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      General
run.sh
g++ -g -02 -std=gnu++17 -static prog.cpp
./a.exe
# compile and test all *.in and *.ans
g++ -g -02 -std=gnu++17 -static prog.cpp
for i in *.in; do
 f=${i%.in}
 ./a.exe < $i > "$f.out"
diff -b -q "$f.ans" "$f.out"
done
Header
// use better compiler options
#pragma GCC optimize("Ofast","unroll-loops")
#pragma GCC target("avx2,fma")
// include everything
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <bits/extc++.h>
#include <sys/resource.h>
// namespaces
using namespace std;
using namespace __gnu_cxx; // rope
using namespace __gnu_pbds; // tree/trie
// common defines
#define fastio

→ ios_base::sync_with_stdio(0);cin.tie(0);
#define nostacklim rlimit RZ; getrlimit(3,&RZ)
 \rightarrow );RZ.rlim cur=-1;setrlimit(3,\&RZ);
#define DEBUG(v) cerr<<__LINE__<<": "<<#v<<" =
\Rightarrow "<<v<<'\n'; #define TIMER
→ cerr<<1.0*clock()/CLOCKS_PER_SEC<<"s\n";
#define ll long long
#define ull unsigned ll
#define i128 _ int128
#define u128 unsigned i128
#define ld long double
// global variables
mt19937 rng((uint32_t)chrono::steady

    _clock::now().time_since_epoch().count());
#define getchar_unlocked() _getchar_nolock()
#define putchar_unlocked(x) _putchar_nolock(x)
void read(unsigned int& n) {
 char c; n = 0;
while ((c=getchar_unlocked())!=' '&&c!='\n')
 n = n * 10 + c - '0';
void read(int& n) {
   char c; n = 0; int s = 1;
   if ((c=getchar_unlocked())=='-') s = -1;
 else n = c - \sqrt{0}:
 while ((c=getchar_unlocked())!=' '&&c!='\n')
 n = n * 10 + c -
void read(ld& n) {
 char c; n = 0;

.ld m = 0, o = 1; bool d = false; int s = 1;

.if ((c=getchar_unlocked())=='-') s = -1;
 else if (c == '.') d = true;
else n = c - '0';
 while ((c=getchar_unlocked())!=' '&&c!='\n')
  if (c == '.') d = true;
```

```
else if (d) { m=m*10+c-'0'; o*=0.1; } else n = n * 10 + c - '0':
 n = s * (n + m * o):
void read(double& n) {
 ld m; read(m); n = m;
void read(float& n) {
  ld m; read(m); n = m;
 void read(string& s) {
 char c; s = ""
 while((c=getchar_unlocked())!=' '&&c!='\n')
bool readline(string& s) {
 char c; s = ""
 while(c=getchar_unlocked()) {
  if (c == '\n') return true;
if (c == EOF) return false;
  s += c:
 return false:
void print(unsigned int n) {
 if (n / 10) print(n / 10);
 putchar unlocked(n % 10 + '0'):
void print(int n) {
 if (n < 0) { putchar_unlocked('-'); n*=-1; }
 print((unsigned int)n);
Common Structs
   n-dimension vectors
// Vec<2, int> v(n, m) = arr[n][m]

// Vec<2, int> v(n, m, -1) default init -1

template<int D, typename T>
struct Vec : public vector < Vec < D-1. T>> {
 template<typename... Args>
 Vec(int n=0, Args... args) : vector<Vec<D-1,
 \rightarrow T>>(n, Vec<D-1, T>(args...)) {}
};
template<typename T>
struct Vec<1, T> : public vector<T> {
   Vec(int n=0. T val=T()) : vector<T>(n. val) {}
   Algorithms
Binary Search
 // search for k in [p,n)
template<typename T>
int binsearch(T x[], int k, int n, int p = 0) {
 for (int i = n; i >= 1; i /= 2)

while (p+i < n && x[p+i] <= k) p += i;

return p; // bool: x[p] == k;
Min/Max Subarray
   max - compare = a < b, reset = a < 0
 // min - compare = a > b, reset = a > 0
// returns {sum, {start, end}}
pair<int, pair<int, int>>
     ContiguousSubarray(int* a, int size,
     bool(*compare)(int, int),
 bool(*reset)(int), int defbest = 0) {
int best = defbest, cur = 0, start = 0, end =
 \rightarrow 0, s = 0;
 for (int i = 0: i < size: i++) {
  cur += a[i];
  if ((*compare)(best, cur)) { best = cur;
    start = s; end = i; }
  if ((*reset)(cur)) { cur = 0; s = i + 1; }
 return {best, {start, end}};
Quickselect
 #define QSNE -999999
```

```
int partition(int arr[], int 1, int r)
 int x = arr[r], i = 1;
for (int j = 1; j <= r - 1; j++)
if (arr[j] <= x)
swap(arr[i++], arr[j]);
 swap(arr[i], arr[r]);
 return i:
// find k'th smallest element in unsorted array,
\rightarrow only if all distinct
int gselect(int arr[], int 1, int r, int k)
 if (!(k > 0 && k <= r - 1 + 1)) return QSNE;
 swap(arr[1 + rng() % (r-1+1)], arr[r]);
 int pos = partition(arr, 1, r);
 if (pos-l==k-1) return arr[pos];
 if (pos-l>k-1) return qselect(arr,l,pos-1,k);
 return qselect(arr, pos+1, r, k-pos+1-1);
|}
|// TODO: compare against std::nth_element()
Saddleback Search
// search for v in 2d array arr[x][y], sorted
→ on both axis
pair<int. int> saddleback search(int** arr. int
\hookrightarrow x, int y, int v) {
int i = x-1, j = 0;
while (i >= 0 && j < y) {
   if (arr[i][j] == v) return {i, j};
  (arr[i][j] > v)? i--: j++;
 return {-1, -1};
Ternary Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
int ternsearch(int a. int b. int (*f)(int)) {
 while (b-a > 4) {
  int m = (a+b)/2;
if (TERNCOMP((*f)(m), (*f)(m+1))) a = m;
  else b = m+1:
for (int i = a+1; i <= b; i++)
if (TERNCOMP((*f)(a), (*f)(i)))
 . a = i;
return a:
#define TERNPREC 0.000001
double ternsearch (double a, double b, double
     (*f)(double)) {
 while (b-a > TERNPREC * 4) {
    double m = (a+b)/2:
  if (TERNCOMP((*f)(m), (*f)(m + TERNPREC))) a
  > = m;
else b = m + TERNPREC;
 for (double i = a + TERNPREC; i <= b; i +=
    TERNPREC)
      if (TERNCOMP((*f)(a), (*f)(i)))
 return a;
Golden Section Search
// < max, > min, or any other unimodal func
#define TERNCOMP(a,b) (a)<(b)
double goldsection(double a, double b, double
 double g(double)) {
  double r = (sqrt(5)-1)/2, eps = 1e-7;
  double x1 = b - r*(b-a), x2 = a + r*(b-a);
  double f1 = f(x1), f2 = f(x2);
 while (b-a > eps)
  while (b-a > eps)

if (TERNCOMP(f2,f1)) {

. b = x2; x2 = x1; f2 = f1;

