

Processes

Processes - Linux

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Processes

Several processes are started at bootstrap

> Automatic

Started at bootstrap, terminated at shut-down. Execute support activities.

- Daemon Processes eg admin
- E-mail applications
- Various control activities, virus scan and others
- · ...

On user request

- Line printer management
- WEB server
- · ...

Processes

- Process identity and process control
 - ID & system calls: pid, getpid, getppid, etc.
- Process creation
 - The creating process is called the parent process, the process created is called the child process
 - It is possible to create a process tree.
 - System calls: fork, exec, system, etc.
- Process synchronization and termination
 - System calls: wait, waitpid, exit.

Process identifier

- Every process has a unique identifier
 - > PID o Process Identifier
- The PID is a non negative integer
 - Although a PID is unique, UNIX reuses the numbers of terminated processes.
 - ➤ PID can be used by concurrent processes for creating unique objects, or temporary filenames
 - For example :
 sprintf(filename, "file-%d", getpid());
 creates a different process-dependent filename

Process identification

- ❖ The first process, PID=0, is a system process
 - The **swapper**, which is responsible for memory management and process scheduling
- The second process, PID=1, is init a daemon executed at user level, but with super-user privileges
 - Becomes the parent of each orphan process, i.e., of a child of a parent process already terminated

Operating Systems

Process identification

```
#include <unistd.h>
pid_t getpid(); // Process ID
pid_t getppid(); // Parent Process ID
```

Other identifiers related to a process

getppid returns the identifier of the parent process

```
pid_t getppid(); // Parent Process ID
```

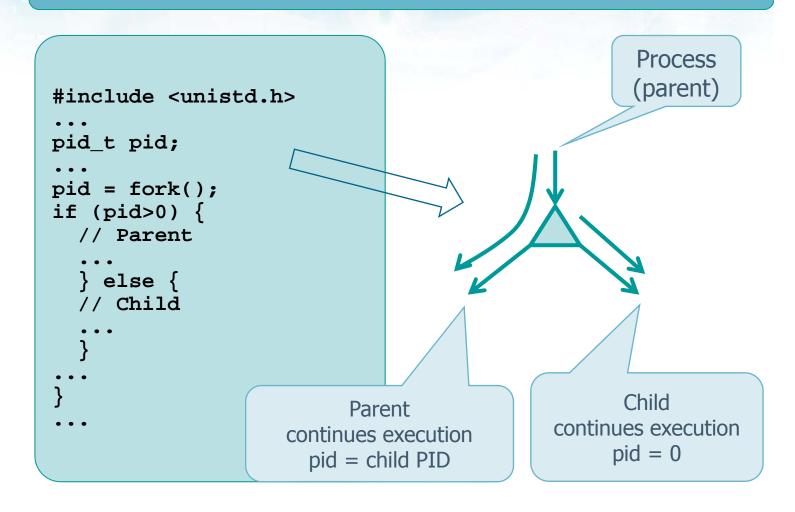
- System call fork() creates a new child process
 - > The child is a copy of the parent excluding the Process ID (PID) returned by fork
 - The parent process receives the child PID
 - A process may have more than one child that can identify on the basis of its PID
 - The child process receives the value 0
 - It can identify its parent by means of the system call getppid
 - Fork is issued once in the parent process, but returns in two different processes, and returns different values to the parent, and to the child.

```
#include <unistd.h>
pid_t fork (void);
```

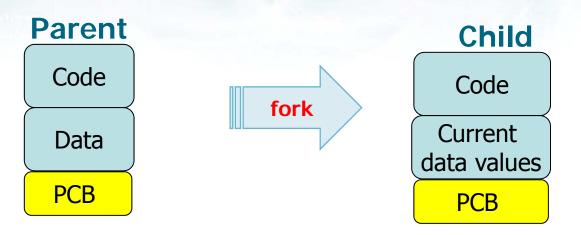
Returned values

- > If fork returns without error
 - Child PID in the parent process
 - Zero in the child process
- > Fork returns -1 in case of error
 - Normally because a limit on the number of allowed process has been reached

```
Process
#include <unistd.h>
                                                    (parent)
pid_t pid;
pid = fork();
switch (pid) {
  case -1:
    // Fork failure
    exit (1);
  case 0:
    // Child
  default:
    // Parent
                                                   Child
                          Parent
                                             continues execution
                     continues execution
                                                  pid = 0
                       pid = child PID
```



