

# **PUBLIC-KEY INFRASTRUCTURE**

UFCFVN-30-M - Computer & Network Security

A Research Coursework

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# TABLE OF CONTENTS

1. LAB TASK	-----	3
1.1 Task 1: RSA Key Cryptanalysis	-----	3
1.2 Task 2: Creating a Certificate Authority (CA)	-----	4
1.3 Task 3: Using your CA to Issue Certificates	-----	6
1.4 Task 4: Deploying Certificates in OpenSSL HTTPS Server	-----	8
1.5 Task 5: Deploying Certificates in Apache HTTPS Server	-----	10
1.6 Task 6: Impersonating a Website (Man-in-the-Middle)	-----	12
1.7 Task 7: Alternative Approaches to (Certificate-Based) PKI (Research Task)	-----	14
1.8 References	-----	16

# 1. LAB TASK

## 1.1 Task 1: RSA Key Cryptanalysis

To find the RSA public key, the following details are provided:

$N=0x8ee32e2bdc9175272b19d9d64667b3869411e62b090bb16c50cabf0b76353ebcda48f9ae02d0e27c75e2d5c0bff5c0bab5519609e96428fd86d968a27c619$

$e=0x11f87$

$p=0x39340499bcb48d0899db332321af58fc9d0c8424d5992fe255da980c4a2e117$


Figure 1 show the procedure for RSA Encryption and how public keys are generated.

### RSA ENCRYPTION (TEXTBOOK VERSION)

- Invented by Rivest, Shamir and Adleman in 1977
- Key Generation:
  - Choose 2 large (e.g. 2048-bit) secret primes  $p$  and  $q$
  - Let  $N = p \cdot q$  we have  $\phi(N) = (p-1) \cdot (q-1)$
  - Choose  $e, d \in \mathbb{Z}_n$  such that  $\text{GCD}(e, \phi(N)) = 1$  and  $e \cdot d \equiv 1 \pmod{\phi(N)}$ 
    - It is common to choose  $e = 3$ ,  $e = 17$  or  $e = 65537$
  - Set  $\text{PK} = (e, N)$  and  $\text{SK} = d$

Figure 1 RSA Encryption

For this task, we use SageMathCell for carrying out the mathematical calculation. Using the above shown formulas, we derive the value of  $d$  (public key) as shown in Figure 2.



Type some Sage code below and press Evaluate.

```
1 p=0x39340499bcb48d0899db332321af58fc9d0c8424d5992fe255da980c4a2e117
2 N=0x8ee32e2bdc9175272b19d9d64667b3869411e62b090bb16c50cabf0b76353ebcda48f9ae02d0e27c75e2d5c0bff5c0bab5519609e96428fd86d968a27c619
3 q=N/p
4 phin=(p-1)*(q-1)
5 e=0x11f87
6 d=(e**1)%phin
7 print(d)
```

Evaluate

212970789819308989496319984031463550864888286220194795883202797339460347843642415393930260585719593497173918429552932948432647895271870832538104832743

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About SageMathCell

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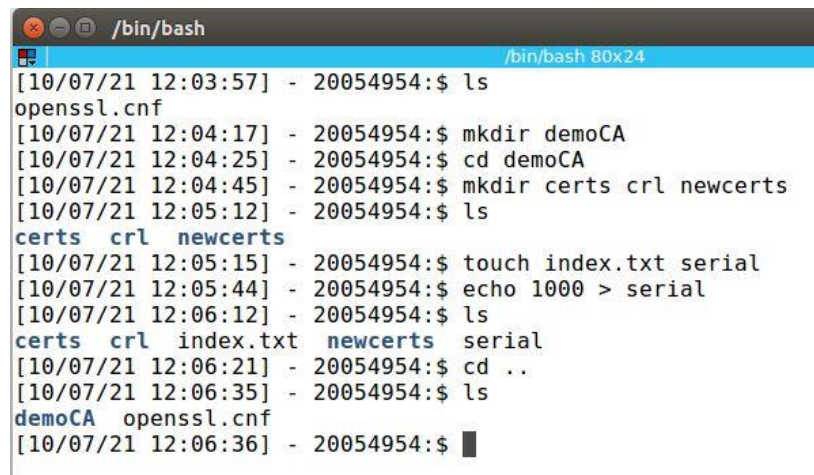
Figure 2 Calculations for finding RSA public key

The RSA public key is:

$D=212970789819308989496319984031463550864888286220194795883202797339460347843642415393930260585719593497173918429552932948432647895271870832538104832743$  (in decimals)

## 1.2 Task 2: Creating a Certificate Authority (CA)

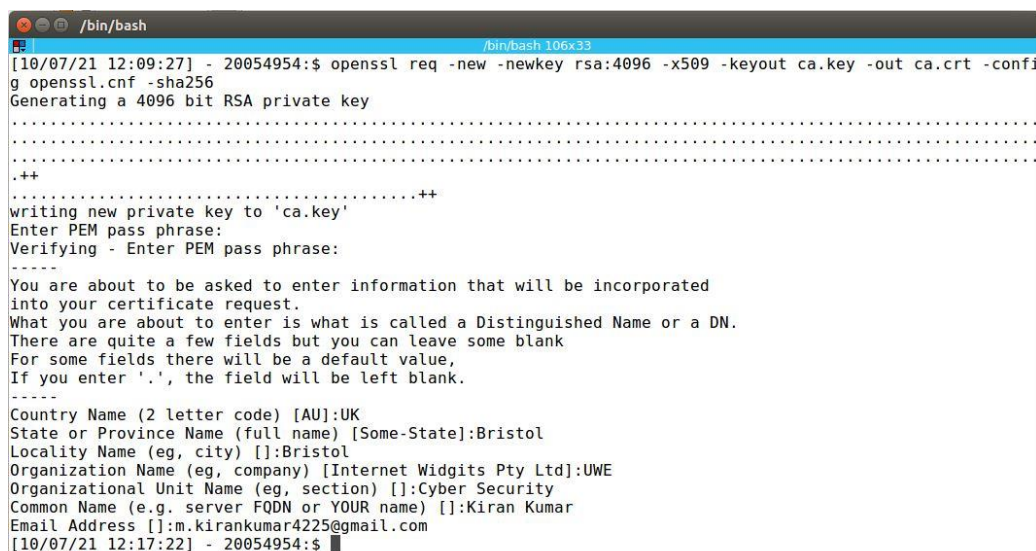
Before creating a Certificate Authority, we need to setup the configuration file (openssl.cnf). To use this file, we need to create some directories and sub-directories. These steps are illustrated in Figure 3.



```
/bin/bash
[10/07/21 12:03:57] - 20054954:$ ls
openssl.cnf
[10/07/21 12:04:17] - 20054954:$ mkdir demoCA
[10/07/21 12:04:25] - 20054954:$ cd demoCA
[10/07/21 12:04:45] - 20054954:$ mkdir certs crt newcerts
[10/07/21 12:05:12] - 20054954:$ ls
certs crt newcerts
[10/07/21 12:05:15] - 20054954:$ touch index.txt serial
[10/07/21 12:05:44] - 20054954:$ echo 1000 > serial
[10/07/21 12:06:12] - 20054954:$ ls
certs crt index.txt newcerts serial
[10/07/21 12:06:21] - 20054954:$ cd ..
[10/07/21 12:06:35] - 20054954:$ ls
demoCA openssl.cnf
[10/07/21 12:06:36] - 20054954:$
```

