

اعمال ریاضی با اعداد

اعمال ریاضی باینری: جمع

• قوانین: مانند جمع دسیمال
 • با این تفاوت که $1+1 = 10$ ← تولید نقلی

$$\begin{array}{r}
 \overset{+1}{1} \overset{+1}{0} \overset{+1}{1} \overset{+1}{1} \overset{+1}{0} \overset{+1}{1} \\
 + 011001 \\
 \hline
 1000110
 \end{array}$$

0+0 = 0(c0) (sum 0 with carry 0) \swarrow

0+1 = 1+0 = 1(c0) \swarrow

1+1 = 0(c1) \swarrow

1+1+1 = 1(c1) \swarrow

Carry	1	1	1	1	1	0
Augend	0	0	1	0	0	1
Addend	0	1	1	1	1	1
Result	1	0	1	0	0	0

سرریز (Overflow)

◀ اگر تعداد بیت ها $n =$ و حاصل جمع $n+1$ بیت نیاز داشته باشد

- ← سرریز

اعمال ریاضی باینری: تفریق

• قوانین:

$0-0 = 1-1 = 0$ (b0) (result 0 with borrow 0) \swarrow

$1-0 = 1$ (b0) \swarrow

$0-1 = 1$ (b1) \swarrow

... \swarrow

X 229
 Y - 46

 $X - Y$ 183

$\begin{array}{r}
010111010 \\
111001101 \\
-001011110 \\
\hline
10110111
\end{array}$

Borrow	1	1	0	0	
Minuend	1	1	0	1	1
Subtrahend	0	1	1	0	1
Result	0	1	1	1	0

روش انجام محاسبات

➤ الگوریتم های اعمال ریاضی مبنای 10 را به خاطر آورید.

➤ آنها را برای مبنای مورد نظر تعمیم دهید.

➤ قانون مبنای مورد نظر را به کار ببرید.

- برای باینری: $1+1=10$

نمایش اعداد

- نمایش اعداد مثبت:

➤ در بیشتر سیستم ها یکسان است.

- نمایش اعداد منفی:

➤ اندازه-علامت (Sign magnitude)

➤ مکمل 1 (Ones complement)

➤ مکمل 2 (Twos complement)

- در بیشتر سیستم ها: مکمل 2

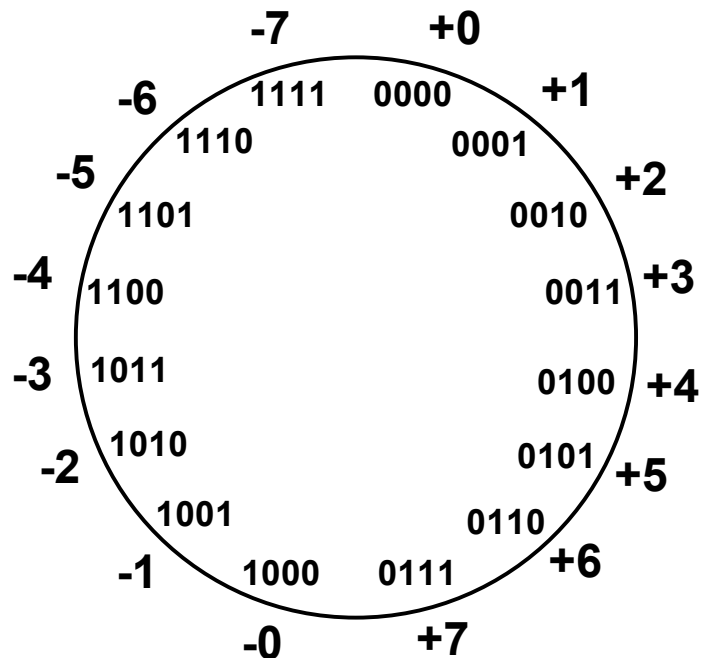
- فرض:

➤ ماشین با کلمه های 4 بیتی:

- ← 16 مقدار مختلف قابل نمایش.

- تقریباً نیمی مثبت، نیمی منفی.

