```
else n = n * 10 + c - '0':
    General
                             7 Graphs
    Algorithms
                              8 2D Geometry
                                                            n = s * (n + m * o):
    Structures
                              9 3D Geometry
                                                           void read(double& n) {
    Strings
                                                            ld m; read(m); n = m;
                              10 Optimization
    Greedy
                              11 Additional
                                                           void read(float& n) {
 ld m: read(m): n = m:
    Math
     General
                                                            void read(string& s) {
                                                            char c; s = "
g++ -g -02 -std=gnu++17 -static prog.cpp
./a.exe
run.sh
                                                            while((c=getchar unlocked())!=' '&&c!='\n')
                                                            bool readline(string& s) {
                                                            char c; s = "";
while(c=getchar unlocked()) {
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp
for i in *.in; do
                                                             if (c == '\n') return true;
if (c == EOF) return false;
s += c;
 f=${i%.in}
 ./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
                                                            return false;
done
                                                            void print(unsigned int n) {
Header
                                                            if (n / 10) print(n / 10);
// use better compiler options
                                                            putchar_unlocked(n % 10 + '0');
#pragma GCC optimize("Ofast","unroll-loops")
#pragma GCC target("avx2,fma")
                                                            void print(int n) {
// include everything
                                                            if (n < 0) { putchar_unlocked('-'); n*=-1; }
 #include <bits/stdc++.h>
                                                            print((unsigned int)n);
#include <bits/extc++.h>
#include <sys/resource.h>
// namespaces
                                                           Common Structs
using namespace std;
                                                               n-dimension vectors
using namespace __gnu_cxx; // rope
                                                               Vec<2, int>v(n, m) = arr[n][m]
using namespace __gnu_pbds; // tree/trie
                                                            // Vec<2, int> v(n, m, -1) default init -1
                                                            template<int D, typename T>
// common defines
#define fastio
                                                            struct Vec : public vector < Vec < D-1, T >> {
                                                              template<typename... Args>

→ ios base::sync with stdio(0);cin.tie(0);
                                                              Vec(int n=0, Args... args) : vector<Vec<D-1,
#define nostacklim rlimit RZ; getrlimit(3,&RZ
                                                            \rightarrow T>>(n, Vec<D-1, T>(args...)) {}
    ):RZ.rlim cur=-1:setrlimit(3.&RZ):
#define DEBUG(v) cerr<< LINE <<": "<<#v<<" =
                                                           template<typename T>
\Rightarrow "<<v<<'\n'; #define TIMER
                                                           struct Vec<1, T> : public vector<T> {
                                                              Vec(int n=0, T val=T()) : vector<T>(n, val)

→ cerr<<1.0*clock()/CLOCKS_PER_SEC<<"s\n";
#define ll long long
#define ull unsigned ll
#define i128 __int128
#define u128 unsigned i128
                                                               {}
                                                                Algorithms
#define ld long double
                                                           Min/Max Subarray
// global variables
                                                              max - compare = a < b, reset = a < 0
mt19937 rng((uint32_t)chrono::steady
                                                            \frac{1}{min} - compare = a > b, reset = a > 0

    clock::now().time since epoch().count());

                                                           // returns {sum, {start, end}}
pair<int, pair<int, int>>
Fast IO
                                                                ContiguousSubarray(int* a, int size,
#ifdef _WIN32
                                                                bool(*compare)(int, int),
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
                                                            bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
                                                            0, s = 0;
for (int i = 0; i < size; i++) {
  cur += a[i];</pre>
void read(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                              if ((*compare)(best, cur)) { best = cur;
  n = n * 10 + c - 0';
                                                            \rightarrow start = s; end = i; }
void read(int& n) {
  char c; n = 0; int s = 1
                                                             if ((*reset)(cur)) { cur = 0; s = i + 1; }
 if ((c=getchar_unlocked())=='-') s = -1;
                                                            return {best, {start, end}}:
 else n = c - '0';
while ((c=getchar_unlocked())!=' '&&c!='\n')
                                                            Quickselect
 n = n * 10 + c - 0';
                                                           #define OSNE -999999
                                                           int partition(int arr[], int 1, int r)
void read(ld& n) {
 char c; n = 0;
ld m = 0, o = 1; bool d = false; int s = 1;
if ((c=getchar_unlocked())=='-') s = -1;
                                                            int x = arr[r], i = 1;
for (int j = 1; j <= r - 1; j++)
...if (arr[j] <= x)
...swap(arr[i++], arr[j]);</pre>
 else if (c == .'.') d = true;
else n = c - '0';
 while ((c=getchar_unlocked())!=' '&&c!='\n') {
                                                            swap(arr[i], arr[r]);
 if (c == '.') d = true;
else if (d) { m=m*10+c-'0'; o*=0.1; }
                                                            return i:
```

```
// find k'th smallest element in unsorted array, void update(int i, int val) {
→ only if all distinct
int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - l + 1)) return QSNE;
swap(arr[1 + rng() % (r-l+1)], arr[r]);
 int pos = partition(arr, 1, r);
if (pos-l==k-1) return arr[pos];
 if (pos-1>k-1) return qselect(arr,1,pos-1,k);
 return qselect(arr, pos+1, r, k-pos+1-1);
// TODO: compare against std::nth_element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted

    on both axis
pair<int, int> saddleback_search(int** arr, int
 \rightarrow x, int y, int v) {
 int i = x-1, j = 0;
 while (i >= 0 && j < y) {
  if (arr[i][j] == v) return {i, j};
  (arr[i][j] > v)? i--: j++;
 return {-1, -1};
Ternary Search
 // < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a, int b, int (*f)(int)) {</pre>
 while (b-a > 4) {
    int m = (a+b)/2;
    if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1:
 for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
   a = i;
 return a;
#define TERNPREC 0.000001
double ternsearch(double a. double b. double
 \leftrightarrow (*f)(double)) {
while (b-a > TERNPREC * 4) {
  double m = (a+b)/2;
  if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
  else b = m + TERNPREC;
 for (double i = a + TERNPREC: i <= b: i +=
     TERNPREC)
      if (TERNCOMP((*f)(a), (*f)(i)))
 return a;
Golden Section Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
double goldsection(double a, double b, double
 while (b-a > eps)
  while (b-a > eps)

if (TERNCOMP(f2,f1)) {

. b = x2; x2 = x1; f2 = f1;

. x1 = b - r*(b-a); f1 = f(x1);
  } else {
   a = x1; x1 = x2; f1 = f2;

x2 = a + r*(b-a); f2 = f(x2);
 return a:
3 Structures
Fenwick Tree
// Fenwick tree, array of cumulative sums -
```

