

CSC3001 Discrete Mathematics

Final Examination

July 25, 2022: 2:00pm - 4:30pm

Name: _____ Student ID: _____

Answer ALL questions in the Answer Book.

Question	Points	Score
1	16	
2	16	
3	16	
4	16	
5	16	
6	16	
7	4	
Total:	100	

1. (16 points) Let $a_1 = 1$. Suppose that, for $n \in \mathbb{N}^+$, we have the recurrence $a_{n+1} = 3a_n \bmod 11$. Find a closed-form expression of the sequence a_n .

Solution:

$a_1 = 1, a_2 = 3, a_3 = 9, a_4 = 5, a_5 = 4, a_6 = 1$. As $a_6 = a_1$, the sequence is periodic thereafter. The closed-form expression is thus

$$a_n = \begin{cases} 1 & n = 1 \pmod{5} \\ 3 & n = 2 \pmod{5} \\ 9 & n = 3 \pmod{5} \\ 5 & n = 4 \pmod{5} \\ 4 & n = 0 \pmod{5} \end{cases}.$$

2. (16 points) Let n be a positive integer. Let b_n be the number of subsets S of $\{1, 2, \dots, n\}$ such that any elements $a, b \in S$ differ by at least 3, that is, $|a - b| \geq 3$. Find a recurrence relation for b_n and compute b_{10} .

Solution:

If n belongs to subset S , then $n - 1$ and $n - 2$ do not belong. The rest of the subset can be chosen in b_{n-3} different ways. If n doesn't belong to subset S , then there are b_{n-1} ways to choose such a subset. Hence

$$b_n = b_{n-1} + b_{n-3}.$$

We have $b_1 = 2, b_2 = 3, b_3 = 4$. Using recursion we get

$$b_4 = 6, b_5 = 9, b_6 = 13, b_7 = 19,$$

$$b_8 = 28, b_9 = 41, b_{10} = 60.$$

3. (16 points) For a list A denote by A^{rev} the reversed list. We recursively define arrays S_n as follows. Let $S_1 = [0, 1]$ be an array of binary strings of length 1. Define

$$S_{n+1} = (S_n \text{ with } 0 \text{ added before each string}) \cup (S_n^{rev} \text{ with } 1 \text{ added before each string}).$$

For example,

$$S_2 = [00, 01, 11, 10],$$

$$S_3 = [000, 001, 011, 010, 110, 111, 101, 100].$$

- (a) Prove that S_n contains all binary strings of length n .

- (b) Prove that each string from S_n differs from its neighbors in at most one bit, and that the first and the last string in S_n also differ in exactly one bit.

Solution:

Part (a) We use induction. Base: S_1 contains all strings of length 1. Step: assume S_n contains all binary strings of length n . Then adding 1 before them will give all binary strings of length $n + 1$ that start with 1. Adding 0 will give all binary strings of length $n + 1$ that start with 0, but their union is all binary strings of length $n + 1$.

Part (b) We use induction. Base: the claim is true for S_1 . Step: assume the claim is true for S_n . Then neighbors within S_n with 0 added before each string also differ by at most one bit. Reversing an array doesn't violate the property hence reversing and adding 1 before each string also keeps the property. We are only left to check two cases: that strings on positions 2^n and $2^n + 1$ differ by exactly one bit and that strings on positions 1 and 2^{n+1} differ by exactly one bit. But we can easily see that string on position 1 is 00...0, string on position 2^{n+1} is 100...0, string on position 2^n is 0100...0, and string on position $2^n + 1$ is 1100...0. Hence we have proven the step.

4. (16 points) Let n be a positive integer. Prove the identity

$$1 \cdot \binom{n}{1} + 2 \cdot \binom{n}{2} + 3 \cdot \binom{n}{3} + \cdots + n \cdot \binom{n}{n} = n \cdot 2^{n-1}.$$

Solution:

Consider all binary strings of length n . Let us count the number of ones in them in two different ways. On the one hand, each of n bits is equal to 1 exactly for 2^{n-1} strings, hence number of ones is $n2^{n-1}$. On the other hand, there are $\binom{n}{k}$ strings with exactly k ones, hence the number of ones is given by

$$1 \cdot \binom{n}{1} + 2 \cdot \binom{n}{2} + 3 \cdot \binom{n}{3} + \cdots + n \cdot \binom{n}{n}.$$

Alternative solution: By the binomial expansion

$$\sum_{k=0}^n \binom{n}{k} x^k = (1+x)^n.$$

Hence taking derivative gives

$$\sum_{k=0}^n \binom{n}{k} k x^{k-1} = n(1+x)^{n-1}.$$

Plugging $x = 1$ gives

$$\sum_{k=0}^n \binom{n}{k} k = n2^{n-1}.$$

5. (16 points) Assume that there are at least 2 vertices in a graph in this question. Prove that the following two definitions of a tree are equivalent:
1. a graph $G = (V, E)$ such that for every two different vertices $u, v \in V$ there is a unique path between u and v ;
 2. a connected graph $G = (V, E)$ with $|V| = |E| + 1$.

Solution:

Let us show that both definitions are equivalent to a connected graph with no cycles.

If G is a connected graph with no cycles then every pair of vertices u, v can be connected by a path P . This path is unique because if there are two different paths P, R , then their union contains a cycle.

If every two vertices in G can be connected by a unique path then G is connected. G also has no cycles C because any two vertices $u, v \in C$ are connected by at least 2 paths.

Let us show by induction on $|V|$ that for connected graphs having no cycles is equivalent to $|V| = |E| + 1$.

Indeed, induction base: $|V| = 2$ - it follows immediately by $|E| = 1$.

Induction step: assume we are given a connected graph with $|V| > 1$. If G has no cycles, then we can find a leaf u (go along any path until we get a vertex of degree 1). Since $G \setminus \{u\}$ also has no cycles we can apply induction base to conclude that $|V| - 1 = (|E| - 1) + 1$, hence $|V| = |E| + 1$.

Now if G has $|V| = |E| + 1$, then from handshaking lemma $\sum_v \deg(v) = 2|E| = 2|V| - 1$ and the fact that $\deg(v) > 0$ (graph is connected) we find that for some vertex v we have $\deg(v) = 1$. It follows that v is a leaf. We conclude by induction that $G \setminus \{v\}$ has no cycles and adding a leaf doesn't add cycles hence G has no cycles.

6. (16 points) Let $m \geq 2$ be a positive integer. Let A_m be the set of all bit strings s of arbitrary length, such that there is a length- m substring that appears in at least 2 different places of s .

For example, $10111011 \in A_4$ as 1011 appears 2 times in the string, $01010 \in A_3$ as 010 appears 2 times in the string, and $0010110 \notin A_3$, as every length-3 substring is distinct.

Let s be a bit string of length at least $m - 1$ that is not in A_m . However if we add any bit to the end of s , it will become an element of A_m . Show that the $(m - 1)$ -th bit of s and the last bit of s are the same.

Solution:

Let the pattern of the last $m - 1$ bits of s be z . Then both $z0$ and $z1$ are in s . If none of $z0$ and $z1$ incidents at the beginning of s , then $0z$ and $1z$ in total appear at least 3 times in s , forcing s to be in A_m by the pigeonhole principle. Therefore s must begin with $z0$ or $z1$, which indicates that s starts and ends with z , which then concludes that the $(m - 1)$ -th bit of s and the last bit of s are both the last bit of z .

7. (4 points) For $p = (x, y) \in \mathbb{R}^2$, define $\|p\|_2 = \sqrt{x^2 + y^2}$.
- (a) (2 points) On a Cartesian coordinate $p = (x, y) \in \mathbb{R}^2$, a unit circle is represented by $\|p\|_2 = 1$. From $(-1, 0)$ on the circle, draw a line $y = k(x + 1)$. This line intersects with the unit circle on $(-1, 0)$ and (x_1, y_1) . Show that if $k \in \mathbb{Q}$ then $x_1, y_1 \in \mathbb{Q}$.
- (b) (2 points) Show that there exists a subset $B \subseteq \mathbb{Q}^2$ such that for every $d \in \mathbb{Q}^+$ there exists a unique unordered pair $p_1, p_2 \in B$ such that $\|p_1 - p_2\|_2 = d$.

Solution:

Part (a) We plug $y = k(x + 1)$ into $x^2 + y^2 = 1$ and obtain $x = \frac{-k^2 \pm 1}{k^2 + 1}$. Then $x_1 = \frac{-k^2 + 1}{k^2 + 1} \in \mathbb{Q}$. Then $y = k(x + 1) \in \mathbb{Q}$.

Part (b) As \mathbb{Q} is countable, we write \mathbb{Q} into $\{d_1, d_2, \dots\}$. Let $P(m)$ be the statement that there exists a subset $B_m \subseteq \mathbb{Q}^2$ such that $|B_m| \leq m + 1$ and for every $d \in \{d_1, d_2, \dots, d_m\}$ there exists a unique unordered pair $p_1, p_2 \in B_m$ with $\|p_1 - p_2\|_2 = d$ and for every $d \in \mathbb{Q} - \{d_1, d_2, \dots, d_m\}$ there is at most one unordered pair $p_1, p_2 \in B_m$ with $\|p_1 - p_2\|_2 = d$. We prove the statement using induction on m .

The base case $m = 1$ can be proved by letting $B_1 = \{(0, 0), (0, d_1)\}$.

When $P(m)$ is true, we specify a subset B_m and investigate d_{m+1} . $P(m)$ dictates that at most one unordered pair with $\|p_1 - p_2\|_2 = d_{m+1}$ exists in B_m . If there is exactly one unordered pair, then letting $B_{m+1} = B_m$ will prove $P(m + 1)$.

Otherwise there is no unordered pair with $\|p_1 - p_2\|_2 = d_{m+1}$. We pick an arbitrary point $p_m \in B_m$ and draw a circle C_m centered at p_m with radius d_{m+1} . From (a) we know that there are infinitely many points p with rational x and y coordinates in the unit circle, and thus infinitely many points $d_{m+1}p + p_m$ with rational x and y coordinates on the circle C_m . Meanwhile, as $|B_m| \leq m + 1$, there are at most $m(m + 1)$ distinct d values such that $\|p_1 - p_2\|_2$ could attain with $p_1, p_2 \in B_m$. As two circles intersect on at most 2 points, there are at most $2m \cdot m(m + 1)$ points

on C_m , adding which to B_m will create some d such that two distinct unordered pairs have the same $\|p_1 - p_2\|_2 = d$. Therefore the number of points we could add to B_m to create B_{m+1} is infinitely many.

The statement follows.

Remark. A real variant of the statement, for $B \subseteq \mathbb{R}^2$ and $d \in \mathbb{R}^+$, can be proved by similar arguments. This is an example of using mathematical induction over reals, known as the *transfinite induction* that we pointed out in the lecture.