

# ساختمان داده و الگوریتم ها

## مبحث یازدهم: معرفی گراف و نمایش گراف

**سجاد شیرعلی شمرضا**

**پاییز 1402**

**شنبه، 20 آبان 1402**

# اطلاع رسانی

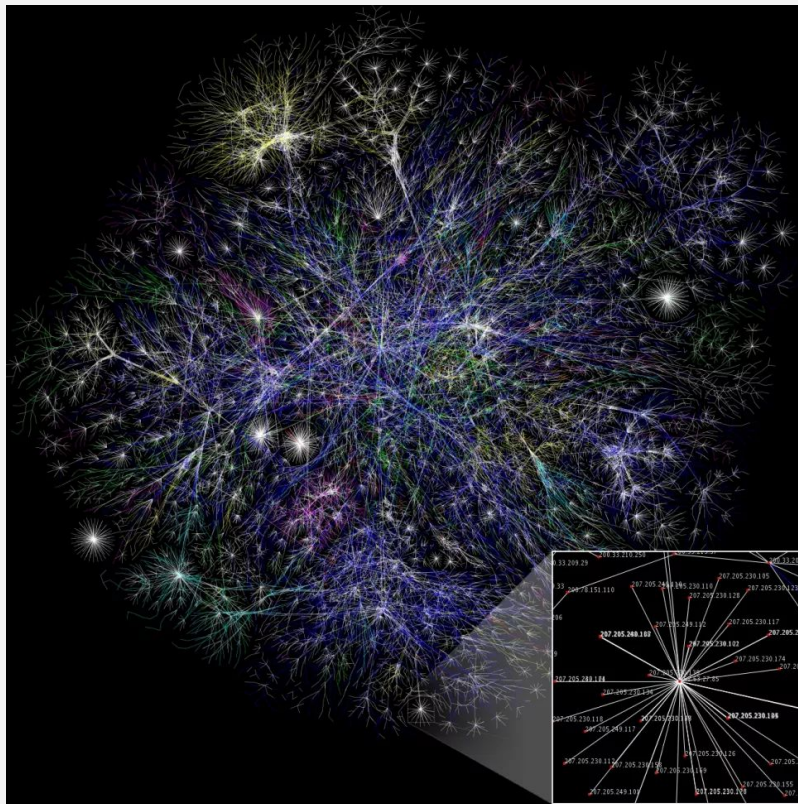
- بخش مرتبط کتاب برای این جلسه: 22

# گراف

**تعریف و نمونه**

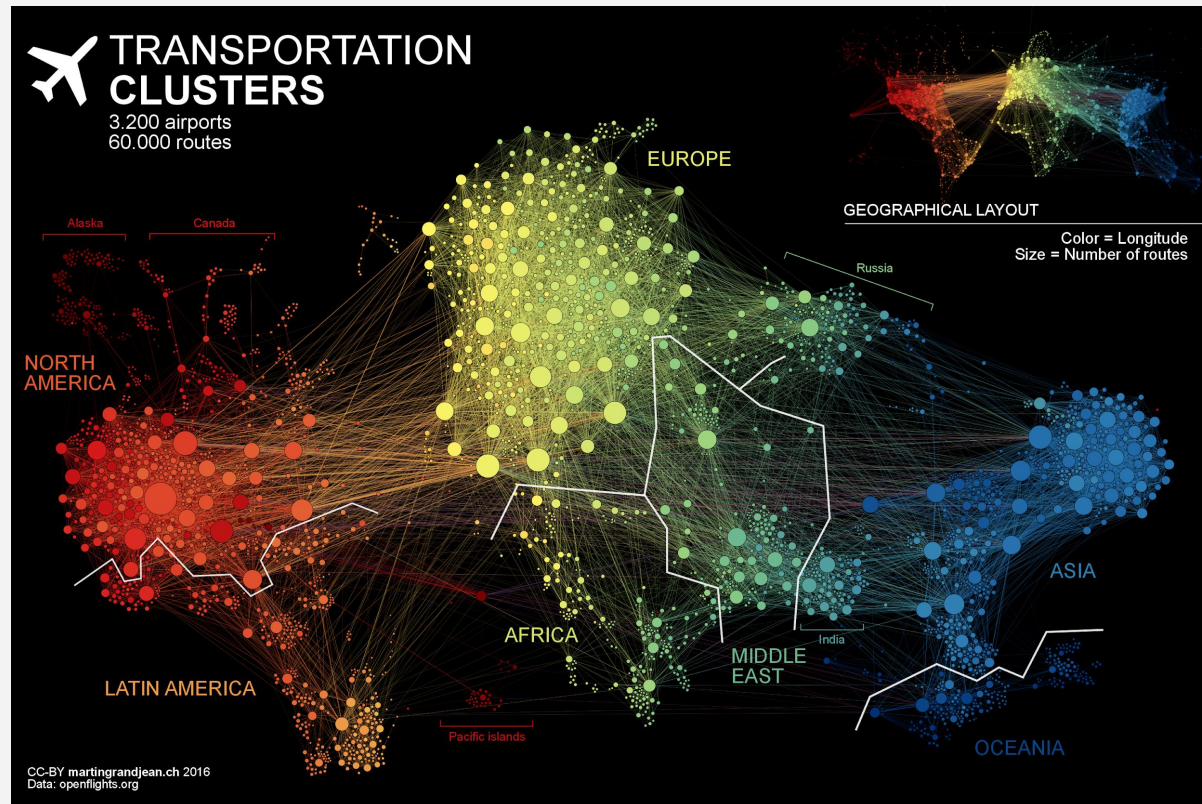
