## Last Lecture

Eugene Wu

### Administrivia

Review session 5/6 2-4PM 480 Mudd

Thank the staff!

Project 2 due 5/9 IOAM. no late days!

Extra credit opportunities

- project 2, recursive queries
- write a tutorial
- draw a concept/topic

## Grading

Normalize all assignments to [0-1]
Weighed sum based on assignment weights
Compute cut-offs, ~ B+ avg
Then add in extra credit

## Agenda

- I. Project Winners
- 2. Modern Databases
- 3. Q&A / Review

### DBMSes in the Wild

#### Classic Relational

\$\$: Oracle, IBM, Microsoft, Teradata, EMC, etc

Free: MySQL, PostgreSQL

#### New Relational

In-Memory, Column-store, Streaming

#### Non-traditional

Search (Google, Bing, Lucene), Scientific, Geographic

#### **NoSQL**

Big Data: Hadoop, Spark, etc

Key-value: Mongo, BerkeleyDB, Cassandra, etc

#### DBMS-as-a-Service

Microsoft Azure, Amazon Redshift/RDS, etc...

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**DBMS-as-a-Service** 

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## Modern Database Systems

Hardware changes affect

Compression is good → Column stores

Large scale aggregation queries

"Data Warehouses"

Memory is cheap → In-memory stores

Transactional systems

## One Size Fits All

#### Traditional Database

Row oriented
Disk based
Tabular data

## One Size Does Not Fits All

Traditional Database
Row oriented
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Columnar

In Memory

Streaming

Scientific

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## Data Warehouses

Store all historical data for future analysis

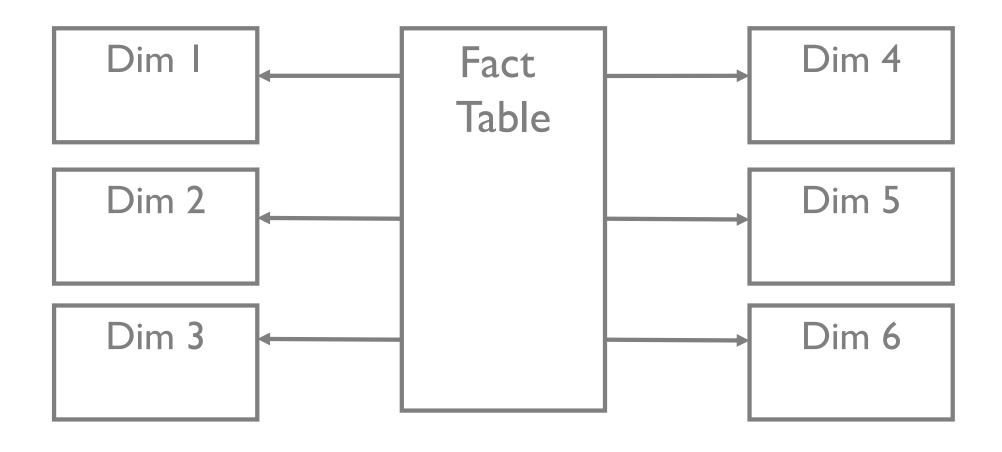
Sales by month over past 20 years

Clicks by youth in texas

Cost by product component

Defacto standard at any company

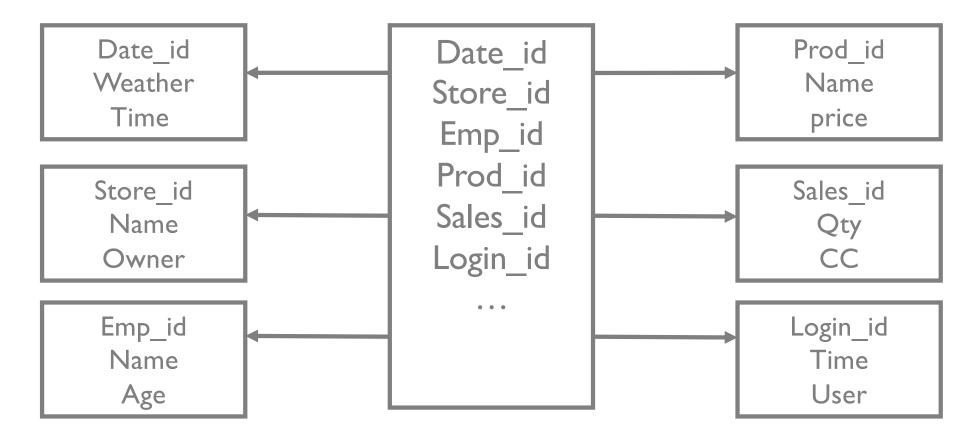
## Star Schema



### Star Schema

Fact table is "fat"

