```
else n = n * 10 + c - '0':
    General
                             7 Graphs
    Algorithms
                              8 2D Geometry
                                                            n = s * (n + m * o):
    Structures
                              9 3D Geometry
                                                           void read(double& n) {
    Strings
                                                            ld m; read(m); n = m;
                              10 Optimization
    Greedy
                              11 Additional
                                                           void read(float& n) {
 ld m: read(m): n = m:
    Math
     General
                                                            void read(string& s) {
                                                            char c; s = "
g++ -g -02 -std=gnu++17 -static prog.cpp
./a.exe
run.sh
                                                            while((c=getchar unlocked())!=' '&&c!='\n')
                                                            bool readline(string& s) {
                                                            char c; s = "";
while(c=getchar unlocked()) {
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp
for i in *.in; do
                                                             if (c == '\n') return true;
if (c == EOF) return false;
s += c;
 f=${i%.in}
 ./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
                                                            return false;
done
                                                            void print(unsigned int n) {
Header
                                                            if (n / 10) print(n / 10);
// use better compiler options
                                                            putchar_unlocked(n % 10 + '0');
#pragma GCC optimize("Ofast","unroll-loops")
#pragma GCC target("avx2,fma")
                                                            void print(int n) {
// include everything
                                                            if (n < 0) { putchar_unlocked('-'); n*=-1; }
 #include <bits/stdc++.h>
                                                            print((unsigned int)n);
#include <bits/extc++.h>
#include <sys/resource.h>
// namespaces
                                                           Common Structs
using namespace std;
                                                               n-dimension vectors
using namespace __gnu_cxx; // rope
                                                               Vec<2, int>v(n, m) = arr[n][m]
using namespace __gnu_pbds; // tree/trie
                                                            // Vec<2, int> v(n, m, -1) default init -1
                                                            template<int D, typename T>
// common defines
#define fastio
                                                            struct Vec : public vector < Vec < D-1, T >> {
                                                              template<typename... Args>

→ ios base::sync with stdio(0);cin.tie(0);
                                                              Vec(int n=0, Args... args) : vector<Vec<D-1,
#define nostacklim rlimit RZ; getrlimit(3,&RZ
                                                            \rightarrow T>>(n, Vec<D-1, T>(args...)) {}
    ):RZ.rlim cur=-1:setrlimit(3.&RZ):
#define DEBUG(v) cerr<< LINE <<": "<<#v<<" =
                                                           template<typename T>
\Rightarrow "<<v<<'\n'; #define TIMER
                                                           struct Vec<1, T> : public vector<T> {
                                                              Vec(int n=0, T val=T()) : vector<T>(n, val)

→ cerr<<1.0*clock()/CLOCKS_PER_SEC<<"s\n";
#define ll long long
#define ull unsigned ll
#define i128 __int128
#define u128 unsigned i128
                                                               {}
                                                                Algorithms
#define ld long double
                                                           Min/Max Subarray
// global variables
                                                              max - compare = a < b, reset = a < 0
mt19937 rng((uint32_t)chrono::steady
                                                            \frac{1}{min} - compare = a > b, reset = a > 0

    clock::now().time since epoch().count());

                                                           // returns {sum, {start, end}}
pair<int, pair<int, int>>
Fast IO
                                                                ContiguousSubarray(int* a, int size,
#ifdef _WIN32
                                                                bool(*compare)(int, int),
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
                                                            bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
                                                            0, s = 0;
for (int i = 0; i < size; i++) {
  cur += a[i];</pre>
void read(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                              if ((*compare)(best, cur)) { best = cur;
  n = n * 10 + c - 0';
                                                            \rightarrow start = s; end = i; }
void read(int& n) {
  char c; n = 0; int s = 1
                                                             if ((*reset)(cur)) { cur = 0; s = i + 1; }
 if ((c=getchar_unlocked())=='-') s = -1;
                                                            return {best, {start, end}}:
 else n = c - '0';
while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                            Quickselect
 n = n * 10 + c - 0';
                                                           #define OSNE -999999
                                                           int partition(int arr[], int 1, int r)
void read(ld& n) {
 char c; n = 0;
ld m = 0, o = 1; bool d = false; int s = 1;
if ((c=getchar_unlocked())=='-') s = -1;
                                                            int x = arr[r], i = 1;
for (int j = 1; j <= r - 1; j++)
...if (arr[j] <= x)
...swap(arr[i++], arr[j]);</pre>
 else if (c == .'.') d = true;
else n = c - '0';
 while ((c=getchar_unlocked())!=' '&&c!='\n') {
                                                            swap(arr[i], arr[r]);
 if (c == '.') d = true;
else if (d) { m=m*10+c-'0'; o*=0.1; }
                                                            return i:
```

```
// find k'th smallest element in unsorted array, void update(int i, int val) {
→ only if all distinct
int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - l + 1)) return QSNE;
swap(arr[1 + rng() % (r-l+1)], arr[r]);
 int pos = partition(arr, 1, r);
if (pos-l==k-1) return arr[pos];
 if (pos-1>k-1) return qselect(arr,1,pos-1,k);
 return qselect(arr, pos+1, r, k-pos+1-1);
// TODO: compare against std::nth_element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted

    on both axis
pair<int, int> saddleback_search(int** arr, int
 \rightarrow x, int y, int v) {
 int i = x-1, j = 0;
 while (i >= 0 && j < y) {
  if (arr[i][j] == v) return {i, j};
  (arr[i][j] > v)? i--: j++;
 return {-1, -1};
Ternary Search
 // < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a, int b, int (*f)(int)) {</pre>
 while (b-a > 4) {
    int m = (a+b)/2;
    if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1:
 for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
   a = i;
 return a;
#define TERNPREC 0.000001
double ternsearch (double a. double b. double
 \leftrightarrow (*f)(double)) {
while (b-a > TERNPREC * 4) {
  double m = (a+b)/2;
  if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
  else b = m + TERNPREC;
 for (double i = a + TERNPREC: i <= b: i +=
     TERNPREC)
      if (TERNCOMP((*f)(a), (*f)(i)))
 return a;
Golden Section Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
double goldsection(double a, double b, double
 while (b-a > eps)
  while (b-a > eps)

if (TERNCOMP(f2,f1)) {

. b = x2; x2 = x1; f2 = f1;

. x1 = b - r*(b-a); f1 = f(x1);
  } else {
   a = x1; x1 = x2; f1 = f2;

x2 = a + r*(b-a); f2 = f(x2);
 return a:
3 Structures
Fenwick Tree
// Fenwick tree, array of cumulative sums -
```