Figure 3 Configuration file setup

Once the setup is done, we create a self-signed certificate for our CA. Figure 4 shows the command for self signing a certificate. The outputs of this command is stored in ca.key and ca.crt files. The ca.key stores the private key of the CA and ca.crt stores the public-key certificate. The details provided for self-signing a certificate are shown in Figure 4.



```
/bin/bash
[10/07/21 12:09:27] - 20054954:$ openssl req -new -newkey rsa:4096 -x509 -keyout ca.key -out ca.crt -config openssl.cnf -sha256
Generating a 4096 bit RSA private key
.....
..++
writing new private key to 'ca.key'
Enter PEM pass phrase:
Verifying - Enter PEM pass phrase:
-----
You are about to be asked to enter information that will be incorporated
into your certificate request.
What you are about to enter is what is called a Distinguished Name or a DN.
There are quite a few fields but you can leave some blank
For some fields there will be a default value,
If you enter '.', the field will be left blank.
-----
Country Name (2 letter code) [AU]:UK
State or Province Name (full name) [Some-State]:Bristol
Locality Name (eg, city) []:Bristol
Organization Name (eg, company) [Internet Widgits Pty Ltd]:UWE
Organizational Unit Name (eg, section) []:Cyber Security
Common Name (e.g. server FQDN or YOUR name) []:Kiran Kumar
Email Address []:m.kirankumar4225@gmail.com
[10/07/21 12:17:22] - 20054954:$
```

Figure 4 Self-signed Certificate creation for CA

Figure 5 and 6 shows the content of the certificate. The subject and issuer are same which proves that it is a self-signed certificate. In Basic Constraints, the CA is set TRUE, this shows that this certificate can sign another certificate. This proves that Certificate Authority has been created successfully.

```
[11/07/21 05:44:13] - 20054954:$ openssl x509 -in ca.crt -noout -text
Certificate:
  Data:
    Version: 3 (0x2)
    Serial Number:
      c5:be:0b:bd:88:57:67:f6
    Signature Algorithm: sha256WithRSAEncryption
    Issuer: C=UK, ST=Bristol, L=Bristol, O=UWE, OU=Cyber Security, CN=Kiran Kumar/emailAddress=m.kirankumar4225@gmail.com
    Validity
      Not Before: Jul 10 16:17:22 2021 GMT
      Not After : Aug  9 16:17:22 2021 GMT
    Subject: C=UK, ST=Bristol, L=Bristol, O=UWE, OU=Cyber Security, CN=Kiran Kumar/emailAddress=m.kirankumar4225@gmail.com
    Subject Public Key Info:
      Public Key Algorithm: rsaEncryption
      Public-Key: (4096 bit)
      Modulus:
        00:c6:d3:b1:91:7b:8c:80:c5:d3:1c:21:ed:c4:1b:
        e0:bc:ae:a9:a3:bb:2c:d5:e7:5e:bc:f5:4b:f9:19:
        c3:45:43:d6:bd:b5:9f:14:da:a9:39:fc:3b:26:37:
        5b:1d:08:0f:b8:bf:02:79:fb:3e:e5:5f:57:6f:44:
        b5:e9:86:c4:0c:83:d2:07:25:61:5e:f3:b7:b9:03:
        0e:8b:68:31:45:ba:44:c7:f3:a3:9b:66:42:cb:ee:
        82:f4:53:de:e3:33:a1:33:a8:54:52:50:6a:ff:ac:
        76:21:0a:c6:6a:f2:cd:9b:ff:9d:71:26:89:bf:9e:
        fd:fe:06:8b:85:af:1f:af:f8:27:37:2d:74:bf:43:
        06:45:c8:27:62:a2:11:ce:b9:42:ab:82:04:95:b9:
        ad:07:3b:cd:a7:13:0d:28:08:d6:0c:ab:d5:e2:03:
        37:ec:e5:58:e3:fa:ff:3b:5d:30:55:09:22:94:26:
        11:29:43:04:6a:5a:c1:ff:43:15:50:18:d4:d4:74:
        60:85:b4:cb:82:50:32:06:c8:e2:32:f0:00:b8:ce:
        a6:54:10:f9:70:ca:c0:fd:19:77:bf:64:a9:63:40:
        ce:e3:1b:39:4f:36:04:94:3e:e0:80:0f:03:c1:23:
        7b:f8:55:94:24:1e:33:6f:f5:0e:6f:b7:a7:f0:4c:
        38:fa:8c:65:ac:de:22:29:97:53:8c:c8:b7:4e:12:
        9b:b6:0c:9e:f5:ea:45:08:9e:27:ea:77:a5:d5:7e:
        b3:77:e6:24:95:e9:bc:d1:02:ba:6a:01:3e:18:b5:
        3a:71:dc:21:01:55:1f:2e:d3:78:da:41:43:54:fd:
        f7:61:b3:99:c7:00:26:a5:09:79:9d:18:d1:3c:2b:
        dc:57:ae:9c:c6:6b:6f:24:fa:0b:df:c6:66:14:38:
        a2:8d:2f:51:47:62:c9:da:f1:3f:74:20:41:c0:bc:
        a4:41:97:24:b2:10:e0:1d:09:eb:94:c1:fe:a0:b7:
        e1:c4:6a:49:fe:a6:cc:29:2f:a5:cc:fb:10:df:6e:
        ed:05:31:f5:e4:d8:d4:90:d7:88:3b:c0:1e:fb:6d:
        f7:b8:51:d5:3d:3b:ae:ae:d9:5c:9f:f2:0a:b6:3f:
```