. x1 = b - r*(b-a); f1 = f(x1);
   a = x1; x1 = x2; f1 = f2;
x2 = a + r*(b-a); f2 = f(x2);
 return a:
```

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3 Structures
Fenwick Tree
// Fenwick tree, array of cumulative sums
\hookrightarrow O(\log n) updates, O(\log n) gets
struct Fenwick { int n: ll* tree:
 void update(int i, int val) {
 .++i;
while (i <= n) {
  tree[i] += val;</pre>
   i += i & (-i);
 Fenwick(int size) {
  n = size;
  tree = new ll[n+1];
for (int i = 1; i <= n; i++)
   .tree[i] = 0;
 Fenwick(int* arr, int size) : Fenwick(size) {
  for (int i = 0; i < n; i++)
...update(i, arr[i]);
 ~Fenwick() { delete[] tree; }
 11 operator[](int i) {
  if (i < 0 || i > n) return 0;
  while (i>0)
  sum += tree[i];
i -= i & (-i);
  return sum;
 ll getRange(int a, int b) { return
    operator[](b) - operator[](a-1); }
Hashtable
// similar to unordered map, but faster
struct chash {
    const uint64_t C = (11)(2e18 * M_PI) + 71;
    ll operator()(11 x) const { return
    builtin bswap64(x*C): }
int main() {
  gp_hash_table<11,int,chash>
\rightarrow hashtable({},{},{},{},{1<<16});
for (int i = 0; i < 100; i++)

hashtable[i] = 200+i;

if (hashtable.find(10) != hashtable.end())

cout << hashtable[10];
Ordered Set
template <typename T>
using oset = tree<T,null type,less<T>,rb tree
    tag, tree order statistics node update>;
template <typename T, typename D> using omap = tree<T,D,less<T>,rb_tree
    _tag,tree_order_statistics_node_update>;
int main()
 oset<int> o_set;
o_set.insert(5); o_set.insert(1);

    o_set.insert(3);
// get second smallest element
 cout << *(o_set.find_by_order(1));</pre>
 // number of elements less than k=4
cout << ' ' << o_set.order_of_key(4) << '\n';</pre>
 // equivalent with ordered map
 omap<int,int> o_map;
o_map[5]=1;o_map[1]=2;o_map[3]=3;
cout << (*(o_map.find_by_order(1))).first;
 cout << ' ' << o_map.order_of_key(4) << '\n';
```

```
Rope
                                                         // print things with prefix "1"
                                                         auto range = trie.prefix_range("1");
// O(\log n) insert, delete, concatenate
                                                         for (auto it = range.first; it !=
int main() {
 // generate rope
                                                         → range.second: it++)
 rope<int> v;
                                                          cout << *it << '
 for (int i = 0; i < 100; i++)
.v.push_back(i);
                                                        Wavelet Tree
 // move range to front
                                                        using iter = vector<int>::iterator;
 rope<int> copy = v.substr(10, 10);
v.erase(10, 10);
                                                        struct WaveletTree {
   Vec<2, int> C; int s;
 v.insert(copy.mutable_begin(), copy);
                                                          // sigma = highest value + 1
                                                         WaveletTree(vector<int>& a. int sigma) :
 // print elements of rope
for (auto it : v) cout << it << "";
                                                            s(sigma), C(sigma*2, 0) {
                                                          build(a.begin(), a.end(), 0, s-1, 1);
                                                         void build(iter b. iter e. int L. int U. int
Segment Tree
                                                          u) {
if (L == U) return
//max(a,b), min(a,b), a+b, a*b, qcd(a,b), a*b
struct SegmentTree {
                                                           int M = (L+U)/2;
 typedef int T;
                                                           C[u].reserve(e-b+1); C[u].push back(0);
 static constexpr T UNIT = INT_MIN;
                                                          for (auto it = b; it != e; ++it)
  C[u].push_back(C[u].back() + (*it<=M));</pre>
 T f(T a, T b) {
 if (a == UNIT) return b;
if (b == UNIT) return a;
                                                           auto p = stable_partition(b, e, [=](int
                                                            i){return i<=M;});
  return max(a,b);
                                                          build(b, p, L, M, u*2);
 int n; vector<T> s;
SegmentTree(int n, T def=UNIT) : s(2*n, def),
                                                          build(p, e, M+1, U, u*2+1);
                                                          // number of occurrences of x in [0,i)
\rightarrow n(n) {}
                                                         int rank(int x, int i) {
   int L = 0, U = s-1, u = 1, M, r;
   while (L != U) {
 SegmentTree(vector<T> arr)

    SegmentTree(arr.size()) {

 for (int i=0:i<arr.size():i++)
                                                           M = (L+U)/2;
r = C[u][i]; u*=2;

    update(i,arr[i]);

                                                           if (x <= M) i = r, U = M;
else i -= r, L = M+1, ++u;
 void update(int pos, T val) {
  for (s[pos += n] = val; pos /= 2;)
   s[pos] = f(s[pos * 2], s[pos*2+1]);
                                                          return i:
                                                          ^{\prime\prime} number of occurences of x in [l,r)
 T query(int b, int e) { // query [b, e)
                                                         int count(int x, int 1, int r) {
  return rank(x, r) - rank(x, 1);
  Tra = UNIT, rb = UNIT;
  for (b+=n, e+=n; b<=); b/=2, e/=2) {
    if (b % 2) ra = f(ra, s[b++]);
    if (e % 2) rb = f(s[--e], rb);
                                                         // kth smallest in [l, r)
int kth(int k, int l, int r) const {
int L = 0, U = s-1, u = 1, M, ri, rj;
  return f(ra, rb):
                                                          while (L != U) {
   M = (L+U)/2;
 T get(int p) { return query(p, p+1); }
                                                           ri = C[u][1]; rj = C[u][r]; u*=2;
                                                           if (k \le rj-ri)^{n}l = ri, r = rj, U = M;
Sparse Table
                                                           else k -= řj-rí, l -= ŕi, r -= ŕj,
template < class T> struct SparseTable {
                                                           L = M+1, ++u;
 vector<vector<T>> m;
                                                           return U:
 SparseTable(vector<T> arr) {
  m.push back(arr);
  for (int k = 1: (1<<(k)) <= size(arr): k++)
                                                         // # elements between [x,y] in [l, r)
                                                         mutable int L, U;
  m.push_back(vector<T>(size(arr)-(1<(k)+1));
                                                         int range(int x, int y, int 1, int r) const {
  for (int i = 0; i < size(arr)-(1<<k)+1; i
                                                          if (y < x \text{ or } r <= 1) return 0;
                                                          L = x; U = y;
 [k][i] = min(m[k-1][i],
                                                          return range(1, r, 0, s-1, 1);
\rightarrow m[k-1][i+(1<<(k-1))]:
}
// min of range [l,r]
                                                         int range(int 1, int r, int x, int y, int u)
                                                         → const {
                                                          if (y < L or U < x) return 0;
if (L <= x and y <= U) return r-l;
T query(int 1, int r) {
 int k = _-lg(r-l+1);
                                                          int M = (x+y)/2, ri = C[u][1], rj = C[u][r];
  return \min(m[k][1], m[k][r-(1<< k)+1]);
                                                          return range(ri, rj, x, M, u*2) + range(1-ri
}
};
                                                            r-rj, M+1, y, u*2+1);
                                                          ^{\prime}// # elements <= x in [l, r]
                                                         int lte(int x, int l, int r) {
  return range(INT_MIN, x, l, r);
typedef trie<string, null_type,

→ trie_string_access_traits<>,

 pat_trie_tag, trie_prefix_search_node_update>

→ trie_type;

int main() {
                                                             Strings
 // generate trie
 trie_type trie;
                                                        Aho Corasick
 for (int i = 0; i < 20; i++)
                                                           range of alphabet for automata to consider
 trie.insert(to string(i)); // true if new,
                                                           MAXC = 26, OFFC = 'a' if only lowercase
\hookrightarrow false if old
```

```
|const int MAXC = 256:
const int OFFC = 0:
struct aho_corasick {
  set<pair<int, int>> out;
  int fail; vector<int> go;
  state() : fail(-1), go(MAXC, -1) {}
 vector<state> s;
  int id = 0;
 aho_corasick(string* arr, int size) : s(1) {
  for (int i = 0; i < size; i++) {
   for (int c : arr[i]) {
   if (s[cur].go[c-OFFC] == -1) {
      s[cur].go[c-OFFC] = s.size();
      s.push back(state());
     cur = s[cur].go[c-OFFC];
   s[cur].out.insert({arr[i].size(), id++});
  for (int c = 0; c < MAXC; c++)
if (s[0].go[c] == -1)
    s[0].go[c] = 0;
  queue int> sq;
for (int c = 0; c < MAXC; c++) {
   if (s[0].so[c] != 0) {
      ...s[s[0].so[c]].fail = 0;</pre>
    sq.push(s[0].go[c]);
  while (sq.size()) {
   int e = sq.front(); sq.pop();
   for (int c = 0; c < MAXC; c++) {
   if (s[e].go[c] != -1) {
      int failure = s[e].fail;
while (s[failure].go[c] == -1)
        failure = s[failure].fail;
      failure = s[failure].go[c];
      s[s[e].go[c]].fail = failure;
      for (auto length : s[failure].out)
s[s[e].go[c]].out.insert(length);
      sq.push(s[e].go[c]);
 // list of {start pos, pattern id}
  vector<pair<int, int>> search(string text)
  vector<pair<int, int>> toret;
  int cur = 0;
  for (int i = 0; i < text.size(); i++) {
  while (s[cur].go[text[i]-OFFC] == -1)
    cur = s[cur].fail;
cur = s[cur].go[text[i]-OFFC];
    if (s[cur].out.size())
    for (auto end : s[cur].out)
. toret.push_back({i - end.first + 1,
     end.second):
  return toret:
Boyer Moore
struct defint { int i = -1; };
vector<int> boyermoore(string txt, string pat)
 vector<int> toret; unordered_map<char, defint>string lcp(string* arr, int n, bool sorted =
 → badchar:
 int m = pat.size(), n = txt.size();
 for (int i = 0; i < m; i++) badchar[pat[i]].i
 \rightarrow = i; int s = 0:
 while (s \leq n - m) {
  int j = m - 1:
```