# GRAPH EXAMPLES

Partial graph of the Internet (in 2005), where each “node” is an IP address, and the “edges” between them reveal connectivity delays (shorter lines = closer IP addresses)



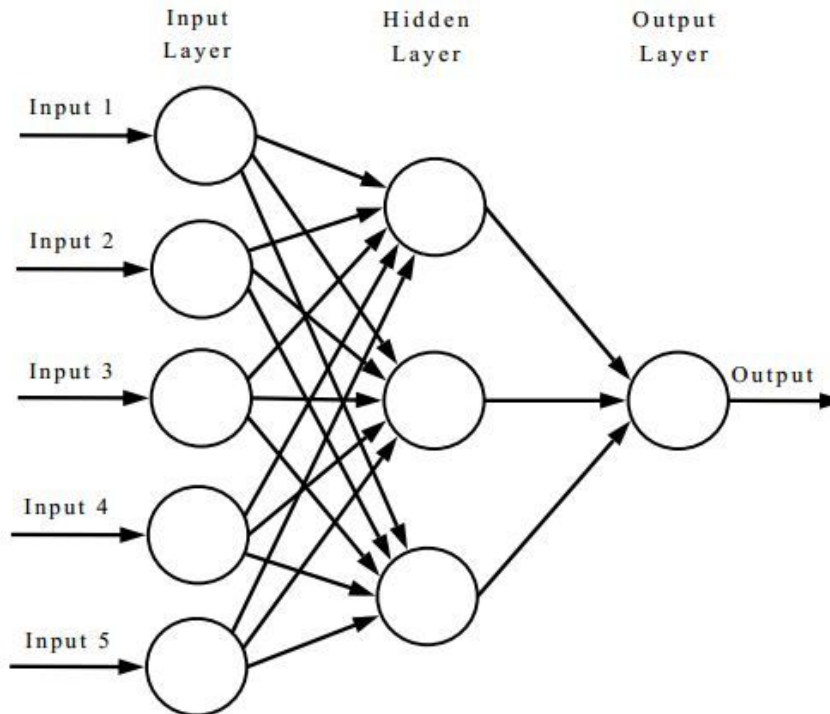
# GRAPH EXAMPLES

Each “node” is an airport,  
and flight routes are  
represented by the “edge”  
in between them