Address space



- Parent and child share their code, but they have a different PCB
- > They also share the value of the data at the time of fork.
- Parent and child may then change the data values independently

- Write a concurrent C program that allows
 - > Creating a child process
 - ➤ Terminating the parent process before the child process, or viceversa
- Output in both cases the Process Identifier of the terminating process and the Process Identifier of its parent.
 - Who is the parent's parent?
 - Who is the child's parent if the parent terminates before the child?

```
tC = atoi (argv[1]);
                                   tF = atoi (argv[2]);
#include <unistd.h>
printf ("Main :
                                   ");
printf ("PID=%d; My parent PID=%d\n",
  getpid(), getppid());
pid = fork();
                                                  Child
if (pid == 0){
  sleep (tC);
  printf ("Child : PIDreturned=%d ", pid);
  printf ("PID=%d; My parent PID=%d\n",
    getpid(), getppid());
                                                  Parent
} else {
  sleep (tF);
  printf ("Parent: PIDreturned=%d ", pid);
  printf ("PID=%d; My parent PID=%d\n",
    getpid(), getppid());
```

> ps

PID TTY TIME CMD

2088 pts/10 00:00:00 bash

2760 pts/10 00:00:00 ps

Shell status (ps: prints process status)

Child waits 2 secs Parent waits 5 secs

> ./u04s01e03-fork 2 5

Main: PID=2813; My parent PID=2088

Child: PIDreturned=0 PID=2814; My parent PID=2813

Parent: PIDreturned=2814 PID=2813; My parent PID=2088

Notice increasing PID values

Child waits 5 secs parent waits 2 secs

> ./u04s01e03-fork 5 2

Main: PID=2815; My parent PID=2088

Parent: PIDreturned=2816 PID=2815; My parent PID=2088

> Child : PIDreturned=0 PID=2816; My parent PID=1

init process PID

Exercise

- Given the following program, draw its
 - Control Flow Graph, CFG
 - > Process generation graph

```
int main () {
  /* fork a child process */
  fork();

  /* fork another child process */
  fork();

  /* fork a last one */
  fork();
}
```

```
Control Flow Graph
int main () {
                                                                          (CFG)
  fork (); // 1
  fork (); // 2
  fork (); // 3
                                                                              \mathbf{F}_{11}
                                                                     \mathbf{F}_{12} \mathbf{F}_{11}
                                                       \mathbf{F}_{21} \mathbf{F}_{1}
```

```
Process generation tree
int main () {
 fork (); // 1
 fork (); // 2
 fork (); // 3
                                          F2
```

Exercise

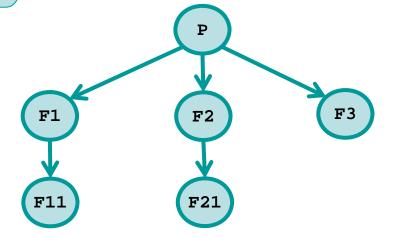
- Given the following program, draw its
 - Control Flow Graph, CFG
 - > Process generation graph

```
pid = fork (); /* call #1 */
if (pid != 0)
  fork (); /* call #2 */
fork (); /* call #3 */
```

Process generation tree

Control Flow Graph (CFG)

• • •



Exercise

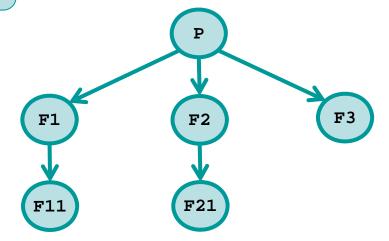
- Given the following program, draw its
 - Control Flow Graph, CFG
 - > Process generation graph

```
pid = fork() /* call #1 */
fork();    /* call #2 */
if (pid != 0)
    fork();    /* call #3 */
```

Process generation tree

Control Flow Graph (CFG)

. . .



