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```

Processes

Signals

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Definition

- A signal is a software interrupt, i.e., an asynchronous notification sent, by the kernel or by another process, to a process to notify it of an event that occurred
 eg kernel
- Signals
 - > can be used as a limited form of inter-process communication
 - allow notify asynchronous events such as the error conditions illustrated for exceptions

1. not otherinformation2. may be lost

Definition

- ❖ A signal is a software interrupt, i.e., an asynchronous notification sent, by the kernel or by another process, to a process to notify it of an event that occurred
- Signals
 - can be used as a limited form of inter-process communication
 - allow notify asynchronous events such as the error conditions illustrated for exceptions

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Signals sent by the exception handlers

Exception	Exception handler	Signal
Divide error	divide_error()	SIGFPE
Debug	debug()	SIGTRAP
Breakpoint	int3()	SIGTRAP
Overflow	overflow()	SIGSEGV
Bounds check	bounds()	SIGSEGV
Invalid opcode	invalid_op()	SIGILL
Segment not present	segment_not_present()	SIGBUS
Stack segment fault	stack_segment()	SIGBUS
General protection	<pre>general_protection()</pre>	SIGSEGV
Page Fault	page_fault()	SIGSEGV
Intel-reserved	None	None
Floating-point error	<pre>coprocessor_error()</pre>	SIGFPE

Terminal signals

- Typing some key combinations at the controlling terminal of a process causes the system to send it a signal:
- Ctrl-C sends an SIGINT signal
 - > by default, this causes the process to terminate.
- Ctrl-Z sends a terminal stop SIGTSTP signal
 - by default, this causes the process to suspend execution.
- Ctrl-\ sends a SIGQUIT signal;
 - by default, this causes the process to terminate and dump core.

Main signals

Name	Description	
SIGABRT	Process abort, generated by system call abort	
SIGALRM	Alarm clock, generated by system call alarm	
SIGFPE	Floating-Point exception	
SIGILL	Illegal instruction	
SIGKILL	Kill (non maskable)	
SIGPIPE	Write on a pipe with no reader	
SIGSEGV	Invalid memory segment access	
SIGCHLD	Child process stopped or exited	
SIGUSR1 SIGUSR2	User-defined signal 1/2 for sync	

You can display the complete list of signals using the shell command kill -1

Operating Systems of continuous del Making:

Signal management

Signal management goes through three phases: signal generation, signal delivery, reaction to a signal

interrpt manage

Signal generation

When the kernel or a process causes an event that generate a signal
 exit(0)

> Signal delivery

- A not yet delivered signal remains pending
- At signal delivery a process executes the actions related to that signal
- The lifetime of a signal is from its generation to its delivery

Signal management

> Reaction to a signal

- To properly react to the asynchronous arrival of a given type of signal, a process must inform the kernel about the action that it will perform when it will receive a signal of that type
- A process may
 - Accept the default behavior (be terminated)
 - Declare to the kernel that it wants to ignore the signals of that type
 - Declare to the kernel that it wants to catch and manage the signals of that type by means of a signal handler function (similarly to the interrupt management)

Signal management

- Precondition to properly handle a received signal for a process is to declare to the kernel if a signal of a given type will be ignored or caught.
- This is done using the system call signal
 - Which instantiates a signal handler

signal does not send a signal!!

send by kill ()

Operating Systems

signal() system call

Arguments

tell what sig to do not send sig

- > sig indicates the type of signal to be caught
 - SIGALRM, SIGUSR1, etc.
- func specifices the address (i.e., pointer) to the function that will be executed when a signal of that type is received by the process
 - This function has a single argument of int type, which indicates the type of signal that will be handled

signal() system call

Return value

- the previous value of the signal handler, i.e., the previous signal handler function
- > SIG_ERR in case of error, errno is set to indicate the cause

Operating Systems

Signal generation

- The kernel generates signals
 - > SIGCHLD, SIGFPE, etc.
- A process can (ask the kernel to) generate a signal by means of the system call
 - kill (and raise) ____
 self kill

kill is misleading, does not kill a process, just send to it a signal

- > alarm
 - Ask the kernel to receive a SIGALRM after a specified amount of time

eg sleep()