Queries access ~6 attrs

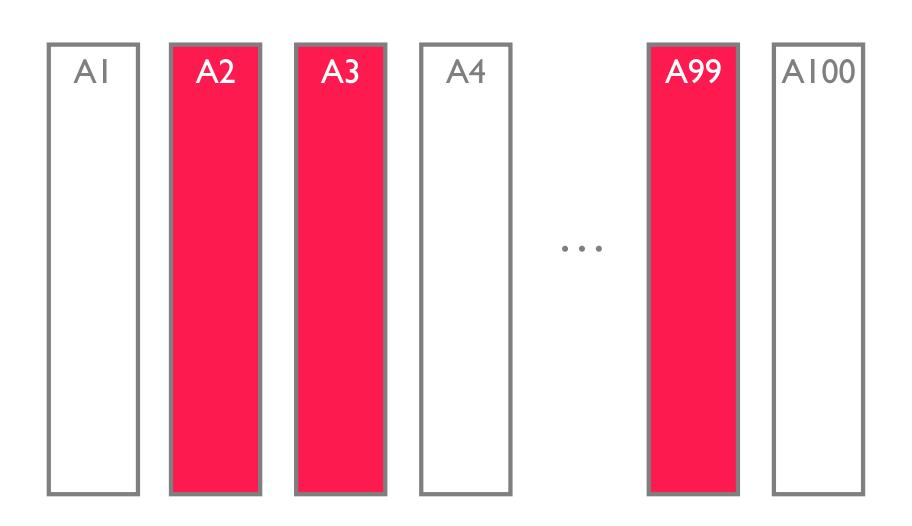


#### 100 attributes

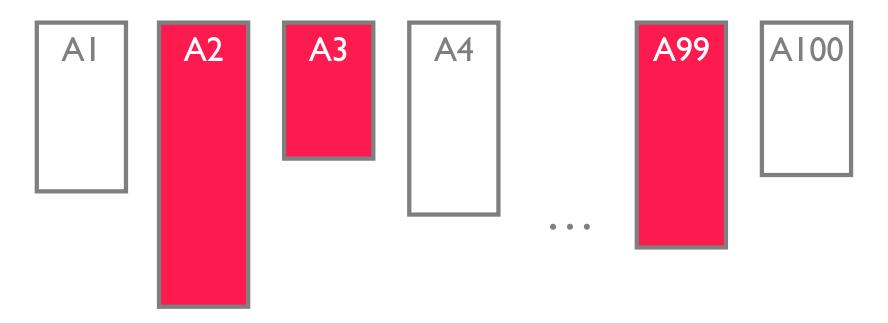
Tuple I Tuple 2 Tuple 3 Tuple 4 Tuple 5

Al	A2	A3	A4		A99	A100
				• • •		

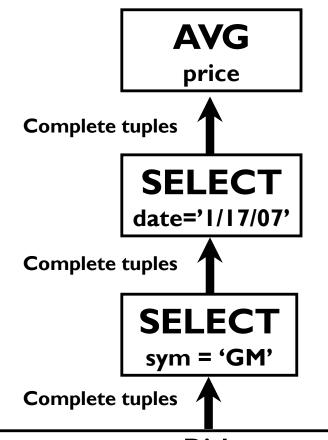
### 16x less data read. Unfair advantage.



# 16x less data read. Unfair advantage.Compression better on single columnExecute on compressed data



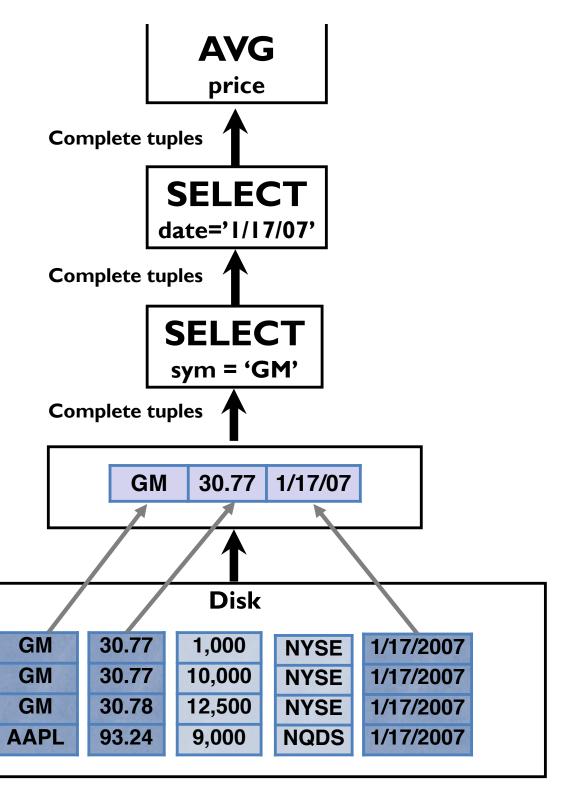
#### Traditional DBMS



		Disk		
GM	30.77	1,000	NYSE	1/17/2007
GM	30.77	10,000	NYSE	1/17/2007
GM	30.78	12,500	NYSE	1/17/2007
AAPL	93.24	9,000	NQDS	1/17/2007

