# ACM-ICPC TEAM REFERENCE DOCUMENT

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## 1 Common

#### 1.1 template

optional, complex, string, vector, list, map, set, deque, queue, stack, bitset, algorithm, sstream, iostream, iomanip, cstdio, cstdlib, ctime, cstring, cmath, cstdarg, cassert, ctime, tuple, unordered\_set, unordered\_map, random, chrono using namespace std;

```
#define ln "\n"
#define pb push_back
#define mp make_pair
#define ins insert
#define sz(x) (int)x.size()
typedef long long ll;
typedef vector<int> vi;
typedef vector<vi> vvi;
typedef pair<int, int> pii;
template<typename T>
ostream& operator<<(ostream& os, vector<T> v) {
      fz(v) {
             os << z << "_{\sqcup}";
      return os:
#ifdef LOCAL
      _LINE___ << "\t" << #x
      #define dbg(x) {}
#endif
#ifdef LOCAL
   #define ass(x) if (!(x)) { cerr << __LINE__ << "\ tassertion_failed:_{\sqcup}" << #x << ln; abort(); }
   #define ass(x) assert(x)
mt19937\_64\ rnd(chrono::steady\_clock::now().time\_since\_epoch
     ().count());
int main()
   ios::sync_with_stdio(0);
   cin.tie(0);
   cout.tie(0)
   #ifdef LOCAL
      freopen("input.txt", "r", stdin);
      freopen("output.txt", "w", stdout);
   #endif
```

#### 1.2 compile

g++ --std=c++20 !.! -DLOCAL -O2 -Wall -Wl,--stack =67108864 -o !.exe

#### 1.3 check

```
:start
gen.exe
sol.exe
copy output.txt answer.txt
NAME.exe
fc output.txt answer.txt
if ERRORLEVEL 1 goto end
goto start
:end
```

#### 1.4 how To Solve Problems

```
1.
2.
3.
4.
     (a)
     (b)
     (c)
     (d)
     (e)
      (f)
     (g)
     (h)
      (i)
      (j)
     (k)
      (1)
    (m)
     (n)
     (o)
                     (
                              )
     (p)
     (q)
```

## 2 Geometry

#### 2.1 structures

```
const ld eps = 1e-6;
struct Point {
    ld x;
    ld v:
    Point operator+(Point B) {
         return Point\{x + B.x, y + B.y\};
     Point operator-(Point B) {
         return\ Point\{x - B.x,\ y - B.y\};
    Point operator*(double t) {
    return Point{x * t, y * t};
     bool operator==(Point B) {
         return abs(x - B.x) \le eps && abs(y - B.y) \le eps;
     Point norm() {
         \begin{array}{l} \operatorname{Id} U = \operatorname{sqrt}(x * x + y * y); \\ \operatorname{return} \operatorname{Point}\{x \ / \ U, \ y \ / \ U\}; \end{array}
    ld len() {
         return sqrt(x * x + y * y);
};
struct Line{
  ll A;
  11 B:
  ll C;
};
struct Rec{
  Line L;
  Seg sg;
Point V;
struct Seg {
    Point a;
    Point b:
    Point to_vec() {
         return Point{b.x - a.x, b.y - a.y};
ld dot(Point p1, Point p2) {
    return p1.x * p2.x + p1.y * p2.y;
```

```
Id dot(Point p1, Point p2, Point p3) {
    return dot(p2 - p1, p3 - p1);
}
ld cross(Point p1, Point p2) {
    return p1.x * p2.y - p2.x * p1.y;
}
ld cross(Point p1, Point p2, Point p3) {
    return cross(p2 - p1, p3 - p1);
}
```

#### 2.2 segment Intersection

```
 \begin{array}{l} {\rm optional} < {\rm Point} > {\rm intersect}({\rm Seg~s1,~Seg~s2})~\{ \\ {\rm if~(fabs(cross(s1.to\_vec(),~s2.to\_vec()))} < {\rm eps})~\{ \\ {\rm return~}\{\}; \\ {\rm }\} \\ {\rm double~t1 = cross((s2.a~-s1.a),~s2.to\_vec())} / {\rm cross(s1.co\_vec(),~s2.to\_vec())}; \\ {\rm double~t2 = cross((s1.a~-s2.a),~s1.to\_vec())} / {\rm cross(s2.co\_vec(),~s1.to\_vec())}; \\ {\rm if~(t1 < -eps~||~t1 > 1 + eps)~\{ \\ {\rm return~}\{\}; \\ {\rm }\} \\ {\rm if~(t2 < -eps~||~t2 > 1 + eps)~\{ \\ {\rm return~}\{\}; \\ {\rm }\} \\ {\rm return~s1.a~+~(s1.b~-s1.a)~*~t1; } \\ \end{array}
```

