```
General
                              7 Graphs
    Algorithms
                              8 2D Geometry
    Structures
                              9 3D Geometry
    Strings
                              10 Optimization
    Greedy
                              11 Additional
    Math
     General
g++ -g -02 -std=gnu++17 -static prog.cpp
./a.exe
run.sh
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp for i i *.in; do
 f=${i%.in}
 ./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
done
Header
// use better compiler options
#pragma GCC optimize("Ofast","unroll-loops")
#pragma GCC target("avx2,fma")
// include everything
 #include <bits/stdc++.h>
#include <bits/extc++.h>
#include <sys/resource.h>
// namespaces
using namespace std;
using namespace __gnu_cxx; // rope
using namespace __gnu_pbds; // tree/trie
// common defines
#define fastio

→ ios base::sync with stdio(0);cin.tie(0);
#define nostacklim rlimit RZ; getrlimit(3,&RZ
    ):RZ.rlim cur=-1:setrlimit(3.&RZ):
#define DEBUG(v) cerr<< LINE <<": "<<#v<<" =
\Rightarrow "<<v<<'\n'; #define TIMER
**define il28 unsigned il28

#define ull unsigned ll
#define il28 unsigned ll
#define il28 unsigned il28
#define ld long double
// global variables
mt19937 rng((uint32_t)chrono::steady

    clock::now().time since epoch().count());

Fast IO
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
void read(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
  n = n * 10 + c - 0';
void read(int& n) {
   char c; n = 0; int s = 1;
   if ((c=getchar_unlocked())=='-') s = -1;
 else n = c - ^{\circ};
while ((c=getchar_unlocked())!=' '&&c!='\n')
 n = n * 10 + c - 0;

n *= s;
void read(ld& n) {
 char c; n = 0;
ld m = 0, o = 1; bool d = false; int s = 1;
if ((c=getchar_unlocked())=='-') s = -1;
 else if (c == .'.') d = true;
else n = c - '0';
 while ((c=getchar_unlocked())!=' '&&c!='\n') {|}
  if (c == '.') d = true;
else if (d) { m=m*10+c-'0'; o*=0.1; }
```

```
else n = n * 10 + c - '0':
 n = s * (n + m * o):
void read(double& n) {
 ld m; read(m); n = m;
void read(float& n) {
 ld m: read(m): n = m:
void read(string& s) {
 char c; s = "
 while((c=getchar unlocked())!=' '&&c!='\n')
bool readline(string& s) {
 char c; s = "";
while(c=getchar unlocked()) {
 if (c == '\n') return true;
if (c == EOF) return false;
s += c;
 return false;
void print(unsigned int n) {
 if (n / 10) print(n / 10);
 putchar_unlocked(n % 10 + '0');
void print(int n) {
 if (n < 0) { putchar_unlocked('-'); n*=-1; }
 print((unsigned int)n);
Common Structs
   n-dimension vectors
// Vec<2, int> v(n, m) = arr[n][m]

// Vec<2, int> v(n, m, -1) default init -1

template<int D, typename T>

struct Vec : public vector<Vec<D-1, T>> {
  template<typename... Args>
  Vec(int n=0, Args... args) : vector<Vec<D-1,
 \rightarrow T>>(n. Vec<D-1. T>(args...)) {}
template<typename T>
struct Vec<1, T> : public vector<T> {
  Vec(int n=0, T val=T()) : vector<T>(n, val)
    {}
     Algorithms
Binary Search
// search for k in [p,n)
template<typename T>
int binsearch(T x[], int k, int n, int p = 0) {
     for (int i = n; i >= 1; i /= 2)
          while (p+i < n \&\& x[p+i] <= k) p += i;
     return p; \frac{1}{bool}: x[p] == k;
Min/Max Subarray
   max - compare = a < b, reset = a < 0
 \frac{1}{2}/ min - compare = a > b. reset = a > 0
// returns {sum, {start, end}}
pair<int, pair<int, int>>
     ContiguousSubarrav(int* a. int size.
    bool(*compare)(int, int).
 bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
 \rightarrow 0, s = 0;
 for (int i = 0; i < size; i++) {
...cur += a[i];
  if ((*compare)(best, cur)) { best = cur;
 > start = s; end = i; }
if ((*reset)(cur)) { cur = 0: s = i + 1: }
 return {best, {start, end}};
Quickselect
```

```
int partition(int arr[], int 1, int r)
  int x = arr[r], i = 1;
 for (int j = 1; j <= r - 1; j++)
. if (arr[j] <= x)
  swap(arr[i++], arr[j]);
 swap(arr[i], arr[r]);
 return i:
// find k'th smallest element in unsorted array
→ only if all distinct
int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - 1 + 1)) return QSNE;
swap(arr[1 + rng() % (r-1+1)], arr[r]);
  int pos = partition(arr, 1, r);
 if (pos-l==k-1) return arr[pos];
 if (pos-1>k-1) return qselect(arr,1,pos-1,k);
 return qselect(arr, pos+1, r, k-pos+1-1);
// TODO: compare against std::nth_element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted
→ on both axis
pair<int, int> saddleback_search(int** arr, int
 \hookrightarrow x, int y, int v) {
 int i = x-1, j = 0;
 while (i >= 0 && j < y) {
   if (arr[i][j] == v) return {i, j};
  (arr[i][i] > v)? i--: j++;
 return {-1, -1};
Ternary Search
// < max, > min, or any other unimodal func #define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a, int b, int (*f)(int)) {
 while (b-a > 4) {
    int m = (a+b)/2;
    if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1:
 for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
 return á;
#define TERNPREC 0.000001
double ternsearch (double a, double b, double
 \rightarrow (*f)(double)) {
while (b-a > TERNPREC * 4) {
  double m = (a+b)/2;
if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
  else b = m + TERNPREC;
 for (double i = a + TERNPREC: i <= b: i +=

→ TERNPREC

      if (TERNCOMP((*f)(a), (*f)(i)))
   .a = i:
 return a;
Golden Section Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)</pre>
double goldsection(double a, double b, double
 double r = (sqrt(5)-1)/2, eps = 1e-7;

double x1 = b - r*(b-a), x2 = a + r*(b-a);

double f1 = f(x1), f2 = f(x2);
  while (b-a > eps)
  while (b-a > eps)

if (TERNCOMP(f2,f1)) {

. b = x2; x2 = x1; f2 = f1;

. x1 = b - r*(b-a); f1 = f(x1);
   a = x1; x1 = x2; f1 = f2; x2 = a + r*(b-a): f2 = f(x2):
 return a:
```

```
3 Structures
```

```
Fenwick Tree
// Fenwick tree, array of cumulative sums
\hookrightarrow O(\log n) updates, O(\log n) gets
struct Fenwick { int n: ll* tree:
 void update(int i, int val) {
 .++i;
while (i <= n) {
  tree[i] += val;</pre>
  i += i & (-i);
 Fenwick(int size) {
 n = size;
 tree = new ll[n+1];
for (int i = 1; i <= n; i++)
  tree[i] = 0;
 Fenwick(int* arr, int size) : Fenwick(size) {
 for (int i = 0; i < n; i++)
...update(i, arr[i]);
 ~Fenwick() { delete[] tree; }
 ll operator∏(int i) {
 if (i < 0 || i > n) return 0;
 while (i>0)
  sum += tree[i];
i -= i & (-i);
 return sum;
ll getRange(int a, int b) { return
    operator[](b) - operator[](a-1); }
Hashtable
```



```
Rope
                                                         // print things with prefix "1"
                                                         auto range = trie.prefix_range("1");
// O(\log n) insert, delete, concatenate
                                                         for (auto it = range.first; it !=
int main() {
 // generate rope
                                                         → range.second: it++)
 rope<int> v;
                                                          cout << *it << '
 for (int i = 0; i < 100; i++)
.v.push_back(i);
                                                        Wavelet Tree
 // move range to front
                                                        using iter = vector<int>::iterator;
 rope<int> copy = v.substr(10, 10);
v.erase(10, 10);
                                                        struct WaveletTree {
   Vec<2, int> C; int s;
 v.insert(copy.mutable_begin(), copy);
                                                          // sigma = highest value + 1
                                                         WaveletTree(vector<int>& a. int sigma) :
