```
General
    Algorithms
    Data Structures
    String
    Math
    Graph
    2D Geometry
    3D Geometry
    Optimization
     General
test.sh
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp
for i in *.in; do
f=${i%.in}
./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
done
Header
// use better compiler options
#pragma GCC optimize("Ofast", "unroll-loops")
#pragma GCC target("avx2, fma")
 // include everything
#include <bits/stdc++.h>
 #include <bits/extc++.h>
#include <sys/resource.h>
// namespaces
using namespace std;
using namespace __gnu_cxx; // rope
using namespace __gnu_pbds; // tree/trie
// common defines
#define fastio
\rightarrow ios_base::sync_with_stdio(0);cin.tie(0); #define nostacklim_rlimit
    RZ; qetrlimit(3, \&RZ); RZ. rlim\ cur=-
insetrlimit(3, €RZ);
#define DEBUG(v) cout«"DEBUG: "«#v«" =
#define ld long double
// global variables
mt19937 rng((uint32_t)chrono::steady

    clock::now().time_since_epoch().count());

Fast IO
void readn(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
  n = n * 10 + c - 0':
void readn(int& n) {
  char c; n = 0; int s = 1;
 if ((c=getchar_unlocked())=='-') s = -1;
 else n = c - '0';
while ((c=getchar_unlocked())!=' '&&c!='\n')
 n = n * 10 + c - 0';

n *= s;
void readn(ld& n) {
 char c; n = 0;
ld m = 0, o = 1; bool d = false; int s = 1;
if ((c=getchar_unlocked())=='-') s = -1;
 else if (c == '.') d = true;
else n = c - '0';
while ((c=getchar_unlocked())!=' '&&c!='\n')

if (c == '.') d = true;
else if (d) { m=m*10+c-'0'; o*=0.1; }
else n = n * 10 + c - '0';
\hat{n} = s * (n + m * o);
```

```
|void readn(double& n) {
 ld m: readn(m): n = m:
void readn(float& n) {
 ld m; readn(m); n = m;
void readn(string& s) {
 char c; s = "";
while((c=getchar_unlocked())!=' '&&c!='\n')
bool readline(string& s) {
 char c: s = "":
 while(c=getchar_unlocked()) {
  if (c == '\n') return true;
if (c == EOF) return false;
 return false:
void printn(unsigned int n) {
 if (n / 10) printn(n / 10):
 putchar_unlocked(n \frac{10}{10} + \frac{10}{10});
void printn(int n) {
  if (n < 0) { putchar_unlocked('-'); n*=-1; }</pre>
 printn((unsigned int)n);
      Algorithms
Min/Max Subarray
// max - compare = a < b, reset = a < 0
// min - compare = a > b, reset = a > 0
// returns {sum, {start, end}}
pair<int, pair<int, int>
     ContiguousSubarray(int* a, int size,
     bool(*compare)(int, int),
 bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
→ 0, s = 0;

for (int i = 0; i < size; i++) {

cur += a[i];
  if ((*compare)(best, cur)) { best = cur;

    start = s; end = i; }
    if ((*reset)(cur)) { cur = 0; s = i + 1; }

 return {best, {start, end}}:
Quickselect
#define QSNE -999999
int partition(int arr[], int l, int r)
int x = arr[r], i = 1;
for (int j = 1; j <= r - 1; j++)
    if (arr[j] <= x)
    swap(arr[i++], arr[j]);</pre>
 swap(arr[i], arr[r]);
// find k'th smallest element in unsorted
→ array, only if all distinct
int qselect(int arr[], int 1, int r, int k)
if (!(k > 0 && k <= r - 1 + 1)) return QSNE;
swap(arr[1 + rng() % (r-1+1)], arr[r]);
int pos = partition(arr, 1, r);
if (pos-1==k-1) return arr[pos];</pre>
 if (pos-1>k-1) return qselect(arr,1,pos-1,k);
 return qselect(arr, pos+1, r, k-pos+1-1);
// TODO: compare against std::nth_element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted
→ on both axis
pair<int, int saddleback_search(int ** arr,
\rightarrow int x, int y, int v) {
