrow t$

We get t at last

c) Set p as a student is known to be cheating

q as a student is good

r as a student will pass

then $\neg p =$ not cheat $\neg r =$ student fail $\neg q =$ not good

$\neg p \wedge \neg r \wedge \neg q =$ not cheat and student fail and not good

i) This statement is if using our set statements, will be if p then q

but will be invalid since $(p \wedge (p \rightarrow \neg r)) = \neg r$ and $\neg r \wedge q \rightarrow \neg q$

the result will be a student isn't good and will not pass

ii) This statement can be represent as if $\neg p$ then $\neg q$. However, there

are no rules to prove and then it would be an invalid statement

iii) This statement can be represent as if q then $\neg p$.

Using (modus ponens) $q \wedge (q \rightarrow \neg p) = \neg p$. We can conclude this

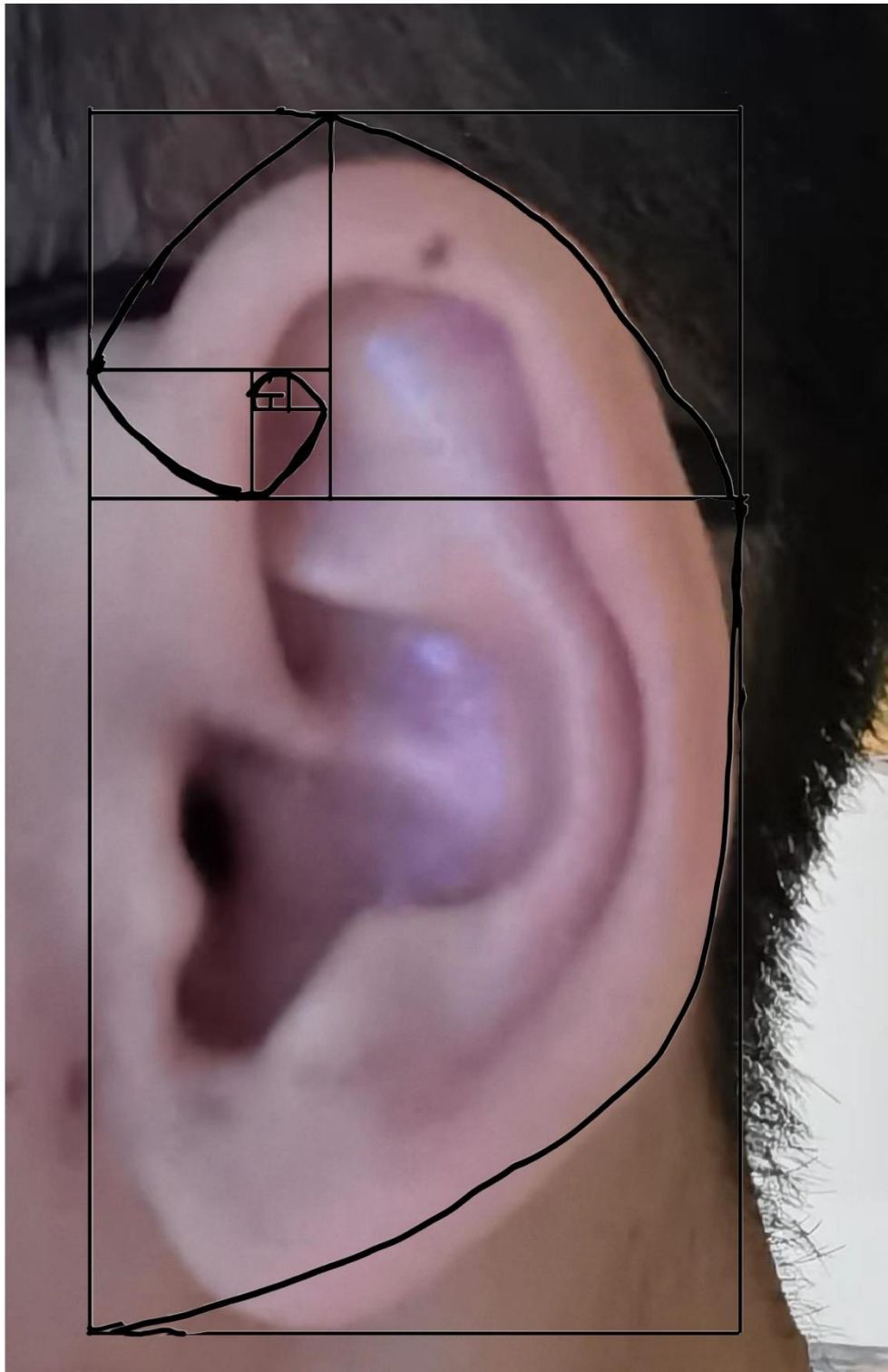
statement is valid

iv) This statement can be represent as if $\neg p$ then $\neg q$. But there

are no rules to prove so we consider it as invalid.

Only iii statement follows logic of statement

Question 6



Picture of my ear with golden spiral(a beautiful gift from my parents)

Question 7

Question 7

a) $D = \begin{bmatrix} 1 & 2 \\ 3 & 6 \end{bmatrix}$ $D^T = \begin{bmatrix} 1 & 3 \\ 2 & 6 \end{bmatrix}$ $E = \begin{bmatrix} 7 & 8 \\ 9 & 10 \end{bmatrix}$

