```
else n = n * 10 + c - '0':
    General
                             7 Graphs
    Algorithms
                              8 2D Geometry
                                                            n = s * (n + m * o):
    Structures
                              9 3D Geometry
                                                           void read(double& n) {
    Strings
                                                            ld m; read(m); n = m;
                              10 Optimization
    Greedy
                              11 Additional
                                                           void read(float& n) {
 ld m: read(m): n = m:
    Math
     General
                                                            void read(string& s) {
                                                            char c; s = "
g++ -g -02 -std=gnu++17 -static prog.cpp
./a.exe
run.sh
                                                            while((c=getchar unlocked())!=' '&&c!='\n')
                                                            bool readline(string& s) {
                                                            char c; s = "";
while(c=getchar unlocked()) {
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp
for i in *.in; do
                                                             if (c == '\n') return true;
if (c == EOF) return false;
s += c;
 f=${i%.in}
 ./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
                                                            return false;
done
                                                            void print(unsigned int n) {
Header
                                                            if (n / 10) print(n / 10);
// use better compiler options
                                                            putchar_unlocked(n % 10 + '0');
#pragma GCC optimize("Ofast","unroll-loops")
#pragma GCC target("avx2,fma")
                                                            void print(int n) {
// include everything
                                                            if (n < 0) { putchar_unlocked('-'); n*=-1; }
 #include <bits/stdc++.h>
                                                            print((unsigned int)n);
#include <bits/extc++.h>
#include <sys/resource.h>
// namespaces
                                                           Common Structs
using namespace std;
                                                               n-dimension vectors
using namespace __gnu_cxx; // rope
                                                               Vec<2, int>v(n, m) = arr[n][m]
using namespace __gnu_pbds; // tree/trie
                                                            // Vec<2, int> v(n, m, -1) default init -1
                                                            template<int D, typename T>
// common defines
#define fastio
                                                            struct Vec : public vector < Vec < D-1, T >> {
                                                              template<typename... Args>

→ ios base::sync with stdio(0);cin.tie(0);
                                                              Vec(int n=0, Args... args) : vector<Vec<D-1,
#define nostacklim rlimit RZ; getrlimit(3,&RZ
                                                            \rightarrow T>>(n, Vec<D-1, T>(args...)) {}
    ):RZ.rlim cur=-1:setrlimit(3.&RZ):
#define DEBUG(v) cerr<< LINE <<": "<<#v<<" =
                                                           template<typename T>
\Rightarrow "<<v<<'\n'; #define TIMER
                                                           struct Vec<1, T> : public vector<T> {
                                                              Vec(int n=0, T val=T()) : vector<T>(n, val)

→ cerr<<1.0*clock()/CLOCKS_PER_SEC<<"s\n";
#define ll long long
#define ull unsigned ll
#define i128 __int128
#define u128 unsigned i128
                                                               {}
                                                                Algorithms
#define ld long double
                                                           Min/Max Subarray
// global variables
                                                              max - compare = a < b, reset = a < 0
mt19937 rng((uint32_t)chrono::steady
                                                            \frac{1}{min} - compare = a > b, reset = a > 0

    clock::now().time since epoch().count());

                                                           // returns {sum, {start, end}}
pair<int, pair<int, int>>
Fast IO
                                                                ContiguousSubarray(int* a, int size,
#ifdef _WIN32
                                                                bool(*compare)(int, int),
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
                                                            bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
                                                            0, s = 0;
for (int i = 0; i < size; i++) {
  cur += a[i];</pre>
void read(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                              if ((*compare)(best, cur)) { best = cur;
  n = n * 10 + c - 0';
                                                            \rightarrow start = s; end = i; }
void read(int& n) {
  char c; n = 0; int s = 1
                                                             if ((*reset)(cur)) { cur = 0; s = i + 1; }
 if ((c=getchar_unlocked())=='-') s = -1;
                                                            return {best, {start, end}}:
 else n = c - '0';
while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                            Quickselect
 n = n * 10 + c - 0';
                                                           #define OSNE -999999
                                                           int partition(int arr[], int 1, int r)
void read(ld& n) {
 char c; n = 0;
ld m = 0, o = 1; bool d = false; int s = 1;
if ((c=getchar_unlocked())=='-') s = -1;
                                                            int x = arr[r], i = 1;
for (int j = 1; j <= r - 1; j++)
...if (arr[j] <= x)
...swap(arr[i++], arr[j]);</pre>
 else if (c == .'.') d = true;
else n = c - '0';
 while ((c=getchar_unlocked())!=' '&&c!='\n') {
                                                            swap(arr[i], arr[r]);
 if (c == '.') d = true;
else if (d) { m=m*10+c-'0'; o*=0.1; }
                                                            return i:
```

```
// find k'th smallest element in unsorted array, void update(int i, int val) {
→ only if all distinct
int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - l + 1)) return QSNE;
swap(arr[1 + rng() % (r-l+1)], arr[r]);
 int pos = partition(arr, 1, r);
if (pos-l==k-1) return arr[pos];
 if (pos-1>k-1) return qselect(arr,1,pos-1,k);
 return qselect(arr, pos+1, r, k-pos+1-1);
// TODO: compare against std::nth_element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted

    on both axis
pair<int, int> saddleback_search(int** arr, int
 \rightarrow x, int y, int v) {
 int i = x-1, j = 0;
 while (i >= 0 && j < y) {
  if (arr[i][j] == v) return {i, j};
  (arr[i][j] > v)? i--: j++;
 return {-1, -1};
Ternary Search
 // < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a, int b, int (*f)(int)) {</pre>
 while (b-a > 4) {
    int m = (a+b)/2;
    if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1:
 for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
   a = i;
 return a;
#define TERNPREC 0.000001
double ternsearch(double a. double b. double
 \leftrightarrow (*f)(double)) {
while (b-a > TERNPREC * 4) {
  double m = (a+b)/2;
  if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
  else b = m + TERNPREC;
 for (double i = a + TERNPREC: i <= b: i +=
     TERNPREC)
      if (TERNCOMP((*f)(a), (*f)(i)))
 return a;
Golden Section Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
double goldsection(double a, double b, double
 while (b-a > eps)
  while (b-a > eps)

if (TERNCOMP(f2,f1)) {

. b = x2; x2 = x1; f2 = f1;

. x1 = b - r*(b-a); f1 = f(x1);
  } else {
   a = x1; x1 = x2; f1 = f2;

x2 = a + r*(b-a); f2 = f(x2);
 return a:
3 Structures
Fenwick Tree
// Fenwick tree, array of cumulative sums -
```

