

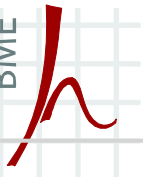
Computer Architectures

1. A card and code based door-lock

Prepared by: **Gábor Horváth**, ghorvath@hit.bme.hu

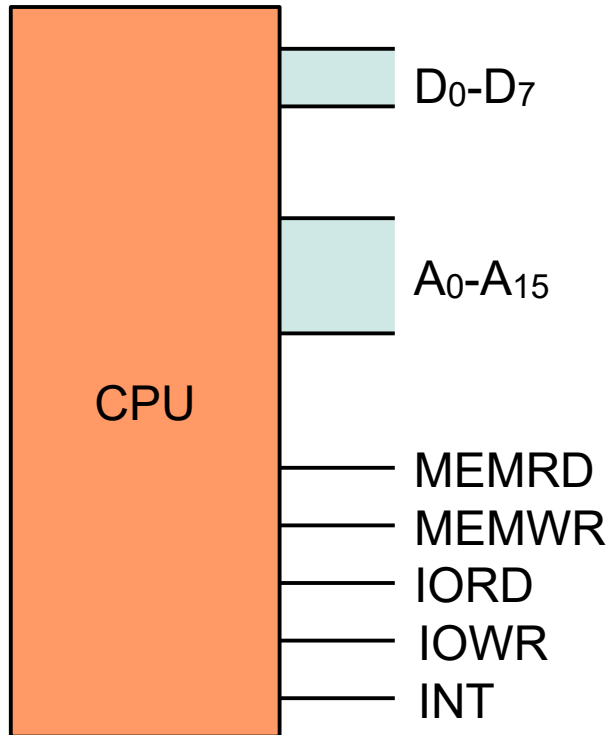
Presented by: **Gábor Lencse**, lencse@hit.bme.hu

- Designing the hardware
 - The building blocks of the system
 - Adding memory to the CPU
 - Adding peripherals to the CPU
- Designing the software
 - State-machine model of the system
 - Flow-charts of the algorithms
 - Implementation



THE HARDWARE

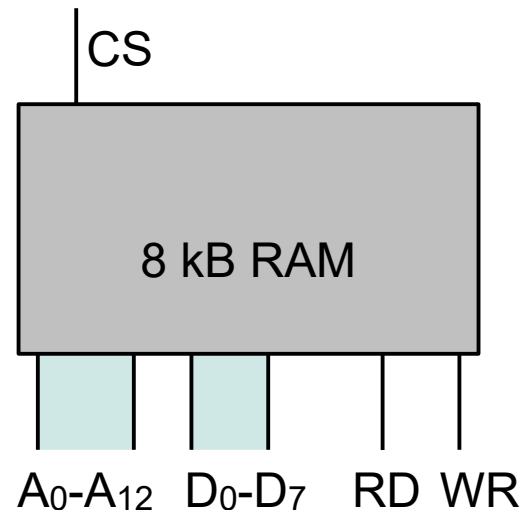
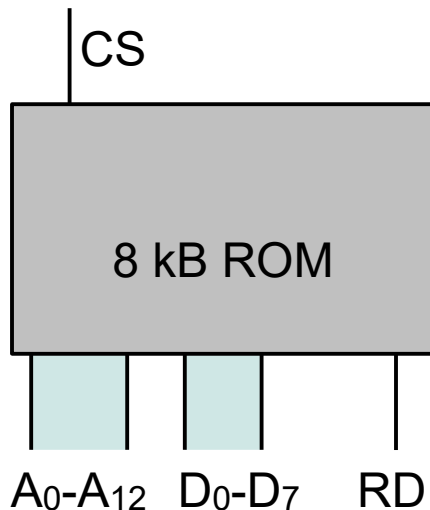
Building blocks of the system



- A 8-bit CPU
 - Data bus: 8 bit wide
 - Address bus: 16 bit wide
 - Multiplexed memory and I/O buses
 - RISC instruction set
 - 3-operand instruction format
 - Starting address: **0000h**
 - Interrupt subroutine address: **1000h**

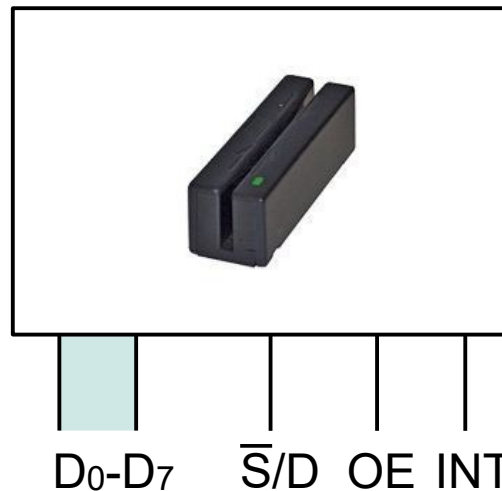
Building blocks of the system

- Memory
 - ROM: to store **instructions** and program constants
 - RAM: to store **program variables** and **stack**
- Let us use 8kB ROM and 8kB RAM
 - 8 bit data bus
 - 13 bit address bus



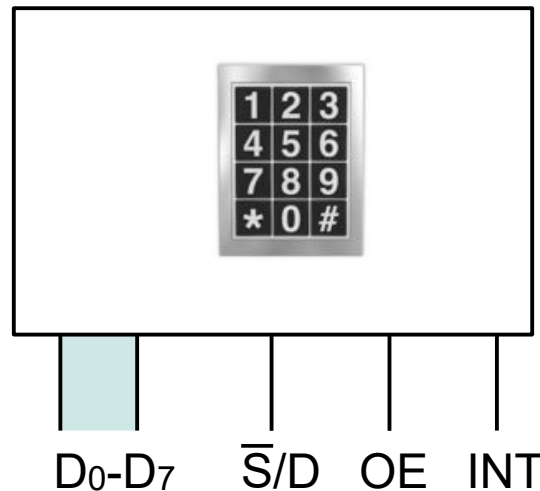
Building blocks of the system

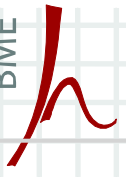
- Card reader
 - OE: output enable
 - It has an 8 bit output: D
 - If $\overline{S}/D = 0$: it gives back an 1, if a new card has arrived
 - If $\overline{S}/D = 1$: it gives back the code (ID) of the last card
 - INT: generates an interrupt, if a new card is recognized



