# **Arboricity and Degeneracy of Graphs**

#### Equivalence, Bounds, and Intuition

#### **Definitions**

**Definition 1.1** (Arboricity). For a graph G = (V, E), the arboricity  $\alpha(G)$  is the minimum number of edge-disjoint forests whose union is G, equivalently:

$$\alpha(G) = \max_{H \subseteq G, |V(H)| \ge 2} \left\lceil \frac{|E(H)|}{|V(H)| - 1} \right\rceil.$$

**Definition 1.2** (Degeneracy). *The* degeneracy  $\kappa(G)$  *is the smallest k such that every subgraph of G has a vertex of degree at most k. Equivalently:* 

$$\kappa(G) = \max_{H \subseteq G} \left( \min_{v \in V(H)} \deg_H(v) \right).$$

## Relationship Between $\alpha$ and $\kappa$

We have the tight bounds:

$$\alpha(G) \leq \kappa(G) \leq 2\alpha(G) - 1.$$

### **Combinatorial Proof that** $\alpha \leq \kappa$

Let  $k = \kappa(G)$ , so G is k-degenerate. Then there exists an ordering of the vertices

$$v_n, v_{n-1}, \ldots, v_1$$

such that each vertex  $v_i$  has at most k neighbors among  $\{v_1, \ldots, v_{i-1}\}$ .

Every edge of any *n*-vertex subgraph *H* is counted exactly once by its later endpoint in this ordering, so

$$|E(H)| \leq k(n-1).$$

By Nash-Williams,

$$\alpha(G) = \max_{H \subset G, |V(H)| > 2} \left\lceil \frac{|E(H)|}{|V(H)| - 1} \right\rceil \le \max_{H} \lceil k \rceil = k,$$

giving  $\alpha(G) \leq \kappa(G)$ .

### **Proof that** $\kappa \leq 2\alpha - 1$

For any subgraph H with n vertices and m edges, the Nash–Williams bound  $m \le (n-1)\alpha$  yields an average degree

$$\overline{d}(H) = \frac{2m}{n} < 2\alpha.$$

Thus

$$\min_{v \in V(H)} \deg_H(v) \leq \overline{d}(H) < 2\alpha \quad \Longrightarrow \quad \min_v \deg_H(v) \leq 2\alpha - 1.$$

Taking the maximum over *H* gives  $\kappa(G) \leq 2\alpha - 1$ .

# **Intuition and Impact**

These inequalities show that arboricity and degeneracy both capture graph sparsity, with arboricity measuring forest decompositions and degeneracy capturing peelability. Despite different definitions, they remain within a factor of two of each other.