#### Tutorial 1

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# Recap: Propositional Logic

Conjunction	^	"AND"	$p \wedge q$
Disjunction	V	"OR"	$p \lor q$
Negation	_	"NOT"	$\neg q$
Implication	$\rightarrow$	"IF, Then"	$p \rightarrow q$
Biconditional	$\leftrightarrow$	"IF AND ONLY IF"	$p \leftrightarrow q$

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Then we can have  $(p \lor q) \to \neg r$ .

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Then we can have  $r \to (p \lor q)$ .

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$$p \land \neg q \land r$$

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$$p \wedge \neg q \wedge r$$

You get an A on the final exam and complete every exercise in the textbook, which helps you obtain an A for the class.

$$(p \land q) \to r$$



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 $q : \mathsf{B}$  is a knight.

A says "B is a knight."

B says "The two of us are of opposite types."

 $p:\mathsf{A}$  is a knight.

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Four combinations:

$$p \wedge q, \ p \wedge \neg q, \ \neg p \wedge q, \ \neg p \wedge \neg q$$

p: A is a knight, q: B is a knight.

Possibilities		A says B is a knight		B says they are opposite types	
p	q	p	q	p	q
Т	Т				
Т	F				
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# Recap: Logical Equivalences

#### **Commutative Laws**

# $p \wedge q \equiv q \wedge p$

$$p\vee q\equiv q\vee p$$

#### **Associative Laws**

$$(p \land q) \land r \equiv q \land (p \land r)$$

$$(p \vee q) \vee r \equiv q \vee (p \vee r)$$

### De Morgan's Laws

$$\neg(p \land q) \equiv \neg p \lor \neg q$$

$$\neg (p \lor q) \equiv \neg p \land \neg q$$

#### **Distributive Laws**

$$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

Prove  $(p \to q) \lor (p \to r) \equiv p \to (q \lor r)$ .

Prove 
$$(p \to q) \lor (p \to r) \equiv p \to (q \lor r)$$
.

$$(p \to q) \lor (p \to r)$$

$$\equiv (\neg p \lor q) \lor (\neg p \lor r)$$

$$\equiv \neg p \lor q \lor \neg p \lor r$$

$$\equiv \neg p \lor (q \lor r)$$

$$\equiv p \to (q \lor r)$$

Prove  $\neg (p \lor (\neg p \land q)) \equiv \neg p \land \neg q$ .

Prove 
$$\neg (p \lor (\neg p \land q)) \equiv \neg p \land \neg q$$
.

$$\neg (p \lor (\neg p \land q))$$

$$\equiv \neg p \land \neg (\neg p \land q)$$

$$\equiv \neg p \land (p \lor \neg q)$$

$$\equiv (\neg p \land p) \lor (\neg p \land \neg q)$$

$$\equiv F \lor (\neg p \land \neg q)$$

$$\equiv \neg p \land \neg q$$

# Recap: Rules of Inference

Modus Ponens	$\begin{array}{c} p \to q \\ p \\ \therefore q \end{array}$	Modus Tollens	$ \begin{array}{c} p \to q \\ \neg q \\ \therefore \neg p \end{array} $
Addition	$\begin{array}{c} p \\ \therefore p \vee q \end{array}$	Simplification	$p \wedge q$ $\therefore q$
Conjunction	$\begin{array}{c} p \\ q \\ \therefore p \wedge q \end{array}$	Resolution	$   \begin{array}{c c}     \neg p \lor r \\     p \lor q \\     \therefore q \lor r   \end{array} $
Transitivity	$ \begin{array}{c} p \rightarrow q \\ q \rightarrow r \\ \therefore p \rightarrow r \end{array} $	Elimination	$ \begin{array}{c} p \lor q \\ \neg p \\ \therefore q \end{array} $

Prove the premises "John works hard", "If John works hard then he isn't having any fun" and "if john isn't having any fun, then he won't make any friends" imply the conclusion "John will not make any friends".

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Prove the premises "John works hard", "If John works hard then he isn't having any fun" and "if john isn't having any fun, then he won't make any friends" imply the conclusion "John will not make any friends".

p: John works hard.

q: John is having any fun.

r: John is making friends.

Assumption:  $p, p \rightarrow \neg q, \neg q \rightarrow \neg r$ .

Conclusion:  $\neg r$ .



### **Exercies**

Assumption:  $p, p \rightarrow \neg q, \neg q \rightarrow \neg r$ .

Conclusion:  $\neg r$ .

Prove:

- $\bigcirc$  p assumption
- ②  $p \rightarrow \neg q$  assumption
- $\bigcirc$   $\neg q$  modus ponens
- $\P$   $\neg q \rightarrow \neg r$  assumption
- $\bigcirc$   $\neg r$  modus ponens

Show that the argument with premises  $(p \wedge t) \to (r \vee s)$ ,  $q \to (u \wedge t)$ ,  $u \to p$ ,  $\neg s$ , q and conclusion  $q \to r$  is valid.

#### Prove:

(1) q assumption

- 8  $p \wedge t$  conjunction
- ②  $q \rightarrow (u \land t)$  assumption
- $(p \land t) \to (r \lor s) \text{ assumption}$
- 3  $u \wedge t$  modus ponens
- $\bigcirc$   $(r \lor s)$  modus ponens

4 u simplification

 $\bigcirc$  ¬s assumption

5 t simplification

 $\bigcirc$  r elimination

(6)  $u \rightarrow p$  assumption

 $Q \rightarrow r$  implication

7 p modus ponens

### Conclusion

- Translate statements into logic.
- Use truth tables to solve logic puzzles.
- Apply rules of inference and logical equivalences.

### Thank You

Thank you for your attention!