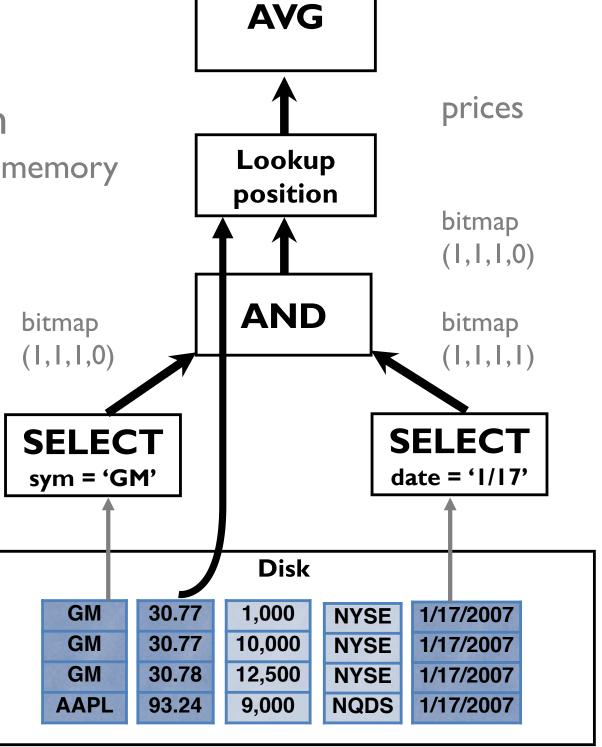
#### Naïve:

Early Materialization Row oriented execution



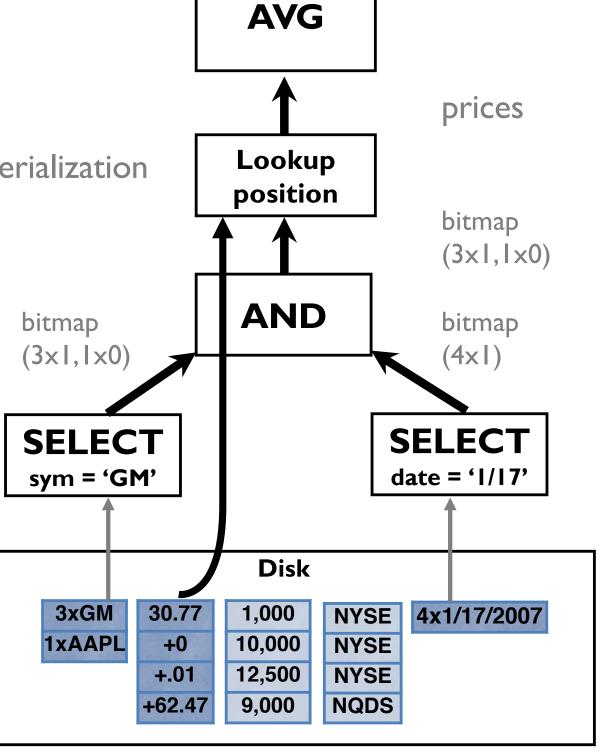
## C-Store Late Materialization

Much less data moving in memory



## C-Store Compression

Only possible w/ late materialization



## Column Stores

Optimized for data warehouses

Store data by attribute/column rather than row

Compression

Compressed query plan execution

50-100x faster than row store

## In-Memory DBMSes

#### Transaction-oriented apps

remove I unit from product move 5 units from org I to org 2 (shopping carts, inventory)

