

EE3025 Assignment-1

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Download all python codes from

https://github.com/Kkuntal990/EE3025-DSP/tree/main/assignment_1/code

and latex-tikz codes from

<https://github.com/srikanth2001/EE3025-DSP/blob/main/Assignment-01/ee18btech11023.tex>

1 PROBLEM

1.1. Defining $x(n)$ and $h(n)$,

$$x(n) = \left\{ \underset{\uparrow}{1}, 2, 3, 4, 2, 1 \right\} \quad (1.1.1)$$

$$h(n) = \left(-\frac{1}{2} \right)^n u(n) + \left(-\frac{1}{2} \right)^{n-2} u(n-2) \quad (1.1.2)$$

1.2. Compute $X(k)$, $H(k)$ and $y(n)$ using FFT and IFFT

2 SOLUTION

2.1. input signal $x(n)$

$$x(n) = \left\{ \underset{\uparrow}{1}, 2, 3, 4, 2, 1 \right\} \quad (2.1.1)$$

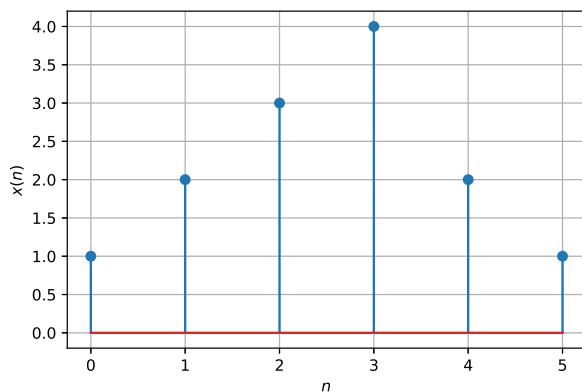


Fig. 2.1: input signal : $x(n)$

2.2. Impulse Response of the System is

$$h(n) = \left(-\frac{1}{2} \right)^n u(n) + \left(-\frac{1}{2} \right)^{n-2} u(n-2) \quad (2.2.1)$$

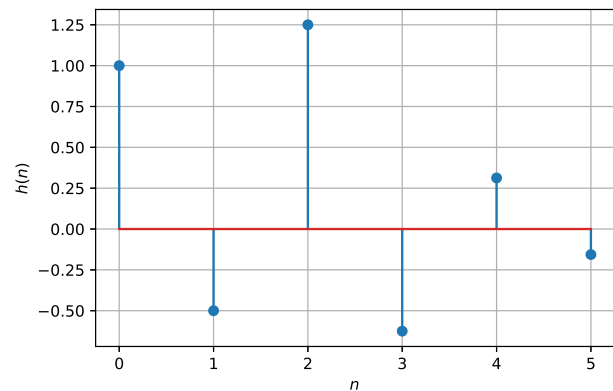


Fig. 2.2: impulse response : $h(n)$

2.3. FFT of the input signal $x(n)$ is

$$X(k) = \sum_{n=0}^{N-1} x(n) e^{-j2\pi kn/N}, \quad k = 0, 1, \dots, N-1 \quad (2.3.1)$$

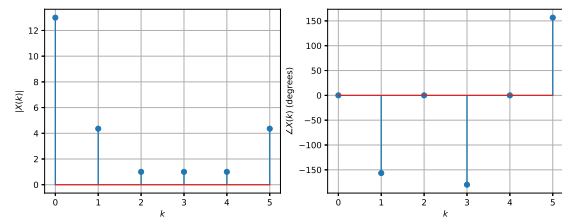
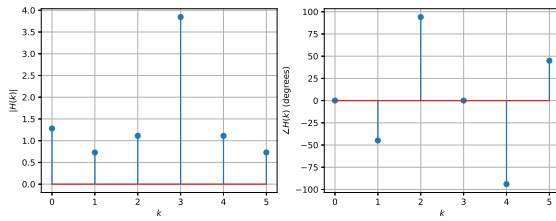


Fig. 2.3: FFT of $x(n)$: $X(k)$

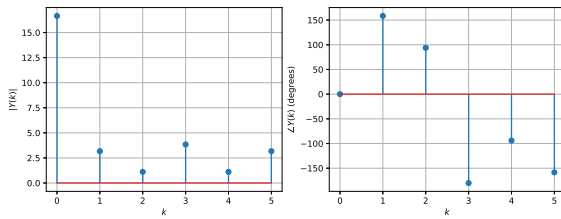
2.4. FFT of the impulse response $h(n)$ is

$$H(k) = \sum_{n=0}^{N-1} h(n) e^{-j2\pi kn/N}, \quad k = 0, 1, \dots, N-1 \quad (2.4.1)$$