Waiting for a signal

- A process can wait for a signal by means of the system call
 - > pause, and any other blocking system call
 - Suspend the process until any signal is received
 - > sleep
 - Suspend the process for a specified amount of time (waits for signal SIGALRM)

Reaction to a signal

- signal system call allows setting three different reactions to the delivery of a signal
 - > Accept the default behavior
 - signal (SIGname, SIG_DFL)
 - Where SIG_DFL is defined in signal.h
 - #define SIG_DFL ((void (*)()) 0
 - Every signal has its own default behavior, defined by the system
 - Most of the default reactions is process termination

Reaction to a signal

> Ignore signal delivery

- signal (SIGname, SIG_IGN)
- Where SIG_IGN is defined in signal.h
 - #define SIG_DFL ((void (*)()) 1
- Some signals cannot be ignored
 - **SIGKILL** and **SIGSTOP** cannot be ignored because the kernel and the superuser would not have the possibility to control all processes
 - Ignoring an illegal memory access, signaled by sigsegv, would produce an undefined process behavior

eg when using & send it to kernel to prevent zombie process

Reaction to a signal

Catch the signal

- signal (SIGname, signal Handler Function)
- where
 - SIGname indicates the signal type
 - signalHandlerFunction is the user defined signal handler function
- The signal handler
 - Is executed when the signal is delivered,
 - When it returns, the process continues with the next instruction, as it happens for interrupts

A signal handler function must be defined for every signal type that must be caugth

```
Signal handler for
#include <signal.h>
                                           signal SIGINT
#include <stdio.h>
#include <unistd.h>
void manager (int sig) {
  printf ("Received signal %d\n", sig);
  // signal (SIGINT, manager);
  return;
                                          Declares the signal
                                              handler
int main() {
  signal (SIGINT, manager);
  while (1) {
    printf ("main: Hello!\n");
    sleep (1);
```

```
Same signal handler
                                         for more than one
void manager (int sig) {
                                            signal type
  if (sig==SIGUSR1)
    printf ("Received SIGUSR1\n");
  else if (sig==SIGUSR2)
    printf ("Received SIGUSR2\n");
  else printf ("Received %d\n", sig);
  return;
                                         Both signal types
int main () {
                                         must be declared
  signal (SIGUSR1, manager);
  signal (SIGUSR2, manager);
```

```
Standard behavior of
         wait
                                                 i=2 PID=3057
                                                 i=1 PID=3056
// signal (SIGCHLD, SIG_IGN);
                                                 i=0 PID=3055
for (i=0; i<3; i++) {
  if (fork() == 0) {
    // child
    sleep (1);
    printf ("i=%d PID=%d\n", i, getpid());
    exit (i);
                                      Wait: ret = 3055 code = 0
                                      Wait: ret = 3056 code = 100
sleep (5);
                                      Wait: ret = 3057 code = 200
for (i=0; i<3; i++) {
  ret = wait (&code);
  printf ("Wait: ret=%d code=%x\n", ret, code);
```

```
Altering the behavior of
                                             Ignore SIGCHLD, sent
        wait
                                             by the kernel to the
                                             parent at the exit of a
signal (SIGCHLD, SIG IGN);
                                             child
for (i=0; i<3; i++) {
  if (fork() == 0) {
    // child
    sleep (1);
    printf ("i=%d PID=%d\n", i, getpid());
    exit (i);
                                        No wait:
                                        Wait: ret = -1 code = 7FFF
                                        Wait: ret = -1 code = 7FFF
sleep (5);
                                        Wait: ret = -1 code = 7FFF
for (i=0; i<3; i++) {
  ret = wait (&code);
  printf ("Wait: ret=%d code=%x\n", ret, code);
```

kill system call

```
#include <signal.h>
int kill (pid_t pid, int sig);
```

- Send signal (sig) to a process or to a group of processes (pid)
 - ➤ A **user** process can send signals only to processes having the same UID
 - > The **superuser** can send signal to any process