#### 2.3 semiplane Intersection

```
ld area
(Point a, Point b, Point c) { return (ld)
1.0 * ((a.x - b.x) * (a.y - c.y) - (a.x - c.x) * (a.y -
             b.y));
bool inside(Point A, Point B, Point C, Point p) {
     if (area(A, B, p) == 0 || area(B, C, p) == 0 || area(A, C, p)
              == 0)
          return false;
     ld v1 = area(A, B, p);
ld v2 = area(B, C, p);
ld v3 = area(A, C, p);
     \operatorname{ld} S = \operatorname{area}(A, B, C);
     v1 = abs(v1);
     v2 = abs(v2);
     v3 = abs(v3);
     S = abs(\hat{S}):
     return abs(S - v1 - v2 - v3) \leq eps;
Point find_bis(Point A, Point B, Point C) {
     auto v = A - B;
auto v2 = C - B;
     v.norm():
     v2.norm():
     return v + v2;
ld get\_acos(double x) {
     return acos(max(-1., min(1., x)));
Line make_line(Point a, Point b) {
     ll A = round(b.y - a.y);
     \begin{array}{l} \text{ll } B = \text{round}(a.x - b.x); \\ \text{ll } C = \text{round}(a.y * b.x - a.x * b.y); \end{array}
     if (A < 0)
     {
          B *= (-1);

C *= (-1);
     ll g = gcd(A, gcd(abs(B), abs(C)));
     A /= g;
     B /= g;
     return Line{A, B, C};
Foint intersect(Line a, Line b) {
    Id d = a.A * b.B - b.A * a.B;
    Id d1 = -a.C * b.B + a.B * b.C;
    Id d2 = -a.A * b.C + a.C * b.A;
     ass(abs(d) > eps);
     Point res = \{d1 / d, d2 / d\};
```

```
ass(on_line(res, a));
          ass(on_line(res, b));
return {d1 / d, d2 / d};
Point get_ort(Point p, Point p2) {
          return {p2.y - p.y, p.x - p2.x};
Point get_ort(Line a) {
          return \{a.A, a.B\};
const ld pi = get \ acos(-1.0);
bool good(deque<Rec> &dq) {
           vector < Point > P;
          vector<Rec> v;
          while (sz(dq)) { v.pb(dq.back());
                     dq.pop_back();
          fi(1, sz(v) - 1) {
                     P.pb(intersect(v[i - 1].L, v[i].L));
          P.pb(intersect(v[0].L, \, v[sz(v) \, \text{-} \, 1].L));
          P.pb(P[0]);
          \begin{aligned} & \text{I.pol} \ [\text{v}], \\ & \text{ld } s = 0, \, \text{s2} = 0; \\ & \text{fi}(1, \, \text{sz}(P) - 1) \, \{ \\ & \text{s} + \text{P}[\text{i}].\text{y} * \text{P}[\text{i} - 1].\text{x}; \\ & \text{s2} + \text{P}[\text{i}].\text{x} * \text{P}[\text{i} - 1].\text{y}; \end{aligned}
          res = abs(s - s2);
          return res > 0 + eps;
bool paral(Line a, Line b) {
return a.A * b.B - a.B * b.A == 0;
bool fun(ll T) {
          if (T = 0)
          return true;
vector<Rec> vec(n + 1);
          set<tuple<ll, ll, ll>> t;
          fi(1, n) {
                     Point p = a[i];
                     Form p = a_{[i]}, Point p = a_{[i]}, Point p = a_{[i]}, p 
                                     p2)};
                     \begin{array}{l} \text{if } (\operatorname{area}(p,p2, \{p.x+1'000'000*\operatorname{vec}[i].V.x, \ p.y+1'000'\\ 000*\operatorname{vec}[i].V.y\}) > 0+\operatorname{eps}) \ \{\\ \operatorname{vec}[i].V = \operatorname{vec}[i].V*-1; \end{array}
                      // vec[i].V.norm();
                     // t.ins({vec[i].L.A, vec[i].L.B, vec[i].L.C});
           // \operatorname{ass}(\operatorname{sz}(t) == n);
           vector < Rec > top, bot;
          fi(1, n) {
                      \begin{array}{c} \text{if (vec[i].V.y} > 0 + \mathrm{eps} \mid\mid (\mathrm{abs(vec[i].V.y}) <= \mathrm{eps} \&\& \\ \mathrm{vec[i].V.x} > 0 + \mathrm{eps})) \; \{ \end{array} 
                               bot.pb(vec[i]);
                     } else {
                               top.pb(vec[i]);
                     }
          vec = vector < Rec > (1);
          fz(bot) vec.pb(z);
reverse(All(top));
          \begin{array}{l} fz(top) \ vec.pb(z); \\ deque < Rec > \ dq; \end{array}
          dbg(vec);
                      while (sz(dq) >= 2) {
                               auto rec = dq.back();
dq.pop_back();
auto rec2 = dq.back();
                               auto P = intersect(rec.L, rec2.L);
                                if (area(vec[i].sg.a, vec[i].sg.b, P) \le eps) {
                                          dq.push_back(rec);
                                          break;
                                } else {
                                          continue;
                     }
```

```
 \begin{array}{l} \mbox{while } (sz(dq)>=2) \ \{ \\ \mbox{auto rec} = dq.front(); \\ \mbox{dq.pop\_front()}; \end{array} 
           auto rec2 = dq.front();
           Point P = intersect(rec.L, rec2.L);
           dbg(area(vec[i].sg.a,\,vec[i].sg.b,\,P));\\
           \begin{array}{l} \text{if } (\operatorname{area}(\operatorname{vec}[i].\operatorname{sg.a},\,\operatorname{vec}[i].\operatorname{sg.b},\,P) <= \operatorname{eps}) \; \{ \\ \operatorname{dq.push\_front}(\operatorname{rec}); \end{array}
                 break;
           } else {
                 continue;
           }
      dq.pb(vec[i]);
}
while (sz(dq) >= 3) {
     auto rec = dq.front();
dq.pop_front();
     auto\ rec2 = \overset{\smile}{dq}.front();
     auto P = intersect(rec.L, rec2.L);
      dbg(P);
      if (area(dq.back().sg.a, dq.back().sg.b, P) <= eps) {
           dq.push_front(rec);
           break;
      } else {
           continue:
}
while (sz(dq) >= 3) {
     auto rec = dq.back();
dq.pop_back();
     auto rec2 = dq.back();
      auto P = intersect(rec.L, rec2.L);
      if \; (area(dq.front().sg.a, \; dq.front().sg.b, \; P) \; <= \; eps) \; \{\\
           dq.push_back(rec);
           break;
      } else {
           continue;
}
dbg(T);
dbg(sz(dq));
return sz(dq) > 2;
```

