Sep22 : Day 3

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```
class Array
{
    private int[]a1;
    private int n;

    public Array(int max)
    {
        a1 = new int[max];
        n=0;
    }
    public void display()
    {
        for(int j=0;j<n;j++)
            System.out.print(a1[j]+" ");
        System.out.println();
    }
    public void insert(int value)
    {
        a1[n] =  value;
        n++;
    }
    public boolean search(int key)
```

Problem statement: Find duplicates in an array

- Given an array a1[] of size N which contains elements from 0 to N-1, you need to find all the elements occurring more than once in the given array.

- **Example 1:**
  - Input:
    - N = 4
    - a[] = {0,3,1,2}
  - Output: -1
  - Explanation: N=4 and all elements from 0 to (N-1 = 3) are present in the given array. Therefore output is -1.

- **Example 2:**
  - Input:
    - N = 5
    - a[] = {2,3,1,2,3}
  - Output: 2 3
  - Explanation: 2 and 3 occur more than once in the given array.

Problem statement: Removing punctuations from a given string

- **Given a string, remove the punctuation from the string if the given character is a punctuation character, as classified by the current C locale. The default C locale classifies these characters as punctuation:**
  - ! " # $ % & ' ( ) * + , - . / : ; ? @ [\ ] ^ _ ` { | } ~

- **Example 1:**
  - Input : %welcome' to @cdacmumbai?<s
  - Output : welcome to cdacmumbai

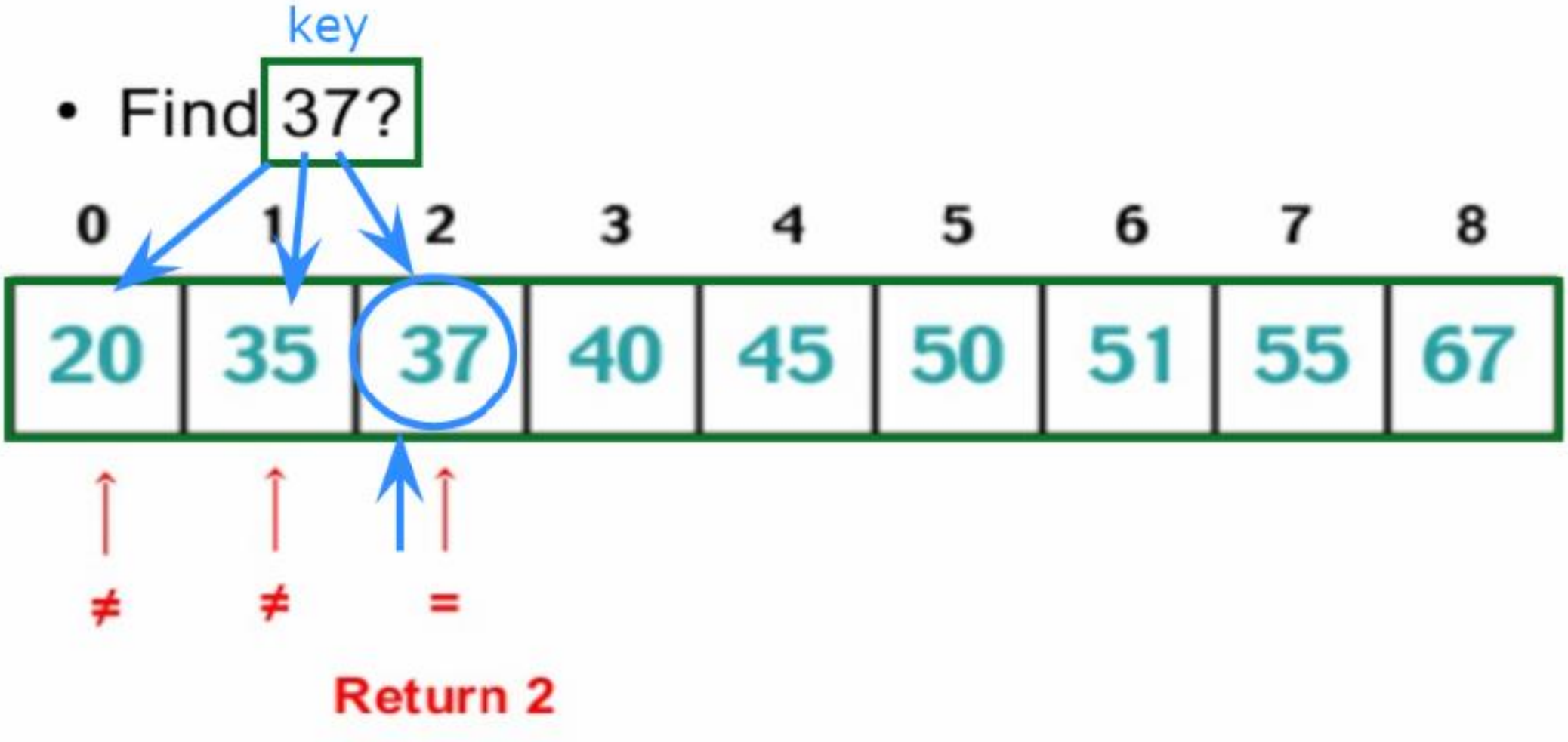
- **Example 2:**
  - Input : Hello!!!, he said ---and went**.