 \hookrightarrow O(log n) updates, O(log n) gets

struct Fenwick { int n; ll* tree;

```
while (i <= n) {
   tree[i] += val;
   i += i & (-i);
 Fenwick(int size) {
  n = size;
  tree = new ll[n+1];
for (int i = 1; i <= n; i++)
   .tree[i] = 0;
 Fenwick(int* arr, int size) : Fenwick(size) {
  for (int i = 0; i < n; i++)
...update(i, arr[i]);
 ~Fenwick() { delete[] tree; }
 ll operator[](int i) {
  if (i < 0 || i > n) return 0;
  \overline{11} \ \overline{sum} = 0;
  while (i>0)
   sum += tree[i];
   i -= i & (-i):
  return sum:
 ll getRange(int a, int b) { return

    operator[](b) - operator[](a-1); }

Hashtable
// similar to unordered map, but faster
struct chash {
    const uint64 t C = (11)(2e18 * M PI) + 71;
 ll operator()(ll x) const { return
    builtin bswap64(x*C); }
int main() {
  gp_hash_table<11,int,chash>
 \rightarrow hashtable({},{},{},{},{1<<16});
 for (int i = 0; i < 100; i++)
hashtable[i] = 200+i;
 if (hashtable.find(10) != hashtable.end())
   cout << hashtable[10];</pre>
Ordered Set
template <typename T>
using oset = tree<T,null_type,less<T>,rb_tree
    _tag,tree_order_statistics_node_update>;
template <typename T, typename D> using omap = tree<T,D,less<T>,rb_tree
    _tag,tree_order_statistics_node_update>;
int main()
 oset<int> o_set;
o_set.insert(5); o_set.insert(1);
 → o_set.insert(3);
// get second smallest element
 cout << *(o set.find by order(1));</pre>
 // number of elements less than k=4
cout << ' ' << o_set.order_of_key(4) << '\n';</pre>
 // equivalent with ordered map
 omap<int,int> o_map;
o_map[5]=1;o_map[1]=2;o_map[3]=3;
 cout << (*(o_map.find_by_order(1))).first;</pre>
 cout << ' ' << o map.order of key(4) << '\n';
Rope
// O(log n) insert, delete, concatenate
int main() {
 // generate rove
 rope<int> v;
 for (int i = 0: i < 100: i++)
  v.push_back(i);
 // move range to front
 rope<int> copy = v.substr(10, 10);
 v.erase(10, 10);
```

```
v.insert(copy.mutable_begin(), copy);
 // print elements of rope
for (auto it : v)
cout << it << "":
Segment Tree
//max(a,b), min(a,b), a+b, a*b, qcd(a,b), a*b
struct SegmentTree {
 typedef int T;
 static constexpr T UNIT = INT MIN:
 T f(T a, T b) {
 if (a == UNIT) return b;
if (b == UNIT) return a;
 return max(a,b);
 Int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
\rightarrow n(n) {}
 SegmentTree(vector<T> arr)

→ SegmentTree(arr.size()) {
 for (int i=0:i<arr.size():i++)

→ update(i.arr[i]):

 void update(int pos, T val) {
  for (s[pos += n] = val; pos /= 2;)
   s[pos] = f(s[pos * 2], s[pos*2+1]);
 T query(int b, int e) { // query [b, e)
 Tra = UNIT, rb = UNIT;

for (b+=n, e+=n; b<e; b/=2, e/=2) {

    if (b % 2) ra = f(ra, s[b++]);

    if (e % 2) rb = f(s[--e], rb);
  return f(ra. rb):
 T get(int p) { return query(p, p+1); }
Sparse Table
template<class T> struct SparseTable {
 vector<vector<T>> m;
SparseTable(vector<T> arr) {
  m.push_back(arr);
  for (int k = 1; (1<<(k)) <= size(arr); k++)
   m.push back(vector<T>(size(arr)-(1<<k)+1)):
   for (int i = 0; i < size(arr)-(1 << k)+1; i
    m[k][i] = min(m[k-1][i],
   m[k-1][i+(1<<(k-1))]:
 }
// min of range [l,r]
T query(int 1, int r) {
  int k = __lg(r-l+1);
  return \min(m[k][1], m[k][r-(1 << k)+1]):
typedef trie<string, null_type,

→ trie_string_access_traits<>,

 pat_trie_tag, trie_prefix_search_node_update>
int main() {
 // generate trie
 trie_type trie;
for (int i = 0; i < 20; i++)
...trie.insert(to_string(i)); // true if new,
\hookrightarrow false if old
 // print things with prefix "1"
 auto range = trie.prefix_range("1");
 for (auto it = range.first; it !=

    range.second; it++)

  cout << *it <<
Wavelet Tree
using iter = vector<int>::iterator;
struct WaveletTree {
```

```
Vec<2, int> C: int s:
 // sigma = highest value + 1
 WaveletTree(vector<int>& a, int sigma) :
    s(sigma), C(sigma*2, 0) {
  build(a.begin(), a.end(), 0, s-1, 1);
 void build(iter b, iter e, int L, int U, int
  if (L == U) return;
  int M = (L+U)/2:
  C[u].reserve(e-b+1); C[u].push_back(0);
  for (auto it = b; it != e; ++it)
    C[u].push_back(C[u].back() + (*it<=M));
  auto p = stable_partition(b, e, [=](int
    i) {return i <= M: }):
  build(b, p, L, M, u*2);
  build(p, e, M+1, U, u*2+1);
 // number of occurences of x in [0,i)
if (x <= M) i = r, U = M;
else i -= r, L = M+1, ++u;
  return i:
 // number of occurrences of x in [l,r)
int count(int x, int l, int r) {
  return rank(x, r) - rank(x, 1);
 // kth smallest in [l, r)
int kth(int k, int l, int r) const {
int L = 0, U = s-1, u = 1, M, ri, rj;
  while (L != U) {
   M = (L+U)/2;
   ri = C[u][1]; rj = C[u][r]; u*=2;

if (k <= rj-ri) 1 = ri, r = rj, U = M;
   else k -= ri-ri. l -= ri. r -= ri.
   L = M+1. ++u:
  return U;
  // # elements between [x,y] in [l, r)
 mutable int L. U:
 int range(int x, int y, int 1, int r) const {
  if (y < x \text{ or } r \le 1) return 0;
  L = x; U = y;
  return range(1, r, 0, s-1, 1);
 int range(int 1, int r, int x, int y, int u)
    const {
  if (y < L or U < x) return 0;
  if (L \le x \text{ and } y \le U) \text{ return } r-1;
  int M = (x+y)/2, ri = C[u][1], rj = C[u][r];
  return range(ri, rj, x, M, u*2) + range(1-ri, Boyer Moore
    r-rj, M+1, y, u*2+1);
 // # elements <= x in [l, r)
int lte(int x, int l, int r) {
  return range(INT_MIN, x, l, r);</pre>
     Strings
Aho Corasick
// range of alphabet for automata to consider
// MAXC = 26, OFFC = 'a' if only lowercase
const int MAXC = 256;
const int OFFC = 0;
struct aho_corasick {
 struct state
  set<pair<int, int>> out;
  int fail; vector<int> go;
  state() : fail(-1), go(MAXC, -1) {}
 vector<state> s;
```

```
int id = 0:
 aho corasick(string* arr, int size) : s(1) {
 for (int i = 0; i < size; i++) {
   int cur = 0;
   for (int c : arr[i]) {
   if (s[cur].go[c-OFFC] == -1) {
     s[cur].go[c-OFFC] = s.size();
     s.push_back(state());
    cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++}):
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
    s[0].go[c] = 0;
  queue<int> sq;
  for (int c = 0; c < MAXC; c++) {
   if (s[0].go[c] != 0) {
    s[s[0].go[c]].fail = 0;
    sq.push(s[0].go[c]);
  while (sq.size()) {
   int e = sq.front(); sq.pop();
  for (int c = 0; c < MAXC; c++) {
   if (s[e].go[c] != -1) {
     int failure = s[e].fail;
while (s[failure].go[c] == -1)
     failure = s[failure].fail;
failure = s[failure].go[c];
     s[s[e].go[c]].fail = failure;
     for (auto length : s[failure].out)
  s[s[e].go[c]].out.insert(length);
      sq.push(\bar{s}[e].go[c]);
 // list of {start pos, pattern id}
 vector<pair<int, int>> search(string text)
 vector<pair<int, int>> toret;
  int cur = 0;
  for (int i = 0; i < text.size(); i++) {</pre>
   while (s[cur].go[text[i]-OFFC] == -1)
    cur = s[cur].fail;
   cur = s[cur].go[text[i]-OFFC];
   if (s[cur].out.size())
    for (auto end : s[cur].out)
. toret.push back({i - end.first + 1,
    end.second});
  return toret:
struct defint { int i = -1: }:
vector<int> boyermoore(string txt, string pat)
 vector<int> toret: unordered map<char, defint>
 → badchar:
 int m = pat.size(), n = txt.size();
 for (int i = 0: i < m: i++) badchar[pat[i]].i string lcp(string* arr. int n. bool sorted =
 \rightarrow = i:
 while (s <= n - m) {
  int j = m - 1;
  while (i \ge 0 \&\& pat[i] == txt[s + i]) i--;
  if (j < 0) {
  .toret.push_back(s);
  s += (s + m < n) ? m - badchar[txt[s +
 → m]].<mark>i</mark> : 1;
} else
  s += max(1, j - badchar[txt[s + j]].i);
return toret:
                                                         int m = a.length(), n = b.length();
                                                         int L[m+1][n+1];
English Conversion
```

```
|const string ones[] = {"", "one", "two",
    "three", "four", "five", "six", "seven",

    "eight", "nine"};
const string teens[] ={"ten", "eleven",
   "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen");
const string tens[] = {"twenty", "thirty",
    "forty", "fifty", "sixty", "seventy",
const string mags[] = {"thousand", "million",
     "billion", "trillion", "quadrillion",
    "quintillion", "sextillion",
    "septillion"};
string convert(int num, int carry) {
if (num < 0) return "negative " +
    convert(-num, 0):
    (num < 10) return ones[num];
(num < 20) return teens[num % 10]
    (num < 100) return tens[(num / 10) - 2] + (num / 10==0?"": " ) + ones[num / 10];
    (num < 1000) return ones[num / 100]
     (num/100==0?"":" ") + "hundred" +
     (num%100==0?"":" ") + convert(num % 100,
return convert(num / 1000, carry + 1) + " " +
    mags[carry] + " " + convert(num % 1000,

⇒ 0);

string convert(int num) {
return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
   vector<int> toret;
 int m = txt.length(), n = pat.length();
 int next[n + 1];
 for (int i = 0; i < n + 1; i++)
 next[i] = 0;
 for (int i = 1; i < n; i++) {
  int j = next[i + 1];
  while (j > 0 && pat[j] != pat[i])
  j = next[j];
 if (j > 0 || pat[j] == pat[i])
next[i + 1] = j + 1;
 for (int i = 0, j = 0; i < m; i++) {
 if (txt[i] == pat[j]) {
  if (++j == n)
   ..toret.push back(i - j + 1);
 } else if (j > 0) {
...j = next[j];
 return toret;
Longest Common Prefix (array)
// longest common prefix of strings in array

    false) {
    if (n == 0) return "";
}

if (!sorted) sort(arr, arr + n);
string r = ""; int v = 0;
 while (v < arr[0].length() && arr[0][v] ==
→ arr[n-1][v])
    r += arr[0][v++];
return r;
Longest Common Subsequence
string lcs(string a, string b) {
```