Building blocks of the system

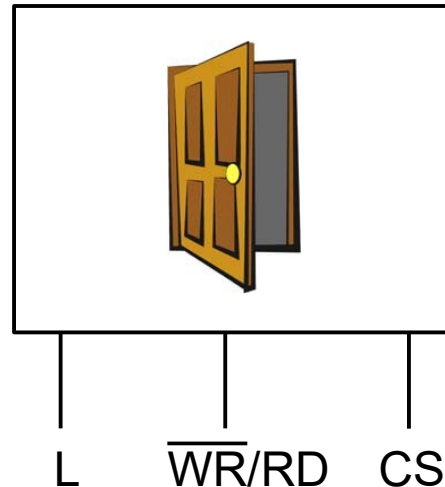
- Keypad
 - OE: output enable
 - It has a 8 bit output D:
 - If $\overline{S}/D = 0$: it gives back an 1 if a button has been pressed
 - If $\overline{S}/D = 1$: it gives back the code of the last key
 - INT: generates an interrupt upon each key press





Building blocks of the system

- Door lock
 - CS: chip select
 - \overline{WR}/RD
 - If $\overline{WR}/RD = 1$: sets L according to the state of the door
 - (open=1, locked=0)
 - If $\overline{WR}/RD = 0$: opens the door. The door automatically locks again after a timeout.



Building blocks of the system

- LED lights
 - Indicate the status of the door (red/green)
- Additional elements:
 - Decoder – for interfacing the memory
 - Comparators – for detecting the addresses of peripherals
 - D flip-flop – for driving the LED lights

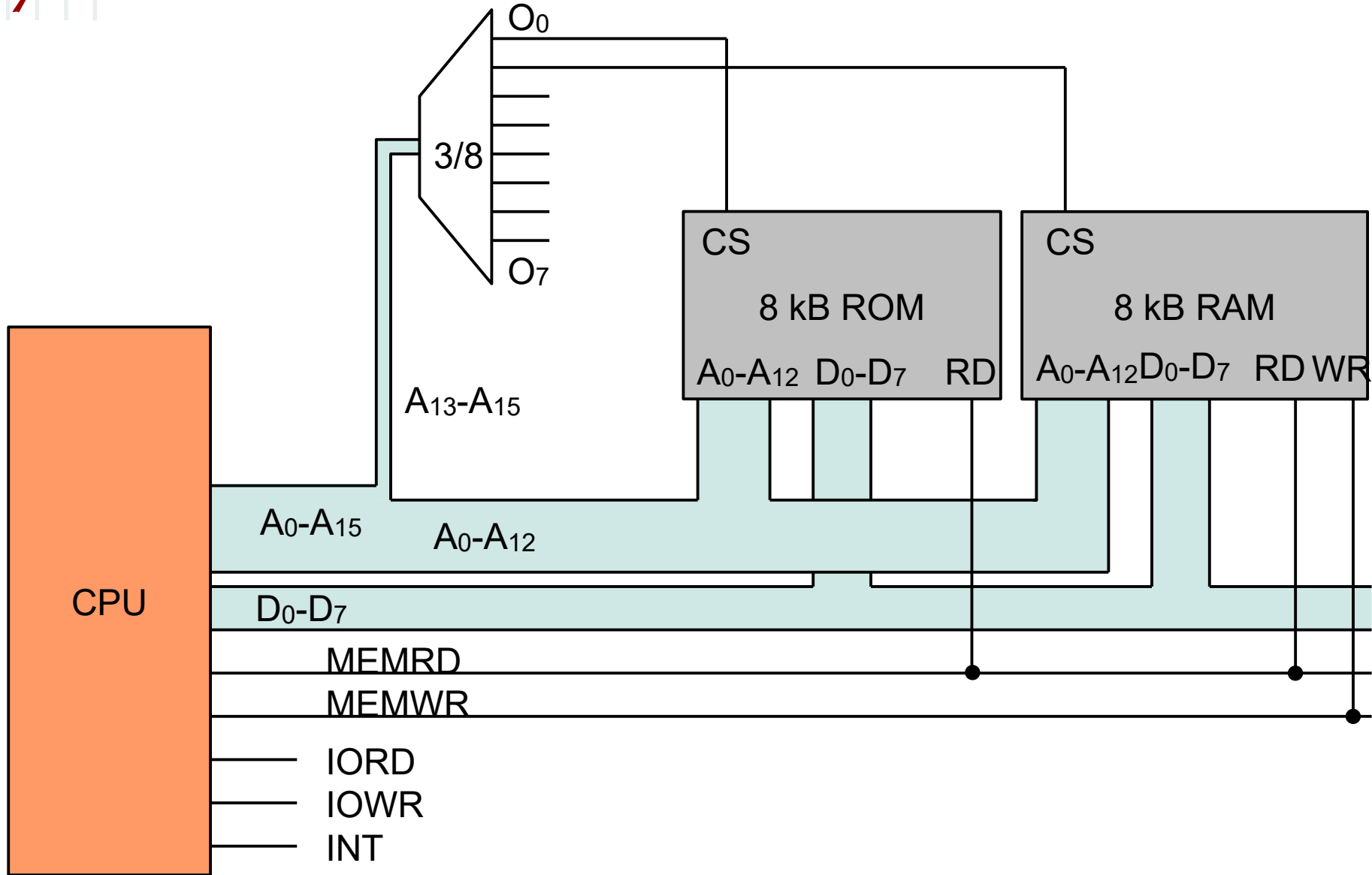
Adding memory to the CPU

- We have 8 kB ROM + 8 kB RAM
- Memory map:

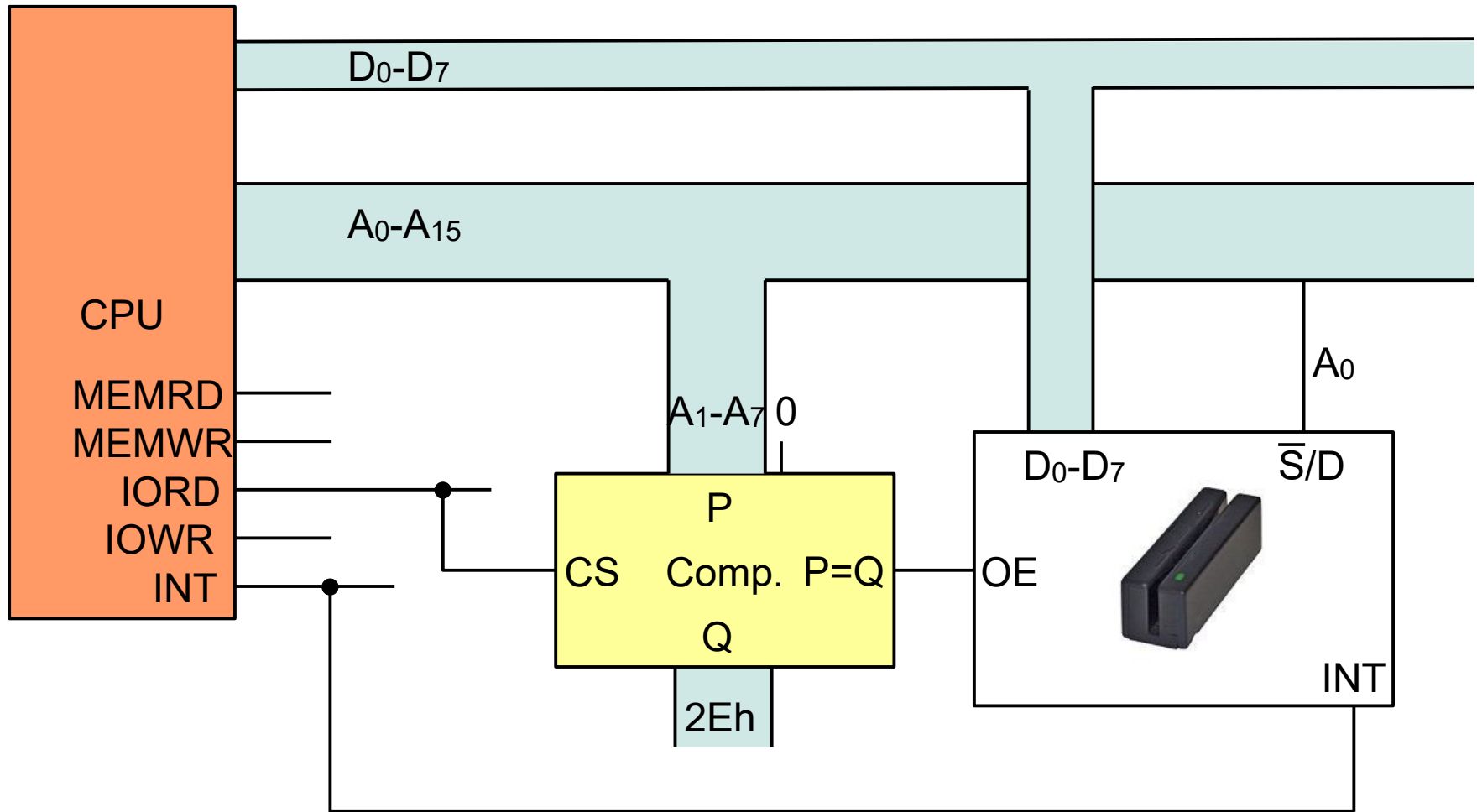
E000h-FFFFh				
C000h-DFFFh				
A000h-BFFFh				
8000h-9FFFh				
6000h-7FFFh				
4000h-5FFFh				
2000h-3FFFh	RAM	010 1	1111	1111
		010 0	0000	0000
0000h-1FFFh	ROM	001 1	1111	1111
		001 0	0000	0000
		000 1	1111	1111
		000 0	0000	0000

- A₁₃-A₁₅ determine which module to use
- A₀-A₁₂ determine the byte position in the selected module