  - Output : Hello he said and went

Problem statement: Program to find the initials of a name.

- Given a string name, we have to find the initials of the name
- Examples 1:
  - Input : Kabhi Haa Kabhi Naa
  - Output : K H K N
    - We take the first letter of all
    - words and print in capital letter.
- Example 2:
  - Input : Mahatma Gandhi
  - Output : M G
- Example 3:
  - Input : Shah Rukh Khan
  - Output : S R K

  - Example 4: your own name

Linear Search

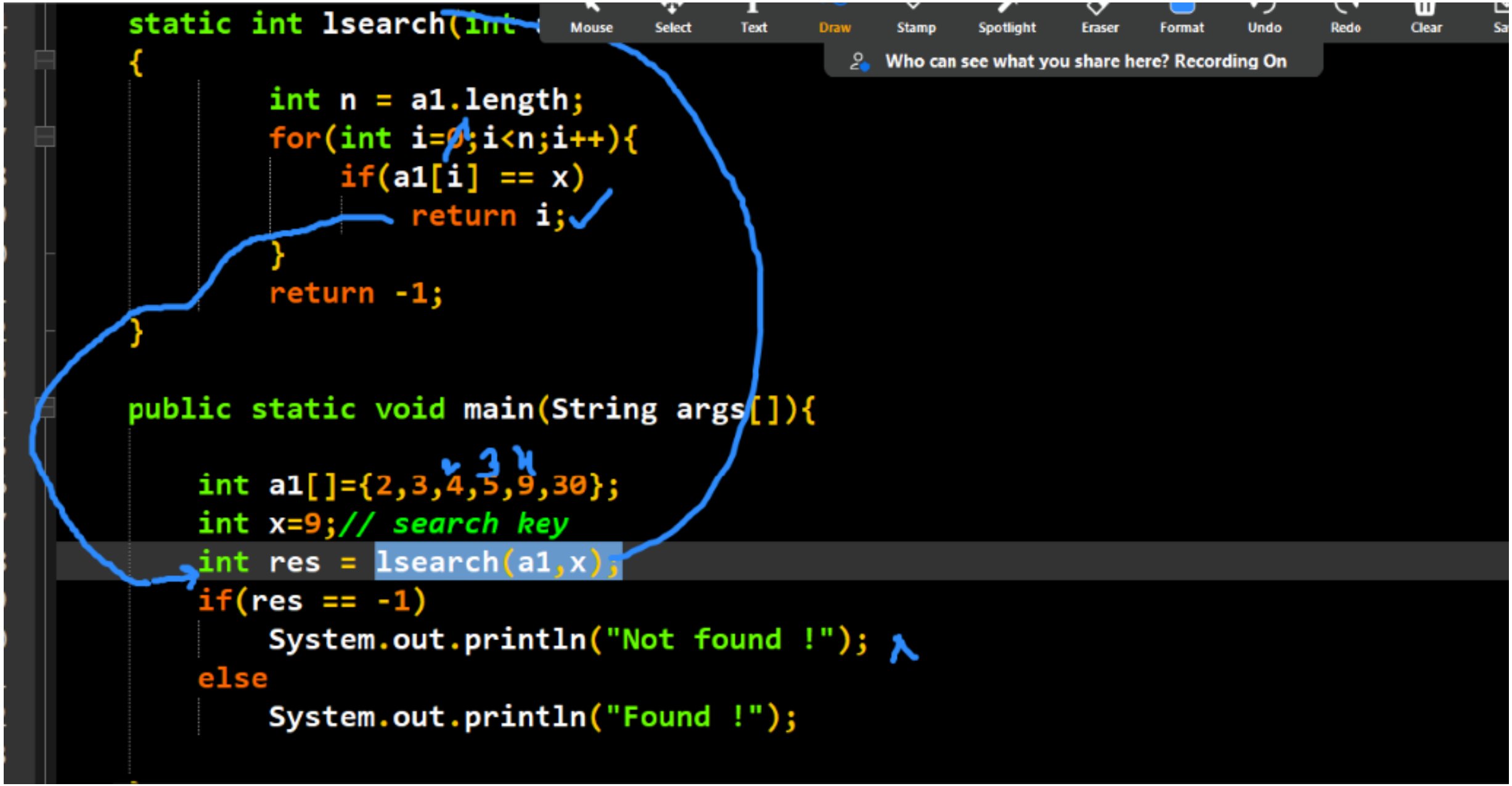


Linear Search

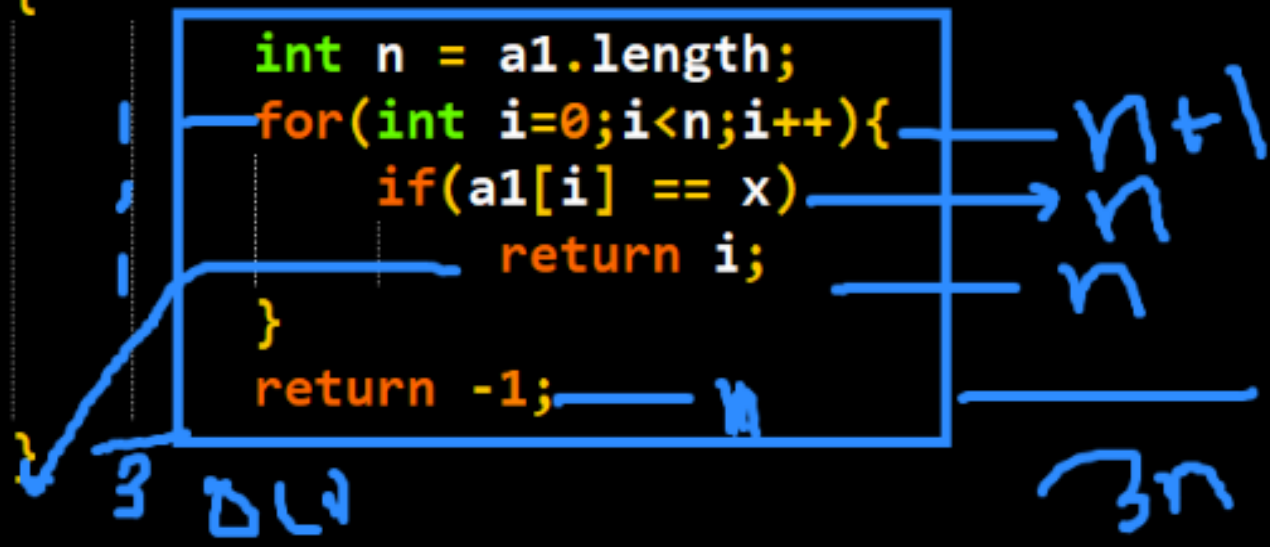
Algorithm

Consider **LA** is a linear array with **N** elements and **K** is a positive integer such that **K<=N**. Following is the algorithm to find an element with a value of ITEM using sequential search.

