# Fundamentos de Programação

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#### Summary

#### Recursive functions

- How recursion works
- The rules for termination
- Recursion vs. repetition
- How to write recursive functions

#### Examples

- Operations on a list
- Towers of Hanoi
- A sorting algorithm (kind of quicksort)

### A crazy idea

What does this function do?

```
sumsq([1, 2, 3]) #-> 14
```

- Does it work on an empty list?
- Can you write it with a generator expression? (Homework!)

```
def sumsq(lst):
    s = 0
    for x in lst:
        s += x**2
    return s
```

- Check out this weird version!
  - It squares first element;
  - Calls sumsq on the rest;
  - · And adds.
  - For empty lst, just return zero.

```
def sumsq2(lst):
    s = 0
    if len(lst) > 0:
        sq0 = lst[0]**2
        s = sq0 + sumsq(lst[1:])
    return s
```

- It is *equivalent* to sumsq in every case, but still calls the original sumsq. Not very useful.
- But if sumsq2 <=> sumsq, why not call itself?

#### Recursive functions

This is what would result.

```
def sumsqR(lst):
    s = 0
    if len(lst) > 0:
        sq0 = lst[0]**2
        s = sq0 + sumsqR(lst[1:])
    return s
```

- This is a recursive function: a function that calls itself.
- Notice that there is no loop instruction, but code gets executed several times, anyway.
- How does it work?

#### How recursion works

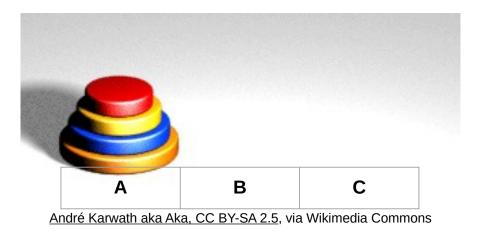
What happens when we call sumsqR([1, 2, 3])?

```
sumsqR([1, 2, 3])
| sumsqR([2, 3])
| sumsqR([3])
| sumsqR([])
| L>0
| L>9 (= 3**2 + 0)
| ->13 (= 2**2 + 9)
| ->14 (= 1**2 + 13)
```

- Notice that at one point, there are 4 frames in memory.
  - 4 variables named lst, 4 named s, 3 named sq0, but all distinct!
- Each frame stores the *local context* of a single function call.
- The frames are stored in the program stack.

### Example: Towers of Hanoi

- The Towers of Hanoi puzzle (Édouard Lucas, 1883).
- Move tower from A to C, using B temporarily.
  - Move only one disk at a time;
  - No disk may be put on top of a smaller disk.



Now solve it in 4 lines of code!



### Example: quicksort

- The quicksort algorithm (by C.A.R. Hoare) goes like this:
  - 1. Pick one of the values in the list (e.g. the first) and store in T.
  - 2. Put values smaller than T into a list L1, the others into a list L2.
  - 3. Sort L1 and L2 (using same algorithm, by the way)
  - 4. Result is L1 + [T] + L2.
- Of course, there's a few more details (the base case).

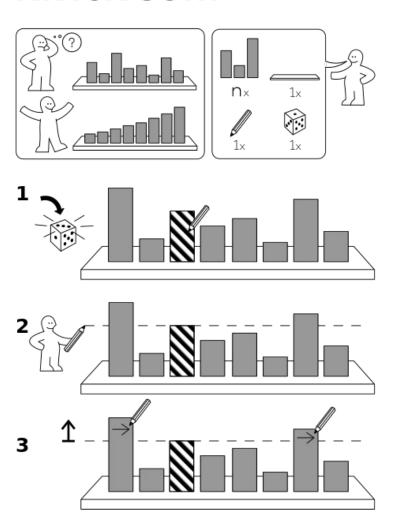
```
def qsorted(lst):
    if len(lst)<=1:  # no need to sort
        return lst[:]  # just return a copy
    T = lst[0]
    L1 = [x for x in lst[1:] if x<T]
    L2 = [x for x in lst[1:] if x>=T]
    return qsorted(L1) + [T] + qsorted(L2)
```

This is simple to understand and quite efficient!

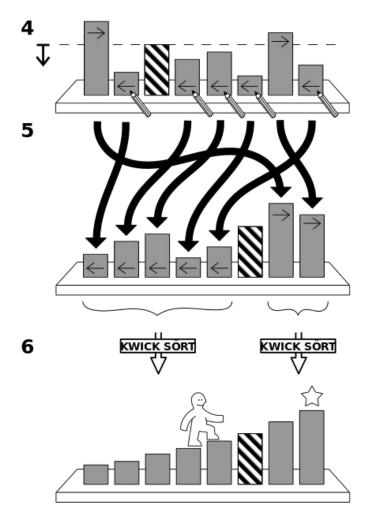
The actual quicksort modifies the list in-place, and is slightly harder to write.

# Example: quicksort

#### **KWICK SÖRT**



idea-instructions.com/quick-sort/ v1.0, CC by-nc-sa 4.0



# The rules for recursion termination



To guarantee that a recursive function **terminates**, it must respect some **rules**:

- 1. There must be cases that can be solved *without* recursive calls. These are called the **base cases**.
  - In sumsqR, the base case is len(lst)==0. In that case, return 0.
- 2. In the other cases, the *context passed* to recursive calls **must always differ** from the *context received*.\*
  - In sumsqR, the argument lst[1:] != lst (always)
- 3. In successive recursive calls, the context must **converge** towards the base cases.
  - In sumsqR, the lst is shortened on each call, until it's empty.

<sup>\*</sup> The *context* is the set of arguments (and global values) that have an impact on the base case / recursive case selection.

### Recursion vs repetition

- Any problem that can be solved by repetition may be solved by recursion, and vice-versa.
- For certain complex problems, recursive solutions are usually more concise and easier to understand.
- Recursive implementations may incur extra time and memory cost because of function calls and stack usage.
- If the problem has a simple iterative solution, that is usually the most efficient, too.

### Writing recursive functions

- To develop a recursive function to solve some problem, follow these guidelines:
  - 1. First, **define** the **arguments** you need, what they **mean**, and the **result** you **expect**, as *rigorously* as possible.
  - 2. Now, **assume** the function will work. Describe how the solution to a problem can be obtained by **modifying** the solutions to **smaller** versions of the problem. This will be the <u>recursive part</u> of the algorithm.
  - 3. Finally, **determine** the **base cases**: which conditions have a trivial solution? This will be the <u>non-recursive part</u> of the algorithm. (Hint: base cases are usually *outside the domain* of the recursive part.)
- While in step 2, you may realize that you need *extra* arguments. Just add them and go back to step 1.