$$D^T \times E = \begin{bmatrix} 1 \times 7 + 3 \times 9 & 1 \times 8 + 3 \times 10 \\ 2 \times 7 + 6 \times 9 & 2 \times 8 + 6 \times 10 \end{bmatrix} = \begin{bmatrix} 7 + 27 & 8 + 30 \\ 14 + 54 & 16 + 60 \end{bmatrix} = \begin{bmatrix} 34 & 38 \\ 68 & 76 \end{bmatrix}$$

$$= \begin{bmatrix} 34 & 38 \\ 68 & 76 \end{bmatrix} = \begin{bmatrix} 34 & 38 \\ 68 & 76 \end{bmatrix}$$

b) Proof: $A^2 = A \cdot A$ $B^2 = B \cdot B$

Left Hand side

$$A^2 = A \cdot A \quad B^2 = B \cdot B$$

$$= \begin{bmatrix} 1 & 0 \\ 2 & -1 \end{bmatrix} \cdot \begin{bmatrix} 1 & 0 \\ 2 & -1 \end{bmatrix} = \begin{bmatrix} 1 & -1 \\ 2 & 3 \end{bmatrix} \cdot \begin{bmatrix} 1 & -1 \\ 2 & 3 \end{bmatrix}$$

$$= \begin{bmatrix} 1 \times 1 + 0 \times 2 & 1 \times 0 + 0 \times 3 \\ 2 \times 1 + (-1) \times 2 & 2 \times 0 + (-1) \times 3 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & -3 \end{bmatrix}$$

$$A^2 - B^2 = \begin{bmatrix} 1 & 0 \\ 0 & -3 \end{bmatrix} - \begin{bmatrix} 3 & -5 \\ 5 & 8 \end{bmatrix} = \begin{bmatrix} -2 & 5 \\ -5 & -7 \end{bmatrix}$$

Right Hand Side

$A - B$

$$= \begin{bmatrix} 1 & 0 \\ 2 & -1 \end{bmatrix} - \begin{bmatrix} 2 & -1 \\ 1 & 3 \end{bmatrix}$$

$$= \begin{bmatrix} 1-2 & 0-(-1) \\ 2-1 & -1-3 \end{bmatrix} = \begin{bmatrix} -1 & 1 \\ 1 & -4 \end{bmatrix}$$

$A + B$

$$= \begin{bmatrix} 1 & 0 \\ 2 & -1 \end{bmatrix} + \begin{bmatrix} 2 & -1 \\ 1 & 3 \end{bmatrix}$$

$$= \begin{bmatrix} 1+2 & 0+(-1) \\ 2+1 & -1+3 \end{bmatrix} = \begin{bmatrix} 3 & -1 \\ 3 & 2 \end{bmatrix}$$

