# ACM-ICPC TEAM REFERENCE DOCUMENT Izhevsk State Technical University (Gogolev, Kismatov, Lykov)

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# 1 Common

# 1.1 template

```
optional, complex, string, vector, list, map, set, deque, queue,
      stack, bitset,
algorithm, sstream, iostream, iomanip, cstdio, cstdlib, ctime,
      cstring, cmath,
cstdarg, cassert, ctime, tuple, unordered_set, unordered_map,
      random, chrono
using namespace std;
#define ln "\n"
#define pb push_back
#define mp make_pair
#define ins insert
\#define sz(x) (int)x.size()
#define All(x) (x).begin(), (x).end()
#define fd(a, b) for (auto i = (a); i <= (b); i++) #define fd(a, b) for (auto i = (a); i >= (b); i--)
#define fx(A) for (auto &x : (A))
typedef long long ll;
typedef vector<int> vi;
typedef vector<vi> vvi;
typedef pair<int, int> pii;
template<typename T>
ostream& operator<<(ostream& os, vector<T> v) {
        fz(v) {
         return os;
#ifdef LOCAL
         \# else
         \#define dbg(x) \{\}
#endif
\#ifdef\ LOCAL
    #define ass(x) if (!(x)) { cerr << _LINE__ << tassertion_failed: _ " << #x << ln; abort(); }
#define ass(x) assert(x) #endif
mt19937\_64\ rnd(chrono::steady\_clock::now().time\_since\_epoch
int main()
    \begin{array}{l} ios::sync\_with\_stdio(0); \\ cin.tie(0); \end{array}
    cout.tie(0)
    \#ifdef\ LOCAL
```

```
freopen("input.txt", "r", stdin); \\ freopen("output.txt", "w", stdout); \\ \#endif \\ \}
```

## 1.2 compile

```
g++ --std=c++20 !.! -DLOCAL -O2 -Wall -Wl,--stack =67108864 -o !.exe
```

#### 1.3 check

```
:start
gen.exe
sol.exe
copy output.txt answer.txt
NAME.exe
fc output.txt answer.txt
if ERRORLEVEL 1 goto end
goto start
:end
```

# 1.4 how to solve problems

- 1. Мат. формулировка
- 2. Введение параметров
- 3. Обращение к теории
- 4. Способы решения:
  - (а) Фиксация параметра
  - (b) Рассмотреть частные случаи
  - (с) Рассмотреть ответ и его свойства
  - (d) Индукция
  - (е) Выдвижение гипотезы
  - (f) Поиск закономерностей
  - (g) Упрощение
  - (h) Предположение
  - (і) Рассмотрение
  - (j) Сведение к другой задаче
  - (k) Сведение к известной задаче
  - (1) Сведение известной задачи к моей
  - (m) Моделирование
  - (n) Обобщение
  - (о) Нахождние контрпримеров
  - (р) Зарандомить (полностью)
  - (q) Найти инвариант

# 2 Geometry

#### 2.1 structures

```
 \begin{array}{l} {\rm const}\ ld\ eps = 1e\text{-}6; \\ {\rm struct}\ Point\ \{ \\ ld\ x; \\ ld\ y; \\ Point\ operator+(Point\ B)\ \{ \\ return\ Point\{x+B.x,\ y+B.y\}; \\ \} \\ Point\ operator-(Point\ B)\ \{ \\ return\ Point\{x-B.x,\ y-B.y\}; \\ \} \\ Point\ operator^*(double\ t)\ \{ \\ return\ Point\{x^*\ t,\ y^*\ t\}; \end{array}
```

```
bool operator==(Point B) {
         return abs(x - B.x) <= eps && abs(y - B.y)<=eps;
     Point norm() {
         ld U = sqrt(x * x + y * y);
         return Point{x / U, y / U};
    ld len() {
         return sqrt(x * x + y * y);
};
struct Line{
  ll\ A;
  ll B;
  ll C;
};
struct Rec{
  Line L:
  Seg sg;
  Point V;
struct Seg {
    Point a;
    Point b:
    Point to vec() {
        return Point{b.x - a.x, b.y - a.y};
ld dot(Point p1, Point p2) {
    \operatorname{return}\ p1.x\ *\ p2.x\ +\ p1.y\ *\ p2.y;
ld dot(Point p1, Point p2, Point p3) {
    return dot(p2 - p1, p3 - p1);
ld cross(Point p1, Point p2) {
    return p1.x * p2.y - p2.x * p1.y;
ld cross(Point p1, Point p2, Point p3) {
    return cross(p2 - p1, p3 - p1);
```

# 2.2 segment intersection

```
 \begin{array}{l} {\rm optional < Point > intersect(Seg~s1,~Seg~s2)~\{} \\ {\rm if~(fabs(cross(s1.to\_vec(),~s2.to\_vec())) < eps)~\{} \\ {\rm return~\{\};} \\ {\rm \}} \\ {\rm double~t1 = cross((s2.a-s1.a),~s2.to\_vec()) / cross(s1.~to\_vec(),~s2.to\_vec());} \\ {\rm double~t2 = cross((s1.a-s2.a),~s1.to\_vec()) / cross(s2.~to\_vec(),~s1.to\_vec());} \\ {\rm if~(t1 < -eps~||~t1 > 1 + eps)~\{} \\ {\rm return~\{\};} \\ {\rm \}} \\ {\rm if~(t2 < -eps~||~t2 > 1 + eps)~\{} \\ {\rm return~\{\};} \\ {\rm \}} \\ {\rm return~s1.a + (s1.b-s1.a)~*~t1;} \\ \end{array}
```