```

int i = x-1, j = 0;

while (i >= 0 && j < v) {

```
if (arr[i][j] == v) return {i, j};
  (arr[i][j] > v)? i-: j++;
 return {-1, -1};
Ternary Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a, int b, int (*f)(int)) {
 while (b-a > 4) {
  int m = (a+b)/2
  if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1;
 for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
 a = i;
return a;
#define TERNPREC 0.000001
double ternsearch (double a, double b, double
 \leftrightarrow (*f)(double)) \{
 while (b-a > TERNPREC * 4) {
  double m = (a+b)/2;
  if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
  for (double i = a + TERNPREC; i <= b; i +=
    TERNPREC)
      if (TERNCOMP((*f)(a), (*f)(i)))
 return á:
3 Data Structures
Fenwick Tree
// Fenwick tree, array of cumulative sums -
 \hookrightarrow O(log n) updates, O(log n) gets
struct Fenwick { int n: ll* tree:
 void update(int i. int val) {
  .++i:
  while (i <= n) {
   tree[i] += val;
   i += i & (-i);
 Fenwick(int size) {
 n = size;
tree = new ll[n+1];
for (int i = 1; i <= n; i++)
tree[i] = 0;</pre>
 Fenwick(int* arr, int size) : Fenwick(size) {
  for (int i = 0; i < n; i++)
...update(i, arr[i]):
 ~Fenwick() { delete[] tree; }
 11 operator[](int i) {
  if (i < 0 \mid | i > n) return 0:
  while (i>0) {
  sum += tree[i];
i -= i & (-i);
  return sum;
 11 getRange(int a, int b) { return

    operator[](b) - operator[](a-1); }

Hashtable
// similar to unordered map, but faster
struct chash { | const uint64_t C = (11)(2e18 * M_PI) + 71;
 ll operator()(ll x) const { return
 \rightarrow __builtin_bswap64(x*C); }
int main() {
   gp_hash_table<11,int,chash>
\rightarrow hashtable({},{},{},{},{1 < 16});
```

```
for (int i = 0; i < 100; i++)
  hashtable[i] = 200+i;
if (hashtable.find(10) != hashtable.end())</pre>
  cout « hashtable[10];
Ordered Set
typedef tree<int,null_type,less<int>,rb_tree
     _tag, tree_order_statistics_node_update>
⇒ ordered
    ordered set:
 lordered_set o_set;
o_set.insert(5); o_set.insert(1);
→ o_set.insert(3);
// get second smallest element
 cout « *(o_set.find_by_order(1)) « '\n';
 // number of elements less than k=4 cout « o_set.order_of_key(4) « '\n';
Rope
// \tilde{O}(\log n) insert, delete, concatenate
int main() {
  // generate rope
 rope<int> v;
for (int i = 0; i < 100; i++)
v.push_back(i);</pre>
  // move range to front
 rope<int> copy = v.substr(10, 10);
 v.erase(10, 10);
 v.insert(copy.mutable_begin(), copy);
 // print elements of rope
 for (auto it : v) cout « it « " ";
Segment Tree
//max(a,b), min(a,b), a+b, a*b, gcd(a,b), a^b
struct SegmentTree {
 typedef int T;
 static constexpr T UNIT = INT MIN;
 T f(T a, T b) {
  if (a == UNIT) return b;
if (b == UNIT) return a;
  return max(a,b);
 int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
 SegmentTree(vector<T> arr)
 SegmentTree(arr.size()) {
for (int i=0;i<arr.size();i++)</pre>
 → update(i,arr[i]);
 void update(int pos, T val) {
  for (s[pos += n] = val; pos /= 2;)
   s[pos] = f(s[pos * 2], s[pos*2+1]);
T query(int b, int e) { // query [b, e) 
T ra = UNIT, rb = UNIT; 
for (b+=n, e+=n; b<e; b/=2, e/=2) { 
if (b % 2) ra = f(ra, s[b++]); 
if (e % 2) rb = f(s[-e], rb);
  return f(ra, rb):
   get(int p) { return query(p, p+1); }
typedef trie<string, null type,
pat trie tag,
int main() {
 // generate trie
 trie_type trie;
 for (int i = 0; i < 20; i++)
  trie.insert(to string(i)): // true if new.