while  $(j \ge 0 \&\& pat[j] == txt[s + j]) j--;$ 

.if (j < 0) {

```
..toret.push_back(s);
   s += (s + m < n) ? m - badchar[txt[s +
   m]].i : 1:
 .} else
   s += \max(1, i - badchar[txt[s + i]].i):
 return toret:
English Conversion
const string ones[] = {"", "one", "two",
    "three", "four", "five", "six", "seven", "eight", "nine"};
const string teens[] ={"ten", "eleven",
    "twelve", "thirteen", "fourteen",
"fifteen", "sixteen", "seventeen",
"eighteen", "nineteen";
const string tens[] = {"twenty", "thirty",
    "forty", "fifty", "sixty", "seventy",
    "eighty", "ninety"};
const string mags[] = {"thousand", "million",
     "billion", "trillion", "quadrillion", "quintillion", "sextillion",
    "septillion"};
string convert(int num, int carry) {
 if (num < 0) return "negative " +

    convert(-num, 0):

     (num < 10) return ones[num];
(num < 20) return teens[num % 10]
 if (num < 100) return tens[(num / 10) - 2] +
     (num\%10==0?"":"") + ones[num\%10]:
 if (num < 1000) return ones[num / 100]
     (num/100==0?"":" ") + "hundred" + (num%100==0?"":" ") + convert(num % 100,
    0);
 return convert(num / 1000, carry + 1) + " " +
     mags[carry] + " " + convert(num % 1000.
    0):
string convert(int num) {
 return (num == 0) ? "zero" : convert(num, 0);
Knuth Morris Pratt
vector<int> kmp(string txt, string pat) {
   vector<int> toret;
 int m = txt.length(), n = pat.length();
 int next[n + 1];
 for (int i = 0; i < n + 1; i++)
  next[i] = 0;
 for (int i = 1; i < n; i++) {
  int j = next[i + 1];
  while (j > 0 && pat[j] != pat[i])
   j = next[j];
  if (j > 0 | pat[j] == pat[i])
   next[i + 1] = j + 1;
 for (int i = 0, j = 0; i < m; i++) {
  if (txt[i] == pat[j]) {
  if (++j == n)
    toret.push_back(i - j + 1);
  } else if (j > 0) {
...j = next[j];
 return toret;
Longest Common Prefix (array)
 // longest common prefix of strings in array
 → false) {
idise; laise; laif (n == 0) return "";
if (!sorted) sort(arr, arr + n);
string r = ""; int v = 0;
 while (v < arr[0].length() && arr[0][v] ==

    arr[n-1][v])
    r += arr[0][v++];

 return r:
```

```
| for (auto c : s) v = (c - 'a' + 1) + v *
Longest Common Subsequence
                                                           → HASHER;
string lcs(string a, string b) {
 int m = a.length(), n = b.length();
                                                          const int MAXN = 1000001;
 int L[m+1][n+1];
 Int L[m+1] [m+1];
for (int i = 0; i <= m; i++) {
    for (int j = 0; j <= n; j++) {
        if (i == 0 || j == 0) L[i][j] = 0;
        else if (a[i-1] == b[j-1]) L[i][j] =
                                                          ull base [MAXN] = \{1\};
                                                          void genBase(int n) {
                                                           for (int i = 1; i \le n; i++)

base[i] = base[i-1] * HASHER:
 \hookrightarrow L[i-1][j-1]+1;
                                                          struct advHash {
   else L[i][j] = \max(L[i-1][j], L[i][j-1]);
                                                           ull v, 1; vector<ull> wip;
                                                           advHash(string& s): v(0) {
 }
// return L[m][n]; // length of lcs
                                                            wip = vector<ull>(s.length()+1);\
                                                             wip[0] = 0;
 string out = "";
                                                             for (int i = 0; i < s.length(); i++)
wip[i+1] = (s[i] - 'a' + 1) + wip[i] *
 int i = m - 1, j = n - 1;
 while (i \ge 0) \& j \ge 0 \{ if (a[i] = b[i]) \}
                                                              HASHER;
                                                            1 = s.length(); v = wip[1];
   out = a[i--] + out;
                                                           ull del(int pos, int len) {
   return v - wip[pos+len]*base[1-pos-len] +
  else if (L[i][j+1] > L[i+1][j]) i--;
                                                              wip[pos]*base[1-pos-len];
  .else j--;
                                                            ull substr(int pos, int len) {
 return out;
                                                            return del(pos+len, (1-pos-len)) -
                                                              wip[pos]*base[len]:
Longest Common Substring
                                                           ull replace(int pos, char c) {
// l is array of palindrome length at that
int manacher(string s, int* 1) {
                                                            return v - wip[pos+1]*base[l-pos-1] + ((c -
                                                                'a' + 1) + wip[pos] *
 int n = s.length() * 2;
                                                              HASHER) *base[1-pos-1];
 for (int i = 0, j = 0, k; i < n; i += k, j =
\rightarrow max(j-k, 0)) {
                                                            ull replace(int pos, string s) {
                                                            // can't increase total string size
 while (i \ge j \&\& i + j + 1 < n \&\& s[(i-j)/2]]
\Rightarrow == s[(i+j+1)/2]) i++;
                                                               wip[pos+s.size()]*base[1-pos-s.size()], c
 1[i] = j;
                                                               wip[pos];
 for (k = 1; i >= k \&\& j >= k \&\& l[i-k] !=
                                                            for (int i = 0; i < s.size(); i++)
    c = (s[i]-'a'+1) + c * HASHER;
    return r + c * base[l-pos-s.size()];</pre>
 \begin{array}{ll} \hookrightarrow & j^-k; k^{++}) \\ \dots & 1[i^+k] = \min(1[i^-k], j^-k); \end{array}
 return *max element(1, 1 + n);
                                                          Subsequence Count
Cyclic Rotation (Lyndon)
                                                            // "banana", "ban" >> 3 (ban, ba..n, b..an)
// simple strings = smaller than its nontrivial
                                                          ull subsequences(string body, string subs) {

→ suffixes

                                                            int m = subs.length(), n = body.length();
// lyndon factorization = simple strings
                                                           if (m > n) return 0;

ull** arr = new ull*[m+1];

for (int i = 0; i <= m; i++) arr[i] = new
\hookrightarrow factorized
// "abaaba" -> "ab", "aab", "a"
vector<string> duval(string s) {
                                                           int n = s.length();
 vector<string> lyndon;
 for (int i = 0; i < n;) {
    int j = i+1, k = i;
    for (; j < n && s[k] <= s[j]; j++)
    if (s[k] < s[j]) k = i;
                                                           for (int i = 1; i <= m; i++)

for (int j = 1; j <= n; j++)

arr[i][j] = arr[i][j-1] + ((body[j-1] ==
                                                              subs[i-1])? arr[i-1][j-1] : 0);
   else k++:
                                                           return arr[m][n]:
  for (; i <= k; i += j - k)
   lyndon.push_back(s.substr(i,j-k));
                                                          Suffix Array + LCP
 return lyndon;
                                                           struct SuffixArray {
}
// lexicographically smallest rotation
                                                           vector<int> sa, 1cp;
                                                           SuffixArray(string& s, int lim=256) {

int n = s.length() + 1, k = 0, a, b;
int minRotation(string s) {
 int n = s.length(); s += s;
                                                            vector<int> x(begin(s), end(s)+1), y(n),
 auto d = duval(s); int i = 0, a = 0;
                                                            \rightarrow ws(max(n, lim)), rank(n);
 while (a + d[i].length() < n) a +=
                                                            sa = lcp = y;
iota(begin(sa), end(sa), 0);

    d[i++].length();

 while (i && d[i] == d[i-1]) a -=
                                                             for (int j = 0, p = 0; p < n; j = max(1, j *

    d[i--].length();

                                                            \rightarrow 2), lim = p) {
return a;
                                                             p = j; iota(begin(y), end(y), n - j);
                                                             for (int i = 0; i < (n); i++)
if (sa[i] >= j)
Hashing
                                                                y[p++] = sa[i] -
#define HASHER 27
                                                              fill(begin(ws), end(ws), 0);
ull basicHash(string s) {
 ull v = 0;
                                                              for (int i = 0; i < (n); i++) ws[x[i]]++;
```

```
|...for (int i = 1; i < (lim); i++) ws[i] +=
 \rightarrow ws[i - 1];
   for (int \bar{i} = n; i--;) sa[--ws[x[v[i]]]] =
    v[i]:
   j]) ? p - 1 : p++;
  for (int i = 1; i < (n); i++) rank[sa[i]] =
 for (int i = 0, j; i < n - 1; lcp[rank[i++]]
  = k)
for (k && k--, j = sa[rank[i] - 1];
s[i + k] == s[j + k]; k++);
 // smallest cyclic shift
int cyclic() { return sa[0]; }
 // longest repeated substring
 pair<int,int> lrs() {
  int length = -1, index = -1;
  for (int i = 0; i < lcp.size(); i++) {
  if (lcp[i] > length) {
  length = lcp[i];
    index = sa[i];
  return {index,length};
 // count distinct substrings, excluding empty
 int distincts() {
  int n = sa.size() - 1, r = n - sa[0];
  for (int i = 1; i < lcp.size(); i++)
  r += (n - sa[i]) - lcp[i - 1];
 return r:
 }
// count repeated substrings. excluding empty
 int repeateds() {
  int r = 0;
for (int i = 1; i < lcp.size(); i++)
  r += \max(lcp[i] - lcp[i-1], 0);
Suffix Tree (Ukkonen's)
struct SuffixTree {
 // n = 2*len+10 or so
enum { N = 50010, ALPHA = 26 };
int toi(char c) { return c - 'a'; }
 t[N][ALPHA],1[N],r[N],p[N],s[N],v=0,q=0,m=2| (1 countries) intervalCover(double L, double R,
 string a;
 void ukkadd(int i, int c) { suff:
  if (r[v]<=q) {
  if (t[v][c]==-1) { t[v][c]=m; l[m]=i;
...p[m++]=v; v=s[v]; q=r[v]; goto suff; }
.v=t[v][c]; q=1[v];
  if (q==-1 || c==toi(a[q])) q++; else {
   l[m+1]=i; p[m+1]=m; l[m]=l[v]; r[m]=q;
   p[m]=p[v]; t[m][c]=m+1; t[m][toi(a[q])]=v;
   1[v]=q; p[v]=m; t[p[m]][toi(a[1[m]])]=m;
v=s[p[m]]; q=1[m];
   while (q < r[m]) = v = t[v][toi(a[q])];
    q+=r[v]-l[v]; }
   if (q==r[m]) s[m]=v; else s[m]=m+2;
   q=r[v]-(q-r[m]); m+=2; goto suff;
 SuffixTree(string a) : a(a) {
 fill(r,r+N,(int)(a).size());
                                                          ull* catalan = new ull[1000000];
 memset(s, 0, sizeof s);
memset(t, -1, sizeof t);
fill(t[1],t[1]+ALPHA,0);
s[0]=1;1[0]=1[1]=-1;r[0]=r[1]=p[0]=p[1]=0;
                                                         void genCatalan(int n, int mod) {
  catalan[0] = catalan[1] = 1;
  for (int i = 2; i <= n; i++) {</pre>
                                                           catalan[i] = 0;
  for(int i=0;i<a.size();i++)
```

