Lecture-Notes-Sep-8-binary-representations

September 12, 2022

1 Lecture 2

- A breif review of Python variables
- Float point real numbers and round-off errors

2 Variable

A placeholder (a block of computer memory) that stores a value. It is called a variable because its value may change (through value assignment).

2.1 Name

Variables are referred to by their names. The names are letters (a-z and A-Z) followed by other letters or digits or the underscore . They are case sensitive, i.e., a and A are two different variables.

2.2 Value

A variable holds a single value. They can be one of the following types * a number * an interger, like 1, 2, 1000 etc. There is effectively no limit on how large the integer can be, as long as the computer has enough memory to hold it. * a float-point number, like 1.3, -10/3, etc, or in scientific notation -5.3×10^{-8} can be written as -5.3e-8 or 5.3E-8, i.e., a value followed by \mathbf{e} or \mathbf{E} followed by an exponent of 10. * a complex number, such as 1+2j (which is commonly known as 1 + 2i) where 1 is the real part and 2 is the imaginary part. * We can use complex(1, 2) to construct it, or simply 1+2J * We can use arithmatic operations such as +, -, *, /, **.

```
[1]: # example what is 5e6 equal to?
print("5e6=", 5e6)
# complex number example
# set the value of the value C to 1+2j
C = 1+2j
print("C=", C)
print("C == complex(1, 2)?", C == complex(1, 2))
# the real part and imaginary part of C can be accessed as C.real and C.imag
print("The real part of C =", C.real)
print("The imaginary part of C =", C.imag)
print("(1+2j)*(2+1j)/(3-1j)=", (1+2j)*(2+1j)/(3-1j))
```

```
5e6= 5000000.0
C= (1+2j)
```

```
C == complex(1, 2)? True
The real part of C = 1.0
The imaginary part of C = 2.0
(1+2j)*(2+1j)/(3-1j)=(-0.499999999999994+1.5j)
```

The last answer is peculiar. If we calculate by hand, we can verify that the true answer is -0.5+1.5j. Why is it a tiny bit off? This is the rounding error, that we will cover later in this lecture

- a string is a string of characters, either single quoted, like 'this is a string', or double quoted "this is another string". Single or double quoted strings are the same as long as the strong of characters are the same.
 - We can compare them, using ==, >, <, >= or <=
 - We can concatenate them using +

```
[2]: s = "it's correct"
    print("s=", s)
    # the operator == compares two values
    print("s == 'test'?", s == 'test')
    print("s + '!' =", s + '!')
```

```
s= it's correct
s == 'test'? False
s + '!' = it's correct!
```

- a boolean, a logical value being either True or False
 - we can compare them using ==, !=
 - we can perform boolean operations and, or, not

```
[3]: # boolean example
a = 2 > 1
print("2 > 1 = ", a)
print("a != True: ", a != True)
x = 3
print("x>0 and x < 10 =", x>0 and x<10)</pre>
```

```
2 > 1 = True
a != True: False
x>0 and x < 10 = True
```

• Other types, such as functions (e.g., print as we have used), lists, dictionaries or classes. We will see them later in this course.

Unlike typed languages, Python variables can hold any type of values.

```
[4]: s=12 # a number

print("s=", s)

s="a" # the same variable, now holds a string

print("s=",s)
```

```
s= 12
```

s= a

3 A Simple Example: Free Falling Object

For a free falling object, given an initial positive x_0 and an initial velocity v_0 at time 0, what is its position x(t) at time t?

$$x(t) = \frac{1}{2}gt^2 + v_0t + x_0$$

```
[5]: # Inputs:
     # the gravitational constant g
     # Python functions are organized in to modules
     # the constant g is defined the scipy model, constants subpackage
     from scipy.constants import g
     # time t
     t = 2
     # initial velocity
     v0 = 2
     # initial positive
     x0 = 1
     # positive
     x = g/2 * t**2 + v0 * t + x0
     print("t =", t, "x =", x)
     # at a different time 3
     print("t =", t, "x =", x, "what?")
     # Why does the value of x does not change?
     # how to we compute x(3)?
```

```
t = 2 x = 24.6133

t = 3 x = 24.6133 what?
```

```
[6]: x = g/2 * t**2 + v0 * t + x0
print("t = 3, x =", x, "correct!")
```

t = 3, x = 51.129925 correct!

4 Based of Numbers

4.1 Base 10

For example,

$$321 = 3 \times 10^2 + 2 \times 10 + 1$$

4.2 Other bases

We can extend this idea to any base:

A number N in base B, with digits $0 \le d_i < B$ for i = 0, ..., m, is

$$N = \sum_{i=0}^{m} d_i B^i$$

We typically write is as $(d_m d_{m-1} \dots d_2 d_1)_B$

Example How do we express $(321)_{10}$ in base 8?