```

```
// print things with prefix "1"
 auto range = trie.prefix_range("1");
 for (auto it = range.first; it !=

→ range.second; it++)

  cout « *it « '
4 String
Aho Corasick
// range of alphabet for automata to consider
// MAXC = 26, OFFC = 'a' if only lowercase
const int MAXC = 256;
const int OFFC = 0;
struct aho_corasick {
  struct state
  set<pair<int, int> out;
  .int fail; vector<int> go
  state() : fail(-1), go(MÁXC, -1) \{\}
 vector<state> s;
 int id = 0;
 aho_corasick(string* arr, int size) : s(1) {
  for (int i = 0; i < size; i++) {
  int cur = 0;
 for (int c : arr[i]) {
    if (s[cur].go[c-OFFC] == -1) {
  s[cur].go[c-OFFC] = s.size();
   ...s.push_back(state());
    cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++});
 for (int c = 0; c < MAXC; c++)
. if (s[0].go[c] == -1)
   ..s[0].go[c] = 0;
  queue<int> sq;
for (int c = 0; c < MAXC; c++) {
   if (s[0].go[c] != 0) {</pre>
  s[s[0].go[c]].fail = 0;
sq.push(s[0].go[c]);
  while (sq.size()) {
  int e = sq.front(); sq.pop();
for (int c = 0; c < MAXC; c++) {
   if (s[e].go[c] != -1) {</pre>
      int failure = s[e].fail;
while (s[failure].go[c] == -1)
      failure = s[failure].fail;
failure = s[failure].go[c];
      s[s[e].go[c]].fail = failure;
      for (auto length : s[failure].out)
  s[s[e].go[c]].out.insert(length);
     sq.push(s[e].go[c]);
 // list of {start pos, pattern id}
 vector<pair<int, int> search(string text)
 vector<pair<int, int» toret;
  int cur = 0:
 for (int i = 0; i < text.size(); i++) {
  while (s[cur].go[text[i]-OFFC] == -1)
  cur = s[cur].fail;
   cur = s[cur].go[text[i]-OFFC];
  if (s[cur].out.size())
  for (auto end : s[cur].out)
  toret.push_back({i - end.first + 1,

→ end.second);
  return toret:
Bover Moore
struct defint { int i = -1; };
vector int boyermoore (string txt, string pat) Longest Common Prefix
```

```
.vector<int> toret; unordered map<char,
   → defint> badchar:
   int m = pat.size(), n = txt.size();
   for (int i = 0; i < m; i++) badchar[pat[i]].i
  \Rightarrow = i;
int s = 0:
   while (s <= n - m) {
int j = m - 1;
      while (j \ge 0 \&\& pat[j] == txt[s + j]) j-;
     if (j < 0) {
        toret.push_back(s);
       s += (s + m < n) ? m - badchar[txt[s + m < n]) ] % m - badchar[txt[s + m < n]) % m - badchar[t
   \rightarrow m]].\dot{i}: 1;
       s += max(1, j - badchar[txt[s + i]].i);
   return toret;
 English Conversion
 const string ones[] = {"", "one", "two",
         "three", "four", "five", "six", "seven", "eight", "nine"};
 const string teens[] ={"ten", "eleven",
const string teens[] ={"ten", "eleven",
    "twelve", "thirteen", "fourteen",
    "fifteen", "sixteen", "seventeen",
    "eighteen", "nineteen"};
const string tens[] = {"twenty", "thirty",
    "forty", "fifty", "sixty", "seventy",
    "eighty", "ninety"};
 const string mags[] = {"thousand", "million",
           "billion", "trillion", "quadrillion", "quintillion", "sextillion",
⇒ "septillion"};
string convert(int num, int carry) {
   if (num < 0) return "negative " +
           convert(-num, 0);
  if (num < 10) return ones[num];
if (num < 20) return teens[num % 10];</pre>
   if (num < 100) return tens[(num / 10) - 2] +
           (num%10==0?"":" ") + ones[num % 10];
   if (num < 1000) return ones[num / 100]
           (num/100==0?"":"") + "hundred" + (num%100==0?"":"") + convert(num % 100,
        0):
   return convert(num / 1000, carry + 1) + " " +
          mags[carry] + " " + convert(num % 1000,
         0);
 string convert(int num) {
  return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
 vector<int> kmp(string txt, string pat) {
   vector<int> toret;
int m = txt.length(), n = pat.length();
   int next[n + 1];
   for (int i = 0; i < n + 1; i++)
   next[i] = 0;
for (int i = 1; i < n; i++) {
     int j = next[i + 1];
      while (j > 0 && pat[j] != pat[i])
       j = next[j];
     if (j > 0 || pat[j] == pat[i])
next[i + 1] = j + 1;
    for (int i = 0, j = 0; i < m; i++) {
     if (txt[i] == pat[j]) {
      if (++j == n)
           toret.push_back(i - j + 1);
 eise if (j
; j = next[j];
; i-;
}
     .} else if _(j > 0) {
   return toret:
```

```
|string lcp(string* arr, int n) {
 if (n == 0) return "";
 sort(arr, arr + n);
string r = ""; int v = 0;
  while (v < arr[0].length() && arr[0][v] ==
 → arr[n-1][v])
∴r += arr[0][v++];
 return r;
Longest Common Subsequence
string lcs(string a, string b) {
  int m = a.length(), n = b.length();
