

PART 1

The Tate Pairing

Definition 3. A *pairing* on an elliptic curve E is a bilinear map $e : A \times B \rightarrow \mathbb{F}_q^*$, where A and B are subgroups of E . We say e is *non-degenerate* when for every $a \in A$, there is at least one $b \in B$ such that $e(a, b) \neq 1$, and vice versa. We say e is *alternating* when $A = B$ and for every $a \in A$ we have $e(a, a) = 1$.

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The Weil Pairing

The iconic pairing, that always helps us when we are dealing with m -torsion $E[m]$,

$$e_m : E[m] \times E[m] \rightarrow \mu_m,$$

where $\mu_m = \{ \zeta \in \overline{\mathbb{F}_q}^* \mid \zeta^m = 1 \}$. The m -Weil pairing is *bilinear*, *non-degenerate*, and *alternating*.

Destructive use: the MOV attack on elliptic curves, reducing elliptic-curve DLOG to \mathbb{F}_q^* -DLOG.

Constructive use* in isogeny-based cryptography: decomposing $K \in E[m]$ into $K = [a]P + [b]Q$ for a given basis P, Q for $E[m]$.

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The Tate-Lichtenbaum Pairing

My favourite pairing, deep connection to division by $[m]$, the crucial insight for this work.

$$t_m : E[m](\mathbb{F}_q) \times E(\mathbb{F}_q)/[m]E(\mathbb{F}_q) \rightarrow \mu_m.$$

The (reduced) m -Tate pairing is *bilinear*. When $\mu_m \subseteq \mathbb{F}_q^*$, it is also *non-degenerate*.

Destructive use: the Frey-Rück attack, using the idea from the MOV attack with the Tate pairing.

Constructive use* in isogeny-based cryptography: finding bases for $E[m]$, computing Sylow- ℓ torsion, navigating isogeny volcanoes, and so on... See [1].

* Many other applications are out of scope for this talk, such as their uses in identity-based cryptography and more generally the whole field of pairing-based cryptography.

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Corollary. Let 2^n divide $p + 1$ for some prime p , and take a supersingular Montgomery curve over \mathbb{F}_{p^2} given by

$$E_A : y^2 = x^3 + Ax^2 + x, \quad \text{with } A \in \mathbb{F}_{p^2}.$$

If a point $P = (x_P, y_P) \in E(\mathbb{F}_{p^2})$ has x_P non-square, then P has order divisible by 2^n