#### 2.4 Line and circle

We say first circle center at point (0,0). Line is (A,B,C). Find  $x_0=\frac{-AC}{A^2+B^2}$ 

 $y_0 = \frac{-BC}{A^2 + B^2}$ 

If dist from  $(x_0, y_0)$  to (0, 0) is bigger than r then answer is zero. Otherwise if it equals to r then answer is  $(x_0, y_0)$ . Otherwise we have 2 solutions:  $d = \sqrt{r^2 - \frac{C^2}{A^2 + B^2}} \ V = \sqrt{\frac{d^2}{A^2 + B^2}} \ a_x = x_0 + BV$   $a_y = y_0 - AV \ b_x = x_0 - BV \ b_y = y_0 + AV$ 

#### 2.5 Two circles

We say first circle center at point (0,0).  $x^2+y^2=r_1^2$   $(x-x_2)^2+(y-y_2)^2=r_2^2$   $x^2+y^2=r_1^2$   $x(-2x_2)+y(-2y_2)+(x_2^2+y_2^2+r_1^2+r_2^2)=0$  Now we have new task to find intersection of line and circle: Ax+By+C=0,  $A=-2x_2$ ,  $B=-2y_2$ ,  $C=x_2^2+y_2^2+r_1^2-r_2^2$ 

#### 2.6 Two common tangets to circles

We have to find all tangent lines to two circles. First circle is at (0,0).

