

C YOUR CODE – INTERMEDIATE CODE GENERATOR

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In partial satisfaction of the requirements for the degree of

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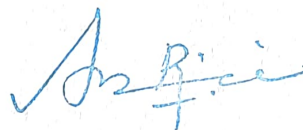
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ABSTRACT

The project is a compiler design project based on the generation of quadruples. The aim of the project is to develop a table which represent the three-address code in a quadruple format, in which the input is a general sentence. The project consists of a code which is meant to generate the quadruple by first tokenizing the given sentence into separate parts and then we apply prefix and postfix notation. The intermediate code generator then takes the stream of tokens and generates intermediate code in the form of three-address code. The quadruple generator is the main focus of the project. It takes the three-address code as input and generates quadruples for each line of code. The quadruples are used as an intermediate representation of the program code during the compilation process. Overall, the project is a comprehensive implementation of a compiler design project based on the generation of quadruples. It provides a practical example of how quadruples can be used as an intermediate representation of the program code during the compilation process.

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CHAPTER 1

1.1 INTRODUCTION

The development of a website that allows developers to input their code and see the output of what the compiler would produce is an excellent tool for developers. As software development becomes increasingly complex, the ability to test code and identify potential errors before deploying it is critical. This tool provides developers with a convenient way to quickly and efficiently test their code without the need to install and configure a complete development environment.

Consider a scenario where a developer is working on a new feature for an application. The feature requires the use of a complex algorithm that the developer has not worked with before. The developer writes the code and attempts to run it, but immediately receives an error message from the compiler. Without the ability to test the code in isolation, the developer must spend significant time trying to isolate the error and determine how to fix it.

However, with the use of the tool we have developed, the developer can simply input the code into the website and select the appropriate part of the compiler to test it. In this case, the developer could select the lexical analyzer to identify and correct any syntax errors. The tool quickly identifies the error and provides a clear message to the developer on how to fix it. With the error corrected, the developer can then run the code again and see the output of what the compiler would produce. This tool saves the developer a significant amount of time and effort in the debugging process.

Additionally, this tool is also useful for teaching programming. Beginners can use the tool to learn the basics of programming without the need to install and configure a development environment. The tool provides an easy-to-use interface that allows users to write code, see the output of what the compiler would produce, and make changes as needed. The feedback provided by the tool is immediate and clear, which helps beginners quickly identify and correct errors.

1.2 PROBLEM STATEMENT

As a software developer, we might have encountered situations where you want to test your code against different compilers, or we might have to compile your code on different platforms. But it can be time-consuming and challenging to set up different compilers and platforms to compile your code manually. This is where the tool you have developed comes in handy. The tool allows developers to input their code and see the output of what the compiler would produce without worrying about installing and configuring different compilers and platforms. With your tool, developers can quickly test their code against different compilers and platforms without leaving their development environment. This tool can also be beneficial for developers who are just starting with programming, as they can see how their code is being compiled and understand the different stages of the compilation process, such as lexical analysis and intermediate code generation. Moreover, the tool can help developers to identify and fix errors in their code during the development phase, making it easier for them to deliver bug-free code. Hence, the tool can save developers a lot of time and effort by providing them with a convenient and efficient way to test their code against different compilers and platforms, and help them deliver high-quality code.

1.3 OBJECTIVES

- Developed a website that allows developers to input their code and see the output of what the compiler would produce.
- Implemented a lexical analyzer, intermediate code generation, and quadruple triple in Python.
- Used ReactJS as frontend and NodeJS with ExpressJS as backend.
- When the server API is called, based on the part of the compiler selected from drop down, that part's shell script is executed and stored into a buffer.
- Utilized Linux file directory to store the buffer and read it for the response to the frontend REST API and display the output on the website.

1.4 HARDWARE REQUIREMENTS

- A server or cloud infrastructure to host the website and the backend logic.
- Sufficient RAM and CPU power to handle multiple user requests simultaneously.
- Sufficient disk space to store the code files and other resources.

1.5 SOFTWARE REQUIREMENTS

- Operating system: Linux or compatible OS.
- Python interpreter installed on the server.
- NodeJS and Express JS installed for the backend.
- A web browser to access the website.

CHAPTER 2

ANATOMY OF A COMPILER

2.1 LEXICAL ANALYSIS

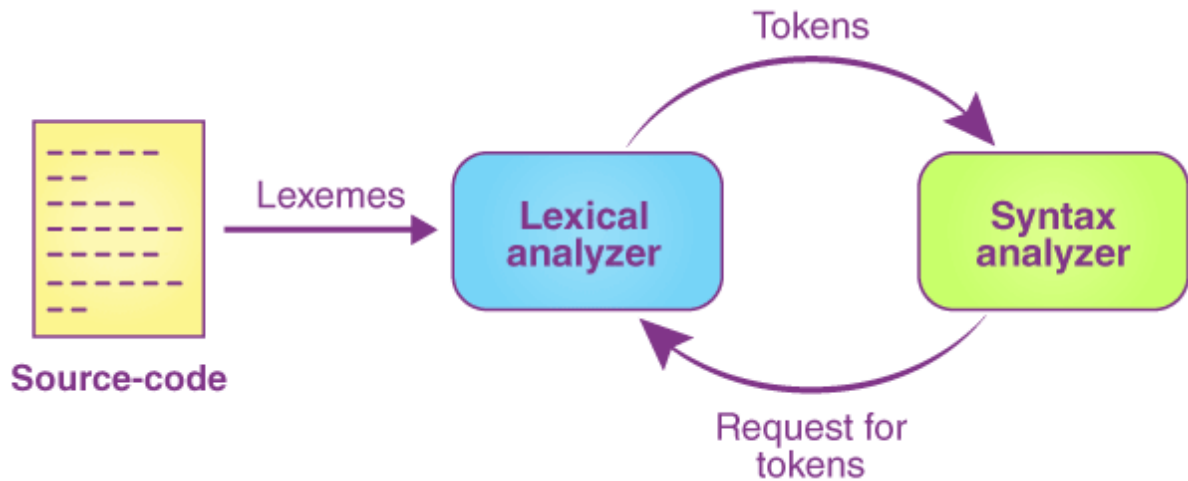


FIG 2.1 Working of Lexical Analyzer

Lexical analysis is the first phase of the compilation process in which the input source code is scanned and analysed to generate a stream of tokens that are subsequently used by the compiler in the later stages of the compilation process.