Figure 5 Self-signed Certificate

```

55:10:b3:3b:af:a3:98:e1:3a:33:0c:93:c0:9e:a8:
4d:9e:0e:cb:3a:45:0d:a2:b7:c6:f8:89:65:66:4d:
30:d1:81:03:a9:b3:d9:22:8e:b4:31:18:0b:c8:56:
46:ae:09
Exponent: 65537 (0x10001)
X509v3 extensions:
  X509v3 Subject Key Identifier:
    69:67:F0:7D:52:1D:87:7D:6A:C3:E4:4F:EC:93:92:9D:0B:FE:B4:4A
  X509v3 Authority Key Identifier:
    keyid:69:67:F0:7D:52:1D:87:7D:6A:C3:E4:4F:EC:93:92:9D:0B:FE:B4:4A

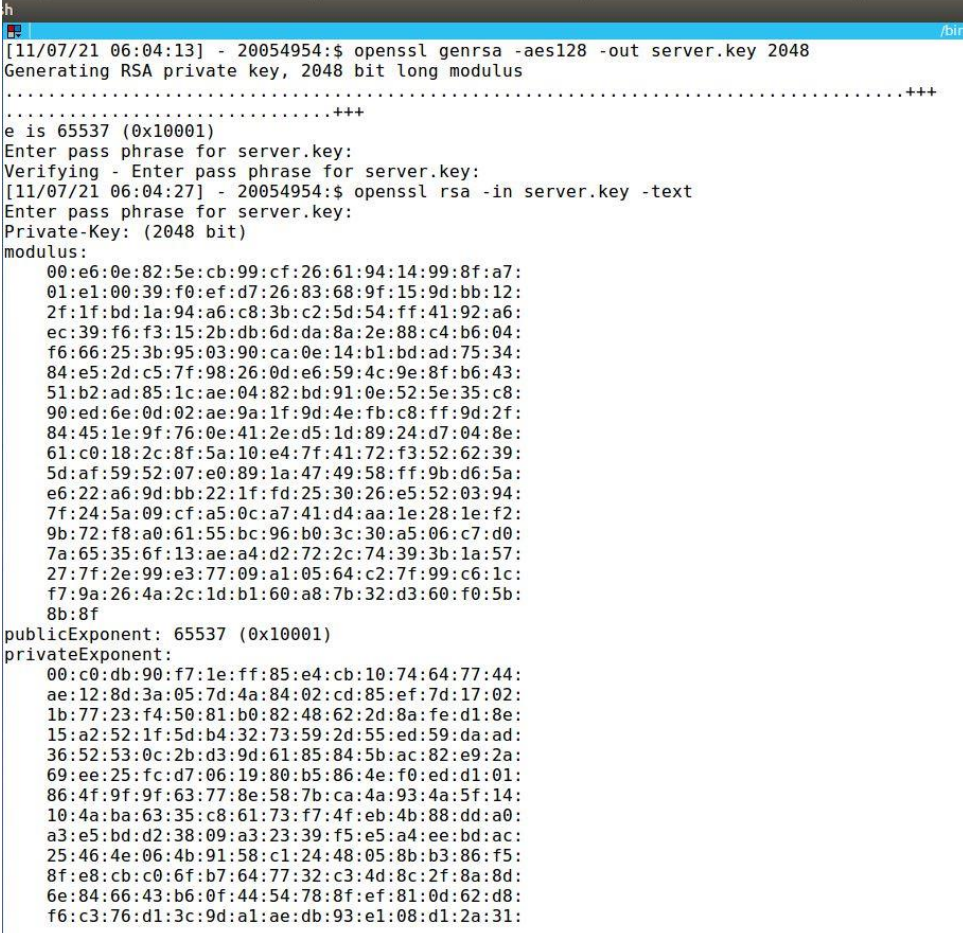
  X509v3 Basic Constraints:
    CA:TRUE
Signature Algorithm: sha256WithRSAEncryption
c6:0b:6a:ad:d0:24:93:33:27:8a:7f:24:02:5d:c2:f5:af:b0:
99:6d:62:71:0a:f4:2c:14:60:5a:f4:96:8d:6b:ec:0f:e7:56:
c4:3f:61:bf:f3:d5:a9:0f:a1:28:31:9e:9a:b2:65:85:8b:8a:
66:e0:7b:fe:99:c9:ab:d7:18:14:86:95:60:0b:a4:65:a6:0c:
93:8c:29:ac:38:7f:4e:a2:39:b4:4e:fe:d5:5e:c6:41:dd:48:
84:5a:71:18:1e:b3:67:e1:23:fc:ba:1a:a5:2b:a1:1f:2f:63:
47:7a:0f:33:e6:a2:6d:a8:21:bd:b8:2a:71:a8:25:58:ac:d3:
d7:5b:91:9b:77:fe:5b:be:ed:9e:bc:d3:06:15:ca:82:98:8a:
fe:e0:7a:2e:ae:c9:0b:99:12:33:f1:2a:b0:62:f6:e7:89:28:
63:97:f8:eb:b1:41:d7:d3:27:e4:57:86:27:d2:0a:ca:6b:cd:
31:f4:b2:4a:cf:08:86:53:62:cc:17:b9:ea:4d:28:70:47:1d:
3e:09:82:79:08:b4:39:75:50:b6:a6:56:4e:a6:c1:b1:dc:5d:
45:0c:ce:a1:9b:f7:08:72:e0:88:ae:04:2d:94:eb:89:fd:83:
da:a4:a4:de:f2:b5:03:63:c8:25:63:a7:48:ae:84:19:90:ca:
d9:b2:f0:48:fe:35:b6:26:30:e4:f8:f3:a7:91:af:ad:5f:9f:
1d:2f:11:6e:d5:3c:c6:0e:ea:d6:f9:92:3e:a2:17:a1:84:99:
8a:ec:d7:b5:79:6b:14:5e:ad:72:41:ea:b1:d6:d3:b0:26:33:
8d:a8:a1:66:fd:03:e5:5f:15:69:1b:fc:1b:e0:92:99:4c:a0:
1c:53:8d:d4:23:04:b0:78:fc:b5:22:14:d6:a7:5e:ca:88:0e:
99:29:5f:06:70:8d:a1:a0:11:9d:33:f8:67:77:08:9b:a9:44:
fc:7f:57:c3:8a:15:64:55:f5:69:7c:1d:78:80:33:b7:27:20:
08:eb:b5:de:25:27:f7:74:db:de:42:dc:12:8b:d7:f5:3c:b2:
ad:47:9a:00:4d:af:e3:0c:9f:0a:d6:b2:31:2c:ba:cf:c2:32:
7c:34:74:06:1b:9f:6f:26:01:70:5a:e1:a8:eb:f6:20:f7:ba:
ab:6c:44:09:51:06:eb:4b:f7:75:f8:6c:34:2c:c9:70:02:a2:
fe:1f:2e:ec:21:b3:e6:4a:3b:43:0d:4f:50:17:ff:a4:a6:75:
6f:f8:09:8f:07:ca:38:33:30:a8:10:00:5a:0a:3d:11:35:f7:
04:df:14:98:92:68:b9:f1:e7:b0:73:b4:92:3a:86:20:4f:7e:
7e:db:a8:86:9f:60:76:65
[11/07/21 05:44:16] - 20054954:$
```