We say first circle center at point (0,0).  $x^2+y^2=r_1^2 (x-x_2)^2+(y-y_2)^2=r_2^2$   $x^2+y^2=r_1^2 x(-2x_2)+y(-2y_2)+(x_2^2+y_2^2+r_1^2+r_2^2)=0$  Now we have new task to find intersection of line and circle:  $Ax+By+C=0,\ A=-2x_2,\ B=-2y_2,\ C=x_2^2+y_2^2+r_1^2-r_2^2$ 

#### 3 Matrix

#### 3.1 Multiply

#### 3.2 Gauss SLAU

```
int gauss (vector < vector < double > > a, vector < double > &
        ans) {
          int n = (int) a.size();
          int m = (int) a[0].size() - 1;
          \begin{array}{l} vector < int > where \ (m, -1); \\ for \ (int \ col=0, \ row=0; \ col < m \ \&\& \ row < n; \ ++col) \ \{ \end{array} 
                    int sel = row;
                    fi(row, n - 1) {
                             if (abs (a[i][col]) > abs (a[sel][col]))
                    if (abs (a[sel][col]) < EPS)
                             continue:
                    fi(col, m) {
                             swap (a[sel][i], a[row][i]);
                    where [col] = row;
                   fi(0, n - 1) {
    if (i != row) {
                                       double c = a[i][col] / a[row][col];
                                        fj(col, m) {
                                                 a[i][j] -= a[row][j] * c;
                             }
                    ++row;
          ans.assign (m, 0);
          fi(0, m - 1) {
                   \begin{array}{c} \text{if (where[i] != -1)} \\ \text{ans[i] = a[where[i]][m] / a[where[i]][i];} \end{array}
          fi(0, n - 1) {
                    double sum = 0;
                    fj(0, m - 1) {
                             \operatorname{sum} \stackrel{\widetilde{}}{+}=\operatorname{ans}[j] * a[i][j];
                    if (abs (sum - a[i][m]) > EPS)
                             return 0;
          fi(0, m - 1) {
                    if (where[i] == -1) return INF;
```

```
} return 1; }
```

## 4 Strings

#### 4.1 Suf LCP

```
string\ s;
vi p, c, rc;

int x = 1;
int gett(int a, int x) {
     if (a > x) return a - x;
     return sz(s) - 1 - (x - a) \% (sz(s) - 1);
int gett2(int a, int x) {
    if\ (a\,+\,x\,<\,sz(s))
         return a + x;
     return 1 + (a + x) \% sz(s);
bool comp1(int a, int b) {
    return (s[a] < s[b]);
int n, S;
vvi e;
void Sort(vector<int> &p) {
    \begin{array}{c} \text{fi}(1, \, \mathbf{n}) \, \{ \\ \text{int } \mathbf{t} = \operatorname{gett}(\mathbf{p}[\mathbf{i}], \, \mathbf{x}); \end{array}
         e[rc[t]].pb(t);\\
     p.clear();
    p.pb(0);
     fi(0, n)
         fz(e[i]) {
              p.pb(z);
         e[i].clear();
    }
void init() {
    n = sz(s) + 1;

s = "#" + s + (char)31;

c = vi(n + 1);
     rc = vi(n + 1);
     fi(0, n) p.pb(i);
     sort(Allf(p), comp1);
     c[1] = 0;
    fi(2, n) {
         if(s[p[i]] == s[p[i-1]]) c[i] = c[i-1];
else c[i] = c[i-1] + 1;
     fi(1, n) {
         rc[p[i]] = c[i];
    e = vvi(n + 1);
    x = 1;
     while (x < n) {
         Sort(p);
         c[1] = 0;
         fi(2, n) {
              if \ (rc[p[i]] == rc[p[i-1]] \ \&\& \ rc[gett2(p[i], \ x)] == rc[
                      gett2(p[i - 1], x)]) {
                   c[i] = c[i-1];
              } else {
                   c[i] = c[i - 1] + 1;
          fi(1, n) {
              rc[p[i]] = c[i];
          x *= 2;
}
```

#### 4.2 LCP

```
\label{eq:continuous_section} \begin{split} & \text{vll lcm}(n+1); \\ & \text{ll } k=0; \\ & \text{fi}(1,\,n) \; \big\{ \\ & \text{ll } x=\text{rev}[i] \text{--}1; \\ \end{split}
```

```
\begin{split} & \text{ll cnt} = \max(k-1,\,(\text{ll})0); \\ & \text{fj}(\text{cnt},\,1e5+10) \; \{\\ & \text{if } (s|i+j] == s[\text{arr}[x]+j]) \; \text{cnt++}; \\ & \text{else break}; \\ & \} \\ & \text{if } (x==0) \; \text{cnt} = 0; \\ & \text{lcm}[x] = \text{cnt}; \\ & k = \text{cnt}; \end{split}
```