#### Data stored in memory

Disk only used for check pointing

No concurrency

Active-active replication for fault-tolerance

### Traditional Database

Indexes queries go faster

Concurrency queries go faster

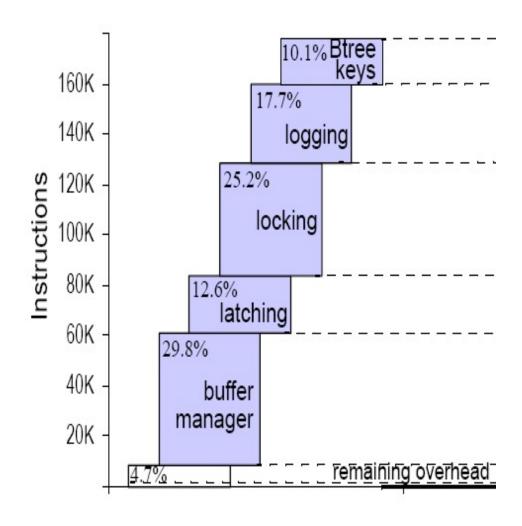
Locking serializability

Logging recovery

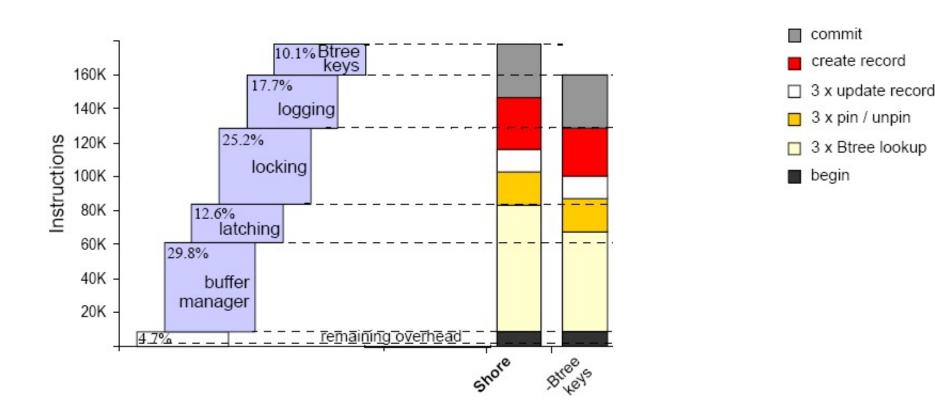
Buffer Manager manage pages in memory

## Results after removing the components (in # instruction)

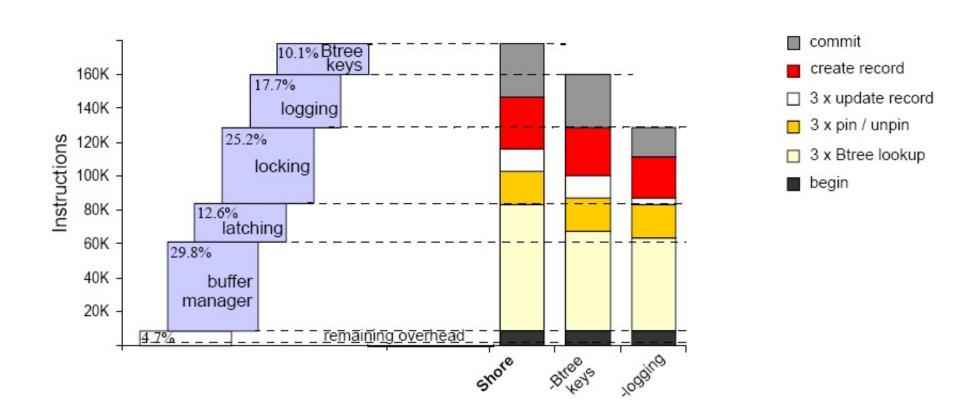
Instruction of useful work is only <2% of a memory resident DB



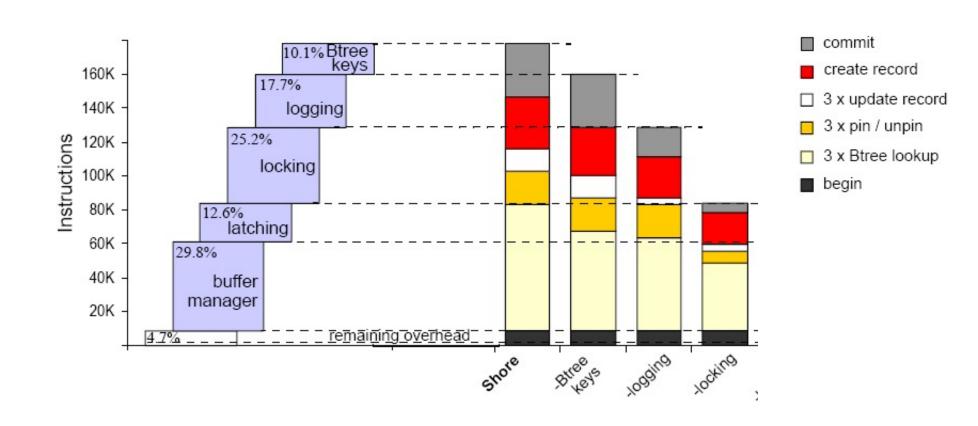
## Effect of removing different components for payment (1/6)



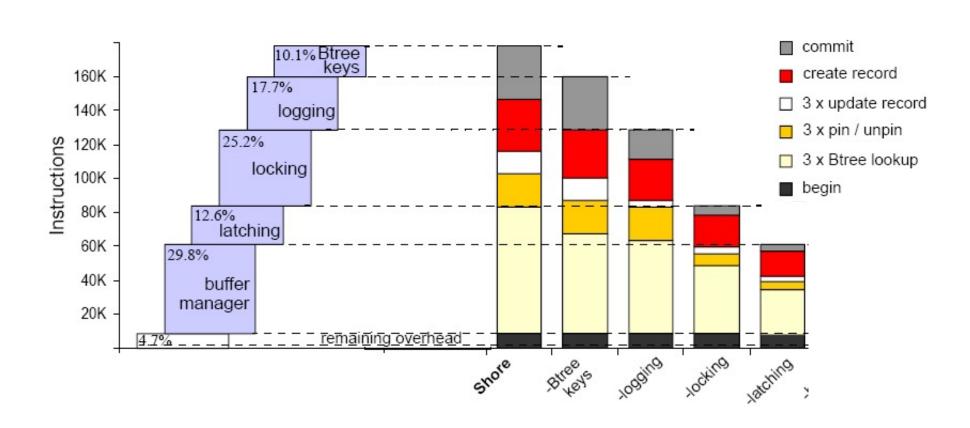
## Effect of removing different components for payment (2/6)



