Stacks Overview

- A stack is a data structure for storing...
 - ...an *ordered* collection of items...
 - ...such that items can only be inserted and removed at one end (called the top)
- A stack is a *last-in-first-out* (LIFO) data structure
 - ◆ Entries are taken out of the stack in the *reverse* order of their insertion
- Another name for a stack is a *pushdown store*

Stacks

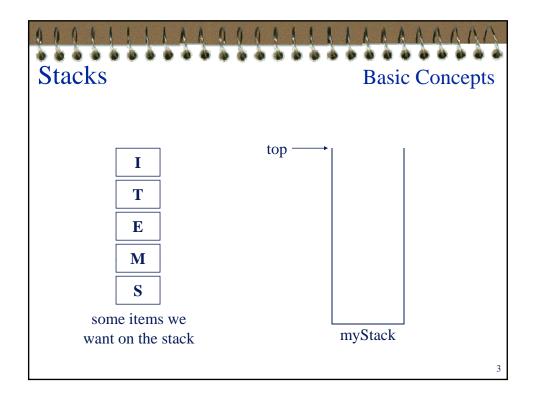
When we say that a stack is "ordered"...

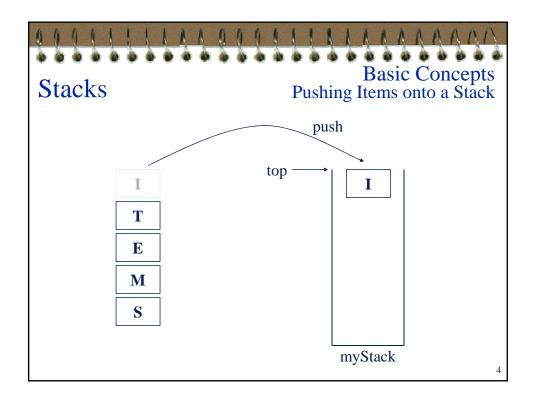
...we mean that there is some positional order that is significant with respect to how items can be accessed

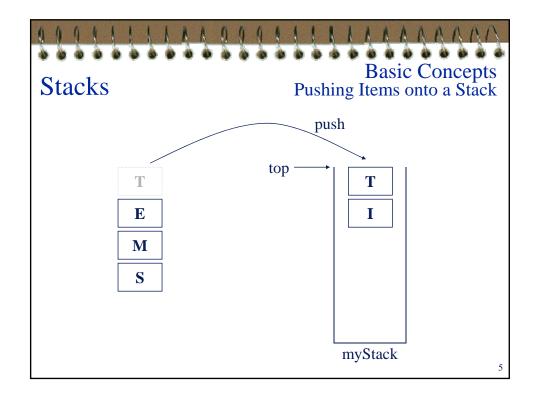
I.e., there is an item that can be accessed first, another second, another third, and so forth

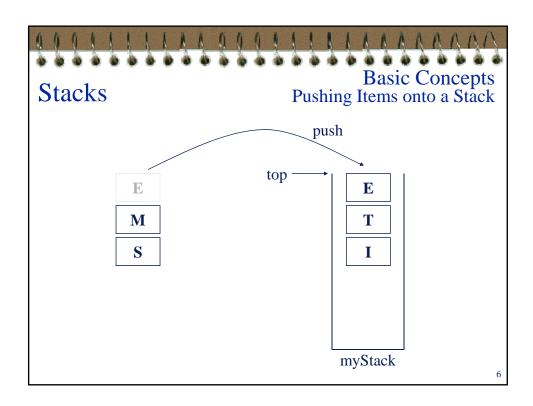
...we don't mean that items are sorted

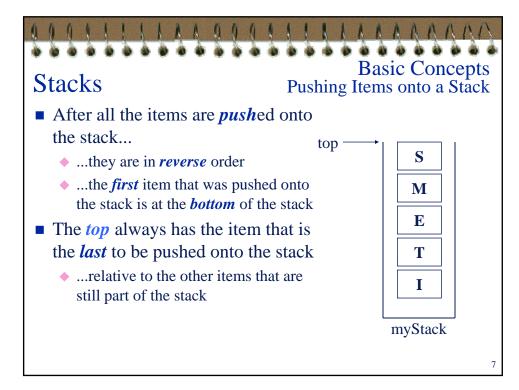
In fact, it is not even necessary that the items are sortable (can be compared)