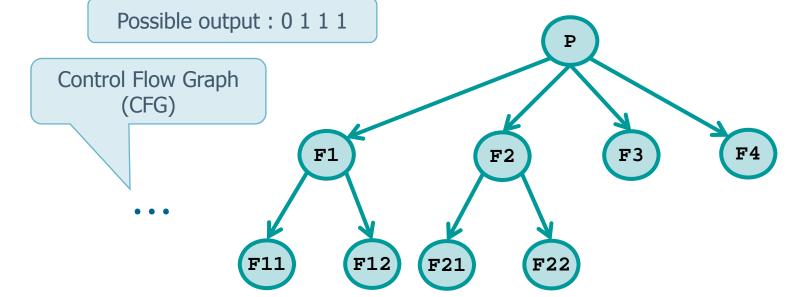
Exercise

- Given the following program, draw its
 - Control Flow Graph, CFG
 - > Process generation graph

```
#include <stdio.h>
int main () {
  int i:
  for (i=0; i<2; i++) {
    printf("i: %d \n", i);
    if (fork()) /* call #1 */
    fork(); /* call #2 */
  }
}</pre>
```

```
for (i=0; i<2; i++) {
  printf("i: %d \n", i);
  if (fork()) // 1
    fork(); // 2
}</pre>
```

Process generation tree



Exercise

- Write a concurrent program that, given n as its argument, generates n children processes
- Each child process outputs its PID and terminates

```
int i, n;

scanf ("%d", &n);
for (i=0; i<n; i++) {
  fork();
  printf ("Proc %d (PID=%d)\n",
      i, getpid());
}

exit (0);</pre>
```

Erroneous solution 1

```
int i, n;

scanf ("%d", &n);
for (i=0; i<n; i++) {
  fork();
  printf ("Proc %d (PID=%d)\n",
      i, getpid());
}

exit (0);</pre>
```

Possible output with n=3

```
Proc 0 (PI D=3188)
Proc 1 (PI D=3188)
Proc 2 (PI D=3188)
Proc 2 (PI D=3191)
Proc 1 (PI D=3190)
Proc 2 (PI D=3190)
Proc 0 (PI D=3189)
Proc 1 (PI D=3189)
Proc 2 (PI D=3189)
Proc 2 (PI D=3192)
Proc 2 (PI D=3192)
Proc 2 (PI D=3193)
Proc 2 (PI D=3193)
Proc 2 (PI D=3193)
Proc 2 (PI D=3195)
```

Erroneous solution 2

```
int i, n;
                                                 Process tree with
                                                     n=3
scanf ("%d", &n);
for (i=0; i<n; i++) {
  if (fork() > 0);
    printf ("Proc %d (PID=%d)\n",
               i, getpid());
exit (0);
                                       F1
                                                    F2
    Generates 7 children processes
                                 F11
                                            F12
                                                   F21
                                F111
```

```
int i, n;
scanf ("%d", &n);
printf ("Start PID=%d\n",
  getpid());
for(i=0; i<n; i++) {
  if (fork() == 0) {
    printf ("Proc %d (PID=%d)\n",
      i, getpid());
    break;
printf ("End PID=%d (PPID=%d)\n",
   getpid(), getppid());
exit(0);
```

```
> ps

PID TTY TIME CMD

088 pts/10 00:00:00 bash

> ./u04s01e06-fork

Star PID=3225

End PID=3225 (PPID=2088)

Proc 2 (PID=3228)

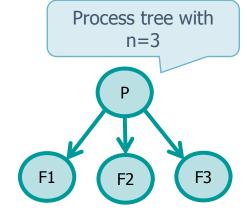
End PID=3228 (PPID=1314)

Proc 1 (PID=3227)

End PID=3227 (PPID=1314)

Proc 0 (PID=3226)

End PID=3226 (PPID=1314)
```



Resources

- The child process is a new entry in the Process Table
- The process resources can be
 - Completely shared among parent and children processes
 - Same address space
 - Partially shared
 - Address spaces partially overlapped
 - > Non shared
 - Separate address spaces

Resources

- In UNIX/Linux parent and child share
 - > The source code
 - > UID, GID, etc.
 - > The open file descriptors (File Description Table)
 - In particular, stdin, stdout, and stderr
 - Concurrent I/O operation implies producing interlaced I/O
 - > The root and the working directory
 - > System resources and their utilization limits
 - > Signal Table
 - > Etc.

