```
else n = n * 10 + c - '0':
    General
                             7 Graphs
    Algorithms
                              8 2D Geometry
                                                            n = s * (n + m * o):
    Structures
                              9 3D Geometry
                                                           void read(double& n) {
    Strings
                                                            ld m; read(m); n = m;
                              10 Optimization
    Greedy
                              11 Additional
                                                           void read(float& n) {
 ld m: read(m): n = m:
    Math
     General
                                                            void read(string& s) {
                                                            char c; s = "
g++ -g -02 -std=gnu++17 -static prog.cpp
./a.exe
run.sh
                                                            while((c=getchar unlocked())!=' '&&c!='\n')
                                                            bool readline(string& s) {
                                                            char c; s = "";
while(c=getchar unlocked()) {
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp
for i in *.in; do
                                                             if (c == '\n') return true;
if (c == EOF) return false;
s += c;
 f=${i%.in}
 ./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
                                                            return false;
done
                                                            void print(unsigned int n) {
Header
                                                            if (n / 10) print(n / 10);
// use better compiler options
                                                            putchar_unlocked(n % 10 + '0');
#pragma GCC optimize("Ofast","unroll-loops")
#pragma GCC target("avx2,fma")
                                                            void print(int n) {
// include everything
                                                            if (n < 0) { putchar_unlocked('-'); n*=-1; }
 #include <bits/stdc++.h>
                                                            print((unsigned int)n);
#include <bits/extc++.h>
#include <sys/resource.h>
// namespaces
                                                           Common Structs
using namespace std;
                                                               n-dimension vectors
using namespace __gnu_cxx; // rope
                                                               Vec<2, int>v(n, m) = arr[n][m]
using namespace __gnu_pbds; // tree/trie
                                                            // Vec<2, int> v(n, m, -1) default init -1
                                                            template<int D, typename T>
// common defines
#define fastio
                                                            struct Vec : public vector < Vec < D-1, T >> {
                                                              template<typename... Args>

→ ios base::sync with stdio(0);cin.tie(0);
                                                              Vec(int n=0, Args... args) : vector<Vec<D-1,
#define nostacklim rlimit RZ; getrlimit(3,&RZ
                                                            \rightarrow T>>(n, Vec<D-1, T>(args...)) {}
    ):RZ.rlim cur=-1:setrlimit(3.&RZ):
#define DEBUG(v) cerr<< LINE <<": "<<#v<<" =
                                                           template<typename T>
\Rightarrow "<<v<<'\n'; #define TIMER
                                                           struct Vec<1, T> : public vector<T> {
                                                              Vec(int n=0, T val=T()) : vector<T>(n, val)

→ cerr<<1.0*clock()/CLOCKS_PER_SEC<<"s\n";
#define ll long long
#define ull unsigned ll
#define i128 __int128
#define u128 unsigned i128
                                                               {}
                                                                Algorithms
#define ld long double
                                                           Min/Max Subarray
// global variables
                                                              max - compare = a < b, reset = a < 0
mt19937 rng((uint32_t)chrono::steady
                                                            \frac{1}{min} - compare = a > b, reset = a > 0

    clock::now().time since epoch().count());

                                                           // returns {sum, {start, end}}
pair<int, pair<int, int>>
Fast IO
                                                                ContiguousSubarray(int* a, int size,
#ifdef _WIN32
                                                                bool(*compare)(int, int),
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
                                                            bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
                                                            0, s = 0;
for (int i = 0; i < size; i++) {
  cur += a[i];</pre>
void read(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                              if ((*compare)(best, cur)) { best = cur;
  n = n * 10 + c - 0';
                                                            \rightarrow start = s; end = i; }
void read(int& n) {
  char c; n = 0; int s = 1
                                                             if ((*reset)(cur)) { cur = 0; s = i + 1; }
 if ((c=getchar_unlocked())=='-') s = -1;
                                                            return {best, {start, end}}:
 else n = c - '0';
while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                            Quickselect
 n = n * 10 + c - 0';
                                                           #define OSNE -999999
                                                           int partition(int arr[], int 1, int r)
void read(ld& n) {
 char c; n = 0;
ld m = 0, o = 1; bool d = false; int s = 1;
if ((c=getchar_unlocked())=='-') s = -1;
                                                            int x = arr[r], i = 1;
for (int j = 1; j <= r - 1; j++)
...if (arr[j] <= x)
...swap(arr[i++], arr[j]);</pre>
 else if (c == .'.') d = true;
else n = c - '0';
 while ((c=getchar_unlocked())!=' '&&c!='\n') {
                                                            swap(arr[i], arr[r]);
 if (c == '.') d = true;
else if (d) { m=m*10+c-'0'; o*=0.1; }
                                                            return i:
```

```
// find k'th smallest element in unsorted array, void update(int i, int val) {
→ only if all distinct
int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - l + 1)) return QSNE;
swap(arr[1 + rng() % (r-l+1)], arr[r]);
 int pos = partition(arr, 1, r);
if (pos-l==k-1) return arr[pos];
 if (pos-1>k-1) return qselect(arr,1,pos-1,k);
 return qselect(arr, pos+1, r, k-pos+1-1);
// TODO: compare against std::nth_element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted

    on both axis
pair<int, int> saddleback_search(int** arr, int
 \rightarrow x, int y, int v) {
 int i = x-1, j = 0;
 while (i >= 0 && j < y) {
  if (arr[i][j] == v) return {i, j};
  (arr[i][j] > v)? i--: j++;
 return {-1, -1};
Ternary Search
 // < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a, int b, int (*f)(int)) {</pre>
 while (b-a > 4) {
    int m = (a+b)/2;
    if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1:
 for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
   a = i;
 return a;
#define TERNPREC 0.000001
double ternsearch(double a. double b. double
 \leftrightarrow (*f)(double)) {
while (b-a > TERNPREC * 4) {
  double m = (a+b)/2;
  if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
  else b = m + TERNPREC;
 for (double i = a + TERNPREC: i <= b: i +=
     TERNPREC)
      if (TERNCOMP((*f)(a), (*f)(i)))
 return a;
Golden Section Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)</pre>
double goldsection(double a, double b, double
 while (b-a > eps)
  while (b-a > eps)

if (TERNCOMP(f2,f1)) {

. b = x2; x2 = x1; f2 = f1;

. x1 = b - r*(b-a); f1 = f(x1);
  } else {
   a = x1; x1 = x2; f1 = f2;

x2 = a + r*(b-a); f2 = f(x2);
 return a:
3 Structures
Fenwick Tree
// Fenwick tree, array of cumulative sums -
```

 \hookrightarrow O(log n) updates, O(log n) gets

struct Fenwick { int n; ll* tree;

```
while (i <= n) {
   tree[i] += val;
   i += i & (-i);
 Fenwick(int size) {
  n = size;
  tree = new ll[n+1];
for (int i = 1; i <= n; i++)
   .tree[i] = 0;
 Fenwick(int* arr, int size) : Fenwick(size) {
  for (int i = 0; i < n; i++)
...update(i, arr[i]);
 ~Fenwick() { delete[] tree; }
 ll operator[](int i) {
  if (i < 0 || i > n) return 0;
  \overline{11} \ \overline{sum} = 0;
  while (i>0)
   sum += tree[i];
   i -= i & (-i):
  return sum:
 ll getRange(int a, int b) { return

    operator[](b) - operator[](a-1); }

Hashtable
// similar to unordered map, but faster