 // print elements of rope
for (auto it : v) cout << it << "";
                                                            s(sigma), C(sigma*2, 0) {
                                                          build(a.begin(), a.end(), 0, s-1, 1);
                                                         void build(iter b. iter e. int L. int U. int
Segment Tree
                                                          u) {
if (L == U) return
//max(a,b), min(a,b), a+b, a*b, qcd(a,b), a*b
struct SegmentTree {
                                                           int M = (L+U)/2;
 typedef int T;
                                                           C[u].reserve(e-b+1); C[u].push back(0);
 static constexpr T UNIT = INT_MIN;
                                                          for (auto it = b; it != e; ++it)
  C[u].push_back(C[u].back() + (*it<=M));</pre>
 T f(T a, T b) {
 if (a == UNIT) return b;
if (b == UNIT) return a;
                                                           auto p = stable_partition(b, e, [=](int
                                                            i){return i<=M;});
  return max(a,b);
                                                          build(b, p, L, M, u*2);
 int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
                                                          build(p, e, M+1, U, u*2+1);
                                                          // number of occurrences of x in [0,i)
\rightarrow n(n) {}
                                                         int rank(int x, int i) {
   int L = 0, U = s-1, u = 1, M, r;
   while (L != U) {
 SegmentTree(vector<T> arr)

    SegmentTree(arr.size()) {

 for (int i=0:i<arr.size():i++)
                                                           M = (L+U)/2;
r = C[u][i]; u*=2;

    update(i,arr[i]);

                                                           if (x <= M) i = r, U = M;
else i -= r, L = M+1, ++u;
 void update(int pos, T val) {
  for (s[pos += n] = val; pos /= 2;)
   s[pos] = f(s[pos * 2], s[pos*2+1]);
                                                          return i:
                                                          ^{\prime\prime} number of occurences of x in [l,r)
 T query(int b, int e) { // query [b, e)
                                                         int count(int x, int 1, int r) {
  return rank(x, r) - rank(x, 1);
  Tra = UNIT, rb = UNIT;
  for (b+=n, e+=n; b<=); b/=2, e/=2) {
    if (b % 2) ra = f(ra, s[b++]);
    if (e % 2) rb = f(s[--e], rb);
                                                         // kth smallest in [l, r)
int kth(int k, int l, int r) const {
int L = 0, U = s-1, u = 1, M, ri, rj;
  return f(ra, rb):
                                                          while (L != U) {
   M = (L+U)/2;
 T get(int p) { return query(p, p+1); }
                                                           ri = C[u][1]; rj = C[u][r]; u*=2;
                                                           if (k \le rj-ri)^{n}l = ri, r = rj, U = M;
Sparse Table
                                                           else k -= řj-rí, l -= ŕi, r -= ŕj,
template < class T> struct SparseTable {
                                                           L = M+1, ++u;
 vector<vector<T>> m;
                                                           return U:
 SparseTable(vector<T> arr) {
  m.push back(arr);
  for (int k = 1: (1<<(k)) <= size(arr): k++)
                                                         // # elements between [x,y] in [l, r)
                                                         mutable int L, U;
  m.push_back(vector<T>(size(arr)-(1<(k)+1));
                                                         int range(int x, int y, int 1, int r) const {
  for (int i = 0; i < size(arr)-(1<<k)+1; i
                                                          if (y < x \text{ or } r <= 1) return 0;
                                                          L = x; U = y;
 [k][i] = min(m[k-1][i],
                                                          return range(1, r, 0, s-1, 1);
\rightarrow m[k-1][i+(1<<(k-1))]:
}
// min of range [l,r]
                                                         int range(int 1, int r, int x, int y, int u)
                                                         → const {
                                                          if (y < L or U < x) return 0;
if (L <= x and y <= U) return r-l;
T query(int 1, int r) {
 int k = _-lg(r-l+1);
                                                          int M = (x+y)/2, ri = C[u][1], rj = C[u][r];
  return \min(m[k][1], m[k][r-(1<< k)+1]);
                                                          return range(ri, rj, x, M, u*2) + range(1-ri
}
};
                                                            r-rj, M+1, y, u*2+1);
                                                          ^{\prime}// # elements <= x in [l, r]
                                                         int lte(int x, int l, int r) {
  return range(INT_MIN, x, l, r);
typedef trie<string, null_type,

→ trie_string_access_traits<>,

 pat_trie_tag, trie_prefix_search_node_update>

→ trie_type;

int main() {
                                                             Strings
 // generate trie
 trie_type trie;
                                                        Aho Corasick
 for (int i = 0; i < 20; i++)
                                                           range of alphabet for automata to consider
 trie.insert(to string(i)); // true if new,
                                                           MAXC = 26, OFFC = 'a' if only lowercase
\hookrightarrow false if old
```

```
|const int MAXC = 256:
const int OFFC = 0:
struct aho_corasick {
  set<pair<int, int>> out;
  int fail; vector<int> go;
  state() : fail(-1), go(MAXC, -1) {}
 vector<state> s;
  int id = 0;
 aho_corasick(string* arr, int size) : s(1) {
  for (int i = 0; i < size; i++) {
   for (int c : arr[i]) {
   if (s[cur].go[c-OFFC] == -1) {
      s[cur].go[c-OFFC] = s.size();
      s.push back(state());
     cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++});
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
    s[0].go[c] = 0;
  queue int> sq;
for (int c = 0; c < MAXC; c++) {
   if (s[0].so[c] != 0) {
      ...s[s[0].so[c]].fail = 0;</pre>
    sq.push(s[0].go[c]);
  while (sq.size()) {
   int e = sq.front(); sq.pop();
   for (int c = 0; c < MAXC; c++) {
   if (s[e].go[c] != -1) {
      int failure = s[e].fail;
while (s[failure].go[c] == -1)
        failure = s[failure].fail;
      failure = s[failure].go[c];
      s[s[e].go[c]].fail = failure;
      for (auto length : s[failure].out)
s[s[e].go[c]].out.insert(length);
      sq.push(s[e].go[c]);
 // list of {start pos, pattern id}
  vector<pair<int, int>> search(string text)
  vector<pair<int, int>> toret;
  int cur = 0;
  for (int i = 0; i < text.size(); i++) {
  while (s[cur].go[text[i]-OFFC] == -1)
    cur = s[cur].fail;
cur = s[cur].go[text[i]-OFFC];
    if (s[cur].out.size())
    for (auto end : s[cur].out)
. toret.push_back({i - end.first + 1,
     end.second):
  return toret:
Boyer Moore
struct defint { int i = -1; };
vector<int> boyermoore(string txt, string pat)
 vector<int> toret; unordered_map<char, defint>string lcp(string* arr, int n, bool sorted =
 → badchar:
 int m = pat.size(), n = txt.size();
 for (int i = 0; i < m; i++) badchar[pat[i]].i
 \rightarrow = i;
int s = 0:
 while (s \leq n - m) {
  int j = m - 1:
```

