

Scheme Coding and Grammars

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Grammars

Consider an expression language with two binary operators "\$" and "#" and two unary prefix operators "*" and "@". It also includes a single alphabetic symbol "X" along with parenthesis "(" and ")". An EBNF expression grammar is given below where the starting non-terminal is <expr> and terminals are highlighted.

```
<expr> ::= <term> { $ <term> }  
<term> ::= <factor> { # <factor> }  
<factor> ::= [ @ | * ] ( X | ( <expr> ) )
```

1. Give an equivalent unambiguous BNF grammar for this expression language.

```
<expr> ::= <term> $ <expr>  
          | <term>  
<term> ::= <factor> # <term>  
          | <factor>  
<factor> ::= @ <symbol>  
            | * <symbol>  
            | <symbol>  
<symbol> ::= X  
            | ( <expr> )
```

2. Give the associativity of the two binary operators and a precedence table for all four operators in your BNF grammar.

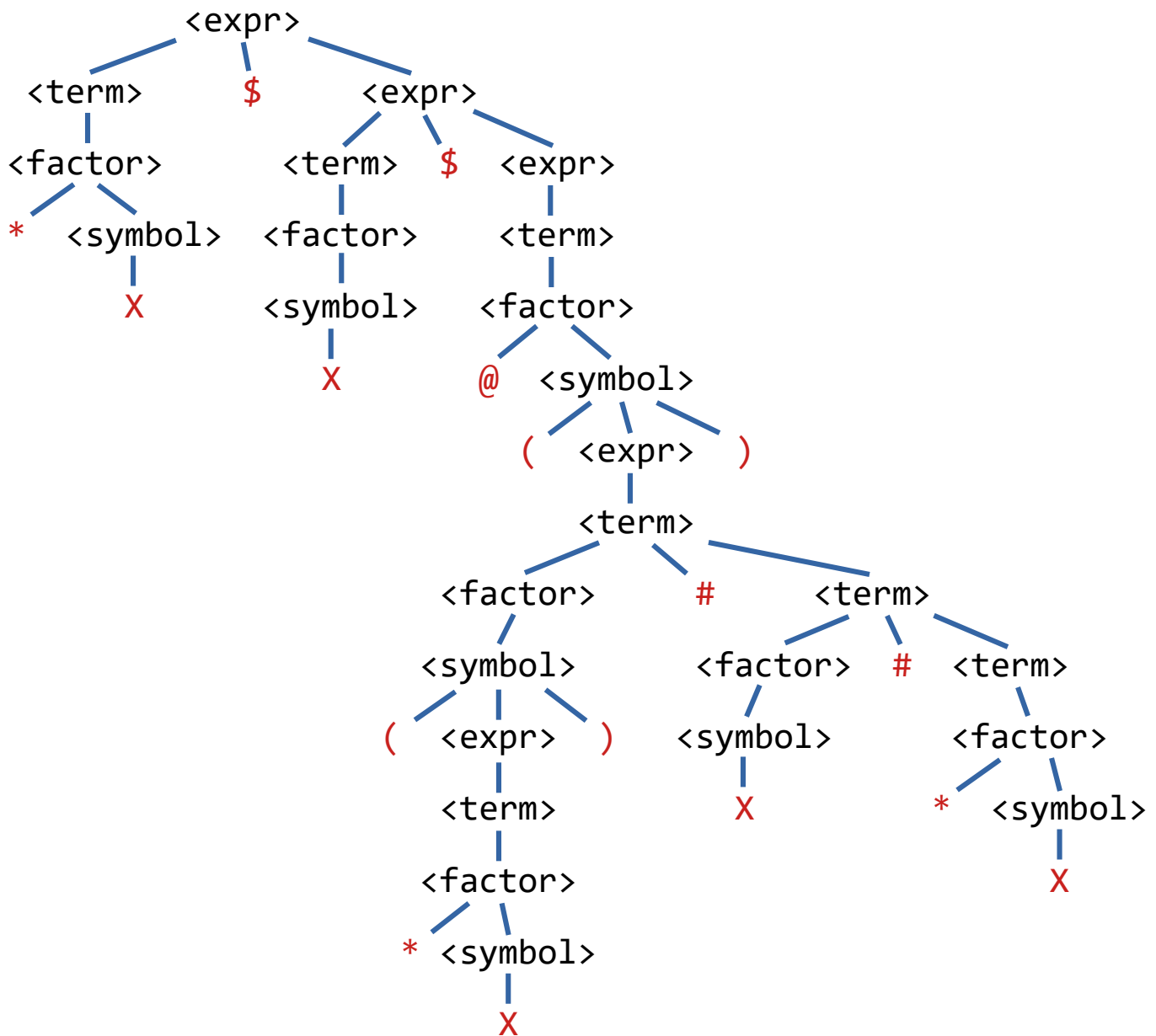
The binary operators \$ and # are right associative because productions #1 and #3 are right recursive with themselves.

