



DATA STRUCTURES AND ALGORITHMS

Sep22 : Day 3

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CDAC Mumbai

Date: 31-10-2022

Day 5: Algorithms & Data Structures

Topics:

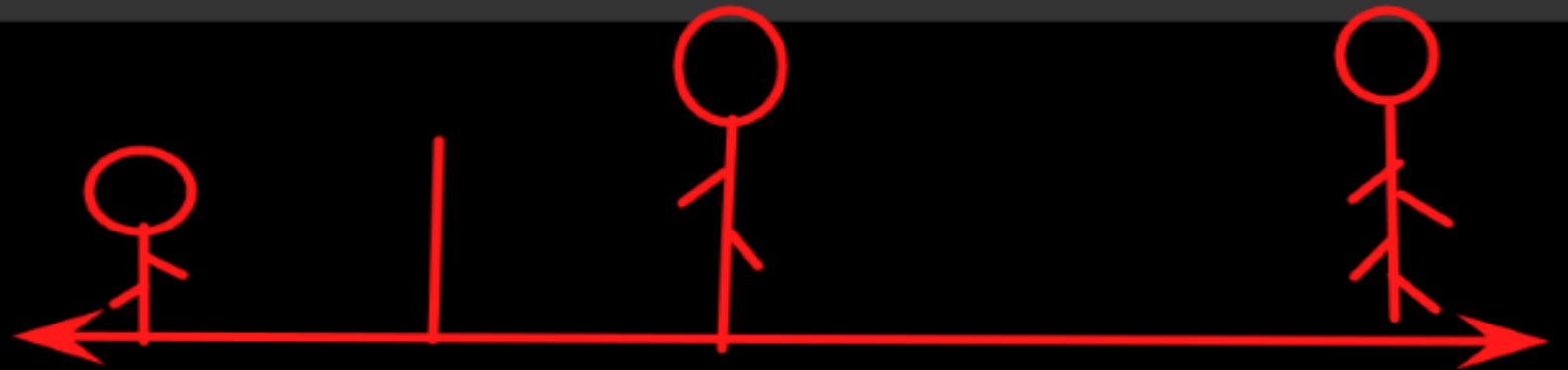
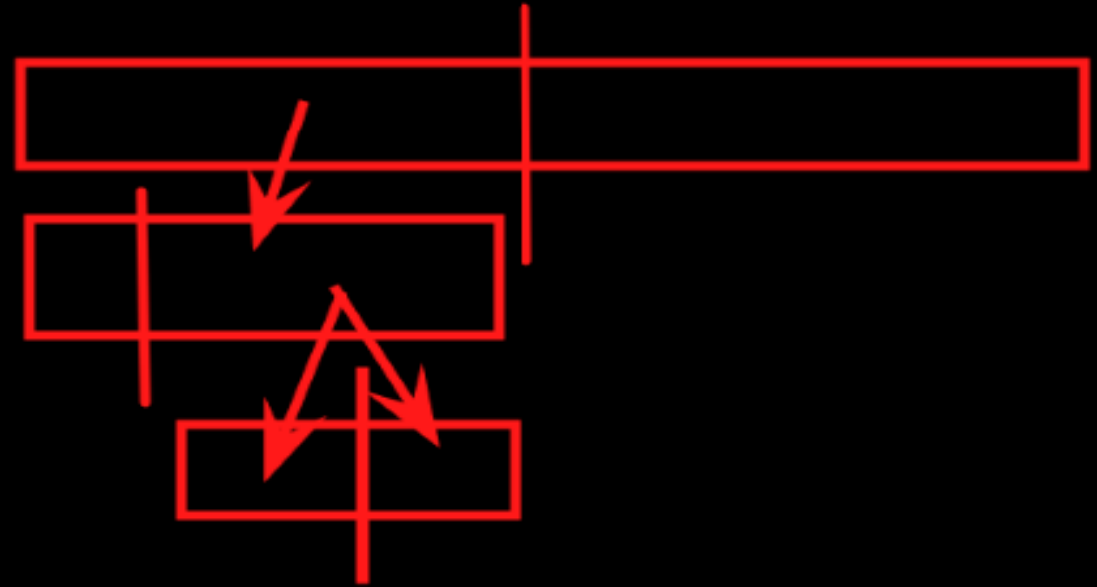
- Sorting
- Heap
- Linked list

Quick Sorting:

-divide and conquer

-3 steps:

- Identify pivot element
- Partition
- recursion



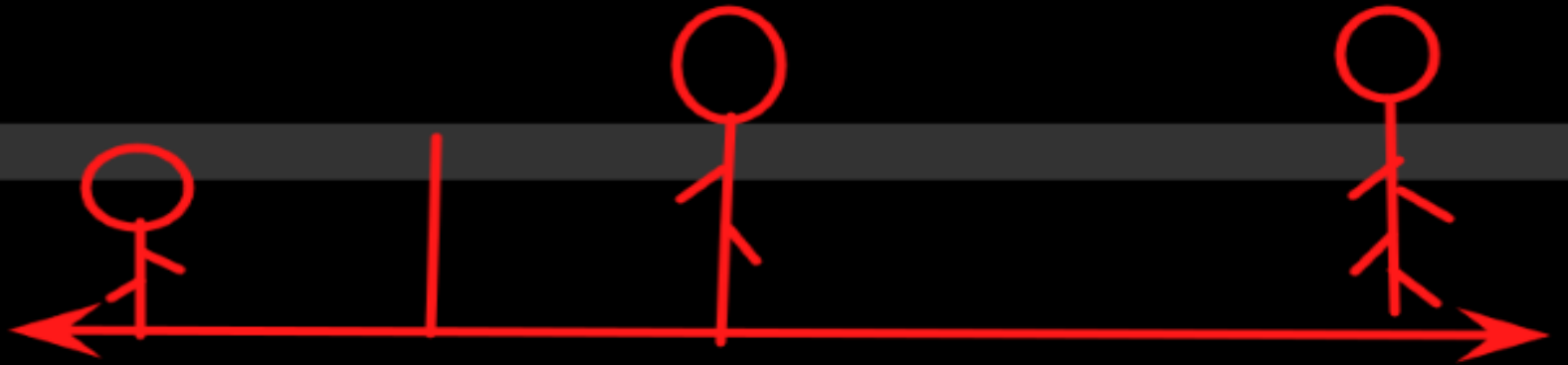
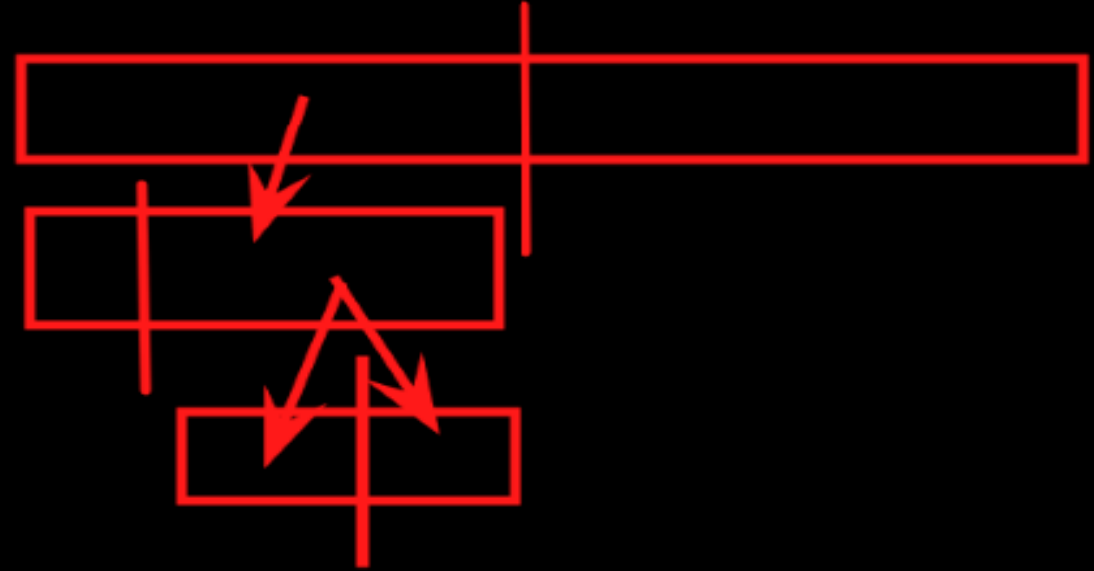
-Partition
-recursion