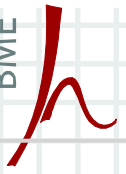
Adding memory to the CPU



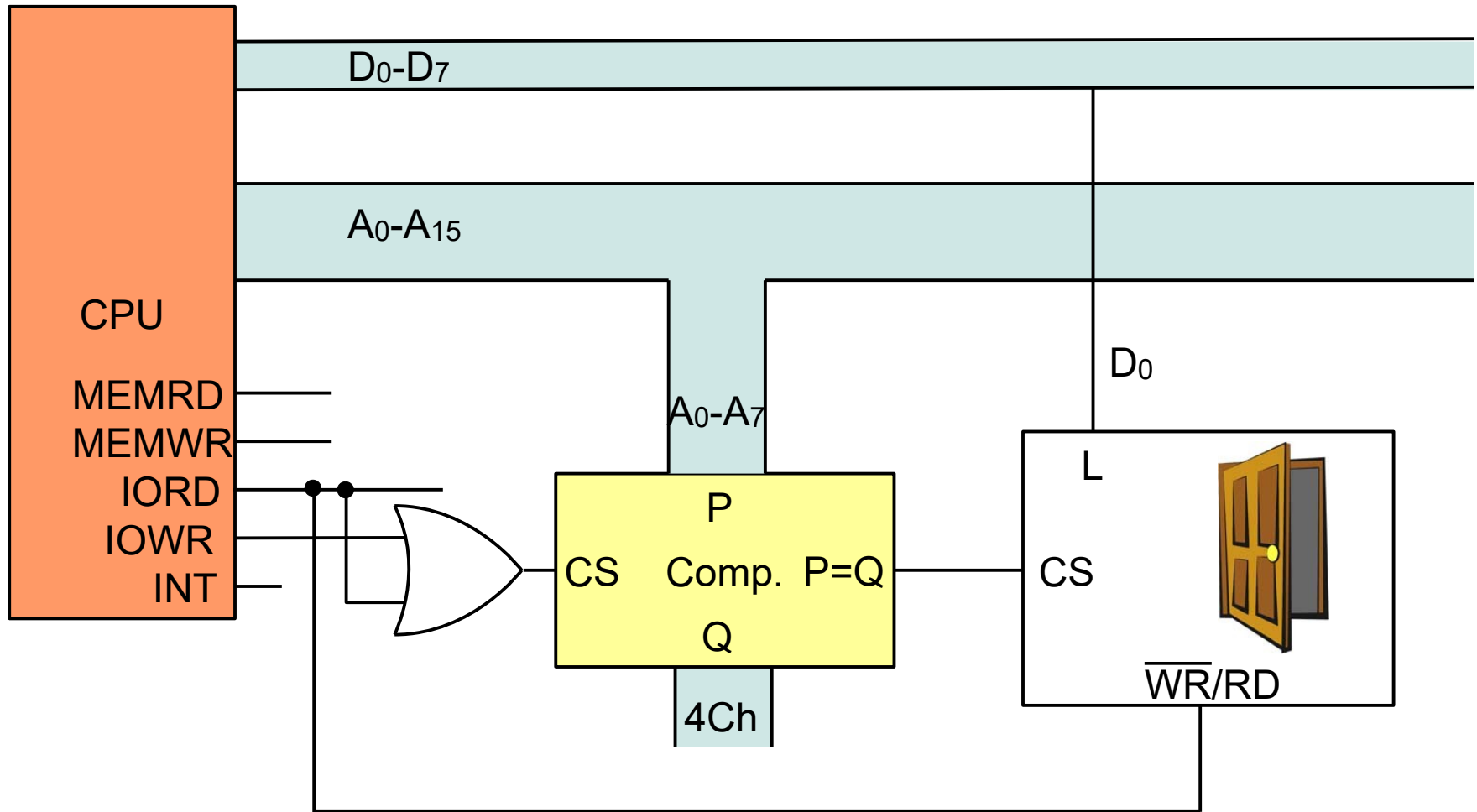
Adding the card reader



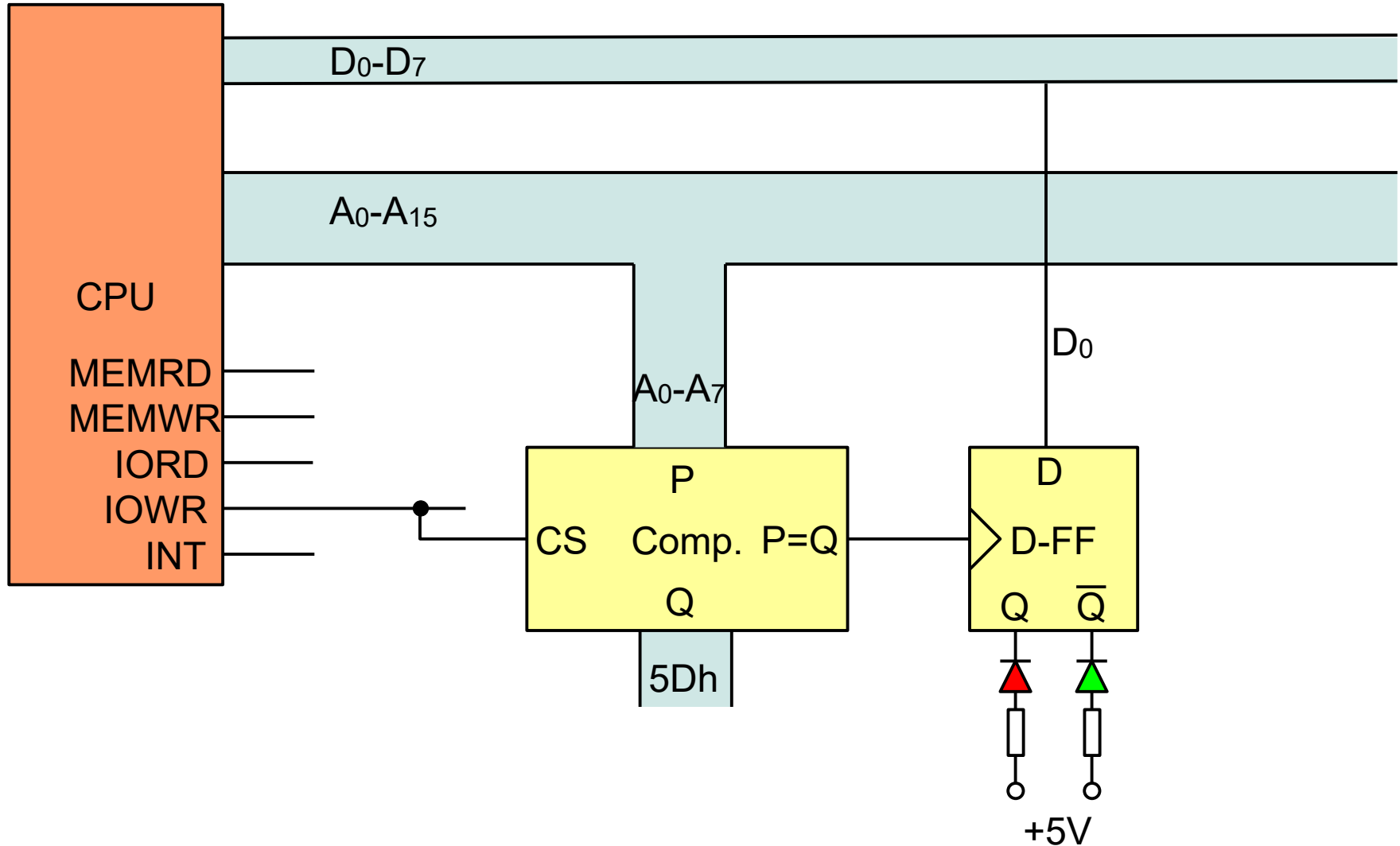
- Adding the keypad: the same, with base address 3Eh