→ ukkadd(i.toi(a[i])):

```
// Longest Common Substring between 2 strings
 // returns {length, offset from first string}
 pair<int, int> best;
int lcs(int node, int i1, int i2, int olen) {
    if (1[node] <= i1 && i1 < r[node]) return 1;
    if (1[node] <= i2 && i2 < r[node]) return 2;
 int mask=0,
 - len=node?olen+(r[node]-1[node]):0;
 for(int c=0; c<ALPHA; c++) if
  (t[node][c]!=-1)
mask |= lcs(t[node][c], i1, i2, len);
 if (mask==3)
  best=max(best, {len,r[node]-len});
 static pair<int, int> LCS(string s, string t)
 SuffixTree
 \rightarrow st(s+(char)('z'+1)+t+(char)('z'+2)):
 st.lcs(0, s.size(), s.size()+t.size()+1, 0);
return st.best;
String Utilities
void lowercase(string& s) {
transform(s.begin(), s.end(), s.begin(),
void uppercase(string& s) {
transform(s.begin(), s.end(), s.begin(),
   ::toupper);
void trim(string &s) {
s.erase(s.begin(),find_if_not(s.begin(),s
     .end(),[](int c){return
   isspace(c);}));
s.erase(find_if_not(s.rbegin(),s.rend(),[](int

    c){return isspace(c);}).base(),s.end());

vector<string> split(string& s, char token) {
    vector<string> v; stringstream ss(s);
    for (string e;getline(ss,e,token);)
        v.push_back(e);
  Greedy
Interval Cover
// L,R = interval [L,R], in = {{l,r}, index}
// does not handle case where L == R
→ vector<pair<pair<double,double,int>> in) {
    int i = 0; pair < double, int > pos = {L,-1};
   vector<int> a;
sort(begin(in), end(in));
    while (pos.first < R) {
        double cur = pos.first;
while (i < (int)in.size() &&</pre>
    in[i].first.first <= cur)</pre>
    max(pos.{in[i].first.second.in[i].second}).
        if (pos.first == cur) return {};
        a.push_back(pos.second);
    return a:
6 Math
Catalan Numbers
```

for  $(int j = i - 1; j >= 0; j--) {$ 

```
catalan[i] += (catalan[j] * catalan[i-j-1]) Discrete Logarithm
    % mod;
                                                           int discretelog(int a. int b. int m) {
   if (catalan[i] >= mod)
                                                            11 n = sqrt(m) + 1, an = 1;
    catalan[i] -= mod:
                                                           for (l1 i = 0; i < n; ++i)
an = (an * a) % m;
unordered_map<ll, 11> vals;
\mathcal{Y}' // TODO: consider binomial coefficient method
                                                            for (11 q = 0, cur = b; q <= n; q++) {
   vals[cur] = q;
                                                             cur = (cur * a) \% m:
Combinatorics (nCr, nPr)
                                                            for (ll p = 1, cur = 1; p <= n; p++) {
// can optimize by precomputing factorials, and
                                                             cur = (cur * an) % m;
 \hookrightarrow fact[n]/fact[n-r]
                                                             if (vals.count(cur)) {
ull nPr(ull n, ull r) {
                                                              int ans = n * p - vals[cur];
                                                              return ans;
 for (ull i = n-r+1: i <= n: i++)
 return v;
                                                            return -1;
ull nPr(ull n, ull r, ull m) {
 ull v = 1;
                                                           Euler Phi / Totient
 for (ull i = n-r+1: i <= n: i++)
 v = (v * i) % m;
                                                           int phi(int n) {
                                                            int r = n;
                                                           for (int i = 2; i * i <= n; i++) {
   if (n % i == 0) r -= r / i;
   while (n % i == 0) n /= i;
úll nCr(ull n. ull r) {
 long double v = 1;
 for (ull i = 1: i <= r: i++)
 v = v * (n-r+i) /i;
                                                            if (n > 1) r = r / n;
 return (ull)(v + 0.001);
                                                            return r:
// requires modulo math
                                                           #define n 100000
// can optimize by precomputing mfac and
                                                           ll phi[n+1];
                                                           void computeTotient() {
\hookrightarrow minv-mfac
                                                            for (int i=1; i<=n; i++) phi[i] = i;
ull nCr(ull n, ull r, ull m) {
 return mfac(n, m) * minv(mfac(k, m), m) % m *
                                                            for (int p=2; p<=n; p++) {
                                                             if (phi[p] == p) {
 \rightarrow minv(mfac(n-k, m), m) % m;
                                                              phi[p] = p-1;
                                                              for (int i = 2*p; i \le n; i += p) phi[i] =
Multinomials
                                                               (phi[i]/p) * (p-1);
11 multinomial(vector<int>& v) {
    ll c = 1, m = v.empty() ? 1 : v[0];
    for(int i = 1; i < v.size(); i++)</pre>
 for (int j = 0; j < v[i]; j++)
...c = c * ++m / (j+1);
                                                           Factorials
                                                           // digits in factorial
 return c;
                                                           #define kamenetsky(n) (floor((n * log10(n /
                                                            \hookrightarrow ME)) + (log10(2 * MPI * n) / 2.0)) + 1)
Chinese Remainder Theorem
                                                            // approximation of factorial
                                                           #define stirling(n) ((n == 1) ? 1 : sqrt(2 *
bool ecrt(ll* r, ll* m, int n, ll& re, ll& mo)
                                                           \hookrightarrow M PI * n) * pow(n / M E, n))
 ill x, y, d; mo = m[0]; re = r[0];
for (int i = 1; i < n; i++) {</pre>
                                                           // natural log of factorial
#define lfactorial(n) (lgamma(n+1))
 d = egcd(mo, m[i], x, y);

if ((r[i] - re) % d != 0) return false;

x = (r[i] - re) / d * x % (m[i] / d);

re += x * mo;
                                                           Prime Factorization
                                                           // do not call directly
                                                           ll pollard rho(ll n. ll s) {
  mo = mo / d * m[i];
                                                            11 x, y;
  re %= mo;
                                                            x = y = rand() \% (n - 1) + 1;
                                                            int head = 1, tail = 2;
while (true) {
 re = (re + mo) \% mo;
 return true;
                                                            while (tide) {
   x = mult(x, x, n);
   x = (x + s) % n;
   if (x == y) return n;
Count Digit Occurences
                                                             11 d = gcd(max(x - y, y - x), n);
/*count(n,d) counts the number of occurences of
                                                            if (1 < \overline{d} \&\& d < n) return d:
                                                             if (++head == tail) y = x, tail <<= 1;
 \rightarrow a digit d in the range [0,n]*/
ll digit_count(ll n, ll d) {
 11 \text{ result} = 0;
 while (n != 0) {
    result += ((n%10) == d ? 1 : 0);
                                                           // call for prime factors
                                                           void factorize(ll n. vector<ll> &divisor) {
                                                            if (n == 1) return;
  n /= 10;
                                                            if (isPrime(n)) divisor.push back(n):
 return result:
                                                            f1 count(11 n, 11 d) {
  if (n < 10) return (d > 0 && n >= d);
  if ((n % 10) != 9) return digit_count(n, d) +
                                                             while (d >= n) d = pollard_rho(n, rand() % (n int josephus(int n, int k) {
                                                             - 1) + 1);
factorize(n / d, divisor);
\hookrightarrow count(n-1, d);
```

return 10\*count(n/10, d) + (n/10) + (d > 0);

factorize(d, divisor);

```
Farey Fractions
// generate 0 <= a/b <= 1 ordered, b <= n
    farey(4) = 0/1 1/4 1/3 1/2 2/3 3/4 1/1
// length is sum of phi(i) for i = 1 to n
vector<pair<int, int>> farey(int n) {
 int h = 0, k = 1, x = 1, y = 0, r;
 vector<pair<int, int>> v;
  v.push back({h, k});
  r = (n-y)/k;
  v += r*k: x' += r*h:
 x = -x; y = -y;

while (k > 1);
 v.push_back({1, 1});
 return v;
Fast Fourier Transform
const double PI = acos(-1):
void fft(vector<cd>& a, bool invert) {
 int n = a.size();
 for (int i = 1, j = 0; i < n; i++) {
  int bit = n >>
  for (; j & bit; bit >>= 1) j ^= bit;
  .j ^= bit;
  if (i < j) swap(a[i], a[i]);
 for (int len = 2; len <= n; len <<= 1) {
   double ang = 2 * PI / len * (invert ? -1 :
  cd wlen(cos(ang), sin(ang));
  for (int i = 0; i < n; i += len) {
   .cd w(1):
   for (int j = 0; j < len / 2; j++) {
    cd u = a[i+j], v = a[i+j+len/2] * w;
    a[i+j] = u + v;
a[i+j+len/2] = u - v;
    .w *= wlen:
 if (invert)
  for (auto\& x : a)
  x /= n;
vector<int> fftmult(vector<int> const& a,
 → vector<int> const& b) {
 vector<cd> fa(a.begin(), a.end()),

    fb(b.begin(), b.end());

 int n = 1 << (32 - _builtin_clz(a.size() +

→ b.size() - 1));
fa.resize(n); fb.resize(n);
 fft(fa. false): fft(fb. false)
 for (int i = 0; i < n; i++) fa[i] *= fb[i];
 fft(fa. true):
 vector<int> toret(n);
 for (int i = 0; i < n; i++) toret[i] =

→ round(fa[i].real());
return toret;
Greatest Common Denominator
ll egcd(ll a, ll b, ll& x, ll& y) {
  if (b == 0) { x = 1; y = 0; return a; }
  ll gcd = egcd(b, a % b, x, y);
 x \stackrel{\smile}{-=} a / b * y;
 swap(x, y);
 return gcd;
Josephus Problem
 // 0-indexed, arbitrary k
 if (n == 1) return 0;
if (k == 1) return n-1;
 if (k > n) return (joséphus(n-1,k)+k)%n;
 int res = josephus(n-n/k,k)-n\%k;
```