for (int i = 0; i <= m; i++) {

```
. for (int j = 0; j <= n; j++) {
...if (i == 0 || j == 0) L[i][j] = 0;
...else if (a[i-1] == b[j-1]) L[i][j] =
                                                        for (int i = 1; i <= n; i++)
base[i] = base[i-1] * HASHER;
                                                       struct advHash {
\hookrightarrow L[i-1][j-1]+1;
...else L[i][j] = max(L[i-1][j], L[i][j-1]);
                                                        ull v, l; vector<ull> wip;
                                                        advHash(string& s): v(0) {
                                                         wip = vector <ull>(s.length()+1);\
.
// return L[m][n]; // length of lcs
                                                          wip[0] = 0;
                                                         for (int i = 0; i < s.length(); i++)
 string out = "";
                                                          wip[i+1] = (s[i] - 'a' + 1) + wip[i] *
 int i = m - 1, j = n - 1;
                                                           HASHER;
 while (i >= 0 && j >= 0) {
                                                         1 = s.length(); v = wip[1];
 if (a[i] == b[j]) {
  .out = a[i--] + out;
                                                        ull del(int pos, int len) {
                                                         return v - wip[pos+len]*base[l-pos-len] +
                                                           wip[pos]*base[1-pos-len]:
  else if (L[i][j+1] > L[i+1][j]) i--;
  else j--;
                                                        ull substr(int pos, int len) {
                                                         return del(pos+len, (1-pos-len)) -
return out;
                                                           wip[pos]*base[len]:
Longest Common Substring
                                                        ull replace(int pos, char c) {
// l is array of palindrome length at that
                                                         return v - wip[pos+1]*base[l-pos-1] + ((c -
                                                            'a' + 1) + wip[pos] *
int manacher(string s, int* 1) {
                                                           HASHER) *base[1-pos-1];
 int n = s.length() * 2;
 for (int i = 0, j = 0, k; i < n; i += k, j =
                                                        ull replace(int pos, string s) {
\rightarrow max(j-k, 0)) {
                                                         // can't increase total string size
 while (i >= j \&\& i + j + 1 < n \&\& s[(i-j)/2]
    == s[(i+j+1)/2]) j++;
                                                           wip[pos+s.size()]*base[l-pos-s.size()], c
 l[i] = j;
                                                           wip[pos];
 for (k = 1; i >= k && j >= k && l[i-k] !=
                                                         for (int i = 0; i < s.size(); i++)
.c = (s[i]-'a'+1) + c * HASHER:
   j-k; k++)
                                                         return r + c * base[1-pos-s.size()];
  1[i+k] = min(1[i-k], j-k);
return *max element(1, 1 + n):
                                                       Subsequence Count
                                                        // "banana", "ban" >> 3 (ban, ba..n, b..an)
Cyclic Rotation (Lyndon)
                                                       ull subsequences(string body, string subs) {
// simple strings = smaller than its nontrivial
                                                        int m = subs.length(), n = body.length();
   suffixes
                                                        if (m > n) return 0;
ull** arr = new ull*[m+1];
// lyndon factorization = simple strings
   factorized
                                                        for (int i = 0; i <= m; i++) arr[i] = new
// "abaaba" -> "ab", "aab", "a"
                                                          ull[n+1];
vector<string> duval(string s) {
                                                        for (int i = 1; i <= m; i++) arr[i][0] = 0;
for (int i = 0; i <= n; i++) arr[0][i] = 1;
 int n = s.length();
vector<string> lyndon;
for (int i = 0; i < n;) {
   int j = i+1, k = i;</pre>
                                                        for (int i = 1; i <= m; i++)
                                                         for (int j = 1; j <= n; j++)
arr[i][j] = arr[i][j-1] + ((body[j-1] ==
 for (; j < n && s[k] <= s[j]; j++)
if (s[k] < s[j]) k = i;
                                                        \rightarrow subs[i-1])? arr[i-1][j-1] : 0);
                                                        return arr[m][n];
   else k++;
  for (; i \leq k; i += j - k)
   lyndon.push_back(s.substr(i,j-k));
                                                       Suffix Array + LCP
                                                       struct SuffixArray {
 return lyndon;
                                                        vector<int> sa, lcp;
}
// lexicographically smallest rotation
                                                        SuffixArray(string&s, int lim=256) {
                                                         int n = s.length() + 1, k = 0, a, b;
.vector<int> x(begin(s), end(s)+1), y(n),
int minRotation(string s) {
int n = s.length(); s += s;
                                                           ws(max(n, lim)), rank(n);
 auto d = duval(s): int i = 0, a = 0:
                                                         sa = lcp = y;
iota(begin(sa), end(sa), 0);
 while (a + d[i].length() < n) a +=

    d[i++].length();

                                                          for (int j = 0, p = 0; p < n; j = max(1, j *
while (i && d[i] == d[i-1]) a -=
                                                           2), lim = p) {

    d[i--].length();

                                                          p = j; iota(begin(y), end(y), n - j);
for (int i = 0; i < (n); i++)
    if [(sa[i] >= j)
return a;
Hashing
                                                             y[p++] = sa[i] - j;
#define HASHER 27
                                                          fill(begin(ws), end(ws), 0);
                                                          for (int i = 0; i < (n); i++) ws[x[i]]++;
for (int i = 1; i < (lim); i++) ws[i] +=
ull basicHash(string s) {
 ull v = 0;
for (auto c : s) v = (c - 'a' + 1) + v *
                                                           ws[i - 1]:

→ HASHER;

                                                          for (int i' = n; i--;) sa[--ws[x[y[i]]]] =
return v;
                                                           y[i];
                                                          swap(x, y); p = 1; x[sa[0]] = 0;
const int MAXN = 1000001;
                                                          for (int i = 1; i < (n); i++) {
   a = sa[i - 1]; b = sa[i];
ull base[MAXN] = {1};
void genBase(int n) {
```

```
x[b] = (y[a] == y[b] && y[a + j] == y[b + |void trim(string &s) 
for (int i = 1; i < (n); i++) rank[sa[i]] =
 → i:
 for (int i = 0, j; i < n - 1; lcp[rank[i++]]
\Rightarrow = k)
...for (k && k--, j = sa[rank[i] - 1];
     s[i + k] == s[j + k]; k++);
Suffix Tree (Ukkonen's)
struct SuffixTree {
 // n = 2*len+10 or so
enum { N = 50010, ALPHA = 26 };
int toi(char c) { return c - 'a'; }
 t[N][ALPHA],1[N],r[N],p[N],s[N],v=0,q=0,m=2|yector<int> intervalCover(double L, double R
 string a;
 void ukkadd(int i, int c) { suff:
  if (r[v]<=q) {
  if (q=-1 || c==toi(a[q])) q++; else {
    l[m+1]=i; p[m+1]=m; l[m]=l[v]; r[m]=q;
    p[m]=p[v]; t[m][c]=m+1; t[m][toi(a[q])]=v;
   i[v]=q; p[v]=m; t[p[m]][toi(a[1[m]])]=m;
v=s[p[m]]; q=1[m];
   while (q < r[m]) \{ v=t[v][toi(a[q])];
   q+=r[v]-l[v]; }
if (q==r[m]) s[m]=v; else s[m]=m+2;
   q=r[v]-(q-r[m]); m+=2; goto suff;
 SuffixTree(string a) : a(a) {
 fill(r,r+N,(int)(a).size());
 memset(s, 0, sizeof s);

memset(t, -1, sizeof t);

fill(t[1],t[1]+ALPHA,0);

s[0]=1;1[0]=1[1]=-1;r[0]=r[1]=p[0]=p[1]=0;
  for(int i=0;i<a.size();i++)
    ukkadd(i.toi(a[i])):
 // Longest Common Substring between 2 strings
 // returns {length, offset from first string}
 pair<int, int> best;
 int lcs(int node, int i1, int i2, int olen) {
  if (l[node] <= i1 && i1 < r[node]) return 1;
  if (l[node] <= i2 && i2 < r[node]) return 2;</pre>
    len=node?olen+(r[node]-1[node]):0;
  for(int c=0; c<ALPHA; c++) if
    (t[node][c]!=-1)
   mask |= lcs(t[node][c], i1, i2, len);
  if (mask==3)
    best=max(best,{len,r[node]-len});
 return mask;
 static pair<int, int> LCS(string s, string t)
  SuffixTree
 \rightarrow st(s+(char)('z'+1)+t+(char)('z'+2));
 st.lcs(0, s.size(), s.size()+t.size()+1, 0); return st.best;
String Utilities
void lowercase(string& s) {
 transform(s.begin(), s.end(), s.begin(),
void uppercase(string& s) {
transform(s.begin(), s.end(), s.begin(),
   ::toupper);
```

```
s.erase(s.begin(),find_if_not(s.begin(),s
               .end(),[](int c){return
           isspace(c);}));
    s.erase(find_if_not(s.rbegin(),s.rend(),[](int
          c){return isspace(c);}).base(),s.end());
 vector<string> split(string& s, char token) {
             vector<string> v; stringstream ss(s);
              for (string e;getline(ss,e,token);)
                          v.push_back(e);
             return v;
         Greedy
  Interval Cover
    // L,R = interval [L,R], in = {{l,r}, index}

// does not handle case where L == R
            vector<pair<double,double>,int>> in) {
             int i = 0; pair<double,int> pos = {L,-1};
             vector<int> a:
               sort(begin(in), end(in));
              while (pos.first < R) {
                           double cur = pos.first;
while (i < (int)in.size() &&</pre>
            in[i].first.first <= cur)</pre>
             max(pos.{in[i].first.second.in[i].second}).
           i++;
                           if (pos.first == cur) return {}:
                          a.push back(pos.second);
             return a:
 6 Math
  Catalan Numbers
 ull* catalan = new ull[1000000];
 void genCatalan(int n, int mod) {
  catalan[0] = catalan[1] = 1;
  for (int i = 2; i <= n; i++) {</pre>
       catalan[i] = 0;
      for (int j = i - 1; j >= 0; j--) {
    catalan[i] += (catalan[j] * catalan[i-j-1])
            % mod;
        if (catalan[i] >= mod)
   catalan[i] -= mod;
  // TODO: consider binomial coefficient method
  Combinatorics (nCr, nPr)
   // can optimize by precomputing factorials, and
           fact[n]/fact[n-r]
 ull nPr(ull n, ull r) {
   .ull v =
   .ull v = 1;
.for (ull i = n-r+1; i <= n; i++)
    return v;
 ull nPr(ull n, ull r, ull m) {
    for (ull i = n-r+1; i \le n; i++)
     v = (v * i) \% m;
    return v:
 ull nCr(ull n, ull r) {
  long double v = 1;
    for (ull i = 1: i <= r: i++)
    v = v * (n-r+i) /i;
return (ull)(v + 0.001);
   // requires modulo math
  // can optimize by precomputing mfac and

→ minv-mfac

 ull nCr(ull n, ull r, ull m) {
  return mfac(n, m) * minv(mfac(k, m), m) % m *
\downarrow \hookrightarrow \min_{k} \min_{k}
```

```
Multinomials
                                                            ..for (int i = 2*p; i<=n; i += p) phi[i] =
(phi[i]/p) * (p-1);
                                                          Factorials
 return c;
                                                          // digits in factorial
                                                          #define kamenetsky(n) (floor((n * log10(n /
Chinese Remainder Theorem
                                                           \rightarrow M_E) + (log10(2 * M_PI * n) / 2.0)) + 1)
bool ecrt(ll* r, ll* m, int n, ll& re, ll& mo)
                                                          // approximation of factorial