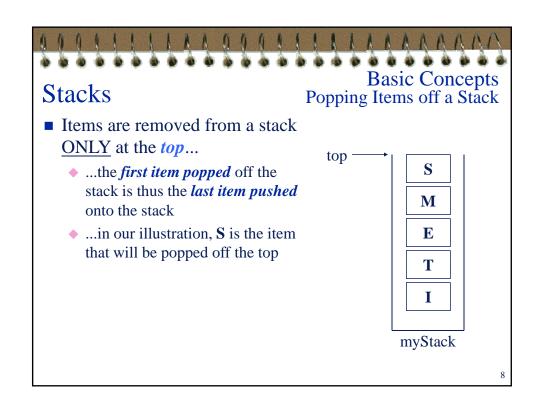


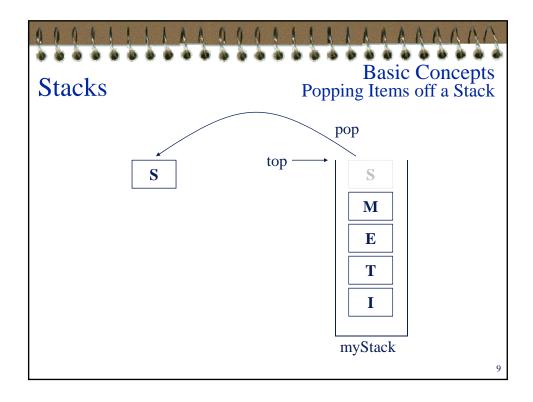


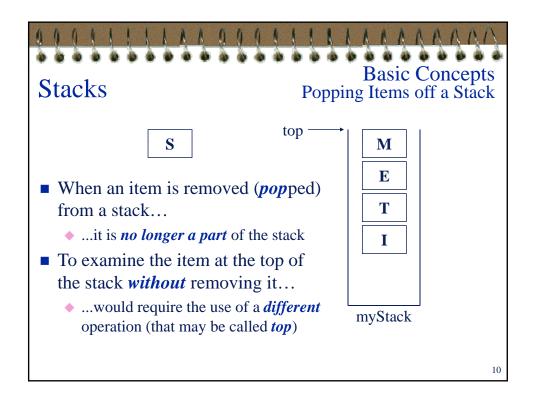


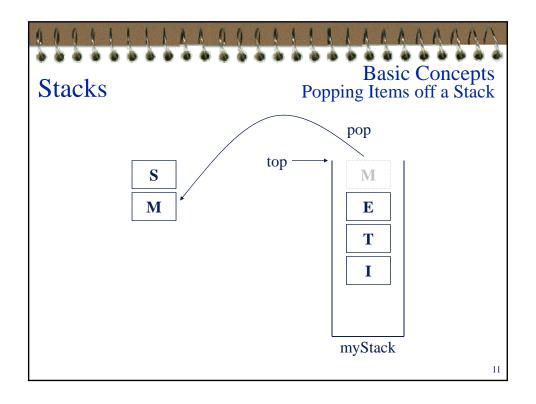


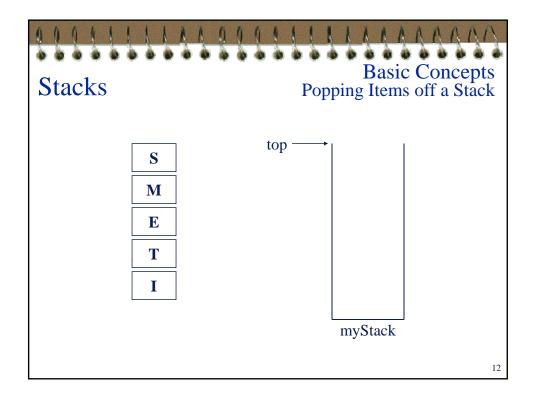














Basic Concepts Accessing the Stack

- At any one time ...
 - ...we can only access the *top item* of the stack
- If you need to access the third item (say) ...
 - ...we must first pop the two that are "above" it
- Textbook authors use as an analogy...
 - ...a PEZ dispenser which holds little rectangular hard candies
 - ...to get the third candy you have to remove the first two from the dispenser

