# **Graph Drawing**

COMP8503

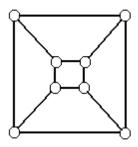
Advanced Topics in Visual Analytics

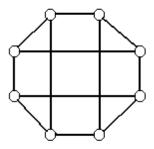
### Reference

- Isabel F. Cruz, Roberto Tamassia. Graph Drawing Tutorial.
- Jesper Nederlof, Roeland Luitwieler. Introduction to Graph Drawing.
- Ioannis G. Tollis, Giuseppe Di Battista, Peter Eades, Roberto Tamassia. Graph Drawing: Algorithms for the Visualization of Graphs
- Graph Drawing, Tom Germano
- Tamassia, R. 1998. Constraints in Graph Drawing Algorithms. Constraints 3, 1 (Apr. 1998), 87-120.
- Eades, P. (1984). A heuristic for graph drawing. Congressus Numerantium, 42:149–160.
- K. Sugiyama & K. Misue. (1995). Graph drawing by magnetic-spring model. J. Visual Lang. Comput. 6(3).
- T. Kamada & S. Kawai. (1989). An algorithm for drawing general undirected graphs. Inform. Process. Lett. 31: 7–15.
- T. Fruchterman & E. Reingold. (1991). Graph drawing by force-directed placement. Softw. Pract. Exp. 21(11): 1129–1164.
- R. Davidson & D. Harel. (1996). Drawing graphics nicely using simulated annealing. ACM Trans. Graph. 15(4): 301–331.
- M. Fiedler. A property of eigenvectors of nonnegative symmetric matrices and its application to graph theory. Czechoslovak Mathematical Journal, 25(100):619–633, 1975.
- Frishman, Y. and Tal, A. 2007. Multi-Level Graph Layout on the GPU. IEEE Transactions on Visualization and Computer Graphics 13, 6 (Nov. 2007), 1310-1319.
- Auber, Y. Chiricota, F. Jourdan, G. Melancon. Multiscale visualization of small world networks. INFOVIS 2003. pp. 75-81.
- Many slides and content are borrowed from those references

## Graphs

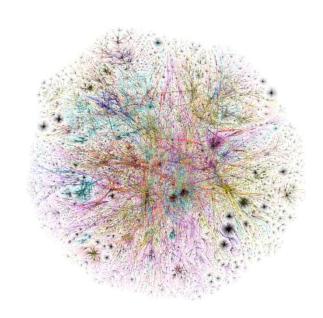
- G = {V, E}, nodes, edges
- Undirected graph, di-graph, weighted graph
- Planar graph drawing





## **Graph Drawings**

- Visualization of graphs/networks
  - Models, algorithms, systems
- Applications
  - Software engineering
  - Database systems
  - Project management
  - Knowledge representation
  - Telecommunications
  - WWW



Internet map from the Opte Project

## **Graph Drawings**

- Readable to users
- Follow drawing conventions
- Satisfy as many aesthetic rules as possible
- Efficient running time

### Outline

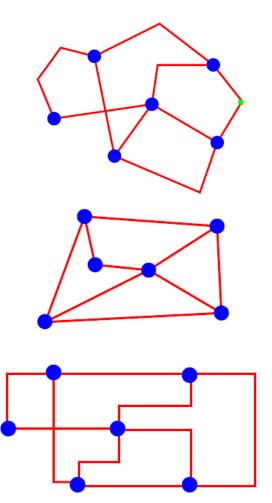
- Drawing Conventions
- Aesthetic Criteria
- Force-directed method
- Multiscale method

## **Drawing Conventions**

Polyline drawings

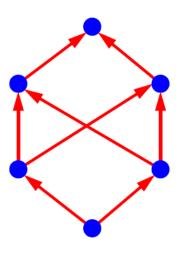
Straight-line drawings

Orthogonal drawings



## **Drawing Conventions**

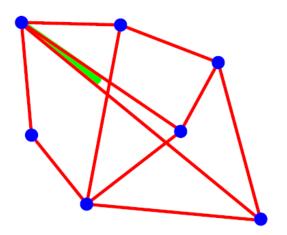
- Planar drawings
  - No crossings allowed
- Upward drawings
  - Drawn as nondecreasing arcs
  - For hierarchical relationships
- Downward drawings
  - Drawn as nonincreasing arcs



## **Drawing Conventions**

- Resolution
  - Smallest distance between vertices
  - Smallest distance between vertices and nonincident edges