return res + ((res<0)?n:res/(k-1));

```
// fast case if k=2, traditional josephus
int josephus(int n) {
return 2*(n-(1<<(32-builtin clz(n)-1)));
Least Common Multiple
#define lcm(a,b) ((a*b)/qcd(a,b))
Modulo Operations
#define MOD \overline{10000000007}
\#define\ madd(a,b,m)\ (a+b-((a+b-m>=0)?m:0))
#define mult(a,b,m) ((ull)a*b%m)
#define msub(a,b,m) (a-b+((a < b)?m:0))
ll mpow(ll b. ll e. ll m) {
11 x = 1:
while (e > 0) {
    if (e % 2) x = (x * b) % m;
    b = (b * b) % m;
 e /= 2;
return x % m:
ull mfac(ull n, ull m) {
ull f = 1;
for (int i = n; i > 1; i--)
 f = (f * i) \frac{\pi}{m}
 return f:
// if m is not quaranteed to be prime
ll minv(ll b, ll m) {
11 \times (0, y) = 0;
if (egcd(b, m, x, y) != 1) return -1;
return (x % m + m) % m;
11 mdiv_compmod(int a, int b, int m) {
if (\underline{gcd(b, m)} != 1) return -1;
return mult(a, minv(b, m), m);
// if m is prime (like 10^9+7)
ll mdiv_primemod (int a, int b, int m) {
return mult(a, mpow(b, m-2, m), m);
// tonelli shanks = sqrt(n) % m, m is prime
ll legendre(ll a, ll m){
if (a % m==0) return 0;
if (m == 2) return 1;
return mpow(a, (m-1)/2, m);
11 msqrt(11 n, 11 m) {
ll s = \_builtin\_ctzll(m-1), q = (m-111)>>s,
\Rightarrow z = rand()%(m-1)+1:
 if (m == 2) return 1;
 if (s == 1) return mpow(n, (m+1)/411, m);
 while (legendre(z,m)!=m-1) z = rand()\%(m-1)+1;
 11 c = mpow(z,q,m), r = mpow(n,(q+1)/2,m), t
 \rightarrow = mpow(n,q,m), M = s;
while (t != 1) {
...ll i=1, ts = (t * t) % m;
...while (ts != 1) i++, ts = (ts * ts) % m;
 lll b = c;
for (int j = 0; j < M-i-1; j++) b = (b * b) %
 r = r * b % m; c = b * b % m; t = t * c % m;
 \rightarrow M = i;
return r;
Modulo Tetration
11 tetraloop(11 a, 11 b, 11 m) {
  if(b == 0 | | a == 1) return 1;
 11 w = tetraloop(a,b-1,phi(m)), r = 1;
 for (;w;w/=2) {
 if (w&1)
  r *= a; if (r >= m) r -= (r/m-1)*m;
 \bar{a} *= a; if (a >= m) a -= (a/m-1)*m;
 return r;
int tetration(int a, int b, int m) {
```

```
if (a == 0 | | m == 1) return ((b+1)&1)%m;
                                                                                                                         |// 1 = even number prime factors, -1 = odd
  return tetraloop(a,b,m) % m;
                                                                                                                           short mu[MAXN] = \{0,1\};
                                                                                                                          void mobius(){
  for (int i = 1; i < MAXN; i++)
    if (mu[i])</pre>
Matrix
template<typename T>
                                                                                                                                for (int' j = i + i; j < MAXN; j += i)
struct Mat : public Vec<2, T> {
                                                                                                                                   mu[j] -= mu[i];
  int w, h;
  Mat(int x, int y) : Vec<2, T>(x, y), w(x),
                                                                                                                           Nimber Arithmetic
 \hookrightarrow h(y) {}
                                                                                                                         #define nimAdd(a,b) ((a)^(b))
ull nimMul(ull a, ull b, int i=6) {
   static const ull M[]={INT_MIN>>32,
  static Mat<T> identity(int n) { Mat<T> m(n,n);
          for (int i=0;i<n;i++) m[i][i] = 1; return

ighthat is in the second in 
                                                                                                                                   M[0]^{M[0]}<16), M[1]^{M[1]}<8),
  Mat<\hat{T}>\& operator+=(const Mat<T>\& m) {
   for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
    (*this)[i][j] += m[i][j];
                                                                                                                                   M[2]^{(M[2] << 4)}, M[3]^{(M[3] << 2)},
                                                                                                                                  M[4]^(M[4]<<1);
                                                                                                                               if (i-- == 0) return a&b;
int k=1<<i:</pre>
     return *this:
                                                                                                                               ull s=nimMul(a,b,i), m=M[5-i],
   Mat<T>& operator-=(const Mat<T>& m) {
                                                                                                                                    t=nimMul(((a^(a>>k))&m)|(s&~m).
    for (int i = 0; i < w; i++)
for (int j = 0; j < h; j++)
                                                                                                                                    ((b^{(b)}) \& m) | (m \& (\sim m >> 1)) << k, i):
                                                                                                                              return ((s^t)&m)<<k|((s^(t>>k))&m);
         (*this)[i][j] -= m[i][j];
     return *this;
                                                                                                                          Permutation
                                                                                                                          //c = array \ size, \ n = nth \ perm, \ return \ index
   Mat<T> operator*(const Mat<T>& m) {
                                                                                                                          vector<int> gen permutation(int c, int n) {
     Mat<T>z(w,m.h);
     for (int i = 0; i < w; i++)
                                                                                                                            vector<int> idx(c), per(c), fac(c); int i;
  ...for (int j = 0; j < h; j++)
....for (int k = 0; k < m.h; k++)
....z[i][k] += (*this)[i][j] * m[j][k];
                                                                                                                            for (i = 0; i < c; i++) idx[i] = i;
                                                                                                                            for (i = 1; i <= c; i++) fac[i-1] = n%i, n/=i;
                                                                                                                            for (i = c - 1; i >= 0; i--
per[c-i-1] = idx[fac[i]],
                                                                                                                              idx.erase(idx.begin() + fac[i]);
  Mat<T> operator+(const Mat<T>& m) { Mat<T>
                                                                                                                            return per;

→ a=*this; return a+=m; }

 Mat<T> operator-(const Mat<T>& m) { Mat<T>
                                                                                                                            // get what nth permutation of vector
                                                                                                                           int get permutation(vector<int>& v) {

    a=*this; return a-=m; }

                                                                                                                           int use = 0, i = 1, r = 0;
for (int e: v) {
   r = r * i++ + __builtin_popcount(use &
  Mat<T>& operator*=(const Mat<T>& m) { return

→ *this = (*this)*m; }

  Mat<T> power(int n) {
                                                                                                                             → -(1<<e));
    .Mat<T> a = Mat<T>::identity(w),m=*this;
                                                                                                                              use |= 1 << e:
     for (;n;n/=2,m*=m) if (n\&1) a *= m:
    return a;
                                                                                                                            return r;
                                                                                                                           Permutation (string/multiset)
Matrix Exponentiation
                                                                                                                          string freq2str(vector<int>& v) {
 // F(n) = c[0]*F(n-1) + c[1]*F(n-2) + ...
// b is the base cases of same length c
                                                                                                                            string s;
                                                                                                                           for (int i = 0; i < v.size(); i++)

for (int j = 0; j < v[i]; j++)

s += (char)(i + 'A');
11 matrix exponentiation(11 n, vector<11> c,
 → vector<11> b) {
   if (nth < b.size()) return b[nth-1];
   Mat<11> a(c.size(), c.size()); ll s = 0;
   for (int i = 0; i < c.size(); i++) a[i][0] =</pre>
                                                                                                                           // nth perm of multiset, n is 0-indexed

    c[i];

                                                                                                                          string gen_permutation(string s, ll n) {
  for (int i = 0; i < c.size() - 1; i++)
                                                                                                                            vector<int> freq(26, 0);
 \hookrightarrow a[i][i+1] = 1;
                                                                                                                            for (auto e : s) freq[e - 'A']++;
  a = a.power(nth - c.size());
                                                                                                                            for (int i = 0; i < 26; i++) if (freq[i] > 0)
  for (int i = 0; i < c.size(); i++)
s += a[i][0] * b[i];
                                                                                                                               freq[i]--; ll v = multinomial(freq);
  return s;
                                                                                                                              if (n < v) return (char)(i+'A') +
                                                                                                                                  gen_permutation(freq2str(freq), n);
Matrix Subarray Sums
                                                                                                                              freq[\overline{i}]++; n -= v;
template < class T> struct MatrixSum {
                                                                                                                            return "";
  Vec<2, T > p;
  MatrixSum(Vec<2, T>& v) {
    p = Vec<2,T>(v.size()+1, v[0].size()+1);
                                                                                                                          Miller-Rabin Primality Test
 for (int i = 0; i < v.size(); i++)
for (int j = 0; j < v[0].size(); j++)
p[i+1][j+1] = v[i][j] + p[i][j+1] +</pre>
                                                                                                                           // Miller-Rabin primality test - O(10 log^3 n)
                                                                                                                         // Miller-Rabin primality test -
bool isPrime(ull n) {
   if (n < 2) return false;
   if (n == 2) return true;
   if (n ½ == 0) return false;
   ull s = n - 1;
   while (s ½ 2 == 0) s /= 2;
   for (int i = 0; i < 10; i++) {
      ull temp = s;
    ull a = rand() ½ (n - 1) + 1;
      ull mod = mpon(a temp n);
      ull mod = mpon(a temp n);

    p[i+1][i] - p[i][i];

  f sum(int u, int l, int d, int r) {
    return p[d][r] - p[d][l] - p[u][r] + p[u][l];
Mobius Function
                                                                                                                               ull mod = mpow(a, temp, n);
const int MAXN = 10000000;
// mu[n] = 0 iff n has no square factors
                                                                                                                               while (temp!=n-1\&\&mod!=1\&\&mod!=n-1) {
```

```
. mod = mult(mod, mod, n);
    temp *= 2:
   if (mod!=n-1&&temp%2==0) return false:
 return true:
Sieve of Eratosthenes
bitset<100000001> sieve;
// generate sieve - O(n log n)
void genSieve(int n) {
    sieve[0] = sieve[1] = 1;
    for (ull i = 3; i * i < n; i += 2)
        if (!sieve[i])</pre>
   for (ull j = i * 3; j <= n; j += i * 2)
...sieve[j] = 1;
// query sieve after it's generated - O(1)
bool querySieve(int n) {
 return n == 2 || (n % 2 != 0 && !sieve[n]);
Compile-time Prime Sieve
const int MAXN = 100000;
template<int N>
struct Sieve {
  bool sieve[N];
  constexpr Sieve() : sieve() {
  sieve[0] = sieve[1] = 1;
  for (int i = 2; i * i < N; i++)
if (!sieve[i])
     for (int j = i * 2; j < N; j += i)
...sieve[j] = 1;
bool isPrime(int n) {
   static constexpr Sieve<MAXN> s;
 return !s.sieve[n]:
Simpson's / Approximate Integrals
// integrate f from a to b, k iterations
// error <= (b-a)/18.0 * M * ((b-a)/2k)^4

// where M = max(abs(f^{^*}(x))) for x in [a,b]

// "f" is a function "double func(double x)"
double Simpsons (double a. double b. int k.
 \rightarrow double (*f)(double)) {
double dx = (b-a)/(2.0*k), t = 0;
 for (int i = 0; i < k; i++)

t += ((i==0)?1:2)*(*f)(a+2*i*dx) + 4 *
 \leftrightarrow (*f)(a+(2*i+1)*dx);
return (t + (*f)(b)) * (b-a) / 6.0 / k;
Common Equations Solvers
// ax^2 + bx + c = 0, find x
vector<double> solveEq(double a, double b,
 → double c) {
vector<double> r;
  double z = \bar{b} * \bar{b} - 4 * a * c;
  if (z == 0)
  r.push_back(-b/(2*a));
  else if (z > 0) {
  r.push back((sgrt(z)-b)/(2*a)):
  r.push\_back((sqrt(z)+b)/(2*a));
 return r:
// ax^3 + bx^2 + cx + d = 0, find x
vector<double> solveEq(double a. double b.
 → double c, double d) {
vector<double> res:
 long double a1 = b/a, a2 = c/a, a3 = d/a;
 long double q = (a1*a1 - 3*a2)/9.0, sq =
 \rightarrow -2*sqrt(q);
 long double r = (2*a1*a1*a1 - 9*a1*a2 +
 \rightarrow 27*a3)/54.0;
long double z = r*r-q*q*q, theta;
 if (z <= 0) {
  theta = a\cos(r/sqrt(q*q*q));
```

```
res.push_back(sq*cos(theta/3.0) - a1/3.0);
     res.push back(sq*cos((theta+2.0*PI)/3.0) -
    \rightarrow a1/3.0):
    res.push back(sq*cos((theta+4.0*PI)/3.0) -
         a1/3.0);
  élse {
   res.push_back(pow(sqrt(z)+fabs(r), 1/3.0));
    res[0] = (res[0] + q / res[0]) *

    ((r<0)?1:-1) - a1 / 3.0;
}
return res;</pre>
 // linear diophantine equation ax + by = c,
          find x and u
 // infinite solutions of form x+k*b/g, y-k*a/g
bool solveEq(ll a, ll b, ll c, ll &x, ll &y, ll
  g = \overline{egcd(abs(a), abs(b), x, y)};
  if (c % g) return false;
 x *= c / g * ((a < 0) ? -1 : 1);

y *= c / g * ((b < 0) ? -1 : 1);

return true;
// m = # equations, n = # variables, a[m][n+1]
 \rightarrow = coefficient matrix
 // a[i][0]x + a[i][1]y + ... + a[i][n]z =
         a[i][n+1]
| \hspace{-0.6cm} \hspace{-0.6cm} \rangle \hspace{-0.6cm} 

→ equation

const double eps = 1e-7:
bool zero(double a) { return (a < eps) && (a >
 \rightarrow -eps); }
vector < double > solveEq(double **a, int m, int
 \rightarrow n) { int cur = 0:
  for (int i = 0; i < n; i++) {
    for (int j = cur; j < m; j++) {
       if (!zero(a[j][i])) {
         double num = a[sat][i] / a[cur][i];
              for (int sot = 0; sot <= n; sot++)
    a[sat][sot] -= a[cur][sot] * num;
          .}
.cur++:
           break
  for (int j = cur; j < m; j++)
    if (!zero(a[j][n])) return vector<double>();
 vector<double> ans(n,0);
for (int i = 0, sat = 0; i < n; i++)
    if (sat < m && !zero(a[sat][i]);
</pre>
  ans[i] = a[sat][n] / a[sat++][i];
return ans:
 // solve A[n][n] * x[n] = b[n] linear equation
// rank < n is multiple solutions, -1 is no

→ solutions
// `alls` is whether to find all solutions, or

 \hookrightarrow any
const double eps = 1e-12;
int solveEq(Vec<2, double>& A, Vec<1, double>&

→ b, Vec<1, double>& x, bool alls=false) {
 int n = A.size(), m = x.size(), rank = 0, br,
  vector<int> col(m); iota(begin(col), end(col),
   → 0);
 for(int i = 0; i < n; i++) {
  double v, bv = 0;</pre>
     for(int r = i; r < n; r++)
   for(int c = i; c < n; c++)
if ((v = fabs(A[r][c])) > bv)
br = r, bc = c, bv = v;
if (bv <= eps) {
      for(int i = i: i < n: i++)
```

```
. if (fabs(b[j]) > eps)
    .return -1;
 swap(A[i], A[br]);
 swap(b[i], b[br]);
 swap(col[i], col[bc]);
 for(int j = 0; j < n; j++)
swap(A[j][i], A[j][bc]);
 bv = 1.0 / A[i][i];
for(int j = (alls)?0:i+1; j < n; j++) {
  if (j != i) {
   double fac = A[j][i] * bv;
   b[j] = fac * b[i];
  for(int k = i+1; k < m; k++)
[k] -= fac*A[i][k];
 rank++;
}
if (alls) for (int i = 0; i < m; i++) x[i] =</pre>
→ -DBL MAX;
for (int i = rank; i--;) {
  bool isGood = true;
 if (alls)
 ifor (int j = rank; isGood && j < m; j++)
if (fabs(A[i][j]) > eps)
 isGood = false;
b[i] /= A[i][i]:
 if (isGood) x[col[i]] = b[i];
 if (!alls)
for(int j = 0; j < i; j++)
   b[i] = A[i][i] * b[i];
return rank:
```

