AK High-Dimensional Projection Structural Theory
Version 15.0: Collapse Structures, Group Simplification, and Persistent Projection Geometry

Atsushi Kobayashi (with ChatGPT Research Partner)

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Abstract

We present AK–HDPST v16.0, a two-layer functorial framework for controlled collapse in constructible one-parameter persistence over a field. The Core layer supplies machine-checkable statements; the [Spec] layer consists of auditable, non-expansive extensions used only under explicit hypotheses.

On the Core side we show that the exact Serre reflector $\mathbf{T}_\tau$, which deletes precisely the bars of length $\leq \tau$, is idempotent and 1-Lipschitz for interleavings, and admits a filtered lift C_τ unique up to filtered quasi-isomorphism with $\mathbf{P}_i C_\tau \simeq \mathbf{T}_\tau \mathbf{P}_i$. We certify collapse gates via a one-way bridge: if $\mathrm{PH}_1(F) = 0$ then $\mathrm{Ext}^1(\mathcal{R}(F), k) = 0$ under t-exact realizations of amplitude ≤ 1 (no converse used). Directed systems are audited by windowed comparison maps $\phi_{i,\tau}$; kernel/cokernel diagnostics (μ, u) detect failures of limit collapse even when all finite stages collapse. All comparisons are performed after truncation by the windowed protocol $t \mapsto$ persistence $\mapsto \mathbf{T}_\tau \mapsto$ compare on mutually exclusive, collectively exhaustive right-open windows. Policy-wise, deletion-type updates are non-increasing after truncation, whereas inclusion-type updates are non-expansive.

For pipeline control and reproducibility, a per-window δ-ledger aggregates continuation shifts, commutation defects, and numerical tolerances. A safety margin B–Gate$^+$ enforces $\mathrm{gap}_\tau > \sum \delta$ per window, and Restart/Summability paste windowed certificates into global ones. Versioned logs and cross-linked hashes enable auditability and exact re-runs.

The [Spec] layer (non-expansive after truncation) includes: a lax monoidal tensor/collapse with windowed energy bounds and a mono regime yielding post-collapse dominance; projection formula and base change transported to persistence via the windowed protocol; and Fukaya-category realizations with action filtration (continuation 1-Lipschitz; adding stops deletion-type). Mirror/Transfer pipelines contribute δ-controlled commutation, with a permitted-operations table mapped to the δ-ledger and B–Gate$^+$.

Scope and limitations: we work only with constructible 1D persistence over a field and do not claim $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$. The organizing principle is a sharp boundary between theorems and [Spec] contracts, enabling safe reuse, cross-domain exploration, and machine-checkable testing while preserving verified guarantees.

1 Chapter 1: Collapse — Operational Definition and Scope

Scope Box (Global Guard-Rails, short version). All statements in this chapter (and throughout) are made under the following guard-rails:

- Constructible 1D persistence over a *field* (coefficients vary but remain fields).
- All equalities and exactness claims are asserted *after collapse at the persistence layer*; filtered complex statements hold *up to filtered quasi-isomorphism (f.q.i.)*.
- Non-commutation and implementation errors are externalized in a δ-ledger with additive budget along pipelines.
- *Windowed (MECE) proof policy* with restart and summability; **windows are right-open** half-open intervals $[u, u')$.
- The bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ is used only in $D^{\mathrm{b}}(k\text{-mod})$ under t-exactness and amplitude ≤ 1; *no converse is claimed*.

Detailed expansions of these guard-rails and their implementable ranges appear in the appendices.

1.0. Windowed proof policy, Overlap Gate, Window Stack (WinFib), and B-Gate$^+$

We formalize the window policy, the *Overlap Gate* for gluing along overlaps (charts × windows), the *Window Stack* as a Grothendieck-fibrational bookkeeping device, and the standard single-layer acceptance gate on the B-side.

Definition 1.1 (Windows (MECE), right-open endpoints, safety margin, and δ-ledger)**.** A *domain window* is a right-open half-open interval $W = [u, u') \subset \mathbb{R}$. A windowing is *MECE* if a family $\{[u_k, u_{k+1})\}_k$ forms a disjoint cover $[u_0, U) = \bigsqcup_k [u_k, u_{k+1})$ with shared endpoints only. The *collapse window* is set by a threshold $\tau > 0$ (fixed within a gate, possibly varying across steps of a pipeline).

For an A-side pipeline with per-step non-commutation/error budgets $\delta_j = \delta_j^{\mathrm{alg}} + \delta_j^{\mathrm{disc}} + \delta_j^{\mathrm{meas}}$, write the *pipeline budget*

$$\Sigma\delta(i, \{\tau_j\}) \ := \ \sum_j \delta_j(i, \tau_j).$$

The *safety margin* gap_τ is determined by the window and collapse policy and must dominate error accumulation. Mandatory logs ensure MECE coverage and additivity $\sum_k \Sigma\delta_k = \Sigma\delta$.

Definition 1.2 (Window restriction (cropping) functor)**.** For a window $W = [u, u')$, define the exact endofunctor $\mathbf{W}_W : \mathsf{Pers}_k^{\mathrm{cons}} \to \mathsf{Pers}_k^{\mathrm{cons}}$ by restricting bars to W (bars fully outside W are deleted; bars intersecting W are cut at the endpoints). On barcodes this is literal cropping; on persistence modules it is precomposition with the inclusion $W \hookrightarrow \mathbb{R}$ and extension by zero outside W. $\mathbf{W}_W$ is 1-Lipschitz for d_{int}.

Definition 1.3 (Overlap Gate)**.** Let $\{(X_\alpha, W_\alpha)\}_{\alpha \in A}$ be a local cover, where X_α is a chart (domain piece) and $W_\alpha = [u_\alpha, u'_\alpha)$ a right-open window. For degree i and a fixed $\tau > 0$, write

$$\mathcal{B}_{\alpha,i} \ := \ \mathbf{T}_\tau \, \mathbf{W}_{W_\alpha}\big(\mathbf{P}_i(F|_{X_\alpha})\big) \ \in \ \mathsf{Pers}_k^{\mathrm{cons}}.$$

For any overlapping pair (α, β) with non-empty overlap

$$\Omega_{\alpha\beta} \ := \ X_\alpha \cap X_\beta \ \times \ \big(W_\alpha \cap W_\beta\big),$$

the *Overlap Gate* $\mathrm{OG}(\alpha, \beta; i, \tau)$ *passes* if

(i) $\mathbf{W}_{W_\alpha \cap W_\beta} \mathcal{B}_{\alpha,i}$ and $\mathbf{W}_{W_\alpha \cap W_\beta} \mathcal{B}_{\beta,i}$ are *isomorphic in* $\mathsf{Pers}_k^{\mathrm{cons}}$ up to budget:

$$d_{\mathrm{int}}\Big(\mathbf{W}_{W_\alpha \cap W_\beta} \mathcal{B}_{\alpha,i},\ \mathbf{W}_{W_\alpha \cap W_\beta} \mathcal{B}_{\beta,i}\Big) \ \le\ \Sigma\delta_{\alpha\beta}(i,\tau),$$

with $\Sigma\delta_{\alpha\beta}$ the overlap-specific budget recorded in the δ-ledger;

(ii) the safety margin on the overlap satisfies $\mathrm{gap}_\tau(\Omega_{\alpha\beta}) > \Sigma\delta_{\alpha\beta}(i,\tau)$;

(iii) tower diagnostics after $\mathbf{T}_\tau$ agree on the overlap (cf. Section 1): $\mu^i_{\mathrm{Collapse}} = u^i_{\mathrm{Collapse}} = 0$ for both sides restricted to $W_\alpha \cap W_\beta$.

A *global Overlap Gate pass* on $(\{X_\alpha\}, \{W_\alpha\}, i, \tau)$ means $\mathrm{OG}(\alpha, \beta; i, \tau)$ passes for all overlapping pairs.

Remark 1.4 (Window Stack (WinFib) as a Grothendieck fibration)**.** Let Win be the small category whose objects are pairs (α, W_α) with W_α right-open, and morphisms $(\alpha, W_\alpha) \to (\beta, W_\beta)$ given by inclusions $X_\alpha \cap W_\alpha \hookrightarrow X_\beta \cap W_\beta$. Define a pseudofunctor

$$\mathcal{S}_i(-,\tau):\ \mathsf{Win}^{\mathrm{op}} \longrightarrow \mathsf{Cat},\ (\alpha, W_\alpha) \longmapsto \Big\{\mathbf{T}_\tau\, \mathbf{W}_{W_\alpha}\big(\mathbf{P}_i(F|_{X_\alpha})\big)\Big\} \subset \mathsf{Pers}_k^{\mathrm{cons}},$$

with reindexing functors induced by cropping and restriction. The associated category of elements $\int \mathcal{S}_i(-,\tau)$ admits a canonical projection

$$\pi:\ \mathsf{WinFib}_{i,\tau}\ :=\ \int \mathcal{S}_i(-,\tau)\ \longrightarrow\ \mathsf{Win},$$

which is a Grothendieck fibration whose cartesian morphisms are induced by inclusions and cropping. The Overlap Gate certifies that objects in fibers glue along overlaps into a global B-side section. Stronger adjunction/commutation claims about π are **[Spec]** and developed in the appendices.

Definition 1.5 (B-Gate$^+$ (standard single-layer gate on the B-side))**.** Fix a right-open window $W = [u, u')$, a collapse threshold $\tau > 0$, and a degree i. We say *B-Gate*$^+$ *passes on* (W, τ, i) if, *after collapse on the B-side*,

(1) $\mathrm{PH}_1(C_\tau F) = 0$;

(2) $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$ (only within the scope stated in the Scope Box);

(3) the tower audit *after* $\mathbf{T}_\tau$ yields $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$ (tail isomorphism on comparison maps ϕ);

(4) the safety margin satisfies $\mathrm{gap}_\tau > \Sigma\delta(i, \{\tau_j\})$.

If a local cover $\{(X_\alpha, W_\alpha)\}$ is used, we additionally require a *global Overlap Gate pass* on $(\{X_\alpha\}, \{W_\alpha\}, i, \tau)$. Optionally (auxiliary), we may require the *auxiliary spectral bars* on the chosen spectral window to vanish after collapse; final gate decisions never rely solely on spectral indicators.

Remark 1.6 (Asymmetric mirror with δ-ledger; soft-commuting policy)**.** Let $U_m, \dots, U_1$ be A-side steps (each deletion-type or ε-continuation), and C_{τ_j} the stepwise collapses. For every i and any fixed B-side collapse threshold τ,

$$d_{\mathrm{int}}\Big(\mathbf{T}_\tau \mathbf{P}_i\big(\mathrm{Mirror}(C_{\tau_m} U_m \cdots C_{\tau_1} U_1 F)\big),\ \mathbf{T}_\tau \mathbf{P}_i\big(C_{\tau_m} U_m \cdots C_{\tau_1} U_1\, \mathrm{Mirror}\, F\big)\Big)\ \le\ \sum_{j=1}^{m} \delta_j(i, \tau_j),$$

uniformly in F. Subsequent 1-Lipschitz post-processing on the persistence layer does not increase the right-hand side. If collapsers T_A, T_B are not provably commuting, we adopt the *soft-commuting* policy: measure

$$\Delta_{\mathrm{comm}} := d_{\mathrm{int}}(T_A T_B M, T_B T_A M).$$

If $\Delta_{\mathrm{comm}} \le \eta$, we accept parallel collapse; else we fix an order and record Δ_{comm} in δ^{alg}.

1.1. Terminology and notation

Fix a base field k. Let Vect_k be finite-dimensional k-vector spaces and $D^{\mathrm{b}}(k\text{-mod})$ the bounded derived category of finite-dimensional k-modules.

Filtered chain complexes. Write $\mathsf{FiltCh}()$ for finite-type (filtered-constructible) filtered chain complexes $F = (C_\bullet, \{F^t C_\bullet\}_{t\in\mathbb{R}})$ over k, with $F^t C_\bullet \subseteq F^{t'} C_\bullet$ for $t \le t'$ and filtered chain maps as morphisms. For each $i \in \mathbb{Z}$ (bounded in homological degree), degreewise homology yields a persistence module with structure maps induced by the inclusions $F^t C_\bullet \hookrightarrow F^{t'} C_\bullet$:

$$\mathbf{P}_i(F)\colon \mathbb{R} \to \mathsf{Vect}_k, \qquad t \longmapsto \mathrm{H}_i(F^t C_\bullet).$$

Constructibility on bounded windows places $\mathbf{P}_i(F)$ in $\mathsf{Pers}_k^{\mathrm{cons}}$. Its barcode is denoted by $\mathrm{PH}_i(F)$. Writing $\mathrm{PH}_1(F) = 0$ means $\mathbf{P}_1(F)$ is the zero object in $\mathsf{Pers}_k^{\mathrm{cons}}$.

Standing convention (constructible range). All persistence arguments take place in $\mathsf{Pers}_k^{\mathrm{cons}} \subset \mathsf{Pers}$ (finite critical set on bounded intervals). Abelianity, Serre subcategories, and exact localizations are used only within $\mathsf{Pers}_k^{\mathrm{cons}}$. When filtered (co)limits are used, they are computed objectwise in $[\mathbb{R}, \mathsf{Vect}_k]$ and then verified to return to $\mathsf{Pers}_k^{\mathrm{cons}}$; we assert equalities *at the persistence layer*. On filtered complexes we retain finite-type hypotheses and note filtered colimits may exit finiteness. We identify $\mathsf{Pers}_k^{\mathrm{ft}}$ with $\mathsf{Pers}_k^{\mathrm{cons}}$ and henceforth write only $\mathsf{Pers}_k^{\mathrm{cons}}$.

Realization functor and test family. A realization is a t-exact functor $\mathcal{R} : \mathsf{FiltCh}(\to) D^{\mathrm{b}}(k\text{-mod})$ of amplitude ≤ 1. We fix the minimal test family $Q = \{k[0]\}$ for Ext^1-vanishing tests.

Interleaving distance. Let d_{int} denote the interleaving (equivalently, bottleneck) distance on $\mathsf{Pers}_k^{\mathrm{cons}}$. All Lipschitz claims use d_{int}.

Epistemic stance. We treat *collapse* as a functorial simplification exposed via higher-dimensional projections and controlled obstruction removal. Operationally, we work with observable indicators (PH_1, Ext^1, $\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}$), an admissible entry gate (the *Collapse Zone*), and a failure typology including *invisible* limit obstructions. We enforce a windowed proof policy and a single-layer judgement after collapse (B-side), with non-commutation externalized via a δ-ledger. Windows are uniformly **right-open**.

1.2. Collapse (operational definition, v15.1)

We define a metrically stable truncation on persistence and a filtered lift used for constructions.

Definition 1.7 (Exact truncation on persistence and abbreviation)**.** For $\tau > 0$, let

$$\mathbf{T}_\tau :\ \mathsf{Pers}_k^{\mathrm{cons}} \longrightarrow (\mathsf{E}_\tau)^\perp$$

be the exact reflective localization (Serre quotient) at the Serre subcategory generated by interval modules of length $\le \tau$, i.e. $\mathbf{T}_\tau$ deletes precisely those finite bars. Here $(\mathsf{E}_\tau)^\perp$ is the full τ-local (orthogonal) subcategory (orthogonal to E_τ for both Hom and Ext^1). Existence and exactness hold in the constructible 1D range (Crawley–Boevey 2015; Chazal–de Silva–Glisse–Oudot 2016). $\mathbf{T}_\tau$ is 1-Lipschitz for d_{int}.

Definition 1.8 (Thresholded collapse lift on filtered complexes)**.** A *thresholded collapse lift* is a functor

$$C_\tau : \ \mathsf{FiltCh}(\longrightarrow)\mathsf{FiltCh}($$

) defined *up to f.q.i.* such that, for every i,

$$\mathbf{P}_i\big(C_\tau(F)\big) \ \cong \ \mathbf{T}_\tau\big(\mathbf{P}_i(F)\big) \quad \text{in } \mathsf{Pers}_k^{\text{cons}}.$$

All (co)limit and pullback compatibilities used in this work are applied *at the persistence layer* via these identifications.

Remark 1.9 (Endpoints and infinite bars)**.** Endpoint conventions are *uniformly right-open* for windows $[u, u')$. Infinite bars and endpoint policies are centralized and referenced consistently; all later references defer to this remark.

Remark 1.10 (Lipschitz stability and scope)**.** $\mathbf{T}_\tau$ is 1-Lipschitz for d_{int}, hence C_τ is 1-Lipschitz *at the persistence level* via $\mathbf{P}_i \circ C_\tau \simeq \mathbf{T}_\tau \circ \mathbf{P}_i$. Compatibilities with filtered colimits and finite pullbacks are used only through $\mathbf{T}_\tau$. Spectral indicators (tails/heat traces) are *not* f.q.i.-invariants; they are controlled by stability under a fixed policy $(\beta, M(\tau), s)$ on an appropriate linearization $L(C_\tau F)$. Spectral quantities are strictly auxiliary.

Definition 1.11 (PH–collapse and Ext–collapse)**.** For $F \in \mathsf{FiltCh}()$:

- *PH–collapse in degree* 1: $\mathrm{PH}_1(F) = 0$.
- *Ext–collapse (tested against* $Q = \{k[0]\}$*)*: $\mathrm{Ext}^1\big(\mathcal{R}(F), k\big) = 0$.

Definition 1.12 (Collapse Zone (binary entry gate))**.** Set

$$\mathrm{CollapseAdmissible}(F) \ :\Longleftrightarrow \quad \mathrm{PH}_1(F) = 0 \ \wedge \ \mathrm{Ext}^1\big(\mathcal{R}(F), k\big) = 0.$$

The *Collapse Zone* is the full subcategory

$$\mathfrak{C} \ := \ \{\, F \in \mathsf{FiltCh}(|)\mathrm{CollapseAdmissible}(F)\,\} \subset \mathsf{FiltCh}(.)$$

Under the hypotheses of the Scope Box (bridge used only in $D^{\mathrm{b}}(k\text{-mod})$ under t-exactness and amplitude ≤ 1), PH–collapse implies Ext–collapse; *no converse is asserted.*

Remark 1.13 (Invariance)**.** $\mathrm{CollapseAdmissible}$ is invariant under filtered quasi-isomorphisms and under equivalences preserving $\mathbf{P}_i(-)$ and $\mathcal{R}(-)$ up to isomorphism in $D^{\mathrm{b}}(k\text{-mod})$.

Abstract functorial viewpoint. Let $\mathsf{Triv} \subset \mathsf{FiltCh}()$ be the full subcategory of objects with vanishing PH_1 and tested Ext^1. If the inclusion $\iota : \mathsf{Triv} \hookrightarrow \mathsf{FiltCh}()$ admits a right adjoint C (existence of a genuine right adjoint is **[Spec]** beyond the implementable range), we call such C a *collapser*; the unit $\eta_F : F \to \iota C(F)$ serves as a canonical collapse morphism. In practice, C_τ (Theorem 1.8) is the operative, metrically stable proxy.

1.3. The entry gate as a measurable predicate

The predicate $\mathrm{CollapseAdmissible}$ separates, by observable tests, objects whose first persistent topology and tested categorical extensions vanish:

- *Measurability:* PH_1 =0 is read off barcodes; Ext^1-vanishing is checked against $\{k[0]\}$ in $D^{\mathrm{b}}(k\text{-mod})$.
- *Functorial simplification:* C_τ erases bars of length $\leq \tau$ and, under the Scope Box hypotheses, preserves admissibility.
- *Cross-domain interface:* the gate supports specifications (geometry, arithmetic, mirror/tropical, PDE) *without* positing cross-domain equivalences.

1.4. Failure landscape and the invisible obstruction $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$

We classify failures of collapse and introduce tower-sensitive diagnostics.

Generic fiber dimension (intuition). For $M \in \mathsf{Pers}_k^{\text{cons}}$, the *generic fiber dimension* is the stabilized rank

$$\text{gdim}(M) \;:=\; \lim_{t \to +\infty} \dim_k M(t),$$

which exists in the constructible setting. After truncation by $\mathbf{T}_\tau$, gdim equals the multiplicity of the infinite bar $I[0, \infty)$ in the barcode.

Observable failure types.

- **Type I (Topological):** $\text{PH}_1(F) eq 0$.
- **Type II (Categorical):** $\text{Ext}^1\big(\mathcal{R}(F), k\big) eq 0$.
- **Type III ([Spec]-level instability):** admissibility is unstable under a prescribed operation (e.g. a specific pullback/filtered colimit) in a way detectable at finite level.

Invisible failure along towers (Type IV). Let $\{F_n\}_{n \in \mathbb{N}}$ be a directed system in $\mathsf{FiltCh}()$ with limit F_∞. For fixed $\tau > 0$ and each degree i, the natural comparison *after truncation by* $\mathbf{T}_\tau$

$$\phi_i : \varinjlim_n \mathbf{T}_\tau\big(\mathbf{P}_i(F_n)\big) \;\longrightarrow\; \mathbf{T}_\tau\big(\mathbf{P}_i(F_\infty)\big)$$

induces the tower-sensitivity invariants

$$\mu^i_{\text{Collapse}} := \text{gdim}\ker(\phi_i), u^i_{\text{Collapse}} := \text{gdim}\,\text{coker}(\phi_i),$$

interpreting gdim as the multiplicity of $I[0, \infty)$ after truncation. Set

$$\mu_{\text{Collapse}} := \sum_i \mu^i_{\text{Collapse}}, u_{\text{Collapse}} := \sum_i u^i_{\text{Collapse}}.$$

We say F_∞ exhibits *invisible failure (Type IV)* if $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) eq (0, 0)$. In that case, all finite layers may appear admissible while the limit fails the gate. Finiteness follows from constructibility and bounded homological degrees.

Purpose and non-identity. $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$ diagnose limit effects of collapse *at the persistence layer*; they are *not* classical Iwasawa invariants. A canonical refinement-limit example exhibits pure cokernel-type failure at fixed τ. They are invariant under f.q.i. and cofinal reindexing.

1.5. Admissible A-side operations and single-layer judgement

We record the permitted A-side steps and the single-layer policy.

Definition 1.14 (Admissible A-side operations)**.** Each A-side step U on filtered complexes is labeled as:

- *Deletion-type (monotone):* induces non-increase of windowed persistence energies and auxiliary spectral counters after applying C_τ. Examples: stop addition/sector shrinking, mollification (low-pass), viscosity increase, threshold lowering.

- *ε-continuation (non-expansive):* is 1-Lipschitz for d_{int}; any drift is bounded by the declared ε.
- *Inclusion-type (stable only):* may increase indicators; only stability (non-expansiveness) is claimed; monotonicity is not asserted.

Every step is followed by a collapse C_{τ_j} and all measurements are taken *only* on the B-side single layer $\mathbf{T}_\tau \mathbf{P}_i$.

Declaration 1.15 (Windowed certificates, Overlap Gate, and pasting)**.** If B-Gate$^+$ (Theorem 1.5) passes on (W, τ, i), we issue a *windowed regularity certificate* for (W, τ): "no failure with no false negative on W." When a local cover $\{(X_\alpha, W_\alpha)\}$ is used, we additionally require a *global Overlap Gate pass* on $(\{X_\alpha\}, \{W_\alpha\}, i, \tau)$. Global statements are obtained by pasting windowed certificates along a MECE partition using the Restart and Summability policies. All cross-domain comparisons are performed after collapse, and any non-commutation is accounted for in the δ-ledger.

1.6. Strengthening notes (fully integrated)

The following reinforcements are fully integrated above and govern all subsequent chapters:

- **Unified endpoint policy.** All windows are *right-open* $[u, u')$. Any earlier "right-inclusive" phrasing is obsolete. Cropping (Theorem 1.2) and Overlap Gate (Theorem 1.3) are formulated for right-open windows uniformly.
- **Overlap-first gluing.** The Overlap Gate is elevated as the primary local-to-global device: global certificates are issued only after all pairwise overlaps pass, with budgets verified against gap_τ on overlaps.
- **Window Stack (WinFib).** Windows and parameters are internalized as a Grothendieck fibration (Theorem 1.4); cartesian morphisms encode type-safe consistency checks on overlaps and inclusions.
- **Guard-rails in a box.** The Scope Box at chapter head shortens (and sharpens) the global guard-rails: constructible 1D, field coefficients, after-collapse layer of truth, δ-ledger, windowed (MECE) proof with restart and summability, and the one-way $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ bridge only in $D^{\mathrm{b}}(k\text{-mod})$.
- **Audit-first policy.** Tower diagnostics ($\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}$) are evaluated *after* $\mathbf{T}_\tau$; invisible failures (Type IV) are treated on equal footing with Types I–III. Certificates require tail isomorphisms (vanishing μ, u).
- **Spectral indicators as auxiliary.** Spectral tails/heat traces are never sole gate criteria; they obey monotonicity for deletion-type steps and stability under ε-continuations, and are always measured after collapse.

References and provenance. Existence, exactness, and interval decomposition in the constructible 1D range follow Crawley–Boevey (2015) and Chazal–de Silva–Glisse–Oudot (2016). Derived and sheaf-theoretic realizations beyond $D^{\mathrm{b}}(k\text{-mod})$ are **[Spec]** and rely on projection formula/base change; they do not enlarge the proven bridge. Strong adjunction/commutation claims for WinFib are also **[Spec]**.

latex

2 Chapter 2: Concrete Model — Finite-Type Filtered Chain Complexes and Thresholded Collapse

2.1. The category $\mathsf{FiltCh}()$ and persistence modules

Fix a field k. Let $\mathsf{Ch}^b(k)$ be the category of bounded chain complexes of finite-dimensional k-vector spaces. A *finite-type filtered chain complex* is a pair

$$F = (C_\bullet, \{F^t C_\bullet\}_{t \in \mathbb{R}})$$

where $F^t C_\bullet \subseteq F^{t'} C_\bullet$ for $t \le t'$, the filtration is exhaustive and left-bounded, and, for each i, the persistence module

$$\mathrm{H}_i(F) : \mathbb{R} \longrightarrow \mathsf{Vect}_k, \qquad t \longmapsto \mathrm{H}_i(F^t C_\bullet)$$

is pointwise finite-dimensional with finitely many critical parameters on compact intervals. Denote by $\mathsf{FiltCh}()$ the category of such F with filtration-preserving chain maps. For each i, let

$$\mathbf{P}_i : \mathsf{FiltCh}(\longrightarrow)\mathsf{Pers}_k^{\mathrm{cons}}$$

be the functor sending $F \mapsto \mathrm{H}_i(F)$. We write $\mathrm{PH}_i(F)$ for the barcode (multiset of intervals) of $\mathbf{P}_i(F)$. Throughout, the interleaving (equivalently, bottleneck) distance on persistence modules is denoted d_{int}.

Standing convention (constructible range and notation). We work inside the constructible subcategory $\mathsf{Pers}_k^{\mathrm{cons}} \subset \mathsf{Pers}$ (finite critical set on bounded intervals, equivalently p.f.d. with finitely many changes on compacts). We identify $\mathsf{Pers}_k^{\mathrm{ft}}$ with $\mathsf{Pers}_k^{\mathrm{cons}}$ by convention. $\mathsf{Pers}_k^{\mathrm{cons}}$ is abelian, admits interval decompositions, and carries a well-defined length. All uses of abelianity, Serre subcategories, and exact localizations are made within $\mathsf{Pers}_k^{\mathrm{cons}}$; see Appendix A for details. For filtered complexes we keep the finite-type hypothesis and record that filtered colimits may exit finiteness (cf. Appendix A). Filtered (co)limits, when used, are computed objectwise in $[\mathbb{R}, \mathsf{Vect}_k]$ and then verified to return to $\mathsf{Pers}_k^{\mathrm{cons}}$; no claim is made outside this regime.

Remark 2.1 (Constructible abelian setting)**.** Within $\mathsf{Pers}_k^{\mathrm{cons}}$, kernels and cokernels are computed pointwise and preserve finiteness; thus $\mathsf{Pers}_k^{\mathrm{cons}}$ is abelian with interval decompositions and Serre localizations by bar length. For each $\tau > 0$, the full subcategory E_τ generated by interval modules of length $\le \tau$ is a hereditary Serre (localizing) subcategory in this constructible 1D setting; see Crawley–Boevey (2015); Chazal–de Silva–Glisse–Oudot (2016).

Remark 2.2 (Windowed proof policy (MECE); right-open endpoints)**.** All statements here are meant to be applied *per window* (cf. Chapter 1). A *domain window* is a right-open half-open interval $[u_k, u_{k+1})$. A windowing is *MECE* if $\bigsqcup_k [u_k, u_{k+1}) = [u_0, U)$, intervals are ordered and adjacent windows meet only in their endpoints. Two mandatory coverage checks are recorded in each run (Appendix G): (i) $\sum_k (u_{k+1} - u_k) = U - u_0$ (total length equals the sum of window lengths), (ii) the number of events (births/deaths counted with multiplicity) on $[u_0, U)$ equals the sum of per-window counts up to rounding tolerance. Collapse thresholds and spectral bins are *fixed on each window* and used only *after collapse* at the persistence layer.

2.2. Thresholded collapse: Serre localization on persistence and filtered lift

We first recall truncation on persistence modules and then lift it to filtered complexes.

Ephemeral part and localization (constructible 1D case). Let $\mathsf{E}_\tau \subset \mathsf{Pers}_k^{\mathrm{cons}}$ be the Serre subcategory generated by interval modules of length $\le \tau$. The reflective localization

$$\mathbf{T}_\tau : \mathsf{Pers}_k^{\mathrm{cons}} \longrightarrow (\mathsf{E}_\tau)^\perp$$

is *exact in the constructible* 1*D range* and 1-Lipschitz for d_{int}. Here $(\mathsf{E}_\tau)^\perp$ denotes the τ-local (orthogonal) subcategory (orthogonal to E_τ for both Hom and Ext^1); in general, "τ-local" is stricter than "τ-torsion-free". We interpret "deleting bars of length $\le \tau$" via this exact localization: for $M \in \mathsf{Pers}_k^{\mathrm{cons}}$,

$$\mathbf{T}_\tau(M) = M/E_\tau(M), \qquad E_\tau(M) \text{ the maximal } \tau\text{-ephemeral subobject.}$$

Remark 2.3 (Endpoints and infinite bars)**.** $\mathbf{T}_\tau$ deletes precisely the *finite* bars of length $\le \tau$; infinite bars are invariant. The choice of open/closed endpoint conventions does not affect $\mathbf{T}_\tau$, since interleaving-equivalence classes (hence bar lengths) are preserved.

Serre reflector: proof sketch (exactness). Exactness of $\mathbf{T}_\tau$ follows from exactness of the Serre quotient $\mathsf{Pers}_k^{\mathrm{cons}} \to \mathsf{Pers}_k^{\mathrm{cons}}/\mathsf{E}_\tau$ and identification of the quotient with $(\mathsf{E}_\tau)^\perp$: localization at a hereditary Serre subcategory is exact; in the constructible 1D barcode setting the quotient is abelian and the orthogonal is reflective. As a left adjoint, $\mathbf{T}_\tau$ preserves colimits; being exact it preserves finite limits/colimits. Shift-commutation yields 1-Lipschitzness for d_{int}. Full details are deferred to Appendix A.

Lifting to filtered complexes. Choose a functor

$$\mathcal{U} : \mathsf{Pers}_k^{\mathrm{cons}} \longrightarrow \mathsf{FiltCh}($$

) such that $\mathbf{P}_i(\mathcal{U}(M)) \cong M$ for M concentrated in degree i, and $\mathcal{U}$ sends finite direct sums of degree-concentrated modules to finite direct sums of elementary filtered complexes built by gluing interval complexes. For $F \in \mathsf{FiltCh}()$ define $C_\tau(F)$ to be any object of $\mathsf{FiltCh}()$ equipped with filtered quasi-isomorphisms in each degree i

$$\mathbf{P}_i\big(C_\tau(F)\big) \xrightarrow{\cong} \mathbf{T}_\tau\big(\mathbf{P}_i(F)\big).$$

Such choices exist and are unique up to filtered quasi-isomorphism; any functorial choice (up to f.q.i.) is a *thresholded collapse functor*. A genuine right adjoint/comonadic presentation at the filtered level is **[Spec]** beyond the implementable range; all exactness claims are asserted at the persistence layer.

2.3. Stability and basic calculus of C_τ

We collect the properties needed later; equalities are in $\mathsf{Pers}_k^{\mathrm{cons}}$ and filtered-level claims are up to f.q.i.

Lemma 2.4 (Shift–commutation and 1-Lipschitz)**.** Let S^ε denote the ε-shift on persistence modules. Then $\mathbf{T}_\tau \circ S^\varepsilon \cong S^\varepsilon \circ \mathbf{T}_\tau$. Consequently $\mathbf{T}_\tau$ preserves ε-interleavings and is 1-Lipschitz for d_{int}.

Proposition 2.5 (Stability, calculus, and (co)limit compatibility)**.** Let $\tau, \sigma > 0$ and $F, G \in \mathsf{FiltCh}()$.

1. *(Non-expansiveness)* For every i, $d_{\mathrm{int}}(\mathbf{P}_i(C_\tau F), \mathbf{P}_i(C_\tau G)) \le d_{\mathrm{int}}(\mathbf{P}_i(F), \mathbf{P}_i(G))$.
2. *(Monotonicity and idempotence)* If $\tau \le \sigma$, then for every i there is a natural epimorphism $\mathbf{P}_i(C_\tau F) \twoheadrightarrow \mathbf{P}_i(\mathsf{C}_\sigma F)$, and $C_\tau \circ \mathsf{C}_\sigma \simeq C_{\max\{\tau,\sigma\}} \simeq \mathsf{C}_\sigma \circ C_\tau$ up to f.q.i. (equivalently, $\mathbf{T}_\tau \circ \mathbf{T}_\sigma = \mathbf{T}_{\max\{\tau,\sigma\}}$ at the persistence layer).
3. *(Exactness on persistence)* Viewing $\mathbf{T}_\tau$ as the localization functor $\mathsf{Pers}_k^{\mathrm{cons}} \to \mathsf{Pers}_k^{\mathrm{cons}}/\mathsf{E}_\tau$ composed with the identification of the quotient with $(\mathsf{E}_\tau)^\perp$, $\mathbf{T}_\tau$ is exact and preserves finite limits/colimits. Consequently, after applying $\mathbf{P}_*$ degreewise to a short exact sequence in $\mathsf{FiltCh}()$, exactness is preserved under $\mathbf{T}_\tau$ at the persistence layer. No claim of short exactness for $\mathbf{P}_i \circ C_\tau$ is made in general.
4. *(Filtered colimits; **[Spec]**)* If $\{F_\lambda\}_{\lambda\in\Lambda}$ is a filtered diagram in $\mathsf{FiltCh}()$ whose colimit is computed degree-wise on chains/filtrations, then for every i, $\mathbf{P}_i(C_\tau(\varinjlim_\Lambda F_\lambda)) \cong \varinjlim_\Lambda \mathbf{P}_i(C_\tau(F_\lambda))$.
5. *(Finite pullbacks; **[Spec]**)* If finite pullbacks in $\mathsf{FiltCh}()$ are computed degreewise and $\mathcal{U}$ preserves finite limits up to f.q.i., then for every pullback square $F \times_H G$, $\mathbf{P}_i(C_\tau(F \times_H G)) \cong \mathbf{P}_i(C_\tau(F) \times_{C_\tau(H)} C_\tau(G))$.

Remark 2.6 (Scope of limit/pullback compatibilities)**.** Compatibilities of C_τ with finite pullbacks and filtered colimits are used *at the persistence layer* via $\mathbf{P}_i \circ C_\tau \simeq \mathbf{T}_\tau \circ \mathbf{P}_i$. **Statements at the filtered-complex level are adopted as [Spec] and hold up to f.q.i.**

Remark 2.7 (Monotonicity vs. stability; deletion- vs. inclusion-type)**.** For *deletion-type* updates (Dirichlet restriction/absorbing boundaries, principal submatrices/Schur complements, positive-semidefinite Loewner contractions, conservative averaging), windowed persistence energies and spectral tails/heat traces computed on $L(C_\tau F)$ are *non-increasing after truncation*. For *inclusion-type* updates we claim only *stability* (non-expansiveness). Spectral indicators are not f.q.i.-invariants; they are controlled by a fixed policy $(\beta, M(\tau), s)$ (see Chapter 11 and Appendix E).

Declaration 2.8 (Spec–C_τ calculus)**.** The filtered lift C_τ realizes $\mathbf{T}_\tau$ degreewise *up to filtered quasi-isomorphism*. Compatibility with filtered colimits and finite pullbacks is used *at the persistence layer* via $\mathbf{P}_i \circ C_\tau \simeq \mathbf{T}_\tau \circ \mathbf{P}_i$. **No stronger claim is made at the filtered-complex level.**

2.4. Windowing (MECE), τ-adaptation, and spectral bins

We formalize the operating policies required for windowed certificates and for the interaction with spectral auxiliaries.

Definition 2.9 (Domain windows (MECE) and coverage checks)**.** A *domain windowing* is a finite or countable family of right-open intervals $\{[u_k, u_{k+1})\}_k$ such that $\bigsqcup_k [u_k, u_{k+1}) = [u_0, U)$ (disjoint union) and $u_k < u_{k+1}$. Coverage checks:

$$\sum_k (u_{k+1} - u_k) = U - u_0, \qquad \#\mathrm{Events}([u_0, U)) = \sum_k \#\mathrm{Events}([u_k, u_{k+1})) \ (\pm \text{ rounding}).$$

Here #Events counts births/deaths (with multiplicity) observed by $\mathbf{P}_i$ on the specified window. Verdicts are issued per window; global statements are obtained by pasting certificates as in Chapter 4.

Definition 2.10 (Collapse thresholds and τ-adaptation)**.** Fix a collapse threshold $\tau > 0$ on each window. A default adaptation ties τ to numerical/filtration resolution:

$$\tau \;=\; \alpha \cdot \max\{\Delta t, \Delta x\} \qquad (\alpha > 0 \text{ fixed per run}),$$

or any comparable rule documented in the run manifest (Appendix G). A *τ-sweep* is a discrete set $\{\tau_\ell\}$ on which $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ and B-Gate$^+$ are evaluated. A *stable band* is a contiguous range of τ-values on which $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$ holds (Appendix J); certificates are issued within stable bands.

Definition 2.11 (Spectral bins and auxiliary bars (aux-bars))**.** For a spectral operator with (ascending) spectrum $(\lambda_m)_{m\geq 1}$, fix a bin width $\beta > 0$ and a range $[a, b]$. Define right-open bins $I_r = [a + r\beta, a + (r+1)\beta)$ for $r = 0, 1, \ldots, \lfloor (b-a)/\beta \rfloor - 1$, and count occupancies $E_r := \#\{m \mid \lambda_m \in I_r\}$. Underflow $\{\lambda_m < a\}$ and overflow $\{\lambda_m \geq b\}$ are recorded. Along a discrete index (e.g. step number), sequences $(E_r(j))_j$ define *auxiliary spectral bars (aux-bars)* as maximal consecutive runs on which $E_r(j) > 0$; their *lifetimes* are measured in the discrete index (rescaled when needed). After applying C_τ on the persistence side, indicators from aux-bars are evaluated as *auxiliary gauges*: *monotone* under deletion-type steps, *stable* under ε-continuations (Appendix E). They never replace the B-side gate.

Remark 2.12 (Separation of roles: persistence vs. spectral auxiliaries)**.** Gate decisions are made on the persistence layer after collapse, never on spectral auxiliaries alone. Aux-bars are used as supportive evidence (e.g. "aux-bars= 0 after C_τ on the spectral window''), with their binning policy $(\beta, [a, b])$ recorded. This maintains invariance guarantees of $\mathbf{T}_\tau$ and the reproducibility of windowed certificates.

2.5. Collapse admissibility and robust variants

Let $\mathcal{R} : \mathsf{FiltCh}(\to) D^{\mathrm{b}}(k\text{-mod})$ be a fixed t-exact realization of amplitude ≤ 1, and fix the minimal test family $Q = \{k[0]\}$.

Definition 2.13 (Admissibility predicate)**.** For $F \in \mathsf{FiltCh}()$, set

$$\mathrm{CollapseAdmissible}(F) \ :\Longleftrightarrow\ \mathrm{PH}_1(F) = 0 \ \wedge\ \mathrm{Ext}^1\big(\mathcal{R}(F), k\big) = 0.$$

Under the bridging assumptions (Chapter 3), $\mathrm{PH}_1(F) = 0 \Rightarrow \mathrm{Ext}^1(\mathcal{R}(F), k) = 0$ in $D^{\mathrm{b}}(k\text{-mod})$.

Definition 2.14 (Robust admissibility at scale ε)**.** Fix $\varepsilon > 0$. We say F is *ε-robustly collapse-admissible* if

$$\mathrm{PH}_1\big(C_\varepsilon(F)\big) = 0 \quad \text{and} \quad \mathrm{Ext}^1\big(\mathcal{R}(C_\varepsilon(F)), k\big) = 0.$$

This weakens the binary gate by discarding bars shorter than ε; ε reflects the admissible noise level.

Remark 2.15 (Usage and guarantees)**.** By Theorem 2.5, C_τ provides a metrically stable simplification compatible with filtered colimits and finite pullbacks (up to f.q.i.), furnishing a robust pre-processing step toward admissibility checks. In practice, one fixes an application-level noise scale ε and tests admissibility on $C_\varepsilon(F)$; the one-way bridge then applies in $D^{\mathrm{b}}(k\text{-mod})$. No global equivalence $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is asserted; only $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ under explicit bridging hypotheses.

Remark 2.16 (Formalizability)**.** The ingredients of this chapter are formalizable: (i) E_τ is a localizing Serre subcategory in the constructible range; (ii) $\mathbf{T}_\tau$ is exact and preserves finite limits/colimits; (iii) Theorem 2.4 implies 1-Lipschitz continuity in d_{int}; and (iv) Theorem 2.5 encodes the (co)limit and pullback compatibility. These can be encoded in Coq/Lean as axiomatized interfaces for the persistence layer, with filtered lifts handled up to f.q.i. (Appendices A–B).

Lemma 2.17 (ε-survival under interleaving perturbations)**.** If $d_{\mathrm{int}}(\mathbf{P}_i(F), \mathbf{P}_i(G)) \leq \varepsilon$, then for every bar b of $\mathbf{P}_i(F)$ and every $\tau_0 > 0$, letting $\ell_{\tau_0}(b)$ denote the $[0, \tau_0]$-clipped length of b, the following holds: if $\ell_{\tau_0}(b) > 2\varepsilon$, then b has a corresponding bar in $\mathbf{T}_{\tau_0}(\mathbf{P}_i(G))$ whose $[0, \tau_0]$-clipped length is at least $\ell_{\tau_0}(b) - 2\varepsilon$. (See Appendix I for the standard matching argument.)

2.6. Local Equivalence on saturation windows (PH–Ext equivalence *per window*)

We elevate the window-local PH–Ext equivalence to a theorem under explicitly verifiable hypotheses; this replaces any global equivalence claims.

Definition 2.18 (Saturation window, length gap, and near-τ non-accumulation)**.** Fix degree $i = 1$, a right-open window $W = [u, u')$, and $\tau > 0$. We say W is a *saturation window at τ* for F if:

- *Event stability:* births/deaths in $\mathbf{P}_1(F)$ restricted to W stabilize (no new events beyond a finite stage for towers, or no significant drift for single objects).
- *Length gap:* the maximal finite bar length inside W is $\leq \tau - \eta$ for some $\eta > 0$.
- *Near-τ non-accumulation:* no sequence of bar lengths in W strictly increases to τ.

We additionally require the post-collapse tail isomorphism at τ: $\mu_{\mathrm{Collapse}} = u_{\mathrm{Collapse}} = 0$ on W for degree 1 (cf. Theorem 2.5 and Chapter 1).

Theorem 2.19 (Local Equivalence on saturation windows)**.** Let $F \in \mathsf{FiltCh}()$. Assume:

1. t-exact realization $\mathcal{R} : \mathsf{FiltCh}(\to) D^{\mathrm{b}}(k\text{-mod})$ of amplitude ≤ 1 (two-term model for $\mathcal{R}(C_\tau F)$);

2. $W = [u, u')$ is a saturation window at $\tau > 0$ in the sense of Theorem 2.18;

3. The tail comparison map $\phi_{1,\tau}$ on W is an isomorphism (i.e. $\mu_{\mathrm{Collapse}} = u_{\mathrm{Collapse}} = 0$ after $\mathbf{T}_\tau$).

Then, *on the window W and at threshold τ*, one has the equivalence

$$\mathrm{PH}_1(C_\tau F) = 0 \quad \Longleftrightarrow \quad \mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0.$$

Proof (sketch). Under (1), $\mathcal{R}(C_\tau F) \in D^{[-1,0]}(k\text{-mod})$ and the edge identification (Appendix C) yields $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) \cong \mathrm{Hom}(H^{-1}(\mathcal{R}(C_\tau F)), k)$. Under the saturation hypotheses (2) and the tail isomorphism (3), the degree-1 persistence on W is determined by stabilized bars, and near-τ accumulation is excluded; after truncation at τ the barcode is insensitive to residual finite fluctuations. Hence, $\mathrm{PH}_1(C_\tau F) = 0$ on W if and only if the stabilized edge group $H^{-1}(\mathcal{R}(C_\tau F))$ vanishes, i.e. $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$. □

Remark 2.20 (Scope and use)**.** Theorem 2.19 is a *window-local* equivalence and replaces any global PH–Ext equivalence claims. It is designed to be verifiable as a gate condition (cf. Chapter 1): event stability, length gap, near-τ non-accumulation, and tail isomorphism ($\mu = u = 0$) are inspected on W. Outside saturation windows, only $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ is asserted (in $D^{\mathrm{b}}(k\text{-mod})$).

2.7. Length Spectrum Operator Λ_{len} (windowed) and its invariance

We formalize the "length spectrum" associated to a barcode and use it as a robust, windowed proxy. This complements but never replaces B-side gates.

Definition 2.21 (Windowed length spectrum operator Λ_{len})**.** Let $M \in \mathsf{Pers}_k^{\mathrm{cons}}$ admit a barcode decomposition $M \cong \bigoplus_j I[b_j, d_j)$. For a right-open window $W = [u, u')$, define the *clipped length* $\ell_W(I[b_j, d_j)) := \max\{0, \min\{d_j, u'\} - \max\{b_j, u\}\}$. The *windowed length spectrum operator* $\Lambda_{\mathrm{len}}(M; W)$ is the diagonal (commutative) endomorphism on $\bigoplus_j k \cdot e_j$ with eigenvalues $\{\ell_W(I[b_j, d_j))\}_j$.

Proposition 2.22 (Isomorphism invariance and identification)**.** *The unordered multiset of eigenvalues of $\Lambda_{\mathrm{len}}(M; W)$ equals the multiset of clipped bar lengths $\{\ell_W(I[b_j, d_j))\}_j$. This multiset is invariant under isomorphisms $M \simeq M'$ in $\mathsf{Pers}_k^{\mathrm{cons}}$. In particular, $\Lambda_{\mathrm{len}}(-; W)$ is well-defined up to isomorphism of bar decompositions.*

Remark 2.23 (Functoriality caveat and intended use)**.** $\Lambda_{\mathrm{len}}(-; W)$ is "natural up to isomorphism" (barcode uniqueness) and not fully functorial under arbitrary morphisms; it is designed for *windowed auditing* and stability checks (e.g. non-increase under deletion-type updates after collapse). For links with classical spectral quantities (heat traces/tails) and sandwich bounds, see Chapter 11 and Appendix E. Λ_{len} is never a sole gate criterion.

2.8. Operating summary for Chapter 2

On each right-open window $[u_k, u_{k+1})$:

- Fix a collapse threshold τ (adapted to resolution; Theorem 2.10) and, if spectral auxiliaries are used, fix a bin policy $(\beta, [a, b])$ (Theorem 2.11).

- Apply A-side steps (deletion-type or ε-continuation), then collapse C_τ, and *measure* solely on the B-side $\mathbf{T}_\tau \mathbf{P}_i$.

- Record the δ-ledger for non-commutation and verify B-Gate$^+$ conditions (Chapter 1); if a local cover is used, verify the Overlap Gate on overlaps.
- Issue a windowed certificate when gates pass; global claims are obtained by pasting windowed certificates along the MECE partition using Restart and Summability (Chapter 4).
- If Theorem 2.19 applies (saturation window), *then and only then* use the local PH–Ext equivalence per window in the gate; otherwise retain the one-way implication.

References and provenance. Serre localization $\mathbf{T}_\tau$ and constructible-barcode abelianity follow Crawley–Boevey (2015) and Chazal–de Silva–Glisse–Oudot (2016). Filtered lifts C_τ are asserted up to f.q.i.; exactness and limit compatibilities are stated at the persistence layer. Local PH–Ext equivalence is a window-local theorem under amplitude ≤ 1, saturation, and tail isomorphism; its proof relies on the edge identification in $D^{\mathrm{b}}(k\text{-mod})$ (Appendix C) and tower diagnostics (Appendix D).

3 Chapter 3: A One-Way Bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ and the Hypothesis Scheme

Scope note (reinforced windowed policy). All statements of this chapter lie within the constructible 1D regime of Chapter 2 and use field coefficients. The implication $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ is proved only in $D^{\mathrm{b}}(k\text{-mod})$ under (B1)–(B3) below. Every claim is issued per domain window (right-open; Chapter 1, Def. 1.0 and Chapter 2, Remark 2.2); gate decisions are taken only on the B-side after collapse, i.e. on single-layer objects $\mathbf{T}_\tau\mathbf{P}_i$ (equivalently $\mathbf{P}_i(C_\tau-)$). Filtered-complex equalities hold only up to filtered quasi-isomorphism (f.q.i.).

3.0. Windowed usage and gate integration

This chapter provides the $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ bridge used by B-Gate$^+$ (Chapter 1, Def. 1.0) after applying the collapse C_τ on a fixed window $W = [u, u')$ and a fixed threshold $\tau > 0$:

- On (W, τ), if $\mathrm{PH}_1(C_\tau F) = 0$, then (under (B1)–(B3)) $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$.
- Together with $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$ at (W, τ) and the safety margin $\mathrm{gap}_\tau > \Sigma\delta$, this yields B-Gate$^+$ passing and thus a windowed certificate on (W, τ) (Chapter 1, Def. 1.0).
- Outside the scope of (B1)–(B3), Ext^1 is not used; B-Gate$^+$ may be evaluated with PH and tower-only parts (cf. Chapter 1).

All cross-domain comparisons (PF/BC, Mirror/Transfer) are performed after collapse, and any non-commutation is recorded in the δ-ledger (Appendix L).

3.1. Bridging Hypotheses (B1–B3)

We fix the notation and conventions of Chapter 2. In particular, k is a field, $\mathsf{FiltCh}(k)$ denotes finite-type (constructible) filtered chain complexes, $\mathbf{P}_i(F)$ is the degree-i persistence module with barcode $\mathrm{PH}_i(F)$, and $\mathcal{R} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}(k\text{-mod})$ is a fixed t-exact realization functor.[1]

(B1) Finite-type over a field. $F \in \mathsf{FiltCh}(k)$ with pointwise finite-dimensional persistence. Filtered (co)limits of constructible persistence modules are computed objectwise in $[\mathbb{R}, \mathsf{Vect}_k]$ and used only under the scope policy of Appendix A (compute in the functor category and verify return to $\mathsf{Pers}_k^{\mathrm{cons}}$); no claim is made outside this regime. Equalities are asserted only at the persistence layer.

[1] We work in cohomological amplitude $[-1, 0]$: amplitude control and t-exactness place $\mathcal{R}(F)$ in $D^{[-1,0]}$ and identify the edge map; see Appendix C.

(B2) **Amplitude ≤ 1 and identification of the H^{-1}-edge.** There is a natural isomorphism

$$H^{-1}\big(\mathcal{R}(F)\big) \;\cong\; \varinjlim_t \, H_1(F^t C_\bullet),$$

and $\mathcal{R}(F) \in D^{[-1,0]}(k\text{-mod})$. Equivalently, $\mathcal{R}(F)$ admits a two-term model

$$\mathcal{R}(F) \;\simeq\; \big[\varinjlim_t H_1(F^t C_\bullet) \xrightarrow{d} \varinjlim_t H_0(F^t C_\bullet)\,\big],$$

concentrated in cohomological degrees $(-1,0)$, functorial in F. Here exactness of filtered colimits in Vect_k and the t-exactness of $\mathcal{R}$ ensure functoriality in F; all statements are confined to $D^{[-1,0]}(k\text{-mod})$ (see Appendix C).

(B3) **Edge identification for degree 1 with $Q = k$.** For any $A \in D^{[-1,0]}(k\text{-mod})$,

$$\mathrm{Ext}^1(A,k) \;\cong\; \mathrm{Hom}\big(H^{-1}(A),k\big),$$

naturally in A, by applying $\mathbf{R}\mathrm{Hom}(-,k)$ to the truncation triangle $\tau_{\leq -1}A \to A \to \tau_{\geq 0}A \to$ and using that k is a field; cf. Appendix C.

Remark 3.1 (On (B2) and the edge identification)**.** We use a two-term realization with amplitude ≤ 1 so that $H^{-1}(\mathcal{R}(F)) \cong \varinjlim_t H_1(F^t C_\bullet)$, relying on exactness of filtered colimits in Vect_k. Over a field, for $A \in D^{[-1,0]}$ one has $\mathrm{Ext}^1(A,k) \cong \mathrm{Hom}(H^{-1}(A),k)$. Proof details and naturality are in Appendix C. The bridge of this chapter is proved only in $D^{\mathrm{b}}(k\text{-mod})$. Uses of other coefficient fields (e.g. Novikov) appear only as **[Spec]** and do not extend the proven bridge.

Remark 3.2 (Outside the bridge domain)**.** Derived geometric/sheaf or symplectic/Floer realizations may be employed as **[Spec]** (Appendix N/O), but the implication $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ is not asserted in those targets. All proofs here remain in $D^{\mathrm{b}}(k\text{-mod})$.

3.1 bis. Window energy and spectral primitives on a fixed window

Definition 3.3 (Window barcode and residual length)**.** Fix a window $W = [u,u')$. For $i \geq 0$, let $\mathcal{B}_i(F;W)$ be the multiset of intervals $I \cap W$ obtained by restricting the barcode of $\mathbf{P}_i(F)$ to W and discarding empty intersections. For $J \in \mathcal{B}_i(F;W)$ write $|J|$ for its length. Given $\tau \geq 0$, define the residual length of J at threshold τ by $|J|_\tau := \max\{\,|J| - \tau,\ 0\,\}$.

Definition 3.4 (Window energy and cumulative spectral content)**.** For degree i and parameters (W,τ), define the window energy

$$E_i(F;W,\tau) \;:=\; \sum_{J \in \mathcal{B}_i(F;W)} |J|_\tau,$$

and the cumulative tail counts

$$C_{i,r}(F;W,\tau) \;:=\; \#\{\, J \in \mathcal{B}_i(F;W) \;:\; |J|_\tau \geq r\,\} \qquad (r \geq 0).$$

We abbreviate $E := E_1$ and $C_r := C_{1,r}$ when $i = 1$.

Proposition 3.5 (Reduction-type monotonicity and layer-cake identity)**.** For each fixed F and W:

1. $E_i(F;W,\tau)$ is finite, piecewise-linear in τ, and nonincreasing in τ.

2. For all $r \geq 0$, $C_{i,r}(F;W,\tau)$ is finite, integer-valued, nonincreasing in both r and τ.

3. Layer-cake identity:
$$E_i(F;W,\tau) \;=\; \int_0^\infty C_{i,r}(F;W,\tau)\,dr.$$

4. Right-derivative in τ exists for all but finitely many τ, with
$$-\frac{\partial}{\partial\tau^+}E_i(F;W,\tau) \;=\; C_{i,0}(F;W,\tau).$$

All statements are functorial at the persistence layer and invariant under f.q.i. of filtered complexes.

Remark 3.6 (Compatibility with collapse; nonexpansive continuation)**.** Let $\mathbf{T}_\tau$ denote the threshold–collapse operator on persistence (Chapter 2). Then for each i,
$$\sum_{J\in\mathcal{B}_i(C_\tau F;W)} |J| \;\leq\; E_i(F;W,\tau),$$
so E_i bounds the total barcode length after collapse. Moreover, the semigroup property $C_{\tau+\varepsilon} \simeq C_\varepsilon \circ C_\tau$ implies the nonexpansive continuation
$$E_i(F;W,\tau+\varepsilon) \;\leq\; E_i(F;W,\tau), \qquad C_{i,r}(F;W,\tau+\varepsilon) \;\leq\; C_{i,r}(F;W,\tau).$$

Proposition 3.7 (Stability of window energy and tails)**.** Let d_{int} be the interleaving distance on persistence modules. Then for fixed W and τ, the maps $F \mapsto E_i(F;W,\tau)$ and $F \mapsto C_{i,r}(F;W,\tau)$ are stable under interleavings: there exists a universal constant L (depending only on the window-complexity bound from Chapter 2) such that for any F, G,
$$\big|E_i(F;W,\tau) - E_i(G;W,\tau)\big| \;\leq\; L\cdot d_{\mathrm{int}}\big(\mathbf{P}_i(F)|_W, \mathbf{P}_i(G)|_W\big),$$
and $C_{i,r}$ is upper semicontinuous in the same metric. Proofs use bottleneck stability and the hinge structure of $|J|_\tau$ (Appendix H).

3.2. The Bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$

Recall that $\mathrm{PH}_1(F) = 0$ means the degree-1 persistence module vanishes; equivalently, $H_1(F^tC_\bullet) = 0$ for all t.

Theorem 3.8 (One-way bridge)**.** Assume (B1)–(B3). If $\mathrm{PH}_1(F) = 0$, then
$$\mathrm{Ext}^1\big(\mathcal{R}(F), k\big) \;=\; 0.$$

Equivalently, under (B2)–(B3) the conclusion reads $H^{-1}\big(\mathcal{R}(F)\big) = 0$.

Proof. $\mathrm{PH}_1(F) = 0$ implies $H_1(F^tC_\bullet) = 0$ for all t, hence $\varinjlim_t H_1(F^tC_\bullet) = 0$. By (B2), $H^{-1}(\mathcal{R}(F)) \cong \varinjlim_t H_1(F^tC_\bullet) = 0$. By (B3), $\mathrm{Ext}^1(\mathcal{R}(F), k) \cong \mathrm{Hom}(H^{-1}(\mathcal{R}(F)), k) = 0$. All steps are natural in F. ◻

Corollary 3.9 (Robust bridge at scale ε)**.** For any $\varepsilon > 0$, if $\mathrm{PH}_1\big(C_\varepsilon(F)\big) = 0$, then
$$\mathrm{Ext}^1\big(\mathcal{R}(C_\varepsilon(F)), k\big) \;=\; 0.$$

In particular, robust admissibility is always tested after truncation (apply C_ε first). All identifications are natural in F and compatible with morphisms in $\mathsf{FiltCh}(k)$.

Proof. Apply Theorem 3.8 to $C_\varepsilon(F)$; the hypotheses are preserved by C_ε (Chapter 2). This uses that C_ε preserves constructibility and that, under the lifting–coherence hypothesis, the t-exact realization $\mathcal{R}$ keeps amplitude ≤ 1 (Appendix B; cf. Chapter 2, §§2.2–2.3). ◻

3.3. Naturality, stability, and windowed gate usage

Proposition 3.10 (Naturality of the edge isomorphisms (B2)–(B3))**.** Under (B1)–(B3), the edge isomorphisms in (B2) and (B3) are natural in F. For any morphism $f : F \to G$ in $\mathsf{FiltCh}(k)$ the diagram

$$\begin{tikzcd} \mathrm{H}^{-1}\big(\mathcal{R}(F)\big) \arrow[r, "\sim"] \arrow[d, "\mathrm{H}^{-1}(\mathcal{R}(f))"'] & \varinjlim_t \mathrm{H}_1(F^t C_\bullet) \arrow[d, "\varinjlim_t \mathrm{H}_1(f^t)"] \\ \mathrm{H}^{-1}\big(\mathcal{R}(G)\big) \arrow[r, "\sim"] & \varinjlim_t \mathrm{H}_1(G^t C_\bullet) \end{tikzcd}$$

commutes, and the identifications in (B3) are functorial in A and compatible with morphisms in $D^{[-1,0]}$.

Remark 3.11 (Stability via thresholded collapse and B-Gate$^+$)**.** By Chapter 2, $\mathbf{T}_\tau$ is 1-Lipschitz for the interleaving distance d_{int} and compatible with filtered colimits at the persistence layer (up to f.q.i.; Appendix B). Hence the premise $\mathrm{PH}_1(C_\varepsilon(F)) = 0$ is metrically stable under interleaving perturbations, and the conclusion of Corollary 3.9 is invariant under functorial choices of C_ε. On a window (W, τ), the quantitative primitives E_1 and $C_{1,r}$ supply monotone, piecewise-linear diagnostics (Proposition 3.5) that are stable (Proposition 3.7) and compatible with gate usage after collapse.

3.4. Survival lemma (ε-survival) and safety margins

We record the windowed ε-survival lemma used operationally; proofs and variants appear in Appendix I.

Lemma 3.12 (ε-survival)**.** Fix a window $W = [u, u']$, a threshold $\tau > 0$, and a safety margin $\mathrm{gap}_\tau > \Sigma\delta$ (the δ-ledger total on W). Let F be constructible and let G be ε-interleaved with F on W. If there exists a bar $J \in \mathcal{B}_1(F; W)$ with residual length $|J|_\tau \geq \varepsilon + \mathrm{gap}_\tau$, then:

1. There is a corresponding bar $J' \in \mathcal{B}_1(G; W)$ with $|J'|_\tau \geq \mathrm{gap}_\tau$.
2. After collapse at τ, the image of J' in $\mathrm{PH}_1(C_\tau G)$ survives B-side gating: it is not removed by tower artifacts and is tracked by the $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$ test.

Equivalently, bars whose residual length exceeds $\varepsilon + \Sigma\delta$ cannot be extinguished by an ε-interleaving and the recorded non-commutation on W.

Corollary 3.13 (Lower bounds via tails)**.** With the notation of Lemma 3.12, if $C_r(F; W, \tau) > 0$ for some $r > \varepsilon + \Sigma\delta$, then $C_0(G; W, \tau) > 0$ and hence $\mathrm{PH}_1(C_\tau G) eq 0$.

Remark 3.14 (Operational reading)**.** The ε-survival lemma quantifies *robust presence* at the PH layer: any bar longer than $\varepsilon + \Sigma\delta$ at residual scale τ must show up after collapse and under ε-perturbations. This feeds directly into windowed diagnostics and prevents false positives when discharging the Ext-part via Theorem 3.8.

3.5. Gate indicators and quantitative linkage

We make explicit how the gate indicators interact in the windowed policy.

Definition 3.15 (Gate indicators on (W, τ))**.** On a fixed (W, τ), the PH-indicator is $\mathrm{PH}_1(C_\tau F)$ (vanishing/nonvanishing); the Ext-indicator is $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k)$ in $D^{\mathrm{b}}(k\text{-mod})$; the collapse indicators are $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ (Appendix D), and the spectral indicators are $\{C_r(F; W, \tau)\}_{r \geq 0}$ and $E(F; W, \tau)$.

Proposition 3.16 (One-way discharge and quantitative certificate)**.** Assume (B1)–(B3). Fix (W, τ) with safety margin $\mathrm{gap}_\tau > \Sigma\delta$.

1. If $\mathrm{PH}_1(C_\tau F) = 0$ and $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0, 0)$, then $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$ and B-Gate$^+$ passes on (W, τ).
2. If $C_r(F; W, \tau) = 0$ for some $r > \varepsilon + \Sigma\delta$, then for every G that is ε-interleaved with F on W, $\mathrm{PH}_1(C_\tau G) = 0$; hence the Ext-part discharges for G by Theorem 3.8.
3. If $E(F; W, \tau) \leq \eta$ with $\eta < \varepsilon + \Sigma\delta$, then necessarily $C_r(F; W, \tau) = 0$ for all $r > \eta$; thus item (2) applies.

Remark 3.17 (Quantitative but non-analytic)**.** The indicators E and C_r are purely topological/persistence-theoretic. We avoid any claim that topological collapse implies analytic regularization. Nonetheless, the monotonicity, piecewise-linearity, and stability of E and C_r offer effective quantitative tools for operational use within the proven bridge domain.

3.6. Scope and limitations

The bridge is strictly one-way: no claim is made that $\mathrm{Ext}^1(\mathcal{R}(F), k) = 0$ implies $\mathrm{PH}_1(F) = 0$. Failure modes (e.g. invisible obstructions beyond the amplitude window) and counterexamples to the converse are documented in Appendix D, including subsection **D.4** (Counterexamples to the converse). Type IV/tower artifacts are detected by $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$; see Appendix D.1–D.3 and Remark A.1 for the generic-fiber (multiplicity of $I[0, \infty)$) interpretation.

3.7. Windowed saturation and domain-restricted triggers [Spec]

The following **[Spec]** items record useful strengthening on fixed windows that may apply in restricted domains; they are not part of the proved bridge.

Declaration 3.18 (Saturation gate: local equivalence [Spec])**.** On a saturation window (Chapter 11), where the event set stabilizes, the inter-window drift is $\leq \eta$, and the edge gap exceeds η, it is **[Spec]**-admissible to adopt a temporary local equivalence:

$$\mathrm{PH}_1(C_{\tau^*} F) = 0 \quad \Longleftrightarrow \quad \mathrm{Ext}^1\big(\mathcal{R}(C_{\tau^*} F), k\big) = 0.$$

This policy is window-local and must be explicitly logged; it does not extend the global one-way bridge outside the saturation conditions.

Declaration 3.19 (Trigger pack (domain-restricted) [Spec])**.** In restricted regimes (e.g. 2D flows, small-data critical spaces, simplified arithmetic models), one may prepare *trigger lemmas* asserting that analytic deviations imply failure of B-Gate$^+$ on the window, e.g.:

$$(\text{blow-up/instability on } W) \;\Rightarrow\; \big(\mathrm{PH}_1(C_\tau F) > 0 \;\text{ or }\; \mu_{\text{Collapse}} > 0 \;\text{ or }\; \text{aux}-\text{bars} > 0\big).$$

Such triggers (when available) supply partial necessity, improving the diagnostic power of the gate. They must be stated with explicit assumptions and logged as **[Spec]**.

3.8. Interaction with PF/BC, Mirror, and the δ-ledger

PF/BC isomorphisms are applied per t, transported to the persistence layer, and then to the collapsed layer (Appendix N). Mirror/Transfer comparisons are performed only after C_τ, and any non-commutation with collapse is recorded in the δ-ledger (Appendix L). The Ext-part of B-Gate$^+$ is always checked on the collapsed object and only within the t-exact, amplitude ≤ 1 scope. No Ext $\Rightarrow$ PH claim is made.

3.9. Formalizability

Hypotheses (B1)–(B3) and Theorem 3.8 admit direct formalization: (B2) via an explicit two-term model for $\mathcal{R}(F)$ and exactness of filtered colimits in Vect_k; (B3) via truncation in $D^{[-1,0]}$ and the edge of the long exact sequence. The primitives E and C_r are definable in the barcode semantics; Proposition 3.5 and Proposition 3.7 reduce to elementary combinatorics plus bottleneck stability (Appendix H). Skeletons and lemma names are listed in Appendix F for Coq/Lean integration. Windowed usage (B-Gate$^+$), the MECE policy, and the δ-ledger are recorded as operational axioms in the formalization stubs (Appendix F), with proofs confined to the persistence layer and the derived category of k-modules.

3.10. Degree specificity

All energy/spectral statements in this chapter are used in degree 1; no claim is made for ($\mathrm{PH}_i \Rightarrow \mathrm{Ext}^i$) with $ieq1$.

4 Chapter 4: Failure Lattice, Local PH–Ext Equivalence, Čech–Ext Gluing, and the Tower-Sensitivity Invariant μ_{Collapse}

Standing hypotheses and scope. We work in the constructible (finite-type) range of persistence (Chapter 2, §2.1), adopt the bridging hypotheses (B1)–(B3) from Chapter 3 with the minimal test family $Q = \{k[0]\}$, and fix a t-exact realization $\mathcal{R} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}(k\text{-mod})$. All uses of filtered (co)limits are asserted *at the persistence layer* (Appendix A). Endpoint conventions and the treatment of infinite bars follow Chapter 2, Remark 2.3. Monotonicity claims for indicators apply only to *deletion-type* updates; inclusion-type updates are *stability-only* (Appendix E). Every claim is *windowed* (Chapter 1, Def. 1.0; Chapter 2, §2.4), and all gate decisions are taken *only* on the B-side after collapse (single layer $\mathbf{T}_\tau \mathbf{P}_i$). We *do not* assert any global equivalence $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$; the *only* core bridge is the one-way implication $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ in $D^{\mathrm{b}}(k\text{-mod})$ under (B1)–(B3) (Chapter 3 and Appendix C).

(B2) (edge identification, recall). There is a natural isomorphism $H^{-1}(\mathcal{R}(F)) \cong \varinjlim_t H_1(F^t C_\bullet)$ and $\mathcal{R}(F) \in D^{[-1,0]}$; see Appendix C.

4.1. Failure lattice and observable vs. invisible modes

We organize collapse failures as follows (cf. Chapter 1, §1.4):

- **Type I (Topological):** $\mathrm{PH}_1(F)eq0$.
- **Type II (Categorical):** $\mathrm{Ext}^1\big(\mathcal{R}(F), k\big)eq0$ (with $Q = \{k[0]\}$).
- **Type III (Functorial/[Spec]):** admissibility unstable under a prescribed operation (e.g. a given pull-back/filtered colimit) at finite level.
- **Type IV (Invisible/tower-level):** all finite layers appear admissible while the limit is not; detected by the tower-sensitivity invariants below.

Types I–II are *observable* on a single object; Type III is *specification-level*; Type IV is *invisible* at finite layers and requires tower diagnostics. We write $I[a,b)$ for the interval module (constructible, p.f.d.) supported on $[a,b)$; endpoint choices do not affect lengths (Chapter 2, Remark 2.3). We emphasize again that *no* global equivalence $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is claimed; only the one-way core bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ under (B1)–(B3) (Chapter 3; Appendix C).

Remark 4.1 (Specification-level failures and functorial calculus)**.** Type III is analyzed using Chapter 2, §2.3: non-expansiveness, shift–commutation (Lemma 2.4), and the persistence-layer (co)limit/pullback compatibilities in Proposition 2.5 (4)(5). Statements at the filtered-complex layer are *adopted as* **[Spec]** and used only up to filtered quasi-isomorphism (Appendix B).

Remark 4.2 (Model towers)**.** Pure-kernel, pure-cokernel, and mixed toy towers, together with the vanishing regime under constructible filtered colimits, are illustrated in Appendix D (D.1–D.3). Counterexamples to the converse $\mathrm{Ext}^1 = 0 \Rightarrow \mathrm{PH}_1 = 0$ are in D.4 (see also Appendix C).

4.2. The tower-sensitivity invariants μ_{Collapse}, u_{Collapse}, and the Defect functor

Generic fiber dimension. As in Chapter 1, §1.4 and Appendix D, Remark A.1, for $M \in \mathsf{Pers}_k^{\text{cons}}$ the *generic fiber dimension* is $\operatorname{gdim}(M) = \lim_{t\to+\infty} \dim_k M(t)$; after truncation by $\mathbf{T}_\tau$, it equals the multiplicity of the infinite bar $I[0,\infty)$ in the barcode.

Comparison map. Fix $\tau > 0$. Let $\{F_n\}_{n\in\mathbb{N}}$ be a directed system in $\mathsf{FiltCh}(k)$ with colimit F_∞. For each degree i, the comparison map of truncated persistence modules is

$$\phi_{i,\tau} : \varinjlim_n \mathbf{T}_\tau\big(\mathbf{P}_i(F_n)\big) \longrightarrow \mathbf{T}_\tau\big(\mathbf{P}_i(F_\infty)\big),$$

where $\mathbf{T}_\tau$ is the exact reflector deleting all bars of length $\le \tau$ (Chapter 2, §2.2) and $\mathbf{P}_i$ is degreewise persistence.

Definition 4.3 (Defect object and tower-sensitivity invariants)**.** For each i and $\tau > 0$, define the *Defect objects* in $\mathsf{Pers}_k^{\text{cons}}$ by

$$\mathsf{Defect}^{\text{ker}}_{i,\tau}(F_\bullet) := \ker(\phi_{i,\tau}), \mathsf{Defect}^{\text{coker}}_{i,\tau}(F_\bullet) := \operatorname{coker}(\phi_{i,\tau}).$$

The *tower-sensitivity invariants* are the generic fiber dimensions

$$\mu_{i,\tau} := \operatorname{gdim}\big(\mathsf{Defect}^{\text{ker}}_{i,\tau}\big), u_{i,\tau} := \operatorname{gdim}\big(\mathsf{Defect}^{\text{coker}}_{i,\tau}\big), \mu_{\text{Collapse}} := \sum_i \mu_{i,\tau},\ u_{\text{Collapse}} := \sum_i u_{i,\tau}.$$

Since complexes are bounded in homological degrees and barcodes are constructible, $\mu_{\text{Collapse}}, u_{\text{Collapse}} < \infty$ on bounded τ-windows.

Proposition 4.4 (Generic dimension equals infinite-bar multiplicity)**.** In the constructible range, for any morphism $\psi : M \to N$ in $\mathsf{Pers}_k^{\text{cons}}$ the generic fiber dimension $\operatorname{gdim}\ker(\psi)$ (resp. $\operatorname{gdimcoker}(\psi)$) equals the multiplicity of $I[0,\infty)$ in the barcode of $\ker(\psi)$ (resp. $\operatorname{coker}(\psi)$). The same holds after truncation by $\mathbf{T}_\tau$.

Proof sketch. Kernels and cokernels are computed pointwise and remain constructible (Chapter 2, Remark 2.1), hence their barcodes are defined. Interval Jordan–Hölder theory identifies the stabilized rank with the $I[0,\infty)$ multiplicity; truncation preserves $I[0,\infty)$ and removes only finite bars (Chapter 2, §2.2). See Appendix D.

Remark 4.5 (Functoriality, invariance, and calculus)**.** The maps $\phi_{i,\tau}$ are natural in the tower, independent of filtered representatives, and invariant under cofinal reindexing (Appendix J). Consequently, $\mathsf{Defect}^{\text{ker/coker}}_{i,\tau}$ and $(\mu_{i,\tau}, u_{i,\tau})$ are invariant under filtered quasi-isomorphisms. Subadditivity under composition, additivity under finite sums, and cofinal invariance are collected in Appendix J.

Proposition 4.6 (Vanishing under constructible filtered colimits)**.** Assume the colimit of the tower $\{F_n\}$ is computed objectwise (chains/filtrations) and that all $\mathbf{P}_i(F_n)$ remain in $\mathsf{Pers}_k^{\text{cons}}$. Then for every $\tau > 0$ and every i, $\phi_{i,\tau}$ is an isomorphism, hence

$$\mathsf{Defect}^{\text{ker}}_{i,\tau} = \mathsf{Defect}^{\text{coker}}_{i,\tau} = 0, \qquad \mu_{\text{Collapse}} = u_{\text{Collapse}} = 0.$$

A complete calculus (subadditivity, additivity, cofinal invariance) is given in Appendix J.

4.3. Local PH–Ext equivalence on saturation windows: proof overview (core)

We recall the window-local equivalence, formally stated in Chapter 2 (Theorem 2.19) and now proved within our guard-rails.

Theorem 4.7 (Local PH–Ext equivalence on saturation windows)**.** Let $F \in \mathsf{FiltCh}(k)$, let $W = [u, u')$ be a right-open window, and fix $\tau > 0$. Assume:

1. *(Amplitude)* t-exact realization $\mathcal{R}$ has amplitude ≤ 1: $\mathcal{R}(C_\tau F) \in D^{[-1,0]}(k\text{-mod})$.
2. *(Saturation/gap)* W is a saturation window at τ (Chapter 2, Def. 2.6): event stability on W, length gap $\leq \tau - \eta$ for some $\eta > 0$, and near-τ non-accumulation.
3. *(Tail isomorphism)* The comparison map $\phi_{1,\tau}$ on W is an isomorphism, i.e. $\mu_{\mathrm{Collapse}} = u_{\mathrm{Collapse}} = 0$ on W for degree 1.

Then, *on the window W and at threshold τ,*

$$\mathrm{PH}_1\big(C_\tau F\big) = 0 \quad \iff \quad \mathrm{Ext}^1\big(\mathcal{R}(C_\tau F), k\big) = 0.$$

Proof sketch. Under (1) we may identify $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) \cong \mathrm{Hom}(H^{-1}(\mathcal{R}(C_\tau F)), k)$ (Appendix C). Under (2) and (3), the degree-1 persistence on W is fully determined by stabilized bars (no near-τ creation/annihilation in the limit), and the tail isomorphism excludes Type IV drift at scale τ. Thus $\mathrm{PH}_1(C_\tau F) = 0$ on W iff $H^{-1}(\mathcal{R}(C_\tau F)) = 0$ on W, i.e. iff $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$. □

Remark 4.8 (Core vs. global)**.** The theorem is *window-local* and sits within the core (no [Spec] clauses). It *does not* assert any global $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$; outside saturation windows we only use $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ (Chapter 3).

4.4. Čech–Ext^1 gluing and the Overlap Gate

We integrate local-to-global gluing via the Overlap Gate (Chapter 1, Def. 1.0) and a Čech-type acyclicity on overlaps.

Definition 4.9 (Čech nerve and local data after collapse)**.** Let $\{X_\alpha\}$ be a cover of the domain, and let $\{W_\alpha = [u_\alpha, u'_\alpha)\}$ be a right-open windowing. For a fixed degree i and threshold $\tau > 0$, set

$$\mathcal{B}_{\alpha,i} \;:=\; \mathbf{T}_\tau\, \mathbf{W}_{W_\alpha}\big(\mathbf{P}_i(F|_{X_\alpha})\big) \;\in\; \mathsf{Pers}_k^{\mathrm{cons}}.$$

Denote higher overlaps by $\mathcal{B}_{\alpha_0\cdots\alpha_p,i}$ via restriction to $X_{\alpha_0\cdots\alpha_p} \times (W_{\alpha_0} \cap \cdots \cap W_{\alpha_p})$. Write $N(\mathcal{U})$ for the Čech nerve of the cover.

Definition 4.10 (Čech–Ext^1–acyclicity (after collapse))**.** We say that the local collapsed data are *Čech–Ext^1–acyclic* on degree 1 if, for each $p \geq 0$ and each nonempty overlap $X_{\alpha_0\cdots\alpha_p}$,

$$\mathrm{Ext}^1\Big(\mathcal{R}\big(C_\tau(F|_{X_{\alpha_0\cdots\alpha_p}})\big),\, k\Big) \;=\; 0,$$

and the associated Čech differential on p-cochains takes values in $\mathrm{Ext}^1(-, k)$ that vanish on all overlaps. Equivalently, the E_1–page of the Čech spectral sequence has Ext^1–row identically zero.

Theorem 4.11 (Overlap Gate with Čech–Ext^1: local-to-global gluing)**.** Fix a degree $i = 1$, a globally right-open windowing, and a threshold $\tau > 0$. Assume:

1. *(Local gates)* On each chart $X_\alpha \times W_\alpha$ the local B–Gate$^+$ passes: $\mathrm{PH}_1(C_\tau F|_{X_\alpha}) = 0$, $\mathrm{Ext}^1(\mathcal{R}(C_\tau F|_{X_\alpha}), k) = 0$, $(\mu, u) = (0, 0)$ on W_α, with safety margin $\mathrm{gap}_\tau > \Sigma\delta_\alpha$.

2. *(Overlap Gate)* For any pair (α, β) with $X_{\alpha\beta} \times (W_\alpha \cap W_\beta) eq \varnothing$, Overlap Gate (Chapter 1, Def. 1.0) passes: the collapsed restrictions agree up to the recorded budget, the safety margin dominates the budget on overlaps, and the tower diagnostics vanish $(\mu, u) = (0, 0)$ on the overlap window.

3. *(Čech–Ext^1–acyclicity)* The local collapsed data are Čech–Ext^1–acyclic in degree 1 (Definition 4.10).

Then the global B–Gate$^+$ passes on $\bigcup_\alpha X_\alpha \times W_\alpha$ at threshold τ: in particular,

$$\mathrm{PH}_1(C_\tau F) = 0, \qquad \mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0, \qquad (\mu, u) = (0, 0),$$

with a safety margin controlled by the minima of the local margins after subtracting the overlap budgets.

Proof sketch. By (1) and (2), the objectwise collapsed local persistence modules glue along overlaps to a global collapsed persistence object in degree 1 (Overlap Gate gives isomorphism on overlaps after collapse). Čech–Ext^1–acyclicity in (3) implies that, after applying the amplitude–≤ 1 realization and restricting to degree (-1)–cohomology, the Čech complex has vanishing H^1; hence the global $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k)$ vanishes. The local PH_1–vanishing and tail isomorphism on overlaps (by (2)) give global $\mathrm{PH}_1(C_\tau F) = 0$ on the glued object. The tower diagnostics vanish globally because they vanish locally and on overlaps, and the Čech differentials preserve the collapsed layer; safety margins subtract the overlap budgets but stay positive by (1)–(2). □

Remark 4.12 (Practical check)**.** In practice one verifies (3) by checking $\mathrm{Ext}^1(\mathcal{R}(C_\tau F|_{X_{\alpha_0 \cdots \alpha_p}}), k) = 0$ for $p = 0, 1$ (Mayer–Vietoris row) under an appropriate acyclicity regime (e.g. affine charts, tame restriction), and records the budget in the δ–ledger. The after-collapse policy ensures that all checks occur on the B–side single layer.

4.5. Type IV: finite admissibility need not pass to the limit

The following precise form of the invisible failure principle works already at the persistence layer and then lifts.

Proposition 4.13 (Type IV: finite-level admissibility may fail at the limit)**.** In the filtered index category $\mathbb{N} \cup \{\infty\}$ with cone apex ∞, there exists a tower $\{F_n\}$ and $\tau > 0$ such that

$$\forall n: \quad \mathrm{PH}_1\big(C_\tau(F_n)\big) = 0 \quad \text{and} \quad \mathrm{Ext}^1\big(\mathcal{R}(C_\tau(F_n)), k\big) = 0,$$

yet

$$\mathrm{PH}_1\big(C_\tau(F_\infty)\big) eq 0 \quad \text{and thus} \quad \mathrm{Ext}^1\big(\mathcal{R}(C_\tau(F_\infty)), k\big) eq 0.$$

Moreover one can arrange $\mu_{\mathrm{Collapse}} = 0$ and $u_{\mathrm{Collapse}} > 0$ (pure cokernel-type mismatch).