Figure 6 Self-signed Certificate



### 1.3 Task 3: Using your CA to Issue Certificates

Firstly, we need to generate RSA public-private key pairs. Figure 7 shows the commands used for key pair generation. To view the key pair, we use the command “openssl rsa -in server.key -text”.



```
h
[11/07/21 06:04:13] - 20054954:$ openssl genrsa -aes128 -out server.key 2048
Generating RSA private key, 2048 bit long modulus
.....+++
e is 65537 (0x10001)
Enter pass phrase for server.key:
Verifying - Enter pass phrase for server.key:
[11/07/21 06:04:27] - 20054954:$ openssl rsa -in server.key -text
Enter pass phrase for server.key:
Private-Key: (2048 bit)
modulus:
 00:e6:0e:82:5e:cb:99:cf:26:61:94:14:99:8f:a7:
 01:e1:00:39:f0:ef:d7:26:83:68:9f:15:9d:bb:12:
 2f:1f:bd:1a:94:a6:c8:3b:c2:5d:54:ff:41:92:a6:
 ec:39:f6:f3:15:2b:db:6d:da:8a:2e:88:c4:b6:04:
 f6:66:25:3b:95:03:90:ca:0e:14:b1:bd:ad:75:34:
 84:e5:2d:c5:7f:98:26:0d:e6:59:4c:9e:8f:b6:43:
 51:b2:ad:85:1c:ae:04:82:bd:91:0e:52:5e:35:c8:
 90:ed:6e:0d:02:ae:9a:1f:9d:4e:fb:c8:ff:9d:2f:
 84:45:1e:9f:76:0e:41:2e:d5:1d:89:24:d7:04:8e:
 61:c0:18:2c:8f:5a:10:e4:7f:41:72:f3:52:62:39:
 5d:af:59:52:07:e0:89:1a:47:49:58:ff:9b:d6:5a:
 e6:22:a6:9d:bb:22:1f:fd:25:30:26:e5:52:03:94:
 7f:24:5a:09:cf:a5:0c:a7:41:d4:aa:1e:28:1e:f2:
 9b:72:f8:a0:61:55:bc:96:b0:3c:30:a5:06:c7:d0:
 7a:65:35:6f:13:ae:a4:d2:72:2c:74:39:3b:1a:57:
 27:7f:2e:99:e3:77:09:a1:05:64:c2:7f:99:c6:1c:
 f7:9a:26:4a:2c:1d:b1:60:a8:7b:32:d3:60:f0:5b:
 8b:8f
publicExponent: 65537 (0x10001)
privateExponent:
 00:c0:db:90:f7:1e:ff:85:e4:cb:10:74:64:77:44:
 ae:12:8d:3a:05:7d:4a:84:02:cd:85:ef:7d:17:02:
 1b:77:23:f4:50:81:b0:82:48:62:2d:8a:fe:d1:8e:
 15:a2:52:1f:5d:b4:32:73:59:2d:55:ed:59:da:ad:
 36:52:53:0c:2b:d3:9d:61:85:84:5b:ac:82:e9:2a:
 69:ee:25:fc:d7:06:19:80:b5:86:4e:f0:ed:d1:01:
 86:4f:9f:9f:63:77:8e:58:7b:ca:4a:93:4a:5f:14:
 10:4a:ba:63:35:c8:61:73:f7:4f:eb:4b:88:dd:a0:
 a3:e5:bd:d2:38:09:a3:23:39:f5:e5:a4:ee:bd:ac:
 25:46:4e:06:4b:91:58:c1:24:48:05:8b:b3:86:f5:
 8f:e8:cb:c0:6f:b7:64:77:32:c3:4d:8c:2f:8a:8d:
 6e:84:66:43:b6:0f:44:54:78:8f:ef:81:0d:62:d8:
 f6:c3:76:d1:3c:9d:a1:ae:db:93:e1:08:d1:2a:31:
```

Figure 7 RSA key pair generation

Now, we create Certificate Signing Request (CSR) with the previously generated server key pair. This CSR is then sent to the CA to generate a certificate for the key and company name.

```

sh
[11/07/21 06:15:49] - 20054954:~$ openssl req -new -key server.key -out server.csr -config openssl.cnf
Enter pass phrase for server.key:
You are about to be asked to enter information that will be incorporated
into your certificate request.
What you are about to enter is what is called a Distinguished Name or a DN.
There are quite a few fields but you can leave some blank
For some fields there will be a default value,
If you enter '.', the field will be left blank.
-----
Country Name (2 letter code) [AU]:UK
State or Province Name (full name) [Some-State]:Bristol
Locality Name (eg, city) []:Bristol
Organization Name (eg, company) [Internet Widgits Pty Ltd]:UWE
Organizational Unit Name (eg, section) []:CS
Common Name (e.g. server FQDN or YOUR name) []:UWEKiranMohanraj.com
Email Address []:.

Please enter the following 'extra' attributes
to be sent with your certificate request
A challenge password []:kiran
An optional company name []:.
[11/07/21 06:18:06] - 20054954:~$ █

