# ACM ICPC REGIONAL 2012

#### 1. Generales

# 1.1. LIS en O(nlgn).

```
vector<int> LIS(vector<int> X) {
   int n = X.size(), L = 0, M[n+1], P[n];
   int lo, hi, mi;

L = 0;
   M[0] = 0;

for(int i=0, j; i<n; i++) {
   lo = 0; hi = L;

   while(lo!=hi) {
      mi = (lo+hi+1)/2;
      if(X[M[mi]] < X[i]) lo = mi;
      else hi = mi-1;
   }

   j = lo;</pre>
```

## 1.2. Problema de Josephus.

```
int survivor(int n, int m) {
  for (int s=0,i=1;i<=n;++i) s = (s+m)%i;</pre>
```

#### 1.3. Contar inversiones.

```
#define MAX_SIZE 100000
int A[MAX_SIZE],C[MAX_SIZE],pos1,pos2,sz;
```

```
P[i] = M[j];

if(j==L || X[i]<X[M[j+1]]) {
    M[j+1] = i;
    L = max(L,j+1);
}

int a[L];

for(int i=L-1,j=M[L];i>=0;i--) {
    a[i] = X[j];
    j = P[j];
}

return vector<int>(a,a+L);
}
```

```
return (s+1);
```

```
long long countInversions(int a, int b) {
   if (a==b) return 0;

int c = ((a+b)>>1);
```

```
long long aux = countInversions(a,c)+countInversions(c+1,b);
pos1 = a; pos2 = c+1; sz = 0;

while(pos1<=c && pos2<=b){
   if(A[pos1]<A[pos2]) C[sz] = A[pos1++];
   else{
      C[sz] = A[pos2++];
      aux += c-pos1+1;
   }
   ++sz;</pre>
```

## 1.4. Números dada la suma de pares.

```
bool solve(int N, int sums[], int ans[]){
  int M = N*(N-1)/2;
  multiset<int> S;
  multiset<int> :: iterator it;

  sort(sums,sums+M);

  for(int i = 2;i<M;++i) {
    if((sums[0]+sums[1]-sums[i])%2!=0) continue;

    ans[0] = (sums[0]+sums[1]-sums[i])/2;
    S = multiset<int>(sums,sums+M);

  bool valid = true;

  for(int j = 1;j<N && valid;++j) {</pre>
```

#### 2.1. Ciclo de Euler.

```
// Las listas de adyacencia se deben ordenar de forma ascendente para
// obtener el ciclo lexicografico minimo deacuerdo a la numeracion
// de las aristas
#define MAX_V 44
#define MAX_E 1995
int N,deg[MAX_V],eu[MAX_E],ev[MAX_E];
list<int> G[MAX_V],L;
```

```
if(pos1>c) memcpy(C+sz,A+pos2,(b-pos2+1)*sizeof(int));
else memcpy(C+sz,A+pos1,(c-pos1+1)*sizeof(int));

sz = b-a+1;
memcpy(A+a,C,sz*sizeof(int));

return aux;
}

ans[j] = (*S.begin())-ans[0];

for(int k = 0;k<j && valid;++k){
    it = S.find(ans[k]+ans[j]);

    if(it==S.end()) valid = false;
    else S.erase(it);
    }
}

if(valid) return true;
}</pre>
```

#### 2. Grafos

```
bool visited[MAX_V];
stack<int> S;
queu<int> Q;

bool connected() {
   int cont = 0;
   Q.push(0);
   memset(visited,false,sizeof(visited));
   visited[0] = true;
```

```
while(!Q.empty()){
      int v = Q.front(); Q.pop();
      ++cont;
      for(list<int>::iterator it = G[v].begin();it!=G[v].end();++it){
            int e = *it;
            int w = eu[e] == v? ev[e] : eu[e];
         if(!visited[w]){
            visited[w] = true;
            Q.push(w);
   return cont == N;
bool eulerian(){
   if(!connected()) return false;
   for (int v = 0; v < N; ++v)
      if(deg[v]&1)
         return false;
   return true;
void take_edge(int v, int w){
   --deg[v]; --deg[w];
   int e = G[v].front();
   G[v].pop_front();
   for(list<int>::iterator it = G[w].begin();it!=G[w].end();++it){
      if(*it==e){
         G[w].erase(it);
         break;
```

```
void euler(int v) {
  while(true) {
      if(G[v].empty()) break;
      int e = G[v].front();
      int w = eu[e] == v? ev[e] : eu[e];
      S.push(e);
      take_edge(v,w);
      v = w;
bool find_cycle(int s) {
  if(!eulerian()) return false;
  int v = s,e;
  L.clear();
   do {
      euler(v);
      e = S.top(); S.pop();
      L.push_back(e);
     v = eu[e] == v? ev[e] : eu[e];
   }while(!S.empty());
  return true;
void print_cycle(int s){
  if(!find_cycle(s)) printf("-1\n");
   else{
      bool first = true;
      reverse(L.begin(), L.end());
      for(list<int>::iterator e = L.begin();e!=L.end();++e){
            if(!first) printf(""");
           first = false;
        printf("%d",1+(*e));
     printf("\n");
```

## 2.2. Euler (Directed graph).

```
int V,E,to[32000],nxt[32000],last[1000],now[1000];
int ans[32000];

void init(){
    memset(last,-1,sizeof(last));
    E = 0;
}

void make_edge(int u, int v){
    to[E] = v; nxt[E] = last[u]; last[u] = E++;
}

// A : vertice inicial
void euler(int A){
    for(int i = 0;i < V;++i)
        now[i] = last[i];

    stack<int> S;
    S.push(A);
```

#### 2.3. Punto de articulación.

```
#define SZ 100
bool M[SZ][SZ];
int N,colour[SZ],dfsNum[SZ],num,pos[SZ],leastAncestor[SZ],parent[SZ];

int dfs(int u) {
   int ans = 0,cont = 0,v;

   stack<int> S;
   S.push(u);

   while(!S.empty()) {
      v = S.top();
      if(colour[v]==0) {
       colour[v] = 1;
       dfsNum[v] = num++;
       leastAncestor[v] = num;
      }

   for(;pos[v]<N;++pos[v]) {
      if(M[v][pos[v]] && pos[v]!=parent[v]) {</pre>
```

```
int cur,sz = 0;
while(!S.empty()) {
    cur = S.top();

    if(now[cur] != -1) {
        S.push(to[now[cur]]);
        now[cur] = nxt[now[cur]];
    }else{
        ans[sz++] = cur;
        S.pop();
    }
}
for(int i = sz - 1;i > 0;--i)
    printf("%d_%d\n",ans[i] + 1,ans[i - 1] + 1)
}
```

```
if(colour[pos[v]]==0) {
    parent[pos[v]]=v;
    S.push(pos[v]);
    if(v==u) ++cont;
    break;
    }else leastAncestor[v]<?=dfsNum[pos[v]];
}

if(pos[v]==N) {
    colour[v] = 2;
    S.pop();

    if(v!=u) leastAncestor[parent[v]]<?=leastAncestor[v];
}

if(cont>1) {
    ++ans;
    printf("%d\n",u);
```

```
for(int i = 0;i<N;++i) {
    if(i==u) continue;
    for(int j = 0;j<N;j++)
        if(M[i][j] && parent[j]==i && leastAncestor[j]>=dfsNum[i]) {
            printf("%d\n",i);
            ++ans;
            break;
        }
}
return ans;
```

## 2.4. Detección de puentes.

```
#define SZ 100
bool M[SZ][SZ];
int N, colour[SZ], dfsNum[SZ], num, pos[SZ], leastAncestor[SZ], parent[SZ];
void dfs(int u){
   int v;
   stack<int> S;
   S.push(u);
   while(!S.empty()){
      v = S.top();
      if(colour[v] == 0) {
          colour[v] = 1;
         dfsNum[v] = num++;
         leastAncestor[v] = num;
      for(;pos[v]<N;++pos[v]){</pre>
         if(M[v][pos[v]] && pos[v]!=parent[v]){
             if(colour[pos[v]] == 0) {
                parent[pos[v]] = v;
               S.push(pos[v]);
             }else leastAncestor[v] <?= dfsNum[pos[v]];</pre>
      if (pos[v] ==N) {
```

```
void Articulation_points() {
   memset(colour, 0, sizeof(colour));
   memset (pos, 0, sizeof (pos));
   memset (parent, -1, sizeof (parent));
   num = 0;
   int total = 0;
   for(int i = 0;i<N;++i) if(colour[i]==0) total += dfs(i);</pre>
  printf("#_Articulation_Points_:_%d\n",total);
         colour[v] = 2;
         S.pop();
         if(v!=u) leastAncestor[parent[v]] <?= leastAncestor[v];</pre>
void Bridge_detection(){
   memset(colour, 0, sizeof(colour));
   memset(pos, 0, sizeof(pos));
  memset(parent,-1, sizeof(parent));
  num = 0;
   int ans = 0;
   for(int i = 0;i<N;i++) if(colour[i]==0) dfs(i);</pre>
   for (int i = 0; i < N; i++)</pre>
      for(int j = 0; j<N; j++)</pre>
         if(parent[j]==i && leastAncestor[j]>dfsNum[i]) {
             printf("%d_-_%d\n",i,j);
             ++ans;
   printf("%d_bridges\n",ans);
```