 \hookrightarrow O(log n) updates, O(log n) gets

struct Fenwick { int n; ll* tree;

```
while (i <= n) {
   tree[i] += val;
   i += i & (-i);
 Fenwick(int size) {
  n = size;
  tree = new ll[n+1];
for (int i = 1; i <= n; i++)
   .tree[i] = 0;
 Fenwick(int* arr, int size) : Fenwick(size) {
  for (int i = 0; i < n; i++)
...update(i, arr[i]);
 ~Fenwick() { delete[] tree; }
 ll operator[](int i) {
  if (i < 0 || i > n) return 0;
  \overline{11} \ \overline{sum} = 0;
  while (i>0)
   sum += tree[i];
   i -= i & (-i):
  return sum:
 ll getRange(int a, int b) { return

    operator[](b) - operator[](a-1); }

Hashtable
// similar to unordered map, but faster
struct chash {
    const uint64 t C = (11)(2e18 * M PI) + 71;
 ll operator()(ll x) const { return
    builtin bswap64(x*C); }
int main() {
  gp_hash_table<11,int,chash>
 \rightarrow hashtable({},{},{},{},{1<<16});
 for (int i = 0; i < 100; i++)
hashtable[i] = 200+i;
 if (hashtable.find(10) != hashtable.end())
   cout << hashtable[10];</pre>
Ordered Set
template <typename T>
using oset = tree<T,null_type,less<T>,rb_tree
    _tag,tree_order_statistics_node_update>;
template <typename T, typename D> using omap = tree<T,D,less<T>,rb_tree
    _tag,tree_order_statistics_node_update>;
int main()
 oset<int> o_set;
o_set.insert(5); o_set.insert(1);
 → o_set.insert(3);
// get second smallest element
 cout << *(o set.find by order(1));</pre>
 // number of elements less than k=4
cout << ' ' << o_set.order_of_key(4) << '\n';</pre>
 // equivalent with ordered map
 omap<int,int> o_map;
o_map[5]=1;o_map[1]=2;o_map[3]=3;
 cout << (*(o_map.find_by_order(1))).first;</pre>
 cout << ' ' << o map.order of key(4) << '\n';
Rope
// O(log n) insert, delete, concatenate
int main() {
 // generate rove
 rope<int> v;
 for (int i = 0: i < 100: i++)
  v.push_back(i);
 // move range to front
 rope<int> copy = v.substr(10, 10);
 v.erase(10, 10);
```

```
v.insert(copy.mutable_begin(), copy);
 // print elements of rope
for (auto it : v)
cout << it << "":
Segment Tree
//max(a,b), min(a,b), a+b, a*b, qcd(a,b), a*b
struct SegmentTree {
 typedef int T;
 static constexpr T UNIT = INT MIN:
 T f(T a, T b) {
 if (a == UNIT) return b;
if (b == UNIT) return a;
 return max(a,b);
 Int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
\rightarrow n(n) {}
 SegmentTree(vector<T> arr)

→ SegmentTree(arr.size()) {
 for (int i=0:i<arr.size():i++)

→ update(i.arr[i]):

 void update(int pos, T val) {
  for (s[pos += n] = val; pos /= 2;)
   s[pos] = f(s[pos * 2], s[pos*2+1]);
 T query(int b, int e) { // query [b, e)
 Tra = UNIT, rb = UNIT;

for (b+=n, e+=n; b<e; b/=2, e/=2) {

    if (b % 2) ra = f(ra, s[b++]);

    if (e % 2) rb = f(s[--e], rb);
  return f(ra. rb):
 T get(int p) { return query(p, p+1); }
Sparse Table
template<class T> struct SparseTable {
 vector<vector<T>> m;
SparseTable(vector<T> arr) {
  m.push_back(arr);
  for (int k = 1; (1<<(k)) <= size(arr); k++)
   m.push back(vector<T>(size(arr)-(1<<k)+1)):
   for (int i = 0; i < size(arr)-(1 << k)+1; i
    m[k][i] = min(m[k-1][i],
   m[k-1][i+(1<<(k-1))]:
 }
// min of range [l,r]
T query(int 1, int r) {
  int k = __lg(r-l+1);
  return \min(m[k][1], m[k][r-(1 << k)+1]):
typedef trie<string, null_type,

→ trie_string_access_traits<>,

 pat_trie_tag, trie_prefix_search_node_update>
int main() {
 // generate trie
 trie_type trie;
for (int i = 0; i < 20; i++)
...trie.insert(to_string(i)); // true if new,
\hookrightarrow false if old
 // print things with prefix "1"
 auto range = trie.prefix_range("1");
 for (auto it = range.first; it !=

    range.second; it++)

  cout << *it <<
Wavelet Tree
using iter = vector<int>::iterator;
struct WaveletTree {
```

```
Vec<2, int> C: int s:
 // sigma = highest value + 1
 WaveletTree(vector<int>& a, int sigma) :
    s(sigma), C(sigma*2, 0) {
  build(a.begin(), a.end(), 0, s-1, 1);
 void build(iter b, iter e, int L, int U, int
  if (L == U) return;
  int M = (L+U)/2:
  C[u].reserve(e-b+1); C[u].push_back(0);
  for (auto it = b; it != e; ++it)
    C[u].push_back(C[u].back() + (*it<=M));
  auto p = stable_partition(b, e, [=](int
    i) {return i <= M: }):
  build(b, p, L, M, u*2);
  build(p, e, M+1, U, u*2+1);
 // number of occurences of x in [0,i)
if (x <= M) i = r, U = M;
else i -= r, L = M+1, ++u;
  return i:
 // number of occurrences of x in [l,r)
int count(int x, int l, int r) {
  return rank(x, r) - rank(x, 1);
 // kth smallest in [l, r)
int kth(int k, int l, int r) const {
int L = 0, U = s-1, u = 1, M, ri, rj;
  while (L != U) {
   M = (L+U)/2;
   ri = C[u][1]; rj = C[u][r]; u*=2;

if (k <= rj-ri) 1 = ri, r = rj, U = M;
   else k -= ri-ri. l -= ri. r -= ri.
   L = M+1. ++u:
  return U;
  // # elements between [x,y] in [l, r)
 mutable int L. U:
 int range(int x, int y, int 1, int r) const {
  if (y < x \text{ or } r \le 1) return 0;
  L = x; U = y;
  return range(1, r, 0, s-1, 1);
 int range(int 1, int r, int x, int y, int u)
    const {
  if (y < L or U < x) return 0;
  if (L \le x \text{ and } y \le U) \text{ return } r-1;
  int M = (x+y)/2, ri = C[u][1], rj = C[u][r];
  return range(ri, rj, x, M, u*2) + range(1-ri, Boyer Moore
    r-rj, M+1, y, u*2+1);
 // # elements <= x in [l, r)
int lte(int x, int l, int r) {
  return range(INT_MIN, x, l, r);</pre>
     Strings
Aho Corasick
// range of alphabet for automata to consider
// MAXC = 26, OFFC = 'a' if only lowercase
const int MAXC = 256;
const int OFFC = 0;
struct aho_corasick {
 struct state
  set<pair<int, int>> out;
  int fail; vector<int> go;
  state() : fail(-1), go(MAXC, -1) {}
 vector<state> s;
```

```
int id = 0:
 aho corasick(string* arr, int size) : s(1) {
 for (int i = 0; i < size; i++) {
   int cur = 0;
   for (int c : arr[i]) {
   if (s[cur].go[c-OFFC] == -1) {
     s[cur].go[c-OFFC] = s.size();
     s.push_back(state());
    cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++}):
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
    s[0].go[c] = 0;
  queue<int> sq;
  for (int c = 0; c < MAXC; c++) {
   if (s[0].go[c] != 0) {
    s[s[0].go[c]].fail = 0;
    sq.push(s[0].go[c]);
  while (sq.size()) {
   int e = sq.front(); sq.pop();
  for (int c = 0; c < MAXC; c++) {
   if (s[e].go[c] != -1) {
     int failure = s[e].fail;
while (s[failure].go[c] == -1)
     failure = s[failure].fail;
failure = s[failure].go[c];
     s[s[e].go[c]].fail = failure;
     for (auto length : s[failure].out)
  s[s[e].go[c]].out.insert(length);
      sq.push(\bar{s}[e].go[c]);
 // list of {start pos, pattern id}
 vector<pair<int, int>> search(string text)
 vector<pair<int, int>> toret;
  int cur = 0;
  for (int i = 0; i < text.size(); i++) {</pre>
   while (s[cur].go[text[i]-OFFC] == -1)
    cur = s[cur].fail;
   cur = s[cur].go[text[i]-OFFC];
   if (s[cur].out.size())
    for (auto end : s[cur].out)
. toret.push back({i - end.first + 1,
    end.second});
  return toret:
struct defint { int i = -1; }:
vector<int> boyermoore(string txt, string pat)
 vector<int> toret: unordered map<char, defint>
 → badchar:
 int m = pat.size(), n = txt.size();
 for (int i = 0: i < m: i++) badchar[pat[i]].i string lcp(string* arr. int n. bool sorted =
 \rightarrow = i:
 while (s <= n - m) {
  int j = m - 1;
  while (j \ge 0 \&\& pat[j] == txt[s + j]) j--;
  if (j < 0) {
  .toret.push_back(s);
  s += (s + m < n) ? m - badchar[txt[s +
 → m]].<mark>i</mark> : 1;
} else
  s += max(1, j - badchar[txt[s + j]].i);
return toret:
                                                         int m = a.length(), n = b.length();
                                                         int L[m+1][n+1];
English Conversion
```

```
|const string ones[] = {"", "one", "two",
    "three", "four", "five", "six", "seven",

    "eight", "nine"};
const string teens[] ={"ten", "eleven",
   "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen");
const string tens[] = {"twenty", "thirty",
    "forty", "fifty", "sixty", "seventy",
const string mags[] = {"thousand", "million",
     "billion", "trillion", "quadrillion",
    "quintillion", "sextillion",
    "septillion"};
string convert(int num, int carry) {
if (num < 0) return "negative " +
    convert(-num, 0):
    (num < 10) return ones[num];
(num < 20) return teens[num % 10]
    (num < 100) return tens[(num / 10) - 2] + (num / 10==0?"": " ) + ones[num / 10];
    (num < 1000) return ones[num / 100]
     (num/100==0?"":" ") + "hundred" +
     (num%100==0?"":" ") + convert(num % 100,
return convert(num / 1000, carry + 1) + " " +
    mags[carry] + " " + convert(num % 1000,

⇒ 0);

string convert(int num) {
return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
   vector<int> toret;
 int m = txt.length(), n = pat.length();
 int next[n + 1];
 for (int i = 0; i < n + 1; i++)
 next[i] = 0;
 for (int i = 1; i < n; i++) {
  int j = next[i + 1];
  while (j > 0 && pat[j] != pat[i])
  j = next[j];
 if (j > 0 || pat[j] == pat[i])
next[i + 1] = j + 1;
 for (int i = 0, j = 0; i < m; i++) {
 if (txt[i] == pat[j]) {
  if (++j == n)
   ..toret.push back(i - j + 1);
 } else if (j > 0) {
...j = next[j];
 return toret;
Longest Common Prefix (array)
// longest common prefix of strings in array

  false) {
  if (n == 0) return "";
}
if (!sorted) sort(arr, arr + n);
string r = ""; int v = 0;
 while (v < arr[0].length() && arr[0][v] ==
→ arr[n-1][v])
    r += arr[0][v++];
return r;
Longest Common Subsequence
string lcs(string a, string b) {
```