$(A - B) \times (A + B)$

$$= \begin{bmatrix} -1 & 1 \\ 1 & -4 \end{bmatrix} \times \begin{bmatrix} 3 & -1 \\ 3 & 2 \end{bmatrix}$$

$$= \begin{bmatrix} -1 \times 3 + 1 \times 3 & -1 \times (-1) + 1 \times 2 \\ 1 \times 3 + (-4) \times 3 & 1 \times (-1) + (-4) \times 2 \end{bmatrix}$$

$$= \begin{bmatrix} -3+3 & 1+2 \\ 3+(-12) & -1+(-8) \end{bmatrix} = \begin{bmatrix} 0 & 3 \\ -9 & -9 \end{bmatrix}$$

LHS = $\begin{bmatrix} -2 & 5 \\ -5 & -7 \end{bmatrix}$ RHS = $\begin{bmatrix} 0 & 3 \\ -9 & -9 \end{bmatrix}$

We can conclude LHS \neq RHS, which means $A^2 - B^2 = (A - B)(A + B)$ is false. $A^2 - B^2$ doesn't equal $(A - B)(A + B)$

$$C^3 = C \cdot C \cdot C$$

$$= \begin{bmatrix} 1 & 0 & 2 \\ 0 & 2 & 1 \\ 2 & 0 & 3 \end{bmatrix} \cdot \begin{bmatrix} 1 & 0 & 2 \\ 0 & 2 & 1 \\ 2 & 0 & 3 \end{bmatrix} \cdot \begin{bmatrix} 1 & 0 & 2 \\ 0 & 2 & 1 \\ 2 & 0 & 3 \end{bmatrix}$$

$$= \begin{bmatrix} 1+0+4 & 0+0+0 & 2+0+6 \\ 0+0+2 & 0+4+0 & 0+2+3 \\ 2+0+6 & 0+0+0 & 4+0+9 \end{bmatrix} = \begin{bmatrix} 5 & 0 & 8 \\ 2 & 4 & 5 \\ 8 & 0 & 13 \end{bmatrix}$$

$$= \begin{bmatrix} 5 & 0 & 8 \\ 2 & 4 & 5 \\ 8 & 0 & 13 \end{bmatrix} \cdot \begin{bmatrix} 1 & 0 & 2 \\ 0 & 2 & 1 \\ 2 & 0 & 3 \end{bmatrix}$$

$$= \begin{bmatrix} 5+0+16 & 0+0+0 & 10+0+24 \\ 2+0+10 & 0+8+0 & 4+4+15 \\ 8+0+26 & 0+0+0 & 16+0+39 \end{bmatrix} = \begin{bmatrix} 21 & 0 & 34 \\ 12 & 8 & 23 \\ 34 & 0 & 55 \end{bmatrix}$$

$$6C^2 = \text{calculated above } C \cdot C \cdot 6$$

$$= \begin{bmatrix} 5 & 0 & 8 \\ 2 & 4 & 5 \\ 8 & 0 & 13 \end{bmatrix} \times 6$$

$$= \begin{bmatrix} 30 & 0 & 48 \\ 12 & 24 & 30 \\ 48 & 0 & 78 \end{bmatrix}$$

$$7C = C \cdot 7$$

$$= \begin{bmatrix} 1 & 0 & 2 \\ 0 & 2 & 1 \\ 2 & 0 & 3 \end{bmatrix} \times 7 = \begin{bmatrix} 7 & 0 & 14 \\ 0 & 14 & 7 \\ 14 & 0 & 21 \end{bmatrix}$$

$$I_3 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$C^3 - 6C^2 + 7C$$

$$= \begin{bmatrix} 21 & 0 & 34 \\ 12 & 8 & 23 \\ 34 & 0 & 55 \end{bmatrix} - \begin{bmatrix} 30 & 0 & 48 \\ 12 & 24 & 30 \\ 48 & 0 & 78 \end{bmatrix} + \begin{bmatrix} 7 & 0 & 14 \\ 0 & 14 & 7 \\ 14 & 0 & 21 \end{bmatrix}$$

$$= \begin{bmatrix} 21-30 & 0-0 & 34-48 \\ 12-12 & 8-24 & 23-30 \\ 34-48 & 0-0 & 55-78 \end{bmatrix} + \begin{bmatrix} 7 & 0 & 14 \\ 0 & 14 & 7 \\ 14 & 0 & 21 \end{bmatrix}$$

