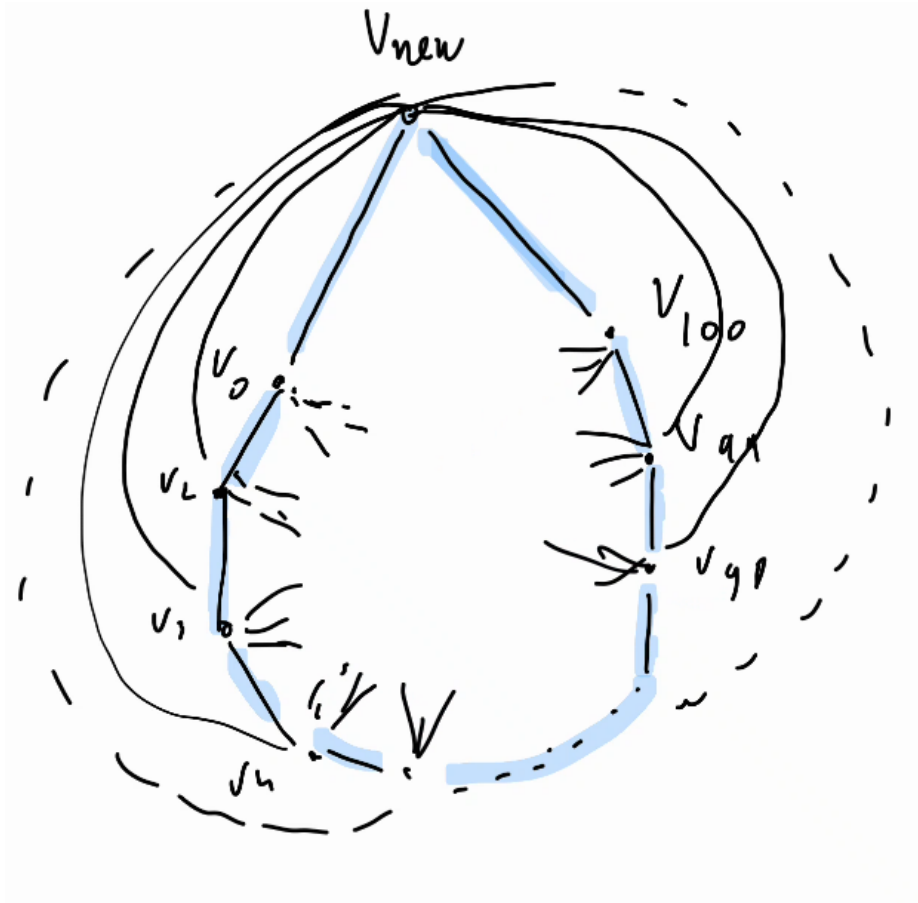


Solution2

2. (KD'2020, Problem 3) Prove that every simple 50-regular graph on 101 vertices has a Hamiltonian cycle. (Out of the 28 submissions, 5 solved this correctly).

Proof:

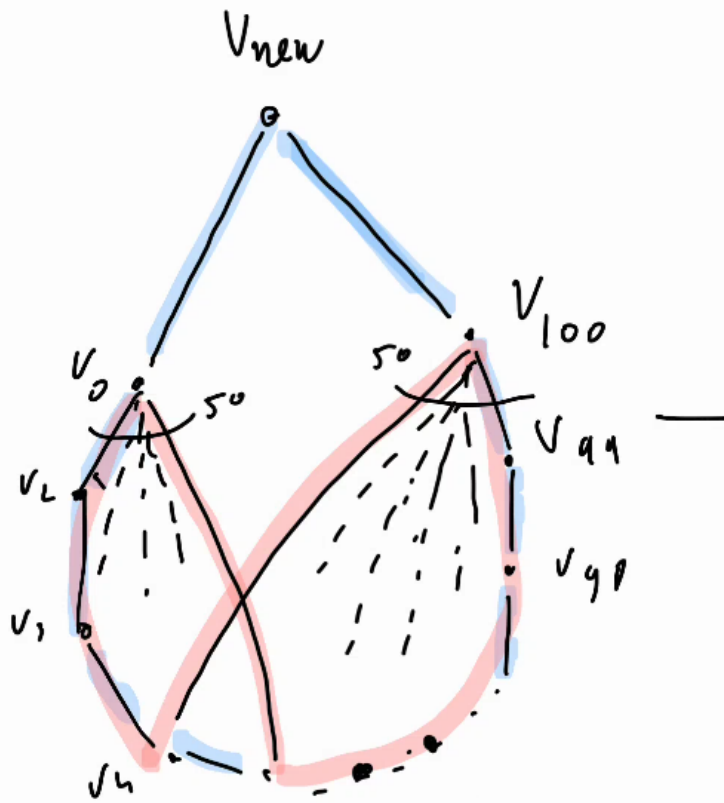
Add vertex and connect it to every other vertex so you get G' with $n=102$ and all the old vertices have $\deg=51$ and the new one 102. Now the Ores condition holds ($51+51=102=n$) $\Rightarrow G'$ is H.



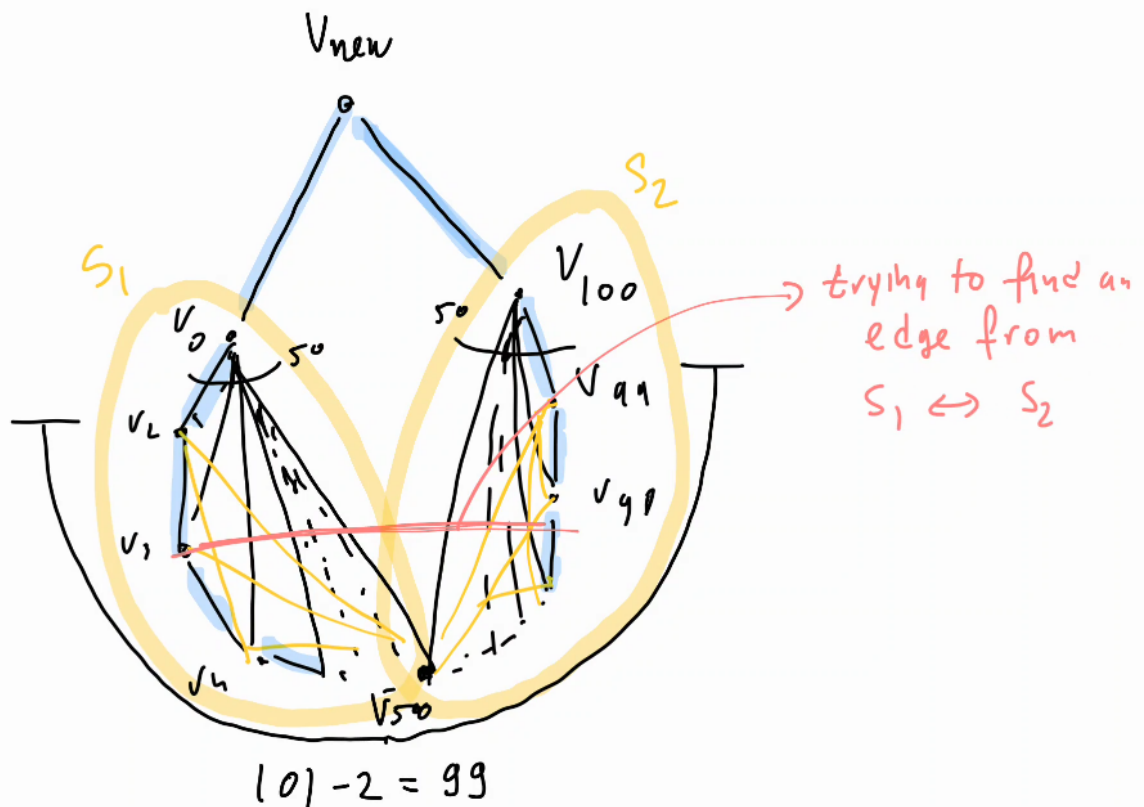
Now we have two options how the edges of v_0 and v_{100} are laid out:

1. They cross each other \Rightarrow immediately H_c .

As mentioned in solution of 1, if one edge crosses you will always find the cut where the 2 sets meet, so you can form the H_c .