Proof (persistence model; then lift). Fix $\tau > 0$ and set $\ell_n = \tau - \frac{1}{n} \uparrow \tau$. Let $M_n := I[0, \ell_n)$ with structure maps $M_n \hookrightarrow M_{n+1}$. Then $\mathbf{T}_\tau(M_n) = 0$ for all n, while $\varinjlim_n M_n \cong I[0, \tau)$. Cone the diagram to the apex module $N := I[0, \infty)$ via $M_n \hookrightarrow I[0, \tau) \hookrightarrow N$, so the colimit of the extended diagram is N. Choose F_n, F_∞ with $\mathbf{P}_1(F_n) \cong M_n$, $\mathbf{P}_1(F_\infty) \cong N$ (Chapter 2, §2.2). Then $\mathbf{T}_\tau(\mathbf{P}_1(F_n)) = 0$ while $\mathbf{T}_\tau(\mathbf{P}_1(F_\infty)) \cong I[0, \infty)$. Thus $\mathrm{PH}_1(C_\tau(F_n)) = 0$ for all n, but $\mathrm{PH}_1(C_\tau(F_\infty)) eq 0$. By (B2)–(B3) this forces $\mathrm{Ext}^1\big(\mathcal{R}(C_\tau(F_\infty)), k\big) eq 0$. Here $\ker(\phi_{1,\tau}) = 0$ and $\mathrm{coker}(\phi_{1,\tau}) eq 0$, so $\mu_{\mathrm{Collapse}} = 0$ and $u_{\mathrm{Collapse}} > 0$. □

$$\begin{tikzcd}
I[0,\tau-\frac{1}{1}) \arrow[r, hook] \arrow[dr] & I[0,\tau-\frac{1}{2}) \arrow[r, hook] & \cdots \arrow[r, hook] \arrow[dr] & I[0,\tau-\frac{1}{n}) \arrow[r, hook] & \cdots \arrow[r, hook] \arrow[dr] & I[0,\tau) \arrow[r, hook] \arrow[dr] & I[0,\infty) \\
 & \mathbf{T}_\tau(\cdot)=0 & & \mathbf{T}_\tau(\cdot)=0 & & \mathbf{T}_\tau(\cdot)=0 & \mathbf{T}_\tau(\cdot)=I[0,\infty)
\end{tikzcd}$$

Figure 1: Type IV intuition (after truncation by $\mathbf{T}_\tau$): finite layers vanish, while the apex produces an infinite bar.

4.6. A natural refinement-limit example (pure cokernel type)

Example 4.14 (Resolution refinement producing a limit infinite bar)**.** Fix $\tau > 0$. Consider a mesh-refinement sequence of scalar fields whose degree-1 persistence at level n has a single bar $[0, \tau - \delta_n)$ with $\delta_n \downarrow 0$. Each finite layer satisfies $\mathbf{T}_\tau(\mathbf{P}_1(F_n)) = 0$. In the refinement limit the coherent structure becomes persistent across all thresholds, yielding an infinite bar in degree 1: $\mathbf{T}_\tau(\mathbf{P}_1(F_\infty)) \cong I[0,\infty)$. This is the natural analogue of Proposition 4.13 in resolution limits (see Appendix D for diagrams). It realizes a *pure cokernel* Type IV failure ($\mu_{\text{Collapse}} = 0$, $u_{\text{Collapse}} > 0$).

Remark 4.15 (When invisible failure is excluded)**.** If the tower satisfies the hypotheses of Proposition 4.6 (and Proposition 2.5 (4) in Chapter 2), then $\phi_{i,\tau}$ is an isomorphism for all i, τ, so no Type IV occurs: for every $\tau > 0$,

$$\big(\forall n:\ \mathrm{PH}_1(C_\tau(F_n)) = 0\big) \implies \mathrm{PH}_1(C_\tau(F_\infty)) = 0,$$

and, under (B1)–(B3), the corresponding Ext^1-vanishing propagates along the tower.

4.7. Restart lemma, summability, and pasting of windowed certificates

We formalize the pasting principle used to obtain global conclusions from windowed certificates.

Definition 4.16 (Per-window safety margin and pipeline budget)**.** Let $\{[u_k, u_{k+1})\}_k$ be a MECE partition of the monitored range (Chapter 2, Def. 2.9). On each window $W_k = [u_k, u_{k+1})$, fix a collapse threshold $\tau_k > 0$ and define the *per-window pipeline budget*

$$\Sigma\delta_k(i) \;:=\; \sum_{j\in J_k} \big(\delta_j^{\text{alg}}(i,\tau_k) + \delta_j^{\text{disc}}(i,\tau_k) + \delta_j^{\text{meas}}(i,\tau_k)\big),$$

where J_k indexes the A-side steps applied before the B-side gate on W_k. The *per-window safety margin* $\mathrm{gap}_{\tau_k} > 0$ is the admissible slack chosen for the gate on W_k (Chapter 1, Def. 1.0).

Lemma 4.17 (Restart lemma (window-to-window inheritance))**.** Fix a degree i. Suppose that on W_k the B-Gate$^+$ passes with safety margin $\mathrm{gap}_{\tau_k} > \Sigma\delta_k(i)$, and that the next window W_{k+1} is reached via a finite number of admissible A-side steps, each either deletion-type or an ε-continuation, followed by collapse $C_{\tau_{k+1}}$ with τ_{k+1} chosen by an adaptation rule (Chapter 2, Def. 2.10). Then there exists $\kappa \in (0,1]$, depending only on the admissible step class and the adaptation policy, such that the safety margin on W_{k+1} satisfies

$$\mathrm{gap}_{\tau_{k+1}} \;\geq\; \kappa\,\big(\mathrm{gap}_{\tau_k} - \Sigma\delta_k(i)\big).$$

In particular, if $\mathrm{gap}_{\tau_k} - \Sigma\delta_k(i) > 0$, then $\mathrm{gap}_{\tau_{k+1}} > 0$ and B-Gate$^+$ can pass on W_{k+1} provided the new budget $\Sigma\delta_{k+1}(i)$ is sufficiently small.

Proof sketch. Deletion-type steps are non-increasing after collapse (Chapter 2, Remark 2.7); ε-continuations are 1-Lipschitz hence introduce a controlled drift proportional to the declared ε. The adaptation of τ preserves the scale of the clipped indicators. Aggregating the drifts yields a multiplicative retention κ of the effective margin.

Definition 4.18 (Summability policy)**.** A run satisfies the *summability policy* if the sequence of per-window budgets on a MECE partition obeys

$$\sum_k \Sigma\delta_k(i) \;<\; \infty$$

for the degrees i on which gates are evaluated. A sufficient design pattern is a geometric decay of step sizes (e.g. $\tau_k = \tau_0\rho^k$, $\beta_k = \beta_0\rho^k$ with $\rho \in (0,1)$) and/or a geometric damping of the number and strength of ε-continuations per window, all recorded in the manifest (Appendix G).

Theorem 4.19 (Pasting windowed certificates)**.** Let $\{[u_k, u_{k+1})\}_k$ be a MECE partition, and suppose that on each window W_k the B-Gate$^+$ passes with $\mathrm{gap}_{\tau_k} > \Sigma\delta_k(i)$. If the summability policy holds (Definition 4.18) and the restart lemma (Lemma 4.17) is applicable at each transition, then the concatenation of windowed certificates yields a global certificate on $\bigcup_k [u_k, u_{k+1})$, i.e. no failure occurs with no false negative on the union of windows for the monitored degrees i.

Proof sketch. By Lemma 4.17, positive safety margin propagates to the next window with a controlled retention factor. Summability of drifts ensures that the cumulative loss of margin remains bounded and does not exhaust the initial slack. The tower audit per window ensures $(\mu, u) = (0,0)$ at each scale and excludes Type IV; the MECE coverage check (Chapter 2, Def. 2.9) guarantees the union covers the monitored range without gaps or double counting.

4.8. Stable bands and τ-sweeps

We record the τ-selection policy used to stabilize the tower audit.

Definition 4.20 (Stable band)**.** For a fixed window W and degree i, a *stable band* is a contiguous range $B \subset (0,\infty)$ such that for all $\tau \in B$ the comparison maps $\phi_{i,\tau}$ are isomorphisms and hence $(\mu_{i,\tau}, u_{i,\tau}) = (0,0)$. A *$\tau$-sweep* is a discrete set $\{\tau_\ell\}$ used to probe $(\mu_{i,\tau_\ell}, u_{i,\tau_\ell})$; a band is declared stable when the sweep detects $(0,0)$ on a consecutive subarray and the outcome persists under refinement of the sweep.

Remark 4.21 (Using stable bands)**.** Windowed certificates are issued within stable bands only. Outside a stable band, the collapse threshold or the pipeline must be redesigned, or the window must be refined. This stabilizes the tower audit across steps and prevents spurious Type IV readings due to poor scale selection.

4.9. Summary

The failure lattice separates observable (Types I–II), specification-level (Type III), and invisible tower effects (Type IV). The Defect objects $\mathrm{Defect}^{\mathrm{ker/coker}}_{i,\tau}$ and the invariants $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ provide principled tower diagnostics: they *vanish* under constructible filtered colimits (Proposition 4.6) and, when positive, certify instability of persistence under passage to limits. The *only* core bridge is $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ in $D^{\mathrm{b}}(k\text{-mod})$; no global equivalence is asserted. Nevertheless, on *saturation windows* with tail isomorphism, we obtain the local equivalence $\mathrm{PH}_1(C_\tau F) = 0 \Leftrightarrow \mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$ (Theorem 4.7). At the level of gluing, the Overlap Gate combined with Čech–Ext^1–acyclicity yields robust local-to-global propagation of B–Gate$^+$ (Theorem 4.11). The Restart lemma and the Summability policy provide a reliable pasting principle (Theorem 4.19) for *windowed certificates*, while stable bands (Definition 4.20) guide τ-selection for tower audits. Counterexamples to converse statements (Ext⇒PH, finite-to-limit admissibility) appear in Appendix C and Appendix D (see especially Remark A.1 and D.4). These invariants and gluing statements are *not* the classical Iwasawa μ, λ–invariants nor their direct analogues; our usage is purely persistence–theoretic and audit–oriented on the B–side single layer after collapse.

5 Chapter 5: Functoriality, Set-Theoretic Coherence, and Formalization Specifications (Proof/Spec)

All adjunction and (co)limit statements in this chapter are made in the *implementable range* and inside $\mathrm{Ho}(\mathsf{FiltCh}(k))$, *up to filtered quasi-isomorphism* (Appendix B). Equalities are asserted *at the persistence layer*. We retain the standing conventions of Chapters 1–4: constructible range, field coefficients, t-exact realization, and the after-truncation policy. Endpoint conventions and infinite bars are as in Chapter 2, Remark 2.3. *Monotonicity claims apply only to deletion-type updates; inclusion-type updates are stability-only* (Appendix E). All statements are *windowed* and gates are evaluated *after collapse* (B-side single layer).

5.1. Exactness, (Co)Limit Behavior, and a Right-Adjoint Collapse in the Implementable Range (up to f.q.i.)

Let k be a field. Recall: $\mathsf{FiltCh}(k)$ is the category of finite-type (constructible) filtered chain complexes; $\mathbf{P}_i$ is degreewise persistence; $\mathbf{T}_\tau$ is the exact truncation deleting all bars of length $\le \tau$ (Chapter 2, §2.2); C_τ is any filtered lift of $\mathbf{T}_\tau$ (Chapter 2, §§2.2–2.3; always *up to f.q.i.*); and $\mathcal{R} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}(k\text{-mod})$ is t-exact. We also keep the minimal test family $Q = \{k[0]\}$.

Persistence-level (reflective) adjunction. Let $\mathsf{Pers}_k^{\mathrm{cons}}$ be the abelian category of constructible k-persistence modules and $\mathsf{Pers}_{k,\tau\text{-tf}}^{\mathrm{cons}} \subset \mathsf{Pers}_k^{\mathrm{cons}}$ the full subcategory of τ-torsion-free objects (no composition factors of length $\le \tau$). As in Chapter 2, §§2.2–2.3:

- $\mathsf{E}_\tau \subset \mathsf{Pers}_k^{\mathrm{cons}}$ (generated by interval modules of length $\le \tau$) is hereditary Serre (localizing).
- The reflector
$$\mathbf{T}_\tau : \mathsf{Pers}_k^{\mathrm{cons}} \longrightarrow \mathsf{Pers}_{k,\tau\text{-tf}}^{\mathrm{cons}}$$
is exact and exhibits a *reflective* adjunction $\mathbf{T}_\tau \dashv \iota_\tau$ with the inclusion $\iota_\tau : \mathsf{Pers}_{k,\tau\text{-tf}}^{\mathrm{cons}} \hookrightarrow \mathsf{Pers}_k^{\mathrm{cons}}$. Consequently, $\mathbf{T}_\tau$ preserves finite limits and colimits and is 1-Lipschitz for d_{int} (Lemma 2.4, Proposition 2.5(1),(3)). As $\mathbf{T}_\tau$ is exact in this abelian setting, it preserves both finite limits and finite colimits.

Filtered-complex level (operational coreflection; implementable range, up to f.q.i.). Define the full subcategory

$$\mathsf{S}_\tau := \Big\{ F \in \mathsf{FiltCh}(k) \;\Big|\; \forall i,\ \mathbf{P}_i(F) \text{ is } \tau\text{-torsion-free and } \operatorname{Ext}^1\big(\mathcal{R}(F), k\big) = 0 \Big\},$$

and let $\mathsf{S}_\tau^{\mathrm{h}}$ be its image in $\mathrm{Ho}(\mathsf{FiltCh}(k))$. For later comparison, recall the "trivial-at-τ" gate

$$\mathsf{Triv}_\tau := \Big\{ F \;\Big|\; \mathrm{PH}_1\big(C_\tau(F)\big) = 0 \text{ and } \operatorname{Ext}^1\big(\mathcal{R}(F), k\big) = 0 \Big\}, \qquad \mathsf{Triv}_\tau^{\mathrm{h}} \subset \mathsf{S}_\tau^{\mathrm{h}}.$$

Proposition 5.1 (Operational collapse as a right adjoint (implementable range; up to f.q.i.; Proof/Spec))**.** Assume (B1)–(B3) (Chapter 3) and the lifting–coherence hypothesis (Appendix B). Then there exists a functor

$$\mathsf{C}_\tau^{\mathrm{comb}} : \mathrm{Ho}(\mathsf{FiltCh}(k)) \longrightarrow \mathsf{S}_\tau^{\mathrm{h}}$$

and a natural transformation $\eta : \mathrm{Id} \Rightarrow \iota\, \mathsf{C}_\tau^{\mathrm{comb}}$ (with $\iota : \mathsf{S}_\tau^{\mathrm{h}} \hookrightarrow \mathrm{Ho}(\mathsf{FiltCh}(k))$ the inclusion) such that, *in this regime*,

1. (Adjunction) $\mathsf{C}_\tau^{\mathrm{comb}}$ is right adjoint to ι:
$$\mathrm{Hom}\big(\iota(G), F\big) \;\cong\; \mathrm{Hom}\big(G, \mathsf{C}_\tau^{\mathrm{comb}}(F)\big) \qquad (G \in \mathsf{S}_\tau^{\mathrm{h}}).$$

2. (Compatibility; persistence layer) For each i, $\mathbf{P}_i\big(\mathsf{C}^{\mathrm{comb}}_\tau(F)\big) \cong \mathbf{T}_\tau\big(\mathbf{P}_i(F)\big)$ in $\mathsf{Pers}^{\mathrm{cons}}_k$.

3. (Compatibility; realization layer) $\mathcal{R}\big(\mathsf{C}^{\mathrm{comb}}_\tau(F)\big) \cong \tau_{\geq 0}\,\mathcal{R}(F)$ in $D^{\mathrm{b}}(k\text{-mod})$.

4. (Soundness) $\mathsf{C}^{\mathrm{comb}}_\tau(F) \in \mathsf{S}^{\mathrm{h}}_\tau$ for all F, and

$$\mathsf{C}^{\mathrm{comb}}_\tau(F) \in \mathsf{Triv}^{\mathrm{h}}_\tau \quad\Longleftrightarrow\quad \mathbf{T}_\tau\big(\mathbf{P}_1(F)\big) = 0 \ \big(\text{i.e. } \mathrm{PH}_1(C_\tau(F)) = 0\big).$$

All equalities above are asserted at the persistence layer, and all filtered-complex statements are in $\mathrm{Ho}(\mathsf{FiltCh}(k))$ *up to f.q.i. only.*

Proof sketch. At persistence level, use $\mathbf{T}_\tau$ (implemented by its lift C_τ) so that $\mathbf{P}_i(C_\tau F) \cong \mathbf{T}_\tau(\mathbf{P}_i F)$. At realization level, use the coreflection $\tau_{\geq 0}$ to enforce amplitude ≤ 1, under which $\mathrm{Ext}^1(-,k) = 0 \iff H^{-1}(-) = 0$ (Chapter 3). Lifting–coherence provides functorial comparison maps $\mathcal{R} \circ C_\tau \Rightarrow \tau_{\geq 0} \circ \mathcal{R}$ (up to f.q.i.), yielding $\mathsf{C}^{\mathrm{comb}}_\tau$ valued in $\mathsf{S}^{\mathrm{h}}_\tau$; triangle identities hold up to f.q.i. in $\mathrm{Ho}(\mathsf{FiltCh}(k))$. □

Corollary 5.2 (Limit/(co)limit behavior of $\mathsf{C}^{\mathrm{comb}}_\tau$ (persistence layer))**.** Under Proposition 2.5(4),(5), $\mathsf{C}^{\mathrm{comb}}_\tau$ preserves finite limits (in particular, finite pullbacks) *up to f.q.i.* at the filtered-complex level; at the *persistence layer* one has, for every degree i,

$$\mathbf{P}_i\big(\mathsf{C}^{\mathrm{comb}}_\tau(\varinjlim_\Lambda F_\lambda)\big) \;\cong\; \varinjlim_\Lambda \mathbf{P}_i\big(\mathsf{C}^{\mathrm{comb}}_\tau(F_\lambda)\big).$$

Hence $\mathsf{C}^{\mathrm{comb}}_\tau$ is 1-Lipschitz at the persistence layer and inherits exactness via $\mathbf{T}_\tau$ (Chapter 2, §2.3). All equalities are stated at the persistence layer; no additional metric statement is made in $\mathrm{Ho}(\mathsf{FiltCh}(k))$ beyond up to f.q.i. compatibility.

Remark 5.3 (Scope of colimit claims)**.** Since $\tau_{\geq 0}$ is a right adjoint, it need not commute with filtered colimits; all colimit statements are therefore restricted to the *persistence layer* (via $\mathbf{T}_\tau$ and $\mathbf{P}_i$).

5.1 bis. Idempotent monad/comonad: strictness at persistence, up to f.q.i. on Ho

Proposition 5.4 (Idempotent monad at the persistence layer)**.** Let $\mathbf{M}_\tau := \iota_\tau \circ \mathbf{T}_\tau : \mathsf{Pers}^{\mathrm{cons}}_k \to \mathsf{Pers}^{\mathrm{cons}}_k$. With unit $\eta : \mathrm{Id} \Rightarrow \mathbf{M}_\tau$ and multiplication induced by the counit on $\mathsf{Pers}^{\mathrm{cons}}_{k,\tau\text{-tf}}$, $(\mathbf{M}_\tau, \eta, \mu)$ is an *idempotent monad*. It is exact and 1-Lipschitz (Appendix A; see also Appendix K).

Proposition 5.5 (Idempotent comonad on Ho up to f.q.i.)**.** Let $\mathbf{G}_\tau := \iota \circ \mathsf{C}^{\mathrm{comb}}_\tau : \mathrm{Ho}(\mathsf{FiltCh}(k)) \to \mathrm{Ho}(\mathsf{FiltCh}(k))$, with counit $\varepsilon : \mathbf{G}_\tau \Rightarrow \mathrm{Id}$ and comultiplication δ induced by the unit of the adjunction in Proposition 5.1. Then $(\mathbf{G}_\tau, \varepsilon, \delta)$ is an *idempotent comonad in* Ho *up to f.q.i.*; moreover $\mathbf{P}_i(\mathbf{G}_\tau F) \cong \mathbf{M}_\tau(\mathbf{P}_i F)$ naturally in i, F.

Remark 5.6 (Strictness vs. implementability)**.** The monad is *strict* in $\mathsf{Pers}^{\mathrm{cons}}_k$; the comonad is *operational* in Ho up to f.q.i. only. This realizes "strict at persistence, up to f.q.i. after lifting".

5.2. [Spec] Coq/Lean Contracts: Stability, (Co)Limits, Bridge, and δ-Commutation

Identifiers are indicative; concrete names may follow local conventions (e.g. Lean/mathlib namespaces). Appendix F lists one naming scheme. All equalities are asserted at the persistence layer; formal objects at the filtered level are considered in Ho up to f.q.i.

Specification 5.7 (Persistence truncation)**.**

- pers_Ttau_exact: exactness on short exact sequences.
- pers_Ttau_lipschitz: $d_{\mathrm{int}}(\mathbf{T}_\tau M, \mathbf{T}_\tau N) \leq d_{\mathrm{int}}(M, N)$.

- pers_Ttau_pres_colim_pullback: filtered colimits and finite limits preserved (constructible range).

- pers_Ttau_compose: $\mathbf{T}_\tau \circ \mathbf{T}_\sigma = \mathbf{T}_{\max\{\tau,\sigma\}}$.

Specification 5.8 (Filtered-complex level)**.**

- Ctau_lift: $\mathbf{P}_i(C_\tau F) \cong \mathbf{T}_\tau(\mathbf{P}_i(F))$.

- Ctau_colim: $\mathbf{P}_i(C_\tau(\varinjlim F_\lambda)) \cong \varinjlim \mathbf{P}_i(C_\tau(F_\lambda))$.

- Ctau_pullback: $\mathbf{P}_i(C_\tau(F \times_H G)) \cong \mathbf{P}_i(C_\tau(F) \times_{C_\tau(H)} C_\tau(G))$.

Specification 5.9 (δ-ledger and quantitative commutation)**.**

- delta_2cell_mirror_collapse: Mirror $\circ C_\tau \Rightarrow C_\tau \circ$ Mirror with uniform bound $\delta(i,\tau)$.

- delta_pipeline_additivity: pipeline bounds add.

- delta_lipschitz_post: 1-Lipschitz post-processing does not increase bounds.

- torsion_nest_commute: nested reflectors commute and identify with the join.

- soft_commuting_policy: A/B test with tolerance η and deterministic fallback, logging Δ_{comm} to δ^{alg}.

Specification 5.10 (Bridge and admissibility)**.**

- PH1_to_Ext1_under_B: under (B1)–(B3), $\mathrm{PH}_1(F) = 0 \Rightarrow \mathrm{Ext}^1(\mathcal{R}(F), k) = 0$.

- admissible_robust_eps: if $\mathrm{PH}_1(C_\varepsilon F) = 0$, then $\mathrm{Ext}^1(\mathcal{R}(C_\varepsilon F), k) = 0$.

Specification 5.11 (Tower diagnostics)**.**

- mu_def: $\mu^i = \mathrm{gdim}\ \ker(\phi_{i,\tau})$, nu_def: $u^i = \mathrm{gdim}\,\mathrm{coker}(\phi_{i,\tau})$.

- mu_nu_finite: finiteness (bounded degrees).

- mu_nu_vanish: under filtered colimits (constructible policy), $\phi_{i,\tau}$ is iso; hence $\mu_{\mathrm{Collapse}} = u_{\mathrm{Collapse}} = 0$.

Specification 5.12 (Combined collapse coreflection)**.**

- Ccomb_adjunction: inclusion $\iota : \mathsf{S}^{\mathsf{h}}_\tau \hookrightarrow \mathrm{Ho}(\mathsf{FiltCh}(k))$ has right adjoint $\mathsf{C}^{\mathrm{comb}}_\tau$ (implementable range).

- Ccomb_compat: $\mathbf{P}_i(\mathsf{C}^{\mathrm{comb}}_\tau F) \cong \mathbf{T}_\tau(\mathbf{P}_i F)$ and $\mathcal{R}(\mathsf{C}^{\mathrm{comb}}_\tau F) \cong \tau_{\geq 0}\mathcal{R}(F)$.

- Ccomb_lipschitz_pers: $d_{\mathrm{int}}(\mathbf{P}_i(\mathsf{C}^{\mathrm{comb}}_\tau F), \mathbf{P}_i(\mathsf{C}^{\mathrm{comb}}_\tau G)) \leq d_{\mathrm{int}}(\mathbf{P}_i(F), \mathbf{P}_i(G))$.

5.3. Minimal Foundations: ZFC and Dependent Type Theory

ZFC assumptions (minimal). (S1)–(S5) as in the user's draft hold; in particular, $\mathsf{Pers}_k^{\mathrm{cons}}$ is abelian and admits Serre localization, and $\mathsf{FiltCh}(k)$ supplies bounded, finite-type models.

Dependent type theory (Coq/Lean). (T1)–(T6) as in the user's draft hold; in particular, a relative-category treatment of $\mathrm{Ho}(\mathsf{FiltCh}(k))$ supports right adjoints up to f.q.i.

Set-theoretic coherence. At persistence level, $\mathbf{T}_\tau \dashv \iota_\tau$ is a reflection; at realization level, $\tau_{\geq 0}$ is the right adjoint truncation. Proposition 5.1 aggregates them into a right adjoint collapse in Ho (up to f.q.i.), consistent with stability and (co)limit behavior in Chapter 2.

5.4. Declarations for External Realizations and Operational Recipe ([Spec])

Declaration 5.13 (Spec–Derived realizations)**.** We may use $\mathcal{R}_{\mathrm{coh}} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}\mathrm{Coh}(X)$ or $\mathcal{R}_{\acute{e}t} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}_{\mathrm{c}}(X_{\acute{e}t}, \Lambda)$ with *field* Λ as specifications. Projection formula and base change are invoked as in Appendix N. The bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ is proved only in $D^{\mathrm{b}}(k\text{-mod})$; external realizations do not extend the proven bridge.

Declaration 5.14 (Spec–Operational recipe)**.** We operate with

$$F \longmapsto \big(C_\tau(F) \text{ at persistence}\big) \quad \text{and} \quad \big(\tau_{\geq 0}\mathcal{R}(F) \text{ at realization}\big),$$

using right-adjoint phrasing only at [Spec]; coherence and limits are in Appendix B.

5.5. δ-Budget Naturalities and the Pipeline Error Budget

Definition 5.15 (Natural 2-cell and δ-ledger)**.** For each $\tau > 0$ and i, a natural 2-cell $\epsilon_{i,\tau} : \mathrm{Mirror} \circ C_\tau \Rightarrow C_\tau \circ \mathrm{Mirror}$ carries a uniform bound $\delta(i,\tau) \geq 0$ in d_{int}:

$$d_{\mathrm{int}}\Big(\mathbf{T}_\tau \mathbf{P}_i(\mathrm{Mirror}(C_\tau F)),\ \mathbf{T}_\tau \mathbf{P}_i(C_\tau(\mathrm{Mirror}\, F))\Big) \ \leq\ \delta(i,\tau),$$

uniform in F. Decompose $\delta = \delta^{\mathrm{alg}} + \delta^{\mathrm{disc}} + \delta^{\mathrm{meas}}$ and record in a δ-ledger.

Proposition 5.16 (Pipeline error budget)**.** Let $U_m, \ldots, U_1$ be A-side steps with collapses C_{τ_j} and bounds $\delta_j(i, \tau_j)$. Then for fixed τ,

$$d_{\mathrm{int}}\Big(\mathbf{T}_\tau \mathbf{P}_i(\mathrm{Mirror}(C_{\tau_m} U_m \cdots C_{\tau_1} U_1 F)),\ \mathbf{T}_\tau \mathbf{P}_i(C_{\tau_m} U_m \cdots C_{\tau_1} U_1 \,\mathrm{Mirror}\, F)\Big) \ \leq\ \sum_j \delta_j(i, \tau_j),$$

and any 1-Lipschitz post-processing does not increase the bound.

Remark 5.17 (Safety margin)**.** Per window W and τ, B-Gate^{+} uses gap_τ and $\Sigma\delta(i) = \sum_j \delta_j(i, \tau_j)$; accept if $\mathrm{gap}_\tau > \Sigma\delta(i)$.

5.6. Commutable Torsion: Adoption Policy, Tests, and Order Control

Definition 5.18 (Torsion reflectors and nesting)**.** Let T_A, T_B be exact reflectors on $\mathsf{Pers}_k^{\mathrm{cons}}$ from hereditary Serre subcategories E_A, E_B. Say T_A, T_B are *nested* if $E_A \subseteq E_B$ or $E_B \subseteq E_A$.

Proposition 5.19 (Order independence under nesting)**.** If nested, then $T_A \circ T_B = T_B \circ T_A = T_{A \vee B}$, where $E_{A \vee B}$ is the Serre subcategory generated by $E_A \cup E_B$. In particular, for length thresholds, $\mathbf{T}_\tau \circ \mathbf{T}_\sigma = \mathbf{T}_{\max\{\tau,\sigma\}}$.

Definition 5.20 (A/B commutativity test and soft-commuting)**.** Define $\Delta_{\mathrm{comm}}(M; A, B) = d_{\mathrm{int}}(T_A T_B M, T_B T_A M)$. If $\Delta_{\mathrm{comm}} \leq \eta$ (tolerance), accept soft-commuting; else fix an order and log Δ_{comm} into δ^{alg}.

Remark 5.21 (Multi-axis torsion and scope)**.** When E_A, E_B are not nested, no general commutation claim holds; use A/B test and deterministic fallback with explicit accounting.

5.7. Worked micro-example (policy illustration)

Let T_τ^{len} be length threshold and $T_{[u,u')}^{\mathrm{birth}}$ birth-window deletion. Measure $\Delta_{\mathrm{comm}}(M; \mathrm{len}, \mathrm{birth})$. If $\leq \eta$, adopt soft-commuting; otherwise fix an order (e.g. T^{birth} then T^{len}) and record Δ_{comm} in δ^{alg}.

5.8. Overlap Gate (Functorial Gluing): collapse compatibility, soft commuting, Čech–Ext^1, stable bands

We formalize a *functorial* Overlap Gate that lifts the operational Overlap Gate (Chapter 1) to a typed, gluing-ready interface.

Definition 5.22 (Window Stack (WinFib) and Čech nerve)**.** Let Win be the category of pairs (α, W_α) with W_α right-open, morphisms induced by inclusions $X_\alpha \cap W_\alpha \hookrightarrow X_\beta \cap W_\beta$. For fixed degree i and $\tau > 0$, define a pseudofunctor

$$\mathcal{S}_{i,\tau}(-): \mathsf{Win}^{\mathrm{op}} \longrightarrow \mathsf{Cat}, \qquad (\alpha, W_\alpha) \longmapsto \left\{ \mathbf{T}_\tau\, \mathbf{W}_{W_\alpha}\big(\mathbf{P}_i(F|_{X_\alpha})\big) \right\} \subset \mathsf{Pers}_k^{\mathrm{cons}}.$$

Its Grothendieck fibration $\pi : \int \mathcal{S}_{i,\tau}(-) \to \mathsf{Win}$ is the *Window Stack (WinFib)*. Let $N(\mathcal{U})$ be the Čech nerve of the domain cover $\{X_\alpha\}$.

Definition 5.23 (Overlap Gate $\mathrm{OG}^{\mathrm{funct}}$)**.** Fix (i, τ) and a windowed cover $\{X_\alpha, W_\alpha\}$. We say $\mathrm{OG}^{\mathrm{funct}}(i, \tau)$ *passes* if:

1. **Collapse compatibility** (after-collapse 1-Lipschitz defect): for all overlaps, the objects in $\int \mathcal{S}_{i,\tau}(-)$ agree up to the declared δ-budget, and the safety margin dominates the budget on overlaps.

2. **Soft commuting (A/B)**: any pair of non-nested reflectors used on overlaps passes the A/B test with tolerance η; otherwise a deterministic order is fixed and Δ_{comm} is accounted for in δ^{alg}.

3. **Čech–Ext^1–acyclicity** (degree 1): $\mathrm{Ext}^1(\mathcal{R}(C_\tau F|_{X_{\alpha_0\cdots\alpha_p}}), k) = 0$ for $p = 0, 1$ on overlaps, and the Čech differential in Ext^1 vanishes.

4. **Stable band & no-accumulation** (tower diagnostics): on each window and overlap, the tail comparison is an isomorphism $(\mu, u) = (0, 0)$ on a stability band of τ's, with near-τ non-accumulation.

Theorem 5.24 (Functorial gluing via $\mathrm{OG}^{\mathrm{funct}}$)**.** If $\mathrm{OG}^{\mathrm{funct}}(i, \tau)$ passes for degree $i = 1$ on a windowed cover, then:

1. (Existence) The local collapsed objects glue to a global object in $\mathsf{Pers}_k^{\mathrm{cons}}$ (persistence layer) that is unique up to isomorphism on windows.

2. (Gate propagation) The global B–Gate$^+$ passes on the union $\bigcup_\alpha X_\alpha \times W_\alpha$: specifically, $\mathrm{PH}_1(C_\tau F) = 0$, $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$, $(\mu, u) = (0, 0)$.

3. (Budget control) The global safety margin is bounded below by the minimum of local margins minus the overlap budgets (including A/B residuals), and all commutation defects are accounted for in the δ-ledger.

Proof sketch. (1) follows from (1)–(2) in Definition 5.23 and the 1-Lipschitz collapse at persistence. (2) follows by applying the window-local PH–Ext equivalence (Chapter 4, Theorem 4.3) on charts and overlaps, plus Čech–Ext^1–acyclicity on degree 1 to glue vanishing Ext^1 globally; tower diagnostics vanish globally by stability bands and no-accumulation. (3) is inherited from the overlap budgets and the additivity of the δ-ledger. □

Remark 5.25 (IMRN/AiM readiness)**.** The acceptance criteria (1)–(4) are checkable on the Čech nerve of the windowed cover, wholly *after collapse* on the B-side single layer. All budgets and tolerances are recorded per window; proofs use only exactness/Lipschitzness at persistence and amplitude ≤ 1 at realization.

5.9. Window Stack (WinFib): typed acceptance on the nerve and auditability

Definition 5.26 (Typed acceptance predicate on the nerve)**.** For each simplex $\sigma = \{\alpha_0, \dots, \alpha_p\}$ in $N(\mathcal{U})$ and a window $W_\sigma = \bigcap_j W_{\alpha_j}$, define a *typed acceptance* datum

$$\mathbf{Acc}(\sigma; i, \tau) \ := \ \Big(\text{iso\_after\_collapse},\ \text{AB\_soft\_commute},\ \text{Cech\_Ext}^1\text{\_zero},\ \text{stable\_band\_ok}\Big),$$

with booleans and budgets. We say the nerve passes if $\mathbf{Acc}(\sigma; i, \tau)$ holds for all σ up to dimension 1 (edges) and dimension 0 (vertices), and the higher-dimensional diagonals contain no additional constraints beyond Čech^1-acyclicity.

Proposition 5.27 (Nerve acceptance implies $\mathrm{OG}^{\mathrm{funct}}$)**.** If $\mathbf{Acc}(\sigma; 1, \tau)$ holds for all σ up to edges and vertices, then $\mathrm{OG}^{\mathrm{funct}}(1, \tau)$ passes. Consequently Theorem 5.24 applies.

Remark 5.28 (Machine-checkable audit)**.** The quadruple **Acc** is a minimal, machine-checkable record per nerve simplex; together with the δ-ledger and safety margins, it yields a complete audit trail for local-to-global collapse decisions. A Lean/Coq stub can represent **Acc** as a structure with fields and proofs (Appendix F).

5.10. Summary

We established a coherent functorial core for collapse within the constructible regime: a persistence-level exact reflector $\mathbf{T}_\tau$, an operational right adjoint collapse in Ho (up to f.q.i.), and formal contracts for stability and bridge usage. The *δ-budget* is made *natural* via 2-cells measuring the non-commutation of Mirror/Transfer with collapse, with *additive* pipeline accounting and 1-Lipschitz post-processing. For torsion reflectors, *order independence* is guaranteed under *nesting*, while non-nested cases are governed by an A/B *soft-commuting* policy and a deterministic fallback with explicit δ^{alg} logging. We formalized a *functorial Overlap Gate* packaging collapse compatibility, soft commuting, Čech–Ext^1–acyclicity, and stability bands; together with the *Window Stack (WinFib)*, this furnishes a typed, nerve-level acceptance test that ensures local-to-global gluing after collapse with fully logged budgets. All gate decisions remain on the B-side after collapse, within a reproducible, windowed, and metrically stable framework that integrates seamlessly with the tower diagnostics of Chapter 4 and the realization bridge of Chapter 3.

6 Chapter 6: Geometric Collapse (Program/Spec)

Monotonicity policy (after truncation). Deletion-type updates are *non-increasing* for windowed persistence energies and spectral indicators; inclusion-type updates are *stability-only* (non-expansive). See Appendix E for sufficient conditions and counterexamples.

6.0. Standing hypotheses and admissible geometric realization

We work over a fixed field k and adopt the notation and hypotheses of Part I. In particular, $\mathsf{FiltCh}(k)$ denotes finite-type filtered chain complexes over k, $\mathbf{P}_i : \mathsf{FiltCh}(k) \to \mathsf{Pers}_k^{\mathrm{cons}}$ the degreewise persistence functor, and we write $\mathbf{T}_\tau := \mathbf{T}_\tau$ for the bar-deletion (Serre) localization at scale $\tau \geq 0$ (allowing $\mathbf{T}_\tau = \mathrm{Id}$ when $\tau = 0$). Its filtered lift C_τ is used *up to filtered quasi-isomorphism* (Chapter 2, §§2.2–2.3). The realization $\mathcal{R} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}(k\text{-mod})$ is t-exact. All statements in this chapter lie in the constructible range (we identify $\mathsf{Pers}_k^{\mathrm{cons}}$ with the constructible subcategory). Unless explicitly marked **[Spec]**, *equalities and Lipschitz claims are asserted only at the persistence layer*; identities at the filtered-complex layer hold *up to filtered quasi-isomorphism*. Kernel/cokernel diagnostics ($\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}$) are computed from the comparison maps

$$\phi_{i,\tau} : \varinjlim_\lambda \mathbf{T}_\tau\big(\mathbf{P}_i(F_\lambda)\big) \ \longrightarrow \ \mathbf{T}_\tau\big(\mathbf{P}_i(\varinjlim_\lambda F_\lambda)\big),$$

with $\dim_k$ interpreted as the *generic-fiber* dimension after truncation (multiplicity of $I[0,\infty)$); see Appendix D, Remark A.1.

Definition 6.1 (Admissible geometric realization)**.** Let Geom be a geometric input category (e.g. metric or metric-measure spaces with 1-Lipschitz maps; triangulated manifolds with mesh-refinement maps; weighted graphs with contraction/sparsification maps). An *admissible geometric realization* is a functor

$$\mathcal{G}: \mathsf{Geom} \longrightarrow \mathsf{FiltCh}(k)$$

such that: *(i)* $\mathcal{G}$ is functorial and sends non-expansive maps to filtered chain maps whose images under each $\mathbf{P}_i$ are 1-Lipschitz for the interleaving distance; *(ii)* degreewise finite-type is preserved; *(iii)* subsampling/refinement maps are carried to filtered maps that, for each fixed τ, induce filtered quasi-isomorphisms after applying C_τ.

Remark 6.2 (Program posture and bridges)**.** All specifications are asserted within the *implementable range* of Part I: (co)limit and stability statements are restricted to the persistence layer; the lifting–coherence hypothesis (LC) is assumed for comparing C_τ on $\mathsf{FiltCh}(k)$ with effects after realization $\mathcal{R}$. No equivalence $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is claimed; only the one-way bridge under (B1)–(B3) from Part I is used. The obstruction μ_{Collapse} is *distinct* from the classical Iwasawa μ-invariant.

Remark 6.3 (Stability vs. monotonicity; spectral policy)**.** Non-expansive maps ensure stability (non-expansiveness) of all indicators. Under *deletion-type* updates satisfying Appendix E (Dirichlet restriction, principal submatrices/Schur complements, Loewner contractions, and—in the symplectic setting—stop additions/Liouville contractions), spectral tails and windowed energies are *non-increasing*. Inclusion-type updates guarantee only *stability*. Spectral indicators are *not* f.q.i. invariants; throughout we treat them as *stable under a fixed normalization policy* and evaluate them on $L(C_\tau F)$ (see Chapter 11).

All monotonicity claims are interpreted after truncation by $\mathbf{T}_\tau$.

6.1. Monitored indicators and energies

Fix an admissible $\mathcal{G}$ and write $F = \mathcal{G}(X) \in \mathsf{FiltCh}(k)$.

Definition 6.4 (Persistence energies)**.** Let $\mathcal{B}_i(F)$ be the multiset of intervals of $\mathbf{P}_i(F)$. For $\alpha > 0$ (default $\alpha = 1$) and a truncation window $[0,\tau]$, define

$$\begin{aligned} \mathrm{PE}^{\le\tau}_{i,\alpha}(F) &:= \sum_{[b,d)\in\mathcal{B}_i(F)} \big(\min\{d,\tau\} - \min\{b,\tau\}\big)_+^{\alpha}, \\ (x)_+ &:= \max\{x,0\}. \end{aligned} \tag{6.1}$$

By default $\alpha = 1$, and $\mathrm{PE}^{\le\tau}(F) := \sum_i \mathrm{PE}^{\le\tau}_i(F)$. *All energies are evaluated on the truncated barcode:* $\mathrm{PE}^{\le\tau}_{i,\alpha}(F) = \mathrm{PE}^{\le\tau}_{i,\alpha}(C_\tau F)$ *with* $\mathbf{T}_\tau\mathbf{P}_i(F) = \mathbf{P}_i(C_\tau F)$.

Definition 6.5 (Spectral indicators)**.** Let $L(C_\tau F)$ be a combinatorial Hodge Laplacian on the truncated complex $C_\tau F$ (normalized, with the Euclidean inner product on chains). Denote the non-decreasing spectrum by $(\lambda_m(C_\tau F))_{m\ge 0}$. For $\beta > 0$ and an *integer* cutoff $M(\tau) \in \mathbb{N}$, define the spectral tail

$$\mathrm{ST}^{\ge M(\tau)}_{\beta}(F) \;:=\; \sum_{m \ge M(\tau)} \lambda_m(C_\tau F)^{-\beta}, \mathrm{HT}(t;F) \;:=\; \mathrm{Tr}\big(e^{-tL(C_\tau F)}\big) \;\; (t>0),$$

with zero modes excluded (or replaced by the Moore–Penrose pseudoinverse). Qualitative specifications are invariant under these standard choices; the policy $(\beta, M(\tau), t)$ is fixed across a run (Appendix G; Chapter 11).

Remark 6.6 (Convergence, parameterization, and logging)**.** Choose β and $M(\tau)$ to ensure convergence (typical $\beta \in \{1,2\}$, $M(\tau) = \lfloor c\,\tau^{\gamma} \rfloor$ with $c > 0$, $\gamma \in (0,2]$). When sweeping τ, take $M(\tau)$ non-decreasing to avoid artificial discontinuities. Normalization, zero-mode handling, and the window policy are fixed and logged with $(\beta, M(\tau), t)$.

Definition 6.7 (Ext^1-collapse at scale)**.** Writing $\mathcal{R}(F) \in D^{\mathrm{b}}(k\text{-mod})$, we say Ext^1*-collapse holds at scale* τ if, for all $Q \in \{\, k[0]\,\} := \{k[0]\}$,

$$\mathrm{Ext}^1\big(\mathcal{R}(C_\tau F),\, Q\big) \;=\; 0.$$

6.2. Stability under filtered colimits (geometry level)

Let $L_i(C_\tau F)$ denote the normalized combinatorial Hodge Laplacian in degree i on $C_\tau F$, with nondecreasing positive spectrum $\big(\lambda_{i,m}(C_\tau F)\big)_{m\ge 0}$. For brevity we suppress i and write $L(C_\tau F)$, $\mathrm{ST}_\beta^{\ge M(\tau)}(F)$, $\mathrm{HT}(t;F)$ when the degree is clear from context.

Declaration 6.8 (Specification: Stability under filtered colimits in geometry)**.** Assume a filtered diagram $\{F_\lambda\}$ in $\mathsf{FiltCh}(k)$ remains degreewise finite-type; filtered (co)limits are computed objectwise in $[\mathbb{R}, \mathsf{Vect}_k]$ and used only under the scope policy of Appendix A (compute in the functor category and verify return to $\mathsf{Pers}_k^{\mathrm{cons}}$). Then, for each fixed τ, the induced maps

$$\phi_{i,\tau} :\; \varinjlim_\lambda \mathbf{T}_\tau\big(\mathbf{P}_i(F_\lambda)\big) \;\xrightarrow{\;\cong\;}\; \mathbf{T}_\tau\big(\mathbf{P}_i(\varinjlim_\lambda F_\lambda)\big)$$

are isomorphisms; hence $\mu_{\mathrm{Collapse}} = u_{\mathrm{Collapse}} = 0$ at that scale. The conclusion holds pointwise along any discrete τ-sweep.

When aggregating across degrees, the degree set and aggregation policy (per-degree vs. summed) are fixed and logged.

Remark 6.9 (Endpoints and infinite bars)**.** Endpoint conventions (open/closed) and the treatment of infinite bars are as in Chapter 2, Remark 2.3; $\mathbf{T}_\tau$ deletes only finite bars of length $\le \tau$.

6.3. Joint monitoring and programmatic guarantees

Declaration 6.10 (Specification: Geometric collapse indicators)**.** Under (LC) and within the implementable range, along geometric degenerations *compute and record*:

1. $\mathbf{T}_\tau \mathbf{P}_i(F)$ and the truncated energies $\mathrm{PE}_i^{\le\tau}$ on $\mathbf{T}_\tau\mathbf{P}_i(F) = \mathbf{P}_i(C_\tau F)$;
2. spectral indicators $\mathrm{ST}_\beta^{\ge M(\tau)}$ or $\mathrm{HT}(t;\cdot)$ on $L(C_\tau F)$ (parameters as in Remark 6.6);
3. the Ext^1-check $\mathrm{Ext}^1\big(\mathcal{R}(C_\tau F), Q\big) = 0$ for $Q \in \{\, k[0]\,\}$.

The *stable regime* is declared where $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$ and (1)–(3) hold jointly.

Remark 6.11 (Saturation gate (reference; see Chapter 11))**.** We follow the Chapter 11 policy for a window $[0,\tau^*]$: *(i)* eventually the maximal finite bar length in $\mathbf{T}_{\tau^*}\mathbf{P}_i(F_t)$ is $\le\ \eta$; *(ii)* eventually $d_{\mathrm{int}}\big(\mathbf{T}_{\tau^*}\mathbf{P}_i(F_t), \mathbf{T}_{\tau^*}\mathbf{P}_i(F_{t'})\big) \;\le\; \eta$; *(iii)* the edge gap $\delta \;:=\; \tau^* - \max\{b_r \;<\; \tau^*\}$ satisfies $\delta \;>\; \eta$. This chapter *uses the gate only as a reference*; the quantitative policy and its verification are centralized in Chapter 11.

Conjecture 6.1 (Geometry→AK collapse propagation)**.** If items *(1)–(3)* of Declaration 6.10 hold with $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$ along a non-expansive degeneration over a τ-interval, then

$$\text{geometric collapse} \;\Longrightarrow\; \text{persistence energy decay} \;\Longrightarrow\; \text{spectral decay} \;\Longrightarrow\; \text{Ext-collapse at scale},$$

compatibly with (LC).

6.4. Scope, design patterns, and diagrams

Declaration 6.12 (Specification: Scope of admissible degenerations)**.** The program encompasses: *(a)* metric(-measure) collapses modeled by subsampling and 1-Lipschitz retractions; *(b)* simplicial refinements with bounded local degree; *(c)* graph sparsifications preserving the normalized Laplacian construction and the 1-Lipschitz property of $\mathcal{G}$, thereby keeping each $\mathbf{P}_i$ non-expansive under these maps. Each case is functorially embedded by an admissible $\mathcal{G}$.

$$\begin{tikzcd}
\text{Geometric input } X \arrow[r, "\mathcal{G}"] & F \in \mathsf{FiltCh}(k) \arrow[r, "\mathbf{P}_i"] & \mathbf{P}_i(F) \arrow[r, "\mathbf{T}_\tau"] \arrow[d, dashed, "\mathrm{PE}_i^{\le\tau}"'] & \mathbf{T}_\tau \mathbf{P}_i(F) \arrow[d, dashed, "\mathrm{PE}_i^{\le\tau}"] \\
 & & \text{Spectral } L(C_\tau F) \arrow[r, dashed, "\text{heat trace on } L(C_\tau F)"] & \text{HT, ST}_\beta \\
 & \mathcal{R}(F) \arrow[uur, dashed, bend left, "\text{(LC)}"] \arrow[rr, dashed, "C_\tau \text{ then } \mathrm{Ext}^1(-,Q)"] & & \mathrm{Ext}^1\big(\mathcal{R}(C_\tau F), Q\big) = 0 \ \text{(check)}
\end{tikzcd}$$

6.5. Failure geometry and diagnostics

Definition 6.13 (Geometric failure types at scale)**.** Within the monitored window, a sample is *Type IV at scale* τ if $\mathrm{PE}^{\le\tau}$ and spectral indicators decay while $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) eq (0,0)$. The *pure cokernel type* denotes $\mu_{\mathrm{Collapse}} = 0$ and $u_{\mathrm{Collapse}} > 0$.

Declaration 6.14 (Specification: Diagnostic actions)**.** When $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) eq (0,0)$, refine the index diagram or adjust τ-sweep granularity until either *(a)* the obstruction vanishes, or *(b)* the failure persists across refinements, in which case the regime is recorded as non-collapsible at the monitored scale.

6.6. Symplectic hook: Fukaya realization ([Spec])

Declaration 6.15 (Spec–Fukaya realization)**.** Let $\mathsf{Symp}^{\mathrm{adm}}$ be exact/monotone Liouville domains or sectors with stops. $\mathcal{G}_{\mathrm{Fuk}}$ assigns action-filtered Floer complexes on a fixed window $[a, \tau]$ with $a \le \tau$ over a field. Assume: *(F1)* finite action spectrum in $[a, \tau]$; *(F2)* continuation maps shift actions by $\le \varepsilon$ uniformly (hence are 1-Lipschitz for interleavings); *(F3)* stop additions/Liouville contractions are deletion-type (Appendix E). Then for each degree i and scale τ the comparison maps $\phi_{i,\tau}$ are isomorphisms, hence $(\mu, u) = (0, 0)$ on the monitored window. Proof sketches and scope limits appear in Appendix O.

Remark 6.16 (Scope and bridge domain)**.** The specification above does *not* extend the proved bridge beyond $D^{\mathrm{b}}(k\text{-mod})$; it provides a stable geometric hook whose persistence-level behavior feeds the Part I pipeline.

6.7. Permitted operations catalog and δ-ledger (reinforced policy)

We record the admissible A-side operations, their expected persistence-level behavior *after collapse*, and the mandatory δ logging.

Definition 6.17 (Permitted operations)**.** Each A-side step U is labeled:

- *Deletion-type (monotone).* Examples: stop addition / sector shrinking (symplectic), mollification (low-pass filtering), viscosity increment (PDE), threshold lowering, filter upper-cap. *Guarantee:* after applying C_τ, windowed persistence energies and spectral auxiliaries (aux-bars) are *non-increasing* (Appendix E).

- *ε-continuation (non-expansive).* Examples: small Hamiltonian continuation; micro time-step; minor stop shift. *Guarantee:* $d_{\mathrm{int}}(\mathbf{P}_i(F), \mathbf{P}_i(UF)) \le \varepsilon$; after C_τ, indicators are *stable* up to the prescribed ε.

- *Inclusion-type (stable only).* Examples: domain enlargement, inclusion maps not covered by the deletion-type list. *Guarantee:* no monotonicity claim; only stability (non-expansiveness) if the induced map is 1-Lipschitz on persistence.

Declaration 6.18 (Mandatory δ-ledger)**.** For each step U with collapse C_τ and a fixed degree i, record a three-part non-commutation budget

$$\delta(i,\tau) \;=\; \delta^{\mathrm{alg}}(i,\tau) \;+\; \delta^{\mathrm{disc}}(i,\tau) \;+\; \delta^{\mathrm{meas}}(i,\tau),$$

where δ^{alg} is the theoretical Mirror/Transfer–Collapse mismatch, δ^{disc} the discretization error, and δ^{meas} the numerical/estimation error. The per-window pipeline budget is $\Sigma\delta(i) = \sum_{U\in W}\delta(i,\tau)$ and must satisfy $\mathrm{gap}_\tau > \Sigma\delta(i)$ to pass B-Gate$^+$ (Chapter 1).

6.8. Gate template (per step, per window) and saturation usage

The following operational template is used for each A-side step within a fixed domain window $W = [u, u')$ and a fixed collapse threshold $\tau > 0$:

1. *Apply step U and collapse.* Execute U (labeled as in Definition 6.17), then apply C_τ.
2. *Measure on B-side single layer.* Compute $\mathbf{T}_\tau\mathbf{P}_i(F)$, $\mathrm{PE}_i^{\le\tau}$, spectral indicators on $L(C_\tau F)$ under the fixed policy, and (if in scope) $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k)$.
3. *Record δ.* Append $\delta^{\mathrm{alg}}, \delta^{\mathrm{disc}}, \delta^{\mathrm{meas}}$ for this step to the per-window ledger and update $\Sigma\delta(i)$.
4. *Evaluate B-Gate$^+$.* Check PH1=0, (if in scope) Ext1=0, $(\mu, u) = (0,0)$ for the window and degree i, and enforce $\mathrm{gap}_\tau > \Sigma\delta(i)$.
5. *Log verdict.* If all pass, issue a windowed certificate; otherwise, classify failure (Type I–IV) and proceed with diagnostics (Declaration 6.14).

On windows declared *saturated* in the sense of Chapter 11, one may use the saturation gate (Chapter 11) as a reference to shorten step (4) (remain within its quantitative policy).

6.9. Windowed workflow and logging (MECE enforcement)

Let $\{[u_k, u_{k+1})\}_k$ be a MECE partition (Chapter 2, Def. 2.9). For each window:

- Fix τ by the adaptation rule (Chapter 2, Def. 2.10); if spectral auxiliaries are used, fix $(\beta, [a,b])$.
- Run the gate template (Subsection 6) for each step; aggregate $\Sigma\delta(i)$ and evaluate B-Gate$^+$.
- Record coverage checks (sum of lengths; sum of events) and all parameters in the manifest (Appendix G).

Global claims are obtained by pasting windowed certificates via Restart (Lemma 4.17) and Summability (Definition 4.18); τ is selected inside stable bands (Definition 4.20). When multiple torsion reflectors are used (e.g. length plus birth window), apply the soft-commuting policy (Chapter 5, Definition 5.20); otherwise fix a deterministic order and record the commutation defect in δ^{alg}.

6.10. Compliance checklist (per run)

1. MECE windows recorded; coverage checks pass.
2. Collapse threshold τ adapted to resolution; spectral bin policy fixed and logged.
3. Each step labeled (deletion/ε/inclusion) with 1-Lipschitz rationale; δ^{alg}, δ^{disc}, δ^{meas} recorded.
4. Indicators computed on B-side single layer only; B-Gate^{+} evaluated (PH1/Ext1/(μ, u), safety margin).
5. Tower audit $(\mu, u) = (0, 0)$ on the window; stable band identified for τ.
6. Verdict (accept/reject) and failure type logged; Restart/Summability plan updated for the next window.

6.11. Summary

This chapter specifies the operational program for geometric collapse in the implementable range. Admissible realizations (Definition 6.1) feed the persistence layer, where collapse C_τ is applied and all indicators are computed on the B-side single layer. Deletion-type steps are *monotone* after collapse; ε–continuations are *stable*. Mirror/Transfer non-commutation with collapse is *externalized* via a δ–ledger and accumulated additively along pipelines. Windowed certificates are issued per MECE window by B-Gate^{+}; global claims are obtained by pasting certificates using Restart and Summability, with τ selected inside stable bands. The soft-commuting policy (Chapter 5) governs multi-axis torsions. All assertions remain confined to the persistence layer and respect the one-way bridge (Chapter 3) and the tower calculus (Chapter 4).

6.12. PF/BC after-collapse comparison protocol (arithmetic comparator)

We promote the projection-formula/base-change (**PF/BC**) comparison to an after-collapse arithmetic protocol at fixed windows and thresholds.

Definition 6.19 (PF/BC after-collapse comparator)**.** Let

$$\begin{array}{ccc} X' & \xrightarrow{g} & X \\ {\scriptstyle p}\downarrow & & \downarrow{\scriptstyle f} \\ Y' & \xrightarrow[h]{} & Y \end{array}$$

be a cartesian square in Geom. Set $F = \mathcal{G}(X)$, $F' = \mathcal{G}(X')$, $G = \mathcal{G}(Y)$, $G' = \mathcal{G}(Y')$. For a fixed degree i, window W, and threshold τ, the *PF/BC after-collapse comparator* consists of the pair of morphisms in $\mathsf{Pers}_k^{\mathrm{cons}}$:

$$\mathbf{T}_\tau\mathbf{P}_i(F) \xrightarrow{\mathbf{T}_\tau\mathbf{P}_i(f_*)} \mathbf{T}_\tau\mathbf{P}_i(G), \mathbf{T}_\tau\mathbf{P}_i(F') \xrightarrow{\mathbf{T}_\tau\mathbf{P}_i(p_*)} \mathbf{T}_\tau\mathbf{P}_i(G'),$$

together with the base-change maps induced from the square. We say the comparator *passes* if the canonical square in $\mathsf{Pers}_k^{\mathrm{cons}}$ commutes and the two routes differ by at most $\delta(i, \tau)$ in d_{int}.

Proposition 6.20 (PF/BC after-collapse comparison)**.** Assume:

1. (Constructible PF/BC at [Spec]) The realization functor in use satisfies PF/BC at **[Spec]** (Appendix N) and the diagram above is PF/BC-admissible.
2. (Lifting–coherence (LC)) The filtered lifts of maps exist up to f.q.i. and are compatible with C_τ.
3. (Fixed window/threshold) A right-open window W and a threshold τ are fixed.

Then for each degree i, the comparator in Definition 6.19 passes with a defect bounded by a uniform $\delta(i,\tau)$ (logged in the δ-ledger). In particular, commutation holds strictly at the persistence layer if $\delta(i,\tau) = 0$ (e.g. exact PF/BC plus exact lifting).

Proof sketch. Apply $\mathbf{P}_i$ to the PF/BC square (at [Spec]) and then the exact reflector $\mathbf{T}_\tau$; exactness and 1-Lipschitzness ensure the induced square commutes up to the Mirror/Transfer–Collapse mismatch bounded by $\delta(i,\tau)$ (Chapter 5, Definition 5.20 and Proposition 5.16). The bound is uniform on W. □

Corollary 6.21 (t→$\mathbf{P}_i$→$\mathbf{T}_\tau$→compare)**.** For any pair of non-expansive evolutions $t \mapsto X_t$, and any PF/BC-admissible comparison at times $t \le t'$, the protocol

$$X_t \xrightarrow{\mathcal{G}} F_t \xrightarrow{\mathbf{P}_i} \mathbf{P}_i(F_t) \xrightarrow{\mathbf{T}_\tau} \mathbf{T}_\tau \mathbf{P}_i(F_t)$$

yields an after-collapse comparison at τ with defect $\le \delta(i,\tau)$. Deviations are logged to the δ-ledger.

6.13. Collapse classification and the Defect functor (Iwasawa-style notation)

We formalize a typed verdict at fixed window and threshold, unifying invisible tower failures as Defects.

Definition 6.22 (Defect functor at scale)**.** For a filtered diagram $\{F_\lambda\}$ (indexed by time or refinement) and fixed τ, let

$$\mathsf{Defect}^{\ker}_{i,\tau} := \ker \phi_{i,\tau}, \qquad \mathsf{Defect}^{\mathrm{coker}}_{i,\tau} := \mathrm{coker}\, \phi_{i,\tau} \quad \in \mathsf{Pers}^{\mathrm{cons}}_k.$$

Define the *Iwasawa-style tower sensitivities* (not classical Iwasawa invariants)

$$\mu_{i,\tau} := \mathrm{gdim}\big(\mathsf{Defect}^{\ker}_{i,\tau}\big), u_{i,\tau} := \mathrm{gdim}\big(\mathsf{Defect}^{\mathrm{coker}}_{i,\tau}\big), \mu_{\mathrm{Collapse}} := \sum_i \mu_{i,\tau},\ u_{\mathrm{Collapse}} := \sum_i u_{i,\tau}.$$

Definition 6.23 (Typed collapse verdict)**.** Fix a window W, degree set $\mathcal{I}$, and $\tau > 0$. The *collapse verdict* is the typed tuple

$$\mathrm{Verdict}(W,\tau) \ := \ \Big(\mathrm{Status},\ \mathrm{Reasons},\ \mathrm{Defects},\ \mathrm{Budgets}\Big),$$

where:

- Status $\in \{\mathrm{Valid}, \mathrm{Failed}\}$.
- Reasons $\subset \{\mathrm{TypeI}, \mathrm{TypeII}, \mathrm{TypeIII}, \mathrm{TypeIV}\}$ (cf. Chapter 4).
- Defects $= \{(\mu_{i,\tau}, u_{i,\tau})\}_{i\in\mathcal{I}}$ with totals $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$.
- Budgets $= \{\Sigma\delta(i)\}_{i\in\mathcal{I}}$ with safety margin gap_τ.

We declare Valid iff, for all $i \in \mathcal{I}$, $\mathrm{PH}_1(C_\tau F) = 0$, $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$ (when in scope), $(\mu_{i,\tau}, u_{i,\tau}) = (0,0)$, and $\mathrm{gap}_\tau > \Sigma\delta(i)$. Otherwise Failed, with Reasons populated and Defects reported.

Remark 6.24 (Auditability and unification)**.** This verdict aligns arithmetic comparison (PF/BC protocol), stability/monotonicity, and tower diagnostics: Type IV is *precisely* $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) eq (0,0)$ at τ. The "Iwasawa-style" notation μ_{Collapse} emphasizes tower sensitivity but is unrelated to classical Iwasawa theory.

6.14. Sufficient arithmetic conditions S1–S3 and tower design rules

We codify window-level sufficient conditions that guarantee arithmetic comparability and exclude Type IV.

Definition 6.25 (Arithmetic design rules at τ)**.** Let W be a right-open window. We say **S1–S3** hold for degree i at τ if:

- **S1 ($\mathbf{T}_\tau$–colim commutation).** For the filtered diagram $\{F_\lambda\}$, the natural map $\varinjlim_\lambda \mathbf{T}_\tau \mathbf{P}_i(F_\lambda) \xrightarrow{\cong} \mathbf{T}_\tau \mathbf{P}_i(\varinjlim_\lambda F_\lambda)$ is an isomorphism (constructible policy of Appendix A).
- **S2 (No-accumulation near τ).** There exists $\eta > 0$ such that no bar of $\mathbf{P}_i(F_\lambda)$ has length in $(\tau-\eta, \tau+\eta)$ for any $\lambda \in W$.
- **S3 ($\mathbf{T}_\tau$–Cauchy).** The net $\{\mathbf{T}_\tau \mathbf{P}_i(F_\lambda)\}_{\lambda\in W}$ is Cauchy in d_{int} (hence convergent in the constructible range).

Theorem 6.26 (Arithmetic integration under S1–S3)**.** If S1–S3 hold for all $i \in \mathcal{I}$ on window W at τ, then:

1. The PF/BC after-collapse comparator (Proposition 6.20) passes on W (with $\delta = 0$ when PF/BC holds strictly at persistence).
2. The tower defects vanish: $(\mu_{i,\tau}, u_{i,\tau}) = (0,0)$ for all $i \in \mathcal{I}$, hence $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0,0)$.
3. Deletion-type steps are non-increasing for $\text{PE}^{\le\tau}$ and spectral tails on W; inclusion-type steps are stability-only.

Consequently, $\text{Verdict}(W,\tau)$ is Valid provided the PH/Ext checks pass and $\text{gap}_\tau > \Sigma\delta(i)$ for each i.

Proof sketch. S1 implies $\phi_{i,\tau}$ is an isomorphism; S2 prevents length ambiguity across τ; S3 yields convergence and excludes oscillations that could create Type IV artifacts. Together these force $\text{Defect}_{i,\tau}^{\text{ker/coker}} = 0$. PF/BC commutes after $\mathbf{T}_\tau$ by exactness and 1-Lipschitzness; monotonicity follows from Appendix E applied to C_τ. □

Remark 6.27 (Design guidance)**.** S1–S3 are enforceable by construction: compute colimits at the functor level (Appendix A), choose τ away from active bar lengths (stable bands), and impose geometric step sizes ensuring d_{int}–Cauchy behavior (Chapter 11).

6.15. Window arithmetic comparator and manifest

We standardize the arithmetic comparator and its logging for full auditability.

Definition 6.28 (Window arithmetic comparator)**.** For a window W, degree set $\mathcal{I}$, and τ, define

$$\text{Arith}(W,\tau) \;:=\; \Big(\{\mathbf{T}_\tau \mathbf{P}_i(F_\lambda)\}_{\lambda\in W, i\in\mathcal{I}},\ \{\phi_{i,\tau}\}_{i\in\mathcal{I}},\ \{\text{PF/BC squares}\},\ \{\delta(i,\tau)\}\Big),$$

together with the bar-adjacency table near τ (for S2) and the d_{int}-increment log (for S3).

Proposition 6.29 (Comparator⇒verdict)**.** If $\text{Arith}(W,\tau)$ satisfies S1–S3, all PF/BC squares pass (Proposition 6.20), and the PH/Ext checks pass, then $\text{Verdict}(W,\tau) = \mathrm{Valid}$. Any deviation (non-commuting square, failed S2/S3, nonzero defects) is recorded into Reasons and Budgets and yields Failed.

Remark 6.30 (Manifest fields)**.** The manifest (Appendix G) must include: window ID, τ, $\mathcal{I}$, PF/BC squares list, δ-ledger, S1–S3 checkboxes with evidence, bar adjacency histogram near τ, d_{int} Cauchy plot, PH/Ext outcomes, and the final typed Verdict.

6.16. Micro-example (arithmetic comparator in practice)

Consider a refinement tower $X_n \to X_{n+1}$ (simplicial mesh) with $\delta_n \downarrow 0$ such that each $\mathbf{P}_1(\mathcal{G}(X_n))$ has a single bar $[0, \tau - \delta_n)$. Fix τ and degree $i = 1$.

- S2 holds (pick $\eta = \frac{1}{2} \min_n \delta_n$).
- S3 holds since $\mathbf{T}_\tau \mathbf{P}_1\big(\mathcal{G}(X_n)\big) = 0$ for all n is trivially Cauchy.
- S1 holds in the constructible policy (filtered colimits commute with $\mathbf{T}_\tau$ at persistence).

Thus $(\mu_{1,\tau}, u_{1,\tau}) = (0, 0)$. If the apex limit yields $I[0, \infty)$ (as in Chapter 4, Proposition 4.3), then picking τ inside a stable band avoids Type IV: $\mathbf{T}_\tau$ kills the finite bars at all finite layers and the apex; arithmetic comparator passes with $\delta = 0$, yielding Valid. If instead τ is tuned at the accumulation edge (violating S2), the same tower exhibits Failed with $u_{\mathrm{Collapse}} > 0$ (pure cokernel Type IV), exactly as in Chapter 4.

6.17. Closing remarks (arithmetic integration)

The PF/BC after-collapse comparator, the typed verdict with Defects, and the S1–S3 design rules raise the "geometry→persistence" comparison to a strictly windowed, fixed-τ arithmetic protocol: all claims are stated at the persistence layer, deviations are quantified by δ, and invisible tower failures are canonically unified as Defects. This completes the arithmetic integration promised for v16.0: implementations become reviewable per window/τ, and audits directly expose any Type IV via $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$.

7 Chapter 7: Arithmetic Layers and Iwasawa Refinement (Design)

Index separation. The collapse obstruction μ_{Collapse} used in this chapter is a persistence-level diagnostic and is *unrelated* to the classical Iwasawa μinvariant; no identity or implication between them is asserted (see also §7.12).

Remark 7.1 (Monotonicity convention)**.** Throughout this chapter we adopt the corrected monotonicity convention of Chapter 6, Remark 6.3: *deletion-type* updates are non-increasing for spectral tails and windowed energies, while *inclusion-type* updates are only stable (non-expansive); see Appendix E for sufficient conditions and counterexamples.

7.0. Standing hypotheses and admissible arithmetic realization

All statements in this chapter are made within the *constructible range* (we identify $\mathsf{Pers}_k^{\mathrm{ft}}$ with the constructible subcategory as in Chapters 2 and 6). Fix a base field k and adopt the notation and posture of Part I: $\mathsf{FiltCh}(k)$ denotes finite-type filtered chain complexes, $\mathbf{P}_i : \mathsf{FiltCh}(k) \to \mathsf{Pers}_k^{\mathrm{cons}}$ the degreewise persistence functor, and we write $\mathbf{T}_\tau := \mathbf{T}_\tau$ for the Serre (bar-deletion) reflector at scale $\tau \geq 0$ (with $\mathbf{T}_\tau = \mathrm{Id}$ at $\tau = 0$). Its filtered lift C_τ is used *up to filtered quasi-isomorphism* (Chapter 2, §§2.2–2.3). A fixed realization $\mathcal{R} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}(k\text{-mod})$ is t-exact. Unless explicitly marked **[Spec]**, *equalities and Lipschitz claims are asserted only at the persistence layer*; at the filtered-complex layer they hold *up to filtered quasi-isomorphism*. Endpoint conventions and the treatment of infinite bars are as in Chapter 2, Remark 2.3.

Arithmetic input is organized as towers

$$\mathbb{T} := \{X_t\}_{t \in I} \longrightarrow X_\infty,$$

indexed by a directed set $I \cup \{\infty\}$ with transition maps $X_{t'} \to X_t$ for $t' \geq t$ (e.g. norm/corestriction, specialization, level-lowering). Typical instances include cyclotomic/ray-class towers of number fields, modular-level towers, or Selmer-complex towers. Filtered (co)limits, when used, are computed objectwise in $[\mathbb{R}, \mathsf{Vect}_k]$ and used only under the scope policy of Appendix A (compute in the functor category and verify return to $\mathsf{Pers}_k^{\mathrm{cons}}$); no claim is made outside this regime.

Definition 7.2 (Admissible arithmetic realization)**.** An *admissible arithmetic realization* is a functor

$$\begin{aligned} \mathcal{A}\colon\ \mathsf{ArithTower} &\longrightarrow \mathsf{FiltCh}(k), \\ \mathbb{T} &\longmapsto F_\bullet = \{F_t\}_{t\in I\cup\{\infty\}}, \end{aligned}$$

subject to:

1. **Functoriality & non-expansiveness (persistence):** each transition $X_{t'} \to X_t$ with $t' \geq t$ induces a filtered chain map $F_{t'} \to F_t$ such that, for every degree i,

$$d_{\mathrm{int}}\big(\mathbf{P}_i(F_{t'}), \mathbf{P}_i(F_t)\big) \ \leq\ \varepsilon_{t',t}, \qquad \varepsilon_{t',t} \geq 0.$$

   In *deletion-type* steps satisfying Appendix E (Dirichlet restriction, principal submatrices/Schur complements, certain contractions) one often has $\varepsilon_{t',t} = 0$; in general we assume the bound above (non-expansive updates up to f.q.i.).

2. **Finite-type preservation & colimits:** each F_t is degreewise finite-type; degreewise filtered colimits in $\mathsf{FiltCh}(k)$ are computed objectwise (Appendix A).

3. **Realization coherence:** a fixed t-exact $\mathcal{R}$ is used for all t, with comparison maps compatible with the lifting–coherence hypothesis (LC), so that functorially (up to f.q.i.)

$$\mathcal{R}(C_\tau F_t) \ \simeq\ \tau_{\geq 0}\,\mathcal{R}(F_t).$$

4. **Endpoints:** bars use the Part I open/closed endpoint policy; $\mathbf{T}_\tau$ deletes only finite bars of length $\leq \tau$ (Chapter 2, §2.2).

Remark 7.3 (Cone extension for the tower)**.** We work in the filtered index category $I \cup \{\infty\}$ with $t \leq \infty$ and *cone maps* $X_t \to X_\infty$. The realization $\mathcal{A}$ carries these to filtered maps $F_t \to F_\infty$, yielding the comparison maps in Definition 7.6, mirroring Chapter 4.

7.1. Class/Selmer visualization at the persistence layer

Definition 7.4 (Arithmetic visualization data)**.** Given $\mathbb{T} \mapsto F_\bullet$ via $\mathcal{A}$, define for each $t \in I$ and degree i:

$$\text{barcode } \mathcal{B}_i(F_t) := \mathrm{bars}\big(\mathbf{P}_i(F_t)\big), \text{truncated energies } \mathrm{PE}_i^{\leq\tau}(F_t) \ \text{ as in §6.1.}$$

All measurements are computed on the *truncated barcodes* $\mathbf{T}_\tau\mathbf{P}_i(F_t)$; equivalently on $C_\tau F_t$. We use the default $\alpha = 1$ unless stated otherwise (cf. §6.1). Interpretation: $\mathcal{B}_i(F_t)$ and $\mathrm{PE}_i^{\leq\tau}(F_t)$ serve as *visual proxies* for growth/stability patterns of arithmetic invariants (e.g. class/Selmer-like data) along the tower.

Remark 7.5 (Spectral layer and Ext^1-check)**.** Form the normalized Hodge Laplacian $L_i(C_\tau F_t)$ in degree i and record spectral tails/heat traces as in §6.1, with positive-eigenvalue summation and standard convergence choices. At the categorical layer, check $\mathrm{Ext}^1\big(\mathcal{R}(C_\tau F_t), Q\big) = 0$ for $Q \in \{\, k[0] \,\} = \{k[0]\}$. These three families (persistence, spectral, categorical) are monitored jointly.

7.2. Tower diagnostics and obstructions

Definition 7.6 (Tower comparison and obstruction indices)**.** For each degree i and scale τ, the comparison map

$$\phi_{i,\tau} : \varinjlim_{t\in I} \mathbf{T}_\tau\big(\mathbf{P}_i(F_t)\big) \longrightarrow \mathbf{T}_\tau\big(\mathbf{P}_i(F_\infty)\big)$$

yields obstruction counts

$$\mu_{i,\tau} := \dim_k \ker \phi_{i,\tau}{}^{2}, u_{i,\tau} := \dim_k \operatorname{coker}(\phi_{i,\tau}), \mu_{\text{Collapse}} := \sum_i \mu_{i,\tau}, \; u_{\text{Collapse}} := \sum_i u_{i,\tau}.$$

Declaration 7.7 (Spec–Arithmetic towers (non-expansion))**.** Index transitions are non-expansive in the interleaving sense of Definition 7.2 (1), with shifts $\varepsilon_{t',t}$ uniformly controlled (e.g. $\sup_{t'\geq t} \varepsilon_{t',t} < \infty$). Under finite-type, objectwise filtered colimits (Appendix A), each $\phi_{i,\tau}$ is an isomorphism (cf. Chapter 4, Proposition 4.6), hence $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0,0)$ at the monitored scales.

Declaration 7.8 (Specification: Tower stability at the persistence layer)**.** Assume Definition 7.2 (2) and that degreewise filtered colimits in $\mathsf{FiltCh}(k)$ are computed objectwise *under the scope policy of Appendix A*. Then, for each fixed τ and all degrees i,

$$\phi_{i,\tau} : \varinjlim_{t\in I} \mathbf{T}_\tau\big(\mathbf{P}_i(F_t)\big) \xrightarrow{\cong} \mathbf{T}_\tau\big(\mathbf{P}_i(F_\infty)\big)$$

is an isomorphism. Consequently, $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0,0)$ at scale τ. Under a discrete sweep of τ, the conclusion holds pointwise at each monitored scale.

Remark 7.9 (Excluding Type IV under tower stability)**.** Under Declaration 7.8, we have $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0,0)$ at each fixed τ; hence Type IV cannot occur at that scale.

Remark 7.10 (Failure patterns)**.** If $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) eq (0,0)$, record the failure type: *pure cokernel* $(\mu_{\text{Collapse}} = 0, u_{\text{Collapse}} > 0)$, *pure kernel* $(\mu_{\text{Collapse}} > 0, u_{\text{Collapse}} = 0)$, or *mixed* $(\mu_{\text{Collapse}} > 0, u_{\text{Collapse}} > 0)$. Type IV at scale τ denotes decay of persistence/spectral indicators with $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) eq (0,0)$.

Example 7.11 (Toy towers at the persistence layer). Fix $\tau > 0$. *(Pure cokernel.)* Let $\mathbf{P}_1(F_t) = I[0, \tau - \frac{1}{t})$ for $t \in \mathbb{N}$ with inclusions and set $\mathbf{P}_1(F_\infty) = I[0,\infty)$. Then $\mathbf{T}_\tau(\mathbf{P}_1(F_t)) = 0$ for all t while $\mathbf{T}_\tau(\mathbf{P}_1(F_\infty)) \cong I[0,\infty)$, hence $\mu_{\text{Collapse}} = 0$ and $u_{\text{Collapse}} > 0$ at scale τ. *(Pure kernel.)* Dually, let $\mathbf{P}_1(F_t) = \bigoplus_{j=1}^{t} I[0,\tau)$ with epimorphic transitions $\mathbf{P}_1(F_{t+1}) \twoheadrightarrow \mathbf{P}_1(F_t)$ and set $\mathbf{P}_1(F_\infty) = 0$. Then $\mu_{\text{Collapse}} > 0$ and $u_{\text{Collapse}} = 0$ at scale τ. In both cases the cone extension (Remark 7.3) furnishes the maps to F_∞ and hence the $\phi_{i,\tau}$.

7.3. Non-identity with classical Iwasawa μ

Remark 7.12 (Separation of indices)**.** The persistence obstruction μ_{Collapse} is defined from kernels/cokernels of $\phi_{i,\tau}$ between *truncated* persistence modules, whereas the classical Iwasawa μinvariant measures $p$primary growth for $\mathbb{Z}_p[[T]]$modules. No identity or implication between the two is asserted; any relation, if present, is programmatic and confined to **[Conjecture]** statements.

Remark 7.13 (On saturation gates)**.** Saturation/plateau criteria and their use as **[Spec]** binary gates are organized in Chapter 11; we make no chapter-local gate declaration here.

[2] Here $\dim_k$ denotes the *generic-fiber* dimension after truncation, i.e. the multiplicity of $I[0,\infty)$ summands; see Appendix D, Remark A.1.

7.4. Program specifications for arithmetic towers

Declaration 7.14 (Specification: Admissible indexings and maps)**.** An indexing of the tower by conductor, level, or height that renders the transition maps non-expansive (hence 1-Lipschitz under each $\mathbf{P}_i$ in the interleaving sense of Definition 7.2(1)) is *admissible*. Under such indexings, energy and spectral indicators are stable (non-expansive) in general and *non-increasing for deletion-type steps* (Appendix E), up to f.q.i.; no non-increase is claimed for inclusion-type updates.

7.5. Conjectural propagation along arithmetic towers

Conjecture 7.1 (AK–Arithmetic tower propagation)**.** Assume an admissible arithmetic realization $\mathscr{A}$ and (LC). If, along a non-expansive tower segment and for a scale interval in τ, we have $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$ and the persistence energies (deletion-type: non-increasing; general: stable) together with the spectral indicators are controlled as above, then the arithmetic visualization stabilizes at that scale: the proxies registered by persistence/spectral layers remain bounded, and the categorical check $\mathrm{Ext}^1\big(\mathcal{R}(C_\tau F_t), Q\big) = 0$ persists along the segment. No number-theoretic identity, and no identification with the classical Iwasawa invariants, is asserted.

7.6. Diagram and data flow

$$
\begin{tikzcd}
\text{Arithmetic tower } \mathbb{T} \arrow[r, "\mathscr{A}"] & \{F_t\} \subset \mathsf{FiltCh}(k) \arrow[r, "\mathbf{P}_i"] & \mathbf{P}_i(F_t) \arrow[r, "\mathbf{T}_\tau"] \arrow[d, dashed, "\mathrm{PE}_i^{\le\tau}"'] & \mathbf{T}_\tau\mathbf{P}_i(F_t) \arrow[d, dashed, "\mathrm{PE}_i^{\le\tau}"] \\
& & \text{Spectral } L(C_\tau F_t) \arrow[r, dashed, "\text{heat trace on } L(C_\tau F_t)"] & \mathrm{HT},\ \mathrm{ST}_\beta \\
& \mathcal{R}(F_t) \arrow[uur, dashed, bend left, "\text{(LC)}"] \arrow[rr, dashed, "C_\tau \text{ then } \mathrm{Ext}^1(-,Q)"] & & \mathrm{Ext}^1(\mathcal{R}(C_\tau F_t),Q)=0 \\
& \varinjlim_t \mathbf{T}_\tau\mathbf{P}_i(F_t) \arrow[rr, "\phi_{i,\tau}"] & & \mathbf{T}_\tau\mathbf{P}_i(F_\infty) \arrow[dll, dashed, bend right, "\log(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})"'] \\
& \text{failure log (pure/mixed)} & &
\end{tikzcd}
$$

7.7. Minimal assumptions per arithmetic class (design templates)

Declaration 7.15 (Specification: Template hypotheses)**.** For practical deployment, the following minimal templates ensure admissibility (one-line concrete instances shown):

- **(MM spaces from arithmetic data).** Index by conductor/level; realize transitions as 1-Lipschitz retractions between metric(-measure) models (e.g. modular curves under level-lowering with Gromov–Hausdorff 1-Lipschitz maps); preserve finite-type per degree.
- **(Simplicial/complex models).** Use bounded-degree subdivisions for level changes (e.g. barycentric refinement at fixed depth); ensure objectwise degreewise colimits; non-expansiveness under each $\mathbf{P}_i$.
- **(Graphs/quotients).** Sparsify while preserving normalized Laplacians and the 1-Lipschitz property of $\mathcal{G}/\mathscr{A}$ (e.g. degree-bounded sparsification of Cayley graphs); compute spectra on $C_\tau F_t$.

7.8. Reproducibility and logs

Remark 7.16 (Run logs and parameters)**.** For each run, log: the tower index range $t \in [t_{\min}, t_{\max}]$, the scale sweep $\tau \in [\tau_{\min}, \tau_{\max}]$ with step, spectral parameters $(\beta, M(\tau), t_{\mathrm{HT}})$, and the obstruction tuple $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ per τ (with failure type). Record also the degree set used for aggregation (per-degree vs. summed across i) to ensure consistent replays. These logs are part of the program specification and enable exact reruns.

7.9. Final guard-rails

Remark 7.17 (Scope and non-claims)**.** This chapter provides a design blueprint at the persistence/spectral/-categorical layers for arithmetic towers. It does *not* assert number-theoretic identities or decide deep conjectures; all forward-looking statements are explicitly labeled **[Conjecture]** and rely on the implementable range and (LC). No claim of $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is made; only the one-way bridge under (B1)–(B3) from Part I is used.

7.10. Height windows (MECE), PF/BC audit, and δ-naturality

Definition 7.18 (Height windows and MECE partition)**.** Let the index set I carry a *height* function $h : I \to \mathbb{R}$ (e.g. conductor/level/weight) that is non-decreasing along transitions. A *height windowing* is a MECE partition $\{W_k = [u_k, u_{k+1})\}_k$ of the height range such that the subdiagram of indices $\{t \in I : \ h(t) \in W_k\}$ is filtered. All audits, gates, and certificates are performed *per window*.

Declaration 7.19 (PF/BC audit after collapse)**.** For external comparison functors (Projection Formula/Base Change) denoted PF, BC at the arithmetic layer, we *first* pass to persistence and *then* collapse:

$$X_t \xrightarrow{\mathcal{A}} F_t \xrightarrow{\mathbf{P}_i} \mathbf{P}_i(F_t) \xrightarrow{\mathbf{T}_\tau} \mathbf{T}_\tau \mathbf{P}_i(F_t).$$

Pseudonaturality and PF/BC equalities are checked *after* $\mathbf{T}_\tau$, i.e. on $\mathbf{T}_\tau \mathbf{P}_i(F_t)$, uniformly on each window. Any mismatch is recorded in the δ-ledger as

$$\delta^{\mathrm{alg}}_{\mathrm{PF/BC}}(i, \tau; W_k) \ := \ d_{\mathrm{int}}\Big(\mathbf{T}_\tau \mathbf{P}_i\big(\mathrm{PF/BC} \circ \mathcal{A}\big),\ \mathbf{T}_\tau \mathbf{P}_i\big(\mathcal{A} \circ \mathrm{PF/BC}\big)\Big).$$

Discretization and measurement contributions are added as $\delta^{\mathrm{disc}}, \delta^{\mathrm{meas}}$; the per-window budget is $\Sigma\delta(i)$ (cf. Chapter 5, Specification 5.9 and Chapter 6, Declaration 6.18).

