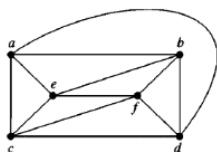
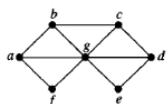


Q.1)

For the given graphs, their respective chromatic numbers are

Max Marks: 1



A

3 and 3

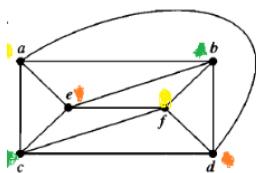
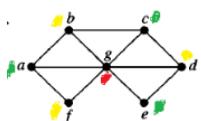
Correct Option

Solution: (A)

Answer: A

Explanation:

We can color both of them with 3 colors so that no two same color is adjacent. but how can we sure that it is not less than 3. As both the graph consist triangle it can't be color with less than 3. So we at least require 3 colors to color them.



B

6 and 5

C

3 and 4

D

4 and 4

Q.2)

 $(\sim q \wedge (p \rightarrow q)) \rightarrow \sim p$  is

Max Marks: 1



A

Tautology

Correct Option

Solution: (A)

Answer: A

Explanation:

$$(\sim q \wedge (p \rightarrow q)) \rightarrow \sim p$$

We know that  $p \rightarrow q \equiv \sim p \vee q$ 

$$(\sim q \wedge (\sim p \vee q)) \rightarrow \sim p$$

Convert  $\wedge \equiv , \vee \equiv +$ 

$$\text{Now, } q' \cdot (p' + q) \rightarrow p'$$

$$(q' \cdot p' + q' \cdot q) \rightarrow p'$$

$$(p' \cdot q' + 0) \rightarrow p'$$

$$p' \cdot q' \rightarrow p'$$

$$p' \cdot p' \rightarrow p'$$

$$(\mu \cdot q) \wedge \mu$$

$$p + q + p' = p + p' + q = 1 + q = 1$$

Now, we can write  $p \vee \sim p \equiv \sim q \vee p \equiv T$   
So, this is a Tautology.

D Contradiction

C Contingency

D Unsatisfiable

Q.3)

Maximum no of edges in a triangle-free, simple planar graph with 10 vertices is \_\_\_\_

Max Marks: 1

Correct Answer

Solution: (16)

Answer: 16

**Explanation:**

Minimum degree in any triangle free graph

$$\delta = 4$$

by using Euler's formula

$$e - n + (k+1) = r \quad \{k = \# \text{f connected components}\}$$

$$e - 10 + 2 = r$$

$$\Rightarrow e - 8 = r$$

As, min degree of region is 4

$$4r \leq 2e$$

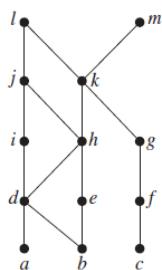
$$\Rightarrow 4(e-8) \leq 2e$$

$$\Rightarrow 4e - 32 \leq 2e$$

$$\Rightarrow e \leq 16$$

Q.4)

Which of the given statements are TRUE for the partial order represented by the following Hasse diagram:



S1 : The maximal elements are l and m.

S2: There is no greatest element.

S3: The GLB of {f, g, h} does not exist.

A S1 and S2

B S2 and S3

C S1 and S3

D S1, S2 and S3

Correct Option

Solution: (D)

Answer: D

**Explanation:**

S1 : l and m are the maximal elements as they have no successor

S2 : There is no greatest elements as the 2 maximal elements are incomparable.  
S3 : Greatest lower bound of  $\{f, g, h\}$  does not exist as they have no lower bounds in common

Q.5)

Given the following conditional statements, which is/are tautology:

Max Marks: 1

S1 :  $[(p \rightarrow q) \wedge (q \rightarrow r)] \rightarrow (p \rightarrow r)$

S2 :  $[(p \vee q) \wedge (p \rightarrow r) \wedge (q \rightarrow r)] \rightarrow r$