#### 2.3 semiplane intersection

```
v3 = abs(v3);
     S = abs(S);
return abs(S - v1 - v2 - v3) \le eps;
Point find bis(Point A, Point B, Point C) {
      auto \vec{v} = \vec{A} - \vec{B};
      auto v2 = C - B;
      v.norm();
      v2.norm();
      \mathrm{return}\ \mathbf{v}\ +\ \mathbf{v2};
ld get_acos(double x) {
      \overline{\text{return acos}}(\max(-1., \min(1., x)));
Line make_line(Point a, Point b) {
      ll A = round(b.v - a.v);
      B = round(a.x - b.x);
      ll C = round(a.y * b.x - a.x * b.y);
      if (A < 0)
           \begin{array}{l} A \ *= (\text{-}1); \\ B \ *= (\text{-}1); \\ C \ *= (\text{-}1); \end{array}
      \hat{l}l g = gcd(A, gcd(abs(B), abs(C)));
     A /= g;
B /= g;
      C /= g;
      return Line{A, B, C};
Point intersect(Line a, Line b) {
    Id d = a.A * b.B - b.A * a.B;
    Id d1 = -a.C * b.B + a.B * b.C;
    Id d2 = -a.A * b.C + a.C * b.A;
      ass(abs(d) > eps);
Point res = {d1 / d, d2 / d};
      ass(on_line(res, a));
     ass(on_line(res, b));
return {d1 / d, d2 / d};
Point get_ort(Point p, Point p2) {
      return {p2.y - p.y, p.x - p2.x};
Point get_ort(Line a) {
      return {a.A, a.B};
const ld pi = get_acos(-1.0);
bool good(deque<Rec> &dq) {
      vector<Point> P;
       vector<Rec> v
      while (sz(dq)) {
            v.pb(dq.back());\\
           dq.pop_back();
      fi(1, sz(v) - 1) {
            P.pb(intersect(v[i-1].L,\,v[i].L));
      P.pb(intersect(v[0].L, v[sz(v) - 1].L));
      ld res = 0;
      P.pb(P[0]);
      ld s = 0, s2 = 0;
     \begin{array}{l} fi(1,\,sz(P)\,\text{-}\,1)\,\,\{\\ s\,+=\,P[i].y\,\,{}^*\,\,P[i\,\text{-}\,1].x;\\ s2\,+=\,P[i].x\,\,{}^*\,\,P[i\,\text{-}\,1].y; \end{array}
      res = abs(s - s2);
      return res > 0 + eps;
bool paral(Line a, Line b) {
return a.A * b.B - a.B * b.A == 0;
bool fun(ll T) {
      if (T == 0)
           return true;
      vector < Rec > vec(n + 1);
      set<tuple<ll, ll, ll>> t;
      fi(1, n) {
            Point p = a[i];
            Point p2 = a[go(i, T + 1, n)];
            vec[i] = Rec\{make\_line(p,\,p2),\,Seg\{p,\,p2\},\,get\_ort(p,\,p2)\}
                    p2)};
            if (area(p, p2, {p.x + 1'000'000 * vec[i].V.x, p.y + 1'000 '000 * vec[i].V.y}) > 0 + eps) { vec[i].V = vec[i].V * -1;
```

```
// vec[i].V.norm();
      // t.ins({vec[i].L.A,vec[i].L.B,vec[i].L.C});
// \operatorname{ass}(\operatorname{sz}(t) == n);
 vector<Rec> top, bot;
fi(1, n) {
     \label{eq:condition} \begin{array}{l} \text{if (vec[i].V.y} > 0 + \text{eps } || \text{ (abs(vec[i].V.y)} <= \text{eps \&\& vec[i].V.x} > 0 + \text{eps)) } \\ \end{array}
          bot.pb(vec[i]);
     } else {
          top.pb(vec[i]);
vec = vector < Rec > (1);
fz(bot) vec.pb(z);
reverse(All(top)):
fz(top) vec.pb(z);
deque<Rec> dq;
dbg(vec);
fi(1, n) {
     while (sz(dq) >= 2) {
          auto rec = dq.back();
          dq.pop_back();
auto rec2 = dq.back();
auto P = intersect(rec.L, rec2.L);
          dbg(P);
          if (area(vec[i].sg.a, vec[i].sg.b, P) <= eps) {
                dq.push_back(rec);
                break;
           \} else \{
                continue;
          }
     }
     \begin{array}{l} \mbox{while } (sz(dq) >= 2) \ \{ \\ \mbox{auto rec} = dq.front(); \end{array}
          dq.pop_front();
auto rec2 = dq.front();
           Point P = intersect(rec.L, rec2.L);
           dbg(P);
           dbg(area(vec[i].sg.a, vec[i].sg.b, P));
          if (area(vec[i].sg.a, vec[i].sg.b, P) \le eps) {
                dq.push_front(rec);
                break;
           } else {
                continue;
     dq.pb(vec[i]);
}
while (sz(dq) >= 3) {
     auto rec = dq.front();
     dq.pop_front();
     auto \ \overline{rec2} = dq.front();
     auto P = intersect(rec.L, rec2.L);
     dbg(P);
     if (area(dq.back().sg.a, dq.back().sg.b, P) <= eps) {
          dq.push_front(rec);
          break;
     \} else \{
          continue;
     }
}
while (sz(dq) >= 3) {
     auto rec = dq.back();
     dq.pop_back();
auto rec2 = dq.back();
     auto P = intersect(rec.L, rec2.L);
     if \; (area(dq.front().sg.a, \, dq.front().sg.b, \, P) <= \, eps) \; \{\\
          dq.push\_back(rec);
          break;
     } else {
          continue;
}
dbg(T);
dbg(sz(dq));
return sz(dq) > 2;
```

#### 2.4 Line and circle

We say first circle center at point (0,0). Line is (A,B,C). Find  $x_0=\frac{-AC}{A^2+B^2}$   $y_0=\frac{-BC}{A^2+B^2}$ 

If dist from  $(x_0, y_0)$  to (0, 0) is bigger than r then answer is zero. Otherwise if it equals to r then answer is  $(x_0, y_0)$ . Otherwise we have 2 solutions:  $d = \sqrt{r^2 - \frac{C^2}{A^2 + B^2}} \ V = \sqrt{\frac{d^2}{A^2 + B^2}} \ a_x = x_0 + BV$  $a_y = y_0 - AV \ b_x = x_0 - BV \ b_y = y_0 + AV$ 

#### 2.5 Two circles

We say first circle center at point (0,0).  $x^2 + y^2 = r_1^2$   $(x - x_2)^2 + (y - y_2)^2 = r_2^2$   $x^2 + y^2 = r_1^2 \ x(-2x_2) + y(-2y_2) + (x_2^2 + y_2^2 + r_1^2 + r_2^2) = 0$  Now we have new task to find intersection of line and circle: Ax + By + C = 0,  $A = -2x_2$ ,  $B = -2y_2$ ,  $C = x_2^2 + y_2^2 + r_1^2 - r_2^2$ 

#### 2.6 Two common tangets to circles

We have to find all tangent lines to two circles. First circle is at (0,0).