#define stirling(n) ((n == 1) ? 1 : sqrt(2 *
 11 x, y, d; mo = m[0]; re = r[0];
for (int i = 1; i < n; i++) {
   d = egcd(mo, m[i], x, y);
   if ((r[i] - re) % d != 0) return false;
   x = (r[i] - re) / d * x % (m[i] / d);
   re += x * mo;
   re -= x * mo;</pre>
                                                          \hookrightarrow M PI * n) * pow(n / M_E, n))
                                                          // natural log of factorial
#define lfactorial(n) (lgamma(n+1))
                                                          Prime Factorization
  mo = mo / d * m[i]:
                                                          // do not call directly
  re %= mo;
                                                          ll pollard_rho(ll n, ll s) {
                                                           .ll x, y;
 re = (re + mo) \% mo;
                                                           x = y = rand() \% (n - 1) + 1;
 return true;
                                                           int head = 1, tail = 2;
while (true) {
                                                           x = mult(x, x, n);
x = (x + s) % n;
if (x == y) return n;
Count Digit Occurences
/*count(n,d) counts the number of occurences of

→ a digit d in the range [0,n]*/

                                                            11 d = _{gcd(max(x - y, y - x), n)};
11 digit count(ll n, ll d) {
                                                            if (1 < \overline{d} \&\& d < n) return d;
                                                            if (++head == tail) y = x, tail <<= 1;
 ll result = 0;
while (n != 0)
 result += ((n\%10) == d?1:0);
  n /= 10;
                                                          // call for prime factors
                                                          void factorize(ll n, vector<ll> &divisor) {
  if (n == 1) return;
 return result;
                                                           if (isPrime(n)) divisor.push_back(n);
11 count(11 n, 11 d) {
   if (n < 10) return (d > 0 && n >= d);
   if ((n % 10) != 9) return digit_count(n, d) +
                                                            while (d >= n) d = pollard_rho(n, rand() % (n | int josephus(int n, int k) {
 \rightarrow count(n-1, d);
                                                           \rightarrow -1) +1);
factorize(n / d. divisor);
 return 10*count(n/10, d) + (n/10) + (d > 0);
                                                            factorize(d, divisor);
Discrete Logarithm
unordered_map<int, int> dlogc;
                                                          Farev Fractions
int discretelog(int a, int b, int m) {
 dlogc.clear();
                                                             generate 0 \le a/b \le 1 ordered. b \le n
 11 \text{ n} = \text{sqrt}(\text{m}) + 1, \text{ an} = 1;
                                                              farey(4) = 0/1 1/4 1/3 1/2 2/3 3/4 1/1
 for (int i = 0; i < n; i++)
an = (an * a) % m;
                                                           // length is sum of phi(i) for i = 1 to n
                                                          vector<pair<int, int>> farey(int n) {
 11 c = an:
                                                           int h = 0, k = 1, x = 1, y = 0, r;
 for (int i = 1; i <= n; i++) {
   if (!dlogc.count(c)) dlogc[c] = i;
                                                           vector<pair<int, int>> v;
                                                           do {
  c = (c * an) \% m;
                                                            v.push_back({h, k});
                                                            r = (n-y)/k;
                                                            y += r*k; x' += r*h;
 for (int i = 0: i <= n: i++) {
                                                           x = -x; y = -y;

while (k > 1);
 if (dlogc.count(c)) return (dlogc[c] * n - i
 \rightarrow + m - 1) % (m-1);
  c = (c * a) \% m;
                                                           v.push_back({1, 1});
                                                           return v:
 return -1;
                                                          Fast Fourier Transform
Euler Phi / Totient
                                                          const double PI = acos(-1);
int phi(int n) {
                                                          void fft(vector<cd>& a, bool invert) {
 int^r = n;
 for (int i = 2; i * i <= n; i++) {
   if (n % i == 0) r -= r / i;
   while (n % i == 0) n /= i;
                                                           int n = a.size();
                                                           for (int i = 1, j = 0; i < n; i++) {
  int bit = n >> 1;
                                                            for (; j & bit; bit >>= 1) j ^= bit;
                                                            j ^= bit:
 if (n > 1) r = r / n;
 return r;
                                                            if (i < j) swap(a[i], a[j]);
 #define n 100000
                                                           for (int len = 2; len <= n; len <<= 1) {
    double ang = 2 * PI / len * (invert ? -1 :
ll phi[n+1];
void computeTotient() {
 for (int i=1; i<=n; i++) phi[i] = i;
                                                            cd wlen(cos(ang), sin(ang));
 for (int p=2; p<=n; p++) {
                                                            for (int i = 0; i < n; i += len) {
 if (phi[p] == p) {
   phi[p] = p-1;
                                                             for (int j = 0; j < len / 2; j++) {
```

```
....cd u = a[i+j], v = a[i+j+len/2] * w;
   a[i+j] = u + v;
a[i+j+len/2] = u - v;
w *= wlen;
 if (invert)
 for (auto\& x : a)
  x /= n;
vector<int> fftmult(vector<int> const& a,

  vector<int> const& b) {
  vector<cd> fa(a.begin(), a.end()),
fb(b.begin(), b.end());
int n = 1 << (32 - _builtin_clz(a.size() +</pre>
 \rightarrow b.size() - 1)):
 fa.resize(n); fb.resize(n);
 fft(fa, false); fft(fb, false)
 for (int i = 0; i < n; i++) fa[i] *= fb[i];
fft(fa, true);</pre>
 vector<int> toret(n):
 for (int i = 0; i < n; i++) toret[i] =
→ round(fa[i].real());
return toret;
Greatest Common Denominator
ll egcd(ll a, ll b, ll& x, ll& y) {
  if (b == 0) { x = 1; y = 0; return a; }
 ll gcd = egcd(b, a % b, x, y);
 x = a / b * y;
 swap(x, y);
return gcd;
Josephus Problem
// 0-indexed, arbitrary k
 if (n == 1) return 0;
if (k == 1) return n-1;
 if (k > n) return (josephus(n-1,k)+k)%n;
 int res = josephus(n-n/k,k)-n\%k;
return res + ((res<0)?n:res/(k-1));
// fast case if k=2, traditional josephus
int josephus(int n) {
return 2*(n-(1<<(32-\_builtin_clz(n)-1)));
Least Common Multiple
#define lcm(a,b) ((a*b)/__gcd(a,b))
Modulo Operations
#define MOD 1000000007
#define madd(a,b,m) (a+b-((a+b-m>=0)?m:0))
#define mult(a,b,m) ((ull)a*b%m)
#define msub(a,b,m) (a-b+((a<b)?m:0))
|11 mpow(11 b, 11 e, 11 m) {
11 \ x = 1:
 while (e'> 0) {
   if (e % 2) x = (x * b) % m;
 b = (b * b) \% m;
  e /= 2;
 return x % m:
ull mfac(ull n, ull m) {
  ull f = 1;
  for (int i = n; i > 1; i--)
 f = (f * i) \frac{\pi}{n}
return f;
// if m is not guaranteed to be prime
ll minv(ll b, ll m) {
.11 x = 0, y = 0;
 if (egcd(b, m, x, y) != 1) return -1;
 return (x % m + m) % m:
Il mdiv_compmod(int a, int b, int m) {
if (\gcd(b, m) != 1) return -1;
```

```
return mult(a, minv(b, m), m);
\frac{1}{2} if m is prime (like 10^{9}+7)
ll mdiv_primemod (int a, int b, int m) {
return mult(a, mpow(b, m-2, m), m);
// tonelli shanks = sqrt(n) % m, m is prime
ll legendre(ll a, ll m){
if (a % m==0) return 0;
 if (m == 2) return 1;
return mpow(a, (m-1)/2, m);
Íl msart(ll n. ll m) {
ll s = builtin ctzll(m-1), q = (m-111)>>s,
   z = rand()\%(m-1)+1;
if (m == 2) return 1;
if (s == 1) return mpow(n,(m+1)/411,m);
 while (legendre(z,m)!=m-1) z = rand()\%(m-1)+1;
 11 c = mpow(z,q,m), r = mpow(n,(q+1)/2,m), t
 \rightarrow = mpow(n,q,m), M = s;
 while (t != 1){
    ll i=1, ts = (t * t) \% m;
  while (ts != 1) i++, ts = (ts * ts) % m;
  11 b = c;
 for (int'j = 0; j < M-i-1; j++) b = (b * b) %
r = r * b \% m; c = b * b \% m; t = t * c \% m;
\rightarrow M = i:
 return r:
Modulo Tetration
11 tetraloop(ll a, ll b, ll m) {
.if(b == 0 ] | a == 1) return 1:
 ll w = tetraloop(a,b-1,phi(m)), r = 1;
 for (:w:w/=2) {
 if (w&1)
  r *= a; if (r >= m) r -= (r/m-1)*m;
 a *= a: if (a >= m) a -= (a/m-1)*m:
 return r:
int tetration(int a, int b, int m) {
 if (a == 0 \mid | m == 1) \text{ return } ((b+1)\&1)\%m;
 return tetraloop(a,b,m) % m;
Matrix
template<typename T>
struct Mat : public Vec<2, T> {
 Mat(int x, int y) : Vec<2, T>(x, y), w(x),
\hookrightarrow h(y) {}
 static Mat<T> identity(int n) { Mat<T> m(n,n);
    for (int i=0:i<n:i++) m[i][i] = 1: return
 Mat<T>& operator+=(const Mat<T>& m) {
  for (int i = 0; i < w; i++)
   for (int j = 0; j < h; j++)

(*this)[i][j] += m[i][j];
  return *this:
 Mat<T>& operator-=(const Mat<T>& m) {
 for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
(*this)[i][j] -= m[i][j];
 return *this;
 Mat<T> operator*(const Mat<T>& m) {
 Mat<T>z(w,m.h);
  for (int i = 0; i < w; i++)
  inf (int j = 0; j < h; j++)
    for (int k = 0; k < m.h; k++)
    z[i][k] += (*this)[i][j] * m[j][k];</pre>
    return z:
 Mat<T> operator+(const Mat<T>& m) { Mat<T>

    a=*this; return a+=m; }

Mat<T> operator-(const Mat<T>& m) { Mat<T>

→ a=*this: return a-=m: }
```

```
Mat<T>& operator*=(const Mat<T>& m) { return
                                                          int get_permutation(vector<int>& v) {
\rightarrow *this = (*this)*m; }
                                                           int use = 0, i = 1, r = 0;
for (int e : v) {
 Mat<T> power(int n) {
                                                            r = r * i++ + __builtin_popcount(use &
  .Mat<T> a = Mat<T>::identity(w), m=*this;
                                                            \rightarrow -(1<<e));
  for (;n;n/=2,m*=m) if (n\&1) a *= m;
                                                            use |= 1 << e;
                                                            return r;
Matrix Exponentiation
                                                           Permutation (string/multiset)
// F(n) = c[0]*F(n-1) + c[1]*F(n-2) + ...
// b is the base cases of same length c
                                                           string freq2str(vector<int>& v) {
ll matrix_exponentiation(ll n, vector<ll> c,
→ vector<ll> b) {
  if (nth < b.size()) return b[nth-1];</pre>
                                                           for (int i = 0; i < v.size(); i++)
                                                            for (int j = 0; j < v[i]; j++)
...s += (char)(i + 'A');
 Mat<11> a(c.size(), c.size()); ll s = 0;
for (int i = 0; i < c.size(); i++) a[i][0] =
                                                           return s:

    c[i];
for (int i = 0; i < c.size() - 1; i++)</pre>
                                                           // nth perm of multiset. n is O-indexed
                                                           string gen_permutation(string s, ll n) {
\hookrightarrow a[i][i+1] = 1;
                                                           vector<int> freq(26, 0);
a = a.power(nth - c.size());
for (int i = 0; i < c.size(); i++)
s += a[i][0] * b[i];</pre>
                                                           for (auto e : s) freq[e - 'A']++;
                                                           for (int i = 0: i < 26: i++) if (freg[i] > 0)
return s;
                                                            freq[i]--; ll v = multinomial(freq);