A

Only S1

B

Only S2

C

Both S1 and S2

Correct Option

Solution: (c)

Answer: C

Explanation:

S1 :  $[(p \rightarrow q) \wedge (q \rightarrow r)] \rightarrow (p \rightarrow r)$

If p, then q AND if q then r would mean if p, then r which is exactly what the RHS is. So, the implication is true. (But RHS implies LHS is not true here as RHS does not say that p should imply q or q should imply r)

S2 :  $[(p \vee q) \wedge (p \rightarrow r) \wedge (q \rightarrow r)] \rightarrow r$

Either p or q is true; if p then r; if q then r. So, either way we have r and RHS is r only. So, again this implication is true.

So, both are tautologies.

D

None of the above

Q.6)

Which of the following is True

Max Marks: 1

- S1. Every satisfiable is a contingency
- S2. Every satisfiable is a tautology
- S3. Every invalid is a contradiction

A

S1 and S2

B

S2 and S3

C

S1 and S3

D

None of the above

Correct Option

Solution: (D)

Explanation:

Tautology: Compound proposition which is always True.

Contradiction: Compound proposition which is always False.

Contingency: Compound proposition which is sometimes True and sometimes False.

Satisfiability: A Compound proposition is satisfiable if there is at least one True in its truth table.

So, Tautology is always satisfiable.

Unsatisfiability: Not even a single True result in its truth table.

So, Contradiction is always unsatisfiable.

Valid: A Compound preposition is valid when it is a tautology.

Invalid: A Compound preposition is invalid when it is either contradiction or contingency.

Important Results:

Tautology is always satisfiable but satisfiable is not always the tautology.

Invalid does not mean that a compound preposition is always False. If a compound preposition is sometimes True and False, then also is said to be invalid.

Summary:

Tautology	Contradiction	Contingency
always TRUE	always FALSE	Sometimes TRUE or FALSE
Satisfiable	Unsatisfiable	Satisfiable
Valid	Invalid	Invalid

Therefore, S1, S2 and S3 are all False. Thus, Option D is correct.

Q.7)

Max Marks: 1

Number of relations  $S$  over set  $\{0,1,2,3\}$  such that  $(x, y) \in S \Rightarrow x = y$  is \_\_\_\_\_

Correct Answer

**Solution:** (16)

**Answer:** 16

**Explanation:**

Only four pairs that satisfy the condition are:  $(0,0), (1,1), (2,2), (3,3)$ .

Now the relation can include or leave each of these 4 pairs. Hence  $2^4$  ways.

For all other 12 possible pairs, the relation must leave them. Hence 1 way.

So no of relations =  $2^4 \times 1 = 16$ .

Note that the empty relation is also valid. In that case the condition is vacuously valid.

Q.8)

What is the largest no of maximal independent set of complete bipartite graph  $K(4,2)$ ?

Max Marks: 1

A 2

B 3

C 4

Correct Option

**Solution:** (c)

**Answer:** C

**Explanation:**

We know that for bipartite complete graph  $K_{m,n}$ , we can partition the vertex set into 2 sets such that one contains  $m$  and other contains  $n$  elements(vertices) and within each of this sets there does not exist any edge between any pair of vertices. So it is nothing but independent set.

So for bipartite complete graph  $K_{m,n}$ , we have 2 independent sets.

So here  $\max(4,2) = 4$ . Hence 4 should be the size of the largest maximal independent set.

D 5

Q.9)

Consider a set  $S = \{2, 3, 4, \dots, 23, 24\}$  and  $R$  is a relation on  $S$  such that  $aRb$  if  $a$  divides  $b$ , then find the number of minimal elements in its hasse diagram

Max Marks: 1

A 9

Correct Option

**Solution:** (A)

**Answer:** A

**Explanation:**

Minimal Element : Let  $(A, R)$  be a poset. Then "a" in  $A$  is a minimal element if there does not exist an element  $b$  in  $A$  such that  $bRa$ .