 $\hookrightarrow$  O(log n) updates, O(log n) gets

struct Fenwick { int n; ll\* tree;

```
while (i <= n) {
   tree[i] += val;
   i += i & (-i);
 Fenwick(int size) {
  n = size;
  tree = new ll[n+1];
for (int i = 1; i <= n; i++)
   .tree[i] = 0;
 Fenwick(int* arr, int size) : Fenwick(size) {
  for (int i = 0; i < n; i++)
...update(i, arr[i]);
 ~Fenwick() { delete[] tree; }
 ll operator[](int i) {
  if (i < 0 || i > n) return 0;
  \overline{11} \ \overline{sum} = 0;
  while (i>0)
   sum += tree[i];
   i -= i & (-i):
  return sum:
 ll getRange(int a, int b) { return

    operator[](b) - operator[](a-1); }

Hashtable
// similar to unordered map, but faster
struct chash {
    const uint64 t C = (11)(2e18 * M PI) + 71;
 ll operator()(ll x) const { return
    builtin bswap64(x*C); }
int main() {
  gp_hash_table<11,int,chash>
 \rightarrow hashtable({},{},{},{},{1<<16});
 for (int i = 0; i < 100; i++)
hashtable[i] = 200+i;
 if (hashtable.find(10) != hashtable.end())
   cout << hashtable[10];</pre>
Ordered Set
template <typename T>
using oset = tree<T,null_type,less<T>,rb_tree
    _tag,tree_order_statistics_node_update>;
template <typename T, typename D> using omap = tree<T,D,less<T>,rb_tree
    _tag,tree_order_statistics_node_update>;
int main()
 oset<int> o_set;
o_set.insert(5); o_set.insert(1);
 → o_set.insert(3);
// get second smallest element
 cout << *(o set.find by order(1));</pre>
 // number of elements less than k=4
cout << ' ' << o_set.order_of_key(4) << '\n';</pre>
 // equivalent with ordered map
 omap<int,int> o_map;
o_map[5]=1;o_map[1]=2;o_map[3]=3;
 cout << (*(o_map.find_by_order(1))).first;</pre>
 cout << ' ' << o map.order of key(4) << '\n';
Rope
// O(log n) insert, delete, concatenate
int main() {
 // generate rove
 rope<int> v;
 for (int i = 0: i < 100: i++)
  v.push_back(i);
 // move range to front
 rope<int> copy = v.substr(10, 10);
 v.erase(10, 10);
```

```
v.insert(copy.mutable_begin(), copy);
 // print elements of rope
for (auto it : v)
cout << it << "":
Segment Tree
//max(a,b), min(a,b), a+b, a*b, qcd(a,b), a*b
struct SegmentTree {
 typedef int T;
 static constexpr T UNIT = INT MIN:
 T f(T a, T b) {
 if (a == UNIT) return b;
if (b == UNIT) return a;
 return max(a,b);
 Int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
\rightarrow n(n) {}
 SegmentTree(vector<T> arr)

→ SegmentTree(arr.size()) {
 for (int i=0:i<arr.size():i++)

→ update(i.arr[i]):

 void update(int pos, T val) {
  for (s[pos += n] = val; pos /= 2;)
   s[pos] = f(s[pos * 2], s[pos*2+1]);
 T query(int b, int e) { // query [b, e)
 Tra = UNIT, rb = UNIT;

for (b+=n, e+=n; b<e; b/=2, e/=2) {

    if (b % 2) ra = f(ra, s[b++]);

    if (e % 2) rb = f(s[--e], rb);
  return f(ra. rb):
 T get(int p) { return query(p, p+1); }
Sparse Table
template<class T> struct SparseTable {
 vector<vector<T>> m;
SparseTable(vector<T> arr) {
  m.push_back(arr);
  for (int k = 1; (1<<(k)) <= size(arr); k++)
   m.push back(vector<T>(size(arr)-(1<<k)+1)):
   for (int i = 0; i < size(arr)-(1 << k)+1; i
    m[k][i] = min(m[k-1][i],
   m[k-1][i+(1<<(k-1))]:
 }
// min of range [l,r]
T query(int 1, int r) {
  int k = __lg(r-l+1);
  return \min(m[k][1], m[k][r-(1 << k)+1]):
typedef trie<string, null_type,

→ trie_string_access_traits<>,

 pat_trie_tag, trie_prefix_search_node_update>
int main() {
 // generate trie
 trie_type trie;
for (int i = 0; i < 20; i++)
...trie.insert(to_string(i)); // true if new,
\hookrightarrow false if old
 // print things with prefix "1"
 auto range = trie.prefix_range("1");
 for (auto it = range.first; it !=

    range.second; it++)

  cout << *it <<
Wavelet Tree
using iter = vector<int>::iterator;
struct WaveletTree {
```

```
Vec<2, int> C: int s:
 // sigma = highest value + 1
 WaveletTree(vector<int>& a, int sigma) :
    s(sigma), C(sigma*2, 0) {
  build(a.begin(), a.end(), 0, s-1, 1);
 void build(iter b, iter e, int L, int U, int
  if (L == U) return;
  int M = (L+U)/2:
  C[u].reserve(e-b+1); C[u].push_back(0);
  for (auto it = b; it != e; ++it)
    C[u].push_back(C[u].back() + (*it<=M));
  auto p = stable_partition(b, e, [=](int
    i) {return i <= M: }):
  build(b, p, L, M, u*2);
  build(p, e, M+1, U, u*2+1);
 // number of occurences of x in [0,i)
if (x <= M) i = r, U = M;
else i -= r, L = M+1, ++u;
  return i:
 // number of occurrences of x in [l,r)
int count(int x, int l, int r) {
  return rank(x, r) - rank(x, 1);
 // kth smallest in [l, r)
int kth(int k, int l, int r) const {
int L = 0, U = s-1, u = 1, M, ri, rj;
  while (L != U) {
   M = (L+U)/2;
   ri = C[u][1]; rj = C[u][r]; u*=2;

if (k <= rj-ri) 1 = ri, r = rj, U = M;
   else k -= ri-ri. l -= ri. r -= ri.
   L = M+1. ++u:
  return U;
  // # elements between [x,y] in [l, r)
 mutable int L. U:
 int range(int x, int y, int 1, int r) const {
  if (y < x \text{ or } r \le 1) return 0;
  L = x; U = y;
  return range(1, r, 0, s-1, 1);
 int range(int 1, int r, int x, int y, int u)
    const {
  if (y < L or U < x) return 0;
  if (L \le x \text{ and } y \le U) \text{ return } r-1;
  int M = (x+y)/2, ri = C[u][1], rj = C[u][r];
  return range(ri, rj, x, M, u*2) + range(1-ri, Boyer Moore
    r-rj, M+1, y, u*2+1);
 // # elements <= x in [l, r)
int lte(int x, int l, int r) {
  return range(INT_MIN, x, l, r);</pre>
     Strings
Aho Corasick
// range of alphabet for automata to consider
// MAXC = 26, OFFC = 'a' if only lowercase
const int MAXC = 256;
const int OFFC = 0;
struct aho_corasick {
 struct state
  set<pair<int, int>> out;
  int fail; vector<int> go;
  state() : fail(-1), go(MAXC, -1) {}
 vector<state> s;
```

```
int id = 0:
 aho corasick(string* arr, int size) : s(1) {
 for (int i = 0; i < size; i++) {
   int cur = 0;
   for (int c : arr[i]) {
   if (s[cur].go[c-OFFC] == -1) {
     s[cur].go[c-OFFC] = s.size();
     s.push_back(state());
    cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++}):
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
    s[0].go[c] = 0;
  queue<int> sq;
  for (int c = 0; c < MAXC; c++) {
   if (s[0].go[c] != 0) {
    s[s[0].go[c]].fail = 0;
    sq.push(s[0].go[c]);
  while (sq.size()) {
   int e = sq.front(); sq.pop();
  for (int c = 0; c < MAXC; c++) {
   if (s[e].go[c] != -1) {
     int failure = s[e].fail;
while (s[failure].go[c] == -1)
     failure = s[failure].fail;
failure = s[failure].go[c];
     s[s[e].go[c]].fail = failure;
     for (auto length : s[failure].out)
  s[s[e].go[c]].out.insert(length);
      sq.push(\bar{s}[e].go[c]);
 // list of {start pos, pattern id}
 vector<pair<int, int>> search(string text)
 vector<pair<int, int>> toret;
  int cur = 0;
  for (int i = 0; i < text.size(); i++) {</pre>
   while (s[cur].go[text[i]-OFFC] == -1)
    cur = s[cur].fail;
   cur = s[cur].go[text[i]-OFFC];
   if (s[cur].out.size())
    for (auto end : s[cur].out)
. toret.push back({i - end.first + 1,
    end.second});
  return toret:
struct defint { int i = -1: }:
vector<int> boyermoore(string txt, string pat)
 vector<int> toret: unordered map<char, defint>
 → badchar:
 int m = pat.size(), n = txt.size();
 for (int i = 0: i < m: i++) badchar[pat[i]].i string lcp(string* arr. int n. bool sorted =
 \rightarrow = i:
 while (s <= n - m) {
  int j = m - 1;
  while (i \ge 0 \&\& pat[i] == txt[s + i]) i--;
  if (j < 0) {
  .toret.push_back(s);
  s += (s + m < n) ? m - badchar[txt[s +
 → m]].<mark>i</mark> : 1;
} else
  s += max(1, j - badchar[txt[s + j]].i);
return toret:
                                                         int m = a.length(), n = b.length();
                                                         int L[m+1][n+1];
English Conversion
```

```
|const string ones[] = {"", "one", "two",
    "three", "four", "five", "six", "seven",

    "eight", "nine"};
const string teens[] ={"ten", "eleven",
   "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen");
const string tens[] = {"twenty", "thirty",
    "forty", "fifty", "sixty", "seventy",
const string mags[] = {"thousand", "million",
     "billion", "trillion", "quadrillion",
    "quintillion", "sextillion",
    "septillion"};
string convert(int num, int carry) {
if (num < 0) return "negative " +
    convert(-num, 0):
    (num < 10) return ones[num];
(num < 20) return teens[num % 10]
    (num < 100) return tens[(num / 10) - 2] + (num / 10==0?"": " ) + ones[num / 10];
    (num < 1000) return ones[num / 100]
     (num/100==0?"":" ") + "hundred" +
     (num%100==0?"":" ") + convert(num % 100,
return convert(num / 1000, carry + 1) + " " +
    mags[carry] + " " + convert(num % 1000,

⇒ 0);

string convert(int num) {
return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
   vector<int> toret;
 int m = txt.length(), n = pat.length();
 int next[n + 1];
 for (int i = 0; i < n + 1; i++)
 next[i] = 0;
 for (int i = 1; i < n; i++) {
  int j = next[i + 1];
  while (j > 0 && pat[j] != pat[i])
  j = next[j];
 if (j > 0 || pat[j] == pat[i])
next[i + 1] = j + 1;
 for (int i = 0, j = 0; i < m; i++) {
 if (txt[i] == pat[j]) {
  if (++j == n)
   ..toret.push back(i - j + 1);
 } else if (j > 0) {
...j = next[j];
 return toret;
Longest Common Prefix (array)
// longest common prefix of strings in array

    false) {
    if (n == 0) return "";
}

if (!sorted) sort(arr, arr + n);
string r = ""; int v = 0;
 while (v < arr[0].length() && arr[0][v] ==
→ arr[n-1][v])
    r += arr[0][v++];
return r;
Longest Common Subsequence
string lcs(string a, string b) {
```

