

DATA LINK CONTROL PROTOCOLS

Requirement & Objectives for Effective data Communication

- Frame synchronization
- Flow Control
- Error Control
- Addressing
- Control & Data on the Same Link
- Link Management

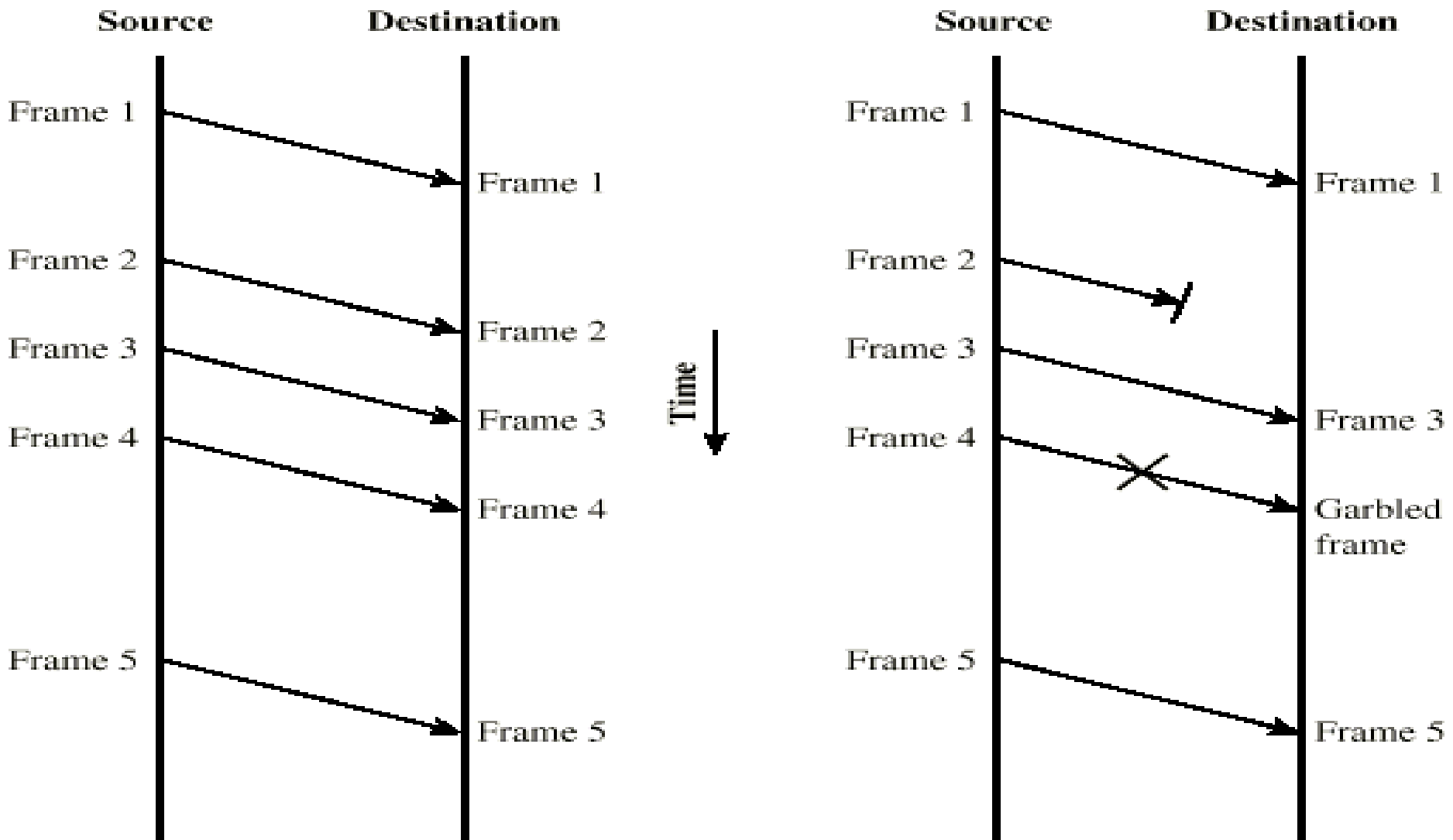
Data Link Layer – Primary Functions

- Flow control – regulate flow of data from sender to receiver
- Error control – detect and correct errors
- Framing – recognizing beginning and end of frames (blocks, packets)

