

Module 6: Exact-Mapping for Entity Resolution

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joint with Brian Kunding and Jerry Reiter

Reading

- ▶ Binette and Steorts (2020)
- ▶ Newcombe et al. (1959)
- ▶ Fellegi and Sunter (1969)
- ▶ Kunding, Reiter, and Steorts (2024)

Notation

Assume two databases A and B (or sometimes called X_1 and X_2 .)

- ▶ The relative size of each database is N_1 and N_2 .
- ▶ Assume there are duplicates across the databases but not within them. This is called a bipartite record linkage assumption.
- ▶ Assume $f = 1, \dots, F$ fields (or attributes).
- ▶ Let L_f denote the number of categories for field f .

Motivation

- ▶ Record pairs that refer to the same entity should be similar.
- ▶ Records pairs that refer to different entities should be dissimilar.

We can compare record pairs using similarity scores (or distance functions).

Examples: Jaccard, Edit, Jaro, Jaro-Winkler.

Comparison Vectors (or Data)

This motivates the comparison vector or comparison data.

Consider

$$\gamma_{ij} = (\gamma_{ij}^1, \dots, \gamma_{ij}^F),$$

where

γ_{ij}^f compares field f for record $i \in A$ and $j \in B$.

Collect all the comparison vectors as

$$\gamma = \{\gamma_{ij}\}_{i=1, j=1}^{N_1, N_2}$$

Comparison Vectors (or Data)

The above notation is used in the literature as it is compact and short!

What do these vectors look like in practice?

Comparison Vectors

Let

$$i = 1, 2, \dots, N_1 \times N_2$$

enumerate the set of all record pairs in $A \times B$.

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- ▶ Each γ_i^j compares the j th field of the records.

Comparison Vectors

How can we visualize the comparison vectors?

$$\gamma_1 = (\gamma_1^{(1)}, \gamma_1^{(2)}, \dots, \gamma_1^{(F)}) \quad (1)$$

$$\gamma_2 = (\gamma_2^{(1)}, \gamma_2^{(2)}, \dots, \gamma_2^{(F)}) \quad (2)$$

$$\vdots \quad (3)$$

$$\gamma_{(N_1 \times N_2)} = (\gamma_{(N_1 \times N_2)}^{(1)}, \gamma_{(N_1 \times N_2)}^{(2)}, \dots, \gamma_{(N_1 \times N_2)}^{(F)}) \quad (4)$$

Let

$$\gamma = (\gamma_1^{(1)}, \gamma_2^{(2)}, \dots, \gamma_{(N_1 \times N_2)}^{(F)})$$

Exact Mapping

We will walk through an approach to create a more efficient representation of the comparison vectors (but not do any dimension reduction).

Notation

- ▶ Let P be the number of exact agreement patterns in γ , which is bounded above by $\prod_{f=1}^F (L_f + 1)$.
- ▶ Consider the following function

$$h_f^{(i,j)} = l_{obs}(\gamma_{ij}^f) 2^{\gamma_{ij}^f + I(f>1) \times \sum_{e=1}^{f-1} (L_e - 1)}, \quad (5)$$

which maps a record pair for a field f to a unique integer.

Summing over fields $f = 1, \dots, F$ for record pair (i, j) results in

$$h^{(i,j)} = \sum_{f=1}^F h_f^{(i,j)}.$$

Notation

- ▶ Enumerate unique hashed agreement patterns from 1 to P .
- ▶ Denote each unique mapped agreement pattern as $h_p = (h_p^1, \dots, h_p^F) \implies$ that when record pair (i, j) has agreement pattern p , we write $\gamma_{ij} = h_p$.
- ▶ Collect all the agreement patterns as $P = \{h_p \mid p \in [P]\}$.

Notation

- ▶ When performing computation, it will be useful to represent a more representative version of h_p , known as a one hot encoding.
- ▶ $e(h_p)$ denotes the $\sum_{f=1}^F L_f$ length vector where the $\ell + \sum_{f=1}^F L_f$ component is 1 when $h_p = \ell$ and otherwise is 0.

Example

Consider five fields with binary agreement patterns (and potential missingness). Suppose that records (5, 7) in have agreement pattern (1, 1, 1, NA, 2), which means

- ▶ agreement in the first three fields.
- ▶ missingness in the fourth field.
- ▶ complete disagreement in the fifth field.

Example

1. Find $h(\gamma_{5,7})$
2. Find $e(h(\gamma_{5,7}))$.

Example

Recall that

$$h_f^{(i,j)} = l_{obs}(\gamma_{ij}^f) 2^{\gamma_{ij}^f + I(f > 1) \times \sum_{e=1}^{f-1} (L_e - 1)} \implies$$

$$h_1^{(5,7)} = 2^1 \tag{6}$$

$$h_2^{(5,7)} = 2^3 \tag{7}$$

$$h_3^{(5,7)} = 2^5 \tag{8}$$

$$h_4^{(5,7)} = 0 \tag{9}$$

$$h_4^{(5,7)} = 2^{10} \tag{10}$$

This implies $h^{(5,7)} = 1066$. Assume this maps to unique integer 42.

Example

How do we create the one hot encoding?

For a binary comparison, we have the following options:

1. (1,0): represents complete agreement of the record pairs.
2. (0,1): represents complete disagreement of the record pairs.
3. (0,0): represents missingness in the record pairs.

Given this,

$$e(h^{(5,7)}) = (1, 0, 1, 0, 1, 0, 0, 0, 0, 1).$$