while $(j \ge 0 \&\& pat[j] == txt[s + j]) j--;$

.if (j < 0) {

```
..toret.push_back(s);
   s += (s + m < n) ? m - badchar[txt[s +
   m]].i : 1:
 .} else
   s += \max(1, i - badchar[txt[s + i]].i):
 return toret:
English Conversion
const string ones[] = {"", "one", "two",
    "three", "four", "five", "six", "seven", "eight", "nine";
const string teens[] ={"ten", "eleven",
    "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen";
const string tens[] = {"twenty", "thirty",
    "forty", "fifty", "sixty", "seventy",

    "eighty", "ninety"};
const string mags[] = {"thousand", "million",
     "billion", "trillion", "quadrillion", "quintillion", "sextillion",
    "septillion"};
string convert(int num, int carry) {
 if (num < 0) return "negative " +

    convert(-num, 0):

     (num < 10) return ones[num];
(num < 20) return teens[num % 10]
 if (num < 100) return tens[(num / 10) - 2] +
     (num\%10==0?"":"") + ones[num\%10]:
 if (num < 1000) return ones[num / 100]
     (num/100==0?"":" ") + "hundred" + (num%100==0?"":" ") + convert(num % 100,
    0);
 return convert(num / 1000, carry + 1) + " " +
    mags[carry] + " " + convert(num % 1000.
    0):
string convert(int num) {
 return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
   vector<int> toret;
 int m = txt.length(), n = pat.length();
 int next[n + 1];
 for (int i = 0; i < n + 1; i++)
  next[i] = 0;
 for (int i = 1; i < n; i++) {
  int j = next[i + 1];
  while (j > 0 && pat[j] != pat[i])
   j = next[j];
  if (j > 0 | pat[j] == pat[i])
   next[i + 1] = j + 1;
 for (int i = 0, j = 0; i < m; i++) {
  if (txt[i] == pat[j]) {
  if (++j == n)
    toret.push_back(i - j + 1);
  } else if (j > 0) {
...j = next[j];
 return toret;
Longest Common Prefix (array)
 // longest common prefix of strings in array
 → false) {
idise; l
if (n == 0) return "";
if (!sorted) sort(arr, arr + n);
string r = ""; int v = 0;
 while (v < arr[0].length() && arr[0][v] ==

    arr[n-1][v])
    r += arr[0][v++];

 return r:
```

```
| for (auto c : s) v = (c - 'a' + 1) + v *
Longest Common Subsequence
                                                           → HASHER;
string lcs(string a, string b) {
 int m = a.length(), n = b.length();
                                                          const int MAXN = 1000001;
 int L[m+1][n+1];
 int L[m+1] [ Lm+1];
for (int i = 0; i <= m; i++) {
    for (int j = 0; j <= n; j++) {
        if (i == 0 || j == 0) L[i][j] = 0;
        else if (a[i-1] == b[j-1]) L[i][j] =</pre>
                                                          ull base [MAXN] = \{1\};
                                                          void genBase(int n) {
                                                           for (int i = 1; i \le n; i++)

base[i] = base[i-1] * HASHER:
 \hookrightarrow L[i-1][j-1]+1;
                                                          struct advHash {
   else L[i][j] = \max(L[i-1][j], L[i][j-1]);
                                                           ull v, 1; vector<ull> wip;
                                                           advHash(string& s): v(0) {
 }
// return L[m][n]; // length of lcs
                                                            wip = vector<ull>(s.length()+1);\
                                                             wip[0] = 0;
 string out = "";
                                                             for (int i = 0; i < s.length(); i++)
wip[i+1] = (s[i] - 'a' + 1) + wip[i] *
 int i = m - 1, j = n - 1;
 while (i \ge 0) \& j \ge 0 \{ if (a[i] = b[i]) \}
                                                              HASHER;
                                                            1 = s.length(); v = wip[1];
   out = a[i--] + out;
                                                           ull del(int pos, int len) {
   return v - wip[pos+len]*base[1-pos-len] +
  else if (L[i][j+1] > L[i+1][j]) i--;
                                                              wip[pos]*base[1-pos-len];
  .else j--;
                                                           ull substr(int pos, int len) {
 return out;
                                                            return del(pos+len, (1-pos-len)) -
                                                              wip[pos]*base[len]:
Longest Common Substring
                                                           ull replace(int pos, char c) {
// l is array of palindrome length at that
int manacher(string s, int* 1) {
                                                            return v - wip[pos+1]*base[l-pos-1] + ((c -
                                                                'a' + 1) + wip[pos] *
 int n = s.length() * 2;
                                                              HASHER) *base[1-pos-1];
 for (int i = 0, j = 0, k; i < n; i += k, j =
\rightarrow max(j-k, 0)) {
                                                           ull replace(int pos, string s) {
                                                            // can't increase total string size
 while (i \ge j \&\& i + j + 1 < n \&\& s[(i-j)/2]]
\Rightarrow == s[(i+j+1)/2]) i++;
                                                               wip[pos+s.size()]*base[l-pos-s.size()], c
 1[i] = j;
                                                               wip[pos];
 for (k = 1; i >= k \&\& j >= k \&\& l[i-k] !=
                                                            for (int i = 0; i < s.size(); i++)
    c = (s[i]-'a'+1) + c * HASHER;
    return r + c * base[l-pos-s.size()];</pre>
 \begin{array}{ll} \hookrightarrow & j^-k; k^{++}) \\ \dots & 1[i^+k] = \min(1[i^-k], j^-k); \end{array}
 return *max element(1, 1 + n);
                                                          Subsequence Count
Cyclic Rotation (Lyndon)
                                                           // "banana", "ban" >> 3 (ban, ba..n, b..an)
// simple strings = smaller than its nontrivial
                                                          ull subsequences(string body, string subs) {
 int m = subs.length(), n = body.length();
// lyndon factorization = simple strings
                                                           if (m > n) return 0;

ull** arr = new ull*[m+1];

for (int i = 0; i <= m; i++) arr[i] = new
\hookrightarrow factorized
// "abaaba" -> "ab", "aab", "a"
vector<string> duval(string s) {
                                                           int n = s.length();
 vector<string> lyndon;
 for (int i = 0; i < n;) {
    int j = i+1, k = i;
    for (; j < n && s[k] <= s[j]; j++)
    if (s[k] < s[j]) k = i;
                                                           for (int i = 1; i <= m; i++)

for (int j = 1; j <= n; j++)

arr[i][j] = arr[i][j-1] + ((body[j-1] ==
                                                              subs[i-1])? arr[i-1][j-1] : 0);
   else k++:
                                                           return arr[m][n]:
  for (; i <= k; i += j - k)
   lyndon.push_back(s.substr(i,j-k));
                                                          Suffix Array + LCP
 return lyndon;
                                                           struct SuffixArray {
}
// lexicographically smallest rotation
                                                           vector<int> sa, 1cp;
                                                           SuffixArray(string& s, int lim=256) {

int n = s.length() + 1, k = 0, a, b;
int minRotation(string s) {
 int n = s.length(); s += s;
                                                            vector<int> x(begin(s), end(s)+1), y(n),
 auto d = duval(s); int i = 0, a = 0;
                                                           \rightarrow ws(max(n, lim)), rank(n);
 while (a + d[i].length() < n) a +=
                                                            sa = lcp = y;
iota(begin(sa), end(sa), 0);

    d[i++].length();

 while (i && d[i] == d[i-1]) a -=
                                                             for (int j = 0, p = 0; p < n; j = max(1, j *

    d[i--].length();

                                                           \rightarrow 2), lim = p) {
return a;
                                                             p = j; iota(begin(y), end(y), n - j);
                                                             for (int i = 0; i < (n); i++)
if (sa[i] >= j)
Hashing
                                                                y[p++] = sa[i] -
#define HASHER 27
                                                              fill(begin(ws), end(ws), 0);
ull basicHash(string s) {
 ull v = 0;
                                                              for (int i = 0; i < (n); i++) ws[x[i]]++;
```

```
|...for (int i = 1; i < (lim); i++) ws[i] +=
 \rightarrow ws[i - 1];
   for (int \bar{i} = n; i--;) sa[--ws[x[v[i]]]] =
    v[i]:
  j]) ? p - 1 : p++;
  for (int i = 1; i < (n); i++) rank[sa[i]] =
 for (int i = 0, j; i < n - 1; lcp[rank[i++]]
  = k)
for (k && k--, j = sa[rank[i] - 1];
s[i + k] == s[j + k]; k++);
 // smallest cyclic shift
int cyclic() { return sa[0]; }
 // longest repeated substring
 pair<int,int> lrs() {
  int length = -1, index = -1;
  for (int i = 0; i < lcp.size(); i++) {
  if (lcp[i] > length) {
  length = lcp[i];
    index = sa[i];
  return {index,length};
 // count distinct substrings, excluding empty
 int distincts() {
  int n = sa.size() - 1, r = n - sa[0];
  for (int i = 1; i < lcp.size(); i++)
  r += (n - sa[i]) - lcp[i - 1];
 return r:
 }
// count repeated substrings. excluding empty
 int repeateds() {
  int r = 0;
for (int i = 1; i < lcp.size(); i++)
  r += \max(lcp[i] - lcp[i-1], 0);
Suffix Tree (Ukkonen's)
struct SuffixTree {
 // n = 2*len+10 or so
enum { N = 50010, ALPHA = 26 };
int toi(char c) { return c - 'a'; }
 t[N][ALPHA],1[N],r[N],p[N],s[N],v=0,q=0,m=2| (1 countries) intervalCover(double L, double R,
 string a;
 void ukkadd(int i, int c) { suff:
  if (r[v]<=q) {
  if (q==-1 || c==toi(a[q])) q++; else {
   l[m+1]=i; p[m+1]=m; l[m]=l[v]; r[m]=q;
   p[m]=p[v]; t[m][c]=m+1; t[m][toi(a[q])]=v;
  1[v]=q; p[v]=m; t[p[m]][toi(a[1[m]])]=m;
v=s[p[m]]; q=1[m];
   while (q < r[m]) = v = t[v][toi(a[q])];
    q+=r[v]-l[v]; }
   if (q==r[m]) s[m]=v; else s[m]=m+2;
   q=r[v]-(q-r[m]); m+=2; goto suff;
 SuffixTree(string a) : a(a) {
 fill(r,r+N,(int)(a).size());
                                                       ull* catalan = new ull[1000000];
 memset(s, 0, sizeof s);
memset(t, -1, sizeof t);
fill(t[1],t[1]+ALPHA,0);
s[0]=1;1[0]=1[1]=-1;r[0]=r[1]=p[0]=p[1]=0;
                                                      void genCatalan(int n, int mod) {
  catalan[0] = catalan[1] = 1;
  for (int i = 2; i <= n; i++) {</pre>
                                                        catalan[i] = 0;
  for(int i=0;i<a.size();i++)
```

→ ukkadd(i.toi(a[i])):

```
// Longest Common Substring between 2 strings
 // returns {length, offset from first string}
 pair<int, int> best;
int lcs(int node, int i1, int i2, int olen) {
    if (1[node] <= i1 && i1 < r[node]) return 1;
    if (1[node] <= i2 && i2 < r[node]) return 2;
 int mask=0,
 - len=node?olen+(r[node]-1[node]):0;
 for(int c=0; c<ALPHA; c++) if
  (t[node][c]!=-1)
mask |= lcs(t[node][c], i1, i2, len);
 if (mask==3)
  best=max(best,{len,r[node]-len});
 static pair<int, int> LCS(string s, string t)
 SuffixTree
 \rightarrow st(s+(char)('z'+1)+t+(char)('z'+2)):
 st.lcs(0, s.size(), s.size()+t.size()+1, 0);
return st.best;
String Utilities
void lowercase(string& s) {
transform(s.begin(), s.end(), s.begin(),
void uppercase(string& s) {
transform(s.begin(), s.end(), s.begin(),
   ::toupper);
void trim(string &s) {
s.erase(s.begin(),find_if_not(s.begin(),s
     .end(),[](int c){return
   isspace(c);}));
s.erase(find_if_not(s.rbegin(),s.rend(),[](int

    c){return isspace(c);}).base(),s.end());

vector<string> split(string& s, char token) {
    vector<string> v; stringstream ss(s);
    for (string e;getline(ss,e,token);)
        v.push_back(e);
  Greedy
Interval Cover
// L,R = interval [L,R], in = {{l,r}, index}
// does not handle case where L == R
→ vector<pair<pair<double,double,int>> in) {
    int i = 0; pair < double, int > pos = {L,-1};
   vector<int> a;
sort(begin(in), end(in));
    while (pos.first < R) {
        double cur = pos.first;
while (i < (int)in.size() &&</pre>
    in[i].first.first <= cur)</pre>
    max(pos.{in[i].first.second.in[i].second}).
        if (pos.first == cur) return {};
        a.push_back(pos.second);
    return a:
6 Math
Catalan Numbers
```