## Graycode Conversions

```
ull graycode2ull(ull n) {
 ull i = 0;
for (; n; n = n >> 1) i ^= n;
 return i;
ull ull2graycode(ull n) {
  return n ^ (n >> 1);
}
```

#### Date Utilities

```
// handles -4799-01-01 to 1465001-12-31 int date2int(int y, int m, int d){
return 1461*(y+4800+(m-14)/12)/4+367*(m-2-(m-14)/12)
    -14)/12*12)/12-3*((v+4900+(m-14)/12)/100)
   /4+d-32075;
pair<int,pair<int,int>> int2date(int x){
 int n,i,j;
 n=4*x/146097;
 x-=(146097*n+3)/4;
i=(4000*(x+1))/1461001;
 x=1461*i/4-31;
 i=80*x/2447:
 return \{100*(n-49)+i+j/11, \{j+2-12*(j/11), \}
\rightarrow x-2447*j/80}};
int dayOfWeek(int y, int m, int d){ //0=sunday
static int cal[]={0,3,2,5,0,3,5,1,4,6,2,4};
 return (y+y/4-y/100+y/400+cal[m-1]+d)\%7;
```

## Unix/Epoch Time

```
// O-indexed month/time, 1-indexed day
// minimum 1970, 0, 1, 0, 0, 0
ull toEpoch(int year, int month, int day, int
→ hour, int minute, int second) {
struct tm t; time_t epoch;

t.tm_year = year - 1900; t.tm_mon = month;

t.tm_mday = day; t.tm_hour = hour;
t.tm min = minute; t.tm sec = second;
```

```
t.tm_isdst = 0; // 1 = daylights savings
epoch = mktime(&t);
return (ull)epoch;
vector<int> toDate(ull epoch) {
time_t e=epoch; struct tm t=*localtime(&e);
return {t.tm vear+1900.t.tm mon.t.tm mdav.t
   .tm_hour,t.tm_min,t.tm_sec};
int getWeekdav(ull epoch) {
time t e=epoch: struct tm t=*localtime(&e):
return t.tm_wday; // 0-6, 0 = sunday
int getDayofYear(ull epoch) {
time_t e=epoch; struct tm t=*localtime(&e);
return t.tm yday; // 0-365
const int months[] =
\rightarrow {31,28,31,30,31,30,31,30,31,30,31};
bool validDate(int year, int month, int day) {
    bool leap = !(year%(year%25?4:16));
    if (month >= 12) return false;
   return day <= months[month] + (leap &&
   month == 1);
```

#### Theorems and Formulae

Montmort Numbers count the number of derangements (permutations where no element appears in its original position) of a set of size n. !0 = 1, !1 = 0, !n = (n+1)(!(n-1))1)+!(n-2)), ! $n = n! \sum_{i=0}^{n} \frac{(-1)^{i}}{i!}$ , ! $n = \left[\frac{n!}{e}\right]$ 

In a partially ordered set, a chain is a subset of elements that are all comparable to eachother. An antichain is a subset where no two are comparable.

|Dilworth's theorem states the size of a max-| $^{\Box}$ imal antichain equals the size of a minimal chain cover of a partially ordered set S. The width of S is the maximum size of an antichain in S, which is equal to the minimum number of chains needed to cover S, or the minimum number of chains such that all elements are in at least one chain.