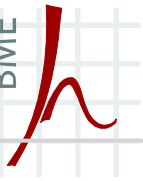


Adding the door lock



Adding the LED lights

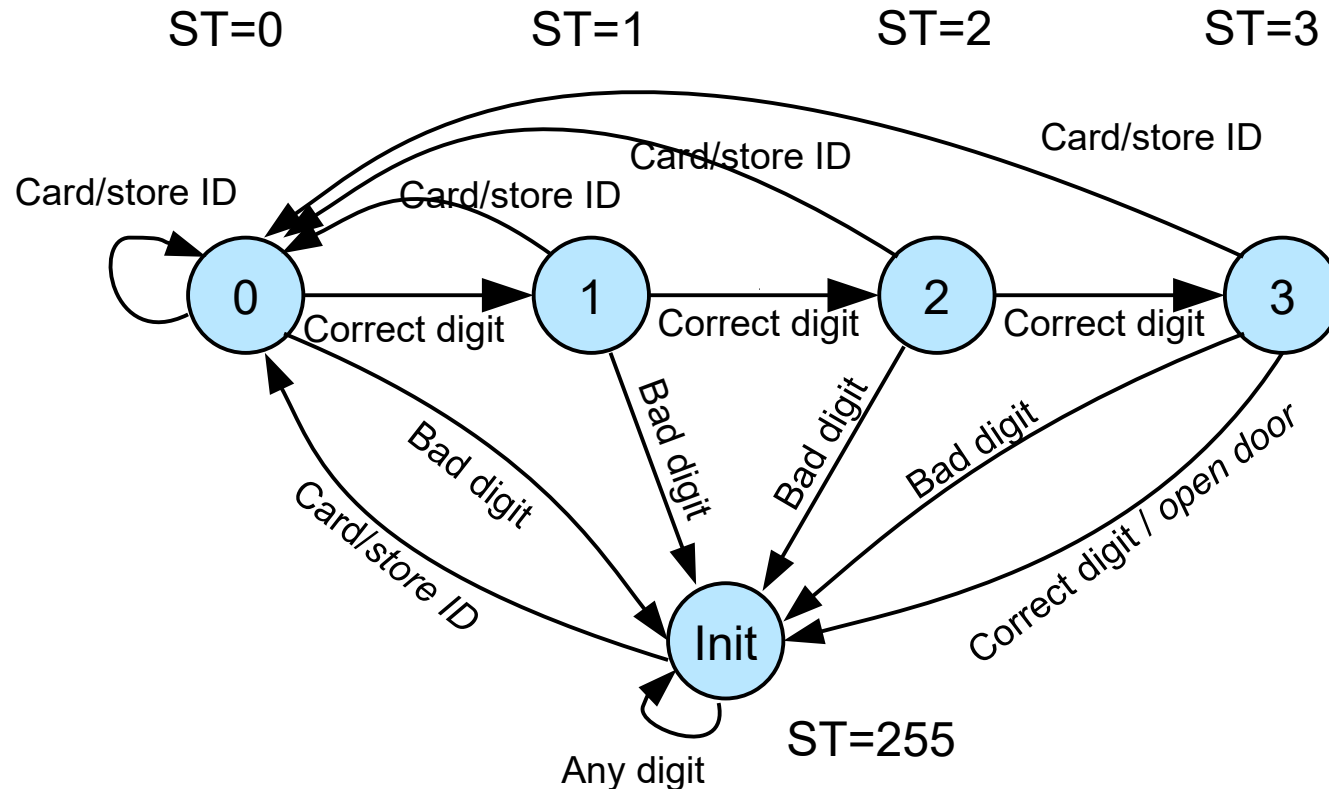




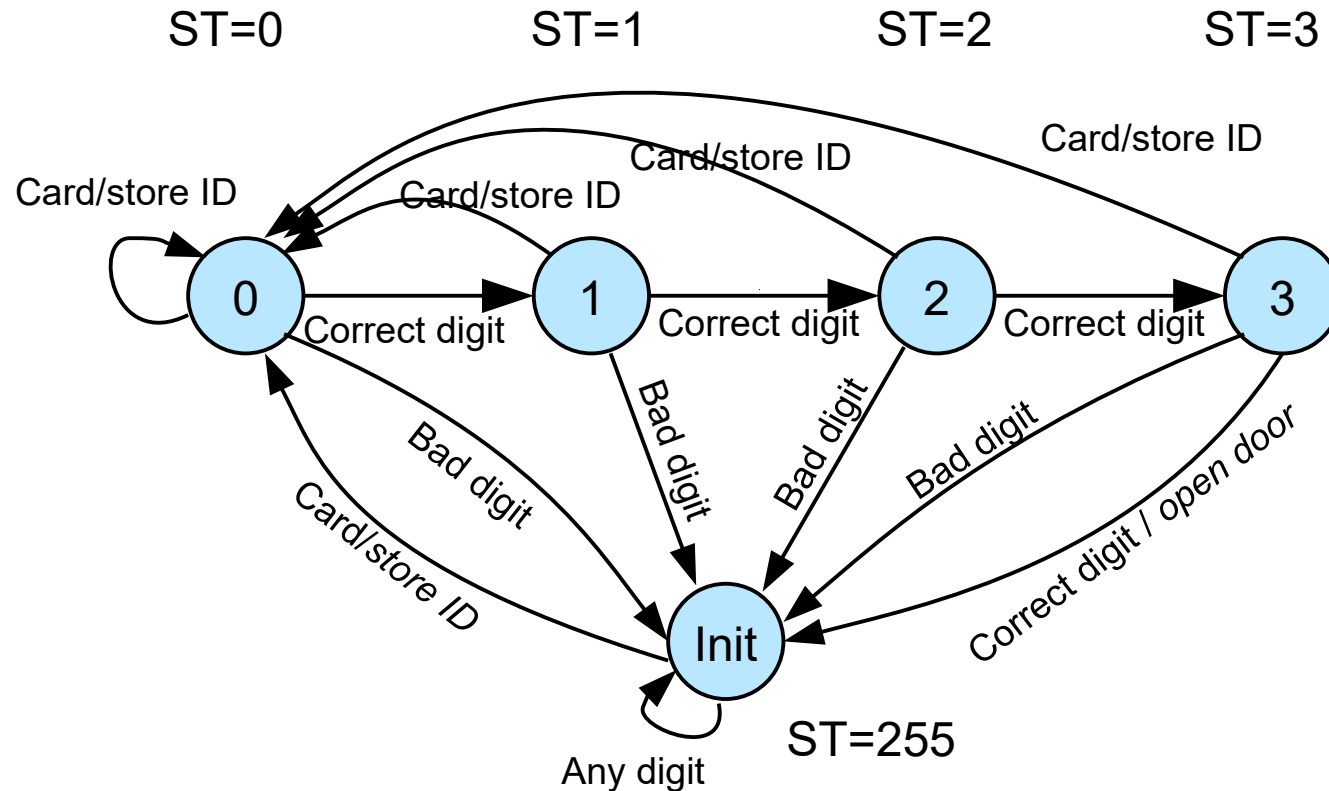
THE SOFTWARE

State-machine of the system

- ST (state) = number of correct code digits typed so far on the keypad

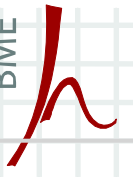


Transitions of the state-machine



- Card event: ST=0, and store the card ID in a local variable
- Correct digit: (if ST!=255 then) ST=ST+1,
 - if now ST==4 then: ST=255 and open door
- Bad digit: ST=255

- The array of 4-digit key codes for each card ID
 - The card ID is 8 bit wide → 256 different IDs are possible
 - An array is used
 - with 256 entries, one for each card ID