The process of lexical analysis is also known as scanning, and the component of the compiler that performs this task is called a lexer or a tokenizer. The primary goal of lexical analysis is to identify the individual lexical units (tokens) of the source code, such as keywords, identifiers, literals, operators, and punctuation marks, and produce a stream of tokens that can be processed by the compiler's parser.

The process of lexical analysis involves several steps. The first step is to read the source code character by character and group them into lexemes, which are the smallest meaningful units of the programming language. Next, the lexer applies a set of rules or regular expressions to identify the lexemes and classify them into different token types.

During this process, the lexer also discards any comments or white spaces that are not significant to the language's syntax. For example, the lexer would ignore any spaces, tabs, or newlines in the source code and focus only on the meaningful tokens that make up the language's syntax.

2.2 INTERMEDIATE CODE GENERATION

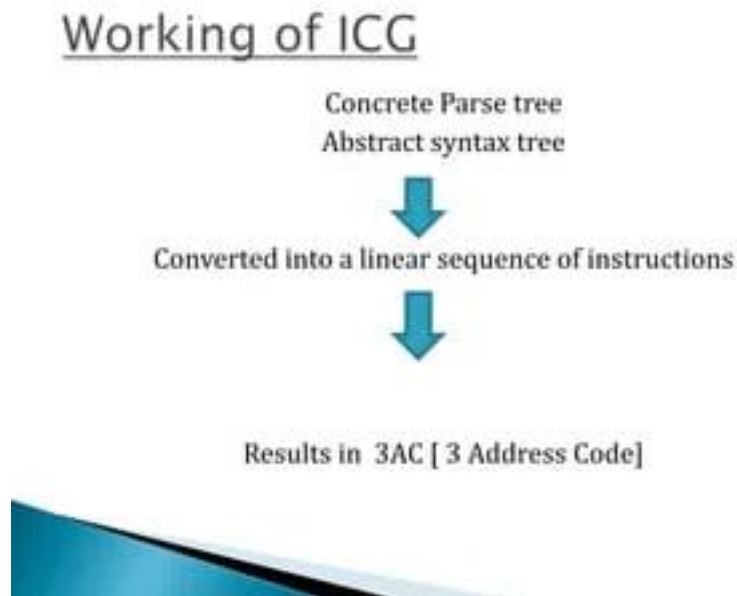


FIG 2.2 Flowchart of Intermediate Code Generator

Intermediate code generation is a crucial phase in the process of compiling a programming language. Its purpose is to translate the source code into an intermediate language representation that is closer to machine code and can be easily optimized and translated into executable code.

During intermediate code generation, the compiler analyses the source code and creates a simplified version of it, typically in the form of a set of instructions or statements in a lower-level language.

The intermediate code is usually designed to be independent of the hardware and operating system on which the code will eventually run, making it easier to port the code to different platforms. It also allows the compiler to perform optimizations that can improve the performance of the generated code, such as dead code elimination and constant folding.

Some examples of intermediate code representations are Three-Address Code (TAC), Quadruple Code, and Intermediate Representation (IR). These representations are usually simpler and more concise than the original source code, making it easier for the compiler to analyse and optimize them. Once the intermediate code is generated, the compiler can then proceed to the next phase, which is code optimization and code generation.

2.3 QUADRUPL

Operator
Source 1
Source 2
Destination

FIG 2.3 Members of Quadruple

In compiler design, a quadruple is a data structure used to represent an executable instruction or operation in an intermediate representation of a program. It contains four fields, namely the operator or instruction to be performed, the addresses of the operands, and the result. The quadruple is named as such because it has four fields. The operator field specifies the operation or instruction to be performed, such as "add", "subtract", "multiply", "divide", "assign", and so on. The operands and result fields contain memory addresses for the variables or memory locations that contain the values involved in the operation. The quadruple is a simple and uniform representation that can be used during code optimization and code generation phases of the compiler. It can be easily translated into assembly language or machine code.

Quadruples are a common intermediate representation in compilers because they provide a compact and uniform representation of executable instructions. They can be used to represent a wide range of operations, including arithmetic and logical operations, memory accesses, and control flow instructions. Quadruples are often generated during the semantic analysis phase of a compiler, where the compiler checks the syntactic and semantic correctness of the input program.

2.4 TRIPLE

Operator
Source 1
Source 2

FIG 2.4 Members of Triple

In compiler design, triples are a common data structure used to represent three-address code, which is an intermediate representation of a program. Three-address code consists of instructions that have at most three operands and one result. The use of three-address code simplifies the analysis and optimization of the program. The triple data structure represents one line of the three-address code and consists of three fields: the operator, the first operand, and the second operand. The third field is used to store the result of the operation.