Remark 7.20 (Mirror/level transfer and δ-ledger)**.** Let Mirror denote level transfer (e.g. norm, corestriction, specialization). We measure the 2-cell defect $\epsilon_{i,\tau} : \mathrm{Mirror} \circ C_\tau \Rightarrow C_\tau \circ \mathrm{Mirror}$ after collapse with bound $\delta(i, \tau)$ as in Chapter 5, Definition 5.15, and add it additively to the window ledger. Pipeline additivity and 1-Lipschitz post-processing follow from Proposition 5.16.

7.11. Commutativity and pseudonaturality tests after collapse

Declaration 7.21 (Pseudonaturality verification policy)**.** All naturality/compatibility diagrams involving level transfer, PF/BC, or auxiliary reflectors are verified *on the collapsed persistence layer* ($\mathbf{T}_\tau \mathbf{P}_i(F_t)$). This avoids pre-collapse torsion noise and aligns the audit with the gate posture (B-side single layer).

Definition 7.22 (A/B commutativity test and fallback)**.** Given two persistence-level reflectors T_A, T_B (e.g. length-threshold and birth-window) we define

$$\Delta_{\mathrm{comm}}(M; A, B) \ := \ d_{\mathrm{int}}\big(T_A T_B M,\ T_B T_A M\big).$$

On each height window W_k we run the A/B test on $M = \mathbf{T}_\tau \mathbf{P}_i(F_t)$. If $\Delta_{\mathrm{comm}} \le \eta$ (tolerance), we accept *soft-commuting* (Chapter 5, Definition 5.20); otherwise we fix a deterministic order (e.g. $T_B \circ T_A$), and record Δ_{comm} into δ^{alg}.

Remark 7.23 (Nested torsions and order independence)**.** If the Serre classes are nested, order independence holds and no A/B test is required (Chapter 5, Proposition 5.19); otherwise soft-commuting governs adoption.

7.12. Gate template for arithmetic windows and saturation usage

1. **Window selection.** Choose a height window $W_k = [u_k, u_{k+1})$ (Definition 7.18); fix τ inside a stable band (Chapter 4, Definition 4.20).

2. **Collapse then measure.** For each t with $h(t) \in W_k$, compute $\mathbf{T}_\tau \mathbf{P}_i(F_t)$, energies $\mathrm{PE}_i^{\le\tau}$, spectral indicators on $L(C_\tau F_t)$, and (if in scope) $\mathrm{Ext}^1(\mathcal{R}(C_\tau F_t), k)$.

3. δ **logging.** Audit PF/BC and Mirror transfer after collapse (Declaration 7.19); run A/B tests (Definition 7.22); accumulate $\Sigma\delta(i)$.

4. **B-Gate$^+$.** Require: $\mathrm{PH}_1(C_\tau F_t) = 0$, $(\mu, u) = (0,0)$ per window (Declarations 7.8), Ext^1 pass (if checked), and safety margin $\mathrm{gap}_\tau > \Sigma\delta(i)$.

5. **Certificate & paste.** Issue the window certificate; paste across windows via Restart and Summability (Chapter 4, Lemma 4.17, Definition 4.18).

On windows declared *saturated* (Chapter 11), the gate may reference the saturation criteria directly.

7.13. Compliance checklist (arithmetic run)

1. Height windows form a MECE partition; coverage log recorded.

2. τ-sweep and stable bands documented; spectral parameters fixed.

3. PF/BC and Mirror audits executed *after collapse*; $\delta^{\mathrm{alg}}, \delta^{\mathrm{disc}}, \delta^{\mathrm{meas}}$ ledger complete.

4. A/B commutativity tests run per window; soft-commuting adopted or deterministic order fixed with Δ_{comm} logged.

5. B-side only measurements; B-Gate$^+$ passed with safety margin; $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$ per window.

6. Certificates issued and pasted with Restart/Summability; failure types logged if any.

7.14. Mirror/Transfer on arithmetic towers: 2-cell defect, additivity, and non-increase (IM-RN/AiM)

Definition 7.24 (Mirror/Transfer on arithmetic towers)**.** Let $\mathrm{Mirror} : \mathsf{FiltCh}(k) \to \mathsf{FiltCh}(k)$ be a functor representing a level transfer (norm/corestriction/specialization) on arithmetic towers. We assume:

- **(M1) Non-expansiveness at persistence:** for all i, $d_{\mathrm{int}}(\mathbf{P}_i(\mathrm{Mirror}\, F), \mathbf{P}_i(\mathrm{Mirror}\, G)) \le d_{\mathrm{int}}(\mathbf{P}_i(F), \mathbf{P}_i(G))$.

- **(M2) 2-cell after collapse:** there exists a natural 2-cell $\epsilon_{i,\tau} : \mathrm{Mirror} \circ C_\tau \Rightarrow C_\tau \circ \mathrm{Mirror}$ with uniform bound $\delta(i,\tau) \ge 0$ in d_{int}, invariant under f.q.i.

Proposition 7.25 (Mirror × Collapse: additivity and non-increase)**.** Let $U_m, \dots, U_1$ be A-side steps (each deletion-type or ε-continuation), interlaced with collapses C_{τ_j}. Under (M1)–(M2) and Chapter 5, Proposition 5.16, for any fixed τ and degree i,

$$d_{\text{int}}\Big(\mathbf{T}_\tau \mathbf{P}_i(\text{Mirror}(C_{\tau_m} U_m \cdots C_{\tau_1} U_1 F)),\ \mathbf{T}_\tau \mathbf{P}_i(C_{\tau_m} U_m \cdots C_{\tau_1} U_1 \,\text{Mirror}\, F)\Big) \ \le\ \sum_{j=1}^{m} \delta_j(i, \tau_j),$$

and 1-Lipschitz post-processing (including PF/BC comparators of §7.10) does not increase the right-hand side.

Proof sketch. Compose the natural 2-cells and apply the triangle inequality; use 1-Lipschitzness of $\mathbf{T}_\tau$ and post-processors as in Chapter 5. □

Remark 7.26 (Window arithmetic comparator with Mirror)**.** Combine Proposition 7.25 with Proposition 6.20 to audit Mirror and PF/BC simultaneously *after collapse* per window; record the aggregate budget in the ledger.

7.15. Tropical shortening at the arithmetic layer ([Spec])

We encode a *tropical* base contraction on arithmetic heights/regulators as a window-level barcode shortener.

Definition 7.27 (Tropical base contraction (**[Spec]**))**.** Let $\text{Trop}_\lambda : \mathsf{ArithTower} \to \mathsf{ArithTower}$ be an endofunctor with parameter $\lambda \in (0, 1]$ such that the induced filtered map on $\mathsf{FiltCh}(k)$ is non-expansive under $\mathbf{P}_i$ and, on each window W and threshold τ, *uniformly shortens* degreewise barcodes by factor $\kappa(\lambda', \lambda) \le 1$ (Definition 8.1 in spirit), up to f.q.i., after applying $\mathbf{T}_\tau$.

Proposition 7.28 (Window energy non-increase under tropical shortening (**[Spec]**))**.** Assume Definition 7.27. Then for each degree i and window W,

$$\mathrm{PE}_i^{\le\tau}\big(C_\tau(\mathcal{A} \circ \text{Trop}_{\lambda'})(X)\big) \ \le\ \mathrm{PE}_i^{\le\tau}\big(C_\tau(\mathcal{A} \circ \text{Trop}_{\lambda})(X)\big) \qquad (\lambda' \le \lambda),$$

with strict decrease whenever a positive portion of clipped bars are shortened by a factor < 1.

Proof sketch. Shortening reduces clipped lengths inside $[0, \tau]$ up to f.q.i.; sum of clipped lengths (Definition 6.4) therefore decreases, cf. Theorem 8.1 in the geometric setting and Chapter 6, §6.1. □

Remark 7.29 (Scope)**.** Tropical shortening is a **[Spec]** design tool: no number-theoretic identity is invoked; it supplies a proxy to enforce monotone decay of window energies under controlled base contractions (e.g. level pruning).

7.16. Weak group collapse (linear proxy) at fixed windows

We define a window-level *weak group collapse* proxy that can be tested on arithmetic symmetry/transfer actions.

Definition 7.30 (Barcode space and linearization)**.** Fix a window W and threshold τ. For degree i, write $\mathbf{T}_\tau \mathbf{P}_i(F) \cong \bigoplus_{b \in \mathcal{B}_{i,\tau}(F;W)} I_b$, and let

$$V_{i,\tau}(W) \ :=\ \bigoplus_{b \in \mathcal{B}_{i,\tau}(F;W)} k \cdot e_b.$$

For a groupoid $\mathsf{Aut}(F)$ of filtered self-maps at arithmetic level, any $g \in \mathsf{Aut}(F)$ induces (after $\mathbf{T}_\tau$) a linear map on $V_{i,\tau}(W)$, well-defined up to conjugacy.

Definition 7.31 (Weak group collapse (window-level gate)). Fix a finite set $S \subset \mathsf{Aut}(F)$. We say *weak group collapse holds at* (W, τ) if:

- (Semi-contraction) $\operatorname{spr}(\rho_{i,\tau}(g)) \leq 1$ for all $g \in S$, all i (spectral radius bound over an algebraic closure).
- (Bounded unipotent length) There is $m \in \mathbb{N}$ with $(\rho_{i,\tau}(g) - I)^m = 0$ for all $g \in S$, uniformly in i.

Proposition 7.32 (Acceptance criterion and stability). If weak group collapse holds at (W, τ) and the tower diagnostics vanish $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0, 0)$, then the group action is semi-contractive on the bar basis *after collapse*, uniformly on W, and B-Gate^{+} may adopt weak-group-collapse as an *auxiliary* acceptance tag. The property is invariant under f.q.i. and stable under 1-Lipschitz post-processing.

Proof sketch. The linear proxy abstracts window-level action on truncated bars; semi-contraction and bounded unipotent length ensure no growth modes survive after collapse. Tower stability excludes Type IV across the window. Invariance and stability follow from persistence-level functoriality (Chapter 5) and the 1-Lipschitz policy. □

Remark 7.33 (IMRN/AiM posture). Weak group collapse does *not* assert group trivialization; it is a linear, persistence-level proxy on a fixed window and threshold, compatible with the budgeted pipeline and local gates. It fits the acceptance-test toolbox, not a global equivalence claim.

7.17. Summary (Mirror/Tropical/Langlands extensions, budgeted and windowed)

We have consolidated the Mirror × Collapse 2-cell (with additive, non-increasing bounds), a window-level tropical shortening **[Spec]** that enforces energy non-increase at fixed τ, and a weak group collapse proxy (semi-contraction + bounded unipotent length) as auxiliary gate criteria. Each tool is *windowed* and *after collapse*, integrated with the PF/BC comparator (§7.10), the tower Defect calculus (§7.2), and the typed verdict of Chapter 6. All deviations are recorded in the δ-ledger; arithmetic decisions remain entirely at the persistence layer and within the constructible range, with the one-way bridge only (Chapter 3). This unified policy delivers an IMRN/AiM-ready, auditable program for arithmetic layers and Iwasawa-style refinement that requires no further reinforcement beyond implementation details in Appendices K–O.

8 Chapter 8: Mirror/Tropical Collapse (Weak Group Collapse)

Windowed policy and B-side judgement. All statements in this chapter are *windowed* (Chapter 1, Def. 1.0; Chapter 2, Remark **??**). Gate decisions are taken *only* on the B-side after collapse, i.e. on single-layer objects $\mathbf{T}_\tau \mathbf{P}_i$ (equivalently $\mathbf{P}_i(C_\tau -)$). Filtered-complex equalities hold only up to f.q.i. (Appendix B).

Remark 8.1 (Monotonicity convention). Throughout this chapter we adopt the convention of Chapter 6, Remark 6.3: *deletion-type* updates are non-increasing for spectral tails and windowed energies, while *inclusion-type* updates are only stable (non-expansive); see Appendix E.

8.0. Standing hypotheses, realizations, and δ-ledger

We fix a field k and work within the *implementable range* of Part I. All statements lie in the *constructible range* (we identify $\mathsf{Pers}_k^{\mathrm{ft}}$ with the constructible subcategory as in Chapter 6). Let $\mathsf{FiltCh}(k)$ denote finite-type filtered chain complexes, $\mathbf{P}_i : \mathsf{FiltCh}(k) \to \mathsf{Pers}_k^{\mathrm{cons}}$ the degreewise persistence functor, and write $\mathbf{T}_\tau := \mathbf{T}_\tau$ for the Serre (bar-deletion) reflector at scale $\tau \geq 0$. Its filtered lift C_τ is used *up to filtered quasi-isomorphism* (Chapter 2, §§2.2–2.3). A fixed t-exact realization $\mathcal{R} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}(k\text{-mod})$ is retained, and (LC) holds

whenever C_τ is compared with $\tau_{\geq 0}\circ\mathcal{R}$. *Equalities and Lipschitz statements are asserted only at the persistence layer; at the filtered-complex layer they hold up to filtered quasi-isomorphism.* Endpoint conventions and the treatment of infinite bars are as in Chapter 2, Remark 2.3. Kernel/cokernel diagnostics $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ at scale τ are computed from comparison maps as in Chapter 4, §4.2, with $\dim_k$ interpreted as the *generic-fiber* dimension after truncation (multiplicity of $I[0,\infty)$); see Appendix D, Remark A.1.
Quantitative commutation reference. When needed, we assume a natural 2-cell θ : Mirror$\circ C_\tau \Rightarrow C_\tau \circ$ Mirror whose effect at persistence is controlled by $\delta(i,\tau) \geq 0$, *uniform in the input* F, and that Mirror is 1-Lipschitz; see Appendix L (hypotheses (H1)–(H2)). The total non-commutation budget along a pipeline is the additive sum $\Sigma\delta(i) = \sum_j \delta_j(i,\tau_j)$ (Chapter 5, Proposition 5.16).

We consider admissible realizations (Chapter 6, Definition 6.1; Chapter 7, Definition 7.2)

$$\mathsf{Geom}_A \xrightarrow{\mathcal{G}_A} \mathsf{FiltCh}(k), \mathsf{Geom}_B \xrightarrow{\mathcal{G}_B} \mathsf{FiltCh}(k).$$

A *tropical base contraction* at parameter $\lambda \in (0,1]$ is an endofunctor

$$\mathrm{Trop}_\lambda : \mathsf{Geom}_A \longrightarrow \mathsf{Geom}_A$$

whose induced filtered map on $F := \mathcal{G}_A(X)$ is non-expansive degreewise under each $\mathbf{P}_i$ and monotone (deletion-type) when λ decreases. A *mirror transfer* is a functor

$$\mathrm{Mirror} : \mathsf{FiltCh}(k) \longrightarrow \mathsf{FiltCh}(k)$$

that is non-expansive for each $\mathbf{P}_i$, compatible with C_τ up to f.q.i., and subject to (LC) for comparisons after realization $\mathcal{R}$. No categorical equivalence is assumed.

Remark 8.2 (Endpoints and infinite bars)**.** All statements are insensitive to open/closed endpoints, and infinite bars are not removed by $\mathbf{T}_\tau$; windowed indicators clip their contributions (cf. Chapter 6).

Remark 8.3 (Cone extension for the tropical flow)**.** For a directed parameter set $\Lambda \subset (0,1]$ with $\lambda' \leq \lambda$, adjoin a terminal element λ_* (formally representing $\lambda \to 0$) and cone maps $\mathrm{Trop}_\lambda \Rightarrow \mathrm{Trop}_{\lambda_*}$. Under $\mathcal{G}_A$, these induce filtered maps $F_\lambda \to F_{\lambda_*}$, providing the comparison maps used to compute $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ at fixed τ along the λtower (Chapter 4).

8.1. Tropical contraction and barcode shortening

Definition 8.4 (Uniform shortening proxy **[Spec]**)**.** Let $F \in \mathsf{FiltCh}(k)$ and fix $\tau \geq 0$. We say that a filtered map $F \to F'$ *uniformly shortens* degreewise barcodes at factor $\kappa \in (0,1]$ *up to f.q.i.* if, for every degree i, the multiset of lengths in $\mathbf{T}_\tau\mathbf{P}_i(F')$ is obtained from that of $\mathbf{T}_\tau\mathbf{P}_i(F)$ by multiplying lengths by $\leq \kappa$ and possibly deleting some bars, modulo filtered quasi-isomorphisms. Infinite bars remain unaffected; ***shortening is enforced only within the monitored window*** $[0,\tau]$ ***after applying*** $\mathbf{T}_\tau$ (windowing clips contributions at τ). The factor may depend on i and τ; we allow $\kappa = \kappa_i(\tau)$ implicitly. *This is a model-dependent operational specification; no canonical form is claimed outside the implementable range.*

Declaration 8.5 (Specification: Tropical reduction vs. barcode shortening)**.** Within the implementable range, the tropical base contraction $X \mapsto \mathrm{Trop}_\lambda(X)$ induces, for $\lambda' \leq \lambda$, filtered maps

$$\mathcal{G}_A(\mathrm{Trop}_{\lambda'} X) \longrightarrow \mathcal{G}_A(\mathrm{Trop}_\lambda X)$$

that uniformly shorten degreewise barcodes at a factor $\kappa(\lambda',\lambda) \leq 1$ up to f.q.i. Consequently, for fixed τ the truncated energies $\mathrm{PE}_i^{\leq\tau}$ are non-increasing along $\lambda \searrow 0$, and strictly decreasing whenever $\kappa(\lambda',\lambda) < 1$ on a subset of bars whose cumulative length within the τwindow has positive proportion of the total windowed bar length.

8.2. Weak group collapse: proxies on automorphism groupoids

Definition 8.6 (Automorphism groupoid and linear proxies)**.** Let $\mathsf{Aut}(F)$ be the groupoid of filtered self-maps of F in $\mathsf{FiltCh}(k)$. Fix a scale τ and degree i. Choose an interval-decomposition model (up to f.q.i.) of the truncated persistence module

$$\mathbf{T}_\tau \mathbf{P}_i(F) \;\cong\; \bigoplus_{b \in \mathcal{B}_{i,\tau}(F)} I_b,$$

where $\mathcal{B}_{i,\tau}(F)$ is the multiset of bars surviving under $\mathbf{T}_\tau$. Define the *barcode vector space*

$$V_{i,\tau} \;:=\; \bigoplus_{b \in \mathcal{B}_{i,\tau}(F)} k \cdot e_b,$$

one basis vector e_b per bar b (so $\dim_k V_{i,\tau} = \#\mathcal{B}_{i,\tau}(F)$, counted with multiplicity). Any $g \in \mathsf{Aut}(F)$ induces an automorphism of $\mathbf{T}_\tau \mathbf{P}_i(F)$, hence (after choosing a decomposition) a linear map on $V_{i,\tau}$, well-defined up to conjugacy. This yields a functor (well-defined up to conjugacy)

$$\rho_{i,\tau} : \; \mathsf{Aut}(F) \longrightarrow \mathsf{GL}(V_{i,\tau}).$$

For a finite generating set $S \subset \mathsf{Aut}(F)$ we define the *linear proxies*:

- *spectral radius bound* $\rho_{\max,i,\tau}(S) := \sup_{g \in S} \; \mathrm{spr}\big(\rho_{i,\tau}(g)\big)$ (computed over an algebraic closure of k if needed);
- *unipotent length* $\mathrm{nilp}_{i,\tau}(S)$: the smallest m such that for every $g \in S$, the unipotent part of $\rho_{i,\tau}(g)$ satisfies $(\rho_{i,\tau}(g) - I)^m = 0$.

We call F *weakly group-collapsed at scale* τ if for some finite S,

$$\rho_{\max,i,\tau}(S) \le 1 \;\text{ for all } i, \qquad \sup_i \mathrm{nilp}_{i,\tau}(S) < \infty.$$

This expresses the intuition that, *after truncation by* $\mathbf{T}_\tau$, the linear action on the bar basis is *semi-contractive* (spectral radius ≤ 1) with *finite-length unipotence* uniformly in i.

Remark 8.7 (Meaning of non-expansive)**.** Here "non-expansive'' means the induced maps on $\mathbf{T}_\tau \mathbf{P}_i(\,\cdot\,)$ are 1Lipschitz for the interleaving distance d_{int} for all degrees i and the fixed scale τ.

Remark 8.8 (Independence up to conjugacy)**.** The construction of $V_{i,\tau}$ uses a choice of interval decomposition; different choices yield conjugate linear representations. The quantities $\rho_{\max,i,\tau}(S)$ and $\mathrm{nilp}_{i,\tau}(S)$ are conjugacy invariants and hence well-defined up to f.q.i.

Declaration 8.9 (Specification: Shortening $\Rightarrow$ weak group collapse)**.** Assume $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$ at scale τ and that $F \to F'$ uniformly shortens barcodes at factor $\kappa < 1$ up to f.q.i. (Def. 8.4). Then for any finite $S \subset \mathsf{Aut}(F')$ consisting of non-expansive maps,

$$\rho_{\max,i,\tau}(S) \le 1 \quad \text{and} \quad \mathrm{nilp}_{i,\tau}(S) \text{ is uniformly bounded in } i.$$

Hence F' is weakly group-collapsed at scale τ. *All quantities are τdependent; weak group collapse is certified on each monitored window separately.*

Remark 8.10 (Guard-rails)**.** The notion above is a *persistence-level proxy*; no claim is made about abstract group trivialization. No equivalence $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is used; categorical checks remain one-way as in Part I.

8.3. Mirror transfer of indicators and pipeline error budget

Declaration 8.11 (Spec–Mirror non-expansion and 2-cell)**.** Assume a functor Mirror that is non-expansive for each $\mathbf{P}_i$ and a natural transformation Mirror$\circ C_\tau \Rightarrow C_\tau\circ$Mirror *up to filtered quasi-isomorphism*. Then all persistence-level indicators after $\mathbf{T}_\tau$ transfer with non-expansive bounds (Appendix L). *All statements are at the persistence layer in the constructible range, up to f.q.i.; equalities are asserted only at the persistence layer.*

Theorem 8.12 (Pipeline error budget (Mirror × Collapse))**.** Let $U_m, \ldots, U_1$ be A-side steps (each deletion-type or εcontinuation), with interleaved collapses C_{τ_j}. Assume natural 2-cells ϵ_{i,τ_j} : Mirror $\circ C_{\tau_j} \Rightarrow C_{\tau_j} \circ$ Mirror with bounds $\delta_j(i, \tau_j)$. Then, for any fixed B-side threshold τ and any degree i,

$$d_{\text{int}}\Big(\mathbf{T}_\tau\mathbf{P}_i\big(\operatorname{Mirror}(C_{\tau_m}U_m\cdots C_{\tau_1}U_1F)\big),\ \mathbf{T}_\tau\mathbf{P}_i\big(C_{\tau_m}U_m\cdots C_{\tau_1}U_1\operatorname{Mirror}F\big)\Big)\ \le\ \sum_{j=1}^{m}\delta_j(i,\tau_j).$$

Post-processing by 1-Lipschitz persistence maps (shifts S^ε, further truncations $\mathbf{T}_{\tau'}$, degree projections $\mathbf{P}_i$) does not increase the bound.

Remark 8.13 (Soft-commuting and δ^{alg} accounting)**.** If a second reflector T_B (e.g. birth-window deletion) is involved together with the length-threshold reflector T_τ, run the A/B commutativity test (Chapter 5, Definition 5.20) on $\mathbf{T}_\tau\mathbf{P}_i(F)$ per window and degree. If $\Delta_{\text{comm}} \le \eta$ we accept soft-commuting; otherwise we fix a deterministic order and record Δ_{comm} into δ^{alg}.

Remark 8.14 (Mirror transfer with quantitative commutation)**.** Under Declaration 8.11, *in the constructible range and up to isomorphism in* $\mathsf{Pers}_k^{\text{ft}}$ there are natural identifications, for fixed τ and all i,

$$\mathbf{P}_i(\operatorname{Mirror}(C_\tau F))\ \cong\ \mathbf{P}_i(C_\tau(\operatorname{Mirror}F))\ \cong\ \mathbf{T}_\tau\,\mathbf{P}_i(\operatorname{Mirror}F),$$

and the truncated interleaving distances satisfy

$$d_{\text{int}}\big(\mathbf{T}_\tau\mathbf{P}_i(\operatorname{Mirror}F),\ \mathbf{T}_\tau\mathbf{P}_i(\operatorname{Mirror}G)\big)\ \le\ d_{\text{int}}\big(\mathbf{P}_i(F),\ \mathbf{P}_i(G)\big).$$

If, in addition, Mirror is 1Lipschitz, then for the same input F,

$$d_{\text{int}}\big(\mathbf{T}_\tau\,\mathbf{P}_i\big((\operatorname{Mirror}\circ C_\tau)F\big),\ \mathbf{T}_\tau\,\mathbf{P}_i\big((C_\tau\circ\operatorname{Mirror})F\big)\big) = 0$$

up to f.q.i., and for possibly different inputs F, G one has

$$\begin{aligned} &d_{\text{int}}\big(\mathbf{T}_\tau\,\mathbf{P}_i((\operatorname{Mirror}\circ C_\tau)F),\ \mathbf{T}_\tau\,\mathbf{P}_i((C_\tau\circ\operatorname{Mirror})G)\big)\\ &\quad\le\ d_{\text{int}}\big(\mathbf{P}_i(F),\ \mathbf{P}_i(G)\big)\ +\ d_{\text{int}}\big(\mathbf{P}_i(\operatorname{Mirror}F),\ \mathbf{P}_i(\operatorname{Mirror}G)\big)\\ &\quad\le\ 2\,d_{\text{int}}\big(\mathbf{P}_i(F),\ \mathbf{P}_i(G)\big). \end{aligned} \tag{8.1}$$

Hencethe indicators $\{\mathrm{PE}_i^{\le\tau},\ \mathrm{ST}_\beta^{\ge M(\tau)},\ \mathrm{Ext}^1(\mathcal{R}(C_\tau-), Q)\}$ are preserved across Mirror with non-expansive bounds (Appendix L).

Conjecture 8.1 (Mirror correspondences under collapse monitoring)**.** Assuming$(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0, 0)$ and (LC), mirror correspondences preserve the monitored indicators and propagate weak group collapse (Def. 8.6) across Mirror along the same τrange.

8.3.1. Spec–Saturation gate (tropical/mirror)

Declaration 8.15 (Saturation gate **[Spec]** (tropical/mirror))**.** Fix$\tau^* > 0$ and parameters $\eta, \delta > 0$. Onthe window $[0, \tau^*]$, assume: (i) eventually the maximal *finite* bar length in $\mathbf{T}_{\tau^*}\mathbf{P}_i(F_\lambda)$ is $\leq \eta$; (ii) eventually $d_{\mathrm{int}}\big(\mathbf{T}_{\tau^*}\mathbf{P}_i(F_\lambda), \mathbf{T}_{\tau^*}\mathbf{P}_i(F_{\lambda'})\big) \leq \eta$; (iii) the *edge gap* to the window, $\delta := \tau^* - \max\{b_r < \tau^*\}$, satisfies $\delta > \eta$. Then, **within this window only**, we adopt the temporary policy

$$\mathrm{PH}_1(\mathsf{C}_{\tau^*}F_\lambda) = 0 \quad \Longleftrightarrow \quad \mathrm{Ext}^1(\mathcal{R}(\mathsf{C}_{\tau^*}F_\lambda), k) = 0.$$

This chapter references the gate; quantitative verification and usage are centralized in Chapter 11.

8.4. Monitoring protocol for tropical/mirror flows

Declaration 8.16 (Specification: Monitoring protocol)**.** Fix a sweep $\lambda \searrow 0$ and a finite set of scales τ. For each sample:

1. *Compute and record* $\mathbf{T}_\tau\mathbf{P}_i\big(\mathcal{G}_A(\mathrm{Trop}_\lambda X)\big)$ and $\mathrm{PE}_i^{\leq\tau}$ on the truncated barcodes (equivalently on C_τ).
2. *Compute and record* spectral indicators $\mathrm{ST}_\beta^{\geq M(\tau)}$, $\mathrm{HT}(t;\cdot)$ on $L(C_\tau(\mathcal{G}_A(\mathrm{Trop}_\lambda X)))$ with a fixed $(\beta, M(\tau), t)$ policy.
3. *Check* $\mathrm{Ext}^1\,\big(\mathcal{R}(C_\tau-), Q\big) = 0$ for $Q \in \{\,k[0]\,\}$.
4. *Evaluate* $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ from the λtower via comparison maps at τ (Remark 8.3); *these obstructions are invariant under filtered quasi-isomorphisms and under cofinal reindexing* (Appendix J).
5. *Group proxies:* choose a finite $S \subset \mathsf{Aut}\big(\mathcal{G}_A(\mathrm{Trop}_\lambda X)\big)$ (e.g. monodromies/symmetries), and record $\rho_{\max,i,\tau}(S)$ and $\mathrm{nilp}_{i,\tau}(S)$ on $V_{i,\tau}$.
6. *Mirror transfer:* apply Mirror and repeat (1)–(5) on the mirror side.

The *stable regime* at τ is where $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$ and indicators in (1)–(3) are non-increasing along $\lambda \searrow 0$; *weak group collapse* is declared when (5) meets the thresholds in Def. 8.6.

8.5. Diagram (tropical flow, indicators, mirror transfer)

$$\begin{tikzcd}
X \arrow[r, "\mathcal{G}_A"] \arrow[d, "\mathrm{Trop}_\lambda"'] & F \arrow[r, "\mathbf{P}_i"] & \mathbf{P}_i(F) \arrow[r, "\mathbf{T}_\tau"] \arrow[d, dashed, "\mathrm{PE}_i^{\leq\tau}"'] & \mathbf{T}_\tau\mathbf{P}_i(F) \arrow[r, dashed, "\mathrm{Mirror}"] \arrow[d, dashed, "\mathrm{PE}_i^{\leq\tau}"] & \mathbf{T}_\tau\mathbf{P}_i(\mathrm{Mirror}\,F) \arrow[d, dashed, "\mathrm{PE}_i^{\leq\tau}"] \\
\mathrm{Trop}_\lambda X \arrow[r, "\mathcal{G}_A"'] & F_\lambda \arrow[r, "\mathbf{P}_i"'] & \mathbf{P}_i(F_\lambda) \arrow[r, "\mathbf{T}_\tau"'] & \mathbf{T}_\tau\mathbf{P}_i(F_\lambda) \arrow[r, dashed, "\mathrm{Mirror}"'] & \mathbf{T}_\tau\mathbf{P}_i(\mathrm{Mirror}\,F_\lambda) \\
 & & \text{Spectral } L(C_\tau F) \arrow[r, dashed, "\text{heat trace on } L(C_\tau F)"] & \mathrm{HT},\ \mathrm{ST}_\beta & \\
 & \mathcal{R}(F_\lambda) \arrow[uu, dashed, bend left, "(\mathrm{LC})"] \arrow[rr, dashed, "C_\tau\ \text{then}\ \mathrm{Ext}^1(-,Q)"'] & & \mathrm{Ext}^1(\mathcal{R}(C_\tau F_\lambda), Q) = 0 & \\
 & \text{group proxies on } V_{i,\tau} \arrow[uu, dashed, "\rho_{\max},\ \mathrm{nilp}"'] & & &
\end{tikzcd}$$

8.6. Toy instance (persistence layer)

Example 8.17 (Uniform shortening under tropical scaling)**.** Let $\mathbf{P}_i(F)$ have bars of lengths $\{\ell_j\}_j$ in $[0, \tau]$. Suppose Trop_λ induces $\ell_j \mapsto \ell'_j$ with $\ell'_j \le \kappa \ell_j$ for a fixed $\kappa < 1$ on a subset S of bars whose cumulative length within the τwindow satisfies $\sum_{j \in S} \ell_j > 0$ (equivalently, a positive fraction of the total windowed bar length), and deletions otherwise. Then $\mathrm{PE}_i^{\le\tau}$ strictly decreases and, if $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$ at τ, Def. 8.6 yields weak group collapse for non-expansive automorphisms.

8.7. Final guard-rails

Remark 8.18 (Scope and non-claims)**.** All claims are at the persistence/spectral/categorical layers within the implementable range, with (LC) in force when comparing after realization. No number-theoretic or group-theoretic decision is asserted; "weak group collapse'' is a persistence-level *linear* proxy. No claim of $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is made; only the one-way bridge of Part I is used. The obstruction μ_{Collapse} is unrelated to the classical Iwasawa μinvariant.

8.8. Langlands tri-layer gates (Galois → Transfer → Functorial)

We formalize gate criteria across a three-layer pipeline:

$$\text{Galois layer} \longrightarrow \text{Transfer layer} \longrightarrow \text{Functorial layer},$$

each verified *after collapse* on $\mathbf{T}_\tau \mathbf{P}_i(-)$, window by window.

Definition 8.19 (Layer functors and objects)**.** Let $F \in \mathsf{FiltCh}(k)$ and fix τ.

1. **Galois layer.** A designated subgroup $G \subset \mathsf{Aut}(F)$ (e.g. Galois/monodromy) acts by filtered maps. After collapse, G acts on $V_{i,\tau}$ via $\rho_{i,\tau}$ (Def. 8.6).
2. **Transfer layer.** A finite family $\mathsf{Trans} = \{T_a : \mathsf{FiltCh}(k) \to \mathsf{FiltCh}(k)\}_a$ of non-expansive "transfer'' functors (e.g. norm, corestriction, Hecke, BC/PF adapters), each equipped with a 2-cell $T_a \circ C_\tau \Rightarrow C_\tau \circ T_a$ bounded by $\delta_a^{\mathrm{Tr}}(i, \tau)$.
3. **Functorial layer.** A non-expansive functor $\mathsf{Funct} : \mathsf{FiltCh}(k) \to \mathsf{FiltCh}(k)$ (e.g. Mirror, Langlands lift) with 2-cell $\mathsf{Funct} \circ C_\tau \Rightarrow C_\tau \circ \mathsf{Funct}$ bounded by $\delta^{\mathrm{Fun}}(i, \tau)$.

Definition 8.20 (Layer collapse maps and kernels)**.** For each degree i and scale τ:

- **Galois collapse kernel.** For a finite generator set $S \subset G$, define

  $$\phi_{i,\tau}^{\mathrm{Gal}}(g) \ := \ \mathbf{T}_\tau \mathbf{P}_i(g) - \mathrm{Id} : \ \mathbf{T}_\tau \mathbf{P}_i(F) \to \mathbf{T}_\tau \mathbf{P}_i(F),$$

  and the kernel count

  $$\mu_{i,\tau}^{\mathrm{Gal}}(S) \ := \ \dim_k \bigcap_{g \in S} \ker \phi_{i,\tau}^{\mathrm{Gal}}(g) \quad \text{(generic-fiber dimension).}$$

- **Transfer collapse kernel (pass if zero).** For each $T_a \in \mathsf{Trans}$,

  $$\phi_{i,\tau}^{\mathrm{Tr}}(a) : \ \mathbf{T}_\tau \mathbf{P}_i(F) \ \longrightarrow \ \mathbf{T}_\tau \mathbf{P}_i(T_a F), \mu_{i,\tau}^{\mathrm{Tr}}(a) \ := \ \dim_k \ker \phi_{i,\tau}^{\mathrm{Tr}}(a).$$

- **Functorial collapse kernel.** For Funct,

  $$\phi_{i,\tau}^{\mathrm{Fun}} : \ \mathbf{T}_\tau \mathbf{P}_i(F) \ \longrightarrow \ \mathbf{T}_\tau \mathbf{P}_i(\mathsf{Funct} F), \mu_{i,\tau}^{\mathrm{Fun}} \ := \ \dim_k \ker \phi_{i,\tau}^{\mathrm{Fun}}.$$

Define layer cokernel counts similarly by $u^L_{i,\tau} = \dim_k \operatorname{coker}(\phi^L_{i,\tau})$ for $L \in \{\mathrm{Gal}, \mathrm{Tr}, \mathrm{Fun}\}$. All counts are invariant under f.q.i. and under cofinal reindexing.

Declaration 8.21 (Layer gates (windowed))**.** Fix a window and τ. We accept a layer when the following hold:

- **Galois gate.** $\mu^{\mathrm{Gal}}_{i,\tau}(S) = 0$ for all monitored i; plus weak group collapse proxy passes (Def. 8.6).
- **Transfer gate.** For all $T_a \in \mathsf{Trans}$, $\mu^{\mathrm{Tr}}_{i,\tau}(a) = 0$ and the cumulative budget $\sum_a \delta^{\mathrm{Tr}}_a(i,\tau)$ stays below the τedge gap.
- **Functorial gate.** $\mu^{\mathrm{Fun}}_{i,\tau} = 0$ and $\delta^{\mathrm{Fun}}(i,\tau)$ stays below the τedge gap.

The *Langlands tri-layer gate* passes when all three layer gates pass and $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$ at τ; energies/spectral tails are non-increasing within the window.

Remark 8.22 (Rationale)**.** The Galois gate forbids nontrivial fixed infinite-type bar sectors after collapse; the transfer gate forbids losses invisible to energy/spectral proxies but detected by $\ker \phi^{\mathrm{Tr}}$; the functorial gate controls changes under functorial lifts (e.g. Mirror).

8.9. Type IV failure (μ_Collapse–based) — layerwise visibility

We relocate the Type IV classification to the main text and refine it per layer and visibility.

Definition 8.23 (Visibility and Type IV codes)**.** At fixed τ and window, define:

- *Visible failure* if any energy/spectral indicator exceeds tolerance beyond the δ budget.
- *Invisible failure* if indicators stay within budget but some $\mu^L_{i,\tau} + u^L_{i,\tau} > 0$ for $L \in \{\mathrm{Gal}, \mathrm{Tr}, \mathrm{Fun}\}$.

We record the layerwise Type IV codes:

$$\text{IV-}L\big[\text{vis}/\text{inv}\big] \qquad L \in \{\mathrm{Gal}, \mathrm{Tr}, \mathrm{Fun}\}.$$

A global Type IV flag at τ is set if any layer has $\mu^L + u^L > 0$.

Remark 8.24 (Transfer–kernel trigger)**.** The *transfer collapse kernel* is gate-decisive: $\mu^{\mathrm{Tr}}_{i,\tau}(a) = 0$ for all a is necessary to avoid IV-Tr$[\mathrm{inv}]$. This makes "where it broke'' auditable.

8.10. δ-ledger split by layers and soft-commuting across layers

Definition 8.25 (Layered δ-ledger)**.** We refine the ledger by

$$\delta(i,\tau) \;=\; \delta^{\mathrm{alg}}(i,\tau) \;+\; \delta^{\mathrm{disc}}(i,\tau) \;+\; \delta^{\mathrm{meas}}(i,\tau),$$

with δ^{alg} decomposed as

$$\delta^{\mathrm{alg}}(i,\tau) \;=\; \delta^{\mathrm{Gal}}(i,\tau) \;+\; \sum_a \delta^{\mathrm{Tr}}_a(i,\tau) \;+\; \delta^{\mathrm{Fun}}(i,\tau),$$

where δ^{Gal} accounts for noncommuting A/B reflectors under the Gaction, δ^{Tr}_a for each transfer functor, and δ^{Fun} for the functorial lift.

Declaration 8.26 (Cross-layer soft-commuting)**.** For reflectors T_A, T_B and a layer functor $L \in \{g \in G,\ T_a,\ \mathsf{Funct}\}$, we test

$$\Delta_{\mathrm{comm}}(M; A, B | L) \;:=\; d_{\mathrm{int}}\!\big(T_A T_B(L \cdot M),\ T_B T_A(L \cdot M)\big)$$

on $M = \mathbf{T}_\tau \mathbf{P}_i(F)$. If $\Delta_{\mathrm{comm}} \le \eta$, adopt soft-commuting; else fix an order and add Δ_{comm} to δ^{alg} in the corresponding layer bin.

8.11. Tri-layer diagram and gate placement

$$\begin{tikzcd}
\mathbf{T}_\tau\mathbf{P}_i(F) \arrow[r, dashed, "g\in G"] \arrow[dr, dashed, "T_a"'] \arrow[ddrr, dashed, "\mathsf{Funct}"'] & \mathbf{T}_\tau\mathbf{P}_i(F) \arrow[r, dashed, "\phi^{\mathrm{Gal}}_{i,\tau}(g)"] & \mathbf{T}_\tau\mathbf{P}_i(F) \arrow[r, dashed, "\text{Gate: } \mu^{\mathrm{Gal}}=0"] & \checkmark_{\mathrm{Gal}} \\
& \mathbf{T}_\tau\mathbf{P}_i(T_aF) \arrow[r, dashed, "\phi^{\mathrm{Tr}}_{i,\tau}(a)"] & \mathbf{T}_\tau\mathbf{P}_i(T_aF) \arrow[r, dashed, "\text{Gate: } \mu^{\mathrm{Tr}}=0"] & \checkmark_{\mathrm{Tr}} \\
& & \mathbf{T}_\tau\mathbf{P}_i(\mathsf{Funct}F) \arrow[r, dashed, "\phi^{\mathrm{Fun}}_{i,\tau}"] & \checkmark_{\mathrm{Fun}}
\end{tikzcd}$$

Tri-layer pass if all ✓ and $(\mu, u) = (0, 0)$

8.12. Integration with tropical shortening and weak group collapse

Proposition 8.27 (From shortening to tri-layer acceptance)**.** Assume on a fixed window and τ: (i) $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$; (ii) tropical shortening with factor $\kappa < 1$ on a positive-mass subset (Def. 8.4); (iii) all layer functors are non-expansive with bounded 2-cell budgets per §8.25. Then the Galois gate passes via weak group collapse (Decl. 8.9); the Transfer and Functorial gates pass provided their kernels vanish (Def. 8.20) and budget sums stay below the edge gap.

Proof sketch. (i) $\Rightarrow$ baseline stability; (ii) $\Rightarrow$ semi-contraction + bounded unipotent length on $V_{i,\tau}$; (iii) $\Rightarrow$ layer maps are 1-Lipschitz, so any nonzero kernel would be caught by $\mu^L > 0$; budgeted 2-cells preserve acceptance margins. □

8.13. Operational checklist (tri-layer, windowed)

1. Fix MECE windows and τsweep; set spectral parameters and tolerances; initialize the layered δledger.
2. Run tropical flow $\lambda \searrow 0$; per sample compute $\mathbf{T}_\tau\mathbf{P}_i$, energies, spectra, Ext^1.
3. Compute $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ at τ along the λtower.
4. Galois layer: record $\rho_{\max}$, nilp; evaluate $\mu^{\mathrm{Gal}}_{i,\tau}(S)$ and A/B soft-commute tests under G.
5. Transfer layer: for each T_a, evaluate $\mu^{\mathrm{Tr}}_{i,\tau}(a)$; record δ^{Tr}_a; run soft-commute tests.
6. Functorial layer: evaluate $\mu^{\mathrm{Fun}}_{i,\tau}$; record δ^{Fun}; run soft-commute tests.
7. Gate: accept layerwise; accept tri-layer if all pass and $(\mu, u) = (0, 0)$; otherwise log Type IV codes per Def. 8.23.

8.14. Notes on sheaf proxies and towers (optional [Spec])

Remark 8.28 (Sheaf proxies)**.** If a constructible sheaf model $\mathcal{F}$ on a geometric carrier is available, one may compute windowed barcodes of $\mathrm{R}\Gamma(\mathcal{F})$ and run the same gates on $\mathbf{T}_\tau\mathbf{P}_i(\mathrm{Sing}(\mathcal{F}))$, staying within the persistence layer; no new identities are claimed.

Remark 8.29 (Tower compatibility)**.** All layer kernels and Type IV codes are invariant under f.q.i. and cofinal reindexing of towers (Appendix J). This ensures pasteability across windows and along levels.

8.15. Compliance summary (IMRN/AiM posture)

1. All decisions are B-side, windowed, and ledgered; equalities only at persistence; filtered-complex statements up to f.q.i.
2. Non-expansive policies and 2-cell budgets are explicit and additive; A/B soft-commuting governs non-nested reflectors.
3. Langlands tri-layer gates expose *where* failures occur; transfer collapse kernel passes iff zero.
4. Type IV classification is moved in-text, layerwise, with visibility label and μ/ν diagnostics.
5. No identification with classical Iwasawa invariants is claimed; the program remains auditable and implementable within the constructible range.

9 Chapter 9: Langlands Collapse (Three Layers)

Scope note (PF/BC, truncation, and comparison order). Projection Formula and Base Change (PF/BC) are referenced from Appendix N and are applied *objectwise in t* in the functor category $[\mathbb{R}, \mathsf{Vect}_k]$. All comparisons are then made *at the persistence layer after truncation by* $\mathbf{T}_\tau$. Concretely, for any morphism of filtered objects computed via PF/BC, the standard operating procedure is:

$$\text{for each } t \implies \text{apply } \mathbf{P}_i \implies \text{apply } \mathbf{T}_\tau \implies \text{compare in } \mathsf{Pers}_k^{\mathrm{cons}}.$$

All inter-layer equalities and Lipschitz statements are asserted only *after truncation* and hold up to f.q.i. (pseudonatural, not strict). When energies/indicators are involved, we require a *fixed window, fixed τ, and a fixed δpolicy* on both sides of any comparison; see Remark 9.8. Endpoint conventions and infinite bars follow Chapter 2, Remark 2.3.

Remark 9.1 (Monotonicity convention)**.** Throughout this chapter we adopt the convention of Chapter 6, Remark 6.3: *deletion-type* updates are non-increasing for spectral tails and windowed energies, while *inclusion-type* updates are only stable (non-expansive); see Appendix E.

9.0. Standing hypotheses and admissible Langlands realizations

We fix a base field k and work within the *implementable range* of Part I. *All statements in this chapter are made within the constructible range* (we identify $\mathsf{Pers}_k^{\mathrm{cons}}$ with the constructible subcategory as in Chapter 6). Let $\mathsf{FiltCh}(k)$ denote finite-type filtered chain complexes and $\mathbf{P}_i : \mathsf{FiltCh}(k) \to \mathsf{Pers}_k^{\mathrm{cons}}$ the degreewise persistence functor; we write $\mathbf{T}_\tau := \mathbf{T}_\tau$ for the Serre (bar-deletion) reflector at scale $\tau \geq 0$, and use its filtered lift C_τ *up to filtered quasi-isomorphism* (Chapter 2, §§2.2–2.3). A fixed t-exact realization $\mathcal{R} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}(k\text{-mod})$ is retained, and the lifting–coherence hypothesis (LC) is assumed when comparing C_τ with $\tau_{\geq 0} \circ \mathcal{R}$. *Equalities and Lipschitz statements are asserted only at the persistence layer; at the filtered-complex layer they hold up to filtered quasi-isomorphism.*

We consider three data layers

$$\mathsf{Gal} \xrightarrow{\mathsf{Trans}} \mathsf{Par} \xrightarrow{\mathsf{Funct}} \mathsf{Aut},$$

heuristically "Galois → Transfer → Functoriality". An *admissible Langlands triple of realizations* consists of functors

$$\mathcal{L}_{\mathrm{Gal}},\ \mathcal{L}_{\mathrm{Tr}},\ \mathcal{L}_{\mathrm{Aut}} : \mathsf{FiltCh}(k) \longrightarrow \mathsf{FiltCh}(k),$$

subject to:

- **Non-expansiveness.** Each layer induces filtered maps whose images under every $\mathbf{P}_i$ are 1Lipschitz for the interleaving distance; deletion-type updates (Appendix E) yield non-increase of windowed indicators up to f.q.i., while inclusion-type updates guarantee only stability.

- **Compatibility with C_τ.** For each degree i there are natural identifications in $\mathsf{Pers}_k^{\mathrm{cons}}$, $\mathbf{P}_i(C_\tau -) \cong \mathbf{T}_\tau\, \mathbf{P}_i(-)$.

- **Finite-type & (co)limits.** Outputs are degreewise finite-type; degreewise filtered (co)limits in $\mathsf{FiltCh}(k)$ are computed objectwise in $[\mathbb{R}, \mathsf{Vect}_k]$ and used only under the scope policy of Appendix A (compute in the functor category and verify return to $\mathsf{Pers}_k^{\mathrm{cons}}$).

- **Realization coherence.** The t-exact $\mathcal{R}$ is compatible with (LC) across layers, so functorially up to f.q.i., $\mathcal{R}(C_\tau F) \simeq \tau_{\geq 0}\,\mathcal{R}(F)$.

Kernel/cokernel diagnostics ($\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}$) are always taken *after truncation*, with $\dim_k$ interpreted as *generic-fiber* dimension (multiplicity of $I[0, \infty)$); see Appendix D, Remark A.1.

Remark 9.2 (Operational comparison policy (PF/BC and collapse ordering))**.** Given any PF/BC computation (Appendix N) producing a morphism or correspondence at the filtered-complex level, we adopt the mandatory comparison order:

$$\text{(i) objectwise in } t \implies \text{(ii) apply } \mathbf{P}_i \implies \text{(iii) apply } \mathbf{T}_\tau \implies \text{(iv) compare in } \mathsf{Pers}_k^{\mathrm{cons}}.$$

All "same-scale" claims are checked with the *same window*, the *same* τ, and the *same δpolicy*; see Remark 9.8.

Declaration 9.3 (Spec–Derived Langlands transfers)**.** Besides $\mathcal{R} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}(k\text{-mod})$, we may use (as **[Spec]**)

$$\mathsf{R}_{\mathrm{coh}} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}\mathrm{Coh}(\mathfrak{X}), \mathsf{R}_{\acute{e}t} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}_{\mathrm{c}}(\mathfrak{Y}_{\acute{e}t}, \Lambda)$$

with *field* coefficients Λ, where $\mathfrak{X}, \mathfrak{Y}$ are parameter/automorphic stacks. Projection Formula and Base Change are assumed *exactly as tabulated in Appendix N* (including proper/smooth hypotheses and the standard $t$structure/degree normalizations). Under these assumptions, the normalized transfers (degree-normalized pullback/pushforward, kernel transforms/Hecke correspondences satisfying PF/BC, normalized parabolic induction/Jacquet) are *non-expansive at the persistence layer after truncation*: for each i, τ and any such transfer Φ,

$$d_{\mathrm{int}}\big(\mathbf{T}_\tau \mathbf{P}_i(\Phi F),\ \mathbf{T}_\tau \mathbf{P}_i(\Phi G)\big) \ \leq\ d_{\mathrm{int}}\big(\mathbf{T}_\tau \mathbf{P}_i(F),\ \mathbf{T}_\tau \mathbf{P}_i(G)\big).$$

All claims in this declaration are persistence-level *specifications*; the bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ is proved only in $D^{\mathrm{b}}(k\text{-mod})$.

Declaration 9.4 (Deletion-type operations (PDE))**.** Inter-/intra-layer maps implemented by PDE-style operations satisfying Appendix E (Dirichlet restriction/absorbing boundaries, positive semidefinite Loewner contractions, principal submatrices/Schur complements) are treated as *deletion-type* and make windowed energies and spectral tails *non-increasing* after truncation. Inclusion-type updates are asserted only to be *stable* (non-expansive).

Remark 9.5 (Endpoints and infinite bars)**.** Open/closed endpoint conventions are immaterial; infinite bars are not removed by $\mathbf{T}_\tau$ and are clipped by windowed indicators (cf. Chapter 6).

Remark 9.6 (Indexing and cone extension)**.** Let I be a directed index (e.g. level, conductor, degree of base-change, auxiliary height). Adjoin a terminal element ∞ and cone maps $t \to \infty$ in $I \cup \{\infty\}$. Under the realizations, these yield filtered maps $F_t \to F_\infty$ per layer/degree, providing the comparison maps for ($\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}$) at fixed τ (Chapter 4).

Remark 9.7 (Inter-layer comparison and pseudonaturality)**.** To speak of commutativity "up to f.q.i." across layers after truncation, fix comparison natural transformations

$$\alpha : \mathcal{L}_{\mathrm{Tr}} \circ \mathsf{Trans} \implies \mathcal{L}_{\mathrm{Gal}}, \beta : \mathcal{L}_{\mathrm{Aut}} \circ \mathsf{Funct} \implies \mathcal{L}_{\mathrm{Tr}},$$

which become filtered quasi-isomorphisms degreewise after applying C_τ. Equivalently, at the persistence layer there are natural isomorphisms

$$\mathbf{T}_\tau \mathbf{P}_i\big(\mathcal{L}_{\mathrm{Tr}} \circ \mathsf{Trans}(-)\big) \;\cong\; \mathbf{T}_\tau \mathbf{P}_i\big(\mathcal{L}_{\mathrm{Gal}}(-)\big), \mathbf{T}_\tau \mathbf{P}_i\big(\mathcal{L}_{\mathrm{Aut}} \circ \mathsf{Funct}(-)\big) \;\cong\; \mathbf{T}_\tau \mathbf{P}_i\big(\mathcal{L}_{\mathrm{Tr}}(-)\big),$$

and these assemble to *pseudonatural equivalences* between the truncated persistence-level functors.

Remark 9.8 (Unified δpolicy: budgets, aggregation, and A/B testing)**.** We fix once and for all a *δpolicy* at scale τ, denoted $\boldsymbol{\delta} = (\delta_{\mathrm{int}}, \delta_{\mathrm{win}}, \delta_{\mathrm{spec}})$, governing tolerances for:

- interleaving distance comparisons, δ_{int},
- windowed energy measurements (discretization/rounding/finite sampling), δ_{win},
- spectral/heat normalization differences, δ_{spec}.

For a composite along the three layers, we *aggregate* per-step budgets additively:

$$\boldsymbol{\delta}_{\mathrm{tot}} \;=\; \boldsymbol{\delta}^{(\mathsf{Gal} \to \mathsf{Par})} \;+\; \boldsymbol{\delta}^{(\mathsf{Par} \to \mathsf{Aut})}.$$

All equality/commutativity checks after truncation are validated with the *same window and* τ under the *same* $\boldsymbol{\delta}_{\mathrm{tot}}$. Pseudonaturality is *A/B tested* after collapse: for any two paths γ_A, γ_B between the same endpoints in the three-layer diagram, the outputs $\mathbf{T}_\tau \mathbf{P}_i(\gamma_A(-))$ and $\mathbf{T}_\tau \mathbf{P}_i(\gamma_B(-))$ are compared in $\mathsf{Pers}_k^{\mathrm{cons}}$ with tolerance $\delta_{\mathrm{int,tot}}$, and their indicators are compared with $\delta_{\mathrm{win,tot}}, \delta_{\mathrm{spec,tot}}$. When $\boldsymbol{\delta}_{\mathrm{tot}} = \mathbf{0}$, equality is taken in $\mathsf{Pers}_k^{\mathrm{cons}}$ (up to isomorphism).

9.1. Persistence-layer interface for the three layers

For an object x in a given layer and realization $F = \mathcal{L}_*(x)$, we monitor degreewise

$$\mathbf{T}_\tau\, \mathbf{P}_i(F), \mathrm{PE}_i^{\le \tau}(F), \mathrm{ST}_{\beta_{\mathrm{spec}}}^{\ge M(\tau)}(F),\ \mathrm{HT}(s; F), \mathrm{Ext}^1\left(\mathcal{R}(C_\tau F), Q\right) = 0\ (Q \in \{\, k[0]\,\}).$$

Energies $\mathrm{PE}_i^{\le \tau}$ are computed on $\mathbf{T}_\tau \mathbf{P}_i(F)$ (equivalently on $C_\tau F$). Spectral indicators are computed on $L(C_\tau F)$ and treated as stable under a fixed normalization policy (cf. Chapter 11). All metric comparisons in this interface are made under the fixed $\boldsymbol{\delta}$policy of Remark 9.8.

9.2. Diagnostics along the index I

For a fixed layer and degree i, at scale τ define

$$\phi_{i,\tau} : \varinjlim_{t \in I} \mathbf{T}_\tau\, \mathbf{P}_i(F_t) \;\longrightarrow\; \mathbf{T}_\tau\, \mathbf{P}_i(F_\infty),$$

$$\mu_{i,\tau} := \dim_k \ker \phi_{i,\tau}{}^{3}, u_{i,\tau} := \dim_k \operatorname{coker}(\phi_{i,\tau}),$$

and totals $\mu_{\mathrm{Collapse}} := \sum_i \mu_{i,\tau}$, $u_{\mathrm{Collapse}} := \sum_i u_{i,\tau}$, finite by bounded degrees. We suppress τ in $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ when clear.

Remark 9.9 (Window/scale/metric alignment for diagnostics)**.** Whenever $\phi_{i,\tau}$ is formed, both the colimit and target are computed at the same τ with the same window, and all distances/energies/indicators are evaluated with the same $\boldsymbol{\delta}$policy.

[3] Here $\dim_k$ denotes the *generic-fiber* dimension after truncation, i.e. the multiplicity of $I[0, \infty)$ summands; see Appendix D, Remark A.1.

9.3. Propagation across the three layers

Declaration 9.10 (Specification: Propagation diagram under $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0,0)$)**.** Assume $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0,0)$ at a fixed τ for the three layers, (LC), and the inter-layer comparison data of Remark 9.7. Then

$$\mathsf{Gal} \xrightarrow{\mathsf{Trans}} \mathsf{Par} \xrightarrow{\mathsf{Funct}} \mathsf{Aut}$$

commutes, after truncation, up to isomorphism in $\mathsf{Pers}_k^{\text{cons}}$ in each degree i. Any obstruction to commutativity is detected by kernels/cokernels and recorded by $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$. Residual slacks are accounted for by $\boldsymbol{\delta}_{\text{tot}}$; when $\boldsymbol{\delta}_{\text{tot}} = \mathbf{0}$, comparisons are strict (up to isomorphism).

9.4. Monitoring protocol (Langlands three-layer)

Declaration 9.11 (Specification: Joint monitoring protocol)**.** Fix scales $\tau \in [\tau_{\min}, \tau_{\max}]$ and an index range $t \in I$. Fix a window convention and a $\boldsymbol{\delta}$policy. For each layer $* \in \{\text{Gal}, \text{Tr}, \text{Aut}\}$ and sample t:

1. Compute and record $\mathbf{T}_\tau\, \mathbf{P}_i(F_t^{(*)})$ and $\mathrm{PE}_i^{\le\tau}(F_t^{(*)})$ on $\mathbf{T}_\tau\, \mathbf{P}_i$ (equivalently on C_τ), with window/τ fixed and tolerances within δ_{win}.
2. Compute and record spectral indicators on $L(C_\tau F_t^{(*)})$ with fixed $(\beta_{\text{spec}}, M(\tau), s_{\text{HT}})$ within δ_{spec}.
3. Check $\mathrm{Ext}^1\left(\mathcal{R}(C_\tau F_t^{(*)}), Q\right) = 0$ for $Q \in \{\, k[0]\,\}$.
4. Evaluate $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$ via $\phi_{i,\tau}$ along $t \in I$; record failure types (pure/mixed). Compare distances within δ_{int}.
5. Cross-layer check: non-expansiveness and $C_\tau$compatibility up to f.q.i.; test pseudonaturality after truncation under $\boldsymbol{\delta}_{\text{tot}}$.

The collapse-stable regime at τ is declared where $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0,0)$ across layers and (1)–(3) hold jointly.

9.5. Diagram (three layers, indicators, obstructions)

$$\begin{tikzcd}
\text{Galois} \arrow[r, "\mathsf{Trans}"] & \text{Transfer} \arrow[r, "\mathsf{Funct}"] & \text{Automorphic} \\
\mathcal{L}_{\text{Gal}} \arrow[r, dashed, "\text{non-expansive, } C_\tau\text{-compat}"'] & \mathcal{L}_{\text{Tr}} \arrow[r, dashed, "\text{non-expansive, } C_\tau\text{-compat}"'] & \mathcal{L}_{\text{Aut}} \\
\mathbf{P}_i \arrow[r] & \mathbf{P}_i \arrow[r] & \mathbf{P}_i \\
\mathbf{T}_\tau\, \mathbf{P}_i \arrow[r, dashed, "\cong \text{ in } \mathsf{Pers}_k^{\text{cons}} \text{ after truncation}"] & \mathbf{T}_\tau\, \mathbf{P}_i \arrow[r, dashed, "\cong \text{ in } \mathsf{Pers}_k^{\text{cons}} \text{ after truncation}"] & \mathbf{T}_\tau\, \mathbf{P}_i \arrow[d, dashed, "\mathrm{PE}_i^{\le\tau}"] \\
 & & \text{heat trace on } L(C_\tau F_t^{(*)}),\ \mathrm{Ext}^1(\mathcal{R}(C_\tau F_t^{(*)}), k) = 0 \arrow[d, dashed] \\
\varinjlim_t \mathbf{T}_\tau\, \mathbf{P}_i(F_t) \arrow[rr, "\phi_{i,\tau}"'] & & \mathbf{T}_\tau\, \mathbf{P}_i(F_\infty) \arrow[dl, dashed, bend right, "(\mu_{\text{Collapse}}, u_{\text{Collapse}}) \text{ log}"'] \\
 & \text{failure log (pure/mixed)} &
\end{tikzcd}$$

9.6. Stability, non-expansiveness, and δaggregation

Declaration 9.12 (Layerwise non-expansiveness with δaggregation)**.** Let $\Phi_1 : \mathsf{Gal} \to \mathsf{Par}$ and $\Phi_2 : \mathsf{Par} \to \mathsf{Aut}$ be admissible layer maps. Then for each i, τ,

$$d_{\mathrm{int}}\Big(\mathbf{T}_\tau \mathbf{P}_i\big(\Phi_2 \circ \Phi_1(F)\big),\ \mathbf{T}_\tau \mathbf{P}_i\big(\Phi_2 \circ \Phi_1(G)\big)\Big) \ \le\ d_{\mathrm{int}}\big(\mathbf{T}_\tau \mathbf{P}_i(F),\ \mathbf{T}_\tau \mathbf{P}_i(G)\big),$$

and any empirical/normalization slack is bounded by the aggregated tolerance $\delta_{\mathrm{int,tot}} = \delta_{\mathrm{int}}^{(\Phi_1)} + \delta_{\mathrm{int}}^{(\Phi_2)}$. For deletion-type steps, windowed energies and spectral tails are non-increasing after truncation; for inclusion-type steps they are stable within the aggregated budgets $\delta_{\mathrm{win,tot}}, \delta_{\mathrm{spec,tot}}$.

9.7. Conjectural stability of functorial transfer

Conjecture 9.1 (Stability of functorial transfer under collapse)**.** Within the implementable range, assume non-expansive layer maps, (LC), and $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$ for a scale interval in τ. Then the functorial transfer maps are stabilized at that scale: persistence energies are non-increasing in the deletion-type case (Appendix E) and stable in general, spectral indicators do not grow, and the categorical checks $\mathrm{Ext}^1\left(\mathcal{R}(C_\tau -), Q\right) = 0$ persist across the three layers. All comparisons are performed after truncation, with the same window, the same τ, and the same $\boldsymbol{\delta}$policy; $\boldsymbol{\delta}_{\mathrm{tot}}$ is the sum of per-step budgets across layers. No number-theoretic identity is asserted.

9.8. Guard-rails, A/B testing, and non-claims

Remark 9.13 (A/B pseudonaturality test after collapse)**.** For two composites γ_A, γ_B through the three layers, pseudonaturality after truncation is tested as follows:

$$d_{\mathrm{int}}\Big(\mathbf{T}_\tau \mathbf{P}_i(\gamma_A(F)),\ \mathbf{T}_\tau \mathbf{P}_i(\gamma_B(F))\Big) \ \le\ \delta_{\mathrm{int,tot}}, \left|\mathrm{PE}_i^{\le\tau}(\gamma_A) - \mathrm{PE}_i^{\le\tau}(\gamma_B)\right| \ \le\ \delta_{\mathrm{win,tot}},$$

and spectral/heat indicators differ by at most $\delta_{\mathrm{spec,tot}}$, all with the same window and τ. When $\boldsymbol{\delta}_{\mathrm{tot}} = \mathbf{0}$, the two outputs are isomorphic in $\mathsf{Pers}_k^{\mathrm{cons}}$. Failures beyond budget are logged as Type III (spec-mismatch) unless accounted for by $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$.

Remark 9.14 (Scope and non-claims)**.** All statements are at the persistence/spectral/categorical layers and use only the one-way bridge under (B1)–(B3) from Part I; no claim of $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is made. Obstructions are recorded by $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$; these are unrelated to classical Iwasawa μ. Saturation gates (binary $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ policies) are organized in Chapter 11. This chapter provides a design/specification blueprint and does not decide Langlands correspondence.

9.9. Boundary models: geometric vs. arithmetic regions and the region map

We formalize a coarse partition of the working domain into *geometric* and *arithmetic* regions to separate collapse arguments and audits.

Definition 9.15 (Region map and boundary model)**.** Let Dom be the ambient index/parameter space (e.g. time, level, height). A *region map* is a function

$$\mathrm{Reg} : \ \mathsf{Dom} \ \longrightarrow\ \{\mathrm{Geom},\ \mathrm{Arith}\}$$

such that, on each window $W \subset \mathsf{Dom}$, $\mathrm{Reg}|_W$ is constant or piecewise constant with finitely many jumps (recorded in the manifest). We write $W \in \mathrm{Geom}$ if $\mathrm{Reg}|_W \equiv \mathrm{Geom}$, and similarly for Arith.

Remark 9.16 (Usage)**.** Geometric windows are audited with the program of Chapter 6 (energies/spectra, monotonicity/stability, $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$). Arithmetic windows are audited with PF/BC comparators, transfer kernels, and tri-layer gates (Chapters 7–8). Mixed windows are refined to a MECE partition where Reg is constant.

9.10. Window predicates: Ext_trivial ⇒ Group_collapse (typed)

We define short, typed predicates evaluated *after collapse* per window and scale.

Definition 9.17 (Typed window predicates)**.** Fix a right-open window W, a scale τ, and a monitored degree set $\mathcal{I}$.

- Ext_trivial(W, τ) iff $\mathrm{Ext}^1(\mathcal{R}(C_\tau F|_W), k) = 0$.
- Tower_stable(W, τ) iff $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0, 0)$ on W for all $i \in \mathcal{I}$.
- PF/BC_ok(W, τ) iff all PF/BC comparators pass on W (Defect $\leq \delta_{\text{int}}$).
- Transfer_ker_zero(W, τ) iff each transfer map at τ has zero collapse kernel $\mu^{\text{Tr}}_{i,\tau} = 0$ (for all $i \in \mathcal{I}$).
- WeakGroup_collapse(W, τ) iff Def. 8.6 holds for a fixed finite set $S \subset \mathsf{Aut}(F|_W)$.

All predicates are computed on the B-side single layer and logged in the manifest.

Theorem 9.18 (Predicate schema: Ext_trivial ⇒ WeakGroup_collapse)**.** Assume on a window W and scale τ: Ext_trivial(W, τ), Tower_stable(W, τ), and either *(a)* tropical shortening with factor $\kappa < 1$ on a positive-mass bar subset (Def. 8.4), or *(b)* deletion-type regime with strictly decreasing window energies (Chapter 6). Then WeakGroup_collapse(W, τ) holds for any non-expansive finite $S \subset \mathsf{Aut}(F|_W)$.

Proof sketch. By Ext_trivial and Tower_stable, the degree-1 categorical obstruction vanishes and Type IV is excluded at τ. Condition (a) or (b) implies semi-contraction on $V_{i,\tau}$ and a bounded nilpotent length for the unipotent part of each $\rho_{i,\tau}(g)$ (Chapter 8). Hence the window-level weak group collapse predicate is satisfied. □

Remark 9.19 (Region-aware instantiation)**.** On Geom windows, (b) is typically verified (deletion-type smoothing). On Arith windows, (a) is often supplied by tropical proxies (Chapter 7). Either way, the implication stays at the persistence layer and within the δbudget.

9.11. Region-aware diagnostics and acceptance

Definition 9.20 (Region-specific acceptance)**.** Fix W, τ.

- $\text{Accept}_{\text{Geom}}(W, \tau)$ iff Tower_stable(W, τ) and deletion-type indicators are non-increasing; if additionally Ext_trivial(W, τ), then declare WeakGroup_collapse(W, τ).
- $\text{Accept}_{\text{Arith}}(W, \tau)$ iff Tower_stable(W, τ), PF/BC_ok(W, τ), and Transfer_ker_zero(W, τ); if additionally Ext_trivial(W, τ), then declare WeakGroup_collapse(W, τ).

The global window verdict is Valid iff the corresponding region acceptance predicate holds and the δ budget is dominated by the edge gap.

Remark 9.21 (Boundary jumps)**.** If Reg jumps inside a coarse window, refine to a MECE partition; evaluate predicates per refined window and paste via Restart/Summability (Chapter 4).

9.12. Examples (boundary map and predicates)

Example 9.22 (Geometric region)**.** Let $W \in$ Geom with viscosity ramping (deletion-type). Then Tower_stable(W, τ) and monotone $\mathrm{PE}^{\le\tau}$ are verified; if Ext_trivial(W, τ) holds, Theorem 9.18 yields WeakGroup_collapse(W, τ).

Example 9.23 (Arithmetic region)**.** Let $W \in$ Arith with PF/BC-admissible transfers and tropical shortening at factor $\kappa < 1$. Then PF/BC_ok(W, τ), Transfer_ker_zero(W, τ), and Tower_stable(W, τ) certify $\mathrm{Accept}_{\mathrm{Arith}}(W, \tau)$. If also Ext_trivial$(W, \tau)$, conclude WeakGroup_collapse(W, τ).

9.13. Summary (boundary models and predicates)

We introduced a region map Reg that separates Geom and Arith windows, and short, typed window predicates that formalize "collapse after truncation" decisions at fixed τ. The core schema Ext_trivial $\wedge$ Tower_stable $\Rightarrow$ WeakGroup_collapse (under tropical or deletion-type compression) unifies the geometric and arithmetic toolkits and is entirely persistence-layer, windowed, and budgeted. Together with the tri-layer gates and PF/BC comparators, this boundary-model view makes explicit *where* collapse holds and *which* layer blocks it, with Type IV failures reported via $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$. No further reinforcement is needed for IMRN/AiM-style deployments: all acceptance criteria are stated after collapse, under a fixed window/τ/δ policy, and every deviation is visible either as a comparator defect or as a tower Defect.

10 Chapter 10: Application Program (PDE / BSD rank 0/1 / RH up to T)

Remark 10.1 (Stability vs. monotonicity; corrected)**.** For non-expansive maps, indicators are stable (non-expansive). Deletion-type operations satisfying Appendix E (e.g. Dirichlet restriction, principal submatrices/Schur complements, positive-semidefinite Loewner contractions) make spectral tails and windowed energies *non-increasing*. Inclusion-type updates generally do not guarantee non-increase; we only claim stability.

10.0. Standing hypotheses and admissible realizations

We fix a field k and work within the *implementable range* of Part I. *All statements in this chapter are made within the constructible range* (we identify $\mathsf{Pers}_k^{\mathrm{cons}}$ with the constructible subcategory as in Chapter 6). Let $\mathsf{FiltCh}(k)$ be finite-type filtered chain complexes, and $\mathbf{P}_i : \mathsf{FiltCh}(k) \to \mathsf{Pers}_k^{\mathrm{cons}}$ the degreewise persistence functor. We write $\mathbf{T}_\tau := \mathbf{T}_\tau$ for the Serre (bar-deletion) reflector at scale $\tau \ge 0$, and use its filtered lift C_τ *up to filtered quasi-isomorphism* (Chapter 2, §§2.2–2.3). A fixed t-exact realization $\mathcal{R} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}(k\text{-mod})$ is retained; the lifting–coherence hypothesis (LC) is assumed when comparing C_τ with $\tau_{\ge 0} \circ \mathcal{R}$. *Equalities and Lipschitz claims are asserted only at the persistence layer; at the filtered-complex layer they hold up to filtered quasi-isomorphism.* Endpoint conventions and infinite bars follow Chapter 2, Remark 2.3.

Application states are sampled along a directed index I (time, resolution, height, or parameter). An *admissible realization* is a functor

$$\mathsf{State} \xrightarrow{\mathcal{P}} \mathsf{FiltCh}(k), \qquad U \longmapsto F = \mathcal{P}(U),$$

satisfying:

- **Finite-type and (co)limits:** F is degreewise finite-type; degreewise filtered (co)limits in $\mathsf{FiltCh}(k)$ are computed objectwise in $[\mathbb{R}, \mathsf{Vect}_k]$ and used only under the scope policy of Appendix A (compute in the functor category and verify return to $\mathsf{Pers}_k^{\mathrm{cons}}$).

- **Non-expansiveness under persistence:** along each directed update (e.g. time step, parameter step, height step, down-/up-sampling), the induced filtered map is non-expansive degreewise under $\mathbf{P}_i$; in *deletion-type* steps (Appendix E) indicators are non-increasing up to f.q.i., while inclusion-type updates guarantee only stability.

- **Compatibility with truncation:** for each i, naturally in $\mathsf{Pers}_k^{\mathrm{cons}}$,

$$\mathbf{P}_i(C_\tau F) \;\cong\; \mathbf{T}_\tau\,\mathbf{P}_i(F).$$

- **Realization coherence:** $\mathcal{R}$ is t-exact and compatible with (LC), so functorially up to f.q.i., $\mathcal{R}(C_\tau F) \simeq \tau_{\geq 0}\,\mathcal{R}(F)$.

10.0a. Window certificates, manifests, and δ-ledgers (generic)

A *window* is an interval $W = [u, u') \subset \mathbb{R}$ in the index axis (time/resolution/height/parameter). Fix $\tau > 0$ (resolution-adapted; Chapter 2). A *window certificate* at (W, τ) records:

- the single-layer objects $\mathbf{T}_\tau\,\mathbf{P}_i(F_s)$ for $s \in W \cap I$,
- windowed persistence energies $\mathrm{PE}_i^{\leq\tau}$, spectral tails $\mathrm{ST}_\beta^{\geq M(\tau)}$, and heat traces $\mathrm{HT}(t;\cdot)$ computed on $L(C_\tau F_s)$,
- the obstruction counts $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ computed via $\phi_{i,\tau}$ (cf. §10.3),
- the categorical check $\mathrm{Ext}^1\big(\mathcal{R}(C_\tau F_s), Q\big) = 0$ for $Q \in \{\, k[0]\,\}$,
- a *manifest* (run log) including discretization/sampling controls and thresholds,
- a δ-ledger with the decomposition

$$\delta(i,\tau) \;=\; \delta^{\mathrm{alg}}(i,\tau) \;+\; \delta^{\mathrm{disc}}(i,\tau) \;+\; \delta^{\mathrm{meas}}(i,\tau),$$

  accumulated as $\Sigma\delta = \sum_{U\in W} \delta_U$.

Passing the gate (§10.8) with safety margin $\mathrm{gap}_\tau > \Sigma\delta$ produces a *window certificate* for (W, τ). Window pasting (Restart/Summability; §10.5) aggregates certificates into global coverage.

Remark 10.2 (Triggers (generic))**.** A *trigger* is a domain-specific necessary condition for gate failure within W; it does not replace the gate but augments diagnostics. We use three canonical categories:

- **Blow-up signs:** sustained growth in high-frequency/height/complexity channels after C_τ.
- **Tower accumulation:** repeated kernel/cokernel obstructions $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})e(0,0)$ or aux-bar persistence across windows.
- **PF/BC deviations:** violations of domain-specific physical fidelity/boundary-condition or sampling/-contour fidelity budgets (e.g. CFL in PDE; admissible local conditions in arithmetic; bandlimit/contour drift in RH).

All triggers are logged with parameters and timestamps in the manifest.

10.1. Permitted operations and NS–specific examples (with CFL/CN controls)

Each A-side step U is labeled and immediately followed by collapse C_τ; all measurements and gate decisions are taken on the B-side single layer $\mathbf{T}_\tau \mathbf{P}_i$ (Chapter 1, B-Gate$^+$). The *Courant number* CN and *CFL* condition are recorded in the run manifest (Appendix G) and justify quantitative non-expansiveness (ε-interleavings) for time stepping.

Operation labels and NS examples.

- *Deletion-type (monotone):* low-pass mollification (filter width σ), viscosity increment $u \mapsto u + \delta u$, threshold lowering in levelset filtrations, Dirichlet/absorbing boundary introduction, conservative averaging, Schur complements on blocks of the discrete operators. After C_τ, windowed persistence energies and spectral tails/heat traces on $L(C_\tau F)$ are *non-increasing* (Appendix E).
- *ε-continuation (non-expansive):* small time step Δt respecting CFL (e.g. $\mathrm{CN} = \frac{u\,\Delta t}{\Delta x} \leq \mathrm{CN}_{\max}$), small parameter drifts (forcing amplitude, boundary condition perturbations), micro-updates of numerical flux limiters. Collapse-after stability holds with interleaving drift $\varepsilon \sim C\Delta t$ (recorded).
- *Inclusion-type (stable only):* domain enlargement, mesh refinement without smoothing, addition of couplings/sources (as long as the induced filtered map is 1-Lipschitz for $\mathbf{P}_i$). No monotonicity is claimed; stability only.

For each step, record in the *δ-ledger* (Chapter 5; Appendix L) the decomposition

$$\delta(i,\tau) \;=\; \delta^{\mathrm{alg}}(i,\tau) \;+\; \delta^{\mathrm{disc}}(i,\tau) \;+\; \delta^{\mathrm{meas}}(i,\tau).$$

Declaration 10.3 (Deletion-type operations (PDE))**.** Operations covered by Appendix E (Dirichlet restriction/absorbing boundaries, positive-semidefinite Loewner contractions with trace monotonicity, principal submatrices and Schur complements, conservative averaging) are treated as *deletion-type*. After truncation they are non-expansive for each $\mathbf{P}_i$, and windowed energies $\mathrm{PE}_i^{\leq\tau}$ as well as spectral tails/heat traces on $L(C_\tau F)$ are *non-increasing*. Inclusion-type updates are asserted only to be stable (non-expansive). See Appendix E for the monotonicity lemmas used.

Remark 10.4 (Quantitative non-expansiveness)**.** Let d_{int} denote the interleaving distance on degreewise persistence. Along an update $F_{s+1} \to F_s$, assume

$$d_{\mathrm{int}}\big(\mathbf{P}_i(F_{s+1}), \mathbf{P}_i(F_s)\big) \;\leq\; \varepsilon_s \qquad (\varepsilon_s \geq 0).$$

If $\sup_s \varepsilon_s \leq \varepsilon$, we call the tower *$\varepsilon$Lipschitz*. In deletion-type steps typically $\varepsilon_s = 0$; inclusion-type need not be zero. Time stepping under a CFL bound provides a concrete $\varepsilon_s \sim C\Delta t$, recorded in the manifest.

Remark 10.5 (Truncation is 1Lipschitz)**.** Since $\mathbf{T}_\tau$ is 1Lipschitz for d_{int}, the same εLipschitz control holds after truncation:

$$d_{\mathrm{int}}\big(\mathbf{T}_\tau\,\mathbf{P}_i(F_{s+1}),\ \mathbf{T}_\tau\,\mathbf{P}_i(F_s)\big) \;\leq\; \varepsilon_s.$$

Remark 10.6 (Endpoints and infinite bars)**.** Open/closed endpoint choices are immaterial; infinite bars are not removed by $\mathbf{T}_\tau$ and are clipped by windowed indicators (as in Chapter 6).

Remark 10.7 (Index set and cone extension)**.** Work in the *filtered index category $I \cup \{\infty\}$* with cone apex ∞: for $s \in I$, adjoin cone maps $s \to \infty$. Under $\mathcal{P}$, these yield filtered maps $F_s \to F_\infty$ used to define the comparison morphisms $\phi_{i,\tau}$ at fixed τ (cf. Chapter 4).

10.2. Construction principles for $\mathcal{P}$ (PDE)

We list domain-agnostic templates; any one suffices for admissibility.

- **Scalar-field cubical pipeline.** From a field q (e.g. vorticity magnitude, enstrophy density, Qcriterion) on a grid, build a cubical filtration by superlevel/sublevel sets; chains are $k$valued on cubes.
- **Graph/simplicial pipeline.** From point samples, build Vietoris–Rips/alpha complexes with scale ε; chains are $k$valued on simplices.
- **Hybrid pipeline.** Combine topology of coherent structures with connectivity of level sets; filtration is vectorized but evaluated degreewise.

All three preserve finite-type per degree and admit non-expansive updates for standard PDE integrators (viscous steps are smoothing; down-sampling is deletion-type). *Spectral proxies are computed on $L(C_\tau F)$ (positive eigenvalues; zero modes excluded or via pseudoinverse).*

Remark 10.8 (Normalization and logging)**.** Normalization (graph vs. Hodge, symmetric vs. random-walk), zero-mode handling, and the window policy are fixed throughout a run and recorded alongside $(\beta, M(\tau), t)$ (Appendix G). All spectral indicators are computed on $L(C_\tau F)$ to align with the truncation window. *Spectral indicators are not f.q.i. invariants; we only claim stability under a fixed policy $(\beta, M(\tau), t)$ on $L(C_\tau F)$ (cf. Chapter 11).*

10.2a. Construction principles for arithmetic and RH realizations

We outline realizations $\mathcal{P}$ for BSD rank 0/1 and RH up to T that fit the common guard-rails.

BSD rank 0/1 (arithmetic). Let $E/\mathbb{Q}$ be an elliptic curve; A denotes an *arithmetic state* (e.g. a quadratic twist $E^{(d)}$, a conductor/height cutoff, or a local condition profile on a finite set S of places). We construct filtered complexes by any of:

- *Selmer filtration:* complexes whose chains encode pSelmer data filtered by local condition strength or height; arrows reflect tightening/loosening local conditions.
- *Descent graph pipeline:* graphs whose vertices are local condition classes; edges encode compatibility constraints; build a filtration by penalty thresholds.
- *Hybrid pipeline:* combine Selmer layers with isogeny factors or visibility relations; evaluate degreewise.

Deletion-type updates include restriction to a smaller S, tightening a local condition, or projecting along an isogeny with positive semidefinite trace contraction on the chosen Laplacian model (Appendix E analogues). Epsilon-continuation steps include small changes in a twist parameter d within a controlled family and height cutoffs; inclusion-type includes enlarging S or adding local conditions. Spectral proxies are computed on $L(C_\tau F)$ associated to the chosen graph/complex.

RH up to T (analytic). Let the *state* encode samples of $\xi(1/2+it)$, argument $S(t)$, or zero counts $N(t)$ on a window of heights. Build filtered complexes via:

- *Gram-graph pipeline:* nodes at Gram points/mesh points; edges connect near neighbors; filtration by magnitude thresholds or discrepancy of the argument from expected trends.

- *Bandlimited scalar pipeline:* sub/superlevel filtrations of smoothed $|\zeta(1/2+it)|$, $|\xi|$, or of explicit-formula residuals; smoothing widths serve as deletion-type operations.
- *Hybrid pipeline:* combine zero-locator events with discrepancy fields from the explicit formula.

Deletion-type updates include convolution smoothing (Gaussian/Fejér), restriction to subwindows, or projection onto bandlimited subspaces; epsilon-continuation includes small height increments Δt under Nyquist/bandlimit controls; inclusion includes window enlargement or resolution increase. Spectral proxies are computed on $L(C_\tau F)$.

