

1. Perform the following steps:

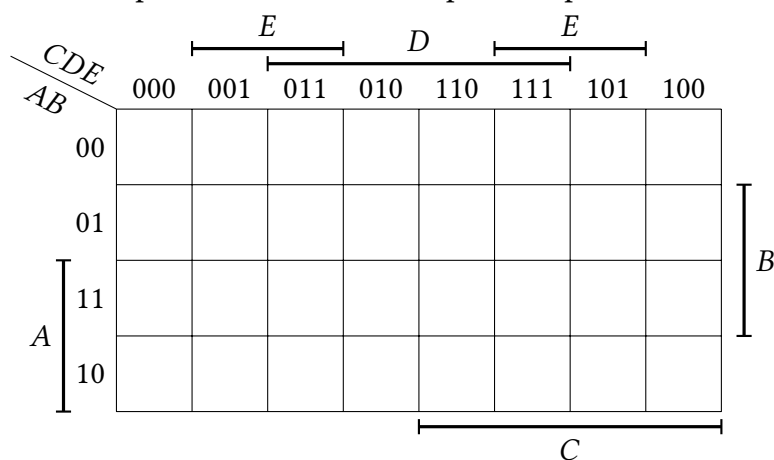
- (a) Calculate the SHA-256 hash  $h$  of the string  $s = \text{"DM Fall 2023 HW3"}$  (without quotes, with all spaces, encoded in UTF-8). Convert hash  $h$  to a 256-bit binary string  $b$  (prepend leading zeros if necessary). Cut the binary string  $b$  into eight 32-bit slices  $r_1, \dots, r_8$ , e.g.  $r_2 = b_{33..64}$ . Xor all slices into a 32-bit string  $d = r_1 \oplus \dots \oplus r_8$ . Compute  $w = d \oplus 0x24d03294$ .

*Hint:* last (least significant) bits of  $h$  are ...01001001, last bits of  $d$  are ...0001.

- (b) Draw the Karnaugh map (use a template below) for a function  $f(A, B, C, D, E)$  defined by the truth table  $w = (w_1 \dots w_{32})$ , where MSB corresponds to  $f(0) = w_1$  and LSB to  $f(1) = w_{32}$ .

- (c) Use K-map to find the minimal DNF and minimal CNF for the function  $f$ .

- (d) Use K-map to find the number of prime implicants, i.e. the size of BCF.



2. For each given function  $f_i$  of 4 arguments, draw the Karnaugh map and use it to find BCF, minimal DNF, and minimal CNF. Additionally, construct ANF (Zhegalkin polynomial) using either the K-map, the tabular ("triangle") method or the Pascal method – use each method at least once.

**Note:** WolframAlpha<sup>1</sup> interprets the query " $n$ -th Boolean function of  $k$  variables" in a reverse manner. In order to employ WolframAlpha properly, manually flip the truth table beforehand, e.g. the correct query for  $f_{10}^{(2)}$  is "5th Boolean function of 2 variables"<sup>2</sup>, which gives  $f_{10}^{(2)} = \neg x_2$ , since  $\text{rev}(1010_2) = 0101_2 = 5_{10}$ .

(a)  $f_1 = f_{47541}^{(4)}$

(b)  $f_2 = \sum m(1, 4, 5, 6, 8, 12, 13)$

(c)  $f_3 = f_{51011}^{(4)} \oplus f_{40389}^{(4)}$

(d)  $f_4 = ABD + \bar{A}\bar{C}D + \bar{B}C\bar{D} + A\bar{C}D$

3. Convert the following formulae to CNF.

(a)  $X \leftrightarrow (A \wedge B)$

(b)  $Z \leftrightarrow \bigvee_i C_i$

(c)  $D_1 \oplus \dots \oplus D_n$

(d)  $\text{majority}(X_1, X_2, X_3)$ <sup>1</sup>

(e)  $R \rightarrow (S \rightarrow (T \rightarrow \bigwedge_i F_i))$

(f)  $M \rightarrow (H \leftrightarrow \bigvee_i D_i)$

4. For each given system of functions  $F_i$ , determine whether it is functionally complete using Post's criterion. For each basis  $F_i$ , use it to represent the majority( $A, B, C$ ) function. Draw a combinational Boolean circuit for each resulting formula.

(a)  $F_1 = \{\wedge, \vee, \neg\}$

(b)  $F_2 = \{f_{14}^{(2)}\}$

(c)  $F_3 = \{\rightarrow, \nrightarrow\}$

(d)  $F_4 = \{1, \leftrightarrow, \wedge\}$

5. Show – without using Post's criterion – that the Zhegalkin basis  $\{\oplus, \wedge, 1\}$  is functionally complete.

<sup>1</sup> Majority function<sup>2</sup> is a Boolean function that is 1 iff the majority (more than half) of the inputs are 1.

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- The diagram shows a logic circuit for a 3-input majority function. It has three inputs: A, B, and C. The circuit consists of several logic gates: AND gates, OR gates, and a final OR gate. The output is labeled  $f(1)$ .

- Binary Decision Diagram <sup>[↗](#)</sup> (BDD) is a representation of a Boolean function as a directed acyclic graph, which consists of *decision* nodes and two *terminal* nodes (0 and 1). Each decision node is labeled by a Boolean variable  $x_i$  and has two child nodes called *low* and *high*. The edge from node to a low (high) child represents an assignment of the value FALSE (TRUE) to variable  $x_i$ . A path from the root node to the 1-terminal (0-terminal) corresponds to an assignment for which the represented Boolean function is true (false). BDD is called *ordered* if variables appear in the same order on all paths from the root. BDD is called *reduced* if it does not contain a node  $v$  with  $\text{high}(v) = \text{low}(v)$ , and there does not exist a pair of nodes  $u, v$  such that the sub-OBDDs rooted in  $u$  and  $v$  are isomorphic.

- (a)  $f_1(x_1, \dots, x_4) = x_1 \oplus x_2 \oplus x_3 \oplus x_4$  (c)  $f_3(x_1, \dots, x_4) = \sum m(1, 2, 5, 12, 15)$   
 (b)  $f_2(x_1, \dots, x_5) = \text{majority}(x_1, \dots, x_5)$  (d)  $f_4(x_1, \dots, x_6) = x_1x_4 + x_2x_5 + x_3x_6$