```

Figure 8 Creating CSR

Figure 9 shows the commands for issuing the certificate where the previously generated CSR is converted into X509 certificate using the CA's private key and public-key certificate.

```

[11/07/21 06:22:13] - 20054954:~$ openssl ca -in server.csr -out server.crt -cert ca.crt -keyfile ca.key -config openssl.cnf
Using configuration from openssl.cnf
Enter pass phrase for ca.key:
Check that the request matches the signature
Signature ok
Certificate Details:
  Serial Number: 4096 (0x1000)
  Validity
    Not Before: Jul 11 10:22:23 2021 GMT
    Not After : Jul 11 10:22:23 2022 GMT
  Subject:
    countryName           = UK
    stateOrProvinceName   = Bristol
    organizationName      = UWE
    organizationalUnitName = CS
    commonName            = UWEKiranMohanraj.com
  X509v3 extensions:
    X509v3 Basic Constraints:
      CA:FALSE
    Netscape Comment:
      OpenSSL Generated Certificate
    X509v3 Subject Key Identifier:
      2F:A3:6D:91:6F:1F:86:30:76:81:BC:2C:1D:73:BC:F3:CC:89:CE:CA
    X509v3 Authority Key Identifier:
      keyid:69:67:F0:7D:52:1D:87:7D:6A:C3:E4:4F:EC:93:92:9D:0B:FE:B4:4A

Certificate is to be certified until Jul 11 10:22:23 2022 GMT (365 days)
Sign the certificate? [y/n]:y

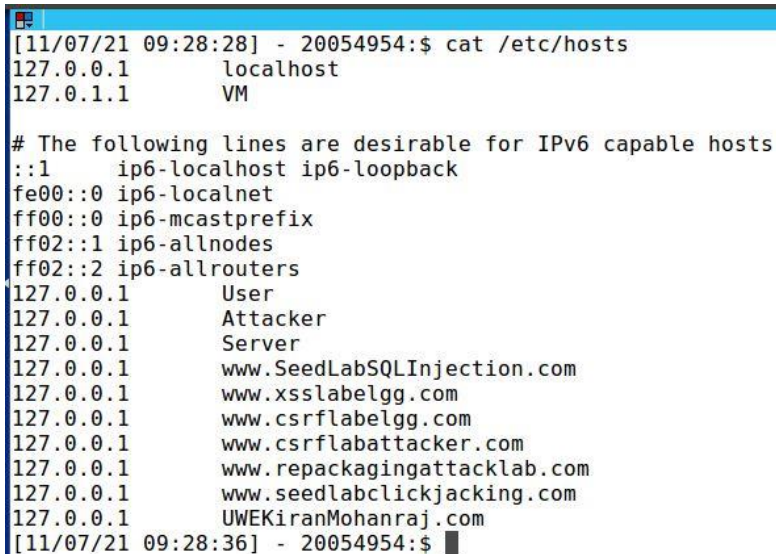
1 out of 1 certificate requests certified, commit? [y/n]y
Write out database with 1 new entries
Data Base Updated
[11/07/21 06:22:42] - 20054954:~$ █

```

Figure 9 Issuing X509 Certificate

## 1.4 Task 4: Deploying Certificates in OpenSSL HTTPS Server

Firstly, the DNS should be configured. We need to add our entry to `/etc/hosts` file so that the machine recognizes the our website without changing the DNS. This entry maps the hostname to the localhost. Figure 10 shows the demonstration of the above steps.

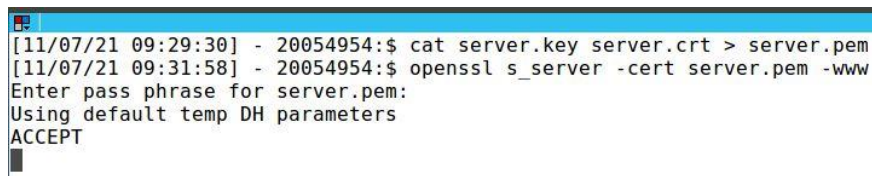


```
[11/07/21 09:28:28] - 20054954:$ cat /etc/hosts
127.0.0.1      localhost
127.0.1.1      VM

# The following lines are desirable for IPv6 capable hosts
::1           ip6-localhost ip6-loopback
fe00::0       ip6-localnet
ff00::0       ip6-mcastprefix
ff02::1       ip6-allnodes
ff02::2       ip6-allrouters
127.0.0.1     User
127.0.0.1     Attacker
127.0.0.1     Server
127.0.0.1     www.SeedLabSQLInjection.com
127.0.0.1     www.xsslabelgg.com
127.0.0.1     www.csrflabelgg.com
127.0.0.1     www.csrfattacklab.com
127.0.0.1     www.repackagingattacklab.com
127.0.0.1     www.seedlabclickjacking.com
127.0.0.1     UWEKiranMohanraj.com
[11/07/21 09:28:36] - 20054954:$
```

Figure 10 Configuring DNS for our hostname

For configuring the web server, we combine the key and certificate into a single file and then launch a web server using `openssl s_server` command. Figure 11 shows that the server is in `ACCEPT` state, indicating server is running.



```
[11/07/21 09:29:30] - 20054954:$ cat server.key server.crt > server.pem
[11/07/21 09:31:58] - 20054954:$ openssl s_server -cert server.pem -www
Enter pass phrase for server.pem:
Using default temp DH parameters
ACCEPT
```

Figure 11 Configuring web server

The port 4433 is the default port which the server can listen. The URL for accessing the server is: [www.UWEKiranMohanraj.com:4433/](http://www.UWEKiranMohanraj.com:4433/)

When we load the URL, we could see an error, which says connection insecure, is displayed in the web browser. In Advanced options, we could see a message that our URL uses invalid security certificates. This is because, the root certificate which we created is not recognised by browser, thus invalidating our certificate.