- Angular resolution
  - Smallest angle formed by two incident edges at a vertex

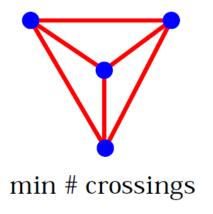


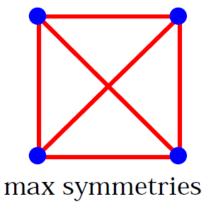
### **Aesthetic Criteria**

- Minimize edge crossings
- Minimize area
- Minimize bends (orthogonal drawing)
- Maximize angular resolution
- Symmetry
- Min. sum / maximum / variance of edge lengths

### **Aesthetic Criteria**

 "In general, one cannot simultaneously optimize two aesthetic criteria"

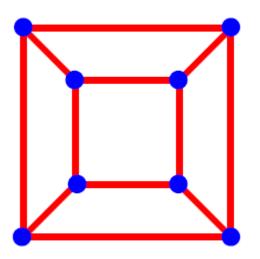


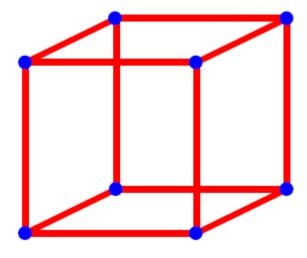


- Complexity
  - Minimizing edge crossing is NP-hard
  - Computing optimal angular resolution is NP-hard

## **Aesthetic Criteria**

Beyond aesthetic criteria





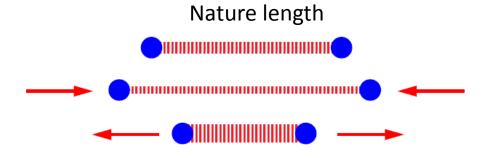
## General Undirected Graph

- Planar straight-line drawings
  - Generates as few edge crossings as possible
- Force-directed method
- Multiscale method

- Define a system of forces acting on the vertices and edges
- Find a minimum energy state by
  - Solving differential equations or
  - Simulating the evolution of the system

- Spring embedder [Eades 1984]
- [Kamada and Kawai 89]
- [Fruchterman and Reingold 90]
- [Davidson and Harel 96]

Replace an edge with a spring of unit length



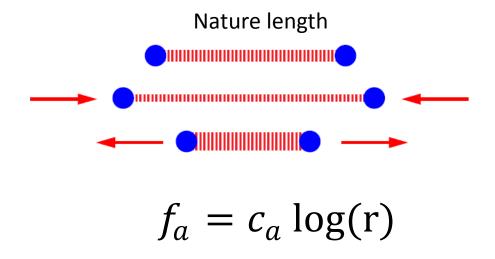
Connect nonadjacent nodes with springs of infinite rest length

Hooke's law

$$f = -\mathbf{k}(\mathbf{x} - \mathbf{x}_0)$$

- f is the force
- k is the factor
- $x_0$  is the nature length
- Force model deviates from Hooke's law

Replace an edge with a spring of unit length



- $f_a$  is the attraction force
- $c_a$  is the attraction factor
- r is the distance between nodes

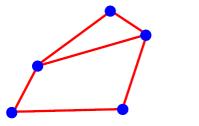
Connect nonadjacent nodes with springs of infinite rest length

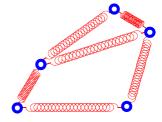
$$f_r = \frac{c_r}{r^2}$$

- $f_r$  is the repulsive force
- $c_r$  is the repulsive factor
- r is the distance between nodes

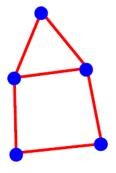
- Initial graph layout randomly
- Update layout iteratively
  - Apply spring forces to connected node pairs
  - Apply spring forces to unconnected node pairs
  - Update layout
  - Until the movements are small enough

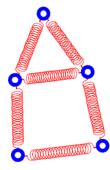
Initial layout and springs





Final configuration





- [Kamada and Kawai 89]
  - For a pair of nodes (u, v), the spring nature length is proportional to d(u, v) which is the distance from u to v
  - Define energy of the system

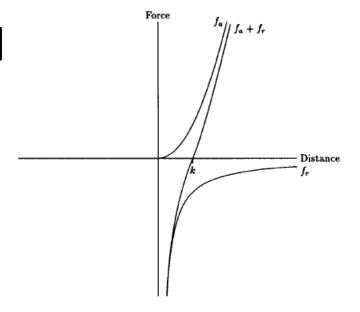
$$\sum k_{u,v}(|p_u - p_v| - d(u,v))^2$$

Reduce energy iteratively by solving PDE equations

- [Fruchterman and Reingold 90]
  - A complex system of forces

$$f_a = \frac{r^2}{k} \qquad f_r = -\frac{k^2}{r}$$

Where 
$$k = C\sqrt{\frac{area}{\#nodes}}$$

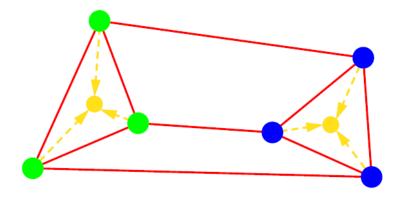


- Control drawings within a boundary
- A "temperate" controlling scheme
  - Temperature from hot to cold and bounds vertex movement

- [Davidson and Harel 96]
  - Consider vertex distributions, edge lengths,
     crossing into a energy function instead of forces
  - Simulated annealing to find solutions
    - Computation costly

- Advantages
  - Simple implementations
  - Easy to add new heuristics, constraints
  - Smooth evolution of layout
  - Supports 3D
  - Often detects and shows symmetries
  - Works well with small graphs

- Add new constraints through means of forces
  - Position constraints with fixed positions or prescribed regions
  - Orientation constraints with "magnetic field"
    - [Sugiyama Misue 95]
  - Group constraints by adding dummy "attractors"



- Disadvantages
  - Slow, running time, convergence
  - Few theoretical supports of drawing quality
  - Difficult to support orthogonal and polyline drawings

- Motivation
  - Force directed method is slow on large graphs
  - The result is sensitive to the initial layout when dealing with large graphs

- Multi-Level Graph Layout on the GPU
  - [Frishman and Tel 07]

- Generate graphs at different spatial scales
  - Partitioning, clustering, coarsening
- Start from the coarsest scale and work back to finest scale
  - How to propagate results from one level to another

- Advantages
  - Reduce computation cost
  - Maintain good quality
  - Less sensitive to initial configurations
- Disadvantages
  - Growing inaccuracy when going from one level to another

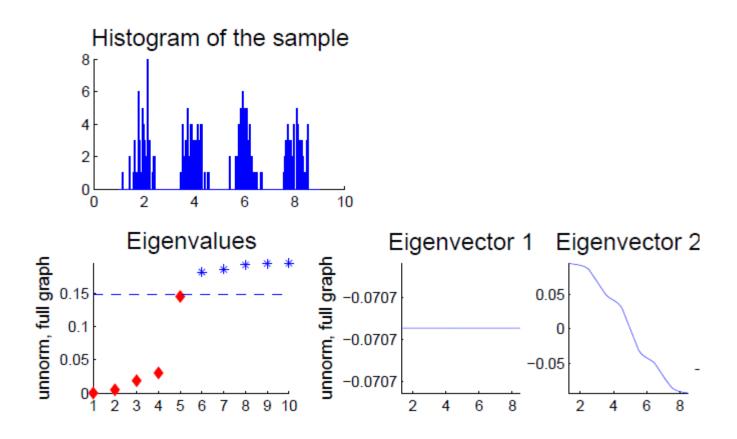
- Graph spectral partitioning
- Multilevel scheme for graph layout
- Accelerate force-directed method

- Graph spectral partition
  - Partition a graph to parts with two properties
    - Similar size <-> balance in layout
    - Minimum cut <-> weakly coupled, independent
- An eigenvector problem [Fiedler 75]
  - Given a graph G, its Laplacian L is defined as

$$\ell_{i,j} := \begin{cases} \deg(v_i) & \text{if } i = j \\ -1 & \text{if } i \neq j \text{ and } v_i \text{ is adjacent to } v_j \\ 0 & \text{otherwise.} \end{cases}$$

- Graph Laplacian L with eigenvalues  $\lambda_0 \le \lambda_1 \le \cdots \le \lambda_{n-1}$ 
  - Positive semi-definite  $\forall i, \lambda_i \geq 0, \quad \lambda_0 = 0$
  - # zero-eigenvalue = # connected component
  - Smallest non-zero eigenvalue => spectral gap, its eigenvector can be used to cluster the graph, namely "spectral clustering"

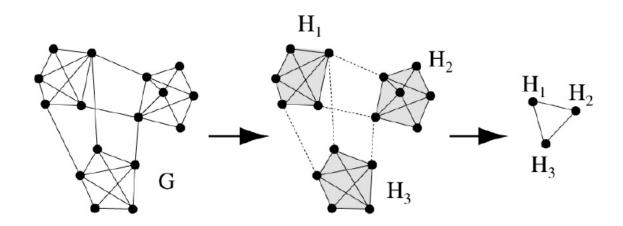
Spectral clustering, a simple example



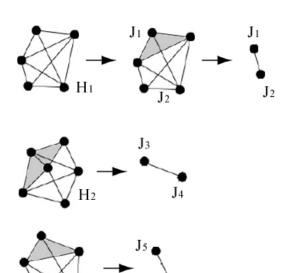
Ulrike von Luxburg, A Tutorial on Spectral Clustering

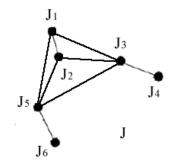
- To compute the eigenvector
  - Power iterations
    - Efficient in computation and memory
    - Slow convergence

- Multilevel representation
  - First level partitioning

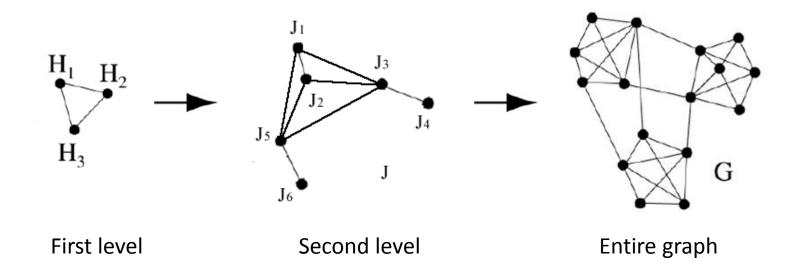


- Multilevel representation
  - Second level partitioning

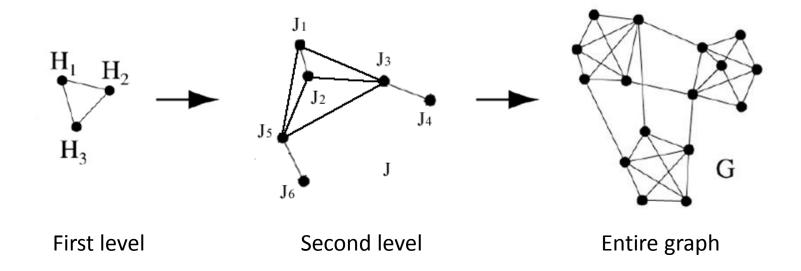




- Multilevel representation
  - Combined



- Compute layout
  - Start from the coarsest level
  - Propagate to finer level, then refine



- Layout propagation
  - Each node in current level is placed at its parent's position in the coarser level
  - Scale positions

$$p_i(x,y) = \sqrt{\frac{|V(L^l)|}{|V(L^{l-1})|}} \cdot p_i(x,y),$$

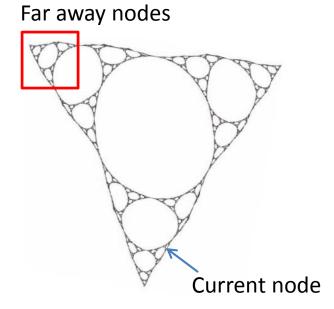
Improve positions

$$p_i = \frac{1}{2} \left( p_i + \frac{1}{degree(i)} \sum_{j \in N(i)} p_j \right).$$

### Accelerate Force-directed Method

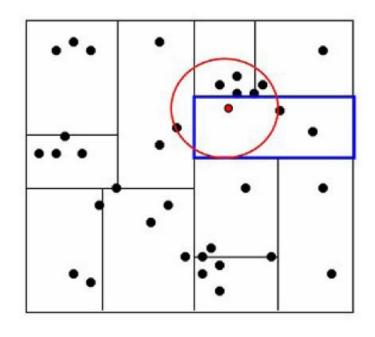
- Bottlenecks
  - Repulsive forces between all pairs of nodes
  - N-body problem, O(|V|^2)

- Solutions
  - For nodes that are far away, use approximations



## Accelerate Force-directed Method

- Kd-tree
  - Spatial partition of nodes
  - Give a particular node
    - Nodes out of a range or its cell are far away nodes
    - Use cell centers to approximate repulsive forces



Region query on a kdtree

### Accelerate Force-directed Method

Parallel implementation

Pseudo code:
For all nodes in parallel
For all other nodes
compute and accumulate repulsive forces

- "Embarrassingly parallel"
- Suits GPU and multi-core architectures

## Summary

- Graph drawing
  - Conventions
  - Aesthetic criteria
- Drawing algorithm for undirected graphs
  - Force directed methods
  - Multiscale methods

## Summary

- Graph drawing
  - Conventions
  - Aesthetic criteria
- Drawing algorithm for undirected graphs
  - Force directed methods
  - Multiscale methods

Questions?