  int L[m+1][n+1]:
 for (int i = 0; i <= m; i++) {
    for (int j = 0; j <= n; j++) {
        if (i == 0 || j == 0) L[i][j] = 0;
        else if (a[i-1] == b[j-1]) L[i][j] =
   L[i-1][j-1]+1;
else L[i][j] = \max(L[i-1][j], L[i][j-1]);
  // return L[m][n]; // length of lcs
 string out = "";
int i = m - 1, j = n - 1;
while (i >= 0 && j >= 0) {
  if (a[i] == b[j]) {
   out = a[i-] + out;
  else if (L[i][j+1] > L[i+1][j]) i-;
  else j-;
 return out;
Longest Common Substring
// l is array of palindrome length at that
int manacher(string s, int* 1) {
int n = s.length() * 2;
 for (int i = 0, j = 0, k; i < n; i += k, j =
 \rightarrow max(j-k, 0)) {
  while (i \ge j \&\& i + j + 1 < n \&\& s[(i-j)/2]
 \rightarrow == s[(i+j+1)/2]) j++;
  .1[i] = j;
  for (k = 1; i >= k && j >= k && l[i-k] !=
 \rightarrow j-k; k++)
  1[i+k] = min(1[i-k], j-k);
 return *max_element(1, 1 + n);
Subsequence Count
    "banana", "ban" » 3 (ban, ba..n, b..an)
ull subsequences(string body, string subs) {
 int m = subs.length(), n = body.length();
 if (m > n) return 0;

ull** arr = new ull*[m+1];

for (int i = 0; i <= m; i++) arr[i] = new
 \rightarrow ull[n+1];
 \rightarrow subs[i-1])? arr[i-1][j-1] : 0);
 return arr[m][n];
5 Math
Catalan Numbers
ull* catalan = new ull[1000000];
void genCatalan(int n, int mod) {
  catalan[0] = catalan[1] = 1;
  for (int i = 2; i <= n; i++) {</pre>
  catalan[i] = 0;
for (int j = i - 1; j >= 0; j-) {
```

catalan[i] += (catalan[i] \* catalan[i-i-1])

```
...if (catalan[i] >= mod)
     ___catalan
...catalan[i]
.}
   // TODO: consider binomial coefficient method
   Combinatorics (nCr, nPr)
   // can optimize by precomputing factorials, and
   \hookrightarrow fact[n]/fact[n-r]
  ull nPr(ull n, ull r) {
    ull v = 1;
     for (ull i = n-r+1; i <= n; i++)
     return v;
   ull nPr(ull n, ull r, ull m) {
     for (ull i = n-r+1; i <= n; i++)
. v = (v * i) % m;
     return v;
   úll nCr(ull n, ull r) {
    long double v = 1;
for (ull i = 1: i <= r: i++)
      v = v * (n-r+i) /i;
     return (ull) (v + 0.001);
  // requires modulo math
   // can optimize by precomputing mfac and
   ull nCr(ull n, ull r, ull m) {
    return mfac(n, m) * minv(mfac(k, m), m) % m *
 \downarrow \rightarrow \min_{k} 
   Chinese Remainder Theorem
   bool ecrt(ll* r, ll* m, int n, ll& re, ll& mo)
     11 \, x, y, d; mo = m[0]; re = r[0];
   mo = mo / d * m[i]:
       re %= mo;
     re = (re + mo) \% mo;
     return true:
   Count Digit Occurences
   /*count(n,d) counts the number of occurences of
   \rightarrow a digit d in the range [0,n]*/
11 digit_count(11 n, 11 d) {
11 result = 0;
  while (n != 0) {
    result += ((n%10) == d ? 1 : 0);
}
       n /= 10;
     return result;
  il count(ll n, ll d) {
    if (n < 10) return (d > 0 && n >= d);
if ((n % 10) != 9) return digit_count(n, d) +
   \rightarrow count(n-1, d);
     return 10*count(n/10, d) + (n/10) + (d > 0):
   Discrete Logarithm
   unordered_map<int, int> dlogc;
   int discretelog(int a, int b, int m) {
     dlogc.clear();
     11 \text{ n} = \text{sqrt(m)} + 1, \text{ an } = 1;
```

for (int i = 0; i < n; i++)
an = (an \* a) % m;

for (int i = 1; i <= n; i++) {
 if (!dlogc.count(c)) dlogc[c] = i;

11 c = an;

```
c = (c * an) % m;
                                                            r = (n-y)/k;
                                                             y += r*k; x += r*h;
                                                           swap(x,h); swap(y,k);
x = -x; y = -y;
while (k > 1);
 c = b:
 for (int i = 0; i <= n; i++) {
   if (dlogc.count(c)) return (dlogc[c] * n - i</pre>
                                                            v.push_back({1, 1});
 \rightarrow + m -1) % (m-1);
 c = (c * a) \% m;
return -1;
                                                           Fast Fourier Transform
                                                           #define cd complex<double>
Euler Phi / Totient
                                                           const double PI = acos(-1);
int phi(int n) {
                                                           void fft(vector<cd>& a, bool invert) {
 int'r = n;
                                                            int n = a.size();
 for (int i = 2; i * i <= n; i++) {
   if (n % i == 0) r -= r / i;
   while (n % i == 0) n /= i;
                                                            for (int i = 1, j = 0; i < n; i++) {
