```
else if (d) { m=m*10+c-'0'; o*=0.1; } else n = n * 10 + c - '0':
    General
                               7 Graphs
    Algorithms
                               8 2D Geometry
                                                               n = s * (n + m * o):
    Structures
                               9 3D Geometry
    Strings
                                                               void read(double& n) {
                               10 Optimization
                                                               ld m; read(m); n = m;
    Greedy
                               11 Python
                                                               void read(float& n) {
  ld m; read(m); n = m;
    Math
                               12 Additional
     General
                                                               void read(string& s) {
                                                               char c; s = "";
while((c=getchar unlocked())!=' '&&c!='\n')
run.sh
g++ -g -02 -std=gnu++17 -static prog.cpp
                                                                s += c:
./a.exe
                                                               bool readline(string& s) {
                                                               char c; s = ""
# compile and test all *.in and *.ans
                                                               while(c=getchar_unlocked()) {
g++ -g -02 -std=gnu++17 -static prog.cpp
                                                                if (c == '\n') return true;
if (c == EOF) return false;
for i in *.in; do
f=${i%.in}
 ./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
                                                               return false:
                                                               void print(unsigned int n) {
                                                               if (n / 10) print(n / 10);
// use better compiler options
                                                               putchar_unlocked(n % 10 + '0'):
#pragma GCC optimize("Ofast", "unroll-loops")
#pragma GCC target("avx2, fma")
// include everything
                                                               void print(int n) {
                                                               if (n < 0) { putchar unlocked('-'); n*=-1; }
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <sys/resource.h>
                                                               print((unsigned int)n);
// namespaces
                                                               Common Structs
using namespace std;
                                                               // n-dimension vectors
// Vec<2, int>v(n, m) = arr[n][m]
// Vec<2, int>v(n, m, -1) default init -1
using namespace __gnu_cxx; // rope
using namespace __gnu_pbds; // tree/trie
// common defines
                                                               template<int D, typename T>
#define fastio
                                                               struct Vec : public vector < Vec < D-1, T >> {
                                                               template<typename... Args>
→ ios_base::sync_with_stdio(0);cin.tie(0);
                                                               Vec(int n=0, Args... args) : vector<Vec<D-1,</pre>
#define nostacklim rlimit RZ: getrlimit(3.&RZ)
                                                               \rightarrow T>>(n, Vec<D-1, T>(args...)) {}
\rightarrow ):RZ.rlim cur=-1:setrlimit(3.\&RZ):
                                                              };
template<typename T>
#define DEBUG(v) cerr<<__LINE__\ << ": " << #v << " =
\rightarrow "<<v<<'\n'; #define TIMER
                                                              struct Vec<1, T>: public vector<T> {
   Vec(int n=0, T val=T()) : vector<T>(n, val) {}
 \begin{array}{l} \hookrightarrow \quad cerr << 1.0*clock()/CLOCKS\_PER\_SEC << "s \ "; \\ \textit{\#define ll long long} \\ \textit{\#define ull unsigned ll} \end{array} 
                                                                  Algorithms
#define i128 _ int128
#define u128 unsigned i128
#define ld long double
                                                               Binary Search
                                                               // search for k in [p,n)
// global variables
                                                               template<typename T>
mt19937 rng((uint32_t)chrono::steady
                                                               int binsearch(T x[], int k, int n, int p = 0) {
                                                               for (int i = n; i >= 1; i /= 2)

. while (p+i < n && x[p+i] <= k) p += i;

. return p; // bool: x[p] == k;

    clock::now().time since epoch().count());

#ifdef _WIN32
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
                                                              Min/Max Subarray
                                                                  max - compare = a < b, reset = a < 0
void read(unsigned int& n) {
                                                               // min - compare = a > b, reset = a > 0
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                              // returns {sum, {start, end}}
pair<int, pair<int, int>>
 n = n * 10 + c - '0';
                                                                    ContiguousSubarray(int* a, int size,
void read(int& n) {
   char c; n = 0; int s = 1;
   if ((c=getchar_unlocked())=='-') s = -1;
                                                                   bool(*compare)(int, int),
                                                               bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
 else n = c - \sqrt{0}:
                                                                \rightarrow 0, s = 0;
 while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                                for (int i = 0; i < size; i++) {
 n = n * 10 + c -
                                                                 cur += a[i];
                                                                 if ((*compare)(best, cur)) { best = cur;
                                                                   start = s; end = i; }
void read(ld& n) {
                                                                if ((*reset)(cur)) { cur = 0; s = i + 1; }
char c; n = 0;

ld m = 0, o = 1; bool d = false; int s = 1;

if ((c=getchar_unlocked())=='-') s = -1;
                                                               return {best, {start, end}};
 else if (c == '.') d = true;
else n = c - '0';
                                                              Quickselect
 while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                               #define QSNE -999999
 if (c == '.') d = true;
```

```
int x = arr[r], i = 1;
                                                               Fenwick Tree
for (int j = 1; j <= r - 1; j++)
if (arr[j] <= x)
swap(arr[i++], arr[j]);
 swap(arr[i], arr[r]);
 return i:
// find k'th smallest element in unsorted array,
→ only if all distinct
int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - 1 + 1)) return QSNE;

swap(arr[1 + rng() % (r-1+1)], arr[r]);

int pos = partition(arr, 1, r);
 if (pos-l==k-1) return arr[pos];
 if (pos-1>k-1) return qselect(arr,1,pos-1,k);
 return qselect(arr, pos+1, r, k-pos+1-1);
// TODO: compare against std::nth_element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted
→ on both axis
pair<int. int> saddleback search(int** arr. int
\hookrightarrow x, int y, int v) {
int i = x-1, j = 0;
while (i >= 0 && j < y) {
  if (arr[i][j] == v) return {i, j};
  (arr[i][j] > v)? i--: j++;
 return {-1, -1};
Ternary Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a. int b. int (*f)(int)) {
 while (b-a > 4) {
                                                               Hashtable
  int m = (a+b)/2;
if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1:
for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
 . a = i;
return a:
#define TERNPREC 0.000001
double ternsearch (double a, double b, double
     (*f)(double)) {
 while (b-a > TERNPREC * 4) {
double m = (a+b)/2;
  if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
  > = m;
else b = m + TERNPREC;
 for (double i = a + TERNPREC; i <= b; i +=
     TERNPREC)
      if (TERNCOMP((*f)(a), (*f)(i)))
 return a;
Golden Section Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
double goldsection(double a, double b, double
 double g(double)) {
  double r = (sqrt(5)-1)/2, eps = 1e-7;
  double x1 = b - r*(b-a), x2 = a + r*(b-a);
  double f1 = f(x1), f2 = f(x2);
 while (b-a > eps)
  while (b-a > eps)

if (TERNCOMP(f2,f1)) {

. b = x2; x2 = x1; f2 = f1;

. x1 = b - r*(b-a); f1 = f(x1);
   a = x1; x1 = x2; f1 = f2;
x2 = a + r*(b-a); f2 = f(x2);
 return a:
```

int partition(int arr[], int 1, int r)

```
3 Structures
```

```
// Fenwick tree, array of cumulative sums
\hookrightarrow O(\log n) updates, O(\log n) gets
struct Fenwick { int n: ll* tree:
 void update(int i, int val) {
 .++i;
while (i <= n) {
  tree[i] += val;</pre>
  i += i & (-i);
 Fenwick(int size) {
 n = size;
 tree = new ll[n+1];
for (int i = 1; i <= n; i++)
  .tree[i] = 0;
 Fenwick(int* arr, int size) : Fenwick(size) {
 for (int i = 0; i < n; i++)
...update(i, arr[i]);
 ~Fenwick() { delete[] tree; }
11 operator[](int i) {
 if (i < 0 || i > n) return 0;
 while (i>0) {
   sum += tree[i];
   i -= i & (-i);
 return sum;
ll getRange(int a, int b) { return
    operator[](b) - operator[](a-1); }
```

# // similar to unordered map, but faster struct chash { const uint64\_t C = (11)(2e18 \* M\_PI) + 71; ll operator()(11 x) const { return builtin bswap64(x\*C); } int main() { gp\_hash\_table<11,int,chash> $\rightarrow$ hashtable({},{},{},{},{1<<16}); for (int i = 0; i < 100; i++) hashtable[i] = 200+i; if (hashtable.find(10) != hashtable.end()) cout << hashtable[i];

