```
else if (d) { m=m*10+c-'0'; o*=0.1; } else n = n * 10 + c - '0':
    General
                              7 Graphs
    Algorithms
                              8 2D Geometry
                                                             n = s * (n + m * o):
    Structures
                              9 3D Geometry
    Strings
                                                            void read(double& n) {
                              10 Optimization
                                                             ld m; read(m); n = m;
    Greedy
                              11 Additional
                                                            void read(float& n) {
  ld m; read(m); n = m;
    Math
     General
                                                             void read(string& s) {
g++ -g -02 -std=gnu++17 -static prog.cpp
./a.exe
                                                             char c: s = ""
                                                             while((c=getchar_unlocked())!=' '&&c!='\n')
                                                              s += c:
                                                            bool readline(string& s) {
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp
                                                             char c: s = ""
                                                             while(c=getchar_unlocked()) {
for i in *.in; do f=${i%.in}
                                                              if (c == '\n') return true;
if (c == EOF) return false;
 ./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
                                                              s += c:
done
                                                             return false:
Header
                                                            void print(unsigned int n) {
// use better compiler options
                                                             if (n / 10) print(n / 10);
#pragma GCC optimize("Ofast","unroll-loops")
#pragma GCC target("avx2,fma")
                                                             putchar unlocked(n % 10 + '0'):
// include everything
                                                             void print(int n) {
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <bits/extc++.h>
#include <sys/resource.h>
                                                             if (n < 0) { putchar_unlocked('-'); n*=-1; }</pre>
                                                             print((unsigned int)n);
// namespaces
using namespace std;
                                                            Common Structs
using namespace __gnu_cxx; // rope
                                                            // n-dimension vectors
// Vec<2, int> v(n, m) = arr[n][m]
// Vec<2, int> v(n, m, -1) default init -1
template<int D, typename T>
using namespace __gnu_pbds; // tree/trie
// common defines
#define fastio
                                                            struct Vec : public vector<Vec<D-1, T>> {
\label{eq:control_state} \rightarrow ios\_base::sync\_with\_stdio(0);cin.tie(0); \\ \textit{\#define nostacklim rlimit}
                                                               template<typename... Args>
                                                               Vec(int n=0, Args... args) : vector<Vec<D-1,
     RZ; getrlimit(3, &RZ); RZ.rlim_cur=-
                                                                T >> (n, Vec < D-1, T > (args...)) {}
\(\Rightarrow\) 1; setrlimit(3,\&RZ); \(\text{#define DEBUG(v) cerr<<_LINE__<<": "<<\#v<<" =
                                                             témplate<typename T>
struct Vec<1, T> : public vector<T> {
                                                               Vec(int n=0, T val=T()) : vector<T>(n, val)
→ cerr<<1.0*clock()/CLOCKS_PER_SEC<<"s\n";
#define ll long long
#define ull unsigned ll
                                                                {}
#define i128 __int128
#define u128 unsigned i128
                                                                 Algorithms
#define ld long double
                                                            Min/Max Subarray
// global variables
                                                             // max - compare = a < b, reset = a < 0
mt19937 rng((uint32_t)chrono::steady
                                                                min - compare = a > b, reset = a > 0

    clock::now().time since epoch().count()):

                                                            // returns {sum, {start, end}}
pair<int, pair<int, int>>
                                                                 ContiguousSubarrav(int* a. int size.
                                                                 bool(*compare)(int, int),
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
                                                             bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
                                                             void read(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                               cur += a[i];
                                                               if ((*compare)(best, cur)) { best = cur;
 n = n * 10 + c - 0';
                                                                start = s; end = i; }
void read(int& n) {
   char c; n = 0; int s = 1;
   if ((c=getchar_unlocked())=='-') s = -1;
                                                              if ((*reset)(cur)) { cur = 0; s = i + 1; }
                                                             return {best, {start, end}};
 else n = c - \sqrt{0}:
 while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                             Quickselect
 n = n * 10 + c -
                                                             #define QSNE -999999
                                                            int partition(int arr[], int 1, int r)
void read(ld& n) {
 char c; n = 0;

.ld m = 0, o = 1; bool d = false; int s = 1;

.if ((c=getchar_unlocked())=='-') s = -1;
                                                              int x = arr[r], i = 1;
                                                             for (int j = 1; j <= r - 1; j++)
if (arr[j] <= x)
 else if (c == '.') d = true;
else n = c - '0';
                                                               swap(arr[i++], arr[j]);
 while ((c=getchar_unlocked())!=' '&&c!='\n') {
                                                             swap(arr[i], arr[r]);
  if (c == '.') d = true;
                                                             return i;
```

```
\frac{1}{1} find k'th smallest element in unsorted array,
\hookrightarrow only if all distinct
int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - 1 + 1)) return QSNE;
swap(arr[1 + rng() % (r-1+1)], arr[r]);
 int pos = partition(arr, 1, r);
 if (pos-l==k-1) return arr[pos];
 if (pos-1>k-1) return qselect(arr,1,pos-1,k);
return qselect(arr, pos+1, r, k-pos+1-1);
// TODO: compare against std::nth element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted
→ on both axis
pair<int, int> saddleback search(int** arr, int
\stackrel{\cdot}{\hookrightarrow} x, int y, int v) {
int i = x-1, j = 0;

while (i >= 0 && j < y) {

if (arr[i][j] == v) return {i, j};
  (arr[i][i] > v)? i--: i++:
 return {-1, -1};
Ternary Search
// < max, > min, or any other unimodal func #define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a, int b, int (*f)(int)) {
 while (b-a > 4) {
  int m = (a+b)/2;
if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1:
 for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
 ...a = i;
return a:
#define TERNPREC 0.000001
double ternsearch (double a, double b, double
 \rightarrow (*f)(double)) {
 while (b-a > TERNPREC * 4) {
   double m = (a+b)/2;
  if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
 → = m;
else b = m + TERNPREC;
 for (double i = a + TERNPREC; i <= b; i +=
    TERNPREC)
      if (TERNCOMP((*f)(a), (*f)(i)))
   a = i;
 return a;
3 Structures
Fenwick Tree
// Fenwick tree, array of cumulative sums -
 \hookrightarrow O(log n) updates, O(log n) gets
struct Fenwick { int n: ll* tree:
  void update(int i, int val) {
  .++i:
  while (i <= n) {
   tree[i] += val;
   i += i & (-i):
 Fenwick(int size) {
  n = size;
  tree = new ll[n+1];
for (int i = 1; i <= n; i++)
   tree[i] = 0:
 Fenwick(int* arr, int size) : Fenwick(size) {
  for (int i = 0; i < n; i++)
update(i, arr[i]);</pre>
```

```
.ll operator[](int i) {
  if (i < 0 || i > n) return 0;
  while (i>0)
   sum += tree[i];
i -= i & (-i);
  return sum:
 ll getRange(int a, int b) { return
    operator[](b) - operator[](a-1); }
Hashtable
 // similar to unordered map, but faster
| struct chash {
| const uint64 t C = (11)(2e18 * M_PI) + 71;
| ll operator()(11 x) const { return
    __builtin_bswap64(x*C); }
int main() {
  gp_hash_table<11,int,chash>
 \rightarrow hashtable({},{},{},{},{1<<16});
 for (int i = 0; i < 100; i++)
. hashtable[i] = 200+i;
.if (hashtable.find(10) != hashtable.end())
. cout << hashtable[10];</pre>
Ordered Set
using oset = tree<T,null_type,less<T>,rb_tree
tag, tree_order_statistics_node_update>; template <typename T, typename D> using omap = tree<T,D,less<T>,rb_tree |
 - _tag,tree_order_statistics_node update>;
int main()
 coset<int> o_set;
o_set.insert(5); o_set.insert(1);
 → o_set.insert(3);
// get second smallest element
 cout << *(o_set.find_by_order(1));</pre>
 // number of elements less than k=4
 cout << ' ' << o set.order of kev(4) << '\n':
  // equivalent with ordered map
 omap<int,int> o_map;
 o_map[5]=1;o_map[1]=2;o_map[3]=3;
 cout << (*(o map.find by order(1))).first;</pre>
 cout << ' ' << o_map.order_of_key(4) << '\n';</pre>
Rope
 // ar{	extsf{O}}(\log n) insert, delete, concatenate
int main() {
 // generate rope
 rope<int> v;
 for (int i = 0; i < 100; i++)
...v.push_back(i);
  // move range to front
 rope<int> copy = v.substr(10, 10);
v.erase(10, 10);
 v.insert(copy.mutable_begin(), copy);
 // print elements of rope
 for (auto it : v) cout << it << " ";
Segment Tree
 //max(a,b), min(a,b), a+b, a*b, gcd(a,b), a\hat{b}
struct SegmentTree {
 typedef int T;
 static constexpr T UNIT = INT_MIN;
 T f(T a, T b) {
    if (a == UNIT) return b;
    if (b == UNIT) return a;
  return max(a,b);
 int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
```

```
SegmentTree(vector<T> arr) :

→ SegmentTree(arr.size()) {
 for (int i=0;i<arr.size();i++)

→ update(i,arr[i]);

 void update(int pos, T val) {
 for (s[pos += n] = val; pos /= 2;)
  s[pos] = f(s[pos * 2], s[pos*2+1]);
 T query(int b, int e) { // query [b, e) }
T ra = UNIT, rb = UNIT;
 for (b+=n, e+=n; b<e; b/=2, e/=2) {
    if (b % 2) ra = f(ra, s[b++]);
    if (e % 2) rb = f(s[--e], rb);
 return f(ra, rb);
 T get(int p) { return query(p, p+1); }
Trie
typedef trie<string, null_type,

→ trie string access traits<>,

 pat_trie_tag, trie_prefix_search_node_update>

→ trie_type;

int main() {
   // generate trie
 trie_type trie;
 for (int i = 0; i < 20; i++)
 trie.insert(to_string(i)); // true if new,
\hookrightarrow false if old
 // print things with prefix "1"
 auto range = trie.prefix_range("1");
for (auto it = range.first; it !=
\hookrightarrow range.second; it++)
 .cout << *it << "
```

```
4 Strings
Aho Corasick
// range of alphabet for automata to consider
// MAXC = 26. OFFC = 'a' if only lowercase
const int MAXC = 256;
const int OFFC = 0:
struct aho_corasick {
 struct state
  set<pair<int, int>> out:
 int fail; vector<int> go;
  state(): fail(-1), go(MAXC, -1) {}
 };
 vector<state> s;
 int id = 0;
 aho_corasick(string* arr, int size) : s(1) {
 for (int i = 0: i < size: i++) {
   int cur = 0;
  .for (int c : arr[i]) {
...if (s[cur].go[c-OFFC] == -1) {
   s[cur].go[c-OFFC] = s.size();
    s.push back(state());
    cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++});
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
   ..s[0].go[\tilde{c}] = 0;
  queue<int> sq;
 for (int c = 0; c < MAXC; c++) {
    if (s[0].go[c] != 0) {
        s[s[0].go[c]].fail = 0;
    sq.push(s[0].go[c]);
  while (sq.size()) {
 int e = sq.front(); sq.pop();
 for (int c = 0; c < MAXC; c++) {
...if (s[e].go[c] != -1) {
```

```
int failure = s[e].fail;
while (s[failure].go[c] == -1)
      failure = s[failure].fail;
failure = s[failure].go[c];
      s[s[e].go[c]].fail = failure;
  for (auto length : s[failure].out)
s[s[e].go[c]].out.insert(length);
     sq.push(s[e].go[c]);
 // list of {start pos, pattern id}
 vector<pair<int, int>> search(string text)
  vector<pair<int, int>> toret;
  int cur = 0;
  for (int i = 0; i < text.size(); i++) {
   while (s[cur].go[text[i]-OFFC] == -1)
    .cur = s[cur].fail;
   cur = s[cur].go[text[i]-OFFC];
   if (s[cur].out.size())
    for (auto end : s[cur].out)
  toret.push_back({i - end.first + 1,
     end.second});
  return toret:
Bover Moore
struct defint { int i = -1; };
vector<int> boyermoore(string txt, string pat)
 vector<int> toret; unordered_map<char, defint> Longest Common Prefix (array)
 → badchar:
 int m = pat.size(), n = txt.size();
 for (int i = 0; i < m; i++) badchar[pat[i]].i
 \rightarrow = i;
int s = 0:
 while (s \leq n - m) {
  int j = m - 1;
  while (j \ge 0) && pat[j] == txt[s + j]) j--;
  if (i < 0) {
   .toret.push back(s);
   s += (s + m < n) ? m - badchar[txt[s +
 \rightarrow mll.i : 1:
  .} else
   s += max(1, j - badchar[txt[s + j]].i);
 return toret;
English Conversion
const string ones[] = {"", "one", "two",
"three", "four", "five", "six", "seven",

"eight", "nine";

const string teens[] ={"ten", "eleven",
   "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen"};
const string tens[] = {"twenty", "thirty",
"forty", "fifty", "sixty", "seventy", 

"eighty", "ninety"};
const string mags[] = {"thousand", "million",
     "billion", "trillion", "quadrillion",
     "quintillion", "sextillion",
string convert(int num, int carry) {
 if (num < 0) return "negative " +
     convert(-num, 0);
     (num < 10) return ones[num];
(num < 20) return teens[num % 10];</pre>
     (\text{num} < 100) \text{ return tens}[(\text{num} / 10) - 2] +
     (num%10==0?"":" ") + ones[num % 10];
     (num < 1000) return ones[num / 100]
     (num/100==0?"":" ") + "hundred" + (num%100==0?"":" ") + convert(num % 100,
```

```
return convert(num / 1000, carry + 1) + " " + |...while (i >= j && i + j + 1 < n && s[(i-j)/2]
     mags[carry] + " " + convert(num % 1000.
    0):
string convert(int num) {
return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
     vector<int> toret;
 int m = txt.length(), n = pat.length();
 int next[n + 1];
for (int i = 0; i < n + 1; i++)
   next[i] = 0;</pre>
 int i = 1; i < n; i++) {
  int j = next[i + 1];
  while (j > 0 && pat[j] != pat[i])
   j = next[j];
  if (j > 0 || pat[j] == pat[i])
  next[i + 1] = i + 1;
 for (int i = 0, j = 0; i < m; i++) {
  if (txt[i] == pat[j]) {
   if (++j == n)
    toret.push_back(i - j + 1);
  .} else if (j > 0) {
  .j = next[j];
 return toret:
// longest common prefix of strings in array
string lcp(string* arr, int n, bool sorted =
false) {
if (n == 0) return "";
 if (!sorted) sort(arr, arr + n);
string r = ""; int v = 0;
 while (v < arr[0].length() && arr[0][v] ==

    arr[n-1][v])
    r += arr[0][v++];

 return r;
Longest Common Subsequence
string lcs(string a, string b) {
 int m = a.length(), n = b.length();
 int L[m+1][n+1];
 for (int i = 0; i <= m; i++) {
    for (int j = 0; j <= n; j++) {
        if (i == 0 || j == 0) L[i][j] = 0;
        else if (a[i-1] == b[j-1]) L[i][j] =
 \rightarrow L[i-1][j-1]+1;
   else L[i][j] = \max(L[i-1][j], L[i][j-1]);
 // return L[m][n]; // length of lcs
 string out = "":
 int i = m - 1, j = n - 1;
while (i >= 0 && j >= 0) {
   if (a[i] == b[j]) {
   out = a[i--] + out;
  else if (L[i][j+1] > L[i+1][j]) i--;
  else j--;
 return out;
Longest Common Substring
// l is array of palindrome length at that
→ index
int manacher(string s. int* 1) {
 int n = s.length() * 2;
 for (int i = 0, j = 0, k; i < n; i += k, j =
```