10 80 70 90 30 20

80 70 90 30 20 90 .

10 20 30 40 79 60 90 50

10 5 7 9 11 12 18 22



Example:

10

16

8

12

15

6

3

9

5



pivote

low

high

cond
pivot > j
pivote < i

10 5 8 12 15 6 3 9 16

10 5 8 9 15 6 3 12 16

10 5 8 9 3 6 15 12 16

j i

pivot = j

6 5 8 9 3 10 15 12 16

6 5 3 9 8

3 5 6 9 8

The following procedure implements quicksort:

QUICKSORT(A, p, r)

```
1  if  $p < r$ 
2       $q = \text{PARTITION}(A, p, r)$ 
3      QUICKSORT( $A, p, q - 1$ )
4      QUICKSORT( $A, q + 1, r$ )
```

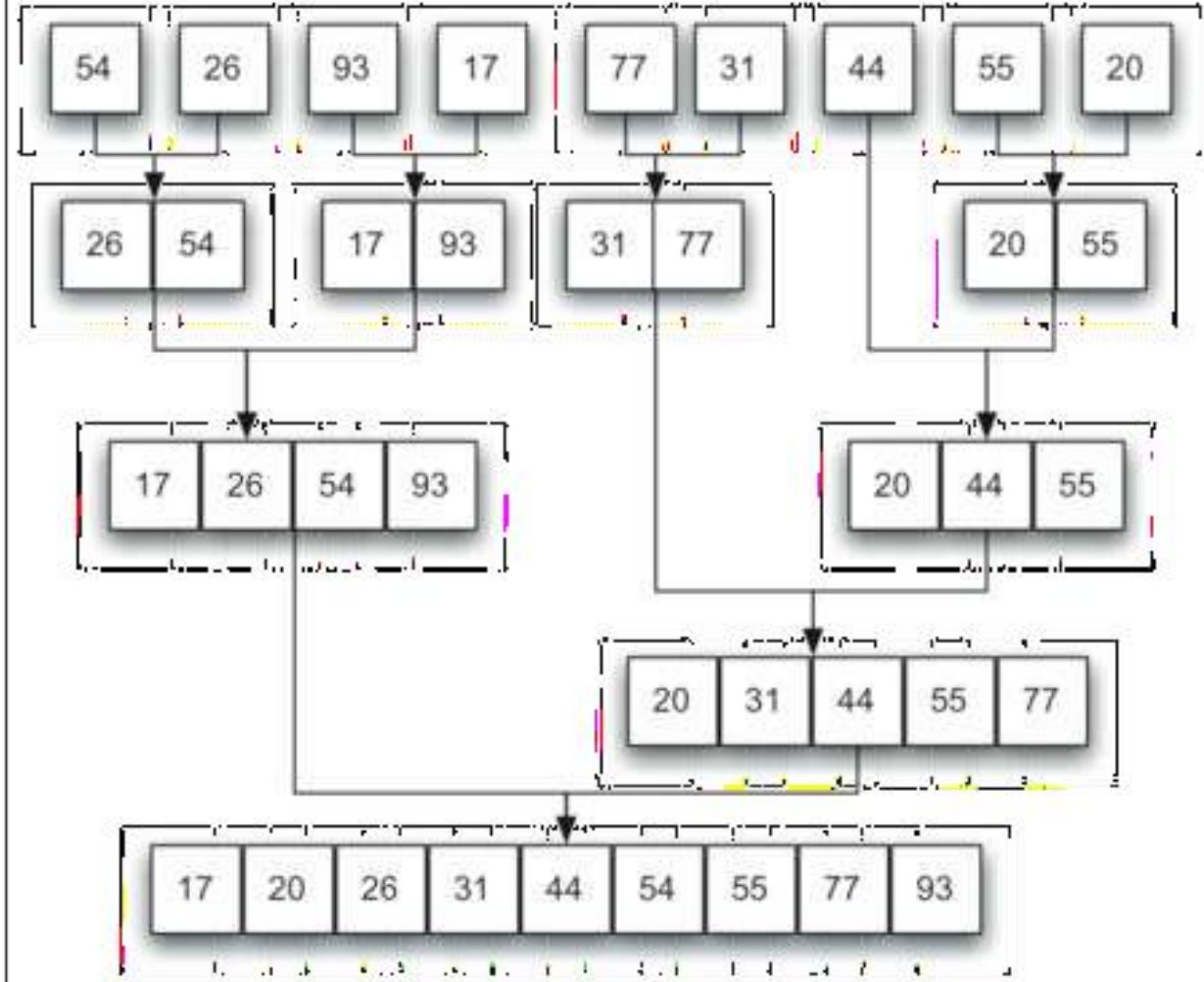
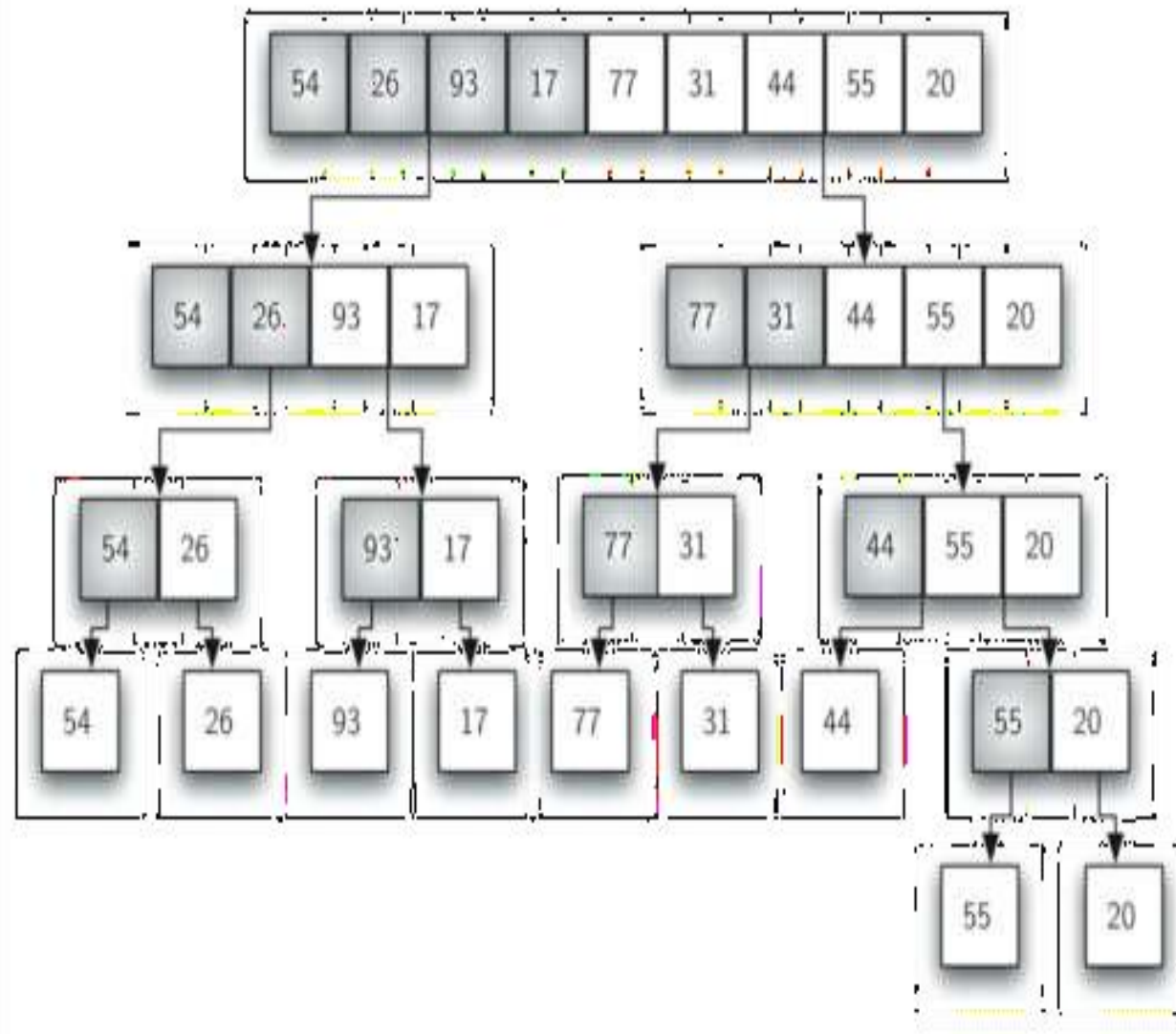
To sort an entire array A , the initial call is QUICKSORT($A, 1, A.length$).