## 2.5. Componentes biconexas (Tarjan).

```
#define MAXN 100000
int V;
vector<int> adj[MAXN];
int dfn[MAXN],low[MAXN];
vector< vector<int> > C;
stack< pair<int, int> > stk;
void cache_bc(int x, int y) {
   vector<int> com;
   int tx,ty;
   do{
      tx = stk.top().first, ty = stk.top().second;
      stk.pop();
      com.push_back(tx), com.push_back(ty);
   }while(tx!=x || ty!=y);
   C.push_back(com);
void DFS(int cur, int prev, int number) {
   dfn[cur] = low[cur] = number;
   for (int i = adj[cur].size()-1;i>=0;--i){
      int next = adj[cur][i];
      if (next==prev) continue;
```

## 2.6. DFS para calcular low iterativo.

```
#define MAXN 100001
#define MAXE 500000

int last[MAXN],nxt[2 * MAXE],to[2 * MAXE],ne = 0;

void add_edge(int &u, int &v) {
   to[ne] = v; nxt[ne] = last[u]; last[u] = ne++;
   to[ne] = u; nxt[ne] = last[v]; last[v] = ne++;
}

int low[MAXN],parent[MAXN],level[MAXN],comp[MAXN];
```

```
if (dfn[next] ==-1) {
         stk.push(make_pair(cur,next));
         DFS (next, cur, number+1);
         low[cur] = min(low[cur], low[next]);
         if(low[next]>=dfn[cur]) cache_bc(cur,next);
      }else low[cur] = min(low[cur],dfn[next]);
void biconn_comp() {
   memset(dfn,-1,sizeof(dfn));
  C.clear();
  DFS(0,0,0);
   int comp = C.size();
  printf("%d\n",comp);
   for (int i = 0; i < comp; ++i) {</pre>
      sort(C[i].begin(),C[i].end());
      C[i].erase(unique(C[i].begin(),C[i].end()),C[i].end());
      int m = C[i].size();
      for(int j = 0; j < m; ++j) printf("%d.", 1 + C[i][j]);</pre>
      printf("\n");
```

```
void dfs(int r) {
    int u,v;

    stack<int> S;

    S.push(r);
    comp[r] = r;
    low[r] = level[r] = 0;
    parent[r] = -1;

while(!S.empty()) {
        u = S.top();
    }
}
```

```
for(int &e = last[u];e != -1;e = nxt[e]) {
    v = to[e];

if(comp[v] != -1 && v != parent[u] && level[u] > level[v]) {
        low[u] = min(low[u],level[v]);
    }else if(comp[v] == -1) {
        S.push(v);
        comp[v] = r;
        low[v] = level[v] = level[u] + 1;
        parent[v] = u;
```

## 2.7. Componentes fuertemente conexas (Tarjan).

```
#define MAX_V 100000
vector<int> L[MAX_V],C[MAX_V];
int V, dfsPos, dfsNum[MAX_V], lowlink[MAX_V], num_scc, comp[MAX_V];
bool in_stack[MAX_V];
stack<int> S;
void tarjan(int v) {
   dfsNum[v] = lowlink[v] = dfsPos++;
   S.push(v); in_stack[v] = true;
   for(int i = L[v].size()-1;i>=0;--i){
      int w = L[v][i];
      if(dfsNum[w] ==-1){
         tarjan(w);
         lowlink[v] = min(lowlink[v],lowlink[w]);
      }else if(in_stack[w]) lowlink[v] = min(lowlink[v], lowlink[w]);
   if(dfsNum[v] == lowlink[v]) {
      vector<int> &com = C[num_scc];
      com.clear();
      int aux;
```

# 2.8. Ciclo de peso promedio mínimo (Karp).

```
#define MAX_V 676
vector< pair<int, int> > L[MAX_V+1];
```

```
break;
      if(last[u] == -1){
        S.pop();
         if(u != r)
            low[ parent[u] ] = min(low[ parent[u] ],low[u]);
      do {
         aux = S.top(); S.pop();
         comp[aux] = num_scc;
         com.push_back(aux);
         in_stack[aux] = false;
      }while (aux!=v);
      ++num_scc;
void build_scc(int _V) {
  V = V;
  memset (dfsNum, -1, sizeof (dfsNum));
  memset(in_stack, false, sizeof(in_stack));
  dfsPos = num_scc = 0;
   for (int i = 0; i < V; ++i)
      if(dfsNum[i]==-1)
         tarjan(i);
int dist[MAX_V+1][MAX_V+2];
```

#### 2.9. Minimum cost arborescence.

```
#define MAX_V 1000
typedef int edge_cost;
edge_cost INF = INT_MAX;

int V,root,prev[MAX_V];
bool adj[MAX_V][MAX_V];
edge_cost G[MAX_V][MAX_V],MCA;
bool visited[MAX_V],cycle[MAX_V];

void add_edge(int u, int v, edge_cost c){
   if(adj[u][v]) G[u][v] = min(G[u][v],c);
   else G[u][v] = c;
   adj[u][v] = true;
}

void dfs(int v){
   visited[v] = true;

for(int i = 0;i<V;++i)
   if(!visited[i] && adj[v][i])</pre>
```

```
if (dist[i][n]!=INT_MAX)
         flag = false;
   if(flag){
      //El grafo es aciclico
      return;
   double ans = 1e15;
   for (int u = 0; u+1 < n; ++u) {
      if (dist[u][n] == INT_MAX) continue;
      double W = -1e15;
      for (int k = 0; k < n; ++k)
         if (dist[u][k]!=INT_MAX)
             W = max(W, (double) (dist[u][n]-dist[u][k])/(n-k));
      ans = min(ans, W);
         dfs(i);
bool check() {
   memset(visited, false, sizeof(visited));
  dfs(root);
   for (int i = 0; i < V; ++i)</pre>
      if(!visited[i])
         return false;
   return true;
int exist_cycle(){
  prev[root] = root;
   for (int i = 0; i < V; ++i) {</pre>
      if(!cycle[i] && i!=root){
```

prev[i] = i; G[i][i] = INF;

```
for (int j = 0; j<V; ++j)</pre>
             if(!cycle[j] && adj[j][i] && G[j][i] <G[prev[i]][i])</pre>
                prev[i] = j;
   }
   for (int i = 0, j; i < V; ++i) {</pre>
      if(cycle[i]) continue;
      memset (visited, false, sizeof (visited));
      j = i;
      while(!visited[j]){
          visited[j] = true;
          j = prev[j];
      if(j==root) continue;
      return j;
   return -1;
void update(int v) {
   MCA += G[prev[v]][v];
   for(int i = prev[v];i!=v;i = prev[i]){
      MCA += G[prev[i]][i];
      cycle[i] = true;
   for (int i = 0; i < V; ++i)</pre>
      if(!cycle[i] && adj[i][v])
         G[i][v] -= G[prev[v]][v];
```

## 2.10. Stable marriage.

```
#define MAX_N 500
int N,pref_men[MAX_N][MAX_N],pref_women[MAX_N][MAX_N];
int inv[MAX_N][MAX_N],cont[MAX_N],wife[MAX_N],husband[MAX_N];
```

```
for(int j = prev[v]; j!=v; j = prev[j]) {
      for(int i = 0; i<V; ++i) {</pre>
         if(cycle[i]) continue;
         if(adj[i][j]){
             if(adj[i][v]) G[i][v] = min(G[i][v],G[i][j]-G[prev[j]][j]);
             else G[i][v] = G[i][j]-G[prev[j]][j];
             adj[i][v] = true;
         if(adj[j][i]){
             if(adj[v][i]) G[v][i] = min(G[v][i],G[j][i]);
             else G[v][i] = G[j][i];
             adj[v][i] = true;
bool min_cost_arborescence(int _root) {
   root = _root;
   if(!check()) return false;
   memset(cycle, false, sizeof(cycle));
  MCA = 0;
   int v;
   while((v = exist_cycle())!=-1)
      update(v);
   for (int i = 0; i < V; ++i)</pre>
      if(i!=root && !cycle[i])
         MCA += G[prev[i]][i];
   return true;
void stable_marriage() {
      for(int i = 0; i<N; ++i)</pre>
             for (int \dot{j} = 0; \dot{j} < N; ++\dot{j})
```

inv[i][pref\_women[i][j]] = j;