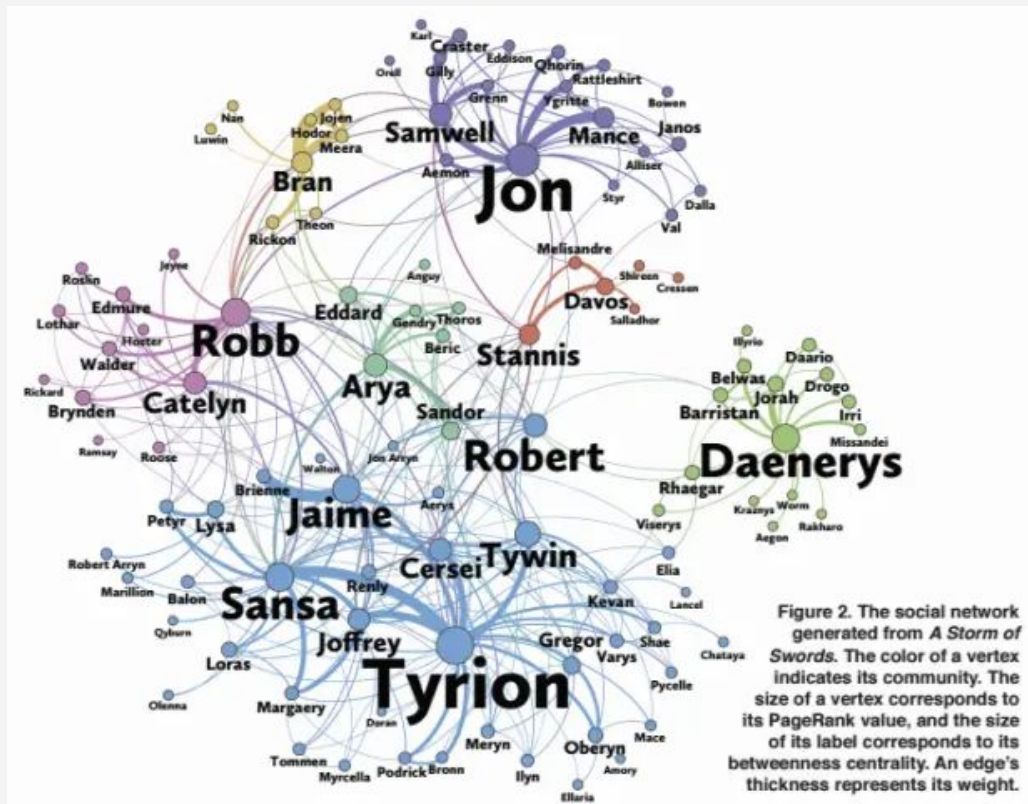
# GRAPH EXAMPLES

Neural networks! Each “node” represents a module of the neural network, and “edge” represent output/input relationships



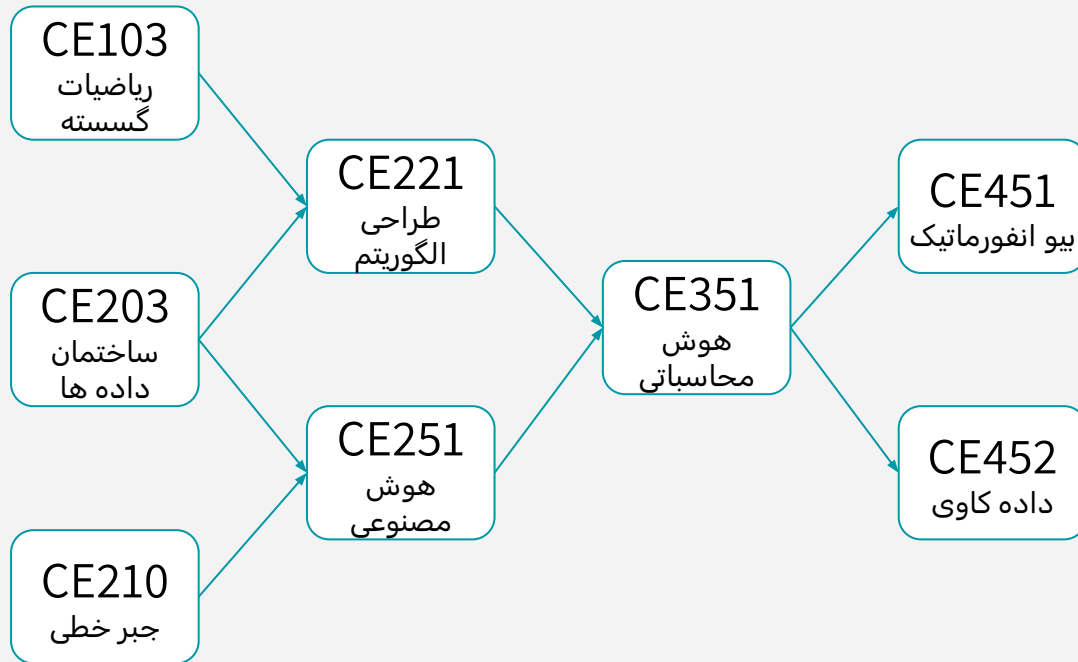
# GRAPH EXAMPLES

Graph of characters in the third book of Game of Thrones, where each “node” is a character, and “edge” reveal frequency of interaction (i.e. 2 names appearing within 15 words of one another).