                                                             int bit = n \gg 1
                                                            for (; j & bit; bit »= 1) j ^= bit;
j ^= bit;
 if (n > 1) r -= r / n;
return r;
                                                             if (i < j) swap(a[i], a[j]);
                                                            for (int len = 2; len <= n; len «= 1) {
   double ang = 2 * PI / len * (invert ? -1 :
 #define n 100000
ll phi[n+1];
void computeTotient() {
                                                             cd wlen(cos(ang), sin(ang));
 for (int i=1; i<=n; i++) phi[i] = i;
                                                             for (int i = 0; i < n; i += len) {
 for (int p=2; p<=n; p++) {
                                                              cd w(1);
  if (phi[p] == p) {
                                                              for (int j = 0; j < len / 2; j++) {
    cd u = a[i+j], v = a[i+j+len/2] * w;
   phi[p] = p-1;
 ...for (int i = 2*p; i<=n; i += p) phi[i] =
                                                               a[i+j] = u + v;
    (phi[i]/p) * (p-1);
                                                               a[i+j+len/2] = u - v;
                                                               w *= wlen;
Factorials
                                                            if (invert)
for (auto& x : a)
// digits in factorial
 #define kamenetsky(n) (floor((n * log10(n / log1)
                                                             . x /= n;
 \hookrightarrow ME)) + (log10(2 * MPI * n) / 2.0)) + 1)
                                                           vector<int> fftmult(vector<int> const& a,
// approximation of factorial
#define stirling(n) ((n == 1) ? 1 : sqrt(2 *

  vector<int> const& b) {
  vector<cd> fa(a.begin(), a.end()),
 \hookrightarrow M PI * n) * pow(n / M E, n))
                                                            → fb(b.begin(), b.end());
 // natural log of factorial
                                                           int n = 1 ≪ (32 - _builtin_clz(a.size() +

→ b.size() - 1));
fa.resize(n); fb.resize(n);
 #define lfactorial(n) (lgamma(n+1))
Prime Factorization
                                                           fft(fa, false); fft(fb, false);
for (int i = 0; i < n; i++) fa[i] *= fb[i];</pre>
 // do not call directly
ll pollard_rho(ll n, ll s) {
                                                            fft(fa, true);
 .11 x, y;
                                                            vector<int> toret(n);
 x = y = rand() \% (n - 1) + 1;
                                                            for (int i = 0; i < n; i++) toret[i] =
 int head = 1, tail = 2;
while (true) {

→ round(fa[i].real());

 x = mult(x, x, n);

x = (x + s) % n;

if (x == y) return n;
                                                            return toret;
                                                           Greatest Common Denominator
  ll d = gcd(max(x - y, y - x), n);
if (1 < d && d < n) return d;
                                                           ll egcd(ll a, ll b, ll& x, ll& y) {
                                                           if (b == 0) { x = 1; y = 0; return a; }
ll gcd = egcd(b, a % b, x, y);
  if (++head == tail) y = x, tail \ll 1;
                                                            x = a / b * y;
 // call for prime factors
                                                            swap(x, y);
void factorize(ll n, vector<ll> &divisor) {
                                                            return gcd;
 if (n == 1) return;
if (isPrime(n)) divisor.push back(n);
                                                           Josephus Problem
  while (d'>= n) d = pollard rho(n, rand() %
                                                           // 0-indexed, arbitrary k
 \hookrightarrow (n - 1) + 1);
                                                           int josephus(int n, int k) {
 factorize(n / d, divisor);
                                                           if (n == 1) return 0;
if (k == 1) return n-1;
  factorize(d, divisor);
                                                            if (k > n) return (josephus(n-1,k)+k)%n;
                                                            int res = josephus(n-n/k,k)-n\%k;
