

On the Schur-Weyl Duality and its Role in Quantum Information

Edward Kim

May 13, 2020

Abstract

The Schur-Weyl Duality describes the decomposition of representations of $GL(d, \mathbb{C})$ and the symmetric group S_n into a direct sum of irreducible tensor representations indexed by partitions described by Young frames. In particular, we are interested in the decomposition of the n -tensored d -dimensional Hilbert space $\mathcal{H}^{\otimes n}$ and the subsequent operators acting on it as described by the two groups above. Calculation of the dimensions of these irreducible representations reduces to a combinatorial problem, allowing us to derive useful asymptotic bounds. These bounds provide elegant proofs for Schumacher's compression rate and the subadditivity properties of von-Neumann entropy.

Contents

1	The Schur-Weyl Duality	2
1.1	Young Tableaux	2
1.1.1	Standard Young Tableau	3
1.1.2	Semi-Standard Young Tableau	4
1.2	Bounds on Dimension of U_λ and V_λ	4
1.3	Schur Polynomials	5
2	Quantum Compression	6
2.1	Keyl-Werner Theorem	6
2.2	Towards Compression	7
3	Subadditivity Properties of von-Neumann entropy	8
3.1	Clebsch-Gordon Decomposition	8
3.2	Subadditivity and SSA	9
4	Conclusion	9

1 The Schur-Weyl Duality

Formal definitions of the Schur-Weyl Duality revolve around the commutation property of the actions of $GL(d, \mathbb{C})$ and S_n for $n, d \geq 1$. First, we define the following basic notions to state the theorem:

Let \mathcal{H} denote a d -dimensional complex Hilbert space with orthonormal basis $\{i\}_{1 \leq i \leq d}$ and let $\mathcal{H}^{\otimes n}$ denote its corresponding n -tensor form. This allows us to define the following representations:

$$R : GL(d, \mathbb{C}) \rightarrow GL(\mathcal{H}^{\otimes n}) :$$

$$R(g) \cdot |\psi_1\rangle \otimes |\psi_2\rangle \otimes \dots \otimes |\psi_n\rangle = g \cdot |\psi_1\rangle \otimes g \cdot |\psi_2\rangle \otimes \dots \otimes g \cdot |\psi_n\rangle, \quad |\psi_i\rangle \in \mathcal{H}$$

and

$$U : S_n \rightarrow GL(\mathcal{H}^{\otimes n}) :$$

$$R(\sigma) \cdot |s_1\rangle \otimes |s_2\rangle \otimes \dots \otimes |s_n\rangle = |s_{\sigma(1)}\rangle \otimes |s_{\sigma(2)}\rangle \otimes \dots \otimes |s_{\sigma(n)}\rangle, \quad s_i \in \{i\}_{1 \leq i \leq d}$$

The Schur-Weyl Duality is based upon the following commutation property: For a subset $A \subseteq GL(\mathcal{H}^{\otimes n})$, define

$$Comm(A) = \{s \in End(\mathcal{H}^{\otimes n}) \mid s \cdot x = x \cdot s, \forall x \in A\}$$

One can check that this is an associative unital \mathbb{C} -algebra. The commutant theorem shown by Schur states that:

Theorem 1. *Let $\mathcal{A} = span R(GL(d, \mathbb{C}))$ and $\mathcal{B} = span U(S_n)$, then $Comm(\mathcal{A}) = \mathcal{B}$ and $Comm(\mathcal{B}) = \mathcal{A}$*

We omit the proof here for the sake of brevity, but the reader can find a proof using the Double Commutant Theorem [1]. A direct corollary is the decomposition of $\mathcal{H}^{\otimes n}$ into tensor products of irreducible representations of $GL(d, \mathbb{C}), S_n$. Here, we borrow our notation and formulation from [2].

$$\mathcal{H}^{\otimes n} = \bigoplus_{\lambda \vdash_d n} U_\lambda \otimes V_\lambda \tag{1}$$

where the U_λ, V_λ refer to irreducible representations of $GL(d, \mathbb{C}), S_n$ respectively and $\lambda \vdash_d n$ refers to the partitions of n having at most d rows. The following block decompositions also hold:

$$R(g) = \bigoplus_{\lambda \vdash_d n} \pi_{U_\lambda}(g) \otimes 1_{V_\lambda} \tag{2}$$

$$U(\sigma) = \bigoplus_{\lambda \vdash_d n} 1_{U_\lambda} \otimes \pi_{V_\lambda}(\sigma) \tag{3}$$

where the maps

$$\pi_{U_\lambda} : GL(d, \mathbb{C}) \rightarrow GL(U_\lambda)$$

$$\pi_{V_\lambda} : S_n \rightarrow GL(V_\lambda)$$

refer to the matrices corresponding to the linear maps describing the action on the corresponding irreducible representations.

1.1 Young Tableaux

We now describe the formal definition of the partitions λ such that $\lambda \vdash_d n$, namely the partitions of n having at most d rows. A *partition* $\lambda = (\lambda_1, \dots, \lambda_d)$ is described as a sequence of non-increasing non-negative integers such that:

$$\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_d, \quad \sum_{i=1}^d \lambda_i = n$$

These partitions will provide the visual groundwork for calculating the dimensions of the irreducible representations present in $U_\lambda \otimes V_\lambda$. The *Young Tableaux* provide a relatively simple combinatorial framework to facilitate these calculations. Suppose we are given a partition as defined above, then one can construct a corresponding *Young frame* as follows: each λ_i represents the number of cells present in the i^{th} row of the Young frame. Furthermore, our rows index from top to bottom, and our cells index from left to right. For example, the Young diagram shown in Figure 1 represents a $(3, 3, 2, 1)$ partition.

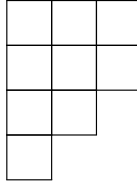


Figure 1: Young diagram of a (3,3,2,1) partition

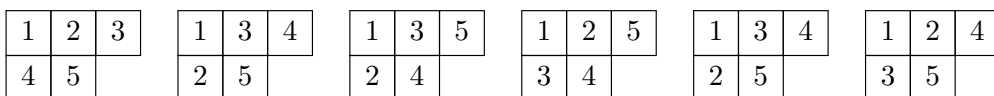
This gives a bijective correspondence between Young frames containing n cells with at most d rows and partitions λ such that $\lambda \vdash_d n$. Hence, we can simply refer to the Young frame λ rather than the partition λ . Once we have determined the proper Young diagram for λ , the dimensions of the irreducible representations are calculated through filling these frames under specific rules respecting the symmetry imposed by the two groups. These rules are enforced via the *standard Young tableau (SYT)* for S_n -representations and the *semi-standard Young tableau (SSYT)* for $GL(d, \mathbb{C})$ -representations.

1.1.1 Standard Young Tableau

Given Young frame λ indexing an irreducible S_n -representation, a filling of λ via the *standard Young tableau* (SYT) obeys the following conditions:

1. Cells are filled with any integer from $\{1, \dots, n\}$ where n will be the total number of cells in the frame. Each integer can only occur once in the tableau i.e can only be of weight one.
2. The rows must be *strictly increasing* from left to right and along columns from top to bottom.

Figure 2 is borrowed from [2] and gives fillings of a $(3, 2)$ frame via this standard Young tableau.

Figure 2: Fillings of a $(3, 2)$ partition via SYT

1.1.2 Semi-Standard Young Tableau

Given Young frame λ indicating an irreducible $GL(d, \mathbb{C})$ -representation, a filling of λ via the *semi-standard Young tableau* (SSYT) follows the following conditions:

1. Cells are filled with any integer from $\{1, \dots, d\}$. These integers can be of arbitrary weight i.e occur more than once or not at all.
2. Rows are *non-decreasing* along rows from left to right and *strictly increasing* along columns from top to bottom.

Figure 3 uses the same frame as Figure 2 except we adhere to the semi-standard Young tableau with integers $\{1, 2, 3\}$

1	2	3
2	3	

1	1	3
2	3	

1	1	2
2	3	

2	2	2
3	3	

1	1	1
3	3	

1	1	1
2	2	

Figure 3: Some fillings of $(3, 2)$ partition with $\{1, 2, 3\}$ under SSYT

A more detailed exposition on the calculation of these Young Tableaux is given in [1], [3].

1.2 Bounds on Dimension of U_λ and V_λ

Based on the Young Tableau filling rules explained in the previous section, we can explicitly calculate $\dim U_\lambda$ and $\dim V_\lambda$ for a given Young frame λ by counting the number of valid fillings possible. For the irreducible $GL(d, \mathbb{C})$ -representation U_λ , we calculate the number of SSYT possible for Young frame λ :

$$\dim U_\lambda = \frac{\prod_{1 \leq i < j \leq d} (\lambda_i - \lambda_j - i + j)}{\prod_{i=1}^{d-1} i!}$$

A simple analysis shows that $(\lambda_j - \lambda_i + i - j) \leq (n + 1)$. Furthermore, there are $d(d - 1)/2$ distinct pairs considered in the numerator's product. This immediately gives us the bound:

$$\dim U_\lambda \leq (n + 1)^{d(d-1)/2} \quad (4)$$

The calculation for the dimension of the irreducible S_n -representation V_λ requires an auxiliary computation called the *Hook-length* of a box. Formally speaking:

$$\dim V_\lambda = \frac{n!}{\prod_i h(i)}$$

where the product walks over all boxes in the Young frame λ . The function value $h(i)$ refers to the hook length of box i which is defined to be the number of cells *below* and to the *right* of box i . We can bound this value by observing that we can simply disregard the vertical cells below each box in a row. Thus, given all boxes in row i , the hook length can be bounded from below:

$$\lambda_i! \leq \prod_{r \in \text{row}_\lambda(i)} h(r)$$

where $\text{row}_\lambda(i)$ is the set of all cells comprising row i of the Young frame λ . By applying this definition for each row, we get the following upper bound:

$$\dim V_\lambda \leq \frac{n!}{\prod_i \lambda_i!} \quad (5)$$

Similarly, we can add the total number of rows to the analysis to get the upper bound:

$$\prod_{r \in \text{row}_\lambda(i)} h(r) \leq (\lambda_i + d)!$$

which subsequently gives the lower bound:

$$\frac{n!}{\prod_i (\lambda_i + d)!} \leq \dim V_\lambda$$

The bounds given for $\dim V_\lambda$ are particularly important for our connection to quantum information. To see this, note that we can reexpress our bounds as follows:

$$\binom{n}{\lambda_1 + d, \lambda_2 + d, \dots, \lambda_d + d} \leq \dim V_\lambda \leq \binom{n}{\lambda_1, \lambda_2, \dots, \lambda_d}$$

where the two bounds are expressed as multinomial coefficients. From classical Shannon entropy, we know that these multinomial coefficients are asymptotic to $2^{nH(\bar{\lambda})}$ where H is the usual Shannon entropy function. Explicitly applying this to our coefficients above shows that:

$$\binom{n}{\lambda_1, \lambda_2, \dots, \lambda_d} \sim 2^{nH(\bar{\lambda})}$$

where $\bar{\lambda} = (\frac{\lambda_1}{n}, \frac{\lambda_2}{n}, \dots, \frac{\lambda_d}{n})$ is the *reduced form* of the Young frame λ . Hence, the bounds above directly yield that:

$$\dim V_\lambda \sim 2^{nH(\bar{\lambda})} \quad (6)$$

1.3 Schur Polynomials

The final tool we need involves taking the trace of the $\pi_{U_\lambda}(g)$ for some $g \in GL(d, \mathbb{C})$. These are determined through the *Schur polynomials*:

$$\text{Tr}(\pi_{U_\lambda}(g)) = s_\lambda(r_1, \dots, r_d) \quad (7)$$

where $\text{spec}(g) = (r_1, \dots, r_d)$. The polynomials s_λ are defined by the following sum over all semi-standard Young Tableau $\Gamma(\lambda)$ over Young frame λ :

$$s_\lambda(x_1, \dots, x_d) = \sum_{t \in \Gamma(\lambda)} x_1^{w_1(t)} x_2^{w_2(t)} \dots x_d^{w_d(t)} \quad (8)$$

where $w_i : \Gamma(\lambda) \rightarrow \{1, \dots, n\}$ denotes the function outputting the weight of integer i assigned in tableau t . In other words, $w_i(t)$ counts the number of occurrences of integer i in tableau t . Using this definition we can bound the Schur polynomials. Without loss of generality, assume that we order $\text{spec}(\rho) = (r_1, \dots, r_d)$ in non-increasing order $r_1 \geq r_2 \geq \dots \geq r_d$.

$$\prod_{i=1}^d r_i^{\lambda_i} \leq s_\lambda(r_1, \dots, r_d) \leq \dim U_\lambda \prod_{i=1}^d r_i^{\lambda_i} \quad (9)$$

The left-hand side simply follows from the definition of the Schur polynomials and the filling rule imposed by the semi-standard Young tableau. As s_λ runs over all such tableau t , it must be at least as large as one such filling obeying the semi-standard Young tableau rule. On the other hand, recalling that the number of semi-standard Young tableau is equal to $\dim U_\lambda$ and noticing that the product gives the maximum value over any such filling of λ under the semi-standard rule gives us the right-hand side.

2 Quantum Compression

2.1 Keyl-Werner Theorem

Perhaps the foundation of our bridge between quantum information and representation theory roots itself in the celebrated Keyl-Werner Theorem. As mentioned in their original paper [4], Keyl and Werner were originally motivated by the problem of seeking a set of observables which could be used to output good approximations of the spectra of an arbitrary density operator $\rho \in \mathcal{D}(\mathcal{H})$ for d -dimensional Hilbert space \mathcal{H} . Roughly speaking, the theorem shows that measuring on the subspaces of the form $U_\lambda \otimes V_\lambda$ where the Young frames of λ are “close” to $\text{spec}(\rho)$ give good estimates of $\text{spec}(\rho)$. The formal proof of the theorem was refined by Hayashi and Matsumoto [5] and corrected slightly by Christandl and Mitchison [6]. We follow the latter’s formulation.

Theorem 2. *Given a density operator ρ with spectrum $r = \text{spec}(\rho)$ and $\epsilon_1 > 0$, define projection operator $P_X = \sum_{\lambda \vdash_d n, \bar{\lambda} \in B_{\epsilon_1}(r)} P_\lambda$ where P_λ is the projection on $U_\lambda \otimes V_\lambda$. Then for $\epsilon_2 > 0$, there exists sufficiently large k such that:*

$$\text{Tr}(P_X \rho^{\otimes k}) > 1 - \epsilon_2$$

where we recall that $\bar{\lambda}$ is reduced partition as first defined in the previous section. To show this, we need a lemma. Assume that $\text{spec}(\rho)$ is ordered in non-increasing order henceforth.

Lemma 1. *Let $\rho \in \mathcal{D}(\mathcal{H})$ and P_λ be the projection on $U_\lambda \otimes V_\lambda$ described by Young frame λ . Then:*

$$\text{Tr}(P_\lambda \rho^{\otimes n}) \leq (n+1)^{d(d-1)/2} \cdot \exp \left\{ -nD(\bar{\lambda} || \text{spec}(\rho)) \right\}$$

where $D(\bar{\lambda} || \text{spec}(\rho))$ refers to the Kullback-Liebler distance between distributions $\bar{\lambda}, \text{spec}(\rho)$

Proof. By Property 2,

$$\rho^{\otimes n} = \bigoplus_{\lambda \vdash_d n} \pi_{U_\lambda}(\rho) \otimes 1_{V_\lambda}$$

We can calculate the probability of projecting on P_λ :

$$\text{Tr}(P_\lambda \rho^{\otimes n}) = \text{Tr}(\pi_{U_\lambda}) \text{Tr}(1_{V_\lambda}) = s_\lambda(r_1, \dots, r_d) \dim V_\lambda$$

Now the bounds derived for the Schur polynomials at (9) and the dimension bounds mentioned at (4) and (5) give us that:

$$\begin{aligned} \text{Tr}(P_\lambda \rho^{\otimes n}) &\leq \dim V_\lambda \dim U_\lambda \prod_i r_i^{\lambda_i} \\ &\leq (n+1)^{d(d-1)/2} \frac{n!}{\prod_i \lambda_i!} \prod_i r_i^{\lambda_i} \\ &\leq (n+1)^{d(d-1)/2} \cdot \exp \left\{ -nD(\bar{\lambda} || \text{spec}(\rho)) \right\} \end{aligned}$$

□

From the above lemma, we arrive at the following corollary.

Corollary 1. Let $P_X = \sum_{\lambda \vdash_d n: \bar{\lambda} \in \mathcal{S}} P_\lambda$. Then:

$$\text{Tr}(P_X \rho^{\otimes n}) \leq (n+1)^{d(d+1)/2} \cdot \exp \left\{ -n \min_{\lambda \vdash_d n, \bar{\lambda} \in \mathcal{S}} D(\bar{\lambda} \| \text{spec}(\rho)) \right\}$$

where $\mathcal{S} \subset [0, 1]^d$.

Proof. Observe the following string of inequalities:

$$\begin{aligned} \text{Tr}(P_X \rho^{\otimes n}) &= \sum_{\lambda \vdash_d n: \bar{\lambda} \in \mathcal{S}} \text{Tr}(P_\lambda \rho^{\otimes n}) \leq \sum_{\lambda \vdash_d n: \bar{\lambda} \in \mathcal{S}} (n+1)^{d(d-1)/2} \cdot \exp \left\{ -n D(\bar{\lambda} \| \text{spec}(\rho)) \right\} \\ &\leq \sum_{\lambda \vdash_d n} (n+1)^{d(d-1)/2} \cdot \exp \left\{ -n \min_{\lambda \vdash_d n, \bar{\lambda} \in \mathcal{S}} D(\bar{\lambda} \| \text{spec}(\rho)) \right\} \\ &\leq (n+1)^d (n+1)^{d(d-1)/2} \cdot \exp \left\{ -n \min_{\lambda \vdash_d n, \bar{\lambda} \in \mathcal{S}} D(\bar{\lambda} \| \text{spec}(\rho)) \right\} \\ &= (n+1)^{d(d+1)/2} \cdot \exp \left\{ -n \min_{\lambda \vdash_d n, \bar{\lambda} \in \mathcal{S}} D(\bar{\lambda} \| \text{spec}(\rho)) \right\} \end{aligned}$$

□

The last inequality comes from the bound

$$|\{\lambda : \lambda \vdash_d n\}| \leq (n+1)^d \quad (10)$$

which just follows from the binomial expansion. We see that by taking our subset \mathcal{S} to be the complement of the ball $\mathcal{B}_{\epsilon_1}(\text{spec}(\rho))$, the probability of projecting onto the subspaces indexed by Young frames which lie outside of this ball tend to zero by the corollary above. This gives us Theorem 2 as required.

2.2 Towards Compression

We now further develop the connection between quantum information and representation theory through an illustrative example concerning the Quantum Compression rate. Recall the setup of the Quantum Compression problem where we have a quantum source describe by ensemble $\{p_X(x), |\psi_x\rangle \langle \psi_x|\}$. This gives us the density operator $\rho = \sum_x p_X(x) |\psi_x\rangle \langle \psi_x|$. The resulting state from n -outputs from the quantum source can thus be described as $\rho^{\otimes n}$. We wish to determine a compression rate R such that we can project onto a subspace of 2^{nR} . Furthermore, we wish for $R < \log d$; that is, ideally the subspace W we wish to project onto is of smaller dimension than d . Finally, this subspace must capture enough of the quantum information encoded into ρ in the sense that:

$$\limsup_{n \rightarrow \infty} \|\rho^{\otimes n} - \Pi_W \rho^{\otimes n} \Pi_W\|_1 = 0$$

We can view this setup as the quantum analogue of the classical Shannon source compression. This connection turns out to be deeper as Schumacher's compression theorem extends the concept of typical sequences in classical Shannon theory to *typical subspaces* [7]. In fact, the KeyL-Werner result tells that these subspaces can be taken to be the subspaces given by *typical Young frames*. We formalize this below. Assume once again that $\text{spec}(\rho)$ is ordered in non-increasing order.

Define the projection operator

$$\Pi(T_\epsilon) = \bigoplus_{\lambda: |\bar{\lambda} - \text{spec}(\rho)| < \epsilon} P_\lambda$$

where P_λ is the projection operator for the subspace $U_\lambda \otimes V_\lambda$. By the KeyL-Werner Theorem (Theorem 2), we know that

$$\lim_{n \rightarrow \infty} \|\rho^{\otimes n} - \Pi(T_\epsilon)\rho^{\otimes n}\Pi(T_\epsilon)\|_1 = 0$$

Furthermore, we use (4), (5) to calculate the dimension of the support of $\Pi(T_\epsilon)$:

$$\text{Tr}(\Pi(T_\epsilon)) \leq \sum_{\lambda: |\bar{\lambda} - \text{spec}(\rho)| < \epsilon} \dim_{U_\lambda} \dim_{Y_\lambda} \leq |\mathcal{I}_{\epsilon, \rho}| (n+1)^{d(d-1)/2} \frac{n!}{\prod_{i=1}^d \lambda_i!} \quad (11)$$

where $\mathcal{I}_{\epsilon, \rho}$ is the set of all Young frames which are ϵ -close to $\text{spec}(\rho)$. Here we can take $(\lambda_1, \dots, \lambda_d)$ to be an ϵ -typical Young frame such that $\prod_{i=1}^d \lambda_i!$ is minimized. Now this set is bounded by the total number of Young frames of n cells having at most d rows given by (10).

Hence by (6), we see that the polynomial factors vanish asymptotically to give way to the exponential factor:

$$\text{Tr}(\Pi(T_\epsilon)) \sim 2^{nH(\bar{\lambda})}$$

which yields that:

$$\text{Tr}(\Pi(T_\epsilon)) \sim 2^{n(S(\rho) + \epsilon)}$$

The derivation aligns with Schumacher's result by noting that the dimension of our Young frame subspace is asymptotic to the von-Neumann entropy $S(\rho)$. As $n \rightarrow \infty$, we observe that we cannot compress better than the rate given by $S(\rho)$. The reader can find a more rigorous approach in [8].

3 Subadditivity Properties of von-Neumann entropy

3.1 Clebsch-Gordon Decomposition

An elegant proof of the subadditivity property of von-Neumann entropy arises from the combination of the *Clebsch-Gordon decomposition* and the Keyl-Werner Theorem. We simply state the decomposition whose proof can be found in [3].

Theorem 3. *Let V_γ, V_ω be two irreducible S_n -representations indexed by partitions $\gamma, \omega \vdash_d n$. Then there exists an S_n -isomorphism:*

$$V_\gamma \otimes V_\omega \cong \bigoplus_{\lambda \vdash n} g_{\lambda\gamma\omega} V_\lambda$$

where $g_{\lambda\gamma\omega}$ are the Kronecker coefficients indicating the multiplicity of V_λ .

Note that there are no restrictions on the number of rows for frame λ . The Kronecker coefficients have another interpretation through the Schur-Weyl duality applied to the analysis of tripartite entanglement [9].

Using this decomposition along with the Keyl-Werner Theorem, Christandl and Mitchison proved the following [6]:

Theorem 4. *Let $\rho_{AB} \in \mathcal{D}(\mathcal{H}_A \otimes \mathcal{H}_B)$ for d -dimensional Hilbert spaces $\mathcal{H}_A, \mathcal{H}_B$. Then there exists a sequence of partitions $(\lambda_i, \gamma_i, \omega_i)$ with $|\lambda_i| = |\gamma_i| = |\omega_i|$ such that $g_{\lambda_i \gamma_i \omega_i} \neq 0$, and*

$$\begin{aligned} \lim_{i \rightarrow \infty} \bar{\lambda}_i &= \text{spec}(\rho_{AB}) \\ \lim_{i \rightarrow \infty} \bar{\gamma}_i &= \text{spec}(\rho_A) \\ \lim_{i \rightarrow \infty} \bar{\omega}_i &= \text{spec}(\rho_B) \end{aligned}$$

3.2 Subadditivity and SSA

Using the above theorem, we can give a simple proof of subadditivity:

Theorem 5. *von-Neumann entropy is subadditive: $S(\rho_{AB}) \leq S(\rho_A) + S(\rho_B)$*

Proof. Theorem 4 gives us a sequence of partitions $(\lambda_i, \gamma_i, \omega_i)$ converging to the spectra of $\rho_{AB}, \rho_A, \rho_B$ respectively. As $g_{\lambda_i \gamma_i \omega_i} \neq 0$, we have that:

$$\dim V_{\lambda_i} \leq \dim V_{\gamma_i} \dim V_{\omega_i}$$

as $i \rightarrow \infty$.

By (5), (6) given above, we have that

$$2^{nS(\rho_{AB})} \leq 2^{nS(\rho_A)S(\rho_B)}$$

where $n = |\lambda_i| = |\gamma_i| = |\omega_i|$. The result follows after taking the logarithm of both sides. \square

In fact, a similar line of thought using recoupling coefficients is used to prove the strong subadditivity property of von-Neumann entropy as well [9]. These recoupling coefficients arise from a direct generalization of the Clebsch-Gordon decomposition introduced above, but the starting ingredients are more-or-less the same.

4 Conclusion

We introduced the Schur-Weyl Duality as a powerful tool allowing us to not only decompose $\mathcal{H}^{\otimes n}$ into tensor products of irreducible representations but also allowing us to utilize the Young Tableaux to easily calculate the dimensions of these irreducible representations. These methods yield deeper insight of the spectra of density operators as seen through the Keyl-Werner result. Instead of thinking about tensor products of eigenspaces of ρ , we can give good approximations through subspaces described by the more intuitive framework given by Young frames. These techniques give us slick proofs of the subadditivity property of von-Neumann entropy and a different interpretation of the Quantum Source Compression rate first derived by Schumacher. However, this is only scratching the bare surface of the deep intertwining of representation theory and quantum information theory. For a more detailed tract on the field, the reader is encouraged to read Aram Harrow's thesis [8]. Matthias Christandl's thesis offers more insight into the role of Kronecker coefficients naturally arising from dimension calculations involving multipartite entangled systems [10].

References

- [1] R. Goodman and N. Wallace, *Symmetry, Representations, and Invariants*. Springer: Graduate Texts in Mathematics, 2009.
- [2] A. Botero, "Quantum information and the representation theory of the symmetric group," *Revista Colombiana de Matematicas*, vol. 50, no. 2, pp. 191–209, 2016.
- [3] W. Fulton and J. Harris, *Representation Theory: A First Course*. Springer: Graduate Texts in Mathematics, 1995.
- [4] M. Keyl and R. Werner, "Estimating the spectrum of density operator," *Physical Review A*, vol. 64, 2001.
- [5] M. Hayashi and K. Matsumoto, "Quantum universal variable-length source coding," *Physical Review A*, Vol.66, 022311, 2002.

- [6] M. Christandl and G. Mitchison, “The spectra of quantum states and the kronecker coefficients of the symmetric group,” *Communications in Mathematical Physics*, vol. 261, pp. 789–797.
- [7] B. Schumacher, “Quantum coding,” *Physical Review A*, vol. 51, pp. 2738–2747, 1995.
- [8] A. W. Harrow, “Applications of coherent classical communication and the schur transform to quantum information theory,” 2005.
- [9] M. Christandl, M. B. Sahinoglu, and M. Walter, “Recoupling coefficients and quantum entropies,” *Annales Henri Poincare*, vol. 19, pp. 385–410.
- [10] M. Christandl, “The structure of bipartite quantum states - insights from group theory and cryptography,” 2006.