Flow Control

- Ensuring the sending entity does not overwhelm the receiving entity
 - Preventing buffer overflow
- Transmission time
 - Time taken to emit all bits of a frame onto the medium (proportional to length of frame)
- Propagation time
 - Time for a bit to traverse the link between sender and receiver

Model of Frame Transmission



(a) Error-free transmission

(b) Transmission with losses and errors

Stop and Wait Flow Control

- Source transmits frame
- Destination receives frame and replies with acknowledgement
- Source waits for ACK before sending next frame
- Destination can stop flow by not sending ACK
- Works well for a few large frames

Fragmentation

- Large block of data may be split into small frames
 - Limited buffer size
 - Longer the transmission, more likely error. If error, must retransmit entire frame
 - Prevents one station occupying medium for long periods
- Stop and wait becomes inadequate⁷

Link Utilization

- $B = R \times (d/V)$

B-length of the link in bits

R – data rate

d – length or distance of the link

V – velocity of propagation

- $a = B/L$

a – variable

L – number of bits in the frame

➤ $a < 1$

Link Utilization

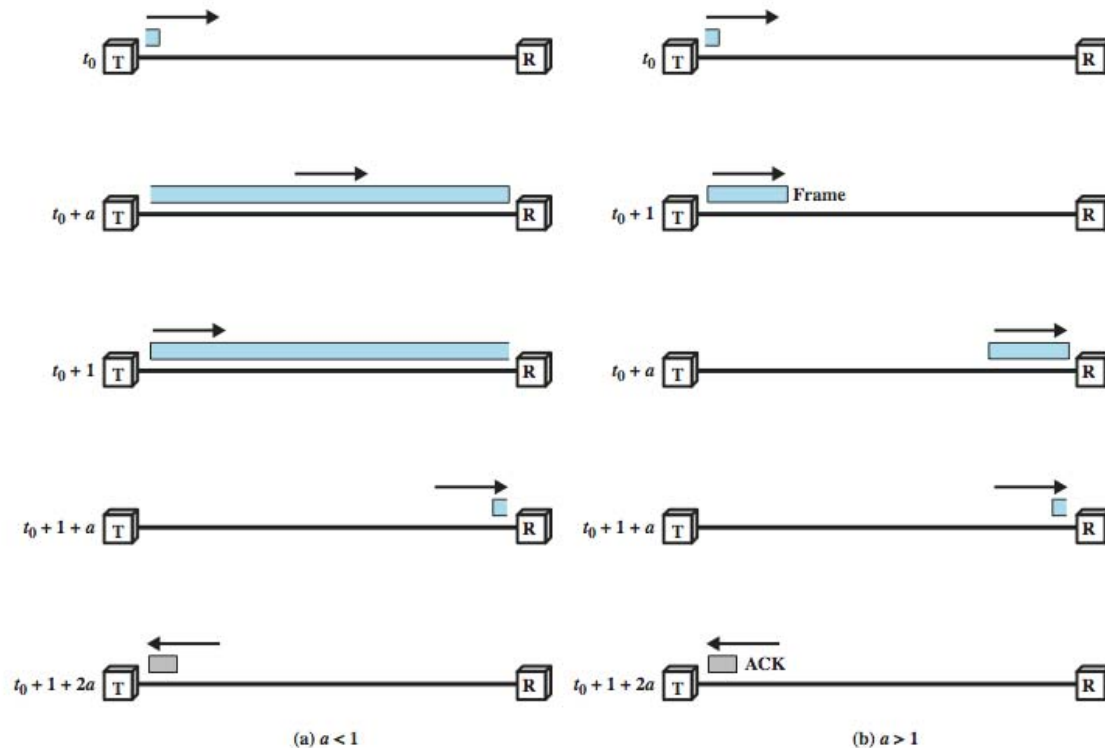


Figure 7.2 Stop-and-Wait Link Utilization (transmission time = 1; propagation time = a)

Drawback- Stop and Wait

- $a < 1$
 - ❖ propagation time is less than transmission time
 - ❖ first bit of the frame has arrived at the destination
before source has completed the transmission of frame