$$= \begin{bmatrix} -9 & 0 & -14 \\ 0 & -16 & -7 \\ -14 & 0 & -23 \end{bmatrix} + \begin{bmatrix} 7 & 0 & 14 \\ 0 & 14 & 7 \\ 14 & 0 & 21 \end{bmatrix} = \begin{bmatrix} -2 & 0 & 0 \\ 0 & -2 & 0 \\ 0 & 0 & -2 \end{bmatrix}$$

$$C^3 - 6C^2 + 7C + kI_3 = 0_3$$

$$= \begin{bmatrix} -2 & 0 & 0 \\ 0 & -2 & 0 \\ 0 & 0 & -2 \end{bmatrix} + \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \times k = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix} \rightarrow \begin{bmatrix} k-2 & 0 & 0 \\ 0 & k-2 & 0 \\ 0 & 0 & k-2 \end{bmatrix} = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

$$1 \times (k) + (-2) = 0$$

$$k = 2$$

The value of k is 2.

Question 8

Question 8		
a) Is-Prime-A(n)	Cost	Frequency
for integer $i \leftarrow 2$ to $n-1$	C_1	At most $n-1$
if $n \bmod i == 0$ then	C_2	At most $n-1$
return False	C_3	At most 1
return True	C_4	At most 1
Is-Prime-B(n)	Cost	Frequency
for integer $i \leftarrow 2$ to $\text{floor}(\sqrt{n})+1$	C_5	At most $\sqrt{n-1} + 1$
if $n \bmod i == 0$ then	C_6	At most $\sqrt{n-1} + 1$
return False	C_7	At most 1
return True	C_8	At most 1
$O(\text{Final complexity for Is-Prime-A}(n)) = C_1 + C_2 + C_3 + C_4 = O(n-1) + O(n-1) + O(1) + O(1) = O(n)$		
$O(\text{Final complexity for Is-Prime-B}(n)) = C_5 + C_6 + C_7 + C_8 = O(\sqrt{n-1} + 1) + O(\sqrt{n-1} + 1) + O(1) + O(1) = O(\sqrt{n})$		
b) Is-Prime-B(n) is more efficient because when n gets larger, $\sqrt{n-1}$ will be much smaller than $(n-1)$. This makes the range of numbers fewer for Is-Prime-B(n) to look for which leads to fewer iterations on finding out the result. So Is-Prime-B(n) is more efficient.		

Question 9

Question 9

Mathematical Induction

$$LHS = 1 - \frac{1}{2} + \frac{1}{3} - \frac{1}{4} + \dots - \frac{1}{2n} \quad RHS = \frac{1}{n+1} + \frac{1}{n+2} + \frac{1}{n+3} + \dots + \frac{1}{2n}$$

Base step: When $n=1$ $LHS = 1 - \frac{1}{2} = \frac{1}{2}$ $RHS = \frac{1}{1+1} = \frac{1}{2}$ $LHS = RHS \forall 1 \in \mathbb{Z}^+$ true

Inductive step: We assume the statement is true for $k=n$

$$\text{then equation } LHS = 1 - \frac{1}{2} + \frac{1}{3} - \frac{1}{4} + \dots + \frac{(-1)^{k+1}}{2k} \quad RHS = \frac{1}{k+1} + \frac{1}{k+2} + \frac{1}{k+3} + \dots + \frac{1}{2k}$$

If we set $n=k+1$ equation will become

$$LHS = 1 - \frac{1}{2} + \frac{1}{3} - \frac{1}{4} + \dots + \frac{(-1)^{k+1}}{2k} + \frac{(-1)^{k+2}}{2k+1} - \frac{1}{2k+2}$$

$$\text{Because } n=k, LHS = RHS = \frac{1}{k+1} + \frac{1}{k+2} + \frac{1}{k+3} + \dots + \frac{1}{2k}$$

$$\text{Thus when } n=k+1 \quad LHS = \frac{1}{k+1} + \frac{1}{k+2} + \frac{1}{k+3} + \dots + \frac{1}{2k} + \frac{1}{2k+1} - \frac{1}{2k+2}$$

$$\begin{aligned} \text{Simplify we get } LHS &= \frac{1}{k+1} + \frac{1}{k+3} + \dots + \frac{1}{2k} + \frac{1}{2k+1} + \left(\frac{1}{2k+1} - \frac{1}{2k+2} \right) \\ &= \frac{1}{k+2} + \frac{1}{k+3} + \dots + \frac{1}{2k} + \frac{1}{2k+1} + \frac{1}{2k+2} \end{aligned}$$