 - each entry is the correct code (4-byte long) of the given card
 - Total memory consumption: $256 \cdot 4 = 1024$ bytes (=400h)
 - The i th code digit of card j is located at
 - array start address + $j \cdot 4 + i$ (j in $\{0..255\}$, i in $\{0..3\}$)
- Local variables:
 - 1-byte integer variable ST
 - 1-byte integer variable CARDID

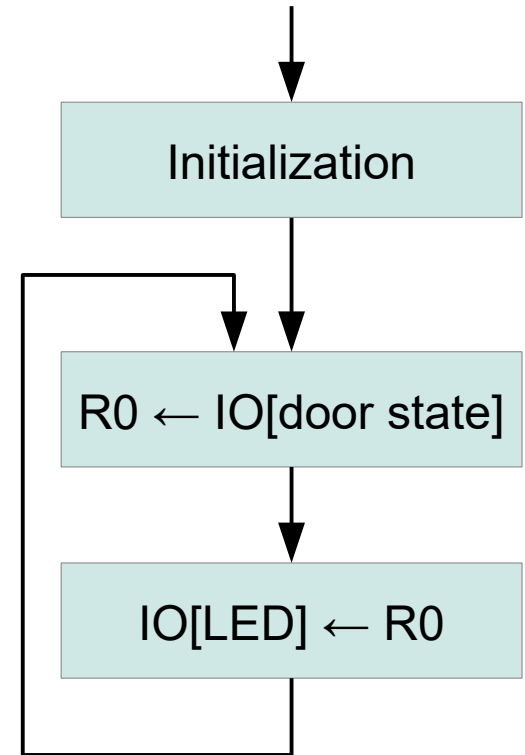


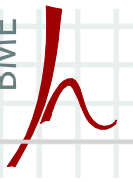
Algorithms

- The program has two parts:
 - Main program
 - Interrupt service routine
- The purpose of the main program:
 - Monitors the state of the door
 - Adjusts the LEDs accordingly
- The purpose of the interrupt service routine
 - Handles card events and key presses
 - Opens the door, if the correct code is given