struct chash {
    const uint64 t C = (11)(2e18 * M PI) + 71;
 ll operator()(ll x) const { return
    builtin bswap64(x*C); }
int main() {
  gp_hash_table<11,int,chash>
 \rightarrow hashtable({},{},{},{},{1<<16});
 for (int i = 0; i < 100; i++)
hashtable[i] = 200+i;
 if (hashtable.find(10) != hashtable.end())
   cout << hashtable[10];</pre>
Ordered Set
template <typename T>
using oset = tree<T,null_type,less<T>,rb_tree
    _tag,tree_order_statistics_node_update>;
template <typename T, typename D> using omap = tree<T,D,less<T>,rb_tree
    _tag,tree_order_statistics_node_update>;
int main()
 oset<int> o_set;
o_set.insert(5); o_set.insert(1);
 → o_set.insert(3);
// get second smallest element
 cout << *(o set.find by order(1));</pre>
 // number of elements less than k=4
cout << ' ' << o_set.order_of_key(4) << '\n';</pre>
 // equivalent with ordered map
 omap<int,int> o_map;
o_map[5]=1;o_map[1]=2;o_map[3]=3;
 cout << (*(o_map.find_by_order(1))).first;</pre>
 cout << ' ' << o map.order of key(4) << '\n';
Rope
// O(log n) insert, delete, concatenate
int main() {
 // generate rove
 rope<int> v;
 for (int i = 0: i < 100: i++)
  v.push_back(i);
 // move range to front
 rope<int> copy = v.substr(10, 10);
 v.erase(10, 10);
```

```
v.insert(copy.mutable_begin(), copy);
 // print elements of rope
for (auto it : v)
cout << it << "":
Segment Tree
//max(a,b), min(a,b), a+b, a*b, qcd(a,b), a*b
struct SegmentTree {
 typedef int T;
 static constexpr T UNIT = INT MIN:
 T f(T a, T b) {
 if (a == UNIT) return b;
if (b == UNIT) return a;
 return max(a,b);
 Int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
\rightarrow n(n) {}
 SegmentTree(vector<T> arr)

→ SegmentTree(arr.size()) {
 for (int i=0:i<arr.size():i++)

→ update(i.arr[i]):

 void update(int pos, T val) {
  for (s[pos += n] = val; pos /= 2;)
   s[pos] = f(s[pos * 2], s[pos*2+1]);
 T query(int b, int e) { // query [b, e)
 Tra = UNIT, rb = UNIT;

for (b+=n, e+=n; b<e; b/=2, e/=2) {

    if (b % 2) ra = f(ra, s[b++]);

    if (e % 2) rb = f(s[--e], rb);
  return f(ra. rb):
 T get(int p) { return query(p, p+1); }
Sparse Table
template<class T> struct SparseTable {
 vector<vector<T>> m;
SparseTable(vector<T> arr) {
  m.push_back(arr);
  for (int k = 1; (1<<(k)) <= size(arr); k++)
   m.push back(vector<T>(size(arr)-(1<<k)+1)):
   for (int i = 0; i < size(arr)-(1 << k)+1; i
    m[k][i] = min(m[k-1][i],
   m[k-1][i+(1<<(k-1))]:
 }
// min of range [l,r]
T query(int 1, int r) {
  int k = __lg(r-l+1);
  return \min(m[k][1], m[k][r-(1 << k)+1]):
typedef trie<string, null_type,

→ trie_string_access_traits<>,

 pat_trie_tag, trie_prefix_search_node_update>
int main() {
 // generate trie
 trie_type trie;
for (int i = 0; i < 20; i++)
...trie.insert(to_string(i)); // true if new,
\hookrightarrow false if old
 // print things with prefix "1"
 auto range = trie.prefix_range("1");
 for (auto it = range.first; it !=

    range.second; it++)

  cout << *it <<
Wavelet Tree
using iter = vector<int>::iterator;
struct WaveletTree {
```

```
Vec<2, int> C: int s:
 // sigma = highest value + 1
 WaveletTree(vector<int>& a, int sigma) :
    s(sigma), C(sigma*2, 0) {
  build(a.begin(), a.end(), 0, s-1, 1);
 void build(iter b, iter e, int L, int U, int
  if (L == U) return;
  int M = (L+U)/2:
  C[u].reserve(e-b+1); C[u].push_back(0);
  for (auto it = b; it != e; ++it)
    C[u].push_back(C[u].back() + (*it<=M));
  auto p = stable_partition(b, e, [=](int
    i) {return i <= M: }):
  build(b, p, L, M, u*2);
  build(p, e, M+1, U, u*2+1);
 // number of occurences of x in [0,i)
if (x <= M) i = r, U = M;
else i -= r, L = M+1, ++u;
  return i:
 // number of occurrences of x in [l,r)
int count(int x, int l, int r) {
  return rank(x, r) - rank(x, 1);
 // kth smallest in [l, r)
int kth(int k, int l, int r) const {
int L = 0, U = s-1, u = 1, M, ri, rj;
  while (L != U) {
   M = (L+U)/2;
   ri = C[u][1]; rj = C[u][r]; u*=2;

if (k <= rj-ri) 1 = ri, r = rj, U = M;
   else k -= ři-ri. l -= ři. r -= ři.
   L = M+1. ++u:
  return U;
  // # elements between [x,y] in [l, r)
 mutable int L. U:
 int range(int x, int y, int 1, int r) const {
  if (y < x \text{ or } r \le 1) return 0;
  L = x; U = y;
  return range(1, r, 0, s-1, 1);
 int range(int 1, int r, int x, int y, int u)
    const {
  if (y < L or U < x) return 0;
  if (L \le x \text{ and } y \le U) \text{ return } r-1;
  int M = (x+y)/2, ri = C[u][1], rj = C[u][r];
  return range(ri, rj, x, M, u*2) + range(1-ri, Boyer Moore
    r-rj, M+1, y, u*2+1);
 // # elements <= x in [l, r)
int lte(int x, int l, int r) {
  return range(INT_MIN, x, l, r);</pre>
     Strings
Aho Corasick
// range of alphabet for automata to consider
// MAXC = 26, OFFC = 'a' if only lowercase
const int MAXC = 256;
const int OFFC = 0;
struct aho_corasick {
 struct state
  set<pair<int, int>> out;
  int fail; vector<int> go;
  state() : fail(-1), go(MAXC, -1) {}
 vector<state> s;
```

```
int id = 0:
 aho corasick(string* arr, int size) : s(1) {
 for (int i = 0; i < size; i++) {
   int cur = 0;
   for (int c : arr[i]) {
   if (s[cur].go[c-OFFC] == -1) {
     s[cur].go[c-OFFC] = s.size();
     s.push_back(state());
    cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++}):
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
    s[0].go[c] = 0;
  queue<int> sq;
  for (int c = 0; c < MAXC; c++) {
   if (s[0].go[c] != 0) {
    s[s[0].go[c]].fail = 0;
    sq.push(s[0].go[c]);
  while (sq.size()) {
   int e = sq.front(); sq.pop();
  for (int c = 0; c < MAXC; c++) {
   if (s[e].go[c] != -1) {
     int failure = s[e].fail;
while (s[failure].go[c] == -1)
     failure = s[failure].fail;
failure = s[failure].go[c];
     s[s[e].go[c]].fail = failure;
     for (auto length : s[failure].out)
  s[s[e].go[c]].out.insert(length);
      sq.push(\bar{s}[e].go[c]);
 // list of {start pos, pattern id}
 vector<pair<int, int>> search(string text)
 vector<pair<int, int>> toret;
  int cur = 0;
  for (int i = 0; i < text.size(); i++) {</pre>
   while (s[cur].go[text[i]-OFFC] == -1)
    cur = s[cur].fail;
   cur = s[cur].go[text[i]-OFFC];
   if (s[cur].out.size())
    for (auto end : s[cur].out)
. toret.push back({i - end.first + 1,
    end.second});
  return toret:
struct defint { int i = -1; }:
vector<int> boyermoore(string txt, string pat)
 vector<int> toret: unordered map<char, defint>
 → badchar:
 int m = pat.size(), n = txt.size();
 for (int i = 0: i < m: i++) badchar[pat[i]].i string lcp(string* arr. int n. bool sorted =
 \rightarrow = i:
 while (s <= n - m) {
  int j = m - 1;
  while (i \ge 0 \&\& pat[i] == txt[s + i]) i--;
  if (j < 0) {
  .toret.push_back(s);
  s += (s + m < n) ? m - badchar[txt[s + m]]
 → m]].<mark>i</mark> : 1;
} else
  s += max(1, j - badchar[txt[s + j]].i);
return toret:
                                                         int m = a.length(), n = b.length();
                                                         int L[m+1][n+1];
English Conversion
```

```
|const string ones[] = {"", "one", "two",
    "three", "four", "five", "six", "seven",

    "eight", "nine"};
const string teens[] ={"ten", "eleven",
   "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen");
const string tens[] = {"twenty", "thirty",
    "forty", "fifty", "sixty", "seventy",
const string mags[] = {"thousand", "million",
     "billion", "trillion", "quadrillion",
    "quintillion", "sextillion",
    "septillion"};
string convert(int num, int carry) {
if (num < 0) return "negative " +
    convert(-num, 0):
    (num < 10) return ones[num];
(num < 20) return teens[num % 10]
    (num < 100) return tens[(num / 10) - 2] + (num / 10==0?"": " ) + ones[num / 10];
    (num < 1000) return ones[num / 100]
     (num/100==0?"":" ") + "hundred" +
     (num%100==0?"":" ") + convert(num % 100,
return convert(num / 1000, carry + 1) + " " +
    mags[carry] + " " + convert(num % 1000,

⇒ 0);

string convert(int num) {
return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
   vector<int> toret;
 int m = txt.length(), n = pat.length();
 int next[n + 1];
 for (int i = 0; i < n + 1; i++)
 next[i] = 0;
 for (int i = 1; i < n; i++) {
  int j = next[i + 1];
  while (j > 0 && pat[j] != pat[i])
  j = next[j];
 if (j > 0 || pat[j] == pat[i])
next[i + 1] = j + 1;
 for (int i = 0, j = 0; i < m; i++) {
 if (txt[i] == pat[j]) {
  if (++j == n)
   ..toret.push back(i - j + 1);
 } else if (j > 0) {
...j = next[j];
 return toret;
Longest Common Prefix (array)
// longest common prefix of strings in array

    false) {
    if (n == 0) return "";
}

if (!sorted) sort(arr, arr + n);
string r = ""; int v = 0;
 while (v < arr[0].length() && arr[0][v] ==
→ arr[n-1][v])
    r += arr[0][v++];
return r;
Longest Common Subsequence
string lcs(string a, string b) {
```

for (int i = 0; i <= m; i++) {

```
.for (int j = 0; j <= n; j++) {
..if (i == 0 || j == 0) L[i][j] = 0;
..else if (a[i-1] == b[j-1]) L[i][j] =</pre>
\hookrightarrow L[i-1][j-1]+1;
...else L[i][j] = max(L[i-1][j], L[i][j-1]);
.
