```
General
                              7 Graphs
    Algorithms
                              8 2D Geometry
    Structures
                              9 3D Geometry
    Strings
                              10 Optimization
    Greedy
                              11 Additional
    Math
     General
g++ -g -02 -std=gnu++17 -static prog.cpp
./a.exe
run.sh
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp for i i *.in; do
 f=${i%.in}
 ./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
done
Header
// use better compiler options
#pragma GCC optimize("Ofast","unroll-loops")
#pragma GCC target("avx2,fma")
// include everything
 #include <bits/stdc++.h>
#include <bits/extc++.h>
#include <sys/resource.h>
// namespaces
using namespace std;
using namespace __gnu_cxx; // rope
using namespace __gnu_pbds; // tree/trie
// common defines
#define fastio

→ ios base::sync with stdio(0);cin.tie(0);
#define nostacklim rlimit RZ; getrlimit(3,&RZ
    ):RZ.rlim cur=-1:setrlimit(3.&RZ):
#define DEBUG(v) cerr<< LINE <<": "<<#v<<" =
\Rightarrow "<<v<<'\n'; #define TIMER
**define il28 unsigned il28

#define ull unsigned ll
#define il28 unsigned ll
#define il28 unsigned il28
#define ld long double
// global variables
mt19937 rng((uint32_t)chrono::steady

    clock::now().time since epoch().count());

Fast IO
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
void read(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
  n = n * 10 + c - 0';
void read(int& n) {
   char c; n = 0; int s = 1;
   if ((c=getchar_unlocked())=='-') s = -1;
 else n = c - ^{\circ};
while ((c=getchar_unlocked())!=' '&&c!='\n')
 n = n * 10 + c - 0;

n *= s;
void read(ld& n) {
 char c; n = 0;
ld m = 0, o = 1; bool d = false; int s = 1;
if ((c=getchar_unlocked())=='-') s = -1;
 else if (c == .'.') d = true;
else n = c - '0';
 while ((c=getchar_unlocked())!=' '&&c!='\n') {|}
  if (c == '.') d = true;
else if (d) { m=m*10+c-'0'; o*=0.1; }
```

```
else n = n * 10 + c - '0':
 n = s * (n + m * o):
void read(double& n) {
 ld m; read(m); n = m;
void read(float& n) {
 ld m: read(m): n = m:
void read(string& s) {
 char c; s = "
 while((c=getchar unlocked())!=' '&&c!='\n')
bool readline(string& s) {
 char c; s = "";
while(c=getchar unlocked()) {
 if (c == '\n') return true;
if (c == EOF) return false;
s += c;
 return false;
void print(unsigned int n) {
 if (n / 10) print(n / 10);
 putchar_unlocked(n % 10 + '0');
void print(int n) {
 if (n < 0) { putchar_unlocked('-'); n*=-1; }
 print((unsigned int)n);
Common Structs
   n-dimension vectors
// Vec<2, int> v(n, m) = arr[n][m]

// Vec<2, int> v(n, m, -1) default init -1

template<int D, typename T>

struct Vec : public vector<Vec<D-1, T>> {
  template<typename... Args>
  Vec(int n=0, Args... args) : vector<Vec<D-1,
 \rightarrow T>>(n. Vec<D-1. T>(args...)) {}
template<typename T>
struct Vec<1, T> : public vector<T> {
  Vec(int n=0, T val=T()) : vector<T>(n, val)
    {}
     Algorithms
Binary Search
// search for k in [p,n)
template<typename T>
int binsearch(T x[], int k, int n, int p = 0) {
     for (int i = n; i >= 1; i /= 2)
          while (p+i < n \&\& x[p+i] <= k) p += i;
     return p; \frac{1}{bool}: x[p] = k;
Min/Max Subarray
   max - compare = a < b, reset = a < 0
 \frac{1}{2}/ min - compare = a > b. reset = a > 0
// returns {sum, {start, end}}
pair<int, pair<int, int>>
     ContiguousSubarrav(int* a. int size.
    bool(*compare)(int, int).
 bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
 \rightarrow 0, s = 0;
 for (int i = 0; i < size; i++) {
...cur += a[i];
  if ((*compare)(best, cur)) { best = cur;
 > start = s; end = i; }
if ((*reset)(cur)) { cur = 0: s = i + 1: }
 return {best, {start, end}};
Quickselect
```

```
int partition(int arr[], int 1, int r)
  int x = arr[r], i = 1;
 for (int j = 1; j <= r - 1; j++)
. if (arr[j] <= x)
  swap(arr[i++], arr[j]);
 swap(arr[i], arr[r]);
 return i:
// find k'th smallest element in unsorted array
→ only if all distinct
int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - 1 + 1)) return QSNE;
swap(arr[1 + rng() % (r-1+1)], arr[r]);
  int pos = partition(arr, 1, r);
 if (pos-l==k-1) return arr[pos];
 if (pos-1>k-1) return qselect(arr,1,pos-1,k);
 return qselect(arr, pos+1, r, k-pos+1-1);
// TODO: compare against std::nth_element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted
→ on both axis
pair<int, int> saddleback_search(int** arr, int
 \hookrightarrow x, int y, int v) {
 int i = x-1, j = 0;
 while (i >= 0 && j < y) {
   if (arr[i][j] == v) return {i, j};
  (arr[i][i] > v)? i--: j++;
 return {-1, -1};
Ternary Search
// < max, > min, or any other unimodal func #define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a, int b, int (*f)(int)) {
 while (b-a > 4) {
    int m = (a+b)/2;
    if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1:
 for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
 return á;
#define TERNPREC 0.000001
double ternsearch (double a, double b, double
 \rightarrow (*f)(double)) {
while (b-a > TERNPREC * 4) {
  double m = (a+b)/2;
if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
  else b = m + TERNPREC;
 for (double i = a + TERNPREC: i <= b: i +=

→ TERNPREC

      if (TERNCOMP((*f)(a), (*f)(i)))
   .a = i:
 return a;
Golden Section Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)</pre>
double goldsection(double a, double b, double
 double r = (sqrt(5)-1)/2, eps = 1e-7;

double x1 = b - r*(b-a), x2 = a + r*(b-a);

double f1 = f(x1), f2 = f(x2);
  while (b-a > eps)
  while (b-a > eps)

if (TERNCOMP(f2,f1)) {

. b = x2; x2 = x1; f2 = f1;

. x1 = b - r*(b-a); f1 = f(x1);
   a = x1; x1 = x2; f1 = f2; x2 = a + r*(b-a): f2 = f(x2):
 return a:
```

```
3 Structures
```

```
Fenwick Tree
// Fenwick tree, array of cumulative sums
\hookrightarrow O(\log n) updates, O(\log n) gets
struct Fenwick { int n: ll* tree:
 void update(int i, int val) {
 .++i;
while (i <= n) {
  tree[i] += val;</pre>
  i += i & (-i);
 Fenwick(int size) {
 n = size;
 tree = new ll[n+1];
for (int i = 1; i <= n; i++)
  tree[i] = 0;
 Fenwick(int* arr, int size) : Fenwick(size) {
 for (int i = 0; i < n; i++)
...update(i, arr[i]);
 ~Fenwick() { delete[] tree; }
 ll operator∏(int i) {
 if (i < 0 || i > n) return 0;
 while (i>0)
  sum += tree[i];
i -= i & (-i);
 return sum;
ll getRange(int a, int b) { return
    operator[](b) - operator[](a-1); }
Hashtable
```



```
Rope
                                                         // print things with prefix "1"
                                                         auto range = trie.prefix_range("1");
// O(\log n) insert, delete, concatenate
                                                         for (auto it = range.first; it !=
int main() {
 // generate rope
                                                         → range.second: it++)
 rope<int> v;
                                                          cout << *it << '
 for (int i = 0; i < 100; i++)
.v.push_back(i);
                                                        Wavelet Tree
 // move range to front
                                                        using iter = vector<int>::iterator;
 rope<int> copy = v.substr(10, 10);
v.erase(10, 10);
                                                        struct WaveletTree {
   Vec<2, int> C; int s;
 v.insert(copy.mutable_begin(), copy);
                                                          // sigma = highest value + 1
                                                         WaveletTree(vector<int>& a. int sigma) :
 // print elements of rope
for (auto it : v) cout << it << "";
                                                            s(sigma), C(sigma*2, 0) {
                                                          build(a.begin(), a.end(), 0, s-1, 1);
                                                         void build(iter b. iter e. int L. int U. int
Segment Tree
                                                          u) {
if (L == U) return
//max(a,b), min(a,b), a+b, a*b, qcd(a,b), a*b
struct SegmentTree {
                                                           int M = (L+U)/2;
 typedef int T;
                                                           C[u].reserve(e-b+1); C[u].push back(0);
 static constexpr T UNIT = INT_MIN;
                                                          for (auto it = b; it != e; ++it)
  C[u].push_back(C[u].back() + (*it<=M));</pre>
 T f(T a, T b) 
 if (a == UNIT) return b;
if (b == UNIT) return a;
                                                           auto p = stable_partition(b, e, [=](int
                                                            i){return i<=M;});
  return max(a,b);
                                                          build(b, p, L, M, u*2);
 int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
                                                          build(p, e, M+1, U, u*2+1);
                                                          // number of occurrences of x in [0,i)
\rightarrow n(n) {}
                                                         int rank(int x, int i) {
   int L = 0, U = s-1, u = 1, M, r;
   while (L != U) {
 SegmentTree(vector<T> arr)

    SegmentTree(arr.size()) {

 for (int i=0:i<arr.size():i++)
                                                           M = (L+U)/2;
r = C[u][i]; u*=2;

    update(i,arr[i]);

                                                           if (x <= M) i = r, U = M;
else i -= r, L = M+1, ++u;
 void update(int pos, T val) {
  for (s[pos += n] = val; pos /= 2;)
   s[pos] = f(s[pos * 2], s[pos*2+1]);
                                                          return i:
                                                          ^{\prime\prime} number of occurences of x in [l,r)
 T query(int b, int e) { // query [b, e)
                                                         int count(int x, int 1, int r) {
  return rank(x, r) - rank(x, 1);
  Tra = UNIT, rb = UNIT;
  for (b+=n, e+=n; b<=); b/=2, e/=2) {
    if (b % 2) ra = f(ra, s[b++]);
    if (e % 2) rb = f(s[--e], rb);
                                                         // kth smallest in [l, r)
int kth(int k, int l, int r) const {
int L = 0, U = s-1, u = 1, M, ri, rj;
  return f(ra, rb):
                                                          while (L != U) {
   M = (L+U)/2;
 T get(int p) { return query(p, p+1); }
                                                           ri = C[u][1]; rj = C[u][r]; u*=2;
                                                           if (k \le rj-ri)^{n}l = ri, r = rj, U = M;
Sparse Table
                                                           else k -= řj-rí, l -= ŕi, r -= ŕj,
template < class T> struct SparseTable {
                                                           L = M+1, ++u;
 vector<vector<T>> m;
                                                           return U:
 SparseTable(vector<T> arr) {
  m.push back(arr);
  for (int k = 1: (1<<(k)) <= size(arr): k++)
                                                         // # elements between [x,y] in [l, r)
                                                         mutable int L, U;
  m.push_back(vector<T>(size(arr)-(1<(k)+1));
                                                         int range(int x, int y, int 1, int r) const {
  for (int i = 0; i < size(arr)-(1<<k)+1; i
                                                          if (y < x \text{ or } r <= 1) return 0;
                                                          L = x; U = y;
 [k][i] = min(m[k-1][i],
                                                          return range(1, r, 0, s-1, 1);
\rightarrow m[k-1][i+(1<<(k-1))]:
}
// min of range [l,r]
                                                         int range(int 1, int r, int x, int y, int u)
                                                         → const {
                                                          if (y < L or U < x) return 0;
if (L <= x and y <= U) return r-l;
T query(int 1, int r) {
 int k = _-lg(r-l+1);
                                                          int M = (x+y)/2, ri = C[u][1], rj = C[u][r];
  return \min(m[k][1], m[k][r-(1<< k)+1]);
                                                          return range(ri, rj, x, M, u*2) + range(1-ri
}
};
                                                            r-rj, M+1, y, u*2+1);
                                                          ^{\prime}// # elements <= x in [l, r]
                                                         int lte(int x, int l, int r) {
  return range(INT_MIN, x, l, r);
typedef trie<string, null_type,

→ trie_string_access_traits<>,

 pat_trie_tag, trie_prefix_search_node_update>

→ trie_type;

int main() {
                                                             Strings
 // generate trie
 trie_type trie;
                                                        Aho Corasick
 for (int i = 0; i < 20; i++)
                                                           range of alphabet for automata to consider
 trie.insert(to string(i)); // true if new,
                                                           MAXC = 26, OFFC = 'a' if only lowercase
\hookrightarrow false if old
```

```
|const int MAXC = 256:
const int OFFC = 0:
struct aho_corasick {
  set<pair<int, int>> out;
  int fail; vector<int> go;
  state() : fail(-1), go(MAXC, -1) {}
 vector<state> s;
  int id = 0;
 aho_corasick(string* arr, int size) : s(1) {
  for (int i = 0; i < size; i++) {
   for (int c : arr[i]) {
   if (s[cur].go[c-OFFC] == -1) {
      s[cur].go[c-OFFC] = s.size();
      s.push back(state());
     cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++});
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
    s[0].go[c] = 0;
  queue int> sq;
for (int c = 0; c < MAXC; c++) {
   if (s[0].so[c] != 0) {
      ...s[s[0].so[c]].fail = 0;</pre>
    sq.push(s[0].go[c]);
  while (sq.size()) {
   int e = sq.front(); sq.pop();
   for (int c = 0; c < MAXC; c++) {
   if (s[e].go[c] != -1) {
      int failure = s[e].fail;
while (s[failure].go[c] == -1)
        failure = s[failure].fail;
      failure = s[failure].go[c];
      s[s[e].go[c]].fail = failure;
      for (auto length : s[failure].out)
s[s[e].go[c]].out.insert(length);
      sq.push(s[e].go[c]);
 // list of {start pos, pattern id}
  vector<pair<int, int>> search(string text)
  vector<pair<int, int>> toret;
  int cur = 0;
  for (int i = 0; i < text.size(); i++) {
  while (s[cur].go[text[i]-OFFC] == -1)
    cur = s[cur].fail;
cur = s[cur].go[text[i]-OFFC];
    if (s[cur].out.size())
    for (auto end : s[cur].out)
. toret.push_back({i - end.first + 1,
     end.second):
  return toret:
Boyer Moore
struct defint { int i = -1; };
vector<int> boyermoore(string txt, string pat)
 vector<int> toret; unordered_map<char, defint>string lcp(string* arr, int n, bool sorted =
 → badchar:
 int m = pat.size(), n = txt.size();
 for (int i = 0; i < m; i++) badchar[pat[i]].i
 \rightarrow = i;
int s = 0:
 while (s \leq n - m) {
  int j = m - 1:
```