#### 4.3 Suf Automaton

```
struct state {
        int len, link;
        map<char, int> next;
const int MAXLEN = 1'000'000;
state st[MAXLEN*2];
int siz, last;
void\ sa\_init()\ \{
        siz = last = 0;

st[0].len = 0;
        st[0].link = -1;
        siz++;
void sa_extend (char c) {
        int cur = siz++
        st[cur].len = st[last].len + 1;
        for (p = last; p != -1 \&\& !st[p].next.count(c); p = st[p].
               link) {
                st[p].next[c] = cur;
        if (p == -1) {
                st[cur].link = 0;
        } else {
                int clone = siz++;

st[clone].len = st[p].len + 1;

st[clone].next = st[q].next;

st[clone].link = st[q].link;
                         for (; p != -1 && st[p].next[c] == q; p = st[p].link)
                                  st[p].next[c] = clone;
                         st[q].link = st[cur].link = clone;
                 }
        last = cur;
```

#### 4.4 pref Function

```
vi prefix_function(string s) {  int \ n = sz(s); \\ vi \ pi(n); \\ fi(1, n - 1) \{ \\ int \ j = pi[i - 1]; \\ while \ (j > 0 \ \&\& \ s[i] \ != s[j]) \ \{ \\ j = pi[j - 1]; \\ \} \\ if \ (s[i] == s[j]) \ j++; \\ pi[i] = j; \\ \} \\ return \ pi; \\ \}
```

### 4.5 Z Function

```
vi z_function (string s) {  \begin{aligned} &\text{int } n = sz(s); \\ &\text{vi } z\ (n); \\ &\text{for (int } i = 1, \, l = 0, \, r = 0; \, i < n; \, i++) \; \{ \\ &\text{if } (i <= r) \; z[i] = \min(r \cdot i + 1, \, z[i \cdot l]); \\ &\text{while } (i + z[i] < n \; \&\& \; s[z[i]] == s[i + z[i]]) \; \{ \\ &z[i]++; \\ \} \\ &\text{if } (i + z[i] \cdot 1 > r) \; \{ \end{aligned}
```

```
\begin{array}{c} l = i; \\ r = i + z[i] - 1; \\ \rbrace \\ \rbrace \\ return \ z; \\ \rbrace \end{array}
```

#### 4.6 Aho Corasick

```
\begin{array}{c} \mathrm{struct} \ \mathrm{Node} \ \{ \\ \mathrm{int} \ \mathrm{leaf} = \text{-}1; \end{array}
         char c = ' \setminus 0';
         Node^* par = nullptr;
         map<char, Node*> nxt;
Node* suf = nullptr;
map<char, Node*> sufs;
         Node(char t, Node* p) {
                 c = t;
                 par = p;
Node* root:
void add(string& s, int j) {
Node* cur = root;
         fi(0, sz(s)- 1) {
                  if(cur->nxt.count(s[i])) {
                           cur = cur->nxt[s[i]];
                  } else {
                           Node* t = new Node(s[i],cur);
                           cur->nxt[s[i]] = t;
                           cur = t;
                  }
         cur->leaf = j;
Node* get_link(Node* a) {
    if(a->suf) return a->suf;
         if(a == root) {
                  a->suf = root;
                  return a->suf;
         if(a->par == root)  {
                 return root;
         char t = a->c;
         Node* cur = get\_link(a->par);
         while(cur != root && !cur->nxt.count(t)) {
                  cur = get_link(cur);
         if(cur->nxt.count(t)) {
                  cur = cur->nxt[t];
         return a->suf = cur;
Node* go(Node* a,char c) {
         Node* cur = a;
         \label{eq:while} while (cur != root \&\& !cur-> nxt.count(c)) \; \{
                  cur = get\_link(cur);
         if(cur->nxt.count(c))
                 cur = cur - > nxt[c];
         return cur;
string s;
ll n;
vector<string> a;
void init() {
         root = new Node('\0',nullptr);
         fi(1,\,n)\,\,\{
                 add(a[i], i);
}
```