 \rightarrow max(i-k, 0)) {

```
\Rightarrow == s[(i+j+1)/2]) j++;
 1[i] = j;
  for (k = 1; i >= k && j >= k && l[i-k] !=
 \rightarrow j-k; k++)
  1[i+k] = min(1[i-k], j-k);
return *max_element(1, 1 + n);
Cyclic Rotation (Lyndon)
// simple strings = smaller than its nontrivial
   suffixes
// lyndon factorization = simple strings
→ factorized
// "abaaba" -> "ab", "aab", "a"
vector<string> duval(string s) {
int n = s.length();
vector<string> lyndon;
for (int i = 0; i < n;) {
 int j = i+1, k = i;

int j = i+1, k = i;

for (; j < n && s[k] <= s[j]; j++)

if (s[k] < s[j]) k = i;
   else k++:
  for (; i \le k; i += j - k)
  lyndon.push back(s.substr(i,j-k));
return lyndon;
// lexicographically smallest rotation
int minRotation(string s) {
int n = s.length(); s += s;
auto d = duval(s); int i = 0, a = 0;
while (a + d[i].length() < n) a +=</pre>
 \rightarrow d[i++].length();
while (i && d[i] == d[i-1]) a -=
→ d[i--].length();
return a;
Subsequence Count
// "banana", "ban" >> 3 (ban, ba..n, b..an)
ull subsequences(string body, string subs) {
int m = subs.length(), n = body.length();
if (m > n) return 0;
 ull** arr = new ull*[m+1];
for (int i = 0; i \le m; i++) arr[i] = new
\hookrightarrow ull[n+1];
for (int i = 1; i <= m; i++) arr[i][0] = 0;
for (int i = 0; i <= n; i++) arr[o][i] = 1;
for (int i = 1; i <= m; i++)
 for (int j = 1; j <= n; j++)
arr[i][j] = arr[i][j-1] + ((body[j-1] ==
\hookrightarrow subs[i-1])? arr[i-1][j-1] : 0);
return arr[m][n]:
Suffix Array + LCP
struct SuffixArray {
vector<int> sa, 1cp;
SuffixArray(string& s, int lim=256) {
   int n = s.length() + 1, k = 0, a, b;
   vector<int> x(begin(s), end(s)+1), y(n),
 \rightarrow ws(max(n, lim)), rank(n);
 sa = lcp = y;
iota(begin(sa), end(sa), 0);
  for (int j = 0, p = 0; p < n; j = max(1, j *
\rightarrow 2), lim = p) {
   p = j; iota(begin(y), end(y), n - j);
  for (int i = 0; i < (n); i++)
if (sa[i] >= j)
y[p++] = sa[i] - j;
```