while $(j \ge 0 \&\& pat[j] == txt[s + j]) j--;$

.if (j < 0) {

```
..toret.push_back(s);
   s += (s + m < n) ? m - badchar[txt[s +
   m]].i : 1:
 .} else
   s += \max(1, i - badchar[txt[s + i]].i):
 return toret:
English Conversion
const string ones[] = {"", "one", "two",
    "three", "four", "five", "six", "seven", "eight", "nine"};
const string teens[] ={"ten", "eleven",
    "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen";
const string tens[] = {"twenty", "thirty",
    "forty", "fifty", "sixty", "seventy",

    "eighty", "ninety"};
const string mags[] = {"thousand", "million",
     "billion", "trillion", "quadrillion", "quintillion", "sextillion",
    "septillion"};
string convert(int num, int carry) {
 if (num < 0) return "negative " +

    convert(-num, 0):

     (num < 10) return ones[num];
(num < 20) return teens[num % 10]
 if (num < 100) return tens[(num / 10) - 2] +
     (num\%10==0?"":"") + ones[num\%10]:
 if (num < 1000) return ones[num / 100]
     (num/100==0?"":" ") + "hundred" + (num%100==0?"":" ") + convert(num % 100,
    0);
 return convert(num / 1000, carry + 1) + " " +
     mags[carry] + " " + convert(num % 1000.
    0):
string convert(int num) {
 return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
   vector<int> toret;
 int m = txt.length(), n = pat.length();
 int next[n + 1];
 for (int i = 0; i < n + 1; i++)
  next[i] = 0;
 for (int i = 1; i < n; i++) {
  int j = next[i + 1];
  while (j > 0 && pat[j] != pat[i])
   j = next[j];
  if (j > 0 | pat[j] == pat[i])
   next[i + 1] = j + 1;
 for (int i = 0, j = 0; i < m; i++) {
  if (txt[i] == pat[j]) {
  if (++j == n)
    toret.push_back(i - j + 1);
  } else if (j > 0) {
...j = next[j];
 return toret;
Longest Common Prefix (array)
 // longest common prefix of strings in array
 → false) {
idise; laise; laif (n == 0) return "";
if (!sorted) sort(arr, arr + n);
string r = ""; int v = 0;
 while (v < arr[0].length() && arr[0][v] ==

    arr[n-1][v])
    r += arr[0][v++];

 return r:
```

```
for (auto c : s) v = (c - 'a' + 1) + v *
Longest Common Subsequence
                                                          → HASHER;
string lcs(string a, string b) {
                                                          return v:
 int m = a.length(), n = b.length();
                                                         const int MAXN = 1000001;
 int L[m+1][n+1];
 for (int i = 0; i <= m; i++) {
    for (int j = 0; j <= n; j++) {
        if (i == 0 || j == 0) L[i][j] = 0;
        ...else if (a[i-1] == b[j-1]) L[i][j] =
                                                         ull base [MAXN] = \{1\};
                                                         void genBase(int n) {
                                                          for (int i = 1; i <= n; i++)
| base[i] = base[i-1] * HASHER;

    L[i-1][j-1]+1;

                                                         struct advHash {
    else L[i][j] = \max(L[i-1][j], L[i][j-1]);
                                                          ull v, 1; vector <ull> wip;
                                                          advHash(string& s): v(0) {
                                                           wip = vector<ull>(s.length()+1);\
 // return L[m][n]; // length of lcs
                                                           wip[0] = 0;
 string out = "";
                                                           for (int i = 0; i < s.length(); i++)
   wip[i+1] = (s[i] - 'a' + 1) + wip[i] *
 int i = m - 1, j = n - 1;
 while (i >= 0 &\check{k} j >= 0) {
                                                             HASHER;
  if (a[i] == b[j]) {
                                                          1 = s.length(); v = wip[1];
   out = a[i--] + out;
                                                          ull del(int pos, int len) {
   return v - wip[pos+len]*base[1-pos-len] +
  else if (L[i][j+1] > L[i+1][j]) i--;
                                                             wip[pos]*base[1-pos-len];
  .else j--;
                                                          ull substr(int pos, int len) {
 return out;
                                                           return del(pos+len, (1-pos-len)) -
                                                             wip[pos]*base[len]:
Longest Common Substring
                                                          ull replace(int pos, char c) {
// l is array of palindrome length at that
                                                           return v - wip[pos+1]*base[l-pos-1] + ((c -
int manacher(string s, int* 1) {
                                                              'a' + 1) + wip[pos] *
 int n = s.length() * 2;
                                                             HASHER) *base[1-pos-1];
 for (int i = \tilde{0}, j = 0, k; i < n; i += k, j =
 \rightarrow max(j-k, 0)) {
                                                          ull replace(int pos, string s) {
                                                           // can't increase total string size
  while (i >= j \&\& i + j + 1 < n \&\& s[(i-j)/2]
 \Rightarrow == s[(i+j+1)/2]) i++;
                                                             wip[pos+s.size()]*base[l-pos-s.size()], c
  1[i] = j;
                                                             wip[pos];
 for (k = 1; i >= k && j >= k && l[i-k] !=
                                                           for (int i = 0; i < s.size(); i++)
c = (s[i]-'a'+1) + c * HASHER:
   j-k; k++)
| 1[i+k] = min(1[i-k], j-k);
                                                           return r + c * base[1-pos-s.size()];
 return *max element(1, 1 + n);
                                                         Subsequence Count
Cyclic Rotation (Lyndon)
                                                          '/ "banana", "ban" >> 3 (ban, ba..n, b..an)
// simple strings = smaller than its nontrivial
                                                         ull subsequences(string body, string subs) {
 int m = subs.length(), n = body.length();
// lyndon factorization = simple strings
                                                          if (m > n) return 0;
\hookrightarrow factorized
                                                          ull** arr = new ull*[m+1];
for (int i = 0; i <= m; i++) arr[i] = new
 // "abaaba" -> "ab", "aab", "a"
vector<string> duval(string s) {
                                                             ull[n+1];
                                                          for (int i = 1; i <= m; i++) arr[i][0] = 0;
for (int i = 0; i <= n; i++) arr[0][i] = 1;
 int n = s.length();
 vector<string> lyndon;
 for (int i = 0; i < n;) {
    int j = i+1, k = i;
    for (; j < n && s[k] <= s[j]; j++)
    if (s[k] < s[j]) k = i;
                                                         for (int i = 1; i <= m; i++)

for (int j = 1; j <= n; j++)

arr[i][j] = arr[i][j-1] + ((body[j-1] ==
                                                             subs[i-1])? arr[i-1][j-1] : 0);
   else k++:
                                                          return arr[m][n]:
  for (; i <= k; i += j - k)
   lyndon.push_back(s.substr(i,j-k));
                                                         Suffix Array + LCP
 return lyndon;
                                                         struct SuffixArray {
}
// lexicographically smallest rotation
                                                          vector<int> sa, lcp;
                                                          SuffixArray(string& s, int lim=256) {
 int n = s.length() + 1, k = 0, a, b;
int minRotation(string s) {
 int n = s.length(); s += s;
                                                           vector<int> x(begin(s), end(s)+1), y(n),
 auto d = duval(s); int i = 0, a = 0;
                                                          \rightarrow ws(max(n, lim)), rank(n);
 while (a + d[i].length() < n) a +=

    d[i++].length();

                                                           iota(begin(sa), end(sa), 0);
 while (i && d[i] == d[i-1]) a -=
                                                           for (int j = 0, p = 0; p < n; j = max(1, j *

    d[i--].length();

return a;
                                                          \rightarrow 2), lim = p) {
                                                            p = j; iota(begin(y), end(y), n - j);
for (int i = 0; i < (n); i++)</pre>
Hashing
                                                             if (sa[i] >= j)
                                                              y[p++] = sa[i] -
 #define HASHER 27
                                                            fill(begin(ws), end(ws), 0);
ull basicHash(string s) {
                                                            for (int i = 0; i < (n); i++) ws[x[i]]++;
 ull v = 0;
```

```
\rightarrow ws[i - 1];
   for (int i = n; i--;) sa[-ws[x[y[i]]]] =
    v[i]:
   j]) ? p - 1 : p++;
  for (int i = 1; i < (n); i++) rank[sa[i]] =
  for (int i = 0, j; i < n - 1; lcp[rank[i++]]
  Suffix Tree (Ukkonen's)
struct SuffixTree {
 // n = 2*len+10 or so
enum { N = 50010, ALPHA = 26 };
int toi(char c) { return c - 'a'; }
 string a;
 void ukkadd(int i, int c) { suff:
  if (r[v]<=q) {</pre>
  p[m]=p[v]; t[m][c]=m+1; t[m][toi(a[q])]=v;
   l[v]=q; p[v]=m; t[p[m]][toi(a[l[m]])]=m;
v=s[p[m]]; q=l[m];
   while (q < r[m]) = v = t[v][toi(a[q])];
    q+=r[v]-1[v]; }
   if (q==r[m]) s[m]=v; else s[m]=m+2;
   q=r[v]-(q-r[m]); m+=2; goto suff;
 SuffixTree(string a) : a(a) {
  fill(r,r+N,(int)(a).size());
  memset(s, 0, sizeof s);

memset(t, -1, sizeof t);

fill(t[1],t[1]+ALPHA,0);

s[0]=1;1[0]=1[1]=-1;r[0]=r[1]=p[0]=p[1]=0;
  for(int i=0;i<a.size();i++)</pre>
    ukkadd(i.toi(a[i])):
 // Longest Common Substring between 2 strings
  // returns {length, offset from first string}
  pair<int, int> best;
 int lcs(int node, int i1, int i2, int olen) {
  if (l[node] <= i1 && i1 < r[node] > return 1;
  if (l[node] <= i2 && i2 < r[node] > return 2;
  int mask=0
   len=node?olen+(r[node]-l[node]):0;
  for(int c=0; c<ALPHA; c++) if
 \rightarrow (t[node][c]!=-1)
\max_{n=0}^{\infty} |= lcs(t[node][c], i1, i2, len);
  if (mask==3)
 → best=max(best,{len,r[node]-len});
  return mask:
 static pair<int, int> LCS(string s, string t)
 \rightarrow st(s+(char)('z'+1)+t+(char)('z'+2));
  st.lcs(0, s.size(), s.size()+t.size()+1, 0);
return st.best;
String Utilities
```