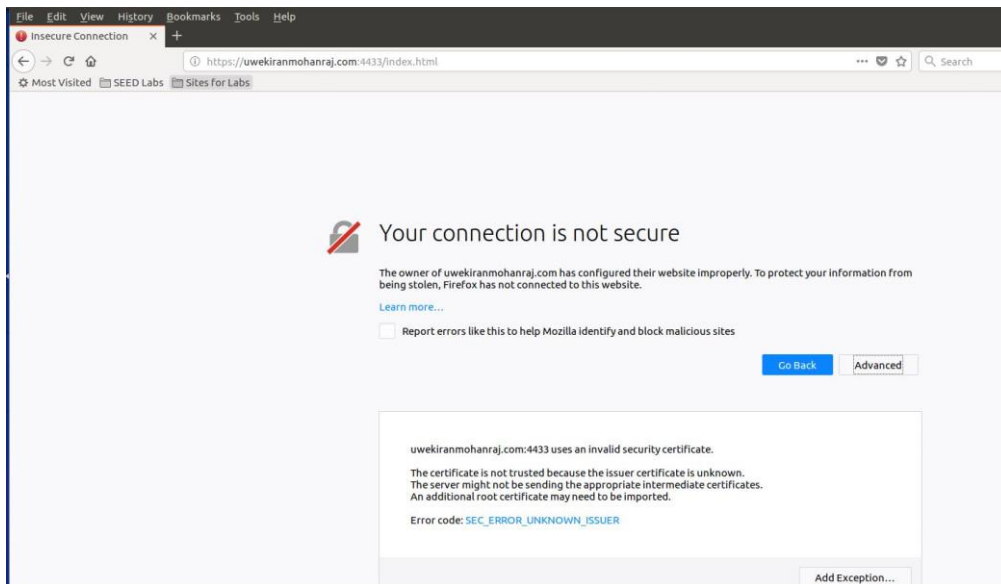


Figure 12 Accessing the web server using our URL

In order to overcome this error, we manually add our certificate to the browser by navigating to certificate manager and importing ca.crt. Figure 13 shows the our root CA has been added as trusted certificates.

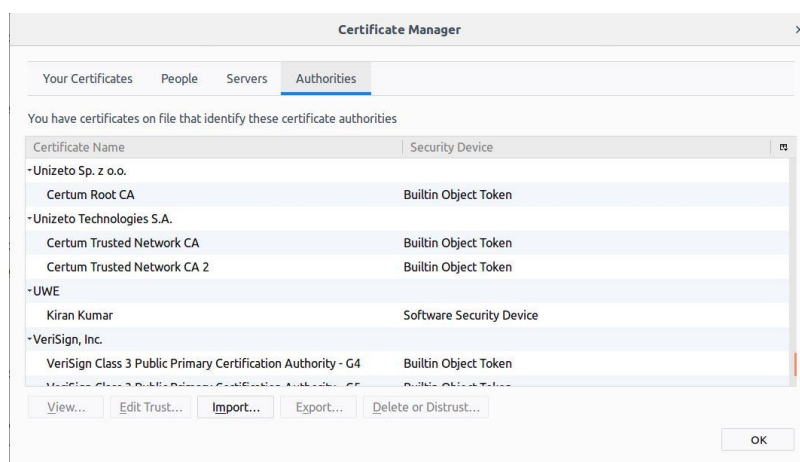


Figure 13 Trusted Certificates of Firefox

Once this setup is done, we try to load our URL in the browser. We could see our index.html being loaded and contents of it has been displayed (as shown in Figure 14).



Figure 14 Successfully accessing our web server

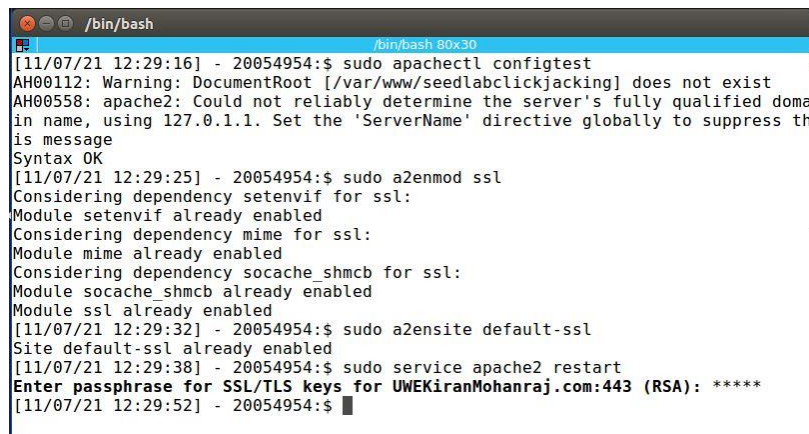
## 1.5 Task 5: Deploying Certificates in Apache HTTPS Server

For this we use Apache web server for deploying the certificates. Firstly, we add a VirtualHost entry to the Apache configuration file (default-ssl.conf). Figure 15 shows our entry added to the configuration. The ServerName entry specifies the name of the website, while the DocumentRoot entry specifies where the files for the website are stored

```
<VirtualHost *:443>
ServerName UWEKiranMohanraj.com
DocumentRoot /var/www/kiran
DirectoryIndex index.html
SSLEngine On
SSLCertificateFile /home/uwe/pki/server.crt
SSLCertificateKeyFile /home/uwe/pki/server.key
</VirtualHost>
<VirtualHost *:80>
ServerName UWEKiranMohanraj.com
DocumentRoot /var/www/kiran
DirectoryIndex index.html
</VirtualHost>
```

Figure 15 VirtualHost entry

Once the setup is done, we run series of commands for starting the apache server. Firstly, the configuration file is tested, then the SSL module is enabled, finally the apache server is configured for HTTPS and restarted. Figure 16 shows the series of commands executed for doing these steps.



```
/bin/bash
[11/07/21 12:29:16] - 20054954:$ sudo apachectl configtest
AH00112: Warning: DocumentRoot [/var/www/seedlabclickjacking] does not exist
AH00558: apache2: Could not reliably determine the server's fully qualified domain name, using 127.0.1.1. Set the 'ServerName' directive globally to suppress this message
Syntax OK
[11/07/21 12:29:25] - 20054954:$ sudo a2enmod ssl
Considering dependency setenvif for ssl:
Module setenvif already enabled
Considering dependency mime for ssl:
Module mime already enabled
Considering dependency socache_shmcb for ssl:
Module socache_shmcb already enabled
Module ssl already enabled
[11/07/21 12:29:32] - 20054954:$ sudo a2ensite default-ssl
Site default-ssl already enabled
[11/07/21 12:29:38] - 20054954:$ sudo service apache2 restart
Enter passphrase for SSL/TLS keys for UWEKiranMohanraj.com:443 (RSA): *****
[11/07/21 12:29:52] - 20054954:$
```

Figure 16 Configuring Apache server

Finally, the web browser is launched for loading the 2 URLs. Figure 17 shows the loading of <https://uwekiranmohanraj.com> and Figure 18 shows the loading of <http://uwekiranmohanraj.com>. The difference that we can notice is that URL with https shows a secure connection whereas the http shows connection insecure. The https URL is secure because the data transferred from the server to the browser is encrypted using SSL, thus creating a secure encrypted connection.