The use of triples is beneficial in optimizing the code because it can simplify the identification of redundant computations and remove them. It can also help to improve code efficiency by allowing the compiler to rearrange code instructions for better performance. For example, the triple data structure can be used to implement common subexpression elimination, where the compiler identifies identical expressions that are computed multiple times and replaces them with a single computation. This optimization can lead to a significant reduction in the number of instructions executed and, therefore, an improvement in program performance.

Triples are often generated during the code generation phase of the compiler, where the compiler translates the intermediate representation of the program into machine code or assembly language. The use of triples can make this translation more straightforward and efficient by simplifying the representation of the program.

CHAPTER 3

ARCHITECTURE AND COMPONENTS

3.1 ARCHITECTURE DIAGRAM

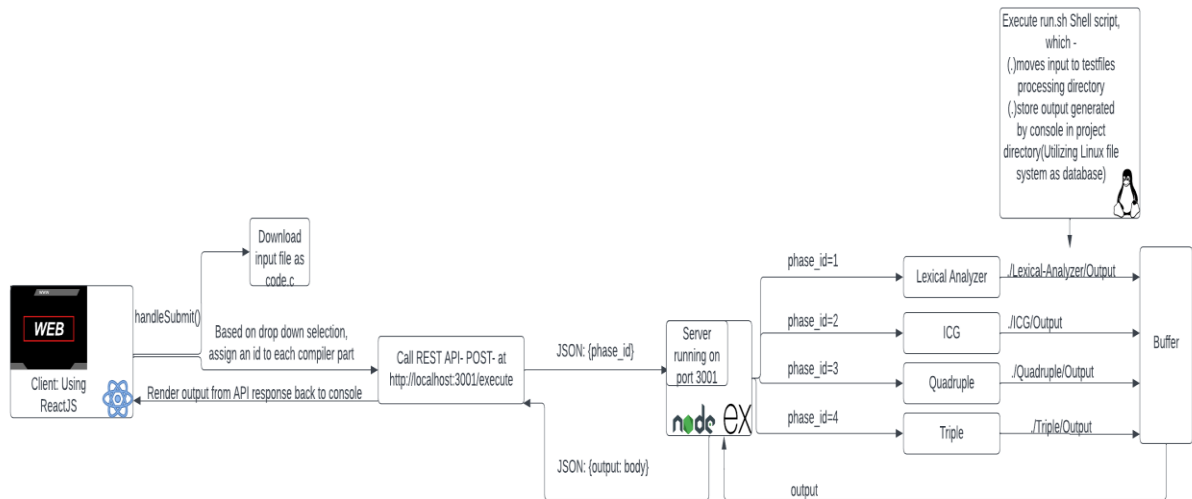


FIG 3.1 Architecture Diagram of the Project

The project consists of several components that work together to achieve its goal. The Front-end UI is responsible for presenting the user interface to the end-user, and it is built using ReactJS, which is a popular JavaScript library for building user interfaces. The REST API serves as the communication interface between the front-end UI and the back-end logic. It is implemented using a RESTful architecture that enables communication between different components using HTTP protocols. On the other hand, the back-end logic consists of several components, including the Lexical Analyzer and the Intermediate Code Generator (ICG). The Lexical Analyzer is responsible for analysing the input source code and generating a stream of tokens, while the ICG is responsible for generating an intermediate code representation of the input source code. The Quadruple and Triple data structures are used to represent the intermediate code generated by the ICG, and they are also implemented as a part of the back-end logic. Overall, the project architecture involves the front-end UI, the REST API, and the back-end logic components, including the Lexical Analyzer, ICG, and data structures. Each component has a specific role and works together to achieve the project's objective.