for (int i = 0; i <= m; i++) {

```
for (int j = 0; j <= n; j++) {
...if (i == 0 || j == 0) L[i][j] = 0;
...else if (a[i-1] == b[j-1]) L[i][j] =
\hookrightarrow L[i-1][j-1]+1;
...else L[i][j] = max(L[i-1][j], L[i][j-1]);
.
// return L[m][n]; // length of lcs
 string out = "";
 int i = m - 1, j = n - 1;
 while (i >= 0 && j >= 0) {
 if (a[i] == b[j]) {
  .out = a[i--] + out;
  else if (L[i][j+1] > L[i+1][j]) i--;
  else j--;
return out;
Longest Common Substring
// l is array of palindrome length at that
int manacher(string s, int* 1) {
int n = s.length() * 2;
                                                          ws[i - 1];
for (int i = 0, j = 0, k; i < n; i += k, j =
\rightarrow max(j-k, 0)) {
                                                          y[i];
 while (i \ge j \&\& i + j + 1 < n \&\& s[(i-j)/2]
\Rightarrow = s[(i+j+1)/2]) j++;
 .1[i] = j;
 for (k = 1; i >= k && j >= k && l[i-k] !=
   i-k: k++)
  1[i+k] = min(1[i-k], j-k);
return *max_element(1, 1 + n);
Cyclic Rotation (Lyndon)
// simple strings = smaller than its nontrivial
// lyndon factorization = simple strings
   factorized
// "abaaba" -> "ab", "aab", "a"
vector<string> duval(string s) {
 int n = s.length();
 vector<string> lyndon;
for (int i = 0; i < n;) {
int j = i+1, k = i;
 for (; j < n && s[k] <= s[j]; j++)
if (s[k] < s[j]) k = i;
                                                       string a;
   else k++;
  for (; i \le k; i += j - k)
                                                        if (r[v]<=q) {
  lyndon.push_back(s.substr(i,j-k));
 return lvndon:
// lexicographically smallest rotation
int minRotation(string s) {
int n = s.length(); s += s;
auto d = duval(s); int i = 0, a = 0; while (a + d[i].length() < n) a +=
   d[i++].length();
 while (i && d[i] == d[i-1]) a -=
\rightarrow d[i--].length();
return a;
Subsequence Count
// "banana", "ban" >> 3 (ban, ba..n, b..an)
ull subsequences(string body, string subs) {
int m = subs.length(), n = body.length();
 if (m > n) return 0;
 ull** arr = new ull*[m+1]:
for (int i = 0; i <= m; i++) arr[i] = new
\hookrightarrow ull[n+1];
for (int i = 1; i <= m; i++) arr[i][0] = 0;
for (int i = 0; i <= n; i++) arr[0][i] = 1;
```

```
for (int i = 1; i <= m; i++)
 for (int j = 1; j <= n; j++)
...arr[i][j] = arr[i][j-1] + ((body[j-1] ==
    subs[i-1])? arr[i-1][j-1] : 0);
return arr[m][n];
Suffix Array + LCP
struct SuffixArray {
vector<int> sa, lcp;
SuffixArray(string& s, int lim=256) {
int n = s.length() + 1, k = 0, a, b;
  vector<int> x(begin(s), end(s)+1), y(n),
   ws(max(n, lim)), rank(n);
 sa = lcp = y;
iota(begin(sa), end(sa), 0);
  for (int j = 0, p = 0; p < n; j = max(1, j *
 \rightarrow 2), lim = p) {
   p = j; iota(begin(y), end(y), n - j);
   for (int i = 0; i < (n); i++); if (sa[i] >= j)
   y[p++] = sa[i] - j;
fill(begin(ws), end(ws), 0);
   for (int i = 0; i < (n); i++) ws[x[i]]++;
   for (int i = 1; i < (lim); i++) ws[i] +=
   for (int i = n: i--:) sa[--ws[x[v[i]]]] =
   swap(x, y); p = 1; x[sa[0]] = 0;
   for (int i = 1; i < (n); i++) {
    a = sa[i - 1]; b = sa[i];
    x[b] = (y[a] == y[b] && y[a + j] == y[b +
   j]) ? p - 1 : p++;
  for (int i = 1; i < (n); i++) rank[sa[i]] =
 for (int i = 0, j; i < n - 1; lcp[rank[i++]]
   for (k \&\& k--, j = sa[rank[i] - 1];

s[i + k] == s[j + k]; k++);
Suffix Tree (Ukkonen's)
struct SuffixTree {
.// n = 2*len*10 or so
enum { N = 50010, ALPHA = 26 };
int toi(char c) { return c - 'a'; }
void ukkadd(int i, int c) { suff:
  if (t[v][c]==-1) { t[v][c]=m; l[m]=i;
    p[m++]=v; v=s[v]; q=r[v]; goto suff; }
    v=t[v][c]; q=1[v];
  if (q==-1 || c==toi(a[q])) q++; else {
    1[m+1]=i; p[m+1]=m; 1[m]=1[v]; r[m]=q;
   p[m] = p[v]; t[m][c] = m+1; t[m][toi(a[q])] = v;
   l[v]=q; p[v]=m; t[p[m]][toi(a[l[m]])]=m;
v=s[p[m]]; q=l[m];
   while (q < r[m]) = v = t[v] [toi(a[q])];
    q+=r[v]-l[v]; }
   if (q==r[m]) s[m]=v; else s[m]=m+2;
   q=r[v]-(q-r[m]); m+=2; goto suff;
 SuffixTree(string a) : a(a) {
 fill(r,r+N,(int)(a).size());
 memset(s, 0, sizeof s);

memset(t, -1, sizeof t);

fill(t[1],t[1]+ALPHA,0);

s[0]=1;1[0]=1[1]=-1;r[0]=r[1]=p[0]=p[1]=0;
  for(int i=0;i<a.size();i++)
    ukkadd(i.toi(a[i])):
```

```
// Longest Common Substring between 2 strings | . . if (catalan[i] >= mod)
                                                    // returns {length, offset from first string}
                                                    pair<int, int> best;
                                                    int lcs(int node, int i1, int i2, int olen) {
  if (1[node] <= i1 && i1 < r[node]) return 1;
  if (1[node] <= i2 && i2 < r[node]) return 2;</pre>
                                                     int mask=0.
                                                    → len=node?olen+(r[node]-l[node]):0;
                                                    for(int c=0: c<ALPHA: c++) if
                                                      (t[node][c]!=-1)
                                                     mask |= lcs(t[node][c], i1, i2, len);
                                                     if (mask==3)
                                                    → best=max(best.{len.r[node]-len}):
                                                    return mask:
                                                    static pair<int, int> LCS(string s, string t)
                                                    \rightarrow st(s+(char)('z'+1)+t+(char)('z'+2));
                                                    st.lcs(0, s.size(), s.size()+t.size()+1, 0);
return st.best;
                                                  String Utilities
                                                   void lowercase(string& s) {
                                                   transform(s.begin(), s.end(), s.begin(),
                                                   void uppercase(string& s) {
                                                   transform(s.begin(), s.end(), s.begin(),
                                                   void trim(string &s) {
                                                   s.erase(s.begin(),find_if_not(s.begin(),s
                                                        .end(),[](int c){return
                                                      isspace(c);}));
                                                    s.erase(find_if_not(s.rbegin(),s.rend(),[](int

    c){return isspace(c);}).base(),s.end());

                                                  vector<string> split(string& s, char token) {
                                                       vector<string> v; stringstream ss(s);
                                                       for (string e;getline(ss,e,token);)
                                                           v.push_back(e);
                                                       return v:
                                                       Greedy
                                                   Interval Cover
                                                   // L,R = interval [L,R], in = {{l,r}, index}