                                                            return res + ((res<0)?n:res/(k-1));
Farey Fractions
 // generate 0 <= a/b <= 1 ordered, b <= n
                                                           /\!/ fast case if k=2, traditional josephus
 // farey(4) = 0/1 1/4 1/3 1/2 2/3 3/4 1/1
                                                           int josephus(int n) {
// length is sum of phi(i) for i = 1 to n
                                                            return 2*(n-(1\ll(32-\_builtin\_clz(n)-1)));
vector<pair<int, int» farey(int n) {
  int h = 0, k = 1, x = 1, y = 0, r;
  vector<pair<int, int» v;</pre>
                                                           Least Common Multiple
                                                           #define lcm(a,b) ((a*b)/_gcd(a,b))
  v.push back({h, k});
```

```
Modulo Operations
#define MOD 1000000007
#define madd(a,b,m) (a+b-((a+b-m>=0)?m:0))
#define mult(a,b,m) ((ull) a*b%m)
#define msub(a,b,m) (a-b+((a < b)?m:0))
| 11 mpow(11 b, 11 e, 11 m) {
 11 x = 1;
 while (e > 0) {
    if (e % 2) x = (x * b) % m;
    b = (b * b) % m;
  e /= 2;
  return x % m:
ull mfac(ull n, ull m) {
ull f = 1;
 for (int i = n; i > 1; i-)

for (f * i) % m;

return f;
// if m is not guaranteed to be prime
\begin{cases} \lim_{x \to 0} (11 & b, 11 & m) \\ 11 & x = 0, y = 0; \end{cases}
 if (egcd(b, m, x, y) != 1) return -1;
return (x % m + m) % m;
ll mdiv_compmod(int a, int b, int m) {
 if (__gcd(b, m) != 1) return -1;
  return mult(a, minv(b, m), m);
// if m is prime (like 10^9+7)
11 mdiv_primemod (int a, int b, int m) {
 return mult(a, mpow(b, m-2, m), m);
Miller-Rabin Primality Test
 // Miller-Rabin primality test - O(10 log^3 n)
bool isPrime(ull n) {
   if (n < 2) return false;
   if (n = 2) return true;
   if (n % 2 == 0) return false;
   ull s = n - 1;
   while (s % 2 == 0) s /= 2;
   for (int i = 0, i < 10, iii)
  for (int i = 0; i < 10; i++) {
  ull temp = s;
   ull a = rand() \% (n - 1) + 1;
   ull mod = mpow(a, temp, n);
   while (temp!=n-1\&\&mod!=1\&\&mod!=n-1)
    mod = mult(mod, mod, n):
   if (mod!=n-1&&temp%2==0) return false;
  return true:
Sieve of Eratosthenes
// generate sieve - O(n log n)
void genSieve(int n) {
 roid gensieve(int in )
sieve[0] = sieve[1] = 1;
for (ull i = 3; i * i < n; i += 2)
...if (!sieve[i])
...for (ull j = i * 3; j <= n; j += i * 2)</pre>
     .sieve[j] = 1;
// query sieve after it's generated - O(1)
bool querySieve(int n) {
 return n' == 2 | | (n \% 2 != 0 \&\& !sieve[n]):
Simpson's / Approximate Integrals
// integrate f from a to b, k iterations
// error \le (b-a)/18.0 * M * ((b-a)/2k)^2
// where M = max(abs(f^{(i)}(x))) for x in [a,b]
 // "f" is a function "double func(double x)"
double Simpsons (double a, double b, int k,
 \rightarrow double (*f)(double)) {
double dx = (b-a)/(2.0*k), t = 0;
 for (int i = 0; i < k; i++)

t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
 \rightarrow (*f)(a+(2*i+1)*dx);
 return (t + (*f)(b)) * (b-a) / 6.0 / k;
```

```
Common Equations Solvers
// ax^2 + bx + c = 0, find x
vector < double > solve Eq (double a, double b,
double c) {
.vector<double> r;
.double z = b * b - 4 * a * c;
 if (z == 0)
 r.push_back(-b/(2*a));
 else if (z > 0) {
   r.push_back((sqrt(z)-b)/(2*a));
  r.push_back((sqrt(z)+b)/(2*a));
// ax^3 + bx^2 + cx + d = 0, find x
vector < double > solve Eq (double a, double b,

    double c, double d) {
    vector<double> res;

 long double a1 = b/a, a2 = c/a, a3 = d/a;
 long double q = (a1*a1 - 3*a2)/9.0, sq =
\rightarrow -2*sqrt(q);
 long double r = (2*a1*a1*a1 - 9*a1*a2 +
\rightarrow 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
 if (z <= 0) {
  theta = acos(r/sqrt(q*q*q));
res.push_back(sq*cos(theta/3.0) - a1/3.0);
  res.push_back(sq*cos((theta+2.0*PI)/3.0) -
 \rightarrow a1/3.0):
  res.push_back(sq*cos((theta+4.0*PI)/3.0) -
   a1/3.0);
 res push_back(pow(sqrt(z)+fabs(r), 1/3.0));
res[0] = (res[0] + q / res[0]) *
   ((r<0)?1:-1) - a1 / 3.0:
 return res;
^{\prime}// m = # equations, n = # variables, a[m][n+1]
\rightarrow = coefficient matrix
// a[i][0]x + a[i][1]y + ... + a[i][n]z =
\begin{vmatrix} \Rightarrow & a[i][n+1] \\ const & double \\ eps = 1e-7; \end{vmatrix}
bool zero(double a) { return (a < eps) && (a >
→ -eps); }
vector < double > solveEq(double **a, int m, int
\rightarrow n) {
if (j != cur) swap(a[j], a[cur]);
     for (int sat = 0; sat < m; sat++) {
  if (sat == cur) continue;