# GRAPH EXAMPLES

CE prerequisites!  
“nodes” are classes  
and an “edge” from  
class A to class B  
means “class B  
depends on class A”





# WHAT ARE GRAPHS USED FOR?

- There are a lot of diverse problems that can be represented as graphs, and we want to answer questions about them
- For example:
  - How do we most efficiently route packets across the internet?
  - Are there natural “clusters” or “communities” in a graph?
  - Which character(s) are least related with \_\_\_\_\_?
  - How should I sign up for classes without violating pre-req constraints?

But first off, some terminology!

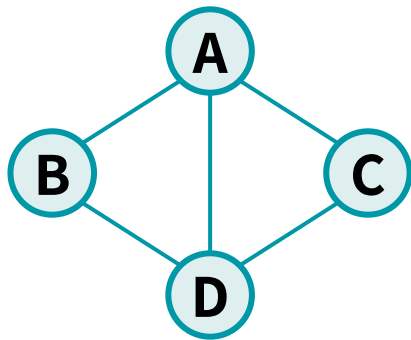
# 2 KINDS OF GRAPHS

We'll deal with both kinds of graphs in this class.

## UNDIRECTED GRAPHS

An undirected graph has  
a set of vertices ( $V$ ) & a set of edges ( $E$ )

Formally,  
 $G = (V, E)$



$V = \{A, B, C, D\}$

$E = \{ \{A, B\}, \{A, C\}, \{A, D\}, \{B, D\}, \{C, D\} \}$

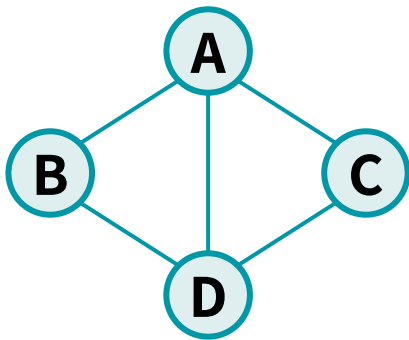
# 2 KINDS OF GRAPHS

We'll deal with both kinds of graphs in this class.

## UNDIRECTED GRAPHS

An undirected graph has  
a set of vertices ( $V$ ) & a set of edges ( $E$ )

Formally,  
 $G = (V, E)$

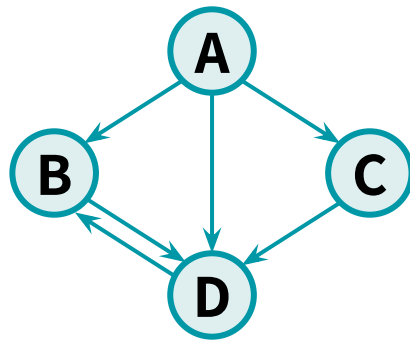


$V = \{A, B, C, D\}$

$E = \{ \{A, B\}, \{A, C\}, \{A, D\}, \{B, D\}, \{C, D\} \}$

## DIRECTED GRAPHS

A directed graph has  
a set of vertices ( $V$ ) & a set of **DIRECTED** edges ( $E$ )



Formally,  
 $G = (V, E)$

$V = \{A, B, C, D\}$

$E = \{ [A, B], [A, C], [A, D], [B, D], [C, D], [D, B] \}$

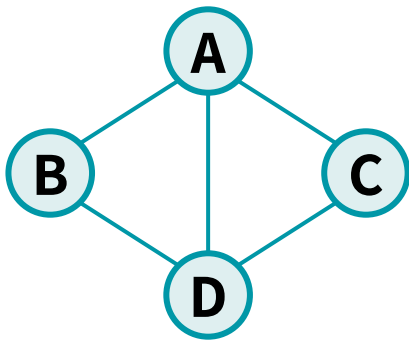
# 2 KINDS OF GRAPHS

We'll deal with both kinds of graphs in this class.