3.2 COMPONENTS DIAGRAMS

3.2.1 LEXICAL ANALYSIS

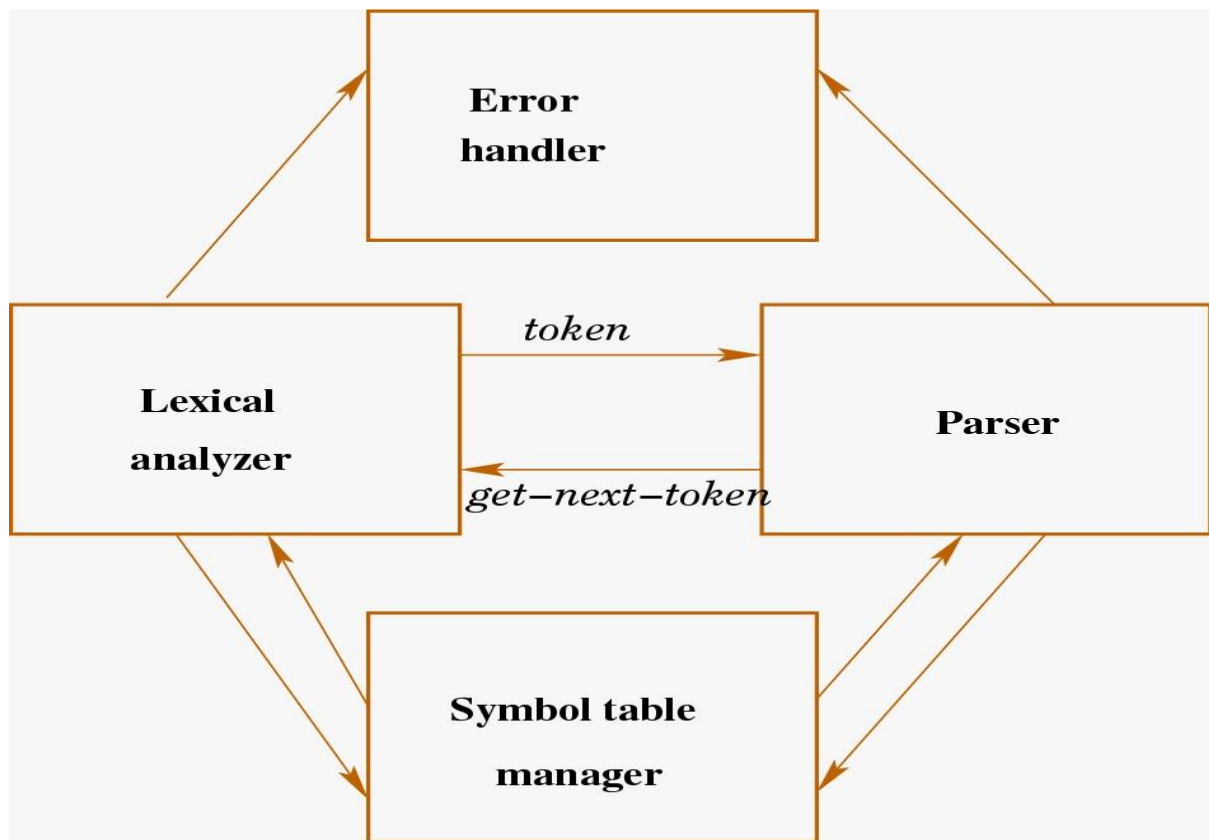


FIG 3.2.1 Component Diagram of Lexical Analyzer

The component diagram of lexical analysis includes four key components: the lexical analyser, error handler, symbol table, and parser. The lexical analyser breaks down the input source code into a sequence of tokens based on predefined rules and passes them to the parser. The error handler detects and reports any errors that occur during the lexical analysis phase, and communicates with both the lexer and the parser to resolve errors. The symbol table maintains a record of all symbols declared in the source code, and is used by both the lexer and parser. The parser analyses the sequence of tokens generated by the lexer and constructs an abstract syntax tree using a formal grammar. Together, these components work to transform the source code into a structured representation that can be further processed by the compiler.

3.2.2 INTERMEDIATE CODE GENERATION

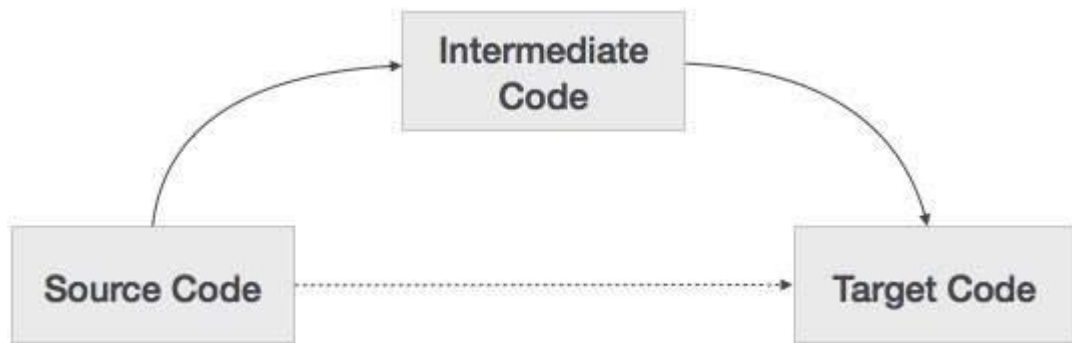


FIG 3.2.2 Component Diagram of Intermediate Code Generator

A component diagram is a type of UML diagram that depicts the system's components and their relationships. In the context of intermediate code generation in a compiler, a component diagram can be used to illustrate the flow of information between the different components of the system, namely the source code, intermediate code, and target code.

The source code is the original program written in a high-level language that the compiler receives as input. The intermediate code is an intermediate representation of the program that is generated during the compilation process. The target code is the final output of the compiler, which is often in the form of machine code or assembly language.

In the component diagram, the source code component would be depicted as the input to the system, with an arrow pointing towards the intermediate code component. The intermediate code component would be the central component of the system, as it represents the intermediate representation of the program that is generated during the compilation process. It would have arrows pointing towards both the source code component and the target code component. This indicates that the intermediate code is generated from the source code and is used to generate the target code.

3.2.3 QUADRUPLER

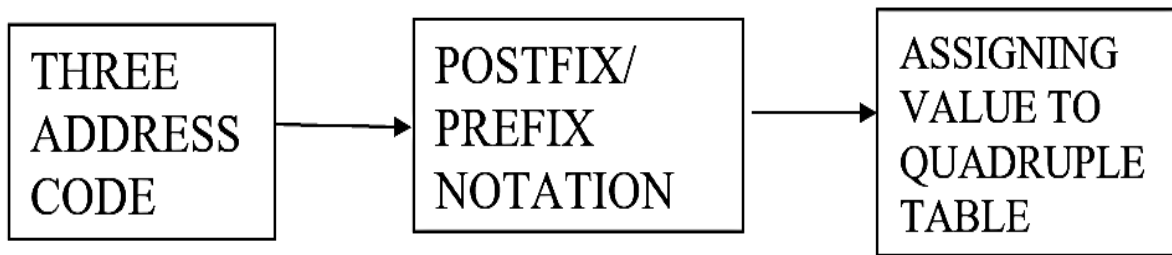


FIG 3.2.3 Component Diagram of Quadruple

In the component diagram, three-address code would be depicted as the input to the system, with arrows pointing towards the postfix/prefix expression component and the quadruple component. The postfix/prefix expression component would be responsible for converting the arithmetic and logical expressions in the three-address code into postfix/prefix notation. The quadruple component would be responsible for generating the quadruples from the postfix/prefix expressions.

The quadruple component may be composed of several sub-components, each responsible for a different aspect of the quadruple generation process, such as parsing the postfix/prefix expressions, building the quadruples, and optimizing the quadruples.