for (int i = 0; i <= m; i++) {

```
for (int j = 0; j <= n; j++) {
...if (i == 0 || j == 0) L[i][j] = 0;
...else if (a[i-1] == b[j-1]) L[i][j] =
\hookrightarrow L[i-1][j-1]+1;
...else L[i][j] = max(L[i-1][j], L[i][j-1]);
.
// return L[m][n]; // length of lcs
 string out = "";
 int i = m - 1, j = n - 1;
 while (i >= 0 && j >= 0) {
 if (a[i] == b[j]) {
  .out = a[i--] + out;
  else if (L[i][j+1] > L[i+1][j]) i--;
  else j--;
return out;
Longest Common Substring
// l is array of palindrome length at that
int manacher(string s, int* 1) {
int n = s.length() * 2;
                                                          ws[i - 1];
for (int i = 0, j = 0, k; i < n; i += k, j =
\rightarrow max(j-k, 0)) {
                                                          y[i];
 while (i \ge j \&\& i + j + 1 < n \&\& s[(i-j)/2]
\Rightarrow = s[(i+j+1)/2]) j++;
 .1[i] = j;
 for (k = 1; i >= k && j >= k && l[i-k] !=
   i-k: k++)
  1[i+k] = min(1[i-k], j-k);
return *max_element(1, 1 + n);
Cyclic Rotation (Lyndon)
// simple strings = smaller than its nontrivial
// lyndon factorization = simple strings
   factorized
// "abaaba" -> "ab", "aab", "a"
vector<string> duval(string s) {
 int n = s.length();
 vector<string> lyndon;
for (int i = 0; i < n;) {
int j = i+1, k = i;
 for (; j < n && s[k] <= s[j]; j++)
if (s[k] < s[j]) k = i;
                                                       string a;
   else k++;
  for (; i \le k; i += j - k)
                                                        if (r[v]<=q) {
  lyndon.push_back(s.substr(i,j-k));
 return lvndon:
// lexicographically smallest rotation
int minRotation(string s) {
int n = s.length(); s += s;
auto d = duval(s); int i = 0, a = 0; while (a + d[i].length() < n) a +=
   d[i++].length();
 while (i && d[i] == d[i-1]) a -=
\rightarrow d[i--].length();
return a;
Subsequence Count
// "banana", "ban" >> 3 (ban, ba..n, b..an)
ull subsequences(string body, string subs) {
int m = subs.length(), n = body.length();
 if (m > n) return 0;
 ull** arr = new ull*[m+1]:
for (int i = 0; i <= m; i++) arr[i] = new
\hookrightarrow ull[n+1];
for (int i = 1; i <= m; i++) arr[i][0] = 0;
for (int i = 0; i <= n; i++) arr[0][i] = 1;
```

```
for (int i = 1; i <= m; i++)
 for (int j = 1; j <= n; j++)
...arr[i][j] = arr[i][j-1] + ((body[j-1] ==
    subs[i-1])? arr[i-1][j-1] : 0);
return arr[m][n];
Suffix Array + LCP
struct SuffixArray {
vector<int> sa, lcp;
SuffixArray(string& s, int lim=256) {
int n = s.length() + 1, k = 0, a, b;
  vector<int> x(begin(s), end(s)+1), y(n),
   ws(max(n, lim)), rank(n);
 sa = lcp = y;
iota(begin(sa), end(sa), 0);
  for (int j = 0, p = 0; p < n; j = max(1, j *
 \rightarrow 2), lim = p) {
   p = j; iota(begin(y), end(y), n - j);
   for (int i = 0; i < (n); i++); if (sa[i] >= j)
   y[p++] = sa[i] - j;
fill(begin(ws), end(ws), 0);
   for (int i = 0; i < (n); i++) ws[x[i]]++;
   for (int i = 1; i < (lim); i++) ws[i] +=
   for (int i = n: i--:) sa[--ws[x[v[i]]]] =
   swap(x, y); p = 1; x[sa[0]] = 0;
   for (int i = 1; i < (n); i++) {
    a = sa[i - 1]; b = sa[i];
    x[b] = (y[a] == y[b] && y[a + j] == y[b +
   j]) ? p - 1 : p++;
  for (int i = 1; i < (n); i++) rank[sa[i]] =
 for (int i = 0, j; i < n - 1; lcp[rank[i++]]
   for (k \&\& k--, j = sa[rank[i] - 1];

s[i + k] == s[j + k]; k++);
Suffix Tree (Ukkonen's)
struct SuffixTree {
.// n = 2*len*10 or so
enum { N = 50010, ALPHA = 26 };
int toi(char c) { return c - 'a'; }
void ukkadd(int i, int c) { suff:
  if (t[v][c]==-1) { t[v][c]=m; l[m]=i;
    p[m++]=v; v=s[v]; q=r[v]; goto suff; }
    v=t[v][c]; q=1[v];
  if (q==-1 || c==toi(a[q])) q++; else {
    1[m+1]=i; p[m+1]=m; 1[m]=1[v]; r[m]=q;
   p[m] = p[v]; t[m][c] = m+1; t[m][toi(a[q])] = v;
   l[v]=q; p[v]=m; t[p[m]][toi(a[l[m]])]=m;
v=s[p[m]]; q=l[m];
   while (q < r[m]) = v = t[v] [toi(a[q])];
    q+=r[v]-l[v]; }
   if (q==r[m]) s[m]=v; else s[m]=m+2;
   q=r[v]-(q-r[m]); m+=2; goto suff;
 SuffixTree(string a) : a(a) {
 fill(r,r+N,(int)(a).size());
 memset(s, 0, sizeof s);

memset(t, -1, sizeof t);

fill(t[1],t[1]+ALPHA,0);

s[0]=1;1[0]=1[1]=-1;r[0]=r[1]=p[0]=p[1]=0;
  for(int i=0;i<a.size();i++)
    ukkadd(i.toi(a[i])):
```

```
// Longest Common Substring between 2 strings | . . if (catalan[i] >= mod)
                                                    // returns {length, offset from first string}
                                                    pair<int, int> best;
                                                    int lcs(int node, int i1, int i2, int olen) {
  if (1[node] <= i1 && i1 < r[node]) return 1;
  if (1[node] <= i2 && i2 < r[node]) return 2;</pre>
                                                     int mask=0.
                                                    → len=node?olen+(r[node]-l[node]):0;
                                                    for(int c=0: c<ALPHA: c++) if
                                                      (t[node][c]!=-1)
                                                     mask |= lcs(t[node][c], i1, i2, len);
                                                     if (mask==3)
                                                    → best=max(best.{len.r[node]-len}):
                                                    return mask:
                                                    static pair<int, int> LCS(string s, string t)
                                                    \rightarrow st(s+(char)('z'+1)+t+(char)('z'+2));
                                                    st.lcs(0, s.size(), s.size()+t.size()+1, 0);
return st.best;
                                                  String Utilities
                                                   void lowercase(string& s) {
                                                   transform(s.begin(), s.end(), s.begin(),
                                                   void uppercase(string& s) {
                                                   transform(s.begin(), s.end(), s.begin(),
                                                   void trim(string &s) {
                                                   s.erase(s.begin(),find_if_not(s.begin(),s
                                                        .end(),[](int c){return
                                                      isspace(c);}));
                                                    s.erase(find_if_not(s.rbegin(),s.rend(),[](int

    c){return isspace(c);}).base(),s.end());

                                                  vector<string> split(string& s, char token) {
                                                       vector<string> v; stringstream ss(s);
                                                       for (string e;getline(ss,e,token);)
                                                           v.push_back(e);
                                                       return v:
                                                       Greedy
                                                   Interval Cover
                                                   // L,R = interval [L,R], in = {{l,r}, index}
t[N][ALPHA],1[N],r[N],p[N],s[N],v=0,q=0,m=2 // does not handle case where L == R vector<int> intervalCover(double L, double R,

    vector<pair<double,double>,int>> in) {
                                                       int i = 0; pair<double,int> pos = {L,-1};
                                                      vector<int> a:
                                                       sort(begin(in), end(in));
                                                       while (pos.first < R) {
                                                           double cur = pos.first;
while (i < (int)in.size() &&</pre>
                                                      in[i].first.first <= cur)
                                                       max(pos,{in[i].first.second,in[i].second})
                                                  if (pos.first == cur) return {};
                                                       return a:
                                                       Math
                                                   Catalan Numbers
                                                  ull* catalan = new ull[1000000];
                                                   void genCatalan(int n, int mod) '{
                                                   catalan[0] = catalan[1] = 1;
for (int_i = 2; i <= n; i++) {</pre>
                                                    catalan[i] = 0;
                                                    for (int j = i - 1; j >= 0; j--) {
    catalan[i] += (catalan[j] * catalan[i-j-1])
```

```
catalan[i] -= mod:
// TODO: consider binomial coefficient method
Combinatorics (nCr. nPr)
 // can optimize by precomputing factorials, and
    fact[n]/fact[n-r]
    nPr(ull n, ull r) {
 for (ull i = n-r+1; i <= n; i++)
v *= i;
 return v;
ull nPr(ull n, ull r, ull m) {
 ull v 🖹
for (ull i = n-r+1; i <= n; i++)
v = (v * i) % m;
return v;
úll nCr(ull n, ull r) {
 long double v = 1;
 for (ull i = 1: i <= r: i++)
 v = v * (n-r+i) /i;