|...for (int i = 1; i < (lim); i++) ws[i] +=

```
void lowercase(string& s) {
                                                 transform(s.begin(), s.end(), s.begin(),
                                                    ::tolower):
                                                 void uppercase(string& s) {
                                                 transform(s.begin(), s.end(), s.begin(),
                                                    ::toupper);
                                                 void trim(string &s) {
                                                 s.erase(s.begin(),find_if_not(s.begin(),s
                                                     .end(),[](int c){return
                                                    isspace(c);}));
                                                 s.erase(find_if_not(s.rbegin(),s.rend(),[](int

    c){return isspace(c);}).base(),s.end());

                                                 vector<string> split(string& s, char token) {
                                                     vector<string> v; stringstream ss(s);
                                                    for (string e;getline(ss,e,token);)
                                                        v.push_back(e);
                                                    return v;
                                                    \mathbf{Greedv}
                                                 Interval Cover

    vector<pair<double,double,int>> in) {

                                                    int i = 0; pair < double, int > pos = {L,-1};
                                                    vector<int> a;
sort(begin(in), end(in));
                                                     while (pos.first < R) {
                                                         double cur = pos.first;
while (i < (int)in.size() &&</pre>
                                                    in[i].first.first <= cur)</pre>
                                                    max(pos, {in[i].first.second,in[i].second}),
                                                    i++;
                                                         if (pos.first == cur) return {};
                                                         a.push back(pos.second):
                                                    return a;
                                                 6 Math
                                                 Catalan Numbers
                                                ull* catalan = new ull[1000000];
                                                 void genCatalan(int n, int mod) {
                                                 catalan[0] = catalan[1] = 1;
                                                 for (int i = 2; i <= n; i++) {
    catalan[i] = 0;
    for (int j = i - 1; j >= 0; j--) {
                                                   catalan[i] += (catalan[j] * catalan[i-j-1])
                                                    % mod;
                                                   if (catalan[i] >= mod)
   catalan[i] -= mod;
                                                 .
// TODO: consider binomial coefficient method
                                                 Combinatorics (nCr. nPr)
                                                 // can optimize by precomputing factorials, and
                                                    fact[n]/fact[n-r]
                                                 ull nPr(ull n, ull r) {
                                                 for (ull i = n-r+1: i <= n: i++)
                                                 ..v *= i;
.return v:
                                                ull nPr(ull n, ull r, ull m) {
                                                 for (ull i = n-r+1; i <= n; i++)
                                                 v = (v * i) \% m;
return v;
                                                 ull nCr(ull n, ull r) {
                                                 long double v = 1;
                                                 for (ull i = 1; i <= r; i++)
                                                 v = v * (n-r+i) /i;
return (ull) (v + 0.001);
```

```
|#define n 100000
// requires modulo math
                                                      ll phi[n+1];
// can optimize by precomputing mfac and
                                                       void computeTotient() {

→ minv-mfac

                                                        for (int i=1: i \le n: i++) phi[i] = i:
ull nCr(ull n, ull r, ull m) {
                                                        for (int p=2; p<=n; p++) {
                                                        if (phi[p] == p) {
 return mfac(n, m) * minv(mfac(k, m), m) % m *
 \rightarrow minv(mfac(n-k, m), m) % m:
                                                         phi[p] = p-1;
for (int i = 2*p; i<=n; i += p) phi[i] =</pre>
                                                           (phi[i]/p) * (p-1);
Multinomials
11 multinomial(vector<int>& v) {
   11 c = 1, m = v.empty() ? 1 : v[0];
 for(int i = 1; i < v.size(); i++)
 for (int j = 0; j < v[i]; j++)
...c = c * ++m / (j+1);
                                                       Factorials
                                                       // digits in factorial
 return c:
                                                       #define kamenetsky(n) (floor((n * log10(n /
                                                       \rightarrow ME)) + (loq10(2 * M_PI * n) / 2.0)) + 1)
                                                       // approximation of factorial
 Chinese Remainder Theorem
bool ecrt(ll* r, ll* m, int n, ll& re, ll& mo) #define stirling(n) ((n == 1) ? 1 : sqrt(2 *
\hookrightarrow M PI * n) * pow(n / M E, n))
                                                       // natural log of factorial
                                                       #define lfactorial(n) (lgamma(n+1))
  d = \gcd(mo, m[i], x, y);
  if ((r[i] - re) % d != 0) return false;
x = (r[i] - re) / d * x % (m[i] / d);
re += x * mo;
                                                       Prime Factorization
                                                       // do not call directly
                                                       ll pollard rho(ll n. ll s) {
  mo = mo / d * m[i];
                                                       .ll x, y;
  re %= mo;
                                                       x = y = rand() % (n - 1) + 1;
int head = 1, tail = 2;
while (true) {
 re = (re + mo) % mo;
return true;
                                                        x = mult(x, x, n);

x = (x + s) \% n;
                                                        if (x == y) return n;
Count Digit Occurences
                                                        11 d = _gcd(max(x - y, y - x), n);
if (1 < d && d < n) return d;
 /*count(n,d) counts the number of occurences of
 \rightarrow a digit d in the range [0,n]*/
                                                        if (++head == tail) y = x, tail <<= 1;
ll digit count(ll n. ll d) {
 ll result = 0:
 while (n != 0) {
result += ((n%10) == d ? 1 : 0);
                                                       // call for prime factors
                                                       void factorize(ll n. vector<ll> &divisor) {
  n /= 10;
                                                        if (n == 1) return;
                                                        if (isPrime(n)) divisor.push back(n);
 return result;
                                                        else {
    ll d = n:
11 count(ll n, ll d) {
                                                        while (\ddot{d} >= n) d = pollard_rho(n, rand() % (n)// 0-indexed, arbitrary k
 if (n < 10) return (d > 0 && n >= d);
if (n % 10) != 9) return digit_count(n, d) +
                                                           - 1) + 1);
                                                         factorize(n'/ d, divisor);
    count(n-1, d);
                                                        factorize(d, divisor);
 return 10*count(n/10, d) + (n/10) + (d > 0):
Discrete Logarithm
                                                       Farev Fractions
int discretelog(int a, int b, int m) {
                                                          generate 0 \le a/b \le 1 ordered, b \le n
    ll n = sqrt(m) + 1, an = 1;
for (ll i = 0; i < n; ++i)
                                                          farey(4) = 0/1 1/4 1/3 1/2 2/3 3/4 1/1