نمایش اعداد



اندازه-علامت:

$\begin{matrix} + \\ \swarrow \\ 0\ 100 = +4 \end{matrix}$
 $\begin{matrix} \searrow \\ 1\ 100 = -4 \\ - \end{matrix}$

High order bit is sign: 0 = positive (or zero), 1 = negative

Three low order bits is the magnitude: 0 (000) thru 7 (111)

Number range for n bits = $[-(2^{n-1} - 1), +(2^{n-1} - 1)]$

Representations for 0

Cumbersome addition/subtraction

Must compare magnitudes to determine sign of result

نمایش اعداد

مکمل 1:

N is positive number, then \bar{N} is its negative 1's complement

$$\bar{N} = (2^n - 1) - N$$

Example: 1's complement of 7

$$2^4 = 10000$$

$$-1 = \underline{00001}$$

$$1111$$

$$-7 = \underline{0111}$$

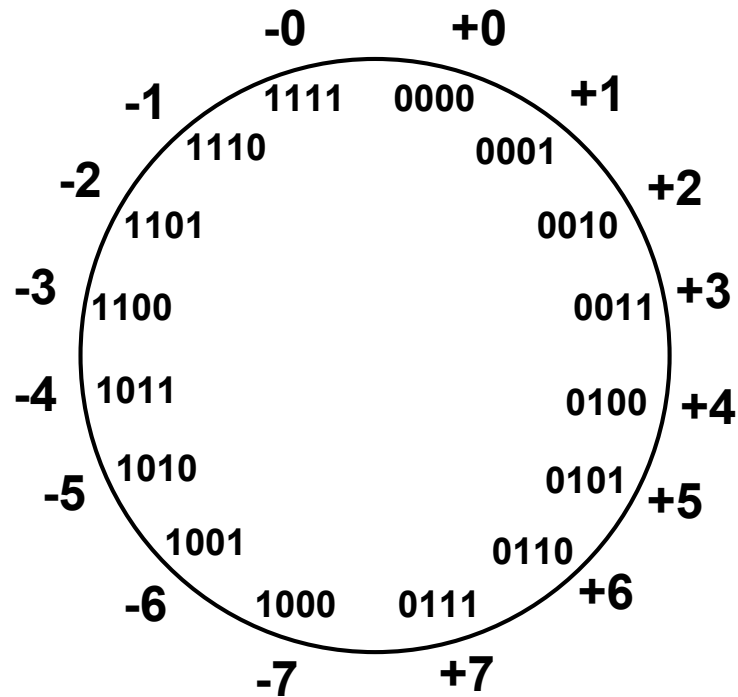
$$1000 = -7 \text{ in 1's comp.}$$

Shortcut method:

simply compute bitwise complement

$$0111 \rightarrow 1000$$

نمایش اعداد



مکمل 1: +
 $0\ 100 = +4$
 $1\ 011 = -4$
 -

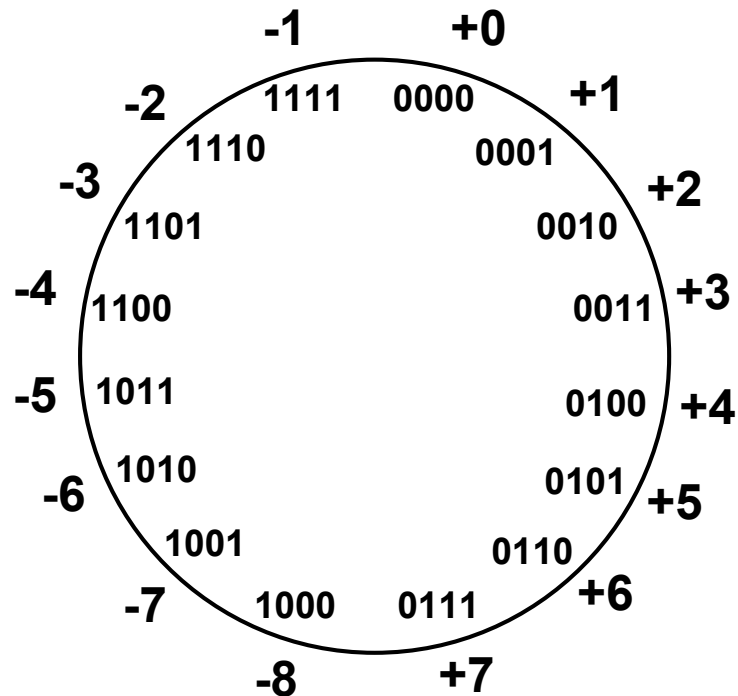
Subtraction implemented by addition & 1's complement