for $(int j = i - 1; j >= 0; j--) {$

```
catalan[i] += (catalan[j] * catalan[i-j-1]) Discrete Logarithm
    % mod;
                                                            int discretelog(int a. int b. int m) {
   if (catalan[i] >= mod)
                                                                 11 n = sqrt(m) + 1, an = 1;
    catalan[i] -= mod:
                                                                for (l1 i = 0; i < n; ++i)
an = (an * a) % m;
unordered_map<11, l1> vals;
\mathcal{Y}' // TODO: consider binomial coefficient method
                                                                 for (11 q = 0, cur = b; q \le n; q++) {
                                                                      vals[cur] = q;
                                                                      cur = (cur * a) \% m:
Combinatorics (nCr, nPr)
                                                                 for (ll p = 1, cur = 1; p <= n; p++) {
// can optimize by precomputing factorials, and
                                                                      cur = (cur * an) % m;
if (vals.count(cur)) {
 \hookrightarrow fact[n]/fact[n-r]
ull nPr(ull n, ull r) {
                                                                           int ans = n * p - vals[cur];
                                                                           return ans;
 for (ull i = n-r+1: i <= n: i++)
 return v;
                                                                 return -1;
ull nPr(ull n, ull r, ull m) {
   ull v = 1;
                                                            Euler Phi / Totient
 for (ull i = n-r+1: i <= n: i++)
 v = (v * i) % m;
                                                            int phi(int n) {
                                                             int'r = n;
                                                            for (int i = 2; i * i <= n; i++) {
   if (n % i == 0) r -= r / i;
   while (n % i == 0) n /= i;
úll nCr(ull n. ull r) {
 long double v = 1;
 for (ull i = 1: i <= r: i++)
 v = v * (n-r+i) /i;
                                                             if (n > 1) r = r / n;
 return (ull)(v + 0.001);
                                                             return r:
// requires modulo math
                                                            #define n 100000
// can optimize by precomputing mfac and
                                                            ll phi[n+1];
                                                            void computeTotient() {
\hookrightarrow minv-mfac
                                                             for (int i=1; i<=n; i++) phi[i] = i;
ull nCr(ull n, ull r, ull m) {
 return mfac(n, m) * minv(mfac(k, m), m) % m *
                                                             for (int p=2; p<=n; p++) {

    minv(mfac(n-k, m), m) % m;

                                                              if (phi[p] == p) {
                                                               phi[p] = p-1;
                                                               for (int i = 2*p; i \le n; i += p) phi[i] =
Multinomials
                                                                (phi[i]/p) * (p-1);
11 multinomial(vector<int>& v) {
    ll c = 1, m = v.empty() ? 1 : v[0];
    for(int i = 1; i < v.size(); i++)</pre>
 for (int j = 0; j < v[i]; j++)
...c = c * ++m / (j+1);
                                                            Factorials
                                                            // digits in factorial
 return c;
                                                            #define kamenetsky(n) (floor((n * log10(n /
                                                             \hookrightarrow ME)) + (log10(2 * MPI * n) / 2.0)) + 1)
Chinese Remainder Theorem
                                                             // approximation of factorial
                                                            #define stirling(n) ((n == 1) ? 1 : sqrt(2 *
bool ecrt(ll* r, ll* m, int n, ll& re, ll& mo)
                                                            \hookrightarrow M PI * n) * pow(n / M E, n))
 ill x, y, d; mo = m[0]; re = r[0];
for (int i = 1; i < n; i++) {</pre>
                                                            // natural log of factorial
#define lfactorial(n) (lgamma(n+1))
  d = egcd(mo, m[i], x, y);

if ((r[i] - re) % d != 0) return false;

x = (r[i] - re) / d * x % (m[i] / d);

re += x * mo;
                                                            Prime Factorization
                                                            // do not call directly
                                                            ll pollard rho(ll n. ll s) {
  mo = mo / d * m[i];
                                                             ll x, y;
  re %= mo;
                                                            x = y = rand() % (n - 1) + 1;
int head = 1, tail = 2;
while (true) {
 re = (re + mo) \% mo;
 return true;
                                                             while (tide) {
   x = mult(x, x, n);
   x = (x + s) % n;
   if (x == y) return n;
Count Digit Occurences
                                                              11 d = \underline{\underline{g}}cd(max(x - y, y - x), n);
/*count(n,d) counts the number of occurences of
                                                             if (1 < \overline{d} \&\& d < n) return d:
                                                              if (++head == tail) y = x, tail <<= 1;
 \rightarrow a digit d in the range [0,n]*/
ll digit_count(ll n, ll d) {
 11 \text{ result} = 0;
 while (n != 0) {
    result += ((n%10) == d ? 1 : 0);
                                                            // call for prime factors
                                                            void factorize(ll n. vector<ll> &divisor) {
                                                             if (n == 1) return;
  n /= 10;
                                                             if (isPrime(n)) divisor.push back(n):
 return result:
                                                             11 count(11 n, 11 d) {
    if (n < 10) return (d > 0 && n >= d);
    if ((n % 10) != 9) return digit_count(n, d) +
                                                              while (d >= n) d = pollard_rho(n, rand() % (n int josephus(int n, int k) {
                                                              - 1) + 1);
factorize(n / d, divisor);
```

factorize(d, divisor);

 \hookrightarrow count(n-1, d);

return 10*count(n/10, d) + (n/10) + (d > 0);

```
Farey Fractions
// generate 0 <= a/b <= 1 ordered, b <= n
   farey(4) = 0/1 1/4 1/3 1/2 2/3 3/4 1/1
// length is sum of phi(i) for i = 1 to n
vector<pair<int, int>> farev(int n) {
 int h = 0, k = 1, x = 1, y = 0, r;
 vector<pair<int, int>> v;
  v.push back({h, k});
  r = (n-y)/k;
  v += r*k: x' += r*h:
 x = -x; y = -y;

while (k > 1);
 v.push_back({1, 1});
 return v;
Fast Fourier Transform
const double PI = acos(-1):
void fft(vector<cd>& a, bool invert) {
 int n = a.size();
 for (int i = 1, j = 0; i < n; i++) {
  int bit = n >>
  for (; j & bit; bit >>= 1) j ^= bit;
  .j ^= bit;
  if (i < j) swap(a[i], a[j]);
 for (int len = 2; len <= n; len <<= 1) {
    double ang = 2 * PI / len * (invert ? -1 :
  cd wlen(cos(ang), sin(ang));
  for (int i = 0; i < n; i += len) {
   .cd w(1):
   for (int j = 0; j < len / 2; j++) {
    cd u = a[i+j], v = a[i+j+len/2] * w;
    a[i+j] = u + v;
a[i+j+len/2] = u - v;
    .w *= wlen:
 if (invert)
  for (auto\& x : a)
  x /= n;
vector<int> fftmult(vector<int> const& a,
 → vector<int> const& b) {
 vector<cd> fa(a.begin(), a.end()),

    fb(b.begin(), b.end());

 int n = 1 << (32 - _builtin_clz(a.size() +

→ b.size() - 1));
fa.resize(n); fb.resize(n);
 fft(fa. false): fft(fb. false)
 for (int i = 0; i < n; i++) fa[i] *= fb[i];
 fft(fa. true):
 vector<int> toret(n);
 for (int i = 0; i < n; i++) toret[i] =
 → round(fa[i].real());
return toret;
Greatest Common Denominator
ll egcd(ll a, ll b, ll& x, ll& y) {
  if (b == 0) { x = 1; y = 0; return a; }
  ll gcd = egcd(b, a % b, x, y);
 x = a / b * y;
 swap(x, y);
 return gcd;
Josephus Problem
 // 0-indexed, arbitrary k
 if (n == 1) return 0;
if (k == 1) return n-1;
 if (k > n) return (joséphus(n-1,k)+k)%n;
 int res = josephus(n-n/k,k)-n\%k;
```