fill(begin(ws), end(ws), 0);

→ ws[i - 1]:

for (int i = 0; i < (n); i++) ws[x[i]]++; for (int i = 1; i < (lim); i++) ws[i] +=

```
. for (int i = n; i--;) sa[-ws[x[y[i]]]] =
                                                      Combinatorics (nCr, nPr)

    y[i];

                                                       // can optimize by precomputing factorials, and
   swap(x, y); p = 1; x[sa[0]] = 0;
                                                           fact[n]/fact[n-r]
   for (int i = 1; i < (n); i++) {
    a = sa[i - 1]; b = sa[i];
    x[b] = (y[a] == y[b] && y[a + j] == y[b +
                                                       ull nPr(ull n, ull r) {
                                                       ull v = 1;
for (ull i = n-r+1; i <= n; i++)
return v;
  for (int i = 1; i < (n); i++) rank[sa[i]] =
                                                      ull nPr(ull n, ull r, ull m) {
                                                        ull v = 1;
                                                       for (ull i = n-r+1; i <= n; i++)
...v = (v * i) % m;
.return v;
 for (int i = 0, j; i < n - 1; lcp[rank[i++]]
for (k \&\& k--, j = sa[rank[i] - 1];
     s[i + k] = s[j + k]; k++);
                                                       ull nCr(ull n. ull r) {
                                                       long double v = 1;
for (ull i = 1: i <= r: i++)
                                                        v = v * (n-r+i) /i;
String Utilities
                                                        return (ull)(v + 0.001):
void lowercase(string& s) {
 transform(s.begin(), s.end(), s.begin(),
                                                       // requires modulo math
// caar{n} optimize by precomputing mfac and
void uppercase(string& s) {
                                                       ull nCr(ull n, ull r, ull m) {
 transform(s.begin(), s.end(), s.begin(),
                                                        return mfac(n, m) * minv(mfac(k, m), m) % m *
minv(mfac(n-k, m), m) \% m:
void trim(string &s) {
                                                       Multinomials
 s.erase(s.begin(),find_if_not(s.begin(),s
                                                      limitinomial(vector<int>& v) {
    ll c = 1, m = v.empty() ? 1 : v[0];
    for(int i = 1; i < v.size(); i++)
        for (int j = 0; j < v[i]; j++)
        c = c * ++m / (j+1);
    }
}</pre>
     .end(), [](int c){return
   isspace(c):})):
 s.erase(find_if_not(s.rbegin(),s.rend(),[](int

→ c){return isspace(c):}).base().s.end()):
                                                        return c:
vector<string> split(string& s, char token) {
    vector<string> v; stringstream ss(s);
                                                       Chinese Remainder Theorem
    for (string e;getline(ss,e,token);)
                                                       bool ecrt(ll* r, ll* m, int n, ll& re, ll& mo)
         v.push back(e);
    return v:
                                                        11 x, y, d; mo = m[0]; re = r[0];
                                                        for (int i = 1; i < n; i++) {
5 Greedy
                                                         d = egcd(mo, m[i], x, y);
                                                        if ((r[i] - re) % d != 0) return false;

x = (r[i] - re) / d * x % (m[i] / d);

re += x * mo;
Interval Cover
// L,R = interval [L,R], in = {{l,r}, index}
// does not handle case where L == R
                                                         mo = mo / d * m[i];
vector<int> intervalCover(double L, double R
                                                        re %= mo;

    vector<pair<pair<double,double>,int>> in)

    int i = 0; pair<double,int> pos = {L,-1};
                                                        re = (re + mo) \% mo;

    vector<int>
a;

                                                        return true:
    sort(begin(in), end(in));
    while (pos.first < R) {
                                                       Count Digit Occurences
         double cur = pos.first;
while (i < (int)in.size() &&</pre>
                                                       /stcount(n,d) counts the number of occurences of
                                                          a digit d in the range [0,n]*/
   in[i].first.first <= cur)</pre>
                                                       11 digit count(ll n, ll d) {
    result += ((n\%10) == d?1:0):
         if (pos.first == cur) return {};
                                                        n /= 10;
         a.push_back(pos.second);
                                                        return result;
    return a;
                                                       ĺl count(ll n, ll d) {
6 Math
                                                        if (n < 10) return (d > 0 \&\& n >= d);
                                                        if ((n % 10) != 9) return digit_count(n, d) +
Catalan Numbers
                                                        \rightarrow count(n-1, d);
ull* catalan = new ull[1000000];
                                                        return 10*count(n/10, d) + (n/10) + (d > 0):
void genCatalan(int n, int mod) {
catalan[0] = catalan[1] = 1;
for (int i = 2; i <= n; i++) {
   catalan[i] = 0;
                                                       Discrete Logarithm
                                                       unordered_map<int, int> dlogc;
  for (int j = i - 1; j \ge 0; j - ) {
                                                       int discretelog(int a, int b, int m) {
   catalan[i] += (catalan[j] * catalan[i-j-1])
                                                        dlogc.clear();
                                                        11 \text{ n} = \text{sqrt}(\text{m}) + 1, \text{ an } = 1;
 if (catalan[i] >= mod)
catalan[i] -= mod;
                                                       for (int i = 0; i < n; i++)
an = (an * a) % m;
                                                        11 c = an:
\gamma'' // TODO: consider binomial coefficient method
                                                       for (int i = 1; i <= n; i++) {
  if (!dlogc.count(c)) dlogc[c] = i;</pre>
```

```
for (int i = 0; i <= n; i++) {
   if (dlogc.count(c)) return (dlogc[c] * n - i</pre>
 \rightarrow + m - 1) % (m-1);
 c = (c * a) \% m:
 return -1;
Euler Phi / Totient
int phi(int n) {
int r = n;

for (int i = 2; i * i <= n; i++) {

   if (n % i == 0) r -= r / i;

   while (n % i == 0) n /= i;
 if (n > 1) r = r / n;
 return r:
#define n 100000
ll phi[n+1];
void computeTotient() {
 for (int i=1; i<=n; i++) phi[i] = i;
 for (int p=2; p<=n; p++) {
  if (phi[p] == p) {
   phi[p] = p-1;
   for (int i = 2*p; i<=n; i += p) phi[i] =
    (phi[i]/p) * (p-1);
Factorials
// digits in factorial
#define kamenetsky(n) (floor((n * log10(n /
 \rightarrow ME)) + (log10(2 * MPI * n) / 2.0)) + 1)
// approximation of factorial
#define stirling(n) ((n == 1) ? 1 : sart(2 *
\hookrightarrow M PI * n) * pow(n / M E, n))
// natural log of factorial
#define lfactorial(n) (lgamma(n+1))
Prime Factorization
// do not call directly
ll pollard_rho(ll n, ll s) {
 .11 x, y;
 x = y = rand() \% (n - 1) + 1;
 int head = 1, tail = 2;
while (true) {
 x = mult(x, x, n);
x = (x + s) % n;
if (x == y) return n;
  11 d = gcd(max(x - y, y - x), n);
if (1 < d && d < n) return d;
  if (++head == tail) y = x, tail <<= 1;
// call for prime factors
void factorize(ll n, vector<ll> &divisor) {
 if (n == 1) return;
 if (isPrime(n)) divisor.push_back(n);
  while (d'>= n) d = pollard_rho(n, rand() % (n)
 \rightarrow -1) + 1);
  factorize(n / d, divisor);
  factorize(d, divisor);
Farev Fractions
    generate 0 \le a/b \le 1 ordered, b \le n
    farey(4) = 0/1 1/4 1/3 1/2 2/3 3/4 1/1
// Jungth is sum of phi(i) for i = 1 to n
vector<pair<int, int>> farey(int n) {
```