                                                       // length is sum of phi(i) for i = 1 to n
    an = (an * a) % m;
unordered map<11, 11> vals:
                                                       vector<pair<int, int>> farey(int n) {
                                                       int h = 0, k = 1, x = 1, y = 0, r;
vector<pair<int, int>> v;
    for (11 q = 0, cur = b; q \le n; q++) {
         vals[cur] = q;
         cur = (cur * a) \% m;
                                                        v.push_back({h, k});
                                                         r = (n-y)/k;
    for (ll p = 1, cur = 1; p <= n; p++) {
                                                         y += r*k; x' += r*h;
         cur = (cur * an) \% m:
                                                        x = -x; y = -y;

while (k > 1);
         if (vals.count(cur)) {
              int ans = n * p - vals[cur];
              return ans;
                                                        v.push_back({1, 1});
     return -1:
                                                       Fast Fourier Transform
Euler Phi / Totient
                                                       #define cd complex<doub
                                                       const double PI = acos(-1);
int phi(int n) {
                                                       void fft(vector<cd>& a, bool invert) {
 int r = n;
                                                        int n = a.size();
 for (int i = 2; i * i <= n; i++) {
   if (n % i == 0) r -= r / i;
                                                        for (int i = 1, j = 0; i < n; i++) {
  int bit = n >> 1:
  while (n % i == 0) n /= i;
                                                         for (; j & bit; bit >>= 1) j ^= bit;
                                                        .j ^= bit:
 if (n > 1) r = r / n;
return r;
                                                         if (i < j) swap(a[i], a[j]);
```

```
for (int len = 2; len <= n; len <<= 1) {
    double ang = 2 * PI / len * (invert ? -1 :
   1);
  cd wlen(cos(ang), sin(ang));
  for (int i = 0; i < n; i += len) {
   cd w(1):
   for (int j = 0; j < len / 2; j++) {
    cd u = a[i+j], v = a[i+j+len/2] * w;
    a[i+j] = u + v;
    a[i+j+len/2] = u - v;
    .w *= wlen;
 if (invert)
 for (auto\& x : a)
  x /= n;
vector<int> fftmult(vector<int> const& a,

    vector<int> const& b) {

 vector<cd> fa(a.begin(), a.end()),
 → fb(b.begin(), b.end());
 int n = 1 \ll (32 - \_builtin\_clz(a.size() +
\rightarrow b.size() - 1));
fa.resize(n); fb.resize(n);
 fft(fa, false); fft(fb, false)
 for (int i = 0; i < n; i++) fa[i] *= fb[i];
 fft(fa, true);
 vector<int> toret(n);
 for (int i = 0; i < n; i++) toret[i] =
 → round(fa[i].real());
 return toret:
Greatest Common Denominator
ll egcd(ll a, ll b, ll& x, ll& y) {
  if (b == 0) { x = 1; y = 0; return a; }
 11 gcd = egcd(b, a \% b, x, y);
 x = a / b * y;
 swap(x, y);
 return gcd;
Josephus Problem
int josephus(int n. int k) {
if (n == 1) return 0;
if (k == 1) return n-1;
 if (k > n) return (josephus(n-1,k)+k)%n;
 int res = josephus(n-n/k,k)-n\%k;
 return res + ((res<0)?n:res/(k-1)):
// fast case if k=2, traditional josephus
int josephus(int n) {
return 2*(n-(1<<(32-builtin clz(n)-1))):
Least Common Multiple
#define lcm(a,b) ((a*b)/qcd(a,b))
Modulo Operations
#define MOD 1000000007
#define madd(a,b,m) (a+b-((a+b-m>=0)?m:0))
#define mult(a,b,m) ((ull)a*b%m)
#define msub(a,b,m) (a-b+((a < b)?m:0))
ll mpow(ll b, ll e, ll m) {
11 x = 1;
 while (e > 0) {
  if (e % 2) x = (x * b) % m;
  b = (b * b) \% m;
  e /= 2;
 return x % m:
ull mfac(ull n, ull m) {
ull f = 1;
for (int i = n: i > 1: i--)
 f = (f * i) \% m;
return f:
```

```
// if m is not guaranteed to be prime
ll minv(ll b, ll m) {
11 min (11 b, 11 m) (11 l m) (11 x = 0, y = 0;

if (egcd(b, m, x, y) != 1) return -1;

return (x % m + m) % m;
11 mdiv_compmod(int a, int b, int m) {
 if (_gcd(b, m) != 1) return -1;
 return mult(a, minv(b, m), m):
 // if m is prime (like 10^9+7)
11 mdiv_primemod (int a, int b, int m) {
 return mult(a, mpow(b, m-2, m), m);
 // tonelli shanks = sqrt(n) % m, m is prime
ll legendre(ll a, ll m){
 if (a % m==0) return 0;
 if (m == 2) return 1;
 return mpow(a, (m-1)/2, m);
11 msqrt(11 n, 11 m) {
 ll s = __builtin_ctzll(m-1), q = (m-111)>>s,
 z = rand()\%(m-1)+1;
 if (m == 2) return 1;
if (s == 1) return mpow(n,(m+1)/411,m);
 while (legendre(z,m)!=m-1) z = rand()\%(m-1)+1;
 11 c = mpow(z,q,m), r = mpow(n,(q+1)/2,m), t
 \rightarrow = mpow(n,q,m), M = s;
 while (t != 1){
    ll i=1, ts = (t * t) % m;
  while (ts != 1) i++, ts = (ts * ts) % m;
  11 b = c;
  for (int'j = 0; j < M-i-1; j++) b = (b * b) %
 r = r * b \% m; c = b * b \% m; t = t * c \% m;
 \rightarrow M = i;
 return r:
Modulo Tetration
11 tetraloop(11 a, 11 b, 11 m) {
 if(b == 0 | | a == 1) return 1;
ll w = tetraloop(a,b-1,phi(m)), r = 1;
 for (;w;w/=2) {
  if (w&1)
  r *= a; if (r >= m) r -= (r/m-1)*m;
  a *= a; if (a >= m) a -= (a/m-1)*m;
 return r:
int tetration(int a, int b, int m) {
  if (a == 0 || m == 1) return ((b+1)&1)%m;
  return tetraloop(a,b,m) % m;
Matrix
template<typename T>
struct Mat : public Vec<2, T> {
 Mat(int x, int y) : Vec<2, T>(x, y), w(x),
 \rightarrow h(v) \{\}
 static Mat<T> identity(int n) { Mat<T> m(n,n);
    for (int i=0;i<n;i++) m[i][i] = 1; return
   m; }
 Mat<\hat{T}>\& operator+=(const Mat<T>\& m) {
  for (int i = 0; i < w; i++)
  for (int j = 0; j < h; j++)
(*this)[i][j] += m[i][j];
  return *this:
 Mat<T>& operator-=(const Mat<T>& m) {
  for (int i = 0; i < w; i++)
   for (int j = 0; j < h; j++)
    (*this)[i][j] -= m[i][j];
  return *this;
 Mat<T> operator*(const Mat<T>& m) {
  Mat < T > z(w,m.h);
  for (int i = 0; i < w; i++)
  for (int j = 0; j < h; j++)
for (int k = 0; k < m.h; k++)
```

```
z[i][k] += (*this)[i][j] * m[j][k];
                                                          for (i = 1; i \le c; i++) fac[i-1] = n\%i, n/=i; | sieve[0] = sieve[1] = 1;
                                                          for (i = c - 1; i >= 0; i--)

per[c-i-1] = idx[fac[i]],
    return z:
 Mat<T> operator+(const Mat<T>& m) { Mat<T>
                                                           idx.erase(idx.begin() + fac[i]);

→ a=*this: return a+=m: }

                                                          return per;
Mat<T> operator-(const Mat<T>& m) { Mat<T>

    a=*this; return a-=m; }

                                                          // get what nth permutation of vector
                                                         int get_permutation(vector<int>& v) {
Mat<T>& operator*=(const Mat<T>& m) { return
                                                          int use = 0, i = 1, r = 0;
for (int e: v) {
   r = r * i++ + __builtin_popcount(use &

→ *this = (*this)*m: }

 Mat<T> power(int n) {
 .Mat<T> a = Mat<T>::identity(w), m=*this;
                                                          \rightarrow -(1<<e));
  for (;n;n/=2,m*=m) if (n\&1) a *= m;
                                                           use |= 1 << e;
                                                          return r;
Matrix Exponentiation
                                                          Permutation (string/multiset)
// F(n) = c[\hat{0}]*F(n-1) + c[1]*F(n-2) + ...
// b is the base cases of same length c
                                                         string freq2str(vector<int>& v) {
ll matrix_exponentiation(ll n, vector<ll> c,
                                                          string s;
                                                          for (int i = 0; i < v.size(); i++)
for (int j = 0; j < v[i]; j++)
s += (char)(i + 'A');
vector<11> b) {
   if (nth < b.size()) return b[nth-1];
   Mat<11> a(c.size(), c.size()); ll s = 0;
   for (int i = 0; i < c.size(); i++) a[i][0] =</pre>

    c[i];

                                                          // nth perm of multiset, n is 0-indexed
 for (int i = 0; i < c.size() - 1; i++)
                                                         string gen permutation(string s, ll n) {
\rightarrow a[i][i+1] = 1;
                                                          vector<int> freq(26, 0);
 a = a.power(nth - c.size());
                                                          for (auto e : s) freq[e - 'A']++;
 for (int i = 0; i < c.size(); i++)
s += a[i][0] * b[i];
                                                          for (int i = 0; i < 26; i++) if (freq[i] > 0)
return s:
                                                           freq[i]--; 11 v = multinomial(freq);