Resources

- In UNIX/Linux parent and child have different
 - > return fork value
 - > PID
 - The parent keeps its PID
 - The child gets a new PID
 - > Data, heap and stack space
 - The initial value of the variables is inherited, but the spaces are completely separated
 - copy-on-write technique is used by the OS
 - New memory is allocated only when one of the processes changes the content of a variable

```
char c, str[10];

c = 'X';
if (fork()) {
    // parent (!=0)
    c = 'F';
    strcpy (str, "parent");
    sleep (5);
} else {
    // child (==0)
    strcpy (str, "child");
}

fprintf(stdout, "PID=%d; PPID=%d; c=%c; str=%s\n",
    getpid(), getppid(), c, str);
```

PID=2777; PPID=2776; c=X; str=child PID=2776; PPID=2446; c=F; str=parent

Output

Process termination

- Five standard methods for process termination
 - > return from main
 - > exit system call
 - > _exit Or _Exit
 - Synonyms defined in ISO C or POSIX
 - Similar effects of exit, but different management of stdio flushing etc.
 - > return from main of the last process thread
 - pthread_exit from the last process thread

Process termination

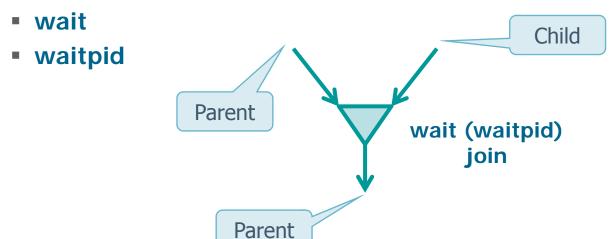
- Three not-normal method for process termination
 - ➤ If a termination signal, or a signal not caught is received
 - > abort
 - Generates the signal sigabort
 - > If the last thread of a process is cancelled

System call wait () and waitpid ()

- When a process terminates (normally or not)
 - > The kernel sends a signal (SIGCHLD) to its parent
 - > For the parent this is an asynchronous event
 - > The parent process may
 - Manage the child termination (and/or the signal)
 - Ignore the event (default)

System call wait () and waitpid ()

- A parent process can manage child termination
 - > By means of system calls



- ➤ Using a signal handler for signal **SIGCHLD**
 - This approach will be introduced in the section devoted to signals

System call wait ()

```
#include <sys/wait.h>
pid_t wait (int *statLoc);
```

- Blocks the calling process if all its children are running (none is already terminated)
 - wait will return as soon as one of its children terminates
- Returns an error if the calling process has not children.

System call wait ()

```
#include <sys/wait.h>
pid_t wait (int *statLoc);
```

- Returns the exit status of a child
 - ➤ If a child process terminates, and the parent process does not call wait, the child exit status remains pending in the kernel Process Control Block

System call wait ()

```
#include <sys/wait.h>
pid_t wait (int *statLoc);
```

The statLoc parameter

Exit status of the child process

- > Is an integer pointer
 - If not NULL collects the exit value of the child
- > The status information are
 - Implementation dependent
 - Recovered using macros in <sys/wait.h> (WIFEXITED, WIFSIGNALED, etc.)

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System call wait ()

```
#include <sys/wait.h>
pid_t wait (int *statLoc);
```

- Returns
 - > The PID of a terminated child

Example

```
pid_t pid, childPid;
int statVal;
pid = fork();
switch (pid) {
  case -1:
    printf ("Fork failed\n"); exit(1);
  case 0:
    // Child
    sleep (10);
  default:
    // Parent
    sleep (05);
```

Example

```
if (pid != 0) {
 // Padre
 childPid = wait (&statVal);
 printf("Child terminated: PID = %d\n", childPid);
 if (WIFEXITED(statVal))
   // WIFEXITED: True if correctly terminated
   // WEXITSTATUS: Takes the 8 returned LSBs (exit)
   printf ("Exit value: %d\n",
      WEXITSTATUS (statVal));
   else
     printf ("Abnormal termination\n");
 exit(code);
```

Zombie processes

- A child process terminated, whose parent is running, but has not et executed wait is in the zombie state
 - ➤ The data segment of the process remains in the process table because the parent could need the child exit status
 - ➤ The child entry is removed only when the parent executes wait
 - Many zombie processes may remain in the system if one or more parents do not execute their wait system call.