int h = 0, k = 1, x = 1, y = 0, r;

vector<pair<int, int>> v;

do {

c = (c * an) % m;

```
r = (n-y)/k;
 y += r*k; x' += r*h;
swap(x,h); swap(y,k);
x = -x; y = -y;
} while (k > 1);
v.push_back({1, 1});
return v:
Fast Fourier Transform
#define cd complex<double>
const double PI = acos(-1);
void fft(vector<cd>& a. bool invert) {
 int n = a.size();
for (int i = 1, j = 0; i < n; i++) {
  int bit = n >> 1;
 for (; j & bit; bit >>= 1) j ^= bit;
 .j ^= biť;
 if (i < j) swap(a[i], a[j]);
 for (int len = 2; len <= n; len <<= 1) {
 double ang = 2 * PI / len * (invert ? -1 :
 cd wlen(cos(ang), sin(ang));
  for (int i = 0; i < n; i += len) {
   cd w(1):
   for (int j = 0; j < len / 2; j++) {
   cd u = a[i+j], v = a[i+j+len/2] * w;
   a[i+j] = u + v;
  a[i+j+len/2] = u - v;
w *= wlen;
 if (invert)
 for (auto\& x : a)
vector<int> fftmult(vector<int> const& a.

    vector<int> const& b) {
vector < cd > fa(a.begin(), a.end()),
fb(b.begin(), b.end());
int n = 1 << (32 - __builtin_clz(a.size() +</pre>

    b.size() - 1));
fa.resize(n); fb.resize(n);
fft(fa, false); fft(fb, false);
for (int i = 0; i < n; i++) fa[i] *= fb[i]; fft(fa, true);
 vector<int> toret(n);
for (int i = 0; i < n; i++) toret[i] =

→ round(fa[i].real());

return toret:
Greatest Common Denominator
ll egcd(ll a, ll b, ll& x, ll& v) {
if (b == 0) \{ x = 1; y = 0; return a; \}
ll gcd = egcd(b, a % b, x, y);
x = a / b * y;
swap(x, y);
return gcd:
Josephus Problem
// O-indexed. arbitrary k
int josephus(int n, int k) {
if (n == 1) return 0;
if (k == 1) return n-1;
if (k > n) return (josephus(n-1,k)+k)%n;
 int res = josephus(n-n/k,k)-n\%k;
return res + ((res<0)?n:res/(k-1)):
\frac{1}{2} fast case if k=2, traditional josephus
int josephus(int n) {
return 2*(n-(1<<(32-\_builtin_clz(n)-1)));
```

