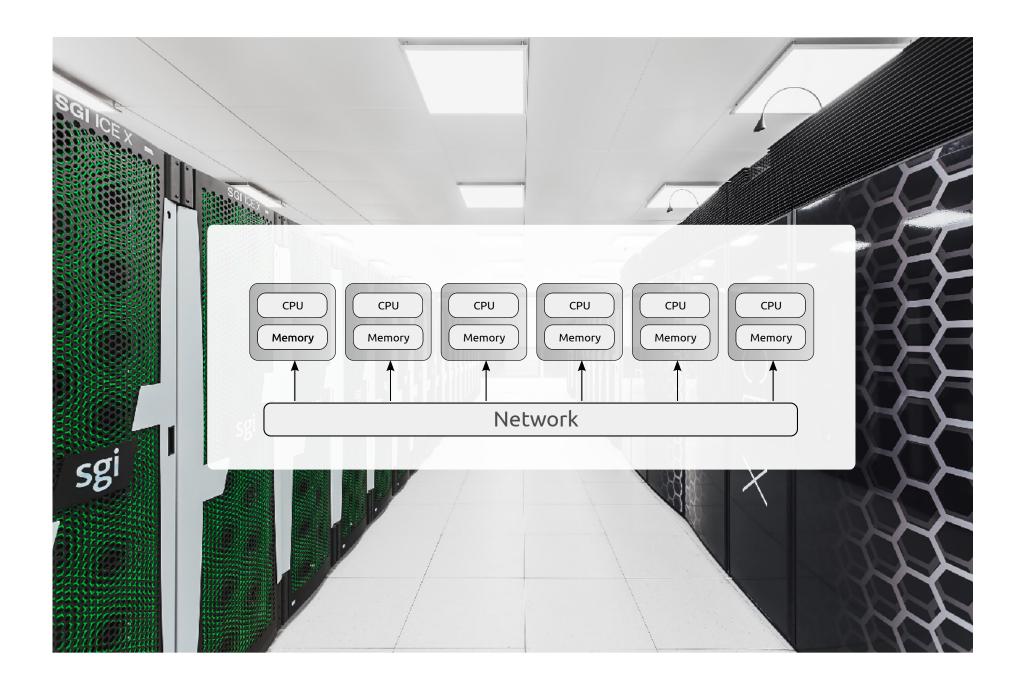
### State-Space Reduction of Non-deterministically Synchronizing Systems Applicable to Deadlock Detection in MPI

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<sup>1</sup> IT4Innovations - National Supercomputing Center, Czech Republic <sup>2</sup> FEI, Technical University of Ostrava, Czech Republic



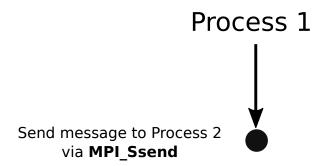


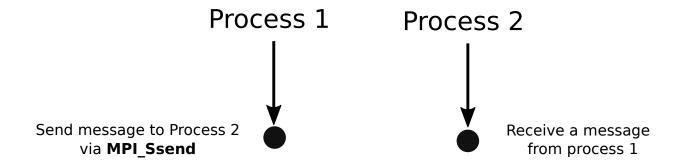
### MPI - Message Passing Interface

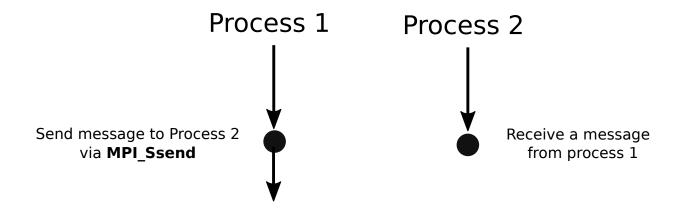
```
MPI Init(&argc, &argv);
int rank;
MPI Comm rank(MPI COMM WORLD, &rank);
int size;
MPI_Comm_size(MPI_COMM_WORLD, &size);
int value;
if (rank == 0) {
    value = 0:
    MPI Send(&value, 1, MPI INT,
            (rank + 1) % size,
            0, MPI COMM WORLD);
for (int t=0; t < 1000; t++) {
    MPI_Recv(&value, 1, MPI INT,
                MPI ANY SOURCE,
                MPI ANY TAG,
                MPI COMM WORLD,
                MPI_STATUS_IGNORE);
    MPI Send(&value, 1, MPI INT,
                (rank + 1) % size,
                0, MPI COMM WORLD);
MPI Finalize();
```