## 5 Graphs

#### 5.1 cutpoints

```
void dfs(ll x, ll p = -1) {
    if (color[x]) return;
    color[x] = 1;
     _time++;
    tin[x] = _time;

tup[x] = _time;
    ll q = 0;
    bool f = 0;
    fy(e[x]) {
if (y == p) continue;
        if (!color[y]) {
             q++;
             dfs(y, x);
             \mathrm{if}\ (\mathrm{tup}[y]>=\mathrm{tin}[x])\ \{
                 f = 1;
             tup[x] = min(tup[x], tup[y]);
        } else {
            tup[x] = min(tup[x], tin[y]);
    if (p == -1) {
    if (q > 1) ans.pb(x);
} else {
         if (f) ans.pb(x);
}
```

## 5.2 bridges

```
 \begin{array}{lll} void \ dfs(ll \ x, \ ll \ p = -1) \ \{ \\ & \ if \ (color[x]) \ return; \\ & \ color[x] = 1; \\ & \ tin[x] = \_time++; \\ & \ tup[x] = tin[x]; \\ & \ for \ (auto \ \&to : e[x]) \ \{ \\ & \ if \ (to == p) \ continue; \\ & \ if \ (color[to]) \ \{ \\ & \ tup[x] = \min(tup[x], \ tin[to]); \\ & \ continue; \\ \} \ else \ \{ \\ & \ dfs(to, \ x); \\ & \ tup[x] = \min(tup[x], \ tup[to]); \\ & \ if \ (tup[to] > tin[x]) \ \{ \\ & \ ans.pb(t[\{to, \ x\}]); \\ & \ dbg(mp(to, \ x)); \\ \} \\ \} \\ \} \end{array} \right\}
```

#### 5.3 min Cost Max Flow

```
\begin{array}{c} bool\ deikstr()\ \{\\ vll\ d(\_k+1,\ INF);\\ set < pll > t; \end{array}
      d[S] = 0;
     \begin{array}{l} a[s] = 0;\\ t.ins(\{d[S], S\});\\ vector < Edge *> p(\_k + 1, NULL);\\ while (sz(t)) \end{array}
           auto [\_, x] = *t.begin();
t.erase(t.begin());
            fy(e[x]) {
                  if (y-)f + 1 \le y-c & d[y-y] > d[x] + y-w {
                        t.erase(\{d[y-y], y-y\});

d[y-y] = d[x] + y-y;
                       t.ins({d[y-y], y-y});

p[y-y] = y;
                  }
           }
     if (d[T] == INF) return false;
int x = T;
      while (x != S)  {
           p[x]->f++;

p[x]->rev->f--;
            x = p[x]->rev->y;
     fi(1,
               k) {
            \overline{\mathrm{fy}}(\mathrm{e}[\mathrm{i}]) {
                 y-y = y-y + d[i] - d[y-y];
     }
```

```
return true;
```

## 6 Data Structures

#### 6.1 fenwick Tree

```
// Linear
vll fen;
void\ add(int\ p,\ ll\ val)\ \{
      for(int i = p; i <= n; i = (i \mid (i + 1))) {
fen[i] += val;
ll sum(int p) {
      ll res = 0;
      for
(int i = p; i >= 0; i = (i & (i + 1)) - 1) {
            res += fen[i];
      return res;
}
// Matrix
vvll fen;
void add(int x, int y, ll val) {
      for(int i = x; i <= n; i = (i \mid (i+1))) {
    for(int j = y; j <= n; j = (j \mid (j+1))) {
        fen[i][j] += val;
      }
fl sum(int x, int y) {
      ll res = 0;
      \begin{array}{l} {\rm for}\,({\rm int}\,\,i=x;\,i>=0;\,i=(i\,\,\&\,\,(i+1))\,\,\cdot\,\,1)\,\,\{\\ {\rm for}\,\,({\rm int}\,\,j=y;\,j>=0;\,j=(j\,\,\&\,\,(j+1))\,\,\cdot\,\,1)\,\,\{\\ {\rm res}\,\,+=\,{\rm fen}[i][j]; \end{array}
             }
      return res;
}
```