// return L[m][n]; // length of lcs
 string out = "";
 int i = m - 1, j = n - 1;
 while (i >= 0 && j >= 0) {
 if (a[i] == b[j]) {
  .out = a[i--] + out;
  else if (L[i][j+1] > L[i+1][j]) i--;
  else j--;
return out;
Longest Common Substring
// l is array of palindrome length at that
int manacher(string s, int* 1) {
 int n = s.length() * 2;
 for (int i = 0, j = 0, k; i < n; i += k, j =
\rightarrow max(j-k, 0)) {
 while (i \ge j \&\& i + j + 1 < n \&\& s[(i-j)/2]
    == s[(i+j+1)/2]) j++;
 l[i] = j;
 for (k = 1; i >= k \&\& j >= k \&\& l[i-k] !=
    i-k; k++)
  1[i+k] = min(1[i-k], j-k);
return *max element(1, 1 + n):
Cyclic Rotation (Lyndon)
// simple strings = smaller than its nontrivial
   suffixes
// lyndon factorization = simple strings
   factorized
// "abaaba" -> "ab", "aab", "a"
vector<string> duval(string s) {
 int n = s.length();
vector<string> lyndon;
for (int i = 0; i < n;) {
   int j = i+1, k = i;</pre>
 for (; j < n && s[k] <= s[j]; j++)
if (s[k] < s[j]) k = i;
   else k++;
  for (; i \le k; i += j - k)
   lyndon.push_back(s.substr(i,j-k));
 return lyndon;
}
// lexicographically smallest rotation
int minRotation(string s) {
int n = s.length(); s += s;
 auto d = duval(s): int i = 0, a = 0:
 while (a + d[i].length() < n) a +=

    d[i++].length();

while (i \&\& d[i] == d[i-1]) a ==

    d[i--].length();

return a;
Hashing
#define HASHER 27
ull basicHash(string s) {
 ull v = 0;
for (auto c : s) v = (c - 'a' + 1) + v *

→ HASHER;

return v;
const int MAXN = 1000001;
ull base[MAXN] = {1};
void genBase(int n) {
```

```
for (int i = 1; i \le n; i++)

base[i] = base[i-1] * HASHER;
struct advHash {
 ull v, l; vector<ull> wip;
 advHash(string& s): v(0) {
 wip = vector<ull>(s.length()+1);\
  wip[0] = 0;
  for (int i = 0; i < s.length(); i++)
   wip[i+1] = (s[i] - 'a' + 1) + wip[i] *
   HASHER;
 1 = s.length(); v = wip[1];
 ull del(int pos, int len) {
  return v - wip[pos+len]*base[l-pos-len] +
    wip[pos]*base[l-pos-len];
 ull substr(int pos, int len) {
  return del(pos+len, (1-pos-len)) -
    wip[pos]*base[len]:
Subsequence Count
 // "banana", "ban" >> 3 (ban, ba..n, b..an)
ull subsequences(string body, string subs) {
 int m = subs.length(), n = body.length();
if (m > n) return 0;
ull** arr = new ull*[m+1];
 for (int i = 0; i \le m; i++) arr[i] = new
for (int i = 1; i <= m; i++) arr[i][0] = 0;
for (int i = 0; i <= n; i++) arr[0][i] = 1;
 for (int i = 1; i <= m; i++)
 for (int j = 1; j <= n; j++)
...arr[i][j] = arr[i][j-1] + ((body[j-1] ==
   subs[i-1])? arr[i-1][j-1] : 0);
 return arr[m][n];
Suffix Array + LCP
struct SuffixArray {
 vector<int> sa, lcp;
 SuffixArray(string& s, int lim=256) {
  int n = s.length() + 1, k = 0, a, b;
  vector<int> x(begin(s), end(s)+1), y(n),
   ws(max(n, lim)), rank(n);
  iota(begin(sa), end(sa), 0);
  for (int j = 0, p = 0; p < n; j = max(1, j *
 \rightarrow 2), \lim_{n \to \infty} p
   p = j; iota(begin(y), end(y), n - j);
   for (int i = 0; i < (n); i++)
if (sa[i] >= j)
     y[p++] = sa[i] - j;
   fill(begin(ws), end(ws), 0);
   for (int i = 0; i < (n); i++) ws[x[i]]++;
for (int i = 1; i < (lim); i++) ws[i] +=
    ws[i - 1];
   for (int i' = n; i--;) sa[--ws[x[v[i]]]] =
    y[i];
   swap(x, y); p = 1; x[sa[0]] = 0;
   for (int i = 1; i < (n); i++) {
   a = sa[i - 1]; b = sa[i];
   x[b] = (y[a] == y[b] && y[a + j] == y[b +
   j]) ? p - 1 : p++;
  for (int i = 1; i < (n); i++) rank[sa[i]] =
  for (int i = 0, j; i < n - 1; lcp[rank[i++]]
  for (k \&\& k--, j = sa[rank[i] - 1];
     s[i + k] == s[j + k]; k++);
Suffix Tree (Ukkonen's)
```

```
struct SuffixTree {
 // n = 2*len+10 or so
enum { N = 50010, ALPHA = 26 };
int toi(char c) { return c - 'a'; }
 void ukkadd(int i, int c) { suff:
  if (r[v]<=q) {
   if (t[v][c]==-1) { t[v][c]=m; l[m]=i;
     p[m++]=v; v=s[v]; q=r[v]; goto suff; }</pre>
   v=t[v][c]; q=1[v];
  if (q==-1 || c==toi(a[q])) q++; else {
   l[m+1]=i; p[m+1]=m; l[m]=l[v]; r[m]=q
    p[m]=p[v]; t[m][c]=m+1; t[m][toi(a[q])]=v;
    l[v]=q; p[v]=m; t[p[m]][toi(a[l[m]])]=m;
   v=s[p[m]]; q=l[m];
   while (q < r[m]) \{ v = t[v][toi(a[q])];
    q+=r[v]-l[v]; }
   if (q==r[m]) s[m]=v; else s[m]=m+2;
q=r[v]-(q-r[m]); m+=2; goto suff;
 SuffixTree(string a) : a(a) {
  fill(r,r+N,(int)(a).size());
 memset(s, 0, sizeof s);
memset(t, -1, sizeof t);
fill(t[1],t[1]+ALPHA,0);
s[0]=1;1[0]=1[1]=-1;r[0]=r[1]=p[0]=p[1]=0;
  for(int i=0;i<a.size();i++)
     ukkadd(i,toi(a[i]));
 .}
.// Longest Common Substring between 2 strings
 // returns {length, offset from first string}
 pair<int, int> best;
 int lcs(int node, int i1, int i2, int olen) {
  if (l[node] <= i1 && i1 < r[node]) return 1;
  if (l[node] <= i2 && i2 < r[node]) return 2;</pre>
  int mask=0,
 → len=node?olen+(r[node]-l[node]):0;
  for(int c=0; c<ALPHA; c++) if
   (t[node][c]!=-1)
mask |= lcs(t[node][c], i1, i2, len);
  if (mask==3)
 → best=max(best,{len,r[node]-len});
  return mask;
 static pair<int, int> LCS(string s, string t)
  SuffixTree
 st(s+(char)('z'+1)+t+(char)('z'+2));
 st.lcs(0, s.size(), s.size()+t.size()+1, 0);
return st.best;
String Utilities
void lowercase(string& s) {
 transform(s.begin(), s.end(), s.begin(),
void uppercase(string& s) {
 transform(s.begin(), s.end(), s.begin(),
 void trim(string &s) {
 s.erase(s.begin(),find_if_not(s.begin(),s
     .end(),[](int c){return
    isspace(c);}));
 s.erase(find_if_not(s.rbegin(),s.rend(),[](int

    c){return isspace(c);}).base(),s.end());

vector<string> split(string& s, char token) {
     vector<string> v; stringstream ss(s);
     for (string e;getline(ss,e,token);)
         v.push_back(e);
     return v:
```

```
5 Greedy
                                                       Interval Cover
                                                        //L,R = interval [L,R], in = \{\{l,r\}, index\}
\rightarrow t[N] [ALPHA],1[N],r[N],p[N],s[N],v=0,q=0,m=2] // does not handle case where L==R string a; vector<int> intervalCover(double L, double R,
                                                          vector<pair<double,double>,int>> in) {
                                                           int i = 0; pair < double, int > pos = {L,-1};
                                                           vector<int> a:
                                                           sort(begin(in), end(in));
                                                           while (pos.first < R) {
                                                                double cur = pos.first;
while (i < (int)in.size() &&
                                                           in[i].first.first <= cur)</pre>
                                                           max(pos,{in[i].first.second,in[i].second}),
                                                                if (pos.first == cur) return {}:
                                                                a.push_back(pos.second);
                                                           return a;
                                                       6 Math
                                                       Catalan Numbers
                                                       ull* catalan = new ull[1000000];
                                                       void genCatalan(int n, int mod) {
                                                       catalan[0] = catalan[1] = 1;

for (int i = 2; i <= n; i++) {

   catalan[i] = 0;

   for (int j = i - 1; j >= 0; j--) {
                                                          catalan[i] += (catalan[j] * catalan[i-j-1])
                                                          % mod;
                                                          if (catalan[i] >= mod)
                                                           catalan[i] -= mod:
                                                       // TODO: consider binomial coefficient method
                                                       Combinatorics (nCr, nPr)
                                                       // can optimize by precomputing factorials, and
                                                          fact[n]/fact[n-r]
                                                       ull nPr(ull n, ull r) {
                                                       ull v =
                                                        for (ull i = n-r+1; i <= n; i++)
                                                       . v *= i;
return v:
                                                      ull nPr(ull n, ull r, ull m) {
                                                       ull v = 1;
for (ull i = n-r+1; i <= n; i++)
                                                        v = (v * i) \% m;
                                                       return v:
                                                       ull nCr(ull n, ull r) {
                                                        long double \dot{v} = 1;
                                                        for (ull i = 1; i <= r; i++)
                                                        v = v * (n-r+i) /i;
                                                       return (ull)(v + 0.001);
                                                        // requires modulo math
                                                       // can optimize by precomputing mfac and

→ minv-mfac

                                                       ull nCr(ull n. ull r. ull m) {
                                                       return mfac(n, m) * minv(mfac(k, m), m) % m *
                                                       \rightarrow minv(mfac(n-k, m), m) % m:
                                                      Multinomials
                                                       ll multinomial(vector<int>& v) {
                                                       11 c = 1, m = v.empty() ? 1 : v[0];
                                                       for(int i = 1; i < v.size(); i++)

for (int j = 0; j < v[i]; j++)

...c = c * ++m / (j+1);
```

return c;

```
Chinese Remainder Theorem
bool ecrt(ll* r, ll* m, int n, ll& re, ll& mo)
11 x, y, d; mo = m[0]; re = r[0];
for (int i = 1; i < n; i++) {
 d = \gcd(mo, m[i], x, y);
 if ((r[i] - re) % d != 0) return false;
x = (r[i] - re) / d * x % (m[i] / d);
 re += x * mo;
.mo = mo / d * m[i];
 re %= mo:
 re = (re + mo) \% mo:
 return true;
Count Digit Occurences
/*count(n,d) counts the number of occurrences of
   a digit d in the range [0,n]*/
11 digit_count(ll n, ll d) {
ll result = 0;
 while (n != 0) {
  result += ((n%10) == d ? 1 : 0);
 n /= 10;
 return result;
il count(ll n, ll d) {
if (n < 10) return (d > 0 \&\& n >= d);
 if ((n % 10) != 9) return digit_count(n, d) +
   count(n-1, d);
return 10*count(n/10, d) + (n/10) + (d > 0);
Discrete Logarithm
unordered_map<int, int> dlogc;
int discretelog(int a, int b, int m) {
.dlogc.clear();
11 n = sqrt(m)+1, an = 1;
for (int i = 0; i < n; i++)
 an = (an * a)'\% m;
 11 c = an;
 for (int i = 1; i <= n; i++) {
    if (!dlogc.count(c)) dlogc[c] = i;
                                                      .do {
 c = (c * an) \% m;
 c = b;
 for (int i = 0; i <= n; i++) {
 if (dlogc.count(c)) return (dlogc[c] * n - i
\rightarrow + m - 1) % (m-1);
 c = (c * a) \% m;
return -1:
Euler Phi / Totient
int phi(int n) {
 int r = n;
for (int i = 2; i * i <= n; i++) {
   if (n % i == 0) r -= r / i;
   while (n % i == 0) n /= i;
if (n > 1) r -= r / n;
return r;
#define n 100000
ll phi[n+1]:
void computeTotient() {
 for (int i=1; i<=n; i++) phi[i] = i;
for (int_p=2; p<=n; p++) {
 if (phi[p] == p) {
  ..phi[p] = p-1;
 ..for (int i = 2*p; i<=n; i += p) phi[i] =
   (phi[i]/p) * (p-1);
Factorials
// digits in factorial
#define kamenetsky(n) (floor((n * log10(n /
\rightarrow ME)) + (log10(2 * MPI * n) / 2.0)) + 1)
```

```
// approximation of factorial
#define stirling(n) ((n == 1) ? 1 : sqrt(2 *
    M PI * n) * pow(n / M E, n))
// natural log of factorial
#define lfactorial(n) (lgamma(n+1))
Prime Factorization
 // do not call directly
11 pollard_rho(ll n, ll s) {
 11 x, y;
 x = y = rand() \% (n - 1) + 1;
 int head = 1, tail = 2;
while (true) {
 white (black, x, n);
x = mult(x, x, n);
x = (x + s) % n;
if (x == y) return n;
  11 d = _{gcd(max(x - y, y - x), n);}
  if (1 < \overline{d} \&\& d < n) return d;
if (++\text{head} == \text{tail}) y = x, tail <<= 1;
 // call for prime factors
void factorize(ll n, vector<ll> &divisor) {
  if (n == 1) return;
 if (isPrime(n)) divisor.push_back(n);
  while (d' \ge n) d = pollard_rho(n, rand() % (n)
 \rightarrow -1) +1);
factorize(n / d, divisor);
  factorize(d, divisor);
 Farey Fractions
    generate 0 \le a/b \le 1 ordered, b \le n
    farey(4) = 0/1 \frac{1}{4} \frac{1}{3} \frac{1}{2} \frac{2}{3} \frac{3}{4} \frac{1}{1}
 // length is sum of phi(i) for i = 1 to n
vector<pair<int, int>> farey(int n) {
 int h = 0, k = 1, x = 1, y = 0, r;
 vector<pair<int, int>> v;
  v.push back({h, k});
  r = (n-y)/k;
  y += r*k; x += r*h;
 swap(x,h); swap(y,k);
x = -x; y = -y;
while (k > 1);
 v.push_back({1, 1});
Fast Fourier Transform
#define cd complex<doubl
const double PI = acos(-1):
void fft(vector<cd>& a, bool invert) {
 int n = a.size();
 for (int i = 1, j = 0; i < n; i++) {
  int bit = n >> 1:
  for (; j & bit; bit >>= 1) j ^= bit;
  i ^= bit:
  if (i < j) swap(a[i], a[j]);
 for (int len = 2; len <= n; len <<= 1) {
  double ang = 2 * PI / len * (invert ? -1 :
   cd wlen(cos(ang), sin(ang));
  for (int i = 0; i < n; i += len) {
    cd w(1);
    for (int j = 0; j < len / 2; j++) {
     cd u = a[i+j], v = a[i+j+len/2] * w;
    a[i+j] = u + v;
a[i+j+len/2] = u - v;
w *= wlen;
 if (invert)
  for (auto& x : a)
   x /= n:
```

```
|vector<int> fftmult(vector<int> const& a,

  vector<int> const& b) {
  vector<cd> fa(a.begin(), a.end()),
 \rightarrow fb(b.begin(), b.end());
int n = 1 << (32 - _builtin_clz(a.size() +

    b.size() - 1));
fa.resize(n); fb.resize(n);
 fft(fa, false); fft(fb, false)
for (int i = 0; i < n; i++) fa[i] *= fb[i];
fft(fa, true);</pre>
 vector<int> toret(n);
 for (int i = 0; i < n; i++) toret[i] =
 → round(fa[i].real());
 return toret;
Greatest Common Denominator
11 egcd(11 a, 11 b, 11% x, 11% y) {
  if (b == 0) { x = 1; y = 0; return a; }
 11 gcd = egcd(b, a \% b, x, y);
 x = a / b * y;
 swap(x, y);
return gcd;
Josephus Problem
// 0-indexed, arbitrary k
int josephus(int n, int k) {
if (n == 1) return 0;
if (k == 1) return n-1;
 if (k > n) return (joséphus(n-1,k)+k)%n;
 int res = josephus(n-n/k,k)-n\%k;
 return res + ((res<0)?n:res/(k-1));
// fast case if k=2, traditional josephus
int josephus(int n) {
return 2*(n-(1<<(32-_builtin_clz(n)-1)));
Least Common Multiple
#define lcm(a,b) ((a*b)/qcd(a,b))
Modulo Operations
#define MOD 1000000007
#define madd(a,b,m) (a+b-((a+b-m>=0)?m:0))
#define mult(a,b,m) ((ull)a*b%m)
#define msub(a,b,m) (a-b+((a<b)?m:0))
ll mpow(ll b, ll e, ll m) {
11 x = 1:
 while (e'> 0) {
  if (e % 2) x = (x * b) % m;
 b = (b * b) \% m:
  e /= 2;
 return x % m:
ull mfac(ull n. ull m) {
 ull f = 1
 for (int i = n; i > 1; i--)
 f = (f * i) \% m;
 return f:
// if m is not quaranteed to be prime
ll minv(ll b. ll m) {
11 x = 0, y' = 0;
 if (egcd(b, m, x, y) != 1) return -1;
 return (x \% m + m) \% m;
il mdiv_compmod(int a, int b, int m) {
 if (__gcd(b, m) != 1) return -1;
 return mult(a, minv(b, m), m);
// if m is prime (like 10^9+7)
11 mdiv_primemod (int a, int b, int m) {
return mult(a, mpow(b, m-2, m), m);
// tonelli shanks = sqrt(n) % m, m is prime
ll legendre(ll a. ll m){
 if (a % m==0) return 0;
 if (m == 2) return 1;