if (n < v) return (char)(i+'A') +</pre>
Matrix Subarray Sums
                                                               gen_permutation(freq2str(freq), n);
template<class T> struct MatrixSum {
                                                             freq[\overline{i}]++; n -= v;
 Vec<2, T> p;
MatrixSum(Vec<2, T>& v) {
                                                           return "":
  p = Vec<2,T>(v.size()+1, v[0].size()+1);
  for (int i = 0; i < v.size(); i++)
                                                           Miller-Rabin Primality Test
 for (int j = 0; j < v[0].size(); j++)
  p[i+1][j+1] = v[i][j] + p[i][j+1] +
                                                           // Miller-Rabin primality test - O(10 log^3 n)
                                                           bool isPrime(ull n) {

    p[i+1][j] - p[i][j];

                                                           if (n < 2) return false;
if (n == 2) return true;
if (n % 2 == 0) return false;
ull s = n - 1;
while (s % 2 == 0) s /= 2;
for (int i = 0; i < 10; i++) {</pre>
 T sum(int u, int 1, int d, int r) {
    return p[d][r] - p[d][1] - p[u][r] + p[u][1];
                                                            ull temp = s;
Mobius Function
                                                             ull a = rand() % (n - 1) + 1;
const int MAXN = 10000000;
                                                             ull mod = mpow(a, temp, n);
// mu[n] = 0 iff n has no square factors
                                                             while (temp!=n-1\&\&mod!=1\&\&mod!=n-1) {
// 1 = even number prime factors, -1 = odd
                                                              mod = mult(mod, mod, n);
short mu[MAXN] = \{0,1\};
                                                              temp *= 2;
void mobius(){
 for (int i = 1; i < MAXN; i++)
                                                             if (mod!=n-1&&temp%2==0) return false;
 if (mu[i])
 for(int'j = i + i; j < MAXN; j += i)
                                                            return true;
    mu[j] -= mu[i];
                                                           Sieve of Eratosthenes
Nimber Arithmetic
                                                           bitset<100000001> sieve;
#define nimAdd(a,b) ((a)^(b))
                                                           // generate sieve - O(n \log n)
ull nimMul(ull a, ull b, int i=6) {
   static const ull M[]={INT_MIN>>32,
                                                           void genSieve(int n) {
                                                           sieve[0] = sieve[1] = 1;
    M[0]^(M[0] << 16), M[1]^(M[1] << 8),
                                                           for (ull i = 3; i * i < n; i += 2)
...if (!sieve[i])
    M[2]^(M[2] << 4), M[3]^(M[3] << 2),
\stackrel{\rightarrow}{\Rightarrow} M[\stackrel{\leftarrow}{4}]^{\sim}(M[\stackrel{\leftarrow}{4}]<<1));
                                                             for (ull j = i * 3; j <= n; j += i * 2)
  if (i-- == 0) return a&b;
                                                               sieve[i] = 1:
  int k=1<<i:
  ull s=nimMul(a,b,i), m=M[5-i],
                                                           \frac{1}{2} query sieve after it's generated - O(1)
    t=nimMul(((a^(a>>k))\&m)|(s\&~m),
                                                           bool quervSieve(int n) {
   ((b^{(b>k)})\&m)|(m\&(~m>>1))<< k, i);
                                                           return n == 2 || (n % 2 != 0 && !sieve[n]);
  return ((s^t)\&m)<\langle k|((s^(t>>k))\&m);
                                                           Compile-time Prime Sieve
Permutation
                                                           const int MAXN = 100000;
// c = array size, n = nth perm, return index
                                                           template<int N>
vector<int> gen permutation(int c, int n) {
                                                           struct Sieve {
 vector<int> idx(c), per(c), fac(c); int i;
                                                           bool sieve[N];
                                                           for (i = 0; i < c; i++) idx[i] = i;
for (i = 1; i <= c; i++) fac[i-1] = n%i, n/=i;
for (i = c - 1; i >= 0; i--)
per[c-i-1] = idx[fac[i]],
                                                              for (int j = i * 2; j < N; j += i)
...sieve[j] = 1;
  idx.erase(idx.begin() + fac[i]);
 return per;
\} // get what nth permutation of vector
```

```
bool isPrime(int n) {
  static constexpr Sieve<MAXN> s;
                                                                  int cur = 0:
                                                                  for (int i = 0; i < n; i++) {
                                                                   for (int j = cur; j < m; j++) {
 return !s.sieve[n]:
                                                                    if (!zero(a[j][i])) {
                                                                    iif (j != cur) swap(a[j], a[cur]);
for (int sat = 0; sat < m; sat++) {
   if (sat == cur) continue;</pre>
Simpson's / Approximate Integrals
// integrate f from a to b, k iterations
// error \langle = (b-a)/18.0 * M * ((b-a)/2k)^2 /

// where M = max(abs(f```(x))) for x in [a,b]

// "f" is a function "double func(double x)"
                                                                       double num = a[sat][i] / a[cur][i];
for (int sot = 0; sot <= n; sot++)
a[sat][sot] -= a[cur][sot] * num;</pre>
double Simpsons (double a, double b, int k,
                                                                      cur++:

    double (*f)(double)) {
    double dx = (b-a)/(2.0*k), t = 0;
    for (int i = 0; i < k; i++)
    t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
}
</pre>
                                                                      break:
 \leftrightarrow (*f)(a+(2*i+1)*dx);
return (t + (*f)(b)) * (b-a) / 6.0 / k;
                                                                  for (int j = cur; j < m; j++)
                                                                   if (!zero(a[j][n])) return vector<double>();
                                                                  vector<double > ans(n,0);
                                                                  for (int i = 0, sat = 0; i < n; i++)
if (sat < m && !zero(a[sat][i]))
ans[i] = a[sat][n] / a[sat++][i];
Common Equations Solvers
// ax^2 + bx + c = 0, find x
vector<double> solveEq(double a, double b,
                                                                  return ans;
 → double c) {
vector<double> r;
                                                                 \frac{1}{2} solve A[n][n] * x[n] = b[n] linear equation
  double z = b * b - 4 * a * c;
if (z == 0)
                                                                 // rank < n is multiple solutions, -1 is no

→ solutions
// `alls` is whether to find all solutions, or

  r.push_back(-b/(2*a));
  else if (z > 0) {
    r.push_back((sqrt(z)-b)/(2*a));
                                                                 \hookrightarrow any
                                                                 const double eps = 1e-12;
int solveEq(Vec<2, double>& A, Vec<1, double>&
  r.push\_back((sqrt(z)+b)/(2*a));

    b, Vec<1, double>& x, bool alls=false) {
    int n = A.size(), m = x.size(), rank = 0, br,
}

 \frac{1}{2} / ax^3 + bx^2 + cx + d = 0, find x
                                                                  vector<int> col(m); iota(begin(col), end(col),
vector<double> solveEq(double a, double b,
  → double c, double d) {
vector<double> res;
                                                                  for(int i = 0; i < n; i++) {
                                                                   double v, bv = 0;
for(int r = i; r < n; r++)
  long double a1 = b/a, a2 = c/a, a3 = d/a;
  long double q = (a1*a1 - 3*a2)/9.0, sq =
                                                                    for(int c = i; c < n; c++)
if ((v = fabs(A[r, [c])) > bv)
  \rightarrow -2*sqrt(q);
 long double r = (2*a1*a1*a1 - 9*a1*a2 +
                                                                   br = r, bc = c, bv = v;
if (bv <= eps) {
  \rightarrow 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
                                                                    for(int j = i; j < n; j++)
if (fabs(b[j]) > eps)
  if (z <= 0) {
   theta = acos(r/sqrt(q*q*q));
                                                                      .return -1;
   res.push_back(sq*cos(theta/3.0) - a1/3.0);
                                                                     break;
   res.push_back(sq*cos((theta+2.0*PI)/3.0) -
                                                                   swap(A[i], A[br]);
                                                                   swap(b[i], b[br]);
  res.push_back(sq*cos((theta+4.0*PI)/3.0) -
                                                                    swap(col[i], col[bc]);
 \rightarrow a1/3.0):
                                                                   for(int j = 0; j < n; j++)

swap(A[j][i], A[j][bc]);

bv = 1.0 / A[i][i];
  res.push_back(pow(sqrt(z)+fabs(r), 1/3.0));
                                                                   for(int j = (alls)?0:i+1; j < n; j++) {
   res[0] = (res[0] + q / res[0]) *
                                                                    if (j != i) {
     ((r<0)?1:-1) - a1 / 3.0;
                                                                      double fac = A[j][i] * bv;
 return res:
                                                                      .b[j] -= fac * b[i];
                                                                     for(int k = i+1; k < m; k++)
A[j][k] -= fac*A[i][k];
\frac{1}{1} linear diophantine equation ax + by = c,
     find x and y
// infinite solutions of form x+k*b/g, y-k*a/g bool solveEq(11 a, 11 b, 11 c, 11 &x, 11 &y, 11
                                                                   rank++;
                                                                  if (alls) for (int i = 0; i < m; i++) x[i] =
  g = egcd(abs(a), abs(b), x, y);
                                                                     -DBL MAX;
 if (c % g) return false;
                                                                  for (int i = rank; i--;) {
   bool isGood = true;
 x *= c / g * ((a < 0) ? -1 : 1);
 y *= c / g * ((b < 0) ? -1 : 1);
                                                                   if (aļls)
                                                                    for (int j = rank; isGood && j < m; j++)
if (fabs(A[i][j]) > eps)
 return true;
\frac{1}{m} = \# equations, n = \# variables, a[m][n+1]
                                                                   isGood = false;
b[i] /= A[i][i];
| \rightarrow | = coefficient matrix
                                                                   if (isGood) x[col[i]] = b[i];
// a[i][0]x + a[i][1]y + ... + a[i][n]z =
                                                                   if (!alls)
for(int j = 0; j < i; j++)
     a[i][n+1]
// find a solution of some kind to linear
                                                                     b[i] -= A[i][i] * b[i];
\rightarrow equation
const double eps = 1e-7;
                                                                  return rank:
bool zero(double a) { return (a < eps) && (a >
 → -eps); }
vector double solve Eq (double **a, int m, int
```