### Ordered Set

```
template <typename T>
using oset = tree<T,null_type,less<T>,rb_tree

→ _tag,tree_order_statistics_node_update>;
template <typename T, typename D>
using omap = tree<T,D,less<T>,rb_tree |
     _tag,tree_order_statistics_node_update>;
 int main()
 oset<int> o_set;
o_set.insert(5); o_set.insert(1);

    o_set.insert(3);
// get second smallest element
  cout << *(o_set.find_by_order(1));</pre>
 // number of elements less than k=4
cout << ' ' << o_set.order_of_key(4) << '\n';</pre>
  // equivalent with ordered map
  omap<int,int> o_map;
 o_map[5]=1;o_map[1]=2;o_map[3]=3;
cout << (*(o_map.find_by_order(1))).first;
 cout << ' ' << o_map.order_of_key(4) << '\n';
```

```
Rope
                                                          // print things with prefix "1"
                                                          auto range = trie.prefix_range("1");
// \bar{O}(\log n) insert, delete, concatenate
                                                         for (auto it = range.first; it !=
int main() {
 // generate rope
                                                         → range.second: it++)
 rope<int> v;
                                                           cout << *it <<
for (int i = 0; i < 100; i++)
.v.push_back(i);
                                                         Wavelet Tree
 // move range to front
                                                         using iter = vector<int>::iterator;
 rope<int> copy = v.substr(10, 10);
                                                        struct WaveletTree {
  Vec<2, int> C; int s;
  // sigma = highest value + 1
 v.erase(10, 10):
 v.insert(copy.mutable_begin(), copy);
                                                          WaveletTree(vector<int>& a. int sigma) :
 // print elements of rope
for (auto it : v) cout << it << "";
                                                             s(sigma), C(sigma*2, 0) {
                                                           build(a.begin(), a.end(), 0, s-1, 1);
                                                          void build(iter b. iter e. int L. int U. int
Segment Tree
                                                           u) {
if (L == U) return
//max(a,b), min(a,b), a+b, a*b, qcd(a,b), a*b
struct SegmentTree {
                                                           \underline{int} \widetilde{M} = (\widetilde{L} + U)/2;
 typedef int T;
                                                           C[u].reserve(e-b+1); C[u].push back(0);
 static constexpr T UNIT = INT_MIN;
                                                           for (auto it = b; it != e; ++it)
  C[u].push_back(C[u].back() + (*it<=M));</pre>
 T f(T a, T b) 
 if (a == UNIT) return b;
if (b == UNIT) return a;
                                                           auto p = stable_partition(b, e, [=](int
                                                            i){return i<=M;});
  return max(a,b);
                                                           build(b, p, L, M, u*2);
 int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
                                                           build(p, e, M+1, U, u*2+1);
                                                          // number of occurrences of x in [0,i)
\rightarrow n(n) {}
                                                         int rank(int x, int i) {
   int L = 0, U = s-1, u = 1, M, r;
   while (L != U) {
 SegmentTree(vector<T> arr)

    SegmentTree(arr.size()) {

 for (int i=0:i<arr.size():i++)
                                                           M = (L+U)/2;
r = C[u][i]; u*=2;

    update(i,arr[i]);

                                                           if (x <= M) i = r, U = M;
else i -= r, L = M+1, ++u;
 void update(int pos, T val) {
  for (s[pos += n] = val; pos /= 2;)
   s[pos] = f(s[pos * 2], s[pos*2+1]);
                                                           return i:
                                                          ^{\prime\prime} number of occurences of x in [l,r)
 T query(int b, int e) { // query [b, e)
                                                         int count(int x, int 1, int r) {
  return rank(x, r) - rank(x, 1);
  Tra = UNIT, rb = UNIT;
  for (b+=n, e+=n; b<= b/=2, e/=2) {
    if (b % 2) ra = f(ra, s[b++]);
    if (e % 2) rb = f(s[--e], rb);
                                                         // kth smallest in [l, r)
int kth(int k, int l, int r) const {
int L = 0, U = s-1, u = 1, M, ri, rj;
  return f(ra, rb):
                                                           while (L != U) {
   M = (L+U)/2;
 T get(int p) { return query(p, p+1); }
                                                            ri = C[u][1]; rj = C[u][r]; u*=2;
                                                            if (k \le rj-ri)^{n}l = ri, r = rj, U = M;
Sparse Table
                                                            else k -= řj-rí, l -= ŕi, r -= ŕj,
template < class T> struct SparseTable {
                                                            L = M+1, ++u;
 vector<vector<T>> m;
                                                           return U:
 SparseTable(vector<T> arr) {
  m.push back(arr);
  for (int k = 1: (1<<(k)) <= size(arr): k++)
                                                          // # elements between [x,y] in [l, r)
                                                         mutable int L, U;
  m.push_back(vector<T>(size(arr)-(1<(k)+1));
                                                          int range(int x, int y, int l, int r) const {
  for (int i = 0: i < size(arr)-(1<<k)+1: i
                                                           if (y < x or r <= 1) return 0;
                                                          L = x; U = y;
 m[k][i] = min(m[k-1][i],
                                                          return range(1, r, 0, s-1, 1);
\rightarrow m[k-1][i+(1<<(k-1))]:
}
// min of range [l,r]
                                                          int range(int 1, int r, int x, int y, int u)
                                                            const {
                                                          if (y < L or U < x) return 0;
if (L <= x and y <= U) return r-1;
T query(int 1, int r) {
 int k = _-lg(r-l+1);
                                                           int M = (x+y)/2, ri = C[u][1], rj = C[u][r];
  return \min(m[k][1], m[k][r-(1<< k)+1]);
                                                           return range(ri, rj, x, M, u*2) + range(1-ri
}
};
                                                            r-rj, M+1, v, u*2+1);
                                                          ^{\prime}// # elements <= x in [l, r)
typedef trie<string, null_type,
                                                         int lte(int x, int 1, int r) {
  return range(INT_MIN, x, 1, r);

→ trie_string_access_traits<>,

 pat_trie_tag, trie_prefix_search_node_update>
Strings
 // generate trie
 trie_type trie;
                                                         Aho Corasick
 for (int i = 0; i < 20; i++)
                                                            range of alphabet for automata to consider
 trie.insert(to string(i)); // true if new,
                                                            MAXC = 26, OFFC = 'a' if only lowercase
\hookrightarrow false if old
```

```
|const int MAXC = 256:
const int OFFC = 0:
struct aho_corasick {
  set<pair<int, int>> out;
  int fail; vector<int> go;
  state() : fail(-1), go(MAXC, -1) {}
 vector<state> s;
  int id = 0;
 aho_corasick(string* arr, int size) : s(1) {
  for (int i = 0; i < size; i++) {
   int cur = 0;
   for (int c : arr[i]) {
   if (s[cur].go[c-OFFC] == -1) {
      s[cur].go[c-OFFC] = s.size();
      s.push back(state());
     cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++});
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
    s[0].go[c] = 0;
  queue<int> sq;
  for (int c = 0; c < MAXC; c++) {
   if (s[0].go[c] != 0) {
    s[s[0].go[c]].fail = 0;
    sq.push(s[0].go[c]);
  while (sq.size()) {
   int e = sq.front(); sq.pop();
   for (int c = 0; c < MAXC; c++) {
   if (s[e].go[c] != -1) {
      in (s[e].go[c] != 1) {
int failure = s[e].fail;
while (s[failure].go[c] == -1)
      failure = s[failure].fail;
failure = s[failure].go[c];
      s[s[e].go[c]].fail = failure;
      for (auto length : s[failure].out)
s[s[e].go[c]].out.insert(length);
      sq.push(s[e].go[c]);
  // list of {start pos, pattern id}
  vector<pair<int. int>> search(string text)
 {
  vector<pair<int, int>> toret;
  int cur = 0:
  for (int i = 0; i < text.size(); i++) {
  while (s[cur].go[text[i]-OFFC] == -1)
     cur = s[cur].fail;
    cur = s[cur].go[text[i]-OFFC];
    if (s[cur].out.size())
    for (auto end : s[cur].out)
toret.push back({i - end.first + 1,
     end.second}):
  .}
.return toret;
Boyer Moore
struct defint { int i = -1; };
vector<int> boyermoore(string txt, string pat)
 vector<int> toret; unordered_map<char, defint>|string | lcp(string* arr, int n, bool sorted =
 → badchar;
 int m = pat.size(), n = txt.size();
 for (int i = 0; i < m; i++) badchar[pat[i]].i
 \rightarrow = i;
int s = 0:
 while (s \leq n - m) {
  int j = m - 1;
```