We say first circle center at point (0,0).  $x^2+y^2=r_1^2$   $(x-x_2)^2+(y-y_2)^2=r_2^2$   $x^2+y^2=r_1^2$   $x(-2x_2)+y(-2y_2)+(x_2^2+y_2^2+r_1^2+r_2^2)=0$  Now we have new task to find intersection of line and circle: Ax+By+C=0,  $A=-2x_2$ ,  $B=-2y_2$ ,  $C=x_2^2+y_2^2+r_1^2-r_2^2$ 

#### 3 Matrix

#### 3.1 Multiply

#### 3.2 Gauss SLAU

```
 \begin{array}{l} {\rm int~gauss~(vector < vector < double > > a,~vector < double > \&~ans} \\ {\rm int~n~=~(int)~a.size();} \\ {\rm int~m~=~(int)~a[0].size()~-1;} \\ {\rm vector < int > where~(m,~-1);} \\ {\rm for~(int~col=0,~row=0;~col < m~\&\&~row < n;~++col)~\{} \\ {\rm int~sel~=row;} \\ {\rm fi(row,~n~-1)~\{} \end{array}
```

```
if\ (abs\ (a[i][col]) > abs\ (a[sel][col])) \\
                  if (abs (a[sel][col]) < EPS)
                           continue;
                  fi(col, m) {
                           swap (a[sel][i], a[row][i]);
                  where [col] = row;
                  fi(0, n - 1) {
                           if (i != row) {
                                    double c = a[i][col] / a[row][col]
                                    fj(col, m) {
                                             a[i][j] -= a[row][j] * c;
                           }
                  ++row;
         ans.assign (m, 0);
         fi(0, m - 1) {
                 fi(0, n - 1) {
                  double sum = 0;
                  fj(0, m - 1) {
                           sum += ans[j] * a[i][j];
                  if (abs (sum - a[i][m]) > EPS)
          \begin{array}{c} fi(0,\,m\,\text{-}\,1)\;\{\\ &\text{if (where[i]==-1) return INF;} \end{array} 
         return 1;
}
```

# 4 Strings

# 4.1 Suf LCP

```
string s;
vi p, c, rc;
int x = 1:
int gett(int a, int x) {
if (a > x) return a - x;
     return sz(s) - 1 - (x - a) \% (sz(s) - 1);
int gett2(int a, int x) {
     if \; (a \, + \, x \, < \, sz(s)) \\
         return a + x:
     return 1 + (a + x) \% sz(s);
bool comp1(int a, int b) {
     return (s[a] < s[b]);
int n. S:
vvi e;
void Sort(vector<int> &p) {
     fi(1, n) {
          int t = gett(p[i], x);
          e[rc[t]].pb(t);
     p.clear();
     p.pb(0);
     fi(0, n) {
          fz(e[i]) {
               p.pb(z);
          e[i].clear();
     }
void init() {
    n = sz(s) + 1;

s = "\#" + s + (char)31;

c = vi(n + 1);
     rc = vi(n + 1);
     fi(0, n) p.pb(i);
     sort(Allf(p), comp1);
     c[1] = 0;
```

```
\begin{array}{l} fi(2,\,n)\;\{\\ if\;(s[p[i]]==s[p[i\text{ - }1]])\;c[i]=c[i\text{ - }1];\\ else\;c[i]=c[i\text{ - }1]+1;\\ . \end{array}
fi(1, n) {
      rc[p[i]] = c[i];
e = vvi(n + 1);
x = 1; while (x < n) {
      Sort(p);
c[1] = 0;
       fi(2, n) {
             if \ (rc[p[i]] == rc[p[i - 1]] \ \&\& \ rc[gett2(p[i], \, x)] == rc \\
                       [gett2(p[i - 1], x)]) {
                    c[i] = c[i-1];
             } else {
                    c[i] = c[i - 1] + 1;
             }
       fi(1, n) {
             \operatorname{rc}[p[i]] = c[i];
          *= 2;
```

#### 4.2 LCP

```
 \begin{array}{l} vll \ lcm(n \ + \ 1); \\ ll \ k \ = \ 0; \\ fl(1, \ n) \ \{ \\ \quad ll \ x \ = \ rev[i] \ - \ 1; \\ \quad ll \ cnt \ = \ max(k \ - \ 1, \ (ll)0); \\ fj(cnt, \ 1e5 \ + \ 10) \ \{ \\ \quad if \ (s[i \ + \ j] \ = \ s[arr[x] \ + \ j]) \ cnt++; \\ \quad else \ break; \\ \} \\ \quad if \ (x \ = \ 0) \ cnt \ = \ 0; \\ \quad lcm[x] \ = \ cnt; \\ k \ = \ cnt; \\ \} \end{array}
```

#### 4.3 Suf automaton

```
struct state {
            int len, link;
            map{<}char,\ int{>}\ next;
const int MAXLEN = 1'000'000;
state st[MAXLEN*2];
int siz. last:
void \ sa\_init() \ \{
           siz = last = 0;

st[0].len = 0;
            st[0].link = -1;
            siz++;
void sa extend (char c) {
            int cur = siz++
            st[cur].len = st[last].len + 1;
            for (p = last; p != -1 \&\& !st[p].next.count(c); p = st[p].
                     link) {
                        st[p].next[c] = cur;
            if (p == -1) {
                        st[cur].link = 0;
            \} \ {\rm else} \ \{
                        \label{eq:continuous_problem} \begin{split} &\inf \; q = st[p].next[c]; \\ &if \; (st[p].len \; + \; 1 = = st[q].len) \; \{ \end{split}
                                     st[cur].link = q;
                         \} else \{
                                     int clone = siz++;
                                     st[clone].len = st[p].len \, + \, 1; \\
                                     st[clone].len = st[p].len + 1;

st[clone].next = st[q].next;

st[clone].link = st[q].link;

for (; p != -1 && st[p].next[c] == q; p

= st[p].link)
                                                 st[p].next[c] = clone;
```

```
st[q].link = st[cur].link = clone; \\ \} \\ last = cur; \\ \}
```