Zombie processes

❖ It the parent process terminates (without executing wait, and the child is still running, the latter is inherited by init the process (PID=1). The child does not become zombie because the system knows that no one is waiting for its exit status.

System call waitpid ()

- If a parent needs to wait a specific child it is better to use waitpid, which
 - > suspends execution of the calling process until a child, specified by *pid* argument, has changed state. By default, **waitpid()** waits only for terminated children.

System call waitpid ()

```
#include <sys/wait.h>

pid_t waitpid (
  pid_t pid,
  int *statLoc,
  int options);
```

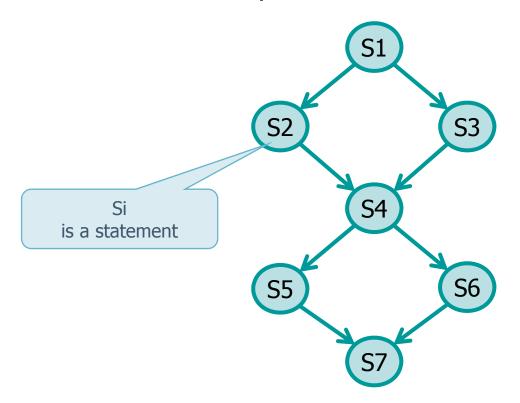
- The parameter pid allows waiting for
 - \triangleright Any child (waitpid==wait) (pid = -1)
 - \rightarrow The child whose PID=pid if >0 (pid >0)
 - Any child whose GID is equal to that of the calling process (pid = 0)
 - ➤ Any child whose GID=abs(pid) (pid < -1)

System call waitpid ()

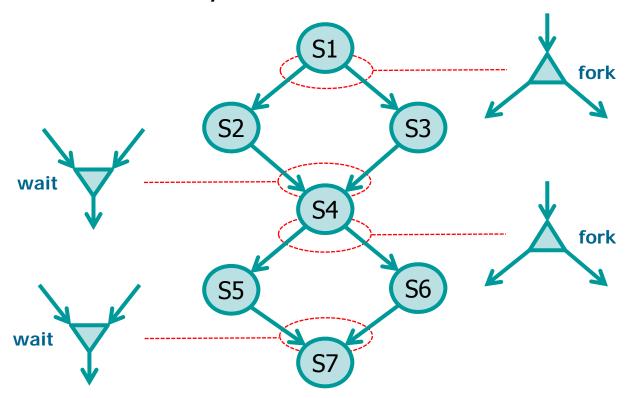
```
#include <sys/wait.h>
pid_t waitpid (
  pid_t pid,
  int *statLoc,
  int options);
```

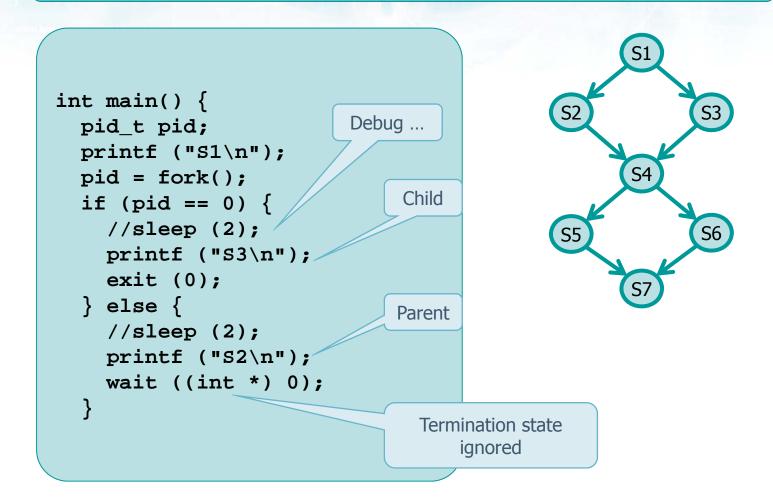
- The options parameter allow additional controls
 - > Default is 0, or is a bitwise OR of constants
 - wnohang, if the child specified by PID is running, the caller does not bock (not blocking version of wait)
 - WCONTINUED, etc.