13

Stacks Stack Errors

- There are *two common errors* that can occur when accessing a stack:
 - ♦ Stack Underflow
 - > The condition resulting from trying to *pop* (remove) an item from an *empty* stack
 - ♦ Stack Overflow
 - > The condition resulting from trying to *push* an item (add an item to) a *full* stack



- In order to avoid a stack underflow...
 - ...a stack ADT should provide a function that tests whether a stack is empty
 - > Appropriately, this would be a public member function in a C++ class
- In order to avoid a stack overflow...
 - ...we should check that the stack is not full before pushing a new item
 - > This implies a *fixed capacity* stack
 - > We could use a function that tests whether a stack is full
 - Implies that such a function is available
 - > We could also compare the *stack capacity* with its *current size*
 - Implies that we have the means to obtain the two values

14

Stacks

Stack Template Class

- Since we are likely to need stacks of different kinds of things...
 - ...it makes sense to implement the stack with a template class
- Two implementations of a stack template class...
 - ...are shown in Section 7.3 of the textbook
 - ...one using a static array, the other using a linked list
- The stack template class implements some general operations required of a stack



Stack Template Class

- The stack template class has the following member functions:
 - ◆ constructor → creates an empty stack
 - void push(const Item& entry)
 - void pop()
 - Item top() const
 - size_type size() const
 - bool empty() const

17

Stacks

Usefulness of Stacks

- Although a stack is a very simple concept, it has several very useful applications → examples:
 - ◆ Run-time stack
 - > When one function calls another, the state (activation record) of the calling function gets pushed onto the run-time stack
 - > When a function ends, control returns to the function that called it by popping the calling function's state off the stack
 - Stacks are also used by compilers to check the syntax of source code



Putting Stacks to Use Reversing a Word

- One simple application of a stack is...
 - ...to reverse the order of the characters in a word
 - > Example: if the user typed in ESIOTROT, the program would output TORTOISE
- While this is a trivial example...
 - ...it demonstrates the common mechanism for putting a stack to use

19

2221111122222111112222222

Stacks

Putting Stacks to Use Reversing a Word

- A pseudocode for reversing a word:
 - While (there are more characters of the word to read)
 - Read a character and push it onto the stack
 - While (the stack is not empty)
 - Output the top character and pop it off the stack
- This is exactly the process...
 - ...portrayed in the illustration shown earlier
 - ...where the word involved is **ITEMS**



Putting Stacks to Use Checking for Balanced Parentheses

- An application that is closely related to evaluating arithmetic (and Boolean) expressions is...
 - ...checking for balanced parentheses
 - ...which means that for every ' (' in an expression there is a matching ')'
- One way of doing it is...
 - ...to use stacks
 - ...but there are alternatives

21

Stacks

Putting Stacks to Use Checking for Balanced Parentheses

- Consider the string (X+Y*(Z+7))*(A+B))
 - Each left parenthesis has a corresponding right parenthesis
- The **is_balanced** function would take a string like this as an argument and...
 - ...return *true* if the parentheses balance
 - ...return false if they don't

Putting Stacks to Use Checking for Balanced Parentheses

■ The algorithm is simple...

```
Scan the characters in the string from left to right if (character read is a left parenthesis)

Push it onto the stack else if (character read is a right parenthesis)

Pop a left parenthesis off the stack else

Do nothing (i.e., simply discard the character read)
```

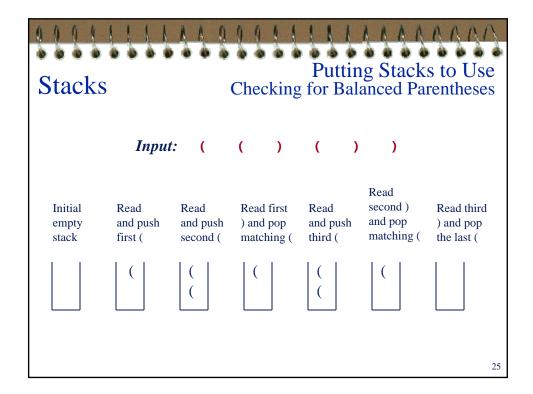
■ (continued)