# 4.4 pref function

```
vi prefix_function(string s) { 
	 int n = sz(s); 
	 vi pi(n); 
	 fi(1, n - 1) { 
	 int j = pi[i - 1]; 
	 while (j > 0 && s[i] != s[j]) { 
	 j = pi[j - 1]; 
	 } 
	 if (s[i] == s[j]) j++; 
	 pi[i] = j; 
	 } 
	 return pi; 
}
```

#### 4.5 Z function

```
 \begin{array}{l} vi \ z\_function \ (string \ s) \ \{ \\ int \ n = sz(s); \\ vi \ z \ (n); \\ for \ (int \ i = 1, l = 0, \ r = 0; \ i < n; \ i++) \ \{ \\ if \ (i < = r) \ z[i] = min(r - i + 1, z[i - l]); \\ while \ (i + z[i] < n \ \&\& \ s[z[i]] = = s[i + z[i]]) \ \{ \\ z[i] + +; \\ \} \\ if \ (i + z[i] - 1 > r) \ \{ \\ l = i; \\ r = i + z[i] - 1; \\ \} \\ return \ z; \\ \} \end{array}
```

# 4.6 Aho Corasick

```
struct Node {
         char c = '\0';
Node* par = nullptr;
map<char, Node*> nxt;
Node* suf = nullptr;
         map<char, Node*> sufs;
          Node(char t, Node* p) {
                   c = t;
                   par = p;
Node* root;
void add(string& s, int j) {
         Node* cur = root;

fi(0, sz(s)-1) {
                   if(cur->nxt.count(s[i])) {
                             cur = cur->nxt[s[i]];
                             Node* t = new Node(s[i],cur);
                             cur->nxt[s[i]] = t;
                             cur = t;
                   }
         cur->leaf = j;
Node* get_link(Node* a) {
         if(a->suf) return a->suf;
         if(a == root) { a->suf = root;
                   return a->suf;
          if(a->par==root) {
                   return root;
```

```
char t = a->c;
          Node* \operatorname{cur} = \operatorname{get\_link}(a->\operatorname{par});
while(\operatorname{cur} != \operatorname{root} \&\& !\operatorname{cur}>\operatorname{nxt.count}(t)) {
                     cur = get\_link(cur);
          if(cur->nxt.count(t)) {
                     cur = cur->nxt[t];
          return a->suf = cur;
Node* go(Node* a,char c) {
          Node* cur = a;
          while(cur != root \&\& !cur->nxt.count(c)) {
                     cur = get_link(cur);
          if(cur->nxt.count(c)) {
                     cur = cur->nxt[c];
           return cur;
}
string s;
ll n;
vector<string> a;
void init() {
          root = new Node('\0',nullptr);
          fi(1, n) {
                     add(a[i], i);
}
```

# 5 Graphs

# 5.1 cutpoints

# 5.2 bridges

```
 \begin{array}{l} \mbox{void } dfs(ll \; x, \; ll \; p = -1) \; \{ \\ \mbox{ if } (\mbox{color}[x]) \; \mbox{return}; \\ \mbox{color}[x] = 1; \\ \mbox{tin}[x] = \mbox{tim} + +; \\ \mbox{tup}[x] = \mbox{tin}[x]; \\ \mbox{for } (\mbox{auto } \& \mbox{to} : \mbox{e}[x]) \; \{ \\ \mbox{ if } (\mbox{to} = \mbox{p}) \; \mbox{continue}; \\ \mbox{if } (\mbox{color}[\mbox{to}]) \; \{ \\ \mbox{ } \mbox{tup}[x] = \mbox{min}(\mbox{tup}[x], \; \mbox{tin}[\mbox{to}]); \\ \mbox{continue}; \\ \mbox{ } \} \; \mbox{else } \{ \\ \mbox{ } \mbox{dfs}(\mbox{to}, \, x); \\ \end{array}
```

```
 \begin{array}{c} tup[x] = min(tup[x], \, tup[to]); \\ if \, (tup[to] > tin[x]) \, \{ \\ & ans.pb(t[\{to, \, x\}]); \\ & dbg(mp(to, \, x)); \\ \} \\ \} \\ \} \\ \end{array} \}
```

#### 5.3 min cost max flow

```
 \begin{array}{l} bool \; deikstr() \; \{ \\ vll \; d(\_k + 1, INF); \\ set < pll > t; \\ d[S] = 0; \\ t.ins(\{d[S], S\}); \\ vector < Edge *> p(\_k + 1, NULL); \\ while \; (sz(t)) \; \{ \\ \quad auto \; [\_, x] = *t.begin(); \\ \quad t.erase(t.begin()); \\ \quad fy(e[x]) \; \{ \\ \quad if \; (y->f+1 <= y->c \; \&\& \; d[y->y] > d[x] + y->w) \\ \quad \{ \\ \quad t.erase(\{d[y->y], y->y\}); \\ \quad d[y->y] = d[x] + y->w; \\ \quad t.ins(\{d[y->y], y->y\}); \\ \quad p[y->y] = y; \\ \} \\ \} \\ \} \\ if \; (d[T] == INF) \; return \; false; \\ int \; x = T; \\ while \; (x != S) \; \{ \\ \quad p[x]->f++; \\ \quad p[x]->rev->f-; \\ \quad x = p[x]->rev->y; \\ \} \\ fi(1, \_k) \; \{ \\ \quad y->w = y->w + d[i] - d[y->y]; \\ \} \\ \} \\ return \; true; \\ \} \\ \end{cases}
```

#### **5.4** flow

```
struct\ Edge\{
               int y;
               int f;
               int c;
               Edge* rev;
int _num[MAX];
int n, m;
\begin{array}{l} \mathrm{int}\ S,T;\\ \mathrm{int}\ K=0; \end{array}
int num(int x) {
               if(\underline{num[x]} == 0) \{
                               \_num[x] = K;
               return _num[x];
vector<vector<Edge*>>e;
vectod \ Vectod \ Edge \ >>e,
void add _ edge(int x, int y, int c) {
            Edge* v = new Edge{y, 0, c, NULL};
            Edge* v2 = new Edge{x, 0, c, v};
            v->rev = v2;
            e[x].pb(v);
               e[y].pb(v2);
 vi color;
int C = 1;
bool dfs(int x) {
    if(x == T) return true;
    if(color[x]) return false;
    color[x] = 1;
                fy(e[x]) {
                               if(y->f + C \le y->c \&\& dfs(y->y))  {
```

```
\begin{array}{c} y\text{-}\!>\!f+=C;\\ y\text{-}\!>\!rev\text{-}\!>\!f-=C;\\ return\ true;\\ \}\\ \}\\ return\ false;\\ \}\\ \\ \text{while}(C>=1)\{\\ C=(1<<30);\\ \text{while}(C>=1)\{\\ color=vi(n+1);\\ \text{while}(dfs(S))\}\\ \\ color=vi(n+1);\\ \}\\ C/=2;\\ \}\\ \end{array}
```