kill system call

```
#include <signal.h>
int kill (pid_t pid, int sig);
```

Arguments

If pid is	Send sig	
>0	To process with PID equal to pid	
==0	To all processes with GID equal to its GID	
<0	To all processes with GID equal to the pid	absolute value of "All process" excludes
==-1	To all processes	a set of system processes

kill system call

```
#include <signal.h>
int kill (pid_t pid, int sig);
```

- Return value
 - > 0 on success
 - \geq -1 on error

raise system call

```
#include <signal.h>
int raise (int sig);
```

- The raise system call allows a process to send a signal to itself
 - raise (sig) is equivalent to
 - kill (getpid(), sig)

Operating Systems

pause system call

```
#include <unistd.h>
int pause (void);
```

- Suspends the calling process until a signal is delivered
- Returns after the completion of the signal handler
 - > returns -1, and errno is set to EINTR

kill -9

alarm system call

```
#include <unistd.h>
unsigned int alarm (unsigned int seconds);
```

- Ask the kernel so send a SIGALRM to the calling process in seconds
 - ➤ The default action for SIGALRM is process termination
 - ➤ A call to alarm(seconds) before the expiration of a previous alarm reschedules the request to the kernel
 - alarm(0) cancels any pending alarm

alarm system call

```
#include <unistd.h>
unsigned int alarm (unsigned int seconds);
```

Return value

- ➤ the number of seconds remaining until the delivery of a previously scheduled alarm
- > zero if there was no previously scheduled alarm.

alarm system call

```
#include <unistd.h>
unsigned int alarm (unsigned int seconds);
```

Warning

- > The signal is generated by the kernel
 - It is possible that the process get the CPU control after some time, depending on the scheduler decisions
- System calls sleep and alarm uses the same kernel timer

Implement system call sleep using system calls alarm and pause

```
include <signal.h>
#include <unistd.h>

static void sig_alrm(int signo) {return;}

unsigned int sleep1(unsigned int nsecs)
{
  if (signal(SIGALRM, sig_alrm) == SIG_ERR)
    return (nsecs);
  alarm (nsecs);
  pause ();
  return (alarm(0));
}

Returns 0, or the remaining time before the delivery if pause returns because another signal has been received
```

Implement system call alarm using system calls fork, signal, kill and pause

```
#include <stdio.h>
#include <unistd.h>
#include <signal.h>

void myAlarm (int sig) {
   printf ("Alarm\n");
}
```

```
int main (void) {
 pid_t pid;
  (void) signal (SIGALRM, myAlarm);
 pid = fork();
 switch (pid) {
    case -1: /* error /*
                                       The child waits
      printf ("fork failed");
                                        and sends
      exit (1);
                                        SIGALRM
    case 0: /* child */
      sleep(5);
      kill (getppid(), SIGALRM);
      exit(0);
                           The parent pauses, and continues
  /* parent */
 pause ();
                          only when it receives the SIGALRM
 exit (0);
                                  sent by the child
```

Signal limitations

- Signals do not convey any information
- The memory of the pending signals is limited
 - Max one signal pending per type
 - Forthcoming signals of the same type are lost
 - > Signals can be ignored
- Signal interrupt functions that must be reentrant
- Produce race conditions
- Some limitations are avoided in POSIX.4

Limited memory

```
static void sigUsr1 (int);
static void sigUsr2 (int);
static void
sigUsr1 (int signo) {
  if (signo == SIGUSR1)
    printf("Received SIGUSR1\n");
  else
   printf("Received wrong SIGNAL\n");
  fprintf (stdout, "sigUsr1 sleeping ...\n");
  sleep (5);
  fprintf (stdout, "... sigUsr1 end sleeping.\n");
 return;
```

Limited memory

```
static void
sigUsr2 (int signo) {
  if (signo == SIGUSR2)
    printf("Received SIGUSR2\n");
  else
    printf("Received wrong SIGNAL\n");

fprintf (stdout, "sigUsr2 sleeping ...\n");
  sleep (5);
  fprintf (stdout, "... sigUsr2 end sleeping.\n");

return;
}
```