## UNDIRECTED GRAPHS

An undirected graph has  
a set of vertices ( $V$ ) & a set of edges ( $E$ )

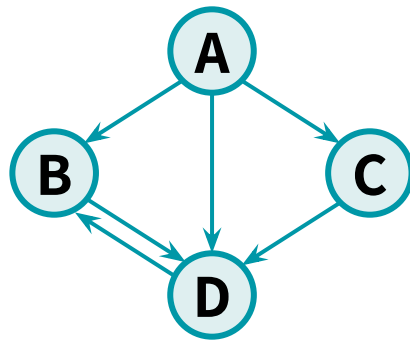
Formally,  
 $G = (V, E)$



The **degree** of vertex D is 3  
Vertex D's **neighbors** are A, B, and C

## DIRECTED GRAPHS

A directed graph has  
a set of vertices ( $V$ ) & a set of **DIRECTED** edges ( $E$ )



Formally,  
 $G = (V, E)$

The **in-degree** of vertex D is 3. The **out-degree** of vertex D is 1.  
Vertex D's **incoming neighbors** are A, B, & C  
Vertex D's **outgoing neighbor** is B

# 2 KINDS OF GRAPHS

We'll deal with both kinds of graphs in this class.

## UNDIRECTED GRAPHS

a set

Formally,  
 $G = (V, E)$

Today, we're only working with ***unweighted*** graphs. These are graphs where edges aren't assigned weights, or all edges are assumed to have the same weight.



The **degree** of vertex D is 3  
Vertex D's **neighbors** are A, B, and C

## DIRECTED GRAPHS

edges (E)

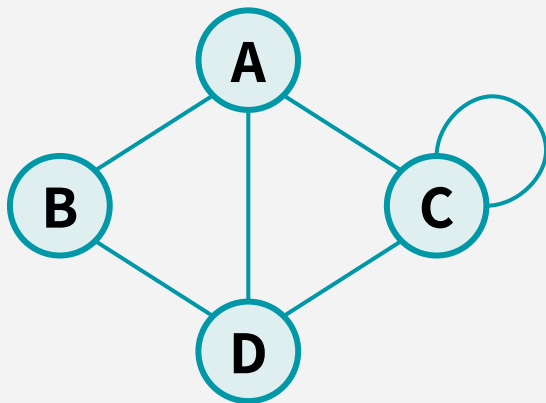
Formally,  
 $G = (V, E)$



The **in-degree** of vertex D is 3. The **out-degree** of vertex D is 1.  
Vertex D's **incoming neighbors** are A, B, & C  
Vertex D's **outgoing neighbor** is B

# GRAPH REPRESENTATIONS

## OPTION 1: **ADJACENCY MATRIX**

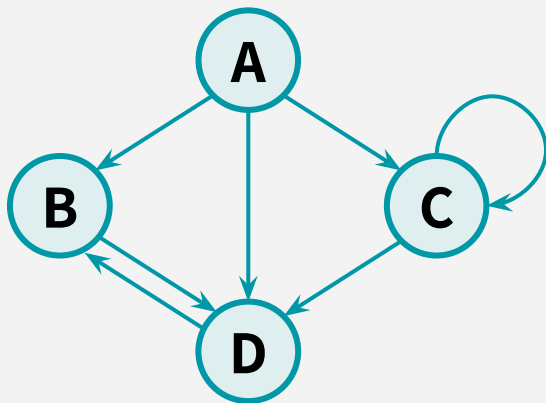


(An undirected graph)

		(destination)			
		A	B	C	D
(source)	A	0	1	1	1
	B	1	0	0	1
	C	1	0	1	1
	D	1	1	1	0

# GRAPH REPRESENTATIONS

## OPTION 1: **ADJACENCY MATRIX**

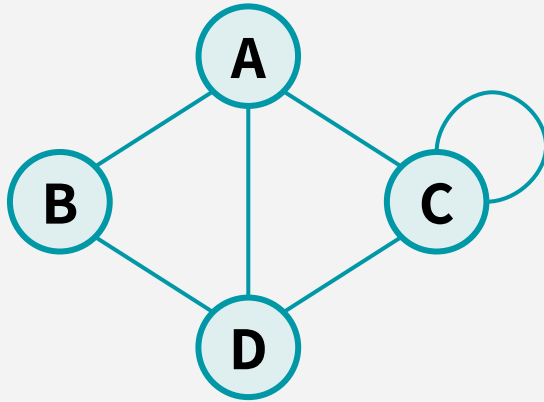


(A directed graph)