Cause $n=k+1$, then first element for right hand side is $\frac{1}{k+2}$

$$RHS = \frac{1}{k+2} + \frac{1}{k+3} + \dots + \frac{1}{2k} + \frac{1}{2k+1} + \frac{1}{2k+2}$$

$$LHS = RHS$$

Thus, we can prove that $\forall n \in \mathbb{Z}^+ \quad 1 - \frac{1}{2} + \frac{1}{3} - \frac{1}{4} + \dots - \frac{1}{2n} = \frac{1}{n+1} + \frac{1}{n+2} + \frac{1}{n+3} + \dots + \frac{1}{2n}$

Question 10

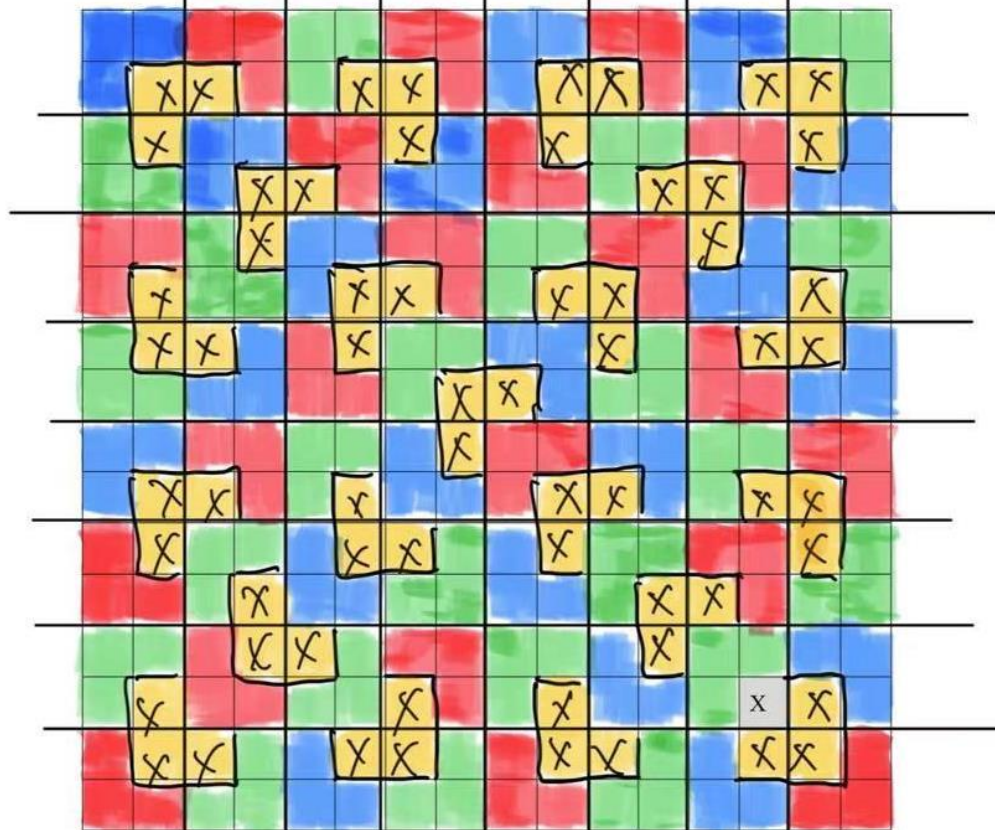
COMP2012 Discrete Mathematics (AY2023-24 Semester 1)

Question 10

[10 marks]

We are referring to the tiling problem in “MI: example 3” in the slides of lecture #5.

Now, we are given the following 16×16 checkboard with a missing square (X).



A triomino looks like one of the followings.