The quadruple component would be depicted as the output of the system, with an arrow pointing towards the three-address code component. This represents the fact that the quadruples are generated from the three-address code.

3.2.4 TRIPLE

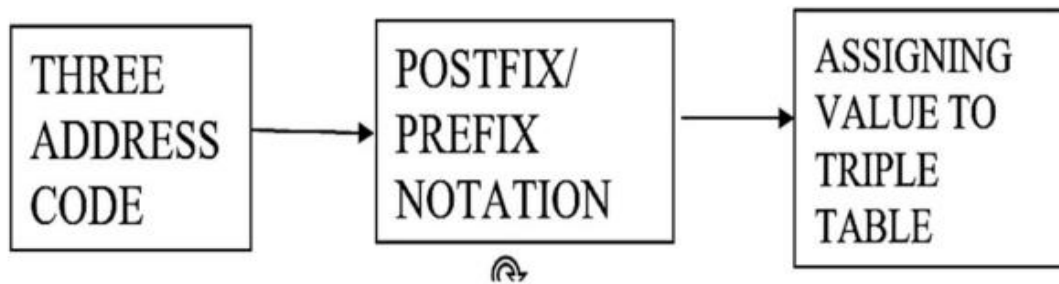


FIG 3.2.4 Component Diagram of Triple

In a component diagram for a triple, the three components of the triple - the operation code, the source operand, and the destination operand - would be represented as separate components, connected by lines that indicate the relationships between them. For example, the operation code component would have a line connecting it to the source operand component, indicating that the operation code specifies the type of operation to be performed on the source operand.

In addition to the three components of the triple, a component diagram for a triple might also include other components that are used in the overall program, such as variables, constants, and control flow structures. These components would also be connected by lines to indicate their relationships.

