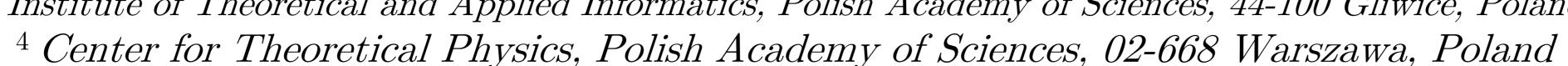
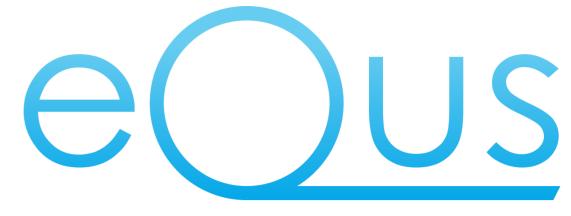
# Coherifying quantum channels

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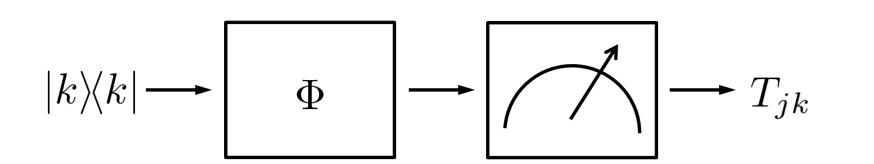




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#### Motivation

#### Classical action of a quantum channel



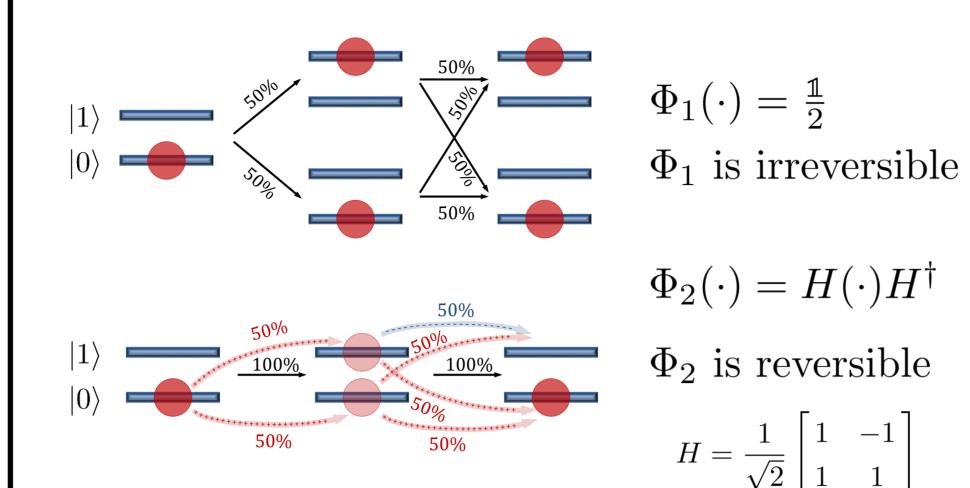
 $\{|k\rangle\}$  - distinguished orthonormal basis

 $T_{jk} = \langle j|\Phi(|k\rangle\langle k|)|j\rangle$  - classical action

What does T tell us about  $\Phi$ ?

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#### Channels with fixed classical action



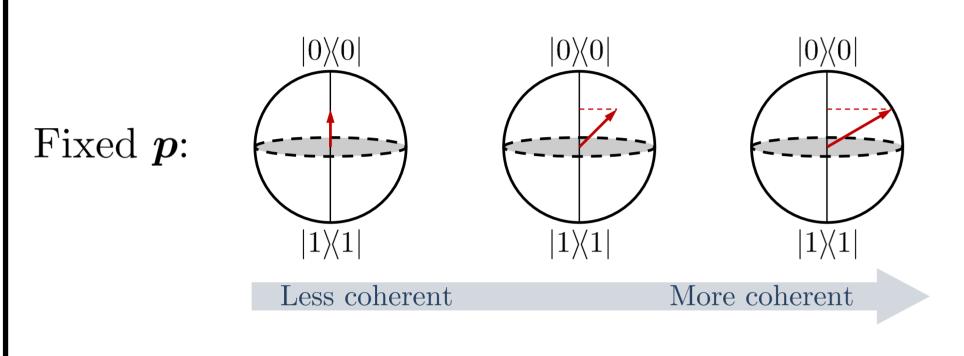
#### Questions

- 1. Is it always possible to explain random stochastic transitions as arising from the underlying deterministic quantum evolution?
- 2. If not, what is the minimal amount of randomness required by quantum theory to explain a given stochastic process?
- 3. And can there exist perfectly distinguishable quantum processes that nevertheless lead to the same classical evolution?

## Setting the scene

#### Coherence of quantum states

 $\langle j|\rho|j\rangle$  - occupations  $p_i$ ,  $\langle j|\rho|k\rangle$  - coherences

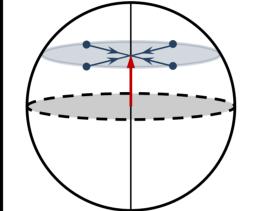


Coherence measures (distance from incoherent states):

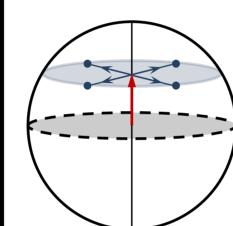
$$C_{\mathrm{e}}(\rho) := S(\rho||\mathcal{D}(\rho)) = S(\boldsymbol{p}) - S(\boldsymbol{\lambda}(\rho))$$
  
 $C_{2}(\rho) := ||\rho - \mathcal{D}(\rho)||_{2} = \boldsymbol{\lambda}(\rho) \cdot \boldsymbol{\lambda}(\rho) - \boldsymbol{p} \cdot \boldsymbol{p}$ 

 $\mathcal{D}$  - decohering channel,  $\lambda(\rho)$  - eigenvalues of  $\rho$ 

#### Decoherence and coherification



Decohering channel  $\mathcal{D}$ :  $\rho \text{ with } \langle j|\rho|j\rangle = p_j \xrightarrow{\mathcal{D}} \rho^{\mathcal{D}} = \text{diag}(\boldsymbol{p})$ 



Coherification C(non-unique inverse of D):  $\rho = \operatorname{diag}(\boldsymbol{p}) \xrightarrow{\mathcal{C}} \rho^{\mathcal{C}} \text{ with } \langle j|\rho|j\rangle = p_j$ 

One can always completely coherify  $\boldsymbol{p}$ :  $\operatorname{diag}(\boldsymbol{p}) \xrightarrow{\mathcal{C}} |\psi\rangle\langle\psi| \text{ with } |\psi\rangle = \sum_{j} \sqrt{p_{j}} e^{i\phi_{j}} |j\rangle$   $C_{e}(|\psi\rangle\langle\psi|) = S(\boldsymbol{p}), \quad C_{2}(|\psi\rangle\langle\psi|) = 1 - \boldsymbol{p} \cdot \boldsymbol{p}$ 

#### Coherence of quantum channels

Choi-Jamiołkowski  $J_{\Phi} = \frac{1}{d}(\Phi \otimes \mathcal{I}) |\Omega\rangle\langle\Omega|$  isomorphism:  $|\Omega\rangle = \sum_{j} |jj\rangle$ 

Classical action on diagonal:  $\langle jk|J_{\Phi}|jk\rangle = \frac{1}{d}T_{jk}$ 

Optimising coherence of  $\Phi \iff$  optimising  $\lambda(J_{\Phi})$ 

Relating randomness of  $\Phi$  with  $\lambda(J_{\Phi})$ :

$$|\psi\rangle \xrightarrow{\Phi} \frac{1}{\sqrt{q_j}} K_j |\psi\rangle$$
 with probability  $q_j$ ,

$$\Phi(\cdot) = \sum_{j} K_{j}(\cdot)K_{j}^{\dagger}, \quad q_{j} = \operatorname{Tr}\left(K_{j} |\psi\rangle\langle\psi|K_{j}^{\dagger}\right)$$

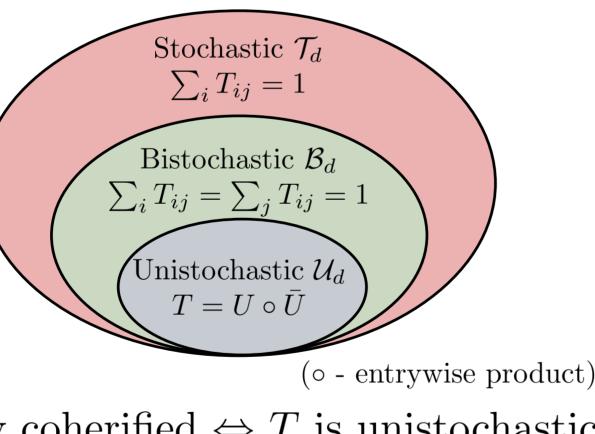
Path probability averaged over all pure states:  $\langle q_j \rangle_{\psi} = \lambda_j(J_{\Phi})$ 

## Limitations on perfect coherification

Can one always completely coherify T?  $T \xrightarrow{\mathcal{C}} |\psi\rangle\langle\psi| \text{ with } |\psi\rangle = \frac{1}{\sqrt{d}} \sum_{j,k} \sqrt{T_{jk}} e^{i\phi_{jk}} |jk\rangle$ 

No! TP condition requires  $\text{Tr}_1(|\psi\rangle\langle\psi|) = \frac{1}{d}$ 



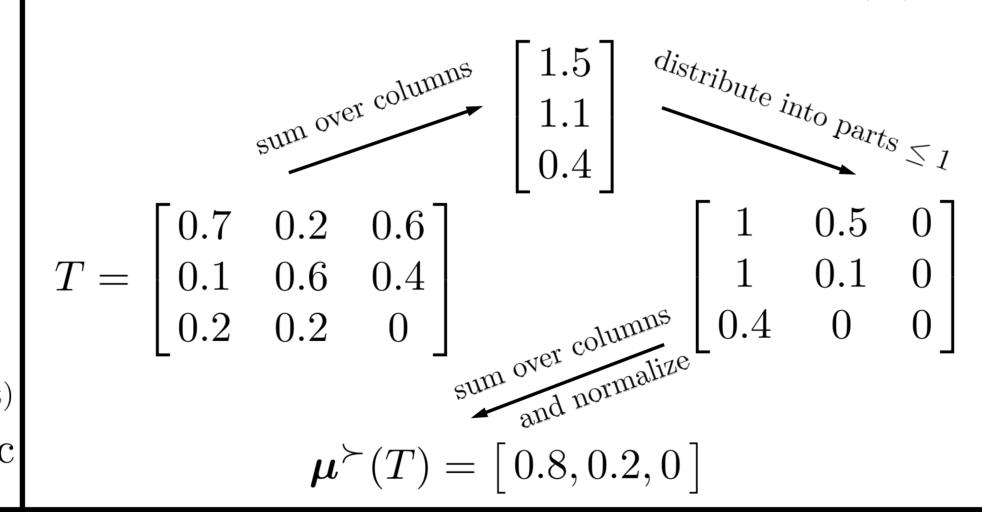


 $\Phi$  can be competely coherified  $\Leftrightarrow T$  is unistochastic

## Coherification upper-bound

Coherifying quantum channels

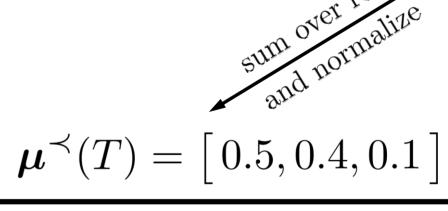
Look for  $\mu^{\succ}(T)$  s. t.  $\forall J_{\Phi}: \mu^{\succ}(T) \succ \lambda(J_{\Phi})$  $\mu^{\succ}$  yields upper bounds for  $C_{\mathrm{e}}(J_{\Phi})$  and  $C_{2}(J_{\Phi})$ Procedure to obtain upper-bounding  $\mu^{\succ}(T)$ :



### Coherification lower-bound

Look for  $\mu^{\prec}(T)$  s. t.  $\exists J_{\Phi}: \mu^{\prec}(T) \prec \lambda(J_{\Phi})$  $\mu^{\prec}$  yields lower bounds for  $C_{\mathrm{e}}(J_{\Phi})$  and  $C_{2}(J_{\Phi})$ Procedure to obtain lower-bounding  $\mu^{\prec}(T)$ :

$$T = \begin{bmatrix} 0.7 & 0.2 & 0.6 \\ 0.1 & 0.6 & 0.4 \\ 0.2 & 0.2 & 0 \end{bmatrix} \xrightarrow{\text{order}} \begin{bmatrix} 0.7 & 0.6 & 0.2 \\ 0.6 & 0.4 & 0.1 \\ 0.2 & 0.2 & 0 \end{bmatrix}$$

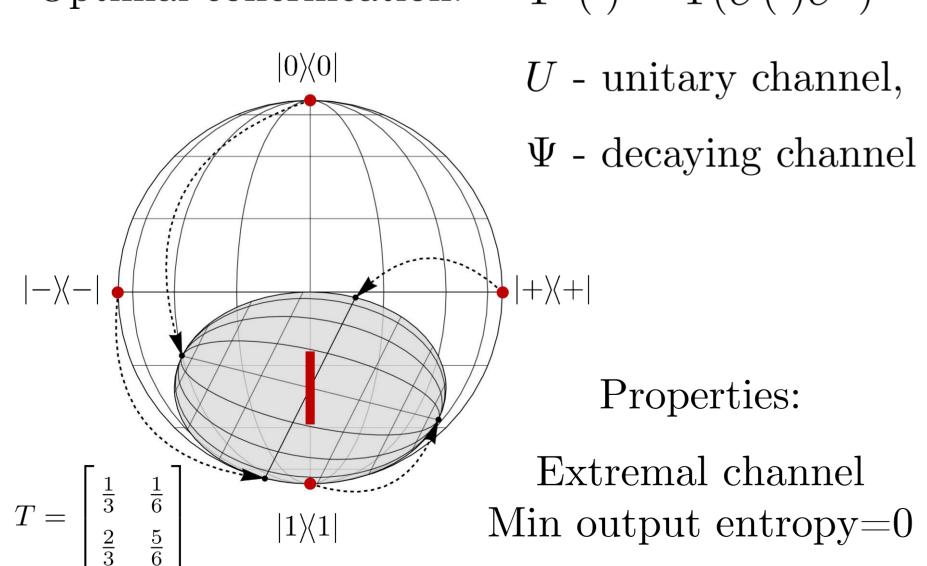


## Example: qubit channel

Classical action:

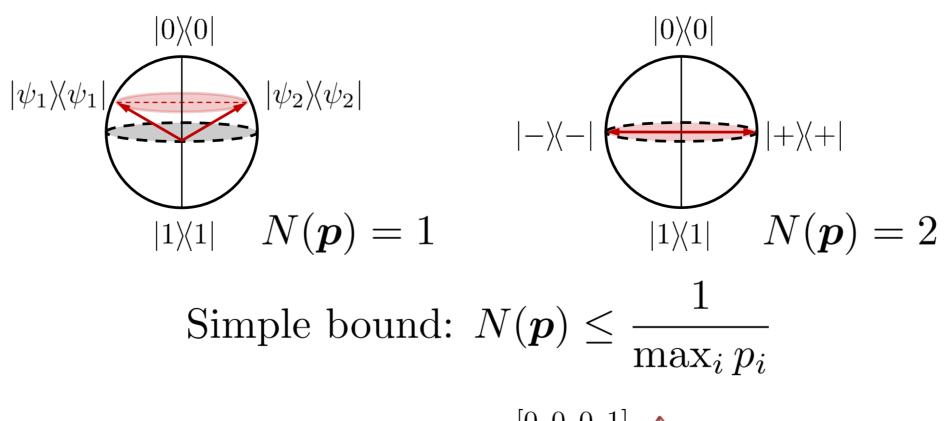
$$T = \begin{bmatrix} a & 1-b \\ 1-a & b \end{bmatrix}$$

Optimal coherification:  $\Phi^{\mathcal{C}}(\cdot) = \Psi(U(\cdot)U^{\dagger})$ 

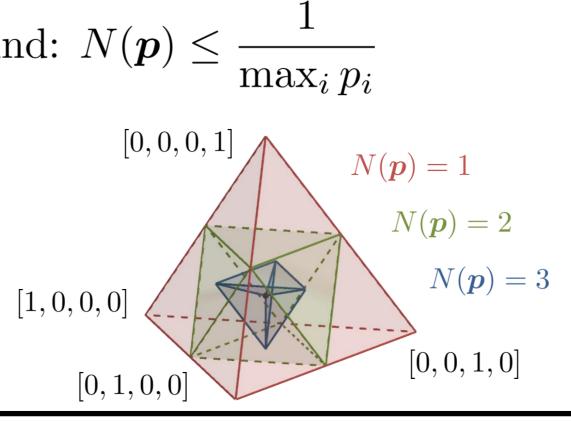


## Quantum states

What is the number  $N(\mathbf{p})$  of perfectly distinguishable states with classical version  $\mathbf{p}$ ?



But things get more complicated. E.g. for d = 4:



# Number of distinct coherifications

What is the number N(T) of perfectly distinguishable channels with classical version T?

Quantum channels

 $1 \leq N(T) \leq d^2$ , both limits achieveable

$1 \leq 1 \vee (1) \leq \alpha$ , both lilling active able		
Classical action	N(T)	Requires entanglement?
Unistochastic	d	No
Unistochastic	$d+1,\ldots,d^2$	Yes
Bistochastic	2	Yes
Circulant	d	No
S.t. $T_{ij} \leq \frac{1}{2}$	2	No

More soon on arXiv! Look for: Distinguishing classically indistinguishable states and channels