while  $(j \ge 0 \&\& pat[j] == txt[s + j]) j--;$ 

.if (j < 0) {

```
. toret.push_back(s);
  ..s += (s + m < n) ? m - badchar[txt[s +
   mll.i : 1:
 .} else
   s += \max(1, i - badchar[txt[s + i]].i):
 return toret:
English Conversion
const string ones[] = {"", "one", "two",
    "three", "four", "five", "six", "seven", "eight", "nine"];
 const string teens[] ={"ten", "eleven",
    "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen");
const string tens[] = {"twenty", "thirty",
"forty", "fifty", "sixty", "seventy",

"eighty", "ninety");

const string mags[] = {"thousand", "million",
     "billion", "trillion", "quadrillion",
     "quintillion", "sextillion",
     "septillion"};
string convert(int num, int carry) {
 if (num < 0) return "negative " +

    convert(-num, 0):

 if (num < 10) return ones[num];
    (num < 20) return teens[num % 10];
(num < 100) return tens[(num / 10) - 2] +
 \(\text{num}\%10==0?"":"\) + ones[num \% 10];
if (num < 1000) return ones[num / 100];
      (num/100==0?"":" ") + "hundred" +
     (num%100==0?"":" ") + convert(num % 100.
 return convert(num / 1000, carry + 1) + " " +
     mags[carry] + " " + convert(num % 1000)
    0):
string convert(int num) {
 return (num == 0) ? "zero" : convert(num, 0);
 Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
   vector<int> toret:
 int m = txt.length(), n = pat.length();
 int next[n + 1];
 for (int i = \bar{0}; i < n + 1; i++)
 next[i] = 0;
for (int i = 1; i < n; i++) {
   int j = next[i + 1];
   while (j > 0 && pat[j] != pat[i])
   j = next[j];
  if (j > 0 | pat[j] == pat[i])
   next[i + 1] = j + 1;
 for (int i = 0, j = 0; i < m; i++) {
  if (txt[i] == pat[j]) {
  .if (++j == n)
    toret.push_back(i - j + 1);
  } else if (j > 0) {
  j = next[j];
 return toret;
Longest Common Prefix (array)
 // longest common prefix of strings in array
 → false) {
 → IAISE) {
  if (n == 0) return "";
  if (!sorted) sort(arr, arr + n);
  string r = ""; int v = 0;
 while (v < arr[0].length() && arr[0][v] ==

    arr[n-1][v])
    r += arr[0][v++];

 return r:
```

```
| for (auto c : s) v = (c - 'a' + 1) + v *
Longest Common Subsequence
                                                             → HASHER;
 string lcs(string a, string b) {
 int m = a.length(), n = b.length();
                                                            const int MAXN = 1000001;
ull base[MAXN] = {1};
  int L[m+1][n+1];
 for (int i = 0; i <= m; i++) {
    for (int j = 0; j <= n; j++) {
        ...if (i == 0 || j == 0) L[i][j] = 0;
        ...else if (a[i-1] == b[j-1]) L[i][j] =
                                                            void genBase(int n) {
                                                             for (int i = 1; i \le n; i++)

base[i] = base[i-1] * HASHER;
 \hookrightarrow L[i-1][j-1]+1;
                                                            struct advHash {
 ull v, l; vector<ull> wip;
    else L[i][j] = \max(L[i-1][j], L[i][j-1]);
                                                             advHash(string& s): v(0) {
                                                              wip = vector<ull>(s.length()+1);\
  // return L[m][n]; // length of lcs
                                                              wip[0] = 0;
  string out = "";
                                                              for (int i = 0; i < s.length(); i++)
wip[i+1] = (s[i] - 'a' + 1) + wip[i] *
  int i = m - 1, j = n - 1;
  while (i \ge 0) \& j \ge 0 \{ if (a[i] = b[i]) \}
                                                                HASHER;
                                                              1 = s.length(); v = wip[1];
   out = a[i--] + out;
                                                             ull del(int pos, int len) {
   return v - wip[pos+len]*base[1-pos-len] +
   else if (L[i][j+1] > L[i+1][j]) i--;
                                                                wip[pos]*base[1-pos-len]:
  .else j--;
                                                             ull substr(int pos, int len) {
 return out;
                                                              return del(pos+len, (1-pos-len)) -
                                                                wip[pos]*base[len]:
Longest Common Substring
                                                             ull replace(int pos, char c) {
 // l is array of palindrome length at that
                                                              return v - wip[pos+1]*base[l-pos-1] + ((c -
int manacher(string s, int* 1) {
                                                                 'a' + 1) + wip[pos] *
 int n = s.length() * 2;
                                                                HASHER) *base[1-pos-1];
 for (int i = \tilde{0}, j = 0, k; i < n; i += k, j =
 \rightarrow max(j-k, 0)) {
                                                             ull replace(int pos, string s) {
  while (i \ge j \&\& i + j + 1 < n \&\& s[(i-j)/2]]
                                                              // can't increase total string size
 \Rightarrow == s[(i+j+1)/2]) i++;
                                                                 wip[pos+s.size()]*base[l-pos-s.size()], c
  .1[i] = j;
                                                                wip[pos];
  for (k = 1; i >= k \&\& j >= k \&\& l[i-k] !=
                                                              for (int i = 0; i < s.size(); i++)
c = (s[i]-'a'+1) + c * HASHER;
  \rightarrow j^-k; k^{++}) 
 \downarrow 1[i+k] = min(1[i-k], j-k); 
                                                              return r + c * base[1-pos-s.size()];
 return *max element(1, 1 + n);
                                                            Subsequence Count
 Cyclic Rotation (Lyndon)
                                                             // "banana", "ban" >> 3 (ban, ba..n, b..an)
 // simple strings = smaller than its nontrivial
                                                            ull subsequences(string body, string subs) {

→ suffixes

                                                             int m = subs.length(), n = body.length();
 // lyndon factorization = simple strings
                                                             if (m > n) return 0;

ull** arr = new ull*[m+1];

for (int i = 0; i <= m; i++) arr[i] = new
 \hookrightarrow factorized
 // "abaaba" -> "ab", "aab", "a"
 vector<string> duval(string s) {
                                                              \rightarrow ull[n+1]:
                                                             for (int i = 1; i <= m; i++) arr[i][0] = 0;
for (int i = 0; i <= n; i++) arr[0][i] = 1;
  int n = s.length();
  vector<string> lyndon;
 for (int i = 0; i < n;) {
    int j = i+1, k = i;
    for (; j < n && s[k] <= s[j]; j++)
    if (s[k] < s[j]) k = i;
                                                             for (int i = 1; i <= m; i++)

for (int j = 1; j <= n; j++)

arr[i][j] = arr[i][j-1] + ((body[j-1] ==
                                                                subs[i-1])? arr[i-1][j-1] : 0);
   else k++:
                                                             return arr[m][n]:
   for (; i <= k; i += j - k)
    lyndon.push_back(s.substr(i,j-k));
                                                            Suffix Array + LCP
 return lyndon;
                                                            struct SuffixArray {
}
// lexicographically smallest rotation
                                                             vector<int> sa, 1cp;
                                                             SuffixArray(string& s, int lim=256) {

int n = s.length() + 1, k = 0, a, b;
int minRotation(string s) {
  int n = s.length(); s += s;
                                                              vector<int> x(begin(s), end(s)+1), y(n),
  auto d = duval(s); int i = 0, a = 0;
                                                             \rightarrow ws(max(n, lim)), rank(n);
  while (a + d[i].length() < n) a +=
                                                              sa = lcp = y;
iota(begin(sa), end(sa), 0);

    d[i++].length();

 while (i && d[i] == d[i-1]) a -=
                                                              for (int j = 0, p = 0; p < n; j = max(1, j *

    d[i--].length();

return a;
                                                             \rightarrow 2), lim = p) {
                                                               p = ', ---
j; iota(begin(y), end(y), n - j);
for (int i = 0; i < (n); i++)
...if (sa[i] >= j)
Hashing
                                                                  y[p++] = sa[i] -
 #define HASHER 27
                                                                fill(begin(ws), end(ws), 0);
ull basicHash(string s) {
 ull v = 0;
                                                                for (int i = 0; i < (n); i++) ws[x[i]]++;
```

```
|...for (int i = 1; i < (lim); i++) ws[i] +=
\rightarrow ws[i - 1];
   for (int i = n; i--;) sa[--ws[x[v[i]]]] =
    v[i]:
  j]) ? p - 1 : p++;
  for (int i = 1; i < (n); i++) rank[sa[i]] =
 for (int i = 0, j; i < n - 1; lcp[rank[i++]]
  // smallest cyclic shift
int cyclic() { return sa[0]; }
 // longest repeated substring
 pair<int,int> lrs() {
  int length = -1, index = -1;
  for (int i = 0; i < lcp.size(); i++) {
  if (lcp[i] > length) {
  length = lcp[i];
    index = sa[i];
  return {index,length};
 // count distinct substrings, excluding empty
 int distincts() {
  int n = sa.size() - 1, r = n - sa[0];
  for (int i = 1; i < lcp.size(); i++)
  r += (n - sa[i]) - lcp[i - 1];
 return r:
 }
// count repeated substrings. excluding empty
 int repeateds() {
  for (int i = 1; i < lcp.size(); i++)
  r += \max(lcp[i] - lcp[i-1], 0);
 return r;
Suffix Tree (Ukkonen's)
struct SuffixTree {
 // n = 2*len+10 or so
enum { N = 50010, ALPHA = 26 };
int toi(char c) { return c - 'a'; }
 t[N] [ALPHA], t[N], t[N], t[N], t[N], t[N], t[N], t[N], t[N], t[N] and t[N] interval Cover (double L, double R,
 string a;
 void ukkadd(int i, int c) { suff:
  if (r[v]<=q) {
  if (t[v][c]==-1) { t[v][c]=m; l[m]=i;
    p[m++]=v; v=s[v]; q=r[v]; goto suff; }
v=t[v][c]; q=1[v];
  if (q==-1 || c==toi(a[q])) q++; else {
   l[m+1]=i; p[m+1]=m; l[m]=l[v]; r[m]=q;
   p[m]=p[v]; t[m][c]=m+1; t[m][toi(a[q])]=v;
   l[v]=q; p[v]=m; t[p[m]][toi(a[l[m]])]=m;
v=s[p[m]]; q=l[m];
   while (q < r[m]) { v = t[v][toi(a[q])];
    q+=r[v]-l[v]; }
   if (q==r[m]) s[m]=v; else s[m]=m+2;
   q=r[v]-(q-r[m]); m+=2; goto suff;
 SuffixTree(string a) : a(a) {
 fill(r,r+N,(int)(a).size());
 memset(s, 0, sizeof s);
memset(t, -1, sizeof t);
fill(t[1],t[1]+ALPHA,0);
s[0]=1;1[0]=1[1]=-1;r[0]=r[1]=p[0]=p[1]=0;
                                                        void genCatalan(int n, int mod) {
  catalan[0] = catalan[1] = 1;
  for (int i = 2; i <= n; i++) {</pre>
                                                         catalan[i] = 0;
  for(int i=0;i<a.size();i++)
```