The main program

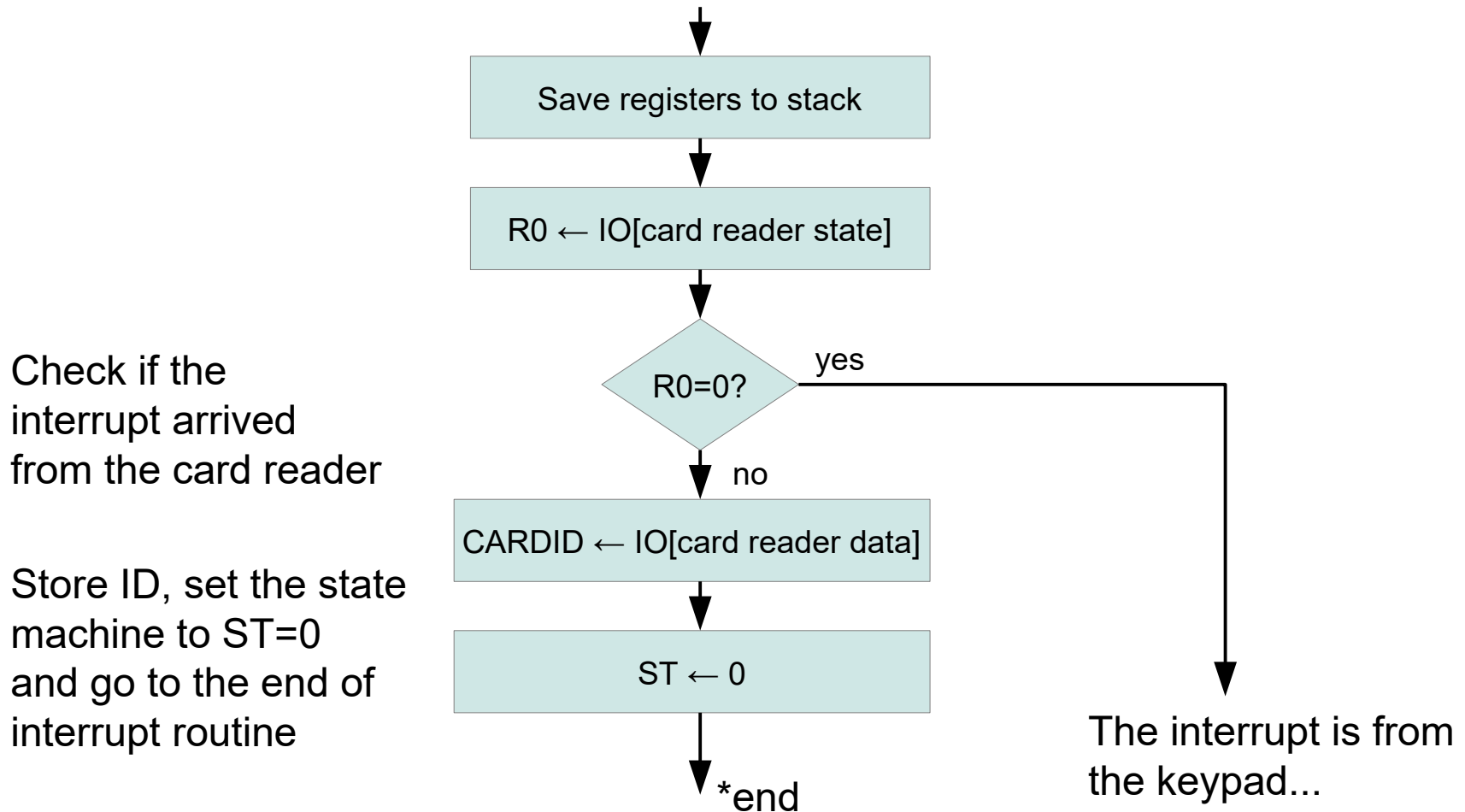
- Initialization:
 - Set stack pointer
 - Set the value of ST (state variable)
 - Enable interrupts
- In an infinite loop:
 - Ask the door for the state
 - Set the LED-s accordingly





The interrupt service routine

- First, we have to find out the interrupt source (polling!)
- Then, we apply the transition on the state machine



The interrupt service routine

The interrupt is from the keypad...

Read keycode from keypad

$R0 \leftarrow \text{IO}[\text{keypad data}]$

Are we in the initial state?

ST=255?

yes

no

Calculate the address of
the correct digit

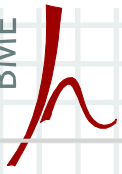
$R1 \leftarrow \text{start of array} + \text{CARDID} * 4 + \text{ST}$