return res + ((res<0)?n:res/(k-1));

```
// fast case if k=2, traditional josephus
int josephus(int n) {
return 2*(n-(1<<(32-builtin clz(n)-1)));
Least Common Multiple
#define lcm(a,b) ((a*b)/qcd(a,b))
Modulo Operations
#define MOD \overline{10000000007}
\#define\ madd(a,b,m)\ (a+b-((a+b-m>=0)?m:0))
#define mult(a,b,m) ((ull)a*b%m)
#define msub(a,b,m) (a-b+((a < b)?m:0))
ll mpow(ll b. ll e. ll m) {
11 x = 1:
while (e > 0) {
    if (e % 2) x = (x * b) % m;
    b = (b * b) % m;
 e /= 2;
return x % m:
ull mfac(ull n, ull m) {
ull f = 1;
for (int i = n; i > 1; i--)
 f = (f * i) \frac{\pi}{n}
 return f:
// if m is not quaranteed to be prime
ll minv(ll b, ll m) {
11 \times (0, y) = 0;
if (egcd(b, m, x, y) != 1) return -1;
return (x % m + m) % m;
11 mdiv_compmod(int a, int b, int m) {
if (\underline{gcd(b, m)} != 1) return -1;
return mult(a, minv(b, m), m);
// if m is prime (like 10^9+7)
ll mdiv_primemod (int a, int b, int m) {
return mult(a, mpow(b, m-2, m), m);
// tonelli shanks = sqrt(n) % m, m is prime
ll legendre(ll a, ll m){
if (a % m==0) return 0:
if (m == 2) return 1;
return mpow(a, (m-1)/2, m);
11 msqrt(11 n, 11 m) {
ll s = __builtin_ctzll(m-1), q = (m-111)>>s.
\Rightarrow z = rand()%(m-1)+1:
 if (m == 2) return 1;
 if (s == 1) return mpow(n, (m+1)/411, m):
 while (legendre(z,m)!=m-1) z = rand()\%(m-1)+1;
 11 c = mpow(z,q,m), r = mpow(n,(q+1)/2,m), t
 \rightarrow = mpow(n,q,m), M = s;
while (t != 1) {
...ll i=1, ts = (t * t) % m;
...while (ts != 1) i++, ts = (ts * ts) % m;
 111 b = c;
for (int j = 0; j < M-i-1; j++) b = (b * b) %
 r = r * b % m; c = b * b % m; t = t * c % m;
 \rightarrow M = i;
return r;
Modulo Tetration
11 tetraloop(11 a, 11 b, 11 m) {
  if(b == 0 | | a == 1) return 1;
 11 w = tetraloop(a,b-1,phi(m)), r = 1;
 for (;w;w/=2) {
 if (w&1)
  r *= a; if (r >= m) r -= (r/m-1)*m;
 \bar{a} *= a; if (a >= m) a -= (a/m-1)*m;
 return r;
int tetration(int a, int b, int m) {
```

```
if (a == 0 | | m == 1) return ((b+1)&1)%m;
                                                                                               |// 1 = even number prime factors, -1 = odd
 return tetraloop(a,b,m) % m;
                                                                                                short mu[MAXN] = \{0,1\};
                                                                                               void mobius(){
  for (int i = 1; i < MAXN; i++)
    if (mu[i])</pre>
Matrix
template<typename T>
                                                                                                    for (int' j = i + i; j < MAXN; j += i)
struct Mat : public Vec<2, T> {
                                                                                                       mu[j] -= mu[i];
 int w, h;
 Mat(int x, int y) : Vec<2, T>(x, y), w(x),
                                                                                                Nimber Arithmetic
 \hookrightarrow h(y) {}
                                                                                               #define nimAdd(a,b) ((a)^(b))
ull nimMul(ull a, ull b, int i=6) {
   static const ull M[]={INT_MIN>>32,
  static Mat<T> identity(int n) { Mat<T> m(n,n);
       for (int i=0;i<n;i++) m[i][i] = 1; return

ighthat is in the second in 
                                                                                                       M[0]^{M[0]}<16), M[1]^{M[1]}<8),
  Mat<\hat{T}>\& operator+=(const Mat<T>\& m) {
  for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
    (*this)[i][j] += m[i][j];
                                                                                                       M[2]^{(M[2] << 4)}, M[3]^{(M[3] << 2)},
                                                                                                      M[4]^(M[4]<<1);
                                                                                                   if (i-- == 0) return a&b;
int k=1<<i:</pre>
    return *this:
                                                                                                   ull s=nimMul(a,b,i), m=M[5-i],
  Mat<T>& operator-=(const Mat<T>& m) {
                                                                                                       t=nimMul(((a^(a>>k))&m)|(s&~m).
   for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
                                                                                                        ((b^{(b)}) \& m) | (m \& (\sim m >> 1)) << k, i):
                                                                                                   return ((s^t)&m)<<k|((s^(t>>k))&m);
       (*this)[i][j] -= m[i][j];
   return *this;
                                                                                                Permutation
                                                                                                //c = array \ size, \ n = nth \ perm, \ return \ index
  Mat<T> operator*(const Mat<T>& m) {
                                                                                                vector<int> gen permutation(int c, int n) {
   Mat<T>z(w,m.h);
   for (int i = 0; i < w; i++)
                                                                                                 vector<int> idx(c), per(c), fac(c); int i;
  ...for (int j = 0; j < h; j++)
....for (int k = 0; k < m.h; k++)
....z[i][k] += (*this)[i][j] * m[j][k];
                                                                                                 for (i = 0; i < c; i++) idx[i] = i;
                                                                                                 for (i = 1; i <= c; i++) fac[i-1] = n%i, n/=i;
                                                                                                 for (i = c - 1; i >= 0; i--
per[c-i-1] = idx[fac[i]],
                                                                                                   idx.erase(idx.begin() + fac[i]);
  Mat<T> operator+(const Mat<T>& m) { Mat<T>
                                                                                                 return per;

→ a=*this; return a+=m; }

 Mat<T> operator-(const Mat<T>& m) { Mat<T>
                                                                                                 // get what nth permutation of vector
                                                                                                int get permutation(vector<int>& v) {

    a=*this; return a-=m; }

                                                                                                 int use = 0, i = 1, r = 0;
for (int e: v) {
   r = r * i++ + __builtin_popcount(use &
 Mat<T>& operator*=(const Mat<T>& m) { return

→ *this = (*this)*m; }

  Mat<T> power(int n) {
                                                                                                  → -(1<<e));
   .Mat<T> a = Mat<T>::identity(w),m=*this;
                                                                                                   use |= 1 << e:
   for (;n;n/=2,m*=m) if (n\&1) a *= m:
   return a;
                                                                                                 return r;
                                                                                                Permutation (string/multiset)
Matrix Exponentiation
                                                                                                string freq2str(vector<int>& v) {
 // F(n) = c[0]*F(n-1) + c[1]*F(n-2) + ...
// b is the base cases of same length c
                                                                                                 string s;
                                                                                                 for (int i = 0; i < v.size(); i++)

for (int j = 0; j < v[i]; j++)

s += (char)(i + 'A');
11 matrix exponentiation(11 n, vector<11> c,
 → vector<11> b) {
   if (nth < b.size()) return b[nth-1];
   Mat<11> a(c.size(), c.size()); ll s = 0;
   for (int i = 0; i < c.size(); i++) a[i][0] =</pre>
                                                                                                 // nth perm of multiset, n is 0-indexed

    c[i];

                                                                                               string gen_permutation(string s, ll n) {
 for (int i = 0; i < c.size() - 1; i++)
                                                                                                 vector<int> freq(26, 0);