- Let us start with the lowest digit, which is the remainder of 321/8?
 - the remainer is calculated by the modulus operator %

```
[7]: # 321 in base 8
N = 321
B = 8
# The last digit
d0 = N % B
print("d0 =", d0)
```

d0 = 1

• For the second last digit, please note that,

$$321 = 40 \times 8 + 1$$

- After we divide 321 by 8, we shift the digits to the right by one position
- Here the quotient 40 (without the remainder) is composed of the remaining digits.
- The second digit is thus the lowest digit of the quotient
- Thus, we can get the last digit of the remainder 40, which is the second last digit of 321

```
[8]: # second last digit
# we get the quotient first, here we use the floor division `//`, which ignoresusthe fractional part.
q = N // B
print("q=", q)

# then d1 is the last digit of q
d1 = q % B
print("d1 =", d1)
```

q= 40d1 = 0

• We then repeat to get the remaining digits one by one

```
[9]: # the thrid last digit
# we get the quotient first
q = q // B
```

```
print("q=", q)
# then d2 is the last digit of q
d2 = q % B
print("d2 =", d2)

# the fourth last digit
# we get the quotient first
q = q // B
print("q=", q)
# then d3 is the last digit of q
d3 = q % B
print("d3 =", d3)
```

q=5 d2 = 5q=0

d3 = 0

• When do we stop?

$$(321)_{10} = 5 \times 8^2 + 0 \times 8 + 1 = (501)_8$$

• This is a tedious process. But each step is repetitive. It is a good candidate for using loops, which we will learn next week!

5 Fractions

The same idea can be extended to fractions, a real number

$$R = \sum_{i=-\infty}^m d_i B^i$$

For example, in base 8,

$$(9.3)_{10} = 1 \times 8 + 1 + 2 \times 8^{-1} + 3 \times 8^{-2} + 1 \times 8^{-3} + 4 \times 8^{-4} + 6 \times 8^{-5} + \dots = (11.2\overline{3146})_8 + (1.2\overline{3146})_8 + (1.2\overline{3146}$$

A finite digit fraction in base 10 may have infinite digits in other bases

5.1 How do we convert a fraction to other bases?

• We have already dealt with the integer part. Lets just consider 0 < R < 1

$$R = \sum_{i = -\infty}^{-1} d_i B^i$$

- Lets start with a example: $(0.12345)_{10} \times 10 = (1.2345)_{10}$
 - we multiply the number by the base
 - the integer part of the product is the first digit

- the other digits are shifted to the left
- In general, multiply R by B, then

$$BR = d_{-1} + \sum_{i=-\infty}^{-1} d_i B^{i+1} = d_{-1} + \sum_{i=-\infty}^{-1} d_{i-1} B^i$$

- So, the integer part of BR is d_{-1}
- The second digit of R becomes the first digit of BR

 $d_1 = 2$ R = 0.3999999999999999999 not equal to 4!

• We repeat the above step to get the first digit of BR, which is d_{-2} , the second digit of R

d_2 = 3 R = 0.199999999999993

• We repeat this process until either R=0 or enough digits is obtained

This is a tedious process This is why we need loops to automate the repetition.

```
[12]: R = (B*R) % 1 # modulus 1 (i.e., % 1) take the fraction part.
    print("R =", R)
    d_3 = int(B*R)
    print("d_3 =", d_3)

R = (B*R) % 1 # modulus 1 (i.e., % 1) take the fraction part.
    print("R =", R)
    d_4 = int(B*R)
    print("d_4 =", d_4)

R = (B*R) % 1 # modulus 1 (i.e., % 1) take the fraction part.
    print("d_8 =", R)
```

```
d_5 = int(B*R)
print("d_5 =", d_5)
```

R = 0.599999999999943

 $d_3 = 4$

R = 0.79999999999545

 $d_4 = 6$

R = 0.399999999996362

 $d_5 = 3$

6 Binary Numbers

- Binary numbers are numbers expressed in base 2
- All modern digital computers use binary numbers internally

7 A fixed point real numbers

- A computer stores real number with a fixed number of binary bits
- How do we express a real number with these digits?
- One choice is to fix the number of digits in both the integer part and the fraction part. This is called the **fixed point format**
- For example, The so called **Q1.14** format uses 1 bit for the integer part and 14 bits for the fractional part, with 1 bit for the sign (+/-). The maximum value it can express is $2-2^{-14}$, the minimum value it can express is $-1+2^{-14}$. The precision is 2^{-14} .
- Of course it cannot express R = 0.3 precisely,

$$(0.3)_{10} = (0.010\overline{011001100110})_2$$

- In **Q1.14** format,

$$(0.3)_{10} \approx (0.01001100110011)_2$$

And the error is

$$(0.3)_{10} - (0.01001100110011)_2 = (0.000000000000000\overline{011001100110})_2$$

This may be a small error, but it may be amplified during subsequent calculations.