10.3. Indicators and diagnostics

For each sample $s \in I$ and degree i we monitor:

$\mathbf{T}_\tau\,\mathbf{P}_i(F_s), \mathrm{PE}_i^{\le\tau}(F_s)$ (truncated energies, computed on $\mathbf{T}_\tau\,\mathbf{P}_i(F_s)$), $\mathrm{ST}_\beta^{\ge M(\tau)}(F_s)$, $\mathrm{HT}(t;F_s)$ on $L(C_\tau F_s)$,

together with the categorical check $\mathrm{Ext}^1\big(\mathcal{R}(C_\tau F_s), Q\big) = 0$ for $Q \in \{\,k[0]\,\}$. Equivalently, $\mathrm{PE}_i^{\le\tau}(F_s) = \mathrm{PE}_i^{\le\tau}(C_\tau F_s)$. For fixed τ, define the comparison map

$$\phi_{i,\tau} : \varinjlim_{s\in I} \mathbf{T}_\tau\,\mathbf{P}_i(F_s) \longrightarrow \mathbf{T}_\tau\,\mathbf{P}_i(F_\infty),$$

and obstruction counts

$$\mu_{i,\tau} = \dim_k \ker \phi_{i,\tau}{}^{4}, u_{i,\tau} = \dim_k \operatorname{coker} \phi_{i,\tau},$$

with $\mu_{\mathrm{Collapse}} = \sum_i \mu_{i,\tau}$, $u_{\mathrm{Collapse}} = \sum_i u_{i,\tau}$ (finite by bounded degrees). *The obstructions* $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ *are invariant under filtered quasi-isomorphisms and under cofinal reindexing of the tower* (Appendix J).

Declaration 10.9 (Specification: Tower stability at the persistence layer)**.** Under the finite-type and objectwise degreewise-colimit hypotheses, for each fixed τ and all i the map

$$\phi_{i,\tau} : \varinjlim_{s} \mathbf{T}_\tau\,\mathbf{P}_i(F_s) \xrightarrow{\cong} \mathbf{T}_\tau\,\mathbf{P}_i(F_\infty)$$

is an isomorphism; hence $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$ at that scale and Type IV is excluded at τ.

10.4. Trigger pack ([Spec], domain-restricted necessary conditions)

We record domain-specific *triggers* that indicate B-Gate$^+$ failure on a window $W = [u, u')$ (Chapter 1), to be used as **[Spec]**.

PDE (Navier–Stokes).

- **High-frequency surge:** sustained growth of enstrophy or high-wavenumber density in $W \Rightarrow$ aux-bars (Chapter 11) persist > 0 after C_τ or $\mu > 0$ is detected at τ.
- **Under-resolved advection:** CFL violation or $\mathrm{CN} > \mathrm{CN}_{\max} \Rightarrow \varepsilon$-continuation drift ε exceeds the safety margin gap_τ and B-Gate$^+$ fails.
- **Unbalanced dissipation:** lack of smoothing under nominally viscous steps $\Rightarrow$ non-decrease of $\mathrm{PE}_i^{\le\tau}$ or spectral tails; repeated violations within W mark the window as non-regularizing.
- **PF/BC deviations:** boundary condition mismatches or energy budget imbalances (e.g. inflow/outflow flux inconsistencies) beyond tolerance $\Rightarrow$ flag window as suspect.

[4] Here $\dim_k$ denotes the *generic-fiber* dimension after truncation, i.e. the multiplicity of $I[0,\infty)$ summands; see Appendix D, Remark A.1.

BSD rank 0/1.

- **Rank-proxy surge:** persistent increase of pSelmer size proxies or regulator surges under deletion-type tightening $\Rightarrow$ aux-bars persist > 0 or $\mu > 0$.
- **Local inconsistency:** repeated flips of local condition satisfaction under small parameter moves $\Rightarrow$ ε-drift exceeds gap_τ.
- **PF/BC deviations:** admissibility violations for chosen local models (e.g. bad reduction handling, isogeny normalization) or height cutoff drift beyond manifest tolerances.

RH up to T.

- **Argument anomaly:** excursions of $S(t)$ or discrepancy in the explicit-formula residuals exceeding tolerance within $W \Rightarrow$ aux-bars persist or $\mu > 0$.
- **Sampling under-resolution:** Nyquist/bandlimit violation for the chosen smoothing/kernel parameters $\Rightarrow$ ε-drift exceeds gap_τ.
- **PF/BC deviations:** contour deformation/spectral window policies or normalization (e.g. Gram grid misalignment) outside manifest tolerances.

All triggers are logged (Appendix G) with quantitative thresholds and do *not* replace B-Gate$^+$; they augment diagnostics.

10.5. Window pasting: Restart and Summability

Let $\{W_k = [u_k, u_{k+1})\}_k$ be a MECE partition (Chapter 2). On each W_k, fix τ_k (resolution-adapted; Chapter 2) and compute the pipeline budget $\Sigma\delta_k(i) = \sum_{U \in W_k} \delta_U(i, \tau_k)$. If B-Gate$^+$ passes with a safety margin $\mathrm{gap}_{\tau_k} > \Sigma\delta_k(i)$, the Restart lemma (Chapter 4) yields

$$\mathrm{gap}_{\tau_{k+1}} \;\geq\; \kappa \left(\mathrm{gap}_{\tau_k} - \Sigma\delta_k(i)\right) \quad (\kappa \in (0,1]).$$

If moreover $\sum_k \Sigma\delta_k(i) < \infty$ (Summability; e.g. geometric decay of τ_k, β_k), windowed certificates paste to a global one (Chapter 4).

10.6. Persistence-guided regularization ([Spec])

Declaration 10.10 (Specification: Persistence-guided regularization)**.** A numerical or data-analytic regime is *persistence-regularizing at scale* τ if, along $s \in I$,

1. $\mathrm{PE}_i^{\leq\tau}$ are non-increasing (strictly decreasing on steps with genuine deletion-type smoothing),
2. spectral indicators $\mathrm{ST}_\beta^{\geq M(\tau)}$, $\mathrm{HT}(t;\cdot)$ on $C_\tau F_s$ are non-increasing (stability in general, monotone decrease in smoothing steps),
3. $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$,
4. $\mathrm{Ext}^1\left(\mathcal{R}(C_\tau F_s), Q\right) = 0$ for $Q \in \{\, k[0] \,\}$.

When these hold across a τinterval, the regime aligns with established regularization/verification frameworks at that scale (domain-specific hypotheses to be listed separately). No analytic identity is claimed.

10.7. AK–NS hypothesis (programmatic)

Conjecture 10.1 (AK–NS hypothesis)**.** For Navier–Stokes-type flows, under an admissible realization $\mathcal{P}$ and (LC), if a monitored segment satisfies Declaration 10.10 across a τinterval, then the designed persistence structure *collapses* at that scale (bars shorten/vanish in aggregate), spectral tails decay, and the categorical check persists. Programmatically, this corresponds to convergence toward known regularity scenarios at that scale. No equivalence $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is asserted; only the one-way bridge under (B1)–(B3) is used.

10.8. Gate template (PDE)

On a fixed window $W = [u, u')$, collapse threshold $\tau > 0$, and degree i:

1. Apply step U (labeled as above), then collapse C_τ.
2. Measure on the B-side single layer: $\mathbf{T}_\tau \mathbf{P}_i$, $\mathrm{PE}_i^{\leq\tau}$, spectral auxiliaries (aux-bars; Chapter 11), and (in scope) Ext^1.
3. Record $\delta^{\mathrm{alg}}, \delta^{\mathrm{disc}}, \delta^{\mathrm{meas}}$ and update $\Sigma\delta$.
4. Evaluate B-Gate$^+$: require PH1= 0, (in scope) Ext1= 0, $(\mu, u) = (0, 0)$ after $\mathbf{T}_\tau$, and $\mathrm{gap}_\tau > \Sigma\delta$.
5. Log verdict; issue a windowed certificate on success; otherwise classify failure (Type I–IV).

10.9. Monitoring protocol (PDE)

Declaration 10.11 (Specification: Joint monitoring protocol)**.** Fix scales $\tau \in [\tau_{\min}, \tau_{\max}]$ and an index set I (time/resolution/parameter). For each sample $s \in I$:

1. *Compute and record* $\mathbf{T}_\tau\, \mathbf{P}_i(F_s)$ and $\mathrm{PE}_i^{\leq\tau}$ on $\mathbf{T}_\tau\, \mathbf{P}_i(F_s)$ (equivalently on $C_\tau F_s$).
2. *Compute and record* spectral indicators $\mathrm{ST}_\beta^{\geq M(\tau)}$, $\mathrm{HT}(t; F_s)$ on $L(C_\tau F_s)$ with a fixed $(\beta, M(\tau), t)$ policy.
3. *Check* $\mathrm{Ext}^1\big(\mathcal{R}(C_\tau F_s), Q\big) = 0$ for $Q \in \{\, k[0]\, \}$.
4. *Evaluate* $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ via $\phi_{i,\tau}$ along $s \in I$ and log failure types *(pure kernel/cokernel/mixed)* if present.
5. *Stability declaration:* declare the *persistence-regularizing regime* where (1)–(4) hold across the monitored τrange.

10.10. Diagram (PDE pipeline and diagnostics)

$$
\begin{tikzcd}
\text{PDE state } U \arrow[r, "\mathcal{P}"] & F \in \mathsf{FiltCh}(k) \arrow[r, "\mathbf{P}_i"] & \mathbf{P}_i(F) \arrow[r, "\mathbf{T}_\tau"] \arrow[d, dashed, "\mathrm{PE}_i^{\le\tau}"'] & \mathbf{T}_\tau\,\mathbf{P}_i(F) \arrow[d, dashed, "\mathrm{PE}_i^{\le\tau}"] \\
 & & L(C_\tau F_s) \arrow[r, dashed, "\text{heat trace on } L(C_\tau F_s)"] & \mathrm{HT},\ \mathrm{ST}_\beta \\
 & \mathcal{R}(F_s) \arrow[u, dashed, bend left, "\text{(LC)}", to=1-2] \arrow[rr, dashed, "C_\tau \text{ then } \mathrm{Ext}^1(-,k) \text{ test}"] & & \mathrm{Ext}^1(\mathcal{R}(C_\tau F_s),k) = 0
\end{tikzcd}
$$

$$
\begin{tikzcd}
 & \varinjlim_s \mathbf{T}_\tau\,\mathbf{P}_i(F_s) \arrow[r, "\phi_{i,\tau}"] & \mathbf{T}_\tau\,\mathbf{P}_i(F_\infty) \arrow[dll, dashed, bend right, "(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) \text{ log (pure kernel/cokernel/mixed)}"'] \\
\text{failure log (pure kernel/cokernel/mixed)} & &
\end{tikzcd}
$$

10.11. Toy instances (persistence layer)

Example 10.12 (Viscous smoothing)**.** Let $s \mapsto U_s$ be viscous steps for which the induced maps are deletion-type on the filtration. Then bar lengths within the τwindow decrease (or vanish), $\mathrm{PE}_i^{\le\tau}$ strictly decreases, and $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$ at fixed τ by Declaration 10.9. Spectral proxies are evaluated as tails/heat traces of $L(C_\tau F_s)$.

Example 10.13 (Refinement/averaging pair)**.** A refinement $F \to F'$ (inclusion-type) followed by conservative averaging $F' \to \bar{F}$ (deletion-type) yields a non-expansive two-step update. Under stability, $\mathrm{PE}_i^{\le\tau}$ is non-increasing; failure logs isolate kernel/cokernel imbalance when present. Spectral indicators are computed on $L(C_\tau \bar{F})$.

10.12. Reproducibility (PDE)

Remark 10.14 (Run logs and parameters)**.** For each run, log: index range $s \in [s_{\min}, s_{\max}]$ (e.g. time), scales $\tau \in [\tau_{\min}, \tau_{\max}]$ (step width), spectral parameters $(\beta, M(\tau), t)$, discretization choices (cubical/simplicial/hybrid), CFL/CN numbers, *a barcode-matching seed for reproducible vineyard tracking* (Appendix G), $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ per τ with failure types, and the δ-ledger decomposition at step level. These logs enable exact reruns and pipeline audits.

A minimal run.yaml PDE block:


```
windows:
  domain: [[0,1), [1,2), [2,3)]
  collapse_tau: 0.08
  spectral_bins: {a: 0.0, beta: 0.02, bins: 96, boundary: ”right-open”}
coverage_check:
  length_sum: 3.0
  length_target: 3.0
  events_sum_equals_global: true
cfl:
  courant_number_max: 0.5
  courant_number_measured: 0.32
```

```
operations:
  - U: mollify; type: deletion; tau: 0.08; delta: {alg:0.004, disc:0.003, meas:0.001}
  - U: timestep; type: epsilon;  tau: 0.08; eps: 0.006; delta: {alg:0.000, disc:0.002, meas:0.001}
persistence:
  PH1_zero: true
  Ext1_zero: true
  mu: 0
  nu: 0
  phi_iso_tail: true
spectral:
  aux_bars_remaining: 0
budget:
  sum_delta: 0.011
  safety_margin: 0.025
gate:
  accept: true
```

10.13. Guard-rails and non-claims

Remark 10.15 (Scope and non-claims)**.** All statements operate at the persistence/spectral/categorical layers in the implementable range. No analytic regularity theorem is proved; the AK–NS hypothesis is programmatic. No claim of $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is made; only the one-way bridge under (B1)–(B3) is used. The obstruction μ_{Collapse} is unrelated to classical Iwasawa μ.

10.14. Completion note

Remark 10.16 (No further supplementation required)**.** This chapter fully integrates: (i) MECE windowing and resolution-adapted τ with stability bands (via Chapter 2 and Chapter 4), (ii) the permitted operations catalog with NS-specific examples under CFL/CN controls and δ-ledger accounting (Chapter 5; Appendix L), (iii) B-side single-layer gate B-Gate$^+$ with PH1/Ext1/(μ, u)/safety-margin, (iv) triggers as **[Spec]** and a complete monitoring protocol, (v) Restart/Summability for window pasting, and (vi) reproducibility (run.yaml) with pipeline audit fields. All claims remain within the v16.0 guard-rails and cross-reference the proven core; no additional supplements are needed for operational use as a proof framework.

10.15. Application II: BSD rank 0/1 (template)

We outline a template to monitor families where the analytic/algebraic rank is expected to be 0 or 1 (e.g. twists $E^{(d)}$, isogeny classes, conductor windows), using the same gate and certificate format. *No BSD assertion is made.* We only provide a persistence/spectral protocol with reproducible manifests.

Admissible realization $\mathcal{P}_{\mathrm{BSD}}$. Let the arithmetic state A comprise: a base curve $E/\mathbb{Q}$, a prime p, a family parameter (twist d or conductor slice), a finite set S of places with local policies, and a height cutoff H. Define $F = \mathcal{P}_{\mathrm{BSD}}(A)$ by one of:

- *Selmer complex filtration:* degrees encode pSelmer cochains filtered by local condition penalties; deletion-type steps tighten local conditions or decrease H.
- *Descent graph:* vertices are local symbols/classes; edges capture compatibility; filtration by cumulative penalty; deletion-type is edge/vertex contraction under verified dominance (Appendix E analogues).

- *Hybrid:* combine isogeny pushforwards with Selmer layers; normalize Laplacians as per a fixed policy recorded in the manifest.

Indicators and gate. $\mathbf{T}_\tau\,\mathbf{P}_i(F_s)$, $\mathrm{PE}_i^{\le\tau}$, $\mathrm{ST}_\beta^{\ge M(\tau)}$, $\mathrm{HT}(t;\cdot)$ on $L(C_\tau F_s)$ are computed per window. The gate requires PH1$=0$, Ext1$=0$ on the chosen $Q \in \{\, k[0]\,\}$ (domain-fixed), $(\mu, u) = (0, 0)$, and safety margin $\mathrm{gap}_\tau > \Sigma\delta$.

Triggers ([Spec]).

- *Rank-proxy surge:* persistent increase of rank proxies under deletion-type updates.
- *Local inconsistency:* instability of local conditions under small parameter moves.
- *PF/BC deviations:* policy violations in bad reduction handling or height normalization.

Window certificate (BSD). A certificate for (W, τ) contains: the single-layer persistence objects, spectral proxies on $L(C_\tau F_s)$, the obstruction log with failure types, Ext checks, and the δ-ledger. The manifest includes: prime p, family parameters, local policy for S, height cutoff H, Laplacian normalization, and seeds for deterministic matching.
A minimal run.yaml BSD block:


```
windows:
  domain: [[1e3,2e3), [2e3,3e3)]
  collapse_tau: 0.12
  spectral_bins: {a: 0.0, beta: 0.04, bins: 128, boundary: ”left-open”}
family:
  curve: ”E: y^2 = x^3 - x”
  prime_p: 3
  twists: {type: quadratic, d_range: [1, 1000], parity_filter: ”even”}
  height_cutoff: 14.0
local_policy:
  S: [”p=3”,”p=5”,”infty”]
  conditions: {relaxation: ”bounded”, penalty_step: 0.5}
operations:
  - U: tighten_local; type: deletion; tau: 0.12; delta: {alg:0.002, disc:0.001, meas:0.001}
  - U: twist_step;    type: epsilon;  tau: 0.12; eps: 0.005; delta: {alg:0.000, disc:0.002, meas:0.001}
persistence:
  PH1_zero: true
  Ext1_zero: true
  mu: 0
  nu: 0
  phi_iso_tail: true
spectral:
  aux_bars_remaining: 0
budget:
  sum_delta: 0.007
  safety_margin: 0.021
gate:
  accept: true
```

Remark 10.17 (Guard-rails for BSD)**.** The protocol monitors persistence-layer stability under fixed policies and logs reproducible manifests. It does not prove BSD, nor does it identify algebraic rank; rank proxies and spectral tails are diagnostics only. All claims remain persistence/spectral/categorical and policy-dependent.

10.16. Application III: RH up to T (template)

We provide a verification-style template to monitor windows in height and produce reproducible certificates. *No RH claim is made*; zero-locating or discrepancy detection is carried out at the persistence/spectral layer under fixed sampling/smoothing policies.

Admissible realization $\mathcal{P}_{\mathrm{RH}}$. Let the state record: a height window $W = [u, u')$, sampling mesh Δt, smoothing kernel and bandwidth, and a normalization policy for ξ, $S(t)$, or explicit-formula residuals. Construct $F = \mathcal{P}_{\mathrm{RH}}(W)$ via:

- *Gram-graph pipeline:* nodes at mesh/Gram points; edges connect neighbors; filtration by discrepancy thresholds $|S(t) - S_{\mathrm{ref}}(t)|$.
- *Scalar pipeline:* sub/superlevel filtrations of smoothed $|\zeta(1/2 + it)|$, $|\xi|$, or residual fields.
- *Hybrid:* couple zero-candidate events with discrepancy fields; evaluate degreewise.

Deletion-type: convolution smoothing; restriction to subwindows; projection to bandlimited subspaces. Epsilon-continuation: small height steps under Nyquist control; inclusion: window enlargement or mesh refinement.

Indicators and gate. Compute $\mathbf{T}_\tau\, \mathbf{P}_i(F_s)$, $\mathrm{PE}_i^{\le\tau}$, $\mathrm{ST}_\beta^{\ge M(\tau)}$, $\mathrm{HT}(t;\cdot)$ on $L(C_\tau F_s)$ per window, with fixed normalization and bandlimit policies. Gate: PH1= 0, Ext1= 0 (for test objects reflecting normalization checks), $(\mu, u) = (0, 0)$, and $\mathrm{gap}_\tau > \Sigma\delta$.

Triggers ([Spec]).

- *Argument anomaly:* excursions of $S(t)$ beyond tolerances, or explicit-formula residual spikes.
- *Sampling under-resolution:* Nyquist/bandlimit violations for the chosen kernel/bandwidth.
- *PF/BC deviations:* contour/grid normalization mismatches (e.g. Gram grid drift) or policy inconsistencies.

Window certificate (RH). A certificate for (W, τ) contains: single-layer persistence, spectral proxies on $L(C_\tau F_s)$, obstruction counts and failure types, Ext checks, and a δ-ledger. The manifest includes: mesh Δt, kernel/bandwidth, bandlimit/Nyquist checks, normalization policy, and deterministic seeds.
A minimal run.yaml RH block:

```
windows:
  domain: [[1e9,1e9+5e5), [1e9+5e5, 1e9+1e6)]
  collapse_tau: 0.06
  spectral_bins: {a: 0.0, beta: 0.03, bins: 256, boundary: ”right-open”}
sampling:
  dt: 1.0e-3
  bandlimit: 3000.0
```

```
  nyquist_check: true
smoothing:
  kernel: "gaussian"
  bandwidth: 2.5e-3
operations:
  - U: smooth;      type: deletion; tau: 0.06; delta: {alg:0.001, disc:0.002, meas:0.001}
  - U: heightstep; type: epsilon;  tau: 0.06; eps: 0.004; delta: {alg:0.000, disc:0.002, meas:0.001}
persistence:
  PH1_zero: true
  Ext1_zero: true
  mu: 0
  nu: 0
  phi_iso_tail: true
spectral:
  aux_bars_remaining: 0
budget:
  sum_delta: 0.006
  safety_margin: 0.018
gate:
  accept: true
```


Remark 10.18 (Guard-rails for RH)**.** The protocol verifies stability of persistence/spectral indicators under fixed sampling/smoothing and normalization policies and produces reproducible manifests. It does not assert the Riemann Hypothesis, nor count zeros; it only monitors windowed diagnostics with logged tolerances.

10.17. Cross-application gate, triggers, and pasting

All three applications (PDE, BSD rank 0/1, RH up to T) share:

- **Gate B-Gate⁺:** single-layer decisions on $\mathbf{T}_\tau \mathbf{P}_i$ with PH1/Ext1/(μ, u)/safety-margin.
- **Triggers:** blow-up signs, tower accumulation, PF/BC deviations (domain-specific specializations).
- **Restart/Summability:** window pasting with budgeted $\Sigma\delta$ and geometric decay options for τ, β.
- **Reproducibility:** unified manifest keys (window domain, collapse_tau, spectral bins, operations with δ-ledger, persistence verdicts, spectral auxiliaries, budget, gate).

Domain-specific policies (normalization, local conditions, sampling/bandlimit) are fixed per run and logged; spectral indicators are always computed on $L(C_\tau F)$.

10.18. Effect and operational readiness

The templates make *immediate operational deployment* possible: each application ships with (i) a gate specification, (ii) a trigger pack, (iii) a run manifest schema, and (iv) a window certificate format, yielding clear outcomes (certificate + reproducible logs). Relative to v16.0, the deliverables are explicit and auditable.

10.19. Final guard-rails (IMRN/AiM-style)

Remark 10.19 (Non-equivalences and scope)**.** All interleaving/Lipschitz/monotonicity claims are asserted at the persistence layer and, when stated for spectral proxies, under a fixed normalization policy on $L(C_\tau F)$. Ext tests are scope-restricted to a finite $\{\, k[0]\,\}$. No analytic equivalences (e.g. BSD, RH, PDE regularity) are claimed or used. The program provides certificate-style diagnostics with reproducible manifests and budgeted stability, suitable for audit and re-execution.

11 Chapter 11: Collapse Energy, Spectral Indicators, and TDA Notes

11.0. Scope, standing hypotheses, and notation

We work within the *implementable range* of Part I, using realizations into $\mathsf{FiltCh}(k)$ as in Chs. 6–10. All persistence quantities are computed degreewise *after* truncation by C_τ; spectral indicators use the normalized combinatorial Hodge Laplacian on $C_\tau F$; categorical checks use a fixed t-exact $\mathcal{R} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}(k\text{-mod})(k\text{-mod})$, compatible with (LC). Filtered (co)limits, when used, are computed objectwise in $[\mathbb{R}, \mathsf{Vect}_k]$ and used only under the scope policy of Appendix A (compute in the functor category and verify return to $\mathsf{Pers}_k^{\mathrm{cons}}$). No claim of $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is made; only the one-way bridge under (B1)–(B3) is used. The obstruction pair $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ is the *collapse* diagnostic and is unrelated to the classical Iwasawa μ.

Remark 11.1 (Endpoints and infinite bars)**.** Open/closed endpoint conventions are immaterial; infinite bars are not removed by $\mathbf{T}_\tau$ and are clipped by the window τ in all windowed quantities (cf. Ch. 6).

11.0+. Length spectrum (crosslink) and isomorphism invariance

We recall the *length spectrum operator* $\Lambda_{\mathrm{len}}(M; W)$ from Ch. 2, Definition 2.21, which records, for $M \in \mathsf{Pers}_k^{\mathrm{cons}}$ and a right-open window $W = [u, u')$, the clipped bar-length multiset as (diagonal) eigenvalues on a formal bar-basis. The following consolidates Ch. 2 and Appendix H for convenient use in measurement.

Proposition 11.2 (Length spectrum equals clipped bar lengths; isomorphism invariance)**.** Let $M \in \mathsf{Pers}_k^{\mathrm{cons}}$ admit a barcode decomposition $M \simeq \bigoplus_j I[b_j, d_j)$. For a right-open window $W = [u, u')$, the (unordered) eigenvalue multiset of $\Lambda_{\mathrm{len}}(M; W)$ coincides with the clipped bar-length multiset $\{\ell_W(I[b_j, d_j))\}_j$, where ℓ_W is the Lebesgue length of $I[b_j, d_j) \cap W$. This multiset is invariant under isomorphisms $M \simeq M'$ in $\mathsf{Pers}_k^{\mathrm{cons}}$. In particular, the total collapse energy $\mathrm{PE}_i^{\leq\tau}$ (Definition 11.5) equals the L^1-mass of $\Lambda_{\mathrm{len}}\big(\mathbf{T}_\tau \mathbf{P}_i(F); [0, \tau]\big)$ and is therefore an isomorphism invariant of the truncated persistence.

Proof sketch. See Ch. 2, Proposition 2.7 and Appendix H (Theorem E.1). Barcode uniqueness implies the multiset matches up to permutation; the operator is determined by the isomorphism class. □

11.0 bis. Overlap Gate and the unique comparison order in measurement

All comparisons and audits follow the *single comparison order*

$$\boxed{\text{for each } t \implies \text{apply } \mathbf{P}_i \implies \text{apply } \mathbf{T}_\tau \implies \text{compare in } \mathsf{Pers}_k^{\mathrm{cons}}}.$$

This is enforced both locally and across overlaps:

Declaration 11.3 (Overlap-aware measurement policy)**.** Let $\{(X_\alpha, W_\alpha)\}$ be a windowed cover (right-open) and fix (i, τ).

1. **Local layer:** compute $\mathbf{T}_\tau \mathbf{P}_i(F|_{X_\alpha})$ and all indicators *after truncation*.

2. **Overlap layer (Overlap Gate, cf. Ch. 1, Def. 1.0 and Ch. 5, Def. 5.8):** on each overlap, require
   - collapse-compatibility up to the recorded δ-budget (Appendix L; additively aggregated),
   - Čech–Ext^1-acyclicity in degree 1 *after truncation*,
   - stability-band condition $(\mu, u) = (0, 0)$ with near-τ non-accumulation.
3. **Global layer:** when all overlaps pass, glue to a global truncated object in $\mathrm{Pers}_k^{\mathrm{cons}}$ (Ch. 5, Thm. 5.8) and perform global audits *after truncation*.

The run manifest (§11.17) must include overlap checks (boolean), Čech–Ext^1 status, and the overlap δ-budgets.

Remark 11.4 (Why this order is unique)**.** Pre-collapse comparisons contaminate measurements by torsion noise and break exactness/metric controls. Truncation $\mathbf{T}_\tau$ is the exact, 1-Lipschitz barrier that makes all downstream comparisons commensurable.

11.1. Collapse energy (windowed persistence energies)

Fix a degree i, a scale $\tau \geq 0$, and an exponent $\alpha > 0$ (default $\alpha = 1$). Let $\mathcal{B}_i(F)$ be the barcode of $\mathbf{P}_i(F)$, and for a bar $b = [b_\ell, b_r)$ write its τ-windowed length

$$\ell_\tau(b) \ := \ \big(\min\{b_r, \tau\} - \min\{b_\ell, \tau\}\big)_+, \qquad (x)_+ := \max\{x, 0\}.$$

Fix a weight function $w_i : \mathcal{B}_i(F) \to [0, \infty)$ (default $w_i \equiv 1$).

Definition 11.5 (Windowed energies)**.** The (weighted) windowed energy at degree i is

$$\mathrm{PE}_i^{\leq\tau}(F; w_i, \alpha) \ := \sum_{b \in \mathcal{B}_i(F)} w_i(b)\, \big(\ell_\tau(b)\big)^\alpha.$$

With the unweighted shorthand $\mathrm{PE}_i^{\leq\tau}(F) := \mathrm{PE}_i^{\leq\tau}(F; 1, \alpha)$ (default $\alpha = 1$), the *collapse energy* vector at τ is

$$\mathbf{CE}^{\leq\tau}(F) \ := \ \big(\mathrm{PE}_i^{\leq\tau}(F)\big)_{i\in\mathbb{Z}}, \qquad \|\mathbf{CE}^{\leq\tau}(F)\|_1 \ := \ \sum_i \mathrm{PE}_i^{\leq\tau}(F).$$

All quantities are computed on the truncated barcode $\mathbf{T}_\tau \mathbf{P}_i(F)$ (equivalently, on $C_\tau F$).

Remark 11.6 (Stability and monotonicity)**.** By Part I (Lemma 2.4, Proposition 2.5), 1-Lipschitz updates under $\mathbf{P}_i$ induce non-expansive changes of $\mathrm{PE}_i^{\leq\tau}$; in *deletion-type* updates, $\mathrm{PE}_i^{\leq\tau}$ is non-increasing up to f.q.i.. Since $\mathbf{T}_\tau$ is 1-Lipschitz for the interleaving distance, the same Lipschitz control carries over after truncation:

$$d_{\mathrm{int}}\big(\mathbf{T}_\tau \mathbf{P}_i(F'),\, \mathbf{T}_\tau \mathbf{P}_i(F)\big) \ \leq \ d_{\mathrm{int}}\big(\mathbf{P}_i(F'),\, \mathbf{P}_i(F)\big).$$

Under finite-type and objectwise degreewise-colimit hypotheses, $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$ at fixed τ excludes Type IV at that scale (Ch. 4).

11.2. Spectral indicators on $C_\tau F$

Let $L(C_\tau F)$ be the normalized combinatorial Hodge Laplacian on $C_\tau F$ (per degree, with the standard Euclidean inner product on chains), and let $\{\lambda_j\}_{j\geq 1}$ denote its positive eigenvalues (zero modes excluded; alternatively treat by Moore–Penrose pseudoinverse).

Definition 11.7 (Spectral tails and heat traces)**.** Fix $\beta > 0$ and a cutoff policy $M(\tau) \in \mathbb{N}$. Define

$$\mathrm{ST}_\beta^{\geq M(\tau)}(F) \;:=\; \sum_{j \geq M(\tau)} \lambda_j^{-\beta}, \qquad \mathrm{HT}(t;F) \;:=\; \sum_{j\geq 1} e^{-t\,\lambda_j}, \quad t > 0.$$

Typical policies: *(i)* $M(\tau) = \lfloor c\,\tau^\gamma \rfloor$ with $c > 0$, $\gamma \in [0,2]$; *(ii)* $t \in [c_1\tau^{-2}, c_2\tau^{-2}]$ with fixed $0 < c_1 \leq c_2$.

Remark 11.8 (Mandatory spectral ordering and norms)**.** For auditability and comparability across runs, it is *mandatory* that

- eigenvalues in spectral artifacts are stored in *ascending* order;
- the chosen matrix norm for numeric tolerances is reported as $\mathrm{norm} \in \{\mathrm{op}, \mathrm{fro}\}$.

These fields are required in $\mathrm{run.yaml}$ (Appendix G) and are rechecked in §11.17.

Remark 11.9 (Stability)**.** Non-expansive filtered maps on $C_\tau F$ induce 1-Lipschitz updates at the persistence layer and stable (non-expansive) spectral responses for $\mathrm{ST}_\beta^{\geq M(\tau)}$, $\mathrm{HT}(t;-)$, provided the policy $(\beta, M(\tau), t)$ is held fixed (cf. Chs. 6–10).

11.3. Auxiliary spectral bars (aux-bars)

We supplement spectral scalars by a *binwise occupancy* summary along a discrete index (e.g. time, step number), evaluated on *collapsed* Laplacians $L(C_\tau F)$.

Definition 11.10 (Spectral bins, occupancies, and aux-bars)**.** Fix a spectral window $[a,b]$ and a bin width $\beta > 0$. Define right-open bins

$$I_r \;:=\; [\,a + r\beta,\; a + (r+1)\beta\,), \qquad r = 0,1,\dots,R-1, \quad R := \left\lfloor \frac{b-a}{\beta} \right\rfloor.$$

For a sample index j (e.g. time step), let $\{\lambda_m(j)\}_{m\geq 1}$ be the positive eigenvalues of $L(C_\tau F_j)$ and define *occupancies*

$$E_r(j) \;:=\; \#\{\,m \;\mid\; \lambda_m(j) \in I_r\,\}, \qquad \text{with under/overflow counts } E_{<a}(j),\; E_{\geq b}(j) \text{ logged explicitly.}$$

For each fixed bin r, define an *auxiliary spectral bar* (aux-bar) as a maximal consecutive run $J \subset \mathbb{Z}$ such that $E_r(j) > 0$ for all $j \in J$. The *lifetime* of the aux-bar is $|J|$ (or a rescaling thereof).

Proposition 11.11 (Cumulative profile monotonicity and stability)**.** Let $C_r(j) := \sum_{s=r}^{R-1} E_s(j)$ be the cumulative profile for the spectral window. Then:

1. *(Deletion-type monotonicity)* Under deletion-type updates (Dirichlet/principal restriction, Schur complement, Loewner contractions; Appendix E), $C_r(j+1) \leq C_r(j)$ for all r.
2. *(Stability under ε-continuations)* If $\|A_{j+1} - A_j\|_{\mathrm{op}} \leq \varepsilon$, then $C_{r+q}(j+1) \leq C_r(j) \leq C_{\max\{0,r-q\}}(j+1)$ with $q = \lceil \varepsilon/\beta \rceil$.

Proof sketch. Appendix E (Propositions B.15, B.16). □

Remark 11.12 (Gate alignment)**.** The cumulative profile and aux-bars are *auxiliary* gauges for B–Gate$^+$ (Declaration **??**). They never override the persistence-layer verdict (PH/Ext/μ–u), but they provide robust, windowed evidence of deletion-type monotonicity or ε-stability.

Remark 11.13 (Mandatory bin policy)**.** Aux-bars must be computed *after* collapse and with a fixed bin policy (a,b,β) per window; boundary is right-open; under/overflow are recorded (Appendix G).

11.4. Categorical check (one-way bridge)

With $\{\, k[0]\,\} = \{k[0]\}$ fixed, we monitor the categorical obstruction

$$\mathrm{Ext}^1\left(\mathcal{R}(C_\tau F), Q\right) = 0 \qquad (Q \in \{\, k[0]\,\}).$$

This check is performed *after truncation* and used only in the one-way direction sanctioned in Part I (under (B1)–(B3)); no converse or equivalence is claimed.

11.5. Collapse diagnostics along towers

For an index category I (time/resolution/parameter), adjoin a terminal element ∞ and cone maps to form comparison morphisms

$$\phi_{i,\tau} : \varinjlim_{t\in I} \mathbf{T}_\tau \mathbf{P}_i(F_t) \longrightarrow \mathbf{T}_\tau \mathbf{P}_i(F_\infty),$$

with

$$\mu_{i,\tau} = \dim_k \ker \phi_{i,\tau}{}^{5}, u_{i,\tau} = \dim_k \operatorname{coker} \phi_{i,\tau}{}^{6},$$

and $\mu_{\mathrm{Collapse}} = \sum_i \mu_{i,\tau}$, $u_{\mathrm{Collapse}} = \sum_i u_{i,\tau}$ (finite by finite-typeness). Under the hypotheses of Part I/Ch. 4, $\phi_{i,\tau}$ are isomorphisms and $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$ at fixed τ.

11.6. Specification: joint monitoring protocol

Declaration 11.14 (Specification: Joint monitoring protocol)**.** Fix a finite sweep $\tau \in [\tau_{\min}, \tau_{\max}]$ and a policy $(\alpha, w_i;\ \beta, M(\tau), t)$. For each sample $t \in I$ and degree i:

1. *Compute and record* $\mathbf{T}_\tau \mathbf{P}_i(F_t)$; evaluate $\mathrm{PE}_i^{\le\tau}(F_t)$ on $\mathbf{T}_\tau \mathbf{P}_i(F_t)$ (equivalently on $C_\tau F_t$).
2. *Compute and record* $\mathrm{ST}_\beta^{\ge M(\tau)}(F_t)$ and $\mathrm{HT}(t; C_\tau F_t)$ using $L(C_\tau F_t)$; compute aux-bars via Definition 11.10 on the same window $[a,b]$ and bin policy β.
3. *Check* $\mathrm{Ext}^1\left(\mathcal{R}(C_\tau F_t), Q\right) = 0$ for $Q \in \{\, k[0]\,\}$.
4. *Evaluate* $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ at each τ via the comparison maps $\phi_{i,\tau}$; log failure type (pure kernel/cokernel/mixed) if $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) eq (0,0)$.
5. *Declare stable regime* at τ when (1)–(3) hold jointly and $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$. Optionally require aux-bars= 0 (Declaration **??**).

All persistence-layer steps are asserted at the persistence layer and are invariant under f.q.i. by construction; spectral and aux-bar steps are not f.q.i. invariants but are *stable* under the fixed policy $(\beta, M(\tau), t)$ on $L(C_\tau F_t)$.

[5] Here $\dim_k$ denotes the *generic fiber* dimension after truncation, i.e. the multiplicity of $I[0,\infty)$ summands; see Appendix D, Remark A.1.

[6] Here $\dim_k$ denotes the *generic fiber* dimension after truncation, i.e. the multiplicity of $I[0,\infty)$ summands; see Appendix D, Remark A.1.

11.7. Noise tolerance and discretization rules

Declaration 11.15 (Specification: noise and discretization policy)**.** Let $\varepsilon > 0$ be the noise scale.

- **Barcode denoising:** apply *ε-clipping* on $\mathbf{T}_\tau \mathbf{P}_i(F)$ by removing bars of length $\le \varepsilon$ within the τ-window. This is stable under bottleneck perturbations $\le \varepsilon$ and preserves f.q.i. invariants.
- **Energy stability:** there exists $C_{i,\tau,\alpha}$ (depending on the number of bars in the window, their maximal multiplicity, and ambient dimension bounds) such that

$$\left|\mathrm{PE}_i^{\le\tau}(F) - \mathrm{PE}_i^{\le\tau}(\tilde{F})\right| \le C_{i,\tau,\alpha}\, \varepsilon^{\min\{1,\alpha\}}$$

  whenever $d_{\mathrm{int}}(\mathbf{T}_\tau \mathbf{P}_i(F), \mathbf{T}_\tau \mathbf{P}_i(\tilde{F})) \le \varepsilon$.
- **Spectral stabilization:** compute spectra on $L(C_\tau F)$; use the same $(\beta, M(\tau), t)$ across runs. Optional averaging over N realizations reduces variance as $N^{-1/2}$. Aux-bar lifetimes are counted with the same binning policy; short lifetimes ≤ 2 frames may be ignored as noise (manifest must record the rule).
- **Resolution rule:** choose sampling so that the minimal feature length resolved is ≥ 3 grid steps; sweep τ on a lattice $\Delta\tau$ satisfying $\Delta\tau \le \frac{1}{2}$ of the minimal resolvable bar length. Here the "bar-length quantum" refers to the minimal resolvable feature length induced by the sampling/grid spacing and the filtration step.

11.8. Categorical check (bridge) and saturation gate

With $\{\, k[0]\,\} = \{k[0]\}$ fixed, we monitor $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), Q) = 0$ after truncation. The one-way bridge is used only under (B1)–(B3). For windows satisfying the *saturation* conditions:

Declaration 11.16 (Saturation gate **[Spec]**)**.** Fix $\tau^* > 0$ and parameters $\eta, \delta > 0$. On the window $[0, \tau^*]$, assume: (i) eventually the maximal *finite* bar length in $\mathbf{T}_{\tau^*}\mathbf{P}_i(F_t)$ is $\le \eta$; (ii) eventually $d_{\mathrm{int}}\big(\mathbf{T}_{\tau^*}\mathbf{P}_i(F_t), \mathbf{T}_{\tau^*}\mathbf{P}_i(F_{t'})\big) \le \eta$; (iii) the *edge gap* to the window, $\delta := \tau^* - \max\{b_r < \tau^*\}$, satisfies $\delta > \eta$. Then, **within this window only**, we adopt the temporary binary policy

$$\mathrm{PH}_1(\mathsf{C}_{\tau^*}F) = 0 \iff \mathrm{Ext}^1(\mathcal{R}(\mathsf{C}_{\tau^*}F), k) = 0.$$

11.9. Data formats, reproducibility, and minimal manifest

Remark 11.17 (Implementation notes and mandatory fields)**.** *Artifacts.* (i) bars.json/h5: list of records $\langle i,\ b_\ell, b_r,\ w\rangle$ per degree; (ii) spec.json/h5: positive eigenvalues of $L(C_\tau F)$ per degree; (iii) aux.json/h5: aux-bar occupancies $E_r(j)$ with bin metadata; (iv) ext.json: boolean for $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), Q)$ with minimal witness; (v) phi.json: ranks of comparison maps and $(\mu_{i,\tau}, u_{i,\tau})$. *Run log.* Store (a) sweep $\tau_{\min}{:}\Delta\tau{:}\tau_{\max}$; (b) policy $(\alpha, w_i;\ \beta, M(\tau), t)$; (c) discretization (grid/complex, step sizes); (d) random seeds; (e) software versions; (f) δ-ledger per step; (g) bin window $[a, b]$, bin width β, under/overflow counts; (h) *mandatory* spectral fields order=ascending, norm $\in \{\mathrm{op}, \mathrm{fro}\}$; (i) overlap checks (overlap_checks), Čech–Ext^1 status (cech_ext1_ok), and tail-isomorphism flag (phi_iso_tail); (j) an optional length_spectrum summary (vector of clipped lengths or hash) for quick audits. *Invariance.* Persistence-layer quantities are taken on $\mathbf{T}_\tau \mathbf{P}_i(-)$ or $C_\tau(-)$ and are invariant under f.q.i.. *Reproducibility.* A single manifest (run.yaml) references all artifacts, declares the tower index set, cone extension, the decision rule for stable regimes, and the overlap status.

A minimal run.yaml block (augmented):

```
windows:
  domain: [[0,1), [1,2), [2,3)]
  collapse_tau: 0.08
  spectral_bins: {a: 0.0, beta: 0.02, bins: 96, boundary: ”right-open”}
coverage_check:
  length_sum: 3.0
  length_target: 3.0
  events_sum_equals_global: true
overlap_checks:
  local_equal_after_collapse: true
  cech_ext1_ok: true
  stability_band_ok: true
spectral_policy:
  order: ”ascending”
  norm: ”op”
operations:
  - U: mollify; type: deletion; tau: 0.08; delta: {alg:0.004, disc:0.003, meas:0.001}
  - U: timestep; type: epsilon;  tau: 0.08; eps:0.006; delta: {alg:0.000, disc:0.002, meas:0.001}
persistence:
  PH1_zero: true
  Ext1_zero: true
  mu: 0
  nu: 0
  phi_iso_tail: true
length_spectrum:
  degree: 1
  tau: 0.08
  eigenvalues: [0.24, 0.51, 0.78]   # optional summary of clipped bar lengths
spectral:
  ST_beta: 2
  ST_M_of_tau: ”floor(0.5 * tau^1.5)”
  HT_t: [0.5*tau^-2, 1.0*tau^-2]
  aux_bars_remaining: 0
budget:
  sum_delta: 0.011
  safety_margin: 0.025
gate:
  accept: true
```


11.10. Compliance checks and unit tests

Declaration 11.18 (Specification: minimal test suite)**.** Every deployment must pass:

- **Stability test.** For synthetic ε-perturbations, verify non-expansiveness of $\mathrm{PE}_i^{\leq\tau}$ and stability of spectral indicators and aux-bars under the fixed policy.

- **Monotone update test.** For a deletion-type update, confirm $\mathrm{PE}_i^{\leq\tau}$ non-increase, spectral tail non-increase, aux-bars non-increase (active bins do not grow), and record $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$ at fixed τ.

- **Cone-extension test.** Verify $\phi_{i,\tau}$ is an isomorphism on a model tower (hence $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0,0)$) and that Type IV is excluded at that τ.

- **Categorical check.** On a curated sample, confirm $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), Q) = 0$ is stable under admissible f.q.i. updates.

11.10 bis. Hooks to Chapter 12 (test harness)

Chapter 12 formalizes the global test harness. In particular:

- **T7 (Saturation gate).** Verifies the quantitative saturation conditions of §11.16 on $[0, \tau^*]$ and logs the temporary local equivalence.

- **T10 (A/B commutativity).** Runs A/B tests after truncation (Appendix K/L) on windowed reflectors; failures trigger a deterministic order with the commutation defect logged as δ^{alg}.

- **T13 (Pipeline δ-budget).** Aggregates Mirror–Collapse and A/B residuals additively and checks post-processing non-increase (§11.17, Appendix L).

Passing T7/T10/T13 is necessary for issuing *windowed certificates* (Ch. 4) and for *pasting* via Restart/Summability (Appendix J).

11.A. Formalization API (Lean/Coq stubs)

All items are stated in the constructible range and, at the filtered–complex layer, *up to filtered quasi-isomorphism*. Identifiers are indicative.

Specification 11.19 (Persistence truncation $\mathbf{T}_\tau$ — exactness and 1-Lipschitz).

- pers_Ttau_exact: $\mathbf{T}_\tau$ is exact on $\mathsf{Pers}_k^{\text{cons}}$ (Serre localization at τ).

- pers_Ttau_lipschitz: $d_{\text{int}}(\mathbf{T}_\tau M, \mathbf{T}_\tau N) \le d_{\text{int}}(M, N)$.

- pers_Ttau_commute_colim_flim: $\mathbf{T}_\tau$ commutes with filtered colimits and finite limits.

Specification 11.20 (Comparison maps ϕ — functoriality and cofinal invariance). For a filtered diagram $\{F_t\}_{t\in I}$ and degree i,

$$\phi_{i,\tau} : \varinjlim_t \mathbf{T}_\tau \mathbf{P}_i(F_t) \longrightarrow \mathbf{T}_\tau \mathbf{P}_i(F_\infty).$$

- phi_natural: natural in (a) morphisms of diagrams; (b) τ via $\mathbf{T}_\tau \Rightarrow \mathbf{T}_\sigma$ for $\tau \le \sigma$.

- phi_cofinal: invariant under cofinal reindexing $I \to I'$.

- phi_sum_prod: additive under finite direct sums; compatible with finite limits (pullbacks) at the persistence layer.

Specification 11.21 (Kernel/cokernel indices (μ, u) — computation and calculus).

- mu_nu_def: $\mu_{i,\tau} = \dim_k \ker \phi_{i,\tau}$, $u_{i,\tau} = \dim_k \operatorname{coker} \phi_{i,\tau}$; $\mu_{\text{Collapse}} = \sum_i \mu_{i,\tau}$, $u_{\text{Collapse}} = \sum_i u_{i,\tau}$ (finite by bounded degree).

- mu_nu_vanish: under degreewise objectwise filtered colimits (Part I/Ch. 4), each $\phi_{i,\tau}$ is an isomorphism, hence $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0,0)$.

- mu_nu_calc: subadditivity under composition, additivity under finite sums, and cofinal invariance (Appendix J).

  Here $\dim_k$ denotes the *generic fiber* dimension after truncation (Appendix D, Remark A.1).

Specification 11.22 (Edge identification and bridge (B2))**.**

- edge_iso_B2: natural isomorphism $H^{-1}(\mathcal{R}(F)) \cong \varinjlim_t H_1(F^t C_\bullet)$ and $\mathcal{R}(F) \in D^{[-1,0]}$.
- ext1_edge_iso: natural isomorphism $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) \cong \mathrm{Hom}(H^{-1}(\mathcal{R}(C_\tau F)), k)$; in particular $H^{-1} = 0$ iff $\mathrm{Ext}^1 = 0$ in this amplitude-≤ 1 range.

Specification 11.23 (Aux-bars (measurement layer))**.**

- aux_bins: parameters (a, b, β); right-open bins $I_r = [a + r\beta, a + (r+1)\beta)$; record under/overflow.
- aux_occupancy: $E_r(j) = \#\{m \mid \lambda_m(j) \in I_r\}$ from the spectrum of $L(C_\tau F_j)$.
- aux_bars: for fixed r, maximal consecutive runs J with $E_r(j) > 0$; lifetime $|J|$ (or rescaled).
- aux_monotone: deletion-type steps $\Rightarrow$ non-increase of active bins and total aux-bar mass; ε-continuations $\Rightarrow$ stability; inclusion-type $\Rightarrow$ stability only.

11.11. Diagram (pipeline and logs)

$$
\begin{tikzcd}
\text{Input } F \arrow[rr, "\mathbf{P}_i"] & & \mathbf{P}_i(F) \arrow[r, "\mathbf{T}_\tau"] \arrow[d, dashed, "\mathrm{PE}_i^{\leq\tau}"] & \mathbf{T}_\tau\mathbf{P}_i(F) \arrow[d, dashed, "\mathrm{PE}_i^{\leq\tau}"] \\
& & L(C_\tau F) \arrow[r, dashed, "\text{heat/aux on } L(C_\tau F)"] & \mathrm{HT},\ \mathrm{ST}_\beta,\ \text{aux-bars} \\
\mathcal{R}(F) \arrow[uu, dashed, bend left, "\text{(LC)}"] \arrow[rrr, dashed, "C_\tau\ \text{ then } \mathrm{Ext}^1(-,Q)"] & & & \mathrm{Ext}^1(\mathcal{R}(C_\tau F), Q) = 0
\end{tikzcd}
$$

$$
\begin{tikzcd}
\varinjlim_t \mathbf{T}_\tau\mathbf{P}_i(F_t) \arrow[rr, "\phi_{i,\tau}"] & & \mathbf{T}_\tau\mathbf{P}_i(F_\infty) \arrow[dll, dashed, bend right, "(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})\ \log"'] \\
\text{failure log (pure kernel/cokernel/mixed)} & &
\end{tikzcd}
$$

11.12. Completion note

Remark 11.24 (No further supplementation required)**.** This chapter fully integrates: (i) collapse energy and spectral indicators on $L(C_\tau F)$; (ii) auxiliary spectral bars (aux-bars) with binning, occupancies, lifetimes, and monotonicity/stability policies; (iii) the joint monitoring protocol with gate usage of aux-bars as optional auxiliary conditions; (iv) noise/discretization policies, including lifetime thresholds and reproducibility; (v) a minimal manifest schema and artifact set; (vi) a formalization stub for persistence, towers, bridge, and aux-bars; (vii) the length-spectrum crosslink (§11.0+) and the Overlap Gate integration (§11.0 bis); (viii) mandatory spectral ordering and norms (§11.8); (ix) hooks to the global test harness (§11.10 bis). All claims remain within the v16.0 guard-rails (B-side single layer, PH1$\to$Ext1 one-way only, MECE windows, δ-ledger, tower diagnostics); no additional supplementation is needed for operational use as a proof/measurement framework.

11.13. Guard-rails

Remark 11.25 (Scope and non-claims)**.** This chapter specifies measurement protocols, auxiliaries, and implementation notes at the persistence/spectral/categorical layers. No analytic regularity, group trivialization, or number-theoretic identity is asserted. All statements respect the guard-rails of Part I; in particular, no claim of $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is made, and μ_{Collapse} differs from classical Iwasawa μ.

12 Chapter 12: Formal Test Suite and Open Problems

Badge policy.

- **[Prop]**: mathematics proved in Part I (core results; cite exact source).
- **[Declaration]**: programmatic specification in the implementable range, verifiable by the test suite in this chapter.
- **[Conjecture]**: forward-looking statement; no claim beyond the stated scope.

12.0. Notation & conventions

- **Constructible range.** We identify $\mathsf{Pers}_k^{\mathrm{cons}}$ with the constructible subcategory of $[\mathbb{R}, \mathsf{Vect}_k]$ and use $\mathsf{Pers}_k^{\mathrm{cons}}$ uniformly (we do not use $\mathsf{Pers}_k^{\mathrm{cons}}$ hereafter).
- **Truncation phrase.** "after applying $\mathbf{T}_\tau$; equivalently on $C_\tau F$" is our standard phrase indicating that a quantity is computed at the persistence layer after truncation (hence equivalently on the filtered lift $C_\tau F$).
- **Generic–fiber dimension.** For a comparison map $\phi_{i,\tau}$ at fixed τ, $\dim_k$ denotes the *generic–fiber dimension after truncation*, i.e. the multiplicity of $I[0,\infty)$ summands in $\mathbf{T}_\tau \mathbf{P}_i(-)$; informally, the $t \to \infty$ stable rank within the τ–window.
- **Spectral ordering and norms.** Positive eigenvalues of $L(C_\tau F)$ are listed in ascending order $\lambda_1 \leq \lambda_2 \leq \cdots$. Matrix/operator norms are denoted by $\|\cdot\|_{\mathrm{op}}$ (operator norm) and $\|\cdot\|_{\mathrm{fro}}$ (Frobenius norm); each test specifies which is used and logs the choice.
- **Obstruction totals and macros.** We write $\mu_{i,\tau} = \dim_k \ker \phi_{i,\tau}$, $u_{i,\tau} = \dim_k \operatorname{coker} \phi_{i,\tau}$, and set the totals $\mu_{\mathrm{Collapse}} := \sum_i \mu_{i,\tau}$, $u_{\mathrm{Collapse}} := \sum_i u_{i,\tau}$ (finite by bounded degree).
- **Endpoints.** Endpoint conventions and the handling of infinite bars follow the global policy (cf. Appendix A, Remark **??**); infinite bars are not removed by $\mathbf{T}_\tau$ and are clipped by the window in all windowed quantities.
- **Non-expansiveness.** We use the spelling "non-expansive"/"non-expansiveness" uniformly.

12.1. Badge inventory (representative items)

Badge	Representative items (label / location)
[Prop]	Stability, idempotence, and exactness of $\mathbf{T}_\tau$ (Prop. 2.5, Ch. 2); Shift–commutation / 1-Lipschitz for $\mathbf{T}_\tau$ (Lemma 2.4, Ch. 2); Operational coreflection $\mathsf{C}_\tau^{\mathrm{comb}}$ on the implementable range (Prop. 5.1, Ch. 5); Tower diagnosis: $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ via cone extension; isomorphism criterion excluding Type IV (Prop. 4.6, Ch. 4).
[Thm]	One-way bridge: $\mathrm{PH}_1(F) = 0 \Rightarrow \mathrm{Ext}^1(\mathcal{R}(F), k) = 0$ under (B1)–(B3) (Thm. 3.8, Ch. 3).
[Declaration]	Ch. 2: (co)limit and pullback compatibility *at the persistence layer only* (after $\mathbf{T}_\tau$); Ch. 6: filtered-colimit stability in geometry; joint indicators; protocol (after truncation); Ch. 7: arithmetic tower stability; non-identity of μ_{Collapse} with Iwasawa μ; Ch. 8: tropical shortening $\Rightarrow$ weak group collapse; mirror transfer *non-expansive after truncation*; Ch. 9: three-layer (Gal$\rightarrow$Trans$\rightarrow$Funct) compatibility as isomorphisms in $\mathsf{Pers}_k^{\mathrm{cons}}$ *after* $\mathbf{T}_\tau\mathbf{P}_i$; Ch. 10: persistence-guided regularization; AK–NS hypothesis (programmatic); Ch. 11: joint monitoring, noise/discretization policy, minimal test suite; central **Saturation gate** (Ch. 11.S).
[Conjecture]	Cross-domain collapse propagation (Chs. 6–10); AK–NS (Ch. 10); mirror-side propagation (Ch. 8); Functorial transfer stability (Ch. 9).

12.2. Formal test suite (unit / integration / regression)

All tests operate at the truncated persistence, spectral (on $L(C_\tau F)$), and categorical layers and are f.q.i. invariant at the persistence layer. A test *passes* iff all stated pass-criteria are met and logs are complete. Pass-criteria must state whether indicators are evaluated *per degree* or *aggregated* across degrees; the choice must be fixed and logged for the run. Spectra use ascending order $\lambda_1 \le \lambda_2 \le \cdots$, and the chosen norm $\|\cdot\|_{\mathrm{op}}$ or $\|\cdot\|_{\mathrm{fro}}$ must be declared and logged.

(T1) Stability under non-expansive updates [Unit]. *Input:* pairs $F \to F'$ with $d_{\mathrm{int}}(\mathbf{P}_i(F), \mathbf{P}_i(F')) \le \varepsilon$.
Assertions: $\left|\mathrm{PE}_i^{\le\tau}(F) - \mathrm{PE}_i^{\le\tau}(F')\right| \le C_{i,\tau,\alpha}\, \varepsilon^{\min\{1,\alpha\}}$ (computed after applying $\mathbf{T}_\tau$; equivalently on $C_\tau F$); spectra of $L(C_\tau F)$ vs. $L(C_\tau F')$ (eigenvalues ascending) satisfy the fixed $(\beta, M(\tau), t)$-policy stability bounds in the declared norm; $\mathrm{Ext}^1(\mathcal{R}(C_\tau -), Q)$ is *stable under admissible f.q.i. updates* (for all $Q \in \{\, k[0] \,\}$); in particular, if $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), Q) = 0$ at baseline, then $\mathrm{Ext}^1(\mathcal{R}(C_\tau F'), Q) = 0$.
Artifacts: bars.json, spec.json, ext.json; run.yaml (norm and spectral policy recorded).

(T2) Monotone update (deletion-/inclusion-type) [Unit]. *Input:* $F \to F'$ monotone.
Assertions: Deletion-type: $\mathrm{PE}_i^{\le\tau}$ and spectral indicators are non-increasing (after $\mathbf{T}_\tau$); spectral eigenvalues are compared in ascending order in the declared norm. In the minimal two-term cone tower (source F, terminal F') the comparison map $\phi_{i,\tau}$ is an isomorphism, hence $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0,0)$ at fixed τ.

Inclusion-type: stability (non-expansiveness) only; no non-increase is claimed.
Artifacts: bars.json, spec.json, ext.json, phi.json; run.yaml (update type, norm choice, and $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$ recorded).

(T3) Filtered-colimit stability [Integration]. *Input:* tower $\{F_\lambda\}_\lambda$.
Assertions: for each fixed τ, comparison maps $\phi_{i,\tau} : \varinjlim_\lambda \mathbf{T}_\tau \mathbf{P}_i(F_\lambda) \xrightarrow{\cong} \mathbf{T}_\tau \mathbf{P}_i(F_{\lambda_*})$; thus $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0,0)$ and Type IV excluded at that scale. *Terminal symbol consistency:* within a run, the same terminal symbol (e.g. F_{λ_*} or F_∞) is used and logged in run.yaml.
Artifacts: phi.json with ranks and $(\mu_{i,\tau}, u_{i,\tau})$; run.yaml (terminal symbol).

(T4) Mirror/tropical pipeline [Integration]. *Input:* X, tropical flow Trop_λ, realization F_λ, mirror functor Mirror.
Assertions: tropical shortening factor $\kappa \leq 1$ implies non-increase of $\text{PE}_i^{\leq\tau}$ (after $\mathbf{T}_\tau$); mirror transfer is *non-expansive after truncation* at the persistence level; group proxies (if used) meet weak-collapse thresholds (Ch. 8). Spectral eigenvalues are reported in ascending order with the declared norm.
Artifacts: bars.json, spec.json, ext.json, phi.json (per λ and on the mirror side); run.yaml.

(T5) Three-layer compatibility [Integration]. *Input:* Gal→Trans→Funct data with comparison natural transformations (Ch. 9).
Assertions: after applying C_τ and $\mathbf{P}_i$, commutativity holds up to isomorphism in $\mathsf{Pers}_k^{\text{cons}}$ for each degree i; indicators consistent across layers; failure logs record type if mismatches occur.
Artifacts: per-layer bars.json, spec.json, ext.json; global phi.json; run.yaml.

(T6) PDE monitoring loop [Regression]. *Input:* index set I (time/resolution/parameter), realization $\mathcal{P}$.
Assertions: protocol of Ch. 11 (Decl. 11.14) holds; stable regime declaration matches (logs $\text{PE}^{\leq\tau}$, spectral indicators and aux-bars, Ext^1, $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$); pass-criteria specify *per-degree vs. aggregated* reporting and are fixed and logged. Repeatability from run.yaml. Spectral eigenvalues ascending; norm choice declared.
Artifacts: bars.json, spec.json, aux.json, ext.json, phi.json over I; run.yaml.

(T7) Saturation gate verification [Integration]. *Input:* a window $[0,\tau^*]$ with candidate saturation (Ch. 11.S).
Assertions: verify the quantitative saturation conditions on the window: (a) maximal *finite* bar length $\leq \eta$; (b) interleaving bound $\leq \eta$ eventually; (c) edge gap $\delta := \tau^* - \max\{b_r < \tau^*\} > \eta$. Confirm the **[Spec]** adoption $\text{PH}_1(\mathsf{C}_{\tau^*}F) = 0 \Leftrightarrow \text{Ext}^1(\mathcal{R}(\mathsf{C}_{\tau^*}F), k) = 0$ is recorded and used only on the saturated window. Spectral reports use ascending ordering and the declared norm.
Artifacts: bars.json, ext.json; run.yaml (saturation parameters, norm, decision).

(T8) ε-clipping regression [Unit]. *Input:* paired runs with unclipped vs. ε-clipped $\mathbf{T}_\tau \mathbf{P}_i$ (Ch. 11.7).
Assertions: stability bound for $\text{PE}_i^{\leq\tau}$ under clipping (theoretical upper bound matched up to tolerance); $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$ are computed on unclipped data and remain identical across the pair; logs explicitly distinguish clipped vs. unclipped reporting.
Artifacts: bars.json (both variants), phi.json; run.yaml.

(T9) MECE window coverage & event accounting [Unit]. *Input:* a MECE windowing $\{[u_k, u_{k+1})\}_k$ and global range $[u_0, U)$.
Assertions: (a) $\sum_k (u_{k+1} - u_k) = U - u_0$; (b) the total number of events (births/deaths counted with multiplicity) on $[u_0, U)$ equals the sum of per-window counts up to rounding tolerance; (c) the collapse threshold τ and spectral bin policy (a, b, β) are *identical* across windows unless explicitly recorded and justified in run.yaml.
Artifacts: run.yaml (windows/coverage_check block), per-window bars.json.

(T10) A/B commutativity test for reflectors [Unit/Integration]. *Input:* two persistence-level reflectors T_A, T_B (e.g. length-threshold and birth-window), tolerance $\eta \geq 0$.
Assertions: compute $\Delta_{\mathrm{comm}}(M; A, B) = d_{\mathrm{int}}(T_A T_B M,\ T_B T_A M)$ on the relevant dataset M. If $\Delta_{\mathrm{comm}} \leq \eta$, adopt *soft-commuting*; else fall back to a fixed order and record Δ_{comm} into δ^{alg} (Appendix L). Pass if the outcome matches the configured policy in run.yaml.
Artifacts: run.yaml (A/B policy: η, order, soft-commuting flag), bars.json before/after.

(T11) Restart experiment & Summability design [Integration/Regression]. *Input:* a MECE window sequence with collapse thresholds τ_k, budgets $\Sigma\delta_k(i)$, and safety margins gap_{τ_k}.
Assertions: (Restart) empirically verify $\mathrm{gap}_{\tau_{k+1}} \geq \kappa(\mathrm{gap}_{\tau_k} - \Sigma\delta_k(i))$ for some $\kappa \in (0, 1]$ (record κ); (Summability) verify $\sum_k \Sigma\delta_k(i) < \infty$ (e.g. geometric decay of τ_k, β_k and bounded step counts); (Pasting) confirm that all B-Gate$^+$ certificates paste to a global certificate.
Artifacts: run.yaml (restart/summability fields), per-window gate logs, cumulative certificate.

(T12) Trigger pack verification (domain-restricted) [Integration]. *Input:* triggers declared for a domain (e.g. PDE, Ch. 10.4).
Assertions: for each trigger, verify that the declared analytical/numerical condition implies a B-Gate$^+$ failure on the window (e.g. sustained enstrophy surge $\Rightarrow$ aux-bars > 0 or $\mu > 0$); record detection rate and false positives (if any). Triggers are **[Spec]** and complement (not replace) B-Gate$^+$.
Artifacts: run.yaml (trigger set, thresholds), aux.json, phi.json, gate verdicts.

(T13) δ-ledger additivity & pipeline budget [Integration]. *Input:* a pipeline of steps $U_m, \ldots, U_1$ with per-step collapses C_{τ_j} and bounds $\delta_j(i, \tau_j)$.
Assertions: verify

$$d_{\mathrm{int}}\Big(\mathbf{T}_\tau \mathbf{P}_i\big(\mathrm{Mirror}(C_{\tau_m} \cdots C_{\tau_1} F)\big),\ \mathbf{T}_\tau \mathbf{P}_i\big(C_{\tau_m} \cdots C_{\tau_1}\, \mathrm{Mirror}\, F\big)\Big) \ \leq\ \sum_{j=1}^{m} \delta_j(i, \tau_j),$$

and that any post-processing by 1-Lipschitz maps does not increase the bound (Appendix L). Pass if the measured defect is $\leq$ the recorded budget.
Artifacts: run.yaml (δ-ledger), bars.json along the pipeline, distance logs.

(T14) Overlap Gate gluing test [Integration]. *Input:* two charts X_1, X_2 with windows W_1, W_2 and overlap $X_{12} = X_1 \cap X_2$, together with a fixed τ and a pair of reflectors T_A, T_B (as in T10).
Assertions: after applying $\mathbf{P}_i$ and $\mathbf{T}_\tau$, verify *local equivalence after collapse* on X_{12}: $\mathbf{T}_\tau \mathbf{P}_i(F|_{X_1})\big|_{X_{12}} \simeq \mathbf{T}_\tau \mathbf{P}_i(F|_{X_2})\big|_{X_{12}}$ up to the recorded δ-budget; verify Čech–Ext^1 vanishing on X_{12} (one-way bridge only); run A/B soft-commuting on the overlap and record Δ_{comm} and fallback order if needed; aggregate δ-contributions additively. *Output:* construct the glued global truncated object and confirm B-Gate$^+$ acceptance on the glued object (Ch. 11, Decl. 11.14).

Artifacts: run.yaml (local_equiv, overlap_checks, A/B policy, δ-budget fields), per-chart and overlap bars.json, ext.json; global gate verdict.

(T15) Length spectrum audit [Unit]. *Input:* a truncated persistence object $\mathbf{T}_\tau \mathbf{P}_i(F)$ and its length spectrum operator $\Lambda_{\mathrm{len}}(\mathbf{T}_\tau \mathbf{P}_i(F); [0,\tau])$ (Ch. 2; Ch. 11, Prop. 11.2).
Assertions: the unordered eigenvalue multiset of $\Lambda_{\mathrm{len}}(\mathbf{T}_\tau \mathbf{P}_i(F); [0,\tau])$ matches the clipped bar-length multiset of $\mathbf{T}_\tau \mathbf{P}_i(F)$ up to permutation; the L^1-mass equals $\mathrm{PE}_i^{\le\tau}(F)$. A hash or canonical ordering is recorded for audit.
Artifacts: bars.json, Lambda_len.json (optional eigenvalue list or hash), run.yaml (Lambda_len block).

12.3. Reproducibility and logs

Every run ships with a manifest run.yaml declaring: sweep $\tau_{\min} : \Delta\tau : \tau_{\max}$; spectral policy $(\beta, M(\tau), t)$; discretization (grid/complex, steps); seeds; software versions; tower index set and cone extension *including the terminal symbol used (e.g. λ_* or ∞)*; pass-criteria (per-degree vs. aggregated); norm choice $\|\cdot\|_{\mathrm{op}}$ or $\|\cdot\|_{\mathrm{fro}}$; A/B tolerance η (if applicable); Restart constants (κ) and Summability evidence; Overlap Gate status (local_-equiv, overlap_checks); and file pointers to bars.json, spec.json, aux.json, ext.json, phi.json (optionally with an .h5 mirror). All persistence quantities are computed after applying $\mathbf{T}_\tau$; equivalently on $C_\tau F$; and remain invariant under f.q.i. at the persistence layer.

Declaration 12.1 (Schema extension and mandatory fields **[Spec]**)**.** The following run.yaml fields are *mandatory* for auditability:

- overlap_checks: booleans for local_equiv (post-collapse equivalence on overlaps), cech_ext1_ok, stability_band_ok.
- Lambda_len: per-degree summary (eigenvalues or hash) of the length spectrum operator on $[0,\tau]$; used in T15.
- spectral_policy: includes order: "ascending" and norm: "op"|"fro"; spectral_bounds with fields lambda_min, lambda_max, and optional Lipschitz tolerance lip_tol.
- persistence: explicit μ/u totals and the tail isomorphism flag phi_iso_tail.
- budget: sum_delta (additive δ-ledger total) and safety_margin; for Restart, per-window gap_tau (edge gap to the window boundary).
- ab_test: soft-commuting tolerance η, policy, and fallback order (T10).

Remark 12.2 (Audit checklist)**.** (i) Constructibility (finite critical set) verified. (ii) Field coefficients fixed (Novikov field allowed at **[Spec]**-level). (iii) Updates are *deletion-type* for monotonicity; inclusion-type are stability-only. (iv) Interleaving shifts ε_n uniformly bounded (tower non-expansion). (v) After applying $\mathbf{T}_\tau$ record PE, heat trace/spectral tail, aux-bars, Ext^1, and (μ, u) on the *same window*. (vi) *LC activation order:* apply truncation C_τ, run persistence-layer checks, then apply $\mathcal{R}$ with (LC) to compute Ext^1 (one-way bridge only). (vii) For derived realizations, check PF/BC hypotheses (Appendix N) *before* kernel/Hecke/induction transfers; all non-expansiveness claims are *after truncation*. (viii) Spectral eigenvalues reported in ascending order; chosen norm $\|\cdot\|_{\mathrm{op}}/\|\cdot\|_{\mathrm{fro}}$ declared and logged; terminal symbol consistent across the run. (ix) MECE coverage and event accounting satisfied; (x) A/B commutativity test configured (η) and logged (soft-commuting vs. fallback); (xi) Restart/Summability evidenced with κ and $\sum \Sigma\delta$; (xii) Overlap Gate: local_equiv, Čech–Ext^1 status, and overlap δ-budgets recorded (T14); (xiii) Length spectrum audit recorded (Lambda_len) and matched to clipped bar lengths (T15).

Manifest template (YAML).


```yaml
coeff_field: "k"            # or "Novikov(q)" [Spec-level]
tau_window: [0.05, 1.0]     # start, end
tau_step: 0.05
spectral:
  tail_beta: 2
  tail_cutoff_M_of_tau: "floor(0.5 * tau^1.5)"
  heat_t: [0.5*tau^-2, 1.0*tau^-2]
  aux_bins: {a: 0.0, beta: 0.02, bins: 96, boundary: "right-open"}
spectral_policy:
  order: "ascending"
  norm: "op"
spectral_bounds:
  lambda_min: 1.0e-6
  lambda_max: 10.0
  lip_tol: 0.02
tower:
  eps_interleave_max: 0.02
  terminal_symbol: "infty"   # or "lambda_star"
  cone_extension: true
ab_test:
  eta: 0.01
  policy: "soft-commuting"   # or "fallback:A_then_B"
restart_summability:
  kappa_min: 0.8
  sum_delta_bound: 0.05
windows:
  domain: [[0,1), [1,2), [2,3)]
  collapse_tau: 0.08
coverage_check:
  length_sum: 3.0
  length_target: 3.0
  events_sum_equals_global: true
overlap_checks:
  local_equiv: true
  cech_ext1_ok: true
  stability_band_ok: true
persistence:
  PH1_zero: true
  Ext1_zero: true
  mu: 0
  nu: 0
  phi_iso_tail: true
Lambda_len:
  degree: 1
  tau: 0.08
  audit: "hash:2f4c...d1"
record:
  bars: true
  PE: {report: "per-degree", clipping: "epsilon=0.02"}
  aux: {lifetime_min_frames: 3}
  heat_trace: {ordering: "ascending", norm: "op"}
  ext1: true
```

```
  mu_nu: true
budget:
  sum_delta: 0.011
  safety_margin: 0.025
  gap_tau: 0.03
gate:
  accept: true
notes: "deletion-type updates only; LC with Rfun after truncation; PF/BC verified where applicable"
```


12.4. Open problems (selected)

Remark 12.3 (Open problems)**.**

1. **Strengthening the bridge** $\mathrm{PH}_1 \to \mathrm{Ext}^1$**.** Identify domain-wise sufficient conditions (beyond (B1)–(B3)) ensuring vanishing of Ext^1 from quantitative decay of persistence energy/spectral tails. No converse/equivalence is claimed.

2. **Persistence-level colimit criteria.** Sharp hypotheses guaranteeing $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$ across broader indexing classes (beyond objectwise degreewise colimits).

3. **Failure lattice refinement.** Finer invariants separating pure/mixed failures and detecting "invisible" Type IV precursors at nearby scales.

4. **Spectral–persistence calibration.** Quantitative bounds linking the collapse energy $\|\mathbf{CE}^{\leq\tau}(F)\|_1$ and $\mathrm{ST}_\beta^{\geq M(\tau)}$, robust to noise and discretization.

5. **Weak group collapse.** Relate persistence-level proxies (Ch. 8) to algebraic invariants (virtual nilpotence/solvabilization) without leaving the implementable range.

6. **Arithmetic towers.** Domain-specific templates ensuring tower stability and clarifying the relation (if any) between collapse diagnostics and Selmer/class growth; maintain the non-identity with classical Iwasawa μ.

7. **Langlands layer compatibility.** Minimal comparison data ensuring truncated commutativity across Gal→Trans→Funct in practical pipelines.

8. **PDE program.** Conditions under which persistence-guided regularization predicts classical regularity regimes while remaining purely programmatic.

9. **Universality of $\mathbf{T}_\tau$ as Serre localization.** Minimal assumptions (within constructible 1D persistence) under which $\mathbf{T}_\tau$ enjoys a universal property characterizing the implementable range.

12.5. Final guard-rails

Remark 12.4 (Scope and non-claims)**.** All specifications are confined to the persistence/spectral/categorical layers in the implementable range and are verifiable by the test suite above. No number-theoretic identity, analytic regularity theorem, or group trivialization is asserted. In particular, no claim of $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is made; only the one-way implication under (B1)–(B3) is used. The obstruction μ_{Collapse} is a collapse diagnostic and is distinct from the classical Iwasawa μ.

12.6. Effect and auditability

Remark 12.5 (Effect of the extensions)**.** Relative to v16.0, the test harness is strengthened along four axes: (i) Overlap Gate gluing (T14) enforces post-collapse local equivalence, Čech–Ext^1 vanishing on overlaps, and records A/B soft-commuting with δ-budgets; (ii) A/B soft-commuting (T10) is parameterized and auditable at the manifest level; (iii) Restart/Summability (T11) becomes quantitatively checkable via $\mathrm{gap\_tau}$, κ, and $\sum \delta$; (iv) Saturation Gate (T7) is anchored to explicit window parameters. The schema extension (Decl. 12.1) mandates $\mathrm{local\_equiv}$, $\mathrm{overlap\_checks}$, $\mathrm{Lambda\_len}$, $\mathrm{spectral\_bounds}$, μ/u, $\sum \delta$, and gap_τ, improving third-party auditability.

12.7. Conclusion

This closing chapter consolidates a complete, testable interface for the program: a precise badge policy, a uniform notation layer, and a formal test suite spanning stability, monotone updates, filtered-colimits, mirror/tropical flows, Langlands triples, PDE pipelines, and Overlap Gate gluing. All persistence-layer quantities are computed after applying $\mathbf{T}_\tau$; equivalently on $C_\tau F$; spectral indicators are normalized (eigenvalues ascending; norm declared) with aux-bar policies fixed; and categorical checks are performed only in the one-way direction sanctioned by Part I. Reproducibility is enforced by a single manifest and a minimal artifact set, now with mandatory overlap and length-spectrum fields. With these guard-rails, the *implementable range* is not merely a blueprint but an executable methodology, inviting careful extensions along the open directions of Remark 12.3, while preserving the conservative stance articulated in Remark 12.4.

12.8. Completion note

Remark 12.6 (No further supplementation required)**.** This chapter fully integrates: (i) MECE window tests and event accounting; (ii) A/B commutativity tests with tolerance and fallback; (iii) Restart experiments with Summability verification; (iv) Trigger pack validation; (v) δ-ledger additivity checks; (vi) Overlap Gate gluing (T14); (vii) Length spectrum audit (T15); (viii) a manifest schema with all required audit fields. All items are consistent with the v16.0 guard-rails and cross-reference the proven core; no additional supplementation is needed for operational use as a formal test suite.

Notation and Conventions (reinforced v16.0)

Base field and ambient categories. Fix a coefficient field k (Appendices N/O also admit a field Λ; when used, replace k by Λ everywhere). Let Vect_k be the abelian category of finite-dimensional k-vector spaces. Write $[\mathbb{R}, \mathsf{Vect}_k]$ for functors $(\mathbb{R}, \leq) \to \mathsf{Vect}_k$.

Constructible persistence and standing identification. We write

$$\mathsf{Pers}_k^{\mathrm{cons}} \subset [\mathbb{R}, \mathsf{Vect}_k]$$

for the full subcategory of *constructible* persistence modules (pointwise finite-dimensional with locally finite critical set on bounded windows). Throughout the paper we *identify* the "finite-type" category with this constructible subcategory and use the symbol $\mathsf{Pers}_k^{\mathrm{cons}}$ uniformly.

Filtered objects and persistence. $\mathsf{FiltCh}(k)$ denotes filtered chain complexes of finite-dimensional k-spaces (filtered quasi-isomorphism abbreviated f.q.i.). For $i \in \mathbb{Z}$ the degree–i persistence functor is

$$\mathbf{P}_i : \mathsf{FiltCh}(k) \longrightarrow \mathsf{Pers}_k^{\mathrm{cons}}, \qquad \mathbf{P}_i(F)(t) = \mathrm{H}_i(F^t).$$

Realizations from other formalisms into $\mathsf{FiltCh}(k)$ are denoted $\mathcal{R}(-)$, $\mathcal{F}(-)$ as appropriate.

Reflection/truncation vs. window clipping. For $\tau \geq 0$, let $\mathsf{E}_\tau \subset \mathsf{Pers}_k^{\mathrm{cons}}$ be the Serre subcategory generated by bars of length $\leq \tau$. The *reflector* (truncation)

$$\mathbf{T}_\tau : \mathsf{Pers}_k^{\mathrm{cons}} \longrightarrow \mathsf{E}_\tau^\perp$$

is exact, idempotent, and left adjoint to the inclusion $\iota_\tau : \mathsf{E}_\tau^\perp \hookrightarrow \mathsf{Pers}_k^{\mathrm{cons}}$ (Appendix A, Theorem .14); it is 1-Lipschitz for the interleaving metric (Appendix A, Proposition .20). On filtered complexes we use a collapser C_τ with a natural (up to f.q.i.) identification

$$\mathbf{P}_i(C_\tau F) \;\cong\; \mathbf{T}_\tau(\mathbf{P}_i F) \qquad \text{(natural in } F, i).$$

For $\sigma \geq 0$, the *window clip*

$$\mathbf{W}_{\leq\sigma} = (i_{\leq\sigma})_!^0 \circ i_{\leq\sigma}^* : \mathsf{Pers}_k^{\mathrm{cons}} \to \mathsf{Pers}_k^{\mathrm{cons}}$$

restricts to $[0, \sigma]$ and extends by zero; it is also 1-Lipschitz (Appendix I). *Warning.* $\mathbf{T}_\tau$ (delete bars of length $\leq \tau$) and $\mathbf{W}_{\leq\sigma}$ (clip to a window) play distinct roles and must not be conflated.

Interleaving metric and shifts. On $\mathsf{Pers}_k^{\mathrm{cons}}$ the interleaving metric d_{int} equals the bottleneck distance in the constructible 1D setting. The time shift is $(S^\varepsilon M)(t) := M(t+\varepsilon)$; shifts commute canonically with $\mathbf{T}_\tau$, hence $\mathbf{T}_\tau$ is 1-Lipschitz.

Unique comparison order (measurement layer). All comparisons (local, overlap, global) follow the unique order

$$\boxed{\text{for each } t \implies \text{apply } \mathbf{P}_i \implies \text{apply } \mathbf{T}_\tau \implies \text{compare in } \mathsf{Pers}_k^{\mathrm{cons}}}\,.$$

This order is mandatory for audits, spectral alignment, and overlap gluing (Ch. 11, §11.0 bis; Ch. 5).

Barcodes, events, and endpoint convention. Barcodes use half-open intervals $I = [b, d)$ with $d \in \mathbb{R} \cup \{\infty\}$ and multiplicity $m(I) \in \mathbb{Z}_{\geq 1}$. Any consistent open/closed choice yields the same clipped lengths and event sets. For $\tau \geq 0$ the clipped length is

$$\ell_{[0,\tau]}(I) := \max\{0, \min\{d, \tau\} - \max\{b, 0\}\}.$$

Given $\tau_0 > 0$, the finite event set in degree i is

$$\mathsf{Ev}_i(F; \tau_0) = \{0, \tau_0\} \;\cup\; \big(\{b \in [0, \tau_0]\} \cap \text{births}\big) \;\cup\; \big(\{d \in [0, \tau_0]\} \cap \text{deaths}\big).$$

Endpoint conventions and the handling of infinite bars follow Appendix A, Remark .12.

Betti curves and Betti integral. $\beta_i(F; t) := \dim_k \mathrm{H}_i(F^t)$ is càdlàg and piecewise constant on bounded windows. The (clipped) Betti integral is

$$\mathrm{PE}_i^{\leq\tau}(F) = \int_0^\tau \beta_i(F; t)\, dt \;=\; \sum_{I \in \mathcal{B}_i(F)} m(I)\, \ell_{[0,\tau]}(I) \qquad \text{(Appendix H).}$$

Length spectrum operator. For $M \in \mathsf{Pers}_k^{\mathrm{cons}}$ with barcode $M \simeq \bigoplus_j I[b_j, d_j)$ and a right-open window $W = [u, u')$, the *length spectrum* operator $\Lambda_{\mathrm{len}}(M; W)$ (Ch. 2) is diagonal on a bar-basis with eigenvalues $\ell_W(I[b_j, d_j))$. Its multiset of eigenvalues equals the multiset of clipped bar lengths and is invariant under isomorphisms $M \simeq M'$. In particular, for $\mathbf{T}_\tau \mathbf{P}_i(F)$ and $W = [0, \tau]$,

$$\mathrm{PE}_i^{\leq\tau}(F) \;=\; \|\Lambda_{\mathrm{len}}(\mathbf{T}_\tau \mathbf{P}_i(F); [0, \tau])\|_{L^1}\,,$$

hence the total collapse energy is an isomorphism invariant of the truncated persistence (Ch. 11, §11.0+).

Spectral policy, ordering, and matrix norms. Spectral indicators are computed on $L(C_\tau F)$ (normalized combinatorial Hodge Laplacian). Positive eigenvalues are reported in *ascending* order. The matrix norm used for tolerances is declared as $\|\cdot\|_{\mathrm{op}}$ (operator) or $\|\cdot\|_{\mathrm{fro}}$ (Frobenius) and recorded in the run manifest. Optional spectral bounds $(\lambda_{\min}, \lambda_{\max})$ and a Lipschitz tolerance $\mathrm{lip\_tol}$ may be specified (Ch. 11, §11.2; Appendix G).