 \hookrightarrow a[i][i+1] = 1;
                                                                                                 for (auto e : s) freq[e - 'A']++;
 a = a.power(nth - c.size());
                                                                                                 for (int i = 0; i < 26; i++) if (freq[i] > 0)
 for (int i = 0; i < c.size(); i++)
s += a[i][0] * b[i];
                                                                                                   freq[i]--; ll v = multinomial(freq);
 return s;
                                                                                                   if (n < v) return (char)(i+'A') +
                                                                                                      gen_permutation(freq2str(freq), n);
Matrix Subarray Sums
                                                                                                   freq[\overline{i}]++; n -= v;
template < class T> struct MatrixSum {
                                                                                                 return "";
  Vec<2, T > p;
  MatrixSum(Vec<2, T>& v) {
   p = Vec<2,T>(v.size()+1, v[0].size()+1);
                                                                                                Miller-Rabin Primality Test
 for (int i = 0; i < v.size(); i++)
for (int j = 0; j < v[0].size(); j++)
p[i+1][j+1] = v[i][j] + p[i][j+1] +</pre>
                                                                                                // Miller-Rabin primality test - O(10 log^3 n)
                                                                                               // Miller-Rabin primality test -
bool isPrime(ull n) {
   if (n < 2) return false;
   if (n == 2) return true;
   if (n ½ == 0) return false;
   ull s = n - 1;
   while (s ½ 2 == 0) s /= 2;
   for (int i = 0; i < 10; i++) {
      ull temp = s;
   ull a = rand() ½ (n - 1) + 1;
   ull mod = mpon(a temp n);
   ull mod = mpon(a temp n);
}</pre>

    p[i+1][i] - p[i][i];

 f sum(int u, int l, int d, int r) {
    return p[d][r] - p[d][l] - p[u][r] + p[u][l];
Mobius Function
                                                                                                   ull mod = mpow(a, temp, n);
const int MAXN = 10000000;
// mu[n] = 0 iff n has no square factors
                                                                                                   while (temp!=n-1\&\&mod!=1\&\&mod!=n-1) {
```

```
. mod = mult(mod, mod, n);
    temp *= 2:
   if (mod!=n-1&&temp%2==0) return false:
 return true:
Sieve of Eratosthenes
bitset<100000001> sieve;
// generate sieve - O(n log n)
void genSieve(int n) {
    sieve[0] = sieve[1] = 1;
    for (ull i = 3; i * i < n; i += 2)
        if (!sieve[i])</pre>
   for (ull j = i * 3; j <= n; j += i * 2)
...sieve[j] = 1;
// query sieve after it's generated - O(1)
bool querySieve(int n) {
 return n == 2 || (n % 2 != 0 && !sieve[n]);
Compile-time Prime Sieve
const int MAXN = 100000;
template<int N>
struct Sieve {
  bool sieve[N];
  constexpr Sieve() : sieve() {
  sieve[0] = sieve[1] = 1;
  for (int i = 2; i * i < N; i++)
if (!sieve[i])
     for (int j = i * 2; j < N; j += i)
...sieve[j] = 1;
bool isPrime(int n) {
   static constexpr Sieve<MAXN> s;
 return !s.sieve[n]:
Simpson's / Approximate Integrals
// integrate f from a to b, k iterations
// error <= (b-a)/18.0 * M * ((b-a)/2k)^4

// where M = max(abs(f^{^*}(x))) for x in [a,b]

// "f" is a function "double func(double x)"
double Simpsons (double a. double b. int k.
 \rightarrow double (*f)(double)) {
double dx = (b-a)/(2.0*k), t = 0;
 for (int i = 0; i < k; i++)

t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
 \leftrightarrow (*f)(a+(2*i+1)*dx);
return (t + (*f)(b)) * (b-a) / 6.0 / k;
Common Equations Solvers
// ax^2 + bx + c = 0, find x
vector<double> solveEq(double a, double b,
 → double c) {
vector<double> r;
  double z = \bar{b} * \bar{b} - 4 * a * c;
  if (z == 0)
  r.push_back(-b/(2*a));
  else if (z > 0) {
  r.push back((sgrt(z)-b)/(2*a)):
  r.push_back((sqrt(z)+b)/(2*a));
 return r:
// ax^3 + bx^2 + cx + d = 0, find x
vector<double> solveEq(double a. double b.
 → double c, double d) {
vector<double> res:
 long double a1 = b/a, a2 = c/a, a3 = d/a;
 long double q = (a1*a1 - 3*a2)/9.0, sq =
 \rightarrow -2*sqrt(q);
 long double r = (2*a1*a1*a1 - 9*a1*a2 +
 \rightarrow 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
 if (z <= 0) {
  theta = a\cos(r/sqrt(q*q*q));
```

```
res.push_back(sq*cos(theta/3.0) - a1/3.0);
 res.push back(sq*cos((theta+2.0*PI)/3.0) -
 \rightarrow a1/3.0):
 res.push back(sq*cos((theta+4.0*PI)/3.0) -
   a1/3.0);
 élse {
 res.push_back(pow(sqrt(z)+fabs(r), 1/3.0));
 res[0] = (res[0] + q / res[0]) *

    ((r<0)?1:-1) - a1 / 3.0;
}
return res;</pre>
// linear diophantine equation ax + by = c,
    find x and u
// infinite solutions of form x+k*b/g, y-k*a/g
bool solveEq(ll a, ll b, ll c, ll &x, ll &y, ll
g = \overline{egcd(abs(a), abs(b), x, y)};
 if (c % g) return false;
x *= c / g * ((a < 0) ? -1 : 1);

y *= c / g * ((b < 0) ? -1 : 1);

return true;
// m = # equations, n = # variables, a[m][n+1]
\rightarrow = coefficient matrix
// a[i][0]x + a[i][1]y + ... + a[i][n]z =
   a[i][n+1]

→ equation

const double eps = 1e-7:
bool zero(double a) { return (a < eps) && (a >
\rightarrow -eps); }
vector < double > solveEq(double **a, int m, int
\rightarrow n) { int cur = 0:
for (int i = 0; i < n; i++) {
 for (int j = cur; j < m; j++) {
  if (!zero(a[j][i])) {
   double num = a[sat][i] / a[cur][i];
     for (int sot = 0; sot <= n; sot++)
    a[sat][sot] -= a[cur][sot] * num;
    .}
.cur++:
    break
 for (int j = cur; j < m; j++)
 if (!zero(a[j][n])) return vector<double>();
vector<double> ans(n,0);
for (int i = 0, sat = 0; i < n; i++)
    if (sat < m && !zero(a[sat][i]);
}</pre>
ans[i] = a[sat][n] / a[sat++][i];
return ans:
// solve A[n][n] * x[n] = b[n] linear equation
// rank < n is multiple solutions, -1 is no

→ solutions
// `alls` is whether to find all solutions, or

\hookrightarrow any
const double eps = 1e-12;
int solveEq(Vec<2, double>& A, Vec<1, double>&

→ b, Vec<1, double>& x, bool alls=false) {
int n = A.size(), m = x.size(), rank = 0, br,
vector<int> col(m); iota(begin(col), end(col),
 → 0);
for(int i = 0; i < n; i++) {
  double v, bv = 0;</pre>
  for(int r = i; r < n; r++)
 for(int c = i; c < n; c++)
if ((v = fabs(A[r][c])) > bv)
br = r, bc = c, bv = v;
if (bv <= eps) {
  for(int i = i: i < n: i++)
```

```
. if (fabs(b[j]) > eps)
     .return -1;
  swap(A[i], A[br]);
swap(b[i], b[br]);
   swap(col[i], col[bc]);
  swap(col[i], col[bd]);
for(int j = 0; j < n; j++)
   swap(A[j][i], A[j][bc]);
bv = 1.0 / A[i][i];
for(int j = (alls)?0:i+1; j < n; j++) {</pre>
   if (j != i) {
     double fac = A[j][i] * bv;
     b[j] = fac * b[i];
   for(int k = i+1; k < m; k++)
A[j][k] -= fac*A[i][k];
  rank++;
 }
if (alls) for (int i = 0; i < m; i++) x[i] =</pre>
 → -DBL MAX;
 for (int i = rank; i--;) {
  bool isGood = true;
  if (alls)
  ifor (int j = rank; isGood && j < m; j++)
if (fabs(A[i][j]) > eps)
  isGood = false;
b[i] /= A[i][i];
if (isGood) x[col[i]] = b[i];
  if (!alls)

for(int j = 0; j < i; j++)

b[j] -= A[j][i] * b[i];
 return rank:
Graycode Conversions
ull graycode2ull(ull n) {
 ull i = 0;
for (; n; n = n >> 1) i ^= n;
return i;
ull ull2graycode(ull n) {
  return n ^ (n >> 1);
Unix/Epoch Time
// O-indexed month/time, 1-indexed day
// minimum 1970, 0, 1, 0, 0, 0
ull toEpoch(int year, int month, int day, int
→ hour, int minute, int second) {
struct tm t; time_t epoch;
 t.tm_year = year - 1900; t.tm_mon = month;
t.tm_mday = day; t.tm_hour = hour;
t.tm_min = minute; t.tm_sec = second;
t.tm_isdst = 0; // 1 = daylights savings
 epoch = mktime(\&t);
 return (ull)epoch;
vector<int> toDate(ull epoch) {
 time_t e=epoch; struct tm t=*localtime(&e):
 return {t.tm_year+1900,t.tm_mon,t.tm_mday,t_
     .tm_hour,t.tm_min,t.tm_sec};
int getWeekday(ull epoch) {
 time_t e=epoch; struct tm t=*localtime(&e);
 return t.tm wday: // 0-6. 0 = sunday
int getDayofYear(ull epoch) {
 time_t e=epoch; struct tm t=*localtime(&e);
 return t.tm yday; // 0-365
const int months [] =
→ {31,28,31,30,31,30,31,31,30,31,30,31};

bool validDate(int year, int month, int day) {

   bool leap = !(year%(year%25?4:16));
     if (month >= 12) return false:
     return day <= months[month] + (leap &&
     month == 1):
```

```
Theorems and Formulae
```