So, for the given poset  $(S, /)$

Minimal elements :  $\{2, 3, 5, 7, 11, 13, 17, 19, 23\} = 9$  elements

B 10

C 12

D 16

Q.10)

If  $G$  is a wheel graph with 6 vertices, then matching number of

Max Marks: 1

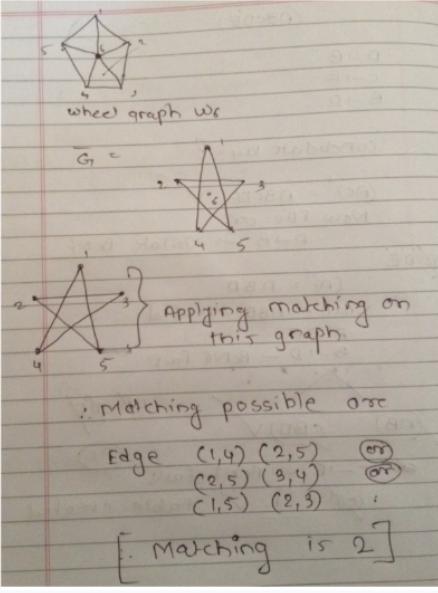
$\bar{G} = \underline{\hspace{2cm}}$ .

Correct Answer

**Solution:** (2)

**Answer:** 2

**Explanation:**



Q.11)

Which of the following statement(s) is/are True?

Max Marks: 2

- S1 : Set of integers on addition operation is Monoid but not group.  
 S2 : Set of integers on subtraction operation is Monoid but not group.  
 S3 : Set of integers on multiplication operation is group but not Abelian group.

<input checked="" type="radio"/> A	S1 only
<input type="radio"/> B	S2 only
<input type="radio"/> C	S1 and S3
<input type="radio"/> D	None of the above

**Correct Option**

**Solution:** (D)  
**Answer:** D  
**Explanation:**  
 S1: It is False  
 Addition of Integer is closed  
 Addition of integer is associative  $((1+2)+3)=(1+(2+3))$   
 Identity property also satisfied  $0+1=1$ , here identity element is 0  
 Inverse of 1 is (-1),  $1+(-1)=0$ , So, inverse also satisfied  
 So, Integer addition will be a group  
 S2: It is also false  
 Associativity also not satisfied  
 $(7 \cdot 2) \cdot 3$  not equal with  $7 \cdot (2 \cdot 3)$   
 i.e.  $(5 \cdot 3)$  not equal with  $(7+1)$   
 but a monoid we have to satisfy associativity and identity property  
 S3: Also False  
 As, Integer multiplication does not satisfy inverse property  
 inverse of 2 will be  $1/2$ , which is not an integer  
 So, Integer multiplication cannot form a group

Q.12)

Max Marks: 2

Consider an undirected graph with  $n$  vertices, vertex 1 has degree 1, while each vertex 2, 3, ...,  $n-1$  has degree 4. The degree of vertex  $n$  is unknown. Which of the following statements must be TRUE?

<input type="radio"/> A	Vertex $n$ has degree 1.
<input type="radio"/> B	Graph is connected.
<input type="radio"/> C	There is a path from vertex 1 to vertex $n$ .

**Correct Option**

**Solution:** (C)  
**Answer:** C  
**Explanation:**  
 A) false.  
 By handshaking lemma you can argue that degree of  $n$  must be odd but it doesn't mean it will be 1 always.  
 For example 1, 4, 4, 4, 4, 3 is a valid degree sequence.  
 B) false.  
 Suppose the graph has 2 components as degree sequence  $G1 = 1, 4, 4, 4, 4, 3$  and  $G2 = 4, 4, 4, 4, 4$ . Which means, it is not always true that graph is connected.  
 C) True  
 Vertex 1 and vertex  $n$  should be a vertex of same component of graph otherwise it will not follow handshaking lemma. Graph is undirected and vertex 1 and  $n$  are connected.

vertex 1 and vertex n should be a vertex of same component of graph. And each vertex must have at least one edge connecting it to other vertices.