→ ukkadd(i.toi(a[i])):

```
// Longest Common Substring between 2 strings
 // returns {length, offset from first string}
 pair<int, int> best;
int lcs(int node, int i1, int i2, int olen) {
    if (1[node] <= i1 && i1 < r[node]) return 1;
    if (1[node] <= i2 && i2 < r[node]) return 2;
  int mask=0,
  > len=node?olen+(r[node]-l[node]):0;
  for(int c=0; c<ALPHA; c++) if
   (t[node][c]!=-1)
mask |= lcs(t[node][c], i1, i2, len);
  if (mask==3)
 → best=max(best.{len.r[node]-len});
 static pair<int, int> LCS(string s, string t)
  SuffixTree
 \rightarrow st(s+(char)('z'+1)+t+(char)('z'+2));
 st.lcs(0, s.size(), s.size()+t.size()+1, 0);
return st.best;
String Utilities
void lowercase(string& s) {
 transform(s.begin(), s.end(), s.begin(),
void uppercase(string& s) {
 transform(s.begin(), s.end(), s.begin(),
   ::toupper);
void trim(string &s) {
s.erase(s.begin(),find_if_not(s.begin(),s
     .end(),[](int c){return
    isspace(c);}));
s.erase(find_if_not(s.rbegin(),s.rend(),[](int

    c){return isspace(c);}).base(),s.end());

vector<string> split(string& s, char token) {
    vector<string> v; stringstream ss(s);
    for (string e;getline(ss,e,token);)
        v.push back(e);
    return v:
    Greedy
Interval Cover
//L,R = interval [L,R], in = \{\{l,r\}, index\}
 ^{\prime\prime} does not handle case where L == R

→ vector<pair<pair<double,double>,int>> in) {
    int i = 0; pair < double, int > pos = {L,-1};
    vector<int> a;
sort(begin(in), end(in));
    while (pos.first < R) {
         double cur = pos.first;
while (i < (int)in.size() &&</pre>
    in[i].first.first <= cur)</pre>
    max(pos.{in[i].first.second.in[i].second}).
   1++:
         if (pos.first == cur) return {};
         a.push_back(pos.second);
    return a:
6 Math
Catalan Numbers
ull* catalan = new ull[1000000];
```

for  $(int j = i - 1; j >= 0; j--) {$ 

```
catalan[i] += (catalan[j] * catalan[i-j-1]) |Discrete Logarithm
    % mod;
                                                          int discretelog(int a, int b, int m) {
   if (catalan[i] >= mod)
                                                           11 n = sqrt(m) + 1, an = 1;
    catalan[i] -= mod:
                                                          for (ll i = 0; i < n; ++i)
an = (an * a) % m;
unordered_map<ll, 11> vals;
)// TODO: consider binomial coefficient method
                                                           for (11 q = 0, cur = b; q <= n; q++) {
   vals[cur] = q;
                                                            cur = (cur * a) \% m:
Combinatorics (nCr, nPr)
                                                           for (ll p = 1, cur = 1; p <= n; p++) {
 // can optimize by precomputing factorials, and
                                                            cur = (cur * an) % m;
 \hookrightarrow fact[n]/fact[n-r]
                                                            if (vals.count(cur)) {
ull nPr(ull n, ull r) {
                                                            int ans = n * p - vals[cur];
 ull v = 1:
                                                             return ans;
 for (ull i = n-r+1; i \le n; i++)
 return v;
                                                           return -1;
ull nPr(ull n, ull r, ull m) {
 .11]] v =
                                                          Euler Phi / Totient
 for (ull i = n-r+1; i \le n; i++)
 .v = (v * i) % m;
return v:
                                                          int phi(int n) {
                                                           int^r = n;
                                                          for (int i = 2; i * i <= n; i++) {
   if (n % i == 0) r -= r / i;
   while (n % i == 0) n /= i;
 úll nCr(ull n. ull r) {
 long double v = 1;
 for (ull i = 1: i <= r: i++)
 v = v * (n-r+i) /i;
return (ull) (v + 0.001);
                                                           if (n > 1) r = r / n:
                                                           return r:
}
// requires modulo math
                                                          #define n 100000
// can optimize by precomputing mfac and
                                                         ll phi[n+1];
                                                          void computeTotient() {
 \hookrightarrow minv-mfac
                                                           for (int i=1; i<=n; i++) phi[i] = i;
ull nCr(ull n, ull r, ull m) {
 return mfac(n, m) * minv(mfac(k, m), m) % m *
                                                           for (int p=2; p<=n; p++) {
 \rightarrow minv(mfac(n-k, m), m) % m;
                                                            if (phi[p] == p) {
                                                             phi[p] = p-1;
                                                             for (int i = 2*p; i<=n; i += p) phi[i] =
Multinomials
                                                              (phi[i]/p) * (p-1);
11 multinomial(vector<int>& v) {
11 c = 1, m = v.empty() ? 1 : v[0];
for(int i = 1; i < v.size(); i++)
    for (int j = 0; j < v[i]; j++)
    c = c * ++m / (j+1);</pre>
                                                          Factorials
                                                          // digits in factorial
 return c;
                                                          #define kamenetsky(n) (floor((n * log10(n /
                                                           \rightarrow ME)) + (log10(2 * MPI * n) / 2.0)) + 1)
Chinese Remainder Theorem
                                                          // approximation of factorial
                                                         #define stirling(n) ((n == 1) ? 1 : sqrt(2 *
bool ecrt(ll* r, ll* m, int n, ll& re, ll& mo)
                                                          \hookrightarrow M PI * n) * pow(n / M E, n))
 ill x, y, d; mo = m[0]; re = r[0];
for (int i = 1; i < n; i++) {</pre>
                                                         // natural log of factorial
#define lfactorial(n) (lgamma(n+1))
  if ((r[i] - re) % d != 0) return false;
x = (r[i] - re) / d * x % (m[i] / d);
re += x * mo;
                                                          Prime Factorization
                                                          // do not call directly
                                                          ll pollard rho(ll n. ll s) {
  mo = mo / d * m[i];
                                                           11 x, y;
  re %= mo;
                                                           x = y = rand() \% (n - 1) + 1;
                                                          int head = 1, tail = 2;
while (true) {
 re = (re + mo) \% mo;
 return true;
                                                           x = mult(x, x, n);
x = (x + s) % n;
                                                            if (x == y) return n;
Count Digit Occurences
                                                           ll d = \_gcd(max(x - y, y - x), n);
if (1 < d && d < n) return d:
 /*count(n,d) counts the number of occurences of
                                                            if (++head == tail) y = x, tail <<= 1;
 \rightarrow a digit d in the range [0,n]*/
ll digit_count(ll n, ll d) {
 ll result = 0:
 while (n != 0)
                                                          // call for prime factors
  result += ((n\%10) == d ? 1 : 0);
                                                          void factorize(ll n. vector<ll> &divisor) {
                                                           if (n == 1) return;
  n /= 10;
                                                           if (isPrime(n)) divisor.push back(n):
 return result:
```

- 1) + 1);
factorize(n / d, divisor);

factorize(d, divisor);

fl count(ll n, ll d) {
 if (n < 10) return (d > 0 && n >= d);
 if ((n % 10) != 9) return digit\_count(n, d) +

return 10\*count(n/10, d) + (n/10) + (d > 0);

 $\hookrightarrow$  count(n-1, d);

```
// generate 0 <= a/b <= 1 ordered, b <= n
                                                     farey(4) = 0/1 1/4 1/3 1/2 2/3 3/4 1/1
                                                  // length is sum of phi(i) for i = 1 to n
                                                 vector<pair<int, int>> farey(int n) {
                                                  int h = 0, k = 1, x = 1, y = 0, r;
                                                   vector<pair<int, int>> v;
                                                   v.push_back({h, k});
                                                   r = (n-y)/k;
                                                   v += r*k: x' += r*h:
                                                  x = -x; y = -y;

while (k > 1);
                                                  v.push_back({1, 1});
                                                  return v;
                                                 Fast Fourier Transform
                                                 #define cd complex<double>
                                                 const double PI = acos(-1):
                                                 void fft(vector<cd>& a, bool invert) {
                                                   int n = a.size();
                                                  for (int i = 1, j = 0; i < n; i++) {
                                                   int bit = n >>
                                                   for (; j & bit; bit >>= 1) j ^= bit;
                                                   .j ^= bit:
                                                   if (i < j) swap(a[i], a[j]);
                                                  for (int len = 2; len <= n; len <<= 1) {
    double ang = 2 * PI / len * (invert ? -1 :
                                                  . 1)
                                                   cd wlen(cos(ang), sin(ang));
                                                   for (int i = 0; i < n; i += len) {
                                                    for (int j = 0; j < len / 2; j++) {
    cd u = a[i+j], v = a[i+j+len/2] * w;
                                                      a[i+j] = u + v;
                                                     a[i+j+len/2] = u - v;
                                                     w = wlen;
                                                  if (invert)
                                                   for (auto& x : a)
                                                    x /= n:
                                                 vector<int> fftmult(vector<int> const& a.
                                                  → vector<int> const& b) {
                                                  vector<cd> fa(a.begin(), a.end()),
                                                  → fb(b.begin(), b.end());
                                                  int n = 1 << (32 - __builtin_clz(a.size() +</pre>
                                                   → b.size() - 1));
                                                  fa.resize(n); fb.resize(n);
                                                  fft(fa, false); fft(fb, false)
                                                  for (int i = 0; i < n; i++) fa[i] *= fb[i];
                                                   fft(fa, true);
                                                   vector<int> toret(n);
                                                  for (int i = 0; i < n; i++) toret[i] =
                                                   → round(fa[i].real());
                                                  return toret:
                                                 Greatest Common Denominator
                                                 11 egcd(11 a, 11 b, 11& x, 11& y) {
  if (b == 0) { x = 1; y = 0; return a; }
                                                  11 gcd = egcd(b, a \% b, x, y);
                                                  x = a / b * y;
                                                  swap(x, y);
                                                  return gcd;
                                                 Josephus Problem
                                                  // 0-indexed, arbitrary k
while (d >= n) d = pollard_rho(n, rand() % (n int josephus(int n, int k) {
                                                  if (n == 1) return 0;
if (k == 1) return n-1;
                                                  if (k > n) return (josephus(n-1,k)+k)%n;
                                                  int res = josephus(n-n/k,k)-n\%k;
                                                  return res + ((res<0)?n:res/(k-1));
```