                                                           if (n < v) return (char)(i+'A') +
Matrix Subarray Sums
                                                             gen_permutation(freq2str(freq), n);
                                                           freq[i]++; n -= v;
template<class T> struct MatrixSum {
 Vec<2, T> p;
                                                          return "":
 .MatrixSum(Vec<2, T>& v) {
. p = Vec<2,T>(v.size()+1, v[0].size()+1);
  for (int i = 0; i < v.size(); i++)
                                                         Miller-Rabin Primality Test
 for (int j = 0; j < v[0].size(); j++)
   p[i+1][j+1] = v[i][j] + p[i][j+1] +
                                                          // Miller-Rabin primality test - O(10 log^3 n)
                                                         bool isPrime(ull n) {

    p[i+1][i] - p[i][i];

                                                          if (n < 2) return false:
                                                          if (n < 2) return false;
if (n = 2) return true;
if (n % 2 == 0) return false;
ull s = n - 1;
while (s % 2 == 0) s /= 2;
for (int i = 0; i < 10; i++) {</pre>
 \tilde{T} sum(int u, int l, int d, int r) {
   return p[d][r] - p[d][l] - p[u][r] + p[u][l];
                                                           ull temp = s;
Mobius Function
                                                            ull a = rand() \% (n - 1) + 1;
const int MAXN = 10000000;
                                                           ull mod = mpow(a, temp, n);
// mu[n] = 0 iff n has no square factors
                                                            while (temp!=n-1\&\&mod!=1\&\&mod!=n-1) {
// 1 = even number prime factors, -1 = odd
                                                             mod = mult(mod, mod, n);
short mu[MAXN] = \{0,1\};
void mobius(){
 for (int i = 1; i < MAXN; i++)
                                                           if (mod!=n-1&&temp%2==0) return false;
 if (mu[i])
 for (int'j = i + i; j < MAXN; j += i)
                                                          return true;
    mu[j] -= mu[i];
                                                         Sieve of Eratosthenes
Nimber Arithmetic
                                                         bitset<100000001> sieve;
#define nimAdd(a,b) ((a)^(b))
                                                          // generate sieve - O(n log n)
ull nimMul(ull a, ull b, int i=6) {
  static const ull M[]={INT_MIN>>32,
                                                         void genSieve(int n) {
                                                          sieve[0] = sieve[1] = 1;
    M[0]^{(M[0] << 16)}, M[1]^{(M[1] << 8)},
                                                          for (ull i = 3; i * i < n; i += 2)
. if (!sieve[i])
   M[2]^(M[2] << 4), M[3]^(M[3] << 2),

\stackrel{\text{ML2J}}{\rightarrow} \underbrace{(\text{M[4]} <<1)}_{\text{if (i--} == 0) \text{ return a\&b;}}

                                                            for (ull j = i * 3; j <= n; j += i * 2)
                                                             sieve[j] = 1;
  int \k=1<<i;
  ull s=nimMul(a,b,i), m=M[5-i],
                                                          // query sieve after it's generated - O(1)
    t=nimMul(((a^(a>>k))&m)|(s&~m),
                                                         bool querySieve(int n) {
    ((b^(b>k))&m)|(m&(\sim m>>1))<< k, i);
                                                          return n == 2 || (n % 2 != 0 && !sieve[n]);
  return ((s^t)\&m)<< k | ((s^(t>>k))\&m);
                                                         Compile-time Prime Sieve
Permutation
                                                         const int MAXN = 100000:
// c = array size, n = nth perm, return index
                                                         template<int N>
vector<int> gen_permutation(int c, int n) {
                                                         struct Sieve
                                                          bool sieve[N]:
vector<int> idx(c), per(c), fac(c); int i;
for (i = 0; i < c; i++) idx[i] = i;
                                                          constexpr Sieve() : sieve() {
```

```
for (int i = 2; i * i < N; i++)
  if (!sieve[i])
    ...for (int j = i * 2; j < N; j += i)
    ...sieve[j] = 1;</pre>
bool isPrime(int n) {
   static constexpr Sieve<MAXN> s;
 return !s.sieve[n];
Simpson's / Approximate Integrals
// integrate f from a to b, k iterations
// error <= (b-a)/18.0 * M * ((b-a)/2k)^4

// where M = max(abs(f```(x))) for x in [a,b]

// "f" is a function "double func(double x)"
double Simpsons (double a, double b, int k,
 \rightarrow double (*f)(double)) {
double dx = (b-a)/(2.0*k), t = 0;
 for (int i = 0; i < k; i++)
t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
\leftrightarrow (*f)(a+(2*i+1)*dx);
return (t + (*f)(b)) * (b-a) / 6.0 / k;
Common Equations Solvers
// ax^2 + bx + c = 0, find x
vector < double > solveEq (double a, double b,
 double c) {
vector<double> r;
double z = b * b - 4 * a * c;
 if (z == 0)
  r.push_back(-b/(2*a));
 else if (z > 0) {
r.push_back((sqrt(z)-b)/(2*a));
  r.push_back((sqrt(z)+b)/(2*a));
 return r;
\frac{1}{2} / ax^3 + bx^2 + cx + d = 0, find x
vector<double> solveEq(double a, double b,
 → double c, double d) {
vector<double> res;
 long double a1 = b/a, a2 = c/a, a3 = d/a;
 long double q = (a1*a1 - 3*a2)/9.0, sq =
 \rightarrow -2*sqrt(q);
 long double r = (2*a1*a1*a1 - 9*a1*a2 +
  \rightarrow 27*a3)/54.0;
 long double z = r*r-q*q*q, theta;
 if (z <= 0) {
  theta = acos(r/sqrt(q*q*q));
res.push_back(sq*cos(theta/3.0) - a1/3.0);
  res.push back(sq*cos((theta+2.0*PI)/3.0) -
  res.push_back(sq*cos((theta+4.0*PI)/3.0) -
  → a1/3.0);
  res.push_back(pow(sqrt(z)+fabs(r), 1/3.0));
  res[0] = (res[0] + q / res[0]) *
 \leftrightarrow ((r<0)?1:-1) - a1 / 3.0;
 return res:
// linear diophantine equation ax + by = c,
\hookrightarrow find x and y
// infinite solutions of form x+k*b/g, y-k*a/g
bool solveEq(ll a, ll b, ll c, ll &x, ll &y, ll
 g = egcd(abs(a), abs(b), x, y);
 if (c % g) return false;
 x *= c / g * ((a < 0) ? -1 : 1);

y *= c / g * ((b < 0) ? -1 : 1);
 return true:
// m = # equations, n = # variables, a[m][n+1]
\hookrightarrow = coefficient matrix
// a[i][0]x + a[i][1]y + ... + a[i][n]z =
\rightarrow a[i][n+1]
```

```
|// find a solution of some kind to linear

→ equation

const double eps = 1e-7;
bool zero(double a) { return (a < eps) && (a >
 → -eps); }
vector < double > solveEq(double **a, int m, int
 \rightarrow n) { int cur = 0:
 for (int i = 0; i < n; i++) {
  for (int j = cur; j < m; j++) {
   if (!zero(a[j][i])) {
    if (j != cur) swap(a[j], a[cur]);
for (int sat = 0; sat < m; sat++) {
   if (sat == cur) continue;</pre>
       double num = a[sat][i] / a[cur][i];
      for (int sot = 0; sot <= n; sot++)
a[sat][sot] -= a[cur][sot] * num;
     cur++:
     break
 for (int j = cur; j < m; j++)
  if (!zero(a[j][n])) return vector<double>();
 vector<double> ans(n,0);
for (int i = 0, sat = 0; i < n; i++)
    if (sat < m && !zero(a[sat][i]))
    ans[i] = a[sat][n] / a[sat++][i];</pre>
 // solve A[n][n] * x[n] = b[n] linear equation
// rank < n is multiple solutions, -1 is no
→ solutions
// `alls` is whether to find all solutions, or
\hookrightarrow any
const double eps = 1e-12;
int solveEq(Vec<2, double>& A, Vec<1, double>&
 \rightarrow b, Vec<1, double>& x, bool alls=false) {
 int n = A.size(), m = x.size(), rank = 0, br,
 → bc;
 vector<int> col(m); iota(begin(col), end(col),
 for(int i = 0; i < n; i++) {
  double v, bv = 0;</pre>
  for(int r = i; r < n; r++)
  for(int c = i; c < n; c++)
    if ((v = fabs(A[r][c])) > bv)
        br = r, bc = c, bv = v;
    if (bv <= eps) {
   for(int j = i; j < n; j++)
    if (fabs(b[j]) > eps)
      return -1:
    break;
  swap(A[i], A[br]);
  swap(b[i], b[br]);
swap(col[i], col[bc]);
  for(int j = 0; j < n; j++)
swap(A[j][i], A[j][bc]);
  bv = 1.0 / A[i][i];
for(int j = (alls)?0:i+1; j < n; j++) {
   if (j != i) {
     double fac = A[j][i] * bv;
     b[j] = fac * b[i];
    for(int k = i+1; k < m; k++)
A[j][k] -= fac*A[i][k];
 rank++:
 if (alls) for (int i = 0; i < m; i++) x[i] =
    -DBL_MAX;
 for (int i = rank; i--;) {
   bool isGood = true;
  if (alls)
   for (int j = rank; isGood && j < m; j++)
...if (fabs(A[i][j]) > eps)
       isGood = false;
  b[i] /= A[i][i];
  if (isGood) x[col[i]] = b[i];
```

```
.if (!alls)
.for(int j = 0; j < i; j++)
..b[j] -= A[j][i] * b[i];</pre>
return rank;
```