```

```
return mpow(a, (m-1)/2, m);
il msqrt(ll n, ll m) {
ll s = builtin ctzll(m-1), q = (m-111)>>s,
  z = rand()\%(m-1)+1;
if (m == 2) return 1;
if (s == 1) return mpow(n,(m+1)/411,m);
while (legendre(z,m)!=m-1) z = rand()\%(m-1)+1;
11 c = mpow(z,q,m), r = mpow(n,(q+1)/2,m), t
= mpow(n,q,m), M = s;
while (t != 1){
 while (t; -1) (t * t) % m;
while (ts != 1) i++, ts = (ts * ts) % m;
 for (int j = 0; j < M-i-1; j++) b = (b * b) %
r = r * b \% m; c = b * b \% m; t = t * c \% m;
\rightarrow M = i;
return r;
Modulo Tetration
11 tetraloop(ll a, ll b, ll m) {
if(b == 0 | a == 1) return 1;
11 w = tetraloop(a,b-1,phi(m)), r = 1;
for (;w;w/=2) {
 if (w&1) {
  r *= a: if (r >= m) r -= (r/m-1)*m:
 a *= a: if (a >= m) a -= (a/m-1)*m:
return r:
int tetration(int a, int b, int m) {
if (a == 0 || m == 1) return ((b+1)&1)%m;
return tetraloop(a,b,m) % m;
Matrix
template<typename T>
struct Mat : public Vec<2, T> {
int w, h;
Mat(int x, int y) : Vec<2, T>(x, y), w(x),
\rightarrow h(v) \{\}
static Mat<T> identity(int n) { Mat<T> m(n,n);
   for (int i=0:i < n:i++) m[i][i] = 1: return

    m; }

.Mat<T>& operator+=(const Mat<T>& m) {
 for (int i = 0; i < w; i++)
 for (int j = 0; j < h; j++)
(*this)[i][j] += m[i][j];
 return *this;
Mat<T>& operator-=(const Mat<T>& m) {
 for (int i = 0; i < w; i++)
  for (int j = 0; j < h; j++)

(*this)[i][j] -= m[i][j];
 return *this;
Mat<T> operator*(const Mat<T>& m) {
 Mat<T>z(w,m.h);
 for (int i = 0; i < w; i++)
  for (int j = 0; j < h; j++)
  for (int k = 0; k < m.h; k++)
z[i][k] += (*this)[i][j] * m[j][k];
   return z:
Mat<T> operator+(const Mat<T>& m) { Mat<T>
→ a=*this: return a+=m: }
Mat<T> operator-(const Mat<T>& m) { Mat<T>

    a=*this; return a-=m; }

Mat<T>& operator*=(const Mat<T>& m) { return
 \rightarrow *this = (*this)*m; }
Mat<T> power(int n) {
 Mat<T> a = Mat<T>::identity(w),m=*this;
 for (;n;n/=2,m*=m) if (n\&1) a *= m;
 return a;
```

```
Matrix Exponentiation
                                                                    Permutation (string/multiset)
// F(n) = c[0]*F(n-1) + c[1]*F(n-2) + ...
                                                                    string freq2str(vector<int>& v) {
// b is the base cases of same length c
ll matrix_exponentiation(ll n, vector<ll> c,
                                                                     string s;
                                                                     for (int i = 0; i < v.size(); i++)
for (int j = 0; j < v[i]; j++)

s += (char)(i + 'A');
return s;
→ vector<11> b) {
if (nth < b.size()) return b[nth-1];</pre>
 Mat<11> a(c.size(), c.size()); 11 s = 0;
for (int i = 0; i < c.size(); i++) a[i][0] =
                                                                     // nth perm of multiset, n is 0-indexed
 \hookrightarrow c[i];
for (int i = 0; i < c.size() - 1; i++)
                                                                    string gen_permutation(string s, ll n) {
                                                                     vector<int> freq(26, 0);
\rightarrow a[i][i+1] = 1;
a = a.power(nth - c.size());
                                                                     for (auto e : s) freq[e - 'A']++;
 for (int i = 0; i < c.size(); i++)
    s += a[i][0] * b[i];
return s;
                                                                     for (int i = 0; i < 26; i++) if (freq[i] > 0)
                                                                       freq[i]--; ll v = multinomial(freq);
                                                                       if (n < v) return (char)(i+'A') +
                                                                         gen_permutation(freq2str(freq), n);
Matrix Subarray Sums
                                                                      freq[\overline{i}]++; n -= v;
template<class T> struct MatrixSum {
 Vec<2, T> p;
                                                                     return "";
 MatrixSum(Vec<2, T>& v) {
  p = Vec<2,T>(v.size()+1, v[0].size()+1);
                                                                    Miller-Rabin Primality Test
  for (int i = 0; i < v.size(); i++)
 ...for (int j = 0; j < v[0].size(); j++)
....p[i+1][j+1] = v[i][j] + p[i][j+1] +
                                                                     // Miller-Rabin primality test - O(10 log^3 n)
                                                                    htter-nath promatity test
bool isPrime(ull n) {
  if (n < 2) return false;
  if (n == 2) return true;
  if (n % 2 == 0) return false;
  ull s = n - 1;
  while (s % 2 == 0) s /= 2;
  factorial false;
  if (n % 2 == 0) s /= 2;</pre>

    p[i+1][j] - p[i][j];

  \stackrel{\mathsf{f}}{\mathsf{I}} \mathtt{sum}( \underset{\mathsf{int}}{\mathsf{int}} \ \mathsf{u}, \ \underset{\mathsf{int}}{\mathsf{int}} \ \mathsf{1}, \ \underset{\mathsf{int}}{\mathsf{int}} \ \mathsf{d}, \ \underset{\mathsf{int}}{\mathsf{int}} \ \mathsf{r}) \ \big\{ \\ \ \mathsf{return} \ \mathsf{p}[\mathsf{d}][\mathsf{r}] \ - \ \mathsf{p}[\mathsf{d}][\mathsf{1}] \ - \ \mathsf{p}[\mathsf{u}][\mathsf{r}] \ + \ \mathsf{p}[\mathsf{u}][\mathsf{1}]; 
                                                                      for (int i = 0; i < 10; i++) {
                                                                       ull temp = s;
                                                                       ull a = rand() % (n - 1) + 1;
Mobius Function
                                                                       ull mod = mpow(a, temp, n);
const int MAXN = 10000000;
                                                                       while (temp!=n-1\&\&mod!=1\&\&mod!=n-1) {
// mu[n] = 0 iff n has no square factors
                                                                        mod = mult(mod, mod, n);
// 1 = even number prime factors, -1 = odd
short mu[MAXN] = {0,1};
                                                                        temp *= 2:
void mobius(){
  for (int i = 1; i < MAXN; i++)</pre>
                                                                      if (mod!=n-1&&temp%2==0) return false;
  if (mu[i])
                                                                     return true;
 ...for (int j = i + i; j < MAXN; j += i)
     mu[j] -= mu[i];
                                                                    Sieve of Eratosthenes
                                                                    bitset<100000001> sieve;
Nimber Arithmetic
                                                                     // generate sieve - O(n log n
#define nimAdd(a,b) ((a)^(b))
                                                                     void genSieve(int n) {
                                                                     sieve[0] = sieve[1] = 1;
for (ull i = 3; i * i < n; i += 2)
    if (!sieve[i])</pre>
ull nimMul(ull a, ull b, int i=6) {
   static const ull M[]={INT_MIN>>32,
     M[0]^(M[0] << 16), M[1]^(M[1] << 8),
                                                                        for (ull j = i * 3; j <= n; j += i * 2)
    M[2]^(M[2] << 4), M[3]^(M[3] << 2),
\stackrel{\rightarrow}{\rightarrow} M[2] (M[2] \times 1), 
\stackrel{\rightarrow}{\rightarrow} M[4] (M[4] \times 1), 
                                                                         .sieve[j] = 1;
   if (i--)=0) return a&b;
                                                                     // query sieve after it's generated - O(1)
  ull s=nimMúl(a,b,i), m=M[5-i],
                                                                    bool querySieve(int n) {
                                                                     return n == 2 || (n % 2 != 0 && !sieve[n]);
     t=nimMul(((a^(a>>k))&m)|(s\&~m).