Implement this Control Flow Graph (CFG) by means of the system calls fork and wait

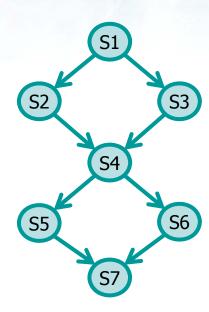


Implement this Control Flow Graph (CFG) by means of the system calls fork and wait

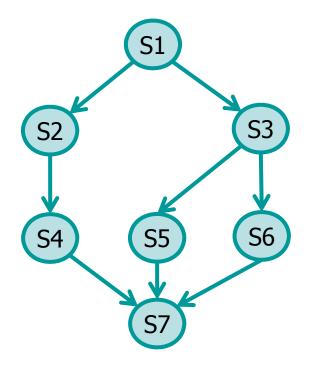




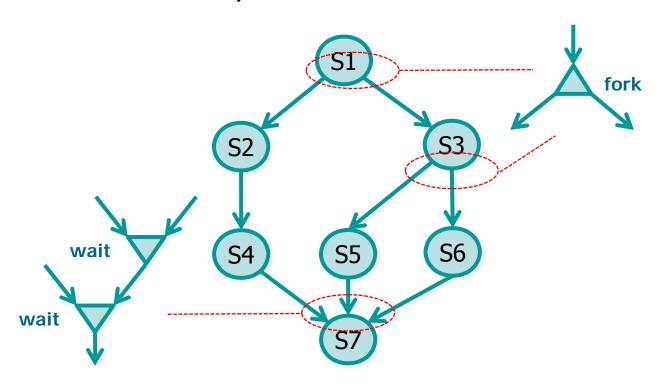
```
printf ("S4\n");
pid = fork();
if (pid == 0) {
  //sleep (2);
 printf ("S6\n");
  exit (0);
} else {
  //sleep (2);
 printf ("S5\n");
  wait ((int *) 0);
printf ("S7\n");
return (0);
```



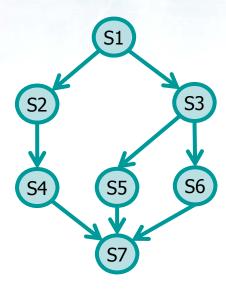
Implement this Control Flow Graph (CFG) by means of the system calls fork and wait



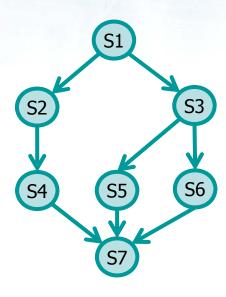
Implement this Control Flow Graph (CFG) by means of the system calls fork and wait



```
int main () {
 pid_t pid;
 printf ("S1\n");
 if ( (pid = fork())==-1 )
   err_sys( "can't fork" );
 if ( pid == 0 ){
   P356();
  } else {
   printf ("S2\n");
   printf ("S4\n");
   while (wait((int *)0)!= pid);
   printf ("S7\n");
   exit (0);
 return (1);
```



```
P356() {
 pid_t pid;
 printf ("S3\n");
  if ( ( pid = fork() ) == -1 )
    err_sys( "can't fork" );
  if (pid > 0 ){
   printf ("S5\n");
   while (wait((int *)0)!=pid );
  } else {
   printf ("S6\n");
    exit (0);
  exit (0);
```



Write a program that

- > Takes as argument an integer value n
- Allocates dynamically an integer vector of dimension n
- > Fills the vector with values reads from the terminal
- Displays the vector content, from the last to the first element, using n-1 processes, each displaying a single element of the vector