Farey Fractions

```
// fast case if k=2, traditional josephus
int josephus(int n) {
return 2*(n-(1<<(32-builtin clz(n)-1)));
Least Common Multiple
#define lcm(a,b) ((a*b)/qcd(a,b))
Modulo Operations
#define MOD 1000000000
#define madd(a,b,m) (a+b-((a+b-m>=0)?m:0))
#define mult(a,b,m) ((ull)a*b%m)
#define msub(a,b,m) (a-b+((a < b)?m:0))
ll mpow(ll b, ll e, ll m) {
11 x = 1:
while (e > 0) {
   if (e % 2) x = (x * b) % m;
   b = (b * b) % m;
 e /= 2;
return x % m:
ull mfac(ull n, ull m) {
ull f = 1;
for (int i = n; i > 1; i--)
f = (f * i) % m;
return f:
// if m is not quaranteed to be prime
ll minv(ll b, ll m) {
11 \times (0, y) = 0;
if (egcd(b, m, x, y) != 1) return -1;
return (x % m + m) % m;
11 mdiv_compmod(int a, int b, int m) {
if (\underline{gcd(b, m)} != 1) return -1;
return mult(a, minv(b, m), m);
\frac{1}{1} if m is prime (like 10^9+7)
11 mdiv_primemod (int a, int b, int m) {
return mult(a, mpow(b, m-2, m), m);
// tonelli shanks = sqrt(n) % m, m is prime
ll legendre(ll a, ll m){
if (a % m==0) return 0;
if (m == 2) return 1;
return mpow(a, (m-1)/2, m);
11 msqrt(11 n, 11 m) {
ll s = __builtin_ctzll(m-1), q = (m-111)>>s,
\Rightarrow z = rand()%(m-1)+1:
if (m == 2) return 1;
if (s == 1) return mpow(n,(m+1)/411,m);
 while (legendre(z,m)!=m-1) z = rand()\%(m-1)+1;
11 c = mpow(z,q,m), r = mpow(n,(q+1)/2,m), t
 \rightarrow = mpow(n,q,m), M = s;
 while (t != 1){
 11 b = c:
 for (int j = 0; j < M-i-1; j++) b = (b * b) \%
r = r * b % m; c = b * b % m; t = t * c % m;
\rightarrow M = i;
return r:
Modulo Tetration
11 tetraloop(11 a, 11 b, 11 m) {
  if(b == 0 | | a == 1) return 1;
11 w = tetraloop(a,b-1,phi(m)), r = 1;
for (;w;w/=2) {
 if (w&1) ·
  r *= a; if (r >= m) r -= (r/m-1)*m;
 \bar{a} *= a; if (a >= m) a -= (a/m-1)*m;
 return r;
int tetration(int a, int b, int m) {
```

```
if (a == 0 | | m == 1) return ((b+1)&1)%m;
                                                                                                  |// 1 = even number prime factors, -1 = odd
 return tetraloop(a,b,m) % m;
                                                                                                   short mu[MAXN] = \{0,1\};
                                                                                                  void mobius(){
  for (int i = 1; i < MAXN; i++)
    if (mu[i])</pre>
Matrix
template<typename T>
                                                                                                        for (int' j = i + i; j < MAXN; j += i)
struct Mat : public Vec<2, T> {
                                                                                                          mu[j] -= mu[i];
 int w, h;
 Mat(int x, int y) : Vec<2, T>(x, y), w(x),
                                                                                                   Nimber Arithmetic
 \hookrightarrow h(y) {}
                                                                                                  #define nimdd(a,b) ((a) ^(b))
ull nimMul(ull a, ull b, int i=6) {
   static const ull M[]={INT_MIN>>32,
  static Mat<T> identity(int n) { Mat<T> m(n,n);
        for (int i=0;i<n;i++) m[i][i] = 1; return

ighthat is in the second in 
                                                                                                          M[0]^{(M[0] << 16)}, M[1]^{(\overline{M}[1] << 8)},
  Mat<\hat{T}>\& operator+=(const Mat<T>\& m) {
  for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
...(*this)[i][j] += m[i][j];
                                                                                                          M[2]^{(M[2] << 4)}, M[3]^{(M[3] << 2)},
                                                                                                         M[4]^(M[4]<<1);
                                                                                                      if (i-- == 0) return a&b;
int k=1<<i:
    return *this:
                                                                                                      ull s=nimMul(a,b,i), m=M[5-i],
  Mat<T>& operator-=(const Mat<T>& m) {
                                                                                                          t=nimMul(((a^(a>>k))&m)|(s&~m).
   for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
                                                                                                           ((b^{(b)}) \& m) | (m \& (\sim m >> 1)) << k, i):
                                                                                                      return ((s^t)&m)<<k|((s^(t>>k))&m);
       (*this)[i][j] -= m[i][j];
    return *this;
                                                                                                   Permutation
                                                                                                  // c = array size, n = nth perm, return index
vector<int> gen_permutation(int c, int n) {
  Mat<T> operator*(const Mat<T>& m) {
    Mat<T>z(w,m.h);
    for (int i = 0; i < w; i++)
                                                                                                    vector<int> idx(c), per(c), fac(c); int i;
  ...for (int j = 0; j < h; j++)
....for (int k = 0; k < m.h; k++)
....z[i][k] += (*this)[i][j] * m[j][k];
                                                                                                    for (i = 0; i < c; i++) idx[i] = i;
for (i = 1; i <= c; i++) fac[i-1] = n%i, n/=i;
                                                                                                    for (i = c - 1; i >= 0; i--)
per[c-i-1] = idx[fac[i]],
                                                                                                      idx.erase(idx.begin() + fac[i]);
  Mat<T> operator+(const Mat<T>& m) { Mat<T>
                                                                                                    return per;

    a=*this; return a+=m; }

 Mat<T> operator-(const Mat<T>& m) { Mat<T>
                                                                                                    // get what nth permutation of vector
                                                                                                   int get permutation(vector<int>& v) {

    a=*this; return a-=m; }

                                                                                                   int use = 0, i = 1, r = 0;
for (int e: v) {
   r = r * i++ + __builtin_popcount(use &
 Mat<T>& operator*=(const Mat<T>& m) { return

→ *this = (*this)*m; }

  Mat<T> power(int n) {
                                                                                                     \rightarrow -(1<<e));
   .Mat<T> a = Mat<T>::identity(w), m=*this;
                                                                                                      use |= 1 << e;
    for (;n;n/=2,m*=m) if (n\&1) a *= m:
   return a;
                                                                                                    return r;
                                                                                                   Permutation (string/multiset)
Matrix Exponentiation
                                                                                                   string freq2str(vector<int>& v) {
 // F(n) = c[0]*F(n-1) + c[1]*F(n-2) + ...
// b is the base cases of same length c
                                                                                                    string s;
                                                                                                   for (int i = 0; i < v.size(); i++)

for (int j = 0; j < v[i]; j++)

s += (char)(i + 'A');
11 matrix exponentiation(11 n, vector<11> c,
 → vector<11> b) {
   if (nth < b.size()) return b[nth-1];
   Mat<11> a(c.size(), c.size()); ll s = 0;
   for (int i = 0; i < c.size(); i++) a[i][0] =</pre>
                                                                                                   // nth perm of multiset, n is 0-indexed

    c[i];

                                                                                                  string gen_permutation(string s, ll n) {
 for (int i = 0; i < c.size() - 1; i++)
                                                                                                    vector<int> freq(26, 0);
 \rightarrow a[i][i+1] = 1;
a = a.power(nth - c.size());
                                                                                                    for (auto e : s) freq[e - 'A']++;
                                                                                                    for (int i = 0; i < 26; i++) if (freq[i] > 0)
 for (int i = 0; i < c.size(); i++)
s += a[i][0] * b[i];
                                                                                                      freq[i]--; ll v = multinomial(freq);
 return s;
                                                                                                      if (n < v) return (char)(i+'A') +
                                                                                                         gen_permutation(freq2str(freq), n);
Matrix Subarray Sums
                                                                                                      freq[\overline{i}]++; n -= v;
template < class T> struct MatrixSum {
                                                                                                    return "";
  MatrixSum(Vec<2, T>& v) {
   p = Vec<2,T>(v.size()+1, v[0].size()+1);
                                                                                                   Miller-Rabin Primality Test
 for (int i = 0; i < v.size(); i++)
for (int j = 0; j < v[0].size(); j++)
p[i+1][j+1] = v[i][j] + p[i][j+1] +</pre>
                                                                                                   // Miller-Rabin primality test - O(10 log^3 n)
                                                                                                 // Multer-Rabin primality test -
bool isPrime(ull n) {
  if (n < 2) return false;
  if (n == 2) return true;
  if (n % 2 == 0) return false;
  ull s = n - 1;
  while (s % 2 == 0) s /= 2;
  for (int i = 0; i < 10; i++) {
    ull temp = s;
    ull a = rand() % (n - 1) + 1;
    ull mod = mnow(a temp n);
}</pre>

    p[i+1][j] - p[i][j];

 f sum(int u, int l, int d, int r) {
    return p[d][r] - p[d][l] - p[u][r] + p[u][l];
Mobius Function
                                                                                                      ull mod = mpow(a, temp, n);
const int MAXN = 10000000;
// mu[n] = 0 iff n has no square factors
                                                                                                      while (temp!=n-1\&\&mod!=1\&\&mod!=n-1) {
```

```
. mod = mult(mod, mod, n);
    temp *= 2:
   if (mod!=n-1&&temp%2==0) return false:
 return true:
Sieve of Eratosthenes
bitset<100000001> sieve;
// generate sieve - O(n log n)
void genSieve(int n) {
    sieve[0] = sieve[1] = 1;
    for (ull i = 3; i * i < n; i += 2)
    if (!sieve[i])</pre>
   for (ull j = i * 3; j <= n; j += i * 2)
...sieve[j] = 1;
// query sieve after it's generated - O(1)
bool querySieve(int n) {
 return n == 2 || (n % 2 != 0 && !sieve[n]);
Compile-time Prime Sieve
const int MAXN = 100000;
template<int N>
struct Sieve {
  bool sieve[N];
  constexpr Sieve() : sieve() {
  sieve[0] = sieve[1] = 1;
  for (int i = 2; i * i < N; i++)
if (!sieve[i])
     for (int j = i * 2; j < N; j += i)
...sieve[j] = 1;
bool isPrime(int n) {
   static constexpr Sieve<MAXN> s;
 return !s.sieve[n]:
Simpson's / Approximate Integrals
// integrate f from a to b, k iterations
// threefact | from a to , k theratton

// error <= (b-a)/18.0 * M * ((b-a)/2k)^4

// where M = max(abs(f```(x))) for x in [a,b]

// "f" is a function "double func(double x)"
double Simpsons (double a. double b. int k.
 \rightarrow double (*f)(double)) {
double dx = (b-a)/(2.0*k), t = 0;
 for (int i = 0; i < k; i++)

t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
    (*f)(a+(2*i+1)*dx);
 return (t + (*f)(b)) * (b-a) / 6.0 / k;
Common Equations Solvers
// ax^2 + bx + c = 0, find x
vector<double> solveEq(double a, double b,
  → double c) {
vector<double> r;
  double z = b * b' - 4 * a * c;
  if (z == 0)
  r.push_back(-b/(2*a));
  else if (z > 0) {
  r.push back((sgrt(z)-b)/(2*a)):
  r.push_back((sqrt(z)+b)/(2*a));
 return r;
// ax^3 + bx^2 + cx + d = 0, find x
vector<double> solveEq(double a. double b.
 → double c, double d) {
vector<double> res:
 long double a1 = b/a, a2 = c/a, a3 = d/a;
 long double q = (a1*a1 - 3*a2)/9.0, sq =
 \rightarrow -2*sqrt(q);
 long double r = (2*a1*a1*a1 - 9*a1*a2 +
 \rightarrow 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
 if (z <= 0) {
  theta = a\cos(r/sqrt(q*q*q));
```

```
res.push_back(sq*cos(theta/3.0) - a1/3.0);
  res.push back(sq*cos((theta+2.0*PI)/3.0) -
  \rightarrow a1/3.0):
 res.push back(sq*cos((theta+4.0*PI)/3.0) -
    a1/3.0);
 élse {
 res.push_back(pow(sqrt(z)+fabs(r), 1/3.0));
 res[0] = (res[0] + q / res[0]) *

    ((r<0)?1:-1) - a1 / 3.0;
}
return res;</pre>
\frac{1}{2} linear diophantine equation ax + by = c,
    find x and y
// infinite solutions of form x+k*b/g, y-k*a/g
bool solveEq(ll a, ll b, ll c, ll &x, ll &y, ll
 g = \tilde{e}gcd(abs(a), abs(b), x, y);
if (c % g) return false;
x *= c / g * ((a < 0) ? -1 : 1);
y *= c / g * ((b < 0) ? -1 : 1);
return true;
^{\prime\prime}/ m = # equations, n = # variables, a[m][n+1]
\rightarrow = coefficient matrix
// a[i][0]x + a[i][1]y + ... + a[i][n]z =
\hookrightarrow a[i][n+1]
// find a solution of some kind to linear

→ equation

const double eps = 1e-7;
bool zero(double a) { return (a < eps) && (a >
\rightarrow -eps); }
vector <double> solveEq(double **a, int m, int
int cur = 0;

for (int i = 0; i < n; i++) {

for (int j = cur; j < m; j++) {
  if (!zero(a[j][i])) {
    if (j!= cur) swap(a[j], a[cur]);
for (int sat = 0; sat < m; sat++) {
   if (sat == cur) continue;
   double num = a[sat][i] / a[cur][i];</pre>
      for (int sot = 0; sot <= n; sot++)
   a[sat][sot] -= a[cur][sot] * num;
    .}
.cur++:
    break
 for (int j = cur; j < m; j++)
 if (!zero(a[j][n])) return vector <double > ();
vector<double> ans(n,0);
for (int i = 0, sat = 0; i < n; i++)
if (sat < m && !zero(a[sat][i]))
ans[i] = a[sat][n] / a[sat++][i];
return ans;
^{\prime}// solve A[n][n] * x[n] = b[n] linear equation
// rank < n is multiple solutions, -1 is no
→ solutions

// `alls` is whether to find all solutions, or
\hookrightarrow any
const double eps = 1e-12;
int solveEq(Vec<2, double>& A, Vec<1, double>&
\rightarrow b, Vec<1, double>& x, bool alls=false) {
int n = A.size(), m = x.size(), rank = 0, br,
\rightarrow bc;
vector<int> col(m); iota(begin(col), end(col),
for(int i = 0; i < n; i++) {
 double v, bv = 0;
for(int r = i; r < n; r++)
 for(int c = i; c < n; c++)
    if ((v = fabs(A[r][c])) > bv)
    br = r, bc = c, bv = v;
    if (bv <= eps) {
   for(int i = i: i < n: i++)
```

```
. if (fabs(b[j]) > eps)
    .return -1;
 swap(A[i], A[br]);
 swap(b[i], b[br]);
 swap(col[i], col[bc]);
 for(int j = 0; j < n; j++)
swap(A[j][i], A[j][bc]);
 bv = 1.0 / A[i][i];
for(int j = (alls)?0:i+1; j < n; j++) {
  if (j != i) {
   double fac = A[j][i] * bv;
   b[j] = fac * b[i];
  for(int k = i+1; k < m; k++)
A[j][k] -= fac*A[i][k];
 rank++;
}
if (alls) for (int i = 0; i < m; i++) x[i] =</pre>
→ -DBL MAX;
for (int i = rank; i--;) {
  bool isGood = true;
 if (aļls)
 for (int j = rank; isGood && j < m; j++)
...if (fabs(A[i][j]) > eps)
 isGood = false;
b[i] /= A[i][i]:
 if (isGood) x[col[i]] = b[i];
 if (!alls)
for(int j = 0; j < i; j++)
   b[j] = A[j][i] * b[i];
return rank:
```