     ((b^(b>>k))&m)|(m&(~m>>1))<< k, i);
  return ((s^t)\&m)<\langle k|((s^(t)>k))\&m);
                                                                    Compile-time Prime Sieve
                                                                    const int MAXN = 100000;
template<int N>
Permutation
                                                                    struct Sieve {
  bool sieve[N];
// c = array size, n = nth perm, return index
vector<int> gen_permutation(int c, int n) {
                                                                      constexpr Sieve() : sieve() {
 vector<int> idx(c), per(c), fac(c); int i;
                                                                       sieve[0] = sieve[1] = 1;
 for (i = 0; i < c; i++) idx[i] = i;
for (i = 1; i <= c; i++) fac[i-1] = n%i, n/=i;
for (i = c - 1; i >= 0; i--)
...per[c-i-1] = idx[fac[i]],
                                                                      for (int i = 2; i * i < N; i++)
  if (!sieve[i])</pre>
                                                                         for (int j = i * 2; j < N; j += i)
...sieve[j] = 1;
   idx.erase(idx.begin() + fac[i]);
                                                                    bool isPrime(int n) {
   static constexpr Sieve<MAXN> s;
// get what nth permutation of vector
int get_permutation(vector<int>& v) {
                                                                     return !s.sieve[n];
 int use = 0, i = 1, r = 0;
 for (int e: v) {
   r = r * i++ + __builtin_popcount(use &
                                                                    Simpson's / Approximate Integrals
 \rightarrow -(1<<e));
                                                                     // integrate f from a to b, k iterations
  use |= 1 << e;
                                                                     // error <= (b-a)/18.0 * M * ((b-a)/2k)^{2}
                                                                     // where M = max(abs(f^{(x)})) for x in [a,b]
 return r;
                                                                     // "f" is a function "double func(double x)"
```

```
double Simpsons (double a, double b, int k,
 \rightarrow double (*f)(double)) {
double dx = (b-a)/(2.0*k), t = 0;
 for (int i = 0; i < k; i++)
t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
(*f)(a+(2*i+1)*dx);
return (t + (*f)(b)) * (b-a) / 6.0 / k;
Common Equations Solvers
// ax^2 + bx + c = 0, find x
vector < double > solveEq (double a, double b,
  double z = b * b - 4 * a * c;
 if (z == 0)
  r.push_back(-b/(2*a));
 else if (z > 0) {
    r.push_back((sqrt(z)-b)/(2*a));
  r.push_back((sqrt(z)+b)/(2*a));
 return r:
\frac{1}{2} / ax^3 + bx^2 + cx + d = 0, find x
vector<double> solveEq(double a, double b,
 → double c, double d) {
vector < double > res;
long double a1 = b/a, a2 = c/a, a3 = d/a;
 long double q = (a1*a1 - 3*a2)/9.0, sq =
 \rightarrow -2*sqrt(q);
 long double r = (2*a1*a1*a1 - 9*a1*a2 +
 \rightarrow 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
 if (z <= 0) {
  theta = acos(r/sqrt(q*q*q));
res.push_back(sq*cos(theta/3.0) - a1/3.0);
  res.push back(sq*cos((theta+2.0*PI)/3.0) -
  - a1/3.0);
.res.push_back(sq*cos((theta+4.0*PI)/3.0) -
\rightarrow a1/3.0);
 else {
  res.push_back(pow(sqrt(z)+fabs(r), 1/3.0));
  res[0] = (res[0] + q / res[0]) *
     ((r<0)?1:-1) - a1 / 3.0:
 return res:
\frac{1}{1} linear diophantine equation ax + by = c,
\hookrightarrow find x and y
// infinite solutions of form x+k*b/g, y-k*a/g bool solveEq(11 a, 11 b, 11 c, 11 &x, 11 &y, 11
 g = egcd(abs(a), abs(b), x, y);
if (c % g) return false;
 x *= c / g * ((a < 0) ? -1 : 1);
 v *= c / g * ((b < 0) ? -1 : 1);
 return true:
// m = # equations. n = # variables. a[m][n+1]
\Rightarrow = coefficient matrix
// a[i][0]x + a[i][1]y + ... + a[i][n]z =
    a[i][n+1]
    find a solution of some kind to linear
\hookrightarrow equation
const double eps = 1e-7:
bool zero(double a) { return (a < eps) && (a >
vector double > solveEq(double **a, int m, int
if (j != cur) swap(a[j], a[cur]);
for (int sat = 0; sat < m; sat++) {
   if (sat == cur) continue;</pre>
```

double num = a[sat][i] / a[cur][i];

for (int sot = 0; sot <= n; sot++

```
_a[sat][sot] -= a[cur][sot] * num;
    .}
.cur++:
    break:
for (int j = cur; j < m; j++)
   if (!zero(a[j][n])) return vector<double>();
vector<double> ans(n,0);
for (int i = 0, sat = 0; i < n; i++)
if (sat < m && !zero(a[sat][i]))
ans[i] = a[sat][n] / a[sat++][i];
return ans;
// solve A[n][n] * x[n] = b[n] linear equation
// rank < n is multiple solutions, -1 is no

    ⇒ solutions
    // `alls` is whether to find all solutions, or
    ...

\hookrightarrow any
const double eps = 1e-12;
int solveEq(Vec<2, double>& A, Vec<1, double>&
\rightarrow b, Vec<1, double>& x, bool alls=false) {
int n = A.size(), m = x.size(), rank = 0, br,
→ bc;
vector<int> col(m); iota(begin(col), end(col),
for(int i = 0; i < n; i++) {
   double v, bv = 0;
   for(int r = i; r < n; r++)</pre>
  for(int c = i; c < n; c++)
if ((v = fabs(A[r][c])) > bv)
br = r, bc = c, bv = v;
  if (bv <= eps) {
  for(int j = i; j < n; j++)
if (fabs(b[j]) > eps)
     return -1:
   break:
  swap(A[i], A[br]);
 swap(b[i], b[br]);
swap(col[i], col[bc]);
  for(int j = 0; j < n; j++)
swap(A[j][i], A[j][bc]);
  bv = 1.0 / A[i][i];
  for(int j = (alls)?0:i+1; j < n; j++) {
   .if (j != i) {
    double fac = A[j][i] * bv;
    .b[j] -= fac * b[i];
    for(int k = i+1; k < m; k++)
     A[i][k] -= fac*A[i][k];
  rank++;
 if (alls) for (int i = 0; i < m; i++) x[i] =
if (alls)
  for (int j = rank; isGood && j < m; j++)
  if (fabs(A[i][j]) > eps)
  isGood = false;
b[i] /= A[i][i];
  if (isGood) x[col[i]] = b[i];
if (!alls)
  for(int j = 0; j < i; j++)
b[j] -= A[j][i] * b[i];
return rank;
Graycode Conversions
ull graycode2ull(ull n) {
ull i = 0;
for (; n; n = n >> 1) i ^= n;
ull ull2graycode(ull n) {
return n ^ (n >> 1);
```

Unix/Epoch Time // O-indexed month/time. 1-indexed day // minimum 1970, 0, 1, 0, 0, 0 ull toEpoch(int year, int month, int day, int → hour, int minute, int second) { struct tm t; time_t epoch; t.tm_year = year - 1900; t.tm_mon = month; t.tm_mday = day; t.tm_hour = hour; t.tm_min = minute; t.tm_sec = second; t.tm_isdst = 0; // 1 = daylights savings epoch = mktime(&t): return (ull)epoch; vector<int> toDate(ull epoch) { time_t e=epoch; struct tm t=*localtime(&e); return {t.tm_year+1900,t.tm_mon,t.tm_mday,t_ .tm hour,t.tm min,t.tm sec}; int getWeekday(ull epoch) { time_t e=epoch; struct tm t=*localtime(&e); return t.tm wday; // 0-6, 0 = sunday int getDayofYear(ull epoch) { time_t e=epoch; struct tm t=*localtime(&e); return t.tm_yday; // 0-365 const int months[] = \rightarrow {31,28,31,30,31,30,31,30,31,30,31}; bool validDate(int year, int month, int day) { bool leap = !(year%(year%25?4:16)); if (month >= 12) return false; return day <= months[month] + (leap &&</pre> month == 1);

Theorems and Formulae

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set of size n. !0 = 1, !1 = 0, !n = (n+1)(!(n-1))1)+!(n-2)), ! $n = n! \sum_{i=0}^{n} \frac{(-1)^i}{i!}$, ! $n = \left[\frac{n!}{e}\right]$

In a partially ordered set, a chain is a subset of elements that are all comparable to eachother. An antichain is a subset where no two are comparable.

Dilworth's theorem states the size of a maximal antichain equals the size of a minimal chain cover of a partially ordered set S. The width of S is the maximum size of an antichain in S, which is equal to the minimum number of chains needed to cover S, or the minimum number of chains such that all elements are in at least one chain.

Rosser's Theorem states the nth prime Floyd Warshall number is greater than n * ln(n) for n > 1.