# 5.5 Kuhn algorithm

#### 5.6 Dinic flow

```
const int MAXN = ...:
struct edge \{
               int a, b, cap, flow;
int n, s, t, d[MAXN], ptr[MAXN], q[MAXN];
vector<edge> e;
vector<int> g[MAXN];
void add_edge (int a, int b, int cap) { edge e1 = { a, b, cap, 0 }; edge e2 = { b, a, 0, 0 }; ...()
               g[a].push_back ((int) e.size());
               e.push_back (e1);
               g[b].\underline{push}\_\underline{back}\ ((int)\ e.size());
               e.push_back (e2);
}
bool bfs() {
               int qh=0, qt=0;
               q[qt++] = s;

memset (d, -1, n * sizeof d[0]);
              \begin{array}{l} d[s] = 0; \\ while \; (qh < qt \; \&\& \; d[t] == \text{-1}) \; \{ \\ & \text{int } v = q[qh++]; \\ & \text{for } (\text{size\_t } i =\! 0; \, i \! <\! g[v]. \text{size}(); \; +\! +\! i) \; \{ \\ & \text{int } id = g[v][i], \\ & \text{to } = e[id]. b; \\ & \text{if } (d[to] == \text{-1} \; \&\& \; e[id]. \text{flow} < e[id]. \text{cap} \\ & ) \; \{ \end{array}
               d[s] = 0:
                                                               egin{aligned} {
m q}[{
m q}t{++}] &= {
m to}; \ {
m d}[{
m to}] &= {
m d}[{
m v}] \,+\,1; \end{aligned}
                               }
               return d[t] != -1;
}
int dfs (int v, int flow) \{
               if (!flow) return 0;
```

```
\begin{array}{c} \mbox{if } (v==t) \mbox{ return flow;} \\ \mbox{for } (; ptr[v] < (int)g[v].size(); ++ptr[v]) \; \{\\ \mbox{int } id = g[v][ptr[v]], \\ \mbox{to } = e[id].b; \\ \mbox{if } (d[to] != d[v] + 1) \mbox{ continue;} \\ \mbox{int } pushed = dfs \; (to, min \; (flow, e[id].cap - e[id]. \\ \mbox{flow})); \\ \mbox{if } (pushed) \; \{\\ \mbox{} e[id].flow += pushed; \\ \mbox{} e[id^-1].flow -= pushed; \\ \mbox{} return \; pushed; \\ \mbox{} return \; pushed; \\ \mbox{} \}\\ \mbox{} \\ \mb
```

# 5.7 dijkstra

```
 \begin{array}{l} vector < int > dijkstra(int\ s) \{ \\ vector < int > d(n+1,\ INF); \\ d[s] = 0; \\ set < pair < int,\ int > > t; \\ t.insert(mp(d[s],\ s)); \\ vector < int > color(n+1,\ 0); \\ while(sz(t)) \left \{ \\ auto[\_,\ x] = *t.begin(); \\ t.erase(t.begin()); \\ if(color[x]) \ continue; \\ color[x] = true; \\ for(auto\ \&[y,\ w]: e[x]) \left \{ \\ if(d[y] > d[x] + w) \left \{ \\ d[y] = d[x] + w; \\ t.insert(mp(d[y],\ y)); \right \} \\ \} \\ return\ d; \\ \} \\ \} \\ return\ d; \\ \\ \end{array}
```

#### 6 Data Structures

# 6.1 fenwick Tree

```
// Linear
vll fen;
void add(int p, ll val) {
      for(int i = p; i \le n; i = (i | (i + 1))) {
            fen[i] += val;
ll sum(int p) {
      ll res = 0;
      for(int i = p; i >= 0; i = (i & (i + 1)) - 1) {
            res += fen[i];
      return res;
}
// Matrix
void add(int x, int y, ll val) {
      \begin{array}{l} \text{for}(\inf \ i = x; \ i <= n; \ i = (i \mid (i+1))) \ \{\\ \text{for}(\inf \ j = y; \ j <= n; \ j = (j \mid (j+1))) \ \{\\ \text{fen}[i][j] \ += val; \end{array}
     }
}
```

```
 \begin{split} & \text{ll sum(int } x, \text{ int } y) \text{ } \{ \\ & \text{ll res} = 0; \\ & \text{for(int } i = x; i >= 0; i = (i \& (i+1)) - 1) \text{ } \{ \\ & \text{ for (int } j = y; j >= 0; j = (j \& (j+1)) - 1) \text{ } \{ \\ & \text{ res} += \text{ fen[i][j]; } \\ & \text{ } \} \\ & \text{ return res; } \} \end{aligned}
```

# 6.2 treap

```
struct Node {
      ll val;
      ll cnt;
     ll y;
Node *l;
      Node *r;
void upd(Node *a) {
      a->cnt = 1;
      if (a->l)
           a->cnt += a->l->cnt;
      if (a->r)
           a->cnt += a->r->cnt;
Node *merge(Node *a, Node *b) {
     if (!a)
           return b:
      if (!b)
           return a;
     if (a->y > b->y) {
 a->r = merge(a->r, b);
            upd(a);
            return a;
      b->l = merge(a, b->l);
      upd(b);
      return b;
pair<Node *, Node *> split(Node *a, ll q) {
     if (!a)
           return \{0, 0\};
      ll left = 0;
      if (a->l)
            \mathbf{left}^{'} += \mathbf{a}\text{-}\mathbf{>}\mathbf{l}\text{-}\mathbf{>}\mathbf{cnt};
      \begin{array}{l} \mathrm{if}\ (\mathrm{left}>=\mathrm{q})\ \{\\ \mathrm{auto}\ [\mathrm{l},\,\mathrm{r}]=\mathrm{split}(\mathrm{a}\text{-}\mathrm{>}\mathrm{l},\,\mathrm{q}); \end{array}
           a\text{-}{>}l=r;
            upd(a);
            return {l, a};
       auto [l, r] = split(a->r, q - left - 1); 
      a->r=1;
      upd(a);
      return \{a, r\};
}
```