Partitioning the array

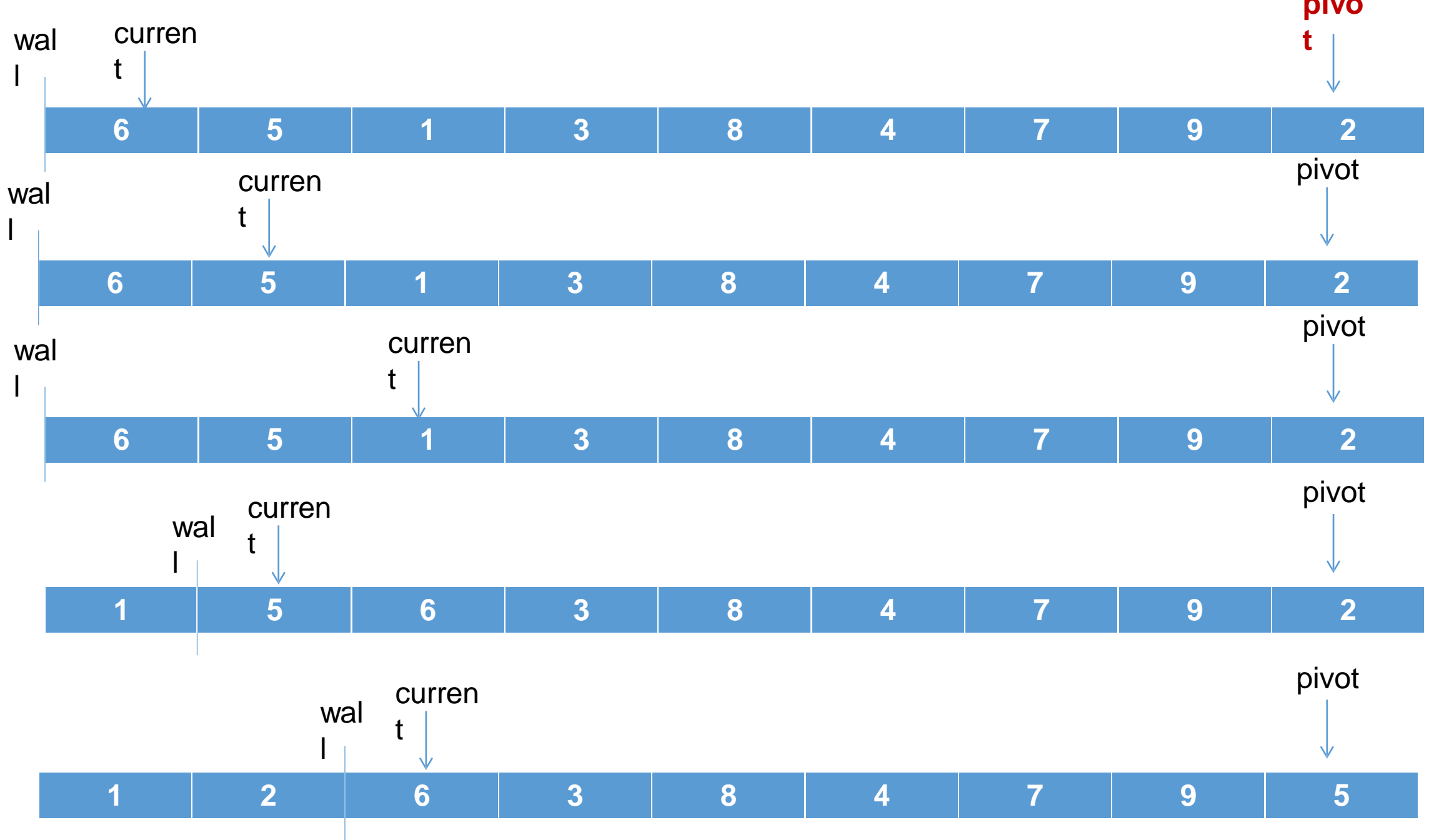
The key to the algorithm is the PARTITION procedure, which rearranges the subarray $A[p..r]$ in place.