• The main limitation of fixed point numbers is the very limited range of numbers it can express.

8 A float point number (simply called a float)

• A float point number has one sign bit S (S=1 is negative), q exponent bits Q, and p precision bits P:

$$S \cdot Q_1 Q_2 \dots Q_q \cdot P_1 P_2 \dots P_{p-1}$$

Note that the leading bits of P is implicit (not stored, but implied from the value of Q)

- -Q = 0: the leading bit is 0
- $-Q \neq 0$: the leading bit is 1

- To express positive or negative exponents, the exponent bits as an unsigned binary integer Q is subtracted by a bias b.
- If every bit of Q is 1, i.e., $Q = (111...1)_2 = 2^q 1$, then Q does not represent an exponent, the number is either $\pm \infty$ or NaN (not a number).
- The maximum exponent Q is $2^q 2$. For none zero Q, the precision bits are always in the form of

$$P=1.P_1P_2\cdots P_{p-1}$$

The leading 1 is implicit, i.e., not stored. So we always have at most p significant bits.

$$R = +P \times 2^{Q-b}$$

• The minimum exponent Q is 0, in this case,

$$P = 0.P_1P_2 \dots P_{p-1}$$

 $R = \pm P \times 2^{-b+1}$

In this case, to get to smaller numbers, we trade off the number of significant bits.

- The smallest positive number it can express is thus $2^{-b+1-(p-1)} = 2^{-b-p+2}$. However, in this case, there is a single significant bit.
- The smallest positive number it can express that keeps p significant bits is 2^{-b+1} .
- The largest positive number it can express is $(1-2^{-p+1}) \times 2^{2^q-2-b}$.

8.1 A double precision float

- In the IEEE 754 standard: p = 53, q = 11, b = 1023.
- The smallest positive number it can express is $2^{-1023-53+2} = 2^{-1074} \approx 5 \times 10^{-324}$. However, in this case, there is a single significant bit.
 - Thus, in Python, $2^{-1075} = 0!$ This is an example of **rounding error**, see below
- The smallest positive number to keep 53 significant bits is 2^{-1022}
- The largest positive number it can express is $(2-2^{-52}) \times 2^{1023} \approx 1.7976931348623157 \times 10^{308}$
 - Any number larger that that is ∞ .

8.1.1 Example

What is the value of the double precision float R where the sign bit S=1 the exponent $Q=(100\ldots0)_2=2^{10}$, and $P=(100\ldots0)_2=2^{51}$

8.2 Rounding error

- A rounding error is an error introduced in the finite-digit representation of real numbers. any number with more than the maximum significant bits that a float point number can represent is rounded.
 - the above example of $2^{-1075} = 0$ is a manifestation.
 - the largest integer with 52 bits of significant bits is $2^{52} 1 = 4,503,599,627,370,495$
 - * if $a/b > (2^{52} 1)/2 \approx 2.251 \times 10^{15}$ (the /2 is to account for rounding up) then, in Python, a b has a rounding error
 - * if $a/b > (2^{52}-1)/4 \approx 1.126 \times 10^{15}$, then b is rounded to 0 when computing a-b, i.e., a-b=a
 - This only depends on the relative magnitude of a and b

```
[13]: # another example of rounding error
a=(2.0**52-1)*1000
b = 0.2*1000
print("a-(a-b)=",a-(a-b))
print("a-b = a is", a-b == a)
print("a-a+b=",a-a+b)
print("a-a+b==a-(a-b) is",(a-a+b)==(a-(a-b)))
```

```
a-(a-b)= 0.0
a-b = a is True
a-a+b= 200.0
a-a+b==a-(a-b) is False
```

The above example shows that, due to rounding errors, in Python (actually, in all computer languages with float point numbers), the association rule may not hold for float numbers!

8.3 The important of Rounding errors

Rounding errors are typically small. It is curious why we worry about such a small error. Some algorithms intrinsically magnifies errors. So affter many iterations, the rounding error may be magnified to an unacceptable large value. We will see such examples in future lectures.

[]	:	
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