Still two representations of 0! This causes some problems

Some complexities in addition

نمایش اعداد

*like 1's comp
except shifted
one position
clockwise*



مکمل 2:

0 100 = + 4

1 100 = - 4

Only one representation for 0

One more negative number than positive number

نمایش اعداد

$$N^* = 2^n - N$$

Example: Twos complement of 7

$$\begin{array}{r} 2^4 = 10000 \\ \text{sub } 7 = \underline{0111} \\ \hline 1001 = \text{repr. of } -7 \end{array}$$

مکمل 2:

Example: Twos complement of -7

$$\begin{array}{r} 2^4 = 10000 \\ \text{sub } -7 = \underline{1001} \\ \hline 0111 = \text{repr. of } 7 \end{array}$$

Shortcut method:

Twos complement = bitwise complement + 1

0111 -> 1000 + 1 -> 1001 (representation of -7)

1001 -> 0110 + 1 -> 0111 (representation of 7)

مکمل 2

- Here's an easier way to compute the 2's complement:

1. Leave all least significant 0's and first 1 unchanged.
2. Replace 0 with 1 and 1 with 0 in all remaining higher significant bits.

Examples:

complement | unchanged
■ N = 1010
 0110
 └─┬─┘
 2's complement

complement | unchanged
N = 01011000
 10101000
 └─┬─┘
 2's complement

جمع و تفریق مکمل 2

If
(carry-in to sign = carry-out)
then ignore carry

if
(carry-in \neq carry-out)
then overflow

$$\begin{array}{r} 4 \quad 0100 \\ + 3 \quad 0011 \\ \hline 7 \quad 0111 \end{array}$$

$$\begin{array}{r} -4 \quad 1100 \\ + (-3) \quad 1101 \\ \hline -7 \quad 11001 \end{array}$$

$$\begin{array}{r} 4 \quad 0100 \\ - 3 \quad 1101 \\ \hline 1 \quad 10001 \end{array}$$

$$\begin{array}{r} -4 \quad 1100 \\ + 3 \quad 0011 \\ \hline -1 \quad 1111 \end{array}$$

Simpler addition scheme makes twos complement the most common choice for integer number systems

جمع و تفریق مکمل 2

Why can the carry-out be ignored?

$-M + N$ when $N > M$:

$$M^* + N = (2^n - M) + N = 2^n + (N - M)$$

Ignoring carry-out is just like subtracting 2^n

$-M + -N$ where $N + M < \text{or} = 2^{n-1}$

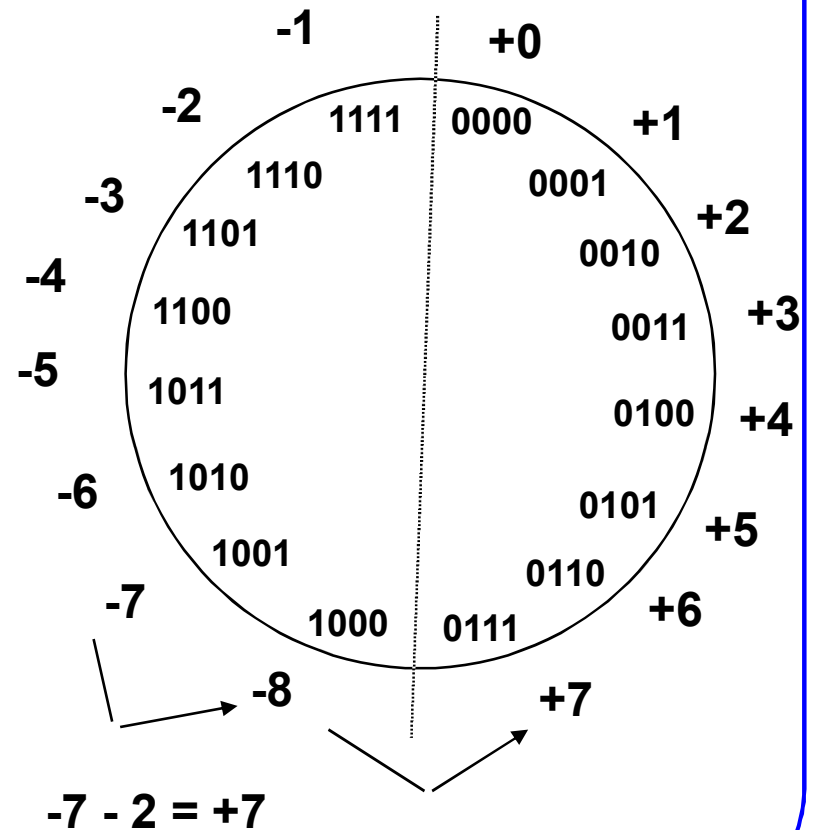
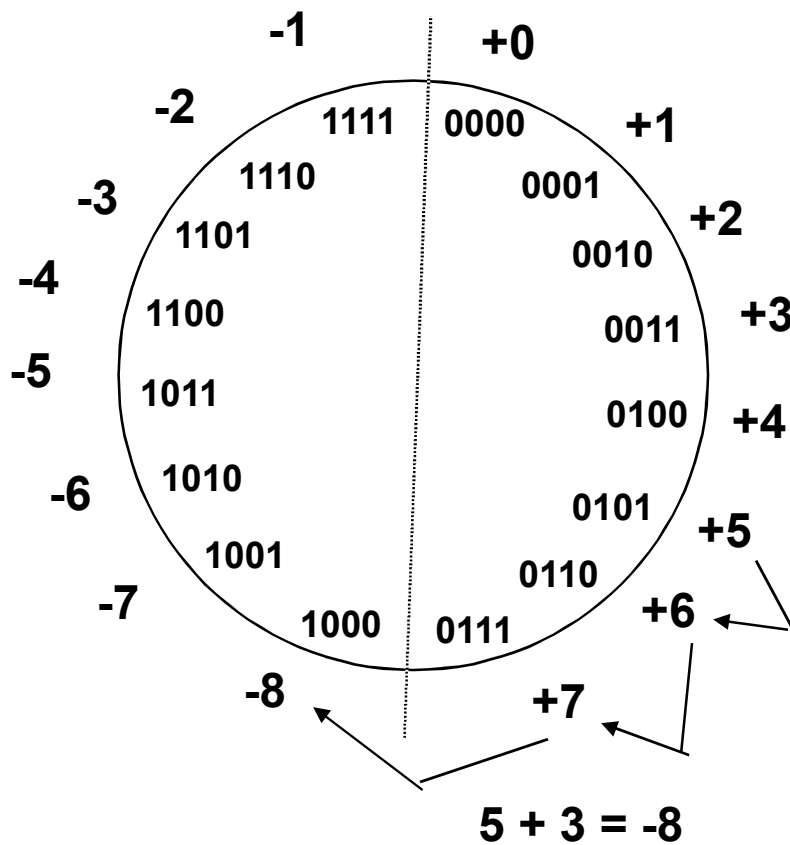
$$\begin{aligned} -M + (-N) &= M^* + N^* = (2^n - M) + (2^n - N) \\ &= 2^n - (M + N) + 2^n \end{aligned}$$

After ignoring the carry, this is just the right twos compl.
representation for $-(M + N)$
or $(M+N)^*$