Limited memory

```
int
main (void) {
  if (signal(SIGUSR1, sigUsr1) == SIG ERR) {
    fprintf (stderr, "Signal Handler Error.\n");
    return (1);
  if (signal(SIGUSR2, sigUsr2) == SIG ERR) {
    fprintf (stderr, "Signal Handler Error.\n");
    return (1);
  while (1) {
    fprintf (stdout, "Before pause.\n");
    pause ();
    fprintf (stdout, "After pause.\n");
  return (0);
                                  The main iterates waiting
                                  signals from shell
```

Limited memory

Shell commands

- > ./pgrm &
- [3] 2636
- > Before pause.
- > kill -USR1 2636
- > Received SIGUSR1
- sigUsr1 sleeping ...
- ... sigUsr1 end sleeping.

After pause.

Before pause.

- > kill -USR2 2636
- > Received SIGUSR2
- sigUsr2 sleeping ...
- ... sigUsr2 end sleeping.

After pause.

Before pause.

Correctly received SIGUSR1

Correctly received SIGUSR2

Observation:

shell command **kill** sends a signal to a process with a specified PID

Limited memory

Two signals sent in sequence

```
> kill -USR1 2636 ; kill -USR2 2636
> Received SIGUSR2
sigUsr2 sleeping ...
... sigUsr2 end sleeping.
Received SIGUSR1
sigUsr1 sleeping ...
... sigUsr1 end sleeping.
After pause.
Before pause.
```

Both are received

The deliver order of the two signal cannot be predicted

Limited memory

```
> kill -USR1 2636 ; kill -USR2 2636 ; kill -USR1 2636
> Received SIGUSR1
sigUsr1 sleeping ...
                                        Three signals sent in
... sigUsr1 end sleeping.
Received SIGUSR2
                                       sequence: two of them
sigUsr2 sleeping ...
                                           are SIGUSR1
... sigUsr2 end sleeping.
After pause.
Before pause.
                                          A SIGUSR1 is lost
> kill -9 2636
              ./pgrm
[3]+ Killed
                             -9 = SIGKILL = Kill
```

Kill a process

Reentrant functions

- The kernel knows where a signal handler returns, but
- The signal handler does not know where it was called, i.e., the control flow was interrupted by the signal

Reentrant functions: Examples

- Suppose a malloc is interrupted, and the signal handler calls another malloc
 - Function malloc manages the list of the free memory regions, which could be corrupted
- Suppose that the execution of a function that uses a static variable is interrupted, but is then called by the signal handler
 - The static variable could be used to store a new value, i.e., it does not remain the same it was before the signal was delivered

Reentrant functions: Conclusions

- The Single UNIX Specification defines the reentrant functions, which can be interrupted without problems
 - > read, write, sleep, wait, etc.
- Most of the I/O standard C are not reentrant
 - > printf, scanf, etc.
 - They use static variables or global variables, thus must be used carefully and being aware of possible problems

Race conditions

Race condition

- ➤ The result of more concurrent processes working on common data depends on the execution order of the processes instructions
- Using signals increases the probability of race conditions.

Race conditions example

- Suppose a process decides to suspend itself for a given number of seconds
- The signal could be delivered before the execution of pause due to a contest switching and scheduling decisions.

```
static void
myHandler (int signo) {
    ...
signal (SIGALARM, myHandler)
alarm (nSec);
pause ();

static void
MyHandler (int signo) {
    delivered before pause
    pause blocks the process
forever because the signal has
    been lost
```

