

## NETAJI SUBHAS UNIVERSITY OF TECHNOLOGY

# **Practical Report**

Database Management Systems

Computer Science Engineering (Internet of Things) Semester~3

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## 1 Introduction

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## 2 Sailors

#### 2.1 Schema

Consider the following relational schema:

```
SAILORS (<u>sid</u>, sname, rating, date_of_birth)

BOATS (<u>bid</u>, bname, color)

RESERVES (<u>sid</u>, bid, date, time_slot)
```

## 2.2 Queries

- 1. Find sailors who've reserved at least one boat
  - (a) Relational Algebra

```
\pi_{sid,sname}(SAILORS \bowtie RESERVES))
```

(b) SQL

```
SELECT sname
FROM SAILORS
WHERE sid IN (
SELECT sid
FROM RESERVES
);
```

- 2. Find names of sailors who've reserved a red or a green boat in the month of March.
  - (a) Relational Algebra

```
\pi_{sname}(SAILORS \bowtie RESERVES \bowtie BOATS) \bowtie 
\sigma_{bname=red \lor bname=green}(\sigma_{date=March}(BOATS \bowtie RESERVES))
```

(b) SQL

```
SELECT sname
1
    FROM SAILORS
2
    WHERE sid IN
3
         (SELECT sid
4
         FROM RESERVES
5
        WHERE bid IN
6
             (SELECT bid
             FROM BOATS
             WHERE bname = 'red' OR bname = 'green')
9
         AND (SELECT extract(month FROM date) FROM RESERVES) = 3)
10
```

- 3. Find names of sailors who've reserved a red and a green boat
  - (a) Relational Algebra

```
\pi_{sname}(SAILORS \bowtie RESERVES \bowtie (\sigma_{color=red}(BOATS))) \cap \pi_{sname}(SAILORS \bowtie RESERVES \bowtie (\sigma_{color=green}(BOATS)))
```

```
SELECT DISTINCT S1.sname
1
    FROM SAILORS S1, RESERVES R1, BOATS B1,
2
    RESERVES R2, BOATS B2
3
    WHERE S1.sid = R1.sid
4
        AND R1.bid = B1.bid
5
        AND S1.sid = R2.sid
6
        AND R2.bid = B2.bid
        AND B1.color = "red"
8
        AND B2.color = "green";
```

- 4. Find SID of sailors who have <u>not</u> reserved a boat after Jan 2018.
  - (a) Relational Algebra

$$\pi_{\rm sid} - \pi_{\rm sid}(SAILORS \bowtie \sigma_{\rm date\_of\_birth > Jan\ 2018}(RESERVES))$$

(b) SQL

```
SELECT sid FROM SAILORS
WHERE sid NOT IN

(SELECT sid FROM RESERVES
WHERE date_of_birth > "2018-01-01")
```

- 5. Find sailors whose rating is greater than that of all the sailors named "John"
  - (a) Relational Algebra

$$\pi_{\text{sid,sname}}(SAILORS) - \pi_{S_2.\text{sid},S_2.\text{sname}}(\sigma_{S_2.\text{rating}} < S.\text{rating}(\rho_{S_2}(SAILORS) \times \rho_S(SAILORS)))$$

(b) SQL

```
SELECT sid, sname FROM SAILORS S1
WHERE S1.rating > ALL
(SELECT S2.rating FROM SAILORS S2
WHERE S2.sname = "John")
```

- 6. Find sailors who've reserved all boats
  - (a) Relational Algebra

$$\pi_{\text{sid.sname}}(\pi_{\text{sid.bid}}(RESERVES) \div \pi_{\text{bid}}(BOATS)) \bowtie SAILORS$$

```
SELECT S.sid, S.sname
1
    FROM SAILORS S
2
    WHERE NOT EXISTS
3
           (SELECT B.bid
           FROM BOATS B
5
           WHERE NOT EXISTS
6
                 (SELECT R.sid, R.bid
                 FROM RESERVES R
8
                 WHERE R.sid = S.sid
9
                        AND R.bid = B.bid))
10
```

- 7. Find name and age of the oldest sailor(s)
  - (a) Relational Algebra

```
\pi_{\text{sname,age}}(\pi_{sid}(SAILORS) - \\ \pi_{S_2.sid}(\sigma_{S_2.age < S.age}(\rho_{S_2}(SAILORS) \times \rho_S(SAILORS)))) \\ \bowtie SAILORS
```

(b) SQL

```
SELECT sname, age FROM SAILORS S1
WHERE S1.date_of_birth > ALL
(SELECT S2.date_of_birth FROM SAILORS S2)
```

8. Find the age of the youngest sailor for each rating with at least 2 such sailors

(a) Relational Algebra

```
\pi_{\text{rating,minage}}(\sigma_{\text{no\_of\_sailors}>1} \\ (\rho_{r(\text{rating,no\_of\_sailors,minage})} \, \mathcal{F} \, (\text{rating,} count(\text{sid}), min(\text{age})) \\ (SAILORS)))
```

(b) SQL

```
SELECT rating, age FROM SAILORS S1
WHERE S1.date_of_birth > ALL AS minage
(SELECT S2.date_of_birth FROM SAILORS S2
WHERE S2.rating = S1.rating)
GROUP BY rating
HAVING COUNT(*) >= 2
```

### 3 Customers

#### 3.1 Schema

Consider the following relational schema:

```
CUSTOMERS (<a href="mailto:cust_num">cust_num</a>, cust_lname, cust_balance)

PRODUCT (<a href="mailto:prod_num">prod_num</a>, prod_name, price)

INVOICE (<a href="mailto:inv_num">inv_date</a>, unit_sold, inv_amount)
```

### 3.2 Queries

Write SQL queries and relational algebraic expression for the following:

- 1. Find the names of the customer who have purchased no item. Set default value of cust\_balance as 0 for such customers.
  - (a) Relational Algebra

```
\pi_{\text{cust\_lname}+\text{""}+\text{cust\_fname}}(\sigma_{\text{cust\_balance}=0}(CUSTOMERS))
```

(b) SQL

```
SELECT concat(cust_lname , " " , cust_fname) as name
FROM CUSTOMERS
WHERE cust_balance = 0
```

2. Write the trigger to update the CUST\_BALANCE in the CUSTOMER table when a new invoice record is entered for the customer.

```
CREATE TRIGGER update_cust_balance

AFTER INSERT ON INVOICE

FOR EACH ROW

BEGIN

UPDATE CUSTOMERS

SET cust_balance = cust_balance + NEW.inv_amount

WHERE cust_num = NEW.cust_num;

END
```

- 3. Find the customers who have purchased more than three units of a product on a day.
  - (a) Relational Algebra

```
\pi_{\text{cust\_lname}+\text{""}+\text{cust\_fname}}(\sigma_{\text{unit\_sold}\geq 3} \ (CUSTOMER \bowtie \sigma_{\text{unit\_sold}} \ (_{\text{cust\_num, inv\_date}>prod\_num} \mathcal{F}_{sum(\text{unit\_sold})}(INVOICE))))
```

```
SELECT concat(cust_lname , " " , cust_fname) as name
FROM CUSTOMERS
WHERE cust_num IN
(
SELECT cust_num
FROM INVOICE
GROUP BY cust_num, inv_date, prod_num
HAVING sum(unit_sold) >= 3
)
```

- 4. Write a query to illustrate Left Outer, Right Outer and Full Outer Join.
  - (a) Left Outer Join

### $CUSTOMER] \bowtie INVOICE$

```
SELECT CONCAT(C.cust_fname, c.cust_lname) as name,
LEFT JOIN INVOICE i
ON C.cust_num=i.cust_num
```

(b) Right Outer Join

#### $CUSTOMER \bowtie [INVOICE]$

```
SELECT CONCAT(C.cust_fname, c.cust_lname) as name,
RIGHT JOIN INVOICE i
ON C.cust_num=i.cust_num
```

(c) Full Outer Join

## $CUSTOMER] \bowtie [INVOICE]$

```
SELECT CONCAT(C.cust_fname, " ", C.cust_lname) as name

LEFT JOIN INVOICE i

ON C.cust_num=i.cust_num

UNION

SELECT CONCAT(C.cust_fname, " ", C.cust_lname) as name, i.inv_amount

RIGHT JOIN INVOICE i

ON C.cust_num=i.cust_num
```

- 5. Count number of products sold on each date.
  - (a) Relational Algebra

```
\pi_{\text{inv\_date},sum(\text{unit\_sold})}(\text{ inv\_date }\mathcal{F}_{sum(\text{unit\_sold})}(INVOICE))
```

```
SELECT inv_date, sum(unit_sold)
FROM INVOICE
GROUP BY inv_date
```

- 6. As soon as customer balance becomes greater than Rs. 100,000, copy the customer\_num in new table called "GOLD\_CUSTOMER"
  - (a) Create table GOLD\_CUSTOMER

```
CREATE TABLE GOLD_CUSTOMER

(
cust_num int,
PRIMARY KEY (cust_num),
FOREIGN KEY (cust_num) REFERENCES CUSTOMERS (cust_num)

)
```

(b) Create a trigger to update the GOLD\_CUSTOMER table when a new invoice record is entered for the customer.

```
CREATE TRIGGER update_gold_customer

AFTER INSERT ON INVOICE

FOR EACH ROW

BEGIN

IF NEW.cust_balance > 100000

AND NEW.cust_num NOT IN (SELECT cust_num FROM GOLD_CUSTOMER) THEN

INSERT INTO GOLD_CUSTOMER VALUES (NEW.cust_num);

END IF;

END
```

7. Add a new attribute CUST\_DOB in customer table

```
ALTER TABLE CUSTOMERS
ADD COLUMN cust_dob date
```

# References

[1] A. Einstein, "Zur Elektrodynamik bewegter Körper. (German) [On the electrodynamics of moving bodies]," *Annalen der Physik*, vol. 322, no. 10, pp. 891–921, 1905.