#### 6.2 dekart Tree

```
struct Node {
     ll val:
      ll cnt;
     ll y;
      Node *1;
     Node *r;
void upd(Node *a) {
     a->cnt = 1:
     if (a->1)
           a->cnt += a->l->cnt;
      if (a->r)
           a->cnt += a->r->cnt;
Node *merge(Node *a, Node *b) {
     if (!a)
           return b;
      if (!b)
           réturn a;
     if (a-y > b-y) \{ a-y = merge(a-y, b);
           upd(a);
           return a;
     b->l = merge(a, b->l);
     upd(b);
     return b;
pair<Node *, Node *> split(Node *a, ll q) {
     if\ (!a)
           return \{0, 0\};
     ll left = 0;
     if (a->l)
     \begin{array}{l} \operatorname{left} += \operatorname{a->l->cnt}; \\ \operatorname{if} (\operatorname{left} >= \operatorname{q}) \; \{ \\ \operatorname{auto} \; [\operatorname{l}, \operatorname{r}] = \operatorname{split}(\operatorname{a->l}, \operatorname{q}); \end{array}
           a->l=r;
           upd(a);
```

```
return\ \{l,\ a\};
           auto [l, r] = split(a->r, q - left - 1);
           a->r = 1;
           upd(a);
           return {a, r};
6.3
                         Long
struct Long \{
           int base = 10000;
           int len = 4:
            vi num = \{1, 0\};
           int operator[](int x) const {
                        if (x > num[0]) return 0;
                        return num[x];
           int& operator[](int x) {
                      if (x > num[0]) {
 num.resize(x + 1, 0);
                                   num[0] = x;
                        return num[x];
           }
Long operator+(const Long& A, const Long& B) {
           Long res;
           int\ n=\max(A[0],\,B[0]);
         \begin{array}{l} \text{fi } i = A \\ \text{fi}(1, n) \; \{ \\ \text{res}[i] = A[i] + B[i]; \\ \text{if } (i > 1) \; \{ \\ \text{res}[i] + = \text{res}[i - 1] \; / \; \text{res.base}; \\ \end{array}
                                   res[i-1] %= res.base;
            \text{while}(\text{res}[\text{res}[0]] >= \text{res.base})  {
                      int m = res[0];

res[m + 1] = res[m] / res.base;
                       res[m] %= res.base;
           return res;
Long operator*(const Long& A, const int& B) {
            Long res:
           int n = A[0];
           fi(1, n) \{
                      \{1, 1\}  \{1, 2\}  \{1, 3\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 4\}  \{1, 
            \text{while}(\text{res}[\text{res}[0]] >= \text{res.base})  {
                      int m = res[0];

res[m + 1] = res[m] / res.base;
                       res[m] %= res.base;
           return res;
Long operator*(const Long& A, const Long& B) {
           Long res;
           fi(1, A[0]) {

if (A[i] == 0) continue;
                      fi (A<sub>[1]</sub> = -0) continue;

fj(1, B<sub>[0]</sub>) {

if (A<sub>[i]</sub> * B<sub>[j]</sub> == 0) continue;

int x = i + j - 1;

res[x] += A<sub>[i]</sub> * B<sub>[j]</sub>;
                                   while(x > 1 && res[x - 1] >= res.base) {
    res[x] += res[x - 1] / res.base;
    res[x - 1] %= res.base;
                        while (res[res[0]] >= res.base) {
                                  int m = res[0];

res[m + 1] += res[m] / res.base;
                                   res[m] %= res.base;
                        }
           return res;
pair<Long, int> operator/(const Long& A, const int B) {
            Long res;
           ll d = 0, m = 0;
```

```
 \begin{array}{l} fdi(A[0],\,1) \,\,\{\\ ll \,\,val \,=\, m \,\,*\,\,res.base \,+\,\,A[i];\\ d \,=\, val \,\,/\,\,B;\\ m \,=\, val \,\,\%\,\,B; \end{array} 
          if (x == 1 \&\& d == 0) continue;
          res[x] = d;
     reverse(Allf(res.num));
     return mp(res, m);
f
void print(const Long& A) {
    printf("%d", A[A[0]]);
    fdi(A[0] - 1, 1) {
        printf("%04d", A[i]);
}
     printf(ln);
Long str_to_Long(const string& s) {
     int n = sz(s);
     Long res;
     int x = 1;
     for
(int i = n - 1; i >= 0; i -= 4) {
          int val = 0;
          for(int j = \max(0, i - 3); j \le i; j++) {
               val = val * 10 + (s[j] - '0');
          res[x] = val;
          x++;
     return res;
Long int_to_Long(int num) {
     Long res;
     int x = 1:
     while(num) {
          int val = num % res.base;
          num /= res.base;
          res[x] = val;
          x++;
     return res:
```