.v.push_back({h, k});

```
Least Common Multiple
                                                        Matrix Exponentiation
                                                           (F(n) = c[0]*F(n-1) + c[1]*F(n-2) + \dots
(b) is the base cases of same length c
#define lcm(a,b) ((a*b)/acd(a,b))
                                                         11 matrix exponentiation(11 n, vector<11> c.
Modulo Operations
                                                         vector<11> b) {
   if (nth < b.size()) return b[nth-1];
   Mat<11> a(c.size(), c.size()); l1 s = 0;
   for (int i = 0; i < c.size(); i++) a[i][0] =</pre>
 #define MOD 1000000007
 \#define\ madd(a,b,m)\ (a+b-((a+b-m>=0)?m:0))
#define mult(a,b,m) ((ull)a*b%m)
#define msub(a,b,m) (a-b+((a<b)?m:0))
                                                          → c[i];
11 mpow(ll b, ll e, ll m) {
                                                          for (int i = 0; i < c.size() - 1; i++)
 11 x = 1;
                                                          \rightarrow a[i][i+1] = 1;
 while (e > 0) {
    if (e % 2) x = (x * b) % m;
    b = (b * b) % m;
                                                          a = \overline{a.power(nth - c.size())};
                                                         for (int i = 0; i < c.size(); i++)
s += a[i][0] * b[i];
return s;
  e /= 2;
 return x % m;
                                                         Nimber Arithmetic
ull mfac(ull n, ull m) {
                                                         #define nimAdd(a,b)
 ull f = 1;
for (int i = n; i > 1; i--)
                                                         ull nimMul(ull a, ull b, int i=6) {
                                                           static const ull M[]={INT_MIN>>32,
 f = (f * i) \% m;
                                                             M[0]^(M[0] << 16), M[1]^(M[1] << 8),
 return f;
                                                             M[2]^(M[2] << 4), M[3]^(M[3] << 2),
                                                            M[4]^{M[4]}<<1);
 // if m is not guaranteed to be prime
11 minv(11 b, 11 m) {
11 x = 0, y = 0;
if (egcd(b, m, x, y) != 1) return -1;
                                                           if (i--==0) return a&b;
                                                           int k=1<<i:
                                                           ull s=nimMúl(a,b,i), m=M[5-i],
                                                             t=nimMul(((a^(a>>k))&m)|(s\&~m),
 return (x % m + m) % m;
                                                             ((b^(b>>k))&m)|(m&(\sim m>>1))<< k, i);
11 mdiv_compmod(int a, int b, int m) {
                                                           return ((s^t)\&m)<\langle k|((s^(t)>k))\&m);
 if (__gcd(b, m) != 1) return -1;
 return mult(a, minv(b, m), m);
                                                         Permutation
                                                         //c = array \ size, \ n = nth \ perm, \ return \ index
 \frac{1}{1} if m is prime (like 10^{9}+7)
                                                         vector<int> gen_permutation(int c, int n) {
11 mdiv_primemod (int a, int b, int m) {
                                                          vector<int> idx(c), per(c), fac(c); int i;
 return mult(a, mpow(b, m-2, m), m);
                                                         for (i = 0; i < c; i++) idx[i] = i; for (i = 1; i <= c; i++) fac[i-1] = n%i, n/=i; double Simpsons (double a, double b, int k,
                                                          for (i = c'-1; i >= 0; i--)
per[c-i-1] = idx[fac[i]],
Matrix
template<typename T>
                                                           idx.erase(idx.begin() + fac[i]);
 struct Mat : public Vec<2, T> {
                                                          return per;
 int w, h;
 Mat(int x, int y) : Vec<2, T>(x, y), w(x),
                                                         // get what nth permutation of vector
 \hookrightarrow h(y) {}
                                                         int get_permutation(vector<int>& v) {
 static Mat<T> identity(int n) { Mat<T> m(n,n);
                                                         int use = 0, i = 1, r = 0;
for (int e: v) {
   r = r * i++ + __builtin_popcount(use &
    for (int i=0:i<n:i++) m[i][i] = 1: return
 \rightarrow -(1<<e));
 .Mat<T>& operator+=(const Mat<T>& m) {
                                                           use |= 1 << e;
  for (int i = 0; i < w; i++)
  for (int j = 0; j < h; j++)
(*this)[i][j] += m[i][j];
                                                          return r;
  return *this;
                                                         Permutation (string/multiset)
 Mat<T>& operator-=(const Mat<T>& m) {
                                                         string freq2str(vector<int>& v) {
  for (int i = 0; i < w; i++)
  for (int j = 0; j < h; j++)
(*this)[i][j] -= m[i][j];
                                                          string s;
                                                          for (int i = 0; i < v.size(); i++)
                                                          for (int j = 0; j < v[i]; j++)
s += (char)(i + 'A');
  return *this;
                                                          return s:
 Mat<T> operator*(const Mat<T>& m) {
  Mat < T > z(w,m.h);
                                                         // nth perm of multiset, n is 0-indexed
  for (int i = 0; i < w; i++)
                                                        string gen_permutation(string s, ll n) {
  for (int j = 0; j < h; j++)
                                                          vector<int> freq(26, 0);
  for (int^*k = 0; k < m.h; k++)

z[i][k] += (*this)[i][j] * m[j][k];
                                                          for (auto e : s) freq[e - 'A']++;
                                                          for (int i = 0; i < 26; i++) if (freq[i] > 0)
    return z:
 |Mat<T> operator+(const Mat<T>& m) { Mat<T>
                                                           freq[i]--; ll v = multinomial(freq);

→ a=*this: return a+=m: }

                                                           if (n < v) return (char)(i+'A') +
 Mat<T> operator-(const Mat<T>& m) { Mat<T>

gen permutation(freg2str(freg), n);

→ a=*this; return a-=m; }

                                                          freq[i]++; n -= v;
 Mat<T>& operator*=(const Mat<T>& m) { return
                                                          return "":

    *this = (*this)*m; }

 Mat<T> power(int n) {
  Mat<T> a = Mat<T>::identity(w),m=*this;
                                                        Miller-Rabin Primality Test
  for (;n;n/=2,m*=m) if (n\&1) a *= m;
                                                         // Miller-Rabin primality test - O(10 log^3 n)
  return à;
                                                        bool isPrime(ull n) {
  if (n < 2) return false;</pre>
```

```
if (n % 2 == 0) return false;
ull s = n - 1;

while (s % 2 == 0) s /= 2;

for (int i = 0; i < 10; i++) {
  ull temp = s;
  ull a = rand() \% (n - 1) + 1;
  ull mod = mpow(a, temp, n);
  while (temp!=n-1\&\&mod!=1\&\&mod!=n-1) {
   mod = mult(mod, mod, n);
   temp *= 2;
  if (mod!=n-1&&temp%2==0) return false;
 return true:
Sieve of Eratosthenes
bitset<100000001> sieve;
// generate sieve - O(n log n)
void genSieve(int n) {
sieve[0] = sieve[1] = 1;
for (ull i = 3; i * i < n; i += 2)
    if (!sieve[i])</pre>
  for (ull j = i * 3; j \le n; j += i * 2)
     sieve[j] = 1;
// query sieve after it's generated - O(1)
bool querySieve(int n) {
 return n == 2 || (n % 2 != 0 && !sieve[n]);
Simpson's / Approximate Integrals
// integrate f from a to b, k iterations // error <= (b-a)/18.0 * M * ((b-a)/2k)^4
// where M = max(abs(f^{(i)}(x))) for x in [a,b] // "f" is a function "double func(double x)"
    double (*f)(double)) {
double dx = (b-a)/(2.0*k), t = 0;

for (int i = 0; i < k; i++)

t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
 \rightarrow (*f)(a+(2*i+1)*dx);
 return (t + (*f)(b)) * (b-a) / 6.0 / k;
Common Equations Solvers
// ax^2 + bx + c = 0, find x
vector<double> solveEq(double a, double b,
 → double c) {
vector<double> r;
 double z = b * b - 4 * a * c;
if (z == 0)
 r.push_back(-b/(2*a));
 else if (z > 0) {
 r.push_back((sqrt(z)-b)/(2*a));
  r.push_back((sqrt(z)+b)/(2*a));
 return r:
\frac{1}{1} = \frac{1}{1} ax^3 + bx^2 + cx + d = 0, find x
vector<double> solveEq(double a, double b,
    double c, double d) {
 vector<double> res;
 long double a1 = b/a, a2 = c/a, a3 = d/a;
 long double q = (a1*a1 - 3*a2)/9.0, sq =
 \rightarrow -2*sqrt(q);
 long double r = (2*a1*a1*a1 - 9*a1*a2 +
 \rightarrow 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
 if (z \le 0) {
  theta = acos(r/sqrt(q*q*q));
  res.push_back(sq*cos(theta/3.0) - a1/3.0);
  res.push back(sq*cos((theta+2.0*PI)/3.0) -
  res.push_back(sq*cos((theta+4.0*PI)/3.0) -
    a1/3.0);
 res.push back(pow(sqrt(z)+fabs(r), 1/3.0));
```