 return (ull)(v + 0.001)
// requires modulo math
// ca\bar{n} optimize by precomputing mfac and

→ minv-mfac

ull nCr(ull n, ull r, ull m) {
return mfac(n, m) * minv(mfac(k, m), m) % m *
\rightarrow minv(mfac(n-k, m), m) % m:
Multinomials
ll multinomial(vector<int>& v) {
    ll c = 1, m = v.empty() ? 1 : v[0];
    for(int i = 1; i < v.size(); i++)</pre>
 for (int j = 0; j < v[i]; j++)
...c = c * ++m / (j+1);
 return c:
Chinese Remainder Theorem
bool ecrt(l1* r. l1* m. int n. l1% re. l1% mo)
 11 x, y, d; mo = m[0]; re = r[0];
 for (int i = 1; i < n; i++) {
   d = egcd(mo, m[i], x, y);
 if ((r[i] - re) % d != 0) return false;

x = (r[i] - re) / d * x % (m[i] / d);

re += x * mo;
  mo = mo / d * m[i];
  re %= mo;
 re = (re + mo) \% mo;
 return true:
Count Digit Occurences
 /*count(n,d) counts the number of occurences of
 \rightarrow a digit d in the range \lceil 0.n \rceil * /
ll digit_count(ll n, ll d) {
 .11 result = 0;
 while (n != 0)
 result += ((n\%10) == d ? 1 : 0);
  n /= 10;
 return result:
11 count(11 n, 11 d) {
    if (n < 10) return (d > 0 && n >= d);
    if ((n % 10) != 9) return digit_count(n, d) +
\hookrightarrow count(n-1, d);
return 10*count(n/10, d) + (n/10) + (d > 0);
```

```
Discrete Logarithm
                                                       Farey Fractions
unordered map<int, int> dlogc;
                                                        // generate 0 \le a/b \le 1 ordered, b \le n
int discretelog(int a, int b, int m) {
                                                           farey(4) = 0/1 1/4 1/3 1/2 2/3 3/4 1/1
dlogc.clear();
ll n = sqrt(m)+1, an = 1;
                                                        // length is sum of phi(i) for i = 1 to n
                                                        vector<pair<int, int>> farey(int n) {
for (int i = 0; i < n; i++)
an = (an * a) % m;
                                                         int h = 0, k = 1, x = 1, y = 0, r;
                                                         vector<pair<int, int>> v;
 11 c = an:
 for (int i = 1; i \le n; i++) {
                                                          v.push back({h, k});
 if (!dlogc.count(c)) dlogc[c] = i;
                                                          r = (n-y)/k;
 c = (c * an) % m;
                                                          y += r*k; x' += r*h;
                                                         swap(x,h); swap(y,k);
x = -x; y = -y;
} while (k > 1);
 for (int i = 0; i <= n; i++) {
 if (dlogc.count(c)) return (dlogc[c] * n - i
                                                         v.push_back({1, 1});
                                                         return v;
\rightarrow + m - 1) % (m-1);
 c = (c * a) \% m;
                                                        Fast Fourier Transform
return -1;
                                                        const double PI = acos(-1):
Euler Phi / Totient
                                                        void fft(vector<cd>& a, bool invert) {
                                                         int n = a.size();
int phi(int n) {
                                                         for (int i = 1, j = 0; i < n; i++) {
  int bit = n >> 1;
 int r = n;
for (int i = 2; i * i <= n; i++) {
    if (n % i == 0) r -= r / i;
    while (n % i == 0) n /= i;
                                                          for (; j & bit; bit >>= 1) j ^= bit;
                                                          j ^= bit;
                                                          if (i < j) swap(a[i], a[j]);
 if (n > 1) r = r / n;
 return r;
                                                         for (int len = 2; len <= n; len <<= 1) {
    double ang = 2 * PI / len * (invert ? -1 :
}
#define n 100000
ll phi[n+1];
                                                          cd wlen(cos(ang), sin(ang));
void computeTotient() {
                                                          for (int i = 0; i < n; i += len) {
for (int i=1; i<=n; i++) phi[i] = i;
                                                           .cd w(1):
for (int p=2; p<=n; p++) {
                                                           for (int j = 0; j < len / 2; j++) {
 if (phi[p] == p) {
                                                           cd u = a[i+j], v = a[i+j+len/2] * w;
 phi[p] = p-1;
for (int i = 2*p; i<=n; i += p) phi[i] =</pre>
                                                           a[i+j] = u + v;
a[i+j+len/2] = u - v;
\rightarrow (phi[i]/p) * (p-1);
                                                           .w *= wlen:
                                                         if (invert)
Factorials
                                                          for (auto& x : a)
// digits in factorial
                                                          x /= n;
#define kamenetsky(n) (floor((n * log10(n /
\hookrightarrow ME)) + (log10(2 * MPI * n) / 2.0)) + 1)
                                                        vector<int> fftmult(vector<int> const& a,
// approximation of factorial
#define stirling(n) ((n == 1) ? 1 : sqrt(2 *
                                                            vector<int> const& b) {
                                                         vector<cd> fa(a.begin(), a.end()),
                                                         → fb(b.begin(), b.end());
\hookrightarrow M PI * n) * pow(n / M E, n))
                                                        // natural log of factorial
#define lfactorial(n) (lgamma(n+1))
                                                         fa.resize(n); fb.resize(n);
Prime Factorization
                                                         fft(fa. false): fft(fb. false)
// do not call directly
                                                         for (int i = 0; i < n; i++) fa[i] *= fb[i];
ll pollard rho(ll n. ll s) {
                                                         fft(fa. true):
                                                         vector<int> toret(n);
x = y = rand() % (n - 1) + 1;
int head = 1, tail = 2;
while (true) {
   x = mult(x, x, n);
   x = (x + s) % n;
   if (x - s) % n;
                                                         for (int i = 0; i < n; i++) toret[i] =
                                                        → round(fa[i].real());
return toret;
  if (x == y) return n;
                                                        Greatest Common Denominator
 ll d = __gcd(max(x - y, y - x), n);
if (1 < d && d < n) return d;
                                                       ll egcd(ll a, ll b, ll& x, ll& y) {
  if (b == 0) { x = 1; y = 0; return a; }
  ll gcd = egcd(b, a % b, x, y);
  if (++head == tail) y = x, tail <<= 1;
                                                         x = a / b * y;
                                                         swap(x, y);
// call for prime factors
                                                         return gcd;
void factorize(ll n. vector<ll> &divisor) {
if (n == 1) return;
 if (isPrime(n)) divisor.push back(n):
                                                        Josephus Problem
                                                        // 0-indexed, arbitrary k
 while (d'>= n) d = pollard_rho(n, rand() % (n|int josephus(int n, int k) {
                                                        if (n == 1) return 0;
if (k == 1) return n-1;

    - 1) + 1);
factorize(n / d, divisor);

    factorize(n / d, divisor);

                                                         if (k > n) return (joséphus(n-1,k)+k)%n;
  factorize(d, divisor);
                                                         int res = josephus(n-n/k,k)-n\%k;
                                                         return res + ((res<0)?n:res/(k-1));
```

```
// fast case if k=2, traditional josephus
int josephus(int n) {
return 2*(n-(1<<(32-builtin clz(n)-1)));
Least Common Multiple
#define lcm(a,b) ((a*b)/qcd(a,b))
Modulo Operations
#define MOD 1000000007
#define madd(a,b,m) (a+b-((a+b-m>=0)?m:0)) #define mult(a,b,m) ((ull)a*b\%m)
#define msub(a,b,m) (a-b+((a<b)?m:0))
ll mpow(ll b, ll e, ll m) {
 while (e' > 0) {
  if (e % 2) x = (x * b) % m;
 b = (b * b) \% m;
 e /= 2;
 return x % m:
ull mfac(ull n, ull m) {
for (int i = n; i > 1; i--)
 f = (f * i) \frac{\overline{\%}}{\%} m;
 return f:
// if m is not guaranteed to be prime
11 minv(11 b, 11 m) {
return (x % m + m) % m;
Il mdiv_compmod(int a, int b, int m) {
 if (_gcd(b, m) != 1) return -1;
 return mult(a, minv(b, m), m);
// if m is prime (like 10^9+7)
ll mdiv_primemod (int a, int b, int m) {
return mult(a, mpow(b, m-2, m), m);
Modulo Tetration
ll tetraloop(ll a, ll b, ll m) {
 if(b == 0 | a == 1) return 1:
 ll w = tetraloop(a,b-1,phi(m)), r = 1;
 for (;w;w/=2) {
 if (w\&1) {
 r *= a; if (r >= m) r -= (r/m-1)*m;
 \bar{a} *= a; if (a >= m) a -= (a/m-1)*m;
 return r:
int tetration(int a, int b, int m) {
  if (a == 0 || m == 1) return ((b+1)&1)%m;
return tetraloop(a,b,m) % m;
Matrix
template<typename T>