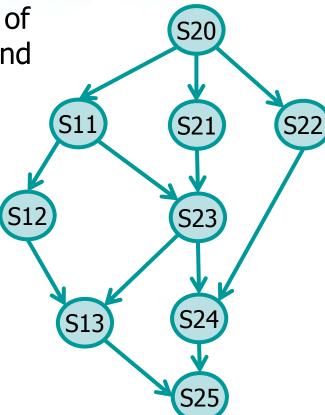
> Hint

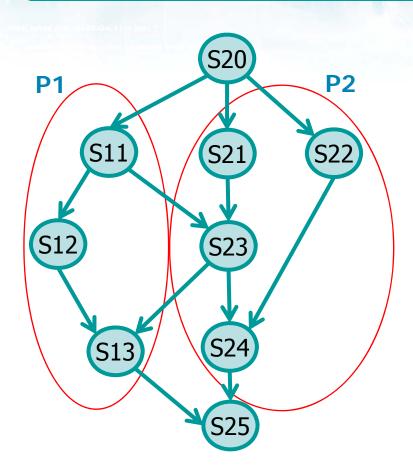
 Synchronize the processes by means of wait system calls

```
int main(int argc, char *argv[]) {
 int i, n, *vet;
 int retValue;
 pid t pid;
 n = atoi (argv[1]);
 vet = (int *) malloc (n * sizeof (int));
 if (vet==NULL) {
   fprintf (stderr, "Allocation Error.\n");
   exit (1);
 fprintf (stdout, "Input:\n");
 for (i=0; i<n; i++) {
   fprintf (stdout, "vet[%d]:", i);
   scanf ("%d", &vet[i]);
```

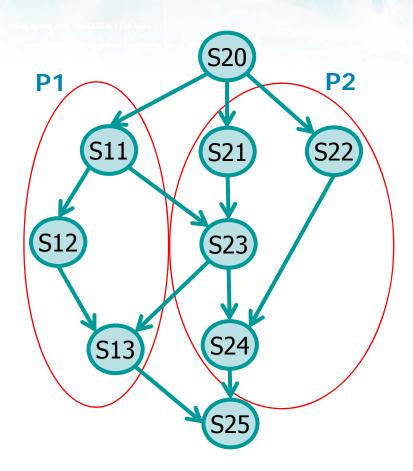
```
fprintf (stdout, "Output:\n");
for (i=0; i<n-1; i++) {
 pid = fork();
  if (pid>0) {
   pid = wait (&retValue);
   break;
  fprintf (stdout, "Run PID=%d\n", getpid());
fprintf (stdout, "vet[%d]:%d - ", i, vet[i]);
fprintf (stdout, "End PID=%d\n", getpid());
exit (0);
```

Implement this Control Flow Graph (CFG) by means of the system calls fork and wait

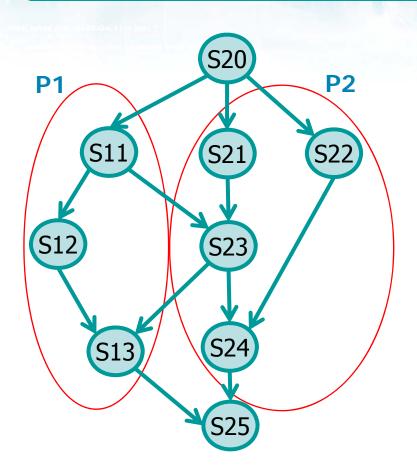




```
main () {
  S20 ();
  pid = fork ();
  if (pid>0) {
    P1 ();
    wait ((int *)0);
  } else {
    P2 ();
  S25 ();
  return;
```



```
P1() {
  S11 ();
 pid = fork ();
  if (pid>0) {
    S12 ();
    wait((int *)0);
  } else {
    ??? To P2 ???;
    exit(0);
  S13 ();
```



```
P2() {
  pid = fork ();
  if (pid>0) {
    S21 ();
    ??? From S1 ???;
    S23 ();
    wait((int *)0);
  } else {
    S22 ();
    exit(0);
  S24 ();
  exit (0);
```

Unfeasible

```
graph
P1() {
                                P2() {
  S11 ();
                                  pid = fork ();
  pid = fork ();
                                   if (pid>0) {
  if (pid>0) {
                                     S21 ();
    S12 ();
                                     ??? From S1 ???;
    wait((int *)0);
                                     s23 ();
  } else {
                                     wait((int *)0);
    ??? To P2 ??//; (S11)
                                     else {
    exit(0);
                                     S22 ();
                                     exit(0);
  s13 ();
                                   S24 ();
                                   exit (0);
```