if (n == 2) return true;

```
res[0] = (res[0] + q / res[0]) * ((r<0)?1:-1)
\rightarrow - a1 / 3.0;
return res;
// linear diophantine equation ax + by = c,
   find x and u
// infinite solutions of form x+k*b/g, y-k*a/g bool solveEq(ll a, ll b, ll c, ll &x, ll &y, ll
g = egcd(abs(a), abs(b), x, y);
if (c % g) return false;
x *= c / g * ((a < 0) ? -1 : 1);
y *= c / g * ((b < 0) ? -1 : 1);
return true:
// m = # equations, n = # variables, a[m][n+1]
\rightarrow = coefficient matrix
// a[i][0]x + a[i][1]y + ... + a[i][n]z =
   a[i][n+1]

    a[i][n+1]
// find a solution of some kind to linear

\hookrightarrow equation
const double eps = 1e-7;
bool zero(double a) { return (a < eps) && (a >
vector<double> solveEq(double **a, int m, int
\hookrightarrow n) {
 int cur = 0;
 for (int i = 0; i < n; i++) {
 for (int j = cur; j < m; j++) {
   if (!zero(a[j][i])) {
    if (j != cur) swap(a[j], a[cur]);
    for (int sat = 0; sat < m; sat++) {
     if (sat == cur) continue;
double num = a[sat][i] / a[cur][i];
     for (int sot = 0; sot <= n; sot++)
[ a[sat][sot] -= a[cur][sot] * num;
    cur++;
    break;
for (int j = cur; j < m; j++)
  if (!zero(a[j][n])) return vector<double>();
 vector < double > ans(n,0);
for (int i = 0, sat = 0; i < n; i++)
    if (sat < m && !zero(a[sat][i]))
    nan[i] = a[sat][n] / a[sat++][i];
    return ans;
// solve A[n][n] * x[n] = b[n] linear equation
// rank < n is multiple solutions, -1 i\bar{s} no

→ solutions
// `alls` is whether to find all solutions, or
\hookrightarrow anu
const double eps = 1e-12;
int solveEq(Vec<2, double>& A, Vec<1, double>&

→ b, Vec<1, double>& x, bool alls=false) {
int n = A.size(), m = x.size(), rank = 0, br,
→ bc;
vector<int> col(m); iota(begin(col), end(col),
for(int i = 0; i < n; i++) {
 double v, bv = 0;
for(int r = i; r < n; r++)
  for(int c = i; c < n; c++)
if ((v = fabs(A[r][c])) > bv)
br = r, bc = c, bv = v;
  if (bv <= eps) {
  for(int j = i; j < n; j++)
if (fabs(b[j]) > eps)
     return -1:
   break;
  swap(A[i], A[br]);
  swap(b[i], b[br]);
  swap(col[i], col[bc]);
  for(int j = 0; j < n; j++)
swap(A[j][i], A[j][bc]);
```

```
bw = 1.0 / A[i][i];
for(int j = (alls)?0:i+1; j < n; j++) {
    if (j != i) {</pre>
    double fac = A[i][i] * bv:
    b[j] = fac * b[i];
    for(int k = i+1; k < m; k++)
A[j][k] -= fac*A[i][k];
  rank++;
 if (alls) for (int i = 0; i < m; i++) x[i] =
→ ¬DBL_MAX;
 for (int i = rank; i--;) {
  bool isGood = true;
 if (aļls)
  for (int j = rank; isGood && j < m; j++)
...if (fabs(A[i][j]) > eps)
 ...isGood = false;

.b[i] /= A[i][i];

.if (isGood) x[col[i]] = b[i];
  if (!alls)
  for(int j = 0; j < i; j++)
b[j] -= A[j][i] * b[i];
 return rank;
Graycode Conversions
ull graycode2ull(ull n) {
for (; n; n = n >> 1) i ^= n;
return i;
```

Unix/Epoch Time

ull ull2graycode(ull n) {
 return n ^ (n >> 1);

```
// 0-indexed month/time, 1-indexed day
// minimum 1970, 0, 1, 0, 0, 0
ull toEpoch(int year, int month, int day, int
→ hour, int minute, int second) {
struct tm t; time_t epoch;
t.tm_year = year - 1900; t.tm_mon = month;
t.tm_mday = day; t.tm_hour = hour;
t.tm_min = minute; t.tm_sec = second;
t.tm_isdst = 0; // 1 = daylights savings
 epoch = mktime(\&t);
return (ull)epoch:
vector<int> toDate(ull epoch) {
time t e=epoch; struct tm t=*localtime(&e);
return {t.tm_year+1900,t.tm_mon,t.tm_mday,t_
   .tm hour.t.tm min.t.tm sec}:
int getWeekday(ull epoch) {
time_t e=epoch; struct tm t=*localtime(&e);
return t.tm wday: // 0-6. 0 = sunday
int getDavofYear(ull epoch) {
time_t e=epoch; struct tm t=*localtime(&e);
return t.tm_yday; // 0-365
const int months[] =
→ {31,28,31,30,31,30,31,31,30,31,30,31};
bool validDate(int year, int month, int day) {
    bool leap = !(vear%(vear%25?4:16));
    if (month >= 12) return false;
    return day <= months[month] + (leap &&
   month == 1):
```

Theorems and Formulae

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set of size n. !0 = 1, !1 = 0, !n = (n+1)(!(n-1))1)+!(n-2)), ! $n = n! \sum_{i=0}^{n} \frac{(-1)^i}{i!}$, ! $n = \lceil \frac{n!}{i!} \rceil$

In a partially ordered set, a chain is a subset of elements that are all comparable to eachother An antichain is a subset where no two are comparable.

Dilworth's theorem states the size of a max-| ⇒ imal antichain equals the size of a minimal chain cover of a partially ordered set S. The width of S is the maximum size of an antichain in S, which is equal to the minimum number of chains needed to cover S, or the minimum number of chains such that all elements are in at least one chain.

Rosser's Theorem states the nth prime 3 number is greater than n * ln(n) for n > 1.

 $(n^{\frac{n+1}{2}})^2$.

Lagrange's Four Square Theorem states & m[k][j] != inf) { of four non-negative integers. This is a special case of the Fermat Polygonal Number FOR(k,n) if (m[k][k] < 0) FOR(i,n) FOR(i,n) Theorem where every positive integer is a if (m[i][k] != inf && m[k][j] != inf) sum of at most n s-gonal numbers. The $nth|_{\gamma}^{-1}$ $m[i][j] = -\inf;$ s-gonal number $P(s, n) = (s - 2)\frac{n(n-1)}{2} + n$

```
7 Graphs
struct edge {
int u,v,w;
edge (int u,int v,int w) : u(u),v(v),w(w) {}
edge (): u(0), v(0), w(0) {}
bool operator < (const edge &e1, const edge
bool operator > (const edge &e1, const edge
\rightarrow &e2) { return e1.w > e2.w; }
```

```
struct subset { int p, rank; };
Eulerian Path
#define edge_list vector<edge>
#define adi sets vector<set<int>>
struct EulerPathGraph {
adj_sets graph; // actually indexes incident
 → edges
edge_list edges; int n; vector<int> indeg;
EulerPathGraph(int n): n(n) {
 indeg = *(new vector<int>(n,0));
 graph = *(new adj_sets(n, set<int>()));
 void add edge(int u, int v) {
 graph[u].insert(edges.size());
indeg[v]++;
  edges.push back(edge(u.v.0)):
bool eulerian_path(vector<int> &circuit) {
 if(edges.size()==0) return false:
  stack<int> st;
int a[] = {-1, -1};
 for(int v=0;v<n;v++) {
  if(indeg[v]!=graph[v].size()) {
   bool b = indeg[v] > graph[v].size();
if (abs(((int)indeg[v])-((int)graph[v])
   .size())) > 1) return
false;
if (a[b] != -1) return false;
```