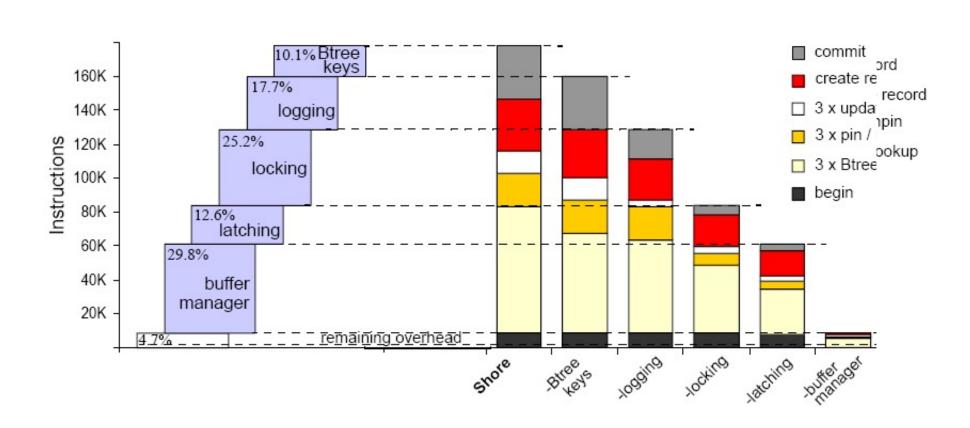
## Effect of removing different components for payment (3/6)



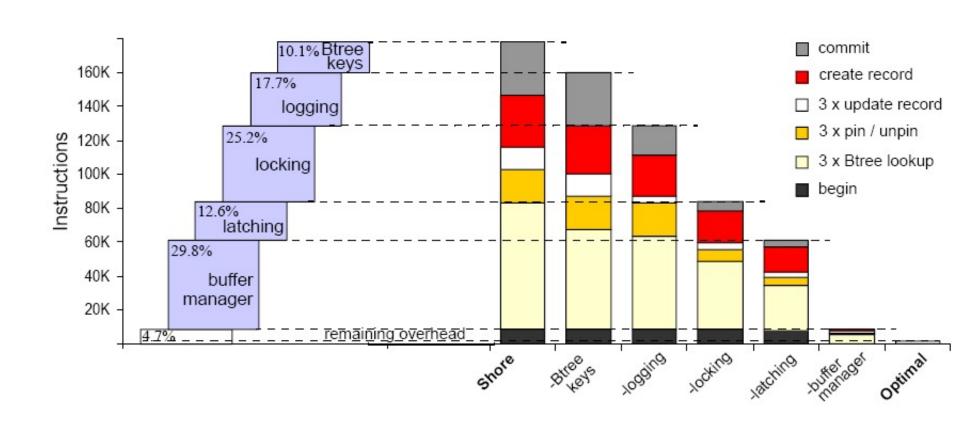
## Effect of removing different components for payment (4/6)



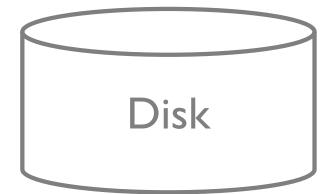
## Effect of removing different components for payment (5/6)



## Effect of removing different components for payment (6/6)



Barebones Executor



In memory DB

Barebones Executor

What about

Parsing

Concurrency?

Recovery?

#### In memory DB

Barebones Executor

**Stored Procedures** 

#### What about

Parsing

Concurrency?

Recovery?

#### Procedure:

#### Query:

p1(10)

In memory DB

Barebones Executor

What about

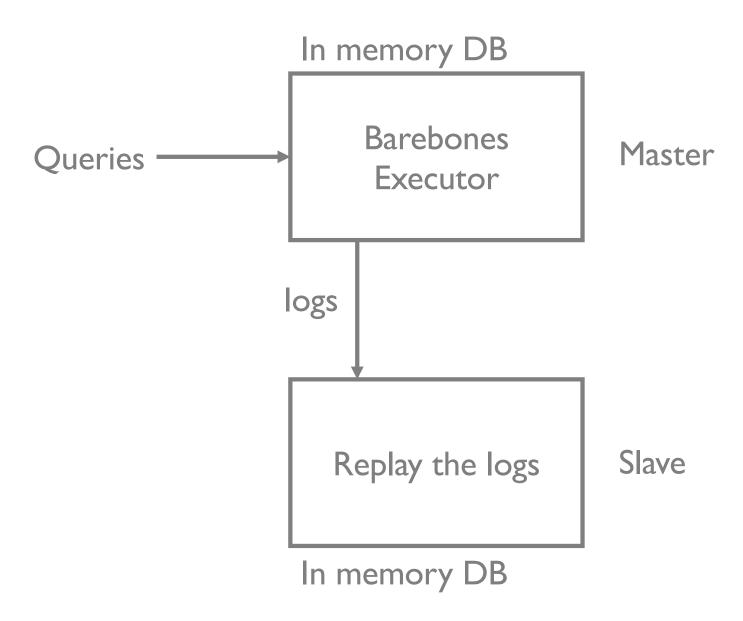
Parsing

Concurrency?