So, there must be a path from 1 to n.

D) false

Not every spanning tree will contain the edge joining 1 to n.

D

Spanning tree will include the edge connecting vertex 1 and n.

Q.13)

Max Marks: 2

Consider the following two statements:

S<sub>1</sub> : All clear explanations are satisfactory.

S<sub>2</sub> : Some excuses are unsatisfactory.

Which one of the following statement(s) follows from S<sub>1</sub> and S<sub>2</sub> as per inference rules of logic?

A

Every excuses are not clear explanations

B

Some excuses are clear explanations

C

Some excuses are not clear explanations

Correct Option

Solution: (c)

Answer: C

Explanation:

Let's define some predicates :

CE(x) : x is a Clear Explanation.

Ex(x) : x is an Excuse.

S(x) : x is Satisfactory.

Now, given statements can be written as follows :

S<sub>1</sub> =  $\forall x(CE(x) \rightarrow S(x))$  (Interpretation : For all X, if X is Clear Explanation then X is Satisfactory).

S<sub>2</sub> =  $\exists x(Ex(x) \wedge \neg S(x))$  (Interpretation : There is some X such that X is an Excuse and X is Not Satisfactory)

So, We are given S<sub>1</sub> and S<sub>2</sub>. And considering these two statements True, we need to show which of the Options is always True.

Using Statement S<sub>2</sub>, you can say that, in the domain, there is at least one X for which Ex(x) is True and S(x) is False. So, Let this X=a and now from S<sub>2</sub> we have Ex(a)=T and S(a)=F

Let's see S<sub>1</sub> for this element a in the domain.

We will find that,

CE(a) → S(a) must be True. But we know that S(a)=F, So, We have CE(a)=F (Because CE(a) → S(a) is True)

So, From S<sub>1</sub> and S<sub>2</sub>, we have Ex(a)=T and CE(a)=F

So, Option C is Correct because it says  $\exists x(Ex(x) \wedge \neg CE(x))$ , which is Always True or we can say that Option C can never be False when S<sub>1</sub> and S<sub>2</sub> are True. Because, for X=a,  $\exists x(Ex(x) \wedge \neg CE(x))$  becomes True.

So, Option C is implied from S<sub>1</sub> and S<sub>2</sub>.

Or You can say that  $(\forall x(CE(x) \rightarrow S(x))) \wedge (\exists x(Ex(x) \wedge \neg S(x))) \rightarrow (\exists x(Ex(x) \wedge \neg CE(x)))$  is a Tautology.

Now, the above method is a Formal way of proving answer. But intuitively, just by looking at the Question, We could say that C is Correct. But, in this direct way, we might make mistakes for complicated questions, So, Knowing the above formal approach gives you assurity of the answer.

D

Some explanations are clear excuses.

Q.14)

Max Marks: 2

From the given statements, which among the following statement(s) is/are TRUE ?

S<sub>1</sub> : ( $\{0, 1, 2, \dots, (m-1)\}$ ,  $+_m$ ) where  $+_m$  stands for "addition-modulo-m"

S<sub>2</sub> : ( $\{0, 1, 2, \dots, m\}$ ,  $+_m$ ) where  $+_m$  stands for "addition-modulo-m".

**A** ONLY S1 is a group.

Correct Option

**Solution:** (A)

**Answer: A**

**Explanation:**

The answer will be A. Only S1 is the group.

Because in the second case, Identity element does not exist.

In the first case identity element is 0. That means for all  $a \in S1$ ,  $a +_m 0 = a$ . because  $a < m$

In the second case, 0 can not be the identity element. For example: for one of the member  $m$  of the set we have  $m +_m 0 = 0$ , It should come  $m$ . That's why S2 is not a group.

However, both S1 and S2 are Semigroup as they satisfy closure and associativity property.

**B**

ONLY S2 is a group.