### Graycode Conversions

```
ull graycode2ull(ull n) {
ull i = 0;
for (; n; n = n >> 1) i ^= n;
 return i;
ull ull2graycode(ull n) {
  return n ^ (n >> 1);
```

### Date Utilities

```
// handles -4799-01-01 to 1465001-12-31
int date2int(int y, int m, int d){
return 1461*(y+4800+(m-14)/12)/4+367*(m-2-(m-14)/12)
   -14)/12*12)/12-3*((v+4900+(m-14)/12)/100)
   /4+d-32075;
pair<int,pair<int,int>> int2date(int x){
 int n,i,j;
 x+=68569
 n=4*x/146097;
x-=(146097*n+3)/4;
i=(4000*(x+1))/1461001;
 x=1461*i/4-31;
 i=80*x/2447:
 return \{100*(n-49)+i+j/11, \{j+2-12*(j/11), \}
\rightarrow x-2447*j/80}};
int dayOfWeek(int y, int m, int d){ //0=sunday
static int cal[]={0,3,2,5,0,3,5,1,4,6,2,4};
return (y+y/4-y/100+y/400+cal[m-1]+d)\%7;
```

### Unix/Epoch Time

```
// O-indexed month/time, 1-indexed day
// minimum 1970, 0, 1, 0, 0, 0
ull toEpoch(int year, int month, int day, int
→ hour, int minute, int second) {
struct tm t; time_t epoch;

t.tm_year = year - 1900; t.tm_mon = month;

t.tm_mday = day; t.tm_hour = hour;
t.tm min = minute; t.tm sec = second;
```

```
t.tm_isdst = 0; // 1 = daylights savings
epoch = mktime(&t);
return (ull)epoch;
vector<int> toDate(ull epoch) {
time_t e=epoch; struct tm t=*localtime(&e);
return {t.tm year+1900,t.tm mon,t.tm mday,t
   .tm_hour,t.tm_min,t.tm_sec};
int getWeekday(ull epoch) {
time t e=epoch; struct tm t=*localtime(&e);
return t.tm_wday; // 0-6, 0 = sunday
int getDayofYear(ull epoch) {
time_t e=epoch; struct tm t=*localtime(&e);
return t.tm_yday; // 0-365
const int months[] =
\leftrightarrow {31,28,31,30,31,30,31,30,31,30,31};
bool validDate(int year, int month, int day) {
    bool leap = !(year%(year%25?4:16));
    if (month >= 12) return false:
   return day <= months[month] + (leap &&
   month == 1):
```

### Theorems and Formulae

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set of size n. !0 = 1, !1 = 0, !n = (n+1)(!(n-1))|1)+!(n-2),  $!n = n! \sum_{i=0}^{n} \frac{(-1)^i}{i!}$ ,  $!n = [\frac{n!}{e}]$ 

In a partially ordered set, a chain is a subset of elements that are all comparable to eachother. An antichain is a subset where no two are comparable.

