```
#include ciding.bb
#include cidi
```

Processes

Process creation

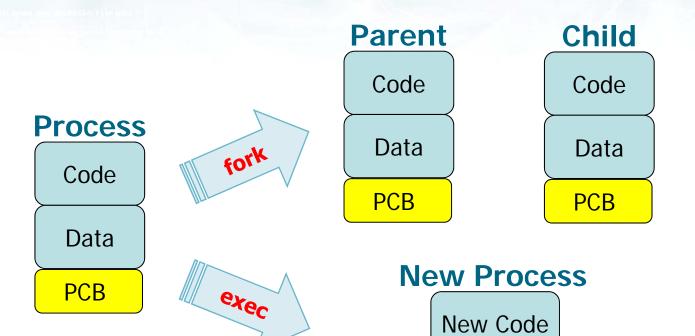
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fork and exec system calls

- System call fork creates a new process duplicating the calling process, then
 - ➤ Parent and child execute **different code sections**Example: a network server duplicates itself at each client request, and the child serves the request while the parent waits for a new client request
 - > Parent and child execute different code
 - Example: a command interpreter (shell)
 - Uses the exec system call

- System call exec substitutes the process code with the executable code of another program
- The new program begins its execution as usual (from main)
- In particular exec
 - Does not create a new process
 - > Substitutes the calling process image with the image of another program.
 - > The process PID does not change
 - fork → duplicates an existent process
 - exec → executes a new program

Address space



New Data

PCB

- 6 versions of exec system call
 - > execl, execlp, execle
 - > execv, execvp, execve

Туре	Action
I (list)	Arguments are a list of strings
v (vector)	Arguments is a vector of strings char **arguments
p (path)	The executable filename is looked for in the directories listed in the environment variable PATH
e (environment)	The last argument is an environment vector envp[] which defines a set of new associations strings name=value

```
#include <unistd.h>
int execl (char *path, char *arg0, ..., (char *)0);
int execlp (char *name, char *arg0, ..., (char *)0);
int execle(char *path, const char *arg0,..., char *envp[]);
int execv (char *path, char *argv[]);
int execvp (char *name, char *arg[]);
int execve (char *path, char *arg[], char *envp[]);
```

❖ The return value is -1 in case of error

Arguments

- > Pathname of the executable file
 - In the "p" versions the complete path is not necessary. The file must be in one of the directories listed in the environment variable PATH (echo \$PATH)
- > Its argument list
 - The first argument is the name of the process (its argv[0)]
 - The other argument of the list are the argument for the executable (argv[1, argv[2], etc).
- Possibly the environment vector

Examples

```
whereis cp: /bin/cp
                                      User defined name
OK
 execl("/bin/cp", "mycp", "./file1", "./file2", NULL);
                                                   Alternative
OK
                                                   termination
 execl("/bin/cp","cp","./file1","./file2",(char*)0);
           Path is missing
NO
 execl("cp","File_copy","./file1","./file2",(char*)0);
OK
             Default path ($PATH)
 execlp("cp", "cp", "./file1", "./file2", (char*)0);
```

- **⋄** exec**v**[p]
 - ➤ Uses a single argument: a vector of strings
 - The vector must be properly initialized

System call exec ()

- exec[lv]e
 - Can provide to the executable a set of environment variables
 - Vector of strings
 - Without "e" the environment of the new process is inherited from the calling process

```
char *env[] = {
   "USER=unknown",
   "PATH=/tmp",
   NULL
};
...
execle (path, arg0, ..., argn, 0, env);
...
execve (path, argv, env);
```

Considerations

- The execed process keeps all open file descriptor (including stdin, stdout, stderr)
 - This allow the process to inherit possible redirections previously set (e.g., by shell)
- Many kernel implement only system call execve
 - ➤ The other versions are macros that use this system call

- Draw the process generation tree of the following C code segment, executed passing as its argument on the command line string "5"
- What does it display?
- Why?

Run with n=5

```
#include <stdio.h>
#include <unistd.h>
int main (int argc, char ** argv) {
 char str[10];
 int n;
 n = atoi(argv[1]) - 1;
 printf ("%d\n", n);
 if (n>0) {
   sprintf (str, "%d", n);
   execl (argv[0], argv[0], str, NULL);
 printf ("End!\n");
 return 1;
```

Solution

```
P(5)
          n=4; printf 4
          exec
P(4)
          n=3; printf 3
          exec
P(3)
          n=2; printf 2
           exec
P(2)
          n=1; printf 1
           exec
P(1)
          n=0; printf 0
          printf End!
```

```
int main (int argc, char ** argv) {
  char str[10];
  int n;
  n = atoi(argv[1]) - 1;
  printf ("%d\n", n);
  if (n>0) {
    sprintf (str, "%d", n);
    execl (argv[0], argv[0], str, NULL);
  }
  printf ("End!\n");
  return 1;
}
```

Output

```
4
3
2
1
0
End!
```

- Draw the process generation tree of the following C code segment
- What does it display?
- Why?

```
#include <stdio.h>
#include <unistd.h>
int main(){
                                        shell command
  int n;
                                           echo
 n=0;
 while (n<3 && fork()){
    if (!fork())
      execlp ("echo", "n++", "n", NULL);
    n++;
    printf ("%d\n", n);
 return (1);
```