PARTITION(A, p, r)

```
1   $x = A[r]$ 
2   $i = p - 1$ 
3  for  $j = p$  to  $r - 1$ 
4      if  $A[j] \leq x$ 
5           $i = i + 1$ 
6          exchange  $A[i]$  with  $A[j]$ 
7  exchange  $A[i + 1]$  with  $A[r]$ 
8  return  $i + 1$ 
```



```
int pi=partition(a1, low, high);  
quicksort(a1, low, pi-1); //Left side  
quicksort(a1, pi+1, high); //Right side  
}  
}  
  
static int partition(int a1[], int low, int high)  
{  
    int pivot=a1[high];  
    int i=(low);  
  
    for(int j=low; j<high; j++)  
    {  
        if(a1[j] < pivot)  
        {  
            i++;  
            swap(a1, i, j);  
        }  
    }  
  
    swap(a1, i+1, high);  
    return (i+1);  
}
```



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```
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Partitioning the array

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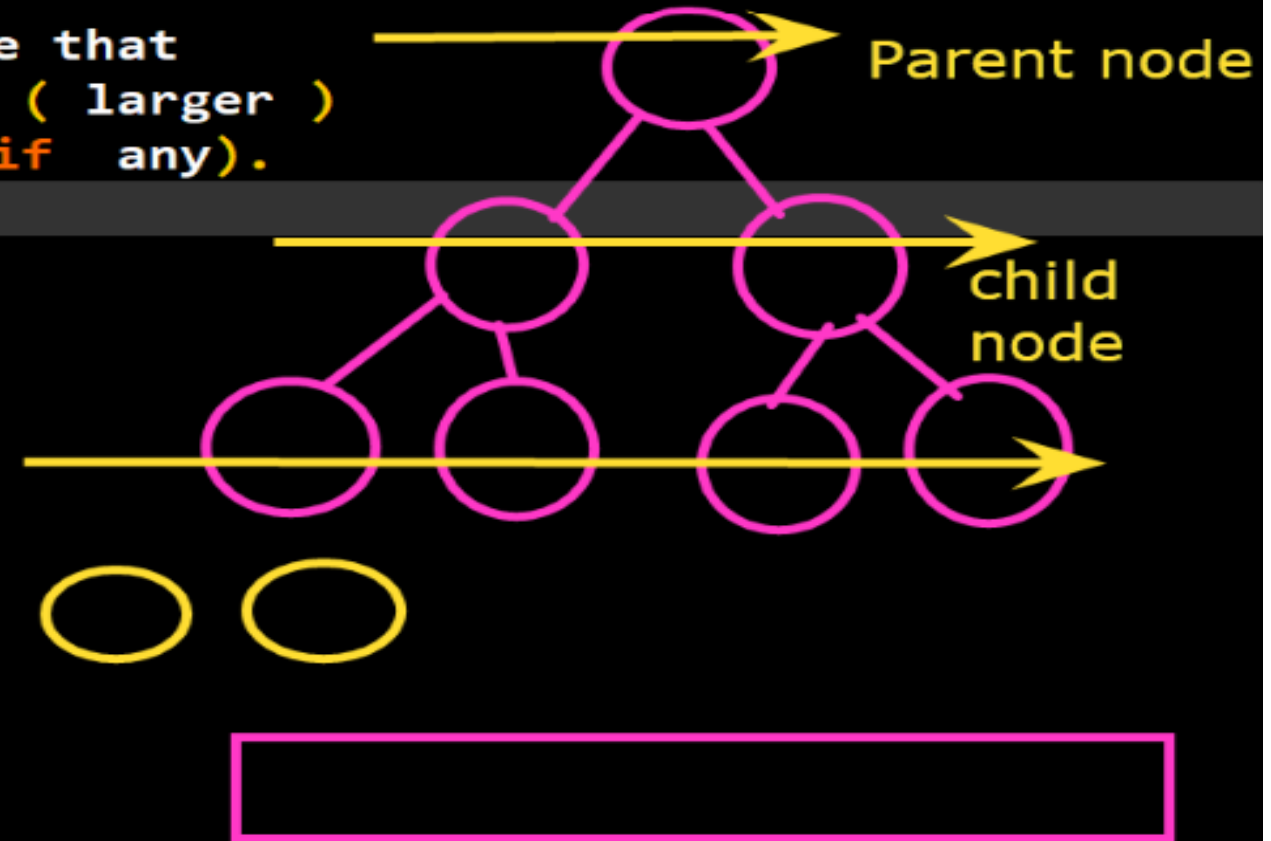
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8  return  $i + 1$ 
```

Heap:

Definition:

A special form of complete binary tree that key value of each node is no smaller (larger) than the key value of its children (if any).



Heap

- **Definition in Data Structure**

- **Heap:** A special form of **complete binary tree** that key value of each node is no smaller (larger) than the key value of its children (if any).

- **Max-Heap: root node has the largest key.**

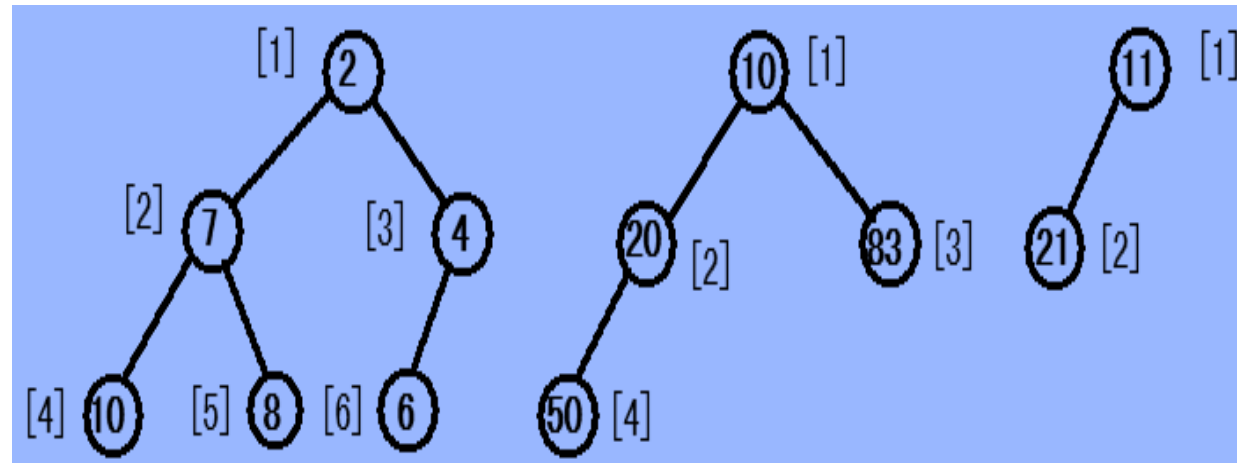
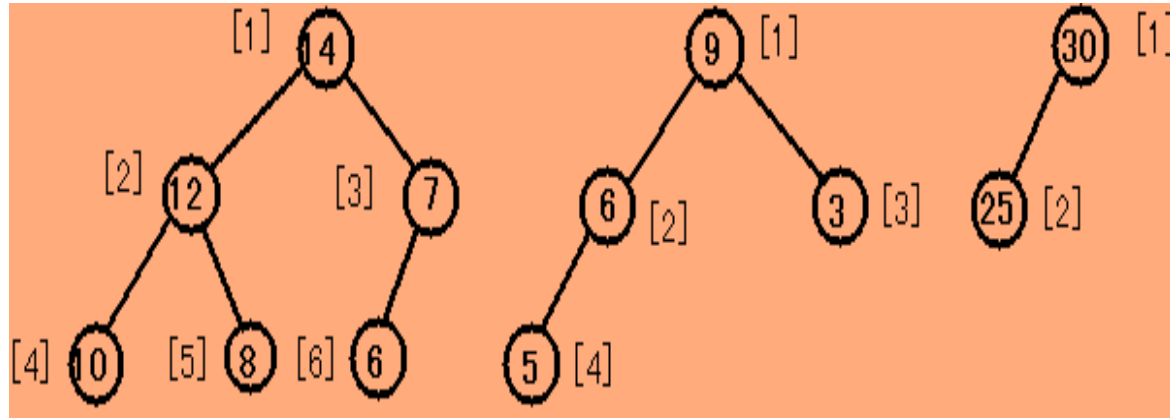
- A **max tree** is a tree in which the key value in each node is **no smaller than** the key values in its children.
- A **max heap** is a **complete binary tree** that is also a max tree.

- **Min-Heap: root node has the smallest key.**

- A **min tree** is a tree in which the key value in each node is **no larger than** the key values in its children.
- A **min heap** is a **complete binary tree** that is also a min tree.

Heap

- **Example:**
 - Max-Heap
 - Min-Heap



Heap:

Definition:

A special form of complete binary tree that key value of each node is no smaller (larger) than the key value of its children (if any).

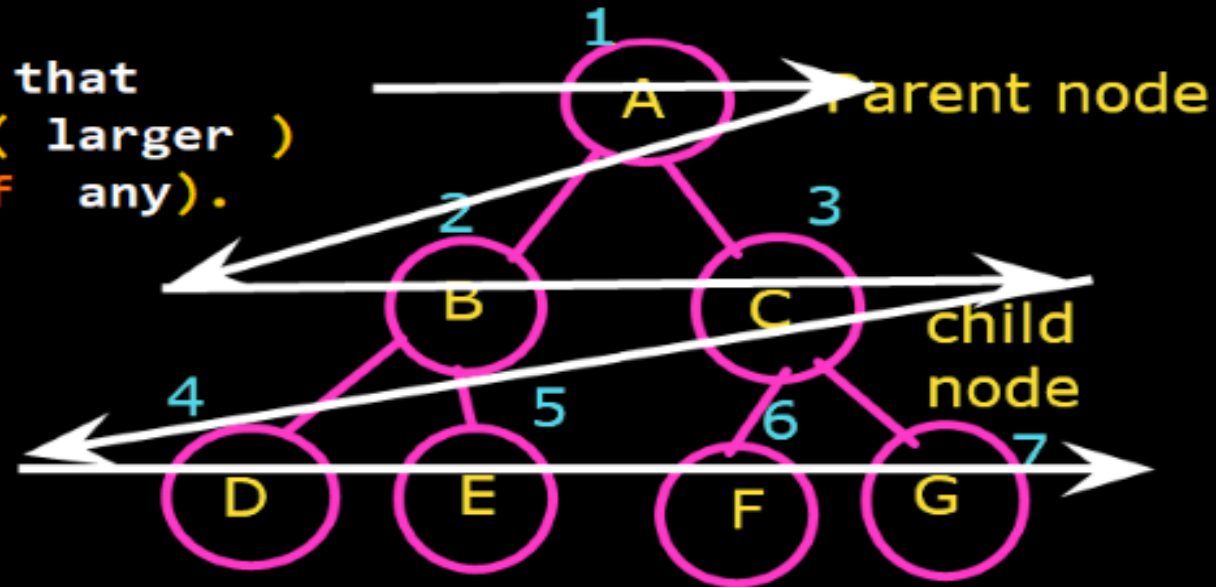
Node=index = i

Left child = $2*i$

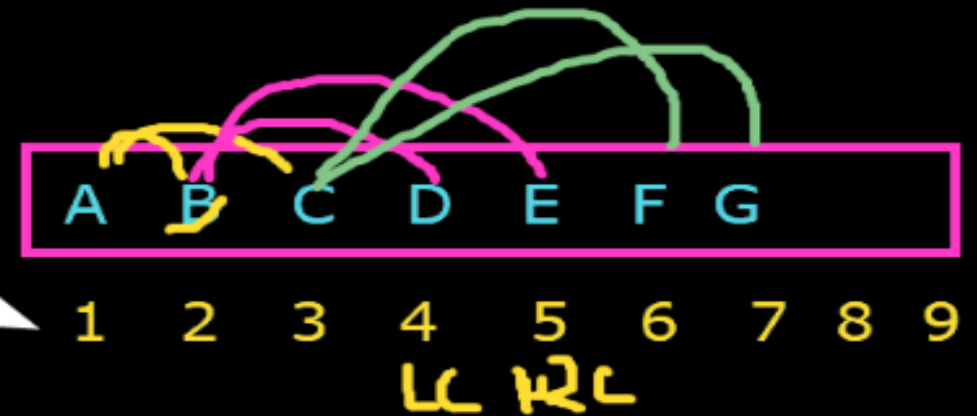
Right child = $2*i+1$

Parent = $i/2$

$i=2$ B
LC=4 D
RC=5 E
P=1



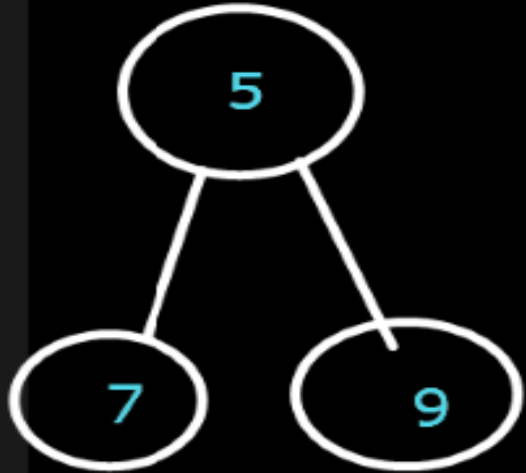
n elements
-int
-char
-float
-string



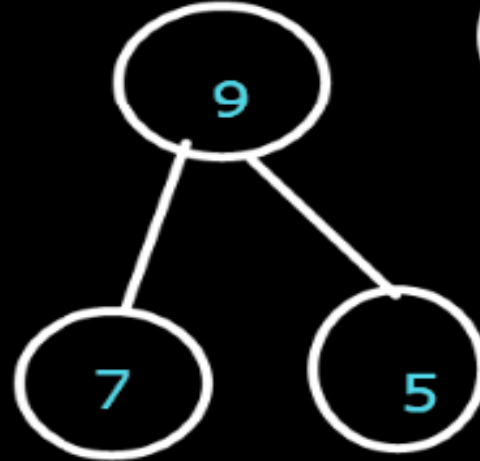
Heap:

Definition:

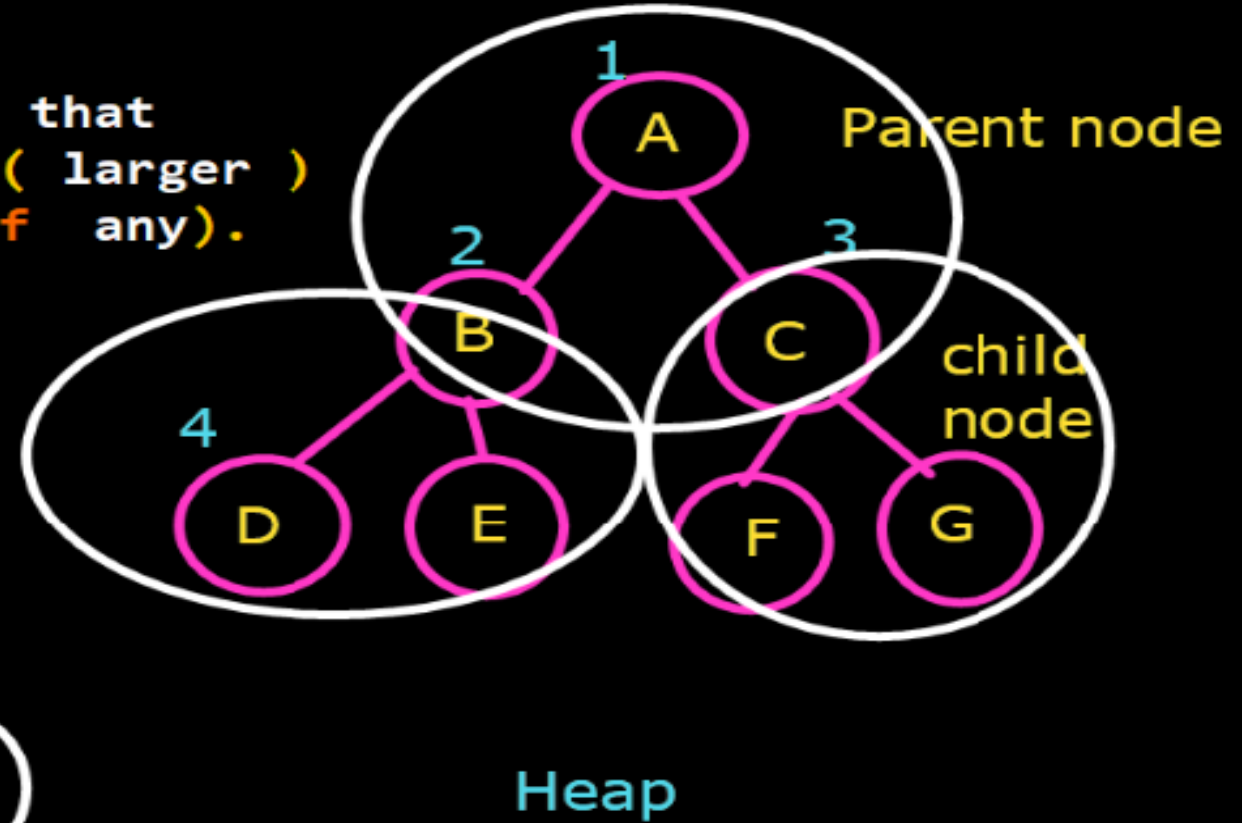
A special form of complete binary tree that key value of each node is no smaller (larger) than the key value of its children (if any).



Min heap



Max heap

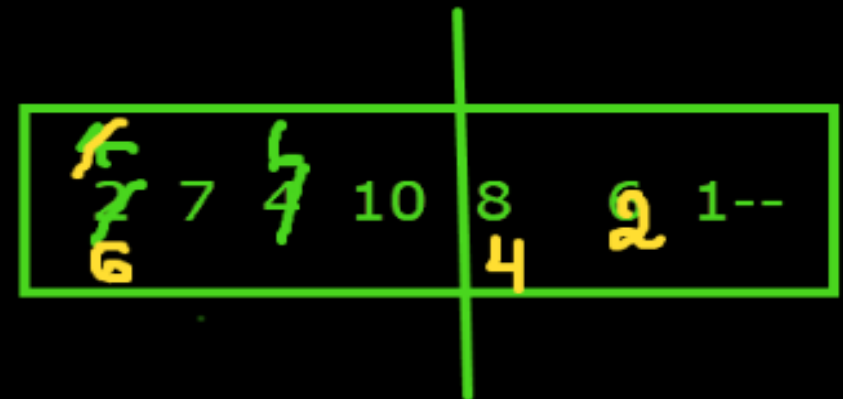
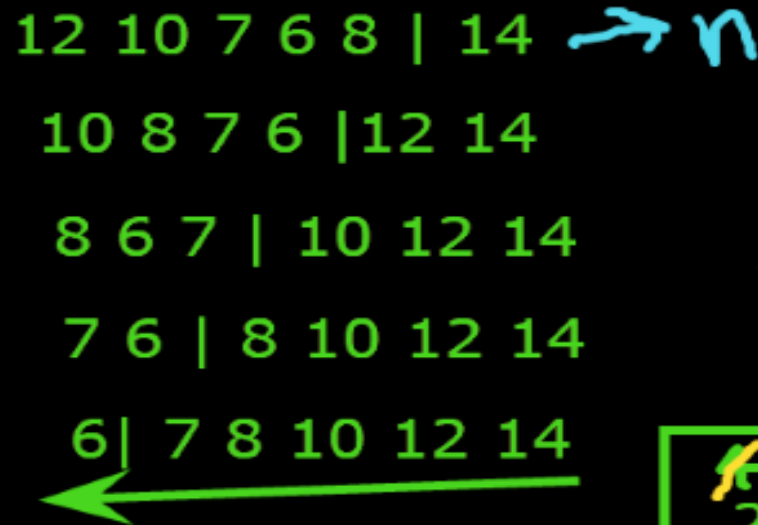


.....

A special form of complete binary tree that
key value of each node is no smaller (larger)
than the key value of its children (if any).

n: elements
Heapify : $O(n)$

2 7 4 10 8 6



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```

    a1[i]=a1[largest];
    a1[largest]=temp;
    heapify(a1, n, largest);