```
1. Start
2. Set J = 0
3. Repeat steps 4 and 5 while J < N
4. IF LA[J] is equal ITEM THEN GOTO STEP 6
5. Set J = J +1
6. PRINT J, ITEM
7. Stop
```


Mouse
Select
Text
Draw
Stamp
Spotlight
Eraser
Format
Undo
Redo
Clear
Who can see what you share here? Recording On
static int lsearch(int
{
int n = a1.length;
for(int i=0;i<n;i++){
if(a1[i] == x)
return i;
}
return -1;
}
public static void main(String args[]){
int a1[]={2,3,4,5,9,30};
int x=9;// search key
int res = lsearch(a1,x);
if(res == -1)
System.out.println("Not found !");
else
System.out.println("Found !");

```java
static int lsearch(int a1[],int x)
{
    int n = a1.length;
    for(int i=0;i<n;i++){
        if(a1[i] == x)
            return i;
    }
    return -1;
}
```



Element Found in
Worst case is O(n)
Best case is O(1)

O (n)

```java
public static void main(String args[]){

    int a1[]={2,3,4,5,9,30};
    int x=2;// search key
    int res = lsearch(a1,x);
    if(res == -1)
        System.out.println("Not found !");
    else
        System.out.println("Found !"+res);
```

Element not found

O(n)

Program 3

Problem: Given an array arr[] of n elements, write a function to search a given element x in arr[].

Examples :

Input : arr[] = {10, 20, 80, 30, 60, 50, 110, 100, 130, 170}
x = 110;
Output : 6
Element x is present at index 6

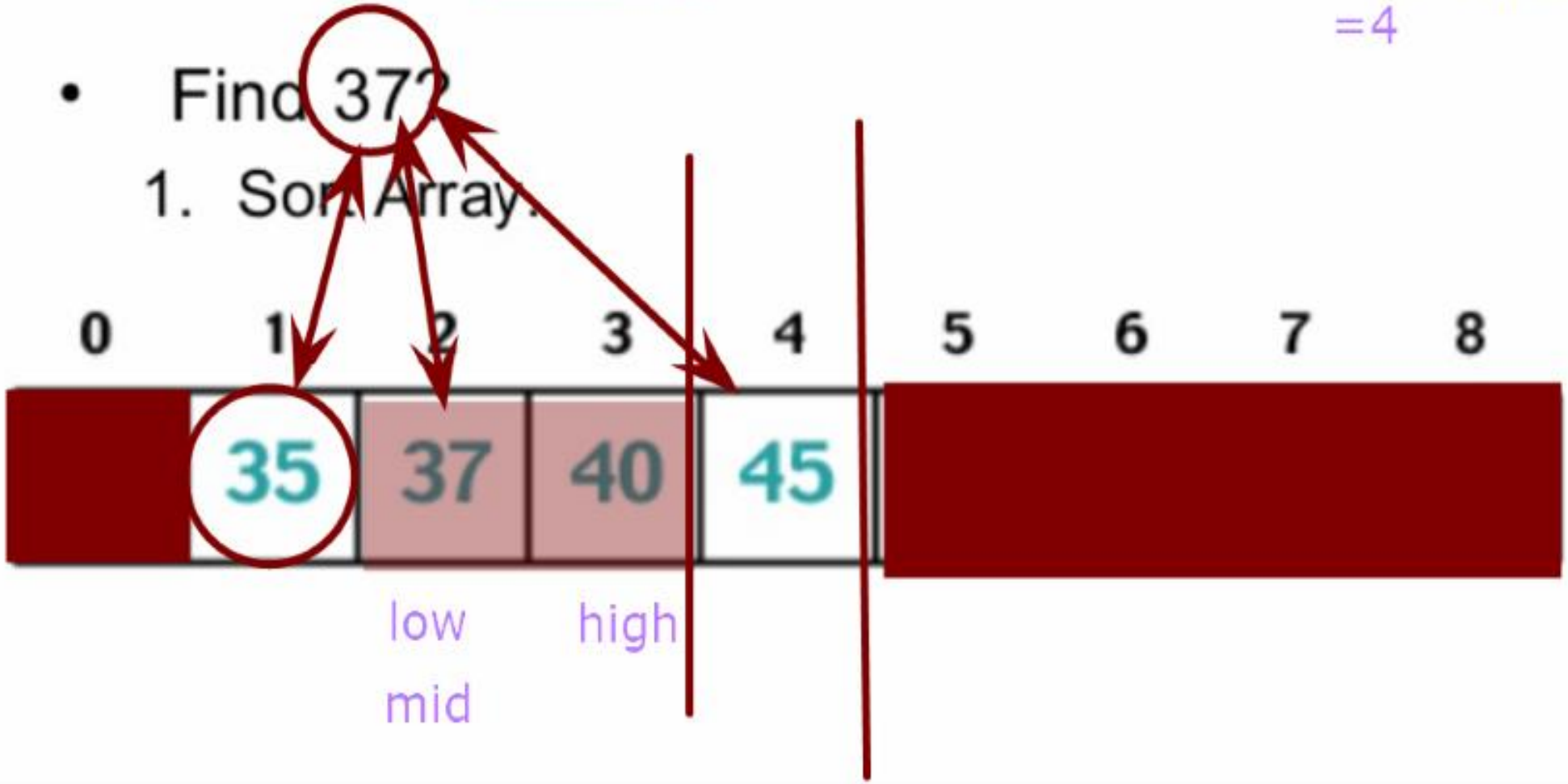
Input : arr[] = {10, 20, 80, 30, 60, 50,
110, 100, 130, 170}
x = 175;
Output : -1
Element x is not present in arr[].

-Sorted order

2 Binary Search

mid=(low+high)/2
=(0+8)/2
=4

- Find 37?
  1. Sort Array



Binary Search

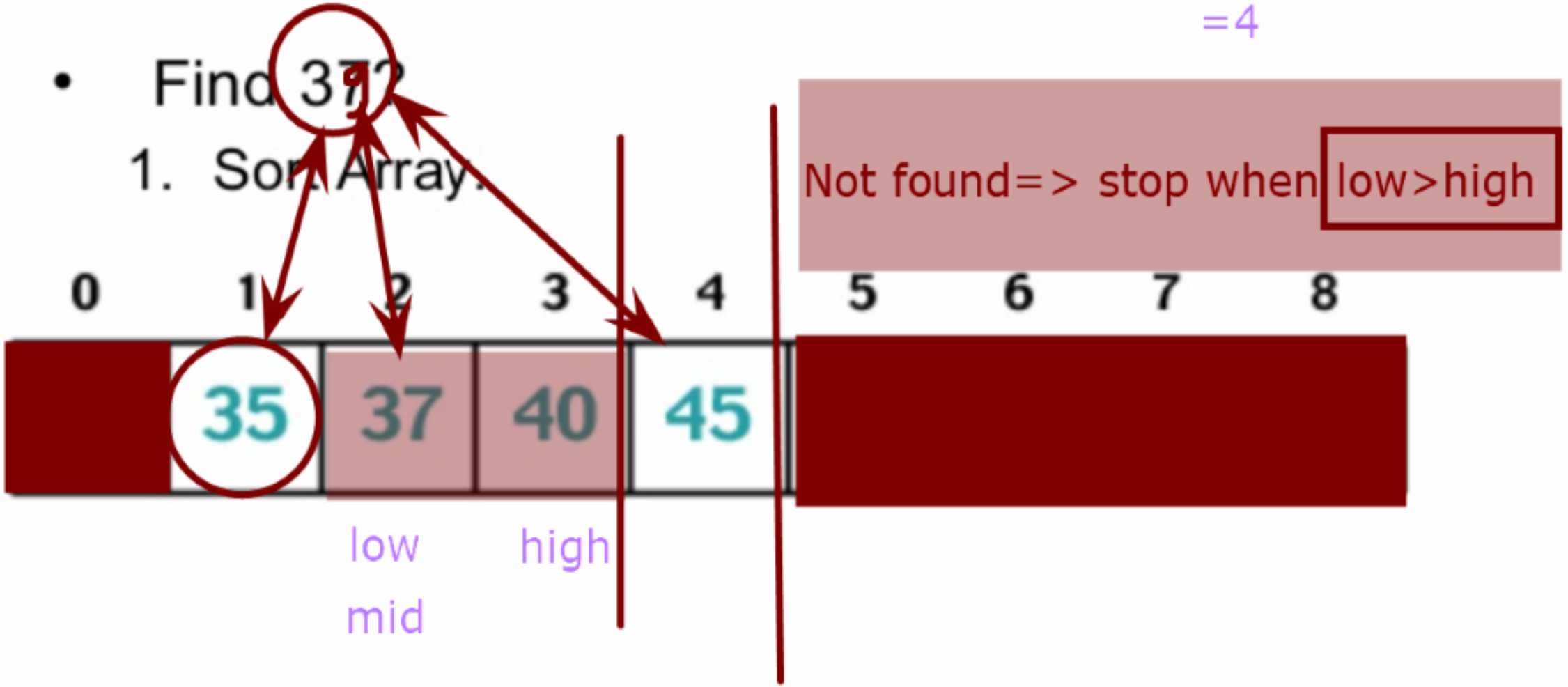
```
Procedure binary_search
    A ← sorted array
    n ← size of array
    x ← value to be searched