23

Stacks

Putting Stacks to Use Checking for Balanced Parentheses

- The parentheses are *balanced*...
 - ...if the stack is *empty* when the **is_balanced** function has read the entire string
- The parentheses are *not balanced*...
 - ...if the is_balanced function produces a stack underflow error during processing
 - > I.e., algorithm dictates a pop but stack is found to be empty
 - ...if there are any *left parentheses left in the stack* (stack not empty) when the **is_balanced** function
 has read and processed the entire string
 - > I.e., stack is not empty when the algorithm ends



Putting Stacks to Use

Stacks

Putting Stacks to Use Evaluating Arithmetic Expressions

■ Suppose we want a program that would evaluate an arithmetic expression such as...

```
((((12 + 9) / 3) + 7.2) * ((6 - 4) / 8))
```

- The expression will consist of...
 - ...non-negative integer or floating-point numbers
 - ◆ ...along with the basic arithmetic operators (+, -, *, /)
 - ...as well as the parentheses themselves



Putting Stacks to Use Evaluating Arithmetic Expressions

Additional Assumptions

- The expression is *formed correctly*
 - (expression operator expression)
 - > An expression can be an integer or floating-point number or another expression
 - > An operator can be +, -, *, /
- Each expression is *fully parenthesized*
 - Each operation has a pair of matching parentheses surrounding its arguments

27

3831111111113831111111111111

Stacks

Putting Stacks to Use Evaluating Arithmetic Expressions

Consider How We'd Do It Manually

- **■** (((6+9)/3) * (6-4))
 - Evaluate the *innermost* expressions
 - (6+9) and (6-4)
- **(**(15/3) * 2)
 - ◆ Evaluate the *innermost* expression > (15 / 3)
- -(5*2)
 - Evaluate the *innermost* expression \rightarrow 10



Putting Stacks to Use Evaluating Arithmetic Expressions

Algorithm Design Considerations

- Trying to keep track of *innermost* expressions is *complicated* and *unnecessary*
- We can simply evaluate the *leftmost of the innermost* expressions each time through
 - ...thus eliminating the need to keep track
- We keep evaluating the *leftmost innermost* expression until we are finished

29

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Stacks

Putting Stacks to Use Evaluating Arithmetic Expressions

Previous Example Revisited

- \blacksquare (((6+9)/3) * (6-4))
- **■** ((15 / 3) * (6 4))
- **(5 * (6 4))**
- **(5 * 2)**
- **1**0



Putting Stacks to Use Evaluating Arithmetic Expressions

Algorithm Design Considerations

- The *end* of the *leftmost of the innermost expressions* is always indicated by a *right parenthesis*
 - Useful for us to find the expression to be evaluated next
- We will use two stacks:
 - One stack contains *numbers* and *intermediate values*
 - > We will refer to it as the *numbers stack*
 - ◆ The other stack contains *operators* from the input
 - > We will refer to it as the *operators stack*

31

Stacks

Putting Stacks to Use Evaluating Arithmetic Expressions

Algorithm Design Considerations

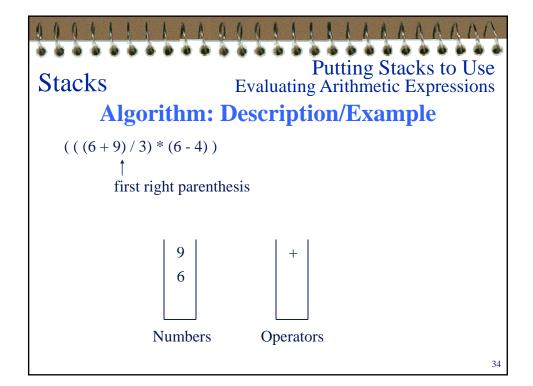
- Since a stack is a LIFO structure...
 - ...it will turn out that the correct two numbers are on top of the numbers stack...
 - ...at the same time that the *appropriate operation* is at the *top of the operators stack*



Putting Stacks to Use Evaluating Arithmetic Expressions

Algorithm: Description/Example

- Start by reading the input *up until the first right* parenthesis
 - ◆ Each *number* along the way is *pushed onto the numbers stack*
 - ◆ Each *operator* along the way is *pushed onto the operators stack*
 - (left parentheses are simply discarded)
 - (each right parenthesis serves as a flag for some
 "computational" actions to be taken to be described)