struct Mat : public Vec<2, T> {
 Mat(int x, int y) : Vec<2, T>(x, y), w(x),
\rightarrow h(v) {}
 static Mat<T> identity(int n) { Mat<T> m(n,n);
    for (int i=0:i<n:i++) m[i][i] = 1: return

    m; }

 Mat<T>& operator+=(const Mat<T>& m) {
 for (int i = 0; i < w; i++)
  for (int j = 0; j < h; j++)
  (*this)[i][j] += m[i][j];
  return *this;
 Mat<T>& operator-=(const Mat<T>& m) {
 for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
(*this)[i][j] -= m[i][j];
  return *this;
```

```
Mat<T> operator*(const Mat<T>& m) {
    Mat < T > z(w,m.h);
    for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
for (int k = 0; k < m.h; k++)
z[i][k] + (*this)[i][j] * m[j][k];
  Mat<T> operator+(const Mat<T>& m) { Mat<T>
 ⇒ a=*this; return a+=m; }
 Mat<T> operator-(const Mat<T>& m) { Mat<T>
 ← a=*this; return a-=m; }
Mat<T>& operator*=(const Mat<T>& m) { return
 \rightarrow *this = (*this)*m: }
  Mat<T> power(int n) {
    Mat<T> a = Mat<T>::identity(w), m=*this;
   for (;n;n/=2,m*=m) if (n\&1) a *=m; return a;
Matrix Exponentiation
// F(n) = c[0]*F(n-1) + c[1]*F(n-2) + \dots
b is the base cases of same length c
ll matrix exponentiation(ll n, vector<ll> c,
procedure | vector<1|> b) {
  if (nth < b.size()) return b[nth-1];
  Mat<1|> a(c.size(), c.size()); ll s = 0;
  for (int i = 0; i < c.size(); i++) a[i][0] =</pre>
 \hookrightarrow a[i][i+1] = 1;
 a = a.power(nth - c.size());
 for (int i = 0; i < c.size(); i++)
s += a[i][0] * b[i];
  return s;
 Matrix Subarray Sums
 template<class T> struct MatrixSum {
  Vec<2, T> p;
  MatrixSum(Vec<2, T>& v) {
    p = Vec<2,T>(v.size()+1, v[0].size()+1);
     for (int i = 0; i < v.size(); i++)
    for (int j = 0; j < v[0].size(); j++)
...p[i+1][j+1] = v[i][j] + p[i][j+1] +
        p[i+1][j] - p[i][i];
 T sum(int u, int l, int d, int r) {
    return p[d][r] - p[d][l] - p[u][r] + p[u][l];
Nimber Arithmetic
 \#define \ nimAdd(a,b) \ ((a)^(b))
ull nimMul(ull a, ull b, int i=6) {
   static const ull M[]={INT_MIN>>32,
          M[0]^{(M[0] << 16)}, M[1]^{(M[1] << 8)},
        M[2]^{M[2]}<<4), M[3]^{M[3]}<<2),
\stackrel{\rightarrow}{\mapsto} \begin{array}{c} M[2] (M[2] < 1); \\ M[4] \cap (M[4] \cap (M[4] < 1); \\ M[4] \cap (M[4] \cap (M[4] < 1); \\ M[4] \cap (M[4] \cap (M[4] \cap (M[4] < 1); \\ M[4] \cap
    if (i-- == 0) return a&b;
     int k=1<<i;</pre>
     ull s=nimMul(a,b,i), m=M[5-i],
          t=nimMul(((a^(a>>k))&m)|(s\&~m),
        ((b^{(b>k)})\&m)|(m\&(~m>>1))<< k, i);
    return ((s^t)\&m)<< k | ((s^(t>>k))\&m);
 Permutation
 // c = array size, n = nth perm, return index
vector<int> gen permutation(int c, int n) {
  vector<int> idx(c), per(c), fac(c); int i;
  for (i = 0; i < c; i++) idx[i] = i;
 for (i = 1; i <= c; i++) fac[i-1] = n%i, n/=i; for (i = c - 1; i >= 0; i--) per[c-i-1] = idx[fac[i]],
    idx.erase(idx.begin() + fac[i]);
```

```
int get_permutation(vector<int>& v) {
 int use = 0, i = 1, r = 0;
for (int e: v) {
   r = r * i++ + __builtin_popcount(use &

    -(1<<e));
..use |= 1 << e;</pre>
 return r;
Permutation (string/multiset)
string freq2str(vector<int>& v) {
 for (int i = 0; i < v.size(); i++)
for (int j = 0; j < v[i]; j++)
s += (char)(i + 'A');
return s;
 string s;
} // nth perm of multiset, n is O-indexed
string gen_permutation(string s, ll n) {
 vector<int> freq(26, 0);
 for (auto e : s) freq[e - 'A']++;
 for (int i = 0; i < 26; i++) if (freq[i] > 0)
  freg[i]--; ll v = multinomial(freg);
  if (n < v) return (char)(i+'A') +

→ gen_permutation(freq2str(freq), n);
freq[i]++; n -= v;
 return "":
Miller-Rabin Primality Test
 // Miller-Rabin primality test - O(10 log^3 n)
// Miller-Rabin primality test -
bool isPrime(ull n) {
   if (n < 2) return false;
   if (n == 2) return true;
   if (n ½ == 0) return false;
   ull s = n - 1;
   while (s % 2 == 0) s /= 2;
   for (int i = 0; i < 10; i++) {
     ull temp = s;
   }
}
   ull a = rand() % (n - 1) + 1;
  ull mod = mpow(a, temp, n);
while (temp!=n-1\&\&mod!=1\&\&mod!=n-1) {
    mod = mult(mod, mod, n);
    temp *= 2;
   if (mod!=n-1&&temp%2==0) return false:
 return true;
 Sieve of Eratosthenes
bitset<100000001> sieve;

// generate sieve - O(n log n)
 void genSieve(int n) {
 void gensieve(int i) {
    sieve[0] = sieve[1] = 1;
    for (ull i = 3; i * i < n; i += 2)
        if (!sieve[i])</pre>
   for (ull j = i * 3; j <= n; j += i * 2)
...sieve[j] = 1;</pre>
// query sieve after it's generated - O(1)
bool querySieve(int n) {
 return n == 2 | | (n \% 2 != 0 \&\& !sieve[n]);
Compile-time Prime Sieve
const int MAXN = 100000;
 template<int N>
struct Sieve {
  bool sieve[N];
  constexpr Sieve() : sieve() {
  sieve[0] = sieve[1] = 1;
 ..for (int i = 2; i * i < N; i++)
...if (!sieve[i])
 for (int j = i * 2; j < N; j += i)
....sieve[j] = 1;
```

```
bool isPrime(int n) {
  static constexpr_Sieve<MAXN> s;
 return !s.sieve[n]:
Simpson's / Approximate Integrals
   integrate f from a to b, k iterations error \langle = (b-a)/18.0 * M * ((b-a)/2k)^2 \rangle
// (0^{-a})/(a^{k}) /4 // where M = max(abs(f)^{**}(x)) for x in [a,b] // "f" is a function "double func(double x)"
double Simpsons (double a, double b, int k,

    double (*f)(double)) {
    double dx = (b-a)/(2.0*k), t = 0;
    for (int i = 0; i < k; i++)
        t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
</pre>
    (*f)(a+(2*i+1)*dx);
 return (t + (*f)(b)) * (b-a) / 6.0 / k;
Common Equations Solvers
// ax^2 + bx + c = 0, find x
vector<double> solveEq(double a, double b,
 → double c) {
vector<double> r:
 double z = b * b
 if (z == 0)
  r.push_back(-b/(2*a));
 else if (z > 0) {
 r.push_back((sqrt(z)-b)/(2*a));
  r.push_back((sqrt(z)+b)/(2*a));
 return r;
\frac{1}{2} / ax^3 + bx^2 + cx + d = 0, find x
vector<double> solveEq(double a, double b,
 double c, double d) {
vector<double> res;
long double a1 = b/a, a2 = c/a, a3 = d/a;
 long double q = (a1*a1 - 3*a2)/9.0, sq =
 \rightarrow -2*sart(a):
 long double r = (2*a1*a1*a1 - 9*a1*a2 +
 \rightarrow 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
 if (z <= 0) {
  theta = acos(r/sqrt(q*q*q));
   res.push_back(sq*cos(theta/3.0) - a1/3.0);
  res.push_back(sq*cos((theta+2.0*PI)/3.0)
  res.push_back(sq*cos((theta+4.0*PI)/3.0) -
    a1/3.0);
 else {
  res.push_back(pow(sqrt(z)+fabs(r), 1/3.0));
  res[0] = (res[0] + q / res[0]) *
     ((r<0)?1:-1) - a1 / 3.0;
 return res;
^{\prime\prime}/ linear diophantine equation ax + by = c,
     find x and y
// infinite solutions of form x+k*b/g, y-k*a/g
bool solveEq(ll a, ll b, ll c, ll &x, ll &y, ll
 g = \overline{egcd(abs(a), abs(b), x, y)};
 if (c % g) return false;
 x *= c / g * ((a < 0) ? -1 : 1);
 y *= c / g * ((b < 0) ? -1 : 1);
return true:
^{\prime}// m = # equations, n = # variables, a[m][n+1]
\Rightarrow = coefficient matrix

// a[i][0]x + a[i][1]y + ... + a[i][n]z =
\stackrel{\hookrightarrow}{\rightarrow} a[i][n+1] // find a solution of some kind to linear
    a[i][n+1]
\rightarrow equation
const double eps = 1e-7;
bool zero(double a) { return (a < eps) && (a >
→ -eps): }
```

```
|vector<double> solveEq(double **a, int m, int | Graycode Conversions
 \rightarrow n) {
 int cur = 0;
for (int i = 0; i < n; i++) {
  for (int j = cur; j < m; j++) {
    if (!zero(a[j][i])) {
     if (j != cur) swap(a[j], a[cur]);
     for (int sat = 0; sat < m; sat++) {
  if (sat == cur) continue;</pre>
      double num = a[sat][i] / a[cur][i];
      for (int sot = 0; sot <= n; sot++)
.a[sat][sot] -= a[cur][sot] * num;
     cur++:
     break:
 for (int j = cur; j < m; j++)
if (!zero(a[j][n])) return vector<double>();
 vector<double> ans(n,0);
 for (int i = 0, sat = 0; i < n; i++)

if (sat < m && !zero(a[sat][i]))

ans[i] = a[sat][n] / a[sat++][i];

return ans;
// solve A[n][n] * x[n] = b[n] linear equation
// rank < n is multiple solutions, -1 is no
    solutions
    'alls' is whether to find all solutions, or
→ anu
const double eps = 1e-12;
int solveEq(Vec<2, double>& A, Vec<1, double>&
 \rightarrow b, Vec<1, double>& x, bool alls=false) {
 int n = A.size(), m = x.size(), rank = 0, br,
 vector < int > col(m); iota(begin(col), end(col), int getDayofYear(ull epoch) {
for(int r = i; r < n; r++)
  for(int c = i; c < n; c++)
    if ((v = fabs(A[r][c])) > bv)
    if (bv <= eps) {
    for(int j = i; j < n; j++)
    if (fabs(b[j]) > eps)
     return -1:
   break:
  swap(A[i], A[br]);
  swap(b[i], b[br]);
  swap(col[i], col[bc]);
  for(int j = 0; j < n; j++)
swap(A[j][i], A[j][bc]);
  bv = 1.0 / A[i][i];
  for(int j = (alls)?0:i+1; j < n; j++) {
   .if (j != i) {
     double fac = A[j][i] * bv;
b[j] -= fac * b[i];
     for(int k = i+1; k < m; k++)
      A[j][k] -= fac*A[i][k];
  rank++:
 if (alls) for (int i = 0; i < m; i++) x[i] =
 → -DBL_MAX;
for (int i = rank; i--;) {
bool isGood = true;
  if (alls)
   for (int j = rank; isGood && j < m; j++)
     if (fabs(A[i][j]) > eps)
  isGood = false;
b[i] /= A[i][i];
  if (isGood) x[col[i]] = b[i];
if (!alls)
   for(int j = 0; j < i; j++)
b[j] -= A[j][i] * b[i];
 return rank;
```

```
ull graycode2ull(ull n) {
ull i = 0;
for (; n; n = n >> 1) i ^= n;
return i:
ull ull2graycode(ull n) { return n ^ (n >> 1);
```