## 2.11. Bipartite matching (Hopcroft Karp).

```
#define MAX_V1 50000
#define MAX_V2 50000
#define MAX_E 150000
int V1, V2, left[MAX_V2], right[MAX_V1];
int E, to[MAX_E], next[MAX_E], last[MAX_V1];
void hk_init(int v1, int v2){
  V1 = v1; V2 = v2; E = 0;
   memset(last,-1, sizeof last);
void hk_add_edge(int u, int v) {
   to[E] = v; next[E] = last[u]; last[u] = E++;
bool visited[MAX_V1];
bool hk_dfs(int u) {
   if(visited[u]) return false;
  visited[u] = true;
   for(int e = last[u], v; e != -1; e = next[e]) {
      v = to[e];
      if(left[v] == -1 || hk_dfs(left[v])){
         right[u] = v;
         left[v] = u;
```

```
if(husband[w]<0 || inv[w][m]<inv[w][husband[w]]) break;
}

dumped = husband[w];
husband[w] = m;
wife[m] = w;
m = dumped;
}
}</pre>
```

```
return true;
   return false;
int hk_match(){
   memset(left,-1,sizeof left);
  memset(right, -1, sizeof right);
  bool change = true;
   while (change) {
      change = false;
      memset (visited, false, sizeof visited);
      for(int i = 0; i < V1; ++i)</pre>
         if(right[i] == -1)
            change |= hk_dfs(i);
   int ret = 0;
   for(int i = 0;i < V1;++i)</pre>
      if(right[i] != -1) ++ret;
   return ret;
```

## 2.12. Algoritmo húngaro.

```
// Maximiza costo del matching
#define MAX V 500
int V, cost[MAX_V][MAX_V];
int lx[MAX_V], ly[MAX_V];
int max_match, xy[MAX_V], yx[MAX_V], prev[MAX_V];
bool S[MAX_V], T[MAX_V];
int slack[MAX_V], slackx[MAX_V];
int q[MAX_V], head, tail;
void init labels() {
   memset(lx,0,sizeof(lx));
   memset(ly,0,sizeof(ly));
   for (int x = 0; x < V; ++x)
      for (int y = 0; y < V; ++y)
         lx[x] = max(lx[x], cost[x][y]);
void update_labels() {
   int x,y,delta = INT_MAX;
   for(y = 0;y<V;++y) if(!T[y]) delta = min(delta,slack[y]);</pre>
   for (x = 0; x < V; ++x) if (S[x]) lx[x] -= delta;
   for(y = 0; y < V; ++y) if(T[y]) ly[y] += delta;
   for(y = 0;y<V;++y) if(!T[y]) slack[y] -= delta;</pre>
void add_to_tree(int x, int prevx) {
   S[x] = true;
   prev[x] = prevx;
   for (int y = 0; y < V; ++y) {
      if(lx[x]+ly[y]-cost[x][y]<slack[y]){
         slack[y] = lx[x]+ly[y]-cost[x][y];
         slackx[y] = x;
void augment(){
   int x,y,root;
```

```
head = tail = 0;
memset(S, false, sizeof(S));
memset(T, false, sizeof(T));
memset (prev, -1, sizeof (prev));
for (x = 0; x < V; ++x) {
   if (xy[x]==-1) {
      q[tail++] = root = x;
      prev[root] = -2;
      S[root] = true;
      break;
for(y = 0;y<V;++y){
   slack[y] = lx[root]+ly[y]-cost[root][y];
   slackx[y] = root;
while (true) {
   while(head<tail) {</pre>
      x = q[head++];
      for (y = 0; y < V; ++y) {
         if(cost[x][y]==lx[x]+ly[y] && !T[y]){
             if(yx[y]==-1) break;
            T[y] = true;
             q[tail++] = yx[y];
             add_to_tree(yx[y],x);
      if(y<V) break;</pre>
   if(y<V) break;</pre>
   update_labels();
   head = tail = 0;
   for(y = 0;y<V;++y) {
      if(!T[y] && slack[y]==0){
```

```
if(yx[y]==-1) {
    x = slackx[y];
    break;
}

T[y] = true;

if(!S[yx[y]]) {
    q[tail++] = yx[y];
    add_to_tree(yx[y],slackx[y]);
    }
}

if(y<V) break;
}
++max_match;</pre>
```

## 2.13. Non bipartite matching.

```
#define MAXN 222
int n;
bool adj[MAXN][MAXN];
int p[MAXN], m[MAXN], d[MAXN], c1[MAXN], c2[MAXN];
int q[MAXN], *qf, *qb;
int pp[MAXN];
int f(int x) {return x == pp[x] ? x : (pp[x] = f(pp[x]));}
void u(int x, int y) \{pp[f(x)] = f(y);\}
int v[MAXN];
void path(int r, int x) {
  if (r == x) return;
   if (d[x] == 0) {
      path(r, p[p[x]]);
      int i = p[x], j = p[p[x]];
      m[i] = j; m[j] = i;
   else if (d[x] == 1) {
      path(m[x], c1[x]);
```

```
for (int cx = x, cy = y, ty; cx!=-2; cx = prev[cx], cy = ty) {
      ty = xy[cx];
      yx[cy] = cx;
      xy[cx] = cy;
int hungarian(){
   int ret = 0;
   \max \ \text{match} = 0;
  memset(xy,-1,sizeof(xy));
  memset (yx, -1, sizeof(yx));
   init_labels();
   for(int i = 0;i<V;++i) augment();</pre>
   for (int x = 0; x < V; ++x) ret += cost[x][xy[x]];
   return ret;
      path(r, c2[x]);
      int i = c1[x], j = c2[x];
      m[i] = j; m[j] = i;
int lca(int x, int y, int r) {
  int i = f(x), j = f(y);
   while (i != j && v[i] != 2 && v[j] != 1){
     v[i] = 1; v[j] = 2;
      if (i != r) i = f(p[i]);
      if (j != r) j = f(p[j]);
   int b = i, z = j;
  if(v[j] == 1) swap(b, z);
   for (i = b; i != z; i = f(p[i])) v[i] = -1;
  v[z] = -1;
   return b;
void shrink_one_side(int x, int y, int b) {
```

```
for(int i = f(x); i != b; i = f(p[i])){
      u(i, b);
      if(d[i] == 1) c1[i] = x, c2[i] = y, *qb++ = i;
bool BFS(int r) {
   for (int i=0; i<n; ++i)</pre>
      pp[i] = i;
   memset(v, -1, sizeof(v));
   memset(d, -1, sizeof(d));
   d[r] = 0;
   qf = qb = q;
   *qb++ = r;
   while(qf < qb) {</pre>
      for(int x=*qf++, y=0; y<n; ++y) {</pre>
         if(adj[x][y] && m[y] != y && f(x) != f(y)){
            if (d[y] == -1) {
               if(m[y] == -1) {
                                      path(r, x);
                                      m[x] = y; m[y] = x;
                                      return true;
                else{
```

# 2.14. Flujo máximo.

```
struct flow_graph{
   int MAX_V,E,s,t;
   int *cap,*to,*next,*last;
bool *visited;

flow_graph() {}

flow_graph(int V, int MAX_E) {
   MAX_V = V; E = 0;
   cap = new int[2*MAX_E], to = new int[2*MAX_E], next = new int[2*MAX_E];
   last = new int[MAX_V], visited = new bool[MAX_V];
   fill(last,last+MAX_V,-1);
}

void clear() {
```

```
p[y] = x; p[m[y]] = y;
                                     d[y] = 1; d[m[y]] = 0;
                                     *qb++ = m[y];
            else if(d[f(y)] == 0){
                              int b = lca(x, y, r);
                              shrink_one_side(x, y, b);
                              shrink_one_side(y, x, b);
   return false;
int match(){
  memset(m, -1, sizeof(m));
  int c = 0;
  for (int i=0; i<n; ++i)</pre>
     if (m[i] == -1)
        if (BFS(i)) c++;
        else m[i] = i;
   return c;
      fill(last, last+MAX_V, -1);
      E = 0;
  void add edge(int u, int v, int uv, int vu = 0){
     to[E] = v, cap[E] = uv, next[E] = last[u]; last[u] = E++;
      to[E] = u, cap[E] = vu, next[E] = last[v]; last[v] = E++;
  int dfs(int v, int f){
     if(v==t || f<=0) return f;
      if(visited[v]) return 0;
      visited[v] = true;
```

for(int e = last[v];e!=-1;e = next[e]) {

```
int ret = dfs(to[e],min(f,cap[e]));

if(ret>0) {
    cap[e] -= ret;
    cap[e^1] += ret;
    return ret;
    }
}

return 0;
}

int max_flow(int source, int sink) {
```

