

CODING PATTERNS FOR TECHNICAL INTERVIEWS AND COMPETITIVE PROGRAMMING

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Let's Connect!

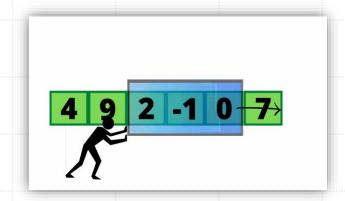
WHAT YOU'LL LEARN TODAY

- Overview of common coding patterns seen in easy/intermediate Leetcode questions for interviews or competitive programming.
- Two Pointer/Sliding Window
- Binary Search
- Breath First Search (BFS) and Depth First Search(DFS)
- Slight overview of Dynamic Programming

TWO POINTERS (SLIDING WINDOW)



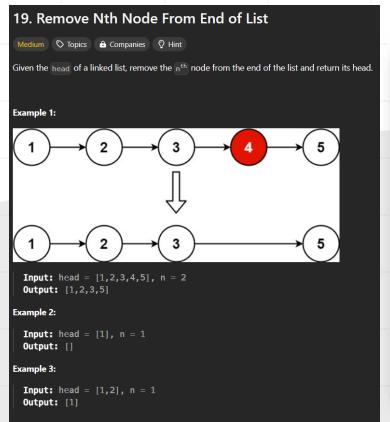
- Refers to using two variables to traverse through a Data Structure (commonly arrays, strings, or linked lists)
- Commonly used to solve problems such as detecting pairs that involve a sorted list
- Usually optimizes solutions to linear time O(n)

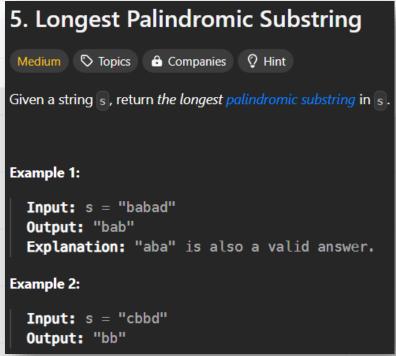


- Common variations of this include:
 - Slow & Fast
 - Left & Right
 - Small & Big
 - You get it...

LET'S PRACTICE!

- 19. Remove Nth node from end of list
- 5. Longest Palindromic
 Substring





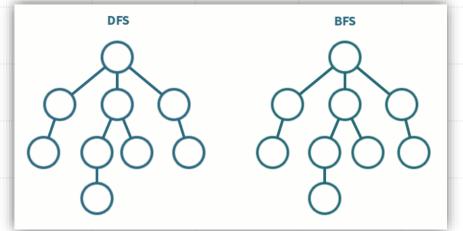
TRAVERSING GRAPHS

Depth First Search (DFS)

- Explores as far down a branch as possible, then backtracking to the next node.
- This backtracking is normally done using a stack or recursion.
- Often used for tree and graph traversal.

Breath First Search (BFS)

- Explore a graph by level, visiting all the nodes at the current height before moving on to the next one
- Useful in tree problems that involve processing nodes by level (or height)



LET'S PRACTICE!

- 100. Same Tree
- 101. Symmetric Tree

Given the roots of two binary trees p and q, write a function to check if they are the same or not.

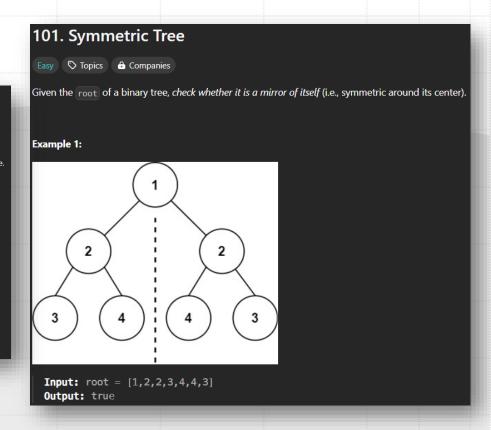
Two binary trees are considered the same if they are structurally identical, and the nodes have the same value.

Example 1:

1
1
1
1
1
Input: p = [1,2,3], q = [1,2,3]
Output: true

100. Same Tree

 Hint: These are not the same



BINARY SEARCH

- Efficient for finding an element in a sorted list by repeatedly dividing the search in half.
- It compares the target value with the middle of the current search range and eliminating half of the search at each step.
- This is ideal for search questions, in which you are given a sorted list.
- It is very time efficient (O(log n)), much faster than a normal O(n) search.



LET'S PRACTICE!

- 69. Sqrt(x)
- (yes, this is a binary search question)



DYNAMIC PROGRAMMING

- Scenario: Imagine solving the typical Fibonacci number problem. It works for small numbers.
- When the input gets too large, this solution is very slow (O(n!))
- Dynamic programming comes in when subproblems like this one overlap. It involves saving solutions to the subproblems that build up to the solution of the main problem

```
public int slowfibonacci(int n)
{
    // Base Case
    if (n <= 1)
        return n;

    // Working recursive fibonacci
    return slowfibonacci(n - 1) + slowfibonacci(n - 2);
}</pre>
```

DYNAMIC PROGRAMMING

- We call this process memoization, and it involves using an extra data structure to save the solutions to the smaller subproblems, this makes solving the big problems faster (Generally O(n))
- There's two types of DP:
 - Top down
 - Bottom up

```
public int coolfibonacci(int n)
{
    // Base case
    if (n <= 1) return n;

    // Create an array to store Fibonacci values up to n
    int[] dp = new int[n + 1];
    dp[0] = 0; // Fibonacci(0) = 0
    dp[1] = 1; // Fibonacci(1) = 1

    // Fill the dp array by iterating from 2 to n
    for (int i = 2; i <= n; i++) {
        dp[i] = dp[i - 1] + dp[i - 2];
    }

    return dp[n]; // Return Fibonacci(n)
}</pre>
```

PRACTICE PROBLEM (NOT LEETCODE)

- Doodoo Dynamics buys long steel rods, cuts them into shorter rods, and sells them.
 Cutting doesn't have a cost and only happens in whole inches
- This table shows the current sale prices for different length rods

length i	1	2	3	4	5	6	7	8	9	10
price p_i	1	5	8	9	10	17	17	20	24	30

- Given a rod of length n, determine the maximum revenue that can be obtained

HINTS

• Which of these rods yields the most revenue?

length i

(e)

	price p_i	1	5	8	9	10	17	17	20	24	30			
9		1		8			5	5	6		8	1		
))							()							
(a)			(t)			(c)				(d)			
1	5	1	5	5	1		5	1	1	1	1	1 1		

(g)

10

(h)

5

(f)

SOLUTION

```
CUT-ROD(p, n)

1 if n == 0

2 return 0

3 q = -\infty

4 for i = 1 to n

5 q = \max(q, p[i] + \text{CUT-ROD}(p, n - i))

6 return q
```