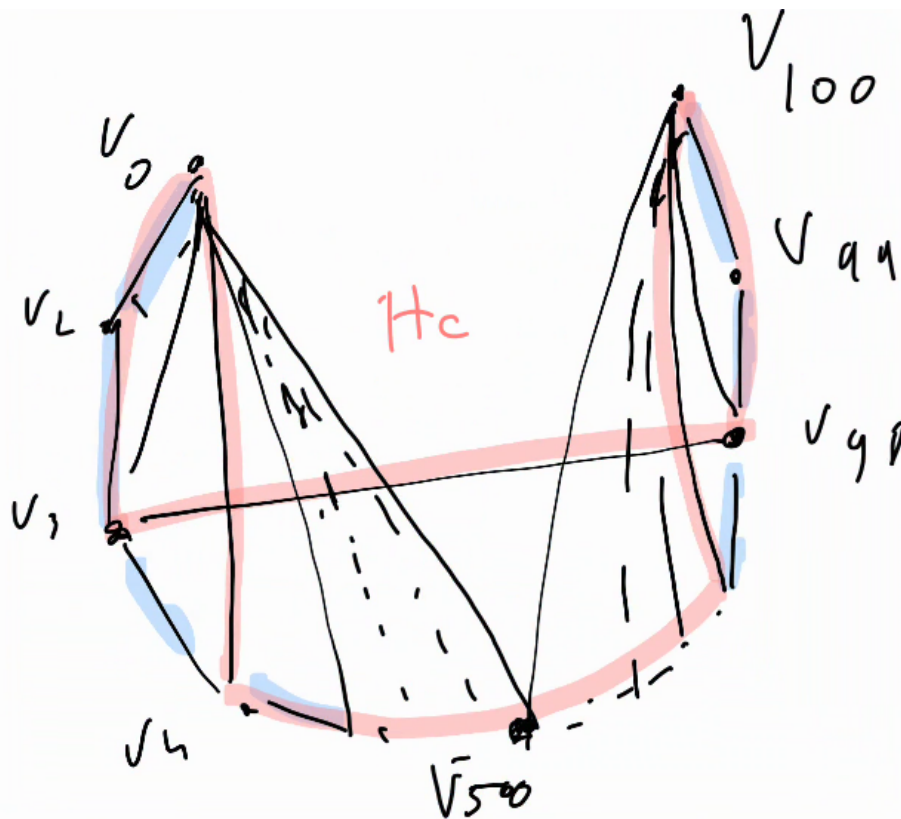


2. The edges are very packed together so they don't cross each other. ($50+50=100-2+1$)
 The edges will met at v_{50} but they wont cross.
 Now we use the 50-regular property. The idea we find and edge from the first subgraph to the second as shown in picture.



We know such an edge should exists because the only way S_1 has no edges going out of

it is if you have a K_{50} , but the vertex v_{50} connects to S_2 which implies that S_1 can not form a $K_{50} \Rightarrow$ there is an edge between v_i and v_j ($0 < i < 50, 50 < j < 100$). Using that edge you can form the H_c as shown below (You can also remove the new vertex we added cause we only needed it to prove that the G has a H_p):



Note:

This is just the Nash-Williams theorem for $k=50$. What we did works generally for any k .

Every k -regular graph on $2k+1$ vertices is Hamiltonian.