

## Assignment 3: Optimization of a City Transportation Network (MST)

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### 1. Objective

The goal of this assignment is to optimize a city transportation network using **Minimum Spanning Tree (MST)** algorithms — **Prim's** and **Kruskal's** — and to compare their performance in terms of total cost, number of operations, and execution time.

The transportation network is represented as a **weighted undirected graph**, where:

- vertices represent city districts;
- edges represent possible roads;
- edge weights represent road construction costs.

### 2. Algorithms Overview

#### Prim's Algorithm

Prim's algorithm grows the MST starting from one vertex and repeatedly adds the cheapest edge connecting the tree to a new vertex.

#### Complexity:

- Using a priority queue:  **$O(E \log V)$**
- Efficient for **dense graphs**

#### Kruskal's Algorithm

Kruskal's algorithm sorts all edges by weight and adds them to the MST in order, skipping edges that would form a cycle (using Union-Find).

#### Complexity:

- **$O(E \log E)$  or  $O(E \log V)$**
- Efficient for **sparse graphs**

### 3. Implementation Summary

Both algorithms were implemented in **Java** with the following classes:

| Package    | File                    | Purpose                     |
|------------|-------------------------|-----------------------------|
| model      | Graph.java, Edge.java   | Custom graph data structure |
| algorithms | Prim.java, Kruskal.java | MST algorithms              |
| util       | JsonIO.java             | JSON input/output           |
| main       | Main.java               | Driver program              |

Graphs are loaded from input.json and results are written to output.json.

#### 4. Input Datasets

Several graph datasets were used to evaluate correctness and performance.

| Dataset | Vertices (V) | Edges (E) | Description               |
|---------|--------------|-----------|---------------------------|
| Small   | 5            | 7         | Simple test for debugging |
| Medium  | 10           | 20        | Moderate graph            |
| Large   | 25           | 70        | Performance test          |

Example JSON for a small graph:

```
{
  "vertices": 5,
  "edges": [
    {"src": 0, "dest": 1, "weight": 2},
    {"src": 0, "dest": 3, "weight": 6},
    {"src": 1, "dest": 2, "weight": 3},
    {"src": 1, "dest": 3, "weight": 8},
    {"src": 1, "dest": 4, "weight": 5},
    {"src": 2, "dest": 4, "weight": 7},
    {"src": 3, "dest": 4, "weight": 9}
  ]
}
```

#### 5. Results

### 5.1 Example Output (Small Graph)

| Metric              | Prim Kruskal |      |
|---------------------|--------------|------|
| Total MST Cost      | 16.0         | 16.0 |
| Operations Count    | 9            | 11   |
| Execution Time (ms) | 1            | 0    |
| MST Edges           | 4            | 4    |

**Both algorithms produced the same MST total cost**, proving correctness.

### 5.2 Medium Graph (10 vertices, 20 edges)

| Metric              | Prim Kruskal |      |
|---------------------|--------------|------|
| Total MST Cost      | 42.0         | 42.0 |
| Operations Count    | 47           | 62   |
| Execution Time (ms) | 2            | 1    |

Observation: Prim’s algorithm used fewer operations due to better priority queue performance on a moderately dense graph.

### 5.3 Large Graph (25 vertices, 70 edges)

| Metric              | Prim Kruskal |       |
|---------------------|--------------|-------|
| Total MST Cost      | 118.0        | 118.0 |
| Operations Count    | 176          | 194   |
| Execution Time (ms) | 4            | 3     |

Observation: Kruskal’s algorithm was slightly faster due to efficient sorting and union-find operations on larger sparse graphs.

## 6. Analysis and Discussion

| Criterion                 | Prim’s Algorithm | Kruskal’s Algorithm |
|---------------------------|------------------|---------------------|
| Implementation complexity | Moderate         | Easy                |

| Criterion | Prim's Algorithm | Kruskal's Algorithm |
|-----------|------------------|---------------------|
|-----------|------------------|---------------------|

|                            |        |                |
|----------------------------|--------|----------------|
| Efficiency on dense graphs | Better | Slightly worse |
|----------------------------|--------|----------------|

|                             |                |        |
|-----------------------------|----------------|--------|
| Efficiency on sparse graphs | Slightly worse | Better |
|-----------------------------|----------------|--------|

|                           |                |            |
|---------------------------|----------------|------------|
| Data structure dependency | Priority Queue | Union-Find |
|---------------------------|----------------|------------|

|                      |                 |              |
|----------------------|-----------------|--------------|
| MST Cost correctness | Same as Kruskal | Same as Prim |
|----------------------|-----------------|--------------|

**Both algorithms produce the same total MST cost**, but their performance depends on graph structure:

- Prim's performs better when many edges exist between vertices (dense graphs);
- Kruskal's performs better when the graph is sparse.

## 7. Conclusions

1. Both Prim's and Kruskal's algorithms correctly compute the Minimum Spanning Tree (MST).
2. Execution times and operation counts vary depending on the graph density.
3. **Prim's algorithm** is more efficient for dense graphs due to its use of a priority queue.
4. **Kruskal's algorithm** is simpler and faster for sparse graphs because it mainly relies on sorting.
5. The total cost of MST is always identical, confirming correctness.