Fig. 2.4: FFT of $h(n)$: $H(k)$

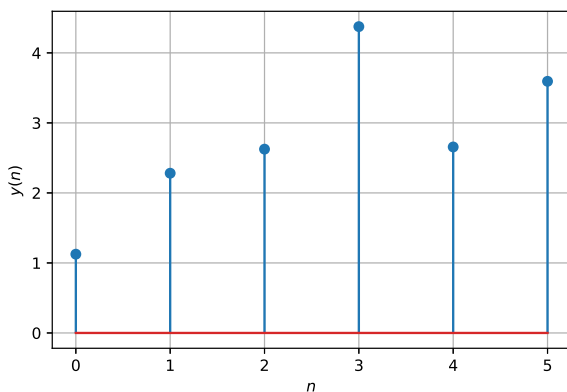
2.5. FFT of output Signal $y(n)$ can be computed as

$$Y(k) = X(k)H(k) \quad (2.5.1)$$

Fig. 2.5: $Y(k) = H(k)X(k)$

2.6. $y(n)$ can be computed by taking IFFT of $Y(k)$

$$y(n) = \frac{1}{N} \sum_{k=0}^{N-1} Y(k) e^{j2\pi nk/N}, \quad k = 0, 1, \dots, N-1 \quad (2.6.1)$$

Fig. 2.6: IFFT of $Y(k)$: $y(n)$

3 PROBLEM

3.1. Wherever possible, express all the above equations as matrix equations.

4 SOLUTION

4.1. FFT of signal $X(n)$

$$X(k) \triangleq W_N^{nk} x(n), \quad k = 0, 1, \dots, N-1 \quad (4.1.1)$$

where $W_N^{nk} = e^{-j2\pi nk/N}$ which can be expressed as:

$$\begin{bmatrix} X(0) \\ X(1) \\ X(2) \\ X(3) \\ X(4) \\ X(5) \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & W_N^1 & W_N^2 & W_N^3 & W_N^4 & W_N^5 \\ 1 & W_N^2 & W_N^4 & W_N^6 & W_N^8 & W_N^{10} \\ 1 & W_N^3 & W_N^6 & W_N^9 & W_N^{12} & W_N^{15} \\ 1 & W_N^4 & W_N^8 & W_N^{12} & W_N^{16} & W_N^{20} \\ 1 & W_N^5 & W_N^{10} & W_N^{15} & W_N^{20} & W_N^{25} \end{bmatrix} \begin{bmatrix} x(0) \\ x(1) \\ x(2) \\ x(3) \\ x(4) \\ x(5) \end{bmatrix} \quad (4.1.2)$$

$$\Rightarrow \begin{bmatrix} X(0) \\ X(1) \\ X(2) \\ X(3) \\ X(4) \\ X(5) \end{bmatrix} = \begin{bmatrix} 1 + 2 + 3 + 4 + 2 + 1 \\ 1 + (2)e^{-j\pi/3} + \dots + (1)e^{-j5\pi/3} \\ 1 + (2)e^{-2j\pi/3} + \dots + (1)e^{-2j5\pi/3} \\ 1 + (2)e^{-3j\pi/3} + \dots + (1)e^{-3j5\pi/3} \\ 1 + (2)e^{-4j\pi/3} + \dots + (1)e^{-4j5\pi/3} \\ 1 + (2)e^{-5j\pi/3} + \dots + (1)e^{-5j5\pi/3} \end{bmatrix} \quad (4.1.3)$$

On solving,

$$\Rightarrow \begin{bmatrix} X(0) \\ X(1) \\ X(2) \\ X(3) \\ X(4) \\ X(5) \end{bmatrix} = \begin{bmatrix} 13 + 0j \\ -4 - 1.732j \\ 1 + 0j \\ -1 + 0j \\ 1 + 0j \\ -4 + 1.732j \end{bmatrix} \quad (4.1.4)$$

4.2. Similarly,

$$\begin{bmatrix} H(0) \\ H(1) \\ H(2) \\ H(3) \\ H(4) \\ H(5) \end{bmatrix} = \begin{bmatrix} h(0) + h(1) + h(2) + h(3) + h(4) + h(5) \\ h(0) + h(1)e^{-j\pi/3} + \dots + h(5)e^{-j5\pi/3} \\ h(0) + h(1)e^{-2j\pi/3} + \dots + h(5)e^{-2j5\pi/3} \\ h(0) + h(1)e^{-3j\pi/3} + \dots + h(5)e^{-3j5\pi/3} \\ h(0) + h(1)e^{-4j\pi/3} + \dots + h(5)e^{-4j5\pi/3} \\ h(0) + h(1)e^{-5j\pi/3} + \dots + h(5)e^{-5j5\pi/3} \end{bmatrix} \quad (4.2.1)$$