Overflow Conditions

سریز

Add two positive numbers to get a negative number
or two negative numbers to get a positive number



سریز

Overflow Conditions

$$\begin{array}{r}
 5 \quad \quad 0111 \\
 \quad \quad 0101 \\
 \hline
 3 \quad \quad 0011 \\
 \hline
 -8 \quad \quad 1000
 \end{array}$$

Overflow

$$\begin{array}{r}
 5 \quad \quad 0000 \\
 \quad \quad 0101 \\
 \hline
 2 \quad \quad 0010 \\
 \hline
 7 \quad \quad 0111
 \end{array}$$

No overflow

$$\begin{array}{r}
 -7 \quad \quad 1000 \\
 \quad \quad 1001 \\
 \hline
 -2 \quad \quad 1110 \\
 \hline
 7 \quad \quad 10111
 \end{array}$$

Overflow

$$\begin{array}{r}
 -3 \quad \quad 1111 \\
 \quad \quad 1101 \\
 \hline
 -5 \quad \quad 1011 \\
 \hline
 -8 \quad \quad 11000
 \end{array}$$

No overflow

Method 1: Overflow when carry in to sign \neq carry out

Method 2: Overflow when $\text{sign}(A) = \text{sign}(B) \neq \text{sign}(\text{result})$

The Binary Multiplication

The diagram illustrates the binary multiplication process. It shows the multiplication of the binary number 101010 by 1011. The multiplier 1011 is shown with its rightmost digit '1' highlighted in a red box. Below the multiplication line, the first partial product is 101010, which is also highlighted in a red box. Below this, the other partial products are shown as 101010 (shaded gray) and 000000. These are then summed to produce the final result 11100110.

$$\begin{array}{r} 101010 \\ \times \quad 1011 \\ \hline 101010 \\ 101010 \\ 000000 \\ + 101010 \\ \hline 11100110 \end{array}$$

Partial Products

ضرب باینری

- Shift-and-add algorithm, as in base 10

M'cand	0	0	0	1	1	0	1
M'plier	0	0	0	0	1	1	0
(1)			0	0	0	0	0
(2)		0	1	1	0	1	
(3)	0	1	1	0	1		
Sum	1	0	0	1	1	1	0

- Check: $13 * 6 = 78$

Binary-Coded Decimal (BCD)

- ◀ A decimal code:
Decimal numbers (0..9)
are coded using 4-bit
distinct binary words
- ◀ Observe that the codes
1010 .. 1111 (decimal
10..15) are NOT
represented (invalid
BCD codes)

□ **TABLE 1-3**
Binary-Coded Decimal (BCD)

Decimal Symbol	BCD Digit
0	0000
1	0001
2	0010
3	0011
4	0100
5	0101
6	0110
7	0111
8	1000
9	1001

Table 1-3 Binary-Coded Decimal (BCD)

Binary-Coded Decimal

- To code a number with n decimal digits, we need $4n$ bits in BCD

e.g. $(365)_{10} = (0011\ 0110\ 0101)_{\text{BCD}}$



- This is different from converting to binary, which is $(365)_{10} = (101101101)_2$
- Clearly, BCD requires more bits. BUT, it is easier to understand/interpret

Application



BCD Addition

Case 1:

$$\begin{array}{r} 0001 \quad 1 \\ 0101 \quad 5 \\ \hline (0) 0110 \quad (0) 6 \end{array}$$

Case 2:

$$\begin{array}{r} 0110 \quad 6 \\ 0101 \quad 5 \\ \hline (0) 1011 \quad (1) 1 \end{array}$$

WRONG!

