

abstract recursive theorem

1.  $\Sigma$  is closed under composition ( $\Sigma$  is a set of functions:  $F(x, y): Z^+ \times Z^+ \rightarrow Z^+$ )
2.  $F(x, x) \in \Sigma$
3.  $\forall f \in \Sigma \exists a \in Z^+ \forall x \in Z^+ E_{F(a, x)} = E_{f(x)}$

- 
4.  $\forall f \in \Sigma \exists i \in Z^+ E_i = E_{f(x)}$

proof)

$$\forall f \in \Sigma \exists a \in Z^+ \forall x \in Z^+ E_{F(a, x)} = E_{f(x)}$$

the function f is arbitrary

- >  $E_{F(a, x)} = E_{f(F(a, x))}$
- >  $E_{F(a, a)} = E_{f(F(a, a))}$
- >  $E_i = E_{f(i)}$

abstract recursive theorem은 sound한 argument이다.

두 집합 A,  $\Sigma$ 에 대하여,

A를 the set of all programs written in Java(C, Python... (Turing Complete Programs))라고 하고,

$\Sigma$ 를 the set of all computable functions  $Z^+ \rightarrow Z^+$ 라고 하면,

즉, 두 집합 A,  $\Sigma$ 를 그렇게 잡아버리면 Sound해진다.

- >  $E_{F(a, x)} = E_{f(x)}$ 이 무조건 발생
- > 정확히 하는 일이 같은 크기가 서로 다른 두 프로그램이 존재한다. -> 이걸 피할 수 없다.

-a quine- :self-pointing program

a quine과 abstract recursive theorem의 관계는 무엇인가

<Chapter 10> : Graphs

$G=(V, E)$

a graph:

1. a directed graph
2. a undirected graph

$V$ : a finite set of nodes

$E$ : a subset of  $V \times V$

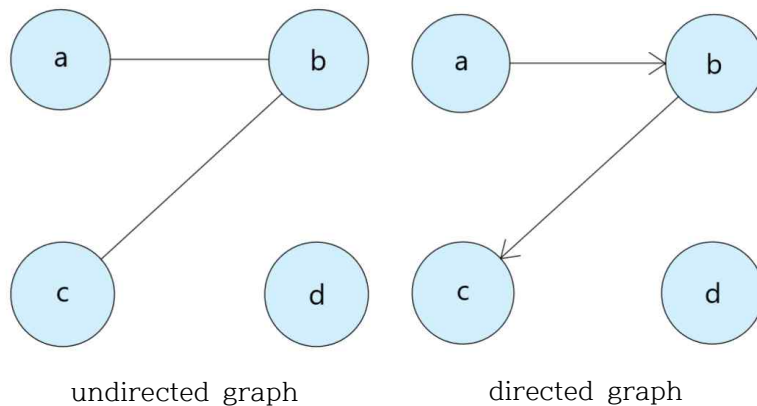
$\rightarrow$  a binary relation on  $V$

(undirected graph의 경우, edge는 a subset of unordered pairs of graph)

-transitive closure-

1. relation
2. directed graph

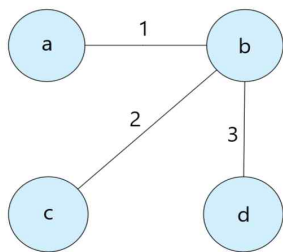
ex)  $V=\{a, b, c, d\}$ ,  $E=\{(a, b), (b, c)\}$



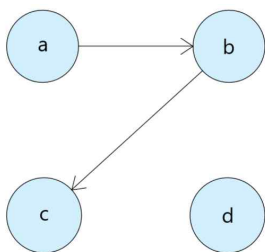
-adjacency/incidence-

Adjacency/Adjacent  $\rightarrow$  both directed & undirected에 사용

Incidence/incident  $\rightarrow$  only for undirected에 사용



edge 1 is incident on(upon) a and b  
 edge 3 is incident on(upon) b and d  
 a and c are adjacent



b is adjacent to a  
 c is adjacent to b

-Representation of graph-

1. Adjacency matrix
2. Adjacency lists

(1) Adjacency matrix