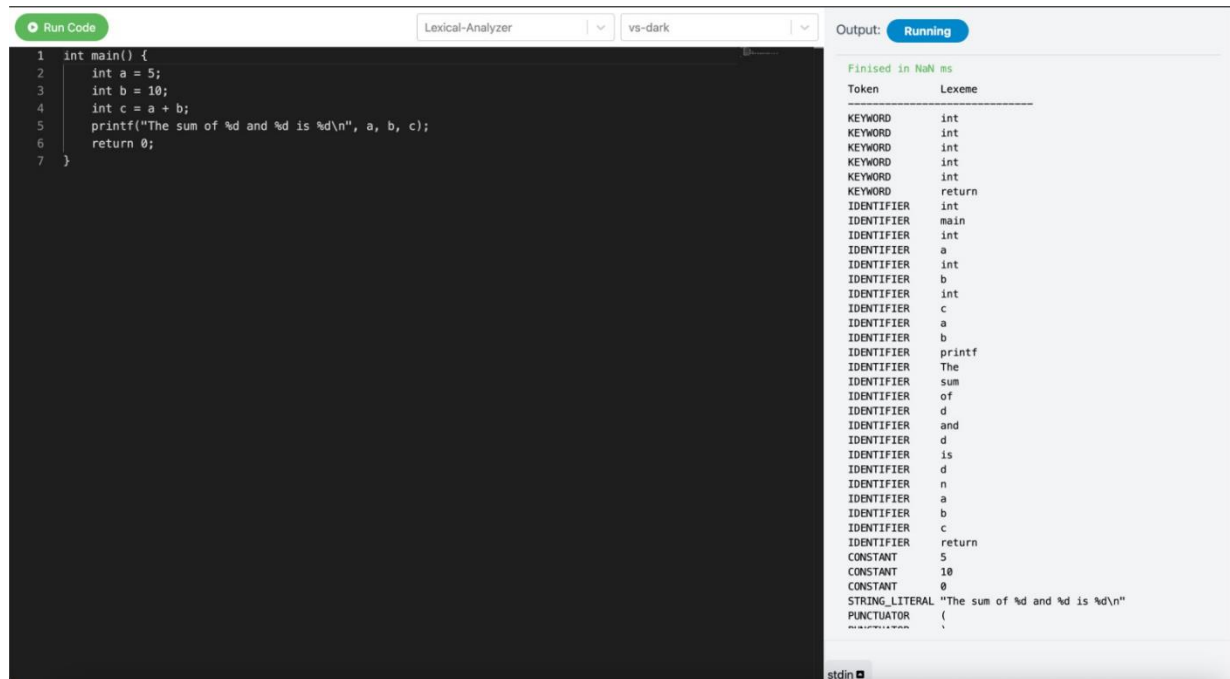
CHAPTER 4

CODING AND TESTING

4.1 LEXICAL ANALYSIS

4.1.1 FRONTEND

LEXICAL ANALYSIS – WITHOUT ERROR



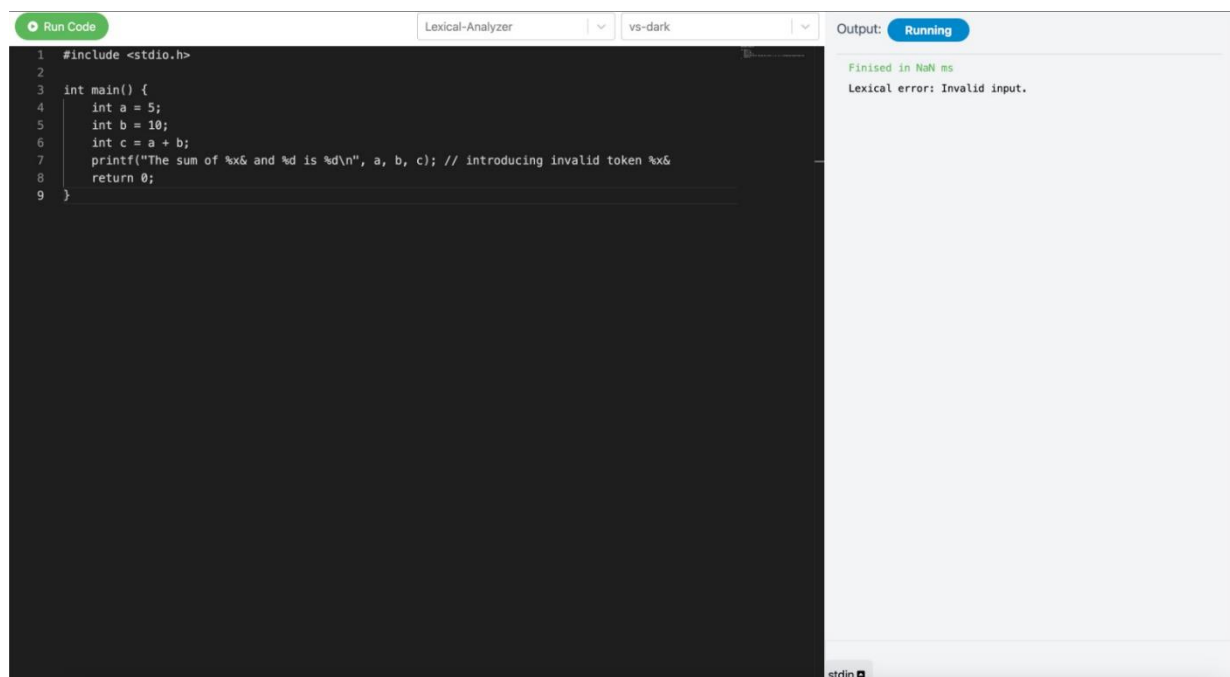
The screenshot shows a web-based Lexical Analyzer interface. On the left, a code editor contains a C program:

```
1 int main() {  
2     int a = 5;  
3     int b = 10;  
4     int c = a + b;  
5     printf("The sum of %d and %d is %d\n", a, b, c);  
6     return 0;  
7 }
```

 The top bar shows 'Run Code', 'Lexical-Analyzer', and 'vs-dark'. The right panel, titled 'Output: Running', shows 'Finised in NaN ms' and a table of tokens and lexemes.

Token	Lexeme
KEYWORD	int
KEYWORD	int
KEYWORD	int
KEYWORD	int
KEYWORD	return
IDENTIFIER	int
IDENTIFIER	main
IDENTIFIER	int
IDENTIFIER	a
IDENTIFIER	int
IDENTIFIER	b
IDENTIFIER	int
IDENTIFIER	c
IDENTIFIER	a
IDENTIFIER	b
IDENTIFIER	printf
IDENTIFIER	The
IDENTIFIER	sum
IDENTIFIER	of
IDENTIFIER	d
IDENTIFIER	and
IDENTIFIER	d
IDENTIFIER	is
IDENTIFIER	d
IDENTIFIER	n
IDENTIFIER	a
IDENTIFIER	b
IDENTIFIER	c
IDENTIFIER	return
CONSTANT	5
CONSTANT	10
CONSTANT	0
STRING_LITERAL	"The sum of %d and %d is %d\n"
PUNCTUATOR	(
PUNCTUATOR)

LEXICAL ANALYSIS – WITH ERROR



The screenshot shows the same Lexical Analyzer interface. The code editor now contains a C program with an invalid token:

```
1 #include <stdio.h>  
2  
3 int main() {  
4     int a = 5;  
5     int b = 10;  
6     int c = a + b;  
7     printf("The sum of %x& and %d is %d\n", a, b, c); // introducing invalid token %x&  
8     return 0;  
9 }
```

 The right panel shows 'Finised in NaN ms' and the message 'Lexical error: Invalid input.'

4.1.2 BACKEND

[illegible]

4.2 INTERMEDIATE CODE GENERATION

4.2.1 FRONTEND

Run Code

ICG

vs-dark

Output: Running

```
1 #include <stdio.h>
2
3 void main()
4 {
5     int x=3,y=2;
6     while(x>y){
7         printf("hello world");
8         x--;
9     }
10 }
11
```

Finished in NaN ms

parser.y:362,27-386,50: warning: unused value: \$3

func begin main

t0 = 3

t1 = 2

L0:

t2 = x > y

IF not t2 GoTo L1

refparam "hello world"

refparam result

call printf, 1

GoTo L0:

L1:

func end

PASSED: ICG Phase

PRINTING SYMBOL TABLE

Symbol Name	Class	Type	Value	I
x	Identifier	int	3	
y	Identifier	int	2	
main	Function	void		
while	Keyword			
int	Keyword			
void	Keyword			
printf	Function			

PRINTING CONSTANT TABLE

constant name	Type
"hello world"	String Constant
2	Number Constant
3	Number Constant

stdin

Run Code

ICG

vs-dark

Output: Running

```
1 #include <stdio.h>
2
3 void main()
4 {
5     int x=3,y=2;
6     while(x>y){
7         printf("hello world");
8         x--;
9     }
10 }
11
```

Line No.