**C**

BOTH S1 AND S2 are groups.

**D**

NEITHER S1 NOR S2 is a group.

**Q.15)**

Max Marks: 2

Let  $G = \{a_1, a_2, \dots, a_{12}\}$  be an Abelian group of order 12. Then the order of the element  $(\prod_{i=1}^{12} a_i)$  is

**A**

1

**B**

2

Correct Option

**Solution:** (B)

**Answer: B**

**Explanation:**

Every element is multiplied with its inverse. The identity element is its own inverse. That leaves us with 11 elements other than identity. So, one element (except the identity element) is its own inverse. The other 10 elements get cancelled out as they are multiplied with their respective inverse elements. So, the order of the element  $\prod_{i=1}^{12} a_i$  is 2.

**C**

8

**D**

12

**Q.16)**

Max Marks: 2

$A = \{1, 2, 3, 4, 5, 6\}$

Given that set A is a group with respect to multiplication mod 7.

G is also a cyclic group, what are the generators of G?

**A**

4 and 6

**B**

2 and 4

**C**

3 and 6

**D**

3 and 5

Correct Option

**Solution:** (D)

**Answer: D**

**Explanation:**

From the given question, we know G is cyclic.

Once we figured out that G is cyclic with  $O(G) = 6$

Then  $\phi(6)$  will give no of generators in G

$$\Rightarrow \phi(6) = \phi(2^1 * 3^1)$$

$$\Rightarrow \phi(6) = \phi(2^1) * \phi(3^1)$$

$$\Rightarrow \phi(6) = [2^1 - 2^0] * [3^1 - 3^0]$$

$$\Rightarrow \phi(6) = 2$$

⇒ And these two gen are 3 and 5

Q.17)

Max Marks: 2

Which of the following statement(s) is/are true?

S1: If  $R_1$  and  $R_2$  are the equivalence relation on  $X$  then  $R_1 \cap R_2^{-1}$  (Inverse of  $R_2$ ) is also an equivalence relation.

S2: If  $R$  is reflexive and transitive relation on  $X$  then  $R \cap R^{-1}$  is an equivalence relation.

A Only S1

B Only S2

C Both S1 and S2

Correct Option

Solution: (c)

Answer: C

Explanation:

Inverse of a relation: If  $(a,b) \in R$ , then  $(b,a) \in R^{-1}$  for all  $a,b \in X$ .

Statement 1:

We know ' $\cap$ ' of two equivalence relations is equivalence. So, we just need to prove  $R_2^{-1}$  is an equivalence relation.

1. If  $(a,a) \in R_2$ , then  $(a,a) \in R_2^{-1}$  also for all  $a \in X$ . So  $R_2^{-1}$  is reflexive.

2. If symmetric pairs  $(a,b) & (b,a) \in R_2$ , then corresponding elements  $(b,a) & (a,b)$  also  $\in R_2^{-1}$ . So  $R_2^{-1}$  is symmetric.

3. If  $(a,b), (b,c) & (a,c) \in R_2$ , then  $(b,a), (c,b) & (c,a) \in R_2^{-1}$  satisfying transitivity property,  $((c,b) & (b,a) \rightarrow (c,a))$ . So  $R_2^{-1}$  is transitive.

Thus  $R_2^{-1}$  is an equivalence relation.

So  $R_1 \cap R_2^{-1}$  is also equivalence.

Statement 2:

1. Since  $R$  is reflexive, so  $R^{-1}$  will also contain reflexive elements. So  $R \cap R^{-1}$  is reflexive.

2. Since  $R$  is not symmetric, then only some but not every symmetric pairs  $(a,b) & (b,a)$  are present in  $R$ .  $R$  contains only  $(a,b)$  but not  $(b,a)$  for some  $a,b \in X$ , then  $R^{-1}$  will contain only  $(b,a)$  but not  $(a,b)$ . So  $R \cap R^{-1}$  contains neither  $(a,b)$  nor  $(b,a)$  thus ensuring no asymmetric elements.  $R$  contains both  $(a,b) & (b,a)$  for some  $a,b \in X$ , then  $R^{-1}$  will also contain both  $(b,a) & (a,b)$ . So  $R \cap R^{-1}$  contains both  $(a,b) & (b,a)$  thus ensuring both symmetric pairs are present. So  $R \cap R^{-1}$  will either contain both symmetric pairs or none at all making it symmetric.

