



**Tecnológico  
de Monterrey**

# Simulation – Basics & Integrals

---

ROBERTO ALEJANDRO CÁRDENAS OVANDO

# Outline

---

- ❖ Taylor Polynomials
- ❖ Approximating integrals with Taylor polynomials
- ❖ Random and Pseudo Random Numbers
- ❖ Riemann Sum

# Taylor Polynomials

---

- ❖ The derivatives are the instantaneous rate of change of a function  $f(x)$  at a given point  $c$
- ❖ Therefore  $f'(x)$  gives us a linear approximation of  $f(x)$  near  $c_i$  for small values of  $\epsilon \in \mathbb{R}$ , we have:

$$f(c + \epsilon) \approx f(c) + \epsilon f'(c)$$

- ❖ If  $f(x)$  has higher derivatives, why stop at the first derivative?



# Taylor Polynomials

---

# Taylor series

---

❖ Let  $f(x)$  be a  $C^n$  polynomial.  $f$  is  $n$ -times continuously differentiable

❖ The  $n$ -th order Taylor polynomial of  $f(x)$  about  $c$  is:

$$T_n(f)(x) = \sum_{k=0}^n \frac{f^{(k)}(c)}{k!} (x - c)^k$$

❖ As Taylor polynomials are approximations of  $f(x)$ , there will be residuals  $R_n$

❖ We want  $R_n(f)(x) \rightarrow 0$

# Taylor series

---

❖ **Theorem:** Suppose  $f(x)$  is  $(n+1)$ -times continuously differentiable. Then,

$$R_n(f)(x) = \int_c^x \frac{f^{(n+1)}(y)}{n!} (x - y)^{n+1} dy$$

❖ This says how much  $Tn(f)(x)$  is off the true value of  $f(x)$ .

# Example

---

❖ Taylor series for:

$$f(x) = e^x \text{ about } 0$$

$$\forall k, f^{(k)}(x) = e^x$$

$$e^x = \sum_{k=0}^{\infty} \frac{f^{(k)}(0)}{k!} (x-0)^k f^{(0)}(0) = e^0 = 1$$

$$f^{(0)}(0) = e^0 = 1 \quad \frac{f^{(2)}(0)}{2!} x^2 = \frac{x^2}{2}$$

$$\frac{f^{(1)}(0)}{1!} x^1 = \frac{x^1}{1} \quad \frac{f^{(3)}(0)}{3!} x^3 = \frac{x^3}{3!}$$

$$\therefore e^x = \sum_{k=0}^{\infty} \frac{x^k}{k!}$$

# Example

---

❖ Function  $f(x) = \cos(x)$  about 0



# Approximating Integrals

---

❖ Approximate the Gaussian curve  $\mu = 0$  and  $\sigma = 1$ :

$$\int_0^{1/3} e^{-x^2} dx$$



# Approximating Integrals

---

# Riemann Sum

---

- ❖ Let a closed interval  $[a,b]$  be partitioned by points:

$$a < x_1 < x_2 < \cdots < x_{n-1} < b$$

- ❖ Where the length of the points are denoted by:

$$\Delta x_1 < \Delta x_2 < \cdots < \Delta x_n$$

- ❖ Let  $X_n^*$  be an arbitrary point in the  $k^{th}$  subinterval. Then:

$$\sum_{k=1}^n f(X_k^*) \Delta x_k$$

- ❖ Is called a Riemann sum for a given function  $f(x)$



# Riemann Sum

---

# Riemann sum

---

❖ The value  $\max(\Delta x_k)$  is called the mesh size



# Riemann sum 2D

---

# Using random numbers

---

❖ Let  $g(x)$  be a function and suppose we wanted  $\theta$  where

$$\theta = \int_0^1 g(x) dx$$

❖ To compute the value of  $\theta$ , note that if  $U$  is uniformly distributed over  $(0,1)$  then we can express  $\theta$  as

$$\theta = E[ g ( U ) ]$$

# Using random numbers

---

- ❖ Independent and identically distributed (iid) random variables have mean  $\theta \rightarrow$  Strong law of large numbers



# Random Numbers

---

- ❖ The building block of a simulation study is the ability to generate random numbers
- ❖ The generated random number will represent an observation from the measured system
- ❖ A random number represents the value of a random variable uniformly distributed an  $(0,1)$

# Pseudo Random Number Generation

---

- ❖ Pseudo Random Numbers (PRN) is a sequence of values
- ❖ They are deterministically generated
- ❖ Have the appearances of being an independent uniform  $(0,1)$  random variables

# Pseudo Random Number Generation

---

- ❖ To generate a PRN starts with an initial value  $X_0$  called the seed
- ❖ The most common approach uses recursion to compute successive values where  $X_n, n \geq 1$ , by letting:

$$X_n = aX_{n-1} \text{ modulo } m$$

Where  $a, m \in \mathbb{N}^+$

- ❖ The quantity  $X_n/m$  is an approximation to the uniform (0,1)
- ❖ The method is called Multiplicative congruential method

# Pseudo Random Number Generation

---

## ❖ Criteria to choose $a$ and $m$ :

1. For any “seed”, the result must “appear” to be a uniform random variable
2. For any seed, the number of values generated before repetition must be large
3. The values can be computed efficiently on a digital computer

# Multiplicative congruential method

---

# Pseudo Random Number Generation

---

## ❖ Another PRN Generator

$$X_n = (aX_{n-1} + C) \text{ modulo } m$$

- ❖ This is called mixed congruential method
- ❖ Mixed = addition + multiplication
- ❖ Quite efficient
  - $m$  bigger than in Multiplicative Congruential method



# Mixed congruential method

---

# Pseudo Random Number Generation

---

❖ Transforming the uniform  $(0,1)$  to  $(a,b)$



# Monte Carlo approach

---

$$\sum_{i=1}^N \frac{g(u_i)}{N} = E[g(u)] = \theta$$

*as  $N \rightarrow \infty$*

❖ This approach is called the Monte Carlo approach

# Monte Carlo approach

---

❖ What happens if the integral goes from  $a$  to  $b$ , instead of  $0,1$



# Monte Carlo approach

---

# Monte Carlo approach

---

❖ Ejemplo

# Monte Carlo approach

---

❖ What happens if we have a multivariate function?

# Monte Carlo approach

---

- ❖ Hence, if we generate  $k$  independent sets, each consisting of  $n$  independent uniform  $(0,1)$  random variables

# Monte Carlo approach

---

✦ Since  $g(U_1^i, U_2^i, \dots, U_P^i)$  for all  $i$  are iid

# HW

---

## ❖ Taylor series: (25 points)

- Handwritten until the 6th-term to get the formula
- Function 1

$$f(x) = \sin(x)$$

- Function 2

If  $i^2 = -1$  compute  $e^{ix}$  about 0



# HW

---

❖ Code: (25 points)

- Riemann sums function

❖ Function (25 points)

$$\int_{-2}^2 e^{x+x^2} dx$$

- Handwritten : Apply Taylor series
- R Code: Riemann sum, and Monte Carlo approach to:

# HW

---

❖ Monte Carlo approach to: (25 points)

$$\int_0^1 \int_0^1 e^{(x+y)^2} dy dx$$