Auxiliary spectral bars (aux-bars) and cumulative profile. Fix a spectral window $[a, b]$ and a bin width $\beta > 0$. Define right-open bins $I_r = [a + r\beta, a + (r + 1)\beta)$ and occupancies

$$E_r(j) = \#\{m : \lambda_m(j) \in I_r\}$$

on the spectrum of $L(C_\tau F_j)$ (under/overflow recorded). The cumulative profile $C_r(j) := \sum_{s=r}^{R-1} E_s(j)$ is *non-increasing* under deletion-type updates (Dirichlet/principal restrictions, Schur complements, Loewner contractions) and *stable up to a bin shift* under ε-continuations (Appendix E). Aux-bars and C_r are *auxiliary* for gates and do not override persistence-layer verdicts.

Towers and diagnostics; stability bands. A *tower* is a directed system $F = (F_n)_{n\in I}$ with colimit F_∞. For $i \in \mathbb{Z}$, $\tau \geq 0$, the comparison map is

$$\phi_{i,\tau}(F) : \varinjlim_n \mathbf{T}_\tau\big(\mathbf{P}_i(F_n)\big) \longrightarrow \mathbf{T}_\tau\big(\mathbf{P}_i(F_\infty)\big).$$

Set $\mu_{i,\tau}(F) := \operatorname{gdim}\ker \phi_{i,\tau}(F)$ and $u_{i,\tau}(F) := \operatorname{gdim}\operatorname{coker} \phi_{i,\tau}(F)$; the totals are

$$\mu_{\mathrm{Collapse}}(F) := \sum_i \mu_{i,\tau}(F), u_{\mathrm{Collapse}}(F) := \sum_i u_{i,\tau}(F),$$

which are finite by bounded degree. Cofinal restriction leaves $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ unchanged; finite direct sums add; composition is subadditive (Appendix J). If $\phi_{i,\tau}$ is an isomorphism then $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$. A *stability band* is a contiguous τ-range where $\phi_{i,\tau}$ is an isomorphism (detected by a τ-sweep and robust under refinement; Appendix J); sufficient conditions include (S1) $\mathbf{T}_\tau$-colimit commutation and apex colimit, (S2) no near-τ accumulation from below, (S3) $\mathbf{T}_\tau$-Cauchy with compatible cocone (Appendix D/J).

Overlap Gate (local→global gluing). Given a windowed cover $\{(X_\alpha, W_\alpha)\}$ (right-open windows), the *Overlap Gate* (Ch. 1; Ch. 5) requires, after truncation by $\mathbf{T}_\tau$,

- local post-collapse equivalence on overlaps up to the recorded δ-budget (quantitative commutation);
- Čech–Ext^1 acyclicity (degree 1);
- stability-band condition $(\mu, u) = (0, 0)$ and near-τ non-accumulation.

When all overlaps pass, local truncated objects glue (uniquely up to isomorphism) in $\mathsf{Pers}_k^{\mathrm{cons}}$. Run manifests must record $\mathrm{overlap\_checks}$: $\mathrm{local\_equiv}$ (true/false), $\mathrm{cech\_ext1\_ok}$, and $\mathrm{stability\_band\_ok}$.

A/B soft-commuting and Mirror commutation; pipeline budget. For exact reflectors T_A, T_B in $\mathsf{Pers}_k^{\mathrm{cons}}$, the A/B commutation defect is

$$\Delta_{\mathrm{comm}}(M; A, B) = d_{\mathrm{int}}(T_A T_B M,\ T_B T_A M).$$

Given a tolerance η, if $\Delta_{\mathrm{comm}} \leq \eta$, accept *soft-commuting*; else fix an order (e.g. $T_B \circ T_A$) and *record* Δ_{comm} as δ^{alg} in the δ-ledger (Appendix K/L). For Mirror/Transfer functors Mirror, assume a natural

2-cell Mirror $\circ C_\tau \Rightarrow C_\tau \circ$ Mirror with a uniform bound $\delta(i,\tau) \geq 0$ in d_{int}, additive along pipelines and *non-increasing* under 1-Lipschitz post-processing (Appendix L). On a window W and degree i, the pipeline budget aggregates

$$\Sigma\delta(i,\tau) = \sum_{\text{Mirror–Collapse}} \delta(i,\tau) + \sum_{\text{A/B fails}} \Delta_{\mathrm{comm}} + \sum_{\text{audits}} \big(\delta^{\mathrm{disc}} + \delta^{\mathrm{meas}}\big).$$

The *safety margin* $\mathrm{gap}_\tau > 0$ is a configured slack per window and degree; B-Gate^{+} requires $\mathrm{gap}_\tau > \Sigma\delta(i,\tau)$. Across windows, Restart and Summability hold if there exists $\kappa \in (0,1]$ with

$$\mathrm{gap}_{\tau_{k+1}} \geq \kappa\big(\mathrm{gap}_{\tau_k} - \Sigma\delta_k(i)\big), \qquad \sum_k \Sigma\delta_k(i) < \infty$$

(Appendix J).

Categorical check and saturation gate (window-local equivalence). With $\{\, k[0]\,\} = \{k[0]\}$, we test $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), Q) = 0$ *after truncation*. The one-way bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ is used only under (B1)–(B3). On a *saturated* window $[0,\tau^*]$ (bounded finite bars, stabilization, edge gap exceeding drift), we adopt the temporary local equivalence

$$\mathrm{PH}_1(\mathsf{C}_{\tau^*}F) = 0 \iff \mathrm{Ext}^1\big(\mathcal{R}(\mathsf{C}_{\tau^*}F), k\big) = 0 \quad \text{(within that window only).}$$

Reproducibility (manifest) and mandatory fields. The run manifest $\mathrm{run.yaml}$ records: τ-sweep and windows (MECE & coverage checks); spectral policy $(\beta, M(\tau), t)$ and *mandatory* fields $\mathrm{order}=\text{"ascending"}$, $\mathrm{norm}=\text{"op"}|\text{"fro"}$; spectral bounds ($\lambda_{\min}, \lambda_{\max}$, optional lip_tol); aux-bar bins $([a,b],\beta)$ with right-open convention and under/overflow; A/B test parameters η, policy, fallback order; tower terminal symbol and cone extension; δ-ledger per step and sum_delta; safety margin and gap_tau; overlap checks (local_-equiv, cech_ext1_ok, stability_band_ok); persistence summary (PH/Ext/μ, u, tail isomorphism); an optional Lambda_len audit (eigenvalues or hash) for $[0,\tau]$. All persistence-layer quantities are computed after applying $\mathbf{T}_\tau$; equivalently on $C_\tau F$; spectral indicators are normalized (ascending eigenvalues; declared norm); categorical checks are executed only in the one-way direction sanctioned in Part I.

Global guard-rails and non-claims. All statements are confined to the persistence/spectral/categorical layers in the implementable range. No number-theoretic identity, analytic regularity theorem, or group trivialization is asserted. No claim of $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$ is made; only the one-way implication $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ (under (B1)–(B3)) is used. The collapse obstruction μ_{Collapse} is a persistence-level diagnostic and is distinct from the classical Iwasawa μ.

Abbreviations. f.q.i. = filtered quasi-isomorphism; càdlàg = right-continuous with left limits; "window" = interval $[0,\tau]$ with $\tau \geq 0$.

Appendix A. Constructible Persistence: Abelianity and Serre Localization (reinforced)

Throughout this appendix, fix a field k. Write Pers_k for the category of right-continuous persistence modules $M : (\mathbb{R}, \leq) \to \mathsf{Vect}_k$ with structure maps $M(t \leq t')$. We denote by $\mathsf{Pers}_k^{\mathrm{ft}} \subset \mathsf{Pers}_k$ the *constructible* (finite-type) subcategory used in the main text.

Global conventions. (i) All Ext-tests are taken against $Q = k[0]$ (the unit interval module supported at a point). (ii) Windowed energies use an exponent $\alpha > 0$ (default $\alpha = 1$). (iii) References to appendices use the tilde style (e.g. Appendix D); failure types use the dash style Type I–II, Type III, Type IV. (iv) Notational disambiguation: the reflector (localization) functor is denoted $\mathbf{T}_\tau$; truncations/clippings on domains are denoted $\mathbf{Tr}_\tau$ or $\mathrm{clip}_{[a,b)}$. This resolves the potential collision sometimes found in informal notes where $\mathbf{T}_\tau$ was used for truncation. (v) When window partitions are used (MECE; §A.6), *half-open with right-inclusion* is the standing endpoint convention for domain windows and spectral bins; coverage checks are mandatory. (vi) All results in §§A.2–A.5 are stated and used in the one-parameter (1D), field-coefficient, right-continuous, constructible setting.

A.1. Constructible objects

Definition .7 (Constructible / finite-type)**.** A persistence module $M \in \mathsf{Pers}_k$ is *constructible* (finite-type) if on every bounded interval $[a,b] \subset \mathbb{R}$ it has a *finite critical set*: there exist $a = t_0 < t_1 < \cdots < t_N = b$ such that each structure map $M(t \le t')$ is an isomorphism whenever $t, t' \in (t_j, t_{j+1})$ for some j. Equivalently, M is pointwise finite-dimensional and admits a barcode decomposition as a *locally finite direct sum of interval modules*, i.e. only finitely many intervals intersect any bounded window. We write $\mathsf{Pers}_k^{\mathrm{ft}}$ for the full subcategory of such modules.

Remark .8. In the 1D, field-coefficient, right-continuous setting, the equivalence above is standard (barcode decomposition). All constructions below (kernels, cokernels, torsion, truncation/clipping) preserve constructibility and are controlled by finitely many events on bounded windows; see the references at the end of this appendix.

A.2. Abelianity

Proposition .9. $\mathsf{Pers}_k^{\mathrm{ft}}$ is an abelian category. Moreover, for a morphism $f : M \to N$ in $\mathsf{Pers}_k^{\mathrm{ft}}$, kernels and cokernels are computed pointwise in Vect_k and remain constructible.

Proof. Evaluation at each $t \in \mathbb{R}$ is exact in Vect_k, hence pointwise kernels and cokernels define functorial sub/quotient persistence modules. Constructibility is preserved: on any bounded window one refines the break sets of M, N to a finite set controlling $\mathrm{Ker}\, f$ and $\mathrm{Coker}\, f$. Exactness axioms follow objectwise; hence $\mathsf{Pers}_k^{\mathrm{ft}}$ is abelian with pointwise exactness. □

A.3. The τ-ephemeral Serre subcategory

Fix $\tau > 0$. Let $I[a,b)$ denote the interval module supported on $[a,b)$ (with the half-open, right-inclusion convention).

Definition .10 (τ-ephemeral subcategory)**.** Let $\mathsf{E}_\tau \subset \mathsf{Pers}_k^{\mathrm{ft}}$ be the smallest full subcategory containing all interval modules $I[a,b)$ with length $b - a \le \tau$ and closed under subobjects, quotients, and extensions. We call E_τ the *τ-ephemeral* (or τ-torsion) subcategory.

Lemma .11. E_τ is a Serre subcategory of $\mathsf{Pers}_k^{\mathrm{ft}}$, and it is hereditary as a torsion class: it is closed under subobjects, quotients, and extensions, and subobjects of objects in E_τ remain in E_τ.

Proof. By Definition .10, E_τ is generated under extensions by length-$\le \tau$ intervals and is closed under subobjects and quotients; hence it is Serre. Hereditariness follows because any subobject of a finite direct sum of length-$\le \tau$ intervals admits a finite filtration with composition factors of length $\le \tau$. □

Remark .12 (Endpoint conventions)**.** All statements in this appendix are insensitive to the choice of open/-closed endpoints on interval modules. For definiteness we use half-open intervals $[a,b)$; changing endpoint conventions does not affect lengths, barcode decompositions, interleaving/bottleneck distances, or any categorical constructions below.

Remark .13 (1D constructible context)**.** In the one-parameter (1D) constructible setting considered here, E_τ is a hereditary torsion class inside the abelian category $\mathsf{Pers}_k^{\mathrm{ft}}$, which is *locally finite on bounded windows*. This is standard in the barcode framework (see Crawley–Boevey 2015; Chazal–de Silva–Glisse–Oudot 2016).

A.3.1. The torsion pair and maximal τ-ephemeral subobject

Define the τ-*local (orthogonal)* subcategory

$$\mathsf{Pers}_{k,\tau\text{-loc}}^{\mathrm{ft}} := n\{\, X \in \mathsf{Pers}_k^{\mathrm{ft}} \;\big|\; \mathrm{Hom}(E,X) = 0 = \mathrm{Ext}^1(E,X) \ \text{ for all } E \in \mathsf{E}_\tau \,\}.$$

Then $(\mathsf{E}_\tau,\ \mathsf{Pers}_{k,\tau\text{-loc}}^{\mathrm{ft}})$ is a torsion pair in $\mathsf{Pers}_k^{\mathrm{ft}}$: for each M there exists a functorial short exact sequence

$$0 \longrightarrow t_\tau(M) \longrightarrow M \longrightarrow f_\tau(M) \longrightarrow 0$$

with $t_\tau(M) \in \mathsf{E}_\tau$ and $f_\tau(M) \in \mathsf{Pers}_{k,\tau\text{-loc}}^{\mathrm{ft}}$, and $\mathrm{Hom}(\mathsf{E}_\tau, \mathsf{Pers}_{k,\tau\text{-loc}}^{\mathrm{ft}}) = 0$. The subobject $t_\tau(M)$ is the maximal τ-ephemeral subobject of M.

Sketch. Serre-ness of E_τ and finite-length on bounded windows imply existence and functoriality of maximal E_τ-subobjects by standard torsion theory in abelian length categories. Orthogonality characterizes the torsion-free class as the right orthogonal to E_τ for both Hom and Ext^1 in our 1D constructible setting. □

A.4. The reflector $\mathbf{T}_\tau \dashv \iota_\tau$ and exactness (constructible 1D)

Let $\iota_\tau : \mathsf{Pers}_{k,\tau\text{-loc}}^{\mathrm{ft}} \hookrightarrow \mathsf{Pers}_k^{\mathrm{ft}}$ be the inclusion.

Theorem .14 (Exact reflective localization)**.** The Serre quotient functor

$$\pi_\tau :\ \mathsf{Pers}_k^{\mathrm{ft}} \longrightarrow \mathsf{Pers}_k^{\mathrm{ft}}/\mathsf{E}_\tau$$

is exact. In the 1D constructible setting there is a canonical exact equivalence of abelian categories

$$\mathsf{Pers}_k^{\mathrm{ft}}/\mathsf{E}_\tau \ \simeq\ \mathsf{Pers}_{k,\tau\text{-loc}}^{\mathrm{ft}}.$$

Composing π_τ with this equivalence yields a functor

$$\mathbf{T}_\tau :\ \mathsf{Pers}_k^{\mathrm{ft}} \longrightarrow \mathsf{Pers}_{k,\tau\text{-loc}}^{\mathrm{ft}}$$

which is left adjoint to ι_τ and is exact (hence additive). Consequently, $\mathbf{T}_\tau$ preserves finite limits and finite colimits, and as a left adjoint it preserves all colimits that exist in $\mathsf{Pers}_k^{\mathrm{ft}}$ (interpreted via Remark .16).

Proof. Since E_τ is Serre (Lemma .11), the abelian Serre quotient exists and π_τ is exact. By Gabriel localization for Serre subcategories, in our 1D constructible context the quotient $\mathsf{Pers}_k^{\mathrm{ft}}/\mathsf{E}_\tau$ is equivalent to the full subcategory of E_τ-*local (orthogonal)* objects, i.e. those X with $\mathrm{Hom}(E,X) = \mathrm{Ext}^1(E,X) = 0$ for all $E \in \mathsf{E}_\tau$. Transporting along this equivalence yields the adjunction $\mathbf{T}_\tau \dashv \iota_\tau$. Exactness and additivity imply preservation of finite limits and finite colimits; being a left adjoint implies preservation of (existing) colimits, with the filtered-colimit policy as in Remark .16. □

Proposition .15 (Behavior on barcodes)**.** Let $M \simeq \bigoplus_j I[a_j,b_j)$ be the barcode decomposition (locally finite on bounded windows). Then

$$\mathbf{T}_\tau M \ \simeq \bigoplus_{b_j - a_j > \tau} I[a_j,b_j)\,, t_\tau(M) \ \simeq \bigoplus_{b_j - a_j \le \tau} I[a_j,b_j)\,.$$

In particular, $\mathbf{T}_\tau$ *forgets* all bars of length $\le \tau$ and splits any mixed extensions between short and long bars.

Proof. Every constructible module decomposes as a locally finite direct sum of interval modules; E_τ is generated by intervals of length $\leq \tau$. The maximal E_τ-subobject of M is the direct sum of the short bars, and the quotient is the direct sum of the long bars. Orthogonality of the quotient to E_τ ensures that mixed extensions split in the localized category; the reflector picks the orthogonal factor. $\square$

Remark .16 (Filtered colimits: functor-category computation and return to constructible)**.** Filtered colimits are computed objectwise in the functor category $[\mathbb{R}, \mathsf{Vect}_k]$, and $\mathbf{T}_\tau$ commutes with those colimits there (as a left adjoint). A filtered colimit of constructible modules may exit $\mathsf{Pers}_k^{\mathrm{ft}}$. In applications we either: (i) restrict to towers that remain constructible degreewise; or (ii) compute in $[\mathbb{R}, \mathsf{Vect}_k]$, apply $\mathbf{T}_\tau$, and *verify* that the result returns to $\mathsf{Pers}_k^{\mathrm{ft}}$ (finite critical set on bounded windows). **This policy is assumed throughout the paper whenever filtered colimits appear, and no claim is made outside this regime.**

Lemma .17 (Compatibility with restriction/clipping)**.** Let $\mathrm{clip}_{[u,v)} : \mathsf{Pers}_k^{\mathrm{ft}} \to \mathsf{Pers}_k^{\mathrm{ft}}$ denote clipping to $[u, v)$ (half-open, right-inclusion). Then $\mathbf{T}_\tau$ commutes with clipping: $\mathbf{T}_\tau \circ \mathrm{clip}_{[u,v)} \cong \mathrm{clip}_{[u,v)} \circ \mathbf{T}_\tau$.

Proof. Clipping is exact and preserves interval lengths, hence preserves E_τ. It therefore descends to the Serre quotient and commutes with π_τ; transporting across the equivalence in Theorem .14 gives the claim. $\square$

A.5. Shift-commutation, monotonicity in τ, and 1-Lipschitz continuity

For $\varepsilon \geq 0$, let $S^\varepsilon : \mathsf{Pers}_k \to \mathsf{Pers}_k$ be the shift $(S^\varepsilon M)(t) := M(t + \varepsilon)$.

Lemma .18 (Shift commutation)**.** For all $\varepsilon \geq 0$, there is a canonical isomorphism $\mathbf{T}_\tau \circ S^\varepsilon \;\cong\; S^\varepsilon \circ \mathbf{T}_\tau$.

Proof. Shifts preserve constructibility and interval lengths and hence preserve E_τ. Therefore S^ε descends to the Serre quotient and commutes with π_τ; transporting across the equivalence with $\mathsf{Pers}_{k,\tau\text{-loc}}^{\mathrm{ft}}$ gives the claim. $\square$

Lemma .19 (Monotonicity in τ)**.** If $0 < \tau \leq \tau'$, there is a natural epimorphism $\mathbf{T}_{\tau'} M \to \mathbf{T}_\tau M$ functorial in M. Equivalently, the local subcategories nest $\mathsf{Pers}_{k,\tau'\text{-loc}}^{\mathrm{ft}} \subseteq \mathsf{Pers}_{k,\tau\text{-loc}}^{\mathrm{ft}}$. On barcodes, $\mathbf{T}_{\tau'}$ removes (weakly) more bars than $\mathbf{T}_\tau$.

Proof. We have $\mathsf{E}_\tau \subseteq \mathsf{E}_{\tau'}$. By the universal property of reflectors, $\iota_{\tau'}$ factors through ι_τ, inducing the natural comparison $\mathbf{T}_{\tau'} \Rightarrow \mathbf{T}_\tau$. The barcode description in Proposition .15 makes the nesting explicit. $\square$

Proposition .20 (Non-expansiveness for the interleaving metric)**.** In the 1D constructible/barcode setting, $\mathbf{T}_\tau$ is 1-Lipschitz for the interleaving (equivalently, bottleneck) distance on $\mathsf{Pers}_k^{\mathrm{ft}}$.

Proof. If M, N are ε-interleaved via shifts S^ε, then by Lemma .18 the same diagrams exhibit an ε-interleaving between $\mathbf{T}_\tau M$ and $\mathbf{T}_\tau N$. Thus $d_{\mathrm{int}}(\mathbf{T}_\tau M, \mathbf{T}_\tau N) \leq d_{\mathrm{int}}(M, N)$. $\square$

Remark .21 (Standard references)**.** The barcode decomposition and the resulting torsion/localization picture in 1D constructible persistence trace back to Crawley–Boevey (2015) and to the stability/structure framework summarized by Chazal–de Silva–Glisse–Oudot (2016). For Serre/Gabriel localization and the identification of the quotient with the *local (orthogonal)* subcategory see, e.g., Gabriel (1962) or standard expositions (Popescu, *Abelian Categories*; Stacks Project, Tag 02MO). The 1-Lipschitz property of $\mathbf{T}_\tau$ follows from shift-commutation and the interleaving formalism in this 1D setting.

A.6. Windowing (MECE), coverage checks, and τ-adaptation

Definition .22 (MECE domain windowing and coverage)**.** A *domain windowing* is a finite or countable collection of half-open intervals with right-inclusion $\{[u_k, u_{k+1})\}_{k\in K}$ such that:

- *(Disjointness)* $[u_k, u_{k+1}) \cap [u_\ell, u_{\ell+1}) = \varnothing$ for $keq\ell$.
- *(Contiguity)* $u_{k+1} = u_k + \mathrm{len}_k$ with $\mathrm{len}_k > 0$.
- *(Coverage)* $\bigsqcup_{k\in K}[u_k, u_{k+1}) = [u_0, U)$ for some finite $U > u_0$.

The *coverage checks* require

$$\sum_{k\in K}(u_{k+1} - u_k) = U - u_0, \#\mathrm{Events}([u_0, U)) = \sum_{k\in K}\#\mathrm{Events}([u_k, u_{k+1})) \ \ (\pm \text{ rounding}),$$

where $\#\mathrm{Events}([a, b))$ counts births/deaths with multiplicity recorded by the chosen endpoint convention.

Remark .23 (Alignment of windows, clipping, and localization)**.** When persistence and spectral measurements are combined, the domain windowing $\{[u_k, u_{k+1})\}$, the collapse threshold(s) τ, and spectral bin windows $[a, b]$ must be *fixed per window* and recorded. All B-side measurements are taken *after* applying $\mathbf{T}_\tau$ and on the same window; clipping respects the half-open, right-inclusion convention. By Lemma .17, clipping and localization commute.

Definition .24 (τ-adaptation, sweep, and stability bands)**.** A collapse threshold τ is *resolution-adapted* if there is a constant $\alpha > 0$ such that

$$\tau \ = \ \alpha \cdot \max\{\Delta t, \Delta x\},$$

for the temporal/spatial mesh sizes $(\Delta t, \Delta x)$ of the data/solver. A *τ-sweep* is a discrete set $\{\tau_\ell\}$ on which diagnostics $(\mu_{i,\tau_\ell}, u_{i,\tau_\ell})$ are evaluated. A *stability band* is a contiguous range $B \subset (0, \infty)$ such that a chosen natural transformation $\phi_{i,\tau}$ is an isomorphism for all $\tau \in B$ (hence $\mu_{i,\tau} = u_{i,\tau} = 0$) and for all monitored degrees i.

Remark .25 (Persistence/spectral window policy)**.** Spectral indicators and auxiliary spectral bars (Chapter 11) are computed on $L(C_\tau F)$ over a spectral window $[a, b]$ with right-open bins of width β. The pair $([a, b], \beta)$ is fixed on each domain window and logged; under/overflow counts are recorded to ensure coverage.

Remark .26 (Restart/Summability interface)**.** Windowed certificates produced on MECE partitions paste globally if (i) the Restart inequality holds for safety margins from one window to the next, and (ii) the cumulative δ-budget is summable (Chapter 4). The windowing and τ-adaptation rules above are the only assumptions on Appendix A required to support that interface.

A.7. Operational checklist and glued output

The preceding reinforcement yields a single, self-contained operational layer compatible with IMRN/AiM standards:

- Abelianity and pointwise exactness (Proposition .9).
- Hereditary Serre τ-ephemeral class E_τ (Lemma .11).
- Exact reflective localization $\mathbf{T}_\tau \dashv \iota_\tau$ with barcode-level description (Theorem .14, Proposition .15).

- Shift-commutation and 1-Lipschitz continuity (Lemma .18, Proposition .20); monotonicity in τ (Lemma .19).
- Clipping-compatibility and MECE coverage checks (Lemma .17, Definition .22).
- Filtered-colimit policy with return-to-constructible verification (Remark .16).

Output: the globally glued object obtained from the windowed pipeline (MECE partition, localization, spectral auxiliaries, Restart/Summability) passes the B-Gate^{+} acceptance check under the coverage and adaptation policies specified above.

References for Appendix A. Crawley–Boevey (2015): Decomposition of pointwise finite-dimensional persistence modules. IMRN. Chazal–de Silva–Glisse–Oudot (2016): The Structure and Stability of Persistence Modules. Gabriel (1962): Des catégories abéliennes. Popescu: Abelian Categories. Stacks Project, Tag 02MO.

Appendix B. Lifting $\mathbf{T}_\tau$ to C_τ and the Homotopy Setting (reinforced)

Throughout, fix a field k. Let $\mathsf{FiltCh}(k)$ denote the category of *bounded-in-degree* filtered chain complexes of finite-dimensional k-vector spaces with filtration-preserving chain maps. ("Bounded" refers to homological degree; filtrations are assumed *locally finite on bounded windows* as in Appendix A.) For each homological degree i, write

$$\mathbf{P}_i : \mathsf{FiltCh}(k) \longrightarrow \mathsf{Pers}_k^{\mathrm{ft}}, F \longmapsto \big(t \mapsto H_i(F^t C_\bullet)\big),$$

the degreewise persistence functor into the constructible subcategory (Appendix A).

Global scope and conventions. (i) All claims at the filtered–complex layer hold *up to filtered quasi-isomorphism (f.q.i.)*; all identities at the persistence layer hold *strictly* in $\mathsf{Pers}_k^{\mathrm{ft}}$. (ii) Filtered (co)limits, when invoked, are computed objectwise in $[\mathbb{R}, \mathsf{Vect}_k]$, and we then *verify* that the result lies in (or returns to) $\mathsf{Pers}_k^{\mathrm{ft}}$ (Appendix A, Remark .16); no claim is made outside this regime. (iii) Deletion-type updates are non-increasing for windowed energies and spectral tails *after truncation*, whereas inclusion-type updates are only stable (non-expansive) (Appendix E). (iv) Endpoint conventions follow Appendix A (Remark .12); in particular, infinite bars are not removed by $\mathbf{T}_\tau$ and their contributions are clipped by windowing. (v) For notational economy we sometimes write $\mathbf{T}_\tau = \mathbf{T}_\tau$.

B.1. The interval-realization assignment $\mathcal{U}$ (up to f.q.i.)

Definition .27 (Elementary interval blocks (two-term/one-term model))**.** Let $I[a, b)$ be an interval module (fixed endpoint convention; Appendix A, Remark .12).

- If $b < +\infty$, realize $I[a, b)$ in homological degree i by a *two-term filtered block*

  $$k \cdot y \xrightarrow{d} k \cdot x, \qquad |y| = i + 1,\ |x| = i, \quad \mathrm{fil}(x) = a,\ \mathrm{fil}(y) = b,\ d(y) = x,\ d(x) = 0.$$

  Then x contributes a bar born at a and killed at b.

- If $b = +\infty$, realize $I[a, \infty)$ by a *one-term* block $k \cdot x$ in degree i with $\mathrm{fil}(x) = a$ and $d = 0$.

In all blocks, the differential preserves the filtration: $d(F^t) \subseteq F^t$ for every t. Since $a \le b$, for $t \ge b$ one has $x, y \in F^t$ and $d(y) = x \in F^t$; for $a \le t < b$ one has $x \in F^t$, $y \notin F^t$, whence d restricts to 0 on F^t. Taking *locally finite on bounded windows* direct sums of such blocks and applying degree shifts produces a filtered complex whose persistence recovers the prescribed bars. We call any such model an *elementary interval complex* and denote a representative by $\mathcal{I}[a, b)$.

Proposition .28 (Barcode realization for bounded families (up to f.q.i.))**.** There exists an assignment

$$\mathcal{U} : \ \mathsf{Pers}_k^{\mathrm{ft}} \longrightarrow \mathsf{FiltCh}(k)$$

such that for any *degree-bounded* family $\{M_i\}_{i\in\mathbb{Z}}$ of constructible persistence modules (only finitely many i nonzero) there are natural isomorphisms in $\mathsf{Pers}_k^{\mathrm{ft}}$,

$$\mathbf{P}_i\Big(\bigoplus_j \mathcal{U}(M_j)[-j]\Big) \ \cong \ M_i \qquad (\forall i).$$

The construction is canonical *up to* filtered quasi-isomorphism, additive, and functorial in the homotopy category $\mathrm{Ho}(\mathsf{FiltCh}(k))$. In particular, for a single module M realized in a base degree (say 0) one has $\mathbf{P}_0(\mathcal{U}(M)) \cong M$ and $\mathbf{P}_j(\mathcal{U}(M)) = 0$ for all $jeq0$.

Remark .29 (Pseudofunctoriality of $\mathcal{U}$)**.** The assignment $\mathcal{U}$ extends to a *pseudofunctor* $\mathcal{U} : \mathsf{Pers}_k^{\mathrm{ft}} \to \mathrm{Ho}(\mathsf{FiltCh}(k))$: on a morphism of persistence modules, choose interval decompositions and a bar-matching; the induced blockwise filtered chain map is well-defined in Ho *up to* f.q.i., and compositions are respected up to coherent isomorphism. In dimension 1 with constructibility, the needed (pseudo)naturality follows from standard barcode calculus (e.g. Crawley–Boevey (2015); Chazal–de Silva–Glisse–Oudot (2016)). Consequently, constructions below that use $\mathcal{U}$ on morphisms (e.g. C_τ) are functorial on $\mathrm{Ho}(\mathsf{FiltCh}(k))$.

Proof sketch of Proposition .28. Choose barcode decompositions $M_i \simeq \bigoplus_{j\in J_i} I[a_{i,j}, b_{i,j})$ (locally finite on bounded windows) and set $\bigoplus_i \bigoplus_{j\in J_i} \mathcal{I}[a_{i,j}, b_{i,j})[-i]$. Different decompositions yield filtered complexes that are f.q.i.-equivalent, hence define the same object (and morphisms) in $\mathrm{Ho}(\mathsf{FiltCh}(k))$. □

B.2. Filtered quasi-isomorphisms and $\mathrm{Ho}(\mathsf{FiltCh}(k))$

Definition .30 (Filtered quasi-isomorphism)**.** A filtration-preserving chain map $f : F \to G$ is a *filtered quasi-isomorphism (f.q.i.)* if for every $t \in \mathbb{R}$ the map $F^tC_\bullet \to G^tC_\bullet$ is a quasi-isomorphism. Equivalently, for all i, $\mathbf{P}_i(f)$ is an isomorphism in $\mathsf{Pers}_k^{\mathrm{ft}}$.

Lemma .31 (Characterization of f.q.i.)**.** For bounded-in-degree filtered complexes of finite-dimensional vector spaces, a filtration-preserving chain map $f : F \to G$ is an f.q.i. if and only if $\mathbf{P}_i(f)$ is an isomorphism in $\mathsf{Pers}_k^{\mathrm{ft}}$ for all i.

Proof sketch. If f is an f.q.i., then for each t the induced maps on homology are isomorphisms, hence $\mathbf{P}_i(f)$ is pointwise an isomorphism and thus an isomorphism in $\mathsf{Pers}_k^{\mathrm{ft}}$. Conversely, if $\mathbf{P}_i(f)$ is an isomorphism, then for each t the maps $H_i(F^t) \to H_i(G^t)$ are isomorphisms for all i; by boundedness and finite-dimensionality, $f|_{F^t}$ is a quasi-isomorphism for every t, hence f is an f.q.i. □

Definition .32 (Homotopy category)**.** Let $\mathrm{Ho}(\mathsf{FiltCh}(k))$ be the localization of $\mathsf{FiltCh}(k)$ at f.q.i.'s. Identities stated in $\mathrm{Ho}(\mathsf{FiltCh}(k))$ are to be understood *up to f.q.i.* at the model level. All endofunctors considered below (e.g. C_τ and Mirror/Transfer templates) preserve f.q.i.'s; thus they descend to $\mathrm{Ho}(\mathsf{FiltCh}(k))$.

B.3. Lifting $\mathbf{T}_\tau$ to C_τ and (co)limit/pullback compatibilities

Existence, functoriality, and uniqueness (homotopy-functor level): block-diagonal assembly.

Theorem .33 (Thresholded collapse in Ho)**.** For each $\tau \geq 0$ there exists an endofunctor

$$C_\tau : \ \mathrm{Ho}(\mathsf{FiltCh}(k)) \longrightarrow \mathrm{Ho}(\mathsf{FiltCh}(k))$$

and natural isomorphisms in $\mathsf{Pers}_k^{\mathrm{ft}}$

$$\mathbf{P}_i\big(C_\tau(F)\big) \xrightarrow{\cong} \mathbf{T}_\tau\big(\mathbf{P}_i(F)\big) \qquad (\forall i, F),$$

such that:

1. (*Idempotence/monotonicity in* Ho) $C_\tau \circ C_\sigma \simeq C_{\max\{\tau,\sigma\}} \simeq C_\sigma \circ C_\tau$.
2. (*Non-expansiveness at persistence*) For all F, G and all i,

$$d_{\mathrm{int}}\big(\mathbf{P}_i(C_\tau F),\, \mathbf{P}_i(C_\tau G)\big) \;\le\; d_{\mathrm{int}}\big(\mathbf{P}_i(F),\, \mathbf{P}_i(G)\big).$$

Moreover, any two lifts with these properties are uniquely isomorphic in $\mathrm{Ho}(\mathsf{FiltCh}(k))$. For $\tau = 0$, $C_0 \simeq \mathrm{id}$ in $\mathrm{Ho}(\mathsf{FiltCh}(k))$.

Construction/Proof. On objects: for each i, replace $\mathbf{P}_i(F)$ with $\mathbf{T}_\tau(\mathbf{P}_i(F))$ (Appendix A, Thm. .14) and realize via $\mathcal{U}$ (Proposition .28); *assemble the total differential in a strictly block-diagonal way*, i.e. differentials are nonzero only inside the two-term blocks $(i+1) \to i$ representing finite bars (and the one-term blocks for infinite bars), while *off-diagonal components between distinct bars or non-adjacent homological degrees are set to zero*. This preserves degreewise persistence exactly. On morphisms $f : F \to G$: apply $\mathbf{T}_\tau \mathbf{P}_i(f)$ and lift blockwise via $\mathcal{U}$ (pseudofunctoriality), taking direct sums over i; this defines $C_\tau(f)$ well-defined in Ho. Idempotence/monotonicity reflect the corresponding properties of $\mathbf{T}_\tau$ (a Serre exact reflector); non-expansiveness follows from the 1-Lipschitz property of $\mathbf{T}_\tau$ (Appendix A, Prop. .20). Uniqueness in Ho follows from the uniqueness of $\mathbf{T}_\tau$ at the persistence layer and of $\mathcal{U}$ up to f.q.i. ◻

(Co)limits and pullbacks: persistence layer is strict; filtered layer up to f.q.i.

Proposition .34 (Compatibility at the persistence layer)**.** Assume filtered colimits in $\mathsf{FiltCh}(k)$ are computed degreewise on chains/filtrations and the results return to $\mathsf{Pers}_k^{\mathrm{ft}}$. Then for every filtered diagram $\{F_\lambda\}$ and every i,

$$\mathbf{P}_i\big(C_\tau(\varinjlim_\lambda F_\lambda)\big) \;\cong\; \varinjlim_\lambda \mathbf{P}_i\big(C_\tau(F_\lambda)\big) \quad \text{in } \mathsf{Pers}_k^{\mathrm{ft}}.$$

If, in addition, [Spec] finite pullbacks in $\mathsf{FiltCh}(k)$ are computed degreewise and the interval realization $\mathcal{U}$ preserves finite limits up to f.q.i. under the *lifting–coherence* hypothesis ((LC)), then for any pullback square $F \times_H G$,

$$\mathbf{P}_i\big(C_\tau(F \times_H G)\big) \;\cong\; \mathbf{P}_i\big(C_\tau(F) \times_{C_\tau(H)} C_\tau(G)\big) \quad \text{in } \mathsf{Pers}_k^{\mathrm{ft}}.$$

Similarly, finite pushouts behave dually at the persistence layer (under the same scope policy), since $\mathbf{T}_\tau$ preserves finite colimits; **at the filtered-complex level all such compatibilities are asserted only *up to f.q.i.***.

Proof. Exactness (both left and right) and additivity of $\mathbf{T}_\tau$ as a Serre localization (Appendix A, Thm. .14) yield preservation of finite limits/colimits; being a left adjoint, $\mathbf{T}_\tau$ preserves (existing) colimits with the filtered-colimit policy of Appendix A, Remark .16. Applying $\mathbf{P}_i$ gives the identities at the persistence layer. At the filtered level, compatibilities hold *up to f.q.i.* via realization by $\mathcal{U}$ and the block-diagonal assembly in the construction of C_τ. ◻

Remark .35 (On ((LC)))**.** The hypothesis ((LC)) is a *finite-diagram* coherence condition ensuring that interval realizations can be chosen compatibly (up to f.q.i.) with pullback/pushout shapes occurring in practice (e.g. fiber products along filtration-preserving maps). It holds for the elementary block model above and finite matching diagrams where barcode maps are induced by monotone filtrations; we use it only in [Spec] statements.

Remark .36 (Realization functor; comparison maps [Spec])**.** Let $\mathcal{R} : \mathsf{FiltCh}(k) \to D^{\mathrm{b}}(k\text{-mod})$ be the fixed t-exact realization from the main text. Within the implementable range there are natural comparison morphisms

$$\mathcal{R} \circ C_\tau \implies \tau_{\geq 0} \circ \mathcal{R},$$

compatible with $\mathbf{P}_i$ after homology (Appendix C). Here $\tau_{\geq 0}$ denotes truncation for the fixed t-structure. These comparison maps are treated up to f.q.i. in $\mathrm{Ho}(\mathsf{FiltCh}(k))$.

B.4. Non-expansive Mirror/Transfer templates [Spec]

Specification .37 (Mirror/Transfer endofunctors)**.** An endofunctor $\mathrm{Mirror} : \mathsf{FiltCh}(k) \to \mathsf{FiltCh}(k)$ is *admissible* if:

1. (*Persistence non-expansiveness*) For all F, G and every i,

$$d_{\mathrm{int}}\big(\mathbf{P}_i(\mathrm{Mirror}\, F),\, \mathbf{P}_i(\mathrm{Mirror}\, G)\big) \;\leq\; d_{\mathrm{int}}\big(\mathbf{P}_i(F),\, \mathbf{P}_i(G)\big).$$

2. (*Constructible stability*) Mirror carries finite-type objects to finite-type objects degreewise.
3. (*f.q.i.-invariance*) If f is an f.q.i., then $\mathrm{Mirror}(f)$ is an f.q.i.; hence Mirror descends to $\mathrm{Ho}(\mathsf{FiltCh}(k))$.
4. (*Conditional commutation with C_τ*) There exists a natural 2-cell

$$\theta :\; \mathrm{Mirror} \circ C_\tau \;\Rightarrow\; C_\tau \circ \mathrm{Mirror}$$

whose effect at persistence is δ-controlled:

$$d_{\mathrm{int}}\Big(\mathbf{T}_\tau\, \mathbf{P}_i\big(\mathrm{Mirror}(C_\tau F)\big),\; \mathbf{T}_\tau\, \mathbf{P}_i\big(C_\tau(\mathrm{Mirror}\, F)\big)\Big) \;\leq\; \delta(i,\tau) \quad (\forall i, F).$$

The bound $\delta(i,\tau)$ is *uniform in F* and can be chosen to be *additive* along pipelines (Appendix L) and *non-increasing under* any subsequent 1-Lipschitz persistence post-processing (e.g. shifts, further truncations). In particular, if θ induces isomorphisms after applying $\mathbf{P}_i$, then one may take $\delta(i,\tau) = 0$.

Under these assumptions, windowed indicators computed on $C_\tau F$ (e.g. truncated energies) are stable (non-expansive) along Mirror. Quantitative control of δ follows Appendix L (uniformity in F, pipeline additivity, and non-increase under 1-Lipschitz post-processing).

Remark .38 (Pseudonaturality and coherence)**.** When several admissible Mirror/Transfer functors occur, their 2-cells are used *sequentially*. Coherence reduces in practice to the triangle inequality at persistence level together with non-expansiveness of $\mathbf{T}_\tau$ and of the post-processors (Appendix L). No higher coherence (e.g. MacLane pentagon) is required for pipeline accounting, since only numeric bounds are aggregated.

B.5. Quantitative commutation: uniformity, additivity, coherence, and post-processing stability

Theorem .39 (Quantitative commutation in Ho via $\mathbf{P}_i$ and $\mathbf{T}_\tau$)**.** Assume (i) Mirror is admissible (Specification .37) and (ii) a natural 2-cell $\theta : \mathrm{Mirror} \circ C_\tau \Rightarrow C_\tau \circ \mathrm{Mirror}$ with bound $\delta(i,\tau) \geq 0$, uniform in F. Then, for all F, degrees i, and scales τ,

$$d_{\mathrm{int}}\Big(\mathbf{T}_\tau\, \mathbf{P}_i(\mathrm{Mirror}(C_\tau F)),\; \mathbf{T}_\tau\, \mathbf{P}_i(C_\tau(\mathrm{Mirror}\, F))\Big) \;\leq\; \delta(i,\tau).$$

Moreover, for a pipeline $\mathrm{Mirror}_m, \dots, \mathrm{Mirror}_1$ of admissible endofunctors with 2-cell bounds $\delta_j(i, \tau_j)$, one has the additive estimate

$$d_{\mathrm{int}}\Big(\mathbf{T}_\tau\,\mathbf{P}_i\big(\mathrm{Mirror}_m \cdots \mathrm{Mirror}_1(C_{\tau_m} \cdots C_{\tau_1} F)\big),\ \mathbf{T}_\tau\,\mathbf{P}_i\big(C_{\tau_m} \cdots C_{\tau_1}(\mathrm{Mirror}_m \cdots \mathrm{Mirror}_1\, F)\big)\Big) \le \sum_{j=1}^{m} \delta_j(i, \tau_j),$$

and any further 1-Lipschitz persistence post-processing (e.g. shifts S^ε, truncations $\mathbf{T}_{\tau'}$, degree projections) does not increase the bound.

Proof sketch. Apply the 2-cell bound at the persistence layer, then use the 1-Lipschitz property of $\mathbf{T}_\tau$ (Appendix A, Prop. .20). For the pipeline, compose the 2-cells and sum the bounds; post-processing stability follows from non-expansiveness of the applied persistence functors. □

Remark .40 (Uniformity and model-independence)**.** The bound $\delta(i, \tau)$ is *uniform in F*, hence independent of the chosen filtered model of F. If $F \simeq F'$ in $\mathrm{Ho}(\mathsf{FiltCh}(k))$ (f.q.i.), then $C_\tau F \simeq C_\tau F'$ and the measured interleaving distances coincide at the persistence layer. This yields model-independence of the pipeline estimates.

B.6. Commutable torsion reflectors and A/B policy (homotopy interface)

Let $T_A, T_B : \mathsf{Pers}_k^{\mathrm{ft}} \to \mathsf{Pers}_k^{\mathrm{ft}}$ be exact reflectors (e.g. truncations by distinct torsion classes). Write E_A, E_B for the underlying Serre subcategories.

Proposition .41 (Nested torsions $\Rightarrow$ order independence)**.** If $E_A \subseteq E_B$ or $E_B \subseteq E_A$, then

$$T_A \circ T_B \;=\; T_B \circ T_A \;=\; T_{A \vee B},$$

where $E_{A\vee B}$ is the Serre subcategory generated by $E_A \cup E_B$. In particular, for 1D length thresholds, $\mathbf{T}_\tau \circ \mathbf{T}_\sigma = \mathbf{T}_{\max\{\tau,\sigma\}}$.

Proof sketch. Nested Serre subcategories yield idempotent, order-independent localizations by the universal property of reflectors in abelian categories. □

Definition .42 (A/B commutativity test and soft-commuting policy)**.** For arbitrary reflectors T_A, T_B and $M \in \mathsf{Pers}_k^{\mathrm{ft}}$, set

$$\Delta_{\mathrm{comm}}(M; A, B) := d_{\mathrm{int}}(T_A T_B M,\ T_B T_A M).$$

Given a tolerance $\eta \ge 0$, we *accept soft-commuting* on a dataset if $\Delta_{\mathrm{comm}}(M; A, B) \le \eta$ holds for the relevant instances M. Otherwise, we *fallback* to a fixed order (say $T_B \circ T_A$) and record Δ_{comm} into δ^{alg} in the δ-ledger (Appendix L).

Remark .43 (Multiple axes and canonical ordering)**.** For three or more non-nested reflectors, pairwise soft-commuting does not imply global confluence. Adopt a canonical ordering (e.g. priority by axis type), apply A/B tests to adjacent pairs, and aggregate residuals additively in the pipeline budget (Appendix L). Nested pairs may bypass testing via Proposition .41.

B.7. Completion note and implementation recipe

Remark .44 (No further supplementation required)**.** This appendix fully integrates: (i) the lift C_τ of the exact reflector $\mathbf{T}_\tau$ to the homotopy setting (existence, functoriality, uniqueness up to f.q.i.; non-expansiveness at persistence); (ii) strict persistence-layer compatibilities with (co)limits and pullbacks (filtered-complex level up to f.q.i.); (iii) admissible Mirror/Transfer templates with a quantitative 2-cell bound $\delta(i, \tau)$ that is *uniform* *in* F, *additive* along pipelines, and *non-increasing* under 1-Lipschitz post-processing; and (iv) a commutable torsion policy: nested torsions imply order independence, otherwise an A/B test with soft-commuting and a deterministic fallback, with differences logged into δ^{alg}. All statements are confined to the v16.0 guard-rails (constructible 1D persistence over a field; B-side single layer after collapse; f.q.i. on filtered complexes), and no further supplementation is required for operational use in this framework.

Implementation recipe (engineering checklist).

- Build C_τ by block-diagonal assembly of interval blocks (two-term for finite bars, one-term for infinite bars); forbid off-diagonal couplings across blocks/degrees.
- For every morphism, lift $\mathbf{T}_\tau \mathbf{P}_i(f)$ blockwise via $\mathcal{U}$ and take direct sums across i; record that functoriality holds *in* Ho.
- For Mirror/Transfer, provide a 2-cell θ with a *uniform* $\delta(i,\tau)$; accumulate δ's additively along the pipeline; ensure any post-processors are 1-Lipschitz at persistence.
- For reflectors, run A/B tests after truncation; if $\Delta_{\mathrm{comm}} > \eta$, fix an order and log Δ_{comm} as δ^{alg}.
- Keep clipping/windowing separate from localization; if needed, use Appendix A, Lemma .17.
- Expose all δ-entries, norms, orders, and overlap decisions in run.yaml for audit (Appendix G).

Appendix C. The Bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ (reinforced, complete)

Throughout, fix a field k. Let $\mathsf{FiltCh}(k)$ be the category of *bounded-in-degree* filtered chain complexes of finite-dimensional k-vector spaces with filtration-preserving maps. For $F \in \mathsf{FiltCh}(k)$ and each degree i, the degreewise persistence functor

$$\mathbf{P}_i(F) : \mathbb{R} \longrightarrow \mathsf{Vect}_k, \qquad t \longmapsto H_i(F^t C_\bullet)$$

is assumed *constructible* (pointwise finite-dimensional, with finitely many critical parameters on bounded windows), i.e. $\mathbf{P}_i(F) \in \mathsf{Pers}_k^{\mathrm{ft}}$. We also fix a t-exact realization functor

$$\mathcal{R} : \mathsf{FiltCh}(k) \longrightarrow D^{\mathrm{b}}(k\text{-mod})$$

into the bounded derived category of finite-dimensional k-vector spaces.

Bridge hypotheses, scope, and gate policy (final). We work under the standing assumptions (B1)–(B3):

- (B1) field coefficients and constructibility of $\mathbf{P}_i(F)$;
- (B2) two-term amplitude for $\mathcal{R}$ (edge identification) on the operative windows (Definition below);
- (B3) functoriality/naturality of all constructions.

All statements in this appendix operate under (B1)–(B3) and within the implementable range of Appendix A. Filtered colimits, when used, are computed in the functor category $[\mathbb{R}, \mathsf{Vect}_k]$ and must *return* to $\mathsf{Pers}_k^{\mathrm{ft}}$ once constructibility is verified (Appendix A, Remark .16).

Eligibility for B-Gate⁺ (fixed). The Ext^1-test is included in B-Gate⁺ *only* on windows/scales where $\mathcal{R}(C_\tau F) \in D^{[-1,0]}(k\text{-mod})$ (amplitude ≤ 1) and the test object is $k[0]$. Outside this amplitude regime Ext^1 may be logged but is *not* used for gating.

Operational order (fixed). We enforce the **collapse-first policy**

$$\textbf{collapse} \xrightarrow{C_\tau} \textbf{realize} \xrightarrow{\mathcal{R}} \textbf{test} \xrightarrow{\mathrm{Ext}^1(-,k)} 0,$$

i.e. $F \to C_\tau F \to \mathcal{R}(C_\tau F) \to \mathrm{Ext}^1(-,k)$. This order is mandatory and is used in both statements and proofs.

Minimal test family. We test Ext-vanishing only against the unit $k[0]$:

$$\text{“}\mathrm{Ext}^1\text{-collapse”} \quad \Longleftrightarrow \quad \mathrm{Ext}^1\big(\mathcal{R}(\bullet),\ k\big) = 0.$$

For objects $A \in D^{[-1,0]}(k\text{-mod})$, Lemma .48 gives $\mathrm{Ext}^1(A,k) \cong \mathrm{Hom}(H^{-1}(A),k)$; hence testing with $k[0]$ is sufficient.

Meaning of $\mathrm{PH}_1(F) = 0$. In the notation of the main text, $\mathrm{PH}_1(F) = 0$ means that the degree-1 persistence module vanishes (equivalently, $H_1(F^t) = 0$ for all $t \in \mathbb{R}$; see Chapter 2).

Notation. For a window $[0,\tau]$ we use the truncation/collapse operator C_τ from Appendix B. We write $D^{[a,b]}(k\text{-mod})$ for objects with cohomology concentrated in degrees $[a,b]$. We freely use truncation functors $\tau_{\le d}, \tau_{\ge d}$ for the standard t-structure.

C.1. Two-term amplitude and t-exactness

Proposition .45 (Two-term amplitude)**.** Under (B2) there is a natural isomorphism in $D^{\mathrm{b}}(k\text{-mod})$

$$\mathcal{R}(F) \ \simeq\ \Big[\, H^{-1}\big(\mathcal{R}(F)\big) \ \xrightarrow{d}\ H^0\big(\mathcal{R}(F)\big) \,\Big],$$

concentrated in cohomological degrees $[-1,0]$. Equivalently, $\mathcal{R}(F) \in D^{[-1,0]}(k\text{-mod})$.

Proof. Since $\mathcal{R}$ is t-exact by (B2), $\tau_{\le 0}\mathcal{R}(F) \xrightarrow{\sim} \mathcal{R}(F) \xleftarrow{\sim} \tau_{\ge -1}\mathcal{R}(F)$. Thus $H^i(\mathcal{R}(F)) = 0$ for $i \notin \{-1,0\}$. Choose any representative two-term complex $E_F = [E^{-1} \to E^0]$ for $\mathcal{R}(F)$ with E^{-1}, E^0 finite-dimensional. The canonical map from E_F to its cohomology model $[H^{-1}(\mathcal{R}(F)) \to H^0(\mathcal{R}(F))]$ is a quasi-isomorphism, natural in F. □

Remark .46. All statements below are invariant under filtered quasi-isomorphism on F and under isomorphism in $D^{\mathrm{b}}(k\text{-mod})$ on $\mathcal{R}(F)$.

C.2. The edge: $H^{-1}(\mathcal{R}(F)) \cong \varinjlim_t H_1(F^t)$ and naturality

Proposition .47 (Edge identification and naturality)**.** Under (B2), for every $F \in \mathsf{FiltCh}(k)$ there is a natural isomorphism

$$H^{-1}\big(\mathcal{R}(F)\big) \ \cong\ \varinjlim_{t\in\mathbb{R}} H_1(F^t C_\bullet).$$

If $f : F \to G$ is a filtration-preserving chain map, then the diagram

$$\begin{array}{ccc} H^{-1}\big(\mathcal{R}(F)\big) & \xrightarrow{\ \sim\ } & \varinjlim_t H_1(F^t C_\bullet) \\ {\scriptstyle H^{-1}(\mathcal{R}(f))}\big\downarrow & & \big\downarrow{\scriptstyle \varinjlim_t H_1(f^t)} \\ H^{-1}\big(\mathcal{R}(G)\big) & \xrightarrow{\ \sim\ } & \varinjlim_t H_1(G^t C_\bullet) \end{array}$$

commutes.

Proof. By Proposition .45, $\mathcal{R}(F)$ is represented by a two-term complex whose degree -1 cohomology is functorial in F. By the construction of $\mathcal{R}$ under (B2) (Appendix B), its degree -1 layer encodes stabilized degree-1 features; concretely, one may model $\mathcal{R}$ via a mapping-telescope of the filtered system $\{F^t\}_{t\in\mathbb{R}}$ truncated to degrees 1 and 0. Over the field k, filtered colimits in Vect_k are exact and computed objectwise in $[\mathbb{R}, \mathsf{Vect}_k]$ (Appendix A, Remark .16). Exactness allows passage of homology through filtered colimits in this bounded, finite-dimensional setting:

$$H^{-1}\big(\mathcal{R}(F)\big) \;\cong\; H_1\Big(\mathrm{Tel}_t\; F^t\Big) \;\cong\; \varinjlim_t H_1(F^t C_\bullet).$$

Naturality in f follows from the functoriality of the telescope construction and homology, and the universal property of colimits. □

C.3. Computing Ext^1 for amplitude $[-1, 0]$

Lemma .48 (Edge lemma for Ext^1)**.** Let $A \in D^{[-1,0]}(k\text{-mod})$. Then there is a natural isomorphism

$$\mathrm{Ext}^1(A, k) \;\cong\; \mathrm{Hom}\big(H^{-1}(A),\, k\big).$$

Proof. Apply $\mathrm{Hom}_{D^b}(-, k[1])$ to the truncation triangle $\tau_{\le -1}A \longrightarrow A \longrightarrow \tau_{\ge 0}A \longrightarrow (\tau_{\le -1}A)[1]$. Since $A \in D^{[-1,0]}$, $\tau_{\le -1}A \simeq H^{-1}(A)[1]$ and $\tau_{\ge 0}A \simeq H^0(A)[0]$. Using semisimplicity of k-vector spaces, $\mathrm{Ext}^1(H^0(A), k) = 0$. Hence

$$\mathrm{Hom}(A, k[1]) \;\cong\; \mathrm{Hom}\big(H^{-1}(A)[1],\, k[1]\big) \;\cong\; \mathrm{Hom}\big(H^{-1}(A), k\big),$$

naturally in A. □

Corollary .49 (Dimension and duality)**.** For any F with $\mathcal{R}(F) \in D^{[-1,0]}$,

$$\dim_k \mathrm{Ext}^1\big(\mathcal{R}(F), k\big) \;=\; \dim_k\, H^{-1}\big(\mathcal{R}(F)\big) \;=\; \dim_k \Big(\varinjlim_t H_1(F^t)\Big),$$

and the canonical pairing $\mathrm{Ext}^1(\mathcal{R}(F), k) \otimes H^{-1}(\mathcal{R}(F)) \longrightarrow k$ is perfect.

Proof. Combine Lemma .48 with Proposition .47. Finite dimensionality yields perfect duality. □

Remark .50 (Base change and choice of the test object)**.** For any field extension $k \subset K$, if $\mathcal{R}(F) \in D^{[-1,0]}$ then

$$\mathrm{Ext}^1\big(\mathcal{R}(F), k\big) = 0 \quad\Longleftrightarrow\quad \mathrm{Ext}^1\big(\mathcal{R}(F) \otimes^{\mathbf{L}}_k K,\, K\big) = 0,$$

since both are equivalent to $H^{-1}(\mathcal{R}(F)) = 0$ and H^{-1} commutes with scalar extension in this range. Moreover, for any finite-dimensional k-vector space V, $\mathrm{Ext}^1(\mathcal{R}(F), V) \cong \mathrm{Hom}(H^{-1}(\mathcal{R}(F)), V)$; hence testing against $k[0]$ is equivalent to testing against $V[0]$ for all V, but in B-Gate$^+$ we fix $V = k$ (minimal unit).

C.4. The bridge and its robust/windowed form

Theorem .51 (Bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$)**.** Let $F \in \mathsf{FiltCh}(k)$. If $\mathrm{PH}_1(F) = 0$ (equivalently, $H_1(F^t) = 0$ for all t), then

$$\mathrm{Ext}^1\big(\mathcal{R}(F),\, k\big) \;=\; 0.$$

Proof. If $\mathrm{PH}_1(F) = 0$ then $\varinjlim_t H_1(F^t) = 0$. By Proposition .47, $H^{-1}(\mathcal{R}(F)) = 0$. Apply Lemma .48 to obtain $\mathrm{Ext}^1(\mathcal{R}(F), k) = 0$. □

Corollary .52 (Robust/windowed bridge)**.** Fix $\tau > 0$. If $\mathrm{PH}_1\big(C_\tau(F)\big) = 0$, then $\mathrm{Ext}^1\big(\mathcal{R}(C_\tau(F)),\ k\big) = 0$.

Proof. Apply Theorem .51 to $C_\tau(F)$. By Appendix B, C_τ preserves the constructible range and $\mathcal{R}(C_\tau(F)) \in D^{[-1,0]}$, so (B2) applies. ◻

Remark .53 (No reverse claim)**.** We do *not* assert $\mathrm{Ext}^1(\mathcal{R}(F), k) = 0 \Rightarrow \mathrm{PH}_1(F) = 0$. See §12 for counterexamples.

C.5. Eligibility and B-Gate⁺ policy

Declaration .54 (Eligibility for using Ext^1 in B-Gate$^+$)**.** Fix a collapse threshold $\tau \geq 0$ and a window $[0, \tau]$. The Ext^1-test $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$ is *admitted into* B-Gate$^+$ if and only if:

1. $\mathcal{R}$ is t-exact and $\mathcal{R}(C_\tau F) \in D^{[-1,0]}(k\text{-mod})$ **(amplitude** ≤ 1**)**;
2. the scope rule of Appendix A (constructibility and return from functor-category colimits) holds; and
3. the test object is $Q = k[0]$ (no other targets are used in gating).

If any condition fails, the Ext^1-value may be logged for diagnostics but is *excluded* from the gate decision on that window.

Remark .55 (Diagnostics outside the eligible regime)**.** Outside amplitude ≤ 1, record a flag ext1_eligible:false and a reason (e.g. amplitude_gt_1). B-Gate$^+$ then relies on the persistence layer ($\mathrm{PH}_1 = 0$), tower isomorphisms $(\mu, u) = (0, 0)$, and the safety margin $\mathrm{gap}_\tau > \Sigma\delta$ only (cf. Chapters 10–12).

C.6. Naturality, exactness, and stability under admissible updates

Proposition .56 (Naturality of the bridge)**.** For any filtration-preserving $f : F \to G$ in $\mathsf{FiltCh}(k)$, the diagram

$$\begin{array}{ccc} \mathrm{Ext}^1\big(\mathcal{R}(F), k\big) & \xrightarrow{\sim} & \mathrm{Hom}\big(H^{-1}(\mathcal{R}(F)), k\big) \\ {\scriptstyle \mathrm{Ext}^1(\mathcal{R}(f),k)}\big\downarrow & & \big\downarrow{\scriptstyle \mathrm{Hom}(H^{-1}(\mathcal{R}(f)),k)} \\ \mathrm{Ext}^1\big(\mathcal{R}(G), k\big) & \xrightarrow{\sim} & \mathrm{Hom}\big(H^{-1}(\mathcal{R}(G)), k\big) \end{array}$$

commutes.

Proof. Combine naturality in Lemma .48 and Proposition .47. ◻

Lemma .57 (Exact triangles and 2-out-of-3)**.** Suppose $F_1 \to F_2 \to F_3 \to F_1[1]$ is a distinguished triangle after realization on a window $[0, \tau]$ with $\mathcal{R}(C_\tau F_i) \in D^{[-1,0]}$ for all i. If two of $\mathrm{Ext}^1(\mathcal{R}(C_\tau F_i), k)$ vanish, then so does the third.

Proof. Apply $\mathrm{Hom}(-, k[1])$ to the triangle in $D^{\mathrm{b}}(k\text{-mod})$ and use exactness of $\mathrm{Hom}(-, k[1])$ on distinguished triangles together with the amplitude restriction, which suppresses higher Ext-terms. ◻

Remark .58 (Stability under admissible updates)**.** When $F \mapsto F'$ is an update that is *non-expansive* at the persistence layer (e.g. deletion-type or ε-continuation post-collapse) and $\mathcal{R}(C_\tau F), \mathcal{R}(C_\tau F') \in D^{[-1,0]}$, then the verdict $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$ persists to F' provided the edge groups remain isomorphic (or vanish) along the filtered colimit; this is immediate from Proposition .47 and Lemma .48. Gating still defers to eligibility (Declaration .54).

C.7. Implementation details, reproducibility, and logging

Remark .59 (Run-time policy and manifest fields)**.** When B-Gate$^+$ is executed, enforce the order $F \to C_\tau F \to \mathcal{R}(C_\tau F) \to \mathrm{Ext}^1(-, k)$ (the **collapse-first** policy). The manifest run.yaml must include:

- ext1_eligible: boolean; true iff $\mathcal{R}(C_\tau F) \in D^{[-1,0]}$;
- ext1_used_in_gate: boolean; true iff ext1_eligible and the gate policy enables it;
- amplitude: reported as $[-1, 0]$ or >1 (if known by static design or a certified check);
- gate_order: fixed as collapse→realize→ext1;
- q_test: fixed as k[0].

Outside eligibility, set ext1_eligible:false, ext1_used_in_gate:false, and continue the gate using persistence/tower/safety criteria only.

Remark .60 (Computational recipe)**.** On an eligible window, computing $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k)$ reduces to:

1. compute $\varinjlim_t H_1((C_\tau F)^t)$ by stabilizing the barcode of degree 1 on $[0, \tau]$ (finite-type);
2. take the linear dual: $\mathrm{Ext}^1 \cong \mathrm{Hom}(\varinjlim_t H_1, k)$.

No derived resolutions are needed in the amplitude $[-1, 0]$ regime.

```
# run.yaml (excerpt)
ext1_eligible: true            # iff amplitude <= 1 was certified
ext1_used_in_gate: true        # iff eligible AND policy permits
amplitude: "[-1,0]"            # else ">1" or "unknown"
gate_order: "collapse→realize→ext1"
q_test: "k[0]"
notes: "PH1=0 implies Ext1=0 used; no reverse claim."
```

C.8. Counterexamples and boundary cases

Example .61 (Failure of the reverse implication $\mathrm{Ext}^1 \Rightarrow \mathrm{PH}_1$)**.** Let F be any filtered chain complex whose degree-1 barcode on $\mathbb{R}$ consists of a single finite bar (a, b) with $a < b < \infty$, and no other degree-1 features. Then $H_1(F^t) eq 0$ for $t \in (a, b)$ (so $\mathrm{PH}_1(F) eq 0$), but $\varinjlim_{t\to+\infty} H_1(F^t) = 0$. If $\mathcal{R}(F) \in D^{[-1,0]}$ (eligible), Corollary .49 gives $\mathrm{Ext}^1(\mathcal{R}(F), k) \cong \mathrm{Hom}(\varinjlim_t H_1(F^t), k) = 0$. Thus $\mathrm{Ext}^1 = 0$ while $\mathrm{PH}_1 eq 0$.

Example .62 (Amplitude breach: diagnostics only)**.** If $\mathcal{R}(C_\tau F) \in D^{[-2,0]}$ with $H^{-2} eq 0$, then Lemma .48 fails. It is possible that $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$ due to cancellations between H^{-2} and H^{-1}, even when $\varinjlim_t H_1((C_\tau F)^t) eq 0$. Hence the Ext^1-test is *ineligible* and must not be used for gating (Remark .55).

Example .63 (Non-constructible tails)**.** If $\mathbf{P}_1(F)$ is not constructible (e.g. infinitely many critical parameters accumulating in a bounded window), filtered colimits in $[\mathbb{R}, \mathsf{Vect}_k]$ may leave $\mathsf{Pers}_k^{\mathrm{ft}}$. The edge identification becomes non-operative for gating; eligibility fails by (B1).

C.9. Additional safeguards and best practices

Remark .64 (Monotonicity across windows)**.** If $0 < \tau \le \tau'$, eligibility at τ' implies eligibility at τ. Moreover, if $\mathrm{Ext}^1(\mathcal{R}(C_{\tau'}F), k) = 0$, then $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$ by functoriality of truncation and Corollary .49.

Remark .65 (Uniformity under base change)**.** By Remark .50, eligibility and the conclusion $\mathrm{Ext}^1 = 0$ are stable under extending scalars $k \subset K$. Thus the gate verdict is field-independent within the amplitude $[-1, 0]$ regime.

Remark .66 (Uniqueness of the test)**.** In the eligible regime, replacing $k[0]$ by any $V[0]$ with $V \in k$-mod yields an equivalent vanishing test. B-Gate$^+$ fixes the unit $k[0]$ to avoid ambiguity.

C.10. Scope and non-claims

Remark .67 (Scope and non-claims)**. The bridge** $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ **is used and proved only in** $D^{\mathrm{b}}(k\textbf{-mod})$**.** The converse implication $\mathrm{Ext}^1(\mathcal{R}(F), k) = 0 \Rightarrow \mathrm{PH}_1(F) = 0$ is *false* in general (Example .61). All filtered colimit uses obey the scope rule stated at the start of this appendix (Appendix A, Remark .16). Derived realizations into other targets (e.g. coherent/étale) appear only as **[Spec]** in the main text and *do not* extend the proved bridge.

Summary of Appendix C (reinforced and final). Under (B1)–(B3), the edge identification $H^{-1}(\mathcal{R}(F)) \cong \varinjlim_t H_1(F^t)$ and the amplitude $[-1, 0]$ model yield the bridge $\mathrm{PH}_1(F) = 0 \Rightarrow \mathrm{Ext}^1(\mathcal{R}(F), k) = 0$. The Ext^1-test is included in B-Gate$^+$ **only** when $\mathcal{R}(C_\tau F) \in D^{[-1,0]}$ (eligibility ≤ 1), and the operational order is **collapse** $\rightarrow$ **realize** $\rightarrow$ **Ext** (**collapse-first**). Outside eligibility, Ext^1 is excluded from gate decisions but may be logged. No reverse implication is claimed, and counterexamples/boundaries are documented in §12. All computations respect constructibility and filtered-colimit scope from Appendix A.

Appendix D. Towers, μ, u, and Examples [Proof/Example] (reinforced)

Throughout, fix a field k. We work in the constructible regime (Appendix A). For each degree $i \in \mathbb{Z}$, let

$$\mathbf{P}_i : \ \mathsf{FiltCh}(()k) \longrightarrow \mathsf{Pers}^{\mathrm{ft}}_k, F \longmapsto \big(t \mapsto H_i(F^t C_\bullet)\big)$$

send a bounded-in-degree filtered chain complex F to its constructible persistence module. The truncation $\mathsf{T}_\tau : \mathsf{Pers}^{\mathrm{ft}}_k \to \mathsf{Pers}^{\mathrm{ft}}_{k,\tau\text{-loc}}$ is exact and 1-Lipschitz (Appendix A). Filtered colimits are computed objectwise in $[\mathbb{R}, \mathsf{Vect}_k]$, with the scope rule of Appendix A, Remark .16. All statements at the filtered–complex layer are *up to filtered quasi-isomorphism (f.q.i.)*; persistence-layer statements take place in $\mathsf{Pers}^{\mathrm{ft}}_k$. All quantities below may depend on the threshold $\tau > 0$; no monotonicity in τ is asserted.

Comparison map and obstruction indices. Let $\{F_n\}_{n\in\mathbb{N}}$ be a directed system in $\mathsf{FiltCh}(()k)$. Let F_∞ be an apex equipped with a cocone $F_n \to F_\infty$ (indexing category $\mathbb{N} \cup \{\infty\}$ with unique morphisms $n \to \infty$). For each i and $\tau > 0$ set

$$\phi_{i,\tau} : \ \mathrm{colim}_{n\in\mathbb{N}}\, \mathsf{T}_\tau\big(\mathbf{P}_i(F_n)\big) \ \longrightarrow \ \mathsf{T}_\tau\big(\mathbf{P}_i(F_\infty)\big),$$

the canonical comparison in $[\mathbb{R}, \mathsf{Vect}_k]$. Define the *tower obstruction indices*

$$\mu_{i,\tau} := \mathrm{gdim}\,\mathrm{Ker}(\phi_{i,\tau}), u_{i,\tau} := \mathrm{gdim}\,\mathrm{coker}(\phi_{i,\tau}), \mu_{\mathrm{Collapse}} := \sum_i \mu_{i,\tau},\ u_{\mathrm{Collapse}} := \sum_i u_{i,\tau}.$$

These sums are finite because complexes are bounded in homological degrees (constructible range). The pair $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ detects *Type IV* (tower-level) failures at scale τ.[7]

[7] For compositions in finite-dimensional linear algebra one has surrogate subadditivity $\dim \mathrm{Ker}(g \circ f) \le \dim \mathrm{Ker} f + \dim \mathrm{Ker} g$ and $\dim \mathrm{coker}(g \circ f) \le \dim \mathrm{coker}\, g + \dim \mathrm{Ker} f$. Appendix J provides the persistence-layer analogue after applying T_τ and taking generic-fiber dimensions.

$$
\begin{tikzcd}
F_1 \arrow[r] \arrow[dr] & F_2 \arrow[r] \arrow[dr] & \cdots \arrow[r] & F_n \arrow[r] \arrow[dr] & \cdots \arrow[r] & F_\infty \\
& \mathbf{P}_i(F_1) \arrow[r] & \mathbf{P}_i(F_2) \arrow[r] & \cdots \arrow[r] & \mathbf{P}_i(F_n) \arrow[r] & \mathbf{P}_i(F_\infty)
\end{tikzcd}
$$

Figure 2: A tower with apex F_∞ and its image under $\mathbf{P}_i$. The comparison $\phi_{i,\tau}$ (defined after applying T_τ) measures the failure of the cocone to exhibit a colimit at scale τ.

Remark A.1 (Generic dimension after truncation)**.** In $\mathsf{Pers}_k^{\mathrm{ft}}$, after applying T_τ the kernel and cokernel of any morphism decompose (noncanonically) as finite direct sums of interval modules. We write $\mathrm{gdim}(-)$ for the *generic fiber* dimension, i.e. the multiplicity of the infinite bar $I[0,\infty)$ in that decomposition. Finite bars contribute zero generic fiber.

Remark A.2 (Invariance of $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$)**.** The indices $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$ are invariant under levelwise f.q.i. replacements of the tower and apex: if $F_n \simeq_{\text{f.q.i.}} F'_n$ and $F_\infty \simeq_{\text{f.q.i.}} F'_\infty$, then $\mathbf{P}_i$ sends these to isomorphisms in $\mathsf{Pers}_k^{\mathrm{ft}}$, hence Ker/coker (and thus μ, u) are unchanged. They are also invariant under cofinal reindexing of the tower, since filtered colimits over cofinal subdiagrams are canonically isomorphic.

D.1. Calculus of defects: generic-fiber interpretation, naturality, and subadditivity

We collect formal properties of the obstruction indices μ, u that are used implicitly throughout the paper and make them explicit for operational use.

Proposition A.3 (Generic-fiber interpretation)**.** Let $f : M \to N$ be a morphism in $\mathsf{Pers}_k^{\mathrm{ft}}$ and fix $\tau > 0$. Then, in the barcode decomposition of $\mathrm{Ker}(\mathsf{T}_\tau f)$ and $\mathrm{coker}(\mathsf{T}_\tau f)$, the multiplicity of $I[0,\infty)$ equals $\mathrm{gdim}\,\mathrm{Ker}(\mathsf{T}_\tau f)$ and $\mathrm{gdim}\,\mathrm{coker}(\mathsf{T}_\tau f)$, respectively. Equivalently, gdim is the generic fiber dimension of the corresponding functor $\mathbb{R} \to \mathsf{Vect}_k$.

Proof. This is standard for constructible 1D persistence over a field: after T_τ, objects are pointwise finite-dimensional and split as finite sums of interval modules; see Appendix A. The generic fiber (dimension on a cofinal ray) counts infinite bars, i.e. copies of $I[0,\infty)$. ◻

Definition A.4 (Morphisms of towers and naturality of ϕ)**.** A *morphism of towers with apex* $(F_\bullet, F_\infty) \to (G_\bullet, G_\infty)$ is a collection of maps $u_n : F_n \to G_n$ and $u_\infty : F_\infty \to G_\infty$ commuting with all structure maps to the apex. Applying $\mathbf{P}_i$, then T_τ, and passing to the filtered colimit yields a commutative square

$$
\begin{tikzcd}
\mathrm{colim}_n\, \mathsf{T}_\tau \mathbf{P}_i(F_n) \arrow[r, "\phi^F_{i,\tau}"] \arrow[d, "\mathrm{colim}\ \mathsf{T}_\tau \mathbf{P}_i(u_n)"'] & \mathsf{T}_\tau \mathbf{P}_i(F_\infty) \arrow[d, "\mathsf{T}_\tau \mathbf{P}_i(u_\infty)"] \\
\mathrm{colim}_n\, \mathsf{T}_\tau \mathbf{P}_i(G_n) \arrow[r, "\phi^G_{i,\tau}"] & \mathsf{T}_\tau \mathbf{P}_i(G_\infty)
\end{tikzcd}
$$

i.e. $\phi_{i,\tau}$ is natural in the tower.

Proposition A.5 (Functoriality and subadditivity under composition)**.** Let $(F_\bullet, F_\infty) \xrightarrow{u} (G_\bullet, G_\infty) \xrightarrow{v} (H_\bullet, H_\infty)$ be morphisms of towers with apex. Then, for each i, τ,

$$
\mu(\phi^H_{i,\tau} \circ \mathrm{colim}\ \mathsf{T}_\tau \mathbf{P}_i(v_n \circ u_n)) \le n\mu(\phi^G_{i,\tau} \circ \mathrm{colim}\ \mathsf{T}_\tau \mathbf{P}_i(u_n)) + \mu(\phi^H_{i,\tau} \circ \mathrm{colim}\ \mathsf{T}_\tau \mathbf{P}_i(v_n)),
$$

$$
u(\phi^H_{i,\tau} \circ \mathrm{colim}\ \mathsf{T}_\tau \mathbf{P}_i(v_n \circ u_n)) \le nu(\phi^H_{i,\tau} \circ \mathrm{colim}\ \mathsf{T}_\tau \mathbf{P}_i(v_n)) + \mu(\phi^G_{i,\tau} \circ \mathrm{colim}\ \mathsf{T}_\tau \mathbf{P}_i(u_n)).
$$

In particular, writing simply $\phi = \phi_{i,\tau}$ for a fixed tower, any factorization of ϕ yields $\mu(\phi \circ \psi) \le \mu(\psi) + \mu(\phi)$ and $u(\phi \circ \psi) \le u(\phi) + \mu(\psi)$.

Proof. Apply T_τ and use exactness together with the standard inequalities for kernels and cokernels of compositions in finite-dimensional linear algebra, then pass to generic fibers via Proposition A.3. □

Proposition A.6 (Additivity on finite direct sums)**.** For two towers $(F_\bullet, F_\infty)$ and $(G_\bullet, G_\infty)$,

$$\mu\big((F \oplus G)_\bullet, (F \oplus G)_\infty\big) = \mu(F_\bullet, F_\infty) + \mu(G_\bullet, G_\infty),$$

$$u\big((F \oplus G)_\bullet, (F \oplus G)_\infty\big) = u(F_\bullet, F_\infty) + u(G_\bullet, G_\infty).$$

Proof. $\mathbf{P}_i$ and T_τ preserve finite direct sums; kernels and cokernels preserve finite direct sums; gdim is additive on direct sums. □

Proposition A.7 (Invariance under f.q.i. and cofinal reindexing)**.** If $F_n \simeq_{\text{f.q.i.}} F'_n$ levelwise and $F_\infty \simeq_{\text{f.q.i.}} F'_\infty$, then μ, u agree for the two towers. If $J \subset \mathbb{N}$ is cofinal, then restricting the tower to J does not change μ, u.

Proof. Follows from $\mathbf{P}_i$ sending f.q.i. to isomorphisms in $\mathsf{Pers}_k^{\text{ft}}$, exactness of T_τ, and the fact that filtered colimits over cofinal subdiagrams are canonically isomorphic. □

D.2. Toy towers: pure kernel / pure cokernel / mixed

Example A.8 (Pure cokernel at a fixed scale)**.** Fix $\tau > 0$ and degree $i = 1$. Let $\mathbf{P}_1(F_n) = I[0, \tau - \frac{1}{n})$ with transition maps the evident inclusions. Let F_∞ satisfy $\mathbf{P}_1(F_\infty) = I[0, \infty)$. Then $\mathsf{T}_\tau(\mathbf{P}_1(F_n)) = 0$ for all n, whereas $\mathsf{T}_\tau(\mathbf{P}_1(F_\infty)) \cong I[0, \infty)$. Hence $\phi_{1,\tau} : 0 \to I[0, \infty)$ has trivial kernel and nontrivial cokernel, so $\mu_{1,\tau} = 0$ and $u_{1,\tau} = 1$ (pure cokernel).

Example A.9 (Pure kernel at a fixed scale)**.** Fix $\tau > 0$. Let $\mathbf{P}_1(F_n) = I[0, \infty)$ for all n, with transition maps the identities (a stationary directed system). Let F_∞ satisfy $\mathbf{P}_1(F_\infty) = 0$, and take the cocone $\mathbf{P}_1(F_n) \to \mathbf{P}_1(F_\infty)$ to be 0 for all n. Then $\mathsf{T}_\tau(\mathbf{P}_1(F_n)) \cong I[0, \infty)$ for all n, so the source of $\phi_{1,\tau}$ is $I[0, \infty)$, while the target is 0. Thus $\phi_{1,\tau} : I[0, \infty) \to 0$ has nontrivial kernel and zero cokernel, hence $\mu_{1,\tau} = 1$, $u_{1,\tau} = 0$ (pure kernel).

Example A.10 (Mixed)**.** Fix $\tau > 0$ and set

$$\mathbf{P}_1(F_n) \;=\; I[0, \tau - \tfrac{1}{n}) \;\oplus\; I[0, \infty),$$

with transition maps the obvious inclusions on the first summand and the identities on the second. Take F_∞ with $\mathbf{P}_1(F_\infty) = I[0, \infty) \oplus 0$, using cocone maps that send the second summand to 0. Then the first summand yields a cokernel contribution exactly as in Example A.8, while the second yields a kernel contribution exactly as in Example A.9. Hence $\mu_{1,\tau} = u_{1,\tau} = 1$ (mixed).

All three persistence-level towers are realizable by filtered complexes via the interval-realization assignment $\mathcal{U}$ (Appendix B), up to f.q.i.; constructibility is preserved.

D.3. When $\phi_{i,\tau}$ is an isomorphism: $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0, 0)$

Proposition A.11 (Isomorphism criterion)**.** Assume:

(i) degreewise filtered colimits in $\mathsf{FiltCh}(()k)$ are computed objectwise on chains and filtrations;

(ii) each $\mathbf{P}_i(F_n)$ lies in $\mathsf{Pers}_k^{\text{ft}}$;

(iii) T_τ commutes with the filtered colimit of $\{\mathbf{P}_i(F_n)\}$ in $[\mathbb{R}, \mathsf{Vect}_k]$, and the result is constructible (Appendix A, Theorem .14);

(iv) the cocone exhibits a colimit at persistence level: the canonical map $\operatorname{colim}_n \mathbf{P}_i(F_n) \xrightarrow{\cong} \mathbf{P}_i(F_\infty)$ is an isomorphism in $[\mathbb{R}, \mathsf{Vect}_k]$.

Then for every i and $\tau > 0$, the comparison map $\phi_{i,\tau}$ is an isomorphism in $\mathsf{Pers}_k^{\mathrm{ft}}$. Consequently $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$.

Remark A.12. Condition (iv) is automatic if F_∞ *is* the colimit of $\{F_n\}$ in a model of filtered complexes for which $\mathbf{P}_i$ is computed objectwise and the scope rule of Appendix A applies; no claim is made beyond that regime.

D.4. Sufficient conditions ensuring $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$

The summability condition $\sum_n d_{\mathrm{int}}\big(\mathbf{P}_i(F_{n+1}), \mathbf{P}_i(F_n)\big) < \infty$ *alone* does not guarantee $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$; see §D.5.1. The following hypotheses are sufficient.

Theorem A.13. Fix i and $\tau > 0$. Each of the following implies that $\phi_{i,\tau}$ is an isomorphism (hence $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$):

(S1) **Commutation and apex colimit:** T_τ commutes with the filtered colimit of $\{\mathbf{P}_i(F_n)\}$ in $[\mathbb{R}, \mathsf{Vect}_k]$, the outcome is constructible, and the cocone exhibits a colimit at persistence level (i.e. Proposition A.11(iv) holds).

(S2) **No τ-accumulation from below:** there exists $\eta > 0$ such that, for all sufficiently large n, no bar in $\mathbf{P}_i(F_n)$ has length in the half-open interval $(\tau - \eta, \tau)$. Equivalently, there is no sequence of bar lengths strictly increasing to τ.

(S3) **T_τ-Cauchy with compatible cocone:** the sequence $\mathsf{T}_\tau(\mathbf{P}_i(F_n))$ is Cauchy in the interleaving metric, and the cocone to $\mathsf{T}_\tau(\mathbf{P}_i(F_\infty))$ identifies the metric limit with the colimit target. (Here we use only the standard completeness/uniqueness of limits for p.f.d. barcodes under the bottleneck/interleaving metric.)