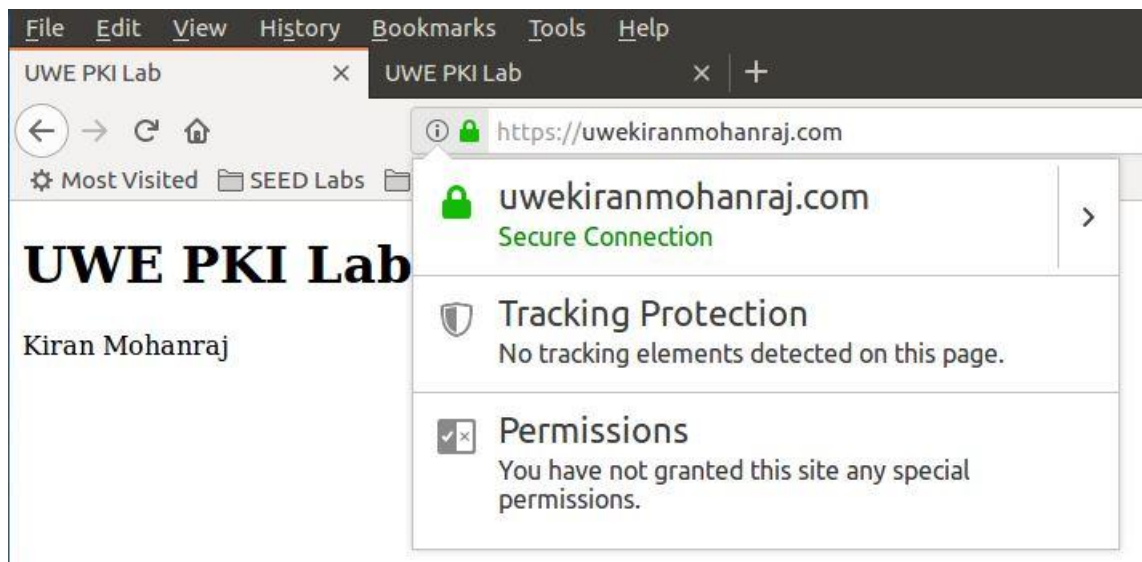


Figure 17 HTTPS secure connection

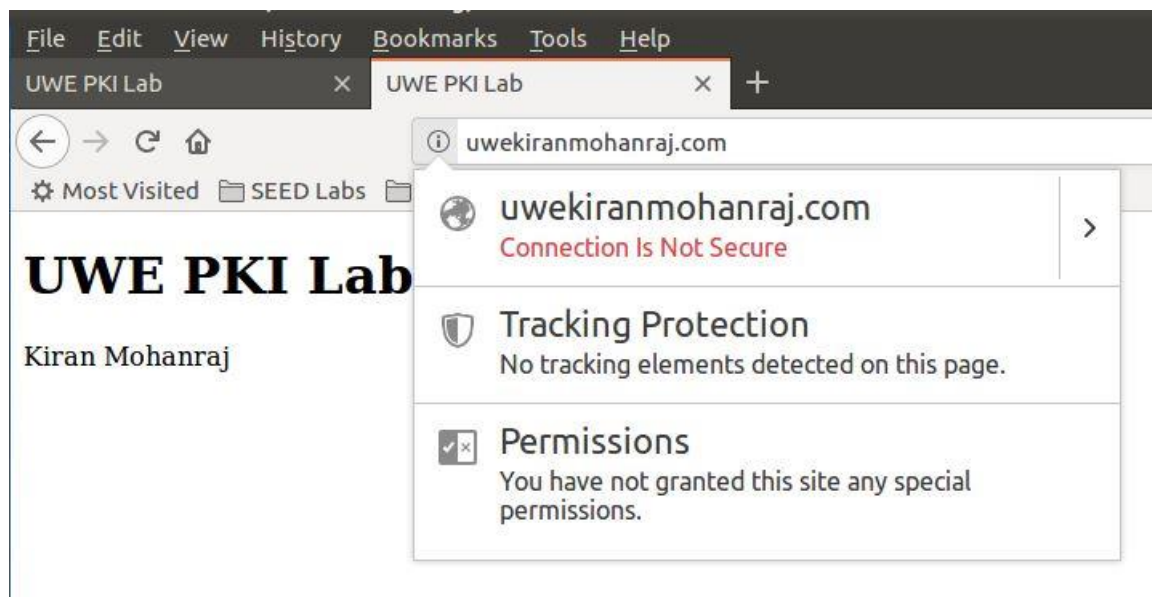


Figure 18 HTTP insecure connection

## 1.6 Task 6: Impersonating a Website (Man-in-the-Middle)

Now, we try to impersonate [www.uwe.ac.uk](http://www.uwe.ac.uk). Firstly, we configure our DNS by adding an entry to the `/etc/hosts` file using the UWE url (Figure 19). Similar to previous task, we UWE's VirtualHost entry to the apache configuration file. Figure 20 shows this the added entry.

```
10 127.0.0.1      User
11 127.0.0.1      Attacker
12 127.0.0.1      Server
13 127.0.0.1      www.SeedLabSQLInjection.com
14 127.0.0.1      www.xsslabelgg.com
15 127.0.0.1      www.csrflabelgg.com
16 127.0.0.1      www.csrfattacklab.com
17 127.0.0.1      www.repackagingattacklab.com
18 127.0.0.1      www.seedlabclickjacking.com
19 127.0.0.1      uwekiranmohanraj.com
20 127.0.0.1      www.uwe.ac.uk
```

Figure 19 Configuring DNS for our hostname

```
<VirtualHost *:443>
  ServerName www.uwe.ac.uk
  DocumentRoot /var/www/kiran
  DirectoryIndex index.html

  SSLEngine On
  SSLCertificateFile /home/uwe/pki/server.crt
  SSLCertificateKeyFile /home/uwe/pki/server.key
</VirtualHost>

<VirtualHost *:80>
  ServerName www.uwe.ac.uk
  DocumentRoot /var/www/kiran
  DirectoryIndex index.html
</VirtualHost>
```

Figure 20 VirtualHost entry

Now we browse our website <https://www.uwe.ac.uk> and see that we get an error saying that the connection is not secure. In more details option, we could see the common name of our UWE url and the certificate does not match, indicating the certificate is not meant for this domain.