Dilworth's theorem states the size of a maximal antichain equals the size of a minimal chain cover of a partially ordered set S. The width of S is the maximum size of an antichain in S, which is equal to the minimum number of chains needed to cover S, or the minimum number of chains such that all elements are in at least one chain.

Rosser's Theorem states the nth prime Floyd Warshall

Lagrange's Four Square Theorem states every natural number is the sum of the squares of four non-negative integers. This is a special case of the Fermat Polygonal Number **Theorem** where every positive integer is  $a \mapsto m[i][j] = -inf$ ; sum of at most n s-gonal numbers. The nths-gonal number  $P(s,n) = (s-2)\frac{n(n-1)}{2} + n$ 

## 7 Graphs

```
struct edge {
int u,v,w;
edge (int u,int v,int w) : u(u),v(v),w(w) {}
edge (): u(0), v(0), w(0) {}
```

```
\rightarrow &e2) { return e1.w < e2.w: }
bool operator > (const edge &e1, const edge
\rightarrow &e2) { return e1.w > e2.w: }
struct subset { int p, rank; };
Eulerian Path
#define edge_list vector<edge>
#define adj_sets vector<set<int>>
struct EulerPathGraph {
adj_sets graph; // actually indexes incident
 → edaes
 edge_list edges; int n; vector<int> indeg;
 EulerPathGraph(int n): n(n) {
 indeg = *(new vector<int>(n,0));
 graph = *(new adj_sets(n, set<int>()));
 void add edge(int u, int v) {
 graph[u].insert(edges.size());
 indeg[v]++;
 edges.push back(edge(u.v.0)):
 bool eulerian_path(vector<int> &circuit) {
 if(edges.size()==0) return false;
  stack<int> st;
 int a[] = {-1};
for(int v=0; v<n; v++) {
   if(indeg[v]!=graph[v].size()) {</pre>
    bool b = indeg[v] > graph[v].size();
    if (abs(((int)indeg[v])-((int)graph[v]
    .size())) > 1) return
    if (a[b] != -1) return false;
   a[b] = v;
  int s = (a[0]!=-1 && a[1]!=-1 ? a[0] :
    (a[0]=-1 & a a[1]=-1 ? edges[0].u : -1));
  if(s==-1) return false;
 while(!st.empty() || !graph[s].empty()) {
   if (graph[s].empty()) {
    circuit.push_back(s); s = st.top();
    st.pop(); }
    int w = edges[*graph[s].begin()].v;
graph[s].erase(graph[s].begin());
    st.push(s); s = w;
 circuit.push_back(s);
 return circuit.size()-1==edges.size();
```

|bool operator < (const edge &e1, const edge

```
number is greater than n*ln(n) for n>1. const l1 inf = 1LL << 62; Nicomachi's Theorem states 1^3+2^3+\ldots+\frac{\#define\ FOR(i,n)\ for\ (int\ i=0;\ i< n;\ i++)}{void\ floydWarshall(Vec<2,\ 11>\&\ m)} {
n^3 = (1+2+...+n)^2 and is equivalent to \inf_{\substack{i=1 \ \text{FOR(i,n) } \text{FOR(i,n) } \text{FOR(i,n) } \text{FOR(i,n) } \text{FOR(i,n) } \text{if (m[i][k] } != \inf_{\substack{i=1 \ \text{FOR(k,n) } \text{FOR(i,n) } \text{FOR(i,n) } \text{ if (m[i][k] } != \inf_{\substack{i=1 \ \text{FOR(k,n) } \text{FOR(i,n) } \text{FOR
                                                                                                                                                                                                                                                                                                                                                                                                                     \rightarrow && m[k][j] != inf)
                                                                                                                                                                                                                                                                                                                                                                                                                           auto newDist = max(m[i][k] + m[k][j], -inf);
                                                                                                                                                                                                                                                                                                                                                                                                                        m[i][j] = min(m[i][j], newDist);
                                                                                                                                                                                                                                                                                                                                                                                                                      FOR(k,n) if (m[k][k] < 0) FOR(i,n) FOR(j,n)
                                                                                                                                                                                                                                                                                                                                                                                                                  if (m[i][k] != inf && m[k][j] != inf)
```

## Minimum Spanning Tree

```
// returns vector of edges in the mst
// graph[i] = vector of edges incident to
→ vertex i
// places total weight of the mst in &total
// if returned vector has size != n-1, there is
vector<edge> mst(vector<vector<edge>> graph,

→ 11 &total) {
```

```
total = 0:
 priority_queue<edge, vector<edge>,

→ greater<edge>> pq;

 vector<edge> MST;
 bitset<20001> marked: // change size as needed
 marked[0] = 1;
 for (edge ep : graph[0]) pq.push(ep);
 while (MST.size()!=graph.size()-1 &&
 → pq.size()!=0) {
  edge e = pq.top(); pq.pop();
int u = e.u, v = e.v, w = e.w;
if(marked[u] && marked[v]) continue;
  else if (marked[u]) swap(u, v);
  for(edge ep : graph[u]) pq.push(ep);
  marked[u] = 1
  MST.push_back(e);
  total += e.w:
 return MST;
Union Find
int uf find(subset* s, int i) {
  if (s[i].p != i) s[i].p = uf_find(s, s[i].p);
 return s[i].p;
void uf_union(subset* s, int x, int y) {
 int xp = uf_find(s, x), yp = uf_find(s, y);
if (s[xp].rank > s[yp].rank) s[yp].p = xp;
 else if (s[xp].rank < s[yp].rank) s[xp].p =
 else { s[yp].p = xp; s[xp].rank++; }
2D Grid Shortcut
#define inbound(x,n) (0<=x\leq x< n)
#define fordir(x, y, n, m) for(auto[dx, dy]:dir)if
\hookrightarrow (inbound(x+dx,n)&Ginbound(y+dy,m))
const pair<int,int> dir[] =
\rightarrow {{1,0},{0,1},{-1,0},{0,-1}};
   2D Geometry
#define point complex<double>
#define EPS 0.0000001
#define sq(a) ((a)*(a))
#define car{b}(a) ((a)*(a)*(a))
double dot(point a, point b) { return

    real(coni(a)*b): }

double cross(point a, point b) { return
 \rightarrow imag(conj(a)*b); }
struct line { point a, b; };
struct circle { point c; double r; };
struct segment { point a, b; };
struct triangle { point a, b, c; };
struct rectangle { point tl, br; };
struct convex_polygon {
 vector<point> points;
 convex_polygon(vector<point> points) :
   points(points) {}
 convex_polygon(triangle a) {
  points.push_back(a.a); points.push_back(a.b);
   points.push back(a.c):
 convex_polygon(rectangle a) {
  points.push_back(a.tl);
    points.push back({real(a.tl).
    imag(a.br)});
  points.push_back(a.br);
    points.push_back({real(a.br),
    imag(a.tl)}):
struct polygon {
 vector <point > points;
 polygon(vector<point> points) :

→ points(points) {}
```

```
point c(0, 0);
 polygon(triangle a) {
 points.push back(a.a); points.push back(a.b);
                                                     for (int i = n - 1, j = 0; j < n; i = j++) {
    points.push back(a.c);
                                                      double v = cross(a.points[i], a.points[j]) /
                                                      → 2;
area += v:
 polygon(rectangle a) {
                                                      c += (a.points[i] + a.points[j]) * (v / 3);
 points.push_back(a.tl);
    points.push_back({real(a.tl),
                                                     c /= area;
    imag(a.br)});
                                                     return {area, c};
  points.push back(a.br);
    points.push_back({real(a.br),
                                                    Intersection
                                                     // -1 coincide, 0 parallel, 1 intersection
    imag(a.tl)});
                                                     int intersection(line a, line b, point& p) {
 polygon(convex_polygon a) {
                                                     if (abs(cross(a.b - a.a, b.b - b.a)) > EPS) {
   p = cross(b.a - a.a, b.b - a.b) / cross(a.b
  for (point v : a.points)
   points.push_back(v);
                                                     \rightarrow - a.a, b.b - b.a) * (b - a) + a;
return 1;
                                                     if (abs(cross(a.b - a.a, a.b - b.a)) > EPS)
   triangle methods
                                                     → return 0;
double area_heron(double a, double b, double
                                                     return -1;
\stackrel{\hookrightarrow}{} c) { if (a < b) swap(a, b);
                                                     // area of intersection
 if (a < c) swap(a, c);
                                                    double intersection(circle a, circle b) {
 if (b < c) swap(b, c);
                                                     double d = abs(a.c - b.c);
if (a > b + c) return -1;
return sqrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)
                                                        (d <= b.r - a.r) return area(a);
(d <= a.r - b.r) return area(b);
                                                     if (d \ge a.r + b.r) return 0;
/16.0);
                                                     double alpha = acos((sq(a.r)' + sq(d) - cos(sq(a.r)')))
                                                        sq(b.r)) / (2 * a.r * d));
double lengthsq(segment a) { return
                                                     double beta = acos((sq(b.r) + sq(d) - sq(a.r))
    sq(real(a.a) - real(a.b)) + sq(imag(a.a)
                                                        /(2 * b.r * d)):
    imag(a.b)); }
                                                     return sq(a.r) * (alpha - 0.5 * sin(2 *
double Tength(segment a) { return
                                                        alpha)) + sq(b.r) * (beta - 0.5 * sin(2 *
    sqrt(lengthsq(a)); }
                                                        beta)):
// circle methods
double circumference(circle a) { return 2 * a.r // -1 outside, 0 inside, 1 tangent, 2

    * M_PI; }

    intersection
int intersection(circle a, circle b,
double area(circle a) { return sq(a.r) * M_PI;
                                                        vector<point>& inter) {
→ }
// rectangle methods
                                                     double d2 = norm(b.c - a.c), rS = a.r + b.r.
                                                     \rightarrow rD = a.r - b.r;
if (d2 > sq(rS)) return -1;
double width(rectangle a) { return

→ abs(real(a.br) - real(a.tl)); }

                                                     if (d2 < sq(rD)) return 0;
double height (rectangle a) { return
                                                     double ca = 0.5 * (1 + rS * rD / d2);

→ abs(imag(a.br) - real(a.tl)); }

                                                     point z = point(ca, sqrt(sq(a.r) / d2 -
double diagonal(rectangle a) { return
                                                        sq(ca));

    sqrt(sq(width(a)) + sq(height(a))); }

                                                     inter.push back(a.c + (b.c - a.c) * z):
double area(rectangle a) { return width(a) *
                                                     if (abs(imag(z)) > EPS) inter.push_back(a.c +
→ height(a): }
                                                         (b.c - a.c) * conj(z));
double perimeter(rectangle a) { return 2 *
                                                     return inter.size():
// check if `a` fit's inside `b
                                                     // points of intersection
// swap equalities to exclude tight fits
                                                     vector<point> intersection(line a, circle c) {
                                                     vector<point> inter;
bool doesFitInside(rectangle a, rectangle b) {
                                                     c.c -= a.a;
a.b -= a.a;
 int x = width(a), w = width(b), v = height(a)

→ h = height(b);

                                                     point m = a.b * real(c.c / a.b);
 if (x > y) swap(x, y);
                                                     double d2 = norm(m - c.c);
 if (w > h) swap(w, h);
                                                     if (d2 > sq(c.r)) return 0;
 if (w < x) return false;
                                                     double l = sqrt((sq(c.r) - d2) / norm(a.b));
 if (y <= h) return true;
                                                     inter.push back(a.a + m + 1 * a.b);
 double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
                                                     if (abs(1) > EPS) inter.push_back(a.a + m - 1
 return sq(a) \le sq(b) + sq(c);

    * a.b);
return inter;
}
// polygon methods
                                                    | Same name of Spine Dates, near per | int snook (int a) { | int b = a & -a, c = a + b; | double intersection(rectangle a, rectangle b) { | return c | ((a ^ c) >> 2) / b; | }
// negative area = CCW, positive = CW
double area(polygon a) {
 double area = 0.0; int n = a.points.size();
                                                     double x1 = max(real(a.tl), real(b.tl)), y1 =
 for (int i = 0, j = 1; i < n; i++, j = (j +
                                                        max(imag(a.tl), imag(b.tl));
                                                     double x2 = min(real(a.br), real(b.br)), y2 =
 area += (real(a.points[j]-a.points[i]))*
                                                        min(imag(a.br), imag(b.br));

    (imag(a.points[j]+a.points[i]));
                                                     return (x2 <= x1 | | y2 <= y1) ? 0 :
return area / 2.0;
                                                        (x2-x1)*(y2-y1);
// get both unsigned area and centroid
                                                    Convex Hull
pair<double, point> area centroid(polygon a) {
                                                    bool cmp(point a, point b) {
  if (abs(real(a) - real(b)) > EPS) return
 int n = a.points.size();
 double area = 0:

    real(a) < real(b):
</pre>
```

```
if (abs(imag(a) - imag(b)) > EPS) return
\rightarrow imag(a) < imag(b);
return false:
convex_polygon convexhull(polygon a) {
sort(a.points.begin(), a.points.end(), cmp);
 vector<point> lower, upper;
 for (int i = 0; i < a.points.size(); i++) {
 while (lower.size() >= 2 &&
   cross(lower.back() - lower[lower.size()
   2], a.points[i] - lower.back()) < EPS)
  lower.pop_back();
 while (upper.size() >= 2 &&
   cross(upper.back() - upper[upper.size()
   2], a.points[i] - upper.back()) > -EPS)
  upper.pop_back();
 lower.push_back(a.points[i]);
 upper.push_back(a.points[i]);
 lower.insert(lower.end(), upper.rbegin() + 1,
   upper.rend());
return convex_polygon(lower);
    3D Geometry
struct point3d {
```

```
double x, y, z;
 point3d operator+(point3d a) const { return
 \rightarrow {x+a.x, y+a.y, z+a.z}; }
 point3d operator*(double a) const { return
 \rightarrow {x*a, y*a, z*a}; }
 point3d operator-() const { return {-x, -y,
 → -z}: }
 point3d operator-(point3d a) const { return
\rightarrow *this + -a; }
 point3d operator/(double a) const { return
 ** *this * (1/a); }
double norm() { return x*x + y*y + z*z; }
 double abs() { return sqrt(norm()); }
 point3d normalize() { return *this /

    this->abs(); }

double dot(point3d a, point3d b) { return
\rightarrow a.x*b.x + a.y*b.y + a.z*b.z; }
point3d cross(point3d a, point3d b) { return
   \{a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,
   a.x*b.v - a.v*b.x}: }
struct line3d { point3d a, b; };
struct plane { double a, b, c, d; } // a*x +
\Rightarrow b*y + c*z + d = 0
struct sphere { point3d c; double r; };
#define sq(a) ((a)*(a))
#define c\bar{b}(a) ((a)*(a)*(a))
double surface(circle a) { return 4 * sq(a.r)
\hookrightarrow M_PI; }
double volume(circle a) { return 4.0/3.0 *
\hookrightarrow cb(a.r) * M PI; }
```