Lagrange's Four Square Theorem states $\rightarrow \&\& m[k][j] = \inf \{$ every natural number is the sum of the squares of four non-negative integers. This is a spe- $\begin{array}{c} \text{cial case of the } \mathbf{Fermat} \ \mathbf{Polygonal} \ \mathbf{Number} \\ \end{array} \\ \begin{array}{c} \mathbf{FOR(k,n)} \ \text{if} \\ \mathbf{fig. (m[k][k]} \\ \end{array} \\ \begin{array}{c} \mathbf{fig. (m[k][k])} \\ \end{array} \\ \begin{array}{c} \mathbf{FOR(i,n)} \\ \end{array} \\ \\ \begin{array}{c} \mathbf{FOR(i,n)} \\ \end{array} \\ \begin{array}{c} \mathbf{$ **Theorem** where every positive integer is a sum of at most n s-gonal numbers. The nth

```
s-gonal number P(s,n) = (s-2)\frac{n(n-1)}{2} + n | Minimum Spanning Tree
   Graphs
struct edge {
 int u,v,w;
edge (int u,int v,int w) : u(u),v(v),w(w) {}
 edge (): u(0), v(0), w(0) {}
bool operator < (const edge &e1, const edge
\rightarrow &e2) { return e1.w < e2.w; }
bool operator > (const edge &e1, const edge
Eulerian Path
#define edge_list vector<edge>
#define_adj_sets vector<set<int>>
struct EulerPathGraph {
 adj_sets graph; // actually indexes incident
 → edaes
 edge_list edges; int n; vector<int> indeg;
 EulerPathGraph(int n): n(n) {
  indeg = *(new vector<int>(n,0));
  graph = *(new adj_sets(n, set<int>()));
 void add edge(int u, int v) {
  graph[u].insert(edges.size());
  indeg[v]++;
  edges.push back(edge(u.v.0)):
 bool eulerian_path(vector<int> &circuit) {
  if(edges.size()==0) return false;
  stack<int> st;
int a[] = {-1, -1};
for(int v=0; v<n; v++) {
   if(indeg[v]!=graph[v].size()) {
        bool b = indeg[v] > graph[v].size();
}
    if (abs(((int)indeg[v])-((int)graph[v]
     .size())) > 1) return
    false;
if (a[b] != -1) return false;
    a[b] = v;
  int s = (a[0]!=-1 && a[1]!=-1 ? a[0] :
    (a[0]==-1 && a[1]==-1 ? edges[0].u : -1));
  if(s==-1) return false;
  while(!st.empty() || !graph[s].empty()) {
   if (graph[s].empty()) {
     circuit.push back(s); s = st.top();
    st.pop(); }
    int w = edges[*graph[s].begin()].v;
    graph[s].erase(graph[s].begin());
    st.push(s); s = w;
 circuit.push_back(s);
  return circuit.size()-1==edges.size();
```

```
number is greater than n * ln(n) for n > 1.

Nicomachi's Theorem states 1^3 + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + ... + 2^3 + 
Nicomachi's Theorem states 1 + 2 + \dots + n^3 = (1 + 2 + \dots + n)^2 and is equivalent to \begin{cases} n + 1 \\ n - 1 \end{cases} and is equivalent to \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} and \begin{cases} n + 1 \\ n - 1 \end{cases} an
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                      m[i][j] = min(m[i][j], newDist);
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                   if (m[i][k] != inf && m[k][j] != inf)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                               \rightarrow m[i][i] = -inf:
```

```
points.push_back(a.br);
                                                          points.push_back({real(a.br),
 // returns vector of edges in the mst
                                                          imag(a.tl)}):
 / graph[i] = vector of edges incident to
→ vertex i
// places total weight of the mst in &total
// if returned vector has size != n-1, there is
                                                     struct polygon {
                                                       vector<point> points;
| → no MST
| vector<edge> mst(vector<vector<edge>> graph.
                                                       polygon(vector<point> points) :

→ 11 &total) {

                                                       → points(points) {}
 total = 0;
priority_queue<edge, vector<edge>,
                                                      polygon(triangle a) {
                                                       points.push_back(a.a); points.push_back(a.b);
 → greater<edge>> pq;
                                                         points.push back(a.c);
 vector<edge> MST;
 bitset<20001> marked; // change size as needed
                                                      polygon(rectangle a) {
 marked[0] = 1;
                                                       points.push_back(a.tl);
 for (edge ep : graph[0]) pq.push(ep); while(MST.size()!=graph.size()-1 &&
                                                          points.push back({real(a.tl).
                                                          imag(a.br)});
 → pq.size()!=0) {
                                                       points.push back(a.br):
 ledge e = pq.top(); pq.pop();
int u = e.u, v = e.v, w = e.w;
if(marked[u] && marked[v]) continue;
                                                          points.push_back({real(a.br),
                                                          imag(a.tl)}):
  else if(marked[u]) swap(u, v);
                                                      polygon(convex_polygon a) {
  for(edge ep : graph[u]) pq.push(ep);
  marked[u] = 1
                                                       for (point v : a.points)
  MST.push_back(e):
                                                         points.push back(v):
  total += e.w:
 return MST:
                                                      // triangle methods
                                                      double area heron(double a, double b, double
                                                      Union Find
int uf find(subset* s, int i) {
  if (s[i].p != i) s[i].p = uf_find(s, s[i].p);
                                                      if (a < c) swap(a, c);
                                                      if (b < c) swap(b, c);
 return s[i].p;
                                                       if (a > b + c) return -1;
                                                      return sqrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)
void uf_union(subset* s, int x, int y) {
 int xp = uf_find(s, x), yp = uf_find(s, y);
if (s[xp].rank > s[yp].rank) s[yp].p = xp;
                                                      // segment methods
 else if (s[xp].rank < s[yp].rank) s[xp].p =
                                                     double lengthsq(segment a) { return
\hookrightarrow yp;
else { s[yp].p = xp; s[xp].rank++; }
                                                          sq(real(a.a) - real(a.b)) + sq(imag(a.a) -
                                                         imag(a.b)); }
                                                      double length(segment a) { return
2D Grid Shortcut

    sqrt(lengthsq(a)); }

#define inbound(x,n) (0<=x\mathcal{E}\mathcal{E}x<n)
                                                      // circle methods
                                                     double circumference(circle a) { return 2 * a.r
#define fordir(x,y,n,m) for(auto[dx,dy]:dir)if

→ * M PI: }

 \rightarrow (inbound(x+dx,n)&\iffsi inbound(y+dy,m))
                                                      double area(circle a) { return sq(a.r) * M_PI;
const pair<int,int> dir[] =
                                                     | → }
|// rectangle methods
\rightarrow {{1,0},{0,1},{-1,0},{0,-1}};
                                                      double width(rectangle a) { return
                                                     → abs(real(a.br) - real(a.tl)); }
double height(rectangle a) { return
   2D Geometry
#define point complex<double>
#define EPS 0.0000001

    abs(imag(a.br) - real(a.tl)); }

#define sq(a) ((a)*(a))
#define cb(a) ((a)*(a)*(a))
                                                     double diagonal (rectangle a) { return
                                                     → sqrt(sq(width(a)) + sq(height(a))); }
double area(rectangle a) { return width(a) *
double dot(point a, point b) { return

    real(conj(a)*b); }

                                                      → height(a): }
double cross(point a, point b) { return
                                                     double perimeter(rectangle a) { return 2 *

    imag(conj(a)*b); }

                                                      struct line { point a, b; };
                                                      // check if `a` fit's inside `b
                                                      // swap equalities to exclude tight fits
struct circle { point c; double r; };
struct segment { point a, b; };
struct triangle { point a, b, c; };
                                                      bool doesFitInside(rectangle a, rectangle b) {
                                                      int x = width(a), w = width(b), y = height(a),
struct rectangle { point tl, br; };
                                                      \rightarrow h = height(b);
struct convex_polygon {
                                                      if (x > y) swap(x, y);
if (w > h) swap(w, h);
 vector<point> points:
 convex polygon(vector<point> points) :
                                                       if (w < x) return false:
    points(points) {}
                                                      if (y <= h) return true;
 convex_polygon(triangle a) {
                                                      double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
  points.push_back(a.a); points.push_back(a.b); return sq(a) <= sq(b) + sq(c);
 → points.push_back(a.c);
                                                      // polygon methods
                                                      // negative area = CCW, positive = CW
 convex_polygon(rectangle a) {
                                                      double area(polygon a) {
  points.push_back(a.tl);
                                                       double area = 0.0; int n = a.points.size();
    points.push_back({real(a.tl),
    imag(a.br)});
```

```
for (int i = 0, j = 1; i < n; i++, j = (j +
    area += (real(a.points[j]-a.points[i]))*
return area / 2.0:
// get both unsigned area and centroid
pair<double, point> area_centroid(polygon a) {
 int n = a.points.size();
 double area = 0;
 point c(0, 0);
 for (int i = n - 1, j = 0; j < n; i = j++) {
 .double v = cross(a.points[i], a.points[i]) /
  c += (a.points[i] + a.points[j]) * (v / 3);
 c /= area;
return {area. c}:
Intersection
// -1 coincide, 0 parallel, 1 intersection
int intersection(line a, line b, point& p) {
if (abs(cross(a.b - a.a, b.b - b.a)) > EPS) {
   p = cross(b.a - a.a, b.b - a.b) / cross(a.b)
\rightarrow - a.a, b.b - b.a) * (b - a) + a;
 return 1:
 if (abs(cross(a.b - a.a. a.b - b.a)) > EPS)

    return 0;

return -1;
}
// area of intersection
double intersection(circle a, circle b) {
double d = abs(a.c - b.c):
 if (d \le b.r - a.r) return area(a);
if (d <= a.r - b.r) return area(b);
if (d >= a.r + b.r) return 0;
 double alpha = acos((sq(a.r)' + sq(d) -
\rightarrow sq(b.r)) / (2 * a.r * d));
double beta = acos((sq(b.r) + sq(d) - sq(a.r))
\rightarrow / (2 * b.r * d));
return sq(a.r) * (alpha - 0.5 * sin(2 *
    alpha)) + sq(b.r) * (beta - 0.5 * sin(2 *
    beta));
}
// -1 outside, 0 inside, 1 tangent, 2
int intersection(circle a, circle b,
→ vector<point>& inter) {
 double d2 = norm(b.c - a.c), rS = a.r + b.r,
\hookrightarrow rD = a.r - b.r;
 if (d2 > sq(rS)) return -1;
 if (d2 < sq(rD)) return 0;
 double ca = 0.5 * (1 + rS * rD / d2);
point z = point(ca, sqrt(sq(a.r) / d2 -
\rightarrow sq(ca)));
 inter.push back(a.c + (b.c - a.c) * z);
 if (abs(imag(z)) > EPS) inter.push_back(a.c +
\rightarrow (b.c - a.c) * conj(z));
return inter.size();
// points of intersection
vector<point> intersection(line a, circle c) {
vector<point> inter;
c.c -= a.a;
a.b -= a.a;
 point m = a.b * real(c.c / a.b);
 double d2 = norm(m - c.c);
 if (d2 > sq(c.r)) return 0;
 double 1 = \operatorname{sqrt}((\operatorname{sq}(c.r) - d2) / \operatorname{norm}(a.b));
 inter.push_back(a.a + m + 1 * a.b);
 if (abs(1) > EPS) inter.push_back(a.a + m - 1
\rightarrow * a.b):
return inter;
// area of intersection
double intersection(rectangle a, rectangle b) { int snoob(int a) {
```

```
double x1 = max(real(a.tl), real(b.tl)), y1 = | int b = a & -a, c = a + b;
                                                       return c | ((a^ c) >> 2) / b:
 → max(imag(a.tl), imag(b.tl));
 double x2 = min(real(a.br), real(b.br)), y2 =
                                                       // example usage
                                                      // example usage
int main() {
   char l1[] = {'1', '2', '3', '4', '
   char l2[] = {'a', 'b', 'c', 'd'};
   int d1 = 5, d2 = 4;
   // prints 12345abcd, 1234a5bcd, ...
 → min(imag(a.br), imag(b.br));
 return (x2 <= x1 | | y2 <= y1) ? 0 :
   (x2-x1)*(y2-y1);
                                                        int min = (1 < < d1) - 1, max = min << d2;
Convex Hull
                                                       for (int i = min; i <= max; i = snoob(i)) {
   int p1 = 0, p2 = 0, v = i;
   while (p1 < d1 || p2 < d2) {
      cout << ((v & 1) ? 11[p1++] : 12[p2++]);
bool cmp(point a, point b) {
  if (abs(real(a) - real(b)) > EPS) return
    real(a) < real(b);
 if (abs(imag(a) - imag(b)) > EPS) return
                                                         v /= 2;
   imag(a) < imag(b);</pre>
 return false;
                                                        cout << '\n':
convex polygon convexhull(polygon a) {
 sort(a.points.begin(), a.points.end(), cmp);
 vector<point> lower, upper;
                                                       Powers
 for (int i = 0; i < a.points.size(); i++) {
...while (lower.size() >= 2 &&
                                                      bool isPowerOf2(11 a) {
                                                       return a > 0 && ! (a & a-1):
    cross(lower.back() - lower[lower.size() -
                                                      bool isPowerOf3(11 a) {
  return a>0&&!(12157665459056928801u11%a);
    2], a.points[i] - lower.back()) < EPS)
   lower.pop_back();
  while (upper.size() >= 2 &&
                                                      bool isPower(ll a, ll b) {
  double x = log(a) / log(b);
    cross(upper.back() - upper[upper.size()
    2], a.points[i] - upper.back()) > -EPS)
                                                       return abs(x-round(x)) < 0.00000000001;
   upper.pop_back();
  lower.push_back(a.points[i]);
                                                      11 Additional
  upper.push_back(a.points[i]);
                                                      Judge Speed
 lower.insert(lower.end(), upper.rbegin() + 1,
                                                       // kattis: 0.50s
// codeforces: 0.421s
    upper.rend());
 return convex polygon(lower);
                                                      // atcoder: 0.455s
                                                      #include <bits/stdc++.h>
                                                      using namespace std;
    3D Geometry
                                                      \frac{1}{1} int v = 1e9/2, p = 1;
                                                      int main() {
struct point3d {
                                                       for (int i = 1; i <= v; i++) p *= i;
 double x, y, z;
 point3d operator+(point3d a) const { return
 \rightarrow {x+a.x, y+a.y, z+a.z}; }
 point3d operator*(double a) const { return
                                                      Judge Pre-Contest Checks
    \{x*a, y*a, z*a\}; \}
                                                           int128 and float128 support?
 point3d operator-() const { return {-x, -y,
                                                       -does extra or missing whitespace cause WA?
   -z}; }
 point3d operator-(point3d a) const { return
                                                       -documentation up to date?
    *this + -a: }
                                                       -printer usage available and functional?
 point3d operator/(double a) const { return
    *this * (1/a); }
                                                       // each case tests a different fail condition
 double norm() { return x*x + y*y + z*z; }
                                                      // try them before contests to see error codes
 double abs() { return sqrt(norm()); }
                                                      struct g { int arr[1000000]; g(){}};
 point3d normalize() { return *this /
    this->abs(); }
                                                      // O=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE
                                                           6=SIGSEGV 7=recursive MLE judge(int n) {
double dot(point3d a, point3d b) { return
 \rightarrow a.x*b.x + a.y*b.y + a.z*b.z; }
                                                           (n == 0) exit(0)
point3d cross(point3d a, point3d b) { return
                                                       if (n == 1) while(1);
if (n == 2) while(1);
a.push_back(g());
    {a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,}
                                                       if (n == 3) while(1) putchar_unlocked('a');
   a.x*b.y - a.y*b.x; }
struct line3d { point3d a, b; };
                                                        if (n == 4) assert(0);
                                                       if (n == 5) 0 / 0;
if (n == 6) *(int*)(0) = 0;
struct plane { double a, b, c, d; } // a*x +
\rightarrow b*y + c*z + d = 0
                                                       return n + judge(n + 1);
struct sphere { point3d c; double r; };
#define sq(a) ((a)*(a))
#define c\bar{b}(a) ((a)*(a)*(a))
                                                      GCC Builtin Docs
double surface(circle a) { return 4 * sq(a.r)
                                                       // 128-bit integer
   M PI; }
                                                        _int128 a;
                                                      unsigned __int128 b;
double volume(circle a) { return 4.0/3.0 *
 \rightarrow cb(a.r) * M_PI; }
                                                      // 128-bit float
                                                       // minor improvements over long double
10 Optimization
                                                        float128 c:
                                                      // log2 floor
                                                      __lg(n);
 // SameNumberOfOneBits, next permutation
                                                      // number of 1 bits
```

// can add ll like popcountll for long longs

```
__builtin_popcount(n);
// number of trailing zeroes
__builtin_ctz(n);
// number of leading zeroes
__builtin_clz(n);
// 1-indemed least significant 1 bit
  builtin ffs(n):
// parity of number
__builtin_parity(n);
Limits
                       \pm 2147483647 \mid \pm 2^{31} - 1 \mid 10^9
int
                                              \frac{1}{2}<sup>32</sup> -\frac{1}{1}<sup>10</sup><sup>9</sup>
uint
                          4294967295
        \pm 9223372036854775807 | \pm 2^{63} - 1 | 10^{18}
11
                                            \frac{1}{2}64 - \frac{1}{1}10<sup>19</sup>
ull
         18446744073709551615
|i128| \pm 170141183460469231... | \pm 2^{\tilde{1}27} - 1 | 10^{38}
|u128| 340282366920938463... | 2^{128} - 1 | 10^{38}
Complexity classes input size (per second):
O(n^n) or O(n!)
                                                       n \leq 10
O(2^n)
                                                       n < 30
O(n^3)
                                                   n < 1000
O(n^2)
                                                  n < 30000
O(n\sqrt{n})
                                                     n < 10^6
O(n \log n)
                                                     n < 10^7
O(n)
                                                     n < 10^9
```