Graycode Conversions ull graycode2ull(ull n) { ull i = 0; for (; n; n = n >> 1) i ^= n; return i; ull ull2graycode(ull n) { return n ^ (n >> 1); Unix/Epoch Time // 0-indexed month/time, 1-indexed day // minimum 1970, 0, 1, 0, 0, 0 ull toEpoch(int year, int month, int day, int → hour, int minute, int second) { t.tm_man = minute; the second; t.tm_year = year - 1900; t.tm_mon = month; t.tm_mday = day; t.tm_hour = hour; t.tm_min = minute; t.tm_sec = second; t.tm_isdst = 0; // 1 = daylights savings epoch = mktime(&t); return (ull)epoch; vector<int> toDate(ull epoch) { time_t e=epoch; struct tm t=*localtime(&e); return {t.tm_year+1900,t.tm_mon,t.tm_mday,t_ .tm hour.t.tm min.t.tm sec}: int getWeekday(ull epoch) { time_t e=epoch; struct tm t=*localtime(&e); return t.tm_wday; // 0-6, 0 = sunday int getDavofYear(ull epoch) { time_t e=epoch; struct tm t=*localtime(&e); return t.tm_yday; // 0-365 const int months[] = Const the models: → {31,28,31,30,31,30,31,30,31,30,31}; bool validDate(int year, int month, int day) { bool leap = !(year%(year%25?4:16)); if (month >= 12) return false; return day <= months[month] + (leap && month == 1):