      double num = a[sat][i] / a[cur][i];
for (int sot = 0; sot <= n; sot++)</pre>
       a[sat][sot] -= a[cur][sot] * num;
     cur++:
     break;
 for (int j = cur; j < m; j++)
  if (!zero(a[j][n])) return vector<double>();
 vector<double> ans(n,0);
 for (int i = 0, sat = 0; i < n; i++
    if (sat < m && !zero(a[sat][i]))
    ans[i] = a[sat][n] / a[sat++][i];
 return ans;
6 Graph
struct edge {
```

edge (int u,int v,int w) : u(u),v(v),w(w) {}

bool operator < (const edge &e1, const edge

edge (): u(0), v(0), w(0) {}

 $\rightarrow$  &e2) { return e1.w < e2.w: }

int u,v,w;

```
bool operator > (const edge &e1, const edge
                                                      return MST;
\rightarrow &e2) { return e1.w > e2.w; }
struct subset { int p, rank; };
                                                      Union Find
                                                      int uf_find(subset* s, int i) {
  if (s[i].p != i) s[i].p = uf_find(s, s[i].p);
Eulerian Path
#define edge_list vector<edge>
#define_adj_sets vector<set<int>
                                                       return s[i].p;
struct EulerPathGraph {
                                                      void uf_union(subset* s, int x, int y) {
adj_sets graph; // actually indexes incident
                                                      int xp = uf_find(s, x), yp = uf_find(s, y);
if (s[xp].rank > s[yp].rank) s[yp].p = xp;

→ edaes

 edge_list edges; int n; vector<int> indeg;
                                                       else if (s[xp].rank < s[yp].rank) s[xp].p =
 EulerPathGraph(int n): n(n) {
                                                      \rightarrow yp;
else { s[yp].p = xp; s[xp].rank++; }
 indeg = *(new vector<int>(n.0)):
 graph = *(new adj_sets(n, set<int>()));
 void add_edge(int u, int v) {
                                                           2D Geometry
  graph[u].insert(edges.size());
                                                      #define point complex<double>
  indeg[v]++;
                                                      double dot(point a, point b) { return
  edges.push_back(edge(u,v,0));

    real(conj(a)*b); }

                                                      double cross (point a, point b) { return
 bool eulerian_path(vector<int> &circuit) {

    imag(conj(a)*b); }

  if(edges.size()==0) return false;
 struct line { point a, b; };
struct circle { point c; double r; };
                                                     struct triangle { point a, b, c; };
                                                      struct rectangle { point tl, br; };
                                                      struct convex_polygon {
 ...if (abs(((int)indeg[v])-((int)graph[v]
                                                       vector<point points;
                                                       convex_polygon(triangle a) {
    .size())) > 1) return
                                                        points.push_back(a.a);
    false:
    if (a[b] != -1) return false;
                                                          points.push_back(a.b);
    a[b] = v;
                                                          points.push_back(a.c);
                                                       .};
                                                       convex_polygon(rectangle a) {
 int s = (a[0]!=-1 && a[1]!=-1 ? a[0] :
                                                       points.push_back(a.tl);
\rightarrow (a[0]==-1 && a[1]==-1 ? edges[0].u : -1));
                                                          points.push_back({real(a.tl),
  if(s==-1) return false;
                                                          imag(a.br)});
  while(!st.empty() || !graph[s].empty()) {
                                                        points.push_back(a.br);
 . if (graph[s].empty()) {
    circuit.push_back(s); s = st.top();
                                                          points.push_back({real(a.br),
\stackrel{\hookrightarrow}{\hookrightarrow} st.pop(); }
                                                          imag(a.tl)});
   else {
   int w = edges[*graph[s].begin()].v;
                                                      #define sq(a) ((a)*(a))
    graph[s].erase(graph[s].begin());
    st.push(s); s = w;
                                                      double circumference(circle a) { return 2 *