    Set lowerBound = 1
    Set upperBound = n

    while x not found
        if upperBound < lowerBound
            EXIT: x does not exists.

        set midPoint = lowerBound + ( upperBound - lowerBound ) / 2

        if A[midPoint] < x
            set lowerBound = midPoint + 1

        if A[midPoint] > x
            set upperBound = midPoint - 1

        if A[midPoint] = x
            EXIT: x found at location midPoint
    end while

end procedure
```

-Sorted order

Binary Search

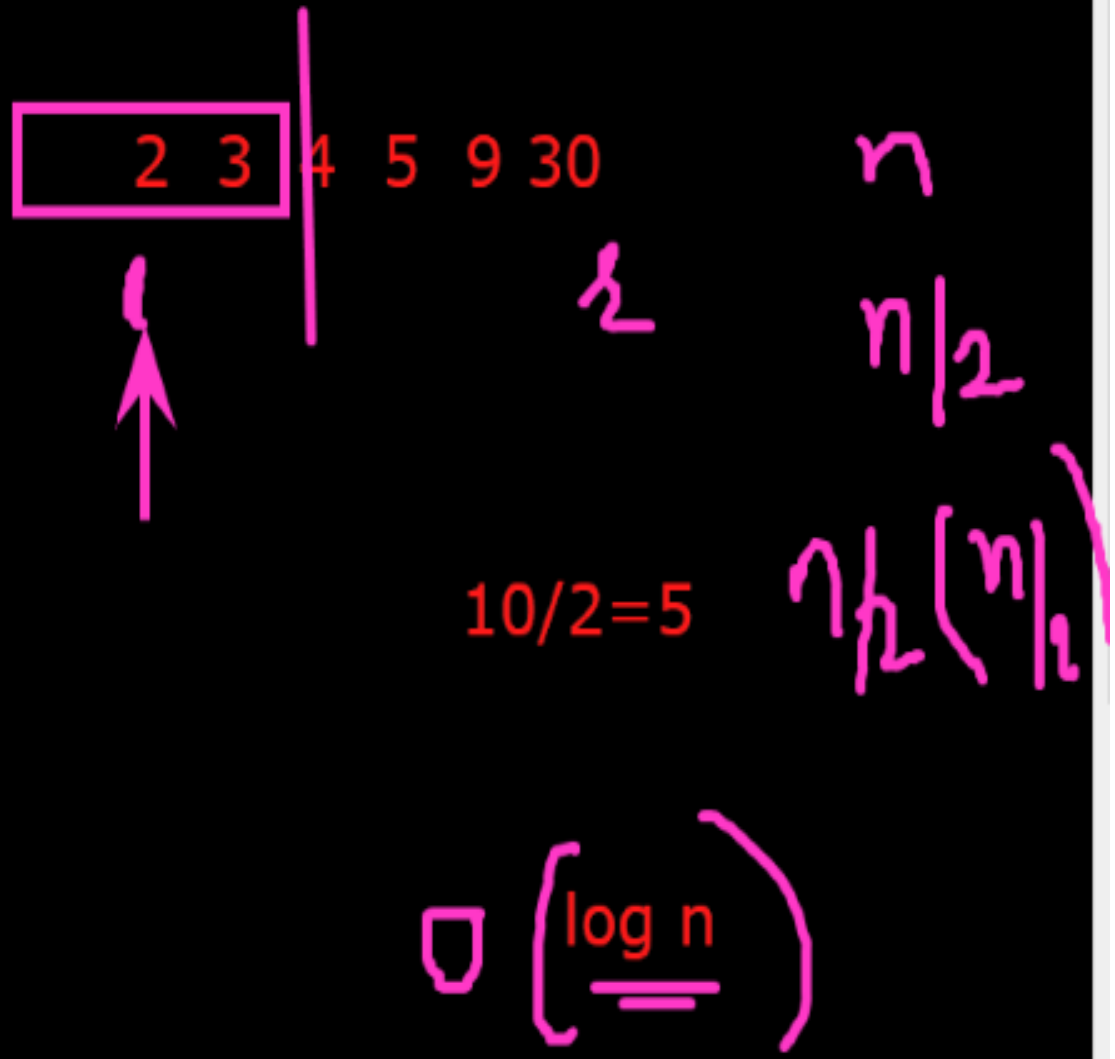
mid=(low+high)/2
=(0+8)/2
=4

- Find 37?
  1. Sort Array



```
static int bsearch(int a1[], int x, int l, int r)
{
    if(r>=l)
    {
        int mid = l+(r-l)/2;

        if(a1[mid] == x)
            return mid;
        if(a1[mid] > x)
            return bsearch(int a1[], int x, int l, int mid-1);

        return bsearch(int a1[], int x, int mid+1, int r);

    }
    return -1;
}
```

```java
static int bsearch(int a1[], int x, int l, int r)
{
    if(r>=l)
    {
        int mid = l+(r-l)/2;

        if(a1[mid] == x)
            return mid;
        if(a1[mid] > x)
            return bsearch(a1,x,l,mid-1);

        return bsearch(a1, x, mid+1, r);

    }
    return -1;
}
```



Algorithms & Data Structure Complexity

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Program:
- Implemented
- Programmer
- Programming language
- H/w or OS
- Testing

Algorithm:
- Design
- Domain Knowledge
- Language
- H/w, OS
- Analysis

Posterior Analysis
- Program
- Dependent PL
- Dependent on H/w
- Time

Priori Analysis
- Algorithms
- Independent of PL
- Independent of H/w
- Time & space

Characteristics of Algorithms:
Input
Output
Finite
unambiguity
Effective
Language independent

Algorithm Complexity
- Time Factor(Time complexity)
- Space Factor(Space complexity)

Asymptotic Notation:

-Best case: minimum time required for execution.

-Average case: average time required for the execution.

-Worst case: maximum time required for the execution.