Graycode Conversions

```
ull gravcode2ull(ull n) {
 for (; n; n = n >> 1) i = n;
return i;
ull ull2graycode(ull n) {
  return n ^ (n >> 1);
```

Unix/Epoch Time

```
// O-indexed month/time, 1-indexed day
// minimum 1970, 0, 1, 0, 0 ull toEpoch(int year, int month, int day, int
→ hour, int minute, int second) {
struct tm t; time_t epoch;
t.tm_year = year = 1900; t.tm_mon = month;
t.tm_mday = day; t.tm_hour = hour;
 t.tm_min = minute; t.tm_sec = second;
 t.tm_isdst = 0; // 1 = daylights savings
 epoch = mktime(&t);
return (ull)epoch;
vector<int> toDate(ull epoch) {
time_t e=epoch; struct tm t=*localtime(&e);
return {t.tm_year+1900,t.tm_mon,t.tm_mday,t_
   .tm hour.t.tm min.t.tm sec}:
int getWeekday(ull epoch) {
time_t e=epoch; struct tm t=*localtime(&e);
return t.tm wday; // 0-6, 0 = sunday
int getDavofYear(ull epoch) {
time t e=epoch; struct tm t=*localtime(&e);
return t.tm yday; // 0-365
const int months[] =
\rightarrow {31,28,31,30,31,30,31,30,31,30,31};
bool validDate(int year, int month, int day) {
    bool leap = !(year%(year%25?4:16));
if (month >= 12) return false;
return day <= months[month] + (leap &&
    month == 1);
```

Theorems and Formulae

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set of size n. !0 = 1, !1 = 0, !n = (n+1)(!(n-1))1)+!(n-2)), ! $n = n! \sum_{i=0}^{n} \frac{(-1)^i}{i!}$, ! $n = \left[\frac{n!}{e}\right]$ In a partially ordered set, a chain is a subset of

elements that are all comparable to eachother An antichain is a subset where no two are comparable.

Dilworth's theorem states the size of a max $\stackrel{\square}{\Rightarrow}$ imal antichain equals the size of a minimal chain cover of a partially ordered set S. The width of S is the maximum size of an antichain in S, which is equal to the minimum number of chains needed to cover S, or the minimum number of chains such that all elements are in at least one chain.

Rosser's Theorem states the nth prime Floyd Warshall number is greater than n * ln(n) for n > 1.