Rosser's Theorem states the *nth* prime Floyd Warshall number is greater than n \* ln(n) for n > 1.

Lagrange's Four Square Theorem states every natural number is the sum of the squares of four non-negative integers. This is a special case of the Fermat Polygonal Number if (m[i][k] != inf && m[k][j]!= inf) **Theorem** where every positive integer is  $a \mapsto m[i][j] = -inf$ ; sum of at most n s-gonal numbers. The nths-gonal number  $P(s,n) = (s-2)\frac{n(n-1)}{2} + n$ 

# 7 Graphs

```
struct edge {
int u,v,w;
edge (int u,int v,int w) : u(u),v(v),w(w) {}
edge (): u(0), v(0), w(0) {}
```

```
|bool operator < (const edge &e1, const edge
\rightarrow &e2) { return e1.w < e2.w: }
bool operator > (const edge &e1, const edge
\rightarrow &e2) { return e1.w > e2.w: }
struct subset { int p, rank; };
Eulerian Path
#define edge_list vector<edge>
#define adj_sets vector<set<int>>
struct EulerPathGraph {
 adj_sets graph; // actually indexes incident
 → edaes
 edge_list edges; int n; vector<int> indeg;
 EulerPathGraph(int n): n(n) {
 indeg = *(new vector<int>(n,0));
  graph = *(new adj_sets(n, set<int>()));
 void add edge(int u. int v) {
  graph[u].insert(edges.size());
  indeg[v]++;
  edges.push back(edge(u.v.0)):
 bool eulerian_path(vector<int> &circuit) {
  if(edges.size()==0) return false;
  stack<int> st;
int a[] = {-1, -1};
for(int v=0;v<n;v++) {
  if(indeg[v]!=graph[v].size()) {
    bool b = indeg[v] > graph[v].size();
    if (abs(((int)indeg[v])-((int)graph[v]
     .size())) > 1) return
    false;
if (a[b] != -1) return false;
    a[b] = v;
  int s = (a[0]!=-1 \&\& a[1]!=-1 ? a[0] :
    (a[0]=-1 & a a[1]=-1 ? edges[0].u : -1):
  if(s==-1) return false:
  while(!st.empty() || !graph[s].empty()) {
   if (graph[s].empty()) {
    circuit.push_back(s); s = st.top();
    st.pop(); }
    int w = edges[*graph[s].begin()].v;
    graph[s].erase(graph[s].begin());
    st.push(s); s = w;
 circuit.push_back(s);
 return circuit.size()-1==edges.size();
```

```
number is greater than n * ln(n) for n > 1.

Nicomachi's Theorem states 1^3 + 2^3 + \dots + \frac{1}{void} floydWarshall (Vec<2, 11>& m) {
n^3=(1+2+...+n)^2 and is equivalent to \inf_{\substack{i=1,\dots,n\\ \text{FOR}(i,n)}} \prod_{\substack{i=1,\dots,n\\ \text
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                    && m[k][j] != inf)
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                              auto newDist = max(m[i][k] + m[k][j], -inf);
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         m[i][j] = min(m[i][j], newDist);
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                       fOR(k,n) if (m[k][k] < 0) FOR(i,n) FOR(j,n)
```

# Minimum Spanning Tree

```
// returns vector of edges in the mst
// graph[i] = vector of edges incident to
→ vertex i
// places total weight of the mst in &total
// if returned vector has size != n-1, there is
vector<edge> mst(vector<vector<edge>> graph,

→ 11 &total) {
```

```
total = 0:
 priority_queue<edge, vector<edge>,

→ greater<edge>> pq;

 vector<edge> MST;
 bitset<20001> marked: // change size as needed
 marked[0] = 1;
 for (edge ep : graph[0]) pq.push(ep);
 while (MST.size()!=graph.size()-1 &&
 \rightarrow pq.size()!=0) {
  edge e = pq.top(); pq.pop();
int u = e.u, v = e.v, w = e.w;
if(marked[u] && marked[v]) continue;
  else if (marked[u]) swap(u, v);
  for(edge ep : graph[u]) pq.push(ep);
  marked[u] = 1
  MST.push_back(e);
  .total += e.w:
 return MST;
Union Find
int uf find(subset* s, int i) {
  if (s[i].p != i) s[i].p = uf_find(s, s[i].p);
 return s[i].p;
void uf_union(subset* s, int x, int y) {
 int xp = uf_find(s, x), yp = uf_find(s, y);
 if (s[xp].rank > s[yp].rank) s[yp].p = xp;
 else if (s[xp].rank < s[yp].rank) s[xp].p =
 else { s[yp].p = xp; s[xp].rank++; }
2D Grid Shortcut
#define inbound(x.n) (0<=x&&x<n)
#define fordir(x, y, n, m) for(auto[dx, dy]:dir) if
\hookrightarrow (inbound(x+dx,n)&\mathcal{G}inbound(y+dy,m))
const pair<int,int> dir[] =
\leftrightarrow \{\{1,0\},\{0,1\},\{-1,0\},\{0,-1\}\};
    2D Geometry
#define point complex<double>
#define EPS 0.0000001
#define sq(a) ((a)*(a))
#define c\overline{b}(a) ((a)*(a)*(a))
double dot(point a, point b) { return

→ real(conj(a)*b); }

double cross(point a, point b) { return
\rightarrow imag(conj(a)*b); }
struct line { point a, b; };
struct circle { point c; double r; };
struct segment { point a, b; };
struct triangle { point a, b, c; };
struct rectangle { point tl, br; };
struct convex_polygon {
 vector<point> points;
 convex_polygon(vector<point> points) :
   points(points) {}
 convex_polygon(triangle a) {
  points.push_back(a.a); points.push_back(a.b);
    points.push_back(a.c);
 convex polygon(rectangle a) {
  points.push_back(a.tl);
    points.push back({real(a.tl).
    imag(a.br)});
  points.push_back(a.br);
    points.push_back({real(a.br),
    imag(a.tl)}):
struct polygon {
 vector <point > points;
 polygon(vector<point> points) :

→ points(points) {}
```

```
double area = 0:
 polygon(triangle a) {
                                                   point c(0, 0);
  points.push back(a.a); points.push back(a.b);
                                                   for (int i = n - 1, j = 0; j < n; i = j++) {
    points.push back(a.c);
                                                    double v = cross(a.points[i], a.points[j])
 polygon(rectangle a) {
                                                    area += v:
  points.push_back(a.tl);
                                                    c += (a.points[i] + a.points[j]) * (v / 3);
    points.push_back({real(a.tl),
    imag(a.br)});
                                                   c /= area;
  points.push back(a.br);
                                                   return {area, c};
    points.push_back({real(a.br),
    imag(a.tl)}):
                                                  Intersection
                                                  // -1 coincide, 0 parallel, 1 intersection
 polygon(convex_polygon a) {
                                                  int intersection(line a, line b, point& p) {
  if (abs(cross(a.b - a.a, b.b - b.a)) > EPS) {
  for (point v : a.points)
   points.push_back(v);
                                                   p = cross(b.a - a.a, b.b - a.b) / cross(a.b)
                                                   \rightarrow -a.a, b.b -b.a) * (b - a) + a;
};
// triangle methods
                                                   return 1;
double area_heron(double a, double b, double
                                                   if (abs(cross(a.b - a.a, a.b - b.a)) > EPS)
\stackrel{\hookrightarrow}{} c) { if (a < b) swap(a, b);
                                                   → return 0;
                                                   return -1;
 if (a < c) swap(a, c);
 if (b < c) swap(b, c);
                                                   // area of intersection
 if (a > b + c) return -1;
                                                  double intersection(circle a, circle b) {
 return sqrt((a+b+c)*(c-a+b)*(c+a-b)*(a+b-c)
                                                   double d = abs(a.c - b.c);
                                                      (d <= b.r - a.r) return area(a)
(d <= a.r - b.r) return area(b)
    /16.0);
                                                   if (d \ge a.r + b.r) return 0:
// segment methods
                                                   double alpha = acos((sq(a.r) + sq(d) -
double lengthsq(segment a) { return
                                                     sq(b.r)) / (2 * a.r * d));
    sq(real(a.a) - real(a.b)) + sq(imag(a.a)
                                                   double beta = acos((sq(b.r) + sq(d) - sq(a.r))
    imag(a.b)); }
                                                     /(2 * b.r * d));
double length(segment a) { return
                                                   return sq(a.r) * (alpha - 0.5 * sin(2 *

    sqrt(lengthsq(a)); }

                                                      alpha) + sq(b.r) * (beta - 0.5 * sin(2 *
// circle methods
                                                      beta));
double circumference(circle a) { return 2 * a.r \Rightarrow
                                                   // -1 outside, 0 inside, 1 tangent, 2
double area(circle a) { return sq(a.r) * M PI;
                                                  intersection intersection (circle a, circle b,
→ }
// rectangle methods
                                                     vector<point>& inter) {
double width(rectangle a) { return
                                                   double d2 = norm(b.c - a.c), rS = a.r + b.r,
  abs(real(a.br) - real(a.tl)); }
                                                   \rightarrow rD = a.r - b.r;
if (d2 > sq(rS)) return -1;
double height (rectangle a) { return

→ abs(imag(a.br) - real(a.tl)); }

                                                   if (d2 < sq(rD)) return 0;
double diagonal(rectangle a) { return
                                                   double ca = 0.5 * (1 + rS * rD / d2):

    sqrt(sq(width(a)) + sq(height(a))); }

                                                   point z = point(ca, sqrt(sq(a.r) / d2 -
double area(rectangle a) { return width(a) *
                                                   \rightarrow sq(ca)):
 → height(a); }
                                                   inter.push back(a.c + (b.c - a.c) * z):
                                                   if (abs(imag(z)) > EPS) inter.push back(a.c +
double perimeter(rectangle a) { return 2 *
   (width(a) + height(a)); }
                                                      (b.c - a.c) * conj(z));
// check if `a` fit's inside `b
                                                   return inter.size();
 // swap equalities to exclude tight fits
                                                   // points of intersection
bool doesFitInside(rectangle a, rectangle b) {
                                                  vector<point> intersection(line a, circle c) {
 int x = width(a), w = width(b), y = height(a)
                                                   vector <point > inter:
\rightarrow h = height(b);
                                                   c.c -= a.a;
a.b -= a.a;
 if (x > y) swap(x, y);
 if (w > h) swap(w, h);
                                                   point m = a.b * real(c.c / a.b);
                                                   double d2 = norm(m - c.c);
 if (w < x) return false;
                                                   if (d2 > sq(c.r)) return 0;
 if (y <= h) return true;
                                                   double l = sqrt((sq(c.r) - d2) / norm(a.b));
 double a=sq(y)-sq(x), b=x*h-y*w, c=x*w-y*h;
                                                   inter.push back(a.a + m + 1 * a.b);
 return sq(a) \le sq(b) + sq(c);
}
// polygon methods
                                                   if (abs(1) > EPS) inter.push_back(a.a + m - 1
                                                   \rightarrow * a.b);
// negative area = CCW, positive = CW
                                                   return inter;
double area(polygon a) {
                                                   // area of intersection
 double area = 0.0; int n = a.points.size();
                                                  double intersection(rectangle a, rectangle b) {
 for (int i = 0, j = 1; i < n; i++, j = (j +
                                                   double x1 = max(real(a.tl), real(b.tl)), y1 =
max(imag(a.tl), imag(b.tl));
                                                   double x2 = min(real(a.br), real(b.br)), y2 =
 min(imag(a.br), imag(b.br));
 return area / 2.0:
                                                   return (x2 \le x1 \mid y2 \le y1) ? 0 :
}
// get both unsigned area and centroid
                                                      (x2-x1)*(y2-y1);
pair<double, point> area_centroid(polygon a) { |}
 int n = a.points.size();
                                                  Convex Hull
```

```
bool cmp(point a, point b) {
 if (abs(real(a) - real(b)) > EPS) return
    real(a) < real(b);
 if (abs(imag(a) - imag(b)) > EPS) return
\rightarrow imag(a) \tilde{\langle} imag(b);
return false:
convex_polygon convexhull(polygon a) {
 sort(a.points.begin(), a.points.end(), cmp);
 vector<point> lower, upper;
 for (int i = 0; i < a.points.size(); i++) {
  while (lower.size() >= 2 &&
    cross(lower.back() - lower[lower.size()
    2], a.points[i] - lower.back()) < EPS)
   lower.pop_back();
  while (upper.size() >= 2 &&
    cross(upper.back() - upper[upper.size()
   2], a.points[i] - upper.back()) > -EPS)
  upper.pop_back();
lower.push_back(a.points[i]);
  upper.push_back(a.points[i]);
 lower.insert(lower.end(), upper.rbegin() + 1,
   upper.rend());
 return convex_polygon(lower);
    3D Geometry
```

```
struct point3d {
 double x, y, z;
 point3d operator+(point3d a) const { return
 \rightarrow {x+a.x, y+a.y, z+a.z}; }
 point3d operator*(double a) const { return
 \rightarrow {x*a, v*a, z*a}: }
 point3d operator-() const { return {-x, -v,
 \rightarrow -z}; }
 point3d operator-(point3d a) const { return
 \rightarrow *this + -a; }
 point3d operator/(double a) const { return
→ *this * (1/a); }
double norm() { return x*x + y*y + z*z; }
 double abs() { return sqrt(norm()); }
 point3d normalize() { return *this /

    this->abs(): }

double dot(point3d a, point3d b) { return
\rightarrow a.x*b.x + a.v*b.v + a.z*b.z: }
point3d cross(point3d a, point3d b) { return
    \{a.y*b.z - a.z*b.y, a.z*b.x - a.x*b.z,
    a.x*b.v - a.v*b.x}: }
struct line3d { point3d a, b; };
struct plane { double a, b, c, d; } // a*x +
\Rightarrow b*y + c*z + d = 0
struct sphere { point3d c; double r; };
\#define\ sq(a)\ ((a)*(a))
#define c\bar{b}(a) ((a)*(a)*(a))
double surface(circle a) { return 4 * sq(a.r)
double volume(circle a) { return 4.0/3.0 *
 \hookrightarrow cb(a.r) * M_PI; }
```