## 2.15. Flujo máximo (Dinic).

```
struct flow_graph{
                                                                                                  while(head < tail) {</pre>
   static const int MAX_V = 500;
                                                                                                     int v = q[head]; ++head;
   static const int MAX_E = 10000;
                                                                                                     for(int e = last[v];e != -1;e = next[e]){
  int E,s,t,head,tail;
                                                                                                        if(cap[e^1] > 0 && dist[to[e]] == -1){
  int cap[2 * MAX_E], to[2 * MAX_E], next[2 * MAX_E], last[MAX_V], dist[MAX_V], q[MAX_V], now[MAX_V];
                                                                                                           q[tail] = to[e]; ++tail;
                                                                                                           dist[to[e]] = dist[v]+1;
   flow_graph(){
     E = 0;
      memset(last,-1,sizeof last);
                                                                                                  return dist[s] != -1;
   void clear(){
     E = 0;
                                                                                               int dfs(int v, int f){
      memset(last,-1,sizeof last);
                                                                                                  if(v == t) return f;
  void add_edge(int u, int v, int uv) {
                                                                                                  for(int &e = now[v];e != -1;e = next[e]){
     to[E] = v, cap[E] = uv, next[E] = last[u]; last[u] = E++;
                                                                                                     if(cap[e] > 0 && dist[to[e]] == dist[v]-1) {
      to[E] = u, cap[E] = 0, next[E] = last[v]; last[v] = E++;
                                                                                                        int ret = dfs(to[e],min(f,cap[e]));
                                                                                                        if(ret > 0){
   bool bfs(){
                                                                                                           cap[e] -= ret;
      memset(dist,-1,sizeof dist);
                                                                                                           cap[e^1] += ret;
     head = tail = 0;
                                                                                                           return ret;
      q[tail] = t; ++tail;
      dist[t] = 0;
```

s = source, t = sink;

x = dfs(s,INT\_MAX);
if(x==0) break;
f += x;

fill (visited, visited+MAX\_V, false);

int f = 0, x;

while(true) {

return f;

};

```
return 0;
}
int max_flow(int source, int sink) {
   s = source; t = sink;
   int f = 0, df;

while(bfs()) {
   for(int i = 0; i <= sink; ++i) now[i] = last[i];
}</pre>
```

## 2.16. Flujo máximo - Costo Mínimo (Succesive Shortest Path).

```
#define MAX V 350
#define MAX_E 2*12500
typedef int cap_type;
typedef long long cost_type;
const cost_type INF = LLONG_MAX;
int V, E, prev[MAX_V], last[MAX_V], to[MAX_E], next[MAX_E];
bool visited[MAX V];
cap_type flowVal, cap[MAX_E];
cost_type flowCost,cost[MAX_E],dist[MAX_V],pot[MAX_V];
void init(int _V) {
   memset(last,-1, sizeof(last));
   V = _{V}; E = 0;
void add_edge(int u, int v, cap_type _cap, cost_type _cost) {
   to[E] = v, cap[E] = \_cap;
   cost[E] = _cost, next[E] = last[u];
  last[u] = E++;
  to[E] = u, cap[E] = 0;
  cost[E] = -_cost, next[E] = last[v];
   last[v] = E++;
// only if there is initial negative cycle
void BellmanFord(int s, int t) {
   bool stop = false;
   for(int i = 0;i<V;++i) dist[i] = INF;</pre>
   dist[s] = 0;
```

```
while(true) {
            df = dfs(s,INT_MAX);
            if(df == 0) break;
             f += df;
      return f;
};
   for(int i = 1;i<=V && !stop;++i){</pre>
      stop = true;
      for (int j = 0; j < E; ++ j) {</pre>
         int u = to[j^1], v = to[j];
         if(cap[j]>0 && dist[u]!=INF && dist[u]+cost[j]<dist[v]) {</pre>
             stop = false;
            dist[v] = dist[u]+cost[j];
   for(int i = 0;i<V;++i) if (dist[i]!=INF) pot[i] = dist[i];</pre>
void mcmf(int s, int t){
   flowVal = flowCost = 0;
   memset (pot, 0, sizeof (pot));
   BellmanFord(s,t);
   while(true) {
      memset (prev, -1, sizeof (prev));
      memset (visited, false, sizeof (visited));
      for(int i = 0;i<V;++i) dist[i] = INF;</pre>
      priority_queue< pair<cost_type, int> > Q;
      Q.push(make_pair(0,s));
      dist[s] = prev[s] = 0;
      while(!Q.empty()){
```

```
int aux = Q.top().second;
Q.pop();

if(visited[aux]) continue;
visited[aux] = true;

for(int e = last[aux];e!=-1;e = next[e]) {
    if(cap[e]<=0) continue;
    cost_type new_dist = dist[aux]+cost[e]+pot[aux]-pot[to[e]];
    if(new_dist<dist[to[e]]) {
        dist[to[e]] = new_dist;
        prev[to[e]] = e;
        Q.push(make_pair(-new_dist,to[e]));
    }
}</pre>
```

## 2.17. Flujo máximo (Dinic + Lower Bounds).

```
struct flow_graph{
  int V,E,s,t;
  int *flow, *low, *cap, *to, *next, *last, *delta;
  int *dist,*q,*now,head,tail;
  flow_graph(){}
   flow_graph(int V, int E) {
      (*this).V = V; (*this).E = 0;
     int TE = 2 \star (E+V+1);
     flow = new int[TE]; low = new int[TE]; cap = new int[TE];
     to = new int[TE]; next = new int[TE];
     last = new int[V+2]; delta = new int[V];
     dist = new int[V+2]; q = new int[V+2]; now = new int[V+2];
  void clear(int V) {
      (*this).V = V; (*this).E = 0;
      fill(last, last+V, -1);
  void add_edge(int a, int b, int 1, int u) {
     to[E] = b; low[E] = 1; cap[E] = u; flow[E] = 0;
     next[E] = last[a]; last[a] = E++;
     to[E] = a; low[E] = 0; cap[E] = 0; flow[E] = 0;
```

```
if (prev[t]==-1) break;
   cap_type f = cap[prev[t]];
   for(int i = t;i!=s;i = to[prev[i]^1]) f = min(f,cap[prev[i]]);
   for(int i = t;i!=s;i = to[prev[i]^1]){
      cap[prev[i]] -= f;
      cap[prev[i]^1] += f;
   flowVal += f;
   flowCost += f*(dist[t]-pot[s]+pot[t]);
   for(int i = 0; i < V; ++i) if (prev[i]!=-1) pot[i] += dist[i];</pre>
   next[E] = last[b]; last[b] = E++;
bool bfs() {
   fill (dist, dist+V+2,-1);
   head = tail = 0;
   q[tail] = t; ++tail;
   dist[t] = 0;
   while(head<tail){</pre>
      int v = q[head]; ++head;
      for(int e = last[v];e!=-1;e = next[e]){
         if(cap[e^1]>flow[e^1] && dist[to[e]]==-1){
            q[tail] = to[e]; ++tail;
            dist[to[e]] = dist[v]+1;
   return dist[s]!=-1;
int dfs(int v, int f) {
```

if(v==t) return f;

```
for(int &e = now[v];e!=-1;e = next[e]){
      if(cap[e]>flow[e] && dist[to[e]] == dist[v]-1) {
         int ret = dfs(to[e],min(f,cap[e]-flow[e]));
         if(ret>0){
            flow[e] += ret;
            flow[e^1] -= ret;
            return ret;
   return 0;
int max_flow(int source, int sink) {
   fill(delta,delta+V,0);
   for(int e = 0; e < E; e += 2) {</pre>
      delta[to[e^1]] -= low[e];
      delta[to[e]] += low[e];
      cap[e] -= low[e];
   last[V] = last[V+1] = -1;
   int sum = 0;
   for(int i = 0;i<V;++i){</pre>
      if(delta[i]>0){
         add_edge(V,i,0,delta[i]);
         sum += delta[i];
      if(delta[i]<0) add_edge(i,V+1,0,-delta[i]);</pre>
   add_edge(sink, source, 0, INT_MAX);
   s = V; t = V+1;
   int f = 0, df;
```

```
fill(flow,flow+E,0);
      while(bfs()){
         for(int i = V+1;i>=0;--i) now[i] = last[i];
         while(true) {
           df = dfs(s,INT_MAX);
            if(df==0) break;
            f += df;
        }
      if(f!=sum) return -1;
      for(int e = 0; e < E; e += 2) {</pre>
         cap[e] += low[e];
         flow[e] += low[e];
         flow[e^1] -= low[e];
         cap[e^1] -= low[e];
      s = source; t = sink;
      last[s] = next[last[s]];
      last[t] = next[last[t]];
      E = 2;
      while(bfs()){
         for(int i = V-1;i>=0;--i) now[i] = last[i];
         while(true) {
            df = dfs(s,INT_MAX);
            if(df==0) break;
            f += df;
      return f;
};
```