                                                      \rightarrow a.r * M_PI; }
                                                      double area(circle a) { return sq(a.r) * M_PI;
  circuit.push_back(s);
  return circuit.size()-1==edges.size();
                                                      double intersection(circle a, circle b) {
  double d = abs(a.c - b.c);
                                                       if (d <= b.r - a.r) return area(a);
                                                      if (d <= a.r - b.r) return area(b);
if (d >= a.r + b.r) return 0;
Minimum Spanning Tree
   returns vector of edges in the mst
                                                       double alpha = acos((sq(a.r) + sq(d) -
// graph[i] = vector of edges incident to
                                                       \Rightarrow sq(b.r)) / (2 * a.r * d));
vertex i
// places total weight of the mst in &total
                                                       double beta = acos((sq(b.r) + sq(d) -
                                                         sq(a.r)) / (2 * b.r * d));
// if returned vector has size != n-1, there is
                                                       return sq(a.r) * (alpha - 0.5 * sin(2 *
vector<edge> mst(vector<vector<edge> graph, 11
                                                          alpha) + sq(b.r) * (beta - 0.5 * sin(2 *
                                                         beta)):
priority_queue<edge, vector<edge>,
                                                      double intersection (rectangle a, rectangle b)

    greater<edge≫ pq;
vector<edge> MST;

                                                       double x1 = max(real(a.tl), real(b.tl)), y1 =
 bitset<20001> marked; // change size as
                                                         max(imag(a.tl), imag(b.tl));

\frac{\text{meeded}}{\text{marked}[0]} = 1;

                                                       double x2 = min(real(a.br), real(b.br)), y2 =
                                                      \rightarrow min(imag(a.br), imag(b.br));
return (x2 <= x1 || y2 <= y1) ? 0 :
for (edge ep : graph[0]) pq.push(ep);
while(MST.size()!=graph.size()-1 &&
                                                          (x2-x1)*(y2-y1);

→ pq.size()!=0) {
  edge e = pq.top(); pq.pop();
 int u = e.u, v = e.v, w = e.w;
if(marked[u] && marked[v]) continue;
else if(marked[u]) swap(u, v);
                                                           3D Geometry
                                                      struct point3d {
                                                       double x, y, z;
  for(edge ep : graph[u]) pq.push(ep);
                                                       point3d operator+(point3d a) const { return
  marked[u] = 1
  MST.push_back(e);
....pusn_back
..total += e.w;
                                                       \rightarrow {x+a.x, y+a.y, z+a.z}; }
                                                       point3d operator*(double a) const { return
```

 $\hookrightarrow$  {x\*a, y\*a, z\*a}; }

```
point3d operator-() const { return {-x, -y,
 \rightarrow -z}; }
 point3d operator-(point3d a) const { return
    *this + -a; }
 point3d operator/(double a) const { return
 \rightarrow *this * (1/a); }
 double norm() { return x*x + y*y + z*z; }
 double abs() { return sqrt(norm()); }
 point3d normalize() { return *this /

    this->abs(): }

double dot(point3d a, point3d b) { return
\rightarrow a.x*b.x + a.y*b.y + a.z*b.z; }
point3d cross(point3d a, point3d b) { return
    \{a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,
    a.x*b.v - a.v*b.x; }
struct line3d { point3d a, b; };
struct plane { double a, b, c, d; } // a*x +
\rightarrow b*y + c*z + d = 0
struct sphere { point3d c; double r; };
#define sq(a) ((a)*(a))
#define cb(a) ((a)*(a)*(a))
double surface(circle a) { return 4 * sq(a.r)
double volume(circle a) { return 4.0/3.0 *
\rightarrow cb(a.r) * M_PI; }
     Optimization
Snoob
 // SameNumberOfOneBits, next permutation
int snoob(int a) {
  int b = a & -a, c = a + b;
  return c | ((a ^ c) » 2) / b;
Powers
bool isPowerOf2(ll a) {
 return a > 0 && !(a & a-1):
bool isPowerOf3(11 a) {
return a>0&&!(12157665459056928801ull%a);
bool isPower(ll a, ll b) {
  double x = log(a) / log(b);
 return abs(x-round(x)) < 0.00000000001;
```