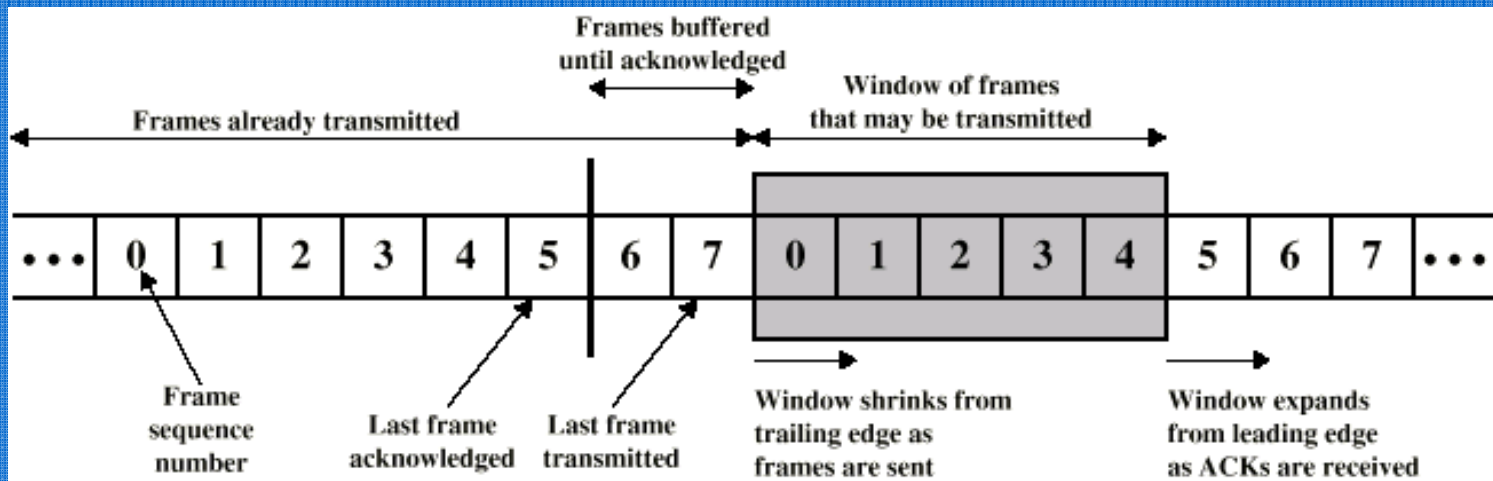
The link is inefficiently utilized

- $a > 1$
 - ❖ Propagation time greater than transmission time
 - ❖ Sender completes the transmission of frame before it reaches the destination

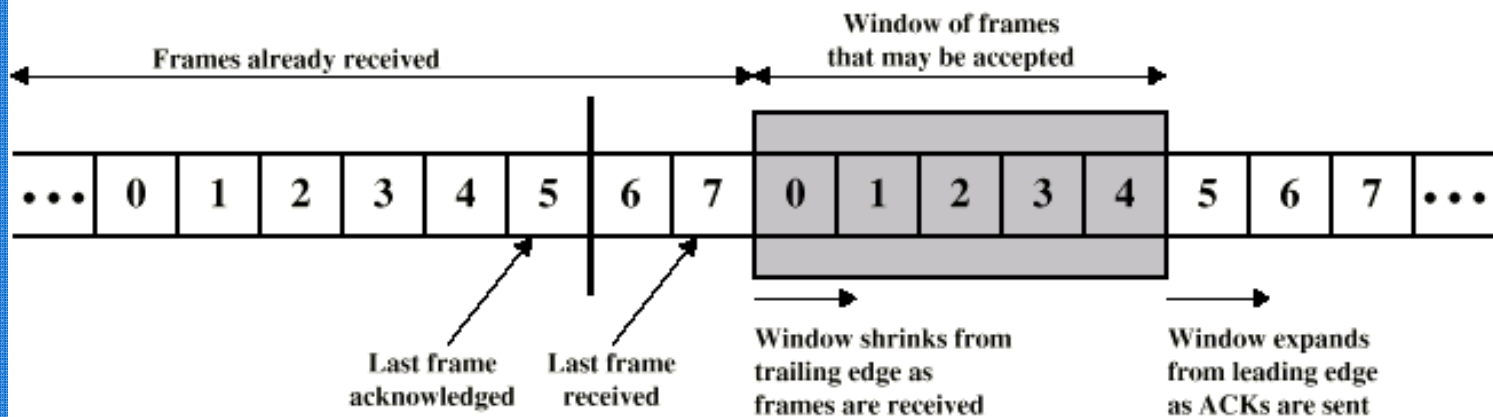
Sliding Windows Flow Control

- Allow multiple frames to be in transit
- Receiver has buffer W long
- Transmitter can send up to W frames without ACK
- Each frame is numbered
- ACK includes number of next frame expected
- ACK can also be used to acknowledge multiple frames