 $(n^{\frac{n+1}{2}})^2$.

every natural number is the sum of the squares of four non-negative integers. This is a specific (k,n) if (m[k][k] < 0) FOR(i,n) FOR(j,n) cial case of the Fermat Polygonal Number if (m[i][k] != inf && m[k][j] != inf) **Theorem** where every positive integer is a sum of at most n s-gonal numbers. The nths-gonal number $P(s,n) = (s-2)\frac{n(n-1)}{2} + n$

Graphs

```
struct edge {
int u, v, w;
edge (int u, int v, int w) : u(u), v(v), w(w) {}
 edge (): u(0), v(0), w(0) {}
};
|bool operator < (const edge &e1, const edge
→ &e2) { return e1.w < e2.w; }
bool operator > (const edge &e1, const edge
```

Eulerian Path

```
#define edge_list vector<edge>
#define adj_sets vector<set<int>>
struct EulerPathGraph {
adj_sets graph; // actually indexes incident
edge_list edges; int n; vector<int> indeg;
EulerPathGraph(int n): n(n) {
 indeg = *(new vector<int>(n.0));
 graph = *(new adj sets(n, set<int>()));
 void add_edge(int u, int v) {
  graph[u].insert(edges.size());
  indeg[v]++;
  edges.push back(edge(u,v,0));
 bool eulerian_path(vector<int> &circuit) {
 if(edges.size()==0) return false:
 stack<int> st;
int a[] = {-1, -1};
for(int v=0;v<n;v++) {
   if(indeg[v]!=graph[v].size()) {</pre>
   bool b = indeg[v] > graph[v].size();
if (abs(((int)indeg[v])-((int)graph[v]_)
     .size())) > 1) return
    false;
if (a[b] != -1) return false;
   a[b] = v;
  int s = (a[0]!=-1 && a[1]!=-1 ? a[0] :
   (a[0]==-1 && a[1]==-1 ? edges[0].u : -1));
  if(s==-1) return false:
  while(!st.empty() || !graph[s].empty()) {
   if (graph[s].empty()) {
    circuit.push_back(s); s = st.top();
   st.pop(); }
   else {
    int w = edges[*graph[s].begin()].v;
    graph[s].erase(graph[s].begin());
    st.push(s); s = w;
 circuit.push_back(s);
```

return circuit.size()-1==edges.size();

```
const 11 inf = 1LL << 62:
                                                                                                                                                                                                                                                                                                                                                        |\#define\ FOR(i,n)\ for\ (int\ i=0;\ i< n;\ i++)
Nicomachi's Theorem states 1^3 + 2^3 + \dots + void floydWarshall (Vec<2, 11>\& m) { <math>n^3 = (1+2+\dots+n)^2 and is equivalent to n^3 = (n^2+2+\dots+n)^2 and is equivalent to n^3 = (n^3+2+\dots+n)^2 and is eq
                                                                                                                                                                                                                                                                                                                                                             \rightarrow && m[k][j] != inf) {
  Lagrange's Four Square Theorem states auto newDist = max(m[i][k] + m[k][j], -inf);
                                                                                                                                                                                                                                                                                                                                                                 m[i][j] = min(m[i][j], newDist);
                                                                                                                                                                                                                                                                                                                                                                                 m[i][j] = -inf;
```

Minimum Spanning Tree

```
// returns vector of edges in the mst
 // araph[i] = vector of edges incident to
// if returned vector has size != n-1, there is
\hookrightarrow no MST
vector<edge> mst(vector<vector<edge>> graph,
total = 0:
 .total = 0;
.priority_queue<edge, vector<edge>,

    greater<edge>> pq;
vector<edge> MST;

 bitset<20001> marked; // change size as needed
marked[0] = 1;
for (edge ep : graph[0]) pq.push(ep);
while(MST.size()!=graph.size()-1 &&
 → pq.size()!=0) {
  edge e = pq.top(); pq.pop();
int u = e.u, v = e.v, w = e.w;
if(marked[u] && marked[v]) continue;
else if(marked[u]) swap(u, v);
   for(edge ep : graph[u]) pq.push(ep);
  marked[u] = 1;
MST.push_back(e);
  total += e.w;
 return MST:
```

Union Find

```
int uf_find(subset* s, int i) {
  if (s[i].p != i) s[i].p = uf_find(s, s[i].p);
 return s[i].p:
void uf_union(subset* s, int x, int y) {
int xp = uf_find(s, x), yp = uf_find(s, y);
if (s[xp].rank > s[yp].rank) s[yp].p = xp;
else if (s[xp].rank < s[yp].rank) s[xp].p =</pre>

    yp;
else { s[yp].p = xp; s[xp].rank++; }
```

2D Grid Shortcut

```
#define inbound(x,n) (0<=x&&x<x)
#define fordir(x,y,n,m) for(auto[dx,dy]:dir)if
\rightarrow (inbound(x+dx,n)\otimes Sinbound(y+dy,m))
const pair<int.int> dir[] =
\rightarrow \{\{1,0\},\{0,1\},\{-1,0\},\{0,-1\}\};
```