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set of size n. !0 = 1, !1 = 0, !n = (n+1)(!(n-1))1+!(n-2), $!n = n! \sum_{i=0}^{n} \frac{(-1)^i}{i!}$, $!n = \left[\frac{n!}{e}\right]$ In a partially ordered set, a chain is a subset of elements that are all comparable to each other. An antichain is a subset where no two are combarable.

Dilworth's theorem states the size of a maximal antichain equals the size of a minimal chain cover of a partially ordered set S. The width of S is the maximum size of an antichain in S, which is equal to the minimum number of chains needed to cover S, or the minimum number of chains such that all elements are in at least one chain.

Rosser's Theorem states the *nth* prime Floyd Warshall number is greater than n * ln(n) for n > 1.

Nicomachi's Theorem states $1^3 + 2^3 + ...$

Lagrange's Four Square Theorem states - && m[k][j] != inf) [every natural number is the sum of the squares of four non-negative integers. This is a special case of the **Fermat Polygonal Number** Theorem where every positive integer is a → m[i][j] = -inf; sum of at most n s-gonal numbers. The nths-gonal number $P(s,n) = (s-2)\frac{n(n-1)}{2} + n$

Graphs

```
struct edge {
int u,v,w;
edge (int u,int v,int w) : u(u),v(v),w(w) {}
edge (): u(0), v(0), w(0) {}
bool operator < (const edge &e1, const edge
\rightarrow &e2) { return e1.w < e2.w; }
bool operator > (const edge &e1, const edge
\rightarrow &e2) { return e1.w > e2.w: }
```

```
struct subset { int p, rank; };
Eulerian Path
#define edge_list vector<edge>
#define adj_sets vector<set<int>>
struct EulerPathGraph {
adj_sets graph; // actually indexes incident
 → edaes
 edge list edges; int n; vector<int> indeg;
EulerPathGraph(int n): n(n) {
 indeg = *(new vector<int>(n.0));
 graph = *(new adj sets(n, set<int>()));
 void add_edge(int u, int v) {
  graph[u].insert(edges.size());
  indeg[v]++:
  edges.push back(edge(u,v,0));
 bool eulerian_path(vector<int> &circuit) {
 if(edges.size()==0) return false;
 stack<int> st;
int a[] = {-1, -1};
```

```
for(int v=0; v<n; v++) {
   if(indeg[v]!=graph[v].size()) {</pre>
                                                                       bool b = indeg[v] > graph[v].size();
if (abs(((int)indeg[v])-((int)graph[v]_
                                                                        .size())) > 1) return
                                                                       false:
                                                                       if (a[b] != -1) return false;
[a[b] = v;
                                                                    int s = (a[0]!=-1 && a[1]!=-1 ? a[0] :
                                                                        (a[0]=-1 & a a[1]=-1 ? edges[0].u : -1);
                                                                     if(s==-1) return false;
                                                                    while(!st.empty() || !graph[s].empty()) {
   if (graph[s].empty()) {
                                                                        circuit.push_back(s); s = st.top();
                                                                       st.pop(); }
                                                                      else {
                                                                       int w = edges[*graph[s].begin()].v;
                                                                        graph[s].erase(graph[s].begin());
                                                                       st.push(s); s = w;
                                                                    circuit.push_back(s);
                                                                    return circuit.size()-1==edges.size():
                                                                const ll inf = 1LL << 62;

#define FOR(i,n) for (int i = 0; i < n; i++)

void floydWarshall(Vec<2, ll>& m) {
Nicomachi's Theorem states 1 2 wold Hoyawai Shart vec 2, 2 int n = m.size();  (n^{\frac{n+1}{2}})^2.  int n = m.size();  [n^{\frac{n+1}{2}}]^2.  FOR(i,n) m[i][i] = min(m[i][i], OLL);  [n^{\frac{n+1}{2}}]^2.  FOR(k,n) FOR(i,n) FOR(j,n) if (m[i][k] != inf
                                                                     auto newDist = max(m[i][k] + m[k][j], -inf);
                                                                    m[i][j] = min(m[i][j], newDist);
                                                                   FOR(k,n) if (m[k][k] < 0) FOR(i,n) FOR(j,n) if (m[i][k] != inf && m[k][j] != inf)
                                                                  // returns vector of edges in the mst
                                                                  // graph[i] = vector of edges incident to
                                                                  → vertex i
// places total weight of the mst in &total
                                                                  // if returned vector has size != n-1, there is
                                                                  vector<edge> mst(vector<vector<edge>> graph.
                                                                   → ll &total) {
                                                                   total = 0;
                                                                   priority_queue<edge, vector<edge>,
                                                                   → greater<edge>> pq;
                                                                    vector<edge> MST;
                                                                   bitset<20001> marked; // change size as needed
                                                                   marked[0] = 1;
                                                                   for (edge ep : graph[0]) pq.push(ep);
while(MST.size()!=graph.size()-1 &&
                                                                    → pq.size()!=0) {`
                                                                    pq.size():-0) \
edge e = pq.top(); pq.pop();
int u = e.u, v = e.v, w = e.w;
if(marked[u] && marked[v]) continue;
                                                                     else if(marked[u]) swap(u, v);
                                                                     for(edge ep : graph[u]) pq.push(ep);
                                                                    marked[u] = 1;
MST.push_back(e);
                                                                    total += e.w:
                                                                   return MST;
```