# 6.3 Long

```
struct Long {
    int base = 10000;
    int len = 4;
    vi num = \{1, 0\};
    int operator[](int x) const {
        if (x > num[0]) return 0;
        return num[x];
    int& operator[](int x) {
        if (x > num[0]) {
            num.resize(x + 1, 0);
            num[0] = x;
        return num[x];
   }
Long operator+(const Long& A, const Long& B) {
    Long res;
    int n = \max(A[0], B[0]);
```

```
 \begin{array}{l} {\rm fi}(1,\,n)\;\{\\ {\rm res}[i]=A[i]+B[i];\\ {\rm if}\;(i>1)\;\{\\ {\rm res}[i]+={\rm res}[i-1]\;/\;{\rm res.base};\\ {\rm res}[i-1]\;\%={\rm res.base};\\ \end{array} 
              \text{while}(\text{res}[\text{res}[0]] >= \text{res.base})  {
                           int m = res[0];

res[m + 1] = res[m] / res.base;
                           res[m] %= res.base;
              return res;
\underline{\text{Long operator*}(\text{const Long\& A, const int\& B)}} \ \{
              Long res;
              int n = A[0];
              fi(1, n) {
                           res[i] = A[i] * B;
                            \begin{array}{l} \mbox{if } (i>1) \ \{ \\ \mbox{res}[i] \ += \mbox{res}[i-1] \ / \mbox{res.base}; \\ \mbox{res}[i-1] \ \%= \mbox{res.base}; \\ \end{array} 
               \text{while}(\text{res}[\text{res}[0]] > = \text{res.base})  {
                           int m = res[0];

res[m + 1] = res[m] / res.base;
                           res[m] %= res.base;
             return res;
Long operator*(const Long& A, const Long& B) {
              Long res
             fi(1, A[0]) {

if (A[i] == 0) continue;
                           \begin{split} & \text{if } (A[i] == 0) \text{ continue;} \\ & \text{fj}(1, B[0]) \ \{ \\ & \text{if } (A[i] * B[j] == 0) \text{ continue;} \\ & \text{int } x = i+j-1; \\ & \text{res}[x] += A[i] * B[j]; \\ & \text{while}(x > 1 \text{ \&& res}[x-1] >= \text{res.base)} \ \{ \\ & \text{res}[x] += \text{res}[x-1] \ / \text{ res.base;} \\ & \text{res}[x-1] \ \% = \text{res.base;} \\ & \text{ } \} \end{split}
                            \text{while}(\text{res}[\text{res}[0]] >= \text{res.base})  {
                                         \begin{array}{ll} \operatorname{int} & \operatorname{res}[0];\\ \operatorname{res}[m+1] + = \operatorname{res}[m] \ / \ \operatorname{res.base};\\ \operatorname{res}[m] \ \% = \operatorname{res.base}; \end{array}
                           }
              return res;
pair{<}Long,\,int{>}\,\,operator/(const\,\,Long\&\,\,A,\,const\,\,int\,\,B)\,\,\{
             Long res;
              ll x = 1;
              ll d = 0, m = 0;
              fdi(A[0], 1) {
                          A[0], A[0]
                           m = val % B;
                           if (x == 1 \&\& d == 0) continue;
                           res[x] = d;
              reverse(Allf(res.num));
             return mp(res, m);
void print(const Long& A) {
             printf("%d", A[A[0]]);
fdi(A[0] - 1, 1) {
    printf("%04d", A[i]);
             printf(ln);
Long str_to_Long(const string& s) {
              int n = \overline{sz}(s);
              Long res;
              int x = 1:
              for(int i = n - 1; i >= 0; i -= 4) {
                            int val = 0;
                            for(int j = \max(0, i - 3); j \le i; j++) {
                                         val = val * 10 + (s[j] - '0');
                            res[x] = val;
                           x++;
              return res;
```

```
Long int_to_Long(int num) {
    Long res;
    int x = 1;
    while(num) {
        int val = num % res.base;
        num /= res.base;
        res[x] = val;
        x++;
    }
    return res;
}
```

## 6.4 persistent treap

```
int next(int l,int r){
             return l + rnd()\%(r - l + 1);
const int MAX = 1e7;
int siz = 1;
vi\ \operatorname{Left}(\operatorname{MAX},\!-1), \operatorname{Right}(\operatorname{MAX},\!-1), \operatorname{Siz}(\operatorname{MAX},\!0), \operatorname{Val}(\operatorname{MAX},\!0);
\begin{array}{c} \mathrm{void} \ \mathrm{upd}(\mathrm{int} \ v) \{ \\ \mathrm{if}(v == \text{-}1) \mathrm{return}; \end{array}
             if(Left[v] != -1)Siz[v] += Siz[Left[v]];
             if(Right[v] != -1)Siz[v] += Siz[Right[v]];
}
bool INIT;
int cop(int v){
             if(INIT)return v;
             int x = siz;
             siz++;
             \begin{array}{l} \operatorname{Siz} + +; \\ \operatorname{Val}[x] = \operatorname{Val}[v]; \\ \operatorname{Left}[x] = \operatorname{Left}[v]; \\ \operatorname{Right}[x] = \operatorname{Right}[v]; \\ \operatorname{Siz}[x] = \operatorname{Siz}[v]; \end{array}
             return x;
}
\begin{array}{c} \mathrm{int\ merg(int\ L,int\ R)\{} \\ \mathrm{if(L==-1)return\ R;} \end{array}
             if(R == -1)return L;
             if(next(\ 1,\ (Siz[L] + Siz[R])*100\ ) <= Siz[L]\ *\ 100)\{
                           L = cop(L);
                           Right[\hat{L}] = merg(Right[L], R);
                           upd(L);
                           return L;
             }else{
                           R = cop(R);
                           Left[R] = merg(L, Left[R]);
                           upd(R);
                           return R:
             }
pair < int, int > split(int v, int k){
             if(Siz[v] <= k)\{
                           return mp(v,-1);
             if(k == 0){
                           return mp(-1,v);
              v = cop(v);
             int siz = 0:
             if(Left[v] \stackrel{\cdot}{!=} -1)siz = Siz[Left[v]];
             if(siz >= k){
                           auto t = \text{split}(\text{Left}[v],k);
                           Left[v] = t.second;
                           upd(v);
                           return mp(t.first,v);
             }else{
                           auto t = \text{split}(\text{Right}[v], k - \text{siz} - 1);
                           Right[v] = t.first;
                           return mp(v,t.second);
             }
}
int root = -1;
void \ seg\_copy(int \ i,int \ j,int \ k)\{
```

```
 \begin{array}{c} \mathrm{auto}\ t = \mathrm{split}(\mathrm{root}, j-1); \\ \mathrm{auto}\ g = \mathrm{split}(\mathrm{root}, i-1); \\ \mathrm{auto}\ h = \mathrm{split}(\mathrm{g.second}, k); \\ \mathrm{auto}\ f = \mathrm{split}(\mathrm{t.second}, k); \\ \mathrm{root} = \mathrm{merg}(\mathrm{t.first}, \mathrm{merg}(\mathrm{h.first}, \mathrm{f.second})); \\ \\ \\ \mathrm{int}\ n, m; \\ \mathrm{vi}\ \mathrm{cnt}, \mathrm{from}, \mathrm{to}; \\ \mathrm{vi}\ \mathrm{cnt}, \mathrm{from}, \mathrm{to}; \\ \mathrm{vi}\ \mathrm{res}; \\ \\ \mathrm{void}\ \mathrm{solve}()\ \{ \\ \mathrm{INIT} = \mathrm{true}; \\ \mathrm{fi}(1, \, n) \{ \\ \mathrm{int}\ v = \mathrm{siz} + +; \\ \mathrm{upd}(v); \\ \mathrm{Val}[v] = \mathrm{i}; \\ \\ \mathrm{root} = \mathrm{merg}(\mathrm{root}, v); \\ \\ \} \\ \mathrm{INIT} = \mathrm{false}; \\ \} \\ \end{array}
```