Precedence Table

Precedence	Operator	Description
3	* @	Unary operator
2	#	Binary operator
1	\$	Binary operator

```
#1.  <expr>      ::= <term> $ <expr>
#2.          | <term>
#3.  <term>      ::= <factor> # <term>
#4.          | <factor>
#5.  <factor>    ::= @ <symbol>
#6.          | * <symbol>
#7.          | <symbol>
#8.  <symbol>    ::= X
#9.          | ( <expr> )
```

3. Using your BNF specification, provide a parse tree for the expression ***X\$X\$@((*X)#X#*X)** .



```

<expr> ::= <term> $ <expr>
          | <term>
<term> ::= <factor> # <term>
          | <factor>
<factor> ::= @ <symbol>
             | * <symbol>
             | <symbol>
<symbol> ::= X
            | ( <expr> )

```

Syntactic Ambiguity

Prove that the following BNF grammar **G** is ambiguous.

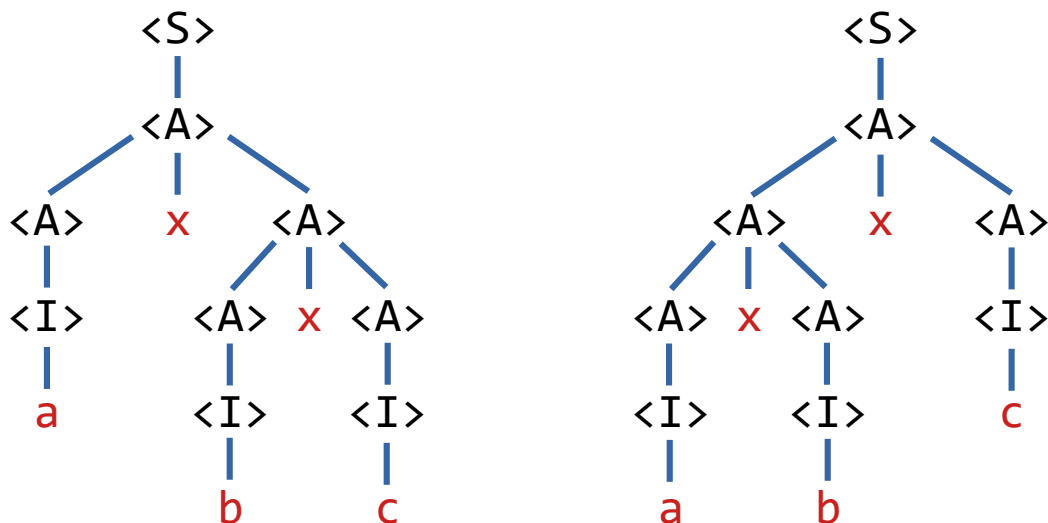
$N = \{ \langle S \rangle, \langle A \rangle, \langle I \rangle \}$
 $T = \{ a, b, c, x \}$
 $P = \{ \langle S \rangle ::= \langle A \rangle$
 $\langle A \rangle ::= \langle A \rangle x \langle A \rangle$
 $\langle A \rangle ::= \langle I \rangle$
 $\langle I \rangle ::= a$
 $\langle I \rangle ::= b$
 $\langle I \rangle ::= c \}$
 $S = \langle S \rangle$

Claim: The BNF grammar **G is ambiguous**

Proof:

Suppose the BNF grammar **G** is NOT ambiguous, that grammar **G** and its language does NOT contains at least one string with two or more distinct parse trees. {

Let **STR** = "axbxc" {



First, there are two valid parse trees for **STR**.

However, this would mean that there is at least one string with two or more distinct parse trees contained in the BNF grammar **G**.

This contradicts the supposed statement which states that grammar **G** is NOT ambiguous.

}

Therefore, the BNF grammar **G** is ambiguous.

}

□