Sliding Window Diagram

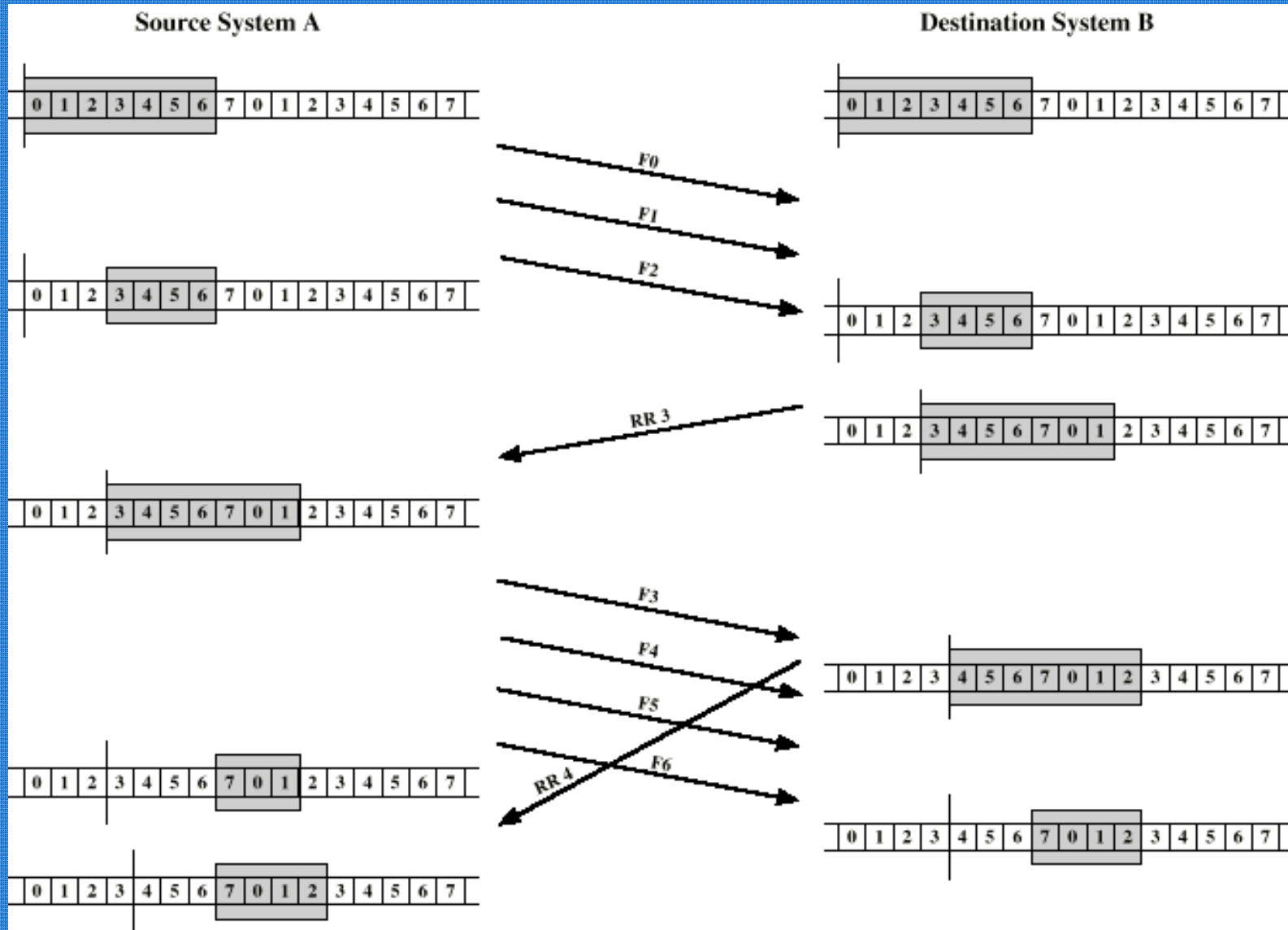


(a) Sender's perspective



(b) Receiver's perspective

Example Sliding Window



- Piggybacking
data and acknowledgement sent together
- RNR –receive not ready (RNR5)
- Sliding window is more efficient than stop and wait flow control
 - transmission link is treated as pipeline filled with frames in the transit