## 7 DP

# 7.1 fft

```
using\ cd = complex{<}double{>};
const double PI = acos(-1);
int reverse(int num, int lg_n) {
      int res=0;
      for (int i = 0; i < lg_n; i++) {
    if (num & (1 << i))
        res |= 1 << (lg_n - 1 - i);
    }
      return res:
void fft(vector<cd> & a, bool invert) {
      int\ n=a.size();
     \begin{array}{l} \mathrm{int}\ \mathrm{lg}\underline{\phantom{}} n = 0;\\ \mathrm{while}\ \overline{((1 << \mathrm{lg}\underline{\phantom{}} n) < n)} \end{array}
           lg_n++;
      for (int i = 0; i < n; i++) {
            if (i < reverse(i, lg_n))
                  swap(a[i], a[reverse(i, lg_n)]);
      for (int len = 2; len <= n; len <<= 1) { double ang = 2 * PI / len * (invert ? -1 : 1);
            cd wlen(cos(ang), sin(ang));
            for (int i = 0; i < n; i += len) {
                  cd w(1);
                   for (int j = 0; j < \text{len } / 2; j++) {
                        cd u = a[i+j], v = a[i+j+len/2] * w;

a[i+j] = u + v;
                         \begin{aligned} a[i+j+len/2] &= u - v; \\ w &*= wlen; \end{aligned} 
                  }
            }
      }
      if (invert) {
            for (cd & x : a) {
                 x /= n;
      }
vector<int> multiply(vector<int>const& a, vector<int>const&
         b) {
      vector<cd> fa(a.begin(), a.end()), fb(b.begin(), b.end());
     \begin{array}{l} int \ n=1; \\ while \ (n < sz(a) + sz(b) \ ) \ \{ \end{array}
           n << = 1:
      fa.resize(n);
      fb.resize(n);
      fft(fa, false);
fft(fb, false);
      for (int i = 0; i < n; i++)
fa[i] *= fb[i];
      fft(fa, true);
```

```
\label{eq:vector} \begin{array}{l} \operatorname{vector} < \operatorname{int} > \operatorname{result}(n); \\ \operatorname{for} \ (\operatorname{int} \ i = 0; \ i < n; \ i++) \\ \operatorname{result}[i] = \operatorname{round}(\operatorname{fa}[i].\operatorname{real}()); \\ \operatorname{for} \ (\operatorname{int} \ i = 0; \ i < n; \ i++) \{ \\ \operatorname{result}[i] = \operatorname{min}(\operatorname{result}[i], 5); \\ \} \\ \operatorname{while}(\operatorname{result}.\operatorname{back}() == 0) \operatorname{result}.\operatorname{pop\_back}(); \\ \operatorname{return} \ \operatorname{result}; \\ \} \end{array}
```

#### 7.2 convex hull trick

```
 \label{eq:struct_line} $$ struct\ Line\ \{$ mutable\ ll\ k,\ m,\ p;$ bool\ operator<(const\ Line\ \&o)\ const\ \{\ return\ k< o.k;\ \}$ 
      bool operator < (ll x) const { return p < x; }
};
struct\ LineContainer: multiset < Line,\ less <>>> \{
      // (for doubles, use \inf = 1/.0, \operatorname{div}(a,b) = a/b) static const ll \inf = INF;
     ll div(ll a, ll b) { // floored division return a / b - ((a \hat{} b) < 0 && a % b);
      bool isect(iterator x, iterator y) {
            if (y == end())
            return x->p = inf, 0;
if (x->k == y->k)
x->p = x->m > y->m ? inf : -inf;
                 x->p = div(y->m - x->m, x->k - y->k);
            \operatorname{return} \ x\text{-}{>}p>=\ y\text{-}{>}p;
      void add(ll k, ll m) {
            auto z = insert(\{k, m, 0\}), y = z++, x = y;
            while (isect(y, z))
                 z = erase(z);
            if (x != begin() \&\& isect(--x, y))
            isect(x, y = erase(y));
while ((y = x) != begin() \&\& (--x)->p >= y->p)
                  isect(x, erase(y));
      ll query(ll x) {
           assert(!empty());
auto l = *lower_bound(x);
return l.k * x + l.m;
};
```