Nesting Count

Params Count

5

99999

-1

5

99999

-1

3

9999

0

6

9999

-1

5

9999

-1

3

9999

-1

7

9999

-1

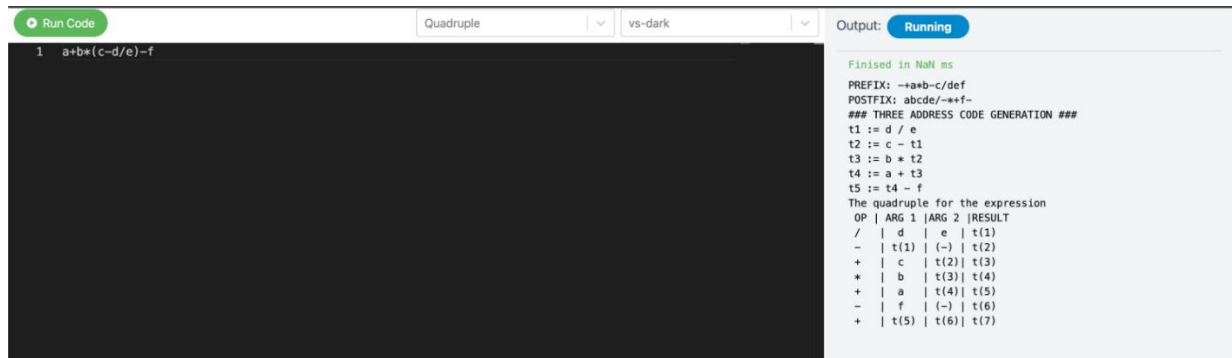
stdin

4.2.2 BACKEND

```
1 OPERATORS = set(['+', '-', '*', '/', '(', ')'])
2 PRI = {'+':1, '-':1, '*':2, '/':2}
3
4 ### INFIX ==> POSTFIX ###
5 def infix_to_postfix(formula):
6     stack = [] # only pop when the coming op has priority
7     output = ''
8     for ch in formula:
9         if ch not in OPERATORS:
10             output += ch
11         elif ch == '(':
12             stack.append('(')
13         elif ch == ')':
14             while stack and stack[-1] != '(':
15                 output += stack.pop()
16             stack.pop() # pop '('
17         else:
18             while stack and stack[-1] != '(' and PRI[ch] <= PRI[stack[-1]]:
19                 output += stack.pop()
20             stack.append(ch)
21     # leftover
22     while stack:
23         output += stack.pop()
24     print(f'POSTFIX: {output}')
25     return output
26
27 ### INFIX ==> PREFIX ###
28 def infix_to_prefix(formula):
29     op_stack = []
30     exp_stack = []
31     for ch in formula:
32         if not ch in OPERATORS:
33             exp_stack.append(ch)
34         elif ch == '(':
35             op_stack.append(ch)
36         elif ch == ')':
37             while op_stack[-1] != '(':
38                 op = op_stack.pop()
39                 a = exp_stack.pop()
40                 b = exp_stack.pop()
```

4.3 QUADRUPLE

4.3.1 FRONTEND



```
1 a+b*(c-d/e)-f
```

Output: Running

Finished in NaN ms

PREFIX: →a+b-c/def
POSTFIX: abcde/-*/-
THREE ADDRESS CODE GENERATION ###
t1 := d / e
t2 := c - t1
t3 := b * t2
t4 := a + t3
t5 := t4 - f

The quadruple for the expression

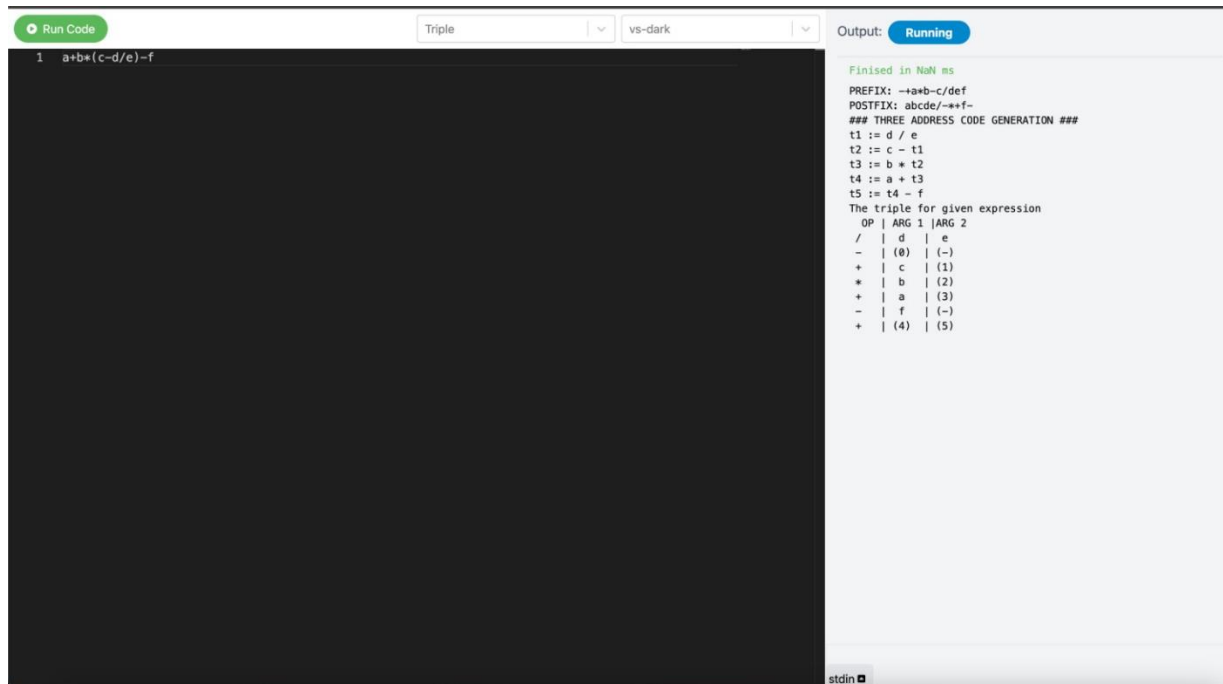
OP	ARG 1	ARG 2	RESULT
/	d	e	t(1)
-	t(1)	c	t(2)
*	t(2)	b	t(3)
+	t(3)	a	t(4)
-	t(4)	f	t(5)