```
if(x==5)
{
    stmt;
}
```

1

1

f(n) = 2 sec

O(2) → constant ← O(1)

```
if(x==5)
{
    stmt;
}
```

```
for(int i=0;i<5;)
{
    SOP("done");
    i++;
}
```

6

5

5

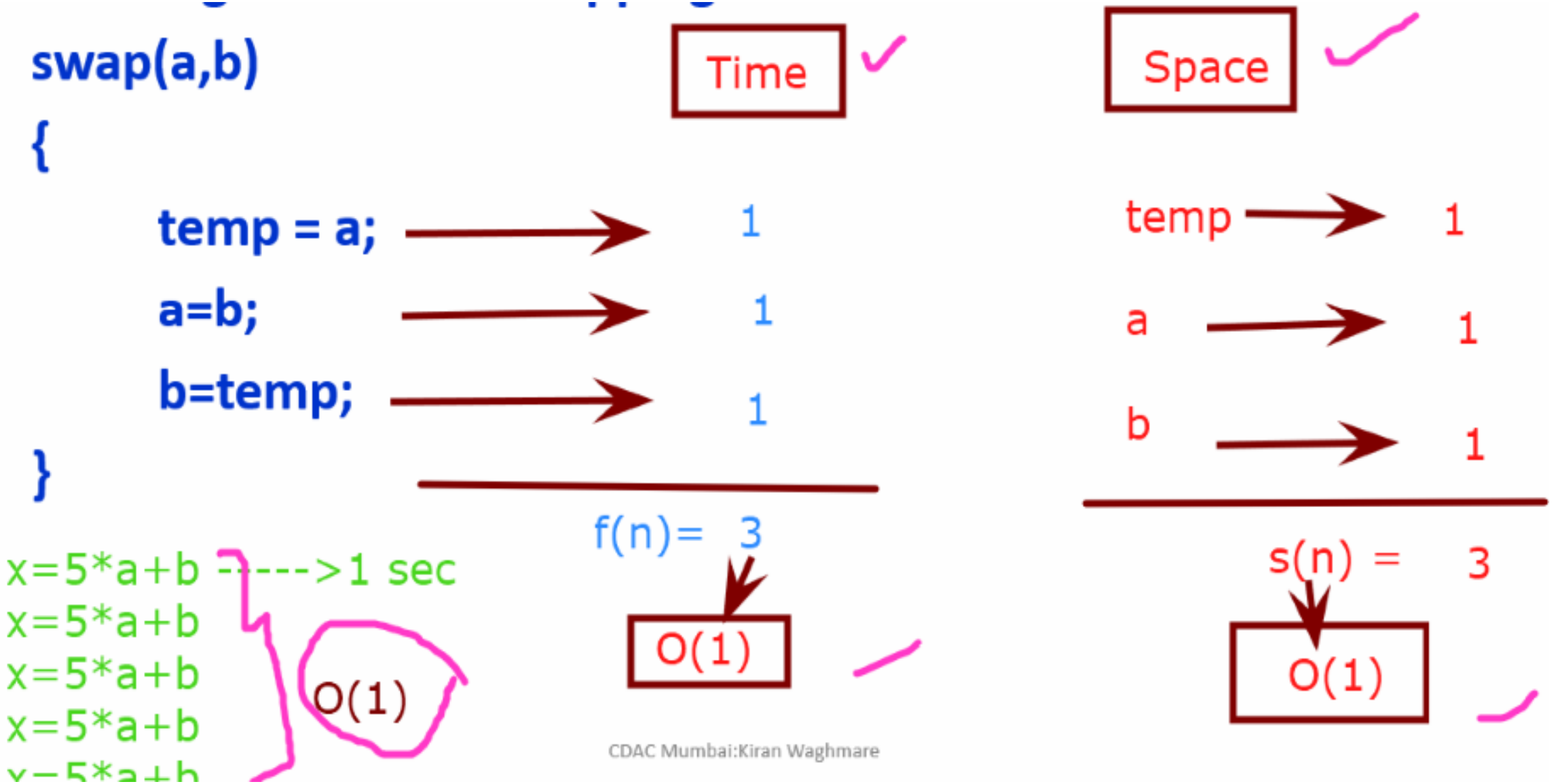
16

i=0,1,2,3,4,5

```
for(int i=0;i<n;i++)
{
    SOP("done");
}
```

n+1

n


swap(a,b)
{
temp = a;
a=b;
b=temp;
}
Time
1
1
1
f(n)= 3
O(1)
Space
temp
a
b
1
1
1
s(n) = 3
O(1)
x=5*a+b ------>1 sec
x=5*a+b
x=5*a+b
x=5*a+b
O(1)
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A[], n=5