```

```

}

```

```

}

```

```

static void heapsort(int a1[])

```

```

{

```

```

    int n = a1.length; ✓

```

```

    for(int i=n/2-1; i>=0; i--)

```

```

        heapify(a1, n, 0);

```

```

    for(int i=n-1; i>=0; i--)

```

```

    {

```

```

        int temp = a1[0];

```

```

        a1[0]=a1[i];

```

```

        a1[i]=temp;

```

```

        heapify(a1, i, 0); //balancing the max heap

```

```

    }

```

```

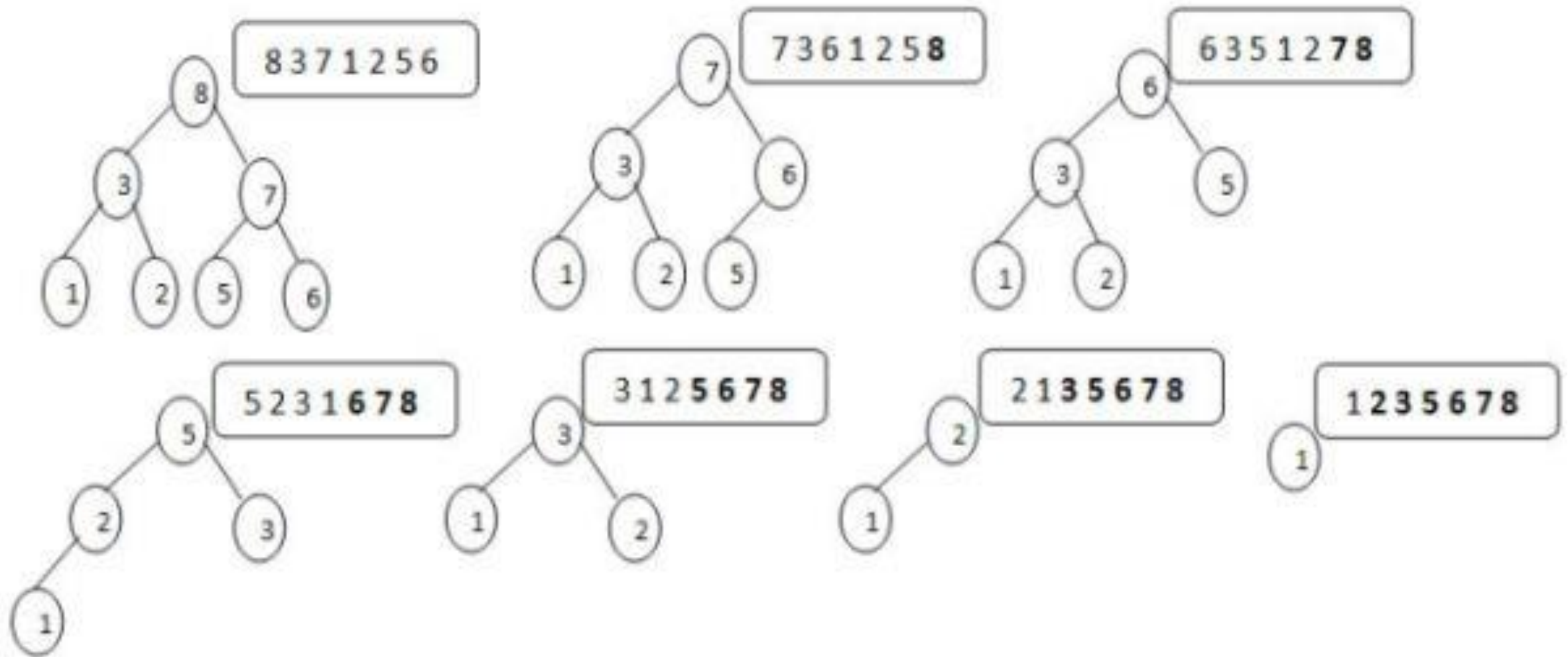
}

```

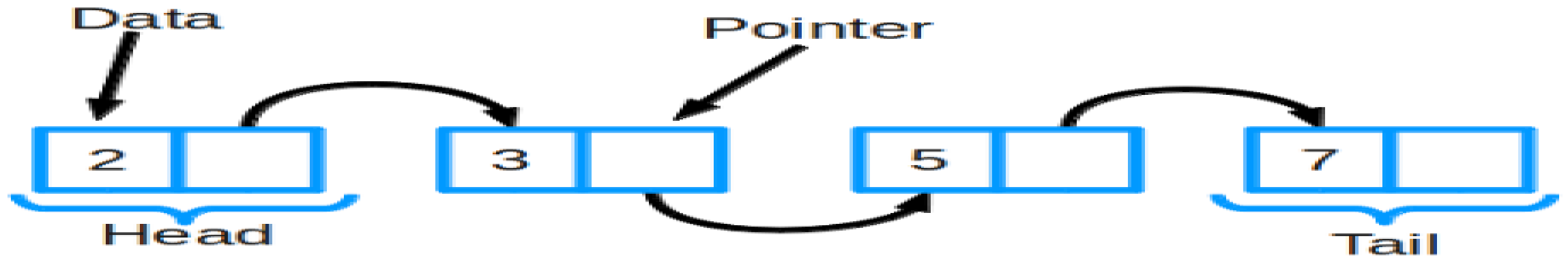
n
 $n-1$



Example:- The fig. shows steps of heap-sort for list (2 3 7 1 8 5 6)



Linked list



```
}
```

Linked list:

- sequence of data structures, which are connected together via links.
- sequence of links
- connections between nodes
- most used data structure

Terms:

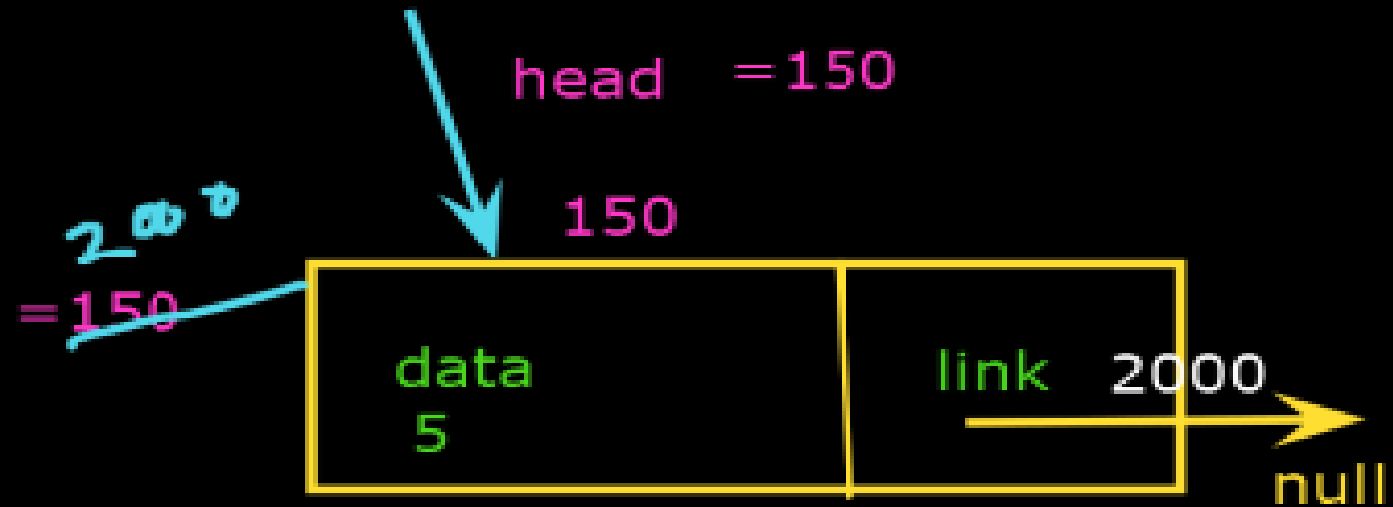
- Link
- Next
- Data
- Linked list



- connections between nodes
- most used data structure

Terms:

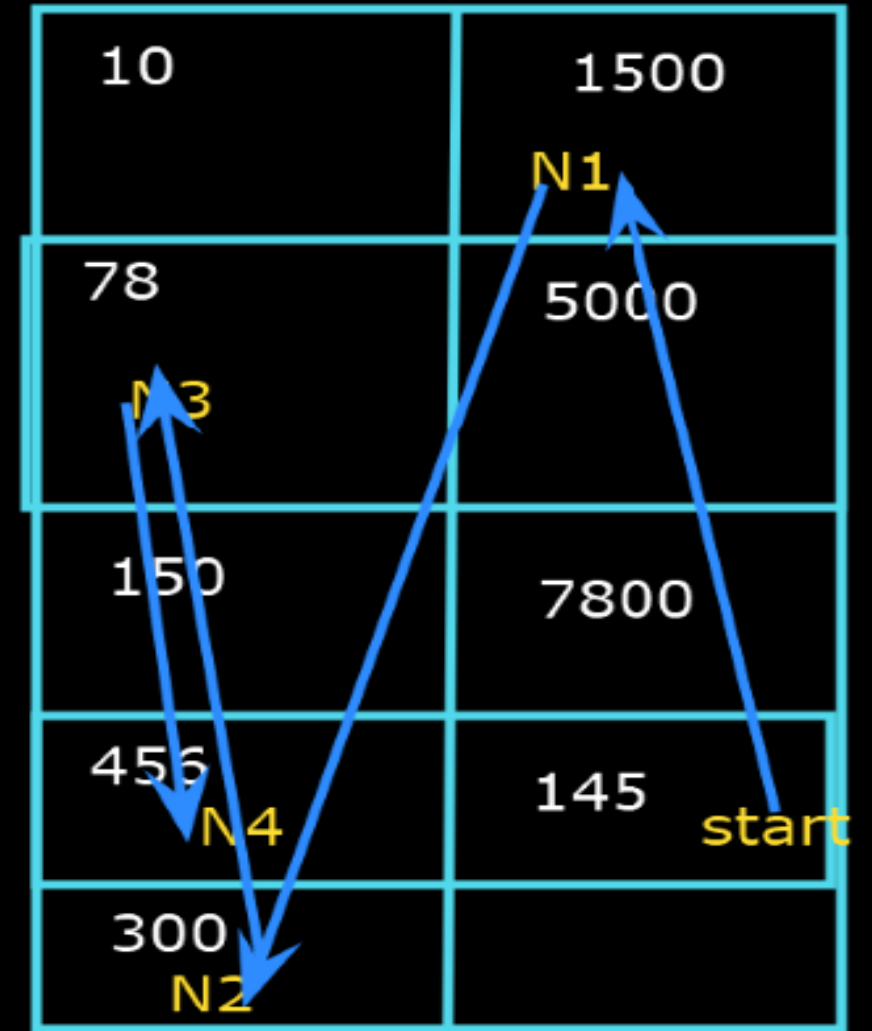
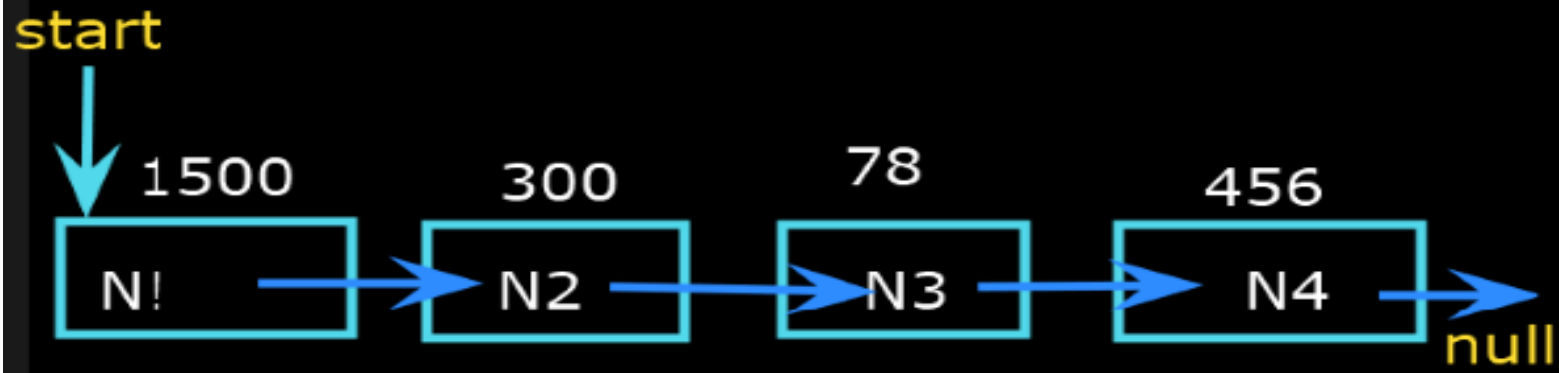
- Link
- Next
- Data
- Linked list



Node

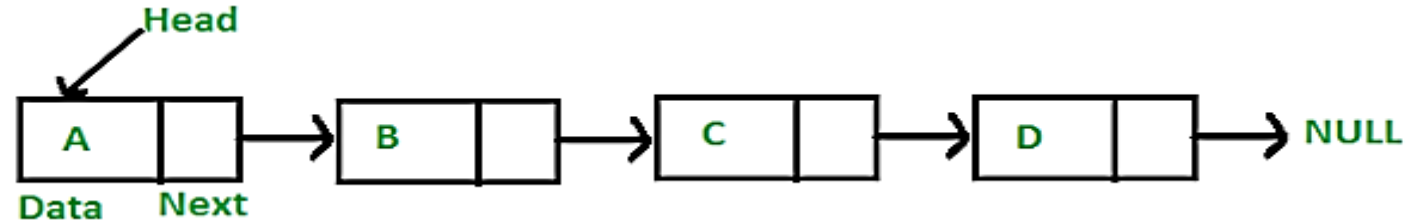


-Linked list



Linked List Representation

- Linked list can be visualized as a chain of nodes, where every node points to the next node.



- As per the above illustration, following are the important points to be considered.
 1. Linked List contains a **link element** called **first**.
 2. Each link carries a **data field(s)** and a **link field** called **next**.
 3. Each link is **linked with its next link** using its **next link**.
 4. **Last link carries a link as null** to mark the end of the list.

}

Linked list:

- sequence of data structures, which are connected together via links.
- sequence of links
- connections between nodes
- most used data structure
- provides lot of flexibility

Terms:

- Link : data=element, link=address
- Next : next is a link: address
- Data : any primitive data types
- Linked list : connection of links :
- First node of linked list = starting node of list
- Last node of linked list = link is null
- chain of nodes....



1. Simple linked list

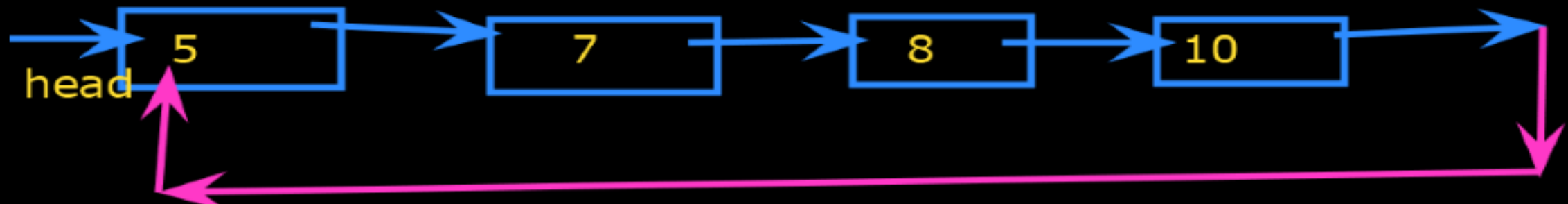
-navigation is in forward direction

2. Doubly linked list

-navigation is in forward and backward direction

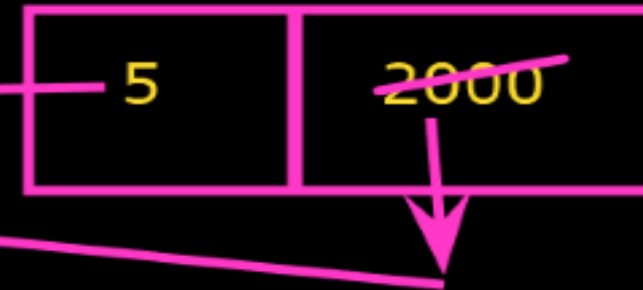
3. Circular linked list

-last item contains link of the first element



Node structure:

```
class Node{  
  
    int data;  
    Node next;  
    Node(int d)  
    {  
        data = d;  
        next = null;  
    }  
}
```



- Last item contains link of the first element

Node structure:

-----singly Linked List-----

```
class Node{
```

```
    int data;
```

```
    Node next;
```

```
    Node(int d)
```

```
{
```

```
    data = d;
```

```
    next = null;
```

```
}
```

```
}
```



Singly Linked list



doubly linked list

Node head;

```
static class Node{
    int data;
    Node next;
    Node(int d)
    {
        data = d;
        next = null;
    }
}

public static void main(String args[])
{
    List1 l1 = new List1();
    l1.head = new Node(11);
    Node second = new Node(22);
    Node third = new Node(33);

    l1.head.next = second;
    second.next = third;
}
```