Case 3:

$$\begin{array}{r} 1000 \quad 8 \\ 1001 \quad 9 \\ \hline (1) 0001 \quad (1) 7 \end{array}$$

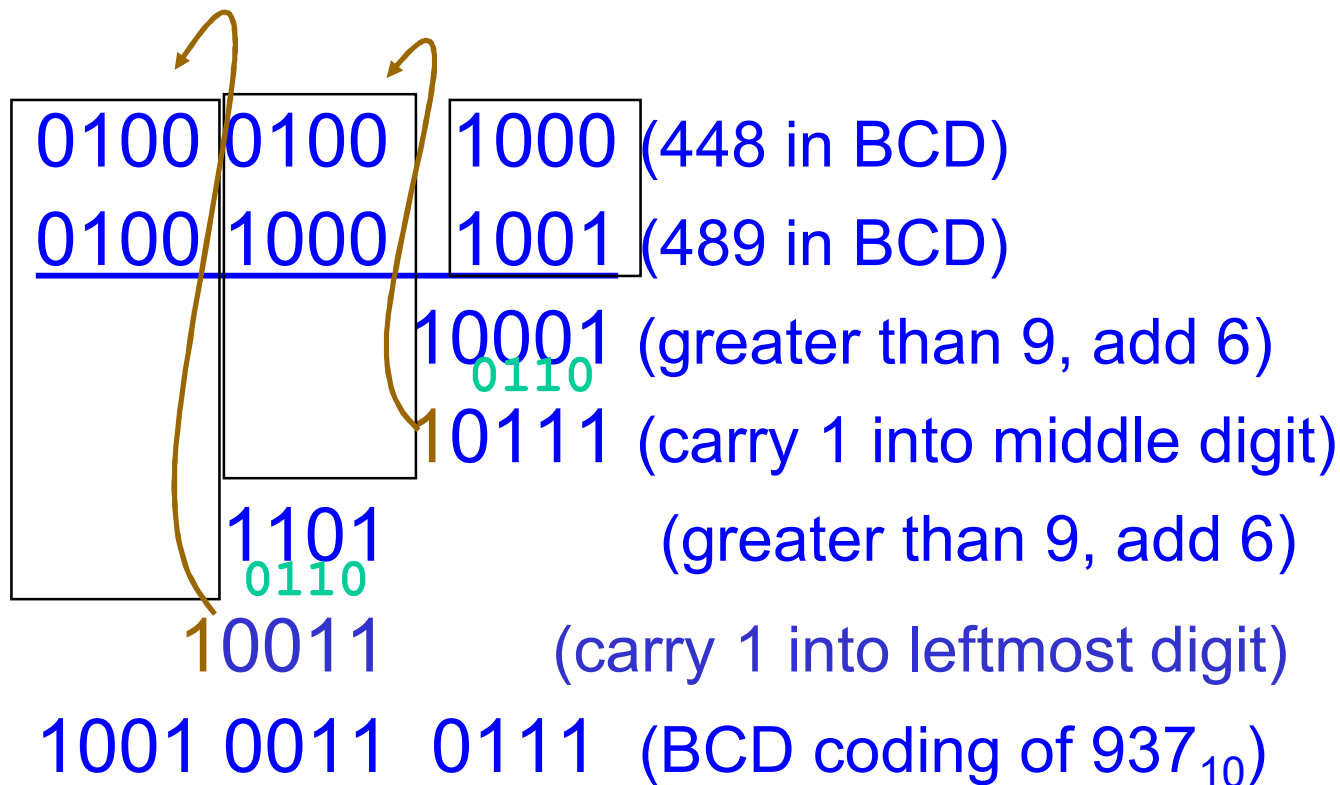
Note that for cases 2 and 3, adding a factor of 6 (0110) gives us the correct result.

BCD Addition (cont.)

- **BCD addition is therefore performed as follows**
 - 1) Add the two BCD digits together using normal binary addition
 - 2) Check if correction is needed
 - a) 4-bit sum is in range of 1010 to 1111
 - b) carry out of MSB = 1
 - 3) If correction is required, add 0110 to 4-bit sum to get the correct result;
→ BCD carry out = 1

BCD Addition (cont.)