[b] = v;

```
int s = (a[0]!=-1 && a[1]!=-1 ? a[0] :
→ (a[0]==-1 && a[1]==-1 ? edges[0].u : -1));
                                                                                                                          if(s==-1) return false;
                                                                                                                           while(!st.empty() || !graph[s].empty()) {
   if (graph[s].empty()) {
                                                                                                                                circuit.push back(s): s = st.top():
                                                                                                                              st.pop(): }
                                                                                                                             else {
                                                                                                                                int w = edges[*graph[s].begin()].v;
                                                                                                                                graph[s].erase(graph[s].begin());
                                                                                                                                st.push(s): s = w:
                                                                                                                           .}
.circuit.push_back(s);
                                                                                                                         return circuit.size()-1==edges.size():
                                                                                                                     Floyd Warshall
                                                                                                                      const 11 inf = 1LL << 62:
Nicomachi's Theorem states 1^3 + 2^3 + ... + \frac{\text{const } 11 \text{ inf}}{\text{#define } FOR(i,n) \text{ for } (int \text{ } i = 0; \text{ } i < n; \text{ } i++)}
n^3 = (1+2+...+n)^2 and is equivalent to value of the property of the property
                                                                                                                       int n = m.size();
FOR(i,n) m[i][i] = min(m[i][i], OLL);
FOR(k,n) FOR(i,n) FOR(j,n) if (m[i][k] != inf
every natural number is the sum of the squares auto newDist = max(m[i][k] + m[k][j], -inf);
                                                                                                                          m[i][j] = min(m[i][j], newDist);
                                                                                                                    Minimum Spanning Tree
                                                                                                                        // returns vector of edges in the mst
                                                                                                                       // graph[i] = vector of edges incident to
                                                                                                                              vertex i places total weight of the mst in Stotal
                                                                                                                       // if returned vector has size != n-1, there is
                                                                                                                      vector<edge> mst(vector<vector<edge>> graph.
                                                                                                                        priority_queue<edge, vector<edge>,
                                                                                                                         → greater<edge>> pq;
                                                                                                                         vector<edge> MST;
                                                                                                                         bitset<20001> marked; // change size as needed
                                                                                                                         marked[0] = 1;
                                                                                                                        for (edge ep : graph[0]) pq.push(ep);
while(MST.size()!=graph.size()-1 &&
                                                                                                                         → pq.size()!=0) {
                                                                                                                           edge e = pq.top(); pq.pop();
                                                                                                                           int u = e.u, v = e.v, w = e.w;
if(marked[u] && marked[v]) continue;
                                                                                                                           else if (marked[u]) swap(u, v);
                                                                                                                           for(edge ep : graph[u]) pq.push(ep);
                                                                                                                           marked[u] = 1;
MST.push_back(e);
                                                                                                                           total += e.w;
                                                                                                                         return MST;
                                                                                                                       Union Find
                                                                                                                      int uf find(subset* s, int i) {
  if (s[i].p != i) s[i].p = uf_find(s, s[i].p);
                                                                                                                        return s[i].p:
                                                                                                                       void uf_union(subset* s, int x, int y) {
                                                                                                                         int xp = uf_find(s, x), yp = uf_find(s, y);
if (s[xp].rank > s[yp].rank) s[yp].p = xp;
                                                                                                                         else if (s[xp].rank < s[yp].rank) s[xp].p =
```