## 2.18. Corte mínimo de un grafo (Stoer - Wagner).

```
#define MAX_V 500
int M[MAX_V][MAX_V], w[MAX_V];
bool A[MAX_V], merged[MAX_V];
int minCut(int n) {
   int best = INT_MAX;
   for(int i=1;i<n;++i) merged[i] = false;</pre>
   merged[0] = true;
   for(int phase=1;phase<n;++phase) {</pre>
      A[0] = true;
      for (int i=1; i<n; ++i) {</pre>
          if(merged[i]) continue;
         A[i] = false;
          w[i] = M[0][i];
      int prev = 0,next;
      for(int i=n-1-phase; i>=0; --i) {
          // hallar siguiente vertice que no esta en A
          next = -1;
          for (int j=1; j<n; ++j)</pre>
             if(!A[j] && (next==-1 || w[j]>w[next]))
```

#### 2.19. Graph Facts.

```
Un grafo es bipartito si y solo si no contiene ciclos de longitud impar.

Todos los arboles son bipartitos.

Las aristas que forman un ciclo, se encuentran en una misma componente biconexa.

Minimum Vertex Cover: para V = (S,T)

DFS desde los vertices que no estan cubiertos por alguna arista del matching, para moverse:

- De izq. a der. usar las aristas que no estan en el matching
```

```
next = j;
         A[next] = true;
         if(i>0){
            prev = next;
            // actualiza los pesos
            for(int j=1; j<n;++j)
               if(!A[j]) w[j] += M[next][j];
      if(best>w[next]) best = w[next];
      // mezcla s y t
      for (int i=0; i<n; ++i) {</pre>
         M[i][prev] += M[next][i];
         M[prev][i] += M[next][i];
      merged[next] = true;
   return best;
- De der. a izq. usar las aristas que estan en el matching
Estan en el vertex cover(independent set):
- De S los no alcanzados (los alcanzados)
- De T los alcanzados(los no alcanzados)
Para usar Teorema de Dilworth colocar tambien aristas que resulten de la transitividad.
Un grafo con grados de vertices iguales a 1 o 2, consiste solo de caminos y ciclos.
```

#### 3.1. Knuth-Morris-Pratt.

```
#define MAX_L 70
int f[MAX_L];

void prefixFunction(string P) {
   int n = P.size(), k = 0;
   f[0] = 0;

   for(int i=1;i<n;++i) {
      while(k>0 && P[k]!=P[i]) k = f[k-1];
      if(P[k]==P[i]) ++k;
      f[i] = k;
   }
}
```

#### 3.2. Aho-Corasick.

```
struct AhoCorasick{
   static const int UNDEF = 0;
   static const int MAXN = 360;
  static const int CHARSET = 26;
  int end, have;
  int tag[MAXN];
   int fail[MAXN];
   int trie[MAXN][CHARSET];
   void init(){
      tag[0] = UNDEF;
      fill(trie[0],trie[0] + CHARSET,-1);
     end = 1;
     have = 0;
   void add(int len, const int* s) {
      int p = 0;
      for(int i = 0; i < len; ++i){</pre>
         if(trie[p][*s] == -1) {
            tag[end] = UNDEF;
            fill(trie[end], trie[end] + CHARSET, -1);
```

```
int KMP(string P, string T){
  int n = P.size(), L = T.size(), k = 0, ans = 0;
   for (int i=0; i<L; ++i) {</pre>
      while (k>0 \&\& P[k]!=T[i]) k = f[k-1];
      if(P[k]==T[i]) ++k;
      if(k==n){
         ++ans;
         k = f[k-1];
   return ans;
            trie[p][*s] = end++;
         p = trie[p][*s];
      tag[p] = (1 \ll have);
      ++have;
   void build() {
      queue<int> bfs;
      fail[0] = 0;
      for(int i = 0;i < CHARSET;++i) {</pre>
         if(trie[0][i] != -1){
            fail[trie[0][i]] = 0;
            bfs.push(trie[0][i]);
         }else{
            trie[0][i] = 0;
```

```
while(!bfs.empty()) {
   int p = bfs.front();
   tag[p] |= tag[fail[p]];
   bfs.pop();

for(int i = 0;i < CHARSET;++i) {
    if(trie[p][i] != -1) {
      fail[trie[p][i]] = trie[fail[p]][i];
    }
}</pre>
```

## 3.3. Algoritmo Z.

```
int next[MAX_P_LEN];
// next[i] : lcp entre la cadena y su sufijo
// a partir del i-esimo caracter

void prefix_kmp(char *P) {
    int L = strlen(P),p = 0,t;

    for(int i = 1; i < L; i++) {
        if(i < p && next[i-t] < p-i) next[i] = next[i-t];
        else {
            int j = max(0, p-i);

            while(i+j < L && P[i+j] == P[j]) ++j;

            next[i] = j;
            p = i + j;
            t = i;
        }
}</pre>
```

#### 3.4. Palíndromos.

```
void manacher(int n, const char s[], int p[]) {
  for (int i = 0, j = 0, k = 0; i <= 2 * (n - 1); ++i) {
    int l = i < k ? min(p[j + j - i], (k - i) / 2) : 0;
    int a = i / 2 - 1, b = (i + 1) / 2 + 1;

  while (0 <= a && b < n && s[a] == s[b]) {
      --a;
      ++b;
      ++1;
    }
}</pre>
```

```
bfs.push(trie[p][i]);
           }else{
            trie[p][i] = trie[fail[p]][i];
};
void LCP(char * P, char *T, int *lcp) {
  int LP = strlen(P),LT = strlen(T);
  int p = 0,t;
  for (int i = 0; i < LT; i++) {</pre>
     if(i 
       int j = max(0, p-i);
       while (i+j < LT && T[i+j] == P[j]) ++j;
       lcp[i] = j;
       p = i + j;
       t = i;
```

```
p[i] = 1;
if(k < 2 * b) {
    j = i;
    k = 2 * b;
}
</pre>
```

# 4. Geometría

## 4.1. Punto y Línea.

```
const double eps = 1e-9;
struct point{
    double x,y;

    point(){}
    point(double _x, double _y) : x(_x), y(_y){}

    double cross(point P){
        return x * P.y - y * P.x;
    }

    bool operator < (const point &p) const{
        if(fabs(x-p.x)>eps) return x<p.x;
        return y>p.y;
    }
};

double cross(point a, point b){
    return a.x * b.y - a.y * b.x;
}
```

# 4.2. Ángulo entre dos vectores.

```
double get_angle(point P1, point P2){
   double sina = P1.y / P1.abs(),cosa = P1.x / P1.abs();
   double sinb = P2.y / P2.abs(),cosb = P2.x / P2.abs();
   double sinc = sinb * cosa - sina * cosb;
   double cosc = cosb * cosa + sina * sinb;
```

## 4.3. Círculos.

```
point get_center(point A, point B, point C) {
   point v1 = (B - A).perp(),v2 = C - A;
   point m1 = (A + B) * 0.5;
   point m2 = (A + C) * 0.5;
```

```
bool polar_cmp(point a, point b) {
   if(a.x >= 0 && b.x < 0) return true;</pre>
   if(a.x < 0 && b.x >= 0) return false;
   if(a.x == 0 && b.x == 0){
      if(a.y > 0 && b.y < 0) return false;</pre>
      if(a.y < 0 && b.y > 0) return true;
   return cross(a,b) > 0;
struct line{
   point p1,p2;
   line(){}
   line(point _p1, point _p2){
      p1 = _p1; p2 = _p2;
      if(p1.x>p2.x) swap(p1,p2);
};
   double x = atan2(sinc,cosc);
   if(x < 0) x += 2 * M_PI;</pre>
   return x;
   double k = (m2 - m1).dot(v2) / v1.dot(v2);
   return m1 + v1 * k;
```

## 4.4. Polígonos.

```
//verdadero : sentido anti-horario, Complejidad : O(n)
bool ccw(const vector<point> &poly) {
   //primero hallamos el punto inferior ubicado mas a la derecha
   int ind = 0,n = poly.size();
   double x = poly[0].x,y = poly[0].y;
   for (int i=1; i < n; i++) {</pre>
      if (poly[i].y>y) continue;
      if (fabs(poly[i].y-y)<eps && poly[i].x<x) continue;</pre>
      ind = i;
      x = poly[i].x;
      y = poly[i].y;
   if (ind==0) return ccw(poly[n-1],poly[0],poly[1]);
   return ccw(poly[ind-1],poly[ind],poly[(ind+1)%n]);
bool isInConvex(vector <Point> &A, const Point &P) {
 int n = A.size(), lo = 1, hi = A.size() - 1;
 if(area(A[0], A[1], P) <= 0) return 0;</pre>
 if(area(A[n-1], A[0], P) <= 0) return 0;</pre>
 while(hi - lo > 1) {
   int mid = (lo + hi) / 2;
   if(area(A[0], A[mid], P) > 0) lo = mid;
   else hi = mid;
 return area(A[lo], A[hi], P) > 0;
```