Race conditions example

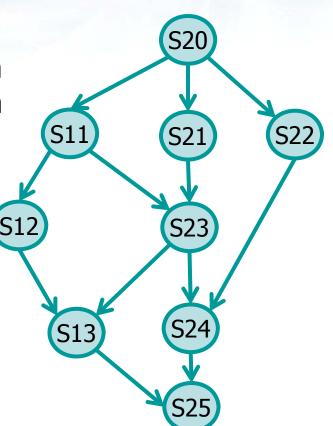
- Suppose two processes P₁ and P₂ decide to synchronize by means of signals
 - ➤ If P₁ (P₂) signal is delivered before P₂ (P₁) executes pause
 - ➤ Process P₂ (P₁) blocks forever waiting a signal

```
P<sub>1</sub>
while (1) {
    ...
    kill (pidP2, SIG...);
    pause ();
}
P<sub>2</sub>
while (1) {
    pause ();
    kill (pidP1, SIG...);
}
```

Exercise

Despite their shortcomings, signals can provide a rough synchronization mechanism

Ignoring the race
conditions (and using
fork, wait, signal,
kill, and pause)
implement this precedence
graph

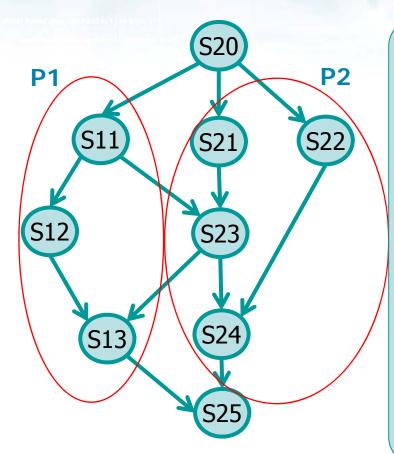


```
static void
sigUsr ( int signo) {
  if (signo== SIGUSR1)
    printf ("SIGUSR1\n");
  else if (signo==SIGUSR2)
    printf ("SIGUSR2\n");
  else
    printf ("Signal %d\n", signo);
  return;
}
```

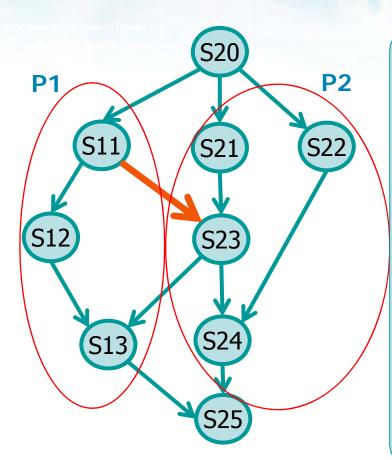
```
int main (void) {
  pid_t pid;

if (signal(SIGUSR1, sigUsr) == SIG_ERR) {
    printf ("Signal Handler Error.\n");
    return (1);
}

if (signal(SIGUSR2, sigUsr) == SIG_ERR) {
    printf ("Signal Handler Error.\n");
    return (1);
}
```

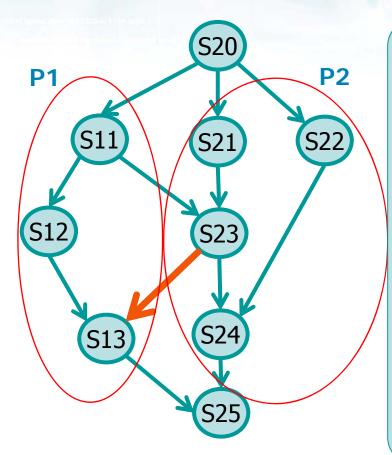


```
printf ("S20\n");
pid = fork ();
if (pid > (pid_t) 0) {
   P1 (pid);
   wait ((int *) 0);
} else {
   P2 ();
   exit (0);
}
printf ("S25\n");
return (0);
}
```



```
void P1 (
   pid_t cpid
) {
   printf ("S11\n");
   sleep (1); //!?
   kill (cpid, SIGUSR1);
   printf ("S12\n");
   pause ();
   printf ("S13\n");

   return;
}
```



```
void P2 ( ){
  if (fork () > 0) {
    printf ("S21\n");
    pause ();
    printf ("S23\n");
    kill (getppid (),
              SIGUSR2);
    wait ((int *) 0);
  } else {
    printf ("S22\n");
    exit (0);
  printf ("S24\n");
  return;
```