t[N][ALPHA],1[N],r[N],p[N],s[N],v=0,q=0,m=2 // does not handle case where L == R vector<int> intervalCover(double L, double R,

    vector<pair<double,double>,int>> in) {
                                                       int i = 0; pair<double,int> pos = {L,-1};
                                                      vector<int> a:
                                                       sort(begin(in), end(in));
                                                       while (pos.first < R) {
                                                           double cur = pos.first;
while (i < (int)in.size() &&</pre>
                                                      in[i].first.first <= cur)
                                                       max(pos,{in[i].first.second,in[i].second})
                                                  if (pos.first == cur) return {};
                                                       return a:
                                                       Math
                                                   Catalan Numbers
                                                  ull* catalan = new ull[1000000];
                                                   void genCatalan(int n, int mod) '{
                                                   catalan[0] = catalan[1] = 1;
for (int_i = 2; i <= n; i++) {</pre>
                                                    catalan[i] = 0;
                                                    for (int j = i - 1; j >= 0; j--) {
    catalan[i] += (catalan[j] * catalan[i-j-1])
```

```
catalan[i] -= mod:
// TODO: consider binomial coefficient method
Combinatorics (nCr. nPr)
 // can optimize by precomputing factorials, and
    fact[n]/fact[n-r]
    nPr(ull n, ull r) {
 for (ull i = n-r+1; i <= n; i++)
v *= i;
 return v;
ull nPr(ull n, ull r, ull m) {
 ull v 🖹
for (ull i = n-r+1; i <= n; i++)
v = (v * i) % m;
return v;
úll nCr(ull n, ull r) {
 long double v = 1;
 for (ull i = 1: i <= r: i++)
 v = v * (n-r+i) /i;
 return (ull)(v + 0.001)
// requires modulo math
// ca\bar{n} optimize by precomputing mfac and

→ minv-mfac

ull nCr(ull n, ull r, ull m) {
return mfac(n, m) * minv(mfac(k, m), m) % m *
\rightarrow minv(mfac(n-k, m), m) % m:
Multinomials
ll multinomial(vector<int>& v) {
    ll c = 1, m = v.empty() ? 1 : v[0];
    for(int i = 1; i < v.size(); i++)</pre>
 for (int j = 0; j < v[i]; j++)
...c = c * ++m / (j+1);
 return c:
Chinese Remainder Theorem
bool ecrt(l1* r. l1* m. int n. l1% re. l1% mo)
 11 x, y, d; mo = m[0]; re = r[0];
 for (int i = 1; i < n; i++) {
   d = egcd(mo, m[i], x, y);
 if ((r[i] - re) % d != 0) return false;

x = (r[i] - re) / d * x % (m[i] / d);

re += x * mo;
  mo = mo / d * m[i];
  re %= mo;
 re = (re + mo) \% mo;
 return true:
Count Digit Occurences
 /*count(n,d) counts the number of occurences of
 \rightarrow a digit d in the range \lceil 0.n \rceil * /
ll digit_count(ll n, ll d) {
 .11 result = 0;
 while (n != 0)
 result += ((n\%10) == d ? 1 : 0);
  n /= 10;
 return result:
11 count(11 n, 11 d) {
    if (n < 10) return (d > 0 && n >= d);
    if ((n % 10) != 9) return digit_count(n, d) +
\hookrightarrow count(n-1, d);
return 10*count(n/10, d) + (n/10) + (d > 0);
```

```
Discrete Logarithm
                                                        Farey Fractions
unordered map<int, int> dlogc;
                                                         // generate 0 \le a/b \le 1 ordered, b \le n
int discretelog(int a, int b, int m) {
                                                            farey(4) = 0/1 1/4 1/3 1/2 2/3 3/4 1/1
dlogc.clear();
ll n = sqrt(m)+1, an = 1;
                                                         // length is sum of phi(i) for i = 1 to n
                                                         vector<pair<int, int>> farey(int n) {
for (int i = 0; i < n; i++)
an = (an * a) % m;
                                                          int h = 0, k = 1, x = 1, y = 0, r;
                                                          vector<pair<int, int>> v;
 11 c = an:
 for (int i = 1; i \le n; i++) {
                                                           v.push back({h, k});
 if (!dlogc.count(c)) dlogc[c] = i;
                                                           r = (n-y)/k;
 c = (c * an) \% m;
                                                           y += r*k; x' += r*h;
                                                          swap(x,h); swap(y,k);
x = -x; y = -y;
} while (k > 1);
 for (int i = 0; i <= n; i++) {
 if (dlogc.count(c)) return (dlogc[c] * n - i
                                                          v.push_back({1, 1});
                                                          return v;
\rightarrow + m -1) % (m-1);
 c = (c * a) \% m;
                                                         Fast Fourier Transform
return -1;
                                                         const double PI = acos(-1):
Euler Phi / Totient
                                                         void fft(vector<cd>& a, bool invert) {
                                                          int n = a.size();
int phi(int n) {
                                                          for (int i = 1, j = 0; i < n; i++) {
  int bit = n >> 1;
 int r = n;
for (int i = 2; i * i <= n; i++) {
    if (n % i == 0) r -= r / i;
    while (n % i == 0) n /= i;
                                                           for (; j & bit; bit >>= 1) j ^= bit;
                                                           j ^= bit;
                                                           if (i < j) swap(a[i], a[j]);
 if (n > 1) r = r / n;
 return r;
                                                          for (int len = 2; len <= n; len <<= 1) {
    double ang = 2 * PI / len * (invert ? -1 :
}
#define n 100000
ll phi[n+1];
                                                           cd wlen(cos(ang), sin(ang));
void computeTotient() {
                                                           for (int i = 0; i < n; i += len) {
for (int i=1; i<=n; i++) phi[i] = i;
                                                            .cd w(1):
for (int p=2; p<=n; p++) {
                                                            for (int j = 0; j < len / 2; j++) {
 if (phi[p] == p) {
                                                            cd u = a[i+j], v = a[i+j+len/2] * w;
 phi[p] = p-1;
for (int i = 2*p; i<=n; i += p) phi[i] =</pre>
                                                            a[i+j] = u + v;
a[i+j+len/2] = u - v;
\rightarrow (phi[i]/p) * (p-1);
                                                            .w *= wlen:
                                                          if (invert)
Factorials
                                                           for (auto& x : a)
// digits in factorial
                                                           x /= n;
#define kamenetsky(n) (floor((n * log10(n /