On solving,

$$\Rightarrow \begin{bmatrix} H(0) \\ H(1) \\ H(2) \\ H(3) \\ H(4) \\ H(5) \end{bmatrix} = \begin{bmatrix} 1.28125 + 0j, \\ 0.51625 - 0.5141875j, \\ -0.078125 + 1.1095625j, \\ 3.84375 + 0j, \\ -0.071825 - 1.1095625j, \\ 0.515625 + 0.5141875j \end{bmatrix} \quad (4.2.2)$$

4.3. Compute $Y(k)$ using Eq (2.5.1)

$$\begin{bmatrix} Y(0) \\ Y(1) \\ Y(2) \\ Y(3) \\ Y(4) \\ Y(5) \end{bmatrix} = \begin{bmatrix} X(0) \cdot H(0) \\ X(1) \cdot H(1) \\ X(2) \cdot H(2) \\ X(3) \cdot H(3) \\ X(4) \cdot H(4) \\ X(5) \cdot H(5) \end{bmatrix} \quad (4.3.1)$$

Solving,

$$\begin{bmatrix} Y(0) \\ Y(1) \\ Y(2) \\ Y(3) \\ Y(4) \\ Y(5) \end{bmatrix} = \begin{bmatrix} 16.6562 + 0j \\ -2.95312 + 1.16372j \\ -0.07812 + 1.10959j \\ -3.84375 - 9.27556j \\ -0.07812 - 1.10959j \\ -2.95312 - 1.16372j \end{bmatrix} \quad (4.3.2)$$

4.4.

$$y(n) \triangleq (W_N^{nk})^* Y(k), \quad n = 0, 1, \dots, N-1 \quad (4.4.1)$$

where $(W_N^{nk})^*$ is conjugate of W_N^{nk} from (4.1.1).

$$\begin{bmatrix} y(0) \\ y(1) \\ y(2) \\ y(3) \\ y(4) \\ y(5) \end{bmatrix} = \begin{bmatrix} 1.125 + 0j \\ 2.28125071 + 0j \\ 2.6250019 - 1.11022302 \times 10^{-16}j \\ 4.37499667 - 1.47104551 \times 10^{-15}j \\ 2.6562481 + 6.10622664 \times 10^{-16}j \\ 3.59375262 - 1.60982339 \times 10^{-15}j \end{bmatrix} \quad (4.4.2)$$

4.5. Properties :

a) symmetry property :

$$W_N^{k+N/2} = -W_N^k$$

b) Periodicity property :

$$W_N^{k+N} = W_N^k$$

c)

$$W_N^2 = W_{N/2}$$

4.6. Using properties to derive FFT from DFT :

$$X(k) = \sum_{n=0}^{N-1} x(n) W_N^{kn}, \quad k = 0, 1, \dots, N-1 \quad (4.6.1)$$

$$= \sum_{n=\text{even}} x(n) W_N^{kn} + \sum_{n=\text{odd}} x(n) W_N^{kn} \quad (4.6.2)$$

$$= \sum_{m=0}^2 x(2m) W_N^{2mk} + \sum_{m=0}^2 x(2m+1) W_N^{(2m+1)k} \quad (4.6.3)$$

using property c, we get,

$$X(k) = \sum_{m=0}^2 x(2m) W_{N/2}^{mk} + W_N^k \sum_{m=0}^2 x(2m+1) W_{N/2}^{mk} \quad (4.6.4)$$

$$= X_1(k) + W_N^k X_2(k) \quad (4.6.5)$$

4.7. • $X_1(k)$ and $X_2(k)$ are 3 point DFTs of $x(2m)$ and $x(2m+1)$, $m=0,1,2$.

• $X_1(k)$ and $X_2(k)$ are periodic, Hence $X_1(k+3) = X_1(k)$ and $X_2(k+3) = X_2(k)$.

• From performing this step once we can see that number of computations have been reduced from N^2 to $\frac{N^2}{2}$.

4.8. Using the above properties recursively we have implemented radix-2 the Fast-Fourier transform algorithm.

Algorithm	$t(N = 128)$	$t(N = 2048)$
DTFT	33.2 ms	7.36 s
FFT	1.54 ms	27.6 ms

• $t(N)$ corresponds to average time of execution for sample size of N .

• We can observe that as we increase N , the difference in execution times is drastically increasing.

• From our implementation of radix-2 FFT we can see that complexity is reduced from $O(n^2)$ to $O(n \log n)$