

a closed semiring  $(S, +, \cdot, 0, 1)$

1. monoid  $(S, +, 0), (S, \cdot, 1)$

-> closed, identity, associative

2.

a quine: a program which copys itself

recursion theorem:

1.  $r(w) = t(\langle R \rangle, w)$

T is a computer program, t is a computable function  $t: \Sigma^* \times \Sigma^* \rightarrow \Sigma^*$

a computable function  $r: \Sigma^* \rightarrow \Sigma^*$

2.  $E_i = E_{f(i)}$

다익스트라

$(S, +, \cdot, 0, 1)$

$S = \mathbb{R}^+ \cup \infty$

$+$  = min

$\cdot$  = +

$0 = \infty$

$1 = 0$

partial order?

1. reflexive

2. anti-symmetric

3. transitive

a binary relation R

$A = \{1, 2, 3\}$

$R = \{(1, 2), (2, 1), (1, 1), (2, 2), (3, 3)\}$  -> not partial order

$R = \{(1, 2), (1, 1), (2, 2), (3, 3)\}$  -> partial order

anti-symmetric?

$\sim$ (if  $aRb$  then  $bRa$ )

$\equiv$  if  $aRb$  then no  $bRa$

fixed point  $\equiv f(x)=x$

NPC의 정의

1.

(a) a set of decision problems  $x$  is in NP

(b) All problems in NP are efficiently reducible to  $x$

2.

If any problem in NPC is solvable efficiently

Then  $P=NP$

$a \in NP$

$b \in P$

$c \in NPC$

relative hardness

$a \geq b$

$a \leq c$

$b \leq c$

$P = \{x | \text{-----}\}$

1.  $x$  is a decision problem

2.  $x$  is efficiently solvable

$NP = \{x | \text{-----}\}$

1.  $x$  is a decision problem

2.  $\forall_{yes-instance\ a} \exists_{certificate\ b}$  Verification ( $a, b$ ) is done efficiently

yes-instance,  $\langle G, k \rangle$  for  $k$ -clique problem