3. Since  $R$  is transitive, so if  $(a,b), (b,c) & (a,c) \in R$  for some  $a,b,c \in X$ , then  $(b,a), (c,b) & (c,a) \in R^{-1}$ . This means  $R \cap R^{-1}$  doesn't contain any of  $(a,b), (b,c) & (a,c)$  or  $(b,a), (c,b) & (c,a)$  elements, so no need to check transitivity property here.

Finally  $R \cap R^{-1}$  is equivalence relation as it is reflexive, symmetric and transitive!

D None of the above.

Q.18)

Max Marks: 2

The proposition  $(P \Rightarrow Q) \wedge (Q \Rightarrow P)$  is a

A Tautology

B Contradiction

C Contingency

Correct Option

Solution: (c)

Answer: C

Explanation:

P	Q	$P \Rightarrow Q$	$Q \Rightarrow P$	$(P \Rightarrow Q) \wedge (Q \Rightarrow P)$
T	T	T	T	T
T	F	F	T	F
F	T	T	F	F
F	F	T	T	F

The given proposition is bidirectional, which is neither tautology nor contradiction.

A proposition which is neither tautology nor contradiction is contingency.

Hence, option C is correct.

D

Absurdity

Q.19)

Let  $f: A \rightarrow B$  and  $g: B \rightarrow C$  be functions.

Which of the following statement(s) is/are true?

Max Marks: 2



- a: If  $g \circ f$  is injective, then  $f$  is injective.
- b: If  $g \circ f$  is surjective, then  $g$  is surjective.

A

a only

B

b only

C

Both a and b

Correct Option

Solution: (C)

Part a)

Suppose  $g \circ f$  is injective. Let  $x, y \in A$  such that

$$\begin{aligned}f(x) = f(y) &\Rightarrow g(f(x)) = g(f(y)) \\&\Rightarrow (g \circ f)(x) = (g \circ f)(y) \\&\Rightarrow x = y \quad (g \circ f \text{ is injective})\end{aligned}$$

Therefore,  $f$  is injective.

Part b)

Suppose  $g \circ f$  is surjective. Then, for any  $z \in C$  there exists  $x \in A$  such that

$$z = (g \circ f)(x) \iff z = g(f(x))$$

Therefore, if we define  $y = f(x)$ , for any  $z \in C$  there exists  $y \in B$  such that  $g(y) = z$ , that is,  $g$  is surjective.

D

None of the above

Q.20)

Which of the following is not a boolean algebra?

Max Marks: 2



- I:  $[D_{110}; /]$
- II:  $[D_{91}; /]$
- III:  $[D_{45}; /]$
- IV:  $[D_{64}; /]$

A

I only

B

II only

C

II and III

D

III and IV

Correct Option

Solution: (D)

Answer: D

Explanation: Let  $n = p_1 p_2 \dots p_k$ , where the  $p_i$  are distinct primes. Then  $D_n$  is a Boolean algebra.

A)  $\{D_{110}, 1\} = \{\underline{2 \times 5 \times 11}\} \checkmark$

• All are distinct prime

B)  $\{D_{91}, 1\} = \{\underline{7, 13}\} \sim$

• All are distinct prime

C)  $\{D_{45}, 1\} = \{\underline{3 \times 3 \times 5}\} \times$

• All are Not Distinct

D)  $\{D_{64}, 1\} = \{\underline{2 \times 2 \times 2 \times 2 \times 2}\} \times$

• All are Not Distinct prime.

close