## 10 Optimization

```
Snoob
```

```
// SameNumberOfOneBits, next permutation
   / example usage
// example usage
int main() {
   char l1[] = {'1', '2', '3', '4', '
   char l2[] = {'a', 'b', 'c', 'd'};
   int d1 = 5, d2 = 4;
   // prints 12345abcd, 1234a5bcd, ...
  int min = (1 < < d1) - 1, max = min < < d2:
  for (int i = min; i <= max; i = snoob(i)) {
  int p1 = 0, p2 = 0, v = i;

while (p1 < d1 || p2 < d2) {

    cout << ((v & 1) ? l1[p1++] : l2[p2++]);
```

```
v /= 2;
 .cout << '\n':
\mathbf{Powers}
bool isPowerOf2(11 a) {
return a > 0 \&\& !(a \& a-1);
bool isPowerOf3(ll a) {
   return a>0&&!(12157665459056928801u11%a):
bool isPower(ll a, ll b)
double x = log(a) / log(b);
return abs(x-round(x)) < 0.00000000001;
```

### 11 Python

### Recursion Limit Removal (Basic)

import sys sys.setrecursionlimit(10\*\*6)

### Recursion Limit Removal (Advanced)

```
# @bootstrap over recursive function
# replace 'return' with 'yield'
# for when sus method does not work
from types import GeneratorType
def bootstrap(f, stack=[]):
def wrappedfunc(*args, **kwargs):
 if stack:
  return f(*args, **kwargs)
 else:
.to = f(*args, **kwargs)
  while True:
if type(to) is GeneratorType:
   . . stačk . append(to)
     to = next(to)
     stack.pop()
     if not stack:
to = stack[-1].send(to)
return to
return wrappedfunc
```

### Python 3 Compatibility

from \_future\_ import division, print\_function if sys.version\_info[0] < 3: from \_builtin import xrange as range from future\_builtins import ascii, filter,  $\rightarrow$  hex, map, oct, zip

## 12 Additional

## Judge Speed

```
// kattis: 0.50s
// codeforces: 0.421s
// atcoder: 0.455s
#include <bits/stdc++.h>
using namespace std;
int v = 1e9/2, p = 1;
int main() {
for (int i = 1: i <= v: i++) p *= i:
cout << p;
```

### Judge Pre-Contest Checks

float128 support? int128 and does extra or missing whitespace cause WA? -documentation up to date? printer usage available and functional?

```
// each case tests a different fail condition
 // try them before contests to see error codes
struct g { int arr[1000000]; g(){}};
vector<g> a;
 // O=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE
if (n == 4) assert(0);
if (n == 5) 0 / 0;
if (n == 6) *(int*)(0) = 0;
 return n + judge(n + 1);
 GCC Builtin Docs
 // 128-bit integer
__int128 a;
unsigned __int128 b;
// 128-bit float
// minor improvements over long double __float128 c; // log2 floor
// log2 floor
__lg(n);
__lg(n);
// number of 1 bits
// can add ll like popcountll for long longs
__builtin_popcount(n);
// number of trailing zeroes
__builtin_ctz(n);
// number of leading zeroes
__builtin_clz(n);
// 1-indexed least significant 1 bit
__builtin_ffs(n);
_builtin_ffs(n);
// parity of number
 __builtin_parity(n);
Limits
                         int
 uint
          \pm 9223372036854775807 | \pm 2^{63} - 1 | 10^{18}
 11
                                                  2^{64} - 1 | 10^{19}
          18446744073709551615
 i128 \mid \pm 170141183460469231... \mid \pm 2^{\overline{1}27} - 1 \mid 10^{38}
                                               \frac{1}{2} \frac{1}{10} \frac{1}{10} \frac{1}{10} \frac{1}{10} \frac{1}{10}
 u128 340282366920938463...
 Complexity classes input size (per second):
 O(n^n) or O(n!)
                                                           n \leq 10
 O(2^n)
                                                          n \leq 30
 O(n^3)
                                                       n \le 1000
 O(n^2)
                                                      n < 30000
                                                         n \le 10^6
 O(n\sqrt{n})
                                                         n \le 10^7
 O(n \log n)
                                                         n < 10^9
 O(n)
```