# State-Space Reduction of Non-deterministically Synchronizing Systems Applicable to Deadlock Detection in MPI

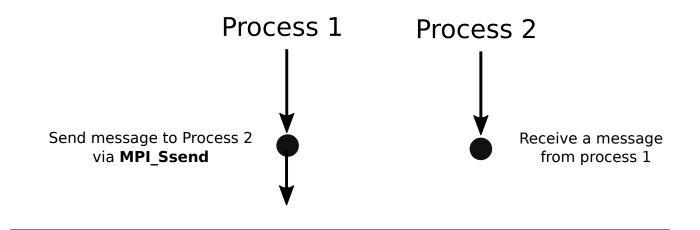
# State-Space Reduction of Non-deterministically Synchronizing Systems Applicable to Deadlock Detection in MPI





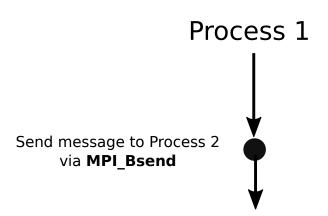


Synchronized send - waits for a matching receive



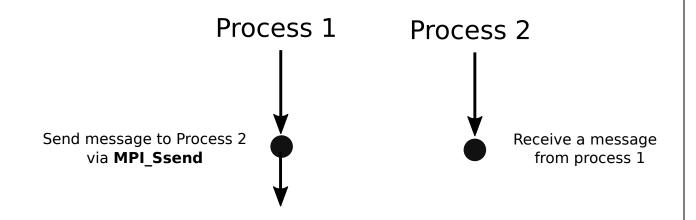
### **MPI\_Bsend**

Buffered send - do not wait for receive



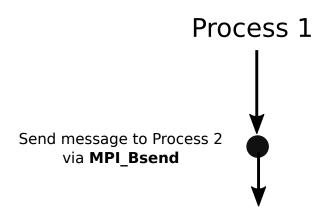
#### **MPI Ssend**

Synchronized send - waits for a matching receive



### MPI\_Bsend

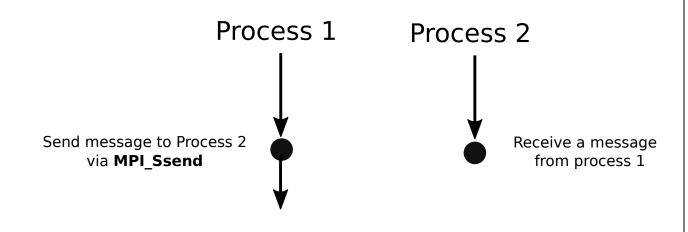
Buffered send - do not wait for receive



MPI\_Send
Standard send

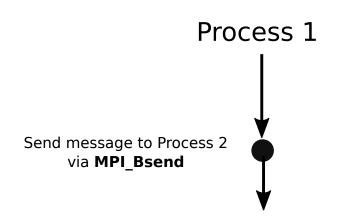
#### **MPI Ssend**

Synchronized send - waits for a matching receive

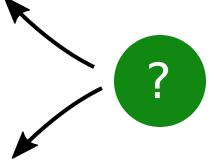


### MPI\_Bsend

Buffered send - do not wait for receive



## MPI\_Send Standard send



Nondeterministic choice

rendezvous / eager

Process 1	Process 2
MPI_Send(to=2) MPI_Recv(from=2)	MPI_Send(to=1) MPI_Recv(from=1)