# 4.5. Convex Hull (Monotone Chain).

```
vector<point> ConvexHull(vector<point> P) {
  sort(P.begin(),P.end());
  int n = P.size(),k = 0;
  point H[2*n];
  for(int i=0;i<n;++i) {</pre>
```

```
bool PointInsidePolygon(const point &P, const vector<point> &poly) {
   int n = poly.size();
   bool in = 0;
   for (int i = 0, j = n-1; i < n; j = i++) {
      double dx = poly[j].x-poly[i].x;
      double dy = poly[j].y-poly[i].y;
      if((poly[i].y<=P.y+eps && P.y<poly[j].y) ||</pre>
          (poly[j].y<=P.y+eps && P.y<poly[i].y))</pre>
         if(P.x-eps<dx*(P.y-poly[i].y)/dy+poly[i].x)</pre>
            in \hat{}=1;
   return in;
//valor positivo : vertices orientados en sentido antihorario
//valor negativo : vertices orientados en sentido horario
double signed_area(const vector<point> &poly) {
 int n = poly.size();
 if(n<3) return 0.0;
 double S = 0.0;
 for (int i=1; i<=n; ++i)</pre>
  S += poly[i%n].x*(poly[(i+1)%n].y-poly[i-1].y);
 S /= 2;
 return S;
      while (k>=2 \&\& !ccw(H[k-2],H[k-1],P[i])) --k;
      H[k++] = P[i];
   for (int i=n-2, t=k; i>=0; --i) {
```

while (k>t && !ccw(H[k-2],H[k-1],P[i])) --k;

```
H[k++] = P[i];
```

## 4.6. Teorema de Pick.

```
A = I + B/2 - 1, donde:
A = Area de un poligono de coordenadas enteras
I = Numero de puntos enteros en su interior
B = Numero de puntos enteros sobre sus bordes

int IntegerPointsOnSegment(const point &P1, const point &P2) {
    point P = P1-P2;
    P.x = abs(P.x); P.y = abs(P.y);
```

## 4.7. Par de puntos más cercano (Sweep Line).

```
#define MAX_N 100000
#define px second
#define py first
typedef pair<long long, long long> point;

int N;
point P[MAX_N];
set<point> box;

bool compare_x(point a, point b) { return a.px<b.px; }

inline double dist(point a, point b) {
   return sqrt((a.px-b.px)*(a.px-b.px)+(a.py-b.py)*(a.py-b.py));
}

double closest_pair() {
   if(N<=1) return -1;</pre>
```

# 4.8. Par de puntos más cercano (Divide and Conquer).

```
void closest_pair(int 1, int r) {
   if(1 == r) return;
```

```
return vector<point> (H,H+k);
  if(P.x == 0) return P.y;
  if(P.y == 0) return P.x;
   return (__gcd(P.x,P.y));
Se asume que los vertices tienen coordenadas enteras. Sumar el valor de esta
funcion para todas las aristas para obtener el numero total de punto en el borde
del poligono.
   sort(P,P+N,compare_x);
  double ret = dist(P[0], P[1]);
  box.insert(P[0]);
   set<point> :: iterator it;
   for (int i = 1, left = 0; i < N; ++i) {</pre>
      while(left<i && P[i].px-P[left].px>ret) box.erase(P[left++]);
      for(it = box.lower_bound(make_pair(P[i].py-ret,P[i].px-ret));
         it!=box.end() && P[i].py+ret>=(*it).py;++it)
            ret = min(ret, dist(P[i],*it));
      box.insert(P[i]);
   return ret;
  int mi = (1 + r) >> 1;
```

int X = p[mi].x;

```
closest_pair(1,mi);
closest_pair(mi + 1,r);
int m = 0;

for(int i = 1;i <= r;++i)
    if(abs(X - p[i].x) <= best)
        aux[m++] = point(p[i].y,p[i].x,p[i].id);

sort(aux,aux + m);

for(int i = 0;i < m;++i){
    int e = i + 1;</pre>
```

# 4.9. Unión de rectángulos (Área).

```
#define MAX_N 10000
struct event {
   int ind;
  bool type;
   event(){};
   event(int ind, int type) : ind(ind), type(type) {};
};
struct point{
   int x,y;
};
int N;
point rects[MAX_N][2];
// rects[i][0] : esquina inferior izquierda
// rects[i][1] : esquina superior derecha
event events_v[2*MAX_N], events_h[2*MAX_N];
bool in_set[MAX_N];
bool compare_x(event a, event b) {
   return rects[a.ind][a.type].x<rects[b.ind][b.type].x;</pre>
bool compare_y (event a, event b) {
   return rects[a.ind][a.type].y<rects[b.ind][b.type].y;</pre>
```

```
while(e < m && aux[e].x - aux[i].x <= best + EPS) {
    double d = dist(aux[i],aux[e]);

    if(d < best) {
        best = d;
        id1 = aux[i].id;
        id2 = aux[e].id;
    }

    ++e;
}</pre>
```

```
long long union_area() {
   int e = 0:
   for(int i = 0;i<N;++i){</pre>
      events_v[e] = event(i,0);
      events_h[e] = event(i,0);
      events_v[e] = event(i,1);
      events_h[e] = event(i,1);
      ++e;
   sort (events_v, events_v+e, compare_x);
   sort(events_h, events_h+e, compare_y);
   memset(in_set, false, sizeof(in_set));
   in set[events v[0].ind] = true;
   long long area = 0;
   int prev_ind = events_v[0].ind, cur_ind;
   int prev_type = events_v[0].type, cur_type;
   for (int i = 1; i < e; ++i) {</pre>
      cur_ind = events_v[i].ind; cur_type = events_v[i].type;
      int cont = 0, dx = rects[cur_ind][cur_type].x-rects[prev_ind][prev_type].x;
      int begin_y;
```

#### 4.10. Geometría 3D.

```
struct XYZ{
    double x,y,z;

XYZ(){}

XYZ(double _x, double _y, double _z) :
    x(_x), y(_y), z(_z){}

void normalize(){
    double r = sqrt(x * x + y * y + z * z);
    x /= r; y /= r; z /= r;
}

XYZ cross(XYZ p){
    return XYZ(y * p.z - z * p.y,z * p.x - x * p.z,x * p.y - y * p.x);
}

double dot(XYZ p){
    return x * p.x + y * p.y + z * p.z;
};

// rotar p con eje de rotacion r
```

```
}

}

in_set[cur_ind] = (cur_type==0);

prev_ind = cur_ind; prev_type = cur_type;
}

return area;
}
```

```
XYZ rotate(XYZ p, XYZ r, double theta) {
  XYZ q(0,0,0);
  double costheta, sintheta;
  r.normalize();
  costheta = cos(theta);
  sintheta = sin(theta);
  q.x += (costheta + (1 - costheta) * r.x * r.x) * p.x;
  q.x += ((1 - costheta) * r.x * r.y - r.z * sintheta) * p.y;
  q.x += ((1 - costheta) * r.x * r.z + r.y * sintheta) * p.z;
  q.y += ((1 - costheta) * r.x * r.y + r.z * sintheta) * p.x;
  q.y += (costheta + (1 - costheta) * r.y * r.y) * p.y;
  q.y += ((1 - costheta) * r.y * r.z - r.x * sintheta) * p.z;
  q.z += ((1 - costheta) * r.x * r.z - r.y * sintheta) * p.x;
  q.z += ((1 - costheta) * r.y * r.z + r.x * sintheta) * p.y;
  q.z += (costheta + (1 - costheta) * r.z * r.z) * p.z;
  return q;
```

## 5. Matemática

#### 5.1. GCD extendido.

```
// a*x + b*y = gcd(a,b)
int extGcd(int a, int b, int &x, int &y) {
   if(b == 0) {
      x = 1;
      y = 0;
      return a;
   }
   int g = extGcd(b,a % b,y,x);
   y -= a / b * x;
```

#### 5.2. Teorema chino del resto.

```
// rem y mod tienen el mismo numero de elementos
long long chinese_remainder(vector<int> rem, vector<int> mod) {
   long long ans = rem[0],m = mod[0];
   int n = rem.size();

   for(int i=1;i<n;++i) {
      int a = modular_inverse(m,mod[i]);
   }
}</pre>
```