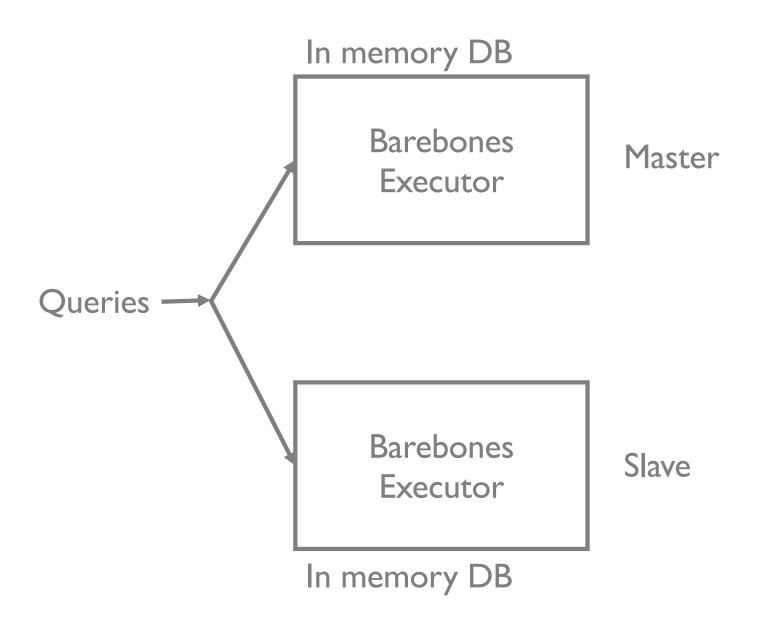
no buffer manager, no concurrency, no locks

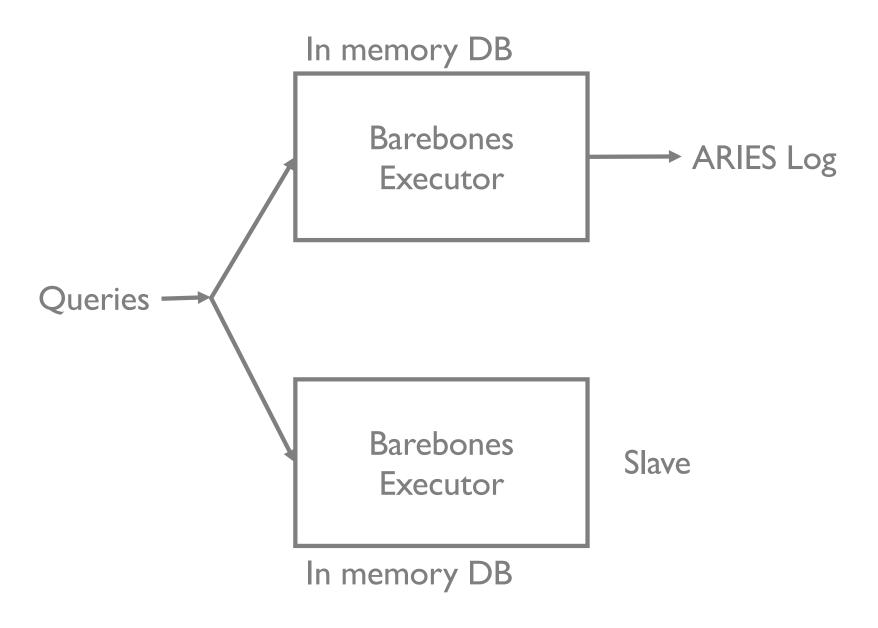
Recovery?

