



Matrix Code Equivalence

combinatorial

Attacks using isometry-invariant substructures

Example: find low-rank codewords in both codes and match them up, construct isometry from this.

or, find peculiar subcodes on both sides, match them up, and construct the isometry between the subcodes

-
- Graph-based algorithm
 - Leon's like algorithm

$$\tilde{\mathcal{O}}(q^{\min(n,m,k)})$$

algebraic

Attacks reducing MCE to solving a system of polynomial equations

Example: write down both generator matrices and add rows in variables of \mathbf{A} and \mathbf{B}

or, use the formulation as tensor isomorphism to get a bilinear system, apply Gröbner techniques

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- direct modelling
 - minor's modelling
 - *improved* modelling

$$\mathcal{O}\left(n^{\omega \frac{n}{4}}\right)$$



Matrix Code Equivalence

equations

$$\mathcal{C}(Ax, By, z) = \mathcal{D}(x, y, T^{-1}z)$$

$$\mathcal{C}(Ax, y, Tz) = \mathcal{D}(x, B^{-1}y, z)$$

$$\mathcal{C}(x, By, Tz) = \mathcal{D}(A^{-1}x, y, z)$$

bilinear system of

- $k(nm - k)$ equations
- $n^2 + m^2$ variables

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