Proof. (S1) is Proposition A.11. For (S2), the gap prevents creation at the apex of new bars of length $> \tau$: every long bar in $\mathsf{T}_\tau(\mathbf{P}_i(F_\infty))$ must appear at some finite stage and stabilize, yielding bijectivity on interval factors. For (S3), completeness of the space of p.f.d. persistence modules up to isometry implies a unique metric limit; the stated compatibility identifies it with the colimit target, so $\phi_{i,\tau}$ is an isometry and hence an isomorphism in $\mathsf{Pers}_k^{\mathrm{ft}}$. $\square$

D.5. A counterexample: $\sum d_{\mathrm{int}} < \infty$ yet $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) eq (0, 0)$

Example A.14 (Summable increments, pure cokernel at the apex)**.** Fix $\tau > 0$ and set $\ell_n = \tau - \sum_{m \geq n} 2^{-m} \uparrow \tau$, so that $\sum_n (\ell_{n+1} - \ell_n) = \sum_n 2^{-n} < \infty$. Let $M_n := I[0, \ell_n)$ with $M_n \hookrightarrow M_{n+1}$ the standard inclusions. Then

$$d_{\mathrm{int}}(M_n, M_{n+1}) = \tfrac{1}{2}(\ell_{n+1} - \ell_n) = 2^{-(n+1)}, \qquad \sum_n d_{\mathrm{int}}(M_n, M_{n+1}) < \infty.$$

Let $\mathbf{P}_1(F_n) = M_n$, and choose an apex with $\mathbf{P}_1(F_\infty) = I[0, \infty)$. For every n, $\mathsf{T}_\tau(M_n) = 0$, while $\mathsf{T}_\tau(\mathbf{P}_1(F_\infty)) = I[0, \infty)$. Thus $\mu_{\mathrm{Collapse}} = 0$, $u_{\mathrm{Collapse}} = 1$ (pure cokernel), despite the summable interleaving distances along the tower.

This shows that $\sum d_{\mathrm{int}} < \infty$ alone is insufficient to force $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$.

D.6. Converse failures and the Type IV catalog

D.6.1. $\mathrm{Ext}^1 = 0$ **does not imply** $\mathrm{PH}_1 = 0$**.** Let $A \in D^{[-1,0]}(k\text{-mod})$ with $H^{-1}(A) = 0$ and $H^0(A)eq0$, e.g. the stalk complex $V[0]$ for a nonzero k-space V. Then $\mathrm{Ext}^1(A, k) \cong \mathrm{Hom}(H^{-1}(A), k) = 0$ by Appendix C (Lemma .48). Choose $F \in \mathsf{FiltCh}(()k)$ with $\mathbf{P}_1(F)eq0$ (e.g. a single finite interval) and $\mathcal{R}(F) \simeq A$; this can be arranged up to f.q.i. using the realization assignment $\mathcal{U}$ (Appendix B). Hence $\mathrm{Ext}^1(\mathcal{R}(F), k) = 0$ while $\mathrm{PH}_1(F)eq0$, refuting the converse of the bridge.

D.6.2. Type IV (pure cokernel) at fixed τ. Example A.8 exhibits $\mu_{\mathrm{Collapse}} = 0$, $u_{\mathrm{Collapse}} > 0$ with $\mathsf{T}_\tau(\mathbf{P}_1(F_n)) = 0$ for all n but $\mathsf{T}_\tau(\mathbf{P}_1(F_\infty))eq0$. Thus finite layers appear admissible while the apex fails.

D.6.3. Type IV (mixed). Example A.10 yields $\mu_{\mathrm{Collapse}} > 0$ and $u_{\mathrm{Collapse}} > 0$ simultaneously, demonstrating that both kernel and cokernel defects can occur in the same tower.

D.6.4. Realization notes. All persistence-level constructions above are realizable by filtered complexes via $\mathcal{U}$ (Appendix B), up to f.q.i.; constructibility is preserved.

D.7. Restart/Summability for window pasting

All persistence-layer statements below are made *after* applying T_τ.

Definition A.15 (Per-window safety margin and pipeline budget)**.** Let $\{W_k = [u_k, u_{k+1})\}_{k\in K}$ be a MECE partition (Appendix A, Definition .22). On each window W_k, fix a collapse threshold $\tau_k > 0$. For a degree i, define the *pipeline budget*

$$\Sigma\delta_k(i) \;:=\; \sum_{U\in W_k} \Big(\delta_U^{\mathrm{alg}}(i,\tau_k) + \delta_U^{\mathrm{disc}}(i,\tau_k) + \delta_U^{\mathrm{meas}}(i,\tau_k)\Big),$$

and the *safety margin* $\mathrm{gap}_{\tau_k}(i) > 0$ as the configured slack for B-Gate⁺ on W_k and degree i.

Lemma A.16 (Restart inequality)**.** Assume that, on window W_k, B-Gate⁺ passes with $\mathrm{gap}_{\tau_k}(i) > \Sigma\delta_k(i)$, and that the transition to W_{k+1} is realized by a finite composition of *deletion-type* steps and ε-continuations (both measured after T_τ). Then there exists $\kappa \in (0, 1]$, depending only on the admissible step class and the τ-adaptation policy, such that

$$\mathrm{gap}_{\tau_{k+1}}(i) \;\geq\; \kappa\big(\mathrm{gap}_{\tau_k}(i) - \Sigma\delta_k(i)\big).$$

Proof sketch. Deletion-type steps are non-increasing for the monitored indicators after T_τ (Appendix E), and ε-continuations are 1-Lipschitz. Aggregating drifts yields the stated retention factor κ. □

Definition A.17 (Summability)**.** A run satisfies *Summability* (on a degree set $I \subset \mathbb{Z}$) if

$$\sum_{k\in K} \Sigma\delta_k(i) \;<\; \infty \qquad (\forall\, i \in I).$$

A sufficient design is a geometric decay of τ_k (hence of spectral/temporal bins) and bounded per-window step counts.

Theorem A.18 (Pasting windowed certificates)**.** Let $\{W_k\}_k$ be MECE, and on each W_k let B-Gate⁺ pass with $\mathrm{gap}_{\tau_k}(i) > \Sigma\delta_k(i)$ for all $i \in I$. If the Restart inequality (Lemma A.16) holds at every transition and Summability (Definition A.17) holds, then the concatenation of windowed certificates yields a global certificate on $\bigcup_k W_k$ for the degrees $i \in I$.

Proof sketch. Iterate Lemma A.16 and sum budgets; Summability ensures that the cumulative loss of the safety margin remains bounded, so positivity of the margin persists. MECE coverage (Appendix A) ensures there are no gaps/overlaps. □

D.8. Stability bands, τ-sweeps, and detection algorithm

Definition A.19 (Stability band via τ-sweep)**.** Fix a window W and degree i. Let $\{\tau_\ell\}_{\ell=1}^{L}$ be an increasing τ-sweep. A contiguous block $\{\tau_a, \dots, \tau_b\}$ is a *stability band* if

$$\mu_{i,\tau_\ell} = u_{i,\tau_\ell} = 0 \quad \text{for all } \ \ell \in \{a, \dots, b\},$$

and the verdict persists upon *refining* the sweep (inserting new τ-values) without introducing μ or u in the band.

Proposition A.20 (Robust detection of stability bands)**.** Assume (S1)–(S3) of Theorem A.13 hold on W. Then any sufficiently fine τ-sweep admits stability bands covering all τ at which $\phi_{i,\tau}$ is an isomorphism; conversely, detecting a stability band by a sweep and its refinement certifies $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0, 0)$ on the band.

Proof sketch. Under (S1)–(S3), $\phi_{i,\tau}$ is an isomorphism on open neighborhoods of the corresponding τ's. A fine sweep samples each neighborhood; refinement eliminates aliasing. The converse follows by definition. □

Remark A.21 (Caveat: non-monotonicity in τ)**.** There is no general monotonicity of $\mu_{i,\tau}$ or $u_{i,\tau}$ in τ. Stability bands may be separated by isolated τ-values where $\phi_{i,\tau}$ fails to be an isomorphism.

Detection algorithm (auditable). Given a sweep $\mathsf{TauSweep} = \{\tau_\ell\}_{\ell=1}^{L}$:

(A1) For each ℓ, compute ϕ_{i,τ_ℓ} and record $(\mu_{i,\tau_\ell}, u_{i,\tau_\ell})$.

(A2) Extract maximal contiguous indices $[a, b]$ with $(\mu, u) = (0, 0)$.

(A3) Refine by inserting midpoints $\tau' = \frac{1}{2}(\tau_\ell + \tau_{\ell+1})$ within each candidate band and recompute (μ, u).

(A4) Accept a band if all refined points also yield $(0, 0)$.

(A5) Emit certificate with hashes of inputs, tower metadata, and the flags indicating which of (S1)–(S3) were used.

A minimal YAML schema for audit is provided in §D.10.

D.9. Implementation guide: APIs, stubs, and tests

This section provides concrete, auditable artifacts for engineering teams. All persistence-layer computations are understood after applying T_τ.

Lean stubs (illustrative).

```
namespace PH
structure Tower ( : Type) :=
  (F    :  →  FiltCh)
  (apex : FiltCh)
  (toApex :   i :   ,  ChainMap (F i) apex)

def P_i (i :  ) : FiltCh → Pers := -- assumed given
def T_tau ( : ) : Pers → Pers := -- exact, 1-Lipschitz
```

```
def phi (i :  ) ( :  ) (T : Tower  ) : PersHom :=
  have src := colim (fun n => T_tau  (P_i i (T.F n)))
  have tgt := T_tau  (P_i i T.apex)
  comparison src tgt -- canonical

def gdim (M : Pers) : Nat := -- multiplicity of I∞[0,) in barcode

def mu (i :  ) ( :  ) (T : Tower  ) : Nat := gdim (kernel (phi i   T))
def nu (i :  ) ( :  ) (T : Tower  ) : Nat := gdim (cokernel (phi i   T))
end PH
```

Listing 1: Lean 4 stubs for towers and obstruction indices

Coq stubs (illustrative).

```
Module PH.
Record Tower := {
  F       : nat -> FiltCh;
  apex    : FiltCh;
  toApex : forall n, ChainMap (F n) apex
}.

Parameter P_i  : Z -> FiltCh -> Pers.
Parameter T_tau : R -> Pers -> Pers. (* exact, 1-Lipschitz *)

Definition phi (i:Z) (tau:R) (T:Tower) : PersHom :=
  let src := colim (fun n => T_tau tau (P_i i (F T n))) in
  let tgt := T_tau tau (P_i i (apex T)) in
  comparison src tgt.

Parameter gdim : Pers -> nat.

Definition mu (i:Z) (tau:R) (T:Tower) : nat := gdim (kernel (phi i tau T)).
Definition nu (i:Z) (tau:R) (T:Tower) : nat := gdim (cokernel (phi i tau T)).
End PH.
```

Listing 2: Coq stubs for towers and obstruction indices

Sample tests.

(T1) **T3 (Filtered-colim stability).** Construct a tower whose apex is the filtered colimit and for which T_τ commutes with colim. Verify $\phi_{i,\tau}$ is an iso and $(\mu, u) = (0, 0)$.

(T2) **T7 (Toy towers).** Instantiate the pure kernel, pure cokernel, and mixed towers of §D.2 and confirm $(\mu, u) = (1, 0), (0, 1), (1, 1)$, respectively.

(T3) **T9 (No τ-accumulation).** Create barcodes with an η-gap below τ; confirm $\phi_{i,\tau}$ is iso.

(T4) **T10 (Cauchy+compatibility).** Build a Cauchy sequence in the bottleneck metric whose limit equals the apex post-T_τ; confirm iso.

(T5) **T11 (Restart+Summability).** Simulate windows and transitions satisfying Lemma A.16 and Definition A.17; verify global certificate via Theorem A.18.

D.10. Audit schema: run.yaml and HDF5 layout

YAML fields (mandatory).


```yaml
phi:
  idx:
    - i: 1
      tau: 0.75
      iso: true          # _{i,} isomorphism?
      mu: 0
      nu: 0
      flags:
        S1: true         # used commutation+apex-colim
        S2: false
        S3: false
      iso_tail:
        passed: true     # tail check on refined sweep
        refinement_levels: 2
windows:
  collapse:
    tau_sweep: [0.5, 0.6, 0.7, 0.75, 0.8]
persistence:
  phi_iso_tail: "strict" # policy for refinement acceptance
tower:
  edges:
    - src: 0; dst: 1; kind: "inclusion"
    - src: 1; dst: 2; kind: "inclusion"
hash:
  inputs: "sha256:..."
  code:   "sha256:..."
```


Listing 3: Minimal audit fields in run.yaml

HDF5 groups (canonical order). We store comparison data under /phi and tower metadata under /tower. A minimal layout is:

Group/Dataset	Contents
/phi/idx/i	integer degrees i
/phi/idx/tau	real thresholds τ
/phi/idx/iso	boolean flags
/phi/idx/mu	nonnegative integers $\mu_{i,\tau}$
/phi/idx/nu	nonnegative integers $u_{i,\tau}$
/phi/idx/flags	bitmask for (S1,S2,S3)
/phi/idx/iso_tail/passed	boolean
/tower/edges/src	integer source indices
/tower/edges/dst	integer target indices
/tower/edges/kind	categorical: inclusion/deletion/epsilon

D.11. Additional formalities: tau-naturality and bandwise certification

We record two further properties that are often used implicitly.

Proposition A.22 (Piecewise constancy off critical thresholds)**.** Fix i and a tower. There exists a finite set $S \subset (0, \infty)$ consisting of bar lengths and their finite sums and differences such that $\mu_{i,\tau}$ and $u_{i,\tau}$ are locally constant on each connected component of $(0, \infty) \setminus S$.

Proof sketch. Within the constructible regime, changes in Ker/coker after T_τ can occur only when τ crosses endpoints of bars that interact with colim/cocone structure; this yields a finite critical set after fixing a finite window of degrees and using boundedness. □

Corollary A.23 (Bandwise certification)**.** If ϕ_{i,τ_0} is an isomorphism at some τ_0 lying in a component of $(0, \infty) \setminus S$, then it is an isomorphism on the whole component.

D.12. Completion note and cross-module conventions

Remark A.24 (No further supplementation required)**.** This appendix now provides: (i) the definition and calculus of the tower obstruction indices (μ, u) (generic fiber dimensions after truncation), (ii) naturality and functoriality of the comparison map ϕ, including subadditivity under composition and additivity under direct sums, (iii) illustrative toy towers (pure kernel/cokernel/mixed) and a counterexample showing that $\sum d_{\mathrm{int}} < \infty$ does not force $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$, (iv) sufficient conditions (S1)–(S3) guaranteeing $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$, (v) a *Restart/Summability* framework to paste windowed certificates into global ones, (vi) a robust τ-sweep procedure and *stability bands* to certify $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}) = (0, 0)$ on contiguous τ-ranges, and (vii) implementation-grade audit schemas (YAML/HDF5), API stubs (Lean/Coq), and test items. All statements are confined to the v16.0 guard-rails (constructible 1D persistence over a field; persistence-layer equalities after truncation; f.q.i. on filtered complexes), and no further supplementation is required for operational use in the proof framework.

Cross-module conventions. Ext-tests are always taken against $k[0]$ (Appendix C): $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$. When windowed energy summaries are referenced elsewhere, the exponent is uniformly $\alpha > 0$ (default $\alpha = 1$). Update monotonicity follows the global rule: *deletion-type* updates are non-increasing for windowed energies and spectral tails after truncation, whereas *inclusion-type* updates are stable (non-expansive); see Appendix E. Type labels follow the global convention *Type I–II / Type III / Type IV*.

Appendix E. Spectral Indicators: Monotonicity, Stability, Counterexamples [Proof/Spec]

For $\tau > 0$, define the *clipped spectrum* $\mathrm{clip}_\tau(H) := (\min\{\lambda_j(H), \tau\})_{j=1}^n$, the *clipped sum*

$$S^{\leq\tau}(H) := \sum_{j=1}^{n} \min\{\lambda_j(H), \tau\},$$

and the *sub-threshold deficit*

$$D^{<\tau}(H) \;:=\; \sum_{j=1}^{n} (\tau - \lambda_j(H))_+, \qquad x_+ := \max\{x, 0\}.$$

We use the operator norm $\|\cdot\|_{\mathrm{op}}$ and Frobenius norm $\|\cdot\|_{\mathrm{fro}}$. All references to filtered colimits follow the scope rule in Appendix A, Remark (A.6). Cross-module conventions (used globally): Ext-tests are taken against $k[0]$ (Appendix C), i.e. $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$; energy exponents are uniform $\alpha > 0$ (default $\alpha = 1$);

type labels use *Type I–II / Type III / Type IV*. When we refer to persistence/filtered complexes, equalities are at the persistence layer (strict in $\mathsf{Pers}_k^{\mathrm{ft}}$); filtered–complex claims are *up to f.q.i.*

Deletion vs. inclusion. When H arises by restricting admissible degrees of freedom, imposing Dirichlet constraints, eliminating internal dofs by shorting (Schur complement), or taking principal submatrices, we call conclusions *deletion-type*. When H is obtained by adding degrees of freedom, couplings, or enlarging a domain, we call them *inclusion-type*. Deletion-type updates admit one-sided monotonicity; inclusion-type updates admit only stability (non-expansive), unless additional order hypotheses are imposed.

E.0. Scope and window policy

All spectral audits in this appendix are *windowed* and performed *after* collapse on the B-side single layer. Concretely, comparisons follow the mandatory order:

$$\boxed{\text{for each } t \implies \text{apply } \mathbf{P}_i \implies \text{apply } \mathbf{T}_\tau \implies \text{compare in } \mathsf{Pers}_k^{\mathrm{ft}}}.$$

When a spectral window $[a, b]$ with bin width $\beta > 0$ is used, bins are *half-open, right-attribution* $I_r = [a + r\beta, a + (r + 1)\beta)$, and eigenvalues at a right boundary are counted in the next bin. Underflows/overflows are recorded. This policy ensures reproducibility and compatibility with Overlap Gate and B–Gate$^+$ (Appendix G; Chapter 1).

E.1. Deletion-type monotonicity (principal/Dirichlet, Schur complement, Loewner)

Proposition B.1 (Principal/Dirichlet restriction: interlacing and counting)**.** Let $A \in \mathbb{R}^{n\times n}$ be Hermitian and B a principal $(n-1)\times(n-1)$ submatrix (obtained, e.g., by pinning a coordinate—Dirichlet restriction). Then Cauchy interlacing holds:

$$\lambda_1(A) \le \lambda_1(B) \le \lambda_2(A) \le \cdots \le \lambda_{n-1}(B) \le \lambda_n(A).$$

In particular, for every $\theta \in \mathbb{R}$, $N_\theta(B) \le N_\theta(A)$.

Proposition B.2 (Schur complement (shorting) monotonicity)**.** Partition $M = \left(\begin{smallmatrix} A & B \\ B^\top & C \end{smallmatrix}\right) \succeq 0$ with $C \succ 0$ and form the Schur complement $S := A - BC^{-1}B^\top$. Then $S \preceq A$. Consequently, for all j and all $\theta \ge 0$, $\lambda_j(S) \le \lambda_j(A), \quad N_\theta(S) \le N_\theta(A)$.

Proposition B.3 (Loewner-order monotonicity)**.** If $0 \preceq A \preceq B$ (Loewner order), then for each j, $\lambda_j(A) \le \lambda_j(B)$ and, for every $\theta \ge 0$, $N_\theta(A) \le N_\theta(B)$.

Remark B.4 (Heat traces and spectral tails)**.** For PSD matrices and $t > 0$, the heat trace $\mathrm{HT}(t; H) = \sum_j e^{-t\lambda_j(H)}$ satisfies: if $A' \preceq A$ (contraction), then $\mathrm{HT}(t; A') \ge \mathrm{HT}(t; A)$; if $A' \succeq A$ (hardening), then $\mathrm{HT}(t; A') \le \mathrm{HT}(t; A)$. Likewise, for spectral tails $\mathrm{ST}_\beta(H) = \sum_{j\ge1} \lambda_j(H)^{-\beta}$ with $\beta > 0$, one has $\mathrm{ST}_\beta(A') \ge \mathrm{ST}_\beta(A)$ under $A' \preceq A$ and the reverse inequality under $A' \succeq A$, provided all $\lambda_j > 0$. In practice, tails are computed on $L(C_\tau F)$ with zero modes removed or handled by pseudoinverses; see Appendix G.

Corollary B.5 (Conservative averaging)**.** If $A_1, \ldots, A_m \succeq 0$ satisfy $A_\ell \preceq A$ for all ℓ, then for any convex combination $\bar{A} := \sum_\ell w_\ell A_\ell$ with $w_\ell \ge 0$, $\sum_\ell w_\ell = 1$, $\bar{A} \preceq A$. Therefore $\lambda_j(\bar{A}) \le \lambda_j(A)$, $N_\theta(\bar{A}) \le N_\theta(A)$ for $\theta \ge 0$, and the heat trace/tail inequalities above apply.

Remark B.6 (Orientation for deletions)**.** Two Loewner orientations occur in practice. *Contractions* (e.g. Schur complements, Kron reduction) produce $A' \preceq A$; *hardening* operations (e.g. some PDE Dirichlet comparisons across different media) may yield $A' \succeq A$. We state deletion-type monotonicities in both orientations.

E.2. Inclusion-type counterexamples

Deletion-type monotonicity does *not* extend naively to inclusion-type operations without additional order hypotheses.

Example B.7 (Neumann/domain inclusion reverses direction)**.** For the Neumann Laplacian on an interval, enlarging the domain decreases the nonzero eigenvalues: on $[0, L]$, the first nonzero Neumann eigenvalue is $(\pi/L)^2$, so passing $L : 1 \to 2$ reduces it from π^2 to $(\pi/2)^2$. Thus any "inclusion $\Rightarrow$ increase" heuristic fails under Neumann-type constraints.

Example B.8 (Indefinite coupling can move eigenvalues both ways)**.** Let $A = I_2 = \mathrm{diag}(1, 1)$ and $B = \left(\begin{smallmatrix} 1 & M \\ M & 1 \end{smallmatrix}\right)$ with $M > 1$. Then B has eigenvalues $1 - M$ and $1 + M$, so for $\theta = 0$, $N_\theta(B) = 1 < N_\theta(A) = 2$, while the top eigenvalue λ_2 increases. Without a Loewner relation ($B - A$ indefinite), no monotone law survives.

Example B.9 (Principal extension lacks a fixed direction)**.** Let $B = [0]$ (eigenvalue 0) and $A = \left(\begin{smallmatrix} 0 & t \\ t & 0 \end{smallmatrix}\right)$ with $teq0$. Going from B to A (adding one dof and a coupling) produces eigenvalues $-|t|$ and $|t|$: the maximum increases to $|t|$, but the minimum decreases to $-|t|$. Hence no uniform increase/decrease holds under inclusion.

These examples justify restricting monotone claims to the deletion/Loewner settings formalized in §E.1.

E.3. Continuity, stability, and truncated functionals

Write $N_{\theta\pm0}(A)$ for the left/right limits at θ (no jump unless θ is an eigenvalue).

Proposition B.10 (Weyl and Hoffman–Wielandt)**.** For Hermitian $A, B \in \mathbb{R}^{n\times n}$,

$$\max_{1\le j\le n} |\lambda_j(A) - \lambda_j(B)| \;\le\; \|A - B\|_{\mathrm{op}}, \Big(\sum_{j=1}^{n} |\lambda_j(A) - \lambda_j(B)|^2\Big)^{1/2} \;\le\; \|A - B\|_{\mathrm{fro}}.$$

Hence $A \mapsto (\lambda_1(A), \dots, \lambda_n(A))$ is 1-Lipschitz from $(\|\cdot\|_{\mathrm{op}})$ into $(\mathbb{R}^n, \|\cdot\|_\infty)$.

Corollary B.11 (Lipschitz stability of clipped spectra)**.** For any $\tau > 0$ and Hermitian A, B,

$$\sum_{j=1}^{n} \Big|\min\{\lambda_j(A), \tau\} - \min\{\lambda_j(B), \tau\}\Big| \;\le\; \sum_{j=1}^{n} |\lambda_j(A) - \lambda_j(B)| \;\le\; \sqrt{n}\,\|A - B\|_{\mathrm{fro}} \;\le\; n\,\|A - B\|_{\mathrm{op}}.$$

Consequently, $S^{\le\tau}$ is $\sqrt{n}$–Lipschitz in $\|\cdot\|_{\mathrm{fro}}$ and n–Lipschitz in $\|\cdot\|_{\mathrm{op}}$.

Proposition B.12 (Semicontinuity of counting indicators)**.** If $A_m \to A$ in operator norm and θ is not an eigenvalue of A, then $N_\theta(A_m) = N_\theta(A)$ for all large m (local constancy). In general,

$$\limsup_{m\to\infty} N_\theta(A_m) \;\le\; N_{\theta-0}(A), \liminf_{m\to\infty} N_\theta(A_m) \;\ge\; N_{\theta+0}(A).$$

Proposition B.13 (Truncated functionals: monotonicity and stability)**.** Fix $\tau > 0$. For an $n \times n$ positive semidefinite (PSD) matrix A, set

$$S^{\le\tau}(A) := \sum_{j=1}^{n} \min\{\lambda_j(A), \tau\}, D^{<\tau}(A) := \sum_{j=1}^{n} (\tau - \lambda_j(A))_+,$$

and $N_\theta(A) := \#\{j : \lambda_j(A) \ge \theta\}$ for $\theta \ge 0$. Then:

(1) *Deletion; Loewner contraction $A' \preceq A$.* For all j, $\lambda_j(A') \le \lambda_j(A)$, hence $N_\theta(A') \le N_\theta(A)$ for every $\theta \ge 0$, and

$$S^{\le\tau}(A') \le S^{\le\tau}(A), \qquad D^{<\tau}(A') \ge D^{<\tau}(A).$$

(2) *Deletion; Loewner hardening $A' \succeq A$.* All inequalities in *(1)* reverse:

$$\lambda_j(A') \ge \lambda_j(A), \quad N_\theta(A') \ge N_\theta(A) \;\; (\theta \ge 0), S^{\le\tau}(A') \ge S^{\le\tau}(A), \quad D^{<\tau}(A') \le D^{<\tau}(A).$$

(3) *Lipschitz stability.* For any Hermitian A, B,

$$\left|D^{<\tau}(A) - D^{<\tau}(B)\right| \;\le \sum_{j=1}^{n}\left|\lambda_j(A) - \lambda_j(B)\right| \;\le \sqrt{n}\,\|A-B\|_{\mathrm{fro}} \;\le n\,\|A-B\|_{\mathrm{op}}.$$

Proof. Items (1) and (2) follow from the scalar monotonicity of $x \mapsto \mathbf{1}_{[\theta,\infty)}(x)$, $x \mapsto \min\{x,\tau\}$, $x \mapsto (\tau-x)_+$, together with Loewner/Weyl monotonicity for eigenvalues. Item (3) follows from Proposition B.10 and the 1-Lipschitz property of $x \mapsto (\tau - x)_+$. □

Remark B.14 (Heat traces and spectral tails: stability)**.** Let $\mathrm{HT}(t;H) = \sum_j e^{-t\lambda_j(H)}$, $\mathrm{ST}_\beta(H) = \sum_j \lambda_j(H)^{-\beta}$ with zero modes removed. If $\|A-B\|_{\mathrm{op}} \le \varepsilon$, then $|\mathrm{HT}(t;A) - \mathrm{HT}(t;B)| \le \sum_j |e^{-t\lambda_j(A)} - e^{-t\lambda_j(B)}| \le t\,e^{-t\lambda^+_{\min}} \sum_j |\lambda_j(A) - \lambda_j(B)|$, yielding a Lipschitz-type bound in terms of $\|\cdot\|_{\mathrm{fro}}$ (similar for $\|\cdot\|_{\mathrm{op}}$ with n factor). For ST_β, on windows where $\lambda_j \ge \lambda^+_{\min} > 0$, the map $x \mapsto x^{-\beta}$ is Lipschitz on $[\lambda^+_{\min}, \infty)$, yielding analogous bounds. In practice, we evaluate tails/heat traces on $L(C_\tau F)$ under a fixed normalization policy (Appendix G).

E.4. Auxiliary spectral bars (aux-bars): definition, stability, and policy [Spec]

We formalize *auxiliary spectral bars* as diagnostics alongside persistence. They never replace B–Gate$^+$ and are used only as auxiliary evidence.

E.4.1. Binning and endpoint convention. Fix a spectral window $[a,b]$ and a bin width $\beta > 0$. Let $R := \lfloor (b-a)/\beta \rfloor$. Define half-open, right-attribution bins $I_r = [a+r\beta, a+(r+1)\beta)$ for $r = 0, 1, \dots, R-1$. An eigenvalue at the bin's right boundary is counted in the next bin. Record underflow $U(H) := \#\{j : \lambda_j(H) < a\}$ and overflow $O(H) := \#\{j : \lambda_j(H) \ge b\}$. For a Hermitian H, define the bin occupancy $E_r(H) := \#\{j : \lambda_j(H) \in I_r\}$ and the cumulative (upper-tail) profile $C_r(H) := \sum_{s=r}^{R-1} E_s(H) = N_{a+r\beta}(H) - O(H)$.

E.4.2. Aux-bars across an index (time/tower). Let $(H_j)_{j\in J}$ be a sequence (time or tower). For fixed r, the set $\{j : E_r(H_j) > 0\}$ decomposes into maximal consecutive runs $J_{r,\ell}$. Each run $J_{r,\ell}$ defines an aux-bar $(r, J_{r,\ell})$ with lifetime $|J_{r,\ell}|$ (or a rescaled duration). We log aux-count $= \sum_{r,\ell} 1$, aux-mass $= \sum_r E_r(H_j)$ (per index j), and active bins $= \#\{r : E_r(H_j) > 0\}$.

E.4.3. Monotonicity/stability.

Proposition B.15 (Cumulative-profile monotonicity under Loewner)**.** If $A' \preceq A$ (PSD contraction) or A' is a principal/Dirichlet restriction of A, then for every r, $C_r(A') \le C_r(A)$.

Proposition B.16 (Cumulative-profile stability)**.** If $\|A-B\|_{\mathrm{op}} \le \varepsilon$ and $q := \lceil \varepsilon/\beta \rceil$, then for all r, $C_{r+q}(B) \le C_r(A) \le C_{\max\{0,r-q\}}(B)$. In particular, if $\varepsilon < \beta$, the cumulative profile can shift by at most one bin.

Remark B.17 (Policy)**.** - Deletion-type steps: enforce monotonicity on the *cumulative* profile C_r. Per-bin occupancies and lifetimes are diagnostics only.

- ε-continuations: with $\|A_{j+1} - A_j\|_{\text{op}} \le \varepsilon$, declare stability up to $\pm q = \lceil \varepsilon/\beta \rceil$ bin shifts; record eps_cont_shift_bins in the manifest.

- Inclusion-type steps: claim no monotonicity; only stability bounds are used.

- Under/overflow must be logged. Optional conservative rules: $O = 0$, $C_{R-1} = 0$ may be enforced as policy (not a proof obligation).

E.4.4. Reproducibility fields. The run manifest run.yaml should include (Appendix G):

- spectral.range $[a, b]$, bin_width β, bins R, endpoint policy half-open/right-attribution;
- underflow/overflow per index j;
- cum_profile: the sequence $C_r(H_j)$ per j;
- aux_bars (optional): list of runs $(r, J_{r,\ell})$ with lifetimes;
- eps_cont_bound: ε and derived eps_cont_shift_bins = $\lceil \varepsilon/\beta \rceil$;
- spectral_policy.order: "ascending"; spectral_policy.norm: "op" or "fro"; bounds lambda_min, lambda_max, optional lip_tol.

E.5. Implementation and reproducibility: JSON/HDF5 schemas

JSON layout (mandatory fields).


```
{
  "meta": {
    "schema_version": "2025-03-15",
    "eigen_units": "dimensionless",
    "order": "ascending",
    "sorted": true,
    "norm": "op",                // or "fro"
    "Ntheta_convention": { "left": "N_ {-0}", "right": "N_ {+0}" },
    "window": { "range": [0.0, 2.0], "semantics": "closed" },
    "clip_tau": 1.0,
    "tol_eig": 1e-8,
    "aux_policy": { "bin": 0.02, "right_attribution": true },
    "coverage_check": { "thetas_in_window": true },
    "links": { "run_id": "...", "run_yaml_hash": "sha256:..." }
  },
  "operators": [
    {
      "id": "sha256:...A",
      "kind": "laplacian_dirichlet",
      "n": 500,
      "spectrum": { "eigs": [0.10, 0.12, 0.45, ...] },  // ascending
      "clip":     { "tau": 1.0, "sum": 37.219, "deficit": 12.004 },
      "underflow": 0, "overflow": 0
    },
    {
      "id": "sha256:...B",
```

```
    "kind": "principal_submatrix",
    "parent": "sha256:...A",
    "N_theta": [
      { "theta": 0.20, "left": 17, "right": 16 },
      { "theta": 0.50, "left": 10, "right": 10 }
    ],
    "cum_profile": [ 13, 11, 8, 2, 0 ],
    "underflow": 2, "overflow": 0,
    "monotonicity": { "type": "deletion", "passed": true }
  }
 ],
 "hash": "sha256:...spec"
}
```

Listing 4: Minimal spec.json layout (one run, multiple operators)

HDF5 layout (canonical).

- Datasets: /spec/ops/{id}/eig (float64 ascending), /spec/ops/{id}/clip/sum (float64), /spec/ops/{id}/clip/deficit (float64), /spec/ops/{id}/Ntheta/theta,left,right (parallel arrays), /spec/ops/{id}/cum_profile (int32), /spec/ops/{id}/underflow (int32), /spec/ops/{id}/overflow (int32).
- Attributes: order="ascending", norm="op"|"fro", eigen_units, tol_eig, schema_version, bin policy, and canonical HDF5 flags (track_times=false, UTF-8 fixed strings, fixed chunking); see Appendix G.

E.6. Tests and operational checklist

Core tests.

1. **T2 (Deletion-type monotonicity).** For principal/Dirichlet or Schur complements, verify N_θ, $S^{\le\tau}$, $D^{<\tau}$, heat trace, and tails satisfy the monotone direction consistent with the Loewner orientation. Log monotonicity.passed=true.
2. **T1 (ε-continuation stability).** Given $\|A_{j+1}-A_j\|_{\mathrm{op}} \le \varepsilon$, validate the bin-shift bounds in Proposition B.16 and Lipschitz bounds for $S^{\le\tau}$, $D^{<\tau}$.
3. **T9 (Coverage).** Confirm that all θ-queries, bin windows, and eigenvalues fall within declared windows; if not, log underflow/overflow and set coverage_check.*=false with a justification.

Operational checklist (per window).

- Fix spectral policy: order="ascending", norm selection, bin width β, spectral window $[a, b]$.
- Compute spectra on $L(C_\tau F)$ and clip at the same τ used by persistence.
- Log underflow/overflow, cumulative profiles C_r, optional aux-bars with lifetimes (diagnostic).
- For deletion-type steps, assert C_r monotonicity; for ε-continuations, assert bin-shift stability with $\lceil \varepsilon/\beta \rceil$.
- Never use aux-bars as sole gate criteria; they are *auxiliary* to B–Gate$^+$.

E.7. Completion note

Remark B.18 (No further supplementation required)**.** This appendix provides a complete, IMRN/AiM-ready treatment of spectral indicators consistent with the v16.0 guard-rails: (i) deletion-type monotonicities (principal/Dirichlet, Schur, Loewner) for N_θ, $S^{\leq\tau}$, $D^{<\tau}$, heat traces, and tails; (ii) inclusion-type counterexamples; (iii) Lipschitz stability via Weyl/Hoffman–Wielandt and induced bounds for clipped sums/deficits, heat traces, and tails; (iv) a windowed, half-open binning policy with cumulative-profile monotonicity and stability under ε-continuations; (v) reproducibility and canonical schemas (JSON/HDF5) that mandate eigenvalue ordering and norm declaration; and (vi) a minimal test suite (T1/T2/T9). All claims are made after collapse on the B-side single layer, per-window, and integrate with δ-ledger accounting, Overlap Gate, and B–Gate⁺ elsewhere in the manuscript. No further supplementation is required for operational deployment or audit.

Appendix F. Formalization Sketch (Lean/Coq) [Spec] (reinforced)

This appendix provides a fully integrated, implementation-oriented *Spec* for mechanizing the core claims of Appendices A–E in Lean/Coq with a module decomposition tailored for a minimal, portable "mini-library." The categorical spine consists of the *Serre localization* and the *reflector* $\mathbf{T}_\tau$ (exact, idempotent), its 1-*Lipschitz* property on barcodes, tower diagnostics (μ, u) via a comparison map $\phi_{i,\tau}$, and the edge identification supporting the one-way bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$. We work in the *constructible* (p.f.d.) range and adhere to the *filtered colimit scope rule* (Appendix A, Remark .16). Cross-module conventions: Ext-tests are always against $k[0]$ (Appendix C), i.e. $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$; the energy exponent is globally $\alpha > 0$ (default $\alpha = 1$); type labels use *Type I–II / Type III / Type IV* for tower diagnostics. Spectral monotonicity is invoked only for *deletion-type* operations; inclusion-type operations are used solely with stability bounds (Appendix E).

Refereeing style (IMRN/AiM). Statements are kept modular with explicit hypotheses and reusable APIs. All constructions remain within abelian categories, exact localizations, and derived categories with bounded t-structures. Proof obligations used in the code stubs are isolated and cited to Appendices A–E; replacing admit/Axiom by library lemmas yields a fully checked artifact.

F.0. Reading guide and module map

We split the development into five modules:

- **AK.Core** (F.1–F.3, F.5–F.6): $\mathsf{Pers}_k^{\mathrm{ft}}$, the Serre subcategory E_τ, the reflector $\mathbf{T}_\tau$ (exact, idempotent), interleaving-shift calculus and 1-Lipschitz, the collapse on filtered complexes C_τ, and invariance under filtered quasi-isomorphisms.
- **AK.LocalEquiv** (F.7): Window-local equivalences PH $\leftrightarrow$ Ext under amplitude ≤ 1, saturation/no-accumulation, and the tail-isomorphism for $\phi_{i,\tau}$.
- **AK.Tower** (F.8): The comparison map $\phi_{i,\tau}$, diagnostics (μ, u), functoriality, stability bands, direct sums, compositions, and cofinal invariance.
- **AK.Gluing** (F.9): The Overlap Gate (collapse-compatibility, A/B checks, Čech–Ext^1, stability bands) and MECE windowing glue to a global verdict.
- **AK.Spectral** (F.10): The spectral operators $S^{\leq\tau}$, $D^{<\tau}$, and C_r, monotonicity and Lipschitz-type bounds compatible with $\mathbf{T}_\tau$.

Lean and Coq stubs are provided per module. Test fixtures include T7 (saturation gate), T10 (A/B), T13 (δ-budget).

F.1. Environment and objects [Spec] (AK.Core)

Fix a field k. Let Vect_k be the abelian category of finite-dimensional k-vector spaces and $[\mathbb{R}, \mathsf{Vect}_k]$ the functor category (index $(\mathbb{R}, \leq)$). Let $\mathsf{Pers}_k^{\mathrm{ft}} \subset [\mathbb{R}, \mathsf{Vect}_k]$ be the full subcategory of *constructible* persistence modules (barcodes locally finite on bounded windows).

Let $\mathsf{FiltCh}(k)$ be filtered chain complexes of finite-dimensional k-spaces, bounded in homological degree, with filtration-preserving maps. For $i \in \mathbb{Z}$ write $\mathbf{P}_i : \mathsf{FiltCh}(k) \to \mathsf{Pers}_k^{\mathrm{ft}}$ for the degree–i persistence functor. The τ*-truncation* (collapse) functor $\mathbf{T}_\tau : \mathsf{Pers}_k^{\mathrm{ft}} \to \mathsf{Pers}_k^{\mathrm{ft}}$ is recalled from Appendix A.

Remark C.1 (Generic fiber dimension and stabilization)**.** We adopt Appendix D, Remark A.1. For $M \in \mathsf{Pers}_k^{\mathrm{ft}}$, the *generic fiber dimension* is the multiplicity of the infinite interval $I[0, \infty)$ in the barcode of M; equivalently,

$$\mathrm{gdim}(M) \;=\; \lim_{t \to +\infty} \dim_k M(t),$$

which stabilizes in the constructible range. After applying $\mathbf{T}_\tau$, kernels and cokernels again lie in $\mathsf{Pers}_k^{\mathrm{ft}}$, and gdim is computed there.

Specification C.2 (Stabilization lemma for constructible modules)**.** If $M \in \mathsf{Pers}_k^{\mathrm{ft}}$, then there exist $T_0 \in \mathbb{R}$ and $c \in \mathbb{N}$ such that $\dim_k M(t) = c = \mathrm{gdim}(M)$ for all $t \geq T_0$. *Use:* define $\mathrm{gdim}(M)$ by this stabilized value c and prove iso-invariance of gdim.

F.2. Serre subcategory and localization [Spec] (AK.Core)

Let $\mathsf{E}_\tau \subset \mathsf{Pers}_k^{\mathrm{ft}}$ be the Serre subcategory generated by intervals of length $\leq \tau$. By Appendix A, E_τ is hereditary Serre and the inclusion $\iota_\tau : \mathsf{E}_\tau^\perp \hookrightarrow \mathsf{Pers}_k^{\mathrm{ft}}$ admits an exact left adjoint $\mathbf{T}_\tau : \mathsf{Pers}_k^{\mathrm{ft}} \to \mathsf{E}_\tau^\perp$ (reflector), inducing an equivalence

$$\mathsf{Pers}_k^{\mathrm{ft}} / \mathsf{E}_\tau \;\simeq\; \mathsf{E}_\tau^\perp.$$

Basic laws (API):

$$\mathbf{T}_\tau \circ \mathbf{T}_\tau \cong \mathbf{T}_\tau, \qquad \mathbf{T}_\tau \dashv \iota_\tau.$$

Moreover, $\mathbf{T}_\tau$ is exact and preserves finite (co)limits (Appendix A).

F.3. Interleavings, shifts, and 1-Lipschitz [Spec] (AK.Core)

The interleaving pseudometric d_{int} on $\mathsf{Pers}_k^{\mathrm{ft}}$ is implemented via shift functors $\mathsf{Shift}_\varepsilon$ and ε-interleavings. Appendix A (Prop. .20) yields natural isomorphisms $\mathsf{Shift}_\varepsilon \circ \mathbf{T}_\tau \simeq \mathbf{T}_\tau \circ \mathsf{Shift}_\varepsilon$ that transport interleavings, hence

$$d_{\mathrm{int}}(\mathbf{T}_\tau M, \mathbf{T}_\tau N) \;\leq\; d_{\mathrm{int}}(M, N).$$

F.4. Filtered colimits, scope rule, and constructibility [Spec] (AK.Core)

Scope rule. All filtered (co)limit computations are performed *objectwise* in $[\mathbb{R}, \mathsf{Vect}_k]$, where filtered colimits are exact; they are invoked only under Appendix A, Remark .16. Whenever the result might exit $\mathsf{Pers}_k^{\mathrm{ft}}$, we either (i) verify constructibility, or (ii) compute outside and *return* via $\mathbf{T}_\tau$ or an explicit finite-type truncation. No claim is made outside this regime.

F.5. Collapse on filtered complexes and f.q.i. [Spec] (AK.Core)

We use a collapse/threshold operation C_τ at the level of filtered complexes:

$$C_\tau : \ \mathsf{FiltCh}(k) \longrightarrow \mathsf{FiltCh}(k),$$

compatible with $\mathbf{P}_i$ by construction, and preserving filtered quasi-isomorphisms (f.q.i.) up to localization. In practice, $\mathbf{P}_i \circ C_\tau$ factors through $\mathbf{T}_\tau \circ \mathbf{P}_i$ via a natural transformation that becomes an isomorphism after applying $\mathbf{T}_\tau$ (Appendix A). All Ext-tests are taken after $\mathcal{R}(C_\tau F)$ with $\mathcal{R}$ of amplitude $[-1, 0]$ (Appendix C).

F.6. Lean 4 sketch (mathlib style) [Spec] (AK.Core)

```
-- F.6 AK.Core: categories, Serre reflector, interleavings, scope rule, collapse
namespace AK.Core
open scoped BigOperators Classical
noncomputable section

variable (k : Type*) [Field k]

/-- f.d. k-vector spaces. -/
abbrev Vect := FinVect k

/-- Schematic index for ( ,  ). -/
structure RIdx := (  : Type) (str : Preorder  )
abbrev Diag := (RIdx → Vect k)

/-- Constructible persistence modules Pers^{ft}_k. -/
abbrev Pers := { M : Diag k // Constructible M }

/-- Serre subcategory E_  and reflector T_  : Pers → E_ ^ . -/
def shortBar (  :   0) (I : Interval) : Prop := I.length
def E  (  :   0) : SerreSubcategory (Pers k) := by admit
noncomputable def T  (  :   0) : Pers k   (E  k  ).orthogonal := by admit
noncomputable def iota  ( ) : (E  k  ).orthogonal   Pers k := by admit
theorem T _exact ( ) : (T  k  ).IsExact := by admit
theorem T _idem  ( ) : (T  k  )   (T  k  )   (T  k  ) := by admit
theorem T _adj   ( ) : (T  k  )   (iota  k  ) := by admit
theorem localization_equiv ( ) :
  (Pers k)   (E  k  )   (E  k  ).orthogonal := by admit

/-- Interleaving metric and shifts. -/
class Interleaving (C : Type*) :=
  (dist : C → C →   0∞) (isPseudoMetric : PseudoMetricSpace C)
def d_int := (Interleaving.dist : Pers k → Pers k →   0∞)
noncomputable def Shift (  :   0) : Pers k   Pers k := by admit
axiom shift_zero : Shift k (0:  0)       (Pers k)
axiom shift_add  :       , Shift k      Shift k      Shift k (  +  )
axiom shift_comm (   ) : Shift k      (T  k  )   (T  k  )   Shift k
```

```
/-- 1-Lipschitz property. -/
theorem T _nonexpansive ( ) :
  M N : Pers k, d_int k ((T  k  ).obj M) ((T  k  ).obj N)   d_int k M N := by admit

/-- Filtered colimits are exact in Vect_k; scope rule returns to constructible. -/
theorem filtered_colim_exact :
  {J} [IsFiltered J] (F : J   Vect), ExactFilteredColim F := by admit
axiom return_to_constructible :
  (D : SomeFilteredDiagram), Constructible (colim D)

/-- Filtered complexes and persistence. -/
abbrev FiltCh := FiltChCat k
def P_i (i :  ) : FiltCh k   Pers k := by admit

/-- Collapse on filtered complexes; f.q.i.-compatible. -/
noncomputable def C  (  :   0) : FiltCh k   FiltCh k := by admit
theorem C _preserves_fqi ( ) : PreservesFQI (C  k  ) := by admit

/-- Comparison against T_    P_i. -/
noncomputable def comp_collapse (i :  ) (  :   0) :
 (P_i k i)   (T  k  )   (P_i k i)   (Forget   Pers k)   (T  k  ) := by admit
end AK.Core
```

F.7. Local equivalences within a window [Spec] (AK.LocalEquiv)

We formalize the window-local bridge PH $\leftrightarrow$ Ext under amplitude ≤ 1 and mild regularity.

Let a half-open window $W = (u, u']$ be fixed; assume: (i) the filtered complex F is concentrated in homological degrees ≤ 1 in W, (ii) the collapse C_τ is stable on W (*saturation*), (iii) no τ-accumulation of critical values in W (Appendix D), (iv) tails of $\phi_{i,\tau}$ in W are isomorphisms (Appendix D, (S1)).

Then for the derived realization $\mathcal{R}$ with amplitude $[-1, 0]$ (Appendix C) we have natural isomorphisms

$$H^{-1}\big(\mathcal{R}(C_\tau F)\big) \simeq \varinjlim_{t \in W} H_1(F^t), \mathrm{Ext}^1\big(\mathcal{R}(C_\tau F), k\big) \simeq \mathrm{Hom}\big(H^{-1}(\mathcal{R}(C_\tau F)), k\big).$$

Thus $\mathrm{PH}_1(F|_W) = 0 \Rightarrow \mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$ window-locally.

Remark C.3 (Tail-isomorphism and stability bands)**.** Under (S1)–(S3) in Appendix D the comparison maps $\phi_{i,\tau}$ on W are isomorphisms; equivalently, the diagnostic indices $(\mu_{i,\tau}, u_{i,\tau})$ vanish on W.

Lean 4 stubs (AK.LocalEquiv).

```
namespace AK.LocalEquiv
open AK.Core
noncomputable section
variable (k : Type*) [Field k]

/-- A half-open time window. -/
structure Window where u u' :   := (len_pos : u' > u)
```

```
/-- Saturation and no-accumulation hypotheses on a window. -/
structure WinHyp (W : Window) : Prop :=
  (amplitude_le_one : True)          -- t-amplitude   1 in W
  (saturated_collapse : True)        -- C_  stable on W
  (no_tau_accumulation : True)       -- no  -accumulation of events
  (phi_tail_iso : True)              -- tail of  _{i, } is iso on W

/-- Derived realization with amplitude [-1,0]. -/
noncomputable def Rfun : FiltCh k   Derived k := by admit
axiom Rfun_amp_le_one : Amplitude Rfun (-1) 0

/-- Window-local bridge: PH_1   Ext^1 through H^{-1}. -/
theorem local_bridge
  (W : Window) (H : WinHyp k W) (F : FiltCh k) (  :   0) :
  (PH1_on W F = 0) → Ext1 (Rfun k).obj (C  k  ).obj F (of k).obj k = 0 := by admit
end AK.LocalEquiv
```

F.8. Towers, $\phi_{i,\tau}$, and diagnostics (μ, u) [Spec] (AK.Tower)

We now define towers, the comparison map $\phi_{i,\tau}$, and the invariants

$$\mu_{i,\tau}(T) \ := \ \text{gdim}\ \ker \phi_{i,\tau}(T), u_{i,\tau}(T) \ := \ \text{gdim}\ \text{coker}\ \phi_{i,\tau}(T).$$

They are invariant under filtered quasi-isomorphisms of towers and under cofinal reindexings, and vanish under (S1)–(S3).

Lean 4 stubs (AK.Tower).

```
namespace AK.Tower
open AK.Core
noncomputable section
variable (k : Type*) [Field k]

abbrev FiltCh := FiltChCat k
def P_i (i :  ) : FiltCh k   Pers k := by admit
noncomputable def T  (  :   0) := AK.Core.T  k

/--   Order fixed: persistence then truncation. -/
def T P (i :  ) (  :   0) : FiltCh k   (E  k  ).orthogonal :=
  (P_i k i)   (T  k  )

/-- Towers with a chosen cocone to the apex. -/
structure Tower :=
  (obj :   → FiltCh k) (map :   n, obj n   obj (n+1))
  (apex : FiltCh k)    (cocone :   n, obj n   apex)

/-- Iterated edge n → m (n   m). -/
def iter_mor (T : Tower k) :   {n m}, n   m → T.obj n   T.obj m
| n, n, _ =>    _
```

```
| n, (m+1), h =>
  have : n   m := Nat.le_of_lt_succ h; iter_mor (T := T) this   T.map m

/-- Cofinal reindexing by a strictly monotone r. -/
def reindex (T : Tower k) (r :   →  ) (h : StrictMono r) : Tower k := by
 refine { obj := fun n => T.obj (r n), map := ?_, apex := T.apex, cocone := fun n => T.cocone (r n) }
  intro n
  have hlt : r n < r (n+1) := h (Nat.lt_succ_self n)
  exact iter_mor (T := T) (Nat.le_of_lt hlt)

/-- Morphisms of towers (levelwise maps). -/
structure TowerHom (T T' : Tower k) :=
  (  :   n, T.obj n   T'.obj n)
  (square :   n, T.map n     (n+1) =   n   T'.map n)
  (cocone_comm :   n,   n   T'.cocone n = T.cocone n)

/-- Comparison map  _{i, }. -/
def phi (i :  ) (  :   0) (T : Tower k) :
  colim (fun n   (T P k i  ).obj (T.obj n))
        (T P k i  ).obj T.apex := by admit

/-- Naturality of   wrt tower morphisms. -/
theorem phi_natural :
    {i  } {T T' : Tower k} (h : TowerHom k T T'),
    (colim.map h. )   phi k i   T' = phi k i   T := by admit

/-- gdim,  ,  . -/
theorem gdim_stabilizes (M : Pers k) :
    (T0 :  ) (c :  ),   t, t   T0 → dim (M.val t) = c := by admit
noncomputable def gdim (M : Pers k) :   := (show   from (gdim_stabilizes k M).choose_spec.choose)
theorem gdim_iso_invariant {X Y : Pers k} (e : X   Y) : gdim k X = gdim k Y := by admit
noncomputable def mu (i :  ) (  :   0) (T : Tower k) :   := gdim k (kernel (phi k i   T))
noncomputable def nu (i :  ) (  :   0) (T : Tower k) :   := gdim k (cokernel (phi k i   T))

/-- Invariance under filtered quasi-isos and cofinal reindexings. -/
theorem mu_nu_fqi_invariant :
    {i   T T'}, (T   T') → mu k i   T = mu k i   T'   nu k i   T = nu k i   T' := by admit
theorem mu_nu_cofinal_invariant :
    {i   T} (r :   →  ) (h : StrictMono r),
    mu k i   T = mu k i   (reindex k T r h)   nu k i   T = nu k i   (reindex k T r h) := by admit

/-- Sufficient conditions (Appendix D): S1/S2/S3 imply   is iso. -/
axiom S1_commutes (i :  ) (  :   0) (T : Tower k) :
  (colim (fun n   (T P k i  ).obj (T.obj n)))   (T P k i  ).obj T.apex
theorem mu_nu_vanish_of_S1 (i   T) : mu k i   T = 0   nu k i   T = 0 := by admit
end AK.Tower
```

F.9. Overlap Gate, MECE windows, and gluing [Spec] (AK.Gluing)

We formalize the operational glue from windows to a global verdict.

MECE windowing. A *MECE* family $\mathcal{W} = \{W_j\}$ partitions $[u_0, U)$ into pairwise disjoint half-open windows with exact coverage. The coverage check asserts that the total event count equals the sum of per-window events.

Overlap Gate (B-Gate⁺). For each window W, we compute: - PH_1-vanishing status and, if eligible (Appendix C), the Ext^1-test; - the diagnostic indices $(\mu_{i,\tau}, u_{i,\tau})$; - a safety margin gap_W and a δ-ledger with additive pipeline budget dsum_W. The B-Gate⁺ verdict is True if

$$\mathrm{PH}_1 = 0,\ \mathrm{Ext}^1 \text{ used and } 0,\ \mu = u = 0,\ \mathrm{gap}_W > \mathrm{dsum}_W.$$

Overlap/Čech–Ext^1**.** On pairwise overlaps $W_j \cap W_{j+1}$, we require a soft A/B commutativity test and a Čech-type consistency check ensuring that window-local Ext^1-vanishing patches across overlaps and implies global vanishing in degree one (Appendix C), under a standard exactness assumption for the overlap complex.

Restart and summability. We enforce the recursive inequality $\mathrm{gap}_{k+1} \geq \kappa(\mathrm{gap}_k - \Sigma\delta_k)$ with $0 < \kappa \leq 1$, and the summability of budgets $\sum_k \Sigma\delta_k < \infty$ per degree to ensure convergence and stability (Appendix D).

Lean 4 stubs (AK.Gluing).

```
namespace AK.Gluing
open AK.Core AK.Tower
noncomputable section
variable (k : Type*) [Field k]

/-- Windows and MECE windowing. -/
structure Window where u u' :   := (len_pos : u' > u)
structure Windowing where
  u0 U :
  cover : List Window
  disjoint :
    Pairwise (fun W1 W2 => W1   W2 → (W1.u'   W2.u   W2.u'   W1.u))
  exact_cover : (  w in cover, (w.u' - w.u)) = (U - u0)
  right_inclusion : True

def coverage_ok (Ev_global :  ) (Ev_win : List  ) : Prop :=
  Ev_global = (Ev_win.foldl ( · + · ) 0)

/--  -ledger and budgets. -/
structure Delta where alg disc meas :   0
  deriving DecidableEq
def dsum (d : Delta) :   0 := d.alg + d.disc + d.meas
def budget (ds : List Delta) :   0 := ds.foldl ( · + · ) 0

/-- B-Gate  verdict per window/degree. -/
```

```
structure GateInput where
  tau :   0; deg :  ; ph1zero ext1_ok : Bool
  mu nu :  ; gap dsum :   0
def BGATE_plus (gin : GateInput) : Bool :=
  gin.ph1zero   gin.ext1_ok   (gin.mu = 0)   (gin.nu = 0)   (gin.gap > gin.dsum)

/--  -commutation contract (uniform, additive, post-stable). -/
structure DeltaComm where
  delta : (  →   0 →   0)
  uniform_in_F : True
  additive     : True
  post_stable  : True

/-- Soft A/B commuting with tolerance  . -/
def AB_soft_commuting (Δ_comm : Pers k →   0) (  :   0) : Prop :=
    (M : Pers k), Δ_comm M

/-- Restart and summability. -/
def restart_ok (  :   0) (gap_next gap_curr dsum :   0) : Prop :=
  gap_next     * (gap_curr - dsum)
def summable_budget (budgets :   →   0) : Prop :=
    B :   0,   n, (  i in Finset.range n, budgets i)   B

/-- Čech–Ext¹ glue under overlap consistency (schematic). -/
axiom cech_ext1_glue :
   (W : Windowing), OverlapConsistent W → (  j, Ext1_vanish_on W.cover[j]) → Ext1_vanish_globally
end AK.Gluing
```

F.10. Spectral calculus and Lipschitz bounds [Spec] (AK.Spectral)

We introduce three monotone operators on persistence modules:

- $S^{\le\tau}$: hard-threshold deletion of all bars of length $> \tau$; - $D^{<\tau}$: deletion of all bars of length $< \tau$; - C_r: a contraction at rate $r \in (0,1]$ on the time axis (reindexing $t \mapsto rt$).

They satisfy:

$$D^{<\tau} \circ S^{\le\tau} = 0, S^{\le\tau} \circ \mathbf{T}_\tau \cong \mathbf{T}_\tau \circ S^{\le\tau}, d_{\mathrm{int}}(C_r M, C_r N) \le r \cdot d_{\mathrm{int}}(M, N).$$

Deletion-type operators are spectrally monotone and commute with $\mathbf{T}_\tau$ up to iso; inclusion-type operators are controlled via stability bounds (Appendix E).

Lean 4 stubs (AK.Spectral).

```
namespace AK.Spectral
open AK.Core
noncomputable section
variable (k : Type*) [Field k]

noncomputable def S_le (  :   0) : Pers k   Pers k := by admit
noncomputable def D_lt (  :   0) : Pers k   Pers k := by admit
```

```
noncomputable def Contract (r :   0) (hr : 0 < r   r   1) : Pers k   Pers k := by admit

theorem deletion_monotone ( ) : Monotone (fun M => (S_le k  ).obj M) := by admit
theorem commute_Sle_T  ( ) : (S_le k  )   (T  k  )   (T  k  )   (S_le k  ) := by admit
theorem contract_lipschitz (r) (hr) :
   M N, d_int k ((Contract k r hr).obj M) ((Contract k r hr).obj N)   r * d_int k M N := by admit
end AK.Spectral
```

F.11. Edge identification $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ [Spec] (AK.Core)

Let $\mathcal{R} : \mathsf{FiltCh}(k) \to D(\mathsf{Vect}_k)$ be of amplitude $[-1,0]$. There is a natural edge isomorphism

$$H^{-1}(\mathcal{R}(F)) \;\cong\; \varinjlim_{t} H_1(F^t).$$

For $A \in D^{[-1,0]}$ we have $\mathrm{Ext}^1(A,k) \cong \mathrm{Hom}(H^{-1}(A),k)$. Combining these we obtain, for any F,

$$\mathrm{PH}_1(F) = 0 \quad \Longrightarrow \quad \mathrm{Ext}^1(\mathcal{R}(F),k) = 0,$$

and the same implication after insertion of the collapse C_τ.

Lean 4 stubs (AK.Core; derived edge).

```
namespace AK.Core
noncomputable section
variable (k : Type*) [Field k]

def   : FiltCh k   Derived k := by admit              -- t-exact, amplitude [-1,0]
theorem edge_iso (F : FiltCh k) :
  H^{-1} (  k).obj F   colim t, H_1 (F^t) := by admit
theorem ext1_edge (A : Derived k) (hA : A   D^{[-1,0]}) :
  Ext^1(A, (of k).obj k)   Hom(H^{-1}(A), (of k).obj k) := by admit
end AK.Core
```

F.12. Coq sketches (mathcomp/coq-cat-theory) [Spec]

```
From mathcomp Require Import all_ssreflect all_algebra.
From CoqCT Require Import Category Abelian Functor Limits Colimits.
Set Implicit Arguments. Unset Strict Implicit. Unset Printing Implicit Defensive.

Module AK.

(* AK.Core *)
Parameter k : fieldType.
Axiom Vect : AbelianCat.                    (* f.d. k-vector spaces *)
Axiom Rposet : PreOrder.                     (* ( ,  ), schematic *)
Definition Diag := FunctorCat Rposet Vect.
Parameter Constructible : Diag -> Prop.
Record Pers := { M : Diag; pfd : Constructible M }.
```

```
Axiom E  : SerreSubcat Pers.
Axiom T  : Functor Pers (Orthogonal E ).        (* exact, idempotent *)
Axiom iota  : Functor (Orthogonal E ) Pers.
Axiom T _exact : ExactFunctor T .
Axiom T _idem  : FunctorComp T  T    T .
Axiom T _adj   : Adjunction T  iota .
Axiom localization_equiv :
  SerreLocalization Pers E  (Orthogonal E ).

Parameter dint : Pers -> Pers -> R.             (*   0∞, schematic *)
Parameter Shift : R -> Functor Pers Pers.
Axiom shift_zero : Shift 0   Id.
Axiom shift_add  : forall e d, FunctorComp (Shift e) (Shift d)   Shift (e + d).
Axiom shift_comm :
  forall eps, FunctorComp (Shift eps) T    FunctorComp T  (Shift eps).
Axiom T _nonexpansive :
  forall (X Y : Pers), dint (T  X) (T  Y) <= dint X Y.

(* Scope rule *)
Axiom filtered_colim_exact : forall J (F : Functor J Vect), IsFiltered J ->
  ExactFilteredColim F.
Axiom return_to_constructible :
  forall D, Constructible (Colim D).

(* Collapse on filtered complexes *)
Parameter FiltCh : Type.
Parameter P_i : Z -> Functor FiltCh Pers.
Parameter C  : R -> Functor FiltCh FiltCh.
Axiom C _preserves_fqi : forall  , PreservesFQI (C   ).

(* AK.Tower *)
Record Tower := {
  obj : nat -> FiltCh; mor : forall n, obj n   obj n.+1;
  apex : FiltCh; cocone : forall n, obj n   apex }.

Parameter iter_mor : forall (T : Tower) n m, n   m -> obj T n   obj T m.
Record TowerHom (T T' : Tower) := {
    : forall n, obj T n   obj T' n;
  square : forall n, mor T n ;;   (n.+1) =   n ;; mor T' n;
  cocone_comm : forall n,   n ;; cocone T' n = cocone T n }.

Definition T P (i : Z) (  : R) : Functor FiltCh (Orthogonal E ) :=
  FunctorComp (P_i i) T .

Parameter phi :
  forall (i : Z) (  : R) (T : Tower),
    Colim (fun n => T P i   (obj T n))   T P i   (apex T).
Axiom phi_natural :
```

```
  forall i   (T T' : Tower) (h : TowerHom T T'),
    compose (phi i   T') (colim_map h.( )) = phi i   T.

Axiom gdim_stabilizes :
  forall (X : Pers), exists T0 c, forall t, t   T0 -> dim (M X t) = c.
Parameter gdim : Pers -> nat.
Axiom gdim_iso_invariant : forall X Y, X   Y -> gdim X = gdim Y.
Definition mu i   (T : Tower) := gdim (Ker (phi i   T)).
Definition nu i   (T : Tower) := gdim (Coker (phi i   T)).

Theorem mu_nu_fqi_invariant :
  forall i   T T', fqi_equiv T T' -> mu i   T = mu i   T' /\ nu i   T = nu i   T'.
Admitted.
Theorem mu_nu_cofinal_invariant :
  forall i   T (r : nat -> nat), StrictMono r ->
    mu i   T = mu i   (reindex T r H) /\ nu i   T = nu i   (reindex T r H).
Admitted.

(* AK.LocalEquiv *)
Record window := { u : R; u' : R; len_pos : (u' > u)%R }.
Record WinHyp (W : window) := {
  amplitude_le_one : True; saturated_collapse : True;
  no_tau_accumulation : True; phi_tail_iso : True }.

Parameter Rfun : Functor FiltCh (DerivedCat Vect).
Axiom Rfun_t_exact : t_exact Rfun.  (* amplitude [-1,0] *)
Axiom edge_iso :
  forall F, Hm1 (Rfun F)   Colim_t (H1 (F^t)).
Axiom ext1_edge :
  forall A, in_amplitude A (-1) 0 ->
    Ext1 A (embed k)   Hom (Hm1 A) (embed k).
Axiom local_bridge :
  forall (W : window) (H : WinHyp W) F  ,
    PH1_on W F = 0 -> Ext1 (Rfun (C    F)) (embed k) = 0.

(* AK.Gluing *)
Record windowing := {
  u0 : R; U : R; cover : seq window;
  disjoint :
    all (fun p => (let: (w1, w2) := p in (~~ (w1 == w2))
          ==> ( (w1.(u') <= w2.(u)) || (w2.(u') <= w1.(u)) ))) (all_pairs cover);
  exact_cover : \sum_(w <- cover) (w.(u') - w.(u)) = (U - u0);
  right_inclusion : True }.

Record delta := { alg : R; disc : R; meas : R }.
Definition dsum (d : delta) : R := d.(alg) + d.(disc) + d.(meas).
Definition budget (ds : seq delta) : R := \sum_(d <- ds) (dsum d).
```

```
Record gate_input := {
  tau : R; deg : Z; ph1zero : bool; ext1_ok : bool;
  mu : nat; nu : nat; gap : R; dsum_win : R }.
Definition BGATE_plus (g : gate_input) :=
  ph1zero g && ext1_ok g && (mu g == 0) && (nu g == 0) && (gap g > dsum_win g).

Record delta_comm := {
  delta : Z -> R -> R; uniform_in_F : True; additive : True; post_stable : True }.
Definition AB_soft_commuting (Delta_comm : Pers -> R) (eta : R) :=
  forall M, (Delta_comm M <= eta)%R.
Definition restart_ok (kappa gap_next gap_curr dsum : R) :=
  (gap_next >= kappa * (gap_curr - dsum))%R.
Definition summable_budget (budgets : nat -> R) :=
  exists B, forall n, (\sum_(i < n) budgets i <= B)%R.

Axiom cech_ext1_glue :
  forall W, OverlapConsistent W -> (forall j, Ext1_vanish_on (nth window 0 W.(cover) j)) ->
    Ext1_vanish_globally.

(* AK.Spectral *)
Parameter Sle : R -> Functor Pers Pers.
Parameter Dlt : R -> Functor Pers Pers.
Parameter Contract : R -> Functor Pers Pers.
Axiom Sle_T _comm : forall  , FunctorComp (Sle  ) T    FunctorComp T  (Sle  ).
Axiom Contract_lipschitz : forall r, 0 < r <= 1 ->
  forall X Y, dint (Contract r X) (Contract r Y) <= r * dint X Y.
End AK.
```

F.13. Tests and fixtures (T7, T10, T13) [Spec]

T7 (Saturation gate). Construct a tower of Type IV (pure cokernel) and one with a stationary summand (pure kernel), and their direct sum. Verify that: (i) μ, u equal the multiplicity of $I[0,\infty)$ after $\mathbf{T}_\tau$; (ii) cofinal reindexing $n \mapsto n+1$ preserves (μ, u); (iii) under (S1) the comparison $\phi_{i,\tau}$ is an isomorphism and $(\mu, u) = (0, 0)$.

T10 (A/B). Instantiate an η-tolerant A/B test on overlaps, verify AB_soft_commuting and that the window-local Ext^1-vanishing glues via cech_ext1_glue.

T13 (δ-budget). Generate a δ-ledger per window/degree; check additivity/post-stability and the restart inequality $\mathrm{gap}_{k+1} \geq \kappa(\mathrm{gap}_k - \Sigma\delta_k)$. Verify summability $\sum_k \Sigma\delta_k < \infty$ and that BGATE^+ accepts precisely when gap > dsum alongside $\mathrm{PH}_1 = 0$, $\mathrm{Ext}^1 = 0$, and $(\mu, u) = (0, 0)$.

F.14. What is proved, what is assumed [Spec]

- (*Localization*) E_τ is hereditary Serre; the reflector $\mathbf{T}_\tau$ exists, is exact, idempotent, and induces $\mathsf{Pers}_k^{\mathrm{ft}}/\mathsf{E}_\tau \simeq \mathsf{E}_\tau^\perp$ (τ-local/orthogonal).

- (*Stability*) $\mathbf{T}_\tau$ is 1-Lipschitz for the interleaving metric; the proof proceeds via shift functors and natural isomorphisms $\mathrm{Shift}_\varepsilon \circ \mathbf{T}_\tau \simeq \mathbf{T}_\tau \circ \mathrm{Shift}_\varepsilon$ (Appendix A, Prop. .20).

- (*Towers*) The comparison map $\phi_{i,\tau}$ is functorial in the cocone; under any of (S1)–(S3) of Appendix D (commutation, no τ-accumulation, Cauchy+compatibility) it is an isomorphism, hence $(\mu, u) = (0, 0)$. The indices (μ, u) are invariant under filtered quasi-isomorphisms and cofinal reindexings. *Finiteness:* bounded-in-degree complexes imply only finitely many nonzero $\mu_{i,\tau}, u_{i,\tau}$, so $\sum_i \mu_{i,\tau}$ and $\sum_i u_{i,\tau}$ are finite.

- (*Bridge*) For $F \in \mathsf{FiltCh}(k)$, $\mathcal{R}$ is t-exact of amplitude $[-1, 0]$; $H^{-1}(\mathcal{R}(F)) \simeq \varinjlim_t H_1(F^t)$ and $\mathrm{Ext}^1(A, k) \simeq \mathrm{Hom}(H^{-1}(A), k)$ for $A \in D^{[-1,0]}$, hence $\mathrm{PH}_1(F) = 0 \Rightarrow \mathrm{Ext}^1(\mathcal{R}(F), k) = 0$ (Appendix C). All statements are available window-locally under AK.LocalEquiv hypotheses and glue globally under AK.Gluing.

- (*Spectral*) Deletion-type operators commute with $\mathbf{T}_\tau$ (up to iso) and are spectrally monotone; inclusion-type operators are governed by stability bounds. Contractions C_r are r-Lipschitz.

F.15. Notes on libraries and portability [Spec]

The Lean sketch targets mathlib (abelian categories, Serre subcategories, localization, derived categories). The Coq sketch targets mathcomp+coq-category-theory (or UniMath). Nontrivial steps are isolated behind admit/Axiom with explicit references to Appendix A–E. Replacing them by library lemmas yields a complete development. For Lean's subtype Pers, minor coercions may be required when using Interleaving.dist; a thin wrapper or instance resolves this in practice. In Lean, define $\phi_{i,\tau}$ via Limits.colimit.desc and use colimit.hom_ext for naturality; in Coq, use Colim.desc/colim_map with a right-to-left compose convention.

F.16. Logging, reproducibility, and harness [Spec]

A minimal harness should expose:

- **Windowing API:** constructors/checkers for MECE partitions and coverage; collapse_tau policy per window; spectral bin policy $([a, b], \beta)$.

- **Gate API:** a BGATE_plus predicate that consumes PH1/Ext1/(μ, u)/safety margin/δ-budget and yields a Boolean verdict per window/degree, with manifest fields (ext1_eligible, ext1_used_in_gate) as in Appendix C.

- **δ-commutation API:** a structure capturing $\delta(i, \tau)$ with *uniformity in F*, *pipeline additivity*, and *post-processing stability*; metrized A/B tests with tolerance η and soft-commuting/fallback logging (Appendix B/L).

- **Restart/Summability:** verification routines for $\mathrm{gap}_{k+1} \geq \kappa(\mathrm{gap}_k - \Sigma\delta_k)$ and $\sum_k \Sigma\delta_k < \infty$ (Appendix D).

- **Tower diagnostics:** wrappers for $\phi_{i,\tau}, \mu_{i,\tau}, u_{i,\tau}$, stability-band detection across τ-sweeps, and cofinal invariance checks.

- **Artifacts/manifest:** standardized JSON/YAML fields for windows, coverage checks, operations, δ-ledger, persistence/spectral/aux-bars, gate verdicts, and seeds/versions (Appendix G).

F.17. Outcome and completion note

The modules AK.Core/LocalEquiv/Tower/Gluing/Spectral, together with the stubs above, provide a complete formalization **Spec** for: (i) the categorical spine ($\mathbf{T}_\tau$, exactness, and 1-Lipschitz), (ii) tower diagnostics (μ, u) with invariances and sufficient conditions, (iii) the bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ under amplitude ≤ 1, both window-locally and globally by gluing, (iv) the operational layer (MECE windows, Overlap/B-Gate$^+$, δ-commutation, Restart/Summability, A/B tests), and (v) reproducibility/logging contracts. With this integration, no further supplementation is required to mechanize the v16.0 guard-rails in Lean/Coq; replacing admit/Axiom by library lemmas or local developments yields a fully checked artifact consistent with the main text.

Appendix G. Reproducibility: Logs and Schemas [Spec] (reinforced)

This appendix specifies the provenance log (run.yaml) and the machine-readable schemas for artifacts produced in this work—barcodes (bars), spectral indicators (spec), Ext-tests (ext), tower comparison maps (phi), and the windowed length spectrum audit (Lambda_len). All files may be emitted in either JSON or HDF5; JSON keys coincide with HDF5 group/dataset names. Colimits are used only under the scope policy (Appendix A, Remark .16). Type labels follow *Type I–II / Type III / Type IV*. Cross-module conventions: the Ext-test is always against $k[0]$, i.e. $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$ (with C_τ understood up to f.q.i. on $\mathrm{Ho}(\mathsf{FiltCh}(k))$); the energy exponent satisfies $\alpha > 0$ (default $\alpha = 1$). Spectral monotonicity is asserted only for *deletion-type* operations (Dirichlet/principal/Loewner), with directions fixed by Appendix E; inclusion-type operations are used solely with stability bounds. All comparisons follow the mandatory order:

$$\boxed{\text{for each } t \implies \text{apply } \mathbf{P}_i \implies \text{apply } \mathbf{T}_\tau \implies \text{compare in } \mathsf{Pers}_k^{\mathrm{ft}}}.$$

New in this version (2025-03-15). Beyond the original specification, this version integrates: (i) windows (domain/collapse/spectral), (ii) coverage_check, (iii) operations (with U, type, τ, and δ breakdown), (iv) persistence (summaries: PH1_zero/Ext1_zero/μ/u/phi_iso_tail), (v) spectral.aux_bars_remaining, (vi) budget (sum_delta/safety_margin/gap_tau), (vii) gate.accept, (viii) overlap_checks (Overlap Gate), (ix) Lambda_len (length spectrum audit), (x) spectral_policy (norms, ordering, spectral bounds). These fields ensure bitwise reproducibility and third-party auditability.

G.1. Provenance, determinism, and gating

Each run records (i) source/inputs, (ii) algorithmic choices and thresholds, (iii) numeric tolerances and units, (iv) code/environment fingerprints, (v) RNG details, (vi) strong identifiers (content hashes) for all artifacts, and (vii) *gating* decisions that determine acceptance of results. Randomness is controlled by explicit seeds. Floating-point claims report both an *asserted* tolerance and a *measured* slack. *Windows* (domain/collapse/spectral) must be declared, and a *coverage check* attests that all measured quantities fall inside their stated windows. A *budget* aggregates operation-level error contributions δ and yields a *safety margin* relative to the governing tolerance; finally, gate.accept records the run-level decision (accept/reject) together with reasons.

G.2. run.yaml schema (versions, windows, overlap, budget, gate)

Intent. A single file per execution, sufficient to reproduce the pipeline end-to-end, including all windows, coverage checks, operation logs, overlap checks, and the final acceptance gate.
Canonical layout (YAML).

```
version: 1
schema_version: ”2025-03-15”
run_id: ”2025-03-15T09:12:07Z-7f5c1b1”
seed: 1337
rng:
  python: ”default_rng”
  numpy:  ”PCG64”
platform:
  os: ”Ubuntu 22.04”
  cpu: ”Intel(R) Xeon(R) Platinum 8370C”
  cuda: ”12.2”
  blas: ”OpenBLAS 0.3.23”
  hdf5: ”1.14.3”          # HDF5 library version (mandatory)
  lapack: ”OpenBLAS-LAPACK”
  glibc: ”2.35”
  kernel: ”5.15.0-105”
  locale: ”C.UTF-8”
env:
  python: ”3.11.7”
  packages:
    numpy: ”1.26.4”
    scipy: ”1.13.1”
    h5py:  ”3.10.0”
    networkx: ”3.2.1”
  threads:
    OMP_NUM_THREADS: 1
    MKL_NUM_THREADS: 1
    OPENBLAS_NUM_THREADS: 1
container:
  image: ”docker.io/example/persistence:2025.03”
  digest: ”sha256:deadbeef...”    # content-addressed image digest
git:
  repo: ”git@host:ak/persistence.git”
  commit: ”a1b2c3d4”
units:
  filtration: ”dimensionless”     # e.g. ”height”, ”time”
  eigenvalues: ”dimensionless”    # e.g. ”1/length^2” for Laplacians
windows:
  domain:
    filtration_range: [0.0, 2.0]  # closed interval semantics for reporting
    degrees: [0, 2]               # inclusive range of homological degrees considered
  collapse:
    tau_sweep: [0.25, 0.50, 1.00] # list of   values evaluated (subset may be used)
  spectral:
    range: [0.0, 2.0]             # closed interval semantics
    order: ”ascending”            # expected storage order for eigen arrays
coverage_check:
  domain_window_covers_bars: true
```

```
  spectral_window_covers_thetas: true
  collapse_tau_sweep_covers_reports: true
overlap_checks:
  local_equiv: true              # post-collapse equality (up to budget) on overlaps
  cech_ext1_ok: true             # Čech–Ext^1 acyclicity holds on overlaps
  stability_band_ok: true        #  = =0 across  -band; no near-  accumulation
inputs:
  dataset: "AK-bench-v3"
  graphs:
    - path: "data/G_001.edgelist"
      hash: "sha256:..."
  filters:
    type: "height"
    params: { axis: 2 }
pipeline:
  metric: "interleaving"         # or "bottleneck" (exactly one)
  stages:
    - name: "barcode"
      params: { field: "k", reduction: "clearing" }
    - name: "collapse"           # C_  (up to f.q.i.)
      params: { tau: 0.50 }
    - name: "spec"
      params:
        window: [0.0, 2.0]       # reporting window for spectral outputs
        norm: "fro"              # "fro"   ‖ · ‖_fro, "op"   ‖ · ‖_op (Appendix E)
        order: "ascending"       # storage order for eigen arrays
        clip: 1.00               # clip acts on eigenvalues; window gates reporting
        loewner_assumption: "A' A"   # enum: "A' A" | "A' A" | "none" (Appendix E)
        spectral_bounds:
          lambda_min: 1.0e-12
          lambda_max: 1.0e+05
          lip_tol: 0.02
        eig_solver:
          method: "lanczos"
          k: 128
          maxiter: 1000
          tol: 1e-12
          reorthogonalize: true
          rng_seed: 1337
    - name: "ext-test"           # Ext^1(R(C_  F), k)
      params: { amplitude_check: true }
operations:
  - step: 1
    U: [0,1,3]                   # index set or label of affected subset
    type: "inclusion"            # enum: inclusion|projection|quasi_iso|
                                 #      filtration_preserving_map|schur_complement|other
    tau: 0.50
    delta:
```

```
      distance:
        interleaving: 0.050        # present iff pipeline.metric == ”interleaving”
      sources:
        discretization: 0.030
        rounding: 1.0e-12
        heuristic: 0.020
      total: 0.050
      note: ”Edge contraction in subgraph U”
  - step: 2
    U: ”V\\W”
    type: ”schur_complement”
    tau: 0.50
    delta:
      distance:
        interleaving: 0.100
      sources:
        elimination: 0.080
        rounding: 0.020
      total: 0.100
persistence:
  PH1_zero: true                    # H_1 extinguished in the reported window
  Ext1_zero: true                   # Ext^1 vanishes (as reported in ext artifact)
  mu: 1                             # generic kernel dimension after truncation
  nu: 0                             # generic cokernel dimension after truncation
  phi_iso_tail:
    passed: false
    i: 1
    tau: 0.50
spectral:
  aux_bars_remaining: 0             # count of spectral aux-bars remaining (diagnostic)
thresholds:
  alpha: 1.0                        #  >0 (default  =1)
  tol:
    distance:
      interleaving: 1e-6            # present iff pipeline.metric == ”interleaving”
      # bottleneck: 1e-6            # present iff pipeline.metric == ”bottleneck”
    eig: 1e-8
    witness: 1e-9
budget:
  sum_delta:
    distance:
      interleaving: 0.150           # sum over operations[*].delta.distance.interleaving
  safety_margin: 0.850              # 1 - sum_delta / thresholds.tol.distance.interleaving
  gap_tau: 0.025                    # edge gap to  -window boundary (used by Restart)
  rationale: ”All deltas accounted for; slack remains >0”
Lambda_len:
  degree: 1
  tau: 0.50
```

```
  audit: ”hash:2f4c...d1”          # or an explicit eigenvalue list (see G.10)
gate:
  accept: true
  reason: ”Coverage ok; safety margin positive; all assertions satisfied”
serialization:
  float_dtype: ”ieee754-f64-le”
  json_sort_keys: true
  hdf5_canonical:
    compression: { algo: ”gzip”, level: 4 }
    shuffle: false
    fletcher32: false
    track_times: false
    fillvalue: 0.0
    string_encoding: ”utf8-fixed”      # fixed-length UTF-8 strings for bitwise reproducibility
    chunk_shapes:
      bars: { i: 4096, birth: 4096, death: 4096, death_is_inf: 4096, mult: 4096 }
      spec_eigs: { eig: 4096 }
      spec_Ntheta: { theta: 512, left: 512, right: 512 }
      phi_idx: { i: 256, tau: 256, iso: 256, mu: 256, nu: 256 }
cache:
  enabled: true
  dir: ”.cache/run_7f5c1b1”
timing:
  wallclock_s: 123.4
  cpu_s: 456.7
  stages:
    barcode: 12.3
    collapse: 4.5
    spec: 80.0
    ext_test: 2.0
status:
  success: true
  errors: []
outputs:
  bars: ”out/bars_7f5c1b1.json”
  spec: ”out/spec_7f5c1b1.json”
  ext:  ”out/ext_7f5c1b1.json”
  phi:  ”out/phi_7f5c1b1.h5”
```


G.3. bars (barcodes) schema

Semantics. A constructible barcode is a multiset of half-open intervals $I = [b, d)$ with degree i. Deaths may be $+\infty$.

JSON layout (with infinity convention, units, and cross-link).


```
{
  ”meta”: {
    ”schema_version”: ”2025-03-15”,
```

```
    "field": "k",
    "filtration_units": "dimensionless",
    "endpoint_convention": "[b,d) (see Chapter~2)",
    "infinity": { "json": "inf" },
    "float_dtype": "ieee754-f64-le",
    "string_encoding": "utf8-fixed",
    "links": {
      "run_id": "2025-03-15T09:12:07Z-7f5c1b1",
      "run_yaml_hash": "sha256:...run"
    }
  },
  "bars": [
    { "i": 0, "birth": 0.0, "death": 0.3,   "mult": 1 },
    { "i": 1, "birth": 0.2, "death": "inf", "mult": 1 }
  ],
  "hash": "sha256:...bars"
}
```

JSON Schema snippet for bars.bars[*].death**.**

```
"bars":{"type":"array","items":{
  "type":"object",
  "properties":{
    "i":{"type":"integer"},
    "birth":{"type":"number"},
    "death":{"oneOf":[{"type":"number"},{"type":"string","enum":["inf"]}]},
    "mult":{"type":"integer","minimum":1}
  },
  "required":["i","birth","death","mult"]
}}
```

HDF5 layout (split representation; no numeric sentinel).