## Unix/Epoch Time

```
// O-indexed month/time, 1-indexed day
// minimum 1970, 0, 1, 0, 0, 0
ull toEpoch(int year, int month, int day, int
 → hour, int minute, int second) {
 t.tm_mday = day; t.tm_hour = hour;
 t.tm_min = minute; t.tm_sec = second;
t.tm_isdst = 0; // 1 = daylights savings
epoch = mktime(&t);
 return (ull)epoch;
vector<int> toDate(ull epoch) {
 time t e=epoch: struct tm t=*localtime(&e):
 return {t.tm year+1900,t.tm mon,t.tm mday,t
   .tm hour,t.tm min,t.tm sec};
int getWeekday(ull epoch) {
 time_t e=epoch; struct tm t=*localtime(&e);
 return t.tm_wday; // 0-6, 0 = sunday
 time_t e=epoch; struct tm t=*localtime(&e);
 return t.tm_yday; // 0-365
const int months[] =

→ {31,28,31,30,31,30,31,31,30,31,30,31};

bool validDate(int year, int month, int day) {
    bool leap = !(year%(year%25?4:16));
    if (month >= 12) return false;
return day <= months[month] + (leap &&</pre>
    month == 1);
```

## Theorems and Formulae

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set of size n. !0 = 1, !1 = 0, !n = (n+1)(!(n-1))|1)+!(n-2),  $!n=n!\sum_{i=0}^{n}\frac{(-1)^{i}}{i!}$ ,  $!n=\left[\frac{n!}{e}\right]$ 

In a partially ordered set, a chain is a subset of elements that are all comparable to eachother. An antichain is a subset where no two are comparable.

**Dilworth's theorem** states the size of a maximal antichain equals the size of a minimal chain cover of a partially ordered set S. The width of S is the maximum size of an antichain in S, which is equal to the minimum number of chains needed to cover S, or the minimum number of chains such that all elements are in at least one chain.

Rosser's Theorem states the nth prime number is greater than n \* ln(n) for n > 1.

```
Nicomachi's Theorem states 1^3 + 2^3 + ... + \begin{vmatrix} const & 11 & inf & = & 1LL & << & 62; \\ \frac{\#define & FOR(i,n) & for & (int & i & = & 0; & i & < & n; & i++)}{n^3} = (1+2+...+n)^2 and is equivalent to void floydWarshall(Vec<2, 11>& m) {
 \begin{array}{c|c} \underline{\text{Lagrange's Four Square Theorem states}} & \overline{\text{FOR}(i,n)} & \underline{\text{min}[i][i]} & \underline{\text{min}(\text{m[i][i]}, \text{OLL})}; \\ \underline{\text{FOR}(k,n)} & \underline{\text{FOR}(k,n)} & \underline{\text{FOR}(i,n)} & \underline{\text{FOR}(j,n)} & \underline{\text{if}} & \underline{\text{m[i][k]}} & \underline{\text{le inf}} \\ \end{array} 
every natural number is the sum of the squares
of four non-negative integers. This is a spe-
cial case of the Fermat Polygonal Number
Theorem where every positive integer is a if (m[i][k] != inf && m[k][j] != inf)
sum of at most n s-gonal numbers. The nth \rightarrow m[i][i] = -inf;
s-gonal number P(s,n) = (s-2)\frac{n(n-1)}{2} + n
7 Graphs
struct edge {
 .int u,v,w;
 edge (int u, int v, int w) : u(u), v(v), w(w) {}
 edge (): u(0), v(0), w(0) {}
bool operator < (const edge &e1, const edge
 \rightarrow &e2) { return e1.w < e2.w; }
bool operator > (const edge &e1, const edge
 \rightarrow &e2) { return e1.w > e2.w; }
struct subset { int p, rank; };
Eulerian Path
#define edge_list vector<edge>
#define adj_sets vector<set<int>>>
struct EulerPathGraph {
 adj_sets graph; // actually indexes incident
 edge_list edges; int n; vector<int> indeg;
  EulerPathGraph(int n): n(n) {
  indeg = *(new vector<int>(n,0));
  graph = *(new adj_sets(n, set<int>()));
  void add_edge(int u, int v) {
   graph[u].insert(edges.size());
   indeg[v]++:
   edges.push_back(edge(u,v,0));
  bool eulerian_path(vector<int> &circuit) {
   if(edges.size()==0) return false;
  stack<int> st;
int a[] = {-1, -1};
for(int v=0; v<n; v++) {
   if(indeg[v]!=graph[v].size()) {
  bool b = indeg[v] > graph[v].size();
if (abs(((int)indeg[v])-((int)graph[v])
      .size())) > 1) return
     false;
if (a[b] != -1) return false;
     .a[b] = v;
```