Read the correct digit
from the memory

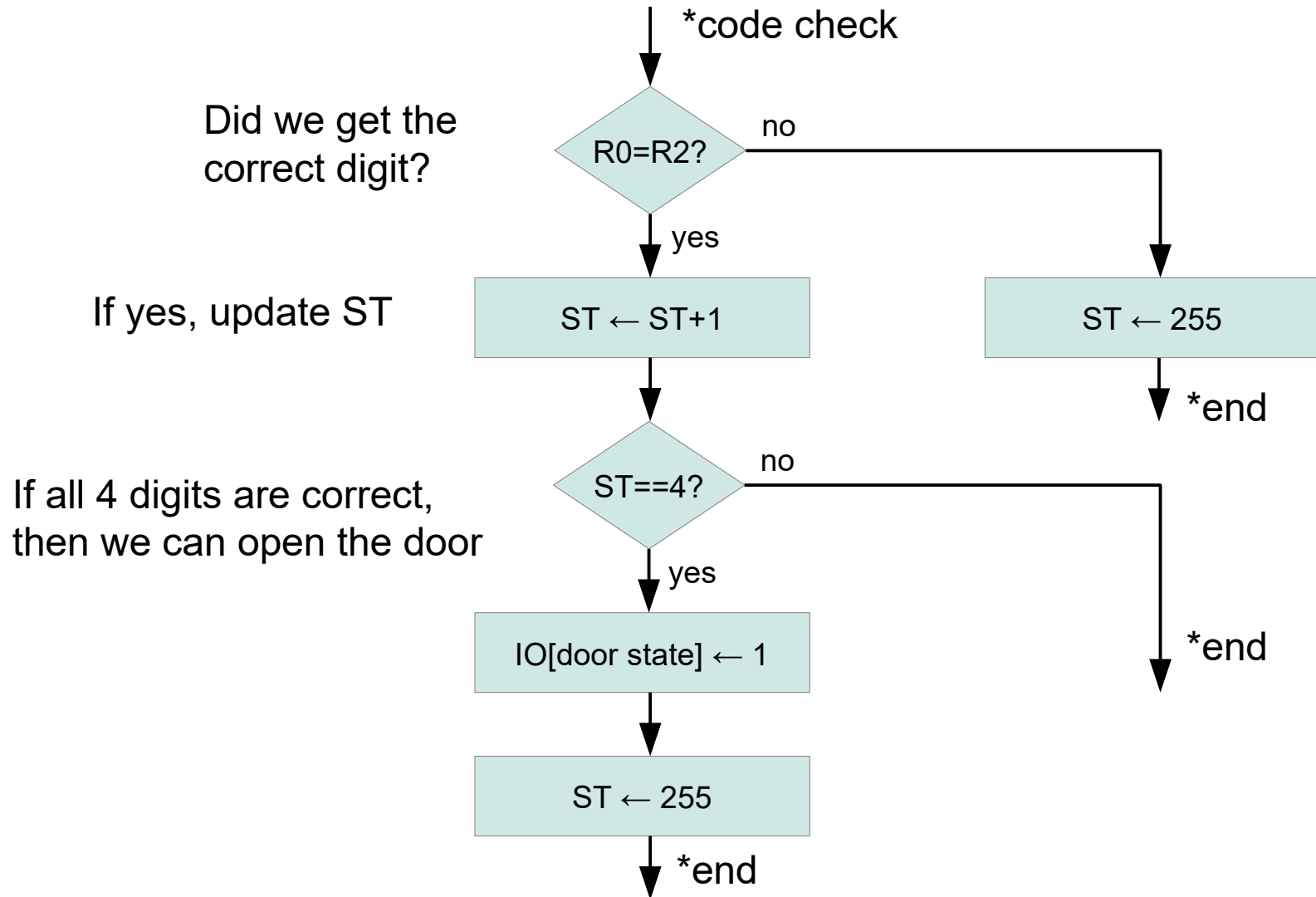
$R2 \leftarrow \text{MEM}[R1]$

*end

*code check

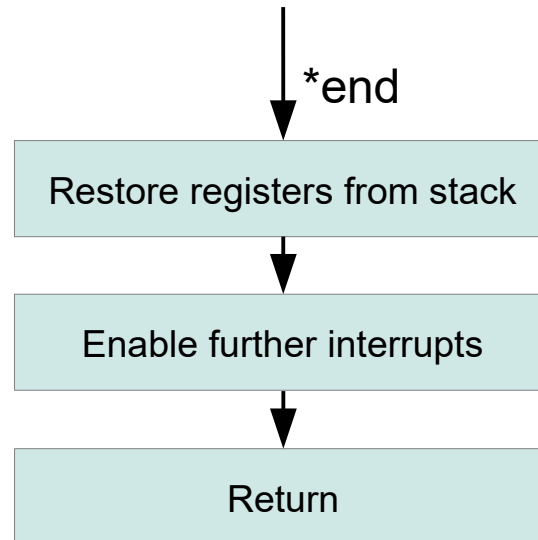


The interrupt service routine



The interrupt service routine

The end of interrupt service routine:



Placement decisions

■ Content of the RAM

- 1-byte integer variable ST: **2000h** (first byte of the RAM)
- 1-byte integer variable CARDID: **2001h** (second byte of the RAM)
- Initial stack pointer: **3FFFh**, the highest address of the RAM, since it grows downwards

■ Content of the ROM

- The program, starting at **0000h**
- The interrupt service routine, starting at **1000h**
- The array of correct codes for each card
 - Size: **400h**
 - Can be placed anywhere, where space is available
 - Let us put it to **500h**
 - The main program is small, it does not go beyond **500h**
 - The end of the array is **8FFh**, which is below the interrupt routine

Assembly implementation

- Implementation of the main program:
- | | |
|----------------------------|----------------------------|
| ORG 0000h | start code at this address |
| SP ← 3FFFh | set stack pointer |
| MEM[2000h] ← 255 | ST=255 (Initial state) |
| EI | enable interrupts |
| label: R0 ← IO[4Ch] | read door status |
| IO[5Dh] ← R0 | update LEDs |
| JUMP label | jump back to reading |

Assembly implementation

- Array of 4-byte codes for each card

ORG 500h

codes: DB 1, 3, 4, 7
DB 5, 7, 2, 9
DB 8, 2, 0, 8
DB 3, 1, 8, 9
...

Assembly implementation

- Implementation of the interrupt service routine:

ORG 1000h

CPU jumps at this address on interrupt

PUSH R0

save R0/R1/R2 register to stack

PUSH R1

PUSH R2

R0 ← IO[2Eh]

read card reader state

JUMP keycode IF R0==0

jump, if IT source is not the reader

R0 ← IO[2Fh]

read card ID

MEM[2001h] ← R0

save to variable CARDID

MEM[2000h] ← 0

ST=0

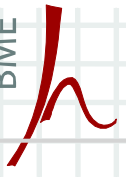
JUMP end

jump to the end of interrupt routine

Assembly implementation

- Implementation of the interrupt service routine:

keycode:	R0 ← I0[3Fh]	R0=new digit
	R1 ← MEM[2000h]	R1=ST
	JUMP end IF R1==255	we are not expecting a digit
	R2 ← MEM[2001h]	R2=CARDID
	R2 ← R2 * 4	R2=4*CARDID
	R2 ← R1 + R2	R2=4*CARDID+ST
	R2 ← MEM[R2+codes]	R2=the correct digit
	JUMP match IF R2==R0	Jump if new digit is correct
	MEM[2000h] ← 255	state machine: ST=225 (init. state)
	JUMP end	and go to the end of interrupt
match:	R1 ← R1+1	If correct, increment ST
	MEM[2000h] ← R1	and save ST to memory
	JUMP end IF R1<4	If this is not the last digit, go to the end
	I0[4Ch] ← 1	If last digit, open the door
	MEM[2000h] ← 255	state machine: ST=225 (init. state)
	JUMP end	Go to the end of interrupt



Assembly implementation

- Implementation of the interrupt service routine:

end:	POP R2	restore registers in a reverse order
	POP R1	
	POP R0	
	EI	enable further interrupts
	RET	return from interrupt

Variations

- Can we solve the problem without RAM?
 - Yes.
 - But we lose the stack
 - No more interrupts!!!
 - Main program continuously polls:
 - The state of the door (to update LEDs)
 - The state of the card reader
 - The state of the keypad
 - We can not put ST and CARDID into the RAM
 - We can store them in registers permanently