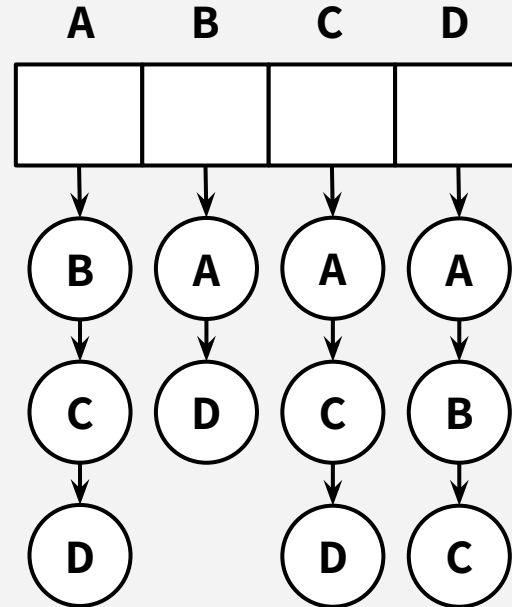
		(destination)			
		A	B	C	D
(source)	A	0	1	1	1
	B	0	0	0	1
	C	0	0	1	1
	D	0	1	0	0

# GRAPH REPRESENTATIONS

## OPTION 2: **ADJACENCY LISTS**



(An undirected graph)

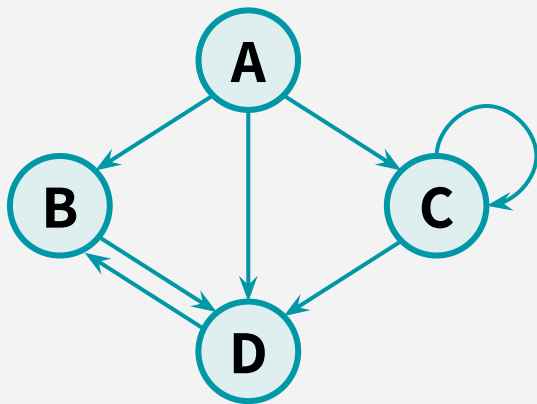


Each list stores a node's neighbors

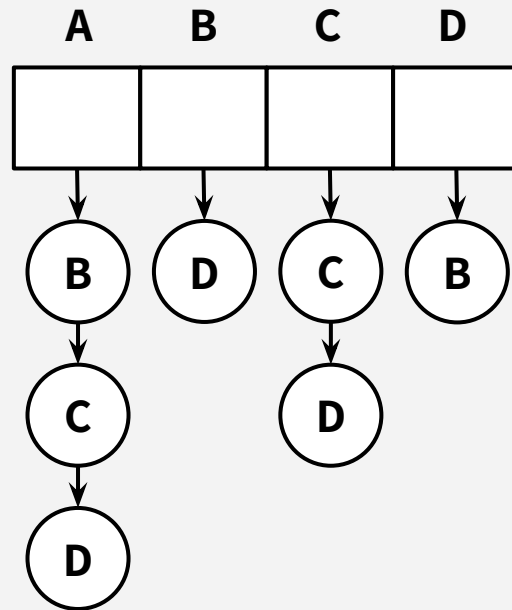


# GRAPH REPRESENTATIONS

## OPTION 2: **ADJACENCY LISTS**



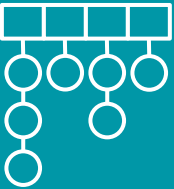
(A directed graph)



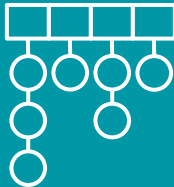
*Tracks outgoing neighbors.*

*(You could also do the same for incoming neighbors as well)*

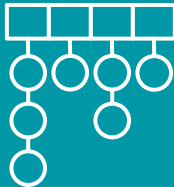
# GRAPH REPRESENTATIONS

For a graph $G = (V, E)$ where $ V  = \mathbf{n}$ , and $ E  = \mathbf{m}$	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$	
<b>EDGE MEMBERSHIP</b> Is $e = \{v, w\}$ in $E$ ?		
<b>NEIGHBOR QUERY</b> Give me $v$ 's neighbors		
<b>SPACE REQUIREMENTS</b>		

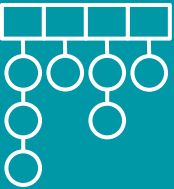
# GRAPH REPRESENTATIONS

For a graph $G = (V, E)$ where $ V  = \mathbf{n}$ , and $ E  = \mathbf{m}$	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$	
<b>EDGE MEMBERSHIP</b> Is $e = \{v, w\}$ in $E$ ?	<b><math>O(1)</math></b>	
<b>NEIGHBOR QUERY</b> Give me $v$ 's neighbors		
<b>SPACE REQUIREMENTS</b>		