# 10 Optimization

```
Snoob
  // SameNumberOfOneBits, next permutation
 int snoob(int a) {
  int b = a & -a, c = a + b;
  return c | ((a ^ c) >> 2) / b;
  '/ example usage
// example usage
int main() {
   char 11[] = {'1', '2', '3', '4', 'char 12[] = {'a', 'b', 'c', 'd'};
   int d1 = 5, d2 = 4;
   // prints 12345abcd, 1234a5bcd, ...
  int min = (1 << d1) -1, max = min << d2;
```

```
for (int i = min; i <= max; i = snoob(i)) {
  int p1 = 0, p2 = 0, v = i;</pre>
 while (p1 < d1 || p2 < d2) {
  cout << ((v & 1) ? 11[p1++] : 12[p2++]);
  v /= 2;
cout << '\n':
```

# $\mathbf{Powers}$ bool isPowerOf2(ll a) { return a > 0 && !(a & a-1); bool isPowerOf3(11 a) { return a>0&&!(12157665459056928801u11%a); bool isPower(ll a, ll b) { double x = log(a) / log(b);return abs(x-round(x)) < 0.00000000001;

## 11 Python

### Recursion Limit Removal (Basic)

import svs svs.setrecursionlimit(10\*\*6)

#### Recursion Limit Removal (Advanced)

```
# @bootstrap over recursive function
# replace 'return' with 'yield'
# for when sys method does not work
from types import GeneratorType
def bootstrap(f, stack=[]):
 def wrappedfunc(*args, **kwargs):
 if stack:
.return f(*args. **kwargs)
  else:

to = f(*args, **kwargs)

while True:

if type(to) is GeneratorType:
     stack append(to)
      to = next(to)
     else:
      stack.pop()
      if not stack:
to = stack[-1].send(to)
return to
return wrappedfunc
```

#### Python 3 Compatibility

```
import sys
from __future__ import division, print_function
if sys.version_info[0] < 3:</pre>
from _builtin_ import xrange as range from future_builtins import ascii, filter,

→ hex, map, oct, zip
```

# 12 Additional

#### Judge Speed

```
kattis: 0.50s
    codeforces: 0.421s
 // atcoder: 0.455s
#include <bits/stdc++.h>
using namespace std;
int v = 1e9/2, p = 1;
int main() {
  for (int i = 1: i <= v: i++) p *= i:</pre>
```

#### Judge Pre-Contest Checks

int128 and float128 support? does extra or missing whitespace cause WA? -documentation up to date? -printer usage available and functional?

```
// each case tests a different fail condition
 // try them before contests to see error codes
struct g { int arr[1000000]; g(){}};
vector<g> a;
 // O=WA 1=TLE 2=MLE 3=OLE 4=SIGABRT 5=SIGFPE
if (n == 4) assert(0);
if (n == 5) 0 / 0;
if (n == 6) *(int*)(0) = 0;
 return n + judge(n + 1);
 GCC Builtin Docs
 // 128-bit integer
__int128 a;
unsigned __int128 b;
// 128-bit float
// minor improvements over long double __float128 c; // log2 floor
__lg(n);
__lg(n);
// number of 1 bits
// can add ll like popcountll for long longs
__builtin_popcount(n);
// number of trailing zeroes
_builtin_ctz(n);
// number of leading zeroes
_builtin_ctz(n);
// number of leading zeroes
_builtin_clz(n);
// 1-indexed_least_significant 1 bit
_builtin_ffs(n).
__builtin_ffs(n);
// parity of number
 __builtin_parity(n);
Limits
                          \begin{array}{c|c} \pm 2147483647 & \pm 2^{31} - 1 | 10^9 \\ 4294967295 | & 2^{32}_{32} - 1 | 10^9_{10} \end{array}
 int
 uint
          \pm 9223372036854775807 | \pm 2^{63} - 1|10^{18}
                                                   \frac{1}{2}^{64} - \frac{1}{10}^{19}
           18446744073709551615
 ull
i128 |\pm 170141183460469231...|\pm 2^{127} - 1|10^{38}
                                                 \frac{1}{2}^{128} - \frac{1}{1}^{10}^{38}
 u128 340282366920938463...
 Complexity classes input size (per second):
 O(n^n) or O(n!)
                                                             n < 10
 O(2^n)
                                                             n \leq 30
 O(n^3)
                                                         n \le 1000
 O(n^2)
                                                       n \le 30000
                                                           n \le 10^6
 O(n\sqrt{n})
 O(n \log n)
                                                           n \le 10^7
 O(n)
                                                           n < 10^9
```