#### 5.3. Número combinatorio.

```
long long comb(int n, int m) {
   if(m>n-m) m = n-m;

  long long C = 1;
    //c^{n}_{ii} -> c^{n}_{ii+1}
   for(int i=0;i<m;++i) C = C*(n-i)/(1+i);
   return C;
}

Cuando n y m son grandes y se pide comb(n,m)%MOD, donde MOD es un numero primo, se puede usar el Teorema de Lucas.

#define MOD 3571

int C[MOD][MOD];

void FillLucasTable() {
   memset(C,0,sizeof(C));</pre>
```

```
return q;
// ASSUME: gcd(a, m) == 1
int modInv(int a, int m) {
   int x,y;
   extGcd(a, m, x, y);
   return (x % m + m) % m;
      int b = modular_inverse(mod[i],m);
      ans = (ans*b*mod[i]+rem[i]*a*m)%(m*mod[i]);
      m \star = mod[i];
   return ans;
   for(int i=0;i<MOD;++i) C[i][0] = 1;</pre>
   for(int i=1;i<MOD;++i) C[i][i] = 1;</pre>
   for (int i=2; i < MOD; ++i)</pre>
      for(int j=1; j<i; ++j)</pre>
         C[i][j] = (C[i-1][j]+C[i-1][j-1]) %MOD;
int comb(int n, int k){
   long long ans = 1;
   while(n!=0){
      int ni = n%MOD,ki = k%MOD;
      n /= MOD; k /= MOD;
      ans = (ans*C[ni][ki])%MOD;
   return (int)ans;
```

#### 5.4. Test de Miller-Rabin.

```
typedef unsigned long long ULL;
ULL mulmod(ULL a, ULL b, ULL c) {
   ULL x = 0, y = a % c;
   while(b > 0){
      if(b & 1) x += y;
      v <<= 1;
      if(x >= c) x -= c;
      if(y >= c) y -= c;
      b >>= 1;
   return x;
ULL pow(ULL a, ULL b, ULL c) {
  ULL x = 1, y = a;
   while(b > 0){
      if(b \& 1) x = mulmod(x,y,c);
      y = mulmod(y, y, c);
      b >>= 1;
   return x;
```

#### 5.5. Polinomios.

```
vector<int> add(vector<int> &a, vector<int> &b) {
   int n = a.size(), m = b.size(), sz = max(n,m);
   vector<int> c(sz,0);

for(int i = 0;i<n;++i) c[i] += a[i];
   for(int i = 0;i<m;++i) c[i] += b[i];

   // mejor no quitar si son reales
while(sz>1 && c[sz-1]==0) {
    c.pop_back();
    --sz;
}
```

```
bool miller_rabin(ULL p, int it) {
   if(p < 2) return false;</pre>
  if(p == 2) return true;
  if((p & 1) == 0) return false;
  ULL s = p - 1;
   while(s % 2 == 0) s >>= 1;
   while (it--) {
      ULL a = rand() % (p-1) + 1, temp = s;
      ULL mod = pow(a,temp,p);
      if (mod == -1 | | mod == 1) continue;
      while(temp != p-1 && mod != p-1) {
         mod = mulmod(mod, mod, p);
         temp <<= 1;
      if (mod != p-1) return false;
   return true;
   return c;
vector<int> multiply(vector<int> &a, vector<int> &b){
   int n = a.size(), m = b.size(), sz = n+m-1;
   vector<int> c(sz,0);
   for (int i = 0; i < n; ++i)</pre>
      for(int j = 0; j<m; ++j)
         c[i+j] += a[i]*b[j];
      // mejor no quitar si son reales
```

**while**(sz>1 && c[sz-1]==0){

```
c.pop_back();
    --sz;
}

return c;
}
bool is_root(vector<int> &P, int r){
    int n = P.size();
```

#### 5.6. Fast Fourier Transform.

```
struct Complex{
   double x,y;
   Complex(){}
   Complex(double _x, double _y):
      x(_x), y(_y){}
   void operator += (Complex &c) {
      x += c.x; y += c.y;
   Complex operator -= (Complex &c) {
      x -= c.x; y -= c.y;
   Complex operator * (Complex &c) {
      return Complex(x * c.x - y * c.y, x * c.y + y * c.x);
};
#define MAXN 262144
Complex A2[MAXN];
void fft(int n, Complex A[], int s){
   int p = __builtin_ctz(n);
   memcpy(A2, A, sizeof(Complex) * n);
   for (int i = 0; i < n; ++i) {</pre>
      int rev = 0;
      for(int j = 0; j < p; ++j) {</pre>
         rev <<= 1;
```

```
long long y = 0;
   for (int i = 0; i < n; ++i) {</pre>
      if(abs(y-P[i])%r!=0) return false;
      y = (y-P[i])/r;
   return y==0;
         rev |= ((i >> j) & 1);
      A[i] = A2[rev];
   Complex w,wn;
   int M = 2, K = 1;
   for(int i = 1;i <= p;++i,M <<= 1,K <<= 1) {</pre>
      wn = Complex(cos(s * 2 * M_PI / M), sin(s * 2 * M_PI / M);
      for(int j = 0; j < n; j += M) {
         w = Complex(1,0);
         for(int 1 = j;1 < K+j;++1){</pre>
            Complex t = w * A[1 + K], u = A[1];
            A[1] += t;
            u -= t;
            A[1 + K] = u;
            w = w * wn;
   if(s == -1)
      for (int i = 0; i < n; ++i)
         A[i].x /= n, A[i].y /= n;;
Complex R[MAXN];
int nR;
```

```
void fft_mult(int nP, Complex P[], int nQ, Complex Q[]) {
    nR = nP + nQ;
    while(_builtin_popcount(nR) > 1)    nR += nR & -nR;

    for(int i = nP;i < nR;++i) P[i] = Complex(0,0);
    for(int i = nQ;i < nR;++i) Q[i] = Complex(0,0);</pre>
```

```
fft (nR,P,1);
fft (nR,Q,1);

for (int i = 0;i < nR;i++) R[i] = P[i] * Q[i];

fft (nR,R,-1);</pre>
```

## 5.7. Stern Brocott.

```
const int MAX_DEN = 3000;
vector<int> Fnum,Fden;

void build(int lnum = 0, int lden = 1, int rnum = 1, int rden = 1) {
   int a = lnum+rnum,b = lden+rden;
   if(b>MAX_DEN) return;
```

```
build(lnum,lden,a,b);
Fnum.push_back(a);
Fden.push_back(b);
build(a,b,rnum,rden);
}
```

#### 6. Estructuras de datos

#### 6.1. Lowest Common Ancestor.

```
#define MAX_N 100000
#define LOG2_MAXN 16

// NOTA : memset (parent, -1, sizeof (parent));
int N, parent [MAX_N], L[MAX_N];
int P[MAX_N] [LOG2_MAXN + 1];

int get_level (int u) {
    if (L[u]!=-1) return L[u];
    else if (parent[u]==-1) return 0;
    return 1+get_level (parent[u]);
}

void init() {
    memset (L, -1, sizeof (L));
    for (int i = 0; i < N; ++i) L[i] = get_level (i);
    memset (P, -1, sizeof (P));

for (int i = 0; i < N; ++i) P[i][0] = parent[i];</pre>
```

```
for(int i = log;i>=0;--i) {
   if(P[p][i]!=-1 && P[p][i]!=P[q][i]) {
      p = P[p][i];
      q = P[q][i];
}
```

## 6.2. Heavy-Light Descomposition.

```
struct HeavyLight{
   static const int MAXN = 100005;
   int N;
   vector<int> E[MAXN];
   int nodedad[MAXN];
   int treesize[MAXN];
   int pos,cntchain;
   int chainleader[MAXN];
   int homechain[MAXN];
   int homepos[MAXN];
   void init(int n){
     N = n;
     for(int i = 0;i < n;++i) E[i].clear();</pre>
     pos = cntchain = 0;
   void add_edge(int u, int v) {
      E[u].push_back(v);
      E[v].push_back(u);
   void explore(int x = 0, int dad = -1){
      nodedad[x] = dad;
     treesize[x] = 1;
      int sz = E[x].size();
```