- **Example: Add 448 and 489 in BCD.**



BCD Negative Number Representation

- **Similar to binary negative number representation but for $r = 10$.**

- ◀ **BCD 9's complement**

- invert each BCD digit (0 → 9, 1 → 8, 2 → 7, 3 → 6, ... 7 → 2, 8 → 1, 9 → 0)

- ◀ **BCD 10's complement**

- $-N \equiv 10^n - N$; 9's complement + 1

Excess-3

• مانند BCD ولی هر رقم +3

◀ جمع سرراست تر

◀ self-complement code

- (مکمل هر رقم = مکمل 9 آن)

Decimal Digit	BCD				Excess-3			
	A	B	C	D	W	X	Y	Z
0	0	0	0	0	0	0	1	1
1	0	0	0	1	0	1	0	0
2	0	0	1	0	0	1	0	1
3	0	0	1	1	0	1	1	0
4	0	1	0	0	0	1	1	1
5	0	1	0	1	1	0	0	0
6	0	1	1	0	1	0	0	1
7	0	1	1	1	1	0	1	0
8	1	0	0	0	1	0	1	1
9	1	0	0	1	1	1	0	0



ASCII character code

- ◀ We also need to represent letters and other symbols → alphanumeric codes
- ◀ ASCII = American Standard Code for Information Interchange. Also known as Western European
- ◀ It contains 128 characters:
 - 94 printable (26 upper case and 26 lower case letters, 10 digits, 32 special symbols)
 - 34 non-printable (for control functions)
- ◀ Uses 7-bit binary codes to represent each of the 128 characters

ASCII Table

		$b_6b_5b_4$ (column)							
$b_3b_2b_1b_0$	Row (hex)	000 0	001 1	010 2	011 3	100 4	101 5	110 6	111 7
0000	0	NUL Null	DLE	SP Space	0	@	P	`	p
0001	1	SOH	DC1	!	1	A	Q	a	q
0010	2	STX	DC2	"	2	B	R	b	r
0011	3	ETX	DC3	#	3	C	S	c	s
0100	4	EOT	DC4	\$	4	D	T	d	t
0101	5	ENQ	NAK	%	5	E	U	e	u
0110	6	ACK	SYN	&	6	F	V	f	v
0111	7	BEL Bell	ETB	'	7	G	W	g	w
1000	8	BS BkSpc	CAN	(8	H	X	h	x
1001	9	HT Tab	EM)	9	I	Y	i	y
1010	A	LF Line Fd	SUB	*	:	J	Z	j	z
1011	B	VT	ESC Escape	+	;	K	[k	{
1100	C	FF	FS	,	<	L	\	l	
1101	D	CR Crg Ret	GS	-	=	M]	m	}
1110	E	SO	RS	.	>	N	^	n	~
1111	F	SI	US	/	?	O	_	o	DEL

Unicode

- Established standard (16-bit alphanumeric code) for international character sets
- Since it is 16-bit, it has 65,536 codes
- Represented by 4 Hex digits
- ASCII is between 0000_{16} .. $007B_{16}$

Unicode Table

<http://www.unicode.org/charts/>

Unicode

062B 1579	ث 062C 1580	ج 062D 1581	ح 062E 1582	خ
0633 1587	س 0634 1588	ش 0635 1589	ص 0636 1590	ض
063B 1595	گ 063C 1596	پ 063D 1597	ئ 063E 1598	ت
0643 1603	ک 0644 1604	ل 0645 1605	م 0646 1606	ن
064B 1611	و 064C 1612	و 064D 1613	و 064E 1614	و
0653 1619	و 0654 1620	و 0655 1621	و 0656 1622	و
065B 1627	و 065C 1628	و 065D 1629	و 065E 1630	و
0663 1635	ز 0664 1636	ذ 0665 1637	ر 0666 1638	ز
066B 1643	ر 066C 1644	ر 066D 1645	ر 066E 1646	ر
0673 1651	ز 0674 1652	ز 0675 1653	ز 0676 1654	ز
067B 1659	پ 067C 1660	ت 067D 1661	ث 067E 1662	پ
0683 1667	ج 0684 1668	چ 0685 1669	خ 0686 1670	چ
068B 1675	ط 068C 1676	ظ 068D 1677	ڈ 068E 1678	ڈ