## Naive approach: Analyze all MPI\_Sends as randezvous

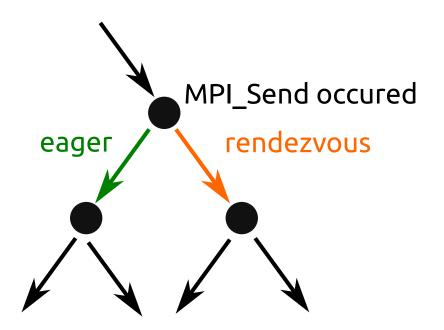
## Naive approach. Analyze all MPI\_Sends as randezvous

### Not complete

## Naive approach. Analyze all MPI\_Sends as randezvous

### Not complete

Approach 2: Encode MPI\_Send choices into state space

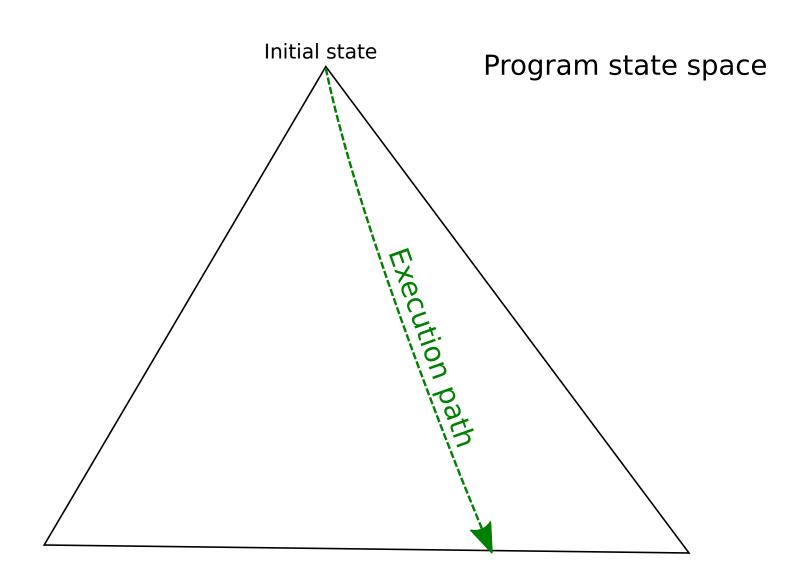


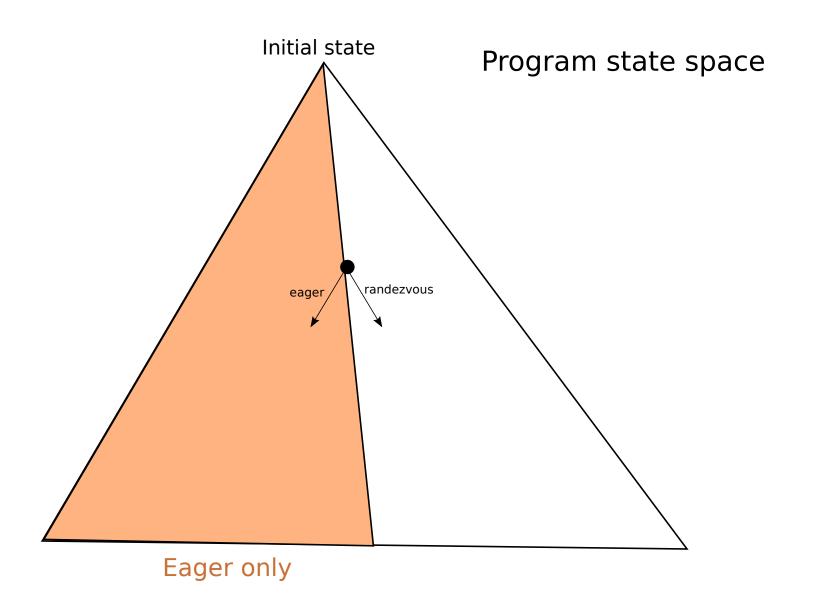


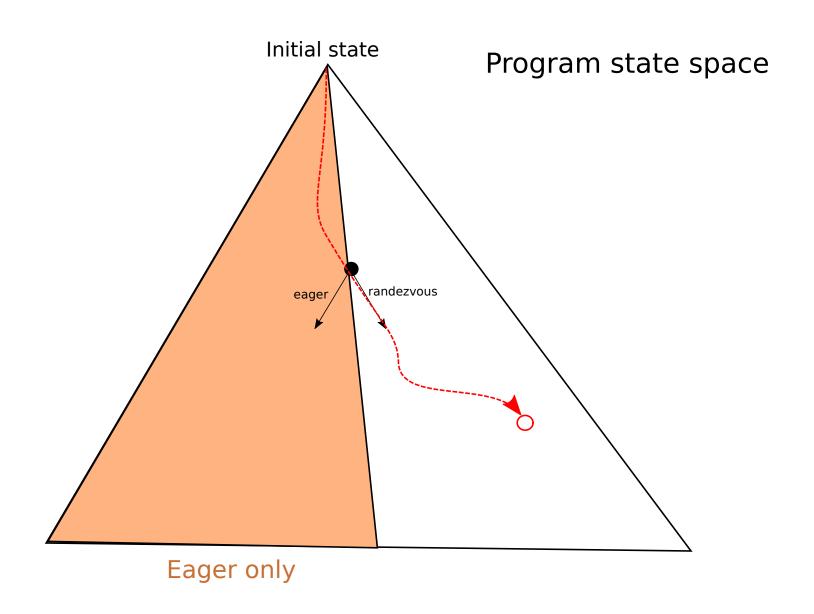
Choices inside MPI\_Sends: no new behavior (except deadlocks)

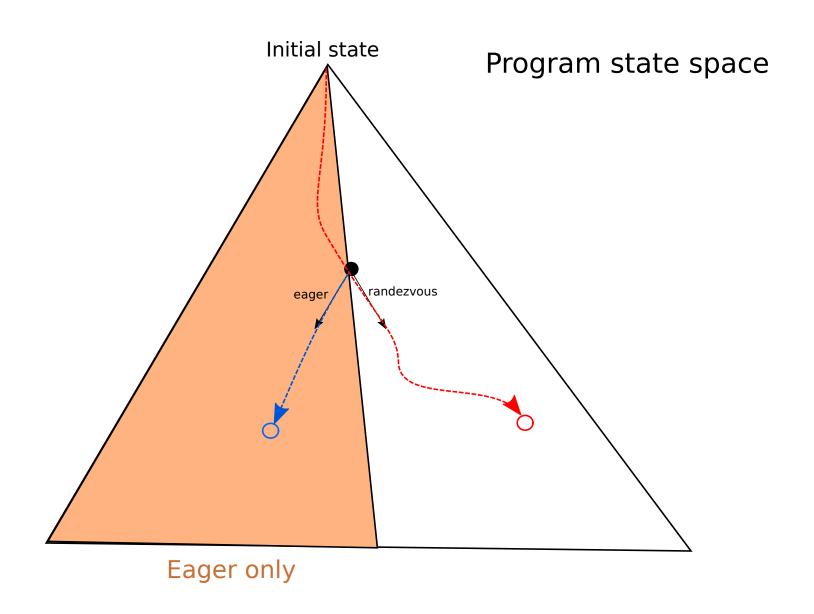
(1) Build a state space while considering only *eager* MPI\_Send