int s = (a[0]!=-1 && a[1]!=-1 ? a[0] :

while(!st.empty() || !graph[s].empty()) {

circuit.push\_back(s); s = st.top();

int w = edges[\*graph[s].begin()].v;

graph[s].erase(graph[s].begin());

return circuit.size()-1==edges.size();

if(s==-1) return false:

st.push(s); s = w;

circuit.push back(s):

if (graph[s].empty()) {

```
int n = m.size();
                                                       \hookrightarrow && m[k][j] != inf) {
                                                       auto newDist = max(m[i][k] + m[k][j], -inf);
                                                       .m[i][j] = min(m[i][j], newDist);
                                                       FOR(k,n) if (m[k][k] < 0) FOR(i,n) FOR(j,n)
                                                      Minimum Spanning Tree
                                                         returns vector of edges in the mst
                                                      // graph[i] = vector of edges incident to
                                                         vertex i places total weight of the mst in Stotal
                                                       // if returned vector has size != n-1. there is
                                                      vector<edge> mst(vector<vector<edge>> graph,
                                                          l<u>l</u> &total) {
                                                       priority_queue<edge, vector<edge>,
                                                       → greater<edge>> pq;
                                                       vector<edge> MST;
                                                       bitset<20001> marked; // change size as needed
                                                       marked[0] = 1;
                                                       for (edge ep : graph[0]) pq.push(ep);
while(MST.size()!=graph.size()-1 &&
                                                          pq.size()!=0) {
                                                        edge e = pq.top(); pq.pop();
                                                        int u = e.u, v = e.v, w = e.w;

if(marked[u] && marked[v]) continue;

else if(marked[u]) swap(u, v);
                                                        for(edge ep : graph[u]) pq.push(ep);
                                                        marked[u] = 1;
MST.push_back(e);
                                                        total += e.w;
                                                       return MST:
                                                      Union Find
                                                      int uf_find(subset* s, int i) {
  if (s[i].p != i) s[i].p = uf_find(s, s[i].p);
                                                       return s[i].p;
                                                      void uf_union(subset* s, int x, int y) {
                                                       int xp = uf_find(s, x), yp = uf_find(s, y);
                                                       if (s[xp].rank > s[yp].rank) s[yp].p = xp;
                                                       else if (s[xp].rank < s[yp].rank) s[xp].p =
                                                       yp;
else { s[yp].p = xp; s[xp].rank++; }
                                                      2D Grid Shortcut
                                                      #define inbound(x,n) (0 <= x \otimes x < n)
                                                      #define fordir(x,y,n,m) for(auto[dx,dy]:dir)if
\rightarrow (a[0]==-1 && a[1]==-1 ? edges[0].u : -1));
                                                           (inbound(x+dx,n) \& inbound(y+dy,m))
                                                      const pair<int,int> dir[] =
                                                       \leftrightarrow \{\{1,0\},\{0,1\},\{-1,0\},\{0,-1\}\};
                                                           2D Geometry
                                                      #define point complex<double>
#define EPS 0.0000001
                                                      #define sq(a) ((a)*(a))
                                                      #define c\overline{b}(a) ((a)*(a)*(a))
                                                      double dot(point a, point b) { return

→ real(conj(a)*b); }

                                                      double cross(point a, point b) { return
                                                       \rightarrow imag(conj(a)*b); }
                                                      struct line { point a, b; };
                                                      struct circle { point c; double r; };
                                                      struct segment { point a, b; };
                                                      struct triangle { point a, b, c; };
```

```
struct rectangle { point tl, br; };
struct convex polygon {
 vector<point> points;
 convex_polygon(vector<point> points) :
 → points(points) {}
 convex_polygon(triangle a) {
  points.push_back(a.a); points.push_back(a.b);
 → points.push_back(a.c);
 convex_polygon(rectangle a) {
  points.push_back(a.tl);
    points.push back({real(a.tl),
    imag(a.br)});
  points.push_back(a.br);
    points.push back({real(a.br),
    imag(a.tl)});
struct polygon {
 vector<point> points;
 polygon(vector point points) :
 → points(points) {}
 polygon(triangle a) {
  points.push_back(a.a); points.push_back(a.b);
   points.push_back(a.c);
 polygon(rectangle a) {
  points.push_back(a.tl);
    points.push_back({real(a.tl),
    imag(a.br)});
  points.push_back(a.br);
    points.push_back({real(a.br),
    imag(a.tl)}):
 polygon(convex_polygon a) {
  for (point v : a.points)
  points.push_back(v);
// triangle methods
double area heron(double a, double b, double
\stackrel{\hookrightarrow}{} c) {
if (a < b) swap(a, b);
 if (a < c) swap(a, c);
 if (b < c) swap(b, c);
 if (a > b + c) return -1:
 return sqrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)
 \rightarrow /16.0);
// seament methods
double lengthsq(segment a) { return
    sq(real(a.a) - real(a.b)) + sq(imag(a.a)
\stackrel{\hookrightarrow}{\hookrightarrow} imag(a.b)); }
double length(segment a) { return

    sqrt(lengthsq(a)); }

// circle methods
double circumference(circle a) { return 2 * a.r // -1 outside, 0 inside, 1 tangent, 2
| → * M PI; }
double area(circle a) { return sq(a.r) * M_PI;
| → }
|// rectangle methods
double width(rectangle a) { return
→ abs(real(a.br) - real(a.tl)): }
double height (rectangle a) { return

    abs(imag(a.br) - real(a.tl)); }

double diagonal(rectangle a) { return

    sqrt(sq(width(a)) + sq(height(a))); }

double area (rectangle a) { return width(a) *
\rightarrow height(a); }
double perimeter(rectangle a) { return 2 *
// check if `a` fit's inside `b
// swap equalities to exclude tight fits
bool doesFitInside(rectangle a, rectangle b) {
```

```
int x = width(a), w = width(b), y = height(a),
\hookrightarrow h = height(b);
 if (x > y) swap(x, y);
 if (w > h) swap(w, h);
 if (w < x) return false;
 if (y <= h) return true;
double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
return sq(a) <= sq(b) + sq(c);</pre>
// polygon methods
// negative area = CCW, positive = CW
double area(polygon a) {
  double area = 0.0: int n = a.points.size():
  for (int i = 0, j = 1; i < n; i++, j = (j - 1)
\rightarrow 1) % n)
    area += (real(a.points[j]-a.points[i]))*
 → (imag(a.points[j]+a.points[i]));
  return area / 2.0
// get both unsigned area and centroid
pair<double, point> area_centroid(polygon a) {
 int n = a.points.size();
 double area = 0:
 point c(0, 0);
 for (int i = n - 1, j = 0; j < n; i = j++) {
    double v = cross(a.points[i], a.points[j]) /
 c += (a.points[i] + a.points[j]) * (v / 3);
 c /= area;
 return {area, c};
Intersection
// -1 coincide. O parallel. 1 intersection
int intersection(line a, line b, point& p) {
if (abs(cross(a.b - a.a, b.b - b.a)) > EPS) {
 p = cross(b.a - a.a, b.b - a.b) / cross(a.b)
 \rightarrow - a.a, b.b - b.a) * (b - a) + a;
 return 1;
 if (abs(cross(a.b - a.a, a.b - b.a)) > EPS)

→ return 0:

return -1:
// area of intersection
double intersection(circle a, circle b) {
double d = abs(a.c - b.c);
if (d <= b.r - a.r) return area(a);</pre>
 if (d <= a.r - b.r) return area(b);
 if (d \ge a.r + b.r) return 0;
 double alpha = acos((sq(a.r) + sq(d) -
\rightarrow sq(b.r)) / (2 * a.r * d));
 double beta = acos((sq(b.r) + sq(d) - sq(a.r))
 \rightarrow / (2 * b.r * d)):
return sq(a.r) * (alpha - 0.5 * sin(2 *
    alpha) + sq(b.r) * (beta - 0.5 * sin(2 *
   beta)):

    intersection
    int intersection(circle a. circle b.

    vector<point>& inter) {

double d2 = norm(b.c - a.c), rS = a.r + b.r,
\rightarrow rD = a.r - b.r:
if (d2 > sq(rS)) return -1;
 if (d2 < sq(rD)) return 0;
 double ca = 0.5 * (1 + rS * rD / d2);
 point z = point(ca, sqrt(sq(a.r) / d2 -
 \rightarrow sq(ca)));
 inter.push_back(a.c + (b.c - a.c) * z);
 if (abs(imag(z)) > EPS) inter.push_back(a.c +
\leftrightarrow (b.c - a.c) * conj(z));
return inter.size();
// points of intersection
vector<point> intersection(line a, circle c) {
vector <point > inter;
```