Theorems and Formulae

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set of size n. !0 = 1, !1 = 0, !n = (n+1)(!(n-1))1)+!(n-2)), ! $n = n! \sum_{i=0}^{n} \frac{(-1)^i}{i!}$, ! $n = \left[\frac{n!}{e}\right]$

In a partially ordered set, a chain is a subset of elements that are all comparable to eachother An antichain is a subset where no two are comparable.

Dilworth's theorem states the size of a maximal antichain equals the size of a minimal chain cover of a partially ordered set S. The width of S is the maximum size of an antichain in S, which is equal to the minimum number of chains needed to cover S, or the minimum number of chains such that all elements are in at least one chain.

Rosser's Theorem states the nth prime number is greater than n * ln(n) for n > 1.

 $(n^{\frac{n+1}{2}})^2$

Lagrange's Four Square Theorem states | FOR(k,n) FOR(i,n) FOR(j,n) if (m[i][k] != inf | convex_polygon(triangle a) { every natural number is the sum of the squares of four non-negative integers. This is a special case of the Fermat Polygonal Number for (k,n) if (m[k][k] < 0) FOR(j,n) Theorem where every positive integer is a lif (m[i][k] != inf && m[k][j] != inf) sum of at most n s-gonal numbers. The $nth \rightarrow m[i][j] = -inf;$ s-gonal number $P(s,n) = (s-2)\frac{n(n-1)}{2} + n$

Graphs struct edge { int u,v,w;
edge (int u,int v,int w) : u(u),v(v),w(w) {} edge (): u(0), v(0), w(0) {} bool operator < (const edge &e1, const edge bool operator > (const edge &e1, const edge

Eulerian Path #define edge_list vector<edge>
#define_adj_sets vector<set<int>>> struct EulerPathGraph { adj_sets graph; // actually indexes incident edges edge_list edges; int n; vector<int> indeg; EulerPathGraph(int n): n(n) { indeg = *(new vector<int>(n,0)); graph = *(new adj_sets(n, set<int>())); void add_edge(int u, int v) { graph[u].insert(edges.size()); indeg[v]++; edges.push back(edge(u,v,0)); bool eulerian_path(vector<int> &circuit) { if(edges.size()==0) return false; stack<int> st;
 int a[] = {-1, -1};
 for(int v=0;v<n;v++) {
 if(indeg[v]!=graph[v].size()) {
 bool b = indeg[v] > graph[v].size();
} if (abs(((int)indeg[v])-((int)graph[v] .size())) > 1) returnfalse;
if (a[b] != -1) return false; a[b] = v;int s = (a[0]!=-1 && a[1]!=-1 ? a[0] :(a[0]=-1 & a a[1]=-1 ? edges[0].u : -1);if(s==-1) return false;
while(!st.empty() || !graph[s].empty()) { if (graph[s].empty()) { circuit.push_back(s); s = st.top(); st.pop(); } else { int w = edges[*graph[s].begin()].v; graph[s].erase(graph[s].begin()); $\bar{s}t.\bar{p}ush(s); s = \bar{w};$

Flovd Warshall

circuit.push_back(s);

Nicomachi's Theorem states $1^3 + 2^3 + ... + \frac{\text{const 11 inf} = 1 \text{LL} << 62;}{\text{#define FOR}(i,n) for (int i = 0; i < n; i++)}$ $n^3 = (1+2+...+n)^2$ and is equivalent to void floydWarshall (Vec<2, 11>& m) { int n = m.size();
FOR(i,n) m[i][i] = min(m[i][i], OLL);

return circuit.size()-1==edges.size();

```
\rightarrow && m[k][j] != inf) {
 auto newDist = max(m[i][k] + m[k][j], -inf);
 m[i][j] = min(m[i][j], newDist);
Minimum Spanning Tree
// returns vector of edges in the mst
// graph[i] = vector of edges incident to
    vertex i
   places total weight of the mst in Stotal
// if returned vector has size != n-1, there is
vector<edge> mst(vector<vector<edge>> graph.
 → 11 &total) {
total = 0:
 priority_queue<edge, vector<edge>,
 → greater<edge>> pq;
 vector<edge> MST;
 bitset<20001> marked; // change size as needed
 marked[0] = 1;
 for (edge ep : graph[0]) pq.push(ep);
while(MST.size()!=graph.size()-1 &&
 → pq.size()!=0) {
 pq.size():-0/(
edge e = pq.top(); pq.pop();
int u = e.u, v = e.v, w = e.w;
if(marked[u] && marked[v]) continue;
else if(marked[u]) swap(u, v);
  for(edge ep : graph[u]) pq.push(ep);
  marked[u] = 1;
MST.push_back(e);
  total += e.w:
 return MST;
Union Find
int uf find(subset* s, int i) {
  if (s[i].p != i) s[i].p = uf_find(s, s[i].p);
return s[i].p;
void uf_union(subset* s, int x, int y) {
int xp = uf_find(s, x), yp = uf_find(s, y);
if (s[xp].rank > s[yp].rank) s[yp].p = xp;
else if (s[xp].rank < s[yp].rank) s[xp].p =</pre>
 \rightarrow yp;
 else { s[yp].p = xp; s[xp].rank++; }
2D Grid Shortcut
#define fordir(x, y, n, m) for(auto[dx, dy]:dir)if
 \rightarrow (inbound(x+dx,n)&\text{Sinbound}(y+dy,m))
const pair<int,int> dir[] =
\rightarrow {{1,0},{0,1},{-1,0},{0,-1}};
    2D Geometry
#define point complex<double>
#define EPS 0.0000001
#define sq(a) ((a)*(a))
#define c\dot{b}(a) ((a)*(a)*(a))
double dot(point a, point b) { return

    real(conj(a)*b);
}
double cross(point a, point b) { return

    imag(conj(a)*b); }
```

struct line { point a, b; };

struct convex_polygon {

vector<point> points;

→ points(points) {}

struct segment { point a, b; };

struct triangle { point a, b, c; };
struct rectangle { point tl, br; };

struct circle { point c; double r; };

