

## Workshop 1: Complexity Notation

**Workshop 1** will take place in week 2. You should prepare solutions, but you don't have to hand them in and they won't get marked.

### Exercise 1 *Complexity Notation*

Solve Exercise 10 and 11 in the book of Mehlhorn/Sanders (Chapter 2.1).

### Exercise 2 *Complexity Notation*

Is it true that if  $f(n) = \Theta(g(n))$  and  $g(n) = \Theta(h(n))$ , then  $h(n) = \Theta(f(n))$ ?

### Exercise 3 *Complexity Notation*

Is it true that if  $f(n) = O(g(n))$  and  $g(n) = O(h(n))$ , then  $h(n) = \Omega(f(n))$ ?

### Exercise 4 *Complexity Notation*

Is it true that a  $\Theta(n^2)$  algorithm always takes longer to run than a  $\Theta(\log n)$  algorithm?

### Exercise 5 *Complexity Notation*

For each pair of functions given below, point out the asymptotic relationships that apply:  
 $f = O(g)$ ,  $f = \Theta(g)$ ,  $f = \Omega(g)$ .

- |   |   |               |
|---|---|---------------|
| • $f(n) = \sqrt{n}$ and $g(n) = \log(n)$      | $f(n) > g(n) \longrightarrow f = \Omega(g)$ | ✓             |
| • $f(n) = 1$ and $g(n) = 2$                   | $f = \Theta(g)$                             | ✓             |
| • $f(n) = 1000 \cdot 2^n$ and $g(n) = 3^n$    | $f = O(g)$                                  | ✓             |
| • $f(n) = 4^{n+4}$ and $g(n) = 2^{2n+2}$      | $f = \Omega(g)$                             | × $\emptyset$ |
| • $f(n) = 5n \log(n)$ and $g(n) = n \log(5n)$ | $f = O(g)$                                  | × $\emptyset$ |
| • $f(n) = n!$ and $g(n) = (n+1)!$             | $f = \Theta(g)$                             | × $\emptyset$ |

### Exercise 6 *Complexity Notation*

Prove that  $n^k = o(c^n)$  for any integer  $k$  and any  $c > 1$ .

## Exercise 1

Ans

Ex.10: a) $n^2 + 10^6 n \in O(n^2)$	✓
b) $n \log(n) \in O(n)$	✗
c) $n \log(n) \in \Omega(n)$	✓
d) $\log(n) \in o(n)$	✓

Ex.11:

### Lemma 6

1.  $c \cdot f(n) = \Theta(f(n))$  for  $\forall c > 0$ .
2.  $f(n) + g(n) = \Omega(f(n))$ .
3.  $f(n) + g(n) = O(f(n))$  if  $g(n) = O(f(n))$ .
4.  $O(f(n)) \cdot O(g(n)) = O(f(n) \cdot g(n))$ .

Proof: 1) For  $n \geq 1$ ,  $c \cdot f(n) \geq f(n)$

For  $n \leq 1$ ,  $c \cdot f(n) \leq f(n)$

Since,  $n=1$  satisfies upper & lower bound,  $\Theta(f(n)) = c \cdot f(n)$

$$2) \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)}$$

### L'Hopital's Rule:

- 1) Take derivative until get out of ' $\frac{0}{0}$ ' or ' $\frac{\infty}{\infty}$ '
- 2) Make sure that after subbing in  $n \rightarrow \infty$  that the ratio is no longer ' $\frac{\infty}{\infty}$ ' or ' $\frac{0}{0}$ '

## Exercise 4

$$\lim_{n \rightarrow \infty} \frac{n^2}{\log(n)} \longrightarrow \frac{2n}{\frac{1}{n}} \longrightarrow 2n^2$$