Example

Program ./pgrm exec itself if it receives as argument 1 or 2

```
n = atoi (argv[1]);
switch (n) {
  case 1:
    printf("#1:PID=%d;PPID=%d\n", getpid(), getppid());
    sleep (n*10);
    execlp ("./pgrm", "./Pgrm", "2", (char *) 0);
    break:
  case 2:
    printf("#2:PID=%d;PPID=%d\n", getpid(), getppid());
    sleep (n*10);
    execlp ("./pgrm", "myPgrm", "3", (char *) 0);
    break;
  default:
    printf("#3:PID=%d;PPID=%d\n", getpid(), getppid());
    sleep (n*10);
    break;
                                       Same pathname but
return (1);
                                     argv[0] (its name) changes
```

Example

```
Run with n=1
    Shell commands (in blue)
                           PID does not change
> ./pgrm 1 &
[2] 2471
#1: PID=2471; PPID=2045
> ps -aux | grep 2471
user 2471 0.0 0.0 4192 352 pts/2 S 19:29 0:00 ./pgrm 1
#2: PID=2471; PPID=2045
> ps -aux | grep 2471
user 2471 0.0 0.0 4192 356 pts/2 S 19:29 0:00 ./Pgrm 2
#3: PID=2471; PPID=2045
> ps -aux | grep 2471
User 2471 0.0 0.0 4192 356 pts/2 S 19:29 0:00 myPgrm 3
[2]+ Exit 1 ./pgrm 1
                                   The name changes
```

UNIX shell skeleton

- Command run in foreground
 - > <command>

```
while (TRUE) {
  write_prompt;
  read_command (command, parameters);
  if (fork() == 0)
    /* Child: Execute command */
    execve (command, parameters);
else
    /* Parent: Wait child */
    wait (&status);
}
```

UNIX shell skeleton

- Command run in background
 - < <command> &

```
while (TRUE) {
  write_prompt;
  read_command (command, parameters);
  if (fork() == 0)
    /* Child: Execute command */
    execve (command, parameters);
/* else */
    /* Parent: does not wait */
    /* wait (&status); */
}
```

Command execution

- It can be useful to execute a shell command from a process
 - ➤ For example for appending a date to a filename or to a file
- System call system solves this problem
 - ➤ It is defined by the standard ISO C and POSIX

system() system call

```
#include <stdlib.h>
int system (const char *string);
```

- System call system
 - ➤ Forks a shell, which execute the string command, while the parent process waits the termination of the shell command
 - > Returns
 - -1 or 127 on error
 - The exit value of the shell that executed the command (with the format of waitpid)

Example

```
system ("date");
...
```

```
system ("ls -laR");
```

```
char str[L];
...
strcpy (str, "ls -la");
system (str);
...
```

system() implementation

```
int system (const char *cmd) {
 pid_t pid;
 int status;
 if (cmd == NULL)
   return(1);
 if ( (pid = fork()) < 0) {
    status = -1;
  } else if (pid == 0) {
   execl("/bin/sh", "sh", "-c", cmd, (char *) 0);
   exit(127);
  } else {
   while (waitpid (pid, &status, 0) < 0)</pre>
      if (errno != EINTR) {
        status = -1;
        break;
 return(status);
```

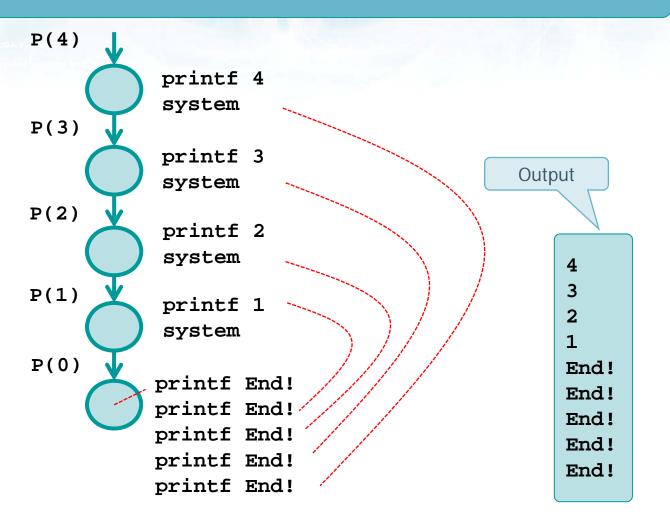
- Draw the process generation tree of the following C program, executed passing as its argument on the command line string "4"
- What does it display?
- Why?

Esercizio

Run with n=4

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
int main(int argc, char ** argv){
  int n;
 char str[10];
 n = atoi (argv[1]);
 if (n>0) {
   printf ("%d\n", n);
   sprintf (str, "%s %d", argv[0], n-1);
   system (str);
 printf("End!\n");
 return (1);
```

Solution



- Draw the process generation tree of the following C code segment
- What does it display?
- Why?

```
#include ...
int main () {
 char str[100];
  int i;
 for (i=0; i<2; i++){
    if (fork()!=0) {
      sprintf (str, "echo system with i=%d", i);
      system (str);
    } else {
      if (fork()==0) {
        sprintf (str, "exec with i=%d", i);
        execlp ("echo", "myPgrm", str, NULL);
 return (0);
```