- Datasets: /bars/i (int32), /bars/birth (float64), /bars/death (float64), /bars/death_is_inf (bool), /bars/mult (int32).
- Attributes: /bars.attrs[field="k"], filtration_units, schema_version, float_dtype, death_encoding="split_scalar_bool", string_encoding="utf8-fixed".

Invariants. Local finiteness on bounded windows (Appendix A). Optional sorting by (i, b, d); multiplicities aggregate identical intervals. Report-window compliance is checked via run.yaml.coverage_check.domain_window_covers_bars.

G.4. spec (spectral indicators) schema

Semantics. Spectral features: clipped sums, counts above/below thresholds with left/right limits, and deletion-type monotonicity diagnostics. Matrices are identified by content hashes.
JSON layout (ascending storage; $N_{\theta\pm0}$; solver params; coverage; cross-link).

```
{
  "meta": {
```

```
    ”schema_version”: ”2025-03-15”,
    ”eigen_units”: ”dimensionless”,
    ”order”: ”ascending”,
    ”sorted”: true,
    ”Ntheta_convention”: { ”left”: ”N_{ -0}”, ”right”: ”N_{ +0}” },
    ”window”: { ”range”: [0.0, 2.0], ”semantics”: ”closed” },
    ”norm”: ”fro”,                    // ”fro”   ‖ · ‖_fro, ”op”   ‖ · ‖_op
    ”clip_tau”: 1.0,
    ”tol_eig”: 1e-8,
    ”loewner_assumption”: ”A' A”,
    ”aux_bars_remaining”: 0,
    ”coverage_check”: { ”thetas_in_window”: true },
    ”string_encoding”: ”utf8-fixed”,
    ”eig_solver”: {
      ”method”: ”lanczos”, ”k”: 128, ”maxiter”: 1000,
      ”tol”: 1e-12, ”reorthogonalize”: true, ”rng_seed”: 1337
    },
    ”links”: {
      ”run_id”: ”2025-03-15T09:12:07Z-7f5c1b1”,
      ”run_yaml_hash”: ”sha256:...run”
    }
  },
  ”operators”: [
    {
      ”id”: ”sha256:...A”,
      ”kind”: ”laplacian_dirichlet”,
      ”n”: 500,
      ”spectrum”: { ”eigs”: [0.10, 0.12, 0.45, ...] },  // ascending
      ”clip”:     { ”tau”: 1.0, ”sum”: 37.219, ”deficit”: 12.004 }
    },
    {
      ”id”: ”sha256:...B”,
      ”kind”: ”principal_submatrix”,
      ”parent”: ”sha256:...A”,
      ”N_theta”: [
        { ”theta”: 0.20, ”left”: 17, ”right”: 16 },
        { ”theta”: 0.50, ”left”: 10, ”right”: 10 }
      ],
      ”monotonicity”: { ”type”: ”deletion”, ”passed”: true }
    }
  ],
  ”hash”: ”sha256:...spec”
}
```


HDF5 layout.

- /spec/ops/{id}/eig (float64, ascending), /spec/ops/{id}/clip/sum (float64), /spec/ops/{id}/clip/deficit (float64), /spec/ops/{id}/Ntheta/theta, /spec/ops/{id}/Ntheta/left, /spec/ops/{id}/Ntheta/right (parallel datasets).

- Attributes: kind, parent, norm ∈ {"fro", "op"} (Appendix E), order="ascending", sorted (bool), eigen_units, tol_eig, schema_version, loewner_assumption ∈ {"A' ⪯ A", "A' ⪰ A", "none"}, aux_bars_remaining (int32), coverage_thetas_in_window (bool), string_encoding="utf8-fixed", and a nested eig_solver group mirroring the JSON.

G.5. ext (Ext-test) schema

Semantics. Outcome of the one-way bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ for $C_\tau F$, with amplitude checks for $\mathcal{R}$ and recorded assumptions.

JSON layout (with assumptions and cross-links).


```
{
  "meta": {
    "schema_version": "2025-03-15",
    "field":"k", "alpha": 1.0,
    "assumptions": {
      "field_is_k": true,
      "constructible_verified": true,
      "t_exact_and_amp_le_1": true
    },
    "string_encoding": "utf8-fixed",
    "links": {
      "run_id": "2025-03-15T09:12:07Z-7f5c1b1",
      "run_yaml_hash": "sha256:...run"
    }
  },
  "tau": 0.50,
  "amplitude": { "ok": true, "range": [-1, 0] },
  "Hminus1": { "dim": 0, "witness_norm": 0.0 },
  "Ext1":    { "dim": 0, "passed": true, "tol": 1e-9, "slack": 0.0 },
  "links":   { "bars": "sha256:...bars", "phi": "sha256:...phi" },
  "hash": "sha256:...ext"
}
```


HDF5 layout.

- Scalars: /ext/tau (float64), /ext/Hminus1/dim (int32), /ext/Ext1/dim (int32), /ext/Ext1/passed (bool), /ext/Ext1/tol (float64), /ext/Ext1/slack (float64).
- Attributes: field="k", alpha=float64, schema_version, assumptions/* (bools as above), string_encoding="utf8-fixed".

G.6. phi (tower comparison) schema

Semantics. Encodes $\phi_{i,\tau}$ for towers, together with (μ, u) as generic-fiber dimensions after truncation, and structural flags for the sufficiency hypotheses (Appendix D, §D.3).

JSON layout (with τ sweep, edge kinds, witnesses, iso-tail, cross-link).


```
{
  "meta": {
```

```
    ”schema_version”: ”2025-03-15”,
    ”definition”: ”phi_{i, }: colim T_  P_i(F_n) → T_  P_i(F_∞)”,
    ”scope”: ”colim in [ ,Vect_k], return-to-constructible policy”,
    ”tau_sweep”: [0.25, 0.50, 1.00],
    ”edge_kinds”: [”inclusion”,”projection”,”quasi_iso”,
                  ”filtration_preserving_map”,”schur_complement”,”other”],
    ”string_encoding”: ”utf8-fixed”,
    ”links”: {
      ”run_id”: ”2025-03-15T09:12:07Z-7f5c1b1”,
      ”run_yaml_hash”: ”sha256:...run”
    }
  },
  ”indices”: [
    {
      ”i”: 1, ”tau”: 0.50,
      ”iso”: false,
      ”mu”: 1, ”nu”: 0,
      ”flags”: { ”S1_commutes”: false, ”S2_noAccum”: true, ”S3_Cauchy”: false },
      ”witness”: { ”ker_generic_dim”: 1, ”coker_generic_dim”: 0 },
      ”iso_tail”: { ”passed”: false }    // mirrors run.yaml.persistence.phi_iso_tail
    }
  ],
  ”tower”: {
    ”nodes”: [
      { ”n”: 0, ”id”: ”sha256:...F0” },
      { ”n”: 1, ”id”: ”sha256:...F1” }
    ],
    ”edges”: [
      { ”src”: 0, ”dst”: 1, ”kind”: ”inclusion” }
    ],
    ”limit”: { ”id”: ”sha256:...Finf” }
  },
  ”hash”: ”sha256:...phi”
}
```


HDF5 layout.

- /phi/idx/i (int32), /phi/idx/tau (float64), /phi/idx/iso (bool), /phi/idx/mu (int32), /phi/idx/nu (int32).
- /phi/idx/flags: S1_commutes, S2_noAccum, S3_Cauchy (bool).
- Optional witnesses: /phi/idx/ker_generic_dim, /phi/idx/coker_generic_dim.
- Optional tail: /phi/idx/iso_tail/passed (bool).
- Optional tower edges: /phi/tower/edges/src,dst (int32), /phi/tower/edges/kind (fixed-length UTF-8 string; attribute string_encoding="utf8-fixed").
- Attributes: schema_version, string_encoding=”utf8-fixed”, tau_sweep (float64 array).

G.7. Lambda_len (windowed length spectrum) schema

Semantics. The length spectrum operator $\Lambda_{\mathrm{len}}(M; [0, \tau])$ is diagonal on bar-basis with eigenvalues the clipped bar lengths on $[0, \tau]$. Its unordered eigenvalue multiset equals the clipped bar-length multiset (Appendix H). The Lambda_len audit records either the eigenvalue list (small instances) or a content hash.
JSON layout.

```
{
  "meta": {
    "schema_version": "2025-03-15",
    "definition": "Lambda_len(P_i(C_  F); [0, ])",
    "degree": 1,
    "tau": 0.50,
    "string_encoding": "utf8-fixed",
    "links": { "bars": "sha256:...bars", "phi": "sha256:...phi", "run_yaml_hash": "sha256:...run" }
  },
  "eigs": [0.24, 0.51, 0.78],      // optional explicit eigenvalues (small cases)
  "hash": "sha256:2f4c...d1"        // required content hash for audit/replay
}
```

HDF5 layout.

- /Lambda_len/meta attributes: schema_version, degree, tau, string_encoding="utf8-fixed", links as above.
- /Lambda_len/eigs (optional; float64 array).
- /Lambda_len/hash (fixed-length UTF-8 string).

G.8. Content hashing and canonical serialization

Each artifact carries a content hash sha256:... over its canonical serialization (JSON with sorted keys; HDF5 with *fixed* dataset/attribute creation order, chunk shapes, compression and filters). All floating datasets are little-endian IEEE-754 double (float64). Cross-file links (bars $\leftrightarrow$ phi $\leftrightarrow$ ext $\leftrightarrow$ spec $\leftrightarrow$ Lambda_len) use these hashes exclusively. *JSON numeric policy:* finite numbers only; positive infinity is encoded as the string "inf" where applicable (see bars.meta.infinity). HDF5 encodes $+\infty$ via the *split representation* /bars/death (float64) + /bars/death_is_inf (bool). *Strings:* all JSON strings and HDF5 string datasets/attributes are *fixed-length* UTF-8 (string_encoding="utf8-fixed") to ensure bitwise reproducibility. *HDF5 canonicalization:* for bit-wise reproducibility, set track_times=false, shuffle=false, fletcher32=false, fillvalue=0.0, compression to GZIP level 4, and use the chunk_shapes recorded in run.yaml; create datasets and attributes in the order shown in this appendix.

G.9. Numeric tolerances, δ-budgets, and audit trail

Every quantitative claim includes:

- **tolerance** (tol) declared in run.yaml;
- **slack** (slack) measured margin to the decision boundary;
- **norm** used for spectral bounds (fro or op), consistent with Appendix E;

- **metric** for persistence distances ($\mathrm{interleaving}$ or $\mathrm{bottleneck}$), with $\mathrm{thresholds.tol.distance}$ carrying the matching key exactly once;

- **budget** aggregation: $\mathrm{operations[*].delta}$ contributions summed into $\mathrm{budget.sum\_delta}$, with $\mathrm{budget.safety\_margin}$ and $\mathrm{budget.gap\_tau}$;

- **solver** details for spectral computations (Lanczos parameters, RNG seed);

- **windows and coverage**: $\mathrm{windows.*}$ declare scopes; $\mathrm{coverage\_check.*}$ record pass/fail;

- **threading** environment and BLAS/LAPACK/HDF5 versions;

- **gate** decision: $\mathrm{gate.accept}$ with $\mathrm{gate.reason}$.

G.10. Tests T14/T15 and reproducibility checklist

T14 (Overlap Gate gluing). On a windowed cover, verify: (i) post-collapse equality (up to budget) on overlaps ($\mathrm{overlap\_checks.local\_equiv{=}true}$), (ii) Čech–Ext^1 acyclicity on overlaps ($\mathrm{cech\_ext1\_ok{=}true}$), (iii) stability band detection ($\mu = u = 0$, $\mathrm{stability\_band\_ok{=}true}$), (iv) A/B soft-commuting with logged residuals, added to the δ-ledger, (v) global gate acceptance with additive budgets.

T15 (Length spectrum audit). Compute $\Lambda_{\text{len}}(\mathsf{T}_\tau \mathbf{P}_i(F); [0, \tau])$ and verify that its eigenvalue multiset equals the clipped bar-length multiset of $\mathsf{T}_\tau \mathbf{P}_i(F)$ (Appendix H). Log either the eigenvalue list or a content hash under $\mathrm{Lambda\_len}$.

Minimal reproducibility checklist.

1. Preserve $\mathrm{run.yaml}$ and all emitted $\mathrm{bars/spec/ext/phi/Lambda\_len}$ files (JSON or HDF5).

2. Confirm $\alpha > 0$ (default $\alpha = 1$) and field k are consistent across files.

3. Verify content hashes and all cross-links resolve; each artifact should carry $\mathrm{meta.links.run\_id}$ and $\mathrm{run\_yaml\_hash}$.

4. Check that $\mathrm{pipeline.metric}$ matches $\mathrm{thresholds.tol.distance}$ (exactly one of $\mathrm{interleaving/bottleneck}$), and that norms/tolerances are consistent.

5. Verify declared **windows** and that $\mathrm{coverage\_check.*}$ is true.

6. Verify eigenvalue *order* metadata: $\mathrm{spec.meta.order{=}"ascending"}$, $\mathrm{spec.meta.sorted}$=true, and HDF5 eigen arrays are non-decreasing.

7. For Dirichlet Laplacians, confirm $\lambda_{\min} > 0$.

8. Ensure each recorded θ used for N_θ lies within $\mathrm{spec.meta.window.range}$; recompute spectral indicators using recorded norm/tolerance/solver settings; check slack ≥ 0.

9. Re-evaluate $\phi_{i,\tau}$ under the scope policy; confirm (μ, u) match generic-fiber counts; check $\mathrm{iso\_tail.passed}$.

10. For Ext-tests, verify amplitude $[-1, 0]$, the assumptions flags in $\mathrm{ext.meta.assumptions}$, and $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$; check $\mathrm{persistence.Ext1\_zero}$.

11. Validate HDF5 canonicalization: chunk shapes, compression, filters, $\mathrm{string\_encoding}$="utf8-fixed", and $\mathrm{track\_times}$=false; death is represented via the split float/bool fields.

12. Inspect $\mathrm{operations}$ and budget: the sum of per-operation δ must match $\mathrm{budget.sum\_delta}$; compute safety margin and $\mathrm{gap\_tau}$; verify $\mathrm{gate.accept}$.

13. Record container details (image & digest) and timing; if unavailable, capture OS and library versions precisely.

Outcome. The versioned schemas above—now with (i) overlap checks for gluing, (ii) a windowed length spectrum audit, (iii) canonical spectral policy (order="ascending", norm="op"|"fro", spectral bounds), (iv) δ-ledger extensions with $\mathrm{gap\_tau}$, and (v) HDF5 canonicalization and fixed-length UTF-8 strings— are sufficient to regenerate all figures and claims in the main text from first principles, within the constructible regime and under the filtered-colimit policy, while making accept/reject criteria explicit and auditable. No further supplementation is required for operational deployment or third-party review.

Appendix H. Betti Integral and Finite τ-Events (Reinforced)

Standing conventions. We work over a field k. All persistence modules are constructible (locally finite on bounded windows); filtered colimits, when used, are taken under the scope policy of Appendix A, Remark .16. Endpoint conventions follow Appendix A, Remark .12; we use half-open bars $[b, d)$ (any consistent choice is immaterial below). Global conventions: Ext^1-tests are always against $k[0]$ (we write $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$, with C_τ understood up to f.q.i. on $\mathrm{Ho}(\mathsf{FiltCh}(k))$); the *energy exponent* satisfies $\alpha > 0$ (default $\alpha = 1$). Tilde references and type dashes (Type I–II / Type III / Type IV) are used uniformly.

H.1. Betti curves and the Betti integral

Let F be a filtered chain complex (or a filtered object realizing a persistence module) with degree-i persistence module $\mathbf{P}_i(F)$. Write its barcode as a locally finite multiset

$$\mathbf{P}_i(F) \;\cong\; \bigoplus_{I \in \mathcal{B}_i(F)} I^{\oplus m(I)}, \qquad I = [b, d) \;\text{ with } d \in \mathbb{R} \cup \{\infty\},\; m(I) \in \mathbb{Z}_{\geq 1}.$$

Define the Betti curve $\beta_i(t) := \dim_k H_i(F^t)$ and the *Betti integral* up to $\tau \geq 0$ by

$$\mathrm{PE}_i^{\leq \tau}(F) \;:=\; \int_0^\tau \beta_i(t)\, dt.$$

Under constructibility, β_i is right-continuous and piecewise constant, and on any bounded window only finitely many bars meet.

Theorem E.1 (Betti integral = clipped barcode mass)**.** For every $\tau \geq 0$,

$$\mathrm{PE}_i^{\leq \tau}(F) \;=\; \sum_{I \in \mathcal{B}_i(F)} m(I) \cdot \lambda(I \cap [0, \tau])\,,$$

where λ is Lebesgue measure and

$$\lambda([b, d) \cap [0, \tau]) = \max\{0, \min\{d, \tau\} - \max\{b, 0\}\} \quad (\text{with } \min\{\infty, \tau\} = \tau).$$

In particular, an infinite bar alive at 0 contributes its clipped length τ.

Proof. By local finiteness, for each bounded window $[0, \tau]$ only finitely many bars I intersect the window, so

$$f(t) \ := \sum_{I \in \mathcal{B}_i(F)} m(I)\, \mathbf{1}_I(t)$$

is a nonnegative measurable function on $[0, \tau]$ given by a finite sum. With the half-open convention $[b, d)$, for all t away from event times (births b and deaths d) we have $\beta_i(t) = f(t)$ and at event times β_i is right-continuous, hence equal to f almost everywhere. Thus, by Tonelli/Fubini on bounded windows (finite sum of nonnegative measurable functions),

$$\int_0^\tau \beta_i(t)\, dt \ = \int_0^\tau f(t)\, dt = \sum_I m(I) \int_0^\tau \mathbf{1}_I(t)\, dt = \sum_I m(I)\, \lambda(I \cap [0, \tau])\,.$$

□

Corollary E.2 (Monotonicity, (a.e.) derivative, piecewise linearity)**.** The map $\tau \mapsto \mathrm{PE}_i^{\leq \tau}(F)$ is nondecreasing, continuous, and piecewise linear on every bounded interval. Its derivative satisfies

$$\frac{d}{d\tau}\, \mathrm{PE}_i^{\leq \tau}(F) = \beta_i(\tau) \quad \text{for a.e. } \tau,$$

and at event points (births/deaths) the right derivative equals $\beta_i(\tau)$ while the left derivative equals $\beta_i(\tau-)$. All breakpoints on $[0, \tau_0]$ lie in $\{0, \tau_0\} \cup \{b \in [0, \tau_0]\} \cup \{d \in [0, \tau_0]\}$.

Remark E.3 (Endpoint and baseline conventions)**.** Changing open/closed endpoint conventions modifies β_i only on a set of measure zero; the integral and the breakpoint set remain unchanged. The baseline 0 is a reference; negative births are allowed and handled by intersecting I with $[0, \tau]$.

Remark E.4 (Energy exponent and α-Betti integral)**.** For $\alpha > 0$, define the α-Betti integral up to $\tau \geq 0$ by

$$\mathrm{PE}_{i,\alpha}^{\leq \tau}(F) \ := \int_0^\tau \big(\beta_i(t)\big)^\alpha\, dt.$$

On each component of $[0, \tau]$ between consecutive event times, β_i is constant, hence $\mathrm{PE}_{i,\alpha}^{\leq \tau}$ is still continuous, nondecreasing, and piecewise linear in τ (with slope $\big(\beta_i(\tau)\big)^\alpha$ on right-open pieces). The case $\alpha = 1$ recovers Theorem E.1. For $\alpha eq 1$ there is no direct "clipped mass" formula, but all algorithmic and verification statements below remain valid verbatim (replace β_i by β_i^α when computing segment slopes).

H.2. Finite τ-events and finite checking sets

Fix $\tau_0 > 0$ and define the finite τ-event set

$$\mathsf{Ev}_i(F; \tau_0) \ := \ \{0, \tau_0\} \ \cup \ \big(\{b \mid [b, d) \in \mathcal{B}_i(F)\} \cap [0, \tau_0]\big) \ \cup \ \big(\{d \mid [b, d) \in \mathcal{B}_i(F)\} \cap [0, \tau_0]\big).$$

By constructibility, $\mathsf{Ev}_i(F; \tau_0)$ is finite.

Proposition E.5 (Finite checking set)**.** Let $g : [0, \tau_0] \to \mathbb{R}$ be continuous and affine on each connected component of $[0, \tau_0] \setminus \mathsf{Ev}_i(F; \tau_0)$ (e.g. a piecewise linear benchmark with breakpoints in Ev_i). Then, for either inequality direction,

$$\mathrm{PE}_i^{\leq \tau}(F) \ \geq \ g(\tau) \quad (\text{resp. } \leq) \quad \text{for all } \tau \in [0, \tau_0]$$

holds if and only if it holds for all $\tau \in \mathsf{Ev}_i(F; \tau_0)$.

Proof. Between consecutive event times, both $\mathrm{PE}_i^{\le\tau}(F)$ and $g(\tau)$ are affine in τ (Corollary E.2). Hence their difference $h(\tau) := \mathrm{PE}_i^{\le\tau}(F) - g(\tau)$ is affine on each closed component $J = [u, v] \subset [0, \tau_0] \setminus \mathsf{Ev}_i(F; \tau_0)$. An affine function on a compact interval attains its extremum at an endpoint, so $h \ge 0$ (resp. $h \le 0$) on J iff $h(u) \ge 0$ and $h(v) \ge 0$ (resp. ≤ 0). Taking the union over all such J plus the singleton event points yields the claim. $\square$

Remark E.6 (Algorithmic evaluation)**.** Algorithm for evaluating $\mathrm{PE}_{i,\alpha}^{\le\tau}$ on $[0, \tau_0]$:

1. Collect all births and deaths intersecting $[0, \tau_0]$; sort to form $\mathsf{Ev}_i(F; \tau_0) = \{0 = t_0 < t_1 < \cdots < t_M = \tau_0\}$.
2. For each segment $[t_j, t_{j+1})$, compute $c_j := \beta_i(t)$ for any $t \in [t_j, t_{j+1})$ (e.g. by counting bars covering that segment), and set the segment slope $s_j := c_j^{\alpha}$.
3. Accumulate linearly: for $\tau \in [t_j, t_{j+1})$,

$$\mathrm{PE}_{i,\alpha}^{\le\tau} \;=\; \sum_{\ell<j} s_\ell\,(t_{\ell+1} - t_\ell) \;+\; s_j\,(\tau - t_j).$$

Complexity: forming/sorting Ev_i is $O(M \log M)$ for M events; the linear pass is $O(M)$ time and $O(M)$ space. This supports streaming if events are pre-sorted. For full reproducibility, record window definition and event counts in the run manifest (Appendix G).

H.3. Consequences for shifts, truncations, and window variation

(i) Equivariance under shifts. For $\varepsilon \in \mathbb{R}$, let the shift S^ε act by $(S^\varepsilon F)^t := F^{t+\varepsilon}$. Then $\beta_i^{S^\varepsilon}(t) = \beta_i(t+\varepsilon)$ and, for $\sigma \ge 0$,

$$\mathrm{PE}_{i,\alpha}^{\le\sigma}(S^\varepsilon F) = \int_0^\sigma \big(\beta_i(t+\varepsilon)\big)^\alpha\,dt = \int_\varepsilon^{\sigma+\varepsilon} \big(\beta_i(u)\big)^\alpha\,du = \mathrm{PE}_{i,\alpha}^{\le\sigma+\varepsilon}(F) - \mathrm{PE}_{i,\alpha}^{\le\varepsilon}(F).$$

Thus $\mathrm{PE}_{i,\alpha}^{\le\sigma}$ is equivariant under simultaneous shifts of both the module and the integration window; it is not invariant if only the module is shifted.

(ii) Truncation monotonicity (deletion-type). Let $\mathbf{T}_{\tau'}$ denote the bar-deletion truncation at scale $\tau' > 0$ on $\mathsf{Pers}_k^{\mathrm{ft}}$ (Appendix A). Then, for every $\sigma > 0$ and every $\alpha > 0$,

$$\mathrm{PE}_{i,\alpha}^{\le\sigma}\big(\mathbf{T}_{\tau'}(\mathbf{P}_i(F))\big) \;\le\; \mathrm{PE}_{i,\alpha}^{\le\sigma}\big(\mathbf{P}_i(F)\big),$$

because $\mathbf{T}_{\tau'}$ deletes precisely the finite bars of length $\le \tau'$ while leaving the others unchanged, and β_i decreases pointwise under deletion. Moreover, $\mathbf{T}_{\tau'}$ is 1-Lipschitz in the interleaving metric (Appendix A, Proposition .20).

(iii) Lipschitz in the window parameter. For $0 \le s \le \tau$ and $\alpha > 0$,

$$\big|\mathrm{PE}_{i,\alpha}^{\le\tau}(F) - \mathrm{PE}_{i,\alpha}^{\le s}(F)\big| = \int_s^\tau \big(\beta_i(t)\big)^\alpha\,dt \;\le\; (\tau - s)\cdot \sup_{t\in[s,\tau]} \big(\beta_i(t)\big)^\alpha,$$

and the supremum is finite on bounded windows by constructibility. Use the same window policy (MECE, right-open) recorded in Appendix G for comparability across runs.

H.4. Stability under interleavings and perturbations

We record a practical bound quantifying stability of PE under barcode perturbations, sufficient for algorithmic and statistical use on bounded windows.

Proposition E.7 (Perturbation bound on bounded windows)**.** Let $\tau_0 > 0$. Suppose two barcodes $\mathcal{B}_i(F)$ and $\mathcal{B}_i(G)$ are δ-matched in the standard bottleneck correspondence: each matched pair $[b,d) \leftrightarrow [b',d')$ satisfies $|b-b'| \le \delta$, $|d-d'| \le \delta$ (with $d=\infty$ allowed), and unmatched bars (if any) have length $\le 2\delta$. Then, for all $\tau \in [0,\tau_0]$ and $\alpha > 0$,

$$\left|\mathrm{PE}^{\le\tau}_{i,\alpha}(F) - \mathrm{PE}^{\le\tau}_{i,\alpha}(G)\right| \;\le\; C_{i,\alpha}(\tau_0,\delta)\cdot\delta,$$

where one can take

$$C_{i,\alpha}(\tau_0,\delta) \;:=\; 2\,N_i([-\delta,\tau_0+\delta])\cdot\max\Bigl\{1,\ \sup_{t\in[-\delta,\tau_0+\delta]}\bigl(\beta_i^F(t)\bigr)^{\alpha-1},\ \sup_{t\in[-\delta,\tau_0+\delta]}\bigl(\beta_i^G(t)\bigr)^{\alpha-1}\Bigr\},$$

and $N_i([a,b])$ denotes the number of bars (counting multiplicity) meeting $[a,b]$ in either barcode. In particular, for $\alpha = 1$,

$$\left|\mathrm{PE}^{\le\tau}_{i}(F) - \mathrm{PE}^{\le\tau}_{i}(G)\right| \;\le\; 2\,\delta\cdot N_i([-\delta,\tau_0+\delta]),$$

which is finite by local finiteness.

Proof sketch. On $[0,\tau]$, $\mathrm{PE}^{\le\tau}_i$ equals the sum of clipped lengths of the bars (Theorem E.1). A δ-perturbation of endpoints changes the clipped length of any matched bar by at most 2δ. Unmatched bars have length $\le 2\delta$ and hence contribute at most 2δ to the clipped length. Summing over bars meeting $[-\delta,\tau_0+\delta]$ gives the $\alpha = 1$ bound. For $\alpha eq 1$, use that on each segment where β_i is constant $= c$, the slope changes by at most $|c^\alpha - c'^\alpha| \le \alpha\max\{c,c'\}^{\alpha-1}|c-c'|$, and $|c-c'|$ is controlled by the number of bar crossings, which is bounded by $N_i([-\delta,\tau_0+\delta])$ on bounded windows. Combining these estimates yields the stated bound. ◻

Remark E.8 (Interpretation)**.** Proposition E.7 gives windowed stability of PE: small bottleneck perturbations of the barcode produce $O(\delta)$ changes in PE, with a constant depending only on the local combinatorics (finite on bounded windows). This is adequate for numerical robustness and cross-run comparisons.

H.5. Implementation notes and numerics

For large barcodes, the following practices improve reproducibility and numerical stability:

1. Event extraction: derive $\mathsf{Ev}_i(F;\tau_0)$ directly from the barcode; for streamed persistence, emit births and deaths as they occur and maintain a running count c_j.
2. Accumulation: use compensated summation (e.g. Kahan) when aggregating segment areas $s_j\,(t_{j+1}-t_j)$ to reduce floating-point error.
3. Types: store event times as 64-bit floats; store counts c_j as 64-bit integers; compute slopes as double-precision floats.
4. Idempotence: with the half-open convention $[b,d)$, repeated evaluation on the same event sequence is bitwise deterministic (assuming fixed sort and tie-break rules).
5. Window policy: record in the manifest (Appendix G) the baseline, window $[0,\tau_0]$, endpoint convention (right-open), and whether negative births are present.

H.6. Testing and validation

We recommend the following minimal test battery.

1. Synthetic bars: verify Theorem E.1 on barcodes with a mix of finite/infinite bars and negative births. Cross-check numerical integration of β_i against clipped-length summation.

2. Endpoint consistency: switch between $[b, d)$ and $(b, d]$ conventions in a controlled harness and verify identical $\mathrm{PE}_{i,\alpha}^{\leq\tau}$ values and identical breakpoint sets.

3. Shift equivariance: for random ε, check $\mathrm{PE}_{i,\alpha}^{\leq\sigma}(S^{\varepsilon}F) = \mathrm{PE}_{i,\alpha}^{\leq\sigma+\varepsilon}(F) - \mathrm{PE}_{i,\alpha}^{\leq\varepsilon}(F)$ up to machine precision.

4. Truncation monotonicity: apply $\mathbf{T}_{\tau'}$ and verify $\mathrm{PE}_{i,\alpha}^{\leq\sigma}$ decreases pointwise in σ.

5. Stability: simulate δ-perturbations of bar endpoints and confirm the bounds in Proposition E.7.

H.7. Variants and generalizations

- Weighted windows: for a nonnegative integrable weight $w \in L^1([0, \tau_0])$, define

$$\mathrm{PE}_{i,\alpha}^{w}(F) := \int_0^{\tau_0} w(t)\,\big(\beta_i(t)\big)^{\alpha}\,dt = \sum_I m(I) \int_{I\cap[0,\tau_0]} w(t)\,\big(\beta_i(t)\big)^{\alpha-1}\,dt,$$

  where the second identity holds by the same Tonelli argument as Theorem E.1. If w is piecewise constant with breaks in $\mathsf{Ev}_i(F; \tau_0)$, finite checking remains valid.

- Alternate baselines: replacing $[0, \tau]$ by $[a, b]$ yields $\mathrm{PE}_i^{[a,b]}(F) = \sum_I m(I)\,\lambda(I \cap [a, b])$, with the same properties and algorithms after shifting time by $-a$.

- Discrete filtrations: if t ranges over a discrete grid $\{t_0 < \cdots < t_M\}$, replace integrals with Riemann sums; all statements adapt with λ replaced by counting measure on grid cells.

Summary. The Betti integral equals the clipped barcode mass (Theorem E.1). Hence $\mathrm{PE}_{i,\alpha}^{\leq\tau}$ is continuous, nondecreasing, and piecewise linear, with breakpoints among births/deaths (Corollary E.2, Remark E.4). Any affine-on-components constraint reduces to a finite event set (Proposition E.5). Under shifts, $\mathrm{PE}_{i,\alpha}^{\leq\sigma}$ obeys an explicit equivariance formula; under deletion-type truncations it is nonincreasing and enjoys 1-Lipschitz stability in interleaving distance (Appendix A, Proposition .20). Bounded-window perturbation bounds (Proposition E.7) provide practical stability. Implementation, numerical, and testing guidelines (Subsections H.5–H.6) ensure reproducibility in line with the global scope policy (Appendix A) and run manifests (Appendix G).

Appendix I. ε-Survival Lemma and Grid-to-Continuum [Proof/Spec]

Standing conventions. We work over a field k. All persistence modules are constructible (locally finite on bounded windows). Any use of filtered colimits follows the scope policy of Appendix A, Remark .16. Global conventions: Ext^1-tests are always against $k[0]$ (so we write $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), k) = 0$); the energy exponent satisfies $\alpha > 0$ (default $\alpha = 1$); tilde references (Appendix˜) and type dashes (Type I–II / Type III / Type IV) are used uniformly. We write d_{int} for the interleaving metric (coinciding with the bottleneck distance in the constructible 1D case).

Remark F.1 (Endpoint conventions are immaterial)**.** We take intervals half-open $[b, d)$ with $d \in \mathbb{R} \cup \{\infty\}$. Any consistent open/closed endpoint choice yields the same clipped lengths and event sets; all statements below are invariant under this choice. Use the same right-open convention as Appendix G for run logs.

Remark F.2 (Relation to cropping in Chapter 2)**.** The window clipping functor $\mathbf{W}_{\leq\tau}$ introduced below implements the restriction to the time window $[0, \tau]$ followed by extension by zero. This is equivalent in purpose to the cropping operator $\mathbf{W}_W$ from Chapter 2 (see also Chapter 2, Definition (Crop)). We keep the present notation to emphasize the functorial construction via restriction and 0-extension.

I.1. Window clipping, nonexpansivity, and ε-survival

Let $\mathsf{Pers}_k^{\mathrm{ft}}$ be the category of constructible persistence modules on $(\mathbb{R}, \leq)$. For $\tau \geq 0$, let

$$i_{\leq\tau} : \ ([0, \tau], \leq) \hookrightarrow (\mathbb{R}, \leq)$$

be the fully faithful inclusion of posets, and write $i_{\leq\tau}^*$ for restriction along $i_{\leq\tau}$. We identify modules on $[0, \tau]$ with modules on $\mathbb{R}$ supported in $[0, \tau]$ via *extension by zero* along $i_{\leq\tau}$; we denote this left Kan extension by $(i_{\leq\tau})_!^0$ ("0-extension"). Define the *window clip* endofunctor

$$\mathbf{W}_{\leq\tau} \ := \ (i_{\leq\tau})_!^0 \circ i_{\leq\tau}^* : \ \mathsf{Pers}_k^{\mathrm{ft}} \longrightarrow \mathsf{Pers}_k^{\mathrm{ft}}.$$

In barcode terms, this coincides with interval intersection $I \mapsto I \cap [0, \tau]$ (discarding empties). For a bar $I = [b, d)$ its clipped length is

$$\ell_{[0,\tau]}(I) := \lambda(I \cap [0, \tau]) = \max\{0, \min\{d, \tau\} - \max\{b, 0\}\},$$

with $\min\{\infty, \tau\} = \tau$. Since only finitely many bars intersect any bounded window, $\mathbf{W}_{\leq\tau}$ preserves constructibility.

Lemma F.3 (Clipping is 1-Lipschitz)**.** For all $M, N \in \mathsf{Pers}_k^{\mathrm{ft}}$ and $\tau \geq 0$,

$$d_{\mathrm{int}}\big(\mathbf{W}_{\leq\tau}M, \ \mathbf{W}_{\leq\tau}N\big) \ \leq \ d_{\mathrm{int}}(M, N).$$

Proof. Let $d_{\mathrm{int}}(M, N) \leq \varepsilon$ be witnessed by an ε-interleaving (φ, ψ) with $\varphi : M \to S^\varepsilon N$ and $\psi : N \to S^\varepsilon M$ satisfying the standard triangle identities. Restricting along $i_{\leq\tau}$ yields an ε-interleaving $(i_{\leq\tau}^*\varphi, i_{\leq\tau}^*\psi)$ between $i_{\leq\tau}^*M$ and $i_{\leq\tau}^*N$, because shifts commute with restriction and the naturality squares and triangle identities restrict pointwise on $[0, \tau]$. Composing with 0-extension $(i_{\leq\tau})_!^0$ produces an ε-interleaving between $\mathbf{W}_{\leq\tau}M$ and $\mathbf{W}_{\leq\tau}N$. Thus $d_{\mathrm{int}}(\mathbf{W}_{\leq\tau}M, \mathbf{W}_{\leq\tau}N) \leq \varepsilon$. Taking the infimum over ε gives the claim. Equivalently in barcode terms (constructible case), the bottleneck matching cost does not increase if one first clips both barcodes to $[0, \tau]$ and then matches, or matches first and clips afterwards. □

Remark F.4 (Shift commutation with clipping)**.** For every $\varepsilon \geq 0$ there is a canonical isomorphism

$$S^\varepsilon \circ \mathbf{W}_{\leq\tau} \ \cong \ \mathbf{W}_{\leq\tau} \circ S^\varepsilon,$$

since restriction along $i_{\leq\tau}$ and 0-extension commute with the time-shift endofunctors. In particular, any ε-interleaving data transport through $\mathbf{W}_{\leq\tau}$ via these identifications.

Lemma F.5 (ε-survival lemma)**.** Let $M, N \in \mathsf{Pers}_k^{\mathrm{ft}}$ with $d_{\mathrm{int}}(M, N) \leq \varepsilon$. Fix $\tau_0 > 0$.

1. **Quantitative form.** For any bar-matching that witnesses $d_{\mathrm{int}}(M, N) \leq \varepsilon$ (not necessarily unique), if a bar I of M is matched to a bar J of N, then

$$\ell_{[0,\tau_0]}(J) \ \geq \ \max\big\{\,\ell_{[0,\tau_0]}(I) - 2\varepsilon, \ 0\,\big\}.$$

2. **Nonvanishing form.** If $\ell_{[0,\tau_0]}(I) > 2\varepsilon$ for some bar I in M, then its matched partner J satisfies $\ell_{[0,\tau_0]}(J) > 0$. Consequently $\mathbf{W}_{\le \tau_0} Neq0$.

3. **Multiplicity lower bound.** If at least r bars of M have clipped length $> 2\varepsilon$, then at least r bars of N remain nonzero after clipping to $[0, \tau_0]$ (counted with multiplicity).

Proof. In the constructible 1D case the interleaving distance equals the bottleneck distance. Under an ε-matching, endpoints of matched bars move by at most ε. Hence the intersection length with $[0, \tau_0]$ can shrink by at most ε at each end, giving (1). Points matched to the diagonal have zero clipped length and thus cannot be matched to bars with clipped length $> 2\varepsilon$; (2) and (3) follow immediately. □

Remark F.6 (Notation)**.** We reserve $\mathbf{T}_\tau$ for the Serre reflection that removes bars of length $\le \tau$ (Appendix A, Theorem .14, and Proposition .20 for its 1-Lipschitz property). Window clipping is denoted $\mathbf{W}_{\le\tau}$. The present statements concern $\mathbf{W}_{\le\tau}$. Appendix H relates clipping to the Betti integral.

I.2. Grid-to-continuum transfer

Let F be a filtered object with degree-i persistence $\mathbf{P}_i(F)$. Suppose F_h is a discretization at mesh h with $\mathbf{P}_i(F_h)$ and a certified bound

$$d_{\text{int}}\big(\mathbf{P}_i(F_h),\ \mathbf{P}_i(F)\big) \ \le\ \varepsilon(h).$$

Theorem F.7 (Grid-to-continuum survival)**.** Fix $\tau_0 > 0$ and $r \in \mathbb{Z}_{\ge 1}$. If $\mathbf{W}_{\le\tau_0}\mathbf{P}_i(F_h)$ contains at least r bars of clipped length $\ge 2\varepsilon(h) + \eta$ for some margin $\eta > 0$, then

$$\mathbf{W}_{\le\tau_0}\mathbf{P}_i(F) \quad \text{contains at least } r \text{ nonzero bars (with multiplicity), each of clipped length } \ \ge \eta.$$

In particular, nonvanishing on the grid with margin $2\varepsilon(h) + \eta$ implies nonvanishing in the continuum window $[0, \tau_0]$ (no false negatives). Record $\varepsilon(h)$, τ_0, and the event counts per window in the run manifest (Appendix G) for reproducibility.

Proof. Apply Lemma F.5 with $M = \mathbf{P}_i(F_h)$, $N = \mathbf{P}_i(F)$, and $\varepsilon = \varepsilon(h)$. Bars with clipped length $> 2\varepsilon(h)$ cannot be matched to the diagonal; at least r such bars survive with length $\ge \eta$, counted with multiplicity. □

I.3. Variants, sharpness, and compatibility

- *Sharpness.* The constant 2 is optimal: a bar with $\ell_{[0,\tau_0]} = 2\varepsilon$ can be shifted by ε at each endpoint so that its clip collapses to length 0.
- *Compatibility with functors.* By Lemma F.3 and Appendix A, Proposition .20, composition with 1-Lipschitz endofunctors (e.g. the reflections $\mathbf{T}_\tau$; mirror/transfer maps under their stated [Spec] hypotheses) preserves the survival inequalities.
- *After-collapse pipeline.* Operationally we compare on the B-side after applying $\mathbf{T}_\tau$ ("after-collapse"). Since $\mathbf{T}_\tau$ is 1-Lipschitz, applying $\mathbf{T}_\tau$ before or after $\mathbf{W}_{\le\tau_0}$ preserves the bounds of Lemma F.5 and Theorem F.7.
- *Towers.* In a tower $F_n \to F_\infty$, if $\varepsilon_n \to 0$ with $d_{\text{int}}(\mathbf{P}_i(F_n), \mathbf{P}_i(F_\infty)) \le \varepsilon_n$, then any uniform margin > 0 propagates from grid to continuum by Theorem F.7 (cf. Appendix D for tower diagnostics).

I.4. Implementation schema and reproducibility (run manifest)

The following fields and checks should be present in the run manifest to certify a grid-to-continuum survival claim at window $[0, \tau_0]$:

- **Metric/tolerance.** Record the metric used and the certified bound: pipeline.metric and thresholds.tol.distance (interleaving or bottleneck, made explicit).
- **Sampling/bandlimit.** Include sampling.dt, bandlimit, nyquist_check, eps_interleave_max (or eps_cont_bound) used for $\varepsilon(h)$.
- **Window coverage.** Set coverage_check.domain_window_covers_bars: true to assert that the declared window matches the computation window.
- **Operations log.** Add an operations entry for the "epsilon-continuation" step with the realized ε, and aggregate into delta.total.
- **Event counts.** Record the number of births/deaths and bar counts per window in the persistence band, after clipping.

A minimal schema example:


```
pipeline:
  metric: interleaving          # or 'bottleneck' with justification
  stages:
    - compute_persistence
    - clip_window               # W_{ }
    - epsilon_continuation      # certify  (h)
    - collapse_short_bars       # optional T_
thresholds:
  tol:
    distance: 0.0125            # certified  (h)
  safety_margin:
    bar_clip_len_min: 2*  +     # enforced margin on grid side
sampling:
  dt: 0.001
  bandlimit: 200.0              # Hz or domain-appropriate
  nyquist_check: true
  eps_interleave_max: 0.0125    # equals thresholds.tol.distance
window:
  tau0: 1.5
  coverage_check:
    domain_window_covers_bars: true
operations:
  - name: epsilon-continuation
    epsilon: 0.0125
    notes: "Certified via stability bound / a priori model"
delta:
  total: 0.0125
persistence:
  degree: 1
```

```
  events:
    window_[0,1.5]:
      births: 7
      deaths: 6
      bars_total: 9
      bars_clip_len_ge_2eps_plus_eta: 3
assertions:
  grid_to_continuum_survival:
    r: 3
    eta: 0.05
    satisfied: true
```

I.5. Minimal tests and regression suite

- **Sharpness test.** Single-bar synthetic: $I = [0, 2\varepsilon)$. Shift endpoints by ε to cause collapse of $\ell_{[0,\tau_0]}$ when $\tau_0 \geq 2\varepsilon$. Confirms optimality of the constant 2.
- **Endpoint policy.** Verify the half-open/right-open convention across the pipeline (clip lengths, event attribution, and logs) matches Appendix G.
- **Tower regression.** Sequence $h_k \to 0$ with $\varepsilon(h_k) \to 0$. Bars on the grid with a uniform margin persist in the continuum; add a TCI regression to check stability of counts and lengths after clipping.
- **Operator separation.** Unit tests that distinguish $\mathbf{W}_{\leq\tau}$ (window clipping) from $\mathbf{T}_\tau$ (short-bar removal), including order-of-application invariance of survival conclusions due to 1-Lipschitzness.

I.6. Formalization stubs (Lean/Coq)

The following stubs (cf. Appendix F, AK.Core/Tower) codify the key interfaces. We keep them minimal to avoid committing to a proof assistant dialect here.

```
-- Lean-style pseudocode

namespace Pers

structure PersModule := -- constructible persistence module over
  (rank_on_window : set   →  )
  -- ... standard fields/constructibility witnesses ...

def W_le (  :   0) (M : PersModule) : PersModule := -- window clipping
  -- restriction to [0,  ] and 0-extension
  sorry

def interleaving_dist (M N : PersModule) :   0∞ := sorry

theorem W_le_lipschitz (  :   0) (M N : PersModule) :
  interleaving_dist (W_le   M) (W_le   N)   interleaving_dist M N := sorry

structure Discretization :=
  (F_h : Type) -- placeholder for filtered object at mesh h
```

```
  (P_i : F_h → PersModule)
  (eps :   0) --  (h)
  -- certificate data ...

theorem eps_survival
  (  :   0) {M N : PersModule}
  (hMN : interleaving_dist M N    )
  ( 0 :   0) :
  -- if bar I in M has clipped length > 2  on [0,  0], then its match in N is nonzero
  sorry

theorem grid_to_continuum
  ( 0 :   0) (r :  ) (    :   0)
  {M N : PersModule}
  (hMN : interleaving_dist M N    )
  (grid_margin : -- at least r bars in W_le  0 M have length   2  +
    sorry) :
  -- conclude: at least r bars in W_le  0 N have length
  sorry

end Pers
```

Reference line for Chapter 4.5. *"By Appendix I (the ε-Survival Lemma, Lemma F.5), any grid-detected bar in the window $[0, \tau_0]$ with clipped length $> 2\varepsilon(h)$ persists in the continuum window; we therefore certify the claim at scale τ_0 from the discretization."*

Summary. Clipping is restriction to $[0, \tau]$ followed by extension by zero to $\mathbb{R}$, hence an endofunctor $\mathbf{W}_{\leq\tau}$ on $\mathsf{Pers}_k^{\mathrm{ft}}$; it preserves constructibility and is 1-Lipschitz (Lemma F.3). Under an ε-interleaving, clipped lengths degrade by at most 2ε; hence grid detections with margin $> 2\varepsilon$ transfer to continuum without false negatives (Lemma F.5, Theorem F.7). These are consistent with the global scope policy and the finite-event structure of Appendices A and H. For implementation, record the metric, certified $\varepsilon(h)$, window coverage, and event counts in the run manifest, and include minimal sharpness/tower tests and optional after-collapse comparisons via the 1-Lipschitz reflection $\mathbf{T}_\tau$.

Appendix J. Calculus of μ, u [Proof + Stability Bands + Window Pasting]

Standing conventions. We work over a field k. All persistence modules are constructible (locally finite on bounded windows). Filtered colimits are computed in the functor category $[\mathbb{R}, \mathsf{Vect}_k]$ under the scope policy of Appendix A, Remark .16, and then (when stated) returned to the constructible range. The reflection $\mathbf{T}_\tau \dashv \iota_\tau$ is exact and 1-Lipschitz (Appendix A, Theorem .14 and Proposition .20). Global conventions: Ext^1 is always against $k[0]$; the energy exponent $\alpha > 0$ (default $\alpha = 1$). All window policies follow the MECE (mutually exclusive–collectively exhaustive), right-open convention used throughout (Appendix G).

Setup and notation. Fix $i \in \mathbb{Z}$. Let $F = (F_n)_{n\in I}$ be a directed system ("tower") of filtered objects for which $\mathbf{P}_i(F_n) \in \mathsf{Pers}_k^{\mathrm{ft}}$, and let F_∞ be its colimit with cocone maps $F_n \to F_\infty$. For $\tau \geq 0$ consider the *comparison map* in the functor category

$$\phi_{i,\tau}(F) : \quad \varinjlim_n \mathbf{T}_\tau\big(\mathbf{P}_i(F_n)\big) \longrightarrow \mathbf{T}_\tau\big(\mathbf{P}_i(F_\infty)\big). \tag{G.1}$$

Define

$$\mu_{i,\tau}(F) \;:=\; \dim_k \ker \phi_{i,\tau}(F), u_{i,\tau}(F) \;:=\; \dim_k \operatorname{coker} \phi_{i,\tau}(F), \tag{G.2}$$

where $\dim_k$ denotes the *generic fiber* dimension in $\mathsf{Pers}_k^{\mathrm{ft}}$ (i.e. the stabilized right-tail dimension $\lim_{t\to+\infty}\dim_k(-)(t)$, equivalently the multiplicity of $I[0,\infty)$ after applying $\mathbf{T}_\tau$; cf. Appendix D, Remark A.1).[8] We also write the *totals*

$$\mu_{\mathrm{Collapse}}(F) \;:=\; \sum_i \mu_{i,\tau}(F), u_{\mathrm{Collapse}}(F) \;:=\; \sum_i u_{i,\tau}(F), \tag{G.3}$$

which are finite because the complexes are bounded in homological degrees (constructible range). All quantities depend on τ; no general monotonicity in τ is asserted.

J.1. Subadditivity under composition

We formalize "composition of towers" via morphisms of directed systems.

Definition G.1 (Morphisms of towers)**.** A morphism $u : F \to G$ consists of maps $u_n : F_n \to G_n$ commuting with the structure maps (in the index category) and inducing a canonical map on colimits $u_\infty : F_\infty \to G_\infty$. Given $u : F \to G$ and $v : G \to H$, the composite $v \circ u : F \to H$ is defined degreewise.

Lemma G.2 (Functoriality of comparison maps)**.** Under the scope policy, each morphism $u : F \to G$ induces a (canonical) morphism of comparison maps

$$\phi_{i,\tau}(u) :\; \phi_{i,\tau}(F) \;\Longrightarrow\; \phi_{i,\tau}(G),$$

natural in both i and τ (here $\Longrightarrow$ indicates a canonical comparison morphism, natural in i, τ). Moreover, for composable $u : F \to G$, $v : G \to H$,

$$\phi_{i,\tau}(v \circ u) \;=\; \phi_{i,\tau}(v) \;\circ\; \phi_{i,\tau}(u).$$

Proof sketch. Apply $\mathbf{P}_i$, then $\mathbf{T}_\tau$, to the maps on systems and pass to filtered colimits in $[\mathbb{R}, \mathsf{Vect}_k]$; exactness of $\mathbf{T}_\tau$ and functoriality of colimits yield naturality and compatibility with composition. □

Theorem G.3 (Subadditivity under composition)**.** Let $u : F \to G$ and $v : G \to H$ be morphisms of towers and fix $\tau \geq 0$. Then

$$\mu_{i,\tau}(v \circ u) \;\leq\; \mu_{i,\tau}(u) \;+\; \mu_{i,\tau}(v), u_{i,\tau}(v \circ u) \;\leq u_{i,\tau}(u) \;+ u_{i,\tau}(v).$$

Proof. By Lemma G.2 we reduce to linear estimates on the generic fiber. In $[\mathbb{R}, \mathsf{Vect}_k]$, kernels and cokernels are computed pointwise; applying $\mathbf{T}_\tau$ preserves exactness. In the constructible range the (right-tail) generic fiber dimension equals the stabilized pointwise dimension. Thus, evaluating the comparison maps at any sufficiently large parameter t (where stabilization holds) and writing $X_t \xrightarrow{f_t} Y_t \xrightarrow{g_t} Z_t$ for the induced linear maps, we have the elementary inequalities

$$\dim\ker(g_t \circ f_t) \;\leq\; \dim\ker f_t \;+\; \dim\ker g_t, \dim\operatorname{coker}(g_t \circ f_t) \;\leq\; \dim\operatorname{coker} f_t \;+\; \dim\ker g_t,$$

and taking these stabilized values yields the stated bounds for the generic fiber dimensions in (G.2).[9] □

[8] For a constructible module M, the generic fiber dimension equals $\dim_k M(t)$ for all sufficiently large t, which stabilizes. After applying $\mathbf{T}_\tau$, it coincides with the multiplicity of $I[0,\infty)$ summands. Kernels/cokernels are taken in $\mathsf{Pers}_k^{\mathrm{ft}}$, where they decompose into finite direct sums of interval modules.

[9] Finite-dimensional linear algebra: for $X \xrightarrow{f} Y \xrightarrow{g} Z$, one has $\ker(g \circ f) \subseteq \ker f \oplus f^{-1}(\ker g)$ and $\operatorname{coker}(g \circ f)$ controlled by $\operatorname{coker} f$ and $\ker g$, giving the stated subadditivity bounds on dimensions.

J.2. Additivity under finite direct sums

Proposition G.4 (Direct-sum additivity)**.** Let $F = F^{(1)} \oplus F^{(2)}$ be the levelwise direct sum of towers (same index category) and similarly for the colimit. Then for every $\tau \geq 0$,

$$\mu_{i,\tau}(F) \; = \; \mu_{i,\tau}\big(F^{(1)}\big) \; + \; \mu_{i,\tau}\big(F^{(2)}\big), u_{i,\tau}(F) \; = u_{i,\tau}\big(F^{(1)}\big) \; + u_{i,\tau}\big(F^{(2)}\big).$$

In particular, the totals satisfy

$$\mu_{\text{Collapse}}(F) \; = \; \mu_{\text{Collapse}}\big(F^{(1)}\big) \; + \; \mu_{\text{Collapse}}\big(F^{(2)}\big), u_{\text{Collapse}}(F) \; = u_{\text{Collapse}}\big(F^{(1)}\big) \; + u_{\text{Collapse}}\big(F^{(2)}\big).$$

Proof. $\mathbf{P}_i$ and $\mathbf{T}_\tau$ preserve finite direct sums; filtered colimits *commute with finite direct sums* in $[\mathbb{R}, \mathsf{Vect}_k]$. Hence $\phi_{i,\tau}(F)$ is block-diagonal with blocks $\phi_{i,\tau}\big(F^{(1)}\big)$ and $\phi_{i,\tau}\big(F^{(2)}\big)$, and kernels/cokernels (hence generic fiber dimensions) add; summing over i gives the identities for $\mu_{\text{Collapse}}, u_{\text{Collapse}}$. □

J.3. Cofinal invariance

Definition G.5 (Cofinal subindexing)**.** Let I be the directed index category for F. A full subcategory $J \subset I$ is cofinal if for every $i \in I$ there exists $j \in J$ with a morphism $i \to j$. The restricted tower $F|_J$ has the same colimit as F in $[\mathbb{R}, \mathsf{Vect}_k]$.

Theorem G.6 (Cofinal invariance)**.** Let $J \subset I$ be cofinal. Then for all $\tau \geq 0$,

$$\mu_{i,\tau}\big(F|_J\big) = \mu_{i,\tau}(F), u_{i,\tau}\big(F|_J\big) = u_{i,\tau}(F),$$

hence also $\mu_{\text{Collapse}}(F|_J) = \mu_{\text{Collapse}}(F)$ and $u_{\text{Collapse}}(F|_J) = u_{\text{Collapse}}(F)$. In particular, passing to any cofinal tail of a sequential tower leaves $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$ unchanged.

Proof. Cofinal restriction does not change colimits in $[\mathbb{R}, \mathsf{Vect}_k]$. Thus the source and target of $\phi_{i,\tau}$ (Eq. (G.1)) remain unchanged, as does $\phi_{i,\tau}$ itself; kernels/cokernels (hence $\mu_{i,\tau}, u_{i,\tau}$, and the totals) agree. □

J.4. Consequences and remarks

- *Triangle-type bounds.* Combining Theorem G.3 across a chain of morphisms yields $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$-control along multi-stage pipelines, useful when towers factor through preprocessing steps (e.g. mirror/transfer under the [Spec] hypotheses of Appendix L).
- *Finite decomposition.* Iterating Proposition G.4 shows that $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$ is additive across any finite direct-sum decomposition of towers, enabling modular bookkeeping.
- *Cofinal tails for convergence.* Theorem G.6 justifies replacing a tower by any convenient tail when verifying sufficient conditions for $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0, 0)$ (Appendix D, §D.3).
- *No general monotonicity in τ.* Absent additional hypotheses, $\tau \mapsto (\mu_{i,\tau}, u_{i,\tau})$ (hence $\tau \mapsto (\mu_{\text{Collapse}}, u_{\text{Collapse}})$) need not be monotone; only the sufficient conditions of Appendix D (e.g. commutation, no τ-accumulation, Cauchy+compatibility) guarantee $(\mu_{\text{Collapse}}, u_{\text{Collapse}}) = (0, 0)$ at fixed τ.
- *Invariance under filtered quasi-isomorphism.* As in Appendix D (and the formalization sketch of Appendix F), $(\mu_{\text{Collapse}}, u_{\text{Collapse}})$ is invariant under replacing each F_n up to f.q.i., since $\mathbf{P}_i$ and $\mathbf{T}_\tau$ respect such replacements and kernels/cokernels (hence generic fiber dimensions) are preserved under isomorphism in $\mathsf{Pers}_k^{\text{ft}}$.

J.5. τ-sweep and stability bands

We formalize the practical, windowed method for selecting scales τ at which tower effects vanish.

Definition G.7 (Stability band for a fixed window and degree)**.** Fix a window (MECE, right-open) and a degree i. A τ-*sweep* is a finite or countable increasing array $\{\tau_\ell\}_{\ell \in L} \subset (0, \infty)$. A contiguous subarray $\{\tau_a, \dots, \tau_b\}$ is a *stability band* if

$$\mu_{i,\tau_\ell}(F) = u_{i,\tau_\ell}(F) = 0 \quad \text{for all } \ell \in \{a, \dots, b\},$$

and the verdict persists under refinement of the sweep (inserting new τ-values) without creating nonzero $\mu_{i,\tau}$ or $u_{i,\tau}$ in the band.

Proposition G.8 (Robust detection under sweep refinement)**.** Assume one of the sufficient conditions of Appendix D, §D.3 holds (e.g. commutation (S1), no τ-accumulation (S2), or $\mathbf{T}_\tau$-Cauchy with compatible cocone (S3)). Then for each i there exists a neighborhood of any τ_0 where $\phi_{i,\tau}$ is an isomorphism, hence $(\mu_{i,\tau}, u_{i,\tau}) = (0, 0)$. A sufficiently fine τ-sweep detects such neighborhoods as stability bands and remains stable under refinement.

Remark G.9 (Aggregating across degrees)**.** Applications often monitor either a fixed finite degree set $I \subset \mathbb{Z}$ or the total $\mu_{\text{Collapse}}, u_{\text{Collapse}}$. A *joint* stability band is the intersection of degreewise bands; it is nonempty whenever each degree admits a band with nontrivial overlap. Report the monitored degree set in the run manifest (Appendix G).

J.6. Window pasting via Restart and Summability

We now integrate a *Restart* inequality and a *Summability* condition to paste windowed certificates into global ones. The δ-ledger (algorithmic/discretization/measurement) is as in Appendix G; Mirror/Transfer commutation defects are treated via Appendix L.

Definition G.10 (Per-window pipeline budget and safety margin)**.** Let $\{W_k = [u_k, u_{k+1})\}_k$ be a MECE partition (right-open). On window W_k, let $\tau_k > 0$ be the selected collapse threshold (possibly from a stability band) and define the *pipeline budget*

$$\Sigma\delta_k(i) \;:=\; \sum_{U \in W_k} \Big(\delta_U^{\text{alg}}(i, \tau_k) + \delta_U^{\text{disc}}(i, \tau_k) + \delta_U^{\text{meas}}(i, \tau_k)\Big),$$

with δ-components recorded per operation (Appendix G). The *safety margin* $\text{gap}_{\tau_k}(i) > 0$ is the configured slack for B-Gate$^+$ on W_k and degree i (Chapter 1).

Lemma G.11 (Restart inequality)**.** Assume that on W_k the B-Gate$^+$ passes with $\text{gap}_{\tau_k}(i) > \Sigma\delta_k(i)$ and that the transition to W_{k+1} is realized by a finite composition of *deletion-type* steps and ε-continuations (both measured after $\mathbf{T}_\tau$). Then there exists $\kappa \in (0, 1]$, depending only on the admissible step class and the τ-adaptation policy, such that

$$\text{gap}_{\tau_{k+1}}(i) \;\geq\; \kappa \,\big(\text{gap}_{\tau_k}(i) - \Sigma\delta_k(i)\big).$$

Proof sketch. Deletion-type steps are nonincreasing for monitored indicators after $\mathbf{T}_\tau$ (Appendix E); ε-continuations are 1-Lipschitz, so the drift is bounded by the declared ε. Aggregating drifts across the finite composition yields a multiplicative retention factor κ. □

Definition G.12 (Summability)**.** A run satisfies *Summability* (on degrees $i \in I$) if

$$\sum_k \Sigma\delta_k(i) \ < \ \infty \qquad (\forall\, i \in I).$$

A sufficient design pattern is geometric decay of thresholds (e.g. $\tau_k = \tau_0\rho^k$, $\rho \in (0,1)$), combined with bounded per-window operation counts, as recorded in Appendix G.

Theorem G.13 (Pasting windowed certificates)**.** Let $\{W_k\}_k$ be MECE. Suppose that on each W_k the B-Gate$^+$ passes with $\mathrm{gap}_{\tau_k}(i) > \Sigma\delta_k(i)$ (for all $i \in I$), that the Restart inequality (Lemma G.11) holds at every transition, and that Summability (Definition G.12) is satisfied. Then the concatenation of per-window certificates yields a global certificate on $\bigcup_k W_k$ for the monitored degrees $i \in I$.

Proof sketch. Iterate Lemma G.11 and sum budgets; Summability ensures that the cumulative loss of safety margin remains bounded, so positivity of the margin persists. MECE coverage (Appendix A) ensures there are no gaps or overlaps; stability bands (if used) fix τ_k within regimes where $(\mu, u) = (0, 0)$. □

J.7. Practical detection, logging, and reproducibility

In practice, one proceeds along a windowed run (Appendix G) as follows:

1. *Select* τ_k on window W_k via a τ-sweep; identify a stability band by Proposition G.8.
2. *Evaluate* (μ, u) at the chosen τ_k and record ϕ_{i,τ_k}-ranks, $\mu_{i,\tau_k}, u_{i,\tau_k}$ in phi artifacts (Appendix G).
3. *Log the pipeline budget* $\Sigma\delta_k(i)$ and $\mathrm{gap}_{\tau_k}(i)$ in run.yaml (Appendix G) and check the Restart inequality.
4. *Paste certificates* across windows using Theorem G.13; ensure that coverage checks and MECE conventions pass (Appendix G).

Mirror/Transfer non-commutation is measured after collapse via a natural 2-cell and contributes additively to $\Sigma\delta$ (Appendix L).

J.8. Edge cases and pitfalls

- *Near-threshold accumulation.* If bars approach length τ from below, stability bands may be thin or absent; operate with finer sweeps and report (μ, u) sensitivity.
- *Metric drift without Summability.* If $\sum_k \Sigma\delta_k(i)$ diverges, the safety margin can be exhausted despite per-window passes; redesign the pipeline (e.g. impose geometric decay).
- *Degree mixing.* When aggregating $\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}}$ across degrees, a band for totals may hide degree-wise failures. Prefer reporting both per-degree and totals (Appendix G).

J.9. Formalization pointers (Lean/Coq) and API alignment

A minimal formalization (Appendix F) provides:

- phi as $\varinjlim \mathbf{T}_\tau \mathbf{P}_i(F_n) \to \mathbf{T}_\tau \mathbf{P}_i(F_\infty)$, with naturality in tower morphisms, composition, and cofinal reindexing.
- mu, nu as generic-fiber dimensions of kernel/cokernel; subadditivity, additivity, and invariance lemmas (Theorems G.3, G.4, G.6).

- Stability-band predicates over discrete τ-sweeps (Definition G.7) and refinement-stability (Proposition G.8).
- Restart/Summability contracts (κ, $\sum \Sigma\delta$) and a pasting theorem as in Theorem G.13. These align with the $\mathrm{run.yaml}$ schema (Appendix G).

J.10. Piecewise constancy off a finite critical set

The following self-contained statement ensures that $\tau \mapsto (\mu_{i,\tau}, u_{i,\tau})$ is piecewise constant away from a finite set of critical scales (per bounded τ-range), hence justifying stability-band detection by refinement.

Proposition G.14 (Finite critical set and piecewise constancy)**.** Fix a tower F, degree i, and a bounded interval $[a, b] \subset (0, \infty)$. There exists a finite set $S \subset [a, b]$ (depending on $F, i, [a, b]$) such that $\mu_{i,\tau}$ and $u_{i,\tau}$ are locally constant on each connected component of $[a, b] \setminus S$.

Proof sketch. In the constructible regime, after applying $\mathbf{T}_\tau$ the barcode of each $\mathbf{P}_i(F_n)$ changes only at finitely many thresholds (finite bar lengths and their finite sums/differences within $[a, b]$). Kernels and cokernels of $\phi_{i,\tau}$ vary only when a bar crosses the τ-cut or when the colimit/cocone structure alters the presence of infinite bars in kernel/cokernel. Thus only finitely many scales in $[a, b]$ can change $\mu_{i,\tau}$ or $u_{i,\tau}$. Between these scales, bar truncations and the induced linear maps (hence generic-fiber dimensions) are constant. □

Corollary G.15 (Band openness)**.** If ϕ_{i,τ_0} is an isomorphism for some $\tau_0 \in (a, b)$, then it is an isomorphism on an open neighborhood $U \subset (a, b)$ of τ_0. In particular, stability bands are unions of (nonempty) open intervals intersected with the sweep.

J.11. Window coherence and after-collapse order

All diagnostics $\phi_{i,\tau}$, $\mu_{i,\tau}$, and $u_{i,\tau}$ are computed *after* applying $\mathbf{T}_\tau$ (B-side single layer), and *with the same window and the same* τ as used by the gate and the δ-ledger (Appendix G). This "for each $t \to \mathbf{P}_i \to \mathbf{T}_\tau \to$compare" order is mandatory for auditability; it guarantees that subsequent 1-Lipschitz post-processing never increases budgets or alters band detection (Appendix A, Appendix L).

J.12. Minimal API sketch (pseudocode)

For engineering integration and tests (T11), we record a minimal pseudocode. It is schematic and omits low-level details (e.g. barcode calls).

```
# Compute (mu, nu) at scale tau for tower F and degree i
def compute_mu_nu(F, i, tau):
    # Persistence after collapse
    Ms = [ T_tau(P_i(F_n), tau) for F_n in F.levels ]
    M_inf = T_tau(P_i(F.apex), tau)
    # Colimit in [R, Vect] (scope policy; return to constructible if needed)
    colim_M = colim(Ms)          # compute objectwise in t
    phi = comparison(colim_M, M_inf)
    mu = generic_fiber_dim(kernel(phi))   # multiplicity of I[0,∞)
    nu = generic_fiber_dim(cokernel(phi))
    return mu, nu
```

```
# Detect stability band on a tau sweep (robust to refinement)
def detect_stability_band(F, i, tau_sweep):
    zero_idxs = [ell for ell,tau in enumerate(tau_sweep)
                 if compute_mu_nu(F,i,tau) == (0,0)]
    bands = maximal_contiguous_subarrays(zero_idxs)
    return bands   # evaluate robustness by inserting midpoints and retesting

# Restart inequality and summability check
def restart_ok(gap_next, gap_curr, dsum, kappa):
    return gap_next >= kappa * (gap_curr - dsum) and gap_curr > dsum

def summable_budget(deltas):
    # deltas: list of per-window sums Σ _k
    return (sum(deltas) < +infty)

# Paste certificates across windows W_k (degree i)
def paste_certificates(windows, i, kappas, budgets, gaps):
    assert len(windows)==len(kappas)==len(budgets)==len(gaps)
    for k in range(len(windows)-1):
        assert restart_ok(gaps[k+1], gaps[k], budgets[k], kappas[k])
    assert summable_budget(budgets)
    return True
```

Summary. Within the constructible range and the filtered-colimit scope of Appendix A, the comparison map $\phi_{i,\tau}$ is functorial in the tower and compatible with composition and finite sums; μ and u (and their totals $\mu_{\text{Collapse}}, u_{\text{Collapse}}$) inherit subadditivity under composition, additivity under finite direct sums, and invariance under cofinal reindexing (and under f.q.i. replacements). Moreover, $\tau \mapsto (\mu_{i,\tau}, u_{i,\tau})$ is piecewise constant off a finite critical set on bounded τ-windows, ensuring robust stability-band detection by sweep refinement. Finally, Restart and Summability provide a principled pasting mechanism for windowed certificates across MECE partitions: with per-window budgets $\Sigma\delta$ and safety margins gap_τ recorded in the run manifest (Appendix G), one obtains a global certificate whenever the per-window gate passes and the restart/-summability conditions hold. All comparisons and budgets are made after collapse on the B-side single layer, with window coherence and 1-Lipschitz post-processing ensuring nonincreasing budgets and reproducible audits.

Appendix K. Idempotent (Co)Monads for Collapse (up to f.q.i.) [Spec + Soft-Commuting Policy]

Standing conventions. We work over a field k. All persistence modules are constructible (locally finite on bounded windows). Filtered (co)limits are computed in the functor category $[\mathbb{R}, \mathsf{Vect}_k]$ under the scope policy of Appendix A, Remark .16; when stated, the result is returned to the constructible range. Reflection $\mathbf{T}_\tau \dashv \iota_\tau$ onto the τ-local (orthogonal) subcategory $(\mathsf{E}_\tau)^\perp \subset \mathsf{Pers}_k^{\text{ft}}$ is exact and 1-Lipschitz (Appendix A, Theorem .14, Proposition .20). Global conventions: Ext^1 is always against $k[0]$; the energy exponent is $\alpha > 0$ (default $\alpha = 1$); windows are MECE and right-open (Appendix G). Distances are measured by the interleaving metric d_{int} (equivalently, bottleneck distance on barcodes for constructible objects).

Throughout, we implicitly identify functors and natural transformations up to filtered quasi-isomorphism (abbrev. f.q.i.) whenever we work in the filtered-homotopy setting; all such identifications are stable under the persistence functors $\mathbf{P}_i$.

K.1. (Persistence layer) the idempotent monad $\iota_\tau \mathbf{T}_\tau$

Let $\iota_\tau : (\mathsf{E}_\tau)^\perp \hookrightarrow \mathsf{Pers}_k^{\mathrm{ft}}$ be the fully faithful inclusion and $\mathbf{T}_\tau : \mathsf{Pers}_k^{\mathrm{ft}} \to (\mathsf{E}_\tau)^\perp$ its left adjoint (Appendix A, Theorem .14). Consider the endofunctor

$$\mathbf{M}_\tau \ := \ \iota_\tau \circ \mathbf{T}_\tau : \ \mathsf{Pers}_k^{\mathrm{ft}} \longrightarrow \mathsf{Pers}_k^{\mathrm{ft}}.$$

Theorem G.16 (Idempotent monad on $\mathsf{Pers}_k^{\mathrm{ft}}$)**.** With unit $\eta : \mathrm{Id} \Rightarrow \iota_\tau \mathbf{T}_\tau$ and multiplication $\mu : \mathbf{M}_\tau^2 = \iota_\tau \mathbf{T}_\tau \iota_\tau \mathbf{T}_\tau \xRightarrow{\iota_\tau\, \varepsilon\, \mathbf{T}_\tau} \iota_\tau \mathbf{T}_\tau = \mathbf{M}_\tau$, induced by the counit $\varepsilon : \mathbf{T}_\tau \iota_\tau \Rightarrow \mathrm{Id}$ on $(\mathsf{E}_\tau)^\perp$, the triple $(\mathbf{M}_\tau, \eta, \mu)$ is a monad on $\mathsf{Pers}_k^{\mathrm{ft}}$. Moreover, μ is a natural isomorphism (idempotence).

Proof. This is standard for reflective subcategories with fully faithful right adjoint. The triangle identities for the adjunction $\mathbf{T}_\tau \dashv \iota_\tau$ yield the monad axioms for $(\iota_\tau \mathbf{T}_\tau, \eta, \mu)$. Since ι_τ is fully faithful, the counit ε is an isomorphism on $(\mathsf{E}_\tau)^\perp$; hence $\mu = \iota_\tau\, \varepsilon\, \mathbf{T}_\tau$ is a natural isomorphism, proving idempotence. ◻

Proposition G.17 (Exactness and non-expansiveness)**.** $\mathbf{M}_\tau$ is exact on $\mathsf{Pers}_k^{\mathrm{ft}}$ and 1-Lipschitz for the interleaving/bottleneck metric:

$$d_{\mathrm{int}}\big(\mathbf{M}_\tau M,\ \mathbf{M}_\tau N\big) \ \le\ d_{\mathrm{int}}(M, N) \qquad (M, N \in \mathsf{Pers}_k^{\mathrm{ft}}).$$

Proof. Exactness follows because both $\mathbf{T}_\tau$ and ι_τ are exact (Appendix A, Theorem .14). The 1-Lipschitz property follows from the fact that $\mathbf{T}_\tau$ commutes with shifts and truncations used in the definition of interleavings, hence it does not increase the minimal interleaving parameter (Appendix A, Proposition .20). ◻

Remark G.18 (Algebras and fixed points)**.** The Eilenberg-Moore category of $\mathbf{M}_\tau$-algebras identifies with $(\mathsf{E}_\tau)^\perp$ via ι_τ. Concretely, $\mathbf{M}_\tau$-algebras $\theta : \mathbf{M}_\tau M \to M$ with the usual coherence are in bijection with objects for which $\eta_M : M \to \mathbf{M}_\tau M$ is an isomorphism, i.e. the essential image of ι_τ. In particular, $\mathbf{M}_\tau$ acts as a *collapse* to the τ-local part, and $\mathbf{M}_\tau$-modules are precisely the τ-collapsed objects.

Example G.19 (Length threshold reflector)**.** In 1-dimensional persistence, $\mathbf{T}_\tau$ kills all summands of lifespan $< \tau$. Then $\mathbf{M}_\tau$ replaces a module by its τ-truncated part. By idempotence, $\mathbf{M}_\tau \mathbf{M}_\tau = \mathbf{M}_\tau$.

K.2. (Ho($\mathsf{FiltCh}(k)$)—implementable range, up to f.q.i.) the idempotent comonad $\iota \circ C_\tau^{\mathrm{comb}}$ [Spec]

[Spec] There exists an implementable subcategory $\mathrm{Ho}(\mathsf{FiltCh}(k))_\tau^{\mathrm{comb}} \subset \mathrm{Ho}(\mathsf{FiltCh}(k))$ ("τ-combinatorial collapses") and a fully faithful inclusion

$$\iota : \ \mathrm{Ho}(\mathsf{FiltCh}(k))_\tau^{\mathrm{comb}} \hookrightarrow \mathrm{Ho}(\mathsf{FiltCh}(k))$$

admitting a right adjoint (coreflector) $C_\tau^{\mathrm{comb}} : \mathrm{Ho}(\mathsf{FiltCh}(k)) \to \mathrm{Ho}(\mathsf{FiltCh}(k))_\tau^{\mathrm{comb}}$, natural up to filtered quasi-isomorphism (Appendix B). Define

$$\mathbf{G}_\tau \ := \ \iota \circ C_\tau^{\mathrm{comb}} : \ \mathrm{Ho}(\mathsf{FiltCh}(k)) \longrightarrow \mathrm{Ho}(\mathsf{FiltCh}(k)).$$

Theorem G.20 (Idempotent comonad on $\mathrm{Ho}(\mathsf{FiltCh}(k))$ up to f.q.i.)**.** Assuming $\iota \dashv C_\tau^{\mathrm{comb}}$, with counit $\varepsilon : \mathbf{G}_\tau \Rightarrow \mathrm{Id}$ and comultiplication $\delta : \mathbf{G}_\tau \xRightarrow{\iota\, \eta\, C_\tau^{\mathrm{comb}}} \mathbf{G}_\tau^2$ (η the unit on $\mathrm{Ho}(\mathsf{FiltCh}(k))_\tau^{\mathrm{comb}}$), the triple $(\mathbf{G}_\tau, \varepsilon, \delta)$ is a comonad on $\mathrm{Ho}(\mathsf{FiltCh}(k))$. Moreover, δ is a natural isomorphism (idempotence). All identities hold in $\mathrm{Ho}(\mathsf{FiltCh}(k))$ and are invariant under filtered quasi-isomorphisms.

Proof. Dual to Theorem G.16. For a coreflective subcategory with fully faithful left adjoint, the comonad axioms follow from the triangle identities. Since η is an isomorphism on $\mathrm{Ho}(\mathsf{FiltCh}(k))^{\mathrm{comb}}_{\tau}$, idempotence follows. Equalities are interpreted in the localized category where f.q.i. are inverted. □

Proposition G.21 (Compatibility with persistence)**.** For any filtered complex F,

$$\mathbf{P}_i\big(\mathbf{G}_\tau F\big) \;\cong\; \mathbf{M}_\tau\big(\mathbf{P}_i(F)\big) \qquad \text{naturally in } i \text{ and } F,$$

and, after $\mathbf{P}_i$,

$$d_{\mathrm{int}}\big(\mathbf{P}_i(\mathbf{G}_\tau F),\ \mathbf{P}_i(\mathbf{G}_\tau G)\big) \;\le\; d_{\mathrm{int}}\big(\mathbf{P}_i(F),\ \mathbf{P}_i(G)\big).$$

Proof. By construction C^{comb}_{τ} lifts $\mathbf{T}_\tau$ (Appendix B), hence $\iota C^{\mathrm{comb}}_{\tau}$ lifts $\iota_\tau \mathbf{T}_\tau$ at the persistence level, yielding the natural isomorphism after applying $\mathbf{P}_i$. Non-expansiveness follows from Proposition G.17 because $\mathbf{P}_i$ is 1-Lipschitz with respect to filtered interleavings and d_{int}. □

Remark G.22 (Scope and usage)**.** The comonad $\mathbf{G}_\tau$ is asserted only on the *implementable range (up to f.q.i.)*, which suffices for algorithms and stability statements. All filtered (co)limit claims comply with Appendix A, Remark .16. In practice, computations are performed on a B-side single layer after collapse, and all claims are made modulo f.q.i.

K.3. Multi-axis torsion reflectors and soft-commuting policy

Beyond the length-threshold reflector $\mathbf{T}_\tau$, one may consider other exact reflectors $T_A, T_B : \mathsf{Pers}^{\mathrm{ft}}_k \to \mathsf{Pers}^{\mathrm{ft}}_k$ arising from hereditary Serre subcategories E_A, E_B (e.g. birth-window deletion, lifespan windowing, support cuts). This subsection specifies order-independence in the nested case and an *operational* "soft-commuting" policy in the general case.

Definition G.23 (Exact reflectors from Serre classes)**.** Let $E \subset \mathsf{Pers}^{\mathrm{ft}}_k$ be a hereditary Serre subcategory. The Serre localization yields an exact reflector $T_E : \mathsf{Pers}^{\mathrm{ft}}_k \to E^{\perp}$ left adjoint to the inclusion $E^{\perp} \hookrightarrow \mathsf{Pers}^{\mathrm{ft}}_k$. We write $T_A := T_{E_A}$, $T_B := T_{E_B}$, and $E_{A\vee B}$ for the Serre subcategory generated by $E_A \cup E_B$.

Proposition G.24 (Nested torsions ⇒ order independence)**.** If $E_A \subseteq E_B$ or $E_B \subseteq E_A$, then

$$T_A \circ T_B \;=\; T_B \circ T_A \;=\; T_{A\vee B}$$

as exact reflectors on $\mathsf{Pers}^{\mathrm{ft}}_k$. In particular, for the 1-dimensional length thresholds, $\mathbf{T}_\tau \circ \mathbf{T}_\sigma = \mathbf{T}_{\max\{\tau,\sigma\}}$.

Proof. In an abelian category, localization at nested Serre subcategories is idempotent and order-independent: if $E_A \subseteq E_B$, then $E_B^{\perp} \subseteq E_A^{\perp}$. The composite $T_A \circ T_B$ (resp. $T_B \circ T_A$) is the reflector onto $E_A^{\perp} \cap E_B^{\perp} = (E_{A\vee B})^{\perp}$. Hence both composites identify with the reflector at the join. The special case for length thresholds follows because the subcategory of intervals of length $< \tau$ is nested in that of length $< \sigma$ when $\tau \ge \sigma$. □

Definition G.25 (A/B commutation defect and soft-commuting policy)**.** For arbitrary exact reflectors T_A, T_B and a dataset $M \in \mathsf{Pers}^{\mathrm{ft}}_k$, define the *commutation defect*

$$\Delta_{\mathrm{comm}}(M; A, B) \;:=\; d_{\mathrm{int}}\big(T_A T_B M,\ T_B T_A M\big).$$

Fix a tolerance $\eta \ge 0$ and a window (MECE, right-open). We say T_A, T_B are *soft-commuting* on the dataset if $\Delta_{\mathrm{comm}}(M; A, B) \le \eta$ holds on the relevant instances (per window). Otherwise we *fallback* to a fixed order (e.g. $T_B \circ T_A$).

Declaration G.26 (Operational policy and δ-ledger integration)**.** Within a window W and for a monitored degree i:

1. *Test.* Measure $\Delta_{\text{comm}}(M; A, B)$ at the persistence layer after applying any prerequisite collapse(s), using the window policy (Appendix G). If $\Delta_{\text{comm}} \le \eta$, adopt *soft-commuting* and allow either order in that window.
2. *Fallback.* If $\Delta_{\text{comm}} > \eta$, fix an order (e.g. $T_B \circ T_A$) and *record* the observed Δ_{comm} as an δ^{alg} entry in the δ-ledger (Appendix G).
3. *Pipeline additivity.* Over a pipeline of such steps across windows, add the recorded defects to the window budgets $\Sigma\delta$ (Appendix G) and consume them in the Restart/Summability calculus (Appendix J).
4. *Reproducibility.* Log η, the chosen order (or the soft-commuting verdict), and the measured Δ_{comm} per window in $\mathrm{run.yaml}$ (Appendix G).

Remark G.27 (Scope and guard-rails)**.** We do *not* assert any a priori commutation for non-nested torsions; order control is purely *operational* via the A/B test. All distances are measured *after* truncation by $\mathbf{T}_\tau$ and on the B-side single layer, consistent with Chapter 1 and Appendix L.

K.4. Interaction with Mirror/Transfer and pipeline δ-budget

Let Mirror be an admissible endofunctor at the filtered level (or on $\mathsf{Pers}_k^{\text{ft}}$) satisfying the quantitative commutation hypotheses of Appendix L: 1-Lipschitz on $\mathbf{P}_i$ and a natural 2-cell Mirror $\circ C_\tau \Rightarrow C_\tau \circ$ Mirror with a uniform bound $\delta(i, \tau) \ge 0$ in d_{int}.

Proposition G.28 (Budget accounting with multiple (co)reflectors and Mirror)**.** In a pipeline that interleaves exact reflectors $T_\bullet$ (including $\mathbf{T}_\tau$) and Mirror/Transfer steps, the total commutation/ordering slack budget on a window is bounded by the *additive* sum of:

- each Mirror–Collapse defect $\delta(i, \tau)$ (Appendix L),
- each A/B commutation defect $\Delta_{\text{comm}}(M; A, B)$ when soft-commuting fails (Declaration G.26),

and is *non-increasing* under any subsequent 1-Lipschitz post-processing (e.g. further truncations, degree projections) at the persistence layer.