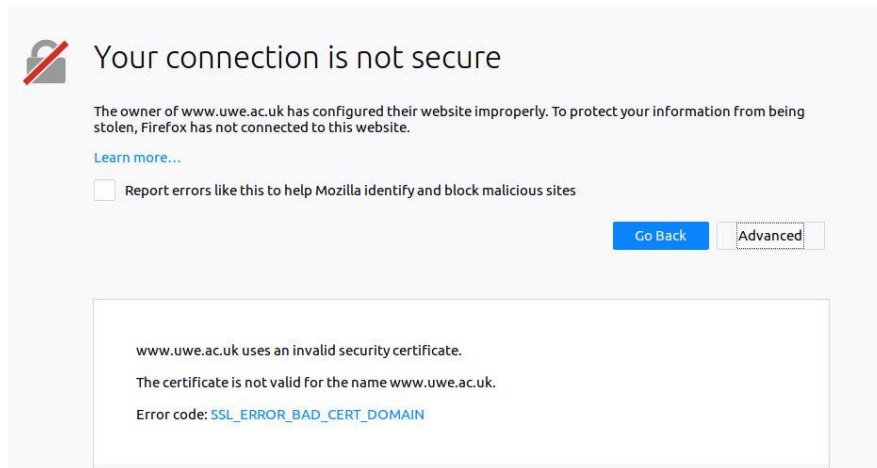


Figure 21 Browsing <https://www.uwe.ac.uk>

This proves that the MITM attack is defeated in the use of public key infrastructure. If the common name matched, then the website would have loaded. However, the common name will match only if the server had a valid certificate from the CA for its domain, which it didn't, in this case.

When we try to browse <http://www.uwe.ac.uk>, the we could see that the page has been redirected to Apache2 Default page (as shown in Figure 22). Thus, proving MITM attack can be implemented.



Figure 22 Browsing <http://www.uwe.ac.uk>



## 1.7 Task 7: Alternative Approaches to (Certificate-Based) PKI (Research Task)

### INTRODUCTION

Public-key infrastructure (PKI) is the basis for using public-key cryptography through key management. PKI helps in authenticating the public key of the user. If the authenticity is compromised, any user can use it to decrypt the messages, leading to problem of repudiation. In certificate-based PKI, a Certification Authority (CA) issues a digital certificate which authenticates the public key of the user. This digital certificate contains the public key and identity of the user. If the user wants to use other entities public key, they must verify its certificate for validating the authenticity of the public key. The major drawback of certificate-based PKI is scalability and certificate management. (Singh *et al.*, 2019).

### HYBRID PKI SCHEMES

To overcome the drawbacks of certificate-based PKI, alternate approaches were proposed. Identity-based PKI (ID-PKI) was proposed by (Shamir, 1985), which used publicly known identity (eg. E-mail address) as public-key and the private key is sent to the owner of the identity (through Email). Although ID-PKI overcome the problem of distributing public key through certificates, it lacks scalability and support for identity management overhead (Singh *et al.*, 2019). To overcome the challenges of the certificate-based PKI and IDPKC, the hybrid approach was introduced which uses the positive features of both these models. In this report, we will be looking three hybrid PKI approaches:

#### ***A) Self- Certificate Scheme:***

Lee (Lee and Kim, 2000) proposed this self-certificate scheme which address the issues of explicit authentication of self-certified key. This self-certificate is user generated for the public key by signing it with private key to the public key. The key pairs of the users are renewed by themselves. The authenticity of the Certificate Authority's certification is maintained without the user interacting with the CA for renewing the key pairs. The certificate-based scheme's revocation method is used by the Certificate Authority to revoke an issued key. The major advantage of this self-certificate scheme is that the user can renew their key pairs without interacting with the Certificate authority, meanwhile having the CA's authentication. The drawback of this scheme is that it has self-certificate overhead similar to the certificate-based PKI overhead (Lee and Kim, 2000).

#### ***B) Certificate based encryption (CBE) Scheme***

Gentry (Gentry, 2003) introduced CBE scheme to provide simple revocation in certificate-based PKI. By exploiting the pairing, this scheme eases out the certificate management of the PKI. This scheme implemented the implicit certification from the IBE (identity-based

encryption) no escrow from public-key certification, thus producing a hybrid PKI scheme. In this model, the user generates their own private key or public key pair and request a certificate from the Certificate authority. For decryption key and proof of certification, the Certificate authority generates implicit certificate by using IBE scheme. This helps in eliminating the third-party queries on the certificate status (Gentry, 2003).

Shao (Shao, 2011) propose a new Certificate-Based Encryption scheme from pairings provably secure in the chosen-key model. This scheme is derived from the CBE scheme of Gentry (Gentry, 2003), after making some improvements. Each client is allowed to choose its partial private key and partial public key independently. To deal with the certificate revocation problem, this paper uses the concept of the time period proposed by Gentry, which is added into certificate information to eliminate the need of the current certification status. The full public key of the client is just its certificate information with a time period. As results, anyone can encrypt messages under the certificate information and the time period so that the receiver can decrypt only if the receiver's partial public key is currently certified. However, if a client entirely trusts the authenticity of a receiver's partial public key, he could use the partial public key of the receiver directly to encrypt messages. Furthermore, this high security level is just at the cost of one more hashing and one more point scalar for encryption, which can be precomputed once for all. This scheme not only overcomes the weakness in the certified-key model of the Gentry scheme, but also raises efficiency.

### ***C) Unified Public Key Infrastructure (UPKI) Scheme***

Lee (Lee, 2010) introduces a new concept called unified public key infrastructure (UPKI) in which both certificate-based and ID-based cryptography are provided to users in a highly combined manner. This scheme assumes the existence of a trusted authority called key generation and certification authority (KGCA) who has the role of both CA and KGC. It checks identification information of user and issues a certificate for a user-chosen public key X. It also issues ID-based partial private key to the user. This paper assumes the existence of multiple KPAs (Key Privacy Agents) who provide key privacy service. In the proposed private key issuing protocol user is authenticated with certificate and user's certified public key X is used to blind the protocol messages such that only the legitimate user can retrieve the ID-based private key.

## **CONCLUSION**

The demand for successfully implemented PKI systems is increasing as many organisations are looking to secure their communications. This report provides various hybrid PKI schemes involving certificates. There is need for some kind of certification is required in order to solve the problem of explicit authentication and scalability related issues. UPKI achieves all these problems except with some communication and complexity overhead.

## 1.8 References

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