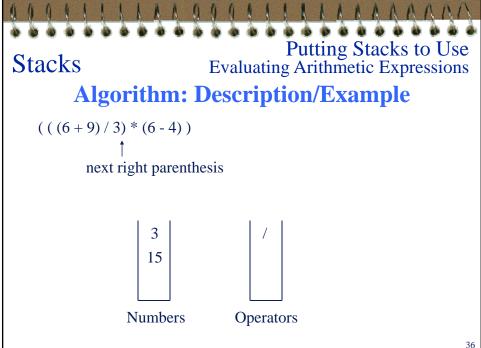


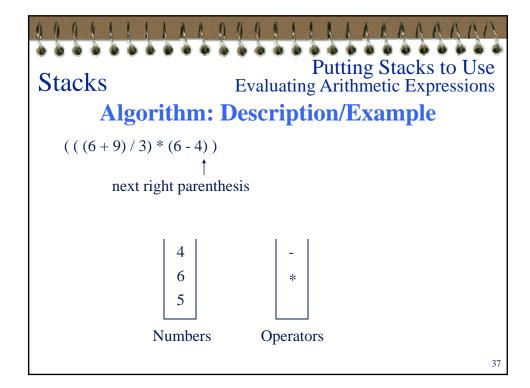


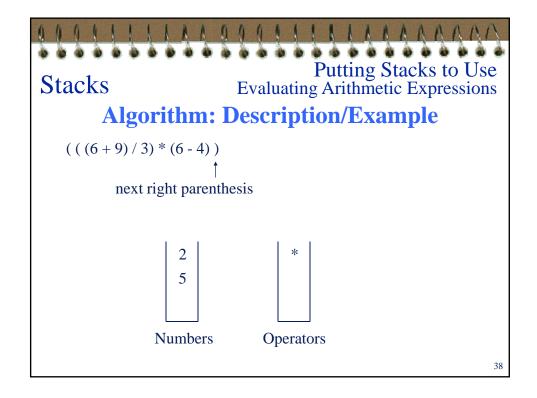
Putting Stacks to Use **Evaluating Arithmetic Expressions**

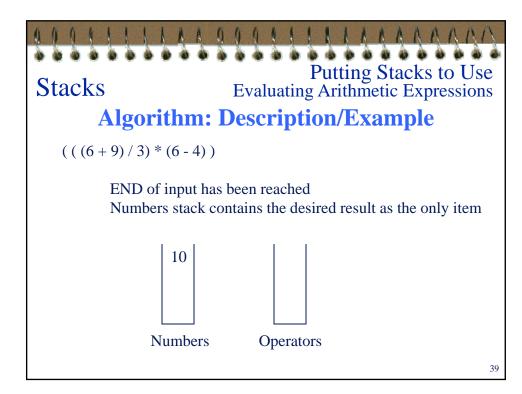
Algorithm: Description/Example

- Each time we see a *right parenthesis*, we combine the top 2 numbers (from the numbers stack) with the top operator (from the operators stack)
 - **♦** Continuing with our example: 6 + 9
 - > Note that the *left* argument should be the *second* number removed from the (numbers) stack → makes a big difference with some operators (e.g., "6 - 9" vs "9 - 6" or "6/9" vs "9/6")
- Perform the operation to get the *intermediate result* (6+9=15), and push it back onto the numbers stack









Putting Stacks to Use Evaluating Arithmetic Expressions

Algorithm: Summary

- When a number is encountered in the input...
 - ...push it onto the numbers stack
- When an operator is encountered in the input...
 - ...push it onto the operators stack
- When a right parenthesis is encountered in the input...
 - ...pop the top two numbers and combine them with the top operation
 - ...evaluate the intermediate result and push it onto the numbers stack



Putting Stacks to Use Evaluating Arithmetic Expressions

Algorithm: Summary

- When a left parenthesis or blank is encountered in the input...
 - ...simply ignore (discard) it
 - NOTE: A more complete algorithm would need to process the left parentheses in some way to ensure that each left parenthesis is balanced by a right parenthesis → our assumption is that the input is *correctly formed*)
- When the END of input is encountered...
 - ...pop the numbers stack to obtain the desired result

41

Textbook Readings

- Chapter 7
 - ♦ Section 7.1
 - Section 7.2
 - Section 7.3