Proof. Concatenate the natural 2-cells along the pipeline and apply the triangle inequality for d_{int}. Each Mirror–Collapse interchange contributes at most $\delta(i, \tau)$. When soft-commuting fails for T_A, T_B, the observed Δ_{comm} accounts for the order selection; adding these yields an upper bound on the overall slack. Since all subsequent steps are 1-Lipschitz (Appendix A, Proposition .20), the bound cannot increase under post-processing. □

Corollary G.29 (Window-wise additivity and stability)**.** If a windowed pipeline uses windows $\{W_j\}_{j\in J}$ (MECE, right-open) and records per-window contributions δ_j^{Mirror} and $\delta_j^{\text{A/B}}$, then

$$\textstyle\sum_{j\in J} \left(\delta_j^{\text{Mirror}} + \delta_j^{\text{A/B}}\right)$$

is a valid global upper bound on the end-to-end slack, and any 1-Lipschitz aggregator on the persistence side preserves this bound.

K.5. Windowed usage and reproducibility ($\mathrm{run.yaml}$ alignment)

For each window W record in $\mathrm{run.yaml}$ (Appendix G):

- the set of reflectors used (e.g. length, birth_window), their order, and whether *soft-commuting* was adopted (ab_test block with tolerance η);
- the measured Δ_{comm} (if any), logged into the δ-ledger under δ^{alg};
- the Mirror/Transfer commutation bounds $\delta(i,\tau)$ (Appendix L) aggregated into $\Sigma\delta$;
- the B-Gate^{+} verdict and the safety margin gap_{τ}, to be consumed by Restart/Summability (Appendix J).

A minimal schema (illustrative only) is:

```
window:
  id: "W01"
  degree: 1
  tau: 0.15
  reflectors:
    - name: length
    - name: birth_window
  ab_test:
    eta: 0.02
    verdict: soft_commuting | fallback
    chosen_order: [birth_window, length]   # if fallback
    Delta_comm: 0.031                      # if measured
  mirror:
    delta:
      i: 1
      tau: 0.15
      value: 0.010
  budget:
    delta_alg: 0.031
    delta_mirror: 0.010
    sum_delta: 0.041
  b_gate_plus:
    passed: true
    gap_tau: 0.12
```

Remark G.30 (Window coherence)**.** A/B tests and Mirror–Collapse commutation checks must be performed on the *same* window (MECE, right-open) and the *same* collapse thresholds τ as logged in run.yaml to ensure consistent budget accounting and pasting (Appendix J).

K.6. Formalization stubs (Lean/Coq) [Spec]

A minimal API (cf. Appendix F) includes:

- reflector_E for a hereditary Serre subcategory E, yielding an exact reflector T_E and the monad $\iota_E T_E$ (idempotent);
- a lemma nested_torsion_order_indep stating $T_A \circ T_B = T_B \circ T_A = T_{A\vee B}$ when $E_A \subseteq E_B$ or $E_B \subseteq E_A$;
- a comm_defect functional producing $\Delta_{\text{comm}}(M;A,B) = d_{\text{int}}(T_A T_B M, T_B T_A M)$ with a policy hook for *soft-commuting* vs. fallback;

- a Mirror–Collapse 2-cell contract mirror_collapse_delta yielding $\delta(i,\tau)$ and the additive pipeline bound under 1-Lipschitz post-processing.

Explicit signatures (schematic):

```
constant reflector_E :
  (E : SerreSubcat (Pers_ft k)) ->
  Σ (T : Pers_ft k   Pers_ft k),
    (is_exact T) × (is_reflector E  T)

theorem nested_torsion_order_indep
  {E_A E_B : SerreSubcat (Pers_ft k)}
  (h : E_A   E_B   E_B   E_A) :
  T_A   T_B   T_{A B}   T_B   T_A   T_{A B}

def comm_defect (M : Pers_ft k)
  (A B : Reflector) :   :=
  d_int ((A   B) M) ((B   A) M)

constant mirror_collapse_delta :
  Π (i :  ) (  :   0),   0
```

K.7. Edge cases and pitfalls

- *Non-exact "reflectors".* The policy assumes exact reflectors (Serre localization). Heuristic truncations without exactness may break monad laws and invalidate stability claims.
- *Inconsistent windows.* Running A/B tests or Mirror–Collapse checks on mismatched windows/τ invalidates budget accounting and the pasting guarantee (Appendix J).
- *Chained non-nested axes.* With T_A, T_B, T_C mutually non-nested, pairwise soft-commuting does not imply global confluence. Adopt a fixed canonical order (e.g. by axis priority), A/B test adjacent pairs, and log residuals.
- *Hidden state across layers.* Always measure Δ_{comm} *after* the B-side collapse and at the monitored degree to avoid spurious layer interactions.
- *Threshold drift.* If τ is adapted online, ensure all δ-ledger entries specify the τ in force at measurement time; do not retroactively amend past entries.

Summary. On the persistence side, collapse is governed by the idempotent monad $\iota_\tau \mathbf{T}_\tau$, exact and 1-Lipschitz. On the filtered–homotopy side (implementable range, up to f.q.i.), collapse is governed by the idempotent comonad $\iota \circ C_\tau^{\text{comb}}$. For *multi-axis* torsion reflectors, order independence holds under nesting; otherwise, an A/B *soft-commuting* policy with a measured defect Δ_{comm} provides operational control, integrating into the additive pipeline δ-budget together with Mirror/Transfer commutation (Appendix L). All steps are carried out *after* collapse on the B-side single layer, logged per window (Appendix G), and pasted globally via Restart/Summability (Appendix J).

Appendix L. Quantitative Commutation for Mirror/Tropical [Spec + Pipeline Budget + A/B Policy]

Standing conventions. We work over a field k; all persistence modules are constructible. Filtered (co)limits are computed in $[\mathbb{R}, \mathsf{Vect}_k]$ under the scope policy of Appendix A, Remark .16, and then returned to the constructible range. The interleaving metric d_{int} (= bottleneck) is used throughout. The truncation $\mathbf{T}_\tau$ is exact and 1-Lipschitz; shift–commutation implies 1-Lipschitz (Appendix A, Proposition .20). On the filtered-complex side we write C_τ; on the persistence side $\mathbf{T}_\tau$. Global conventions: Ext^1 is always against $k[0]$; the energy exponent satisfies $\alpha > 0$ (default $\alpha = 1$); windows are MECE and right–open (Appendix G). All comparisons follow the unique order "for each $t \to \mathbf{P}_i \to \mathbf{T}_\tau \to$compare," on the B-side single layer after collapse (Chapter 1, Appendix N).

L.1. Hypotheses

[Spec] Let Mirror be a functor on filtered complexes (or, via $\mathbf{P}_i$, a persistence-layer endofunctor). Assume:

(H1) 1**-Lipschitz of** Mirror**.** For every degree i,

$$d_{\text{int}}\Big(\mathbf{P}_i(\text{Mirror}\,F),\ \mathbf{P}_i(\text{Mirror}\,G)\Big) \ \le\ d_{\text{int}}\Big(\mathbf{P}_i(F),\ \mathbf{P}_i(G)\Big) \qquad \text{for all filtered complexes } F, G.$$

(No extra assumption is needed for C_τ: by Appendix B one has $\mathbf{P}_i(C_\tau F) \cong \mathbf{T}_\tau(\mathbf{P}_i F)$, and $\mathbf{T}_\tau$ is 1-Lipschitz by Appendix A, Proposition .20.)

(H2) δ**-controlled natural** 2**-cell.** There exists a natural 2-cell

$$\theta:\ \text{Mirror}\circ C_\tau \ \Rightarrow\ C_\tau \circ \text{Mirror}$$

(interpreted up to f.q.i. in $\mathrm{Ho}(\mathsf{FiltCh}(k))$; cf. Appendix B) such that, *uniformly in* F, for all i, τ,

$$d_{\text{int}}\Big(\mathbf{P}_i(\text{Mirror}(C_\tau F)),\ \mathbf{P}_i(C_\tau(\text{Mirror}\,F))\Big) \ \le\ \delta(i,\tau).$$

Remark H.1 (Interpretation)**.** (H2) measures, at the persistence layer, the failure of Mirror and C_τ to commute; the control $\delta(i,\tau) \ge 0$ may depend on (i,τ) but not on F. It is an external assumption; no intrinsic value of δ is claimed without additional structure.

Remark H.2 (Strict commutation as a special case)**.** If $\delta(i,\tau) = 0$ (e.g. θ induces isomorphisms after applying $\mathbf{P}_i$), then Mirror and C_τ commute up to isomorphism at the truncated persistence layer in degree i and window τ.

L.2. Conditional bound

Theorem H.3 (Quantitative commutation under (H1)–(H2), uniform in F)**.** Under *(H1)* and *(H2)*, for all filtered complexes F, all i, and all τ,

$$d_{\text{int}}\Big(\mathbf{T}_\tau\,\mathbf{P}_i(\text{Mirror}(C_\tau F))\ ,\ \mathbf{T}_\tau\,\mathbf{P}_i(C_\tau(\text{Mirror}\,F))\Big) \ \le\ \delta(i,\tau),$$

and the bound is *uniform in* F.

Proof sketch. By (H2), $d_{\mathrm{int}}\big(\mathbf{P}_i(\mathrm{Mirror}(C_\tau F)), \mathbf{P}_i(C_\tau(\mathrm{Mirror}\, F))\big) \le \delta(i,\tau)$. Applying $\mathbf{T}_\tau$, which is 1-Lipschitz (Appendix A, Proposition .20), yields the stated inequality. Since $\mathbf{T}_\tau$ is exact, *no additional stability loss* is introduced for downstream kernel/cokernel diagnostics.[10] □

Corollary H.4 (Stability under additional 1-Lipschitz post-processing)**.** If Φ is any endofunctor on persistence modules with $d_{\mathrm{int}}(\Phi X, \Phi Y) \le d_{\mathrm{int}}(X,Y)$, then for all i, τ, F,

$$d_{\mathrm{int}}\Big(\Phi\, \mathbf{T}_\tau\, \mathbf{P}_i(\mathrm{Mirror}(C_\tau F))\ ,\ \Phi\, \mathbf{T}_\tau\, \mathbf{P}_i(C_\tau(\mathrm{Mirror}\, F))\Big) \ \le\ \delta(i,\tau).$$

Corollary H.5 (Two-stage pipeline; additive control)**.** Let $\mathrm{Mirror}_1, \mathrm{Mirror}_2$ satisfy *(H1)* and admit 2-cells with controls δ_1, δ_2 as in *(H2)*. Then for all i, τ, F,

$$d_{\mathrm{int}}\Big(\mathbf{T}_\tau\, \mathbf{P}_i(\mathrm{Mirror}_2\, \mathrm{Mirror}_1(C_\tau F))\ ,\ \mathbf{T}_\tau\, \mathbf{P}_i(C_\tau(\mathrm{Mirror}_2\, \mathrm{Mirror}_1\, F))\Big) \ \le\ \delta_1(i,\tau) + \delta_2(i,\tau).$$

Proof sketch. Insert the intermediate term $\mathbf{T}_\tau \mathbf{P}_i(\mathrm{Mirror}_2(C_\tau\, \mathrm{Mirror}_1\, F))$ and apply the triangle inequality. Use Theorem H.3 for Mirror_1 with the 1-Lipschitz post-processing $\Phi = \mathbf{T}_\tau \mathbf{P}_i\, \mathrm{Mirror}_2$ to bound the first leg by δ_1, and Theorem H.3 for Mirror_2 to bound the second leg by δ_2.

Corollary H.6 (Strict commutation and tower diagnostics)**.** If $\delta(i,\tau) = 0$ for all i, τ (i.e. Mirror and C_τ commute up to isomorphism at the truncated persistence layer), then for any tower $\{F_n\} \to F_\infty$ and any i, τ, the comparison maps computed after either order agree up to isomorphism:

$$\phi_{i,\tau}\big(\mathrm{Mirror}(C_\tau F_\bullet)\big) \ \cong\ \phi_{i,\tau}\big(C_\tau(\mathrm{Mirror}\, F_\bullet)\big).$$

Consequently, the tower obstruction indices are preserved,

$$\mu_{i,\tau}\big(\mathrm{Mirror}(C_\tau F_\bullet)\big) = \mu_{i,\tau}\big(C_\tau(\mathrm{Mirror}\, F_\bullet)\big), u_{i,\tau}\big(\mathrm{Mirror}(C_\tau F_\bullet)\big) = u_{i,\tau}\big(C_\tau(\mathrm{Mirror}\, F_\bullet)\big),$$

and likewise for the totals $(\mu_{\mathrm{Collapse}}, u_{\mathrm{Collapse}})$.

Remark H.7 (Typical sources of $\delta = 0$)**.** If Mirror preserves the filtration and there is a filtration–natural isomorphism $\mathrm{Mirror} \circ C_\tau \cong C_\tau \circ \mathrm{Mirror}$ (e.g. Mirror is induced by a reindexing monotone endomorphism of $(\mathbb{R}, \le)$ that fixes the window cut, or by a direct-sum decomposition functor that respects sublevel sets), then $\delta(i,\tau) = 0$ for all i, τ.

Remark H.8 (Uniformity across degrees)**.** If Mirror admits a degreewise 1-Lipschitz constant independent of i on a fixed finite set of degrees $I \subset \mathbb{Z}$, then the commutation control may be chosen $\delta(i,\tau) \le \delta_I(\tau)$ uniformly for $i \in I$.

L.3. Counterexample notes (necessity of assumptions)

Remark H.9 (Why the hypotheses are needed)**.** Without a δ-controlled natural 2-cell (H2), individual non-expansiveness of Mirror and C_τ does *not* constrain the distance between the two composites: triangle inequalities do not relate $\mathrm{Mirror} \circ C_\tau$ and $C_\tau \circ \mathrm{Mirror}$ in the absence of a comparison map. Concretely, one can build (Appendix D) towers whose bar lengths accumulate from below at τ (Type IV phenomena). Collapsing first may erase the near-τ mass, while mirroring first may convert it into persistent features that survive collapse; the gap can be made arbitrarily large despite each functor being 1-Lipschitz on its own. Thus no quantitative bound of the form in Theorem H.3 is available in general unless (H2) (or an equivalent commutation control) is imposed.

Remark H.10 (Scope)**.** All claims are made under the constructible scope and filtered-(co)limit policy of Appendix A. When Mirror is only defined up to filtered quasi-isomorphism, (H2) is interpreted in $\mathrm{Ho}(\mathsf{FiltCh}(k))$, and distances are taken after applying $\mathbf{P}_i$.

[10] Exactness does not turn a small distance bound into equality of kernels/cokernels; it ensures that truncation itself adds no extra error to such diagnostics.

L.4. Pipeline error budget and windowed accounting

We integrate quantitative commutation into a *windowed* error budget that is additive along pipelines, uniform in the input F, and non-increasing under any subsequent 1-Lipschitz post-processing at the persistence layer.

Definition H.11 (Per-step commutation bounds and window budget)**.** Consider a pipeline on a window $W = [u, u')$ consisting of steps

$$\Pi \ := \ \Big(\ \cdots \ \to \ \mathrm{Mirror}_j \ \to \ C_{\tau_j} \ \to \ T_{A_j/B_j} \ \to \ \cdots \ \Big),$$

where Mirror_j are Mirror/Transfer steps, C_{τ_j} are collapses (filtered lift of $\mathbf{T}_{\tau_j}$), and T_{A_j/B_j} are exact reflectors on $\mathsf{Pers}_k^{\mathrm{ft}}$ (Appendix K). For each Mirror–Collapse pair, let $\delta_j(i, \tau_j)$ be the (H2) bound on degree i. For each pair of reflectors T_A, T_B that are applied in some order with an *A/B test* (Appendix K), let $\Delta_{\mathrm{comm}}(M; A, B)$ be the measured commutation defect (Definition G.25); when $\Delta_{\mathrm{comm}} > \eta$, we fallback to a fixed order and record Δ_{comm} as an δ^{alg} contribution. The *window budget* is the additive sum of all such contributions on W:

$$\Sigma\delta_W(i) \ := \ \sum_j \delta_j(i, \tau_j) \ + \sum_{\text{A/B tested pairs}} \Delta_{\mathrm{comm}}(M; A, B) \quad \text{(when exceeded tolerance)}.$$

Theorem H.12 (Pipeline budget: additivity, uniformity, and post-processing stability)**.** Fix a window W, a monitored degree i, and a pipeline Π as in Definition H.11. For any filtered input F,

$$d_{\mathrm{int}}\Big(\mathbf{T}_\tau \, \mathbf{P}_i(\Pi_{\mathrm{lhs}}(F)) \ , \ \mathbf{T}_\tau \, \mathbf{P}_i(\Pi_{\mathrm{rhs}}(F))\Big) \ \le \ \Sigma\delta_W(i),$$

where the left/right versions aggregate Mirror–Collapse orderings and A/B reflector orderings according to the window's operational choices. The bound:

- is *additive* in the per-step Mirror–Collapse bounds $\delta_j(i, \tau_j)$ and the reflector A/B defects Δ_{comm};
- is *uniform in F* (each δ_j is uniform by (H2), and Δ_{comm} are measured at the persistence layer independent of representatives up to f.q.i.);
- is *non-increasing* under any subsequent 1-Lipschitz post-processing on the persistence layer (e.g. further truncations $\mathbf{T}_{\tau'}$, degree projections $\mathbf{P}_i$, shifts S^{ε}).

Proof sketch. Compose the natural 2-cells for Mirror–Collapse steps and measure each A/B reflector defect in d_{int}; apply the triangle inequality to obtain additivity. Uniformity follows from (H2) being uniform in F and from the fact that A/B defects are computed *after* applying $\mathbf{P}_i$ (hence f.q.i. invariant). Post-processing stability comes from the 1-Lipschitz property of the applied persistence functors (Appendix A). □

Remark H.13 (run.yaml alignment and B-Gate$^+$)**.** Record each $\delta_j(i, \tau_j)$ and Δ_{comm} per window in the δ-ledger (Appendix G); aggregate into $\Sigma\delta_W(i)$ and compare to the safety margin gap_τ for B-Gate$^+$ on the same window and degree (Chapter 1). This integrates with Restart/Summability for pasting (Appendix J).

L.5. Operational A/B policy (test ⇒ fallback ⇒ δ_{alg} accounting)

We adopt the *soft-commuting* policy of Appendix K at the persistence layer after collapse, with explicit windowed logging.

Definition H.14 (A/B commutativity test)**.** Given exact reflectors T_A, T_B (Appendix K) and a dataset $M \in \mathsf{Pers}_k^{\mathrm{ft}}$ on window W, define

$$\Delta_{\mathrm{comm}}(M; A, B) \ := \ d_{\mathrm{int}}\big(T_A T_B M, \ T_B T_A M\big).$$

Fix a tolerance $\eta \ge 0$. If $\Delta_{\mathrm{comm}} \le \eta$, we say T_A, T_B are *soft-commuting* on W.

Declaration H.15 (Windowed A/B procedure)**.** On a window W and degree i:

1. *Measure* $\Delta_{\text{comm}}(M; A, B)$ on $\mathbf{T}_\tau \mathbf{P}_i$ (same window and τ as the budget).
2. *Accept* soft-commuting if $\Delta_{\text{comm}} \le \eta$; otherwise *fallback* to a fixed order (e.g. $T_B \circ T_A$).
3. *Record* Δ_{comm} as an δ^{alg} contribution if soft-commuting fails; add it to $\Sigma\delta_W(i)$.
4. *Log* η, the decision (soft-commute or fixed order), and the numerical Δ_{comm} in run.yaml (Appendix G).

Proposition H.16 (Combined bound with A/B policy)**.** For a window W with the A/B procedure of Declaration H.15, the pipeline bound in Theorem H.12 holds with

$$\Sigma\delta_W(i) \;=\; \sum_j \delta_j(i, \tau_j) \;+\; \sum_{\text{A/B fails}} \Delta_{\text{comm}}(M; A, B),$$

and is non-increasing under any subsequent 1-Lipschitz persistence post-processing.

Proof sketch. Immediate from Theorem H.12 by including the A/B residuals as per-step additive costs. □

Remark H.17 (Best practices)**.** Use a canonical ordering for multiple axes (e.g. priority by axis type) and apply A/B tests only to adjacent pairs; record all policies and tolerances. For nested torsions (Appendix K, Proposition G.24), skip the test (order independence holds).

L.6. Edge cases, pitfalls, and window coherence

- *Mismatched windows or thresholds.* All commutation bounds must be computed on the *same* window and collapse thresholds τ as used for B-Gate$^+$ and for the δ-ledger; otherwise budget accounting and pasting (Appendix J) become invalid.
- *Degree mixing.* When reporting a total over degrees (e.g. $\mu_{\text{Collapse}}, u_{\text{Collapse}}$), a small A/B residual in one degree can be obscured. Keep both per-degree and aggregated logs (Appendix G).
- *Cascaded non-commutation.* Pairwise soft-commuting does not imply global commutation across three or more axes; use a fixed canonical order and log all residuals.

L.7. Minimal pseudocode (pipeline budget and ledger)

The following schematic code illustrates how to accumulate windowed commutation budgets in a pipeline; it is consistent with the run.yaml schema in Appendix G.

```
# Mirror–Collapse commutation ledger for one window W, degree i
def window_budget(steps, i, tau):
    ”””
    steps: list of pipeline steps on window W, each tagged as:
          - (”mirror”, Mirror_j, tau_j, delta_ij)   #  _j(i,  _j)
          - (”reflectors”, T_A, T_B, eta, Delta)   # A/B with tolerance  , measured Δ
          - (”post”, functor)                      # 1-Lipschitz post-processing
    ”””
    Sigma = 0.0
    for s in steps:
        if s[0] == ”mirror”:
            _, _, tau_j, delta_ij = s
```

```
        Sigma += delta_ij
    elif s[0] == "reflectors":
        _, T_A, T_B, eta, Delta = s
        if Delta > eta:
            Sigma += Delta          # fallback; record into ^{alg}
        else:
            pass                    # soft-commuting; no cost added
    elif s[0] == "post":
        pass                        # 1-Lipschitz => no increase
  return Sigma

# Gate check (B-side, after collapse), same window and
def gate_passed(Sigma_delta, gap_tau):
  return (gap_tau > Sigma_delta)
```

Summary. If Mirror is 1-Lipschitz and there is a natural 2-cell controlling its commutator with C_τ by $\delta(i, \tau)$ (uniform in F), then truncation preserves this control:

$$d_{\mathrm{int}}\Big(\mathbf{T}_\tau \mathbf{P}_i(\mathrm{Mirror}\, C_\tau F),\ \mathbf{T}_\tau \mathbf{P}_i(C_\tau\, \mathrm{Mirror}\, F)\Big)\ \le\ \delta(i, \tau).$$

For multi-stage pipelines the controls *add*, are *uniform in F*, and are *non-increasing* under 1-Lipschitz post-processing (Theorem H.12). Operationally, A/B tests for non-nested reflectors provide a *soft-commuting* decision per window; failures fall back to a fixed order and are logged as δ^{alg} contributions (Declaration H.15, Proposition H.16). All items are computed *after* collapse on the B-side single layer, logged per window (run.yaml, Appendix G), and pasted globally via Restart/Summability (Appendix J).

Appendix M. (Optional) Lax Monoidal Compatibility [Spec + Windowed Usage + Budget Integration]

Status. All claims marked [Spec] are windowed, reproducible contracts intended for operational use. They are compatible with the standard 1D constructible persistence setting and do not assert strict monoidality beyond the stated hypotheses. Full derivations are included so that no further reinforcement is required.

M.0. Standing conventions and scope

Ground field and categories. Fix a field k. Let $\mathsf{Pers}_k^{\mathrm{c}}$ denote the category of constructible (1-parameter) persistence modules $M : \mathbb{R} \to \mathsf{Vect}_k$, i.e. functors that are pointwise finite dimensional and have locally finite sets of critical parameters (equivalently, barcodes locally finite). We write Vect_k for vector spaces over k. Filtered chain complexes will be denoted $F = (F^t)_{t\in\mathbb{R}}$, each F^t a bounded-below chain complex of finite-dimensional k-vector spaces, with nondecreasing inclusions in t and locally finite changes.

Interleaving metric and truncation. We use the interleaving metric d_{int} on constructible 1D persistence modules; in this setting d_{int} coincides with the bottleneck distance on barcodes. We assume the existence of: (i) an exact, functorial, filtration-truncation endofunctor $\mathbb{T} : \mathsf{Pers}_k^{\mathrm{c}} \to \mathsf{Pers}_k^{\mathrm{c}}$ for each $\tau > 0$ ("collapse" or "truncation" at scale τ), and (ii) a chain-level operation C_τ producing a filtered complex $C_\tau F$ whose persistent homology is $\mathbb{T}$ of that of F. We assume $\mathbb{T}$ is 1-Lipschitz with respect to d_{int} and exact on short exact sequences of persistence modules. (These properties hold for many standard truncation/collapse operators; we take them here as standing assumptions to avoid external cross-references.)

Windowing and measurability. Fix a MECE, right-open windowing of $[0,\infty)$ into intervals $\mathcal{W} = \{[w_j, w_{j+1})\}_{j\in J}$ with $0 = w_0 < w_1 < \cdots$, locally finite on bounded ranges. For a filtered complex F and degree i, the Betti function $\beta_i(F;t) := \dim_k H_i(F^t)$ is piecewise constant with locally finite jumps; hence it is measurable and integrable on bounded windows.

Energy functional. For a window size $\sigma > 0$ and degree i,

$$\mathrm{PE}_i^{\le\sigma}(F) \;:=\; \int_0^{\sigma} \beta_i(F;t)\,dt.$$

and similarly $\mathrm{PE}_i^{\le\sigma}(M)$ for a persistence module M, via any filtered presentation. All integrals are Lebesgue integrals; by constructibility on bounded windows, they reduce to finite sums.

Pointwise tensor. On the filtered-complex side and on persistence modules we use the pointwise tensor:

$$(M\otimes N)(t) := M(t)\otimes_k N(t), \qquad (F\otimes G)^t := F^t\otimes_k G^t.$$

All tensor products are over k.

[Spec] scope. Results labeled [Spec] are reproducible windowed contracts that depend only on the hypotheses stated herein. They interoperate with budget accounting (Section M.5) and with τ-sweeps (Section M.4). No assertion of strict monoidality is made beyond the lax statements we prove.

M.1. Hypotheses and the laxator

We work under the following hypotheses.

(M1) Exact, constructible tensor. The pointwise tensor $\otimes$ is exact and biadditive, and preserves constructibility on bounded windows: for any $M, N \in \mathsf{Pers}_k^{\mathrm{c}}$, the set of critical parameters of $M\otimes N$ on a bounded window is locally finite, and $\dim_k(M(t)\otimes_k N(t)) < \infty$. Likewise for filtered complexes F, G.

(M2) Künneth over a field (pointwise). For each $t\in\mathbb{R}$ and each $i\in\mathbb{Z}$,

$$H_i\big((F\otimes G)^t\big) \;\cong\; \bigoplus_{p+q=i} H_p(F^t)\otimes_k H_q(G^t),$$

naturally in (F,G) and t. Over a field, Tor-terms vanish.

(M3) Lax compatibility for collapse. There is a natural transformation (the laxator) in the homotopy setting (up to filtered quasi-isomorphism)

$$\lambda_{\tau,F,G} :\; C_\tau(F\otimes G) \;\Longrightarrow\; C_\tau F \;\otimes\; C_\tau G,$$

natural in F, G and τ.

We will occasionally strengthen (M3) to:

(**M3⁺**) For all $t\in\mathbb{R}$ and all degrees i, the induced map in homology

$$H_i\big(\lambda^t_{\tau,F,G}\big) :\; H_i\big((C_\tau(F\otimes G))^t\big) \;\longrightarrow\; H_i\big((C_\tau F\otimes C_\tau G)^t\big)$$

is a monomorphism (equivalently, ranks do not decrease under λ at each (t,i)).

Remark H.18 (Intervals under tensor)**.** For interval modules over k,

$$k_{[a,b)} \otimes k_{[c,d)} \;\cong n \begin{cases} k_{[\max\{a,c\},\ \min\{b,d\})}, & \text{if } [a,b) \cap [c,d) eq \emptyset, \\ 0, & \text{otherwise.} \end{cases}$$

Thus tensor intersects lifespans; at the barcode level this is the Künneth rule over a field.

Remark H.19 (Preservation of constructibility)**.** Since on any bounded window only finitely many bars of M and N meet, the critical set of $M \otimes N$ is contained in the finite union of the critical sets of M and N; hence it is locally finite. Pointwise finite dimensionality is preserved because $\dim_k(M(t) \otimes_k N(t)) = \dim_k M(t) \cdot \dim_k N(t) < \infty$.

Remark H.20 (Monotonicity scope)**.** Tensor is neither purely inclusion-type nor deletion-type in general. We therefore confine monotonicity statements to windowed energy upper bounds and to the (M3$^+$) regime, where the laxator is pointwise monomorphic in homology.

M.2. Windowed energy via overlap integrals

Definition H.21 (Betti function and windowed energy)**.** For a filtered complex F and degree i, set $\beta_i(F;t) := \dim_k H_i(F^t)$ for $t \in \mathbb{R}$. For $\sigma > 0$ define the clipped Betti integral (windowed energy)

$$\mathrm{PE}_i^{\le \sigma}(F) \;:=\; \int_0^{\sigma} \beta_i(F;t)\, dt.$$

For a persistence module M we define $\beta_i(M;t)$ and $\mathrm{PE}_i^{\le \sigma}(M)$ analogously, using any filtered presentation (well-defined by exactness).

Theorem H.22 (Convolution-type upper bound; equality under Künneth)**.** Assume *(M1)–(M2)*. Then for all filtered complexes F, G, all degrees i, and all windows $\sigma > 0$,

$$\mathrm{PE}_i^{\le \sigma}(F \otimes G) \;\le\; \sum_{p+q=i} \int_0^{\sigma} \beta_p(F;t)\, \beta_q(G;t)\, dt,$$

where the sum $\sum_{p+q=i}$ is finite due to bounded homological range. Consequently,

$$\mathrm{PE}_i^{\le \sigma}(F \otimes G) \;\le\; \sum_{p+q=i} \Big(\sup_{t \in [0,\sigma]} \beta_q(G;t) \Big)\, \mathrm{PE}_p^{\le \sigma}(F),$$

and symmetrically with $(F,p) \leftrightarrow (G,q)$. If, moreover, the pointwise Künneth isomorphism of *(M2)* holds, then equality holds:

$$\mathrm{PE}_i^{\le \sigma}(F \otimes G) \;=\; \sum_{p+q=i} \int_0^{\sigma} \beta_p(F;t)\, \beta_q(G;t)\, dt.$$

Proof. By (M2), for each t the homology of $(F \otimes G)^t$ decomposes as $\bigoplus_{p+q=i} H_p(F^t) \otimes H_q(G^t)$, whence

$$\beta_i(F \otimes G;t) \;\le\; \sum_{p+q=i} \beta_p(F;t)\, \beta_q(G;t),$$

with equality when Künneth isomorphism holds. The functions involved are measurable and piecewise constant on bounded windows. Integrate both sides over $[0,\sigma]$ and apply Tonelli/Fubini; finiteness follows from constructibility. The second inequality is obtained by extracting $\sup_{t\in[0,\sigma]} \beta_q(G;t)$ from the integral. □

Remark H.23 (Barcode-level interpretation)**.** In the interval-decomposable case (1D constructible over a field), $\beta_p(F;t)$ counts the number of p-bars alive at t. The integrand $\beta_p(F;t)\,\beta_q(G;t)$ counts ordered pairs of alive bars; the overlap integral sums the lengths of pairwise intersections, matching tensor-as-intersection (Remark H.18).

M.3. Collapse, lax monoidality, and energy dominance

Proposition H.24 (Collapsed convolution bound)**.** *Assume (M1)–(M3). For any $\tau, \sigma > 0$ and all degrees i,*

$$\mathrm{PE}_i^{\leq\sigma}\big(C_\tau F \otimes C_\tau G\big) \;\leq\; \sum_{p+q=i} \int_0^\sigma \beta_p(C_\tau F;t)\,\beta_q(C_\tau G;t)\,dt.$$

Equality holds when the pointwise Künneth isomorphism holds for $(C_\tau F, C_\tau G)$.

Proof. Apply Theorem H.22 to the filtered complexes $C_\tau F$ and $C_\tau G$. ◻

Theorem H.25 (Energy dominance under a monomorphic laxator)**.** *Assume (M1)–(M3) and (M3$^+$). Then for all $\tau, \sigma > 0$ and all degrees i,*

$$\mathrm{PE}_i^{\leq\sigma}\big(C_\tau(F\otimes G)\big) \;\leq\; \mathrm{PE}_i^{\leq\sigma}\big(C_\tau F \otimes C_\tau G\big).$$

Consequently, combining with Proposition H.24,

$$\mathrm{PE}_i^{\leq\sigma}\big(C_\tau(F\otimes G)\big) \;\leq\; \sum_{p+q=i} \int_0^\sigma \beta_p(C_\tau F;t)\,\beta_q(C_\tau G;t)\,dt.$$

Proof. By (M3$^+$), for each t the map $H_i(\lambda^t_{\tau,F,G})$ is injective, hence

$$\beta_i\big(C_\tau(F\otimes G);t\big) \;\leq\; \beta_i\big(C_\tau F\otimes C_\tau G;t\big)$$

for all t. Integrate over $[0,\sigma]$. ◻

Remark H.26 (Necessity of monomorphy)**.** Without (M3$^+$), $\lambda_{\tau,F,G}$ could a priori decrease ranks at specific (t,i), and one cannot guarantee PE-nonincrease along λ. In that regime, only the (symmetric) convolution bounds of Theorem H.22 and Proposition H.24 are safe.

M.4. Interaction with τ-sweeps and stability bands

Definition H.27 (τ-sweep and stability band)**.** For a fixed degree i and window $[0,\sigma]$, a τ-sweep probes $\tau > 0$ to identify stability bands where the functor $\mathbb{T}$ is locally constant in τ (e.g. no bars of length in $(\tau-\varepsilon, \tau+\varepsilon)$ intersect $[0,\sigma]$). Within a stability band, small τ-variations do not alter $\beta_i(\mathbb{T}M;t)$ on $[0,\sigma]$.

Proposition H.28 (Robustness within stability bands)**.** *Fix i and $\sigma > 0$. Within a stability band for τ, the inequality*

$$\mathrm{PE}_i^{\leq\sigma}\big(C_\tau(F\otimes G)\big) \;\leq\; \mathrm{PE}_i^{\leq\sigma}\big(C_\tau F \otimes C_\tau G\big)$$

of Theorem H.25 is robust under small perturbations of τ: both sides remain constant in τ on the band.

Proof. On a stability band, $\beta_i(C_\tau(\cdot);t)$ is locally constant in τ for all $t \in [0,\sigma]$, hence the integrals defining the PE functionals are constant. ◻

M.5. Budget integration and quantitative gaps [Spec]

In operational pipelines, one may wish to track commutation defects and measurement tolerances.

Definition H.29 (Budget ledger)**.** Fix a window $[0, \sigma]$. A budget ledger records nonnegative error terms δ_ℓ (from, e.g., numerical integration tolerance, discretization, or measured interleaving gaps) and aggregates $\Sigma\delta := \sum_\ell \delta_\ell$. Downstream 1-Lipschitz operations (e.g. $\mathbb{T}$, homology, degree projections) do not increase $\Sigma\delta$.

Proposition H.30 (Integrating a measured laxity gap)**.** Suppose a metric comparison is available at the persistence layer:

$$d_{\mathrm{int}}\Big(\mathbb{T}\mathbf{P}_i\big(C_\tau(F \otimes G)\big),\ \mathbb{T}\mathbf{P}_i\big(C_\tau F \otimes C_\tau G\big)\Big) \ \le\ \delta_{\mathrm{lax}}$$

for some degree i and window $[0, \sigma]$. Then δ_{lax} may be recorded in the budget ledger for $[0, \sigma]$. Subsequent 1-Lipschitz processing preserves this bound.

Remark H.31 (When no metric control is available)**.** In the absence of such metric control, rely on the energy inequalities (Theorem H.22, Theorem H.25) and budget only measurement and A/B residuals arising from the pipeline.

M.6. Edge cases and pitfalls

- Non-exact "tensor" surrogates. If the operation used in place of $\otimes$ is not exact and biadditive, Künneth may fail and the convolution bound becomes invalid.
- Unverified Künneth. If (M2) is not verified in the context at hand, use the inequality form of Theorem H.22; do not claim equality.
- Missing monomorphy. Without (M3$^+$), the comparison $C_\tau(F \otimes G) \to C_\tau F \otimes C_\tau G$ need not be energy-nonincreasing.
- Window mismatch. All measurements (Betti integrals, laxator rank checks) must use the same MECE, right-open windows and the same τ; otherwise splicing across windows may fail.

M.7. Worked examples and tests

Example H.32 (Single-interval bars)**.** Let $\mathbf{P}_p(F) = k_{[a,b)}$, $\mathbf{P}_q(G) = k_{[c,d)}$ be the only nontrivial bars (other degrees vanish). Then by Remark H.18,

$$\mathbf{P}_{p+q}(F \otimes G) \ \cong\ k_{[\max\{a,c\},\, \min\{b,d\})},$$

and for any $\sigma > 0$,

$$\mathrm{PE}^{\le\sigma}_{p+q}(F \otimes G) \ =\ \lambda\Big([\max\{a, c\},\ \min\{b, d\}) \cap [0, \sigma]\Big),$$

where $\lambda(\cdot)$ denotes Lebesgue measure (length).

Example H.33 (Collapsed intervals)**.** If $b - a \le \tau$ or $d - c \le \tau$, then $C_\tau F$ or $C_\tau G$ kills the corresponding bar. The right-hand side of the convolution bound reduces accordingly; if both bars are killed, then under (M3$^+$) the left-hand side $\mathrm{PE}^{\le\sigma}_{p+q}(C_\tau(F \otimes G))$ must vanish.

Proposition H.34 (Test model for (M3$^+$): finite direct sums of intervals)**.** Let $F \simeq \bigoplus_r k_{I_r}[-p_r]$ and $G \simeq \bigoplus_s k_{J_s}[-q_s]$ be finite direct sums of interval modules (shifted in homological degree). Suppose C_τ removes all summands with length $\leq \tau$ and acts as the identity on others. Define $\lambda_{\tau,F,G}$ on summands by the canonical inclusions

$$C_\tau(k_{I_r} \otimes k_{J_s}) \hookrightarrow C_\tau k_{I_r} \otimes C_\tau k_{J_s},$$

extended additively. Then (M3$^+$) holds, and for all $i, \sigma > 0$,

$$\mathrm{PE}_i^{\leq\sigma}\big(C_\tau(F \otimes G)\big) \leq \mathrm{PE}_i^{\leq\sigma}\big(C_\tau F \otimes C_\tau G\big).$$

Proof. On intervals, tensor is intersection; C_τ either kills short intervals or leaves them unchanged. If either factor is killed, the target summand on the right is 0, and the source summand on the left is also 0 (intersection with a killed factor is empty), so the induced map is injective. If both factors survive, C_τ acts as identity and the inclusion is literal. Additivity preserves injectivity. Energy nonincrease follows by integrating ranks of monomorphisms degreewise. □

M.8. Formal underpinnings: constructibility, measurability, exactness

For completeness we record the basic facts used throughout.

Lemma H.35 (Constructibility preserved by tensor)**.** Assume (M1). If $M, N \in \mathsf{Pers}_k^{\mathrm{c}}$, then $M \otimes N \in \mathsf{Pers}_k^{\mathrm{c}}$. If F, G are constructible filtered complexes, then $F \otimes G$ is constructible.

Proof. See Remark H.19. For filtered complexes, apply the same argument degreewise; bounded homological range on windows is preserved by tensor over a field. □

Lemma H.36 (Measurability and finiteness of energy)**.** For constructible F, each $\beta_i(F; \cdot)$ is piecewise constant with locally finite jumps. Hence $\mathrm{PE}_i^{\leq\sigma}(F) < \infty$ for all $\sigma > 0$.

Proof. Constructibility implies only finitely many sublevel changes on bounded windows. The Betti function is a finite sum of characteristic functions of intervals; integrability follows. □

Lemma H.37 (Exactness and 1-Lipschitzness of $\mathbb{T}$)**.** Assume $\mathbb{T}$ is exact on short exact sequences and 1-Lipschitz for d_{int}. Then for any morphism $f : M \to N$ in $\mathsf{Pers}_k^{\mathrm{c}}$,

$$d_{\mathrm{int}}(\mathbb{T}M, \mathbb{T}N) \leq d_{\mathrm{int}}(M, N).$$

Proof. This is a standing property of $\mathbb{T}$. In barcode terms, $\mathbb{T}$ is realized by interval truncation which is nonexpansive for bottleneck distance; exactness follows from functoriality and additivity on intervals. □

M.9. Operational checklist (windowed, reproducible) [Spec]

For each experiment window $[0, \sigma]$:

- Record the tensor context: pairs (F, G), monitored degrees i, windows $[0, \sigma]$, and collapse thresholds τ.
- State whether Künneth (M2) is assumed/verified on the chosen windows; if not, use the $\leq$ form of Theorem H.22.

- State whether (M3) and optionally (M3$^+$) are assumed/verified on the chosen windows; e.g., check monomorphy of $H_i(\lambda^t_{\tau,F,G})$ at sampled t.
- Fix numerical tolerances for spectra/Betti integrations and any clipping thresholds; enter them in the budget ledger as δ-terms.
- If a measured laxity gap δ_{lax} is available (Proposition H.30), record it and include it in $\Sigma\delta$.

M.10. Summary of contracts [Spec]

Under exact pointwise tensor and a collapse-compatible laxator $C_\tau(F \otimes G) \Rightarrow C_\tau F \otimes C_\tau G$, the clipped Betti integral of a tensor admits robust, windowed upper bounds expressed by overlap integrals of Betti curves (Theorem H.22, Proposition H.24). With the additional monomorphy hypothesis (M3$^+$), $C_\tau(F \otimes G)$ is energy-dominated by $C_\tau F \otimes C_\tau G$ on each window (Theorem H.25). These contracts are stable under τ-sweeps within stability bands (Proposition H.28), are compatible with budget accounting (Proposition H.30), and suffice for operational, windowed use without further supplementation.

M.11. Formalization blueprint (Lean/Coq) [Spec]

A minimal API includes:

- An exact, biadditive tensor on $\mathsf{Pers}^{\mathrm{c}}_k$ preserving constructibility; a lemma $\mathrm{kun\_neth}$ yielding
$$\beta_i(F \otimes G; t) = \sum_{p+q=i} \beta_p(F;t)\,\beta_q(G;t)$$
over a field.
- A natural transformation $\mathrm{laxator}$ $\lambda_{\tau,F,G} : C_\tau(F \otimes G) \Rightarrow C_\tau F \otimes C_\tau G$.
- An optional hypothesis $\mathrm{laxator\_mono}$ encoding (M3$^+$) degreewise in homology.
- Overlap-integral bounds $\mathrm{PE\_conv\_bound}$ and $\mathrm{PE\_collapsed\_bound}$ corresponding to Theorem H.22 and Proposition H.24, with window parameters explicit.
- Hooks to record optional metric gaps as δ-terms for a pipeline budget.

Appendix N. Projection Formula and Base Change [Spec + Windowed Protocol + Budget Integration]

Standing conventions. We work over a coefficient *field* Λ (e.g. a base field k or, at [Spec]-level, a Novikov field), and all statements below are phrased uniformly for Λ. All persistence modules are constructible (locally finite on bounded windows). Filtered (co)limits are computed objectwise in $[\mathbb{R}, \mathsf{Vect}_\Lambda]$ under the scope policy of Appendix A, Remark .16, and then (when stated) returned to the constructible range. The interleaving metric d_{int} (= bottleneck in the constructible 1D setting) is used throughout. Truncation $\mathbf{T}_\tau$ is exact and 1-Lipschitz (Appendix A, Proposition .20); we keep the filtered-complex functor C_τ *up to f.q.i.* and write $\mathbf{P}_i$ for degree–i persistence. Global conventions: Ext^1 is always taken against $\Lambda[0]$ (so we write $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), \Lambda) = 0$); the energy exponent satisfies $\alpha > 0$ (default $\alpha = 1$); windows are MECE and right–open (Appendix G). References to "infinite bars/generic dimension" point to Appendix D, Remark A.1. Monotonicity claims follow the global policy: deletion-type only (nonincreasing), inclusion-type merely stable/nonexpansive (Appendix E).

N.1. Hypotheses (PF/BC layer) and normalizations

We fix a class of filtered spaces and maps $f : X \to Y$ for which the usual six–functor formalism is available on $D^b_c(\mathrm{Shv}_\Lambda(-))$ (bounded derived category of constructible Λ–sheaves), and adopt:

(N0) **Coefficients.** Λ is a field; all objects have *finite Tor-dimension* (no Tor corrections).

(N1) **Finiteness/constructibility.** All sheaves are constructible; the standard t-structure is used; per-object (co)homology is finite dimensional.

(N2) **Proper/smooth hypotheses.** Projection formula (PF) and base change (BC) are taken under the usual hypotheses:

- PF: for f *proper*, $Rf_*(A \otimes^{\mathbf{L}} f^*B) \simeq Rf_*A \otimes^{\mathbf{L}} B$.
- BC: for a Cartesian square with f proper (or smooth with the appropriate $f^!$ variant), $Lg^*Rf_*A \simeq Rf'_*Lg'^*A$.

(N3) **Degree normalization and objectwise evaluation.** We use the *cohomological* convention on D^b_c and evaluate realizations *objectwise in t*:

$$\mathcal{R}(F)^t \cong \mathcal{R}(F^t), \qquad \mathbf{P}_i(F)(t) \;\cong\; H_i(F^t) \;\cong\; H^{-i}\big(\mathcal{R}(F^t)\big).$$

Hence $\mathbf{P}_i$ reads off the $(-i)$-th cohomology sheaf along the filtration. Any geometric shift from $f^!$ (smooth case) is *absorbed* by this bookkeeping.

(N4) **Tensor.** The tensor is pointwise in t: $(A \otimes^{\mathbf{L}} B)^t \cong A^t \otimes^{\mathbf{L}} B^t$, exact over the field Λ; constructibility is preserved (Appendix H justifies Tonelli on bounded windows).

Remark I.1 (Scope and return to constructible)**.** All PF/BC comparisons below are formed in the derived category, computed objectwise in t, and then passed to the persistence layer via $\mathbf{P}_i$. Any filtered colimit is taken in $[\mathbb{R}, \mathsf{Vect}_\Lambda]$ under Appendix A's scope policy and *returned to* $\mathsf{Pers}^{\mathrm{ft}}_\Lambda$ (by verification of constructibility or by applying $\mathbf{T}_\tau$).

N.2. Projection formula / base change at the persistence layer

Let $f : X \to Y$ and a Cartesian square

$$\begin{array}{ccc} X' & \xrightarrow{g'} & X \\ {\scriptstyle f'}\downarrow & & \downarrow{\scriptstyle f} \\ Y' & \xrightarrow{g} & Y \end{array}$$

satisfy (N0)–(N2). For filtered complexes F on X and G on Y, write $\mathcal{R}(F), \mathcal{R}(G)$ for their realizations in D^b_c, computed objectwise in t.

Theorem I.2 (PF/BC transported to $\mathbf{P}_i$ and $\mathbf{T}_\tau$ [Spec])**.** Under (N0)–(N4) the following canonical isomorphisms hold, *natural in* (i, τ, f, g, F, G) and *up to f.q.i.* on the filtered–complex side; they are asserted *after truncation by* $\mathbf{T}_\tau$:

$$\text{(PF)} \quad \mathbf{T}_\tau\, \mathbf{P}_i\big(Rf_*(\mathcal{R}(F) \otimes^{\mathbf{L}} f^*\mathcal{R}(G))\big) \cong \mathbf{T}_\tau\, \mathbf{P}_i\big(Rf_*\mathcal{R}(F) \otimes^{\mathbf{L}} \mathcal{R}(G)\big)$$

$$\text{(BC)} \quad \mathbf{T}_\tau\, \mathbf{P}_i\big(Lg^*Rf_*\mathcal{R}(F)\big) \cong \mathbf{T}_\tau\, \mathbf{P}_i\big(Rf'_*Lg'^*\mathcal{R}(F)\big).$$

Proof sketch. PF and BC are canonical isomorphisms in D_c^b under (N0)–(N2). Evaluate objectwise in t (N3), identify $\mathbf{P}_i$ with $(-i)$-cohomology in t, and then apply $\mathbf{T}_\tau$. Exactness of $\mathbf{T}_\tau$ (Appendix A, Theorem .14) preserves short exact sequences induced by PF/BC on cohomology sheaves; hence isomorphisms descend to the persistence layer *after truncation*. Naturality in (f, g, F, G) follows from naturality of PF/BC; naturality in (i, τ) is clear from functoriality of $\mathbf{P}_i$ and $\mathbf{T}_\tau$. ◻

Corollary I.3 (Compatibility with collapse at the filtered-complex layer)**.** Assume in addition that C_τ realizes $\mathbf{T}_\tau$ after applying $\mathbf{P}_i$ (Appendix B) up to f.q.i. Then the PF/BC isomorphisms of Theorem I.2 hold with $\mathcal{R}(C_\tau F)$ in place of $\mathbf{T}_\tau\,\mathbf{P}_i(\mathcal{R}(F))$, *after truncation by* $\mathbf{T}_\tau$ *and up to f.q.i. on filtered complexes*.

Remark I.4 (What is *not* claimed)**.** We do *not* assert global Lipschitz control for PF/BC beyond the 1-Lipschitz behavior of $\mathbf{T}_\tau$ (Appendix A) and any additional [Spec] commutation controls (Appendix L). PF/BC are exact identities at the sheaf level; any persistence–level discrepancy indicates violation of hypotheses or implementation drift and must be logged as such (see §N.3).

N.3. Windowed protocol and reproducible audit

All PF/BC audits are *windowed*. The mandatory comparison order is:

$$\boxed{\text{for each } t \implies \text{apply } \mathbf{P}_i \implies \text{apply } \mathbf{T}_\tau \implies \text{compare in } \mathsf{Pers}^{\mathrm{ft}}_\Lambda}.$$

Use the *same* MECE, right–open windows and the *same* τ as the rest of the run (Appendix G). In run.yaml, record per window:

- the PF/BC hypothesis set used (proper/smooth, finite Tor, degree normalization);
- the functors and objects compared (e.g. $Rf_*(\mathcal{R}(F) \otimes f^*\mathcal{R}(G))$ vs. $Rf_*\mathcal{R}(F) \otimes \mathcal{R}(G)$);
- the verdict (isomorphism detected) and, if any numerical/non–ideal drift is observed *after truncation*, its breakdown into δ^{disc} and δ^{meas} together with tolerance(s);
- the window and τ used for the comparison, matching those used by B–Gate$^+$ and the δ–ledger (Appendix G).

When PF/BC hypotheses are satisfied, the algebraic contribution δ^{alg} is 0 by design; only discretization/measurement components may be nonzero.

N.4. Budget integration and window pasting

PF/BC isomorphisms themselves contribute $\delta^{\mathrm{alg}} = 0$ when the hypotheses hold. However, in a *pipeline* that also includes Mirror/Transfer steps and non–nested reflectors (Appendix L/K), the δ–budget on a window W is:

$$\Sigma\delta_W(i) \;=\; \sum_{\text{Mirror–Collapse steps}} \delta_j(i, \tau_j) \;+\; \sum_{\text{A/B fails}} \Delta_{\mathrm{comm}}(M; A, B) \;+\; \sum_{\text{PF/BC}} \big(\delta^{\mathrm{disc}} + \delta^{\mathrm{meas}}\big).$$

with $\delta^{\mathrm{disc}}, \delta^{\mathrm{meas}}$ for PF/BC typically small or null. This budget is *additive*, *uniform in F* for Mirror–Collapse (Appendix L), and *non–increasing* under any subsequent 1-Lipschitz persistence post–processing (Appendix L, Theorem H.12). Window pasting then follows from Restart/Summability (Appendix J): log $\Sigma\delta_W(i)$ and gap_τ per window and verify $\mathrm{gap}_\tau > \Sigma\delta_W(i)$.

N.5. Ext–tests under change of functor / coefficients

PF/BC isomorphisms transport Ext^1–tests along canonical identifications:

Proposition I.5 (Portability of the Ext^1–test (sheaf layer))**.** Under (N0)–(N2) and Theorem I.2, for any PF/BC isomorphism $A \xrightarrow{\sim} B$ in D^b_c, there is a natural isomorphism

$$\mathrm{Ext}^1(A, \Lambda) \xrightarrow{\sim} \mathrm{Ext}^1(B, \Lambda).$$

In particular, if $\mathrm{Ext}^1(\mathcal{R}(C_\tau F), \Lambda) = 0$, then Ext^1 also vanishes for any PF/BC partner of $\mathcal{R}(C_\tau F)$.

Remark I.6 (Bridge stays one–way)**.** The one–way bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ (Appendix C) is unchanged: PF/BC *transport* the test across equivalent sheaf–theoretic descriptions. No converse or new implication is claimed.

N.6. Implementation notes and checkpoints

- **Finite windows / constructibility.** On bounded t–windows, bar events are finite (Appendix H); PF/BC are computed objectwise in t and preserved by $\mathbf{T}_\tau$.
- **Exactness bookkeeping.** Reductions to persistence use only: (i) PF/BC hold in D^b_c; (ii) $\mathbf{P}_i$ reads off $(-i)$–cohomology in t; (iii) $\mathbf{T}_\tau$ is exact (Appendix A, Theorem .14) and 1–Lipschitz (Appendix A, Proposition .20); (iv) filtered colimits respect the scope policy (Appendix A, Remark .16).
- **Proper/smooth reminder.** We use the cohomological convention, and any $f^!$–induced shifts in the smooth case are absorbed by (N3). PF/BC are invoked under the proper/smooth hypotheses (N2).
- **Window coherence.** All PF/BC audits must use the *same* windows and τ as those employed by B–Gate$^+$ and the δ–ledger (Appendix G); otherwise budget accounting and pasting (Appendix J) become invalid.

N.7. Formalization stubs (Lean/Coq) [Spec]

A minimal API (cf. Appendix F) includes:

- pf_iso: $Rf_*(A \otimes f^*B) \cong Rf_*A \otimes B$ under properness; bc_iso: $Lg^*Rf_*A \cong Rf'_*Lg'^*A$ for Cartesian squares (and smooth variants with $f^!$);
- to_pers: objectwise evaluation in t plus $\mathbf{P}_i$ extraction and $\mathbf{T}_\tau$ application to transport PF/BC to $\mathsf{Pers}^{\mathrm{ft}}_\Lambda$;
- pfbc_pers_nat: naturality of these isomorphisms in (i, τ, f, g, F, G);
- hooks to log any residual numeric slack (post–truncation) as δ^{disc}, δ^{meas} for the window budget (Appendix L).

Summary. Under the standard PF/BC hypotheses over a field Λ, projection formula and base change descend—via objectwise evaluation in t, $\mathbf{P}_i$, and the exact truncation $\mathbf{T}_\tau$—to canonical, *natural* isomorphisms at the persistence layer, uniformly in (i, τ, f, g, F, G). Comparisons must follow the windowed protocol "for each $t \to$ persistence $\to$ collapse $\to$ compare," with the *same* windows and τ as the rest of the run; any residual numerical drift (post–truncation) is accounted for in the δ–ledger and aggregated in the pipeline budget. These audits integrate seamlessly with Mirror/Transfer commutation (Appendix L), multi–axis reflectors (Appendix K), and Restart/Summability pasting (Appendix J), while keeping the one–way bridge $\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ intact (Appendix C).

Appendix O. Fukaya Realization & Stability [Spec + Permitted Ops + δ-Ledger + B-Gate^{+}] (Reinforced)

Standing conventions. We work over a coefficient *field* Λ (e.g. a ground field k or a Novikov field). All persistence modules are constructible (locally finite on bounded windows). Filtered (co)limits are computed objectwise in $[\mathbb{R}, \mathsf{Vect}_\Lambda]$ under the scope policy of Appendix A, Remark .16, and then (when stated) returned to the constructible range. The interleaving metric d_{int} (which agrees with the bottleneck distance in the constructible 1D setting) is used throughout. Truncation $\mathbf{T}_\tau$ is exact and 1-Lipschitz (Appendix A, Proposition .20). Global conventions: Ext^1 is always taken against $\Lambda[0]$; the energy exponent satisfies $\alpha > 0$ (default $\alpha = 1$). References to "generic fiber dimension / infinite bars" point to Appendix D, Remark A.1. Monotonicity claims follow the global policy: *deletion-type only* (nonincreasing), inclusion-type merely stable/nonexpansive (Appendix E). Windows are MECE and right–open; bars are half-open $[b, d)$ (Appendix H/G).

Remark J.1 (Right–open windows and half-open bars)**.** Right–open windows and half-open $[b, d)$ bars enforce MECE coverage and avoid double-counting at endpoints; events at a right boundary are attributed to the subsequent window.

Definition J.2 (Filtered chain model and persistence)**.** A filtered chain complex over Λ means a chain complex $C_\bullet$ equipped with an exhaustive, increasing filtration $\{F^t C_\bullet\}_{t\in\mathbb{R}}$ by subcomplexes such that the differential preserves the filtration and continuation data act by filtered maps. For each homological degree i, the degree–i persistence module is $t \mapsto H_i(F^t C_\bullet) \in \mathsf{Vect}_\Lambda$. Constructibility on bounded windows means only finitely many jumps (break times) occur and vector-space ranks are finite on bounded intervals.

Definition J.3 (Deletion-type morphism)**.** A morphism $M \to N$ in Pers_Λ is *deletion-type* on a window $W \subset \mathbb{R}$ if, after restricting to W and post-composing with the truncation $\mathbf{T}_\tau$ (for any $\tau \geq 0$), the induced map can only shorten or remove existing bars and cannot create new bars on W. Equivalently, all deletion-type monotone indicators of Appendix E are nonincreasing under this morphism.

O.1. Realization functor and hypotheses

[Spec] Fix a Liouville/Weinstein sector (X, λ) with a (possibly empty) system of *stops*. Write $\mathsf{Fuk}(X; \text{stops})$ for a wrapped/exact/monotone Fukaya-type category for which Floer-theoretic chain models admit an *action filtration*. We adopt the convention that the *action value is the filtration parameter t*, increasing with larger action (so sublevel sets F^t mean action $\leq t$). We package the chain-level construction into a realization functor

$$\mathcal{F} : \text{(geometric input)} \longrightarrow \mathsf{FiltCh}(\Lambda),$$

natural in continuation data and stop operations, with degree–i persistence $\mathbf{P}_i\big(\mathcal{F}(-)\big) \in \mathsf{Pers}_\Lambda^{\text{ft}}$ (constructible on bounded windows). We assume:

(O0) **Coefficients/admissibility.** Λ is a field; in the monotone/exact cases and with admissible almost complex structures, action and index filtrations are well defined; continuation solutions have finite energy.

(O1) **Constructibility (action window).** On every bounded action window $[0, \sigma]$ the Floer complexes have finitely many generators and finitely many break times; hence $\mathbf{P}_i(\mathcal{F}(-))$ is constructible. This local finiteness equally holds with Novikov coefficients on bounded windows.

(O2) **Continuation shift bound.** Any continuation map for a homotopy of data with controlled action shift ε induces a filtered chain map whose filtration increase is $\leq \varepsilon$.

(O3) **Stop operations are deletion-type.** Adding a stop or shrinking a sector removes generators and/or increases differentials in a way that corresponds to a *deletion-type* operation at the persistence layer: in any fixed action window no new bars are created.

(O4) **Up to filtered quasi-isomorphism.** Chain models are considered up to filtered quasi-isomorphism; all claims are invariant under f.q.i. by exactness of $\mathbf{T}_\tau$ and the scope policy.

Remark J.4 (Action filtration normalization)**.** Our choice "action value equals filtration parameter" fixes the direction of filtration; monotone time reparametrizations that preserve order act by reindexings and, after normalization, give isometries in d_{int}.

O.2. Stability: continuation and stops

Theorem J.5 (Continuation 1-Lipschitz)**.** Under (O2), for any two realizations related by a continuation with action shift ε,

$$d_{\mathrm{int}}\Big(\mathbf{P}_i(\mathcal{F}_0),\ \mathbf{P}_i(\mathcal{F}_1)\Big)\ \le\ \varepsilon, d_{\mathrm{int}}\Big(\mathbf{T}_\tau\mathbf{P}_i(\mathcal{F}_0),\ \mathbf{T}_\tau\mathbf{P}_i(\mathcal{F}_1)\Big)\ \le\ \varepsilon$$

for all $i, \tau \ge 0$.

Proof sketch. A filtered chain map whose filtration increase is $\le\ \varepsilon$ commutes up to shift with the time-translation endofunctors, hence yields an ε-interleaving at the persistence level. The first inequality follows. Applying $\mathbf{T}_\tau$ preserves the bound by its 1-Lipschitz property (Appendix A, Proposition .20). □

Proposition J.6 (Deletion-type monotonicity for stops)**.** Under (O3), adding a stop or shrinking a sector induces, on any window and after $\mathbf{T}_\tau$, a deletion-type morphism: for every i and $\tau \ge 0$,

$$\mathbf{T}_\tau\, \mathbf{P}_i\big(\mathcal{F}_{\text{with stop}}\big)\ \preceq\ \mathbf{T}_\tau\, \mathbf{P}_i\big(\mathcal{F}_{\text{without stop}}\big),$$

and all deletion-type monotone indicators (Appendix E) are nonincreasing under this operation.

Remark J.7 (Inclusion-type caution)**.** Operations that enlarge the admissible region or remove stops can increase features. We only assert nonexpansivity via continuation when an explicit shift control is available; no deletion-type monotonicity is claimed in that direction.

Corollary J.8 (Compositionality of shift bounds)**.** If $\mathcal{F}_0 \to \mathcal{F}_1 \to \mathcal{F}_2$ are continuations with shifts $\varepsilon_1, \varepsilon_2$, then $d_{\mathrm{int}}(\mathbf{P}_i(\mathcal{F}_0), \mathbf{P}_i(\mathcal{F}_2)) \le \varepsilon_1 + \varepsilon_2$.

O.3. Towers, comparison map, and diagnostics

Let $F = (F_n)_{n\in I}$ be a directed system of geometric inputs (e.g. refining Hamiltonians/perturbations or nested stop systems) with colimit F_∞. Apply $\mathcal{F}$ and $\mathbf{P}_i$ to obtain a tower in $\mathsf{Pers}^{\mathrm{ft}}_\Lambda$. For $\tau \ge 0$ consider the comparison map (Appendix J)

$$\phi_{i,\tau}(F):\quad \varinjlim_n\ \mathbf{T}_\tau\big(\mathbf{P}_i(\mathcal{F}(F_n))\big)\ \longrightarrow\ \mathbf{T}_\tau\big(\mathbf{P}_i(\mathcal{F}(F_\infty))\big).$$

Definition J.9 (Sufficient tower hypotheses)**.** Assume: (S1) commutation up to vanishing error with truncation and reindexing; or (S2) no τ-accumulation of break times; or (S3) Cauchy in d_{int} with compatible structure maps. Each is made precise in Appendix D, §D.3 and ensures stability of barcodes under the colimit and truncation.

Theorem J.10 (When the comparison is an isomorphism)**.** If the tower admits continuation controls $\varepsilon_n \to 0$ with $d_{\mathrm{int}}\big(\mathbf{P}_i(\mathcal{F}(F_n)), \mathbf{P}_i(\mathcal{F}(F_\infty))\big) \leq \varepsilon_n$ and satisfies any of (S1)/(S2)/(S3), then $\phi_{i,\tau}(F)$ is an isomorphism for all $\tau \geq 0$. Consequently,

$$\mu_{i,\tau}(F) = u_{i,\tau}(F) = 0,$$

where (μ, u) denote the generic fiber dimensions of the kernel/cokernel (Appendix D, Remark A.1).

Proof sketch. Combine Theorem J.5 with the tower criteria of Appendix D, §D.3 to control the interleavings along the system and its colimit; then apply the comparison formalism of Appendix J to deduce that $\phi_{i,\tau}$ is an isomorphism. Vanishing of (μ, u) follows. □

Corollary J.11 (Grid $\Rightarrow$ continuum survival)**.** In discretization towers (mesh $h \to 0$) with certified continuation bounds $\varepsilon(h) \to 0$, any bar detected in a fixed window $[0, \tau_0]$ whose $\mathbf{T}_{\tau_0}$-clipped length exceeds $2\varepsilon(h)$ persists in the limit (Appendix I, Theorem I:F.7).

O.4. Permitted-operations table (windowed, post-collapse) and δ-ledger

All comparisons follow the protocol "for each $t \to$ persistence $\to$ collapse $\mathbf{T}_\tau \to$ compare," on MECE, right–open windows and a fixed τ (Appendix G/N). The following table summarizes permitted operations, their type, quantitative contracts (after collapse), and how to record them in the δ–ledger.

Operation	Type	Quantitative contract after $\mathbf{T}_\tau$ and ledger entry
Add stop / shrink sector	Deletion	Nonincreasing deletion indicators; no new bars on the window; ledger: $\delta^{\mathrm{alg}} = 0$, record $\delta^{\mathrm{disc}}, \delta^{\mathrm{meas}}$ if any
Continuation (tame homotopy)	Shift	$d_{\mathrm{int}} \leq \varepsilon$ (Theorem J.5); ledger: $\delta^{\mathrm{alg}} = \varepsilon$
Hamiltonian change (bounded drift)	Shift	$d_{\mathrm{int}} \leq \varepsilon$; ledger: $\delta^{\mathrm{alg}} = \varepsilon$
Almost complex structure change (tame)	Shift	$d_{\mathrm{int}} \leq \varepsilon$; ledger: $\delta^{\mathrm{alg}} = \varepsilon$
Regrading / Maslov shift	Bookkeeping	Isometry (degree reindexing); ledger: $\delta^{\mathrm{alg}} = 0$
Monotone time reparametrization	Reindex	Isometry after reindex normalization; ledger: $\delta^{\mathrm{alg}} = 0$
Mirror/Transfer post-processing	External	As audited (Appendix L); if commutation defect $\delta(i, \tau)$, ledger: $\delta^{\mathrm{alg}} = \delta(i, \tau)$
Non-nested reflectors (if used)	External	A/B test or soft-commuting fallback (Appendix K/L); ledger: $\delta^{\mathrm{alg}} = \Delta_{\mathrm{comm}}$ if fallback

Definition J.12 (δ–ledger and budgets)**.** Per window $W = [u, u')$ and degree i, define the aggregate budget

$$\Sigma\delta_W(i) \;:=\; \sum_{\text{continuations}} \varepsilon \;+\; \sum_{\text{Mirror/Transfer}} \delta(i, \tau) \;+\; \sum_{\text{A/B fails}} \Delta_{\mathrm{comm}} \;+\; \sum_{\text{audits}} (\delta^{\mathrm{disc}} + \delta^{\mathrm{meas}}),$$

where "audits" include PF/BC checks (Appendix N) and numerical tolerances (Appendix G).

O.5. B-Gate$^+$, restart, and summability (window pasting)

We adopt B-Gate$^+$ with a per-window safety margin $\mathrm{gap}_\tau(i) > 0$ computed after $\mathbf{T}_\tau$. On window W and degree i, the gate *passes* if

$$\mathrm{gap}_\tau(i) \;>\; \Sigma\delta_W(i).$$

Across consecutive windows $(W_k)_k$, assume: (i) transitions are finite compositions of deletion-type steps and ε-continuations (measured post-collapse), and (ii) Summability holds $\sum_k \Sigma\delta_{W_k}(i) < \infty$ (Appendix J). Then the Restart inequality of Appendix J yields for some $\kappa \in (0, 1]$,

$$\mathrm{gap}_{\tau_{k+1}}(i) \;\geq\; \kappa\Big(\mathrm{gap}_{\tau_k}(i) - \Sigma\delta_{W_k}(i)\Big),$$

so positivity of the margin persists along the pipeline. Per-window certificates therefore paste to a global certificate on $\bigcup_k W_k$ (Appendix J, Theorem J:G.13).

Remark J.13 (Choice of κ)**.** The constant κ accounts for uniform losses at window transitions (e.g. finite alignment overheads or reindexing coercions). Under exact commutation, one can take $\kappa = 1$.

O.6. Windowed usage and run.yaml alignment

Record in $\mathrm{run.yaml}$ per window and degree:

- the operation sequence (stops/sector changes, continuations) with quantitative parameters (ε, thresholds τ, sweep settings);
- the δ–ledger entries and their sum $\Sigma\delta_W(i)$;
- the B-Gate$^+$ safety margin $\mathrm{gap}_\tau(i)$ and pass/fail verdict;
- any external steps (Mirror/Transfer, reflectors) with A/B policy data (η, Δ_{comm}) (Appendix K/L);
- constructibility checks: upper bounds on generators and event counts (Appendix H).

All diagnostics (μ, u, comparison maps $\phi_{i,\tau}$) are computed *after* truncation $\mathbf{T}_\tau$ and logged with the same window and τ.

O.7. Failure modes and audit checklist

Failure modes (outside our scope).

- **Loss of filtration control.** Non-exact data or bubbling may invalidate (O2); interleaving bounds then fail.
- **Near-threshold accumulation.** Type IV behavior (Appendix D) may produce bar-length accumulation at τ; comparison maps need not stabilize without (S2)/(S3).
- **Inclusion-type operations.** Removing stops or enlarging sectors can increase features; only stability (not monotonicity) is claimed, and only when continuation control is present.

Audit checklist (runtime verifications).

1. Record continuation shift bounds ε and certify d_{int}-nonexpansivity (Appendix A).
2. Verify constructibility on each window $[0, \sigma]$ (finite generators/events; Appendix H) and log per-window counts (Appendix G).
3. For stop additions/sector shrinkage, mark operation as deletion-type and apply Appendix E indicators.
4. For towers, log ε_n, check (S1)/(S2)/(S3) where applicable, and compute (μ, u); $\phi_{i,\tau}$ iso $\Rightarrow$ $(\mu, u) = (0, 0)$ (Appendix J).
5. If external functors are used, run Mirror/Transfer commutation audits and A/B tests (Appendix L/K) and bookkeep residuals in $\Sigma\delta_W(i)$.

O.8. Formalization stubs (Lean/Coq) [Spec]

A minimal API (cf. Appendix F) includes:

- fukaya_realize: returns an action-filtered chain model $\mathcal{F}$ up to f.q.i., with constructibility on bounded windows (O1).
- cont_eps: encodes continuation maps with filtration increase $\leq \varepsilon$ and yields $d_{\mathrm{int}} \leq \varepsilon$ (Theorem J.5).
- stop_delete: produces deletion-type morphisms for stops/sector-shrink (Proposition J.6).
- tower_phi_iso: checks sufficient criteria so that $\phi_{i,\tau}$ is an isomorphism and $(\mu, u) = (0, 0)$ (Appendix D/J).
- Hooks for δ–ledger entries and B-Gate$^+$ checks; restart/summability contracts (Appendix J).

O.9. Examples and regression tests

Example J.14 (Exact wrapped setting with a new stop)**.** Let $X = T^*Q$ with its standard exact form and consider a wrapped setting with a stop at infinity. Adding an additional stop supported on a Legendrian subset removes Reeb chords crossing the stop. On any fixed action window, no new generators appear and differentials can only increase; hence the induced map is deletion-type (Proposition J.6). Deletion indicators (Appendix E) weakly decrease; regression test: the windowed barcode total variation does not increase after $\mathbf{T}_\tau$.

Example J.15 (Continuation bound from Hamiltonian drift)**.** Suppose H_0, H_1 are cofinal Hamiltonians with $\sup_{x,t}(H_1 - H_0) \leq \varepsilon$ on the relevant support and homotoped by a tame path. The action change along continuation solutions is bounded by ε; thus the induced filtered map has filtration increase $\leq \varepsilon$, and $d_{\mathrm{int}} \leq \varepsilon$ (Theorem J.5). Regression test: pairwise barcode bottleneck distance never exceeds the certified ε.

Example J.16 (Grid-to-continuum)**.** Discretize a time-dependent Floer datum with mesh h and certified continuation bound $\varepsilon(h) = O(h^\alpha)$. For any fixed τ_0, bars of length $> 2\varepsilon(h)$ in $\mathbf{T}_{\tau_0}$-collapse survive in the limit (Corollary J.11). Regression test: survival rates converge monotonically after accounting for $\Sigma\delta_W(i)$.

O.10. Summary

Floer-theoretic realizations with action filtration yield constructible persistence (O1). Continuation with shift ε is 1-Lipschitz at the persistence level (O2 $\Rightarrow$ Theorem J.5). Adding stops or shrinking sectors is deletion-type and hence nonincreasing for all deletion indicators (Proposition J.6). A windowed, post-collapse permitted-ops table prescribes how to assign δ–ledger entries; B-Gate$^+$ requires $\mathrm{gap}_\tau > \Sigma\delta$ per window and pastes via Restart/Summability (Appendix J). Under standard tower hypotheses (Appendix D), the comparison map $\phi_{i,\tau}$ is an isomorphism and $(\mu, u) = (0, 0)$ (Theorem J.10); grid-to-continuum survival follows (Appendix I). All items respect the MECE/right–open window policy, are evaluated "$t \to \mathbf{P}_i \to \mathbf{T}_\tau \to$compare," and integrate with the pipeline budget and A/B policy (Appendix L/K) in a reproducible run.yaml workflow (Appendix G).

Concluding Remarks and Acknowledgments

Coefficients and standing scope. Throughout, unless explicitly stated otherwise, coefficients are taken in a *field*. In the provable Core (including the bridge in Appendix C) the target of realizations is $D^{\mathrm{b}}(k\text{-mod})$. In **[Spec]** appendices that use sheaf-theoretic or Fukaya-categorical realizations we write the coefficient field as Λ. This is a notational convenience only: all persistence-layer statements remain within the constructible 1D regime over a field. We do not leave this regime. In particular, all filtered (co)limits are computed *objectwise* in $[\mathbb{R}, \mathsf{Vect}_\Lambda]$ under the scope policy (Appendix A) and returned to the constructible subcategory by verification or by applying $\mathbf{T}_\tau$.

What is proved and formalized (Core). Within constructible 1D persistence over a field, the following are established and arranged for machine-checkability and formalization.

- *Exact truncation and filtered lift.* The Serre reflector $\mathbf{T}_\tau$ deletes precisely the finite bars of length $\le \tau$, is exact, idempotent, and 1-Lipschitz for the interleaving metric; it admits a filtered lift C_τ (unique up to filtered quasi-isomorphism) with $\mathbf{P}_i(C_\tau F) \cong \mathbf{T}_\tau(\mathbf{P}_i F)$ (Appendix A/B).
- *One-way bridge $PH_1 \Rightarrow Ext^1$.* Under a t-exact realization of amplitude ≤ 1, $\mathrm{PH}_1(F) = 0$ implies $\mathrm{Ext}^1(\mathcal{R}(F), k) = 0$. No converse or global equivalence is asserted (Appendix C).
- *Tower diagnostics.* For the windowed comparison map $\phi_{i,\tau} \colon \varinjlim \mathbf{T}_\tau \mathbf{P}_i(F_n) \to \mathbf{T}_\tau \mathbf{P}_i(F_\infty)$, the generic fiber dimensions $\mu_{i,\tau} := \operatorname{gdim} \ker \phi$ and $u_{i,\tau} := \operatorname{gdim} \operatorname{coker} \phi$ detect invisible limit failures (Type IV). These enjoy subadditivity under composition, additivity under finite sums, and cofinal invariance; they vanish when $\phi_{i,\tau}$ is an isomorphism (Appendix D/J).
- *Windowed indicators and calculus.* The Betti integral equals the clipped barcode mass on $[0, \tau]$; ε-survival provides grid-to-continuum robustness without false negatives; deletion-type spectral indicators are non-increasing after collapse, while general operations are non-expansive (Appendix H/I/E).
- *Windowed pasting.* B–Gate$^+$ enforces a per-window safety margin $\mathrm{gap}_\tau > \sum \delta$; Overlap Gate glues local certificates after collapse; Restart/Summability pastes windowed certificates into global ones (Appendix J). All comparisons are performed after truncation by the protocol $t \to \mathbf{P}_i \to \mathbf{T}_\tau \to$ compare on mutually exclusive, collectively exhaustive right-open windows.

What is specified and how it is audited ([Spec]). All **[Spec]** items are *non-expansive after truncation*, shipped with explicit preconditions, and audited by windowed diagnostics (μ, u) together with δ-budgeting.

- *Mirror/Transfer commutation with collapse.* A natural 2-cell $\mathrm{Mirror} \circ C_\tau \Rightarrow C_\tau \circ \mathrm{Mirror}$ with a uniform bound $\delta(i, \tau)$ (in the interleaving metric at the persistence layer) gives δ-controlled commutation; bounds add along pipelines and do not increase under 1-Lipschitz post-processing (Appendix L).
- *Multi-axis reflectors and A/B soft-commuting.* For exact reflectors arising from hereditary Serre subcategories, order independence holds in the nested case; otherwise an A/B commutativity test with tolerance η and deterministic fallback is used. Residuals Δ_{comm} are logged as δ^{alg} (Appendix K).
- *Lax monoidal tensor/collapse.* Windowed energy admits overlap-integral bounds; in a pointwise-mono regime (M3$^+$) one obtains post-collapse dominance $\mathrm{PE}(\, C_\tau(F \otimes G)\,) \le \mathrm{PE}(\, C_\tau F \otimes C_\tau G\,)$. Optional metric gaps may be integrated into the δ-budget (Appendix M).
- *Projection formula and base change.* PF/BC are transported to the persistence layer by the windowed protocol $t \to \mathbf{P}_i \to \mathbf{T}_\tau \to$ compare. Any residual numerical slack (post-truncation) is budgeted as $\delta^{\mathrm{disc}}, \delta^{\mathrm{meas}}$ (Appendix N).
- *Fukaya realizations.* Action filtration yields constructible persistence on bounded windows; continuation is 1-Lipschitz; stop addition is deletion-type; tower stability and grid-to-continuum survival follow under standard hypotheses (Appendix O).

All **[Spec]** statements are used only under their stated hypotheses, are auditable per window via (μ, u) and the δ-ledger, and never override Core guarantees.

Operational pipeline (end-to-end). We summarize the practical path from input to audit-ready outputs.

- *Windows and after-collapse policy.* Partition time/parameter into MECE right-open windows; fix τ per window (optionally via τ-sweep and stability bands). All measurements and comparisons are made on the B-side single layer after applying $\mathbf{T}_\tau$.
- *Overlap Gate.* On overlaps, require post-collapse equivalence up to budget; verify Čech–Ext^1 acyclicity in degree 1 and stability-band conditions; then glue.
- *B–Gate$^+$.* Per window and degree i, require $\mathrm{PH}_1 = 0$, (eligible) $\mathrm{Ext}^1 = 0$, $(\mu, u) = (0, 0)$, and a safety margin $\mathrm{gap}_\tau > \sum \delta$.
- *Restart/Summability.* Across windows, enforce a restart inequality $\mathrm{gap}_{\tau_{k+1}} \geq \kappa(\mathrm{gap}_{\tau_k} - \sum \delta_k)$ for some $\kappa \in (0, 1]$ and summability $\sum_k \sum \delta_k < \infty$; paste local certificates into global ones.
- *δ-ledger.* Aggregate algorithmic (δ^{alg}: Mirror–Collapse defects $\delta(i, \tau)$, A/B residuals Δ_{comm}), discretization (δ^{disc}), measurement (δ^{meas}) per window; compare $\sum \delta$ to gap_τ.
- *Reproducibility.* Emit versioned artifacts for bars/spec/ext/phi/Lambda_len with cross-linked hashes. Fix spectral ordering (ascending) and matrix norm (operator/Frobenius). Use canonical JSON/HDF5 serialization (fixed endianness, fixed-length UTF-8 strings, creation-order stability) for bitwise re-runs (Appendix G).

Reproducibility, formalization, and tests. Appendix G specifies the manifest (run.yaml) and artifact schemas with canonical serialization and cross-linked hashes; the δ-ledger records per-window budgets, including Mirror–Collapse $\delta(i, \tau)$ and A/B Δ_{comm}. Appendix F sketches a Lean/Coq formal development for the Core: Serre localization and $\mathbf{T}_\tau$; 1-Lipschitz continuity; comparison maps and (μ, u); the bridge's edge identification; and API stubs for budget accounting, A/B policy, PF/BC transport, and lax monoidal contracts. A minimal test suite (Appendix G/12) validates stability, deletion-type monotonicity, filtered-colimit behavior, Mirror/tropical pipelines, three-layer correspondences, saturation gates, A/B soft-commuting, and Restart/Summability pasting.

Scope and guard-rails. All statements obey the following guard-rails.

- *Regime.* Constructible 1D persistence over a field; filtered (co)limits computed objectwise in $[\mathbb{R}, \mathsf{Vect}]$ and returned to $\mathsf{Pers}^{\mathrm{ft}}$ by verification or $\mathbf{T}_\tau$.
- *Collapse layer.* $\mathbf{T}_\tau$ is exact/idempotent/1-Lipschitz; shift-commutation holds; deletion-type vs. inclusion-type policy applies after truncation (monotonicity vs. non-expansiveness).
- *Single-layer judgement.* All diagnostics and budgets are evaluated after collapse on $\mathbf{T}_\tau \mathbf{P}_i$; PF/BC and Mirror/Transfer comparisons follow the windowed protocol $t \to \mathbf{P}_i \to \mathbf{T}_\tau \to$ compare.
- *Bridge and non-claims.* The bridge remains strictly one-way ($\mathrm{PH}_1 \Rightarrow \mathrm{Ext}^1$ under amplitude ≤ 1); no claim of $\mathrm{PH}_1 \Leftrightarrow \mathrm{Ext}^1$. The collapse diagnostic μ_{Collapse} is unrelated to classical Iwasawa μ.

Limitations. We do not assert metric Lipschitz control beyond $\mathbf{T}_\tau$ (and any explicitly stated δ-controls). Spectral indicators are not f.q.i.-invariants; they are evaluated under fixed policies, with deletion-type monotonicity and general non-expansiveness bounds after collapse. We do not extend the bridge, strong adjunctions, or strict monoidality beyond the proven/declared domains. PF/BC and lax monoidal statements remain windowed **[Spec]** contracts; all budgets are computed after collapse on the B-side single layer.

Future directions.

- *Quantitative links between persistence energies and spectral tails.* Under robust hypotheses and verified window policies, strengthen two-sided bounds, sensitivity analysis, and noise models.
- *Tower stability and diagnostics.* Broaden verifiable criteria for $(\mu, u) = (0, 0)$ beyond (S1)–(S3), refine Type IV precursors across scales, and automate stability-band detection on τ-sweeps.
- *Domain templates.* Provide standardized templates (arithmetic/Langlands/PDE/Fukaya) with δ-controls, PF/BC audits, A/B policies, and richer manifests for end-to-end reviews.
- *Formalization.* Expand shift/interleaving libraries, PF/BC transports, survival lemmas, pipeline budget calculus, and soft-commuting policies in proof assistants; extend test harnesses for regression and coverage metrics.

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Final note. The consistent separation between the provable Core and the **[Spec]** extensions is the organizing principle of this work: it enables safe reuse, cross-domain exploration, and reproducible evaluation while preserving mathematically verified guarantees. Windowed B-side judgement (collapse first), explicit δ-budgets, A/B soft-commuting, and Restart/Summability pasting provide an end-to-end methodology that is conservative in scope and extensible within the v16.0 guard-rails.