\hookrightarrow ME)) + (log10(2 * MPI * n) / 2.0)) + 1)
                                                         vector<int> fftmult(vector<int> const& a,
// approximation of factorial
#define stirling(n) ((n == 1) ? 1 : sqrt(2 *
                                                            vector<int> const& b) {
                                                          vector<cd> fa(a.begin(), a.end()),
                                                          → fb(b.begin(), b.end());
\hookrightarrow M PI * n) * pow(n / M E, n))
                                                         int n = 1 << (32 - _builtin_clz(a.size() +

→ b.size() - 1));
// natural log of factorial
#define lfactorial(n) (lgamma(n+1))
                                                          fa.resize(n); fb.resize(n);
Prime Factorization
                                                          fft(fa, false); fft(fb, false)
// do not call directly
                                                          for (int i = 0; i < n; i++) fa[i] *= fb[i];
ll pollard rho(ll n. ll s) {
                                                          fft(fa. true):
                                                          vector<int> toret(n);
x = y = rand() % (n - 1) + 1;
int head = 1, tail = 2;
while (true) {
   x = mult(x, x, n);
   x = (x + s) % n;
   if (x - s) % n;
                                                          for (int i = 0; i < n; i++) toret[i] =
                                                         → round(fa[i].real());
return toret;
  if (x == y) return n;
                                                         Greatest Common Denominator
 ll d = __gcd(max(x - y, y - x), n);
if (1 < d && d < n) return d;
                                                        ll egcd(ll a, ll b, ll& x, ll& y) {
  if (b == 0) { x = 1; y = 0; return a; }
  ll gcd = egcd(b, a % b, x, y);
  if (++head == tail) y = x, tail <<= 1;
                                                          x = a / b * y;
                                                          swap(x, y);
// call for prime factors
                                                          return gcd;
void factorize(ll n. vector<ll> &divisor) {
if (n == 1) return;
 if (isPrime(n)) divisor.push back(n):
                                                         Josephus Problem
                                                         // 0-indexed, arbitrary k
 while (d'>= n) d = pollard_rho(n, rand() % (n|int josephus(int n, int k) {
                                                         if (n == 1) return 0;
if (k == 1) return n-1;

    - 1) + 1);
factorize(n / d, divisor);

    factorize(n / d, divisor);

                                                          if (k > n) return (josephus(n-1,k)+k)%n;
  factorize(d, divisor);
                                                          int res = josephus(n-n/k,k)-n\%k;
                                                          return res + ((res<0)?n:res/(k-1));
```

```
// fast case if k=2, traditional josephus
int josephus(int n) {
return 2*(n-(1<<(32-builtin clz(n)-1)));
Least Common Multiple
#define lcm(a,b) ((a*b)/qcd(a,b))
Modulo Operations
#define MOD 1000000007
#define madd(a,b,m) (a+b-((a+b-m>=0)?m:0)) #define mult(a,b,m) ((ull)a*b\%m)
#define msub(a,b,m) (a-b+((a<b)?m:0))
ll mpow(ll b, ll e, ll m) {
 while (e' > 0) {
  if (e % 2) x = (x * b) % m;
  b = (b * b) \% m;
  e /= 2;
 return x % m:
ull mfac(ull n, ull m) {
 for (int i = n; i > 1; i--)
 f = (f * i) \frac{\overline{\%}}{\%} m;
 return f:
// if m is not guaranteed to be prime
11 minv(11 b, 11 m) {
return (x % m + m) % m;
Il mdiv compmod(int a, int b, int m) {
 if (_gcd(b, m) != 1) return -1;
 return mult(a, minv(b, m), m);
// if m is prime (like 10^9+7)
ll mdiv_primemod (int a, int b, int m) {
 return mult(a, mpow(b, m-2, m), m);
Modulo Tetration
ll tetraloop(ll a, ll b, ll m) {
 if(b == 0 | a == 1) return 1:
 ll w = tetraloop(a,b-1,phi(m)), r = 1;
 for (;w;w/=2) {
 if (w\&1) {
 r *= a; if (r >= m) r -= (r/m-1)*m;
  \bar{a} *= a; if (a >= m) a -= (a/m-1)*m;
 return r:
int tetration(int a, int b, int m) {
   if (a == 0 | | m == 1) return ((b+1)&1)%m;
 return tetraloop(a,b,m) % m;
Matrix
template<typename T>
struct Mat : public Vec<2, T> {
 Mat(int x, int y) : Vec<2, T>(x, y), w(x),
\rightarrow h(v) {}
 static Mat<T> identity(int n) { Mat<T> m(n,n);
    for (int i=0:i < n:i++) m[i][i] = 1: return

    m; }

 Mat<T>& operator+=(const Mat<T>& m) {
 for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
(*this)[i][j] += m[i][j];
  return *this;
 Mat<T>& operator-=(const Mat<T>& m) {
 for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
(*this)[i][j] -= m[i][j];
  return *this;
```

```
Mat<T> operator*(const Mat<T>& m) {
   Mat < T > z(w,m.h);
  for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
for (int k = 0; k < m.h; k++)
z[i][k] + (*this)[i][j] * m[j][k];
 Mat<T> operator+(const Mat<T>& m) { Mat<T>
 ⇒ a=*this; return a+=m; }
 Mat<T> operator-(const Mat<T>& m) { Mat<T>
 ← a=*this; return a-=m; }
Mat<T>& operator*=(const Mat<T>& m) { return
 \rightarrow *this = (*this)*m: }
 Mat<T> power(int n) {
  Mat<T> a = Mat<T>::identity(w), m=*this;
  for (;n;n/=2,m*=m) if (n\&1) a *=m; return a;
Matrix Exponentiation
// F(n) = c[0]*F(n-1) + c[1]*F(n-2) + ...
// b is the base cases of same length c
ll matrix_exponentiation(ll n, vector<ll> c,
 → vector<ll> b) {
    if (nth < b.size()) return b[nth-1];
    Mat<ll> a(c.size(), c.size()); ll s = 0;
    for (int i = 0; i < c.size(); i++) a[i][0] =

    c[i];

 for (int i = 0; i < c.size() - 1; i++)
 \hookrightarrow a[i][i+1] = 1;
 a = a.power(nth - c.size());
 for (int i = 0; i < c.size(); i++)
s += a[i][0] * b[i];
  return s;
 Matrix Subarray Sums
 template < class T> struct MatrixSum {
 Vec<2, T> p;

MatrixSum(Vec<2, T>& v) {
    p = Vec<2,T>(v.size()+1, v[0].size()+1);
   for (int i = 0; i < v.size(); i++)
for (int j = 0; j < v[0].size(); j++)
p[i+1][j+1] = v[i][j] + p[i][j+1] +
     p[i+1][i] - p[i][i];
 T sum(int u, int l, int d, int r) {
    return p[d][r] - p[d][l] - p[u][r] + p[u][l];
 Mobius Function
const int MAXN = 10000000;
// mu[n] = 0 iff n has no square factors
 // 1 = even number prime factors, -1 = odd
 short mu[MAXN] = \{0, 1\};
 void mobius(){
 for (int i = 1; i < MAXN; i++)
  for (int' j = i + i; j < MAXN; j += i)
...mu[j] -= mu[i];
 Nimber Arithmetic
 #define nimAdd(a,b) ((a)^(b))
 ull nimMul(ull a, ull b, int i=6) {
   static const ull M[]={INT_MIN>>32,
     M[0]^{(M[0] << 16)}, M[1]^{(M[1] << 8)},
    M[2]^(M[2] << 4), M[3]^(M[3] << 2),
if (i-- == 0) return a&b;
int k=1<<i;
   ull s=nimMul(a,b,i), m=M[5-i],
     t=nimMul(((a^(a>>k))&m)|(s&-m),
 \stackrel{\hookrightarrow}{\rightarrow} ((b^{(b>k)})\&m) \mid (m\&(\sim m>>1)) << k, i);
   return ((s^t)\&m)<\langle k|((s^(t)>k))\&m);
```

```
Permutation
// c = array size, n = nth perm, return index
vector<int> gen_permutation(int c, int n) {
 vector<int> idx(c), per(c), fac(c); int i;
 for (i = 0; i < c; i++) idx[i] = i;