convex_polygon(vector<point> points) :

```
points.push back(a.a); points.push back(a.b);
   points.push back(a.c);
 convex_polygon(rectangle a) {
 points.push_back(a.tl);
    points.push back({real(a.tl).
   imag(a.br)});
 points.push_back(a.br);
    points.push_back({real(a.br),
    imag(a.tl)}):
struct polygon {
 vector<point> points;
 polygon(vector<point> points) :
 → points(points) {}
 polygon(triangle a) {
 points.push_back(a.a); points.push_back(a.b);
   points.push back(a.c);
polygon(rectangle a) {
 points.push back(a.tl);
    points.push back({real(a.tl),
   imag(a.br)});
 points.push_back(a.br);
   points.push_back({real(a.br),
    imag(a.tl)});
polygon(convex_polygon a) {
 for (point v : a.points)
  points.push_back(v);
// triangle methods
double area_heron(double a, double b, double
\hookrightarrow c) {
if (a < b) swap(a, b);
if (a < c) swap(a, c);
 if (b < c) swap(b, c);
if (a > b + c) return -1;
return sqrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)
( /16.0);
// segment methods
double lengthsq(segment a) { return
    sq(real(a.a) - real(a.b)) + sq(imag(a.a) -
   imag(a.b)); }
double length(segment a) { return

    sqrt(lengthsq(a)); }

   circle methods
double circumference(circle a) { return 2 * a.r

→ * M PI; }

double area(circle a) { return sq(a.r) * M PI:
→ }
// rectangle methods
double width(rectangle a) { return
⇒ abs(real(a.br) - real(a.tl)); }
double height(rectangle a) { return

→ abs(imag(a.br) - real(a.tl)); }

double diagonal (rectangle a) { return

    sqrt(sq(width(a)) + sq(height(a)));
]
double area (rectangle a) { return width(a) >
\rightarrow height(a); }
double perimeter(rectangle a) { return 2 *
   (width(a) + height(a)); }
// check if `a` fit's inside `b
// swap equalities to exclude tight fits
bool doesfitInside(rectangle a, rectangle b) {
int x = width(a), w = width(b), y = height(a),
\rightarrow h = height(b);
if (x > y) swap(x, y);
if (w > h) swap(w, h):
```

```
if (w < x) return false;
 if (y <= h) return true;
 double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
return sq(a) \le sq(b) + sq(c);
// polygon methods
// negative area = CCW, positive = CW
double area(polygon a) {
  double area = 0.0; int n = a.points.size();
  for (int i = 0, j = 1; i < n; i++, j = (j +
    area += (real(a.points[j]-a.points[i]))*
return area / 2.0;
// get both unsigned area and centroid
pair<double, point> area_centroid(polygon a) {
 int n = a.points.size();
double area = 0;
point c(0, 0);
 for (int i = n - 1, j = 0; j < n; i = j++) {
 double v = cross(a.points[i], a.points[j]) /
 area += v:
  c += (a.points[i] + a.points[j]) * (v / 3);
 c /= area:
return {area, c};
Intersection
// -1 coincide, 0 parallel, 1 intersection
int intersection(line a, line b, point& p) {
if (abs(cross(a.b - a.a, b.b - b.a)) > EPS) {
 p = cross(b.a - a.a, b.b - a.b) / cross(a.b)
\rightarrow - a.a, b.b - b.a) * (b - a) + a;
 return 1:
 if (abs(cross(a.b - a.a, a.b - b.a)) > EPS)

→ return 0;

return -1;
}
// area of intersection
double intersection(circle a, circle b) {
double d = abs(a.c - b.c);
if (d <= b.r - a.r) return area(a);
if (d <= a.r - b.r) return area(b);</pre>
 if (d \ge a.r + b.r) return 0;
 double alpha = acos((sq(a.r)' + sq(d) -
\rightarrow sg(b.r)) / (2 * a.r * d)):
double beta = acos((sq(b.r) + sq(d) - sq(a.r))
\rightarrow / (2 * b.r * d));
return sq(a.r) * (alpha - 0.5 * sin(2 *
    alpha)) + sq(b.r) * (beta - 0.5 * sin(2 *
    beta));
// -1 outside, 0 inside, 1 tangent, 2
intersection
int intersection(circle a, circle b,

    vector<point>& inter) {

 double d2 = norm(b.c - a.c), rS = a.r + b.r,
\rightarrow rD = a.r - b.r;
if (d2 > sq(rS)) return -1;
 if (d2 < sq(rD)) return 0;
 double ca = 0.5 * (1 + rS * rD / d2):
 point z = point(ca, sqrt(sq(a.r) / d2 -
\rightarrow sq(ca)));
 inter.push back(a.c + (b.c - a.c) * z);
 if (abs(imag(z)) > EPS) inter.push_back(a.c +
\rightarrow (b.c - a.c) * conj(z));
return inter.size();
// points of intersection
vector<point> intersection(line a. circle c) {
vector<point> inter;
c.c -= a.a;
a.b -= a.a;
 point m = \hat{a}.b * real(c.c / a.b);
 double d2 = norm(m - c.c);
```

```
if (d2 > sq(c.r)) return 0;
 double l = sqrt((sq(c.r) - d2) / norm(a.b));
inter.push_back(a.a + m + 1 * a.b);
if (abs(1) > EPS) inter.push_back(a.a + m - 1 \rightarrow cb(a.r) * M_PI; }
\rightarrow * a.b);
return inter:
double intersection(rectangle a, rectangle b) {
// area of intersection
double x1 = max(real(a.tl), real(b.tl)), y1 =
 → max(imag(a.tl), imag(b.tl));
double x2 = min(real(a.br), real(b.br)), y2 =
   min(imag(a.br), imag(b.br));
return (x2 <= x1 || y2 <= y1) ? 0 :
   (x2-x1)*(y2-y1);
Convex Hull
bool cmp(point a, point b) {
  if (abs(real(a) - real(b)) > EPS) return
    real(a) < real(b);
 if (abs(imag(a) - imag(b)) > EPS) return
   imag(a) < imag(b);</pre>
return false;
convex_polygon convexhull(polygon a) {
sort(a.points.begin(), a.points.end(), cmp);
 vector<point> lower, upper;
 for (int i = 0; i < a.points.size(); i++) {
 while (lower.size() >= 2 &&
    cross(lower.back() - lower[lower.size() -
    2], a.points[i] - lower.back()) < EPS)
  lower.pop_back();
  while (upper.size() >= 2 &&
    cross(upper.back() - upper[upper.size() -
   2], a.points[i] - upper.back()) > -EPS)
  upper.pop_back();
  lower.push back(a.points[i]);
  upper.push_back(a.points[i]);
 lower.insert(lower.end(), upper.rbegin() + 1,
   upper.rend());
return convex_polygon(lower);
    3D Geometry
struct point3d {
```

```
double x, y, z;
point3d operator+(point3d a) const { return
 \rightarrow {x+a.x, y+a.y, z+a.z}; }
point3d operator*(double a) const { return
 \rightarrow {x*a, y*a, z*a}; }
point3d operator-() const { return {-x, -y,
   -z}; }
point3d operator-(point3d a) const { return
   *this + -a: }
point3d operator/(double a) const { return
   *this * (1/a); }
double norm() { return x*x + y*y + z*z; }
double abs() { return sqrt(norm()); }
point3d normalize() { return *this /
    this->abs(); }
double dot(point3d a, point3d b) { return
\rightarrow a.x*b.x + a.y*b.y + a.z*b.z; }
point3d cross(point3d a, point3d b) { return
   \{a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,
\Rightarrow a.x*b.y - a.y*b.x}; }
struct line3d { point3d a, b; };
struct plane { double a, b, c, d; } // a*x +
\rightarrow b*y + c*z + d = 0
struct sphere { point3d c; double r; };
#define sq(a) ((a)*(a))
#define c\bar{b}(a) ((a)*(a)*(a))
```

```
double surface(circle a) { return 4 * sq(a.r) * GCC Builtin Docs

M_PI; }
double volume(circle a) { return 4.0/3.0 *
10 Optimization
 // SameNumberOfOneBits, next permutation
int snoob(int a) {
int b = a & -a, c = a + b;
return c | ((a \hat{c}) >> 2) / b;
// example usage
int main() {
   char l1[] = {'1', '2', '3', '4', '5', char l2[] = {'a', 'b', 'c', 'd'};
   int d1 = 5, d2 = 4;
   // prints 12345abcd, 1234a5bcd, ...
 int min = (1 << d1) -1, max = min << d2;
 for (int i = min; i <= max; i = snoob(i)) {
  int p1 = 0, p2 = 0, v = i;
  while (p1 < d1 || p2 < d2) {
    cout << ((v & 1) ? 11[p1++] : 12[p2++]);
   .v /= 2;
  cout << '\n';
Powers
bool isPowerOf2(ll a) {
 return a > 0 \&\& !(a \& a-1);
bool isPowerOf3(11 a) {
   return a>0&&!(12157665459056928801ull%a);
                                                           O(n^3)
bool isPower(ll a, ll b) {
  double x = log(a) / log(b);
                                                           O(n^2)
 return abs(x-round(x)) < 0.00000000001;
                                                          O(n\sqrt{n})
                                                           O(n \log n)
11 Additional
                                                           O(n)
Judge Speed
// kattis: 0.50s
  / codeforces: 0.421s
// atcoder: 0.455s
#include <bits/stdc++.h>
using namespace std;
\inf v = 1e9/2, p = 1;
int main() {
  for (int i = 1; i <= v; i++) p *= i;</pre>
 cout << p;
Judge Pre-Contest Checks
     int128 and float128 support?
-does extra or missing whitespace cause WA?
-documentation up to date?
-printer usage available and functional?
// each case tests a different fail condition
// try them before contests to see error codes
struct g { int arr[1000000]; g(){}};
vector<ğ> a;
// O=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE
⇒ 6=SIGSEGV 7=recursive MLE
int judge(int n) {
```

if (n == 0) exit(0);

if (n == 4) assert(0);
if (n == 5) 0 / 0;

return n + judge(n + 1);

if (n == 6) * (int*)(0) = 0:

if (n == 1) while(1);
if (n == 2) while(1) a.push_back(g());

if (n == 3) while(1) putchar_unlocked('a');

```
// 128-bit integer
 int128 a;
unsigned __int128 b;
// 128-bit float
// minor improvements over long double
__float128 c;
// log2 floor
__lg(n);
// number of 1 bits
// can add il like popcountll for long longs
__builtin_popcount(n);
// number of trailing zeroes
__builtin_ctz(n);
// number of leading zeroes
_builtin_clz(n);
// 1-indemed least significant 1 bit
__builtin_ffs(n);
// parity of number
__builtin_parity(n);
Limits
                        \pm 2147483647 \mid \pm 2^{31} - 1 \mid 10^9
int
                                               \bar{2}^{32} - \bar{1}|10^9
                          4294967295
uint
        \pm 9223372036854775807 | \pm \overline{2}^{63} - \overline{1}|\overline{10}^{18}
                                               2^{64} - 1|10^{19}
ull
         18446744073709551615
|i128| \pm 170141183460469231... | \pm 2^{\overline{1}27} - 1 | 10^{38}
|u128| 340282366920938463... | 2^{128} - 1 | 10^{38}
Complexity classes input size (per second):
O(n^n) or O(n!)
                                                        n < 10
O(2^n)
                                                       n < 30
```

n < 1000

n < 30000

 $n < 10^{6}$

 $n < 10^7$

 $n < 10^9$