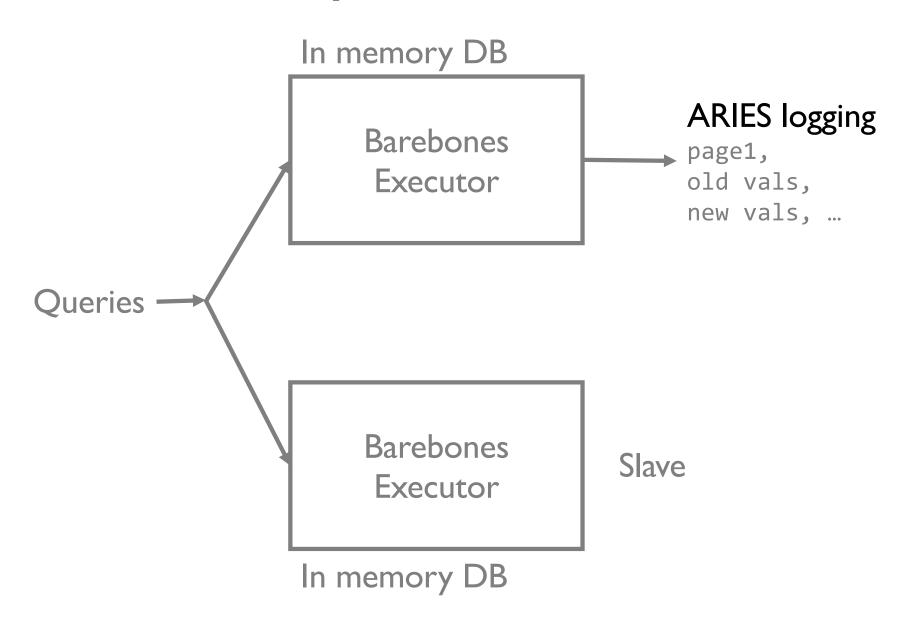
## Log Shipping

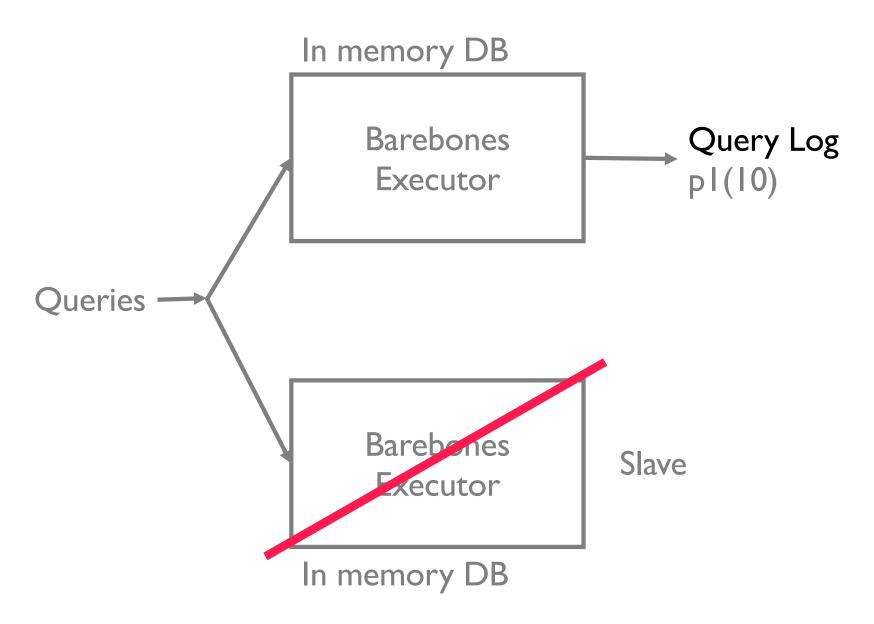


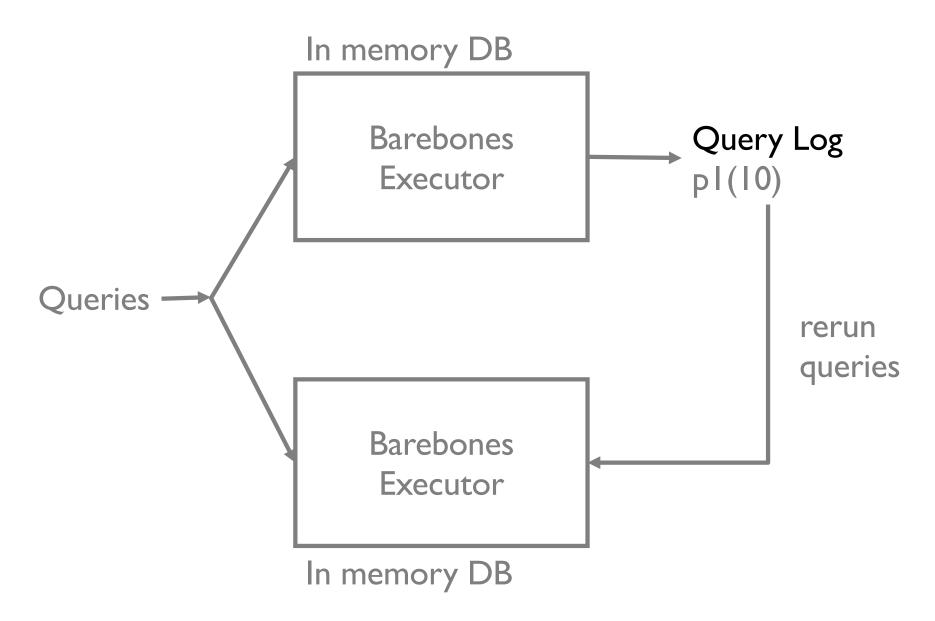
#### Active-Active











#### One Size Does Not Fits All

Traditional Database
Row oriented
Disk based
Tabular data

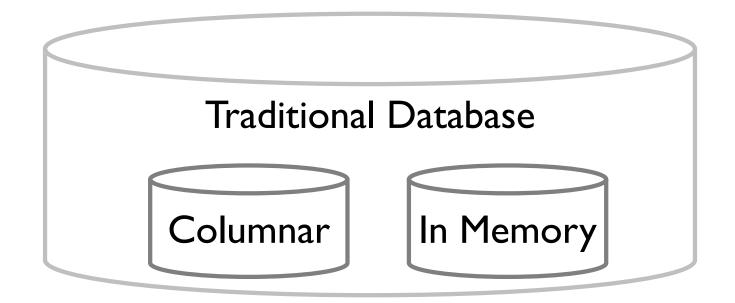
Columnar

In Memory

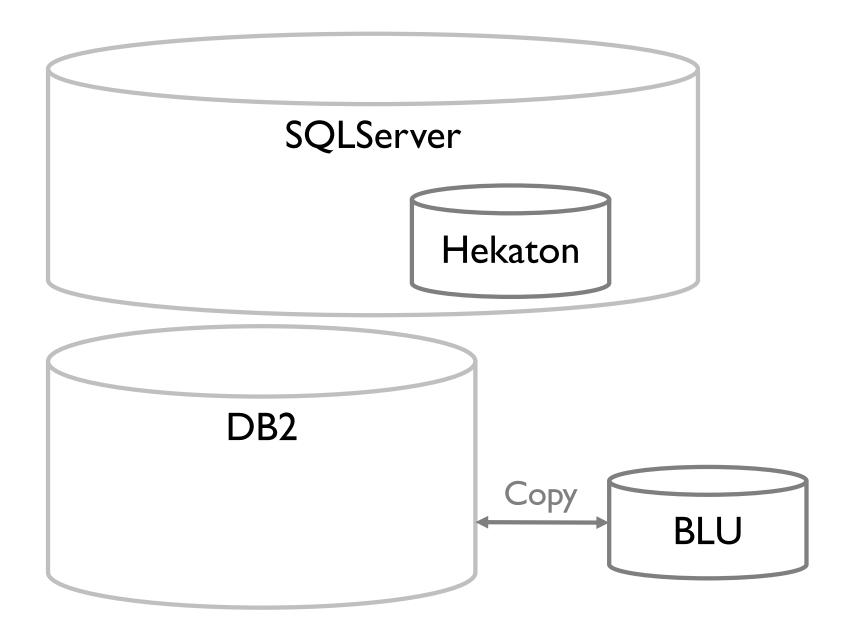
Streaming

Scientific

#### One Size Does Not Fits All



### One Size Does Not Fits All



# OK Let's Step Back

Discussed the how:

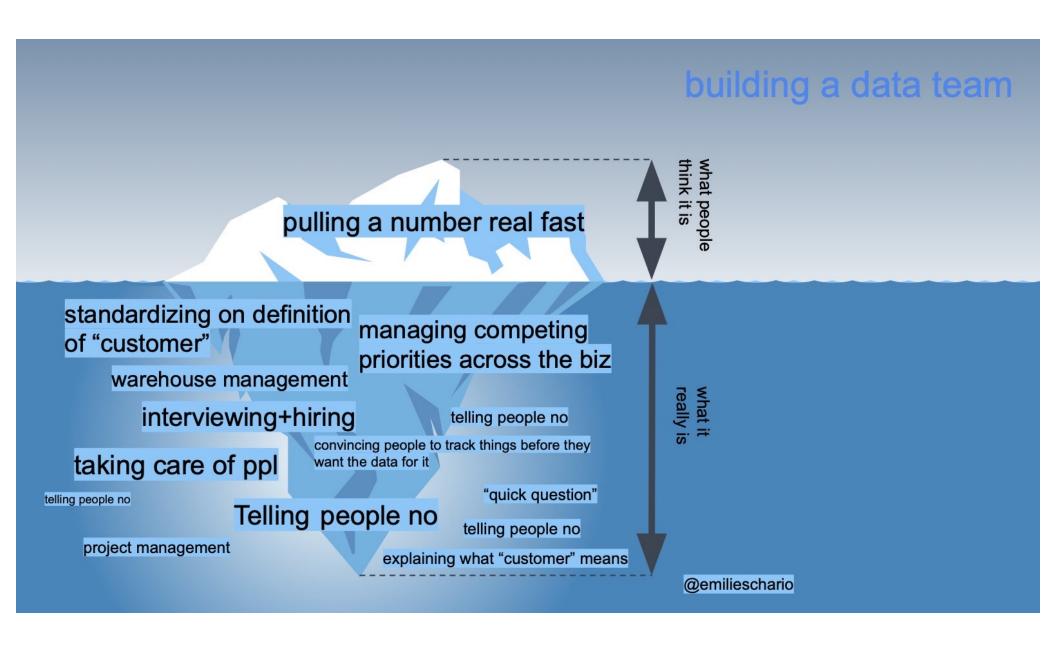
how to model data needed by an application

how to query databases

how databases execute queries

how databases run fast, correctly

To what end?



