## $\phi: E \to E'$

**Definition 2.** The *kernel*  $\ker \varphi$  are the points  $P \in E$  that are mapped to infinity  $\mathcal{O}' \in E'$ . For all isogenies we will care about, the degree  $\deg \varphi$  is equal to the size of  $\ker \varphi$ .

$$\ker \varphi := \{ P \in E \mid \varphi(P) = \mathcal{O}' \}.$$

The kernel is *cyclic* if every point in  $\ker \varphi$  is a multiple of some *generator*  $K \in \ker \varphi$ . Then, for  $n = \deg \varphi$ ,

$$\ker \varphi = \{ K, [2]K, [3]K, ..., [n]K \}$$

**Theorem 1.** If the kernel  $\ker \varphi$  is cyclic, with generator  $K=(x_K,y_K)\in E(\mathbb{F}_q)$ , then it's easy to compute a smooth degree  $\varphi$ : We can compute E' and  $\varphi(P)$  for  $P\in E(\mathbb{F}_q)$ .

## IN PRACTICE

- computing isogenies often **bottleneck** in efficiency of isogeny-based cryptography
- isogenies of degree  $2^n$  are **most efficient**, for largest n with rational points of order  $2^n$
- can describe such an isogeny by a single point  $K \in E[2^n](\mathbb{F}_q)$
- we **choose** our field  $\mathbb{F}_q$  and curve E such that we have many of such points/isogenies



The Study of Vermeer (1964)