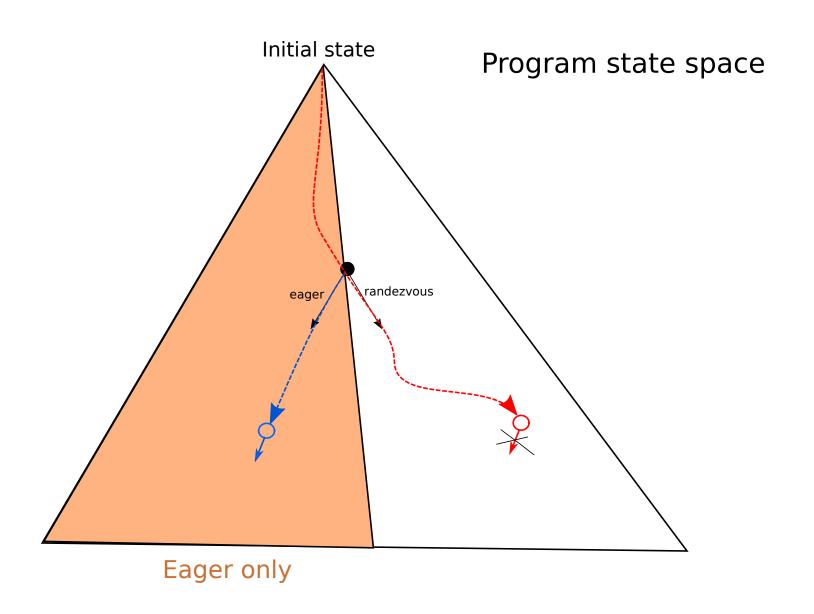
(2) Find missing deadlocks





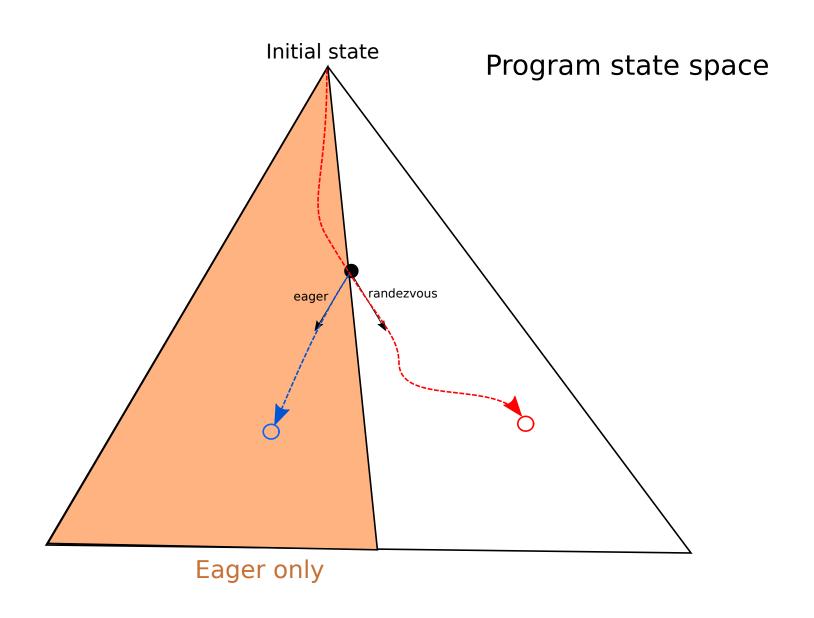


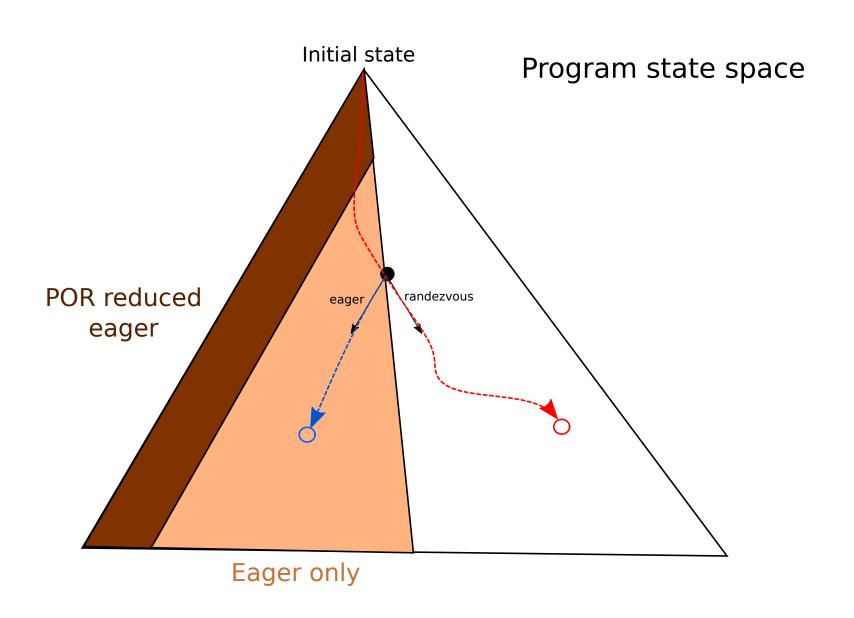


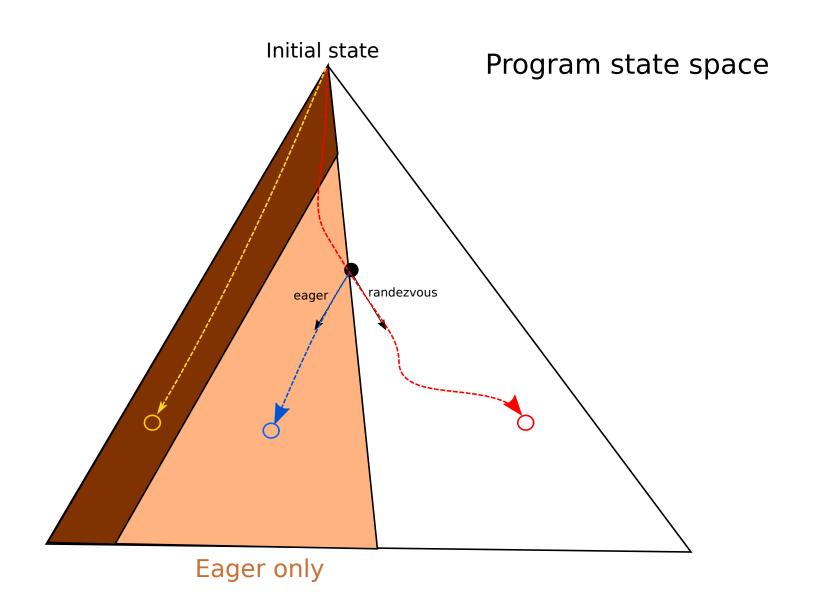


(1) Build a **parial-order reduced** state space while considering only *eager* MPI\_Send

(2) Find missing deadlocks







### Input

- **R** a system obtained by applying any\* POR method on a system **T** any POR that preserves Mazurkiewicz Traces
- I the independent set used in reduction
- C set of 'candidates' that may nondetermistically synchronize

### **Output**

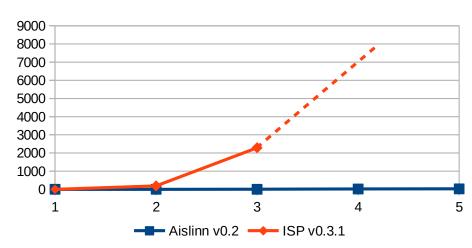
Is there a deadlock in **T** while considering synchronizations of **C**?

### http://verif.cs.vsb.cz/aislinn/

Benchmark: Workers

# of verified processes	2	3	4	5	6
Aislinn v0.2	0.7	1.1	4.3	17.6	26.2
ISP v0.3.1	2.1	191.9	2287.2	>7200	>7200

\_\_\_execution times in seconds



### Aislinn

Approach 3 compared to Approach 2 (with optimizations) **3x - 15x** speedup