for (i = 1; i <= c; i++) fac[i-1] = n\( i, n/=i; \)
for (i = c - 1; i >= 0; i--)
per[c-i-1] = idx[fac[i]],
  idx.erase(idx.begin() + fac[i]);
// get what nth permutation of vector
int get_permutation(vector<int>& v) {
 int use = 0, i = 1, r = 0;
for (int e: v) {
   r = r * i++ + __builtin_popcount(use &
 \rightarrow -(1<<e));
  use = 1 << e;
 return r:
Permutation (string/multiset)
string freq2str(vector<int>& v) {
 string s;
 for (int i = 0; i < v.size(); i++)
for (int j = 0; j < v[i]; j++)
s += (char)(i + 'A');
return s;
 // nth perm of multiset, n is 0-indexed
string gen_permutation(string s, 11 n) {
 vector<int> freq(26, 0);
 for (auto e : s) freq[e - 'A']++;
 for (int i = 0; i < 26; i++) if (freq[i] > 0)
 ← {
 freq[i]--; ll v = multinomial(freq);
  if (n < v) return (char)(i+'A') +

→ gen_permutation(freq2str(freq), n);
freq[i]++; n -= v;
 return "";
Miller-Rabin Primality Test
 // Miller-Rabin primality test - O(10 log^3 n)
bool isPrime(ull n) {
bool isfrime(ull n) {
    if (n < 2) return false;
    if (n = 2) return true;
    if (n % 2 == 0) return false;
    ull s = n - 1;
    while (s % 2 == 0) s /= 2;
    for (int i = 0; i < 10; i++) {</pre>
  ull temp = s;
  ull a = rand() \% (n - 1) + 1;
  ull mod = mpow(a, temp, n);
while (temp!=n-1&&mod!=1&&mod!=n-1) {
   mod = mult(mod, mod, n);
    temp *= 2;
   if (mod!=n-1&&temp%2==0) return false;
 return true;
Sieve of Eratosthenes
bitset<100000001> sieve;
 // generate sieve - O(n log n)
void genSieve(int n) {
 void generate (int in )
sieve[0] = sieve[1] = 1;
for (ull i = 3; i * i < n; i += 2)
    if (!sieve[i])
    for (ull j = i * 3; j <= n; j += i * 2)</pre>
     .sieve[j] = 1;
}
// query sieve after it's generated - O(1)
bool querySieve(int n) {
return n == 2 || (n % 2 != 0 && !sieve[n]);
```

```
Compile-time Prime Sieve
const int MAXN = 100000:
template<int N>
struct Sieve {
 bool sieve[N];
 constexpr Sieve() : sieve() {
   sieve[0] = sieve[1] = 1;
  for (int i = 2; i * i < N; i++)
if (!sieve[i])
    for (int j = i * 2; j < N; j += i)
...sieve[j] = 1;
bool isPrime(int n) {
   static constexpr Sieve<MAXN> s;
 return !s.sieve[n]:
Simpson's / Approximate Integrals
   integrate f from a to b, k iterations
// error <= (b-a)/18.0 * M * ((b-a)/2k)^4

// where M = max(abs(f^{```}(x))) for x in [a,b]

// "f" is a function "double func(double x)"
double Simpsons (double a, double b, int k,
 → double (*f)(double)) {
double dx = (b-a)/(2.0*k), t = 0;

for (int i = 0; i < k; i++)

t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
 \leftrightarrow (*f)(a+(2*i+1)*dx);
return (t + (*f)(b)) * (b-a) / 6.0 / k;
Common Equations Solvers
// ax^2 + bx + c = 0, find x
vector<double> solveEq(double a. double b.
 double c) {
vector<double> r;
double z = b * b - 4 * a * c;
 if (z == 0)
  r.push_back(-b/(2*a));
 else if (z > 0) {
  r.push_back((sqrt(z)-b)/(2*a));
  r.push_back((sqrt(z)+b)/(2*a));
 return r:
 \frac{1}{2} ax^3 + bx^2 + cx + d = 0, find x
vector < double > solveEq (double a, double b,

    double c, double d) {
    vector<double> res;

 long double a1 = b/a, a2 = c/a, a3 = d/a:
 long double q = (a1*a1 - 3*a2)/9.0, sq =
 \rightarrow -2*sqrt(q);
 long double r = (2*a1*a1*a1 - 9*a1*a2 +
 \rightarrow 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
 if (z <= 0) {
  theta = acos(r/sqrt(q*q*q));
  res.push_back(sq*cos(theta/3.0) - a1/3.0);
  res.push_back(sq*cos((theta+2.0*PI)/3.0)
  res.push_back(sq*cos((theta+4.0*PI)/3.0) -
    a1/3.0);
  res.push_back(pow(sqrt(z)+fabs(r), 1/3.0));
  res[0] = (res[0] + q / res[0]) *
    ((r<0)?1:-1) - a1 / 3.0:
// linear diophantine equation ax + by = c,
    find x and y
// infinite solutions of form x+k*b/g, y-k*a/g bool solveEq(11 a, 11 b, 11 c, 11 &x, 11 &y, 11
 g = egcd(abs(a), abs(b), x, y);
 if (c % g) return false;
 x *= c / g * ((a < 0) ? -1 : 1);
```

```
y *= c / g * ((b < 0) ? -1 : 1);
return true;
// m = # equations, n = # variables, a[m][n+1]
 \Rightarrow = coefficient matrix
// a[i][0]x + a[i][1]y + ... + a[i][n]z =
\rightarrow a[i][n+1]
// find a solution of some kind to linear
 \rightarrow equation
const double eps = 1e-7;
bool zero(double a) { return (a < eps) && (a >
vector double > solveEq(double **a, int m, int
  \rightarrow n) { int cur = 0;
 for (int i = 0; i < n; i++) {
   for (int j = cur; j < m; j++) {
      if (!zero(a[j][i])) {
     if (j != cur) swap(a[j], a[cur]);
     for (int sat = 0; sat < m; sat++) {
  if (sat == cur) continue;
       double num = a[sat][i] / a[cur][i];
       for (int sot = 0; sot <= n; sot++)
    a[sat][sot] -= a[cur][sot] * num;
     cur++;
     break:
  .}
  for (int j = cur; j < m; j++)
  if (!zero(a[j][n])) return vector<double>();
  vector<double> ans(n,0);
for (int i = 0, sat = 0; i < n; i++)
    if (sat < m && !zero(a[sat][i]))
    ans[i] = a[sat][n] / a[sat++][i];</pre>
  return ans:
 // solve A[n][n] * x[n] = b[n] linear equation
 // rank < n is multiple solutions, -1 is no
     `alls` is whether to find all solutions, or
G anu
const double eps = 1e-12;
int solveEq(Vec<2, double>& A, Vec<1, double>&
 \rightarrow b, Vec<1, double>& x, bool alls=false) {
 int n = A.size(), m = x.size(), rank = 0, br,
  vector<int> col(m); iota(begin(col), end(col),
  → 0);
 for(int j = i; j < n; j++)
if (fabs(b[j]) > eps)
      return -1:
    break:
   swap(A[i], A[br]);
swap(b[i], b[br]);
   swap(col[i], col[bc]);
   for(int j = 0; j < n; j++)
    swap(A[j][i], A[j][bc]);
  bw = 1.0 / A[i][i];
for(int j = (alls)?0:i+1; j < n; j++) {
   if (j != i) {</pre>
     double fac = A[j][i] * bv;
      b[j] = fac * b[i];
     for(int k = i+1; k < m; k++)
A[i][k] -= fac*A[i][k];
  rank++;
  if (alls) for (int i = 0; i < m; i++) x[i] =
  → -DBL_MAX;
  for (int i = rank; i--;) {
```

```
bool isGood = true:
 if (alls)
  for (int j = rank; isGood && j < m; j++)
  if (fabs(A[i][j]) > eps)
 isGood = false;
b[i] /= A[i][i];
 if (isGood) x[col[i]] = b[i];
if (!alls)
    for(int j = 0; j < i; j++)
    b[j] -= A[j][i] * b[i];</pre>
.}
return rank;
Graycode Conversions
ull graycode2ull(ull n) {
ull ull2graycode(ull n) {
  return n ^ (n >> 1);
Unix/Epoch Time
// O-indexed month/time, 1-indexed day
 // minimum 1970, 0, 1, 0, 0, 0
ull toEpoch(int year, int month, int day, int
→ hour, int minute, int second) {
t.tm_mday = day; t.tm_hour = hour;
t.tm_min = minute; t.tm_sec = second;
t.tm_isdst = 0; // 1 = daylights savings
epoch = mktime(&t);
return (ull)epoch;
vector<int> toDate(ull epoch) {
 time_t e=epoch; struct tm t=*localtime(&e);
return {t.tm year+1900,t.tm mon,t.tm mday,t
   .tm_hour,t.tm_min,t.tm_sec};
int getWeekday(ull epoch) {
 time_t e=epoch; struct tm t=*localtime(&e);
return t.tm_wday; // 0-6, 0 = sunday
int getDayofYear(ull epoch) {
time_t e=epoch; struct tm t=*localtime(&e);
return t.tm_yday; // 0-365
const int months[] =
\rightarrow {31,28,31,30,31,30,31,30,31,30,31};
bool validDate(int year, int month, int day) {
    bool leap = !(year%(year%25?4:16));
   if (month >= 12) return false;
return day <= months[month] + (leap &&
   month == 1);
```