## 7 DP

#### 7.1 fft

```
using cd = complex<double>;
const double PI = acos(-1);
int reverse(int num, int lg_n) {
     int res = 0;
    for (int i = 0; i < lg_n; i++) {
    if (num & (1 << i))
        res |= 1 << (lg_n - 1 - i);
     return res;
void fft(vector<cd> & a, bool invert) {
     int n = a.size();
    int lg_n = 0;
while ((1 << lg_n) < n)
lg_n++;
     for (int i = 0; i < n; i++) {
         if \; (i < reverse(i, lg\_n))
              swap(a[i],\;a[reverse(i,\;lg\_n)]);
    for (int len = 2; len <= n; len <<= 1) { double ang = 2 * PI / len * (invert ? -1 : 1);
          cd wlen(cos(ang), sin(ang));
          for (int i = 0; i < n; i += len) {
              cd w(1);
              for (int j = 0; j < \text{len } / 2; j++) {
                   cdu = a[i+j], v = a[i+j+len/2] * w;
a[i+j] = u + v;
a[i+j+len/2] = u - v;
                   w *= wlen;
              }
         }
     }
    if (invert) {
```

```
for (cd & x : a) {
              x /= n;
    }
vector<int> multiply(vector<int>const& a, vector<int>const&
     vector <\!\!cd\!\!> fa(a.begin(),\,a.end()),\,fb(b.begin(),\,b.end());
    \begin{array}{l} int\ n=1;\\ while\ (n< sz(a)+sz(b)\ )\ \{ \end{array}
         n <<= 1;
     fa.resize(n);
    fb.resize(n);
    fft(fa, false);
fft(fb, false);
    for (int i = 0; i < n; i++)
fa[i] *= fb[i];
    fft(fa, true);
     vector<int> result(n);
    for (int i = 0; i < n; i++)

result[i] = round(fa[i].real());

for( int i = 0; i < n; i++){
       result[i] = min(result[i],5);
     while(result.back() == 0)result.pop_back();
    return result:
}
```

#### 7.2 convex Hull Trick

```
 \begin{array}{l} struct\ Line\ \{\\ mutable\ ll\ k,\ m,\ p;\\ bool\ operator<(const\ Line\ \&o)\ const\ \{\ return\ k< o.k;\ \} \end{array} 
                  bool operator<(ll x) const { return p < x; }
struct LineContainer : multiset<Line, less<>>> {
                   // (for doubles, use \inf = 1/.0, \operatorname{div}(a,b) = a/b)
                 // (to dotted the first three first three
                 bool is
ect(iterator x, iterator y) \{
                                  \mathrm{if}\;(\mathrm{y} == \mathrm{end}(\mathrm{y}))
                                                  return x->p = \inf_{x \to \infty} 0;
                                  if (x->k == y->k)

x->p = x->m > y->m? inf : -inf;
                                                   x\text{-}{>}p=\mathrm{div}(y\text{-}{>}m\text{ - }x\text{-}{>}m,\;x\text{-}{>}k\text{ - }y\text{-}{>}k);
                                   \mathrm{return}\ x\text{-}\!\!>\!\!\mathrm{p}>=\mathrm{y}\text{-}\!\!>\!\!\mathrm{p};
                 void add(ll k, ll m) { auto z = insert(\{k, m, 0\}), y = z++, x = y;
                                    while (isect(y, z))
                                 | z = erase(z);
| if (x != begin() && isect(--x, y))
| isect(x, y = erase(y));
| while ((y = x) != begin() && (--x)->p >= y->p)
                                                     isect(x, erase(y));
                 ll query(ll x) {
                                   assert(!empty());
auto l = *lower_bound(x);
return l.k * x + l.m;
};
```

## 8 Comb

#### 8.1 basic

$$\begin{split} Cat(n,k) &= \frac{1}{n+1}C(2n,n) \\ N(n,k) &= \frac{1}{n}C(n,k)C(n,k-1) \\ &\sum_{n=0}^{\inf} \sum_{k=1}^{n} N(n,k)z^n t^k = \frac{1+z(t-1)-\sqrt{1-2z(t+1)+z^2(t-1)^2}}{2z} \\ S(n,k) &= \sum_{j=0}^{k-1} (-1)^j C(k,j) \cdot (m-j)^n \ O(k\log n) \end{split}$$