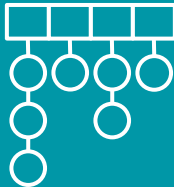
# GRAPH REPRESENTATIONS

<p>For a graph <math>G = (V, E)</math>          where <math> V  = \mathbf{n}</math>, and <math> E  = \mathbf{m}</math></p>	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$	
<p><b>EDGE MEMBERSHIP</b>          Is <math>e = \{v, w\}</math> in <math>E</math>?</p>	<p><b><math>O(1)</math></b></p>	<p><b><math>O(\deg(v))</math> or  <math>O(\deg(w))</math></b></p>
<p><b>NEIGHBOR QUERY</b>          Give me <math>v</math>'s neighbors</p>		
<p><b>SPACE REQUIREMENTS</b></p>		

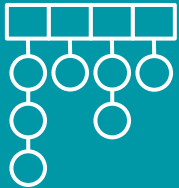
# GRAPH REPRESENTATIONS

<p>For a graph <math>G = (V, E)</math>          where <math> V  = \mathbf{n}</math>, and <math> E  = \mathbf{m}</math></p>	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$	
<p><b>EDGE MEMBERSHIP</b>          Is <math>e = \{v, w\}</math> in <math>E</math>?</p>	<p><b><math>O(1)</math></b></p>	<p><b><math>O(\deg(v))</math> or  <math>O(\deg(w))</math></b></p>
<p><b>NEIGHBOR QUERY</b>          Give me <math>v</math>'s neighbors</p>	<p><b><math>O(n)</math></b></p>	
<p><b>SPACE REQUIREMENTS</b></p>		

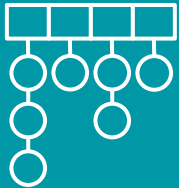
# GRAPH REPRESENTATIONS

<p>For a graph <math>G = (V, E)</math>  where <math> V  = \mathbf{n}</math>, and <math> E  = \mathbf{m}</math></p>	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$	
<p><b>EDGE MEMBERSHIP</b>  Is <math>e = \{v, w\}</math> in <math>E</math>?</p>	<p><b><math>O(1)</math></b></p>	<p><b><math>O(\deg(v))</math> or  <math>O(\deg(w))</math></b></p>
<p><b>NEIGHBOR QUERY</b>  Give me <math>v</math>'s neighbors</p>	<p><b><math>O(n)</math></b></p>	<p><b><math>O(\deg(v))</math></b></p>
<p><b>SPACE REQUIREMENTS</b></p>		

# GRAPH REPRESENTATIONS

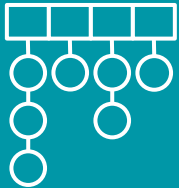
For a graph $G = (V, E)$ where $ V  = \mathbf{n}$ , and $ E  = \mathbf{m}$	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$	
<b>EDGE MEMBERSHIP</b> Is $e = \{v, w\}$ in $E$ ?	$O(1)$	$O(\deg(v))$ or $O(\deg(w))$
<b>NEIGHBOR QUERY</b> Give me $v$ 's neighbors	$O(n)$	$O(\deg(v))$
<b>SPACE REQUIREMENTS</b>	$O(n^2)$	

# GRAPH REPRESENTATIONS

For a graph $G = (V, E)$ where $ V  = \mathbf{n}$ , and $ E  = \mathbf{m}$	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$	
<b>EDGE MEMBERSHIP</b> Is $e = \{v, w\}$ in $E$ ?	$O(1)$	$O(\deg(v))$ or $O(\deg(w))$
<b>NEIGHBOR QUERY</b> Give me $v$ 's neighbors	$O(n)$	$O(\deg(v))$
<b>SPACE REQUIREMENTS</b>	$O(n^2)$	$O(n + m)$



# GRAPH REPRESENTATIONS

For a graph $G = (V, E)$ where $ V  = n$ , and $ E  = m$	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$	
<b>EDGE MEMBERSHIP</b> Is $e = \{v, w\}$ in $E$ ?	$O(1)$	$O(\deg(v))$ or $O(\deg(w))$
<b>NEIGHBOR QUERY</b> Give me $v$ 's neighbors	$O(n)$	$O(\deg(v))$
<b>SPACE REQUIREMENTS</b>	$O(n^2)$	$O(n + m)$

Generally, better for  
sparse graphs  
(where  $m \ll n^2$ ).

**We'll assume this  
representation,  
unless otherwise  
stated.**



سوال؟