#### Production ML

#### An on-call engineer's biggest nightmare

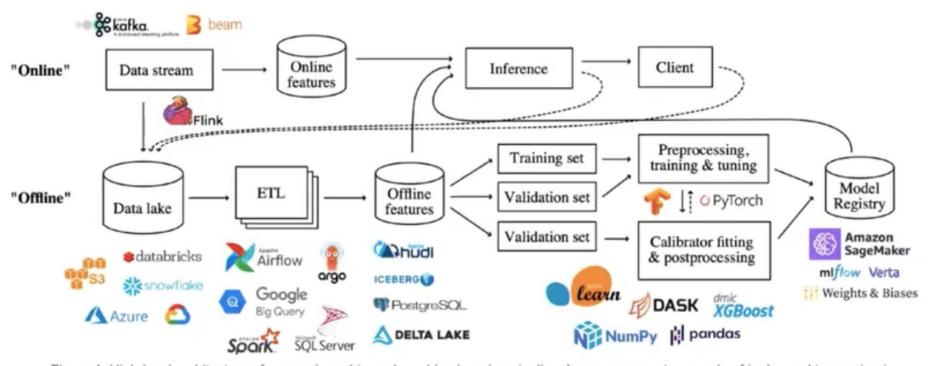


Figure 1: High-level architecture of a generic end-to-end machine learning pipeline. Logos represent a sample of tools used to construct components of the pipeline, illustrating heterogeneity in the tool stack. Shankar et al. 2021

https://www.facebook.com/Engineering/videos/1578607659138164/

4xxx: "here are facts about data management"

- 4111: basics
- 4112: gory DBMS internals

6xxx: "how people figured out those facts"

- 6111: info extraction and web
- 6113: class and modern DB research
- 6998: systems for human data interaction

PhD: "we don't understand anything. Plz help"

### Thanks!

Please fill out courseworks survey cool.