2D Geometry

```
|#define point complex<double>
#define EPS 0.0000001
#define sa(a) ((a)*(a))
#define c\bar{b}(a) ((a)*(a)*(a))
double dot(point a, point b) { return

    real(conj(a)*b);
}
double cross(point a, point b) { return

    imag(conj(a)*b); }

struct line { point a, b; };
struct circle { point c; double r; };
struct segment { point a, point b; };
struct triangle { point a, b, c; };
struct rectangle { point tl, br; };
struct convex_polygon {
 vector<point points;
 convex_polygon(vector<point> points) :
 → points(points) {}
 convex_polygon(triangle a) {
  points.push_back(a.a); points.push_back(a.b);
    points.push back(a.c);
 convex polygon(rectangle a) {
  points.push_back(a.tl);
    points.push_back({real(a.tl),
    imag(a.br)});
  points.push_back(a.br);
    points.push_back({real(a.br),
    imag(a.tl)}):
struct polygon {
 vector <point > points;
 polygon(vector point points) : points(points)
 polygon(triangle a) {
  points.push_back(a.a); points.push back(a.b);
    points.push back(a.c):
 polygon(rectangle a) {
  points.push_back(a.tl);
    points.push back({real(a.tl).
    imag(a.br)});
  points.push back(a.br):
    points.push back({real(a.br).
    imag(a.tl)}):
 polygon(convex_polygon a) {
  for (point v : a.points)
   points.push_back(v);
   triangle methods
double area heron(double a, double b, double c)
 \overrightarrow{if} (a < b) swap(a, b);
 if (a < c) swap(a, c);
 if (b < c) swap(b, c);
 if (a > b + c) return -1;
return sqrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)
→ /16.0):
// seament methods
double lengthsq(segment a) { return
    sq(real(a.a) - real(a.b)) + sq(imag(a.a) -
   imag(a.b)); }
double length(segment a) { return

    sgrt(lengthsq(a)); }

// circle methods
double circumference(circle a) { return 2 * a.r

→ * M PI: }

double area(circle a) { return sq(a.r) * M_PI;
→ }
// rectangle methods
```

```
double width(rectangle a) { return

→ abs(real(a.br) - real(a.tl)); }
double height (rectangle a) { return

→ abs(imag(a.br) - real(a.tl)); }

double diagonal(rectangle a) { return

    sqrt(sq(width(a)) + sq(height(a))); }

double area(rectangle a) { return width(a) *
→ height(a); }
double perimeter(rectangle a) { return 2 *
   (width(a) + height(a)); }
// check if `a` fit's inside `b
// swap equalities to exclude tight fits
bool doesFitInside(rectangle a, rectangle b) {
int x = width(a), w = width(b), y = height(a),
\rightarrow h = height(b);
if (x > y) swap(x, y);
if (w > h) swap(w, h);
if (w < x) return false;
if (y <= h) return true;</pre>
double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
return sq(a) \le sq(b) + sq(c);
// polygon methods
// negative area = CCW, positive = CW
double area(polygon a) {
  double area = 0.0; int n = a.points.size();
 for (int i = 0, j = 1; i < n; i + +, j = (j - 1)

    1) % n)
area +=

    (real(a.points[j]-a.points[i]))*(imag(a | )
    .points[j]+a.points[i]));
 return area / 2.0;
// get both unsigned area and centroid
pair<double, point> area_centroid(polygon a) {
 int n = a.points.size():
double area = 0;
point c(0, 0);
for (int i = n - 1, j = 0; j < n; i = j++) {
 double v = cross(a.points[i], a.points[j]) /
 area += v:
 c += (a.points[i] + a.points[j]) * (v / 3);
c /= area;
.return {area, c};
Intersection
// -1 coincide, 0 parallel, 1 intersection
int intersection(line a, line b, point& p) {
if (abs(cross(a.b - a.a, b.b - b.a)) > EPS) {
p = cross(b.a - a.a, b.b - a.b) / cross(a.b)
\rightarrow a.a, b.b - b.a) * (b - a) + a;
 return 1;
if (abs(cross(a.b - a.a, a.b - b.a)) > EPS)

→ return 0;

return -1:
// area of intersection
double intersection(circle a, circle b) {
double d = abs(a.c - b.c);
 if (d <= b.r - a.r) return area(a);
if (d <= a.r - b.r) return area(b);
if (d >= a.r + b.r) return 0;
double alpha = acos((sq(a.r) + sq(d) -
\rightarrow sq(b.r)) / (2 * a.r * d));
double beta = acos((sq(b.r) + sq(d) - sq(a.r))
\rightarrow / (2 * b.r * d)):
return sq(a.r) * (alpha - 0.5 * sin(2 *
    alpha)) + sq(b.r) * (beta - 0.5 * sin(2 *
   beta)):
// -1 outside, 0 inside, 1 tangent, 2
int intersection (circle a, circle b,
→ vector<point>& inter) {
```

```
double d2 = norm(b.c - a.c), rS = a.r + b.r,
\rightarrow rD = a.r - b.r;
if (d2 > sq(rS)) return -1;
 if (d2 < sq(rD)) return 0;
 double ca = 0.5 * (1 + rS * rD / d2)
 point z = point(ca, sqrt(sq(a.r) / d2 -
   sq(ca))):
 inter.push_back(a.c + (b.c - a.c) * z);
 if (abs(imag(z)) > EPS) inter.push_back(a.c +
   (b.c - a.c) * conj(z));
 return inter.size();
 // points of intersection
vector<point> intersection(line a, circle c) {
vector ont inter;
c.c = a.a;
a.b = a.a;
 point m = a.b * real(c.c / a.b);
 double d2 = norm(m - c.c);
 if (d2 > sq(c.r)) return 0;
 double l = sqrt((sq(c.r) - d2) / norm(a.b));
 inter.push_back(a.a + m + 1 * a.b);
 if (abs(1) > EPS) inter.push_back(a.a + m - 1
 return inter:
 / area of intersection
max(imag(a.tl), imag(b.tl));
 double x2 = min(real(a.br), real(b.br)), y2 =
   min(imag(a.br), imag(b.br));
 return (x2 <= x1 | | y2 <= y1) ? 0 :
   (x2-x1)*(y2-y1);
Convex Hull
bool cmp(point a, point b) {
 if (abs(real(a) - real(b)) > EPS) return
 → real(a) < real(b);
if (abs(imag(a) - imag(b)) > EPS) return
   imag(a) < imag(b);</pre>
 return false:
convex_polygon convexhull(polygon a) {
 sort(a.points.begin(), a.points.end(), cmp);
 vector<point> lower, upper;
 for (int i = 0; i < a.points.size(); i++) {
  while (lower.size() >= 2 &&
    cross(lower.back() - lower[lower.size()
   2], a.points[i] - lower.back()) < EPS)
   lower.pop_back();
  while (upper.size() >= 2 &&
    cross(upper.back() - upper[upper.size() -
   2], a.points[i] - upper.back()) > -EPS)
   upper.pop_back();
  lower.push_back(a.points[i]);
  upper.push_back(a.points[i]);
 lower.insert(lower.end(), upper.rbegin() + 1,
   upper.rend());
 return convex polygon(lower);
    3D Geometry
struct point3d {
 double x, y, z;
 point3d operator+(point3d a) const { return
```

```
\rightarrow {x+a.x, y+a.y, z+a.z}; }
point3d operator*(double a) const { return
\rightarrow {x*a, y*a, z*a}; }
point3d operator-() const { return {-x, -y,
→ -z}: }
point3d operator-(point3d a) const { return
\rightarrow *this + -a; }
point3d operator/(double a) const { return
\rightarrow *this * (1/a): }
```

```
double norm() { return x*x + y*y + z*z; }
 double abs() { return sqrt(norm()); }
 point3d normalize() { return *this /

    this->abs(); }

double dot(point3d a, point3d b) { return
\rightarrow a.x*b.x + a.y*b.y + a.z*b.z; }
point3d cross(point3d a, point3d b) { return
    \{a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,
\stackrel{\Rightarrow}{\Rightarrow} a.x*b.y - a.y*b.x}; }
struct line3d { point3d a, b; };
struct plane { double a, b, c, d; } // a*x +
\Rightarrow b*y + c*z + d = 0
struct sphere { point3d c; double r; };
#define sq(a) ((a)*(a))
#define cb(a) ((a)*(a)*(a))
double surface(circle a) { return 4 * sq(a.r)
double volume(circle a) { return 4.0/3.0 *
\hookrightarrow cb(a.r) * M_PI; }
10 Optimization
```

```
Snoob
// SameNumberOfOneBits, next permutation
// example usage
int main() {
   char 11[] = {'1', '2', '3', '4', 'char 12[] = {'a', 'b', 'c', 'd'};
   int di = 5, d2 = 4;
   // prints 12345abcd, 1234a5bcd, ...
  int min = (1 << d1) - 1, max = min << d2;
  for (int i = min; i <= max; i = snoob(i)) {
  int p1 = 0, p2 = 0, v = i;

while (p1 < d1 || p2 < d2) {

    cout << ((v & 1) ? 11[p1++] : 12[p2++]);
   v /= 2;
   cout << '\n';
```

Powers bool isPowerOf2(ll a) { return a > 0 && ! (a & a-1);bool isPowerOf3(11 a) { return a>0&&!(12157665459056928801u11%a);

bool isPower(ll a, ll b) {
 double x = log(a) / log(b); return abs(x-round(x)) < 0.00000000001;

11 Additional

Judge Speed

```
kattis: 0.50s
// codeforces: 0.421s
// atcoder: 0.455s
#include <bits/stdc++.h>
using namespace std;
int v = 1e9/2, p = 1;
int main() {
for (int i = 1; i <= v; i++) p *= i;
cout << p;
```

Judge Pre-Contest Checks

int128 and float128 support? -does extra or missing whitespace cause WA? -documentation up to date? -printer usage available and functional?

```
// each case tests a different fail condition
// try them before contests to see error codes
struct g { int arr[1000000]; g(){}};
// O=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE
    6=SIGSEGV 7=recursive MLE judge(int n) {
    (n == 0) exit(0)
if (n == 1) while(1);
if (n == 2) while(1) a.push_back(g());
if (n == 3) while(1) putchar_unlocked('a');
if (n == 4) assert(0);
if (n == 5) 0 / 0;
if (n == 6) *(int*)(0) = 0;
 return n + judge(n + 1);
```

GCC Builtin Docs

```
// 128-bit integer
   int128 a;
unsigned __int128 b;
// 128-bit float
 // minor improvements over long double
__float128 c;
// log2 floor
 __lg(n);
// number of 1 bits
 // can add ll like popcountll for long longs
__builtin_popcount(n);
// number of trailing zeroes
__builtin_ctz(n);
__Dulitin_cuz(n),
// number of leading zeroes
__builtin_clz(n);
// 1-indexed least significant 1 bit
 builtin ffs(n);
 // parity of number
 __builtin_parity(n);
```

```
Limits
                        \pm 2147483647 \mid \pm 2^{31} - 1 \mid 10^9
int
                                              \frac{1}{2}32 - \frac{1}{1}|\tilde{1}\tilde{0}^9|
                          4294967295
uint
         \pm 9223372036854775807 | \pm \overline{2}^{63} - 1 | 10^{18}
11
                                               |\overline{2}^{64} - 1| 10^{19}
ull
         18446744073709551615
|i128| \pm 170141183460469231... | \pm 2^{127} - 1 | 10^{38}
                                              \frac{1}{2}^{128} - 1|10^{38}
u128 340282366920938463...
Complexity classes input size (per second):
```

 $O(n^n)$ or O(n!) $O(2^n)$ n < 30 $O(n^{3})$ n < 1000 $O(n^2)$ n < 30000 $O(n\sqrt{n})$ $n < 10^6$ $n < 10^7$

 $O(n \log n)$ $n < 10^9$ O(n)