Union Find

return s[i].p;

```
void uf_union(subset* s, int x, int y) {
                                                     int xp = uf_find(s, x), yp = uf_find(s, y);
                                                     if (s[xp].rank > s[yp].rank) s[yp].p = xp;
else if (s[xp].rank < s[yp].rank) s[xp].p =
                                                     else { s[yp].p = xp; s[xp].rank++; }
                                                     2D Grid Shortcut
                                                     #define inbound(x.n) (0<=x86x<n)
                                                    #define fordir(x, y, n, m) for(auto[dx, dy]:dir)if
                                                     (inbound(x+dx,n)&Ginbound(y+dy,m))
                                                    const pair<int,int> dir[] =
                                                     \leftrightarrow \{\{1,0\},\{0,1\},\{-1,0\},\{0,-1\}\};
                                                         2D Geometry
                                                    #define point complex<double>
#define EPS 0.0000001
                                                     #define sq(a) ((a)*(a))
                                                     #define cb(a) ((a)*(a)*(a))
                                                    double dot(point a, point b) { return

    real(conj(a)*b);
}
                                                    double cross (point a, point b) { return

    imag(conj(a)*b); }

                                                    struct line { point a, b; };
                                                    struct circle { point c; double r; };
                                                    struct segment { point a, b; };
                                                    struct triangle { point a, b, c; };
struct rectangle { point tl, br; };
                                                     struct convex_polygon {
                                                     vector<point points;
                                                     convex_polygon(vector<point> points) :
                                                     → points(points) {}
                                                     convex_polygon(triangle a) {
                                                      points.push_back(a.a); points.push_back(a.b);
                                                      points.push_back(a.c);
                                                     convex_polygon(rectangle a) {
                                                      points.push_back(a.tl);
                                                         points.push back({real(a.tl).
                                                        imag(a.br)}):
                                                      points.push_back(a.br);
                                                         points.push_back({real(a.br),
                                                        imag(a.tl)});
                                                     struct polygon {
                                                     vector <point > points;
                                                     polygon(vector point points) :

→ points(points) {}

                                                     polygon(triangle a) {
                                                      points.push_back(a.a); points.push_back(a.b);
                                                        points.push back(a.c);
                                                     polygon(rectangle a) {
                                                      points.push_back(a.tl);
                                                         points.push_back({real(a.tl),
                                                        imag(a.br)});
                                                      points.push_back(a.br);
                                                         points.push back({real(a.br),
                                                        imag(a.tl)});
                                                     polygon(convex_polygon a) {
                                                      for (point v : a.points)
                                                       points.push_back(v);
                                                     // triangle methods
                                                    double area_heron(double a, double b, double
int uf find(subset* s, int i) {
  if (s[i].p_!= i) s[i].p = uf_find(s, s[i].p);
                                                    \hookrightarrow c) {
if (a < b) swap(a, b);
                                                     if (a < c) swap(a, c);
                                                     if (b < c) swap(b, c);
```

```
if (a > b + c) return -1;
                                                  double d = abs(a.c - b.c);
                                                  if (d <= b.r - a.r) return area(a);
 return sqrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)
                                                      (d \le a.r - b.r) return area(b);
                                                  if (d \ge a.r + b.r) return 0;
}
// segment methods
                                                   double alpha = acos((sq(a.r)' + sq(d) -
                                                  \rightarrow sq(b.r)) / (2 * a.r * d));
double lengthsq(segment a) { return
                                                  double beta = acos((sq(b.r) + sq(d) - sq(a.r)))
    sq(real(a.a) - real(a.b)) + sq(imag(a.a)
                                                     /(2 * b.r * d)):
   imag(a.b)): }
                                                  return sq(a.r) * (alpha - 0.5 * sin(2 *
double length(segment a) { return
                                                      alpha)) + sq(b.r) * (beta - 0.5 * sin(2 *

    sqrt(lengthsq(a)); }

                                                     beta)):
   circle methods
double circumference(circle a) { return 2 * a.r.
                                                    -1 outside, 0 inside, 1 tangent, 2
double area(circle a) { return sq(a.r) * M PI;
                                                  int intersection(circle a, circle b,
\rightarrow } // rectangle methods
                                                  → vector<point>& inter) {
                                                  double d2 = norm(b.c - a.c), rS = a.r + b.r,
double width(rectangle a) { return
                                                  \rightarrow rD = a.r - b.r;

→ abs(real(a.br) - real(a.tl)); }

                                                  if (d2 > sq(rS)) return -1;
double height (rectangle a) { return
                                                  if (d2 < sq(rD)) return 0;

→ abs(imag(a.br) - real(a.tl)); }

                                                  double ca = 0.5 * (1 + rS * rD / d2);
double diagonal (rectangle a) { return
                                                  point z = point(ca, sqrt(sq(a.r) / d2 -

    sqrt(sq(width(a)) + sq(height(a))); }

                                                     sq(ca)));
double area(rectangle a) { return width(a) *
                                                  inter.push_back(a.c + (b.c - a.c) * z);
                                                  if (abs(imag(z)) > EPS) inter.push_back(a.c
→ height(a); }
double perimeter(rectangle a) { return 2 *
                                                      (b.c - a.c) * conj(z));
   (width(a) + height(a)); }
                                                  return inter.size():
// check if `a` fit's inside `b
                                                  // points of intersection
// swap equalities to exclude tight fits
                                                 vector<point> intersection(line a, circle c) {
bool doesFitInside(rectangle a, rectangle b) {
                                                  vector<point> inter;
c.c -= a.a;
int x = width(a), w = width(b), y = height(a),
\rightarrow h = height(b);
                                                  a.b -= a.a;
if (x > y) swap(x, y);
if (w > h) swap(w, h);
                                                  point m = a.b * real(c.c / a.b);
double d2 = norm(m - c.c);
                                                  if (d2 > sq(c.r)) return 0;
 if (w < x) return false;
                                                  double 1 = sqrt((sq(c.r) - d2) / norm(a.b));
 if (v <= h) return true;
                                                   inter.push back(a.a + m + 1 * a.b);
 double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
                                                  if (abs(1) > EPS) inter.push_back(a.a + m - 1 10 Optimization
 return sq(a) \le sq(b) + sq(c);
}
// polygon methods

    * a.b);
return inter:
// negative area = CCW, positive = CW
                                                  // area of intersection
double area(polygon a) {
  double area = 0.0; int n = a.points.size();
  for (int i = 0, j = 1; i < n; i++, j = (j - 1)
                                                     max(imag(a.tl), imag(b.tl));
                                                  double x2 = min(real(a.br), real(b.br)), y2 =
    area += (real(a.points[j]-a.points[i]))*|
                                                     min(imag(a.br), imag(b.br));
return (x2 \le x1 \mid | y2 \le y1) ? 0 :
  return area / 2.0;
                                                     (x2-x1)*(y2-y1);
// get both unsigned area and centroid
pair<double, point> area_centroid(polygon a) {
                                                 Convex Hull
 int n = a.points.size():
                                                  bool cmp(point a, point b) {
  if (abs(real(a) - real(b)) > EPS) return
 double area = 0:
 point c(0, 0);
                                                      real(a) < real(b);
 for (int i = n - 1, j = 0; j < n; i = j++) {
                                                   if (abs(imag(a) - imag(b)) > EPS) return
 double v = cross(a.points[i], a.points[j]) /
                                                     imag(a) < imag(b);</pre>
                                                  return false;
  areá += v:
  c += (a.points[i] + a.points[j]) * (v / 3);
                                                  convex_polygon convexhull(polygon a) {
                                                  sort(a.points.begin(), a.points.end(), cmp);
 c /= area;
                                                  vector<point> lower, upper;
 return {area, c};
                                                  for (int i = 0; i < a.points.size(); i++) {
                                                   while (lower.size() >= 2 &&
                                                      cross(lower.back() - lower[lower.size() -
Intersection
// -1 coincide, 0 parallel, 1 intersection
                                                     2], a.points[i] - lower.back()) < EPS)
int intersection(line a, line b, point& p) {
                                                    lower.pop_back();
                                                   while (upper.size() >= 2 &&
 if (abs(cross(a.b - a.a, b.b - b.a)) > EPS) {
 p = cross(b.a - a.a, b.b - a.b) / cross(a.b)
                                                      cross(upper.back() - upper[upper.size() -
\rightarrow - a.a, b.b - b.a) * (b - a) + a;
                                                     2], a.points[i] - upper.back()) > -EPS)
 return 1:
                                                    upper.pop_back();
                                                   lower.push_back(a.points[i]);
 if (abs(cross(a.b - a.a. a.b - b.a)) > EPS)
                                                    upper.push_back(a.points[i]);

    return 0;

return -1;
                                                   lower.insert(lower.end(), upper.rbegin() + 1,
}
// area of intersection
                                                  → upper.rend());
double intersection(circle a, circle b) {
                                                  return convex polygon(lower);
```

```
3D Geometry
                                                      struct point3d {
                                                       double x, y, z;
                                                       point3d operator+(point3d a) const { return
                                                       \rightarrow {x+a.x, y+a.y, z+a.z}; }
                                                       point3d operator*(double a) const { return
                                                       \rightarrow {x*a, y*a, z*a}; }
                                                       point3d operator-() const { return {-x, -y,
                                                       \rightarrow -z}; }
                                                       point3d operator-(point3d a) const { return
                                                       \rightarrow *this + -a; }
                                                       point3d operator/(double a) const { return
                                                       \rightarrow *this * (1/a); }
                                                       double norm() { return x*x + y*y + z*z; }
                                                       double abs() { return sqrt(norm()); }
                                                       point3d normalize() { return *this /
                                                      |

this->abs(); }
                                                      double dot(point3d a, point3d b) { return
                                                      \rightarrow a.x*b.x + a.y*b.y + a.z*b.z; }
                                                      point3d cross(point3d a, point3d b) { return
                                                          \{a.v*b.z - a.z*b.v, a.z*b.x - a.x*b.z,
                                                          a.x*b.y - a.y*b.x; }
                                                      struct line3d { point3d a, b; };
                                                      struct plane { double a, b, c, d; } // a*x +
                                                      \hookrightarrow b*u + c*z + d = 0
                                                      struct sphere { point3d c; double r; };
                                                      #define sq(a) ((a)*(a))
                                                      #define c\bar{b}(a) ((a)*(a)*(a))
                                                      double surface(circle a) { return 4 * sq(a.r)
                                                      double volume(circle a) { return 4.0/3.0 *
                                                      \rightarrow cb(a.r) * M PI; }
                                                      Snoob
                                                      // SameNumberOfOneBits, next permutation
                                                      int snoob(int a) {
// example usage
                                                     int main() {
  char l1[] = {'1', '2', '3', '4',
   char l2[] = {'a', 'b', 'c', 'd'};
  int di = 5, d2 = 4;
  // prints 12345abcd, 1234a5bcd, ...
                                                       int min = (1 < < d1) - 1, max = min < < d2;
                                                       for (int i = min; i <= max; i = snoob(i)) {
   int p1 = 0, p2 = 0, v = i;
                                                        while (p1 < d1 || p2 < d2) {
                                                        cout << ((v & 1) ? l1[p1++] : l2[p2++]);
                                                        v /= 2;
                                                        .cout << '\n';
                                                      Powers
                                                      bool isPowerOf2(ll a) {
return a > 0 && !(a & a-1);
                                                      bool isPowerOf3(11 a) {
return a>0&&!(12157665459056928801u11%a);
                                                      bool isPower(ll a, ll b) {
  double x = log(a) / log(b);
                                                       return abs(x-round(x)) < 0.00000000001:
                                                      11 Additional
                                                      Judge Speed
                                                     // kattis: 0.50s
                                                        / codeforces: 0.421s
```

// atcoder: 0.455s

```
|#include <bits/stdc++.h>
using namespace std;
int v = 1e9/2, p = 1;
int main() {
  for (int i = 1; i <= v; i++) p *= i;</pre>
 cout << p;
Judge Pre-Contest Checks
     int128 and
                        float128 support?
 -does extra or missing whitespace cause WA?
 documentation up to date?
 printer usage available and functional?
   each case tests a different fail condition
 // try them before contests to see error codes
struct g { int arr[1000000]; g(){}};
vector<g> a;
// O=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE
→ 6=SIGSEGV 7=recursive MLE int judge(int n) {
    (n == 0) exit(0);
    (n == 1) while(1);

(n == 2) while(1) a.push_back(g());

(n == 3) while(1) putchar_unlocked('a');
 if (n == 4) assert(0);
if (n == 5) 0 / 0;
if (n == 6) *(int*)(0) = 0;
 return n + judge(n + 1);
```

GCC Builtin Docs // 128-bit integer

```
__int128 a;
unsigned __int128 b;
// 128-bit float
// minor improvements over long double
__float128 c;
// log2 floor
__lg(n);
// number of 1 bits
 // can add ll like popcountll for long longs
builtin_popcount(n);
// number of trailing zeroes
builtin_ctz(n);
// number of leading zeroes
builtin_clz(n);
// 1-indexed least significant 1 bit
__builtin_ffs(n);
// parity of number
 __builtin_parity(n);
```

Limits

```
\pm 2147483647 \mid \pm 2^{31} - 1 \mid 10^9
int
                                                \bar{2}^{32} - \bar{1}|10^9
                           4294967295
uint
         \pm 9223372036854775\overline{807}|\pm \overline{2}^{63}-1|\overline{10}^{18}
11
                                                 \overline{2}^{64} - 1|10^{19}
          18446744073709551615
|i128| \pm 170141183460469231... | \pm 2^{\tilde{1}27} - 1 | 10^{38}
u128 340282366920938463...
```

Complexity classes input size (per second):

	r (r)
$O(n^n)$ or $O(n!)$	$n \le 10$
$O(2^n)$	$n \leq 30$
$O(n^3)$	$n \le 1000$
$O(n^2)$	$n \le 30000$
$O(n\sqrt{n})$	$n \le 10^6$
$O(n \log n)$	$n \le 10^7$
O(n)	$n < 10^9$