else {

```
c.c -= a.a;
                                                       #define sq(a) ((a)*(a))
#define cb(a) ((a)*(a)*(a))
                                                                                                               GCC Builtin Docs
 a.b -= a.a;
                                                                                                                // 128-bit integer
 point m = a.b * real(c.c / a.b):
                                                        double surface(circle a) { return 4 * sq(a.r) * int128 a;
 double d2 = norm(m - c.c);
                                                                                                               unsigned __int128 b;
                                                          M PI; }
 if (d2 > sq(c.r)) return 0;
                                                                                                               // 128-bit float
                                                       double volume(circle a) { return 4.0/3.0 *
 double 1 = \operatorname{sqrt}((\operatorname{sq}(c.r) - d2) / \operatorname{norm}(a.b));
                                                                                                                // minor improvements over long double
                                                         \rightarrow cb(a.r) * M_PI; }
 inter.push_back(a.a + m + 1 * a.b);
                                                                                                                 float128 c;
 if (abs(1) > EPS) inter.push_back(a.a + m - 1
                                                        10 Optimization
                                                                                                               77 log2 floor
\rightarrow * a.b);
                                                                                                               __lg(n);
// number of 1 bits
return inter;
                                                        Snoob
                                                                                                               // can add ll like popcountll for long longs
                                                         // SameNumberOfOneBits, next permutation
// area of intersection
                                                       int snoob(int a) {
  int b = a & -a, c = a + b;
  return c | ((a ^ c) >> 2) / b;
                                                                                                               __builtin_popcount(n);
// number of trailing zeroes
double intersection(rectangle a, rectangle b) {
 double x1 = max(real(a.tl), real(b.tl)), v1 =
                                                                                                                _builtin_ctz(n);

→ max(imag(a.tl), imag(b.tl));
                                                                                                               // number of leading zeroes
                                                       // example usage int main() { char l1[] = {'1
double x2 = min(real(a.br), real(b.br)), y2 =
                                                                                                               builtin_clz(n);
// 1-indexed least significant 1 bit

→ min(imag(a.br), imag(b.br));
                                                        char 11[] = {'1', '2', '3', '4', char 12[] = {'a', 'b', 'c', 'd'};
int d1 = 5, d2 = 4;
return (x2 <= x1 || y2 <= y1) ? 0 :
                                                                                                                 _builtin_ffs(n);
                                                                                                               // parity of number
   (x2-x1)*(y2-y1);
                                                         // prints 12345abcd, 1234a5bcd, ...
                                                                                                                __builtin_parity(n);
                                                         int min = (1 << d1) - 1, max = min << d2;
Convex Hull
                                                                                                               Limits
                                                         for (int i = min; i <= max; i = snoob(i)) {
bool cmp(point a, point b) {
                                                         int p1 = 0, p2 = 0, v = i;
while (p1 < d1 || p2 < d2) {
                                                                                                                                    \pm 2147483647 \mid \pm 2^{31} - 1 \mid 10^9
if (abs(real(a) - real(b)) > EPS) return
                                                                                                               _{\rm int}

    real(a) < real(b);
</pre>
                                                                                                                                                       \bar{2}^{32} - \bar{1}|10^9
                                                                                                                                      4294967295
                                                          cout \langle ((v \& 1) ? 11[p1++] : 12[p2++]);
                                                                                                               uint
if (abs(imag(a) - imag(b)) > EPS) return
                                                                                                                       \pm 9223372036854775807 | \pm \overline{2}^{63} - \overline{1}|\overline{10}^{18}

    imag(a) < imag(b);
</pre>
                                                                                                                                                       2^{64} - 1|10^{19}
                                                                                                               ull
                                                                                                                       18446744073709551615
                                                          cout << '\n';
return false;
                                                                                                                      |\pm 170141183460469231...|\pm 2^{\overline{1}27} - 1|10^{38}
convex_polygon convexhull(polygon a) {
                                                                                                                |\underline{\mathbf{u}}128| \quad 340282366920938463... | \mathbf{\bar{2}}^{128} - \mathbf{\bar{1}}| \mathbf{\bar{1}} \mathbf{\bar{0}}^{38} 
 sort(a.points.begin(), a.points.end(), cmp);
                                                        Powers
                                                                                                               Complexity classes input size (per second):
 vector<point> lower, upper;
                                                        bool isPowerOf2(11 a) {
for (int i = 0; i < a.points.size(); i++) {
  while (lower.size() >= 2 &&
                                                        return a > 0 \&\& !(a \& a-1);
                                                                                                               O(n^n) or O(n!)
                                                                                                                                                              n < 10
                                                                                                               O(2^n)
                                                                                                                                                              n < 30
    cross(lower.back() - lower[lower.size() -
                                                        bool isPowerOf3(11 a) {
return a>0&&!(12157665459056928801u11%a);
   2], a.points[i] - lower.back()) < EPS)
                                                                                                               O(n^3)
                                                                                                                                                            n < 1000
  lower.pop_back();
                                                                                                               |O(n^2)|
                                                                                                                                                          n \le 30000
                                                        bool isPower(ll a, ll b) {
  double x = log(a) / log(b);
  while (upper.size() >= 2 &&
    cross(upper.back() - upper[upper.size()
                                                                                                                O(n\sqrt{n})
                                                                                                                                                             n < 10^{6}
                                                        return abs(x-round(x)) < 0.00000000001;
   2], a.points[i] - upper.back()) > -EPS)
                                                                                                               O(n \log n)
                                                                                                                                                             n < 10^7
   upper.pop back();
                                                                                                                                                             n < 10^9
                                                                                                               |O(n)|
                                                        11 Additional
  lower.push_back(a.points[i]);
  upper.push_back(a.points[i]);
                                                        Judge Speed
 lower.insert(lower.end(), upper.rbegin() + 1.
                                                          kattis: 0.50s
                                                           codeforces: 0.421s
→ upper.rend());
                                                        // atcoder: 0.455s
return convex_polygon(lower);
                                                        #include <bits/stdc++.h>
                                                        using namespace std:
                                                        int v = 1e9/2, p = 1;
    3D Geometry
                                                        int main() {
struct point3d {
                                                        for (int i = 1; i <= v; i++) p *= i;
 double x, y, z;
                                                        cout << p;
 point3d operator+(point3d a) const { return
\rightarrow {x+a.x, y+a.y, z+a.z}; }
                                                        Judge Pre-Contest Checks
point3d operator*(double a) const { return
\hookrightarrow {x*a, y*a, z*a}; }
                                                            int128 and float128 support?
point3d operator-() const { return {-x, -y,
                                                         does extra or missing whitespace cause WA?
\hookrightarrow -z}; }
                                                        documentation up to date?
.point3d operator-(point3d a) const { return
\rightarrow *this + -a; }
                                                        printer usage available and functional?
point3d operator/(double a) const { return

→ *this * (1/a); }

                                                        // each case tests a different fail condition
 double norm() { return x*x + y*y + z*z; }
                                                        // try them before contests to see error codes
 double abs() { return sqrt(norm()); }
                                                        struct g { int arr[1000000]; g(){}};
 point3d normalize() { return *this /
                                                        vector<g> a;

    this->abs(): }

                                                        // O=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE
                                                       → 6=SIGSEGV 7=recursive MLE int judge(int n) {
double dot(point3d a, point3d b) { return
                                                        if (n == 0) exit(0)
\rightarrow a.x*b.x + a.v*b.v + a.z*b.z: }
                                                            (n == 1) while(1);
(n == 2) while(1) a.push_back(g());
point3d cross(point3d a, point3d b) { return
    \{a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,
                                                        if (n == 3) while(1) putchar_unlocked('a');
\stackrel{\hookrightarrow}{\Rightarrow} a.x*b.y - a.y*b.x}; }
                                                            (n == 4)   assert(0);
(n == 5)   0  / 0;
struct line3d { point3d a, b; };
struct plane { double a, b, c, d; } // a*x +
                                                        if (n == 6) * (int*)(0) = 0:
\rightarrow b*u + c*z + d = 0
                                                        return n + judge(n + 1);
struct sphere { point3d c; double r; };
```