

Lec-1-计算机基本组成

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Four Functions

组成

冯·诺依曼计算机 (Von Neumann Machine)

哈佛结构计算机的特点

总结

摩尔定律 (The Era of Moore's Law)

原因

未来发展方向

Four Functions

- Data movement
- Data storage
- Data processing
- Control

组成

- I/O
- CPU
 - Control Unit ---- 控制单元
 - Registers ---- 寄存器
 - Arithmetic and Logic Unit(ALU)----逻辑控制单元
- Memory

冯·诺依曼计算机 (Von Neumann Machine)

1. **存储结构**: 指令和数据共享同一存储器。
2. **总线设计**: 指令和数据共享同一总线。这意味着在任何时刻, CPU只能读取指令或数据, 不能同时进行两者的访问。
3. **执行顺序**: 指令通常按顺序存放和执行, 但在特定条件下可以根据运算结果或设定的条件改变执行顺序。
4. **灵活性**: 由于程序和数据存储在单一内存中, 程序可以动态地修改或存取数据。
5. **硬件复杂性**: 结构相对简单, 易于实现和扩展。
6. **应用场景**: 广泛应用于通用计算机、PC、服务器等。
7. **Storage Structure**: Instructions and data share the same memory.
8. **Bus Design**: Instructions and data share the same bus. This means that at any given moment, the CPU can only read instructions or data, but not both simultaneously.
9. **Execution Order**: Instructions are typically stored and executed in sequence, but under certain conditions, the execution order can be changed based on the results of computations or predefined conditions.

10. **Flexibility:** Since programs and data are stored in the same memory, programs can dynamically modify or access data.
11. **Hardware Complexity:** The structure is relatively simple, making it easy to implement and expand.
12. **Application Scenarios:** Widely used in general-purpose computers, PCs, servers, etc.

哈佛结构计算机的特点

1. **存储结构:** 指令和数据存储在不同的存储器中。
2. **总线设计:** 采用独立的指令总线 and 数据总线，允许指令和数据同时被访问。
3. **并行处理:** 由于指令和数据可以并行访问，哈佛结构能够提高处理速度和效率。
4. **硬件复杂性:** 需要双存储器和双总线设计，增加了硬件成本和复杂性。
5. **灵活性:** 由于指令存储器通常是只读的，程序无法直接修改指令存储器中的内容。
6. **应用场景:** 常用于嵌入式系统、数字信号处理器（DSP）等对性能要求较高的应用。

1. **Storage Structure:** Instructions and data are stored in separate memory units.
2. **Bus Design:** Independent instruction and data buses are used, allowing simultaneous access to instructions and data.
3. **Parallel Processing:** Since instructions and data can be accessed in parallel, the Harvard architecture can improve processing speed and efficiency.
4. **Hardware Complexity:** Requires a dual-memory and dual-bus design, which increases hardware costs and complexity.
5. **Flexibility:** As the instruction memory is typically read-only, programs cannot directly modify the contents of the instruction memory.
6. **Application Scenarios:** Commonly used in embedded systems, digital signal processors (DSPs), and other applications with high performance requirements.

总结

- **性能:** 哈佛结构通过并行访问指令和数据，避免了冯·诺依曼架构中的总线冲突问题，从而提高了执行效率。
- **灵活性与成本:** 冯·诺依曼架构由于其简单性和灵活性，适用于通用计算，而哈佛架构则在专用计算领域表现出色。

摩尔定律 (The Era of Moore's Law)

Double (every 18-24 months):

- The number of transistors per chip
- Computer performance : higher packing density means shorter electrical paths, giving higher performance
- CPU clock speed •Power and cooling requirements are reduced (until 2013)
- One of the above, at constant cost (until 2016)

原因

Dennard Scaling (登纳德缩放定律) 是1974年由Robert Dennard提出的, 它与摩尔定律共同指导了集成电路行业多年。Dennard Scaling的核心观点是, 随着晶体管尺寸的缩小, 其功率密度保持不变, 从而使芯片的功率与芯片面积成正比

Dennard (1974) observed that voltage and current should be proportional to the linear dimensions of a transistor

- Therefore, as transistors shrank, so did required voltage and current
- power is proportional to the area of the transistor.
- 因此, 随着晶体管尺寸的缩小, 所需的电压和电流也随之减小.
- 功率与晶体管的面积成比例.

未来发展方向

performance / watt

operations / Joules (Joules = watts x second)