```
sum(A,n)
{
    s=0;                    1
    for(i=0;i<n;i++)        n+1
    {
        s=s+A[i];           n
    }
    return s;               1
}
```

Time	Space

A → n (n)
n → 1
s → 1
i → 1

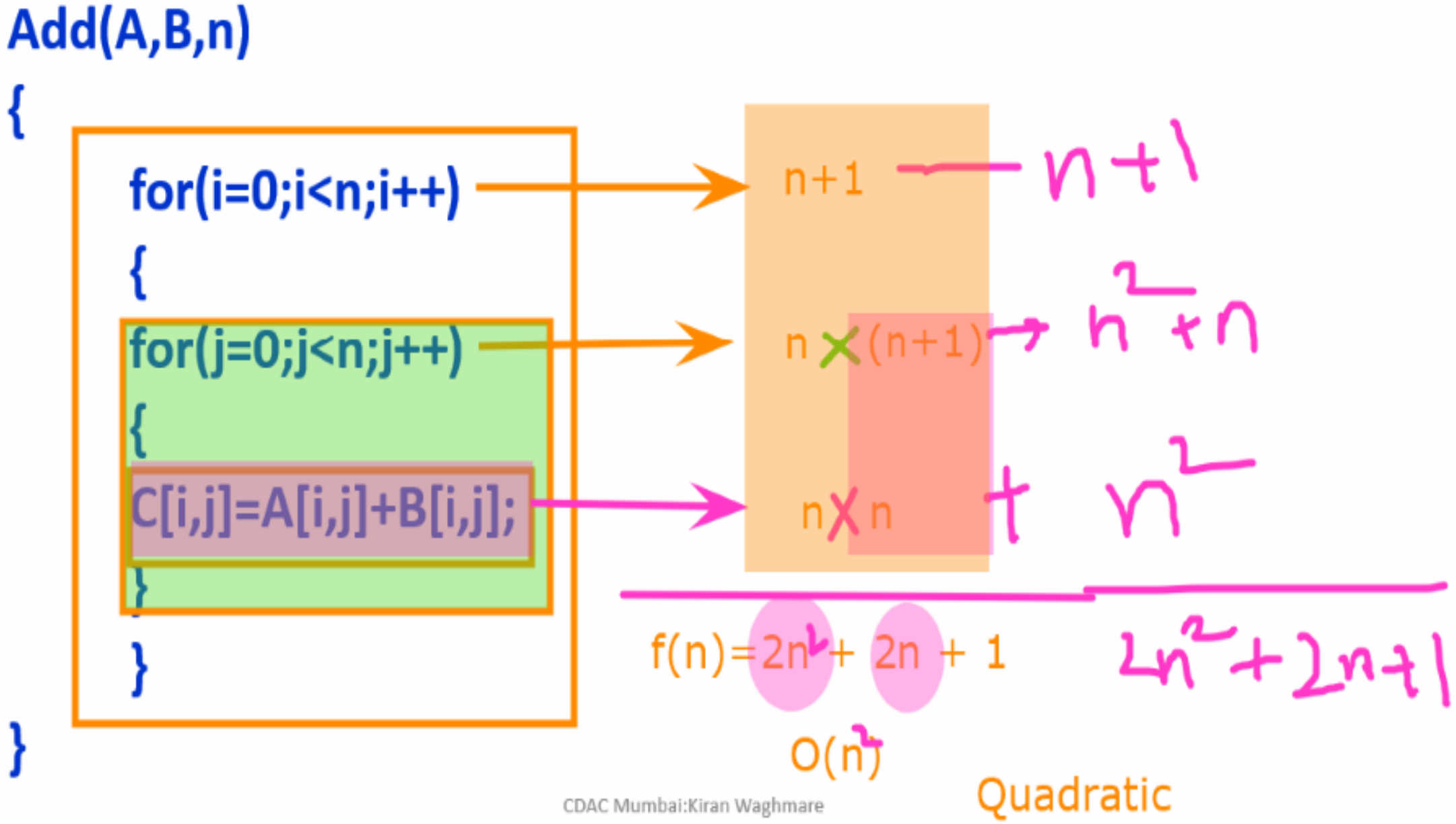
s(n)= n+3

O(n)

f(n)= 2n+3

O(n)

Linear Complexity


Add(A,B,n)
{
for(i=0;i<n;i++)
{
for(j=0;j<n;j++)
{
C[i,j]=A[i,j]+B[i,j];
}
}
}
n+1
n+1
n x (n+1)
n^2+n
n x n
n^2
f(n)=2n^2+ 2n + 1
2n^2+2n+1
O(n^2)
Quadratic
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Ex 5:	Ex 6:	Ex 7:

for(i=0;i<n;i++)
{
stmt;
}
O(n)
	for(i=n;i>0;i--)
	{
	stmt;
	}
	O(n)
		for(i=1;i<n;i+2)
		{
		stmt; — n/2
		}
		O(n)

Ex 8:	**Ex 9:**	**Ex 10:**

for(i=1;i<n;i=i+20)
{
stmt; — n/20
}
O(n)
	for(i=0;i<n;i++)
	{
	for(j=0;j<n;j++)
	{
	stmt;
	}
	}
	O(n²)
		for(i=0;i<n;i++)
		{
		for(j=0;j<i;j++)
		{
		stmt;
		}
		}
		O(n²)

The order of growth for all time complexities are indicated in the graph below: