```
else n = n * 10 + c - '0':
    General
                              7 Graphs
    Algorithms
                              8 2D Geometry
                                                             n = s * (n + m * o):
    Structures
                              9 3D Geometry
                                                            void read(double& n) {
    Strings
                                                             ld m; read(m); n = m;
                              10 Optimization
    Greedy
                              11 Additional
                                                            void read(float& n) {
 ld m: read(m): n = m:
    Math
     General
                                                            void read(string& s) {
                                                             char c; s = "
g++ -g -02 -std=gnu++17 -static prog.cpp
./a.exe
run.sh
                                                             while((c=getchar unlocked())!=' '&&c!='\n')
                                                            bool readline(string& s) {
                                                             char c; s = "";
while(c=getchar unlocked()) {
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp for i i *.in; do
                                                              if (c == '\n') return true;
if (c == EOF) return false;
s += c;
f=${i%.in}
f=${i%.in}
./a.exe < $i > "$f.out"
.diff -b -q "$f.ans" "$f.out"
                                                             return false;
done
                                                            void print(unsigned int n) {
Header
                                                             if (n / 10) print(n / 10);
// use better compiler options
                                                             putchar_unlocked(n % 10 + '0');
#pragma GCC optimize("Ofast","unroll-loops")
#pragma GCC target("avx2,fma")
                                                            void print(int n) {
// include everything
                                                             if (n < 0) { putchar_unlocked('-'); n*=-1; }
 #include <bits/stdc++.h>
                                                             print((unsigned int)n);
#include <bits/extc++.h>
#include <sys/resource.h>
// namespaces
                                                            Common Structs
using namespace std;
                                                                n-dimension vectors
using namespace __gnu_cxx; // rope
                                                                Vec<2, int>v(n, m) = arr[n][m]
using namespace __gnu_pbds; // tree/trie
                                                             // Vec<2, int> v(n, m, -1) default init -1
                                                            template<int D, typename T>
// common defines
#define fastio
                                                            struct Vec : public vector < Vec < D-1, T >> {
                                                               template<typename... Args>

→ ios base::sync with stdio(0);cin.tie(0);
                                                               Vec(int n=0, Args... args) : vector<Vec<D-1,
#define nostacklim rlimit RZ; getrlimit(3,&RZ
                                                             \rightarrow T>>(n, Vec<D-1, T>(args...)) {}
    ):RZ.rlim cur=-1:setrlimit(3.&RZ):
#define DEBUG(v) cerr<< LINE <<": "<<#v<<" =
                                                            template<typename T>
\Rightarrow "<<v<<'\n'; #define TIMER
                                                            struct Vec<1, T> : public vector<T> {
                                                               Vec(int n=0, T val=T()) : vector<T>(n, val)

→ cerr<<1.0*clock()/CLOCKS_PER_SEC<<"s\n";
#define ll long long
#define ull unsigned ll
#define i128 __int128
#define u128 unsigned i128
                                                                {}
                                                                 Algorithms
#define ld long double
                                                            Min/Max Subarray
// global variables
                                                               max - compare = a < b, reset = a < 0
mt19937 rng((uint32_t)chrono::steady
                                                             \frac{1}{min} - compare = a > b, reset = a > 0
                                                            // returns {sum, {start, end}}
pair<int, pair<int, int>>

    clock::now().time since epoch().count());

Fast IO
                                                                 ContiguousSubarray(int* a, int size,
#ifdef _WIN32
                                                                 bool(*compare)(int, int),
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
                                                             bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
                                                             0, s = 0;
for (int i = 0; i < size; i++) {
  cur += a[i];</pre>
void read(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                               if ((*compare)(best, cur)) { best = cur;
  n = n * 10 + c - 0';
                                                             \rightarrow start = s; end = i; }
void read(int& n) {
   char c; n = 0; int s = 1;
   if ((c=getchar_unlocked())=='-') s = -1;
                                                              if ((*reset)(cur)) { cur = 0; s = i + 1; }
                                                             return {best, {start, end}}:
 else n = c - '0';
while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                            Quickselect
 n = n * 10 + c - 0;

n *= s;
                                                            #define OSNE -999999
                                                            int partition(int arr[], int 1, int r)
void read(ld& n) {
 char c; n = 0;
ld m = 0, o = 1; bool d = false; int s = 1;
if ((c=getchar_unlocked())=='-') s = -1;
                                                             int x = arr[r], i = 1;
for (int j = 1; j <= r - 1; j++)
...if (arr[j] <= x)
...swap(arr[i++], arr[j]);</pre>
 else if (c == '.') d = true;
else n = c - '0';
while ((c=getchar_unlocked())!=' '&&c!='\n') {
                                                             swap(arr[i], arr[r]);
  if (c == '.') d = true;
else if (d) { m=m*10+c-'0'; o*=0.1; }
                                                             return i:
```

```
|// find k'th smallest element in unsorted array,
→ only if all distinct
int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - l + 1)) return QSNE;
swap(arr[1 + rng() % (r-l+1)], arr[r]);
 int pos = partition(arr, 1, r);
if (pos-l==k-1) return arr[pos];
 if (pos-1>k-1) return qselect(arr,1,pos-1,k);
 return qselect(arr, pos+1, r, k-pos+1-1);
|}
|// TODO: compare against std::nth_element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted
→ on both axis
pair<int, int> saddleback_search(int** arr, int
 \stackrel{\cdot}{\hookrightarrow} x, int y, int v) {
 int i = x-1, j = 0;
while (i >= 0 && j < y) {
  if (arr[i][j] == v) return {i, j};
  (arr[i][i] > v)? i--: i++:
 return {-1, -1}:
 Ternary Search
 // < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a, int b, int (*f)(int)) {
 while (b-a > 4) {
  int m = (a+b)/2
  if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1:
  for (int i = a+1; i <= b; i++)
  if (TERNCOMP((*f)(a), (*f)(i)))
 ...a = i;
return a:
#define TERNPREC 0.000001
double ternsearch (double a, double b, double
 (*f)(double)) {
while (b-a > TERNPREC * 4) {
    double m = (a+b)/2;
    if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
  = m;
else b = m + TERNPREC;
  for (double i = a + TERNPREC; i <= b; i +=
    TERNPREC)
      if (TERNCOMP((*f)(a), (*f)(i)))
    a = i:
 return á;
     Structures
Fenwick Tree
// Fenwick tree, array of cumulative sums -
 \rightarrow O(log n) updates, O(log n) gets
struct Fenwick {
 int n; ll* tree;
  void update(int i, int val) {
  .++i;
  while (i <= n) {
   tree[i] += val;
   i += i & (-i);
 Fenwick(int size) {
  | n = size;
| tree = new | l[n+1];
| for (int i = 1; i <= n; i++)
| tree[i] = 0;
  Fenwick(int* arr, int size) : Fenwick(size) {
  for (int i = 0; i < n; i++)
update(i, arr[i]);
  ~Fenwick() { delete[] tree; }
 ll operator[](int i) {
```

```
.if (i < 0 || i > n) return 0;
.ll sum = 0;
  ++i;
  while (i>0)
  sum += trée[i];
   i = i & (-i);
  return sum:
 ll getRange(int a, int b) { return
operator[](b) - operator[](a-1); };
Hashtable
 // similar to unordered map, but faster
struct chash {
    const uint64_t C = (11)(2e18 * M_PI) + 71;
 ll operator()(ll x) const { return
    builtin bswap64(x*C); }
int main() {
  gp_hash_table<11,int,chash>
 \rightarrow hashtable({},{},{},{},{1<<16});
 for (int i = 0; i < 100; i++)
hashtable[i] = 200+i;
 if (hashtable.find(10) != hashtable.end())
  cout << hashtable[10];
 Ordered Set
template <typename T>
using oset = tree<T,null_type,less<T>,rb_tree

tag, tree_order_statistics_node_update>;
template <typename T, typename D>
using omap = tree<T,D,less<T>,rb_tree |
    _tag,tree_order_statistics_node_update>;
int main()
 oset<int> o_set;
 o set.insert(5); o set.insert(1);
 \rightarrow o set.insert(3);
 // get second smallest element
 cout << *(o set.find by order(1));</pre>
 // number of elements less than k=4
 cout << ' ' << o_set.order_of_key(4) << '\n';
 // equivalent with ordered map
 omap<int.int> o map:
 o_map[5]=1;o_map[1]=2;o_map[3]=3;
 cout << (*(o_map.find_by_order(1))).first;
cout << ' ' << o_map.order_of_key(4) << '\n';</pre>
Rope
 // O(log n) insert, delete, concatenate
int main() {
 // generate rope
 rope<int> v;
 for (int i = 0; i < 100; i++)
  v.push back(i):
 // move range to front
 rope<int> copy = v.substr(10, 10);
 v.erase(10, 10);
 v.insert(copy.mutable_begin(), copy);
 // print elements of rope
 for (auto it : v) cout << it << " ":
Segment Tree
 //max(a,b), min(a,b), a+b, a*b, gcd(a,b), a\hat{b}
struct SegmentTree {
 typedef int T;
 static constexpr T UNIT = INT_MIN;
 T f(T a, T b) {
    if (a == UNIT) return b;
    if (b == UNIT) return a;
  return max(a,b);
 int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
```

```
SegmentTree(vector<T> arr) :

→ SegmentTree(arr.size()) {
 for (int i=0;i<arr.size();i++)

→ update(i,arr[i]);

 void update(int pos, T val) {
 for (s[pos += n] = val; pos /= 2;)
  s[pos] = f(s[pos * 2], s[pos*2+1]);
 T query(int b, int e) { // query [b, e) }
T ra = UNIT, rb = UNIT;
 for (b+=n, e+=n; b<e; b/=2, e/=2) {
    if (b % 2) ra = f(ra, s[b++]);
    if (e % 2) rb = f(s[--e], rb);
 return f(ra, rb);
 T get(int p) { return query(p, p+1); }
Trie
typedef trie<string, null_type,

→ trie string access traits<>,

 pat_trie_tag, trie_prefix_search_node_update>

→ trie_type;

int main() {
   // generate trie
 trie_type trie;
 for (int i = 0; i < 20; i++)
 trie.insert(to_string(i)); // true if new,
\hookrightarrow false if old
 // print things with prefix "1"
 auto range = trie.prefix_range("1");
for (auto it = range.first; it !=
\hookrightarrow range.second; it++)
 .cout << *it << "
4 Strings
```

```
Aho Corasick
// range of alphabet for automata to consider
// MAXC = 26. OFFC = 'a' if only lowercase
const int MAXC = 256;
const int OFFC = 0:
struct aho_corasick {
 struct state
  set<pair<int, int>> out:
 int fail; vector<int> go;
  state(): fail(-1), go(MAXC, -1) {}
 };
 vector<state> s;
 int id = 0;
 aho_corasick(string* arr, int size) : s(1) {
 for (int i = 0: i < size: i++) {
   int cur = 0;
  .for (int c : arr[i]) {
...if (s[cur].go[c-OFFC] == -1) {
   s[cur].go[c-OFFC] = s.size();
    s.push back(state());
    cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++});
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
   ..s[0].go[\tilde{c}] = 0;
  queue<int> sq;
 for (int c = 0; c < MAXC; c++) {
    if (s[0].go[c] != 0) {
        s[s[0].go[c]].fail = 0;
    sq.push(s[0].go[c]);
  while (sq.size()) {
 int e = sq.front(); sq.pop();
 for (int c = 0; c < MAXC; c++) {
...if (s[e].go[c] != -1) {
```

```
int failure = s[e].fail;
while (s[failure].go[c] == -1)
      failure = s[failure].fail;
failure = s[failure].go[c];
      s[s[e].go[c]].fail = failure;
  for (auto length : s[failure].out)
s[s[e].go[c]].out.insert(length);
     sq.push(s[e].go[c]);
 // list of {start pos, pattern id}
 vector<pair<int, int>> search(string text)
  vector<pair<int, int>> toret;
  int cur = 0;
  for (int i = 0; i < text.size(); i++) {
   while (s[cur].go[text[i]-OFFC] == -1)
    .cur = s[cur].fail;
   cur = s[cur].go[text[i]-OFFC];
   if (s[cur].out.size())
    for (auto end : s[cur].out)
  toret.push_back({i - end.first + 1,
     end.second});
  return toret:
Bover Moore
struct defint { int i = -1; };
vector<int> boyermoore(string txt, string pat)
 vector<int> toret; unordered_map<char, defint> Longest Common Prefix (array)
 → badchar:
 int m = pat.size(), n = txt.size();
 for (int i = 0; i < m; i++) badchar[pat[i]].i
 \rightarrow = i;
int s = 0:
 while (s \leq n - m) {
  int j = m - 1;
  while (j \ge 0) && pat[j] == txt[s + j]) j--;
  if (i < 0) {
   .toret.push back(s);
   s += (s + m < n) ? m - badchar[txt[s +
 \rightarrow mll.i : 1:
  .} else
   s += max(1, j - badchar[txt[s + j]].i);
 return toret;
English Conversion
const string ones[] = {"", "one", "two",
"three", "four", "five", "six", "seven",

"eight", "nine";

const string teens[] ={"ten", "eleven",
   "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen"};
const string tens[] = {"twenty", "thirty",
"forty", "fifty", "sixty", "seventy", 

"eighty", "ninety"};
const string mags[] = {"thousand", "million",
     "billion", "trillion", "quadrillion",
     "quintillion", "sextillion",
string convert(int num, int carry) {
 if (num < 0) return "negative " +
     convert(-num, 0);
     (num < 10) return ones[num];
(num < 20) return teens[num % 10];</pre>
     (\text{num} < 100) \text{ return tens}[(\text{num} / 10) - 2] +
     (num%10==0?"":" ") + ones[num % 10];
     (num < 1000) return ones[num / 100]
     (num/100==0?"":" ") + "hundred" + (num%100==0?"":" ") + convert(num % 100,
```

```
return convert(num / 1000, carry + 1) + " " + |...while (i >= j && i + j + 1 < n && s[(i-j)/2]
     mags[carry] + " " + convert(num % 1000.
    0):
string convert(int num) {
return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
     vector<int> toret;
 int m = txt.length(), n = pat.length();
 int next[n + 1];
for (int i = 0; i < n + 1; i++)
   next[i] = 0;</pre>
 int i = 1; i < n; i++) {
  int j = next[i + 1];
  while (j > 0 && pat[j] != pat[i])
   j = next[j];
  if (j > 0 || pat[j] == pat[i])
  next[i + 1] = i + 1;
 for (int i = 0, j = 0; i < m; i++) {
  if (txt[i] == pat[j]) {
   if (++j == n)
    toret.push_back(i - j + 1);
  .} else if (j > 0) {
  .j = next[j];
 return toret:
// longest common prefix of strings in array
string lcp(string* arr, int n, bool sorted =
false) {
if (n == 0) return "";
 if (!sorted) sort(arr, arr + n);
string r = ""; int v = 0;
 while (v < arr[0].length() && arr[0][v] ==

    arr[n-1][v])
    r += arr[0][v++];

 return r;
Longest Common Subsequence
string lcs(string a, string b) {
 int m = a.length(), n = b.length();
 int L[m+1][n+1];
 for (int i = 0; i <= m; i++) {
    for (int j = 0; j <= n; j++) {
        if (i == 0 || j == 0) L[i][j] = 0;
        else if (a[i-1] == b[j-1]) L[i][j] =
 \rightarrow L[i-1][j-1]+1;
   else L[i][j] = \max(L[i-1][j], L[i][j-1]);
 // return L[m][n]; // length of lcs
 string out = "":
 int i = m - 1, j = n - 1;
while (i >= 0 && j >= 0) {
   if (a[i] == b[j]) {
   out = a[i--] + out;
  else if (L[i][j+1] > L[i+1][j]) i--;
  else j--;
 return out;
Longest Common Substring
// l is array of palindrome length at that
→ index
int manacher(string s. int* 1) {
 int n = s.length() * 2;
 for (int i = \overline{0}, j = 0, k; i < n; i += k, j =
```

 \rightarrow max(i-k, 0)) {

```
for (k = 1; i >= k && j >= k && l[i-k] !=
 \rightarrow j-k; k++)
  1[i+k] = min(1[i-k], j-k);
return *max_element(1, 1 + n);
Cyclic Rotation (Lyndon)
// simple strings = smaller than its nontrivial
    suffixes
// lyndon factorization = simple strings
→ factorized
// "abaaba" -> "ab", "aab", "a"
vector<string> duval(string s) {
int n = s.length();
vector<string> lyndon;
for (int i = 0; i < n;) {
 int j = i+1, k = i;

int j = i+1, k = i;

for (; j < n && s[k] <= s[j]; j++)

if (s[k] < s[j]) k = i;
   else k++:
  for (; i \le k; i += j - k)
  lyndon.push back(s.substr(i,j-k));
return lyndon;
// lexicographically smallest rotation
int minRotation(string s) {
int n = s.length(); s += s;
auto d = duval(s); int i = 0, a = 0;
while (a + d[i].length() < n) a +=</pre>
 \rightarrow d[i++].length();
while (i && d[i] == d[i-1]) a -=
→ d[i--].length();
return a;
Subsequence Count
// "banana", "ban" >> 3 (ban, ba..n, b..an)
ull subsequences(string body, string subs) {
int m = subs.length(), n = body.length();
if (m > n) return 0;
 ull** arr = new ull*[m+1];
for (int i = 0; i \le m; i++) arr[i] = new
\hookrightarrow ull[n+1];
for (int i = 1; i <= m; i++) arr[i][0] = 0;
for (int i = 0; i <= n; i++) arr[o][i] = 1;
for (int i = 1; i <= m; i++)
 for (int j = 1; j <= n; j++)
arr[i][j] = arr[i][j-1] + ((body[j-1] ==
\hookrightarrow subs[i-1])? arr[i-1][j-1] : 0);
return arr[m][n]:
Suffix Array + LCP
struct SuffixArray {
vector<int> sa, 1cp;
SuffixArray(string& s, int lim=256) {
   int n = s.length() + 1, k = 0, a, b;
   vector<int> x(begin(s), end(s)+1), y(n),
 \rightarrow ws(max(n, lim)), rank(n);
 sa = lcp = y;
iota(begin(sa), end(sa), 0);
  for (int j = 0, p = 0; p < n; j = max(1, j *
\rightarrow 2), lim = p) {
   p = j; iota(begin(y), end(y), n - j);
  for (int i = 0; i < (n); i++)
if (sa[i] >= j)
y[p++] = sa[i] - j;
   fill(begin(ws), end(ws), 0);
  for (int i = 0; i < (n); i++) ws[x[i]]++;
for (int i = 1; i < (lim); i++) ws[i] +=
```

 $\Rightarrow == s[(i+j+1)/2]) j++;$

1[i] = j;

→ ws[i - 1]:

```
. for (int i = n; i--;) sa[-ws[x[y[i]]]] =
                                                          void uppercase(string& s) {
\hookrightarrow y[i];
                                                           transform(s.begin(), s.end(), s.begin(),
   ::toupper);
                                                          void trim(string &s) {
                                                           s.erase(s.begin(),find_if_not(s.begin(),s
.end(),[](int c){return
                                                              isspace(c);}));
                                                           isspace(c);}));
s.erase(find_if_not(s.rbegin(),s.rend(),[](int) | 11 multinomial(vector<int>& v) {
            | 11 c = 1, m = v.empty() ? 1 : v[0];
  for (int i = 1; i < (n); i++) rank[sa[i]] =
                                                              c){return isspace(c);}).base(),s.end());
 for (int i = 0, j; i < n - 1; lcp[rank[i++]]
                                                           vector<string> split(string& s, char token) {
                                                               vector<string> v; stringstream ss(s);
for (k \&\& k--, j = sa[rank[i] - 1];
     s[i + k] = s[j + k]; k++);
                                                               for (string e;getline(ss,e,token);)
                                                               v.push_back(e);
return v;
Suffix Tree (Ukkonen's)
                                                               Greedy
struct SuffixTree {
  enum { N = 50000, ALPHA = 26 };
  int toi(char c) { return c - 'a'; }
                                                          Interval Cover
                                                           //L,R = interval [L,R], in = \{\{l,r\}, index\}
ting at t[N] [ALPHA],1[N],r[N],p[N],s[N],v=0,q=0,m=2 // does not handle case where L == R vector<int> intervalCover(double L, double R
 string a;
                                                           → vector<pair<pair<double.double>.int>> in)
 void ukkadd(int i, int c) { suff:
 void ukkadd(int 1, int c) { suff:
   if (r[v] <=q) {
    if (t[v][c]==-1) { t[v][c]=m; l[m]=i;
        p[m++]=v; v=s[v]; q=r[v]; goto suff; }
   v=t[v][c]; q=l[v];</pre>
                                                               int i = 0; pair < double, int > pos = {L,-1};
                                                              vector<int> a;
sort(begin(in), end(in));
                                                               while (pos.first < R) {
                                                                    double cur = pos.first;
while (i < (int)in.size() &&</pre>
  if (q==-1 | c==toi(a[q])) q++; else {
  in[i].first.first <= cur)
                                                               max(pos,{in[i].first.second,in[i].second})
   i++;
                                                                    if (pos.first == cur) return {};
                                                                    a.push back(pos.second):
\rightarrow q+=r[v]-l[v]; }
   if (q=r[m]) s[m]=v; else s[m]=m+2;
                                                               return a;
   q=r[v]-(q-r[m]); m+=2; goto suff;
                                                               Math
 SuffixTree(string a) : a(a) {
  fill(r,r+N,(int)(a).size());
                                                          Catalan Numbers
                                                          ull* catalan = new ull[1000000];
 memset(s, 0, sizeof s);

memset(t, -1, sizeof t);

fill(t[1],t[1]+ALPHA,0);

s[0]=1;1[0]=1[1]=-1;r[0]=r[1]=p[0]=p[1]=0;
                                                          void genCatalan(int n, int mod) {
  catalan[0] = catalan[1] = 1;
                                                           for (int i = 2; i <= n; i++) {
    catalan[i] = 0;
    for (int j = i - 1; j >= 0; j--) {
 for(int i=0;i<a.size();i++)

    ukkadd(i,toi(a[i]));
                                                             catalan[i] += (catalan[j] * catalan[i-j-1])
                                                               % mod;
 // Longest Common Substring between 2 strings
                                                             if (catalan[i] >= mod)
  catalan[i] -= mod;
 // returns {length, offset from first string}
 pair<int, int> best;
int lcs(int node, int i1, int i2, int olen) {
   if (1[node] <= i1 && i1 < r[node]) return 1;
   if (1[node] <= i2 && i2 < r[node]) return 2;</pre>
                                                           // TODO: consider binomial coefficient method
  int mask=0
c len=node?olen+(r[node]-l[node]):0;
...for(int c=0; c<ALPHA; c++) if</pre>
                                                           Combinatorics (nCr. nPr)
                                                           // can optimize by precomputing factorials, and
\hookrightarrow (t[node][c]!=-1)
                                                           \hookrightarrow fact[n]/fact[n-r]
   mask |= lcs(t[node][c], i1, i2, len);
  if (mask==3)
                                                           ull nPr(ull n, ull r) {
→ best=max(best,{len,r[node]-len});
                                                           for (ull i = n-r+1: i <= n: i++)
 return mask;
                                                           v *= i;
return v:
 static pair<int, int> LCS(string s, string t)

    SuffixTree

                                                          ull nPr(ull n, ull r, ull m) {
→ st(s+(char)('z'+1)+t+(char)('z'+2));

st.lcs(0, s.size(), s.size()+t.size()+1, 0);

return st.best;
                                                           for (ull i = n-r+1; i <= n; i++)
                                                            v = (v * i) \% m;
                                                           return v;
                                                          ull nCr(ull n, ull r) {
long double v = 1;
String Utilities
                                                           for (ull i = 1; i <= r; i++)
void lowercase(string& s) {
                                                            v = v * (n-r+i) /i;
transform(s.begin(), s.end(), s.begin(),
                                                           return (ull)(v + 0.001);
   ::tolower);
                                                           // requires modulo math
```

```
|// can optimize by precomputing mfac and

→ minv-mfac

ull nCr(ull n, ull r, ull m) {
 return mfac(n, m) * minv(mfac(k, m), m) % m
    minv(mfac(n-k, m), m) \% m:
Multinomials
 for(int i = 1; i < v.size(); i++)

for (int j = 0; j < v[i]; j++)

...c = c * ++m / (j+1);
 return c:
Chinese Remainder Theorem
bool ecrt(ll* r, ll* m, int n, ll& re, ll& mo)
 11 x, y, d; mo = m[0]; re = r[0];
 for (int i = 1; i < n; i++) {
  d = egcd(mo, m[i], x, y);

if ((r[i] - re) % d != 0) return false;

x = (r[i] - re) / d * x % (m[i] / d);

re += x * mo;
  mo = mo / d * m[i];
  re %= mo;
 re = (re + mo) % mo;
return true;
Count Digit Occurences
/*count(n,d) counts the number of occurences of
 \rightarrow a digit d in the range \lceil 0.n \rceil * /
ll digit_count(ll n, ll d) {
 ll result = 0:
 while (n != 0) {
  result += ((n%10) == d ? 1 : 0);
  n /= 10:
 return result:
11 count(11 n, 11 d) {
   if (n < 10) return (d > 0 && n >= d);
 if ((n % 10) != 9) return digit_count(n, d) +
 \rightarrow count(n-1, d);
 return 10*count(n/10, d) + (n/10) + (d > 0):
Discrete Logarithm
unordered map<int, int> dlogc:
 int discretelog(int a. int b. int m) {
 dlogc.clear():
 11 \ \tilde{n} = \operatorname{sqrt}(\tilde{m}) + 1, an = 1;
 for (int i = 0; i < n; i++)
an = (an * a) % m;
  11 c = an;
 for (int i = 1; i <= n; i++) {
   if (!dlogc.count(c)) dlogc[c] = i;
  c = (c * an) \% m;
  c = b
 for (int i = 0; i <= n; i++) {
  if (dlogc.count(c)) return (dlogc[c] * n - i</pre>
 \rightarrow + m - 1) % (m-1):
  c = (c * a) \% m;
 return -1:
Euler Phi / Totient
int phi(int n) {
 int^r = n;
 for (int i = 2: i * i <= n: i++) {
  if (n % i == 0) r -= r / i;
while (n % i == 0) n /= i;
 if (n > 1) r = r / n;
 return r;
#define n 100000
```

```
|ll phi[n+1];
void computeTotient() {
 for (int i=1; i<=n; i++) phi[i] = i;
 for (int p=2; p<=n; p++) {
  .if (phi[p] == p) {
   phi[p] = p-1;
for (int i = 2*p; i<=n; i += p) phi[i] =
    (phi[i]/p) * (p-1);
Factorials
// digits in factorial
#define kamenetsky(n) (floor((n * log10(n /
\hookrightarrow ME)) + (log10(2 * MPI * n) / 2.0)) + 1)
// approximation of factorial
#define stirling(n) ((n == 1) ? 1 : sqrt(2 *
\hookrightarrow M PI * n) * pow(n / M E, n))
// natural log of factorial
#define lfactorial(n) (lgamma(n+1))
Prime Factorization
// do not call directly
ll pollard_rho(ll n, ll s) {
ll x, y;
x = y = rand() % (n - 1) + 1;
 int head = 1, tail = 2;
 while (true) {
 x = mult(x, x, n);

x = (x + s) \% n;
  if (x == y) return n;
  11 d = _gcd(max(x - y, y - x), n);
if (1 < d && d < n) return d;
  if (++head == tail) y = x, tail <<= 1;
// call for prime factors
void factorize(ll n, vector<ll> &divisor) {
 if (n == 1) return;
 if (isPrime(n)) divisor.push back(n);
  while (d'>= n) d = pollard rho(n, rand() % (n
 - 1) + 1);
factorize(n / d, divisor);
factorize(d, divisor);
Farev Fractions
   generate 0 \le a/b \le 1 ordered. b \le n
   farey(4) = 0/1 \ 1/4 \ 1/3 \ 1/2 \ 2/3 \ 3/4 \ 1/1
 // length is sum of phi(i) for i = 1 to n
vector<pair<int, int>> farev(int n) {
 int h = 0, k = 1, x = 1, y = 0, r;
vector<pair<int, int>> v;
 do {
 v.push back({h, k});
  r = (n-y)/k;
  y += r*k; x' += r*h;
  swap(x,h); swap(y,k);
 x = -x; y = -y;
} while (k > 1);
 v.push_back({1, 1});
 return v:
Fast Fourier Transform
#define cd complex<double>
const double PI = acos(-1):
void fft(vector<cd>& a, bool invert) {
 int n = a.size();
 for (int i = 1, j = 0; i < n; i++) {
  .int bit = n >> 1;
  .for (; j & bit; bit >>= 1) j ^= bit;
.j ^= bit;
  if (i < j) swap(a[i], a[j]);
 for (int len = 2; len <= n; len <<= 1) {
```

```
.double ang = 2 * PI / len * (invert ? -1 :
                                                       \overline{11} x = 0, y = 0;
 → 1):
                                                         if (egcd(b, m, x, y) != 1) return -1;
  cd wlen(cos(ang), sin(ang));
                                                         return (x % m + m) % m;
  for (int i = 0; i < n; i += len) {
  cd w(1);
   for (int j = 0; j < len / 2; j++) {
    cd u = a[i+j], v = a[i+j+len/2] * w;
                                                        11 mdiv_compmod(int a, int b, int m) {
                                                         if (__gcd(b, m) != 1) return -1;
                                                         return mult(a, minv(b, m), m);
   a[i+j] = u + v;
  a[i+j+len/2] = u - v;
                                                        // if m is prime (like 10^9+7)
   w = wlen;
                                                        ll mdiv_primemod (int a, int b, int m) {
                                                         return mult(a, mpow(b, m-2, m), m);
 if (invert)
                                                        Modulo Tetration
 for (auto& x : a)
  x /= n;
                                                        ll tetraloop(ll a, ll b, ll m) {
                                                         if(b == 0 | | a == 1) return 1;
ll w = tetraloop(a,b-1,phi(m)), r = 1;
vector<int> fftmult(vector<int> const& a.
                                                         for (:w:w/=2) {

    vector<int> const& b) {

                                                          if (w&1) {
    r = a; if (r >= m) r -= (r/m-1)*m;
 vector<cd> fa(a.begin(), a.end()),

    fb(b.begin(), b.end());

 int n = 1 << (32 - __builtin_clz(a.size() +</pre>
                                                          a *= a: if (a >= m) a -= (a/m-1)*m:
 \rightarrow b.size() - 1));
 fa.resize(n); fb.resize(n);
                                                         return r:
 fft(fa, false); fft(fb, false);
 for (int i = 0; i < n; i++) fa[i] *= fb[i];
                                                        int tetration(int a, int b, int m) {
  if (a == 0 || m == 1) return ((b+1)&1)%m;
 fft(fa, true);
                                                         return tetraloop(a,b,m) % m;
 vector<int> toret(n);
 for (int i = 0; i < n; i++) toret[i] =

→ round(fa[i].real());
                                                        Matrix
 return toret:
                                                        template<typename T>
                                                        struct Mat : public Vec<2. T> {
                                                         int w, h;
Greatest Common Denominator
                                                         Mat(int x, int y) : Vec<2, T>(x, y), w(x),
ll egcd(ll a, ll b, ll& x, ll& y) {
    if (b == 0) { x = 1; y = 0; return a; }
    ll gcd = egcd(b, a % b, x, y);
                                                         static Mat<T> identity(int n) { Mat<T> m(n,n);
 x = a / b * y;
                                                            for (int i=0;i<n;i++) m[i][i] = 1; return
 swap(x, y);
                                                            m; }
                                                         Mat<\hat{T}>\& operator+=(const Mat<T>\& m) {
 return gcd;
                                                          for (int i = 0; i < w; i++)
                                                           for (int j = 0; j < h; j++)
  (*this)[i][j] += m[i][j];
Josephus Problem
                                                          return *this:
// 0-indexed, arbitrary k
int josephus(int n. int k) {
                                                         Mat<T>& operator = (const Mat<T>& m) {
 if (n == 1) return 0;
if (k == 1) return n-1;
                                                          for (int j = 0; j < h; j++)
for (int j = 0; j < h; j++)
(*this)[i][j] -= m[i][j];
 if (k > n) return (joséphus(n-1,k)+k)%n;
 int res = josephus(n-n/k,k)-n\%k;
                                                          return *this;
 return res + ((res<0)?n:res/(k-1)):
} // fast case if k=2, traditional josephus
                                                         Mat<T> operator*(const Mat<T>& m) {
                                                          Mat<T> z(w,m.h);
int josephus(int n) {
                                                          for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
 return 2*(n-(1<<(32-__builtin_clz(n)-1)));
                                                            for (int k = 0; k < m.h; k++)
z[i][k] += (*this)[i][j] * m[j][k];
Least Common Multiple
                                                             return z:
#define lcm(a,b) ((a*b)/qcd(a,b))
                                                         Mat<T> operator+(const Mat<T>& m) { Mat<T>
Modulo Operations
                                                            a=*this: return a+=m: }
                                                         Mat<T> operator-(const Mat<T>& m) { Mat<T>
#define MOD 1000000007
#define madd(a,b,m) (a+b-((a+b-m>=0)?m:0))
#define mult(a,b,m) ((ull)a*b/m)
#define msub(a,b,m) (a-b+((a<b)?m:0))
                                                            a=*this; return a-=m; }
                                                         Mat<T>& operator*=(const Mat<T>& m) { return
                                                           *this = (*this)*m; }
                                                         Mat<T> power(int n) {
11 mpow(ll b, ll e, ll m) {
                                                          Mat<T> a = Mat<T>::identity(w), m=*this;
for (;n;n/=2,m*=m) if (n&1) a *= m;
 .11 x = 1;
 Matrix Exponentiation
 return x % m;
                                                         /\!/ F(n) = c[\hat{0}]*F(n-1) + c[1]*F(n-2) + ...

/\!/ b is the base cases of same length c
ull mfac(ull n, ull m) {
                                                        |ll matrix_exponentiation(ll n, vector<ll> c,
 for (int i = n; i > 1; i--)

  vector<11> b) {
  if (nth < b.size()) return b[nth-1];
}
</pre>
 f = (f * i) \% m;
 return f;
                                                         Mat<11> a(c.size(), c.size()); ll s = 0;
for (int i = 0; i < c.size(); i++) a[i][0] =
} // if m is not guaranteed to be prime
```

```
for (int i = 0; i < c.size() - 1; i++)
 \rightarrow a[i][i+1] = 1:
 a = a.power(nth - c.size());
 for (int i = 0; i < c.size(); i++)

s += a[i][0] * b[i];
 return s;
Nimber Arithmetic
#define nimAdd(a,b) ((a) \hat{}(b))
ull nimMul(ull a, ull b, int i=6) {
   static const ull M[]={INT_MIN>>32,
     M[0]^{(M[0] << 16)}, M[1]^{(M[1] << 8)},
     M[2]^(M[2] << 4), M[3]^(M[3] << 2),
   M[4]^{(M[4]<<1)};
  if (i-- == 0) return a&b;
  int k=1<<i:
  ull s=nimMul(a,b,i), m=M[5-i],
     t=nimMul(((a^(a>>k))&m)|(s\&~m),
  ((b^(b>>k))&m)|(m&(~m>>1))<<k,i);
return((s^t)&m)<<k|((s^(t>>k))&m);
Permutation
// c = array size, n = nth perm, return index
vector<int> gen_permutation(int c, int n) {
 vector<int> idx(c), per(c), fac(c); int i;
 for (i = 0; i < c; i++) idx[i] = i;
for (i = 1; i <= c; i++) fac[i-1] = n%i, n/=i;
for (i = c - 1; i >= 0; i--)
per[c-i-1] = idx[fac[i]],
  idx.erase(idx.begin() + fac[i]);
 return per;
// get what nth permutation of vector
int get permutation(vector<int>& v) {
 int use = 0, i = 1, r = 0;
 for (int e: v) {
   r = r * i++ + __builtin_popcount(use &
 \rightarrow -(1<<e));
  use |= 1 << e:
 return r:
Permutation (string/multiset)
string freq2str(vector<int>& v) {
 for (int i = 0; i < v.size(); i++)
  for (int j = 0; j < v[i]; j++)
s += (char)(i + 'A');
 return s;
// nth perm of multiset, n is O-indexed
string gen_permutation(string s, ll n) {
 vector<int> freq(26, 0);
 for (auto e : s) freq[e - 'A']++;
 for (int i = 0; i < 26; i++) if (freq[i] > 0)
  freq[i]--; ll v = multinomial(freq);
if (n < v) return (char)(i+'A') +</pre>
  → gen_permutation(freq2str(freq), n);
  freq[i]++; n-= v;
 return "":
Miller-Rabin Primality Test
// Miller-Rabin primality test - O(10 log^3 n)
bool isPrime(ull n) {
  if (n < 2) return false;
  if (n == 2) return true;
  if (n % 2 == 0) return false;</pre>
 ull s = n - 1;
while (s % 2 == 0) s /= 2;
 for (int i = 0; i < 10; i++) {
  ull temp = s;
  ull a = rand() % (n - 1) + 1;
  ull mod = mpow(a, temp, n);
  while (temp!=n-1\&\&mod!=1\&\&mod!=n-1) {
   mod = mult(mod, mod, n);
```

```
.temp *= 2;
  if (mod!=n-1&&temp%2==0) return false:
.}
.return true;
Sieve of Eratosthenes
bitset<100000001> sieve;
// generate sieve - O(n log n)
void genSieve(int n) {
  sieve[0] = sieve[1] = 1;
| Sieve[j] - 1, | for (ull i = 3; i * i < n; i += 2) | if (!sieve[i]) | | for (ull j = i * 3; j <= n; j += i * 2) |
    sieve[j] = 1;
// query sieve after it's generated - O(1)
bool quervSieve(int n) {
return n == 2 || (n % 2 != 0 && !sieve[n]);
Simpson's / Approximate Integrals
   integrate f from a to b, k iterations
// error <= (b-a)/18.0 * M * ((b-a)/2k)^4
// where M = max(abs(f^{*})^*(x))) for x in [a,b] // "f" is a function "double func(double x)"
double Simpsons (double a, double b, int k,
 double (*f)(double)) {
  double dx = (b-a)/(2.0*k), t = 0;
for (int i = 0; i < k; i++)

t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
\hookrightarrow (*f)(a+(2*i+1)*dx);
return (t + (*f)(b)) * (b-a) / 6.0 / k;
Common Equations Solvers
// ax^2 + bx + c = 0, find x
vector<double> solveEq(double a, double b,

    double c) {
    vector<double> r;
}
 double z = b * b - 4 * a * c;
if (z == 0)
 r.push_back(-b/(2*a));
 else if (z > 0) {
 r.push back((sqrt(z)-b)/(2*a));
 r.push_back((sqrt(z)+b)/(2*a));
.}
.return r;
\frac{1}{2} / ax^3 + bx^2 + cx + d = 0, find x
vector<double> solveEq(double a, double b,
double c, double d) {
.vector<double> res;
.long double a1 = b/a, a2 = c/a, a3 = d/a;
 long double q = (a1*a1 - 3*a2)/9.0, sq =
 \rightarrow -2*sqrt(q);
 long double r = (2*a1*a1*a1 - 9*a1*a2 +
\stackrel{\frown}{\rightarrow} 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
 if (z \le 0) {
  theta = acos(r/sqrt(q*q*q));
  res.push_back(sq*cos(theta/3.0) - a1/3.0);
  res.push_back(sq*cos((theta+2.0*PI)/3.0) -
 \rightarrow a1/3.0):
 res.push back(sq*cos((theta+4.0*PI)/3.0) -
 \rightarrow a1/3.0):
 res.push_back(pow(sqrt(z)+fabs(r), 1/3.0));
 res[0] = (res[0] + q / res[0]) *
\rightarrow ((r<0)?1:-1) - a1 / 3.0;
return res;
\frac{1}{1} linear diophantine equation ax + by = c.
   find x and y
// infinite solutions of form x+k*b/g, y-k*a/g
bool solveEq(11 a, 11 b, 11 c, 11 &x, 11 &y, 11
```

```
g = egcd(abs(a), abs(b), x, y);
 if (c % g) return false;
x *= c / g * ((a < 0) ? -1 : 1);

y *= c / g * ((b < 0) ? -1 : 1);
 return true;
}
// m = # equations, n = # variables, a[m][n+1]
\rightarrow = coefficient matrix
// a[i][0]x + a[i][1]y + ... + a[i][n]z =
    a[i][n+1]
\stackrel{\longleftrightarrow}{/\!\!/} find \ a \ solution \ of \ some \ kind \ to \ linear
\hookrightarrow equation
const double eps = 1e-7;
bool zero(double a) { return (a < eps) && (a >
→ -eps); }
vector < double > solve Eq (double **a, int m, int
 int cur = 0;
for (int i = 0; i < n; i++) {
    for (int j = cur; j < m; j++) {
        if (!zero(a[j][i])) {
  ...if (j != cur) swap(a[j], a[cur]);
     for (int sat = 0; sat < m; sat++) {
  ...if (sat == cur) continue;
  double num = a[sat][i] / a[cur][i];
for (int sot = 0; sot <= n; sot++)
a[sat][sot] -= a[cur][sot] * num;
     }
cur++
    break
 for (int j = cur; j < m; j++)
  if (!zero(a[j][n])) return vector<double>();
 vector<double ans(n,0);
for (int i = 0, sat = 0; i < n; i++)
if (sat < m && !zero(a[sat][i]))
ans[i] = a[sat][n] / a[sat++][i];
return ans;
// solve A[n][n] * x[n] = b[n] linear equation
// rank < n is multiple solutions, -1 is no
\stackrel{\textstyle \longrightarrow}{\nearrow} solutions alls is whether to find all solutions, or
\hookrightarrow any
const double eps = 1e-12:
int solveEq(Vec<2, double>& A, Vec<1, double>&

→ b, Vec<1, double>& x, bool alls=false) {
int n = A.size(), m = x.size(), rank = 0, br,
vector<int> col(m); iota(begin(col), end(col),
c   0);
for(int i = 0; i < n; i++) {
  double v, bv = 0;</pre>
  for(int \dot{r} = i; \dot{r} < n; \dot{r} + +)
  for(int c = i; c < n; c++)
    if ((v = fabs(A[r][c])) > bv)
    br = r, bc = c, bv = v;
    if (bv <= eps) {
   for(int j = i; j < n; j++)
if (fabs(b[j]) > eps)
      .return -1;
  swap(A[i], A[br]);
swap(b[i], b[br]);
  swap(col[i], col[bc]);
  for(int j = 0; j < n; j++)
    swap(A[j][i], A[j][bc]);
bv = 1.0 / A[i][i];
for(int j = (alls)?0:i+1; j < n; j++) {</pre>
  if (j != i) {
   ..double fac = A[j][i] * bv;
     .b[i] -= fac * b[i];
   for(int k = i+1; k < m; k++)
A[j][k] -= fac*A[i][k];
rank++;
```

```
-DBL_MAX;
 for (int i = rank; i--:)
  bool isGood = true:
  if (alls)
  for (int j = rank; isGood && j < m; j++)
   if (fabs(A[i][j]) > eps)
  isĜood = false;
b[i] /= A[i][i];
  if (isGood) x[col[i]] = b[i];
if (!alls)
  for(int j = 0; j < i; j++)
b[j] -= A[j][i] * b[i];
return rank:
Gravcode Conversions
ull gravcode2ull(ull n) {
 ull i = 0;
 for (; n; n = n >> 1) i ^= n; return i;
ull ull2graycode(ull n) { return n ^ (n >> 1);
Unix/Epoch Time
// O-indexed month/time, 1-indexed day
/// minimum 1970, 0, 1, 0, 0, 0
ull toEpoch(int year, int month, int day, int
 → hour, int minute, int second) {
 struct tm t; time_t epoch;
t.tm_year = year - 1900; t.tm_mon = month;
t.tm_mday = day; t.tm_hour = hour;
t.tm_min = minute; t.tm_sec = second;
t.tm_isdst = 0; // 1 = daylights savings
 epoch = mktime(&t);
 return (ull)epoch;
vector<int> toDate(ull epoch) {
 time t e=epoch: struct tm t=*localtime(&e):
 return {t.tm_year+1900,t.tm_mon,t.tm_mday,t_
    .tm hour, t.tm min, t.tm sec};
```

Theorems and Formulae

int getWeekdav(ull epoch) {

int getDayofYear(ull epoch) {

return t.tm_yday; // 0-365

const int months[] =

month == 1):

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set \rightarrow of size n. !0 = 1, !1 = 0, !n = (n+1)(!(n-1))|1)+!(n-2), $!n = n! \sum_{i=0}^{n} \frac{(-1)^i}{i!}$, $!n = [\frac{n!}{e}]$

time t e=epoch; struct tm t=*localtime(&e);

time_t e=epoch; struct tm t=*localtime(&e);

→ {31,28,31,30,31,30,31,30,31,30,31}; bool validDate(int year, int month, int day) { bool leap = !(year%(year%25?4:16));

if (month >= 12) return false; return day <= months[month] + (leap &&

return t.tm wday; // 0-6, 0 = sunday

In a partially ordered set, a chain is a subset of elements that are all comparable to eachother An antichain is a subset where no two are comparable.

Dilworth's theorem states the size of a maximal antichain equals the size of a minimal

if (alls) for (int i = 0; i < m; i++) x[i] = |chain cover of a partially ordered set S. The graph[s].erase(graph[s].begin()); width of S is the maximum size of an antichain in S, which is equal to the minimum number of chains needed to cover S, or the minimum number of chains such that all elements are in at least one chain.

Rosser's Theorem states the nth prime number is greater than n * ln(n) for n > 1.

Nicomachi's Theorem states $1^3 + 2^3 + ... + const 11$ inf = 1LL << 62; $n^3 = (1 + 2 + ... + n)^2$ and is equivalent to void floydWarshall (Vec<2, 11>& m) {

every natural number is the sum of the squares of four non-negative integers. This is a special case of the Fermat Polygonal Number **Theorem** where every positive integer is a sum of at most n s-gonal numbers. The $nth \mapsto m[i][j] = -inf;$ s-gonal number $P(s,n) = (s-2)\frac{n(n-1)}{2} + n$

```
7 Graphs
struct edge {
```

int u,v,w; edge (int u,int v,int w) : u(u),v(v),w(w) {} edge (): u(0), v(0), w(0) {} |}; |bool operator < (const edge &e1. const edge bool operator > (const edge &e1, const edge

```
Eulerian Path
#define edge_list vector<edge>
#define adj sets vector<set<int>>
struct EulerPathGraph {
 adj_sets graph; // actually indexes incident
 → edaes
 edge_list edges; int n; vector<int> indeg;
 EulerPathGraph(int n): n(n) {
 indeg = *(new vector<int>(n,0));
 graph = *(new adj_sets(n, set<int>()));
 void add_edge(int u, int v) {
  graph[u].insert(edges.size());
  indeg[v]++;
  edges.push_back(edge(u,v,0));
 bool eulerian_path(vector<int> &circuit) {
  if(edges.size()==0) return false;
 stack<int> st;
int a[] = {-1, -1};
for(int v=0;v<n;v++)
   if(indeg[v]!=graph[v].size()) {
    bool b = indeg[v] > graph[v].size();
if (abs(((int))indeg[v])-((int)graph[v])
     .size())) > 1) return
    false;
if (a[b] != -1) return false;
    a[b] = v;
 int s = (a[0]!=-1 && a[1]!=-1 ? a[0] :
→ (a[0]==-1 && a[1]==-1 ? edges[0].u : -1));
```

while(!st.empty() || !graph[s].empty()) {

circuit.push back(s); s = st.top();

int w = edges[*graph[s].begin()].v;

if(s==-1) return false:

⇒ st.pop(); }

else {

if (graph[s].empty()) {

```
st.push(s); s = w;
circuit.push back(s):
return circuit.size()-1==edges.size();
```

Flovd Warshall

```
\frac{(n\frac{n+1}{2})^2}{\text{Lagrange's Four Square Theorem states}} \begin{vmatrix} \inf n = \min(\text{size}(); \\ \text{FOR}(i,n) & \text{m[i][i]} = \min(\text{m[i][i]}, \text{OLL}); \\ \text{FOR}(k,n) & \text{FOR}(i,n) & \text{FOR}(j,n) & \text{if } (\text{m[i][k]} & \text{inf } (\text{m[i][k]}) \end{vmatrix} = \inf(n\frac{n+1}{2})^2 \cdot (n\frac{n+1}{2})^2 \cdot (
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                       \rightarrow && m[k][j] != inf)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                              auto newDist = max(m[i][k] + m[k][j], -inf);
m[i][j] = min(m[i][j], newDist);
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         FOR(k,n) if (m[k][k] < 0) FOR(i,n) FOR(j,n)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                       if (m[i][k] != inf && m[k][j] != inf)
```

Minimum Spanning Tree

```
// returns vector of edges in the mst
// graph[i] = vector of edges incident to
   vertex i places total weight of the mst in Stotal
// if returned vector has size != n-1, there is
vector<edge> mst(vector<vector<edge>> graph,
priority_queue<edge, vector<edge>.

→ greater<edge>> pq;

vector<edge> MST;
bitset<20001> marked: // change size as needed
 marked[0] = 1;
for (edge ep : graph[0]) pq.push(ep);
while(MST.size()!=graph.size()-1 &&
 → pq.size()!=0) {
 pq.size():-0/
edge e = pq.top(); pq.pop();
int u = e.u, v = e.v, w = e.w;
if(marked[u] && marked[v]) continue;
else if(marked[u]) swap(u, v);
  for(edge ep : graph[u]) pq.push(ep);
  marked[u] = 1
  MST.push back(e):
  total += e.w:
 return MST;
```

Union Find

```
int uf_find(subset* s, int i) {
  if (s[i].p != i) s[i].p = uf_find(s, s[i].p);
  return s[i].p;
void uf_union(subset* s, int x, int y) {
int xp = uf_find(s, x), yp = uf_find(s, y);
 if (s[xp].rank > s[yp].rank) s[yp].p = xp;
else if (s[xp].rank < s[yp].rank) s[xp].p =
else { s[yp].p = xp; s[xp].rank++; }
```

2D Grid Shortcut

```
#define inbound(x,n) (0<=x\mathcal{E}\mathcal{E}x<n)
#define fordir(x, y, n, m) for(auto[dx, dy]:dir)if
\hookrightarrow (inbound(x+dx,n)&\mathcal{G}inbound(y+dy,m))
const pair<int,int> dir[] =
\rightarrow {{1,0},{0,1},{-1,0},{0,-1}};
```

```
double width(rectangle a) { return
    2D Geometry
                                                     → abs(real(a.br) - real(a.tl)); }
#define point complex<double>
                                                    double height(rectangle a) { return
#define EPS 0.0000001

    abs(imag(a.br) - real(a.tl)); }

#define sq(a) ((a)*(a))
                                                    double diagonal(rectangle a) { return
#define cb(a) ((a)*(a)*(a))

    sgrt(sg(width(a)) + sg(height(a))); }

double dot(point a, point b) { return
                                                    double area (rectangle a) { return width(a) *

→ real(conj(a)*b); }

                                                    → height(a); }
double cross(point a, point b) { return
                                                    double perimeter(rectangle a) { return 2 *

    imag(conj(a)*b); }

                                                        (width(a) + height(a)); }
struct line { point a, b; };
                                                    // check if `a` fit's inside `b
struct circle { point c; double r; };
                                                    // swap equalities to exclude tight fits
struct segment { point a, point b; };
                                                    bool doesfitInside(rectangle a, rectangle b) {
struct triangle { point a, b, c; };
struct rectangle { point tl, br; };
                                                     int x = width(a), w = width(b), y = height(a),
                                                     \rightarrow h = height(b):
struct convex_polygon {
  vector<point> points;
                                                     if (x > y) swap(x, y);
if (w > h) swap(w, h);
 convex_polygon(vector<point> points) :
                                                        (w < x) return false;</pre>

→ points(points) {}
                                                     if (y <= h) return true;
 convex_polygon(triangle a) {
                                                     double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
 points.push_back(a.a); points.push_back(a.b);
                                                     return sq(a) \le sq(b) + sq(c):

→ points.push_back(a.c);

                                                       polygon methods
 convex_polygon(rectangle a) {
                                                       negative area = CCW, positive = CW
 points.push_back(a.tl);
                                                    double area(polygon a) {
    points.push back({real(a.tl).
                                                      double area = 0.0; int n = a.points.size();
   imag(a.br):
                                                      for (int i = 0, j = 1; i < n; i++, j = (j +
  points.push_back(a.br);
                                                        1) % n)
    points.push_back({real(a.br),
                                                        area += (real(a.points[j]-a.points[i]))*|
    imag(a.tl)}):
                                                        (imag(a.points[j]+a.points[i]));
};
                                                      return area / 2.0:
struct polygon {
                                                       get both unsigned area and centroid
vector<point> points;
                                                    pair<double, point> area_centroid(polygon a) {
polygon(vector point points) :
                                                     int n = a.points.size();
→ points(points) {}
                                                     double area = 0:
polygon(triangle a) {
                                                     point c(0, 0);
                                                     for (int i = n - 1, j = 0; j < n; i = j++) {
   double v = cross(a.points[i], a.points[j]) /
 points.push_back(a.a); points.push_back(a.b);
   points.push back(a.c);
 polygon(rectangle a) {
                                                      c += (a.points[i] + a.points[j]) * (v / 3);
 points.push_back(a.tl);
    points.push back({real(a.tl),
    imag(a.br)});
                                                     return {area, c};
  points.push_back(a.br);
                                                    Intersection
    points.push back({real(a.br),
                                                    // -1 coincide, 0 parallel, 1 intersection int intersection(line a, line b, point& p) {
    imag(a.tl)});
                                                     if (abs(cross(a.b - a.a, b.b - b.a)) > EPS) {
 polygon(convex_polygon a) {
                                                     p = cross(b.a - a.a, b.b - a.b) / cross(a.b)
  for (point v : a.points)
                                                     \rightarrow -a.a, b.b - b.a) * (b - a) + a; return 1;
   points.push_back(v);
                                                     if (abs(cross(a.b - a.a. a.b - b.a)) > EPS)
   triangle methods
                                                     → return 0;
double area_heron(double a, double b, double
\rightarrow c) { if (a < b) swap(a, b):
                                                    // area of intersection
if (a < c) swap(a, c);
                                                    double intersection(circle a, circle b) {
 if (b < c) swap(b, c);
                                                    double d = abs(a.c - b.c);
if (d <= b.r - a.r) return area(a);
if (d <= a.r - b.r) return area(b);
 if (a > b + c) return -1:
return sqrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)
                                                     if (d \ge a.r + b.r) return 0;
   /16.0);
                                                     double alpha = acos((sq(a.r) + sq(d) -
                                                     \rightarrow sq(b.r)) / (2 * a.r * d));
// segment methods
                                                     double beta = acos((sq(b.r) + sq(d) - sq(a.r))) \rightarrow \{x+a.x, y+a.y, z+a.z\}; \}
double lengthsq(segment a) { return
                                                        /(2 * b.r * d));
    sq(real(a.a) - real(a.b)) + sq(imag(a.a) -
                                                     return sq(a.r) * (alpha - 0.5 * sin(2 *
   imag(a.b)); }
                                                        alpha) + sq(b.r) * (beta - 0.5 * sin(2 *
double length(segment a) { return
                                                        beta));

    sqrt(lengthsq(a)); }

// circle methods
                                                       -1 outside, 0 inside, 1 tangent, 2
double circumference(circle a) { return 2 * a.r
                                                    int intersection circle a, circle b,

→ * M_PI; }

double area(circle a) { return sq(a.r) * M_PI;
                                                        vector<point>& inter) {
\rightarrow } // rectangle methods
                                                     double d2 = norm(b.c - a.c), rS = a.r + b.r,
                                                    \rightarrow rD = a.r - b.r:
```

```
if (d2 > sq(rS)) return -1;
 if (d2 < sq(rD)) return 0;
 double ca = 0.5 * (1 + rS * rD / d2);
point z = point(ca, sqrt(sq(a.r) / d2 -
 \rightarrow sq(ca))):
 inter.push_back(a.c + (b.c - a.c) * z);
 if (abs(imag(z)) > EPS) inter.push_back(a.c
   (b.c - a.c) * conj(z));
 return inter.size():
// points of intersection
vector<point> intersection(line a, circle c) {
 vector<point> inter;
c.c -= a.a;
a.b -= a.a;
 point m = a.b * real(c.c / a.b);
 double d2 = norm(m - c.c);
 if (d2 > sq(c.r)) return 0;
 double 1 = \operatorname{sqrt}((\operatorname{sq}(c.r) - d2) / \operatorname{norm}(a.b));
 inter.push_back(a.a + m + 1 * a.b);
 if (abs(1) > EPS) inter.push back(a.a + m - 1
 \rightarrow * a.b):
 return inter:
// area of intersection
double x2 = min(real(a.br), real(b.br)), y2 =
 → min(imag(a.br), imag(b.br));
return (x2 <= x1 || y2 <= y1) ? 0 :
 \rightarrow (x2-x1)*(y2-y1);
Convex Hull
bool cmp(point a, point b) {
  if (abs(real(a) - real(b)) > EPS) return
    real(a) < real(b);
 if (abs(imag(a) - imag(b)) > EPS) return

    imag(a) < imag(b);
</pre>
 return false;
convex_polygon convexhull(polygon a) {
 sort(a.points.begin(), a.points.end(), cmp);
 vector<point> lower, upper;
 for (int i = 0; i < a.points.size(); i++) {
  while (lower.size() >= 2 &&
    cross(lower.back() - lower[lower.size()
    2], a.points[i] - lower.back()) < EPS)
   lower.pop back();
  while (upper.size() >= 2 &&
    cross(upper.back() - upper[upper.size() -
    2], a.points[i] - upper.back()) > -EPS)
   upper.pop_back();
  lower.push_back(a.points[i]);
  upper.push back(a.points[i]);
 lower.insert(lower.end(), upper.rbegin() + 1,
    upper.rend());
 return convex polygon(lower);
     3D Geometry
struct point3d {
 double x, y, z;
 point3d operator+(point3d a) const { return
 point3d operator*(double a) const { return
 \rightarrow {x*a, v*a, z*a}: }
 point3d operator-() const { return {-x, -y,
 point3d operator-(point3d a) const { return
 \rightarrow *this + -a; }
 point3d operator/(double a) const { return
 \rightarrow *this * (1/a); }
 double norm() { return x*x + y*y + z*z; }
 double abs() { return sqrt(norm()); }
```

```
.point3d normalize() { return *this /
this->abs(); }
double dot(point3d a, point3d b) { return
 \rightarrow a.x*b.x + a.y*b.y + a.z*b.z; }
point3d cross(point3d a, point3d b) { return
    \{a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,
\stackrel{\hookrightarrow}{\Rightarrow} a.x*b.y - a.y*b.x}; }
struct line3d { point3d a, b; };
struct plane { double a, b, c, d; } // a*x +
 \Rightarrow b*y + c*z + d = 0
struct sphere { point3d c; double r; };
#define sq(a) ((a)*(a))
#define cb(a) ((a)*(a)*(a))
double surface(circle a) { return 4 * sq(a.r) *
double volume(circle a) { return 4.0/3.0 *
 \hookrightarrow cb(a.r) * M PI; }
```

10 Optimization

Snoob

```
// SameNumberOfOneBits, next permutation
// example usage
int main() {
    char l1[] = {'1', '2', '3', '4', '
    char l2[] = {'a', 'b', 'c', 'd'};
    int d1 = 5, d2 = 4;
    // prints 12345abcd, 1234a5bcd, ...
  int min = (1 < < d1) - 1, max = min << d2;
 for (int i = min; i <= max; i = snoob(i)) {
    int p1 = 0, p2 = 0, v = i;
    while (p1 < d1 || p2 < d2) {
        cout << ((v & 1) ? 11[p1++] : 12[p2++]);
      v /= 2;
   .cout << '\n';
```

```
bool isPowerOf2(11 a) {
  return a > 0 && !(a & a-1);
bool isPowerOf3(11 a) {
return a>0&&!(12157665459056928801ull%a);
bool isPower(ll a, ll b) {
  double x = log(a) / log(b);
  return abs(x-round(x)) < 0.00000000001;</pre>
```

11 Additional

Judge Speed

```
// kattis: 0.50s
// codeforces: 0.421s
// atcoder: 0.455s
#include <bits/stdc++.h>
using namespace std;
int v = 1e9/2, p = 1;
int main() {
    for (int i = 1; i <= v; i++) p *= i;
```

Judge Pre-Contest Checks

int128 and float128 support? -does extra or missing whitespace cause WA? -documentation up to date? -printer usage available and functional?

```
// each case tests a different fail condition
 // try them before contests to see error codes
struct g { int arr[1000000]; g(){}};
vector<g> a;
 // O=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE
if (n == 4) assert(0);
if (n == 5) 0 / 0;
if (n == 6) *(int*)(0) = 0;
 return n + judge(n + 1);
 GCC Builtin Docs
 // 128-bit integer
__int128 a;
unsigned __int128 b;
// 128-bit float
// minor improvements over long double __float128 c; // log2 floor
__lg(n);
__lg(n);
// number of 1 bits
// can add ll like popcountll for long longs
__builtin_popcount(n);
// number of trailing zeroes
_builtin_ctz(n);
// number of leading zeroes
_builtin_ctz(n);
// number of leading zeroes
_builtin_clz(n);
// 1-indexed_least_significant 1 bit
_builtin_ffs(n).
__builtin_ffs(n);
// parity of number
 __builtin_parity(n);
Limits
                          \begin{array}{c|c} \pm 2147483647 & \pm 2^{31} - 1 | 10^9 \\ 4294967295 | & 2^{32}_{32} - 1 | 10^9_{10} \end{array}
 int
 uint
          \pm 9223372036854775807 | \pm 2^{63} - 1|10^{18}
                                                   \frac{1}{2}^{64} - \frac{1}{10}^{19}
           18446744073709551615
 ull
i128 |\pm 170141183460469231...|\pm 2^{127} - 1|10^{38}
                                                 \frac{1}{2}^{128} - \frac{1}{1}^{10}^{38}
 u128 340282366920938463...
 Complexity classes input size (per second):
 O(n^n) or O(n!)
                                                             n < 10
 O(2^n)
                                                             n \leq 30
 O(n^3)
                                                         n \le 1000
 O(n^2)
                                                       n \le 30000
                                                           n \le 10^6
 O(n\sqrt{n})
 O(n \log n)
                                                           n \le 10^7
 O(n)
                                                           n < 10^9
```