Theorems and Formulae

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set of size n. !0 = 1, !1 = 0, !n = (n + 1)(!(n - 1) + !(n - 2)), $!n = n! \sum_{i=0}^{n} \frac{(-1)^{i}}{i!}$, $!n = [\frac{n!}{e}]$

In a partially ordered set, a chain is a subset of elements that are all comparable to eachother. An antichain is a subset where no two are comparable.

Dilworth's theorem states the size of a maximal antichain equals the size of a minimal chain cover of a partially ordered set S. The width of S is the maximum size of an antichain in S, which is equal to the minimum number

of chains needed to cover S, or the minimum circuit.push_back(s); number of chains such that all elements are in return circuit.size()-1==edges.size(); at least one chain.

Rosser's Theorem states the nth prime Floyd Warshall number is greater than n * ln(n) for n > 1.

Lagrange's Four Square Theorem states - && m[k][j] != inf) every natural number is the sum of the squares of four non-negative integers. This is a special case of the Fermat Polygonal Number **Theorem** where every positive integer is a sum of at most n s-gonal numbers. The nths-gonal number $P(s,n) = (s-2)\frac{n(n-1)}{2} + n$ Minimum Spanning Tree

7 Graphs

```
struct edge {
int u,v,w;
edge (int u,int v,int w) : u(u),v(v),w(w) {}
edge (): u(0), v(0), w(0) {}
bool operator < (const edge &e1, const edge
\rightarrow &e2) { return e1.w < e2.w; }
bool operator > (const edge &e1, const edge
```

```
struct subset { int p, rank; };
Eulerian Path
#define edge_list vector<edge>
#define_adj_sets_vector<set<int>>
struct EulerPathGraph {
adj_sets graph; // actually indexes incident

→ edges

edge_list edges; int n; vector<int> indeg;
 EulerPathGraph(int n): n(n) {
 indeg = *(new vector<int>(n,0));
 graph = *(new adj_sets(n, set<int>()));
 void add_edge(int u, int v) {
 graph[u].insert(edges.size());
  indeg[v]++;
 edges.push back(edge(u,v,0));
 bool eulerian_path(vector<int> &circuit) {
 if(edges.size()==0) return false;
 stack<int> st;
int a[] = {-1, -1};
for(int v=0;v<n;v++)
  if(indeg[v]!=graph[v].size()) {
    bool b = indeg[v] > graph[v].size();
   if (abs(((int)indeg[v])-((int)graph[v])
    .size())) > 1) return
   false;
if (a[b] != -1) return false;
   a[b] = v;
 int s = (a[0]!=-1 && a[1]!=-1 ? a[0] :
\rightarrow (a[0]==-1 && a[1]==-1 ? edges[0].u : -1));
 if(s==-1) return false;
  while(!st.empty() || !graph[s].empty()) {
  if (graph[s].empty()) {
    circuit.push_back(s); s = st.top();
   st.pop(): }
   int w = edges[*graph[s].begin()].v;
   graph[s].erase(graph[s].begin());
```

st.push(s); s = w;

```
Nicomachi's Theorem states 1^3 + 2^3 + ... + \frac{\text{#define FOR}(i, n) for (int }{\text{i = 0; i < n; i++)}} n^3 = (1+2+...+n)^2 and is equivalent to int n = m \text{ size}(1).
                                                                      const ll inf = 1LL << 62:
                                                                      int n = m.size();
FOR(i,n) m[i][i] = min(m[i][i], OLL);
FOR(k,n) FOR(i,n) FOR(j,n) if (m[i][k] != inf
                                                                        auto newDist = max(m[i][k] + m[k][j], -inf);
                                                                        m[i][j] = min(m[i][j], newDist);
                                                                       FOR(k,n) if (m[k][k] < 0) FOR(i,n) FOR(j,n)
                                                                       if (m[i][k] != inf && m[k][j] != inf)
                                                                       \hookrightarrow m[i][i] = -inf;
```

```
returns vector of edges in the mst
 // graph[i] = vector of edges incident to
     vertex i
    places total weight of the mst in Stotal
 // if returned vector has size != n-1, there is
vector<edge> mst(vector<vector<edge>> graph.
 → ll &total) {
total = 0;
priority_queue<edge, vector<edge>,

    greater<edge>> pq;
vector<edge> MST;

 bitset<20001> marked; // change size as needed
for (edge ep : graph[0]) pq.push(ep);
while(MST.size()!=graph.size()-1 &&
    pq.size()!=0) {
  edge e = pq.top(); pq.pop();
 int u = e.u, v = e.v, w = e.w;
if(marked[u] && marked[v]) continue;
else if(marked[u]) swap(u, v);
  for(edge ep : graph[u]) pq.push(ep);
  marked[u] = 1;
MST.push_back(e);
  total += e.w:
 return MST:
```

Union Find

```
int uf find(subset* s, int i) {
  if (s[i].p != i) s[i].p = uf_find(s, s[i].p);
 return s[i].p;
void uf_union(subset* s, int x, int y) {
int xp = uf_find(s, x), yp = uf_find(s, y);
if (s[xp].rank > s[yp].rank) s[yp].p = xp;
else if (s[xp].rank < s[yp].rank) s[xp].p =</pre>
 \rightarrow yp;
else { s[yp].p = xp; s[xp].rank++; }
```

2D Grid Shortcut

```
#define inbound(x,n) (0 <= x + x < n)
#define fordir(x,y,n,m) for(auto[dx,dy]:dir)if
    (inbound(x+dx,n)\otimes Sinbound(y+dy,m))
const pair<int,int> dir[] =
\rightarrow \{\{1,0\},\{0,1\},\{-1,0\},\{0,-1\}\};
```

2D Geometry

```
#define point complex double>
#define EPS 0.0000001
#define sq(a) ((a)*(a))
#define c\bar{b}(a) ((a)*(a)*(a))
double dot(point a, point b) { return
→ real(coni(a)*b): }
```

```
double cross(point a, point b) { return

    imag(conj(a)*b); }

struct line { point a, b; };
struct circle { point c; double r; };
struct segment { point a, b; };
struct triangle { point a, b, c; };
struct rectangle { point tl, br; };
struct convex_polygon {
 vector<point points;
 convex_polygon(vector<point> points) :
    points(points) {}
  convex polygon(triangle a) {
  points.push_back(a.a); points.push_back(a.b);
    points.push_back(a.c);
 convex_polygon(rectangle a) {
  points.push_back(a.tl);
     points.push_back({real(a.tl),
    imag(a.br)});
  points.push_back(a.br);
     points.push back({real(a.br),
     imag(a.tl)});
struct polygon {
 vector point points;
 polygon(vector<point> points) :
 → points(points) {}
 polygon(triangle a) {
  points.push_back(a.a); points.push_back(a.b);
  → points.push back(a.c):
 polygon(rectangle a) {
  points.push_back(a.tl);
    points.push_back({real(a.tl),
    imag(a.br)});
  points.push_back(a.br);
     points.push_back({real(a.br),
    imag(a.tl)}):
 polygon(convex_polygon a) {
  for (point v : a.points)
   points.push_back(v);
// triangle methods
double area heron(double a, double b, double
 \hookrightarrow c) {
if (a < b) swap(a, b);
 if (a < c) swap(a, c);
 if (b < c) swap(b, c):
 if (a > b + c) return -1;
 return sgrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)
// segment methods
double lengthsq(segment a) { return
     sq(real(a.a) - real(a.b)) + sq(imag(a.a)
\stackrel{\hookrightarrow}{\hookrightarrow} imag(a.b)); }
double length(segment a) { return

    sqrt(lengthsq(a)); }

// circle methods
double circumference(circle a) { return 2 * a.r.
double area(circle a) { return sq(a.r) * M PI;
| → }
|// rectangle methods
double width(rectangle a) { return
```