8 2D Geometry

```
#define point complex < double >
#define EPS 0.0000001
#define sq(a) ((a)*(a))
#define c\bar{b}(a) ((a)*(a)*(a))
double dot(point a, point b) { return

→ real(conj(a)*b);
}
double cross(point a, point b) { return

    imag(conj(a)*b); }

struct line { point a, b; };
struct circle { point c; double r; };
```

```
struct segment { point a, b; };
struct triangle { point a, b, c; };
struct rectangle { point tl, br; };
struct convex_polygon {
 vector<point points;
 convex_polygon(vector<point> points) :

    points(points) {}

 convex_polygon(triangle a) {
  points.push back(a.a); points.push back(a.b);
    points.push back(a.c);
 convex polygon(rectangle a) {
  points.push back(a.tl):
    points.push back({real(a.tl).
   imag(a.br)});
  points.push back(a.br):
    points.push back({real(a.br).
    imag(a.tl)}):
struct polygon {
 vector <point > points:
 polygon(vector<point> points) :
 → points(points) {}
 polygon(triangle a) {
  points.push_back(a.a); points.push_back(a.b);
    points.push back(a.c);
polygon(rectangle a) {
  points.push back(a.tl);
    points.push back({real(a.tl),
    imag(a.br)});
  points.push_back(a.br);
    points.push back({real(a.br),
    imag(a.tl)});
 polygon(convex_polygon a) {
  for (point v : a.points)
   points.push_back(v);
 // triangle methods
double area_heron(double a, double b, double
 \hookrightarrow c) {
if (a < b) swap(a, b);
 if (a < c) swap(a, c);
 if (b < c) swap(b, c);
 if (a > b + c) return -1;
return sqrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)
// segment methods
double lengthsq(segment a) { return
    sq(real(a.a) - real(a.b)) + sq(imag(a.a) -
    imag(a.b)); }
double length(segment a) { return
    sqrt(lengthsq(a)); }
    circle methods
double circumference(circle a) { return 2 * a.r
double area(circle a) { return sq(a.r) * M PI:
| \xrightarrow{} \}
// rectangle methods
double width(rectangle a) { return

→ abs(real(a.br) - real(a.tl)); }

double height(rectangle a) { return

→ abs(imag(a.br) - real(a.tl)); }

double diagonal (rectangle a) { return

    sqrt(sq(width(a)) + sq(height(a))); }

double area (rectangle a) { return width(a)
 → height(a); }
double perimeter(rectangle a) { return 2 *
 // check if `a` fit's inside `b
```

```
// swap equalities to exclude tight fits
                                                   // points of intersection
bool doesFitInside(rectangle a, rectangle b) { |vector<point> intersection(line a, circle c) { | → b*y + c*z + d = 0
                                                     vector point inter;
 int x = width(a), w = width(b), y = height(a),
                                                     c.c -= a.a;
a.b -= a.a;

→ h = height(b);

 if (x > y) swap(x, y);
if (w > h) swap(w, h);
                                                     point m = a.b * real(c.c / a.b);
                                                     double d2 = norm(m - c.c):
 if (w < x) return false;
if (y <= h) return true;</pre>
                                                     if (d2 > sq(c.r)) return 0;
                                                     double l = sqrt((sq(c.r) - d2) / norm(a.b));
 double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
                                                     inter.push back(a.a + m + 1 * a.b);
 return sq(a) \le sq(b) + sq(c);
                                                     if (abs(1) > EPS) inter.push_back(a.a + m - 1
                                                     \rightarrow * a.b);
// polygon methods
                                                     return inter:
// negative area = CCW. positive = CW
                                                                                                        Snoob
double area(polygon a) {
                                                     // area of intersection
                                                    double intersection(rectangle a, rectangle b) {
    int snoob(int a) {
        int b = a & -a, c = a + b;
        return c | ((a ^ c) >> 2) / b;
    }
}
  double area = 0.0; int n = a.points.size();
  for (int i = 0, j = 1; i < n; i++, j = (j - 1)
\hookrightarrow 1) % n)
                                                       max(imag(a.tl), imag(b.tl));
    area += (real(a.points[j]-a.points[i]))*
                                                     double x2 = min(real(a.br), real(b.br)), y2 =
    (imag(a.points[j]+a.points[i]));
                                                       min(imag(a.br), imag(b.br));
                                                     return (x2 <= x1 | | y2 <= y1) ? 0 :
  return area / 2.0;
                                                        (x2-x1)*(y2-y1);
// get both unsigned area and centroid
pair<double, point> area_centroid(polygon a) {
                                                    Convex Hull
 int n = a.points.size();
                                                    bool cmp(point a, point b) {
 double area = 0;
 point c(0, 0);
                                                     if (abs(real(a) - real(b)) > EPS) return
 for (int i = n - 1, j = 0; j < n; i = j++) {
                                                        real(a) < real(b);
 .double v = cross(a.points[i], a.points[i]) /
                                                     if (abs(imag(a) - imag(b)) > EPS) return
                                                                                                          v /= 2;
                                                        imag(a) < imag(b):
  area += v:
                                                     return false:
  c \leftarrow (a.points[i] + a.points[j]) * (v / 3);
                                                    convex_polygon convexhull(polygon a) {
 c /= area:
                                                     sort(a.points.begin(), a.points.end(), cmp);
 return {area, c};
                                                                                                        Powers
                                                     vector<point> lower, upper;
                                                     for (int i = 0; i < a.points.size(); i++) {
                                                      while (lower.size() >= 2 &&
Intersection
                                                        cross(lower.back() - lower[lower.size()
// -1 coincide, 0 parallel, 1 intersection
                                                        2], a.points[i] - lower.back()) < EPS)
int intersection(line a, line b, point& p) {
                                                       lower.pop_back();
 if (abs(cross(a.b - a.a, b.b - b.a)) > EPS) {
                                                      while (upper.size() >= 2 &&
 p = cross(b.a - a.a, b.b - a.b) / cross(a.b)
                                                        cross(upper.back() - upper[upper.size()
\rightarrow -a.a, b.b -b.a) * (b - a) + a;
                                                        2], a.points[i] - upper.back()) > -EPS)
  return 1:
                                                       upper.pop back():
 if (abs(cross(a.b - a.a, a.b - b.a)) > EPS)
                                                      lower.push_back(a.points[i]);
→ return 0:
                                                      upper.push_back(a.points[i]);
 return -1:
                                                     lower.insert(lower.end(), upper.rbegin() + 1,
// area of intersection
                                                        upper.rend()):
double intersection(circle a, circle b) {
                                                     return convex_polygon(lower);
 double d = abs(a.c - b.c);
if (d <= b.r - a.r) return area(a);</pre>
 if (d <= a.r - b.r) return area(b);
                                                        3D Geometry
 if (d >= a.r + b.r) return 0;
double alpha = acos((sq(a.r) + sq(d) -
                                                    struct point3d {
   sq(b.r)) / (2 * a.r * d));
                                                     double x, y, z;
                                                                                                         cout << p;
 double beta = acos((sq(b.r) + sq(d) - sq(a.r))
                                                     point3d operator+(point3d a) const { return
\rightarrow / (2 * b.r * d)):
                                                     \rightarrow {x+a.x, y+a.y, z+a.z}; }
 return sq(a.r) * (alpha - 0.5 * sin(2 *
                                                     point3d operator*(double a) const { return
    alpha)) + sq(b.r) * (beta - 0.5 * sin(2 *
                                                       \{x*a, y*a, z*a\}; \}
    beta)):
                                                     point3d operator-() const { return {-x, -y,
}
// -1 outside, 0 inside, 1 tangent, 2
                                                       -z}; }
                                                     point3d operator-(point3d a) const { return
\stackrel{\hookrightarrow}{\text{int}} intersection(circle a, circle b,
                                                     \rightarrow *this + -a; }
                                                     point3d operator/(double a) const { return

    vector<point>& inter) {

                                                     * *this * (1/a); }
double norm() { return x*x + y*y + z*z; }
 double d2 = norm(b.c - a.c), rS = a.r + b.r,
\rightarrow rD = a.r - b.r;
                                                     double abs() { return sqrt(norm()); }
 if (d2 > sq(rS)) return -1;
                                                     point3d normalize() { return *this /
 if (d2 < sq(rD)) return 0;
                                                       this->abs(): }
 double ca = 0.5 * (1 + rS * rD / d2);
 point z = point(ca, sqrt(sq(a.r) / d2 -
                                                    double dot(point3d a, point3d b) { return
                                                    \rightarrow a.x*b.x + a.y*b.y + a.z*b.z; }
 inter.push_back(a.c + (b.c - a.c) * z);
                                                    point3d cross(point3d a, point3d b) { return
 if (abs(imag(z)) > EPS) inter.push_back(a.c +
                                                       \{a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,
\rightarrow (b.c - a.c) * coni(z)):
return inter.size();
                                                       a.x*b.y - a.y*b.x; }
```

struct line3d { point3d a, b; };

```
struct plane { double a, b, c, d; } // a*x +
struct sphere { point3d c; double r; };
#define sq(a) ((a)*(a))
#define c\bar{b}(a) ((a)*(a)*(a))
double surface(circle a) { return 4 * sq(a.r)
double volume(circle a) { return 4.0/3.0 *
\hookrightarrow cb(a.r) * M PI: }
10 Optimization
 // SameNumberOfOneBits, next permutation
int main() {
   char 11[] = {'1', '2', '3', '4', '5'};
   char 12[] = {'a', 'b', 'c', 'd'};
   int d1 = 5, d2 = 4;
   // prints 12345abcd, 1234a5bcd, ...
 int min = (1 < < d1) - 1, max = min << d2;
 for (int i = min; i <= max; i = snoob(i)) {
  int p1 = 0, p2 = 0, v = i;
  while (p1 < d1 || p2 < d2) {
cout << ((v & 1) ? l1[p1++] : l2[p2++]);
  cout << '\n':
bool isPowerOf2(ll a) {
return a > 0 && !(a & a-1);
bool isPowerOf3(11 a) {
 return a>0&&!(12157665459056928801ull%a);
bool isPower(ll a, ll b) {
  double x = log(a) / log(b);
 return abs(x-round(x)) < 0.00000000001;
11 Additional
Judge Speed
    kattis: 0.50s
   codeforces: 0.421s
// atcoder: 0.455s
#include <bits/stdc++.h>
using namespace std;
int v = 1e9/2, p = 1;
int main() {
  for (int i = 1; i <= v; i++) p *= i;</pre>
Judge Pre-Contest Checks
     int128 and
                        float 128 support?
 -does extra or missing whitespace cause WA?
-documentation up to date?
-printer usage available and functional?
// each case tests a different fail condition
// try them before contests to see error codes
struct g { int arr[1000000]; g(){}};
vector<g> a;
// O=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE
→ 6=SIGSEGV 7=recursive MLE int judge(int n) {
 if (n == 0) exit(0); if (n == 1) while(1);
 if (n == 2) while(1) a.push_back(g());
 if (n == 3) while(1) putchar_unlocked('a');
```

if (n == 4) assert(0);

```
if (n == 5) 0 / 0;
 if (n == 6) * (int*)(0) = 0;
return n + judge(n + 1);
GCC Builtin Docs
 // 128-bit integer
__int128 a;
unsigned __int128 b;
// 128-bit float
// minor improvements over long double
 float128 c;
// log2 floor
__lg(n);
   number of 1 bits
// can add ll like popcountll for long longs
__builtin_popcount(n);
// number of trailing zeroes
__builtin_ctz(n);
// number of leading zeroes
__builtin_clz(n);
 7/1-indexed least significant 1 bit
__builtin_ffs(n);
// parity of number
__builtin_parity(n);
Limits
                       \pm 2147483647 \mid \pm 2^{31} - 1 \mid 10^9
int
                                            \frac{1}{2}<sup>32</sup> -\frac{1}{1}<sup>10</sup><sup>9</sup>
uint
                         4294967295
        \pm 922337203\overline{6854775807} | \pm \overline{2}^{63} - \overline{1}| \overline{10}^{18}
11
                                           \frac{1}{2}64 - \frac{1}{1}10<sup>19</sup>
         18446744073709551615
ull
|128| \pm 170141183460469231... | \pm 2^{127} - 1 |10^{38}|
                                          \frac{1}{2}^{128} - \frac{1}{10}^{10}^{38}
u128 340282366920938463...
Complexity classes input size (per second):
O(n^n) or O(n!)
                                                     n < 10
O(2^n)
                                                     n < 30
O(n^3)
                                                  n < 1000
O(n^2)
                                                n < 30000
```

 $n < 10^6$

 $n < 10^7$

 $n < 10^9$

 $O(n\sqrt{n})$

O(n)

 $O(n \log n)$