#### 8 Comb

# 8.1 basic

```
Cat(n,k) = \frac{1}{n+1}C(2n,n)
N(n,k) = \frac{1}{n}C(n,k)C(n,k-1)
\sum_{n=0}^{\inf} \sum_{k=1}^{n} N(n,k)z^{n}t^{k} = \frac{1+z(t-1)-\sqrt{1-2z(t+1)+z^{2}(t-1)^{2}}}{2z}
S(n,k) = \sum_{j=0}^{k-1} (-1)^{j}C(k,j) \cdot (m-j)^{n} O(k\log n)
```

#### 9 Math

#### 9.1 Simpson

```
double a, b; const int n = 1000*1000; double s = 0; double h = (b - a) / N; fi(0, n) { double x = a + h * i; s += f(x) * ((i=0 || i==N) ? 1 : ((i&1)==0) ? 2 : 4);
```

```
\label{eq:special} \begin{array}{l} s \\ s \end{array} *= h \ / \ 3;
```

#### 9.2 2 SAT

```
ll n;
vvi e;
vi color;
vvi e2;
vvi num; int id = 1;
\begin{array}{c} \mathrm{int} \ \mathrm{num}(\mathrm{int} \ v, \mathrm{int} \ t) \{ \\ \mathrm{if}(!\_\mathrm{num}[v][t] \ )\_\mathrm{num}[v][t] = \mathrm{id} + +; \end{array}
            return _num[v][t];
vvi e2r;
vi u;
vi ord;
void dfs2(int v){
            if(u[v])return;
            u[v] = 1;
            \begin{array}{c} u_{\lfloor v_{\rfloor} - - \rfloor} \\ fz(e2[v]) \{ \\ dfs2(z); \end{array}
            ord.pb(v);
int comp = 0;
vi component;
void dfs3(int v){
            if(component[v])return;
component[v] = comp;
            fz(e2r[v]){
                        dfs3(z);
vi ans:
bool can(){
            ord.clear();
            ans.clear();
            e2r = vvi(4*n + 1);
            comp = 1;
            fi(1,4*n){
                         \begin{array}{c} )\{\\ \mathrm{fz}(\mathrm{e2[i]})\{\\ \mathrm{e2r[z].pb(i)}; \end{array} 
            u = vi(4*n + 1);
           fi(1,4*n){ dfs2(i);
            dbg("S");
            component = vi(4*n + 1);
            fdi(sz(ord) -1,0){
                        if(!component[ord[i]]){
                                    dfs3(ord[i]);
                                     comp++;
            dbg("D");
            fi(1,2*n){
                        if(component[num(i,1)] == component[num(i,2)]
                        if(component[num(i,1)] > component[num(i,2)])
                                 ans.pb(i);
            return true;
}
```

#### 9.3 gcd

```
\begin{array}{l} \mbox{int gcd (int a, int b, int \& x, int \& y) \{} \\ \mbox{if } (a == 0) \; \{ \\ \mbox{$x = 0$; $y = 1$;} \\ \mbox{return b;} \\ \mbox{\}} \\ \mbox{int $x1$, $y1$;} \\ \mbox{int $d = \gcd(b\%a, a, $x1$, $y1$);} \end{array}
```

```
 \begin{array}{c} x = \, y1 \, \cdot \, (b \, / \, a) \, * \, x1; \\ y = \, x1; \\ return \, d; \\ \} \end{array}
```

#### 9.4 Diofant

```
void shift_solution (int & x, int & y, int a, int b, int cnt) {
         x + = cnt * b;
         y -= cnt * a;
}
int find_all_solutions (int a, int b, int c, int minx, int maxx,
      int miny, int maxy) {
         int x, y, g;
         if (! find_any_solution (a, b, c, x, y, g))
                 return 0;
         \mathbf{a} /= g; b /= g;
         \begin{array}{l} {\rm int~sign\_a = a {>} 0~?~+1:-1;} \\ {\rm int~sign\_b = b {>} 0~?~+1:-1;} \end{array}
         shift\_solution (x, y, a, b, (minx - x) / b);
         if (x < minx)
                  shift_solution (x, y, a, b, sign_b);
         if\ (x>maxx)
                  return 0;
         int\ lx1=x;
         shift\_solution~(x,~y,~a,~b,~(maxx~-~x)~/~b);
         if (x > maxx)
                 shift_solution (x, y, a, b, -sign_b);
         int rx1 = x;
         shift\_solution (x, y, a, b, - (miny - y) / a);
         if (y < miny)
                 shift_solution (x, y, a, b, -sign_a);
         if (y > maxy)
                 return 0;
         int\ lx2=x;
         shift\_solution~(x,~y,~a,~b,~-~(maxy~-~y)~/~a);
         if (y > maxy)
                 shift_solution (x, y, a, b, sign_a);
         int rx2 = x;
         if\ (lx2>rx2)
                 swap (lx2, rx2);
         int\ lx=max\ (lx1,\ lx2);
         int rx = min (rx1, rx2);
         return (rx - lx) / abs(b) + 1;
}
```

#### 9.5 chinese theorem

$$\sum_{i=1}^{n} a_i \left(\frac{M}{m_i}\right)^{\varphi(m_i)}$$

# 9.6 joseph

```
\begin{array}{l} \text{int joseph (int } n, \text{ int } k) \; \{ \\ & \text{ if } (n = = 1) \text{ return } 0; \\ & \text{ if } (k = = 1) \text{ return } n\text{--}1; \\ & \text{ if } (k > n) \text{ return (joseph (n-1, k) + k) } \% \; n; \\ & \text{ int cnt } = \; n \; / \; k; \\ & \text{ int res } = \text{ joseph (n - cnt, k)}; \\ & \text{ res -= n } \% \; k; \\ & \text{ if } (\text{res } < 0) \text{ res } + = \; n; \\ & \text{ else res } + = \text{ res } / \; (k \text{ - 1}); \\ & \text{ return res;} \\ \} \end{array}
```

# 9.7 discret log

```
int solve (int a, int b, int m) {
    int n = (int) sqrt (m + .0) + 1;
    int cur;
    int an = 1;
    fi(0, n - 1) {
        an = (an * a) % m;
    }
    map<int,int> vals;

    cur = an;
    fi(1, n) {
        if (!vals.count(cur))
            vals[cur] = i;
        cur = (cur * an) % m;
    }
    cur = b;
    fi(0, n) {
        if (vals.count(cur)) {
            if (ans < m) return ans;
        }
        cur = (cur * a) % m;
    }
    return -1;
}</pre>
```