4.3.2 BACKEND

```
88
89 def Quadruple(pos):
90     stack = []
91     op = []
92     x = 1
93     for i in pos:
94         if i not in OPERATORS:
95             stack.append(i)
96         elif i == '-':
97             op1 = stack.pop()
98             stack.append("t(%s)" % x)
99             print("{0:^4s} | {1:^4s} | {2:^4s}|{3:4s}".format(
100                 i, op1, "(-)", " t(%s)" % x))
101             x = x+1
102             if stack != []:
103                 op2 = stack.pop()
104                 op1 = stack.pop()
105                 print("{0:^4s} | {1:^4s} | {2:^4s}|{3:4s}".format(
106                     "+", op1, op2, " t(%s)" % x))
107                 stack.append("t(%s)" % x)
108                 x = x+1
109         elif i == '=':
110             op2 = stack.pop()
111             op1 = stack.pop()
112             print("{0:^4s} | {1:^4s} | {2:^4s}|{3:4s}".format(i, op2, "(-)", op1))
113         else:
114             op1 = stack.pop()
115             op2 = stack.pop()
116             print("{0:^4s} | {1:^4s} | {2:^4s}|{3:4s}".format(
117                 i, op2, op1, " t(%s)" % x))
118             stack.append("t(%s)" % x)
119             x = x+1
120
121
122 print("The quadruple for the expression ")
123 print(" OP | ARG 1 |ARG 2 |RESULT ")
124 Quadruple(pos)
125
```

4.4 TRIPLE

4.4.1 FRONTEND



```
1 a+b*(c-d/e)-f
```

Output: Running

Finished in NaN ms

PREFIX: --a+b-c/def
POSTFIX: abcde/-+*f-

THREE ADDRESS CODE GENERATION

t1 := d / e
t2 := c - t1
t3 := b * t2
t4 := a + t3
t5 := t4 - f

The triple for given expression

OP	ARG 1	ARG 2
/	d	e
-	(t1)	(-)
+	c	(1)
*	b	(2)
+	a	(3)
-	f	(-)
+	(4)	(5)

4.4.2 BACKEND

```
def Triple(pos):  
    stack = []  
    op = []  
    x = 0  
    for i in pos:  
        if i not in OPERATORS:  
            stack.append(i)  
        elif i == '-':  
            op1 = stack.pop()  
            stack.append("(%s)" % x)  
            print("{0:^4s} | {1:^4s} | {2:^4s}".format(i, op1, "(-)"))  
            x = x+1  
            if stack != []:  
                op2 = stack.pop()  
                op1 = stack.pop()  
                print("{0:^4s} | {1:^4s} | {2:^4s}".format("+", op1, op2))  
                stack.append("(%s)" % x)  
                x = x+1  
        elif i == '=':  
            op2 = stack.pop()  
            op1 = stack.pop()  
            print("{0:^4s} | {1:^4s} | {2:^4s}".format(i, op1, op2))  
        else:  
            op1 = stack.pop()  
            if stack != []:  
                op2 = stack.pop()  
                print("{0:^4s} | {1:^4s} | {2:^4s}".format(i, op2, op1))  
                stack.append("(%s)" % x)  
                x = x+1  
    print("The triple for given expression")  
    print(" OP | ARG 1 | ARG 2 ")  
    Triple(pos)
```

CHAPTER 5

RESULT

The implementation of the phases of a compiler, including the lexical analyser, intermediate code generation, quadruples, and triples, has been successful in our project. The integration of modern web technologies, such as React JS and Node JS, has made the compiler more accessible to a wider range of users. Our compiler can effectively translate source code into executable code, making it a useful tool for developers and programmers. Through the development process, we encountered various challenges, such as ensuring the accuracy of the intermediate code generation process, but we were able to overcome these challenges through careful planning and testing. Overall, our project demonstrates our proficiency in compiler development and our ability to apply the concepts and tools learned in class to real-world applications. We believe that our compiler has the potential to be a valuable resource for the programming community and we are excited to see how it will be used in the future.

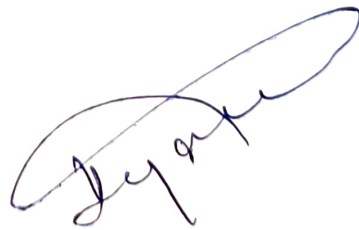
CHAPTER 6

CONCLUSION

In conclusion, the development of a compiler that implements the various phases of the compilation process has been a challenging and rewarding experience. Through the implementation of the lexical analyser, intermediate code generation, quadruples, and triples, we have gained a deeper understanding of the inner workings of compilers and the importance of each phase in the compilation process. The integration of modern web technologies has made the compiler more user-friendly and accessible to a wider range of users. We believe that our compiler has the potential to be a valuable resource for developers and programmers, and we look forward to seeing how it will be used in the future.

The development process has also taught us valuable lessons about software engineering and project management. We learned the importance of careful planning, testing, and documentation in ensuring the success of a project. Additionally, we discovered the importance of communication and collaboration in a team setting, and how effective teamwork can lead to more efficient and effective outcomes.

Overall, our project has been a valuable learning experience that has allowed us to apply the concepts and tools learned in class to a real-world application. We are proud of what we have accomplished and look forward to applying our newfound knowledge to future projects and endeavours.

A handwritten signature in blue ink, appearing to read 'Jeyaraj', is located in the bottom right corner of the page.