## 6.3. **Treap.**

```
typedef long long ptype;
```

```
return parent[p];
      for(int i = 0;i < sz;++i){</pre>
         if(E[x][i] != dad) {
            explore(E[x][i], x);
            treesize[x] += treesize[ E[x][i] ];
  void heavy_light(int x = 0, int dad = -1, int k = -1, int p = 0) {
     if(p == 0){
         k = cntchain++;
         chainleader[k] = x;
      homechain[x] = k;
      homepos[x] = pos++;
      int mx = -1, sz = E[x].size();
      for (int i = 0; i < sz; ++i)
         if(E[x][i]] = dad \&\& (mx == -1 || treesize[E[x][i]] > treesize[E[x][mx]]))
            mx = i;
      if(mx != -1) heavy_light(E[x][mx], x, k, p + 1);
      for(int i = 0; i < sz; ++i)</pre>
         if(E[x][i] != dad && i != mx)
            heavy_light(E[x][i], x, -1, 0);
};
```

```
ptype seed = 47;
ptype my_rand(){
   seed = (seed * 279470273) % 4294967291LL;
   return seed;
typedef struct node * pnode;
struct node {
 int x, y, cnt;
 pnode L, R;
 node() {}
 node(int x, int y): x(x), y(y), cnt(1), L(NULL), R(NULL) {}
pnode T;
inline int cnt(pnode &it) {
 return it ? it->cnt : 0;
inline void upd_cnt(pnode &it) {
 if (it) {
  it->cnt = cnt(it->L) + cnt(it->R) + 1;
// Split Treap
void split(pnode t, int x, pnode &L, pnode &R) {
 if (!t) L = R = NULL;
 else
  if (x < t->x)
    split (t->L, x, L, t->L), R = t;
    split (t->R, x, t->R, R), L = t;
   upd cnt(t);
// Split Implicit Treap
void split(pnode t, pnode &L, pnode &R, int key) {
 if (!t) L = R = NULL;
 else
   int cntL = cnt(t->L);
  if (key <= cntL)</pre>
```

```
split (t->L, L, t->L, key), R = t;
    split (t->R, t->R, R, key - cntL - 1), L = t;
   upd_cnt(t);
// For Treap & Implicit Treap
void merge (pnode &t, pnode L, pnode R) {
 if (!L) t = R;
 else if(!R) t = L;
 else if (L->V > R->V)
  merge (L->R, L->R, R), t = L;
  merge (R->L, L, R->L), t = R;
 upd_cnt(t);
// Combines 2 treaps
pnode unite (pnode 1, pnode r) {
 if (!1 || !r) return 1? 1: r;
 if (1->y > r->y) swap (1, r);
 pnode lt, rt;
 split (r, 1->x, lt, rt);
 1->L = unite(1->L, lt);
 1->R = unite(1->R, rt);
 return 1;
// Find in Treap
bool find (pnode &t, int x) {
 if(!t) return 0;
 else if (t->x == x) return 1;
 else return find (x < t->x ? t->L: t->R, x);
// Erase from Treap
void erase (pnode &t, int x) {
 if (t-> x == x)
  merge (t, t->L, t->R);
  erase (x < t->x ? t->L: t->R, x);
// Insert into Treap
void insert(pnode &t, pnode it) {
```

```
if (!t) t = it;
else if (it->y > t->y)
    split (t, it->x, it->L, it->R), t = it;
else insert (it->x < t->x ? t->L: t->R, it);
}

// Insert into Treap and return the # of greater elements
int insert(pnode &t, pnode it) {
    int ret = 0;
    if (!t) t = it;
else if (it->y > t->y)
    split (t, it->x, it->L, it->R), t = it, ret = cnt(t->R);
else if (it->x < t->x)
```

```
ret = 1 + cnt(t->R) + insert(t->L, it);
else
  ret = insert(t->R, it);
upd_cnt(t);
return ret;
}

// Safely insert into Treap
void insert(int x)
{
  if(!find(T, x))
    insert(T, new node(x, rand()));
}
```

## 7. Matrices

#### 8. Mathematical facts

8.1. **Números de Catalan.** están definidos por la recurrencia:

$$C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i}$$

$$C_n = \frac{1}{n+1} {2n \choose n} = {2n \choose n} - {2n \choose n+1}$$

8.2. Números de Stirling de primera clase. son el número de permutaciones de n elementos con exactamente k ciclos disjuntos.

8.3. Números de Stirling de segunda clase. son el número de particionar un conjunto de n elementos en k subconjuntos no vacíos.

$${n \brace k} = k {n-1 \brace k} + {n-1 \brace k-1}$$

Además:

$${n \brace k} = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^n$$

8.4. **Números de Bell.** cuentan el número de formas de dividir n elementos en subconjuntos.

$$\mathcal{B}_{n+1} = \sum_{k=0}^{n} \binom{n}{k} \mathcal{B}_k$$

X	0	1	2	3	4	5	6	7	8	9	10
$\mathcal{B}_x$	1	1	2	5	15	52	203	877	4.140	21.147	115.975

8.5. **Derangement.** permutación que no deja ningún elemento en su lugar original

$$!n = (n-1)(!(n-1)+!(n-2)); !1 = 0, !2 = 1$$

$$!n = n! \sum_{k=0}^{n} \frac{(-1)^k}{k!}$$

8.6. Números armónicos.

$$H_n = \sum_{k=1}^n \frac{1}{k}$$

$$\frac{1}{2n+1} < H_n - \ln n - \gamma < \frac{1}{2n}$$

 $\gamma = 0.577215664901532860606512090082402431042159335...$ 

8.7. Número de Fibonacci.  $f_0 = 0, f_1 = 1$ :

$$f_n = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2}\right)^n - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2}\right)^n$$

$$f_{n+1}^2 + f_n^2 = f_{2n+1}, f_{n+2}^2 - f_n^2 = f_{2n+2}$$

$$f_n = \sum_{j=0}^{\lfloor \frac{n}{2} \rfloor} \binom{n-j}{j}$$

8.8. Sumas de combinatorios.

$$\sum_{i=n}^{m} \binom{i}{n} = \binom{m+1}{n+1}$$

$$\sum_{i=0}^{k} \binom{n}{i} \binom{m}{k-i} = \binom{n+m}{k}$$

8.9. Funciones generatrices. Una lista de funciones generatrices para secuencias útiles:

$(1,1,1,1,1,1,\ldots)$	$\frac{1}{1-z}$
$(1,-1,1,-1,1,-1,\ldots)$	$\frac{1}{1+z}$
$(1,0,1,0,1,0,\ldots)$	$\frac{1}{1-z^2}$
$(1,0,\ldots,0,1,0,1,0,\ldots,0,1,0,\ldots)$	$\frac{1}{1-z^2}$
$(1,2,3,4,5,6,\ldots)$	$\frac{1}{(1-z)^2}$
$(1, \binom{m+1}{m}, \binom{m+2}{m}, \binom{m+3}{m}, \dots)$	$\frac{1}{(1-z)^{m+1}}$
$(1,c,\binom{c+1}{2},\binom{c+2}{3},\ldots)$	$\frac{1}{(1-z)^c}$
$(1,c,c^2,c^3,\ldots)$	$\frac{1}{1-cz}$
$(0,1,\frac{1}{2},\frac{1}{3},\frac{1}{4},\ldots)$	$\ln \frac{1}{1-z}$

Truco de manipulación:

$$\frac{1}{1-z}G(z) = \sum_{n} \sum_{k < n} g_k z^n$$

8.10. The twelvefold way. ¿Cuántas funciones  $f: N \to X$  hay?

N	X	Any $f$	Injective	Surjective
dist.	dist.	$x^n$	$(x)_n$	$x!\binom{n}{x}$
indist.	dist.	$\binom{x+n-1}{n}$	$\binom{x}{n}$	$\binom{n-1}{n-x}$
dist.	indist.	$\binom{n}{1} + \ldots + \binom{n}{x}$	$[n \le x]$	$\binom{n}{k}$
indist.	indist.	$p_1(n) + \dots p_x(n)$	$[n \le x]$	$p_x(n)$

Where  $\binom{a}{b} = \frac{1}{b!}(a)_b$  and  $p_x(n)$  is the number of ways to partition the integer n using x summands.

8.11. **Teorema de Euler.** si un grafo conexo, plano es dibujado sobre un plano sin intersección de aristas, y siendo v el número de vértices, e el de aristas y f la cantidad de caras (regiones conectadas por aristas, incluyendo la región externa e infinita), entonces

$$v - e + f = 2$$

8.12. **Burnside's Lemma.** Si X es un conjunto finito y G es un grupo de permutaciones que actúa sobre X, sean  $S_x = \{g \in G : g * x = x\}$  y  $Fix(g) = \{x \in X : g * x = x\}$ . Entonces el número de órbitas está

dado por:

$$N = \frac{1}{|G|} \sum_{x \in X} |S_x| = \frac{1}{|G|} \sum_{g \in G} |Fix(g)|$$

8.13. Ángulo entre dos vectores. Sea  $\alpha$  el ángulo entre  $\vec{a}$  y  $\vec{b}$ :

$$\cos\alpha = \frac{\vec{a} \cdot \vec{b}}{\|\vec{a}\| \|\vec{b}\|}$$

8.14. Proyección de un vector. Proyección de  $\vec{a}$  sobre  $\vec{b}$ :

$$\operatorname{proy}_{\vec{b}} \vec{a} = (\frac{\vec{a} \cdot \vec{b}}{\vec{b} \cdot \vec{b}}) \vec{b}$$

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