abs(real(a.br) - real(a.tl)); }

double height (rectangle a) { return

→ abs(imag(a.br) - real(a.tl)): }

double diagonal (rectangle a) { return

sgrt(sg(width(a)) + sg(height(a))); }

```
|double area(rectangle a) { return width(a) *
 → height(a); }
double perimeter(rectangle a) { return 2 *

→ (width(a) + height(a)); }

// check if `a` fit's inside `b`

// swap equalities to exclude tight fits
bool doesfitInside(rectangle a, rectangle b) {
 int x = width(a), w = width(b), y = height(a),
 \hookrightarrow h = height(b);
 if (x > y) swap(x, y);
if (w > h) swap(w, h);
 if (w < x) return false;
 if (y <= h) return true;
 double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
 return sq(a) \le sq(b) + sq(c):
 // polygon methods
 // negative area = CCW, positive = CW
double area(polygon a) {
  double area = 0.0; int n = a.points.size(); for (int i = 0, j = 1; i < n; i++, j = (j +
→ 1) % n)
area += (real(a.points[j]-a.points[i]))*

    (imag(a.points[j]+a.points[i]));
  return area / 2.0:
.
// get both unsigned area and centroid
pair < double, point > area centroid (polygon a) {
 int n = a.points.size();
 double area = 0:
 point c(0, 0);
 for (int i = n - 1, j = 0; j < n; i = j++) {
  double v = cross(a.points[i], a.points[j])
 \rightarrow area += v:
  c += (a.points[i] + a.points[j]) * (v / 3);
 c /= area;
 return {area, c};
Intersection
// -1 coincide, 0 parallel, 1 intersection
int intersection(line a, line b, point& p) {
 if (abs(cross(a.b - a.a. b.b - b.a)) > EPS) {
 p = cross(b.a - a.a, b.b - a.b) / cross(a.b)
 \rightarrow - a.a, b.b - b.a) * (b - a) + a;
  return 1:
 if (abs(cross(a.b - a.a, a.b - b.a)) > EPS)
→ return 0;
 return -1;
 // area of intersection
double intersection(circle a, circle b) {
 double d = abs(a.c - b.c);
if (d <= b.r - a.r) return area(a);
if (d <= a.r - b.r) return area(b);</pre>
 if (d >= a.r + b.r) return 0;
double alpha = acos((sq(a.r) + sq(d) -
 \rightarrow sq(b.r)) / (2 * a.r * d));
 double beta = acos((sq(b.r) + sq(d) - sq(a.r))
 \rightarrow / (2 * b.r * d));
 return sq(a.r) * (alpha - 0.5 * sin(2 *
    alpha) + sq(b.r) * (beta - 0.5 * sin(2 *
|}
|// -1 outside, 0 inside, 1 tangent, 2
intersection intersection (circle a, circle b,
vector<point>& inter) {
double d2 = norm(b.c - a.c), rS = a.r + b.r,
 \rightarrow rD = a.r - b.r;
 if (d2 > sq(rS)) return -1;
 if (d2 < sq(rD)) return 0;
 double ca = 0.5 * (1 + rS * rD / d2):
 point z = point(ca, sqrt(sq(a.r) / d2
 \rightarrow sq(ca));
```

```
|point3d cross(point3d a, point3d b) { return
                                                                                                           if (n == 1) while(1);
if (n == 2) while(1) a.push_back(g());
 inter.push_back(a.c + (b.c - a.c) * z);
if (abs(imag(z)) > EPS) inter.push_back(a.c +
                                                         \{a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,
                                                                                                           if (n == 3) while(1) putchar_unlocked('a');
\rightarrow (b.c - a.c) * conj(z));
                                                        a.x*b.y - a.y*b.x; }
                                                                                                           if (n == 4) assert(0):
return inter.size();
                                                     struct line3d { point3d a, b; };
                                                                                                           if (n == 5) 0 / 0;
if (n == 6) *(int*)(0) = 0;
                                                     struct plane { double a, b, c, d; } // a*x +
// points of intersection
                                                     \Rightarrow b*y + c*z + d = 0
                                                                                                           return n + judge(n + 1);
vector<point> intersection(line a, circle c) {
                                                     struct sphere { point3d c; double r; };
vector < point > inter;
c.c -= a.a;
                                                     #define sq(a) ((a)*(a))
a.b -= a.a:
                                                     #define c\bar{b}(a) ((a)*(a)*(a))
                                                                                                          GCC Builtin Docs
 point m = a.b * real(c.c / a.b);
                                                     double surface(circle a) { return 4 * sq(a.r)
                                                                                                          // 128-bit integer
 double d2 = norm(m - c.c);
                                                                                                          __int128 a;
unsigned __int128 b;
                                                     \rightarrow M PI: }
 if (d2 > sq(c.r)) return 0;
                                                     double volume(circle a) { return 4.0/3.0 *
                                                                                                          // 128-bit float
 double l = sqrt((sq(c.r) - d2) / norm(a.b));
                                                     \hookrightarrow cb(a.r) * M PI; }
 inter.push_back(a.a + m + 1 * a.b);
                                                                                                           // minor improvements over long double
if (abs(1) > EPS) inter.push_back(a.a + m - 1 10 Optimization
                                                                                                            _float128 c;
                                                                                                          7/ log2 floor
\rightarrow * a.b):
                                                                                                          __lg(n);
return inter:
                                                                                                          // number of 1 bits
                                                     // SameNumberOfOneBits, next permutation
// area of intersection
                                                                                                          // can add ll like popcountll for long longs
                                                     int snoob(int a) {
double intersection (rectangle a, rectangle b) { int b = a & -a, c = a + b;
                                                                                                          __builtin_popcount(n);
// number of trailing zeroes
                                                     return c | ((a ^ c) >> 2) / b;
double x1 = max(real(a.tl), real(b.tl)), y1 =

→ max(imag(a.tl), imag(b.tl));
                                                                                                            _builtin_ctz(n);
                                                      // example usage
                                                                                                          // number of leading zeroes
double x2 = min(real(a.br), real(b.br)), y2 =
                                                     int main() {
    char l1[] = {'1', '2', '3', '4', char l2[] = {'a', 'b', 'c', 'd'};
    int d1 = 5, d2 = 4;
                                                                                                          _builtin_clz(n);
// 1-indexed least significant 1 bit

→ min(imag(a.br), imag(b.br));
                                                                                          '5'};
return (x2 \le x1 \mid | y2 \le y1) ? 0 :
                                                                                                            builtin ffs(n);
    (x2-x1)*(y2-y1);
                                                                                                          7/ parity of number
                                                      // prints 12345abcd, 1234a5bcd, ...
                                                                                                           _builtin_parity(n);
                                                      int^* min = (1 << d1) - 1, max = min << d2;
                                                     for (int i = min; i <= max; i = snoob(i)) {
   int p1 = 0, p2 = 0, v = i;
   while (p1 < d1 || p2 < d2) {
Convex Hull
                                                                                                          Limits
bool cmp(point a, point b) {
if (abs(real(a) - real(b)) > EPS) return
                                                                                                                             \pm 2147483647 \mid \pm 2^{31} - 1 \mid 10^9
                                                                                                          _{
m int}
                                                        cout << ((v & 1) ? l1[p1++] : l2[p2++]);

→ real(a) < real(b);
if (abs(imag(a) - imag(b)) > EPS) return
                                                                                                                               4294967295
                                                                                                          uint
                                                       v /= 2;
                                                                                                                 \pm 9223372036854775807 | \pm \overline{2}^{63} - \overline{1}|\overline{10}^{18}

→ imag(a) < imag(b);
</p>
                                                       cout << '\n';
return false:
                                                                                                          ull
                                                                                                                 18446744073709551615
                                                                                                          |i128| \pm 170141183460469231... | \pm 2^{127} - 1 | 10^{38}
convex_polygon convexhull(polygon a) {
                                                                                                          |u128| 340282366920938463... | 2^{128} - \bar{1}|\bar{1}0^{38}|
sort(a.points.begin(), a.points.end(), cmp);
                                                     Powers
 vector<point> lower, upper;
                                                     bool isPowerOf2(ll a) {
                                                                                                          Complexity classes input size (per second):
for (int i = 0; i < a.points.size(); i++) {</pre>
                                                      return a > 0 && !(a & a-1);
                                                                                                          O(n^n) or O(n!)
 while (lower.size() >= 2 &&
    cross(lower.back() - lower[lower.size() -
                                                     bool isPowerOf3(11 a) {
                                                                                                          |O(2^n)|
                                                      return a>0&&! (12157665459056928801ull%a);
   2], a.points[i] - lower.back()) < EPS)
                                                                                                          O(n^3)
   lower.pop_back();
                                                     bool isPower(ll a, ll b) {
  double x = log(a) / log(b);
                                                                                                          O(n^2)
  while (upper.size() >= 2 &&
    cross(upper.back() - upper[upper.size()
                                                      return abs(x-round(x)) < 0.00000000001;
                                                                                                          O(n\sqrt{n})
   2], a.points[i] - upper.back()) > -EPS)
                                                                                                          O(n \log n)
   upper.pop_back();
                                                     11 Additional
 lower.push_back(a.points[i]);
                                                                                                          O(n)
 upper.push_back(a.points[i]);
                                                     {f Judge\ Speed}
                                                        kattis: 0.50s
 lower.insert(lower.end(), upper.rbegin() + 1,
                                                        codeforces: 0.421s
→ upper.rend());
                                                     // atcoder: 0.455s
return convex_polygon(lower);
                                                     #include <bits/stdc++.h>
                                                     using namespace std;
                                                     int v = 1e9/2, p = 1;
    3D Geometry
                                                     int main() {
                                                      for (int i = 1; i <= v; i++) p *= i;
struct point3d {
                                                      cout << p;
double x, y, z;
point3d operator+(point3d a) const { return
\rightarrow {x+a.x, y+a.y, z+a.z}; }
                                                     Judge Pre-Contest Checks
point3d operator*(double a) const { return
                                                         int128 and float128 support?
\hookrightarrow {x*a, v*a, z*a}: }
point3d operator-() const { return {-x, -y,
                                                      does extra or missing whitespace cause WA?
                                                      documentation up to date?
point3d operator-(point3d a) const { return
                                                      printer usage available and functional?
\rightarrow *this + -a; }
point3d operator/(double a) const { return
\rightarrow *this * (1/a); }
                                                      // each case tests a different fail condition
double norm() { return x*x + y*y + z*z; }
                                                     // try them before contests to see error codes
 double abs() { return sqrt(norm()); }
                                                     struct g { int arr[1000000]; g(){}};
                                                     vector<g> a;
point3d normalize() { return *this /
                                                     // O=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE

    this->abs(); }

                                                    ⇒ 6=SIGSEGV 7=recursive MLE
int judge(int n) {
double dot(point3d a, point3d b) { return
                                                      if (n == 0) \text{ exit}(0);
\rightarrow a.x*b.x + a.y*b.y + a.z*b.z; }
```

 $\overline{2}^{32} - 1|10^9$

 $2^{64} - 1|10^{19}$

n < 10

n < 30

n < 1000

 $n < 10^{6}$

 $n \le 10^{7}$

 $n < 10^9$

n < 30000